

Faculty of Computer Science School of Software Engineering

SQLITE RDBMS EXTENSION FOR DATA INDEXING USING B-TREE MODIFICATIONS

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OUTLINE

- B-tree and modifications
- SQLite and extensions
- B-tree modifications C++ library and research conducted
- Motivation, existing solutions and main goals
- Algorithm of selecting the index structure
- Technologies used
- Implementation
- Usage example and experiment conducted

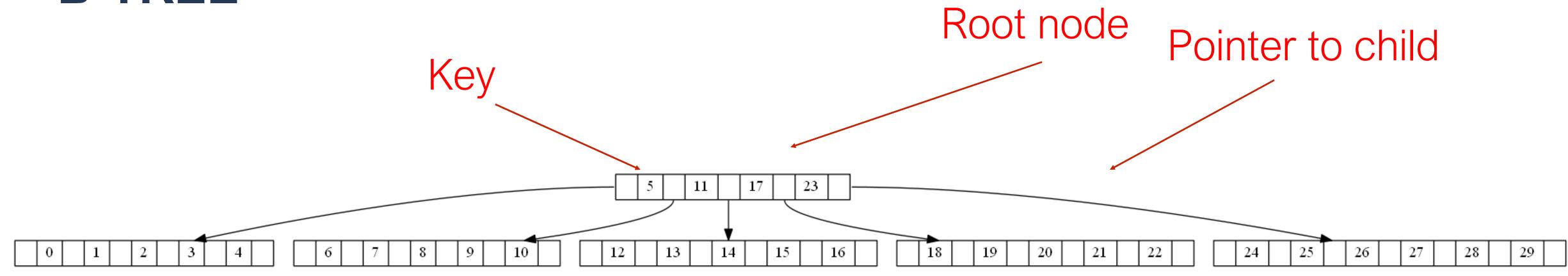


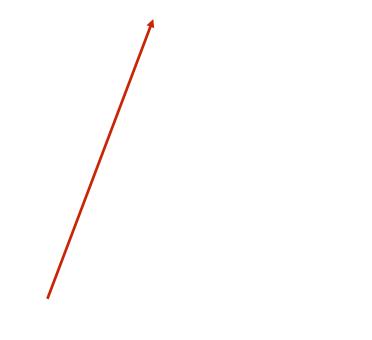
B-TREE

- Balanced search tree [1]
- Nodes may contain more than 1 key and more than 2 pointers to the children nodes [1]
- If some node contains k keys than it contains k+1 pointers to the children nodes [1]
- **B-tree order** is such a t number that (k is the count of keys in the node):
 - ✓ for each non-root node: $t 1 \le k \le 2t 1$ [1]
 - ✓ for root node in the non-empty tree: $1 \le k \le 2t 1$ [1]
 - ✓ for root node in the empty tree: k = 0 [1]
- B-tree height is $O(log_t n)$, where t is tree order and n is the count of keys in the tree [1]
- Usually used as the data index [1]



B-TREE





The B-tree example, t = 6

Leaf node



B-TREE OPERATIONS

Searching

- ✓ Time complexity $O(tlog_t n)$ [1]
- ✓ Memory usage O(t) [1]
- ✓ Disk operations count $O(log_t n)$ [1]

Nodes split (the part of insertion)

- ✓ Time complexity O(t) [1]
- ✓ Memory usage O(t) [1]
- ✓ Disk operations count O(1) [1]

Insertion

- ✓ Time complexity $O(tlog_t n)$ [1]
- ✓ Memory usage $-O(tlog_t n)$ for simple recursion and O(t) for tail recursion or loop [1]
- ✓ Disk operations count $O(log_t n)$ [1]

Deletion

- ✓ Time complexity $O(tlog_t n)$ [1]
- ✓ Memory usage $-O(tlog_t n)$ for simple recursion and O(t) for tail recursion or loop [1]
- ✓ Disk operations count $O(log_t n)$ [1]



B-TREE MODIFICATIONS

B⁺-tree

- ✓ Only leaf nodes contain real keys (real data), other nodes contain router keys [2]
- ✓ Deletion is probably faster than in B-tree

B*-tree

- ✓ Each node (except of the root node) is filled at least by 2/3 not 1/2 [3]
- ✓ Keys insertion in B*-tree is expected to be faster than in B-tree

• B*+-tree

- ✓ Developed previously by author of this work [4]
- ✓ Combines the main B+-tree and B*-tree features together [4]

^{[3] &}quot;B*-tree." NIST Dictionary of Algorithms and Data Structures. Available: https://xlinux.nist.gov/dads/HTML/bstartree.html (accessed Dec. 24, 2018).



