Binary Sequences with Minimum Peak Sidelobe Level up to Length 68

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Abstract: Results of an exhaustive search for minimum peak sidelobe level binary sequences are presented. Several techniques for efficiency implementation of search algorithm are described. A table of number of non-equivalent optimal binary sequences with minimum peak sidelobe (MPS) level up to length 68 is given. This number can be used in prediction of the longest length for a given sidelobe level of binary sequences. The examples of optimal binary MPS sequences having high merit factor are shown.

Exhaustive search, binary, minimum peak sidelobe, sequences, aperiodic, autocorrelation function, merit factor, optimal, peak sidelobe

Introduction

Binary sequences with low autocorrelation sidelobe levels are useful in different applications: radar, communication systems, information security, synchronization and so on. The main point in radar is to gain the signal-to-noise ratio benefits of a long pulse along the range resolution of a short pulse. This requires longer and longer sequences, with low aperiodic autocorrelation sidelobes. There are no theoretical methods to generate such sequences, so they have been produced by computer searches. The computational complexity of such searches increases with length: a search for N=71 is approximately twice as difficult as that for N=70 that is why computational time doubles.

This paper adds to available knowledge for record length of binary MPS sequence and provides a number of non-equivalence classes for each lengths up to $N = 68$. Despite the computational challenges, progress has been made over time due to improvements in both computational resources and search methods.

Lindner [1] in 1975 did an exhaustive search for binary MPS sequences up to $N = 40$. Cohen et al. [2] in 1990 continued up to $N = 48$. Coxson and Ruso [3] executed an exhaustive search of binary MPS sequences for $N = 64$. So binary MPS sequences are only known for the lengths up to $N = 48$ and for the length $N = 64$.

Apart from known results of global exhaustive search of binary MPS sequences, there are some useful results of local search of binary sequences with the low aperiodic autocorrelation. Kerdock et al. [4] in 1986 found binary sequences for lengths $N = 51,69,88$ with $PSL = 3,4,5$ respectively. These are still the best. Elders-Boll et al. [5] in 1997 found best known binary PSL sequences with $PSL = 4$ for the lengths from $N = [49;61]$. Coxson and Ruso [3] in 2004 continued the list of best known binary PSL sequences with $PSL = 4$ up to $N = 70$. Nunn and Coxson [6] in 2008 found best known binary PSL sequences with $PSL = 4$ up to $N = 82$ and with $PSL = 5$ for $N = [83,105]$.

We have made an exhaustive search of binary MPS sequences during 4 months implementation 1 supercomputer Flagman RX240T8.2 on the base of 8 NVIDIA TESLA C2059 with 3584 parallel graphical processors and on the base of 2 processors Intel Xeon X5670 (up to Six-Core) and using CUDA compilation. Our algorithm is based on a concept of Mertens [7], the branch-and-bound algorithm, using new assembler instructions for calculation of aperiodic autocorrelation function and "package" regime.

Preliminaries

A binary sequence of length N is an N-tuple $A = (a_0, a_1, ..., a_{N-1})$ where each $a_n \in \{-1, 1\}$, $n = 0, 1, ..., N-1$. The aperiodic autocorrelation of A at shift τ defined as

$$C_\tau = \sum_{n=0}^{N-1-\tau} a_n \cdot a_{n+\tau}. \quad (1)$$

There are two principal measures of level of sidelobe level. The primary measure is the peak sidelobe level (PSL):

$$PSL(C) = \max_{1 \le \tau \le N-1} |C_\tau|. \quad (2)$$

For optimal binary sequences by PSL criteria the peak sidelobe has to be minimum:

$$MPS = \min_A PSL. \quad (3)$$

A secondary measure, is the merit factor (MF):

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$$MF(C) = \frac{N^2}{2\sum\limits_{\tau=1}^{N-1}[C_\tau]^2}. \qquad (4)$$

PSL affects the maximum of self interference of the sequence and merit factor determine average interference. There are three operations that preserve peak sidelobe level in binary codes: reversal: $R(a_n) = a_{N-1-n}$, negation: $N(a_n) = -a_n$, alternating sign $S(a_n) = (-1)^n a_n$. The sequences obtained within such transformations will be formed class of equivalence.

