

(ii)  $t = 9 - \lambda$   
(iii)  $\lambda(i)$

4. On  
9x6,  
repres.  
 $\times 0.3$

5. A  
instru

• 20 °

• 30 °

•

•

(1)

th

(2)

as

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# Blue Book

EXAMINATION BOOK

Box No. \_\_\_\_\_

NAME Yu Cheng YC4845  
SUBJECT CS4 ECE 6913  
CLASS \_\_\_\_\_  
SECTION B  
INSTRUCTOR \_\_\_\_\_  
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for example,  
converting 6 to 1

translate 9x6 by shifting the bit string

Yu Cheng

YC4845

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problem 1:

a) CPU ~~per~~ performance equation

$$\text{Execution Time} = IC \times CPI \times T_{cycle}$$

CPU ~~per~~ clock cycles = (Instruction number \* CPI) / clock frequency

$$IC_{P_1} \times CPI_{P_1} \times T_{cycle P_1} = IC_{P_2} \times CPI_{P_2} \times T_{cycle P_2} = IC_{P_3} \times CPI_{P_3} \times T_{cycle P_3}$$

$$\frac{IN \times 1}{5 \times 10^3} = \frac{IN \times 2}{5 \times 10^3} = \frac{IN \times 3}{5 \times 10^3}$$

∴ The ratio's should be  $P_1 : P_2 : P_3 = 1 : 2 : 3$

b)

~~Performance~~: ~~CP~~:

$$\frac{IN_{P_1} \times 1}{CF_1} = \frac{IN_{P_2} \times 2}{CF_2} = \frac{IN_{P_3} \times 3}{CF_3} = \frac{IN_{P_1} \times 1}{5 \times 10^3}$$

$$\frac{1.5 \times 1}{CF_1} = \frac{1.25 \times 2}{CF_2} = \frac{1 \times 3}{CF_3}$$

$$\therefore CF_1 = 7.5 \times 10^3 \text{ Hz}$$

$$CF_2 = 6.025 \times 10^3 \text{ Hz}$$

$$CF_3 = 5 \times 10^3 \text{ Hz}$$

and add instead of an actu  
we can calculate  $9 \times 6$  by  
that result]. Show the best  
8-bit unsigned

yc 4895

Problem 2:

a) add  $x_7 \times t_0$ ~~addi x8, x10, x7~~~~addi~~slli x10, ~~20~~, 16srli x10, x10, ~~20~~, 16~~slli t1, t0, 20~~~~slli x10, ~~t0~~, 32~~

$$\frac{1}{2}x10 = \frac{1}{2}x10 + 16$$

b)

mv t0, x10

mv t1, x10

mv t2, x11

li t3, 0

sum-loop: add t3, t3, t1.

addi t1, t1, 1

④ bft t1, t2, sum-loop

sum

sfai t3 t3 2

mv x10, t3

$$g = A[C[16] + B[32]] \cdot 6, 7$$

$$(iii) A[i] = 4B[8i-81] + 4C[32i+32]$$

4. One possible performance enhancement is to do a shift and add instead of an actual multiplication. [Since  $9 \times 6$ , for example, can be written  $(2 \times 2 \times 2 + 1) \times 6$ , we can calculate  $9 \times 6$  by shifting the bit string 6 to the left three times and then adding 6 to that result]. Show the best way to calculate  $9 \times 6$  using shifts and adds/subtracts. Assume both  $A$  and  $B$  are 32-bit integers.