SQLITE

- Popular open-source embedded relational DBMS
- Written in the C language
- Uses the B-tree as the default index
- SQLite extensions are the dynamically linked libraries [5]

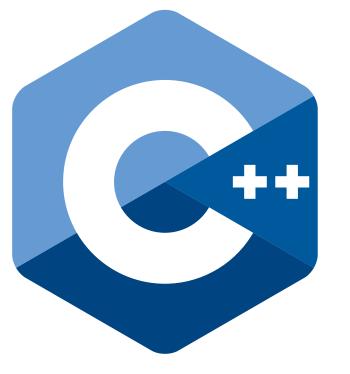




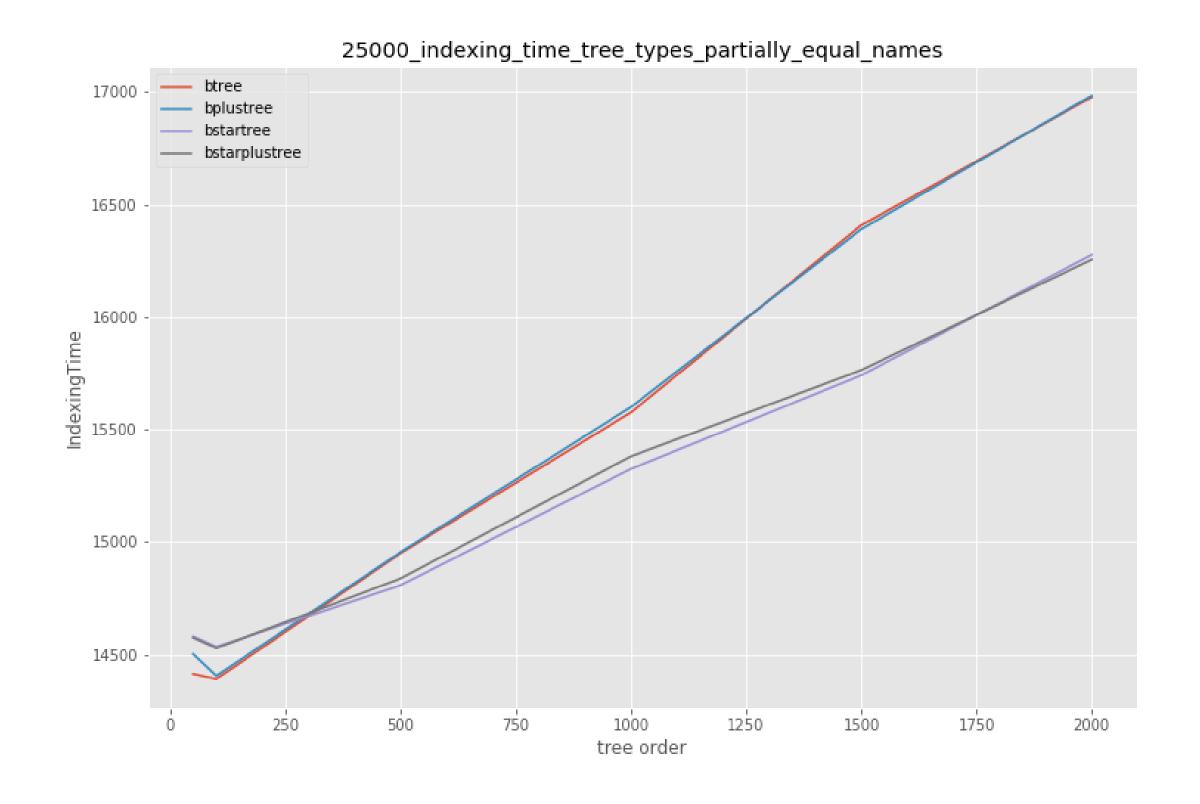


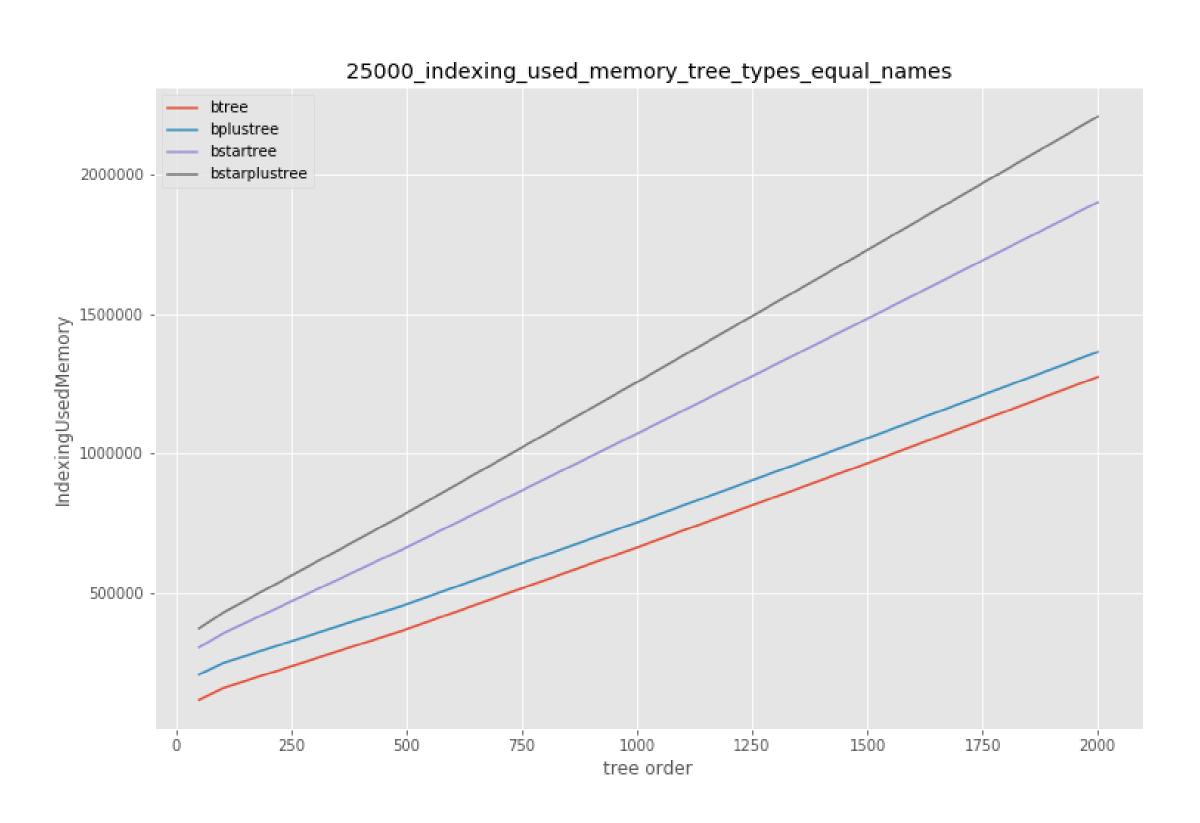
B-TREE MODIFICATIONS C++ LIBRARY

- Developed previously [4]
- Contains B-tree, B+-tree, B*-tree and B*+-tree implementations
- In the current work connected to the SQLite as the run-time loadable extension

