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-CENG312-ASSIGNMENT 3 IP, ICMP, NAT, DHCP-

a.) When we look at the packet traffic "Fragmentation needed" ICMP messages are triggered in case amount of data that should be carried by datagram exceeds the maximum amount of data can be carried within the datagram.

b.) When endpoints receive the "Fragmentation needed" message endpoints can divide or in other words fragment datagram into smaller IPv4 packets. To identify fragements 16-bit field is used as an identification field and this identification field is copied to fragmented packet header so that on the destination side receiver will be able to reassemble the fragments that belong to a particular frame. In the packet traffic there is a problem related to fragmentation, because even though endpoint with IP address 62.177.254.141 has received ICMP message "Destination unreachable (Fragmentation needed)" three times it was not able to perform fragmentation. (Figure 1 and 2)

No. Time	Source	Destination	Protocol I	Length Calculated window size	e Info
1 2005-10-12 05:11:48,326590	62.177.254.141	62.177.254.1	DNS	78	Standard query 0x0e69 A scsc.msg.yahoo.com
2 2005-10-12 05:11:48,327329	100.100.100.10	0 62.177.254.141	ICMP	70	Destination unreachable (Fragmentation needed)
3 2005-10-12 05:11:49,327191	62.177.254.141	62.177.254.1	DNS	78	Standard query 0x0e69 A scsc.msg.yahoo.com
4 2005-10-12 05:11:50,327263	62.177.254.141	62.177.254.1	DNS	78	Standard query 0x0e69 A scsc.msg.yahoo.com
5 2005-10-12 05:11:52,327711	62.177.254.141	62.177.254.1	DNS	78	Standard query 0x0e69 A scsc.msg.yahoo.com
6 2005-10-12 05:11:54,329648	80:05:5d:7d:1		ARP	60	Who has 62.177.254.141? Tell 62.177.254.1
7 2005-10-12 05:11:54,329673		80:05:5d:7d:1	ARP	42	62.177.254.141 is at 08:00:46:f4:3a:09
8 2005-10-12 05:11:56,328401	62.177.254.141		DNS	78	Standard query 0x0e69 A scsc.msg.yahoo.com
9 2005-10-12 05:11:56,329135		0 62.177.254.141		70	Destination unreachable (Fragmentation needed)
10 2005-10-12 05:12:03,329401	62.177.254.141		DNS	78	Standard query 0xb16b A scsc.msg.yahoo.com
11 2005-10-12 05:12:03,330124		0 62.177.254.141		70	Destination unreachable (Fragmentation needed)
12 2005-10-12 05:12:04,329919	62.177.254.141		DNS	78	Standard query 0xb16b A scsc.msg.yahoo.com
13 2005-10-12 05:12:05,330114	62.177.254.141		DNS	78	Standard query 0xb16b A scsc.msg.yahoo.com
14 2005-10-12 05:12:07,330481	62.177.254.141		DNS	78	Standard query 0xb16b A scsc.msg.yahoo.com
15 2005-10-12 05:12:09,593433	80:05:5d:7d:1		ARP	60	Who has 62.177.254.141? Tell 62.177.254.1
16 2005-10-12 05:12:09,593453		80:05:5d:7d:1	ARP	42	62.177.254.141 is at 08:00:46:f4:3a:09
17 2005-10-12 05:12:11,331277	62.177.254.141		DNS	78	Standard query 0xb16b A scsc.msg.yahoo.com
18 2005-10-12 05:12:11,332056	100.100.100.10	0 62.177.254.141	ICMP	70	Destination unreachable (Fragmentation needed)
Frame 2: 70 bytes on wire (560		:1b:a0). Dst: So	nv f4:3a:09	(08:00:46:f4:3a:09)	
	b:a0 (80:05:5d:7d c: 100.100.100.100.10 bytes (5): 0x00 (DSCP: CSO Services Codepoir tion Notification) lot set Not set	0, Dst: 62.177.2 , ECN: Not-ECT) t: Default (0)	54.141	Also when we look	it we can t, normally in order to cion. This also is a problem

Figure 1: Fragmentation problem in the traffic

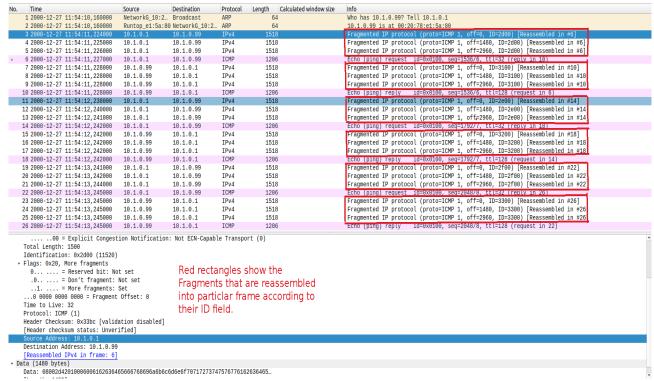


Figure 2: Reassembly operation on fragments is done according to id fields

<u>2.)</u>

- **a.)** In the packet capture Info column, there is an information saying that "Fragmented IP protocol" with fragment's offset number. (Figure 3)
- **b.)** Purpose of IP fragmentation is to divide large datagrams into smaller packets which does not fit into a single datagram previously. Later, fragments will be reassembled at the destination according to identification fields.
- **c.)** When we look at the line numbers 3,4, and 5 fragments with the ID=2d00 are all reassembled into number #6 frame. In the lines 7,8, and 9 fragments with ID=3100 are all reassembled into number #10 frame. In the lines 11,12, and 13 fragments with ID=2e00 are all reassembled into number #14 frame. In the lines 15,16, and 17 fragments with ID=3200 are all reassembled into number #18 frame. In the lines 19,20, and 21 fragments with ID=2f00 are all reassembled into number #22 (Figure 3 and Figure 4)

frame. In the lines 23,24, and 25 with ID=3300 are all reassembled into number #26.

c.) If we look at the line numbers 3,4, and 5, fragments which have ID=2d00 belong to

to frame with number #6. Each Fragment has a total length of 1500, so total length is 4500. We can also compute the original packet size by using offset values. First packet's offset is 0 and second packet's offset is 1480, so data length of first packet should be 1480. Second packet offset is 1480 and third packet offset is 2960, so data length of second packet is 1480. Third packet also has the data length 1480, so total data length is 4440 and also each fragment has 20 byte header field so total 60 bytes for headers. Indeed 4440+60 = 4500 is the original IP packet length which conforms to total length information in Wireshark bottom panel. (Figure 4 and Figure 6)

d.) For the fragments whose ID's are 2d00, 2e00, and 2f00 Time to live field is 32 whereas, fragments whose ID's are 3100,3200, and 3300 have Time to live field value 128. So TTL value is different for some fragments. Reason that's why TTL value for some fragments is longer than others can be attributed to importance of fragment, may be some fragments some fragments are given more importance due to role of data they carry. (Figure 5)

