A Binary Search Tree Implementation

Chapter 25

Data Structures and Abstractions with Java, 4e, Global Edition Frank Carrano

Binary Search Tree (BST)

- For each node in a binary search tree
 - Node's data is greater than all data in node's left subtree
 - Node's data is less than all data in node's right subtree
- Every node in a binary search tree is the root of a binary search tree

Binary Search Tree

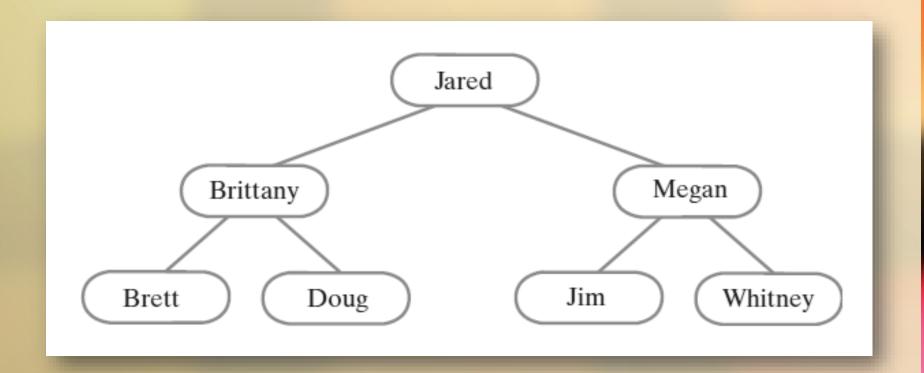


FIGURE 25-1 A binary search tree of names

Interface for the BST

```
package TreePackage;
import java.util.Iterator;
public interface SearchTreeInterface<T extends Comparable<? super T>>
   /** Searches for a specific entry in this tree.
      @param entry An object to be found.
      @return True if the object was found in the tree. */
  public boolean contains(T entry);
   /** Retrieves a specific entry in this tree.
      @param entry An object to be found.
      @return Either the object that was found in the tree or
               null if no such object exists. */
  public T getEntry(T entry);
   /** Adds a new entry to this tree, if it does not match an existing
      object in the tree. Otherwise, replaces the existing object with
                         minumen
```

Interface for the BST

```
/** Adds a new entry to this tree, if it does not match an existing
   object in the tree. Otherwise, replaces the existing object with
   the new entry.
   @param newEntry An object to be added to the tree.
   @return Either null if newEntry was not in the tree already, or
             the existing entry that matched the parameter newEntry
             and has been replaced in the tree. */
public T add(T newEntry);
/** Removes a specific entry from this tree.
   @param entry An object to be removed.
   @return Either the object that was removed from the tree or
            null if no such object exists. */
public T remove(T entry);
/** Creates an iterator that traverses all entries in this tree.
   @return An iterator that provides sequential and ordered access
            to the entries in the tree. */
public Iterator<T> getInorderIterator();
```

LISTING 25-1 An interface for a search tree

Understanding the Specifications

- Methods will use return values instead of exceptions to indicate whether an operation has failed.
- Pay attention to
 - the entry's equals method
 - the entry's compareTo method

Understanding the Specifications

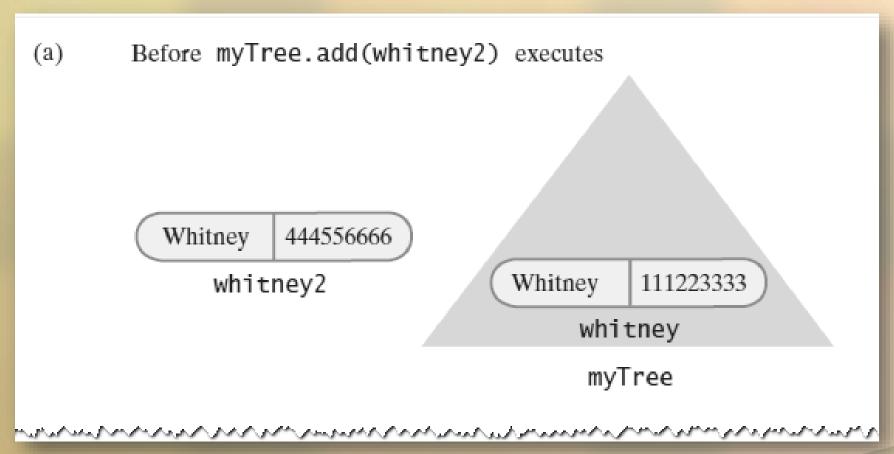


FIGURE 25-2 Adding an entry that matches an entry already in a binary search tree

