

## Problem 1

1. **Inductive hypothesis:** At the end of each  $i^{\text{th}}$  loop, the variable *out* became the sum of all elements of  $A[1 \dots i]$ .

**Base case** ( $i = 0$ ): The induction hypothesis is trivially true.

**Induction step:** Let the variable  $out_i$  be the version of the variable at the end of  $i^{\text{th}}$  iteration. By the inductive hypothesis,  $out_i$  became the sum of all the elements in  $A[1 \dots i]$ . At the end of the loop  $(i + 1)^{\text{th}}$ , the variable  $out_{i+1}$  became  $out_i + A[i + 1]$ , which is the sum of elements in  $A[1 \dots i + 1]$ .

**Wrap up:** When  $i = n$ , the inductive hypothesis implies that the returned variable *out* became the sum of all elements in  $A$ .

2. **Inductive hypothesis:** For every complete binary tree  $T$  with height  $h$  and number  $x$ , the function `search()` returns "yes" if the tree contains  $x$ , and "no" otherwise.

**Base case** ( $h = 0$ ): The induction hypothesis is trivially true.

**Induction step:** Assuming the inductive hypothesis for a non-empty input  $T_h$  with height  $h$ , consider  $T_{h+1}$  with height  $h + 1$ . The function starts with the root  $T_{h+1}.val$ .

- (a)  $T_h$  is not empty, so  $T_{h+1}$  is not empty — the function won't return "no" at the beginning.
- (b) If  $T_{h+1}.val = x$ , the function returns "yes".
- (c) If  $T_{h+1}.val > x$ , the function calls `search( $T_{h+1}.l, x$ )` where  $\text{height}(T_{h+1}.l) = h$ . By the inductive hypothesis, it returns "yes" iff  $T_{h+1}.l$  contains  $x$ , and "no" otherwise.
- (d) If  $T_{h+1}.val < x$ , the function calls `search( $T_{h+1}.r, x$ )` where  $\text{height}(T_{h+1}.r) = h$ . Similarly, it returns "yes" iff  $T_{h+1}.r$  contains  $x$ , and "no" otherwise.

**Wrap up:** The inductive hypothesis matches exactly the statement we set out to prove.