We are interested to find all non-equivalent classes of binary MPS sequences for each length N from the range $N = [49,68]$.

Effective global algorithm for exhaustive search of binary PSL sequences

We modified all achievements of Nunn and Coxson algorithm which based on the main idea of Mertens's branch-and-bound algorithm as described below.

Base idea of branch-and-bound algorithm. Let us consider the binary sequence from two opposite sides of the length $N/2$. We called them left and right half-sequences. In the first step we ? choose left bit a_0 and right bit a_{N-1}. They can be considered like new code $a_0\, a_{N-1}$. This pair can be "00", "01", "10" and "11". We present the algorithm in graph form, where possible pairs of opposite bits can be shown like leaves of the tree:
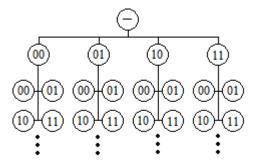
Fig. 1 Set of sequences presented in tree manner

We can check the PSL level of autocorrelation during each step. If on the next step the PSL is not valued all branch of the tree can be excluded from searching. Also due to operations that preserve peak sidelobe level in binary codes we can say that on the first step one pair $a_0 = 0$, $a_{N-1} = 0$ is formed whole class of

equivalence for all other possible values a_0 a_{N-1}. So we the only one branch instead of for. On the second? step we have only 3 non-equivalent branches

$$0,0,a_2,...,a_{N-3},0,0, \qquad 0,0,a_2,...,a_{N-3},1,0, \qquad 0,1,a_2,...,a_{N-3},1,0.$$

The search space can be reduced by excluding equivalent branches, as shown in Fig.2.
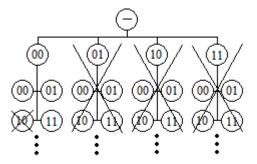
Fig.2. Excluding of the equivalent sequences

Our modifications.

1. We used recursive implementation of our algorithm (Fig.3) using inline options for all external operations.

2. Our main idea is to use new assembler instructions for computing autocorrelation function of binary sequences. We can find side lobes of aperiodic autocorrelation using XOR operation. To determine the level of sidelobe we have to calculate the numbers of zeros and units for each shift of sequences. New Intel processors have microarchitecture Intel Core of version SSE4.2 which operating with the set of command on low level. For example C/C++ Microsoft compiler has function __popcnt64 of intrin library and also compilersGCC and G++ has function _mm_popcnt_u64 of smmintrin library for calculation the number of units in binary sequences by 1 cycle.

3. For excluding equivalent sequences we used reverse transformation for two bytes at the time instead of each bit. All possible reverses are stored in static massive with 65536 different bit variations.

4. Also we realized parallel computing in multiprocessor system for all set of non-equivalent sequences separately each from other. We implemented our algorithm on CUDA SDK using function __popcll() for calculating number of unit bits.

5. Finally we used "package" regime to find some binary PSL sequences with the lengths $N, N+2, N+4,....,$ because cross correlation functions for left and right parts of the sequences with lengths $N, N+2, N+4,....,$ are the same.
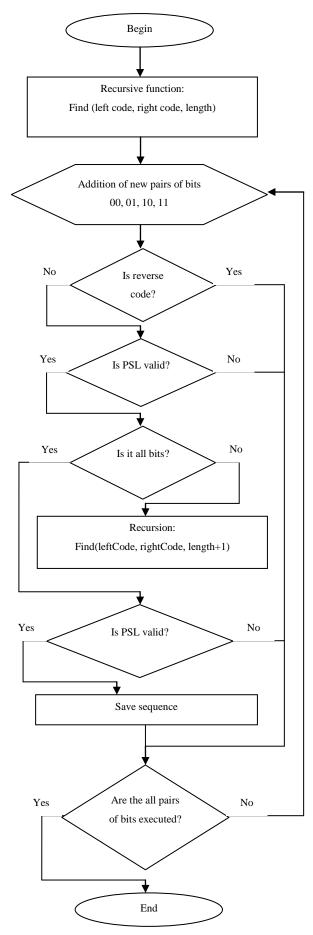
Fig.3 Block-scheme of recursive algorithm