Understanding the Specifications

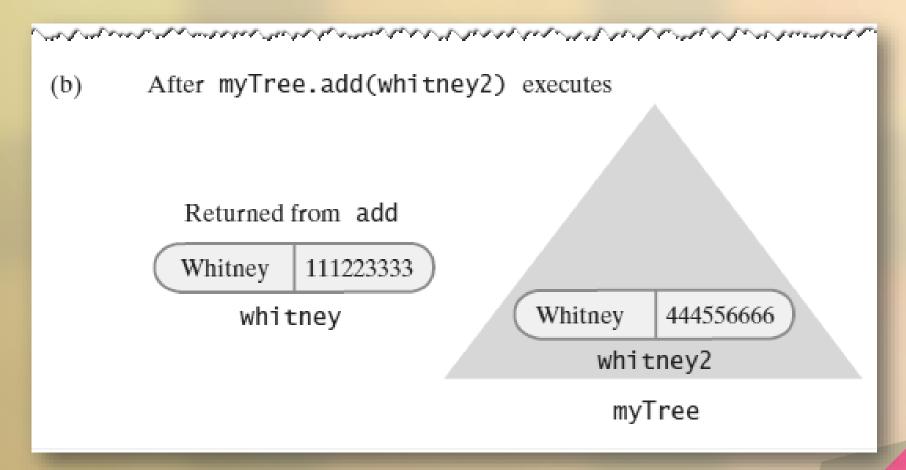


FIGURE 25-2 Adding an entry that matches an entry already in a binary search tree

If Duplicate Entries Are Allowed

- If any entry e has a duplicate entry d, we require that d occur in the right subtree of e's node
- For each node in a binary search tree:
 - Data in a node is greater than data in node's left subtree
 - Data in a node is less than or equal to data in node's right subtree

If Duplicate Entries Are Allowed

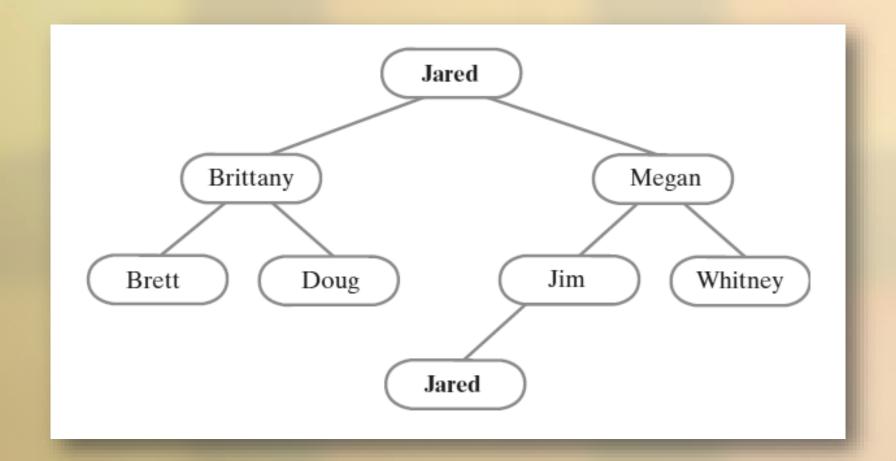


FIGURE 25-3 A binary search tree with duplicate entries

Searching and Retrieving

```
Algorithm bstSearch(binarySearchTreeRoot, desiredObject)
// Searches a binary search tree for a given object.
// Returns true if the object is found.

if (binarySearchTreeRoot is null)
    return false
else if (desiredObject == object in binarySearchTreeRoot)
    return true
else if (desiredObject < object in binarySearchTreeRoot)
    return bstSearch(left child of binarySearchTreeRoot, desiredObject)
else
    return bstSearch(right child of binarySearchTreeRoot, desiredObject)</pre>
```

Algorithm that describes actual implementation more closely

```
public T getEntry(T entry)
   return findEntry(getRootNode(), entry);
} // end getEntry
private T findEntry(BinaryNode<T> rootNode, T entry)
   T result = null;
   if (rootNode != null)
  Suran arman ara maran ara mana maran ara maran ara
      T rootEntry = rootNode.getData();
       if (entry.equals(rootEntry))
          result = rootEntry;
       else if (entry.compareTo(rootEntry) < 0)</pre>
          result = findEntry(rootNode.getLeftChild(), entry);
      else
          result = findEntry(rootNode.getRightChild(), entry);
   } // end if
   return result:
} // end findEntry
```

Searching and Retrieving

```
public boolean contains(T entry)
{
    return getEntry(entry) != null;
} // end contains
```

Method contains can simply call getEntry to see whether a given entry is in the tree:

Adding an Entry

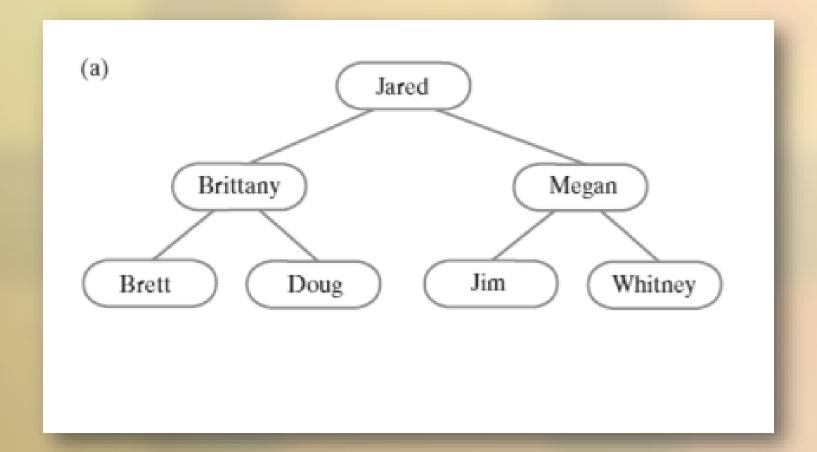


FIGURE 25-4 (a) A binary search tree;