No.	Time	Source	Destination	Protocol	Length	Calculated window size	Info	
		11:54:10,160000 NetworkG_10:22:1		ARP	6		Who has 10.1.0.99? Tell 10.1.0.1	
		11:54:10,160000 Runtop_e1:5a:80	NetworkG_10:		6		10.1.0.99 is at 00:20:78:e1:5a:80	
		11:54:11,224000 10.1.0.1	10.1.0.99	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=0 ID=2d00) [Reassembled in #6]	
		11:54:11,225000 10.1.0.1	10.1.0.99	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=1480 ID=2d00) [Reassembled in #6]	
		11:54:11,226000 10.1.0.1	10.1.0.99	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=2960 ID=2d00) [Reassembled in #6]	
		11:54:11,227000 10.1.0.1	10.1.0.99	ICMP	120		Echo (ping) request id=0x0100, seq=1536/6, ttl=32 (reply in 10)	
		11:54:11,228000 10.1.0.99	10.1.0.1	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=0, ID=3100) [Reassembled in #10]	
		11:54:11,228000 10.1.0.99	10.1.0.1	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=1480, ID=3190) [Reassembled in #10]	
		11:54:11,228000 10.1.0.99	10.1.0.1	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=2960, ID=3100) [Reassembled in #10]	
		11:54:11,228000 10.1.0.99	10.1.0.1	ICMP	120		Echo (ping) reply id=0x0100, seq=1536/6, ttl=128 (request in 6)	
		11:54:12,238000 10.1.0.1	10.1.0.99	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=0, ID=2e00) [Reassembled in #14]	
		11:54:12,240000 10.1.0.1	10.1.0.99	IPv4	151	-	Fragmented IP protocol (proto=ICMP 1, off=1480, ID=2e00) [Reassembled in #14]	
		11:54:12,241000 10.1.0.1	10.1.0.99	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=2960, ID=2e00) [Reassembled in #14]	
		11:54:12,242000 10.1.0.1	10.1.0.99	ICMP	120		Echo (ping) request id=0x0100, seq=1792/7, ttl=32 (reply in 18)	
		11:54:12,242000 10.1.0.99	10.1.0.1	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=0, ID=3200) [Reassembled in #18]	
		11:54:12,242000 10.1.0.99	10.1.0.1	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=1480, ID=3200) [Reassembled in #18]	
		11:54:12,242000 10.1.0.99	10.1.0.1	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=2960, ID=3200) [Reassembled in #18]	
		11:54:12,242000 10.1.0.99	10.1.0.1	ICMP	120		Echo (ping) reply id=0x0100, seq=1792/7, ttl=128 (request in 14)	
		11:54:13,241000 10.1.0.1	10.1.0.99	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=0, ID=2f00) [Reassembled in #22]	
		11:54:13,242000 10.1.0.1	10.1.0.99	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=1480, ID=2f00) [Reassembled in #22]	
		11:54:13,244000 10.1.0.1	10.1.0.99	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=2960, ID=2f00) [Reassembled in #22]	
		11:54:13,245000 10.1.0.1	10.1.0.99	ICMP	120		Echo (ping) request id=0x0100, seq=2048/8, ttl=32 (reply in 26)	
		11:54:13,245000 10.1.0.99	10.1.0.1	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=0, ID=3300) [Reassembled in #26]	
		11:54:13,245000 10.1.0.99	10.1.0.1	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=1480, ID=3300) [Reassembled in #26]	
		11:54:13,245000 10.1.0.99	10.1.0.1	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=2960, ID=3300) [Reassembled in #26]	
	26 2000-12-27	11:54:13,245000 10.1.0.99	10.1.0.1	ICMP	120		Echo (ping) reply id=0x0100, seq=2048/8, ttl=128 (request in 22)	
*	Flags: 0x20,	More fragments						
		= Reserved bit: Not set	Fach fra	gment with	1D=2	d00		
		= Don't fragment: Not set		_				
		= More fragments: Set	nas the	data part l	ength .	1480		
		00 1000 = Fragment Offset: 1480	hytes S	the total	data n	art length		
	Time to Live							
	Protocol: IC		of these	three frag	ments	is 4440.		
		sum: 0x3303 [validation disabled]	Alen eac	h fragmen	t had a	in haadar		
		ksum status: Unverified]		•				
	Source Addre		field of 2	20 bytes, s	o in toi	tal 4440 + 3*20	0 = 4500	
		Address: 10.1.0.99	chould b	a tha land	th of o	riginal ID packet	.	
-		IPv4 in frame: 6]	Should b	e the leng	111 01 0	riginal IP packet	Li	
	ata (1480 <mark>byte</mark>							
		5465666768696a6b6c6d6e6f707172737	47576776162636	465666768696a6	6b6c6d			
	[Length: 148	9]						,

