Algorithmen der Bioinformatik I WS 2017/2018

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Most important implementation: CLUSTAL W

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CLUSTAL W: improving the sensitivity of progressive multiple sequence alignment through sequence weighting, position-specific gap penalties and weight matrix choice

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Received July 12, 1994; Revised and Accepted September 23, 1994

45.933 citations in the scientific literature (Web of Science)





Figure: Des Higgins, Dublin (http://businessetc.thejournal.ie/)

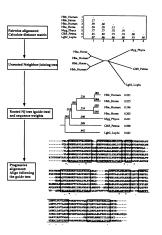


Figure: Progressive Alignment in CLUSTAL W (Thompson et al, 1994)



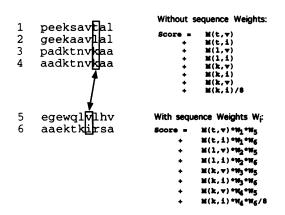


Figure: Sequence weighting in profile alignment (Thompson et al, 1994)



Gap penalties in CLUSTAL W

- Use affine-linear gap penalties with gap opening penalty (GOP) and gap extension penalty (GEP)
- Initial values of GOP and GEP specified by user
- During progressive alignment, GOP and GEP modified depending on
 - Substitution matrix
 - Similarity between sequences
 - Length of sequences
 - Differences in sequence lengths
 - Local sequence composition (hydrophilic or hydrophobic amino acid residues)
 - Existing gaps
 - Position in sequence: end gaps get lower penalty



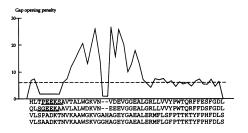
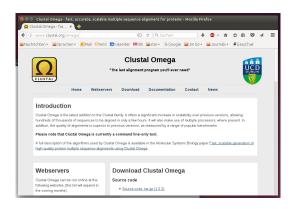


Figure: Variable gap penatlies depending on sequence composition, existing gaps etc. (Thompson *et al*, 1994)



CLUSTAL Ω



Latest version of *CLUSTAL*: align up to 100.000 sequences. Based on fast clustering algorithm



1792–1797 Nucleic Acids Research, 2004, Vol. 32, No. 5 DOI: 10.1093/nar/ekh340

MUSCLE: multiple sequence alignment with high accuracy and high throughput

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ABSTRACT

We describe MUSCLE, a new computer program for creating multiple alignments of protein sequences. Elements of the algorithm include fast distance estimation using kmer counting, progressive alignment using a new profile function we call the log-expectation score, and refinement using tree-dependent restricted partitioning. The speed and

variant on this strategy is used by T-Coffee (5), which aligns profiles by optimizing a score derived from local and global alignments of all pairs of input sequences. Misalignments by progressive methods are sometimes readily apparent (Fig. 1), motivating further processing (refinement). For a recurreview of multiple alignment methods, see Notredame (6). Here we describe MUSCLE (multiple sequence comparison by log-expectation), a new computer program for multiple protein sequence alignment.



Most time-consuming step in progressive alignment: Calculation of *guide tree*.

To calculate distances for N sequences of length ℓ :

 $O(N^2)$ pairwise alignments calculated, total complexity

$$O(N^2 \cdot \ell^2)$$

Calculating guide tree with Neighbour-Joining takes

$$O(N^3)$$

time



Strategy of MUSCLE (1):

- Calculate pairwise similarities of sequences using k-mer occurrences (k-mer = word of length k)
- Turn similarity values into distance values $d_{i,j}$
- Calculate guide tree using UPGMA
- Progressive alignment



To calculate distances $d_{i,j}$ for sequences X_i, X_j of length ℓ_i, ℓ_j using k-mer frequencies:

Define for k-mer τ :

- $n_i(\tau)$ = frequency of τ in X_i
- $n_j(\tau)$ = frequency of τ in X_j

and 'k-mer similarity' as

$$F_{i,j} = \frac{\sum_{\tau} \min \left\{ n_i(\tau), n_j(\tau) \right\}}{\min \left\{ \ell_i, \ell_j \right\} - k + 1}$$

Distance between X_i and X_i defined as

$$d_{i,j} = 1 - F_{i,j}$$



In general: *NJ* produces phylogenetically more accurate trees than *LIPGMA* since it can deal with different mutation rates

But: test runs showed that *UPGMA* may be superior to *NJ* to construct guide trees.