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Adding an Entry

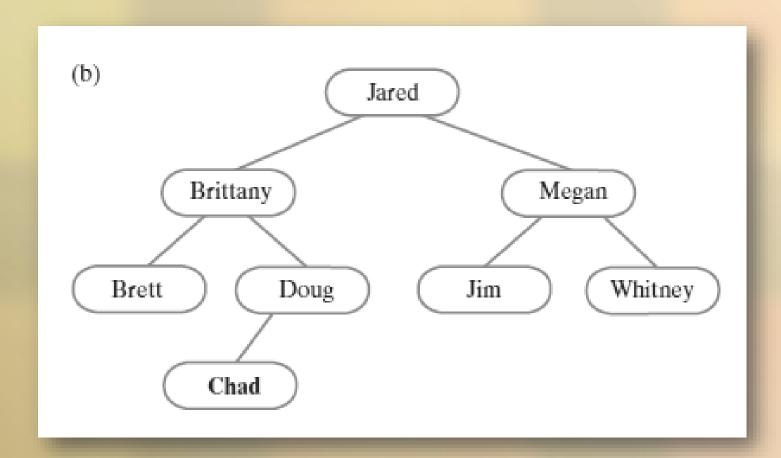


FIGURE 25-4 (b) the same tree after adding Chad

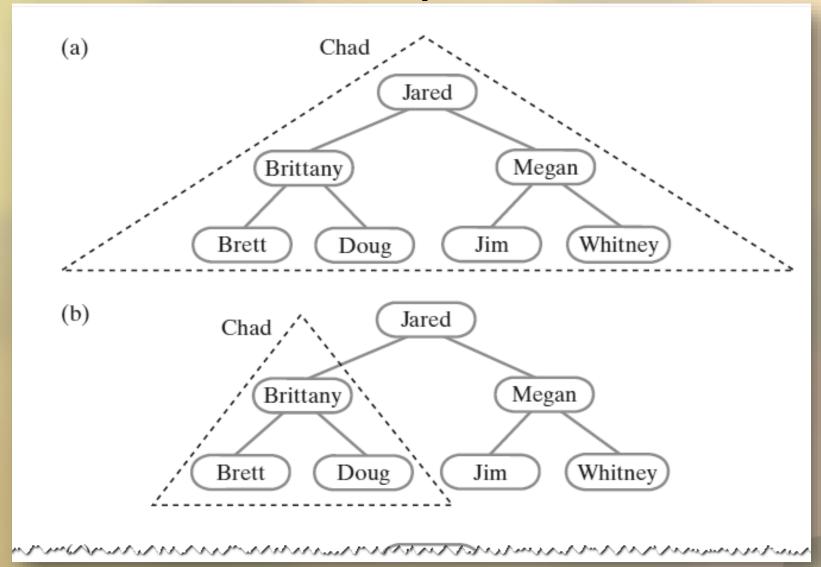


FIGURE 25-5 Recursively adding *Chad* to smaller subtrees of a binary search tree

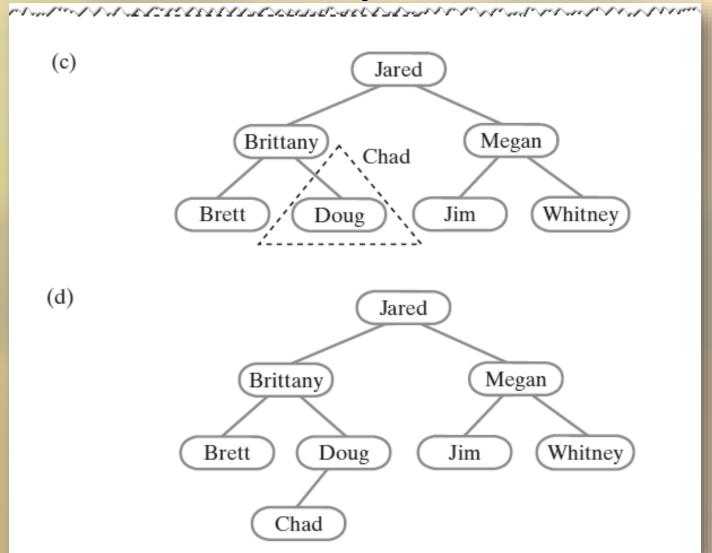


FIGURE 25-5 Recursively adding *Chad* to smaller subtrees of a binary search tree

```
Algorithm add(binarySearchTree, newEntry)
// Adds a new entry to a binary search tree.
// Returns null if newEntry did not exist already in the tree. Otherwise, returns the
// tree entry that matched and was replaced by newEntry.

result = null
if (binarySearchTree is empty)
    Create a node containing newEntry and make it the root of binarySearchTree
else
    result = addEntry(binarySearchTree, newEntry)

return result;
```

Handle the addition to an empty binary search tree as a special case

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```
Algorithm addEntry(binarySearchTree, newEntry)
// Adds a new entry to a binary search tree that is not empty.
// Returns null if newEntry did not exist already in the tree. Otherwise, returns the
// tree entry that matched and was replaced by newEntry.
result = null
if (newEntry matches the entry in the root of binarySearchTree)
   result = entry in the root
   Replace entry in the root with newEntry
else if (newEntry < entry in the root of binarySearchTree)
   if (the root of binarySearchTree has a left child)
       result = addEntry(left subtree of binarySearchTree, newEntry)
   else
       Give the root a left child containing newEntry
else // newEntry > entry in the root of binarySearchTree
```

Recursive algorithm for adding a new entry

