DORAM: What is? Why Care? What Did we Do?

Maliciously-secure DORAM with the best-known semi-honest asymptotics.
Fastest DORAM for large (proj: $> 2^{21}$)

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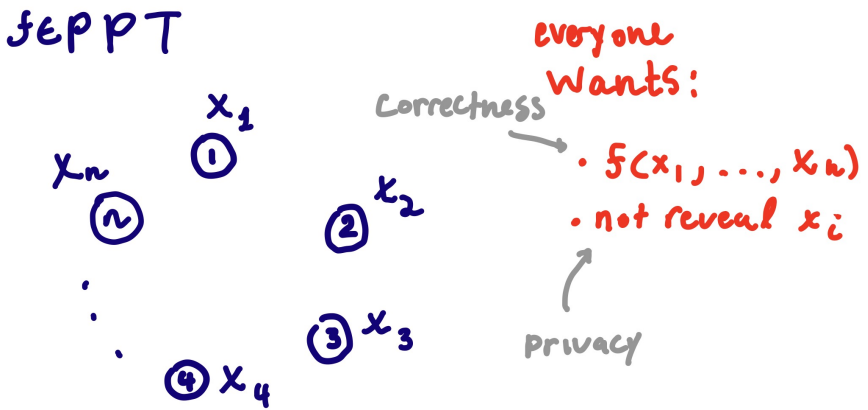
Overview

1. Preliminaries
2. DORAM Uses
3. Our Contributions
4. Future Work

Prelemenaries: What is Secure Multi-Party Computation (MPC)?

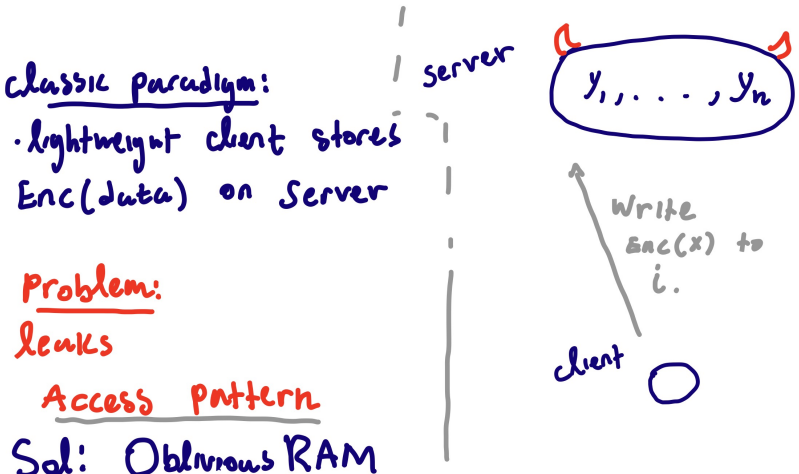


Prelemenaries: Semi-Honest vs Malicious Security in MPC

Semi-Honest security: if everyone follows the protocol precisely, security is gureateed.

Malicous Security: no matter what anyone does, security is gurenteed.

Prelemenaries: What is ORAM?



Prelemenaries: Bottlenecks of ORAM

- $O(\log N)$ interactive sequencial rds per r/w query
- Each ORAM serves a Single client
- Complex client
- Amortized (bad phases)

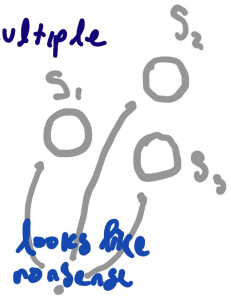
(we simplify here, there have been works attacking this, largely Still issue)

Prelemenaries: What is **D**ORAM?

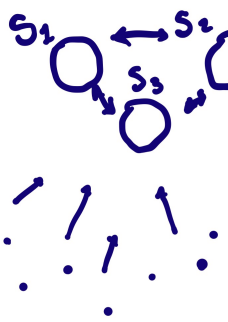
Distributed Oblivious RAM is a Relaxation of the problem.

Add: trust assumption, multiple servers, Some are honest.

Ask: Can we do better?



DORAM Uses: ORAM-replacement supporting many extremely-lightweight clients



! Why Care? !

- lots of clients!
  - ↳ extremely simple
    - no install needed
- 3x overhead over trusting client
- 1 rd per query (instead of $O(\log N)$)

instead of $O(\log N)$ Amortized

DORAM Uses: MPC in the random access model RAM-MPC

Circuits are a pain! ! Why care? !
+ circuits are less efficient

DORAM enables RAM-MPC
(write python* instead of circom)

Our Contributions: Theoretical Contributions

Let N be the number of elements stored in the DORAM, D be the payload size, κ be the computational security parameter.

Theorem (DORAM, informal)

There exists a (3,1)-MPC maliciously secure MPC scheme on bits or fields with $O(|C|)$ communication and computation complexity which achieves $O((\kappa + D) \log N)$ amortized communication complexity and amortized computation complexity per random access.

only semi-honest was known.

Theorem (RAM-MPC, Informal)

There exists a (3,1)-DORAM scheme which achieves $3(\log N + D)$ client query communication complexity, $O((\kappa + D) \log N)$ amortized $\{$communication,computation$\}$ complexity between the servers per query, and **non-amortized** *1 round of interaction per query for client.*

only semi-honest was known.

Our Contributions: Practical Comparison with Existing DORAM

Takeaway (mostly projected): for practical databse sizes and good networks, we improve on previous and concurrent semi-honest works by a factor of 10x-100x. [1]
Expect benchmarks on Saturday, but roughly, at $N = 2^{20}$ we able to get while state of the art can do 600 queries (semi-honest) while state of the art can do 1 query per second (semi-honest)

[1] our tests are not precise (yet)!

Our Contributions: Open Source EMP-Toolkit code

EMP Toolkit (158 GitHub stars) – we're workin on it!

Thanks Everyone!

Demo! (with DORAM chant)

goto
`https://matanshtepel.com/DORAM/HackLodge_demo.html`