Figure 3: Calculating original packet length with help of offset information

No.	Time	Source	Destination	Protocol	Length	Calculated window size	Info
	1 2000-12-27	11:54:10,160000 NetworkG_10:22:1	b Broadcast	ARP	64		Who has 10.1.0.99? Tell 10.1.0.1
	2 2000-12-27	11:54:10,160000 Runtop_e1:5a:80	NetworkG_10:22:	1b ARP	64		10.1.0.99 is at 00:20:78:e1:5a:80
	3 2000-12-27	11:54:11,224000 10.1.0.1	10.1.0.99	IPv4	1518		Fragmented IP protocol (proto=ICMP 1, off=0, ID=2d00) [Reassembled in #6]
	4 2000-12-27	11:54:11,225000 10.1.0.1	10.1.0.99	IPv4	1518		Fragmented IP protocol (proto=ICMP 1, off=1480, ID=2d00) [Reassembled in #6
	5 2000-12-27	11:54:11,226000 10.1.0.1	10.1.0.99	IPv4	1518		Fragmented IP protocol (proto=ICMP 1, off=2960, ID=2d00) [Reassembled in #6]
•	6 2000-12-27	11:54:11,227000 10.1.0.1	10.1.0.99	ICMP	1206		Echo (ping) request id=0x0100, seq=1536/6, ttl=32 (reply in 10)
		11:54:11,228000 10.1.0.99	10.1.0.1	IPv4	1518		Fragmented IP protocol (proto=ICMP 1, off=0, ID=3100) [Reassembled in #10]
		11:54:11,228000 10.1.0.99	10.1.0.1	IPv4	1518		Fragmented IP protocol (proto=ICMP 1, off=1480, ID=3100) [Reassembled in #10]
		11:54:11,228000 10.1.0.99	10.1.0.1	IPv4	1518		Fragmented IP protocol (proto=ICMP 1, off=2960, ID=3100) [Reassembled in #10]
		11:54:11,228000 10.1.0.99	10.1.0.1	ICMP	1206		Echo (ping) reply id=0x0100, seq=1536/6, ttl=128 (request in 6)
		11:54:12,238000 10.1.0.1	10.1.0.99	IPv4	1518		Fragmented IP protocol (proto=ICMP 1, off=0, ID=2e00) [Reassembled in #14]
	12 2000-12-27	11:54:12,240000 10.1.0.1	10.1.0.99	IPv4	1518		Fragmented IP protocol (proto=ICMP 1, off=1480, ID=2e00) [Reassembled in #14]
		11:54:12,241000 10.1.0.1	10.1.0.99	IPv4	1518		Fragmented IP protocol (proto=ICMP 1, off=2960, ID=2e00) [Reassembled in #14]
		11:54:12,242000 10.1.0.1	10.1.0.99	ICMP	1206		Echo (ping) request id=0x0100, seq=1792/7, ttl=32 (reply in 18)
		11:54:12,242000 10.1.0.99	10.1.0.1	IPv4	1518		Fragmented IP protocol (proto=ICMP 1, off=0, ID=3200) [Reassembled in #18]
		11:54:12,242000 10.1.0.99	10.1.0.1	IPv4	1518		Fragmented IP protocol (proto=ICMP 1, off=1480, ID=3200) [Reassembled in #18]
		11:54:12,242000 10.1.0.99	10.1.0.1	IPv4	1518		Fragmented IP protocol (proto=ICMP 1, off=2960, ID=3200) [Reassembled in #18]
		11:54:12,242000 10.1.0.99	10.1.0.1	ICMP	1206		Echo (ping) reply id=0x0100, seq=1792/7, ttl=128 (request in 14)
		11:54:13,241000 10.1.0.1	10.1.0.99	IPv4	1518		Fragmented IP protocol (proto=ICMP 1, off=0, ID=2f00) [Reassembled in #22]
		11:54:13,242000 10.1.0.1	10.1.0.99	IPv4	1518		Fragmented IP protocol (proto=ICMP 1, off=1480, ID=2f00) [Reassembled in #22]
		11:54:13,244000 10.1.0.1	10.1.0.99	IPv4	1518		Fragmented IP protocol (proto=ICMP 1, off=2960, ID=2f00) [Reassembled in #22]
		11:54:13,245000 10.1.0.1	10.1.0.99	ICMP	1206		Echo (ping) request id=0x0100, seq=2048/8, ttl=32 (reply in 26)
		11:54:13,245000 10.1.0.99	10.1.0.1	IPv4	1518		Fragmented IP protocol (proto=ICMP 1, off=0, ID=3300) [Reassembled in #26]
		11:54:13,245000 10.1.0.99	10.1.0.1	IPv4	1518		Fragmented IP protocol (proto=ICMP 1, off=1480, ID=3300) [Reassembled in #26]
		11:54:13,245000 10.1.0.99	10.1.0.1	IPv4	1518		Fragmented IP protocol (proto=ICMP 1, off=2960, ID=3300) [Reassembled in #26]
	26 2000-12-27	11:54:13,245000 10.1.0.99	10.1.0.1	ICMP	1206	5	Echo (ping) reply id=0x0100, seq=2048/8, ttl=128 (request in 22)
→ Fr	ame 4: 1518 b	oytes on wire (12144 bits), 1518	bytes captured (1	2144 bits)	on interf	ace unknown, id 0	
		c: NetworkG_10:22:1b (00:00:65:1					
+ Ir	ternet Protoc	col Version 4, Src: 10.1.0.1, Dst	: 10.1.0.99				
	0100 = 1						
	0101 =	Header Length: 20 bytes (5)					
*	Differentiat	ed Services Field: 0x00 (DSCP: CS	0, ECN: Not-ECT)				
	0000 00	= Differentiated Services Codepoi	int: Default (0)				
		= Explicit Congestion Notification	on: Not ECN-Capab	e Transpor	t (0)		
	Total Length	: 1500					
	Identificati	on: 0x2d00 (11520)	Total langti	ofosc	h fragn	ont	
*	Flags: 0x20,	More fragments	Total length				
	0	= Reserved bit: Not set	with ID=2d	00 is 15	i00 byte	es and	
	.0	= Don't fragment: Not set					
	1	= More fragments: Set	each fragm	ень наѕ	an nea	ider neid	
	0 0101 110	00 1000 = Fragment Offset: 1480	of 20 bytes				
	* * * * * * * * * * * * * * * * * * *		wytou				14

Figure 4: Total length and header length fileds of fragments

No.	Time	Source	Destination	Protocol	Length	Calculated window size	Info
		11:54:10,160000 NetworkG_10:22:1		ARP	6	4	Who has 10.1.0.99? Tell 10.1.0.1
	2 2000-12-27	11:54:10,160000 Runtop_e1:5a:80	NetworkG_10:22	:1b ARP	6		10.1.0.99 is at 00:20:78:e1:5a:80
	3 2000-12-27	11:54:11,224000 10.1.0.1	10.1.0.99	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=0, ID=2d00) [Reassembled in #6]
	4 2000-12-27	11:54:11,225000 10.1.0.1	10.1.0.99	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=1480, ID=2d00) [Reassembled in #6]
	5 2000-12-27	11:54:11,226000 10.1.0.1	10.1.0.99	IPv4	151	3	Fragmented IP protocol (proto=ICMP 1, off=2960, ID=2d00) [Reassembled in #6]
	6 2000-12-27	11:54:11,227000 10.1.0.1	10.1.0.99	ICMP	120	6	Echo (ping) request id=0x0100, seq=1536/6, ttl=32 (reply in 10)
		11:54:11,228000 10.1.0.99	10.1.0.1	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=0, ID=3100) [Reassembled in #10]
		11:54:11,228000 10.1.0.99	10.1.0.1	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=1480, ID=3100) [Reassembled in #10]
	9 2000-12-27	11:54:11,228000 10.1.0.99	10.1.0.1	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=2960, ID=3100) [Reassembled in #10]
	10 2000-12-27	11:54:11,228000 10.1.0.99	10.1.0.1	ICMP	120	5	Echo (ping) reply id=0x0100, seq=1536/6, ttl=128 (request in 6)
	11 2000-12-27	11:54:12,238000 10.1.0.1	10.1.0.99	IPv4	151	3 /	Fragmented IP protocol (proto=ICMP 1, off=0, ID=2e00) [Reassembled in #14]
		11:54:12,240000 10.1.0.1	10.1.0.99	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=1480, ID=2e00) [Reassembled in #14]
	13 2000-12-27	11:54:12,241000 10.1.0.1	10.1.0.99	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=2960, ID=2e00) [Reassembled in #14]
		11:54:12,242000 10.1.0.1	10.1.0.99	ICMP	120		Echo (ping) request id=0x0100, seq=1792/7, ttl=32 (reply in 18)
		11:54:12,242000 10.1.0.99	10.1.0.1	IPv4	151	3 / /	Fragmented IP protocol (proto=ICMP 1, off=0, ID=3200) [Reassembled in #18]
		11:54:12,242000 10.1.0.99	10.1.0.1	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=1480, ID=3200) [Reassembled in #18]
		11:54:12,242000 10.1.0.99	10.1.0.1	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=2960, ID=3200) [Reassembled in #18]
	18 2000-12-27	11:54:12,242000 10.1.0.99	10.1.0.1	ICMP	120		Echo (ping) reply id=0x0100, seq=1792/7, ttl=128 (request in 14)
	19 2000-12-27	11:54:13,241000 10.1.0.1	10.1.0.99	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=0, ID=2f00) [Reassembled in #22]
		11:54:13,242000 10.1.0.1	10.1.0.99	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=1480, ID=2f00) [Reassembled in #22]
		11:54:13,244000 10.1.0.1	10.1.0.99	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=2960, ID=2f00) [Reassembled in #22]
		11:54:13,245000 10.1.0.1	10.1.0.99	ICMP	120		Echo (ping) request 1d=0x0100, seq=2048/8, ttl=32 (reply in 26)
		11:54:13,245000 10.1.0.99	10.1.0.1	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=0, ID=3300) [Reassembled in #26]
		11:54:13,245000 10.1.0.99	10.1.0.1	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=1480, ID=3300) [Reassembled in #26]
		11:54:13,245000 10.1.0.99	10.1.0.1	IPv4	151		Fragmented IP protocol (proto=ICMP 1, off=2960, ID=3300) [Reassembled in #26]
	26 2000-12-27	11:54:13,245000 10.1.0.99	10.1.0.1	ICMP	120	· // /	Echo (ping) reply id=0x0100, seq=2048/8, ttl=128 (request in 22)
	Total Length:	: 1500				///	
		on: 0x2d00 (11520)				/	
*	Flags: 0x20,	More fragments			- M	//	
		= Reserved bit: Not set			- 1	-	
		= Don't fragment: Not set					
	1 =	= More fragments: Set					
		00 0000 = Fragment Offset: 0					
	Time to Live:						
	Protocol: ICM		Ero o	monto	haca II	No are	
		sum: 0x33bc [validation disabled]		ments w			
		ksum status: Unverified]	2d0	0.2e00. a	ind 2f0	Orespectively	
	Source Addres					eld value of 32	
		Address: 10.1.0.99	nave	e rime to	ilive ne	eid value of 32	
		IPv4 in frame: 6]					
▼ Da	ata (1480 byte	es)					