Results of exhaustive search of binary MPS sequences

Our results are presented in Table 1. There are number of binary PSL sequences, which means that PSLi sequences have exactly $PSL = i$, not less. Synthesized sequences are available on our website [7].

Table 1. Size of set non-equivalent PSL sequences

Length	Size of set non-equivalent PSL sequences				
N	PSL1	PSL2	PSL3	PSL4	PSL5
2	1	0	0	0	0
3	1	1	0	0	0
4	1	1	1	0	0
5	1	3	1	1	0
6	0	4	4	1	1
7	1	7	5	5	1
8	0	8	12	8	6
9	0	10	23	20	29
10	0	5	46	35	30
11	1	7	53	97	52
12	0	16	87	133	152
13	1	11	126	287	246
14	0	9	152	486	583
15	0	13	223	800	1050
16	0	10	361	1173	2176
17	0	4	307	2243	3490
18	0	2	339	3025	7205
19	0	1	419	4661	11645
20	0	3	625	6245	21456
21	0	3	505	9826	32539
22	0	0	378	11840	58331
23	0	0	515	16533	86812
24	0	0	858	20673	148583
25	0	1	436	29794	206762
26	0	0	242	31205	329356
27	0	0	388	40193	469454
28	0	2	624	49884	753204
29	0	0	284	63059	966451
30	0	0	86	59506	1390617
31	0	0	251	71546	processing
32	0	0	422	89190	2894816
33	0	0	139	98644	processing
34	0	0	51	84636	4567602
35	0	0	111	98331	processing
36	0	0	161	118624	8951507
37	0	0	55	119053	processing
38	0	0	17	89067	11788025

Table 1. Size of set non-equivalent PSL sequences (continuing)

39	0	0	30	101808	processing
40	0	0	57	118731	22333659
41	0	0	15	112039	processing
42	0	0	4	72716	24453952
43	0	0	12	83417	processing
44	0	0	15	98334	44270683
45	0	0	4	82538	processing
46	0	0	1	47331	41354620
47	0	0	1	54896	processing
48	0	0	4	64424	74010972
49	0	0	0	49088	processing
50	0	0	0	25169	57294359
51	0	0	1	28249	processing
52	0	0	0	33058	processing
53	0	0	0	23673	processing
54	0	0	0	10808	processing
55	0	0	0	11987	processing
56	0	0	0	15289	processing
57	0	0	0	9476	processing
58	0	0	0	4026	processing
59	0	0	0	4624	processing
60	0	0	0	5542	processing
61	0	0	0	3246	processing
62	0	0	0	1212	processing
63	0	0	0	1422	processing
64	0	0	0	1859	processing
65	0	0	0	1003	processing
66	0	0	0	324	processing
67	0	0	0	414	processing
68	0	0	0	491	processing

The results of an exhaustive search of binary MPS sequences up to length $N = 68$ are presented in Table 2. Also in the Table 2 there are the highest level of MF between binary MPS sequences and examples of such sequences in hexadecimal format.