Possible reason: Alignment most accurate if sequences closely related



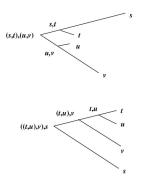


Figure: Wrong branching order in UPGMA tree (Edgar, 2004)



Strategy of MUSCLE (2):

- Iterative improvement of alignment:
 - Use Kimura distances to calculate new guide tree (Kimura distance estimates number of substitutions based on observed mismatches in alignment)
 - ► For nodes that are different in new tree, re-calculate progressive alignment



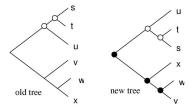


Figure: Progressive alignment newly calculated only if branching order is different: black nodes in new tree (Edgar, 2004)



Strategy of MUSCLE (3):

- Refinement of MSA
 - Partition sequences into two groups
 - Re-align profiles from the two groups
 - Accept new alignment, if score is improved



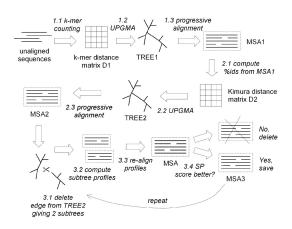


Figure: Multiple alignment with MUSCLE (Edgar, 2004)



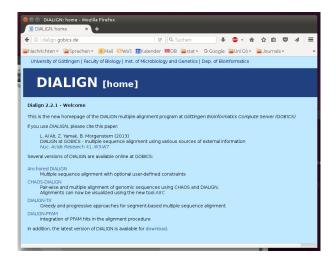


Figure: DIALIGN at GOBICS



Idea: combine global and local alignment.

- Search for local pairwise similarities
- Compose alignments as consistent set of local pairwise alignments
- Ignore non-related parts of sequences
- No gap penalty

Local pairwise gap-free alignment called fragment



```
S_1 Y I A V L F A E D S_2 L A C V I F G S S_3 P W D D V T F D A E
```

Figure: MSA composed of fragments, i.e. gap-free pairwise alignments



```
S_1 Y I A V L F A E D S_2 L A C V I F G S S_3 P W D D V T F D A E
```

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Figure: MSA composed of fragments, i.e. gap-free pairwise alignments



```
S_1 Y I A V L F A E D

S_2 L A C V I F G S

S_3 P W D D V T F D A E
```

Figure: MSA composed of *fragments*, *i.e.* gap-free pairwise alignments



```
S_1 Y I A - V L F - A E D

S_2 - L A C V I F - G S -

S_3 P W D D V T F D A E -
```

Figure: MSA as consistent set of fragments



```
S_1 y I A - V L F - A E d

S_2 - L A c V I F - G s -

S_3 p w d d V T F d A E -
```

Figure: Resulting MSA: non-aligned positions shown in lower case



• Weight score of fragment f:

$$w(f) = -\log Pr(f)$$

Pr(f) = probability of occurrence of fragment f in random sequences of same length

- Score of alignment: sum of weight scores of fragments in alignment - no gap penalty!
- Optimization problem: Find consistent set of fragments with max total weight, i.e. set of fragments that fits into one single multiple alignment.

