

CENG111 Concepts in Computer Engineering

Chapter 3 Operating Systems



Chapter 3 Operating Systems

3.1 The History of Operating Systems

3.2 Operating System Architecture

3.3 Coordinating the Machine's Activities

3.4 Handling Competition Among Processes

3.5 Security



Computer System Structure

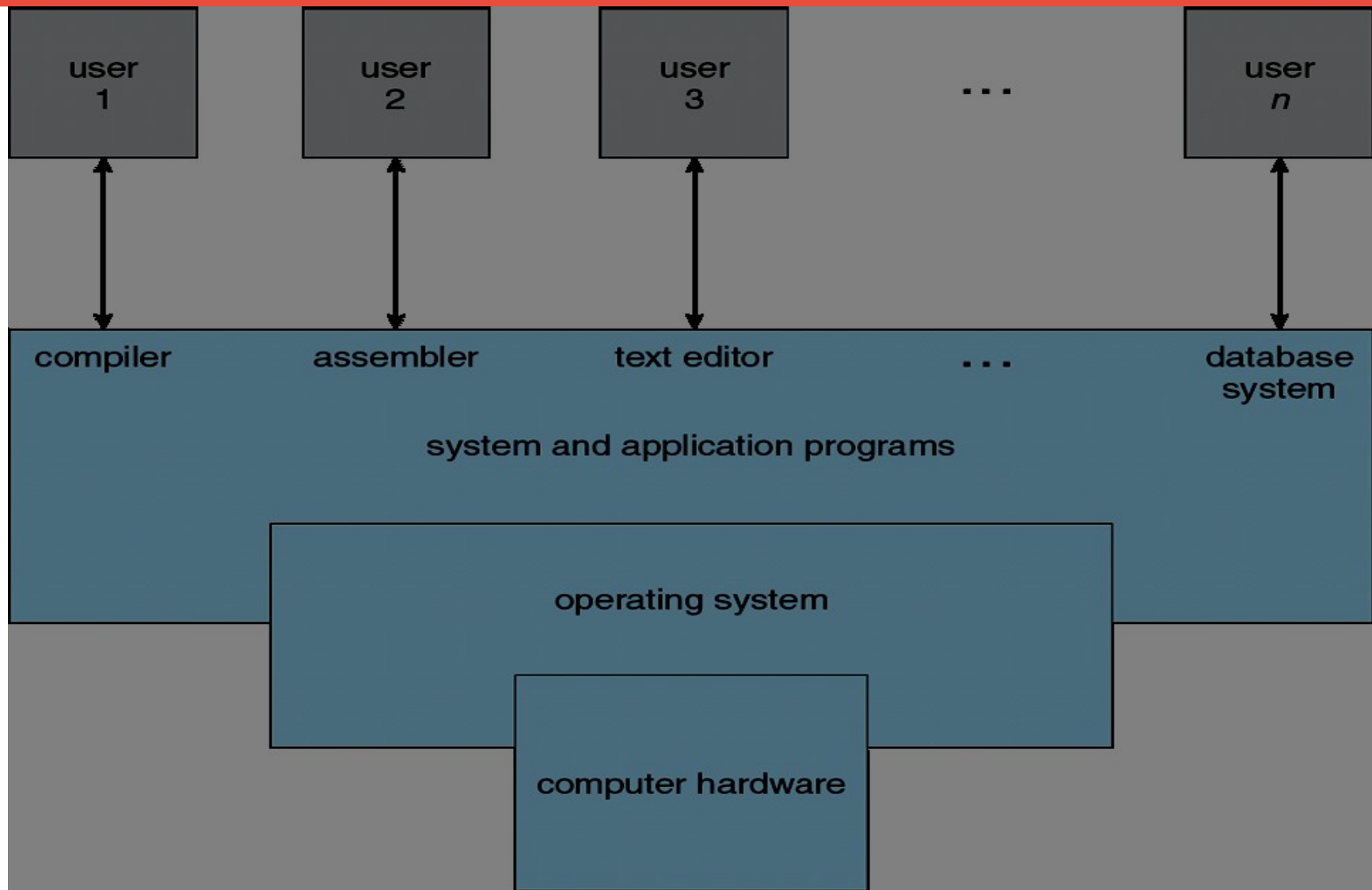
Hardware - provides basic computing resources
CPU, memory, I/O devices

