

A thick black L-shaped frame is positioned on the left and right sides of the slide, framing the main title and chapter information.

# **BACS1024**

# **INTRODUCTION TO**

# **COMPUTER SYSTEMS**

## **Chapter 3: Floating Point Representation**

# 0. Overview

1. Signed & Unsigned Number Representation
2. Floating-point Number Representation & Standard
3. Bitwise Logical Operations

# **1. Signed & unsigned Num Representation**

# 1. Signed & Unsigned Number Representation

## ■ Integer representation

### □ Unsigned integer

- ❖ hold a positive value, and no negative value
- ❖ uses the most significant bit (MSB) as a part of the value
- ❖ E.g.: **Unsigned** integer:  $1111\ 1111_2 = +255_{10}$

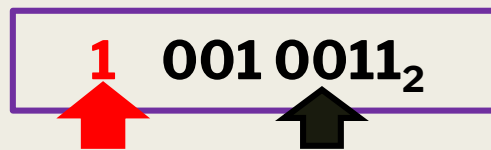
### □ Signed integer

- ❖ hold both positive and negative numbers
- ❖ uses the most significant bit (MSB) to identify if the number is positive or negative. **0** indicates a **positive** while **1** indicates a **negative**.
- ❖ E.g.: **Signed** integer:  $0111\ 1111_2 = +127_{10}$   
 $1111\ 1111_2 = -127_{10}$

However, this made calculation difficult.

# 1. Signed & Unsigned Number Representation

- In **Signed-and-magnitude representation**



**sign bit (MSB)**

**Magnitude bits**

- This made calculation difficult.

□ E.g.: In **signed** integer:  $0111\ 1111_2 = 2^6 + 2^5 + 2^4 + 2^3 + 2^2 + 2^1 + 2^0$   
 $= +127_{10}$   
 $1111\ 1111_2 = -2^7 + 2^6 + 2^5 + 2^4 + 2^3 + 2^2 + 2^1 + 2^0$   
 $= -1_{10}$   
 $\neq -127_{10}$

- Solution: **Two's Complement**

# 1. Signed & Unsigned Number Representation

## ■ Two's Complement

- 2 steps: 1.) One's complement

Perform inversion. i.e. change 1's to 0's and 0's to 1's

2.) Adding 1 to the result of One's complement

- As the computer stores data in byte basis, therefore **8 bit system** is used.

- E.g.:

$$10_{10} = 0000\ 1010_2$$

$$\text{One's complement} = 1111\ 0101_2$$

$$\text{Add } 1_2 = \underline{\hspace{2cm}} 1_2 +$$

$$\text{Two's complement} = 1111\ 0110_2$$

$$= -(2^7) + 2^6 + 2^5 + 2^4 + 2^2 + 2^1$$

$$= -10_{10}$$

# 1. Signed & Unsigned Number Representation

## ■ Two's Complement

- ❑ Uses: To represent negative value & perform subtraction.
- ❑ E.g.:  $30_{10} - 10_{10}$   
 $= 30_{10} + (-10_{10})$

$$30_{10} = 0001\ 1110_2$$

$$10_{10} = 0000\ 1010_2$$

$$\text{One's complement} = 1111\ 0101_2$$

$$\text{add } 1_2 = \underline{\quad\quad\quad 1_2 +}$$

$$\text{Two's complement} = 1111\ 0110_2$$
$$= -10_{10}$$

### Continued:

#### Decimal

$$\begin{array}{r} 30_{10} \\ + (-10)_{10} \\ \hline 20_{10} \end{array}$$

#### Binary

$$\begin{array}{r} 0001\ 1110_2 \\ + 1111\ 0110_2 \\ \hline (1)0001\ 0100_2 \\ = 20_{10} \end{array}$$

# 1. Signed & Unsigned Number Representation

## ■ **Overflow**

- ☐ Occurs when the result of an arithmetic operation does not fit into the fixed number of bits available for the result.
- ☐ Occur only when both operands have the same sign.
- ☐ Detected by the fact that the sign of the result is opposite of both operands.
- ☐ Stored in overflow flag (OF)

## ■ **Carry**

- ☐ Occurs when the result of an arithmetic operation exceeds the fixed number of bits allocated, without regard to the sign.
- ☐ The carry bit is ignored in single precision 2's complement addition and subtraction.
- ☐ Detected when extra '1' bit is generated.
- ☐ Stored in carry flag (CF)



# 1. Signed & Unsigned Number Representation

## ■ Overflow

□ E.g.:

$$\begin{array}{rcl} 64_{10} & 0100\ 0000_2 & \\ + 65_{10} & + 0100\ 0001_2 & \\ \hline 129_{10} & 1000\ 0001_2 & = -127_{10} \end{array}$$

In **8 bit system**, data ranges from:  
-128 to +127.

$$64_{10} + 65_{10} = 129_{10} \text{ which is } > \mathbf{127}$$

**i.e.** 129 does not fit into the number of bits available.

Therefore, **OVERFLOW** occur.

