

USE YOUR IMAGINATION™

Blue Book

EXAMINATION BOOK

Box No. Ø

NAME Yu Cheng

SUBJECT FCE b913

CLASS _____

SECTION section B

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11" x 8.5" 8 LEAVES 16 PAGES

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Problem 1:

1.1

Data forwarding alone doesn't suffice here because by the time the first ALU instruction completes the second ALU (execute) step, the second ALU is already in the first ALU step.

" a Nop between the 2 instructions, as well as data forwarding between the second ALU step and the decode step

CC0 CC1 CC2 CC3 CC4 CC5 CC6 CC7 CC8
ADD x1x2x3 IF ID ALU1 ALU2 MEM WB
ADD x5x4x1 IF ID ALU1 ALU2 MEM WB

2903

2903

(SIX) 0 . dX b1

(SIX) 8 11X b1

2903

2903

11X 01X SIX bbb

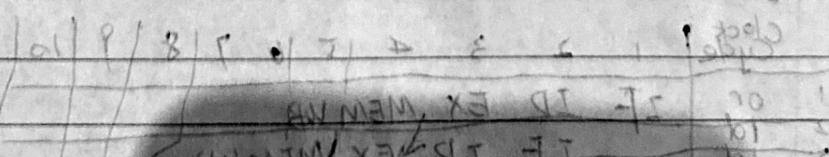
2903 SIX bbb

2903

2903

11X 01X SIX bbb

5.6



aluring controller

IF ID EX MEM NR

909891

Problem 2:

2.1

Hazards identified:

or x13 x12 x11

ld x10, 0(x13) Ex to 1st RAW Hazard

ld x11, 8(x13) Ex to 2nd RAW Hazard

add x12 x10 x11 MEM to 1st RAW [load-use-data] & MEM to 2nd RAW Hazard

subi x13 x12 16 Ex to 1st RAW Hazard

No PS introduced to resolve Hazards:

or x13 x12 x11

NOPs

NOPs

ld x10, 0(x13) Ex to 1st RAW Hazard resolution with 2 NOPs

ld x11, 8(x13) Ex to 2nd RAW Hazard resolved as well from above
2 NOPs

1 NOPs

add x12 x10 x11 MEM to 1st RAW [load-use-data] & MEM to 2nd RAW Hazard's resolved with 2 NOPs

NOPs

NOPs

subi x13 x12, 16 Ex to 1st only RAW Hazard resolved with 2 NOPs

2.2

Clock Cycle	1	2	3	4	5	6	7	8	9	10
1 or	ZF	ID	EX	MEM WB						
2 ld										
3 ld										
4 NOP										
5 add										
6 subi										

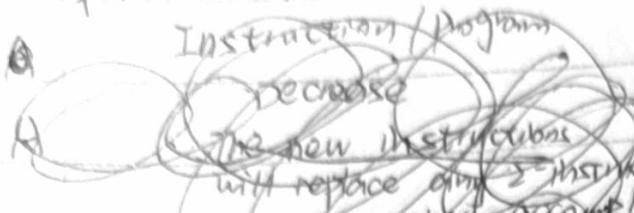
Mandatory NOP for which no forwarding solution possible

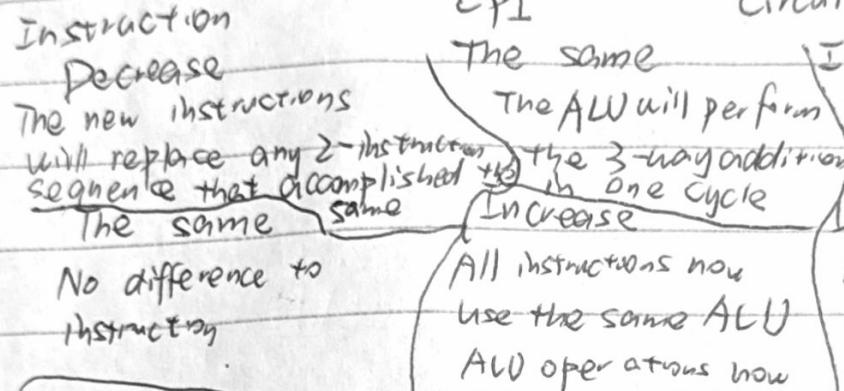
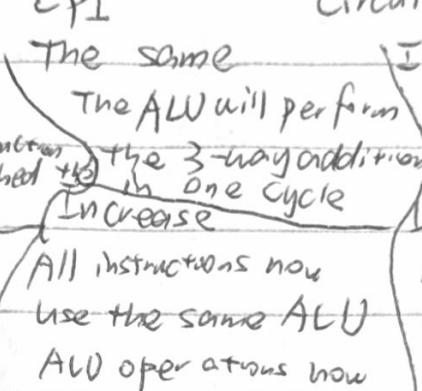
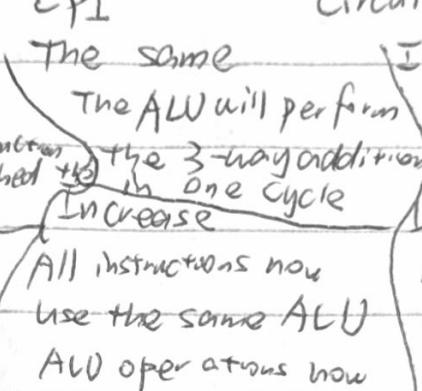
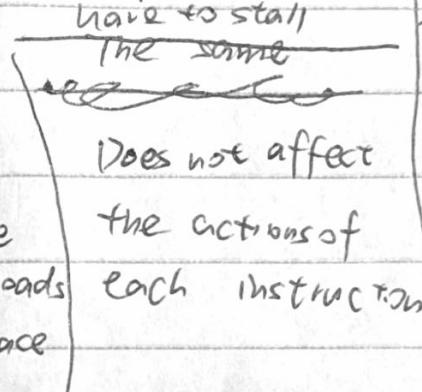
IF ID EX (MEM WB)