Operating system - Controls and coordinates use of hardware among various applications and users

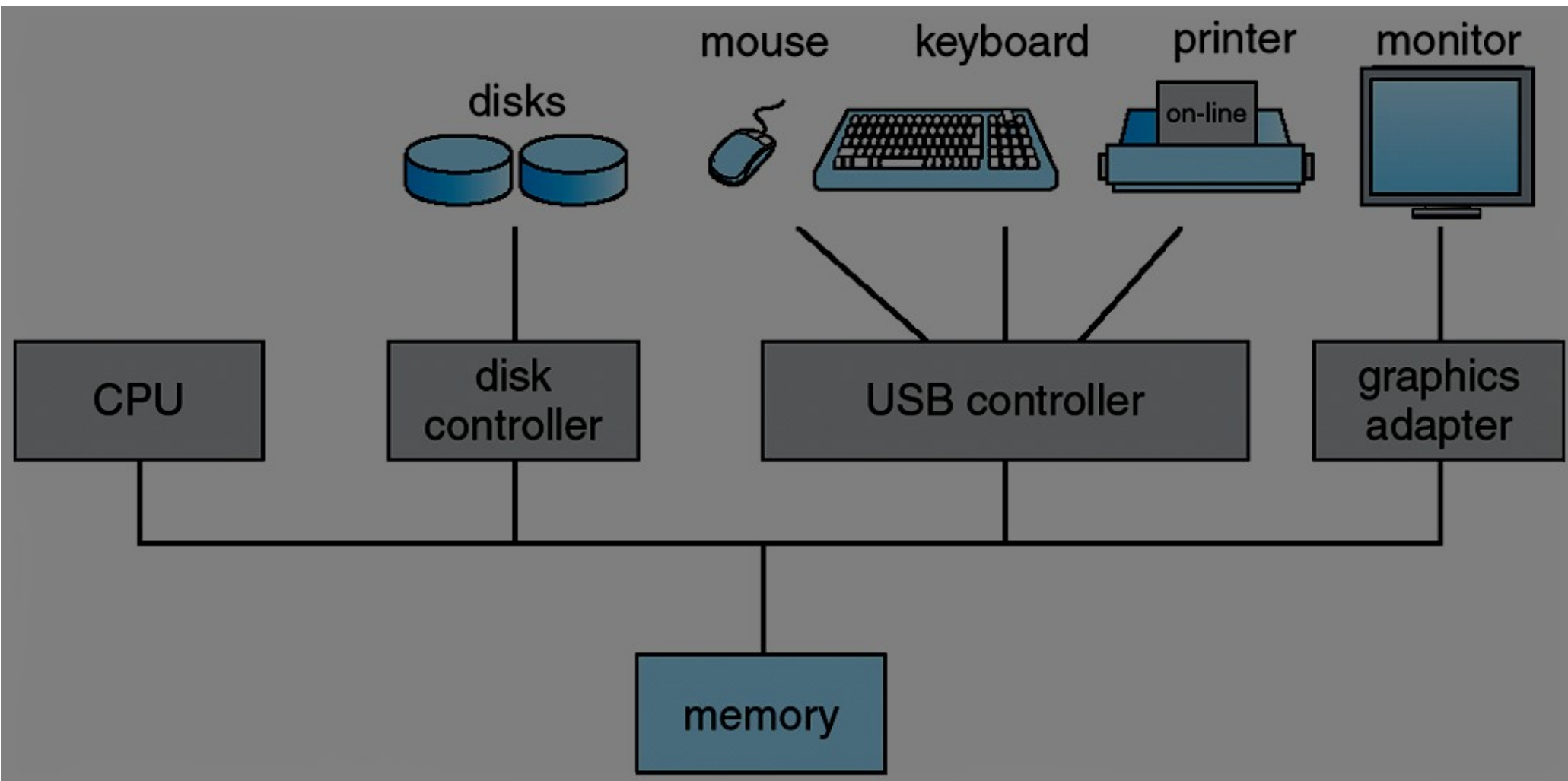
Application programs - define the ways in which the system resources are used to solve the computing problems of the users; Word processors, compilers, web browsers, database systems, video games