```
else // newEntry > entry in the root of binarySearchTree
{
   if (the root of binarySearchTree has a right child)
     result = addEntry(right subtree of binarySearchTree, newEntry)
   else
     Give the root a right child containing newEntry
}
return result
```

Recursive algorithm for adding a new entry

```
public T add(T newEntry)
   T result = null;
   if (isEmpty())
      setRootNode(new BinaryNode<>(newEntry));
   else
      result = addEntry(getRootNode(), newEntry);
   return result;
} // end add
```

```
// Adds newEntry to the nonempty subtree rooted at rootNode.
private T addEntry(BinaryNode<T> rootNode, T newEntry)
   assert rootNode != null;
   T result = null;
   int comparison = newEntry.compareTo(rootNode.getData());
   if (comparison == 0)
   {
      result = rootNode.getData();
      rootNode.setData(newEntry);
   else if (comparison < 0)</pre>
      if (rootNode.hasLeftChild())
         result = addEntry(rootNode.getLeftChild(), newEntry);
      else
         rootNode.setLeftChild(new BinaryNode<>(newEntry));
```

```
else
      assert comparison > 0;
      if (rootNode.hasRightChild())
         result = addEntry(rootNode.getRightChild(), newEntry);
      else
         rootNode.setRightChild(new BinaryNode<>(newEntry));
   } // end if
   return result;
} // end addEntry
```

```
Algorithm addEntry(binarySearchTree, newEntry)
  // Adds a new entry to a binary search tree that is not empty.
  // Returns null if newEntry did not exist already in the tree. Otherwise, returns
  // tree entry that matched and was replaced by newEntry.
  result = null
  currentNode = root node of binarySearchTree
  found = false
  while (found is false)
     if (newEntry matches the entry in currentNode)
         found = true
         result = entry in currentNode
         Replace entry in currentNode with newEntry
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```

```
else if (newEntry < entry in currentNode)
   if (currentNode has a left child)
      currentNode = left child of currentNode
   else
      found = true
      Give currentNode a left child containing newEntry
else // newEntry > entry in currentNode
```

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                              else // newEntry > entry in currentNode
                                                           if (currentNode has a right child)
                                                                                          currentNode = right child of currentNode
                                                           else
                                                                                          found = true
                                                                                          Give currentNode a right child containing newEntry
return result
```

```
private T addEntry(T newEntry)
   BinaryNodeInterface<T> currentNode = getRootNode();
   assert currentNode != null;
   T result = null;
   boolean found = false;
   while (!found)
      T currentEntry = currentNode.getData();
      int comparison = newEntry.compareTo(currentEntry);
      if (comparison == 0)
      { // newEntry matches currentEntry;
         // return and replace currentEntry
         found = true;
         result = currentEntry;
         currentNode.setData(newEntry);
```