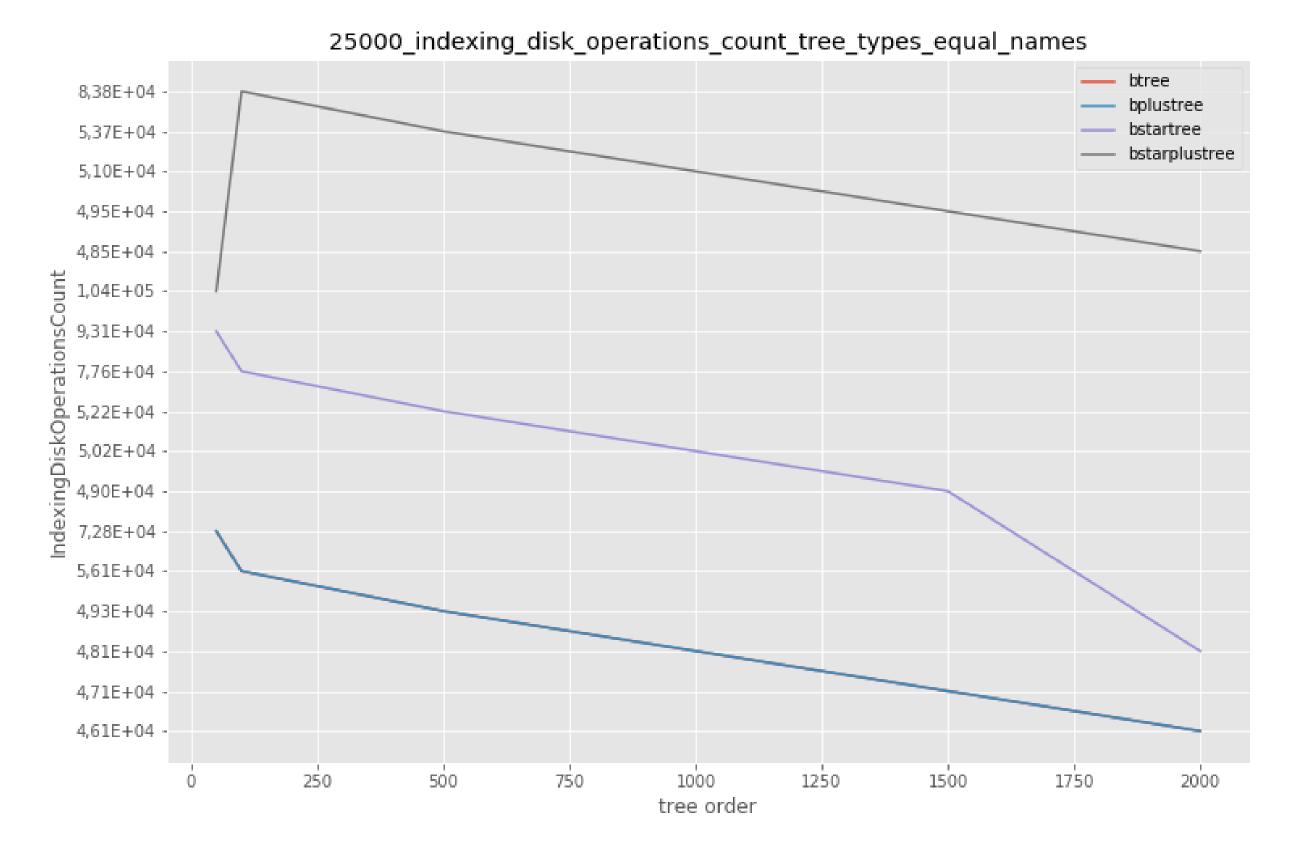




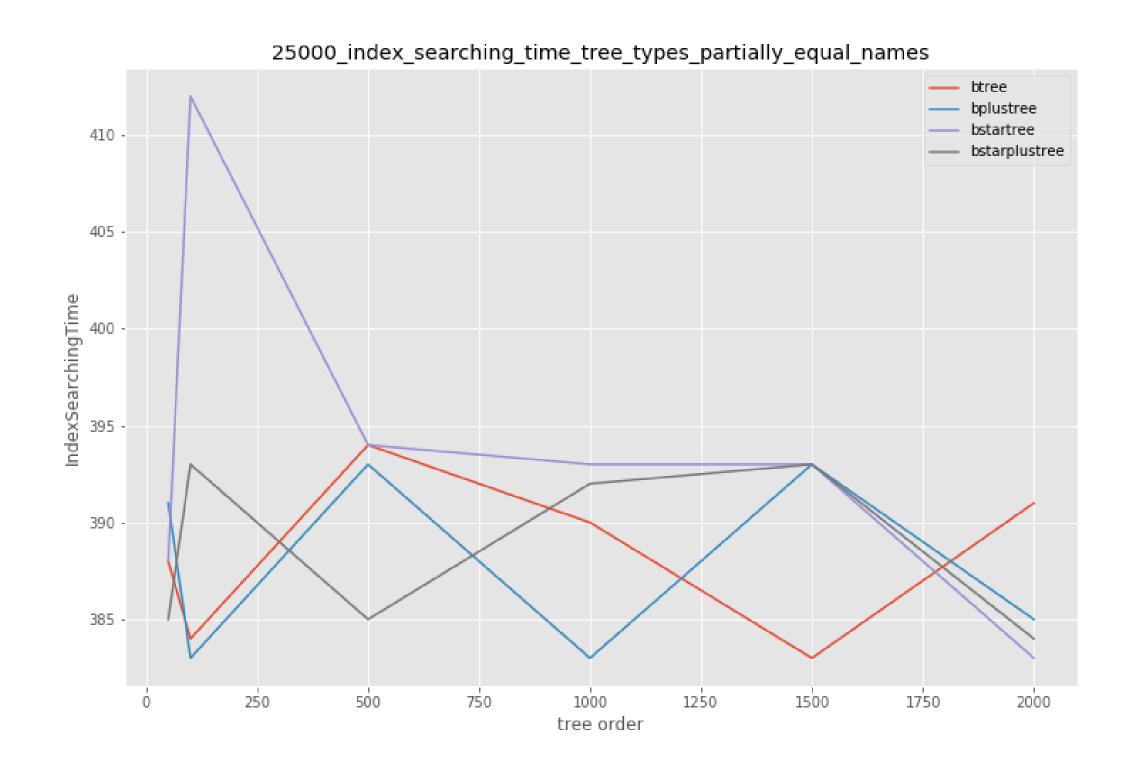


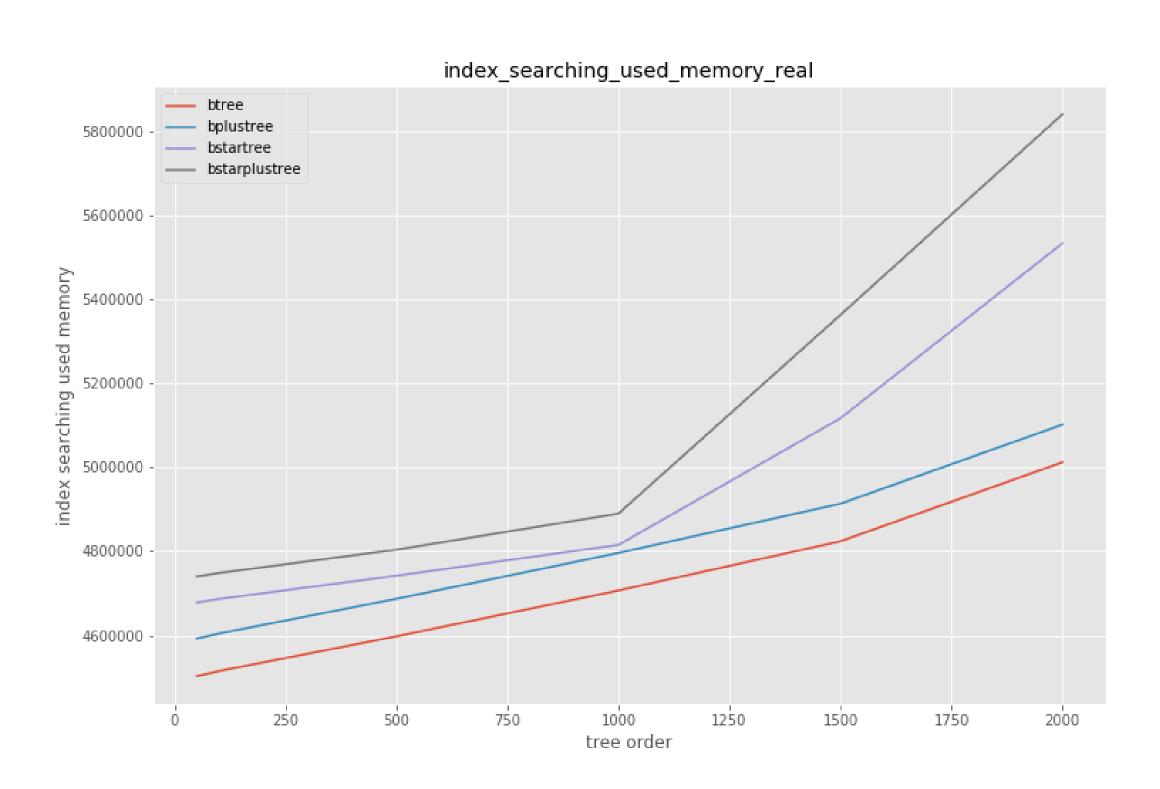




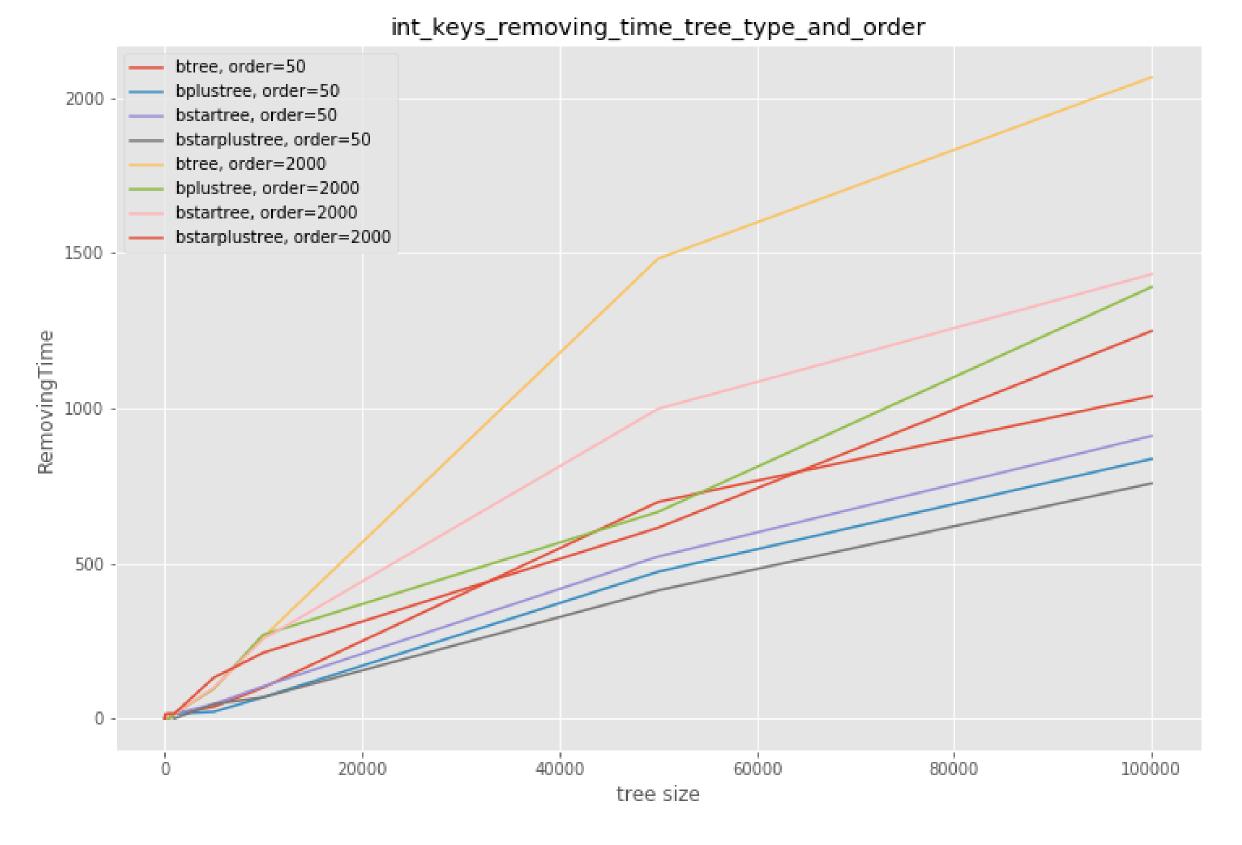














MOTIVATION AND EXISTING SOLUTIONS

Motivation

- ✓ In the current time, the processed data amounts are continuously growing
- ✓ It is necessary to develop new efficient approaches to data indexing in the DBMSs
- ✓ SQLite contains a small number of index structures by default it is relevant to add new ones

Existing solutions

- ✓ B-tree is the default index structure in the SQLite.
- ✓ There are some SQLite extensions (for example, the extension adding the R-tree indexing)
- ✓ No extensions with B+-tree, B*-tree, and B*+-tree were found



THE MAIN GOALS OF THE WORK

- To add B-tree modifications such as B+-tree, B*-tree and B*+-tree to SQLite
- To develop and implement an algorithm that would allow selecting the indexing data structure (B-tree, B+-tree, B*-tree or B*+-tree) when a user manipulates a table



ALGORITHM OF SELECTING THE INDEX STRUCTURE

- Selects from the B-tree and its modifications (B+-tree, B*-tree and B*+-tree)
- Runs at the start of each table operation (search, insertion, updating, deletion)
- Performs the index structure rebuilding only on each 1000-th operation and only for the first 10000 operations
- Index structure selection depends on the ratios of the numbers of operations of different types (search, insert, delete) on the tree
- The tree order of the B-trees and their modifications used in the SQLite extension developed in this work equals 750