Users; People, machines, other computers

Computer System Structure



Computer System Organization



Top-Level Components

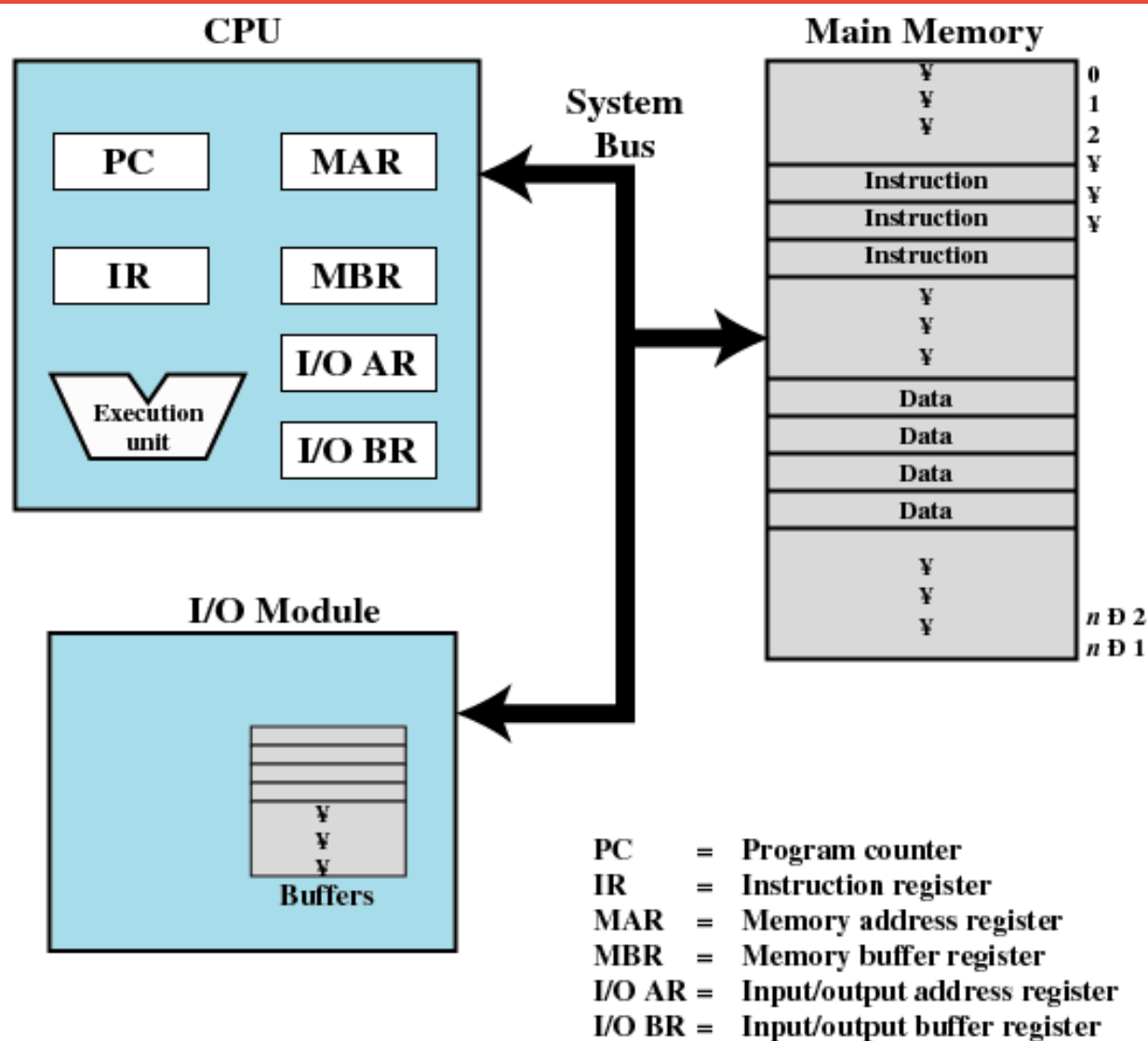


Figure 1.1 Computer Components: Top-Level View

Central Processing Unit (CPU)

Processor (CPU) to perform computations:

Executes a sequence of stored instructions called a program, program is kept in the main memory

Control unit : responsible for deciding which instruction in a program should be executed