```
ćűrréntNode. setLettChild(new BiháryNode<45(newEntry));
         } // end if
      else
         assert comparison > 0;
         if (currentNode.hasRightChild())
            currentNode = currentNode.getRightChild();
         else
            found = true;
            currentNode.setRightChild(new BinaryNode<+>(newEntry));
         } // end if
      } // end if
   } // end while
   return result:
} // end addEntry
```

Removing an Entry Whose Node Is a Leaf

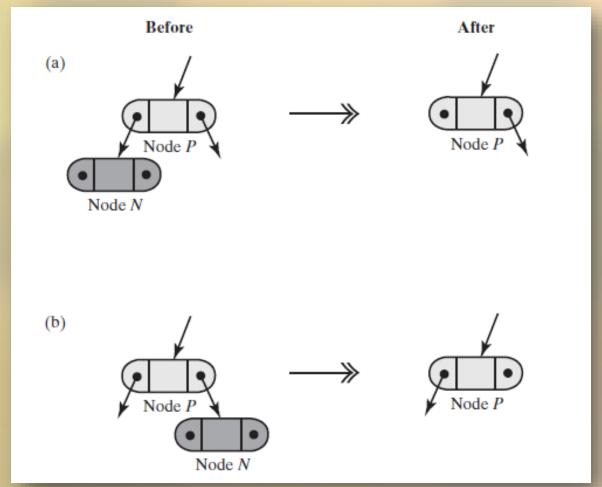


FIGURE 25-6 (a) Two possible configurations of a leaf node N; (b) the resulting two possible configurations after removing node

Removing an Entry Whose Node Has One Child

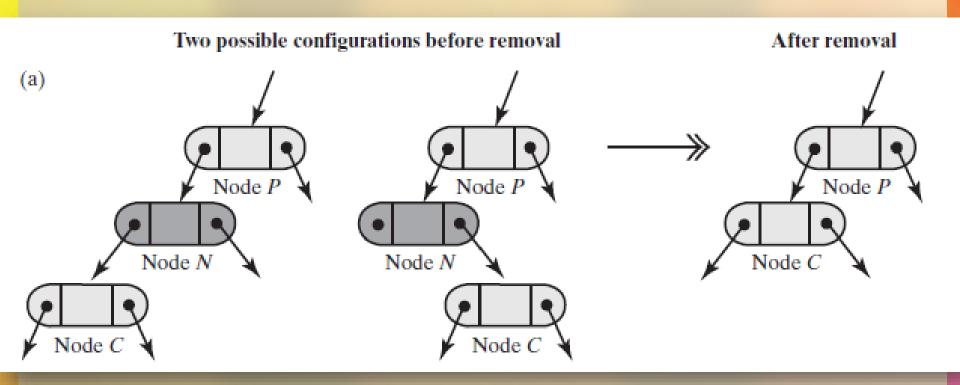


FIGURE 25-7 Removing a node N from its parent node P when N has one child and is itself (a) a left child

Removing an Entry Whose Node Has One Child

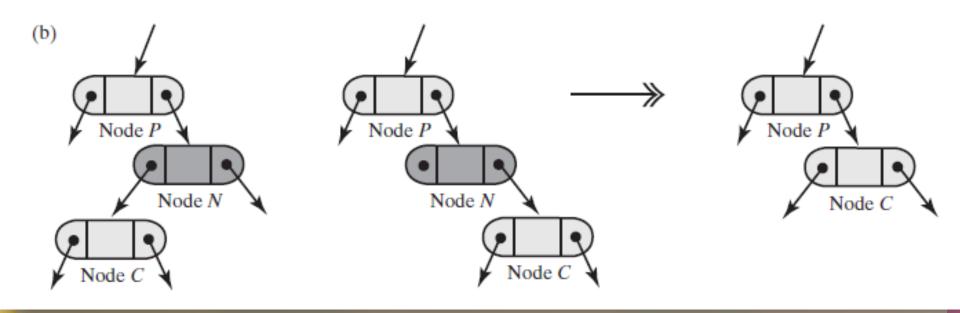
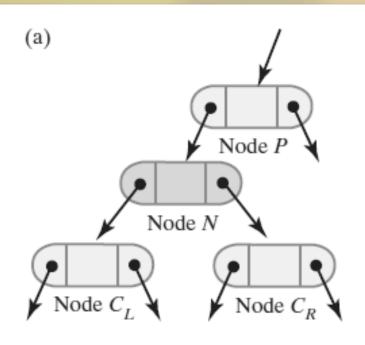
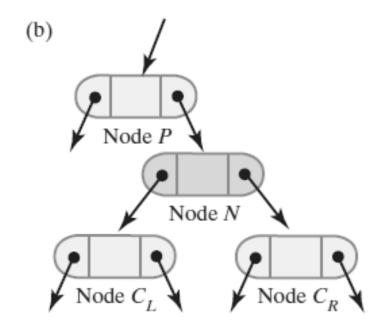