Figure 5: Time to live fields of fragments with ID's 2d00,2e00,2f00 is 32

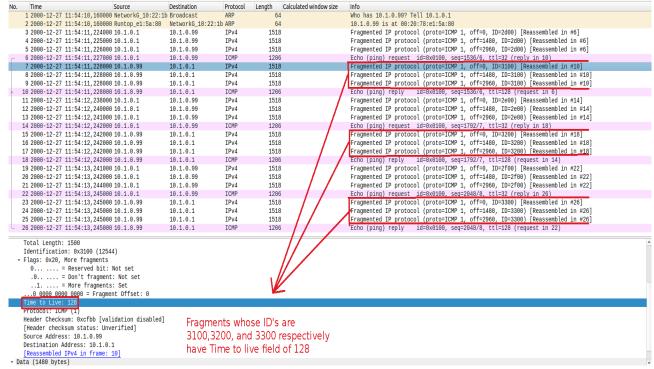


Figure 6: Time to live fields of fragments with ID's 3100,3200,3300 is 128

3.) If we look at the line number 8 we can see in the packet comments section that "IPv6 packets are encapsulated in IPv4" this is the method of tunelling and encapsulation where IPv6 datagrams are embedded inside IPv4 datagram to sent to destination address. Also when we look at the Internet protocol version 4 part below the packet capture table we can see that Protocol field is also IPv6.