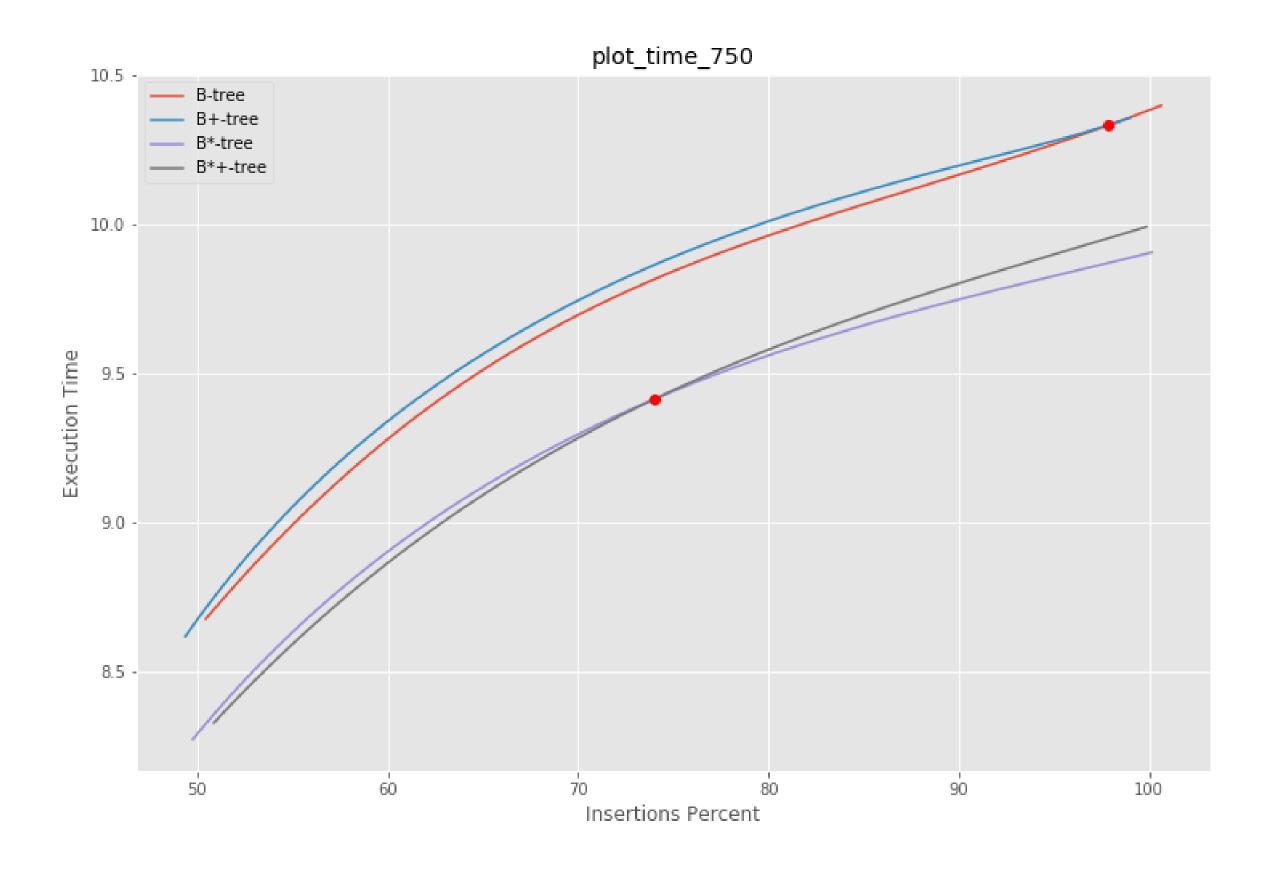
ALGORITHM OF SELECTING THE INDEX STRUCTURE

- 1. If the current total count of the operations on the tree is equal to 0 or more than 10000 or not a multiple of 1000, then exit the algorithm, otherwise go to step 2.
- 2. If the current count of the modifying operations on the tree (key insertions, key deletions) is less than 10 % of the current total count of the operations on the tree, then exit the algorithm, otherwise go to step 3.
- 3. If the current count of the key insertion operations is more than p=73.97% of the current count of the modifying operations on the tree, then select the B*-tree as the index structure and go to step 5, otherwise go to step 4.

- 4. Select the B*+-tree as the index structure and go to step 5.