Since  $129_{10} \neq -127_{10}$ , this is **invalid**.

## ■ Carry

□ E.g.:

$$\begin{array}{rcl} 106_{10} & 0110\ 1010_2 & \\ + (-2)_{10} & + 1111\ 1110_2 & \\ \hline 104_{10} & (1)0110\ 1000_2 & = 104_{10} \end{array}$$

In **8 bit system**, extra 1 bit is generated. Carry occur.

Carry ignore

## **2. Floating-point Number Representation & Standard**

## 2. Floating-point Num Representation & Standard

- **Exponential notation**

- ❑ For decimal numbers, floating point is represented in **scientific notation**.

- ❑ E.g.:

- 12345

- =  $12345 \times 10^0$

- =  $0.12345 \times 10^5$

- =  $0.0012345 \times 10^7$

- =  $123450000 \times 10^{-4}$

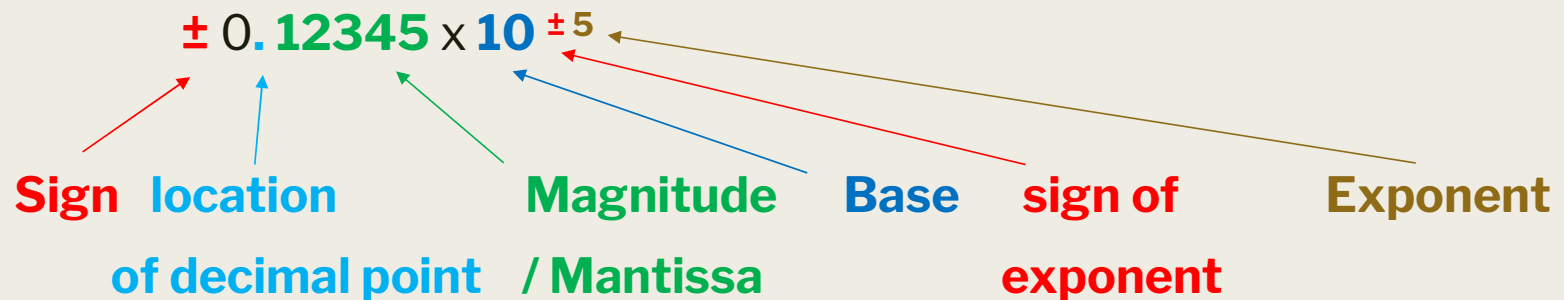
- ❑ Allows a range of overly large & very small numbers to be represented with a few digits.

## 2. Floating-point Num Representation & Standard

- **Exponential notation**

- ❑ A number is represented by the combination of 6 specifications.

- ❑ E.g.: To represent  $12345_{10}$



## 2. Floating-point Num Representation & Standard

- **Floating Point Representation**

- ☐ Floating point numbers will be stored and manipulated in the computer using standard predefined format, usually in 8 bits basis.
- ☐ There are 2 key format applied: **SEEMMMMM** & **IEEE754** notation
- ☐ The base of exponent and location of the binary point are standardize as part of the format. Therefore, they are not required to be stored at all.

## 2. Floating-point Num Representation & Standard

- **Floating Point Representation - SEEMMMMM**

- In **SEEMMMMM** format,

- ❖ 1 digit for sign
    - ❖ 2 digits for exponent
    - ❖ 5 digits for mantissa, the decimal point location is assumed to be located at the beginning of mantissa.

- **Excess-N** notation

- ❖ **N** is the chosen middle value
    - ❖ E.g.: Excess-50 allow a magnitude ranges as follow:

$$0.00001 \times 10^{-50} < \text{Number} < 0.99999 \times 10^{49}$$

| Excess   | -48 | -49 | 50 | 51 | 52 |
|----------|-----|-----|----|----|----|
| Exponent | -2  | -1  | 0  | 1  | 2  |

## 2. Floating-point Num Representation & Standard

### ▪ Floating Point Representation - SEEMMMMM

#### ❑ Conversion:

❖ The base is 10.