Table 2. Results of exhaustive search of binary MPS sequences

Length	PSL	MF	Optimal or best known by MF?	Sequence	Size of set
2	1	2	yes	0	1
3	1	4,5	yes	1	1
4	1	4	yes	2	1

Table 2. Results of exhaustive search of binary MPS sequences (continuing)

5	1	6,25	yes	02	1
6	2	2,571	yes	02	4
7	1	8,167	yes	0D	1
8	2	4	yes	1A	8
9	2	3,375	yes	02C	10
10	2	3,846	yes	02C	5
11	1	12,1	yes	0ED	1
12	2	7,2	yes	0A6	16
13	1	14,083	yes	00CA	1
14	2	5,158	yes	00CA	9
15	2	4,891	no	0329	13
16	2	4,571	no	1DDA	10
17	2	4,516	yes	0192B	4
18	2	6,48	yes	0168C	2
19	2	4,878	no	0EEDA	1
20	2	5,263	no	04D4E	3
21	2	6,485	no	005D39	3
22	3	6,205	yes	013538	378
23	3	5,628	yes	084BA3	515
24	3	8	yes	31FAB6	858
25	2	7,102	no	031FAB6	1
26	3	7,511	no	07015B2	242
27	3	9,851	yes	0F1112D	388
28	2	7,84	yes	4B7770E	2
29	3	6,782	yes	04B7770E	284
30	3	7,627	yes	03F6D5CE	86
31	3	7,172	yes	07E736D5	251
32	3	7,111	no	01E5AACC	422
33	3	8,508	yes	003CB5599	139
34	3	8,892	yes	0CC01E5AA	51
35	3	7,562	no	00796AB33	111
36	3	6,894	no	3314A083E	161
37	3	6,985	no	031D5AD93F	55
38	3	8,299	yes	003C34AA66	17
39	3	6,391	no	13350BEF3C	30
40	3	7,407	yes	2223DC3A5A	57
41	3	7,504	no	038EA520364	15
42	3	8,733	yes	04447B874B4	4

Table 2. Results of exhaustive search of binary MPS sequences (continuing)

43	3	6,748	no	005B2ACCE1C	12
44	3	6,286	no	202E2714B96	15
45	3	6,575	no	02AF0CC6DBF6	4
46	3	6,491	no	03C0CF7B6556	1
47	3	7,126	no	069A7E851988	1
48	3	6,128	no	24AC8847B87C	4
49	4	8,827	yes	05E859E984451	49088
50	4	8,17	yes	038FE23225492	25169
51	3	7,517	no	0E3F88C89524B	1
52	4	8,145	yes	05FB6D5D9D8E3	33058
53	4	7,89	no	00FF66EAE96B1C	23673
54	4	7,327	no	043B48A28793B3	10808
55	4	7,451	no	1658A2BC0A133B	11987
56	4	8,167	yes	0C790164F6752A	15289
57	4	7,963	no	01B4DE3455B93BF	9476
58	4	8,538	yes	008D89574E1349E	4026
59	4	8,328	no	1CAD63EFF126A2E	4624
60	4	8,108	no	119D01522ED3C34	5542
61	4	7,563	no	0024BA568EB83731	3246
62	4	8,179	yes	000C67247C59568B	1212
63	4	9,587	yes	1B3412F0501539CE	1422
64	4	9,846	yes	26C9FD5F5A1D798C	1859
65	4	8,252	no	04015762C784EC369	1003
66	4	7.751	no	03FEF2CCB0B8CAC54	324
67	4	7.766	no	073C2FADC44255264	414
68	4	8.438	no	562B8CA48E0C9027E	491

Conclusion

Optimal binary MPS sequences are updated for lengths 49 to 63 and 65 to 68. Also the number of non-equivalent PSL1, PSL2, PSL3, SPL4 and PSL5 sequences are found for lengths 2 to 68. The number of MPS and PSL sequences are useful for prediction of longest sequences possible for a given PSL. Such optimal sequences are highly sought after in radar and cryptography. These sequences were considered from MF criteria and the MPS sequences with highest MF were identified. The number of PSL is sequences rapidly increased for increasing PSL. For example for the length 48 the number of nonequivalent binary sequences with PSL3 is 4, PSL4 is 64424, PSL5 is 74010972.

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Preparing for

INTERNATIONAL WORKSHOP ON CODING AND CRYPTOGRAPHY, WCC 2013, April 15-19, 2013, Bergen, Norway.

Date: 14 December 2012.