FIGURE 25-7 Removing a node N from its parent node P when N has one child and is itself (b) a right child





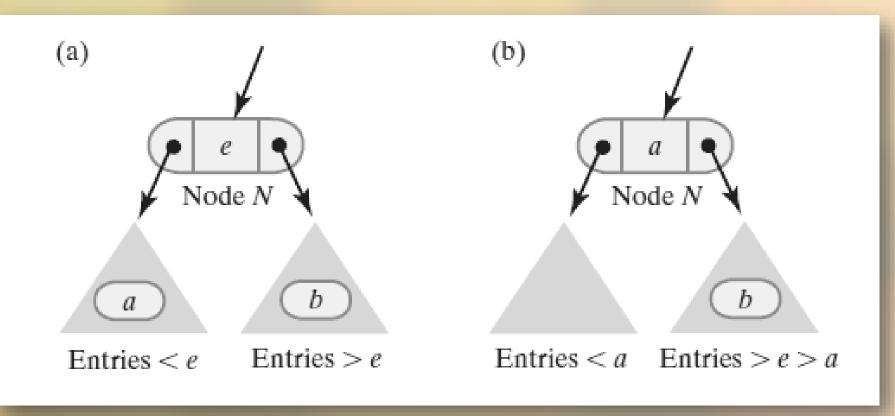


FIGURE 25-9 Node N and its subtrees: (a) the entry a is immediately before the entry e, and b is immediately after e; (b) after deleting the node that contained a and replacing e

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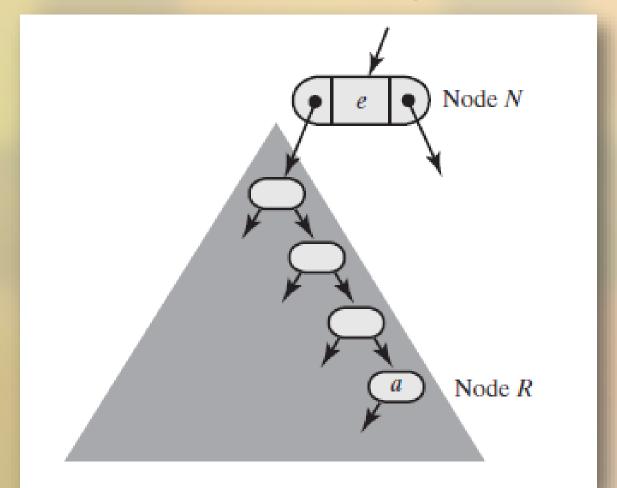


FIGURE 25-10 The largest entry a in node N's left subtree occurs in the subtree's rightmost node R

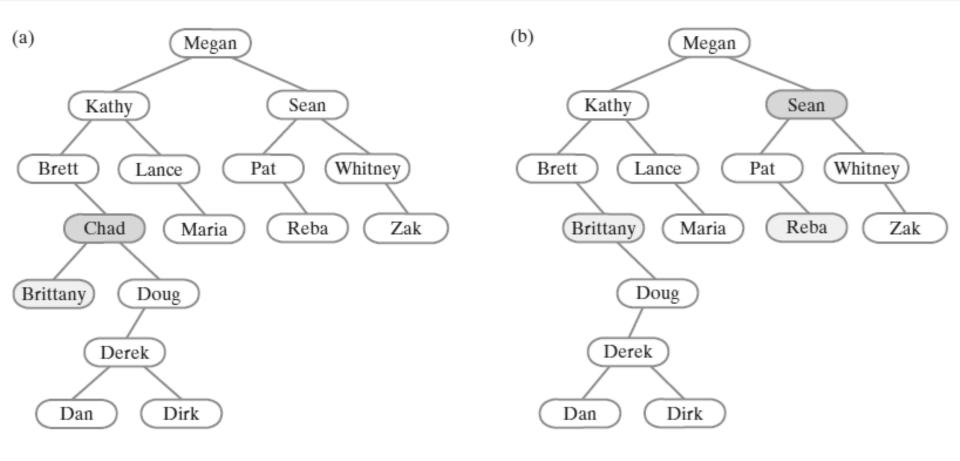


FIGURE 25-11 (a) A binary search tree; (b) after removing Chad

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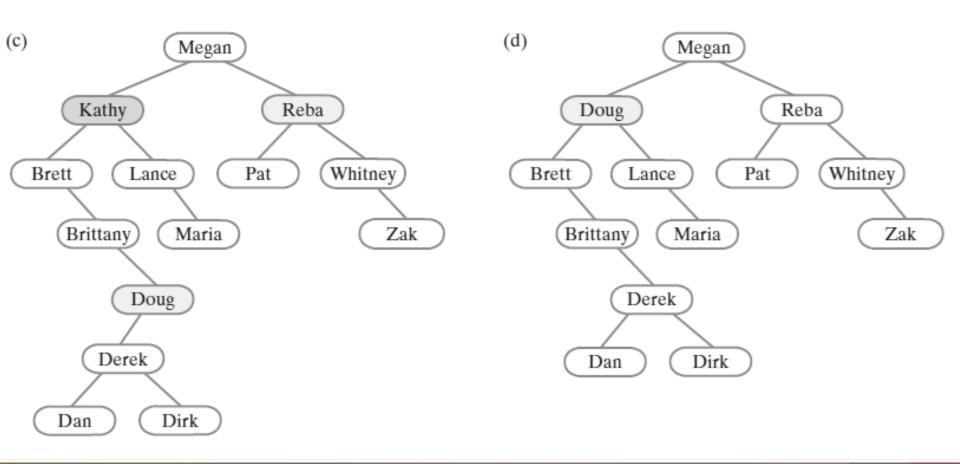


FIGURE 25-11 (c) after removing Sean; (d) after removing Kathy

Removing an Entry in the Root

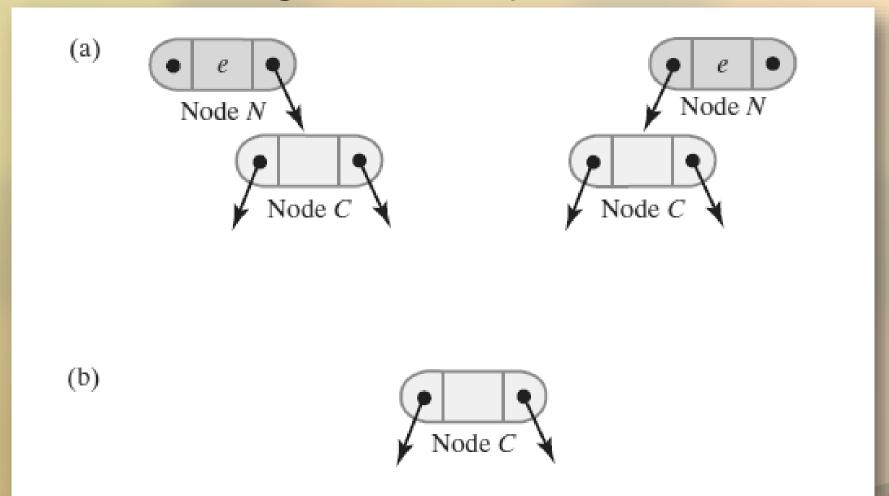


FIGURE 25-12 (a) Two possible configurations of a root that has one child; (b) after removing the root