- 5. If the new index structure was selected in the steps 3-4, then rebuild the existing index structure into the new selected index structure saving all the data stored in the existing index structure.

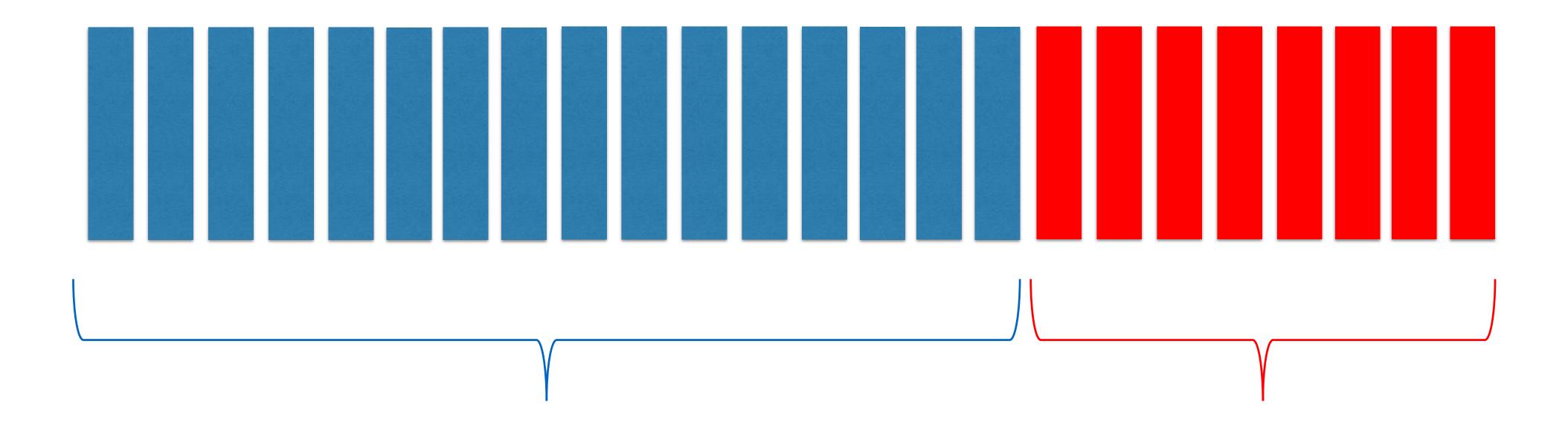


ALGORITHM OF SELECTING THE INDEX STRUCTURE





TREE KEY STRUCTURE



primary key value

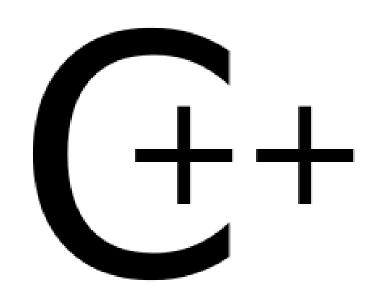
row id (8 bytes)



TECHNOLOGIES USED





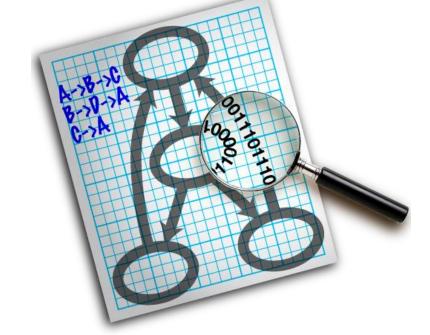














IMPLEMENTATION

- The C API for the B-tree modifications library is implemented as dynamically shared library using the extern "C" { ... } C++ statement
- The SQLite extension registers the *btrees_mods* module for creating the virtual tables
- Virtual table is any table created using such a module
- The extension also registers the *btreesModsVisualize*, *btreesModsGetTreeOrder*, *btreesModsGetTreeType* functions



