Problem 2

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Algorithm:
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\begin{array}{l} i \leftarrow 1 \\ \text{while houses[i]} == \text{``keep going'':} \\ i = i * 2 \\ \text{return bsearch}(\frac{i}{2}, \, \text{i}) \end{array}
```

Runtime:

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The bounds of h will always be 2^k \le h \le 2^{k+1} where k = \lfloor log_2(h) \rfloor.
The range binary searched will be (2^{k+1} - 2^k), so the total runtime would be... log_2(2^{k+1} - 2^k) \le log(k) + C \to O(log(h)).
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This means the algorithm runs in O(log(h)) time.

Correctness:

Hypothesis: Returns the index to where houses will return "Party's here!" without making changes to houses.

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Base Case: h = 1, the alg. returns 1. \checkmark Induction Step:
Assume: \forall i < n, it returns i when i = h.
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The function considers some range [i, 2i] where $i = \lfloor log_2(h) \rfloor$. We also know this will return the correct result as, we know h >= i and h <= 2i. Meaning we have not throw out the correct h and do not have to consider any previous numbers.

This proves that our algorithm with always find the correct h according to our hypothesis.