```
Algorithm remove(binarySearchTree, entry)
oldEntry = null
if (binarySearchTree is not empty)
   if (entry matches the entry in the root of binarySearchTree)
      oldEntry = entry in root
      removeFromRoot(root of binarySearchTree)
   else if (entry < entry in root)
      oldEntry = remove(left subtree of binarySearchTree, entry)
   else // entry > entry in root
      oldEntry = remove(right subtree of binarySearchTree, entry)
return oldEntry
```

```
// Removes an entry from the tree rooted at a given node.
// rootNode is a reference to the root of a tree.
// entry is the object to be removed.
// oldEntry is an object whose data field is null.
// Returns the root node of the resulting tree; if entry matches
                                        an entry in the tree, oldEntry's data field is the entry
                                     that was removed from the tree; otherwise it is null.
private BinaryNode<T> removeEntry(BinaryNode<T> rootNode, T entry,
                                                                                                                             ReturnObject oldEntry)
           if (rootNode != null)
                      T rootData = rootNode.getData();
                      int comparison = entry.compareTo(rootData);
                      if (comparison == 0) // entry == root entry
                                 oldEntry.set(rootData);
                                 rootNode = removeFromRoot(rootNode);
                      else if (comparison < 0) // entry < root entry</pre>
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```

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       else if (comparison < 0) // entry < root entry</pre>
          BinaryNode<T> leftChild = rootNode.getLeftChild();
          BinaryNode<T> subtreeRoot = removeEntry(leftChild, entry, oldEntry);
          rootNode.setLeftChild(subtreeRoot);
       else
                                   // entry > root entry
          BinaryNode<T> rightChild = rootNode.getRightChild();
          rootNode.setRightChild(removeEntry(rightChild, entry, oldEntry));
       } // end if
   } // end if
   return rootNode:
} // end removeEntry
```

```
Algorithm removeFromRoot(rootNode)
// Removes the entry in a given root node of a subtree.
if (rootNode has two children)
   largestNode = node with the largest entry in the left subtree of rootNode
   Replace the entry in rootNode with the entry in largestNode
   Remove largestNode from the tree
else if (rootNode has a right child)
   rootNode = rootNode's right child
else
   rootNode = rootNode's left child // Possibly null
// Assertion: If rootNode was a leaf, it is now null
return rootNode
```

```
// Removes the entry in a given root node of a subtree.
// rootNode is the root node of the subtree.
// Returns the root node of the revised subtree.
private BinaryNode<T> removeFromRoot(BinaryNode<T> rootNode)
   // Case 1: rootNode has two children
   if (rootNode.hasLeftChild() && rootNode.hasRightChild())
      // Find node with largest entry in left subtree
      BinaryNode<T> leftSubtreeRoot = rootNode.getLeftChild();
      BinaryNode<T> largestNode = findLargest(leftSubtreeRoot);
      // Replace entry in root
      rootNode.setData(largestNode.getData());
      // Remove node with largest entry in left subtree
     rootNode_setLeftChild(removeLargest(leftSubtreeRoot));
```

```
rootNode.setData(largestNode.getData());
      // Remove node with largest entry in left subtree
      rootNode.setLeftChild(removeLargest(leftSubtreeRoot));
   } // end if
   // Case 2: rootNode has at most one child
   else if (rootNode.hasRightChild())
      rootNode = rootNode.getRightChild();
   else
      rootNode = rootNode.getLeftChild();
   // Assertion: If rootNode was a leaf, it is now null
   return rootNode;
} // end removeEntry
```

```
// Finds the node containing the largest entry in a given tree.
// rootNode is the root node of the tree.
// Returns the node containing the largest entry in the tree.
private BinaryNode<T> findLargest(BinaryNode<T> rootNode)
{
   if (rootNode.hasRightChild())
     rootNode = findLargest(rootNode.getRightChild());
   return rootNode;
} // end findLargest
```

```
// Removes the node containing the largest entry in a given tree.
// rootNode is the root node of the tree.
// Returns the root node of the revised tree.
private BinaryNode<T> removeLargest(BinaryNode<T> rootNode)
   if (rootNode.hasRightChild())
      BinaryNode<T> rightChild = rootNode.getRightChild();
      rightChild = removeLargest(rightChild);
      rootNode.setRightChild(rightChild);
  else
      rootNode = rootNode.getLeftChild();
   return rootNode;
} // end removeLargest
```

```
Algorithm remove(entry)
result = null
currentNode = node that contains a match for entry
parentNode = currentNode's parent
if (currentNode != null) // That is, if entry is found
   result = currentNode's data (the entry to be removed from the tree)
   // Case 1
   if (currentNode has two children)
      // Get node to remove and its parent
       nodeToRemove = node containing entry's inorder predecessor; it has at most one child
       parentNode = nodeToRemove's parent
       Copy entry from nodeToRemove to currentNode
       currentNode = nodeToRemove
      // Assertion: currentNode is the node to be removed; it has at most one child
      // Assertion: Case 1 has been transformed to Case 2
   // Case 2: currentNode has at most one child
   Delete currentNode from the tree
return result
```