Arithmetic and logic unit (ALU) : responsible for executing the actual instructions

Register : quickly accessible location available to CPU

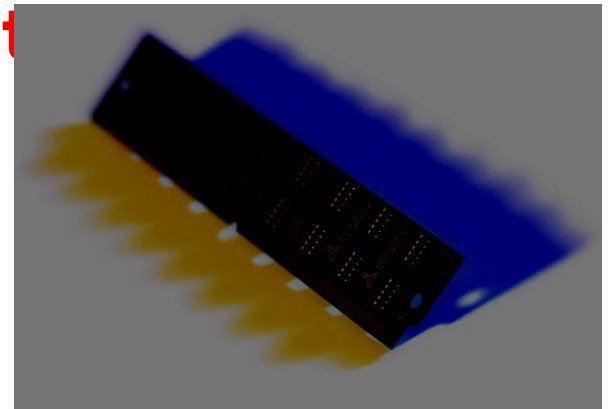
Main Memory

Collection of locations, each of which is capable of storing both instructions and data

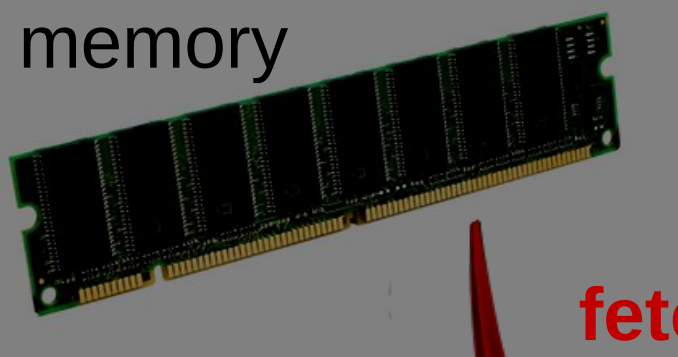
Every location consists of an address, which is used to access the location, and the contents of the location