o. Time	Source	Destination	Protocol		Calculated window size Info
1 2011-05-30 20:29:16,097280		2607:f0d0:200		106	
2 2011-05-30 20:29:16,294291		2002:1806:add		106	
3 2011-05-30 20:29:16,295000		2607:f0d0:200		94	
4 2011-05-30 20:29:16,296339		2607:f0d0:200		758	
5 2011-05-30 20:29:16,506624		2002:1806:add		94	
6 2011-05-30 20:29:16,507834		2002:1806:add		498	
7 2011-05-30 20:29:16,704527		2607:f0d0:200		94	
8 2011-05-30 20:29:19,758768		2607:f0d0:200		778	
9 2011-05-30 20:29:19,963078		2002:1806:add		502	
10 2011-05-30 20:29:20,167766		2607:f0d0:200	TCP	94	
11 2011-05-30 20:29:21,702919	fe80::5083:65		DHCPv6	147	
12 2011-05-30 20:29:22,702904	fe80::5083:65		DHCPv6	147	
13 2011-05-30 20:29:24,703065	fe80::5083:65		DHCPv6	147	
14 2011-05-30 20:29:28,703206	fe80::5083:65		DHCPv6	147	
15 2011-05-30 20:29:34,973212		2002:1806:add		94	
16 2011-05-30 20:29:34,973527		2607:f0d0:200		94	
17 2011-05-30 20:29:36,703973	fe80::5083:65	ff02::1:2	DHCPv6	147	Solicit XID: 0xc37c88 CID: 0001000114294f3cd48564a3b133
Notice that these IPv6 packets Expert Info (Comment/Commen [Notice that these IPv6 p. [Severity level: Comment] [Group: Comment]	nt): Notice that ackets are encaps	these IPv6 packe ulated in IPv4 -	ts are en the pack	capsulated ket is rout	sed on the IPv4 header through an IPv4 network. in IPv4 - the packet is routed based on the IPv4 header through an IPv4 network.] ed based on the IPv4 header through an IPv4 network.]
Figure 1. Expert Info (Comment/Comment [Notice that these IPv6 pp [Severity level: Comment] [Group: Comment] Frame 8: 778 bytes on wire (6224 Ethernet II, Src: HewlettP_a7:bf	nt): Notice that ackets are encaps bits), 778 bytes :a3 (d4:85:64:a7	these IPv6 packe ulated in IPv4 - s captured (6224 :bf:a3), Dst: Ca	ts are en the pack bits) on dant_31:bl	icapsulated ket is rout interface b:c1 (00:0	in IPv4 - the packet is routed based on the IPv4 header through an IPv4 network.] ed based on the IPv4 header through an IPv4 network.] unknown, id 0 1:5c:31:bb:c1)
Notice that these IPv6 packets [Expert Info (Comment/Commer [Notice that these IPv6 pp. [Severity level: Comment] [Group: Comment] Frame 8: 778 bytes on wire (6224	nt): Notice that ackets are encaps bits), 778 bytes :a3 (d4:85:64:a7	these IPv6 packe ulated in IPv4 - s captured (6224 :bf:a3), Dst: Ca	ts are en the pack bits) on dant_31:bl	icapsulated ket is rout interface b:c1 (00:0: When w	in IPv4 - the packet is routed based on the IPv4 header through an IPv4 network.] de based on the IPv4 header through an IPv4 network.] unknown, id 0 Lisc:31:bb:c1) e look at the get request
- Notice that these IPv6 packets - [Expert Info [Comment/Commer] [Notice that these IPv6 pu [Severity level: Comment] Frame 8: 778 bytes on wire (6224 Ethernet II, Src: HewlettP_A7:bf Internet Protocol Version 4, Src 0100 Version: 4	ht): Notice that ackets are encaps bits), 778 byte: :a3 (d4:85:64:a7 : 24.6.173.220, [these IPv6 packe ulated in IPv4 - s captured (6224 :bf:a3), Dst: Ca	ts are en the pack bits) on dant_31:bl	icapsulated ket is rout interface b:c1 (00:0: When w	in IPv4 - the packet is routed based on the IPv4 header through an IPv4 network.] de based on the IPv4 header through an IPv4 network.] unknown, id 0 Lisc:31:bb:c1) e look at the get request
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- Notice that these IPv6 packets - [Expert Info (Comment/Commen [Notice that these IPv6 pn [Severity level: Comment] - Frame 8: 778 bytes on wire (6224 Ethernet II, Src: Hewlettp_aribf Internet Protocol Version 4, Src 0100 = Version: 4 0101 = Header Length: 20 - Differentiated Services Field: Total Length: 764 - Total Length: 764	bits), 778 byte: :a3 (d4:85:64:a7: : 24.6.173.220, [bytes (5) 0x00 (DSCP: CS0	these IPv6 packersulated in IPv4 - s captured (6224:bf:a3), Dst: Caust: 192.88.99.1	ets are en the pack bits) on dant_31:bl	interface b:c1 (00:00 When w ine we and prot	in IPV4 - the packet is routed based on the IPV4 header through an IPV4 network.] ed based on the IPV4 header through an IPV4 network.] unknown, id 0 1.55:31:15b:c1) e look at the get request can see a packet comment occl information which s that IPV6 datagrams are
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- Notice that these IPv6 packets - [Expert Info (Comment/Commer) - [Notice that these IPv6 p [Severity level: Comment] - [Group: Comment] - Frame 8: 778 bytes on wire (6224 - Ethernet II, Src: HewlettP_a7:bf - Internet Protocol Version 4, Src - 1080 Version: 4 0191 = Header Length: 20 - Differentiated Services Field: - Total Length: 764 - Identification: 0x2d40 (11584) - Flags: 0x80 0 0000 0000 0000 = Fragment - Ime to Live: 128 - Protocol: 19v6 (43) - Header Checksum: Dxxxxxxxxxxxxxxxxxxxxxxxxxxxxxxxxxxxx	ht): Notice that ackets are encaps bits), 778 byte: a3 (d4:85:64:a7: 24.6.173.220, l bytes (5) 0x00 (DSCP: CS0 Offset: 0	these IPv6 packersulated in IPv4 - s captured (6224:bf:a3), Dst: Caust: 192.88.99.1	ets are en- the pack bits) on dant_31:bl	interface b:c1 (00:0: When w ine we and prot ndicate	in IPV4 - the packet is routed based on the IPV4 header through an IPV4 network.] ed based on the IPv4 header through an IPv4 network.] unknown, id 0 1:50:31:10b:c1) e look at the get request can see a packet comment occl information which is that IPv6 datagrams are lated inside IPv4 packets
- Notice that these IPv6 packets - [Expert Info (Comment/Commer [Notice that these IPv6 pn [Severity level: Comment] Frame 8: 778 bytes on wire (622 Ethernet II, 5rc: HewlettP_a7:bf Internet Protocol Version 4, 5rc 0100 Version; 4 0101 = Header Length: 20 Differentiated Services Field: Total Length: 764 Identification: 0x2d40 (11584) Flags: 0x00 0 0000 0000 0000 = Fragment Time to Live: 128 Frotocools 1000 (45) Header Checksum xxxxxxxxxxxxxxxxxxxxxxxxxxxxxxxxxxxx	ht): Notice that ackets are encaps this), 778 byte: bis), 778 byte: bis, 31 (44:85:64:a7: 24.6:173.226, bytes (5) 0x00 (DSCP: CS0 Offset: 0	these IPv6 packersulated in IPv4 - s captured (6224:bf:a3), Dst: Caust: 192.88.99.1	ets are en- the pack bits) on dant_31:bl	interface b:c1 (00:0: When w ine we and prot ndicate	in IPV4 - the packet is routed based on the IPV4 header through an IPV4 network.] ed based on the IPv4 header through an IPv4 network.] unknown, id 0 1:50:31:10b:c1) e look at the get request can see a packet comment occl information which is that IPv6 datagrams are lated inside IPv4 packets
- Notice that these IPv6 packets - [Expert Info (Comment/Commer) - [Notice that these IPv6 p [Severity level: Comment] - [Group: Comment] - Frame 8: 778 bytes on wire (6224 - Ethernet II, Src: HewlettP_a7:bf - Internet Protocol Version 4, Src - 1080 Version: 4 0191 = Header Length: 20 - Differentiated Services Field: - Total Length: 764 - Identification: 0x2d40 (11584) - Flags: 0x80 0 0000 0000 0000 = Fragment - Ime to Live: 128 - Protocol: 19v0 (43) - Header Checksum: Dxxxxxxxxxxxxxxxxxxxxxxxxxxxxxxxxxxxx	ht): Notice that accets are encaps bits), 778 byte: a3 (d4:85:64:a7:24.6.173.220, i bytes (5) exe0 (DSCP: CS0 offset: 0 lation disabled] ified]	these IPv6 packe ulated in IPv4 - s captured (6224 bf:a3), Dst: Ca sst: 192.88.99.1 , ECN: Not-ECT)	ets are en- the pack bits) on dant_31:bl	interface bic1 (00:0 When w ine we and profindicate encapsusefore s	in IPV4 - the packet is routed based on the IPV4 header through an IPv4 network.] ed based on the IPv4 header through an IPv4 network.] unknown, id 6 elook at the get request can see a packet comment occol information which s that IPv6 datagrams are lated inside IPv4 packets ent to destination
Notice that these IPv6 packets - [Expert Info [Comment/Comment] [Severity Info [Comment] [Severity Level: Comment] [Frame 8: 778 bytes on wire (6224 Ethernet II, 5rc: HewlettP_a7:bf Internet Protocol Version 4, 5rc 0:109 Version: 4 0:101 = Header Length: 20 Differentiated Services Field: Total Length: 704 Identification: 0x2d40 (11584) Flags: 0x00 0:0000 00000 = Fragment Time to Live: 128 Protocous IPvos (45) Header Checksum xxxxxxxx [Header Checksum status: Univer Source Address: 24.6.173.220 Destination Address: 192.88.90 Destination Address: 192.88.00 Internet Protocol Version 6, 5rc	ht): Notice that ackets are encaps bits), 778 byte: :a3 (d4:85:64:a7 : 24.6.173.226, bytes (5)	these IPv5 packeuluated in IPv4 - s captured (6224 stbf:a3), Dst: Cac Dst: 192.88.99.1 , ECN: Not-ECT)	bits are en. the pack bits) on dant_31:bl c i i c t t t t t t t t t t t t t t t	interface bic1 (00:0 When w ine we and prot ndicate encapsu before s	in IPv4 - the packet is routed based on the IPv4 header through an IPv4 network.] unknown, id 0 in:5c:31:bb:c1) e look at the get request can see a packet comment ocol information which s that IPv6 datagrams are lated inside IPv4 packets ent to destination
- Notice that these IPv6 packets - [Expert Info (Comment/Commer) - [Notice that these IPv6 p [Severity level: Comment] - [Group: Comment] - Frame 8: 778 bytes on wire (6224 - Ethernet II, Src: HewlettP_a7:bf - Internet Protocol Version 4, Src - 1080 Version: 4 0191 = Header Length: 20 - Differentiated Services Field: - Total Length: 764 - Identification: 0x2d40 (11584) - Flags: 0x80 0 0000 0000 0000 = Fragment - Ime to Live: 128 - Protocol: 19v0 (43) - Header Checksum: Dxxxxxxxxxxxxxxxxxxxxxxxxxxxxxxxxxxxx	ht): Notice that ackets are encaps bits), 778 byte: :a3 (d4:85:64:a7 : 24.6.173.226, bytes (5)	these IPv5 packeuluated in IPv4 - s captured (6224 stbf:a3), Dst: Cac Dst: 192.88.99.1 , ECN: Not-ECT)	bits are en. the pack bits) on dant_31:bl c i i c t t t t t t t t t t t t t t t	interface bic1 (00:0 When w ine we and prot ndicate encapsu before s	in IPv4 - the packet is routed based on the IPv4 header through an IPv4 network.] unknown, id 0 itsc::31:bb::01) e look at the get request can see a packet comment occil information which s that IPv6 datagrams are lated inside IPv4 packets ent to destination