```
public T remove(T entry)
  T result = null;
  // Locate node (and its parent) that contains a match for entry
  NodePair pair = findNode(entry);
  BinaryNode<T> currentNode = pair.getFirst();
   BinaryNode<T> parentNode = pair.getSecond();
   if (currentNode != null) // Entry is found
      result = currentNode.getData(); // Get entry to be removed
      // Case 1: currentNode has two children
      if (currentNode.hasLeftChild() && currentNode.hasRightChild())
        // Replace entry in currentNode with the entry in another node
         // that has at most one child; that node can be deleted
```

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```
// Get node to remove (contains inorder predecessor; has at
         // most one child) and its parent
         pair = getNodeToRemove(currentNode);
         BinaryNode<T> nodeToRemove = pair.getFirst();
         parentNode = pair.getSecond();
         // Copy entry from nodeToRemove to currentNode
         currentNode.setData(nodeToRemove.getData());
         currentNode = nodeToRemove;
         // Assertion: currentNode is the node to be removed; it has at
                       most one child
         // Assertion: Case 1 has been transformed to Case 2
      } // end if
      // Case 2: currentNode has at most one child; delete it
      removeNode(currentNode, parentNode);
   } // end if
   return result;
} // end remove
```

```
private NodePair findNode(T entry)
{
   NodePair result = new NodePair();
   boolean found = false;
   . . .
   if (found)
      result = new NodePair(currentNode, parentNode);
      // Located entry is currentNode.getData()
   return result;
} // end findNode
```

The private method findNode

```
// Find the inorder predecessor by searching the left subtree; it will be the largest
// entry in the subtree, occurring in the node as far right as possible
leftSubtreeRoot = left child of currentNode
rightChild = leftSubtreeRoot
priorNode = currentNode
while (rightChild has a right child)
{
    priorNode = rightChild
    rightChild = right child of rightChild
}
// Assertion: rightChild is the node to be removed and has no more than one child
```

Pseudocode for the private method getNodeToRemove

```
private NodePair getNodeToRemove(BinaryNode<T> currentNode)
   // Find node with largest entry in left subtree by
   // moving as far right in the subtree as possible
   BinaryNode<T> leftSubtreeRoot = currentNode.getLeftChild();
   BinaryNode<T> rightChild = leftSubtreeRoot;
   BinaryNode<T> priorNode = currentNode;
   while (rightChild.hasRightChild())
      priorNode = rightChild;
      rightChild = rightChild.getRightChild();
   } // end while
   // rightChild contains the inorder predecessor and is the node to
   // remove; priorNode is its parent
   return new NodePair(rightChild, priorNode);
} // end getNodeToRemove
```

```
private void removeNode(BinaryNode<T> nodeToRemove, BinaryNode<T> parentNode)
  BinaryNode<T> childNode;
   if (nodeToRemove.hasLeftChild())
      childNode = nodeToRemove.getLeftChild();
  else
      childNode = nodeToRemove.getRightChild();
   // Assertion: If nodeToRemove is a leaf, childNode is null
   assert (nodeToRemove.isLeaf() && childNode == null) ||
          !nodeToRemove.isLeaf();
   if (nodeToRemove == getRootNode())
      setRootNode(childNode);
   else if (parentNode.getLeftChild() == nodeToRemove)
      parentNode.setLeftChild(childNode);
  else
      parentNode.setRightChild(childNode);
} // end removeNode
```

Efficiency of Operations

- For tree of height h
 - The operations add, remove, and getEntry are O(h)
- If tree of n nodes has height h = n
 - These operations are O(n)
- Shortest tree is full
 - Results in these operations being O(log n)

Efficiency of Operations

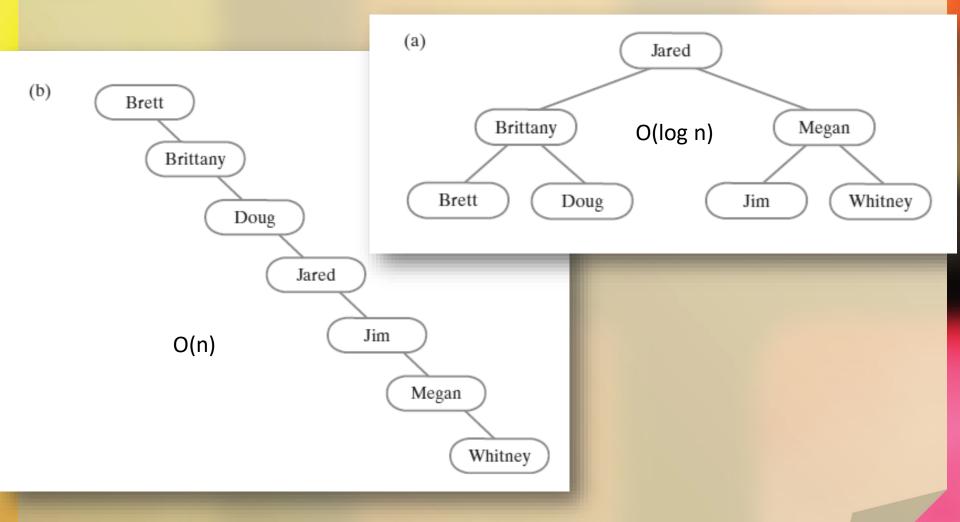


FIGURE 25-13 Two binary search trees that © 2016 Pearson Education, Ltd. All rights reserved. contain the same data

The Importance of Balance

- Do not need a full binary search tree to get O(log n) performance
- Complete tree will also give us O(log n) performance.

Order in Which Nodes Are Added

- Order in which you add entries to a binary search tree affects the shape of the tree
- If you add entries into an initially empty binary search tree, do not add them in sorted order.

End

Chapter 25