❖ The implied decimal point is at the beginning of the mantissa.

❖ Assume that:

✓ Excess-50 is applied.

✓ A 0 represents positive and 5 represents negative.

$$05324567 = 0.24567 \times 10^3 = 245.67$$

$$54810000 = -0.10000 \times 10^{-2} = -0.00100000$$

$$55555555 = -0.5555 \times 10^5 = -55555$$

| Excess   | -48 | -49 | 50 | 51 | 52 |
|----------|-----|-----|----|----|----|
| Exponent | -2  | -1  | 0  | 1  | 2  |

## 2. Floating-point Num Representation & Standard

- **Floating Point Representation - SEEMMMMM**

- ❑ **Normalization**

- ❖ Shift number left by increasing the exponent until leading zeros are eliminated.

- ❑ **Steps:** (Convert decimal number into SEEMMMMM format)

- 1. Provide number with exponent (0 if not yet specified)
    2. Increase / decrease exponent to shift decimal point to proper position.
    3. Decrease exponent to eliminate leading zeros on mantissa.
    4. Correct precision by adding 0's or discarding / rounding least significant digits



## 2. Floating-point Num Representation & Standard

- **Floating Point Representation - SEEMMMMM**

- ❑ **Normalization and formatting**

- ❑ E.g.:

- ❖ Given 246.8035, normalize it and represent it in SEEMMMMM format.

- ❑ Steps:

1. Add exponent :  $246.8035 \times 10^0$
2. Position decimal point :  $0.2468035 \times 10^3$
3. Already normalized, no adjustment is required.
4. Trim mantissa to 5 digits :  $0.24680 \times 10^3$
5. Convert the number : 05324680

- ❖ **Assume** that:
      - ✓ Excess-50 is applied.
      - ✓ A 0 represents positive and 5 represents negative.

## 2. Floating-point Num Representation & Standard

- **Floating Point Representation - SEEMMMMM**

- ❑ **Normalization and formatting**

- ❑ E.g.:

- ❖ Given  $-1255 \times 10^{-3}$ , normalize it and represent it in SEEMMMMM format.

- ❑ Steps:

1. The number is already in exponential form.
2. Position decimal point :  $-0.1255 \times 10^1$
3. Already normalized, no adjustment is required.
4. Trim mantissa to 5 digits :  $-0.12550 \times 10^1$
5. Convert the number : 55112550

- ❖ **Assume** that:

- ✓ Excess-50 is applied.
      - ✓ A 0 represents positive and 5 represents negative.

## 2. Floating-point Num Representation & Standard

- **Floating Point Representation - SEEMMMMM**

- ❑ **Addition and subtraction**

- ❖ Exponent and mantissa treated separately.
    - ❖ Exponents of numbers must agree.
      - ✓ Align decimal points
      - ✓ Least significant digits may be lost
    - ❖ Overflow of the most significant digit may occurs.
    - ❖ Number must be shifted right and the exponent incremented to accommodate overflow.

## 2. Floating-point Num Representation & Standard

- **Floating Point Representation - SEEMMMMM**

- **Addition and subtraction**

- ❖ E.g.: Add the 2 floating-point numbers.

$$\begin{array}{r} 05199520 \\ + 04967850 \\ \hline \end{array}$$

Align exponents

$$= 05199520$$

By adding 2 zeros in front to mantissa

$$= \underline{0510067850} +$$

Add mantissa, (1) indicates a carry

$$= (1)0019850$$

Carry requires right shift of exponent

$$= 0\ 52\ 10019(850)$$

Round

$$= 0\ 52\ 10020$$

- ❖ **Assume** that:
  - ✓ Excess-50 is applied.
  - ✓ A 0 represents positive and 5 represents negative.

## 2. Floating-point Num Representation & Standard

- **Floating Point Representation - SEEMMMMM**

- **Addition and subtraction**

- ❖ E.g.: **Check result.**

$$\begin{array}{r} 05199520 \\ + 04967850 \\ \hline \end{array}$$

$$05199520 = 0.99520 \times 10^1$$

$$04967850 = 0.67850 \times 10^{-1}$$

In sign-magnitude form

$$= 9.520$$

$$= 0.067850 + 10.019850$$

$$= 0.1001985 \times 10^2$$

- ❖ **Assume** that:
  - ✓ Excess-50 is applied.
  - ✓ A 0 represents positive and 5 represents negative.