Yu Cheng - yc9845

Problem 3:

i) ~~add~~,  $x_6, x_7$ . ~~slli~~  $x_7, 128$   
~~add~~  $x_6, x_7, x_6$   
~~# load C[128]~~  
~~lw x7 \$12(x7)~~  
~~add~~  $x_8, x_7, x_5$   
~~add~~  $x_8, x_4, x_6, x_7$

~~lw x~~  
~~sub x8, x9, x4~~

ii) ~~lw x \$128(x6)~~  
~~lw x7 \$64(x7)~~  
~~add x24, x6, x7~~  
~~slli x34, x24, 2~~  
~~add x4, x5, x24~~  
 $li x_7, 16$   
 $li x_6, 32$   
 $add x_7, x_7, x_6$   
 $mul x_5, x_7, x_5$   
 $sub x_8, x_9, x_5$

iii) ~~add x7, x10, x10~~  $\# x_7 = 8i$   
~~addi x7, x7, -81~~  $\# x_7 = 8i - 81$   
~~slli x7, x7, 2~~  
~~add x7, x7, x6~~  $\# x_7 = \text{address of } B[8i-81]$

~~slli x7, x10, 2~~  
~~add x7, x7, x5~~  $\# x_7 = \text{address of } A[i]$   
~~lw x24 0(x24)~~

fgi 8  
Add instead of an actual multiplication. Show the best way to calculate  $9 \times 6$  by shifting left. Show the best way to calculate  $8 \times 8$ .

addi x7, x12, x12 #  $x7 = 32$

addi x7, x7, 32 #  $x7 = 32 + 32$

slli x7, x7, 2

add x7, x7, x5 #  $x5 = \text{address of C} [32 + 32]$

slli x4, x8, 2

add x4, x4, x5 #  $x4 = \text{address of AC}$

AC = 32 + 32

AC = 64

AC = 128

AC = 256

AC = 512

AC = 1024

AC = 2048

AC = 4096

AC = 8192

AC = 16384

AC = 32768

AC = 65536

AC = 131072

AC = 262144

AC = 524288

Created with Scanner Pro

Tu Cheng ycs8x5

\* Problem 4:

$$0 \times 2E_{hex} = 46 \Rightarrow 00101110 \text{ in binary}$$

$$0 \times 3D_{hex} = 61 \Rightarrow 0011101 \text{ in binary}$$

$$46 \times 61 = 2806 = 0xAFF616$$

$$61 = (32+16+8+4+1)$$

$$46 \times 61 = 46 \times (32+16+8+4+1)$$

l1 to  $0 \times 2E$

l1 t1  $0 \times 3D$

l1 t2

l1  $\times 0 46$

~~s1 l1 x11, x10, 5~~

~~s1 l1 x12, x10, 4~~

~~add x11, x11, x12~~

~~s1 l1 x12, x10, 3~~

~~add x11, x11, x12~~

~~s1 l1 x12, x10, 2~~

~~add x11, x11, x12~~

~~s1 l1 t3 to 4~~

~~add t2 t2 t3~~

~~s1 l1 t3 to 3~~

~~add t2 t2 t3~~

~~s1 l1 t3 to 2~~

~~add t2 t2 t3~~

Problem 5:

1) 142 instructions  $\Rightarrow$  8 bits ( $128 < 142 < 256$ )

32 registers  $\Rightarrow$  5 bits ; 16 bit immediates

1 reg in, 1 reg out :  $8 + 5 + 5 = 18$  bits  $\Rightarrow 2k$  bits

2 reg in, 1 reg out :  $8 + 5 + 5 + 5 = 23$  bits  $\Rightarrow 2k$  bits

1 reg in, 1 reg out, 1 imm :  $8 + 5 + 5 + 16 = 34$  bits  $\Rightarrow 4k$  bits

1 imm in, 1 reg out :  $8 + 16 + 5 = 29$  bits  $\Rightarrow 3k$  bits

(2)

Assuming it has 100 instruction  
Since the largest instruction type requires 8-bit instructions

Thus memory used by fixed-length encoding =  $40 \times 100 = 4000$  bits  
whereas, each instruction type in the variable encoding will  
use the number of bits from 1.

$$(100 \times 1.2 \times 2k + 0.3 \times 2k + 0.25 \times 40 + 0.25 \times 32) = 3200 \text{ bits}$$