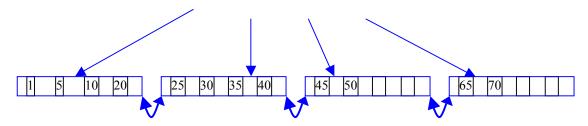
EECS 484: Final Sample Exam #1

Nam	Please <i>print</i> your name and unic ne:	name below.	
Unio	qname:		
2. 3. 3. 3. 4. 4. Honor C	You have 80 minutes to complete The questions vary widely in the question at the expense of an easone of the questions on this expour explanation, in addition to At the end of the exam, sign the Code Pledge: I have neither given	xam require you to provide an explanation. You the correctness of your answer.	avoid getting stuck on a hard will be graded on the clarity o
the Hono	Signature:		
			(For Instructor's Use
		Q1: Indexes (20 points)	
		Q2: Query Evaluation. (25 points)	
		Q3: Transactions & Recovery (20 points)	
		TOTAL	

Question 1: (20 points) Indexes

25	42	65		



(a) (4 points) Consider the B+-tree shown above. Draw the tree after inserting the key value 32*. Assume the convention that the left pointer points to values that are strictly less than the key value. Assume also that each node in this B+-tree can hold at most four data entries (i.e., the order of the tree is 2).

(b) (4 points) Using your answer from part (a), draw the tree after additionally inserting key value 15*.

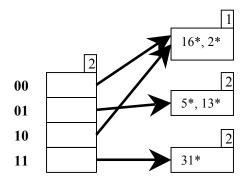
(c) (4 points) Consider a relation R(a,b,c,d) containing a million records. Suppose that R is stored in a heap file with unclustered indexes. The indexes use Alternative 2, where data entries are (value, rid) pairs. Suppose also that the records in R are randomly ordered. Assume that a is a candidate key for R, with values in the range 0 to 999,999. For each of the following queries, choose the approach that would likely require the fewest I/Os to complete the query. The approaches are as follows:

- i) Scan through the entire heap file containing R.
- ii) Use a B+-tree index on R.a.
- iii) Use a hash index on R.a.

Find all tuples in R. _____ Find all tuples in R such that a < 50. _____ Find all tuples in R such that a = 50. _____ Find all tuples in R such that a > 50 and a < 100. _____

(d) (8 points) Consider the following extendible hash index. Assume that the hash function is $h(K) = K \mod 32$, and that at most 2 data entries fit on a page.

Re-draw the index to show how it would appear after inserting the key values 8* and 22* (in that order).



Question 2: (25 points) Query Evaluation & Optimization

a) (15 points) Consider the join R $\bowtie_{R.a = S.b} S$, and suppose you are given the following information about the relations to be joined. For all of the following questions, assume that the cost metric is the total number of page I/Os and that the cost of writing out the result should be uniformly ignored. For all questions, show your work to receive full credit.

Relation R contains 10,000 tuples and has 10 tuples per page Relation S contains 2,000 tuples and has 10 tuples per page Attribute S.b is the primary key for S Both relations are stored as unsorted heap files Neither relation has any indexes 52 buffer pages are available

What is the cost of joining R and S using a page-oriented nested loops join?

What is the cost of joining R and S using Grace hash join?

Suppose you chose to join R and S using a <u>block-oriented nested loops</u> join, and that you want to do this using \leq 1200 I/Os. What is the minimum number of buffer pages necessary? Explain your answer.

b) (3 points) Consider a relational schema including A(a1,a2,a3) and B(b1, b2, b3), and the following relational algebra expression over these relations. To which of the given expressions is it equivalent? (Select all that apply.)

$$\pi_{a3} \; \big(\sigma_{b2\,=\,5} \; \big(A \; \bowtie_{a1\,=\,b1} \; B\big)\big)$$

i.
$$\pi_{a3} \left(\sigma_{b2} = 5 \left(\sigma_{a1} = b_1 \left(\sigma_{b2} = 5 \vee b_2 = 10 \right) \left(A \times B \right) \right) \right)$$

ii.
$$\pi_{a3} \left(\sigma_{b2=5} \left(\pi_{a3} \left(A \bowtie_{a1=b1} B \right) \right) \right)$$

iii.
$$\pi_{a3} ((\sigma_{b2=5} B) \bowtie_{a1=b1} A))$$

Answer:

c) (4 points) Consider the four-way join $R \bowtie_{R.a = S.b} S \bowtie_{S.c = T.d} T \bowtie_{S.e = V.f} V$, where each relation (R,S,T,V) is the same size, and the result size for each two-way join is the same. Consider also the following two join plans (one left-deep and one bushy). Which plan would you choose? Explain your answer in terms of what the query engine must do for each.



d) (3 points) *Briefly* explain what role the system catalogs play in query optimization.

Ou	estio	n 3: (20	points) T	ransactions & Recovery						
				roduced by following the strict two-phase locking (Strict 2PL) protocol is guaranteed to have						
	which of the following properties? (Select all that apply.)									
	i. The schedule must be serializable.									
	ii.	The sc	hedule mu	st be recoverable.						
	iii.	The sc	hedule mu	st avoid cascading aborts.						
	iv.	The sc	hedule mu	st not result in a deadlock.						
	Answer:									
1. \	(1	intal Ca		fellowing (incomplete) schoolule involving transcriptions T1 T2 and T2. December schoolule						
				following (incomplete) schedule involving transactions T1, T2, and T3. Does this schedule						
T1	avoic		ng aborts?	Explain.						
R(X)	T2	Т3	-						
W(X	_			-						
W (2	1)		D(V)	-						
			R(X) W(X)	-						
		D(V)	W(A)	-						
		R(Y)		-						
		W(Y)								
		_] Yes	□ No						
_	1 .	. [l i es	□ No						
Exp	lanat	ion:								
c) (4	l noir	nts) Con	sider the f	ollowing (complete) schedule involving transactions T1, T2, and T3, and consider a system						
				phase locking (2PL). Could this schedule possibly arise given this locking protocol? Explain.						

T1	T2	T3
R(X)		
W(X)		
	R(Y)	
		R(X)
	W(Y)	
		W(X)
		Commit
Commit		
	Abort	

		Yes			No
	_			_	

Explanation:

	,
d) (3 points) <i>Briefly</i> explain the difference between a system crash and a media failure.	
d) (5 points) Consider a system that implements the ARIES logging and recovery protocol. If the system fails repeated during recovery, what is the maximum number of log records that can be written (as a function of the number of update and other log records written before the crash) before restart completes successfully? Explain.	.ly