## 2. Floating-point Num Representation & Standard

- **Floating Point Representation - SEEMMMMM**

- ❑ **Multiply and divide**

- ❖ Mantissa: Multiplied or divided
    - ❖ Exponent: Added or subtracted and adjusted excess value since added twice

- ❖ E.g.:

- ✓ Assume that two number with exponent 3, each representing 53.
    - ✓ Adding the two exponent:  $53 + 53 = 106$
    - ✓ Since 50 (excess-50) is added twice, subtract:  $106 - 50 = 56$

- ❖ Normalization necessary to:

- ✓ Restore location of decimal point
    - ✓ Maintain precision of the result.

- ❖ **Assume** that:
  - ✓ Excess-50 is applied.
  - ✓ A 0 represents positive and 5 represents negative.

## 2. Floating-point Num Representation & Standard

- **Floating Point Representation - SEEMMMMM**

- **Multiply and divide**

- ❖ E.g.: multiply two floating pint numbers

05220000  
x 04712500

Add exponent, subtract offset =  $75 + 47 - 50 = 49$

Multiply mantissa =  $0.20000 \times 0.12500$   
= 0.025000000

Normalized the result = 04825000

- ❖ **Assume** that:
  - ✓ Excess-50 is applied.
  - ✓ A 0 represents positive and 5 represents negative.

## 2. Floating-point Num Representation & Standard

- **Floating Point Representation - SEEMMMMM**

- **Multiply and divide**

- ❖ E.g.: **Check result**

$$05220000 = 0.20000 \times 10^2 = 20$$

$$04712500 = 0.12500 \times 10^{-3} = 0.000125$$

$$\begin{aligned} \text{Multiply} &= 20 \times 0.000125 \\ &= 0.025000000 \times 10^{-1} \end{aligned}$$

$$\text{Normalizing and rounding} = 0.25000 \times 10^{-2}$$

- ❖ **Assume** that:
  - ✓ Excess-50 is applied.
  - ✓ A 0 represents positive and 5 represents negative.



## 2. Floating-point Num Representation & Standard

- **Floating Point Representation – IEEE754**

## ❑ Typical floating point format in computer

- ❖ Consists of 32 bits, divided into:

✓ 1 bit of sign

✓ 8 bits of exponent, excess-127 notation, base 2

✓ 23 bits of mantissa

bit 0 1 8 9

31



## 2. Floating-point Num Representation & Standard

- **Floating Point Representation – IEEE754**

- ❑ Normalized numbers must always start with a 1, the leading bit is not stored, but is instead implied.
- ❑ This bit is located to the left of the implied binary position. So, numbers are normalized to the form 1.MMMMMMM...

| Precision    | Single (32-bit)         | Double (64-bit)           |
|--------------|-------------------------|---------------------------|
| Sign         | 1 bit                   | 1 bit                     |
| Exponent     | 8 bits                  | 11 bits                   |
| Implied base | 2                       | 2                         |
| Range        | $2^{-126}$ to $2^{127}$ | $2^{-1022}$ to $2^{1023}$ |
| mantissa     | 23                      | 52                        |
| Notation     | Excess-127              | Excess-1023               |

## 2. Floating-point Num Representation & Standard

- **Floating Point Representation – IEEE754**

- ❑ E.g.: Convert  $253.75_{10}$  to binary floating point form.

$$253.75_{10}$$

$$1111\ 1101.11_2$$

$$\begin{aligned}\text{Therefore, } 253.11_{10} &= 1111\ 1101.11_2 \\ &= 1.111\ 110111 \times 2^{+7}\end{aligned}$$

|   |          |               |
|---|----------|---------------|
| 0 | 10000110 | 11111011100.. |
|---|----------|---------------|

|      |          |          |
|------|----------|----------|
| Sign | Exponent | Mantissa |
| 1    | 8        | 23 bits  |

## 2. Floating-point Num Representation & Standard

- **Floating Point Representation - Conversion**

- ❑ E.g.: Convert of fraction number

- ❑  $16.5_{10} \rightarrow \text{base } 2$

|       |       |       |          |          |          |
|-------|-------|-------|----------|----------|----------|
| $2^2$ | $2^1$ | $2^0$ | $2^{-1}$ | $2^{-2}$ | $2^{-3}$ |
| 4     | 2     | 1     | 1/2      | 1/4      | 1/8      |