Figure 7: IPv6 packets encapsulated inside IPv4 datagrams

- **a.)** We can observe two ICMP message types as a result of executing traceroute command from terminal. First one is Type 11 which gives the time exceeded message when a router receives a datagram with TTL of 0 or 1. So by using this TTL value infinite routing loops are prevented. Routers can not forward a datagram that has a TTL of 0 or 1 and packet is dropped. ICMP message of code 0 indicates Time-to-live became equal to 0 during transit not during reassembly. Second is Type 3 ICMP message which indicates that the router cannot find the destination network. Code 3 which says "port uncreachable" indicates that trasport layer protocol we are attempting to communicate is not available on the arriving side.
 - **b.)** As a result of executing traceroute google.com command, intermediate network devices are same in both cases because when we look at the hop IP addresses, in the Wireshark and in the output of traceroute they are same. It is intuitive to get same intermediate device IP addresses because we have captured packets during the same traceroute execution. (Figure 10)
 - **c.)** Routers with IP addresses given inside blue rectangle in the screenshot below have have been sent 3 packets. Routers with the IP addresses given inside the red rectangle have been sent 1 packet. (Figure 8 and Figure 9)

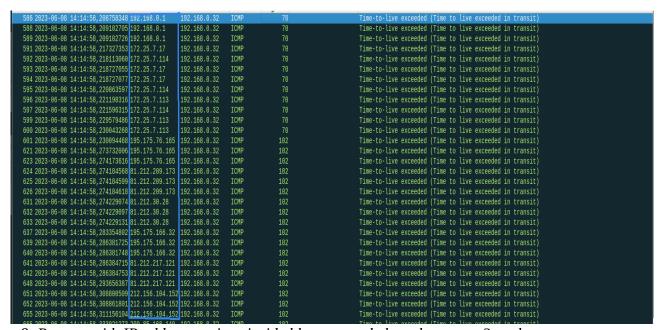


Figure 8: Routers with IP addresses given inside blue rectangle have been sent 3 packets

665 2023-06-08 14:14:58,333021273 209.85.168.140	192.168.0.32 ICMP	102	Time-to-live exceeded (Time to live exceeded in transit)
667 2023-06-08 14:14:58,334399051 72.14.222.58	192.168.0.32 ICMP	102	Time-to-live exceeded (Time to live exceeded in transit)
669 2023-06-08 14:14:58,335308721 74.125.51.44	192.168.0.32 ICMP	110	Time-to-live exceeded (Time to live exceeded in transit)
670 2023-06-08 14:14:58,343530413 142.251.61.242	192.168.0.32 ICMP	102	Time-to-live exceeded (Time to live exceeded in transit)
671 2023-06-08 14:14:58,344754740 142.251.92.65	192.168.0.32 ICMP	110	Time-to-live exceeded (Time to live exceeded in transit)
672 2023-06-08 14:14:58,345201851 <mark>14</mark> 2.251.227.252	192.168.0.32 ICMP	110	Time-to-live exceeded (Time to live exceeded in transit)
674 2023-06-08 14:14:58,349824901 <mark>2</mark> 09.85.254.243	192.168.0.32 ICMP	102	Time-to-live exceeded (Time to live exceeded in transit)
684 2023-06-08 14:14:58,365977146 142.251.92.2	192.168.0.32 ICMP	70	Time-to-live exceeded (Time to live exceeded in transit)
685 2023-06-08 14:14:58,366764923 <mark>14</mark> 2.251.92.67	192.168.0.32 ICMP	110	Time-to-live exceeded (Time to live exceeded in transit)
690 2023-06-08 14:14:58,390165271 <mark>1</mark> 42.250.187.174	192.168.0.32 ICMP	70	Destination unreachable (Port unreachable)
693 2023-06-08 14:14:58,401546910 <mark>1</mark> 42.251.243.28	192.168.0.32 ICMP	182	Time-to-live exceeded (Time to live exceeded in transit)
711 2023-06-08 14:14:58,521040648 142.250.187.174	192.168.0.32 ICMP	70	Destination unreachable (Port unreachable)
712 2023-06-08 14:14:58,521191387 142.251.243.28	192.168.0.32 ICMP	182	Time-to-live exceeded (Time to live exceeded in transit)
713 2023-06-08 14:14:58,523857302 142.251.242.230	192.168.0.32 ICMP	110	Time-to-live exceeded (Time to live exceeded in transit)

Figure 9: Routers with IP address given inside the red rectangle have been sent 1 packet

```
| Paddresses given inside the blue rectangle match with the | Paddresses given inside the blue rectangle match with the | Paddresses given inside the blue rectangles in Wireshark. | Paddresses given inside the blue rectangles in Wireshark. | Paddresses given inside the blue rectangles in Wireshark. | Paddresses given inside the blue rectangles in Wireshark. | Similarly, IPaddresses given inside the red rectangle match with the | Paddresses given inside the red rectangle match with the | Paddresses given inside the red rectangle match with the | Paddresses given inside the red rectangle match with the | Paddresses given inside the red rectangle match with the | Paddresses given inside the red rectangle match with the | Paddresses given inside the red rectangle in | Paddresses given inside the red rectangle match with the | Paddresses given inside the red rectangle in | Paddresses given inside the red rectangle match with the | Paddresses given inside the red rectangle in | Paddresses given inside the red rectangle match with the | Paddresses given inside the red rectangle in | Paddresses given | Pad
```

Figure 10: traceroute command output and the hops along with their IP addresses and packets they received

NAT SECTION

1.) NAT protocol stands for Network address translation which allows all devices inside local area network to share just one IP address as far as outer networks are

considered, so all datagrams leaving from local network share same NAT IP address, but they are distinguished by different port numbers. The mapping between datagram source IP address and NAT IP address, is implemented in the form of NAT table. NAT protocol also helps IPv4 address exhaustion because just one IP address is needed from ISP for all devices in local network.