OS sees and manages memory

Programs/data need to be brought to



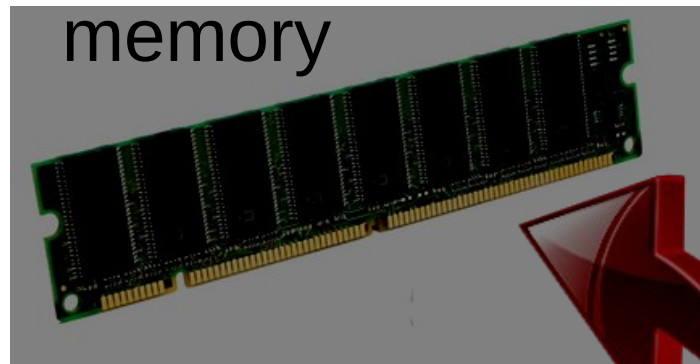
memory



fetch/read



CPU



write/store



CPU

Instruction Cycle

The processor fetches the instruction from memory

Program counter (PC) holds address of the instruction to be fetched next

Program counter is incremented after each fetch

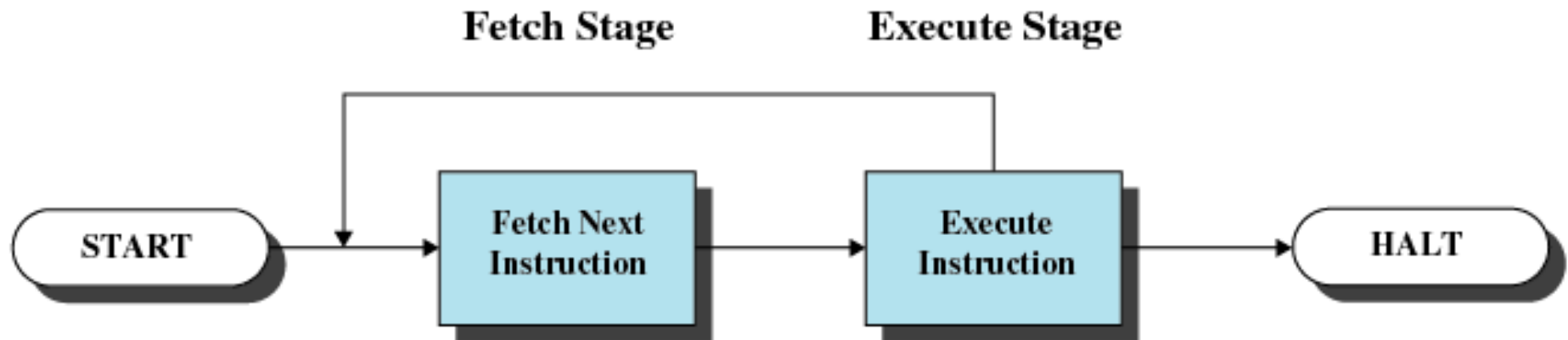
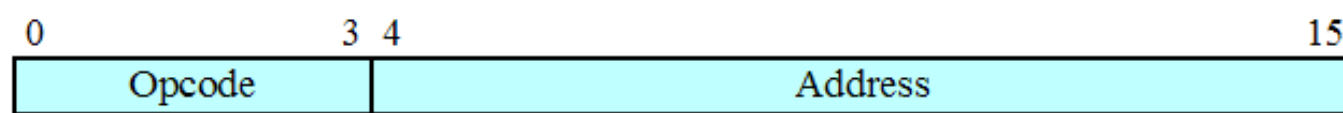
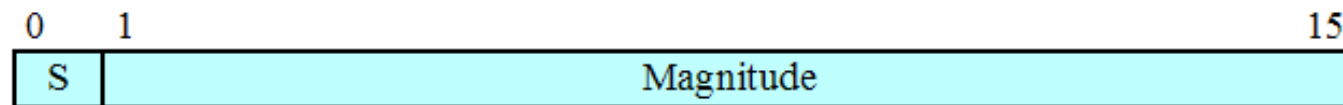


Figure 1.2 Basic Instruction Cycle

Characteristics of a Hypothetical Machine



(a) Instruction format



(b) Integer format

Program Counter (PC) = Address of instruction
Instruction Register (IR) = Instruction being executed
Accumulator (AC) = Temporary storage