- ❑  $16.5_{10} = 16_{10} + 0.5_{10}$   
     $= 10000_2 + 0.1_2$   
     $= 10000.1_2$

## 2. Floating-point Num Representation & Standard

- **Floating Point Representation - Conversion**

- ❑ E.g.: Convert of fraction number

- ❑  $16.5_{10} \rightarrow \text{base } 2$

|       |       |       |          |          |          |
|-------|-------|-------|----------|----------|----------|
| $2^2$ | $2^1$ | $2^0$ | $2^{-1}$ | $2^{-2}$ | $2^{-3}$ |
| 4     | 2     | 1     | 1/2      | 1/4      | 1/8      |

- ❑  $16.5_{10} = 16_{10} + 0.5_{10}$   
 $= 10000_2 + 0.1_2$   
 $= 10000.1_2$

**OR,**

$0.5 \times 2 = \mathbf{1.00}$



(Reading the integer only, in top-down direction)

## 2. Floating-point Num Representation & Standard

**X 2 because it need to  
change to base 2.**

Since it is multiplied using  
fractional numbers,  
therefore it must add **0.** in  
front of the answer.

$0.110101_2$



### ▪ **Floating Point Representation - Conversion**

❑ E.g.: Convert of fraction number

❑  $0.828125_{10} \rightarrow \text{base 2}$

$$0.828125 \times 2 = \mathbf{1.656250}$$

$$0.\mathbf{656250} \times 2 = \mathbf{1.312500}$$

$$0.\mathbf{312500} \times 2 = \mathbf{0.625000}$$

$$0.625000 \times 2 = \mathbf{1.250000}$$

$$0.250000 \times 2 = \mathbf{0.500000}$$

$$0.500000 \times 2 = \mathbf{1.000000}$$

Therefore,  $0.828125_{10} = 0.110101_2$

## 2. Floating-point Num Representation & Standard

- **Floating Point Representation - Conversion**

- ❑ E.g.: Convert of fraction number
- ❑  $0.828125_{10} \rightarrow$  base 2
- ❑ To validate it:

|         | $2^0$ | $2^{-1}$ | $2^{-2}$ | $2^{-3}$ | $2^{-4}$ | $2^{-5}$  | $2^{-6}$ |
|---------|-------|----------|----------|----------|----------|-----------|----------|
| Decimal | 0     | $1/2$    | $1/4$    | $1/8$    | $1/16$   | $1/32$    | $1/64$   |
|         | 0     | 0.5      | 0.25     | 0.125    | 0.0625   | $0.03125$ | 0.015625 |
| Binary  | 0     | 1        | 1        | 0        | 1        | 0         | 1        |

Therefore,  $0.828125_{10} = 0.110101_2$

## 2. Floating-point Num Representation & Standard

- **Floating Point Representation - Conversion**

- ❑ E.g.: Convert of fraction number

- ❑  $11.110011_2 \rightarrow \text{base } 10$

| Binary          | 1     | 1     | . | 1        | 1        | 0        | 0        | 1        | 1        |
|-----------------|-------|-------|---|----------|----------|----------|----------|----------|----------|
| Decimal         | $2^1$ | $2^0$ | . | $2^{-1}$ | $2^{-2}$ | $2^{-3}$ | $2^{-4}$ | $2^{-5}$ | $2^{-6}$ |
|                 | 2     | 1     | . | 1/2      | 1/4      | 1/8      | 1/16     | 1/32     | 1/64     |
|                 | 2     | 1     | . | 0.5      | 0.25     | 0.125    | 0.0625   | 0.03125  | 0.015625 |
|                 | +     | +     | + | +        | +        | +        | +        | +        |          |
| $3.796875_{10}$ |       |       |   |          |          |          |          |          |          |

Therefore,  $11.110011_2 = 3.796875_{10}$



## 2. Floating-point Num Representation & Standard

- **Floating Point Representation - Conversion**

- ❑ E.g.: Convert of fraction number

- ❑  $10011.10111_2 \rightarrow \text{base}_{16}$

$10011.10111_2$

0001 0011 . 1001 1000<sub>2</sub>

1

3.