2.) When we look at the lan-packets screenshot below, scenario might be as follows: host with IP address 10.142.154.239 and destination with IP address 8.8.8.8 are located in the same local area network and host is trying to establish a TCP connection with server whose IP address is 8.8.8.8. When we look at the wan-packets screenshot below, scenario might be as follows: host with IP address 136.102.83.11 and destination with IP address 8.8.8.8 are located in different networks and host is trying to establish a TCP connection with server whose IP address is 8.8.8.8. This scenarios can be related to NAT protocol as follows: Actual IP address of host may be 10.142.154.239 inside it's local area network but let's say as in WAN scenario server with IP address 8.8.8.8 is located outside the LAN and any host when connecting to any server in LAN is using IP address 10.142.154.139, but when server for which host is going to establish connection is outside the LAN, there is a NAT router which translates IP address of host from 10.142.154.239 to 136.102.83.11. Thus when server sends response to host it sees it's NAT IP which is 136.102.83.11 not it's actual local network IP address which is 10.142.154.239. (Figure 11 and Figure 12)

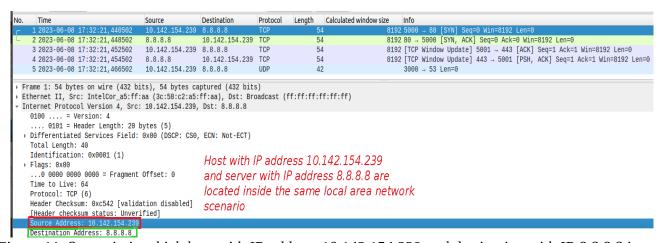


Figure 11: Scenario in which host with IP address 10.142.154.239 and destination with IP 8.8.8.8 in same LAN

No	. Time	Source	Destination	Protocol	Length	Calculated window six	ze	Info		
	1 2023-06-08 17:32:21,440502	136.102.83.11	8.8.8.8	TCP	5	4	8192	5000 → 80 [SYN] Se	q=0 Win=8192 Len=0	
L	2 2023-06-08 17:32:21,448502	8.8.8.8	136.102.83.11	TCP	5	4	8192	80 → 5000 [SYN, AC	K] Seq=0 Ack=0 Win=	=8192 Len=0
	3 2023-06-08 17:32:21,452502	136.102.83.11	8.8.8.8	TCP	5	4				Seq=1 Ack=1 Win=8192 Len=0
	4 2023-06-08 17:32:21,454502	8.8.8.8	136.102.83.11	TCP	5	4] 443 → 5001 [PSH,	ACK] Seq=1 Ack=1 Win=8192 Len=0
	5 2023-06-08 17:32:21,466502	136.102.83.11	8.8.8.8	UDP	4	2		3000 → 53 Len=0		
L										
1	Frame 1: 54 bytes on wire (432 b.									
ľ	Ethernet II, Src: IntelCor_a5:ff			adcast (f	T:TT:TT:	rr:rr:rr)				
ľ	Internet Protocol Version 4, Src	: 136.102.83.11,	DST: 8.8.8.8							
1	0100 = Version: 4	hutaa (E)								
	0101 = Header Length: 20 • Differentiated Services Field:		ECN: Not ECT)							
1	0000 00 = Differentiated S									
1	00 = Explicit Congest			e Transnor	t (a)					
	Total Length: 40	cion Mocificacion	i. Not Lon-oupabl	c ii anspoi	(0)					
	Identification: 0x0001 (1)									
	▼ Flags: 0x00									
П	0 = Reserved bit: No	ot set	Host with IP	addres.	s 136.	102.83.11				
П	.0 = Don't fragment:	Not set	and server v	vith ID -	ddroc	0000				
L	0 = More fragments:	Not set								
П	0 0000 0000 0000 = Fragment		are located .	in diffei	rent lo	cal are netwo	rks			
П	Time to Live: 64		scenario							
П	Protocol: TCP (6)		Sceriario							
П	Header Checksum: 0x8f4e [valid	lation disabled]								
	[Header checksum status: Unver	ified]								
	Source Address: 136.102.83.11									
	Destination Address: 8.8.8.8									

Figure 12: Scenario in which host with IP address 136.102.83.11 and destination with IP 8.8.8.8 in different LANs

DHCP SECTION

- **1.)** ipconfig /release command will inform the DHCP server that we don't want to use previously assigned IP address any larger and we release it. After releasing the previously assigned IP address we can request new IP address from DHCP server with ipconfig /renew command. Note: As I am using **UBUNTU** distribution of Linux, I have used **dhclient** program for releasing and obtaining new IP address with DHCP which is equivalent to ipconfig program on Windows.
- **2.)** When we look at the packet capture on Wireshark 1 DHCP packet was captured for the release command, 4 DHCP packets were captured for first renew command. 2 DHCP packets were captured for second renew command.
- 3.) If we look at the Wireshark packet capture when when I made two new requests respectively to DHCP server both newly assigned IP address in both situation were same as we can see in the below screenshots. In both situations requested IP address was 192.138.0.1, but this may not always be the case, IP address can differ from previous one.
- **4.)** When we look at the packets as a result of first renew command, we can see

four different DHCP message types which are DISCOVER, OFFER, REQUEST, and ACK respectively. DHCP DISCOVER message is sent by client to locate DHCP server when client attempts to connect network for the first time. DHCP OFFER message is sent by server in response to DHCP discover message of client, and it caries configuration information along with the new IP address. DHCP REQUEST message is sent by server. Finally DHCP ACK message is sent by server to acknowledge DHCP REQUEST message sent from DHCP client. After receiving DHCP ACK message client obtains configuration parameters including new IP address.

5.) For first dhclient command (which is equivalent to piconfig /renew on Windows) common option fileds to all packets are as follows: subnet mask which defines the subnet IP of computer that belongs to particular LAN, IP address of DHCP server, and DHCP message type option which indicates the sending purpose of DHCP packet that I have explained detaily in the previous question (4) it can be of type DISCOVER, OFFER, REQUEST, ACK, etc.. to indicate what client wants from DHCP server and DHCP servers response.