Algorithm:

Calculate all pairwise optimal alignments (= chains of fragments)

 $\mathcal{M}_1 :=$ set of fragments from pairwise optimal alignments

- Calculate overlap weights for fragments from M₁ depending on weights of overlapping fragments.
- Sort set M₁ of fragments according to 'overlap weights'
- ► Greedily select consistent subset M₂ of M₁
- Repeat iteratively, given consistency constraints imposed by M₂



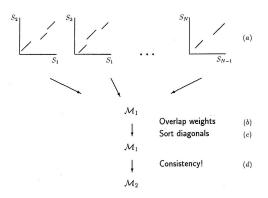
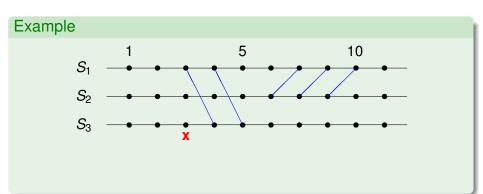


Figure: Greedy algorithm to calculate MSA, iteratively applied

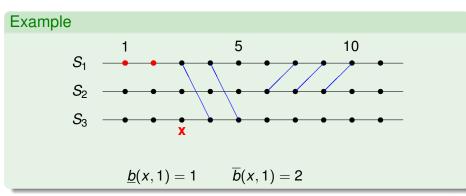


To decide if new fragment is consistent: use *consistency bounds* $\underline{b}(x, i)$ and $\overline{b}(x, i)$





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Results for DIALIGN:

- Better than other methods for *locally* related sequences
- Inferior on globally related sequences



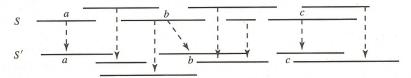
One-	dimen	sional c	haining	g prob	lem:			
		-			C. The Company of	-	 	
					•		 - -	
R				A				

Find a best chain of weighted intervals ('fragments')

For input set of n fragments, best chain can be found in $O(n \cdot \log n)$ time by *dynamic programming*



Two-dimensional chaining problem:



Find best chain of two-dimensional fragments

Also solvable in $O(n \cdot \log n)$ time.



In sequence alignment, number n of fragments can be large. Space-efficient algorithm possible:

- Go column-wise through DP matrix
- For *i*-th column, calculate arrays S[i] and L[i] with score and last fragment in longest chain up to (i, j) for each j.
- For new column i, calculate weight w(f) for each fragment f starting in i. Calculate total weight W(f) of optimal chain ending in f and its 'predecessor' P(f) using arrays S[i-1] and L[i-1]:

$$W(f) = w(f) + W[i-1, j-1]$$

 $P(f) = P[i-1, j-1]$



- For new fragment ending in column i', update list E[i'] of fragments ending in column i'
- After fragments starting in column i have been processed: calculate S[i] and L[i] from S[i-1], L[i-1] and E[i]. Delete S[i-1], L[i-1].
- Maintain fragment f* in which best chain so far ends
- Finally: start trace back at f*



Evaluation of Multiple Protein Alignment Software

```
.NLFVALYDfvasqdntlsitkGEKLRVLgynhn.....gE
1aboA
             kGVIYALWDyepqnddelpmkeGDCMTIIhrede.....deiE
1ycsB
1pht
             gYOYRALYDvkkereedidlhlGDILTVNkgslvalgfsdggearpeeiG
1 i hvA
              .NFRVYYRDsrd.....pvwkGPAKLLWkq.....eG
1vie
              .drvrkksga.....peG
1aboA
        36
             WCEAOt..knggGWVPSNYITPVN.....
1ycsB
             WWWAR1..ndkeGYVPRNLLGLYP.....
             WLNGYnettgerGDFPGTYVEYIGrkkisp
1pht
1ihvA
             AVVIQd..nsdiKVVPRRKAKIIRd.....
1vie
             YAVESeahpgsvOIYPVAALERIN.....
Key
alpha helix RED
beta strand GREEN
                                 BAIIBASE
core blocks UNDERSCORE
                                 Reference alignments
```

Figure: Reference alignment with core blocks from BAliBASE



Evaluation of Multiple Protein Alignment Software

Benchmark alignments contain reliable core blocks

For Evaluation:

- Sum-of-Pairs (SP) score: ration of correctly aligned pairs of positions in core blocks
- Total-column (TC) score: ration of correctly aligned columns in core blocks



Evaluation of Multiple Protein Alignment Software

Results:

- DIALIGN best method for local MSA
- Outperformed by other methods on weakly but globally related sequences