9

8<sub>16</sub>

Additional Zero added

## 2. Floating-point Num Representation & Standard

- **Floating Point Representation - Conversion**

- ❑ E.g.: Convert of fraction number

- ❑  $39.B8_{16} \rightarrow \text{base } 10$

- $$= (3 \times 16^1) + (9 \times 16^0) + (.) + (B \times 16^{-1}) + (8 \times 16^{-2})$$

- $$= 48_{10} + 9_{10} + (.) 0.6875_{10} + 0.03125_{10}$$

- $$= 57.71875_{10}$$

## 2. Floating-point Num Representation & Standard

- **Floating Point Representation - Conversion**

- ❑ E.g.: Convert of fraction number

- ❑  $4F5.09_{16} \rightarrow \text{base}_2$

4 F 5 . 0 9<sub>16</sub>

0100 1111 0101 . 0000 1001<sub>2</sub>

Additional Zero added

# **3. Bitwise Logical Operations**

# 3. Bitwise logical operations

- Bitwise logical operations are including:

| Logical operation                                  | Symbol                            | Example       |                        |
|--|-----------------------------------|---------------|------------------------|
| AND= Yield TRUE if both operands are TRUE          | $(\bullet)$ , $(\wedge)$          | $A \bullet B$ | $A \wedge B$           |
| OR = Yield TRUE if either/both operand is/are TRUE | $(+)$ , $(\vee)$                  | $A + B$       | $A \vee B$             |
| NOT = Inverts the value of its operand             | $(')$ , $(\neg)$                  | $A'$          | $\neg A$               |
| XOR = exclusive disjunction                        | $(\oplus)$ , $(\underline{\vee})$ | $A \oplus B$  | $A \underline{\vee} B$ |

### 3. Bitwise logical operations

- Bitwise Operator precedence

| Expression        | Order of Operations |
|-------------------|---------------------|
| $A + B'$          | NOT, then OR        |
| $(A + B)'$        | OR, then NOT        |
| $A + (B \cdot C)$ | AND, then OR        |

- Relative order of precedence: **NOT** > **XOR** > **AND** > **OR**

# 3. Bitwise logical operations

- Bitwise logical operations
- E.g.: **AND**

|     |   |   |   |   |
|-----|---|---|---|---|
|     | 1 | 1 | 0 | 0 |
| AND | 1 | 0 | 1 | 0 |
|     | 0 | 0 | 0 | 0 |

- Bitwise logical operations
- E.g.: **OR**

|    |   |   |   |   |
|----|---|---|---|---|
|    | 1 | 1 | 0 | 0 |
| OR | 1 | 0 | 1 | 0 |
|    | 1 | 1 | 1 | 0 |

# 3. Bitwise logical operations

- Bitwise logical operations
- E.g.: **NOT**

|    |   |   |
|----|---|---|
| OR | 1 | 0 |
|    | 0 | 1 |

- Bitwise logical operations
- E.g.: **XOR**

|     |   |   |   |   |
|-----|---|---|---|---|
|     | 1 | 1 | 0 | 0 |
| XOR | 1 | 0 | 1 | 0 |
|     | 0 | 1 | 1 | 0 |



# **Chapter Review**

# Chapter Review

## 1. Signed and unsigned number representation

- ❑ Signed: Positive only
- ❑ Unsigned: Positive (0) /negative (1)

## 2. Floating-point Num Representation & Standard

- ❑ Sign-magnitude notation  
 $\pm 0.12345 \times 10^{\pm 5}$
- ❑ SEEMMMMM notation  
Excess-50, 1=+ve, 2=-ve  
25512345

- ❑ IEEE754 notation
  - ❖ Single precision (1:8:23)
  - ❖ Double precision (1:11:1023)
- ❑ Floating-point number conversion
  - ❖ B / O / H  $\rightarrow$  D:  $2/8/16^{-n}$
  - ❖ D  $\rightarrow$  B / O / H: Mantissa x base
  - ❖ B  $\rightarrow$  O / H: Collapse
  - ❖ H / O  $\rightarrow$  B: Expand

## 3. Bitwise logical operations

| Operator | Symbol                               | Example       |                        |
|----------|--------------------------------------|---------------|------------------------|
| AND      | ( $\bullet$ ), ( $\wedge$ )          | $A \bullet B$ | $A \wedge B$           |
| OR       | ( $+$ ), ( $\vee$ )                  | $A + B$       | $A \vee B$             |
| NOT      | ( $'$ ), ( $\neg$ )                  | $A'$          | $\neg A$               |
| XOR      | ( $\oplus$ ), ( $\underline{\vee}$ ) | $A \oplus B$  | $A \underline{\vee} B$ |