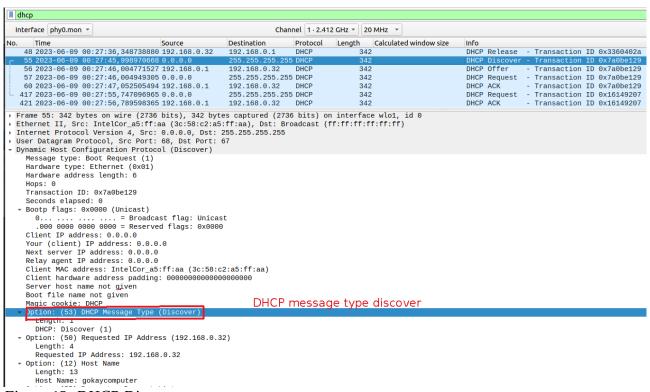


Figure 13: DHCP Discover message

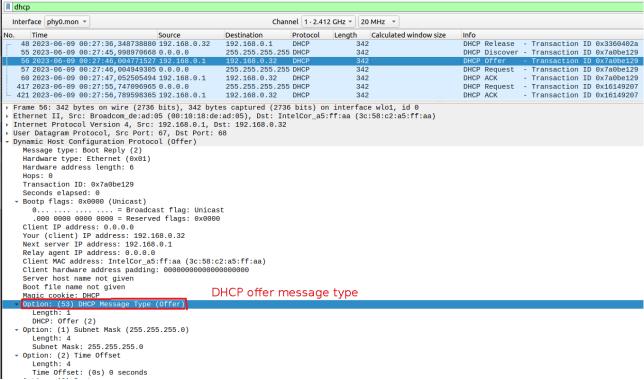


Figure 14: DHCP Offer message type

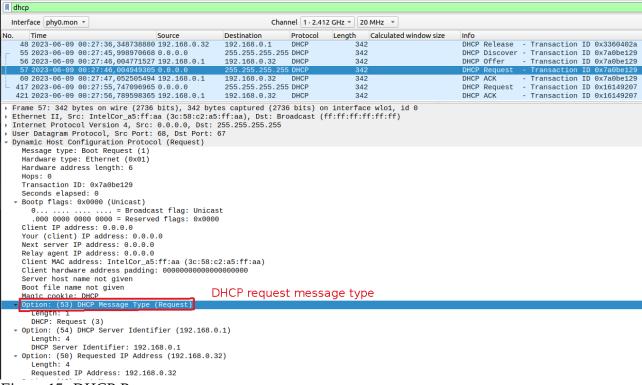


Figure 15: DHCP Request message type

```
ackay@ackaycomputer:—$ sudo dhclient -v -r wlo1
xilled old client process
Internet Systems Consortium DHCP Client 4.4.1
Copyright 2004-2018 Internet Systems Consortium.
All rights reserved.
For info, please visit https://www.isc.org/software/dhcp/

Listenting on LPF/wlo1/3c:58:c2:a5:ff:aa
Sending on Scket/falback
DHCPRELEASE of 192.168.0.32 on wlo1 to 192.168.0.1 port 67 (xid=0x629a7f54)
gokay@gokaycomputer:—$ sudo dhclient -v wlo1
Internet Systems Consortium DHCP Client 4.4.1
Copyright 2004-2018 Internet Systems Consortium.
All rights reserved.
For info, please visit https://www.isc.org/software/dhcp/

Listening on LPF/wlo1/3c:58:c2:a5:ff:aa
Sending on Scket/falback
DHCPDISCOVER on wlo1 to 255.255 is 255 port 67 interval 3 (xid=0x7a0be129)
DHCPOFER of 192.168.032 from 192.168.0.1
DHCPREQUEST for 192.16 jo 32 on wlo1 to 255.255.255 port 67 (xid=0x29e10b7a)
DHCPAFER of 192.168.0.32 from 192.168.0.1 (xid=0x7a0be129)
bound to 192.168.0.32 - renewal in 243054 seconds.
gokay@gokaycomputer:—$ sudo dhclient -v wlo1
Internet Systems Consortium DHCP Client 4.4.1
Copyright 2004-2018 Internet Systems Consortium.
All rights reserved.
For info, please visit https://www.isc.org/software/dhcp/

Listening on LPF/wlo1/3c:58:c2:a5:ff:aa
Sending on LPF/wlo1/3c:58:c2:a5:ff:aa
Sending on LPF/wlo1/3c:58:c2:a5:ff:aa
Sending on Scket/falback
DHCPACK of 192.168.0.32 on wlo1 to 255.255.255.255 port 67 (xid=0x7921416)
DHCPACK of 192.168.0.32 on wlo1 to 255.255.255.255 port 67 (xid=0x7921416)
DHCPACK of 192.168.0.32 on wlo1 to 255.255.255.255 port 67 (xid=0x7921416)
DHCPACK of 192.168.0.32 - renewal in 235556 seconds.
gokay@gokaycomputer:—$ |
```

Figure 16: dhclient command equivalent to ipconfig in Wİndows for releasing IP address and obtaining new IP address

dhcp							
Interface phy0.mon 🕶		Ch	annel 1 · 2.41	20 MHz	•		
o. Time	Source	Destination	Protocol	Length Calcu	lated window size	Info	
48 2023-06-09 00:27:36,348738880		192.168.0.1	DHCP	342		DHCP Release	- Transaction ID 0x3360402
55 2023-06-09 00:27:45,998970668		255.255.255.2		342			- Transaction ID 0x7a0be12
56 2023-06-09 00:27:46,004771527		192.168.0.32	DHCP	342		DHCP Offer	- Transaction ID 0x7a0be12
57 2023-06-09 00:27:46,004949305		255.255.255.2		342			- Transaction ID 0x7a0be12
60 2023-06-09 00:27:47,052505494		192.168.0.32	DHCP	342			- Transaction ID 0x7a0be12
417 2023-06-09 00:27:55,747096965		255.255.255.2		342			- Transaction ID 0x1614926
421 2023-06-09 00:27:56,789598365	192.168.0.1	192.168.0.32	DHCP	342		DHCP ACK	- Transaction ID 0x1614920
Frame 60: 342 bytes on wire (2736 Ethernet II, Src: Broadcom_de:ad: Internet Protocol Version 4, Src: User Datagram Protocol, Src Port: Dynamic Host Configuration Protoco Message type: Boot Reply (2) Hardware type: Ethernet (0x01) Hardware address length: 6 Hops: 0 Transaction ID: 0x7a0be129 Seconds elapsed: 0 * Bootp flags: 0x0000 (Unicast) 0 Broadca .000 0000 0000 0000 = Reserve Client IP address: 0.0.0 Your (client) IP address: 192.16 Relay agent IP address: 192.16 Relay agent IP address: 1ntelCor_a5 Client MAC address: IntelCor_a5 Client hardware address padding Server host name not given Boot file name not given Boot file name not given Madic_cookie: DHCP	55 (06:10:18:de 192.168.0.1, I 67, Dst Port: bl (ACK) st flag: Unica d flags: 0x000 68.0.32 .0.1 :ff:aa (3c:58:	e:ad:05), Dst: 1 Dst: 192.168.0.5 68 st 0	ntelCor <u>a</u> 5	:ff:aa (3c:58:c			
Length: 1	(ACK)						
DHCP: ACK (5)							
- Option: (1) Subnet Mask (255.25	5.255.0)						
Length: 4							
Subnet Mask: 255.255.255.0							
→ Option: (2) Time Offset							
Length: 4							
Time Offset: (0s) 0 seconds							

Figure 17: DHCP ACK message for second dhclient command (same as ipconfig /renew)