IF ID EX (MEM WB)

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problem 3:

A 
Instruction / Program
decrease
The new instructions will replace any 2-instruction sequence that accomplished the same, such as ADD rs1, rs2, rd and ADD rs3, rd, rd.

A 
Instruction
Decrease
The new instructions will replace any 2-instruction sequence that accomplished the same.
B 
CPI
The same
The ALU will perform the 3-way addition in one cycle.
Circuit complexity
decrease
3-operand ALU is more complex than a 2-operand one.
B 
CPI
The same
All instructions now use the same ALU.
ALU operations now have to stall.
Circuit complexity
decrease
Now there is one less adder / ALU cycle.
C 
CPI
The same or decrease
If the more registers enable the compiler to avoid loads and stores, Decrease
Otherwise, the same
Circuit complexity
increase
More registers complicate the register file.

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problem 4:

41

~~I-Mem is read: $140\text{ pJ} + 2 \times 70\text{ pJ} + 60\text{ pJ} = 340\text{ pJ}$~~

4.2

$$140P] + 2 \times 70P] + 60P] + 140P] = 480P]$$

4.3

$\overline{J-400} \rightarrow 2000 \text{ years} = 160 \text{ B.C.}$ $J = 700 \text{ A.D.}$

$$I - Mem + 2 \text{ registers} = 140 PJ + 2 \times 70 PJ = 280 PJ$$

Problem

~~11 - 2017.02.01.05~~

En el valle de la Cuenca del Río Magdalena se han establecido 1000 km² de plantaciones de café y 100 km² de cultivo de cacao.

$$\partial S_1 = \Gamma(S_1) + \Gamma(-S_1)$$

ANSWER

9484

Problem 5:

① representing $-1/4$:

$$\frac{1}{4} = 0.25 \Rightarrow$$

$$-2 + 127 = 125$$

So the bit pattern =

Sign Exponent Mantissa
1 0111101 ~~000000000000~~ since 0.25 is represented as .100000000000

Thus in binary: 1011110100000000000000000000000000000000

② representing -1/3:

The sign bit is 1 because $-1/3$ is negative

$\frac{1}{3}$ binary is a repeating fraction, approximating 1.0101010101...
binary which can be represented as 1.0101010101 multiplied by 2^{-1}
 $-1 + 127 = 126$

So the bit pattern:

Sign Exponent Mantissa
 | 0111110 010101010101010 which will fill up the Mantissa

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problem 6

	Sign	Exponent	Fraction	Decimal value
The number 0	0	00000	0000000000	0
Charge of an electron	0	00000	0000000000	6.10352×10^{-19} Coulombs
smallest positive normalized number largest	0	10000	1000000000	6.10352×10^{-5}
smallest positive normalized number	0	11110	1111111111	6.552×10^{-5}
smallest positive de-normalized number	0	00000	0000000000	5.96×10^{-5}
largest positive denormalized number	0	00000	1111111111	6.097×10^{-5}
distance b/w Earth and Neptune in inches	0	11111	0000000000	$+00$

charge of an electron is too smaller than the smallest number it can represent in half-formate. Thus it will be under flow to 0 in this format.

The Decimal value of smallest positive normalized number is $2^{-10} \approx 6.10352 \times 10^{-5}$

The Decimal value of largest positive normalized number is $2^{15}(2-2^{-10}) = 6.552$

The Decimal value of smallest positive denormalized number is $2^{-10} \approx 5.96 \times 10^{-5}$

The Decimal value of largest positive denormalized number is $(1-2^{-10})2^{-10} \approx 6.09756 \times 10^{-5}$

distance b/w Earth and Neptune in inches is much larger than the maximum representable number in half-precision

so it will overflow to $+00$