IMPLEMENTATION

Method		Purpose
btreesModsCreate(sqlite3*, void*, int,	const	Creates a new table.
char* const*, sqlite3_vtab**, char**)		
btreesModsUpdate(sqlite3_vtab*,	int,	Inserts, deletes or updates a value of a row in
sqlite3_value**, sqlite_int64*)		the table.
btreesModsFilter(sqlite3_vtab_cursor*,	int,	Searches for a row in the table.
const char*, int, sqlite3_value**)		



IMPLEMENTATION

Method	Purpose		
btreesModsVisualize(sqlite3_context*, int,	Outputs the graphical representation of the		
sqlite3_value**)	table's index structure (tree) into the GraphViz		
	DOT file.		
btreesModsGetTreeOrder(sqlite3_context*,	Outputs the order of the tree used as the		
int, sqlite3_value**)	table's index structure.		
btreesModsGetTreeType(sqlite3_context*,	Outputs the type of the tree (1 – B-tree, 2 – B+-		
int, sqlite3_value**)	tree, 3 - B*-tree, 4 - B*+-tree) used as the		
	table's index structure.		



USAGE EXAMPLE

```
SQLite version 3.26.0 2018-12-01 12:34:55
Enter ".help" for usage hints.
Connected to a transient in-memory
Use ".open FILENAME" to reopen on a persistent database.
sqlite> .load ./btrees_mods
sqlite> CREATE VIRTUAL TABLE btt USING btrees_mods(id INTEGER PRIMARY KEY, a INTEGER, b TEXT);
sqlite> INSERT INTO btt VALUES (4, 2, "ABC123");
sqlite> INSERT INTO btt VALUES (7, 3, "def");
sqlite> SELECT * FROM btt WHERE id = 4;
4 2 ABC123
sqlite> SELECT * FROM btt WHERE id = 7;
7 | 3 | def
sqlite> SELECT * FROM btt WHERE id = 4 OR id = 7;
4 2 ABC123
7 3 def
sqlite> .tables
btrees_mods_idxinfo btt
                                           btt_real
sqlite> SELECT * FROM btt_real;
4 2 ABC123
sqlite> SELECT * FROM btrees_mods_idxinfo;
btt | 1 | 0 | id | INTEGER | 4 | tree_18291557263097.btree
sqlite> DROP TABLE btt;
sqlite> .tables
btrees_mods_idxinfo
sqlite> SELECT * FROM btrees_mods_idxinfo;
sqlite> .exit
```



EXPERIMENT CONDUCTED USING THE SQLITE EXTENSION

Operation on the table	Total execution time (ms)	Mean execution time per row (ms)
Table creation	20	-
First 500 rows insertion	10301	20.6
Next 500 rows insertion	10322	20.6
1001st row insertion (including the B+-tree into the B*-tree rebuilding)	40	40
Next 499 rows insertion	9386	18.8
Last 500 rows insertion	9032	18.1



EXPERIMENT CONDUCTED USING THE SQLITE EXTENSION

Operation on the table	Total execution time (ms)	Mean execution time per row (ms)
First 500 rows deletion	11558	23.1
Next 500 rows deletion	10708	21.4
1001st row insertion (including the	62	62
B*-tree into the B*+-tree rebuilding)		
Next 499 rows deletion	9418	18.9
Last 500 rows deletion	8863	17.7
1000 rows insertion	18890	18.9
Next 5000 rows insertion	92395	18.5
(including the B*+-tree into the B*-		
tree rebuilding)		



EXPERIMENT CONDUCTED USING THE SQLITE EXTENSION

- The key insertion on the B*-tree was faster than on the B+-tree
- The key deletion on the B*+-tree was faster than on the B*-tree
- The key insertion on the B*-tree is slightly faster than on the B*+-tree
- The search on the table takes about 1 ms on all the B-tree modifications considered in this work



SUMMARY

- B-trees modifications library is connected to the SQLite as an extension using developed in this work C API
- Research conducted using this library is presented
- Algorithm of selecting the index structure is developed and implemented
- Experiment on the trees performance is conducted using the developed SQLite extension



REFERENCES

- [1] R. Bayer and E. McCreight, "Organization and maintenance of large ordered indexes," *Acta Informatica*, vol. 1, no. 3, pp. 173 189, 1972.
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- [4] A. Rigin, "On the Performance of Multiway Trees in the Problem of Structured Data Indexing," (in Russian), coursework, Dept. Soft. Eng., HSE, Moscow, Russia, 2018.
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Thank you for your attention!

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