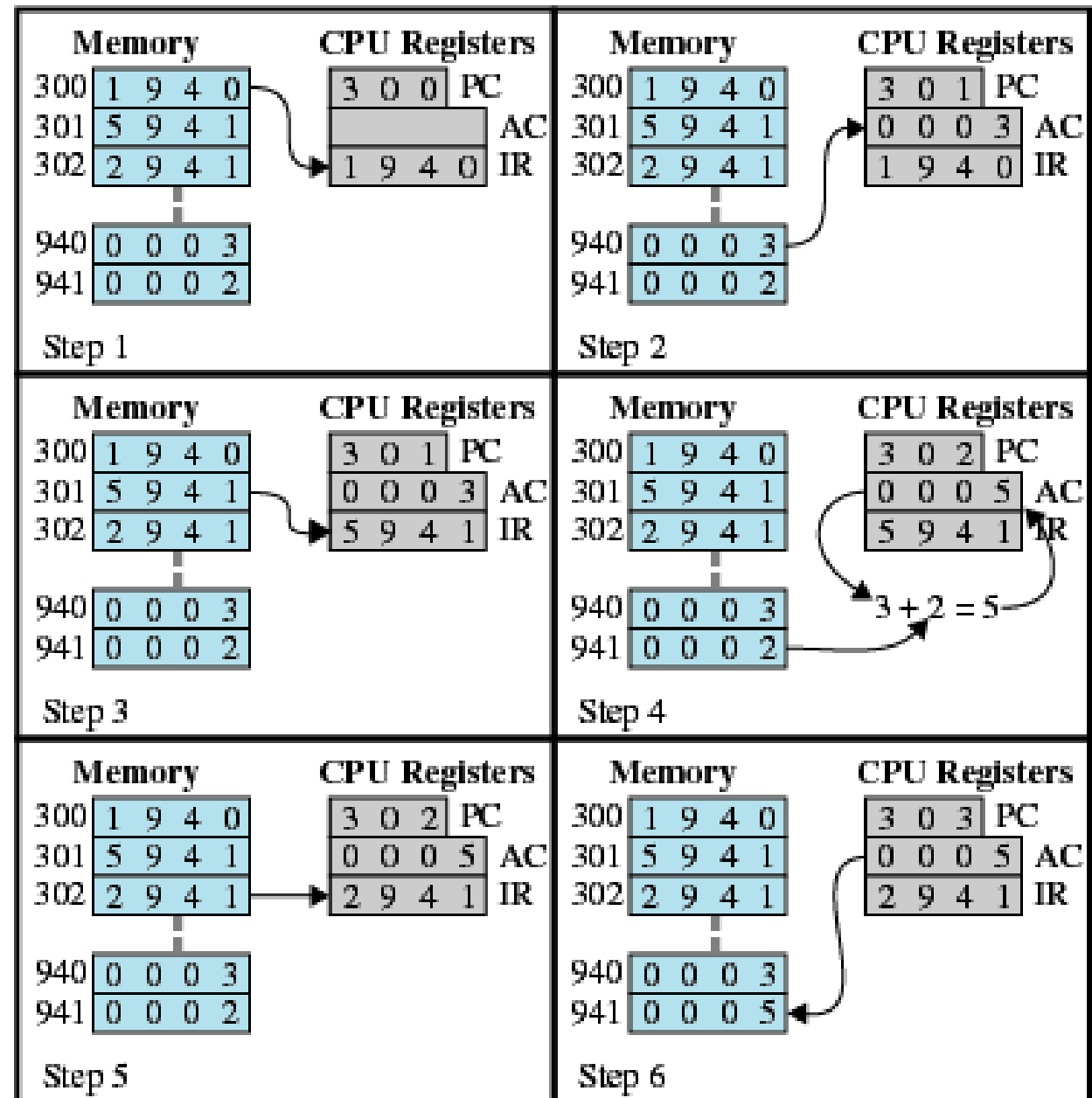
(c) Internal CPU registers

0001 = Load AC from Memory
0010 = Store AC to Memory
0101 = Add to AC from Memory

(d) Partial list of opcodes

Example of Program Execution

0001 – Load AC from Memory
0010 – Store AC to Memory
0101 – Add to AC from Memory



I/O Devices

Large variety, varying speeds: Disk, monitor, printer

Each has a controller:

- Hides low-level details from OS**

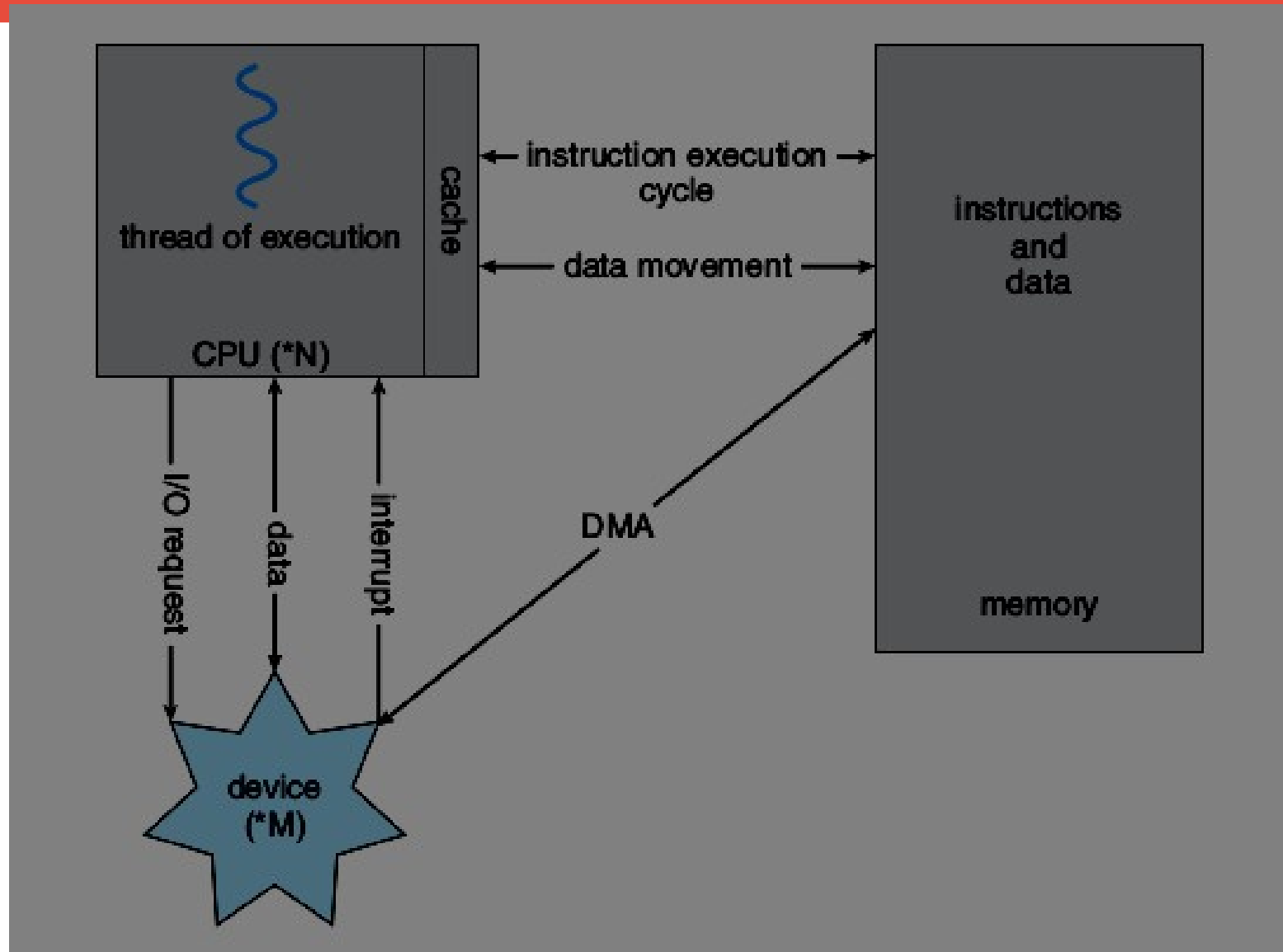
- Manages data flow between device and CPU/memory**

Interconnects

Flow of data and information in the system

Switches map a fixed number of inputs to outputs, total number of ports on a switch

How a Computer Works?



Why to Learn How OS works?

Knowing how operating systems work is a fundamental and critical to anyone who is a serious software developer

When you write a program and it runs too slow, but you see nothing wrong with your code, where else will you look for a solution

Understand **system performance**

Operating System

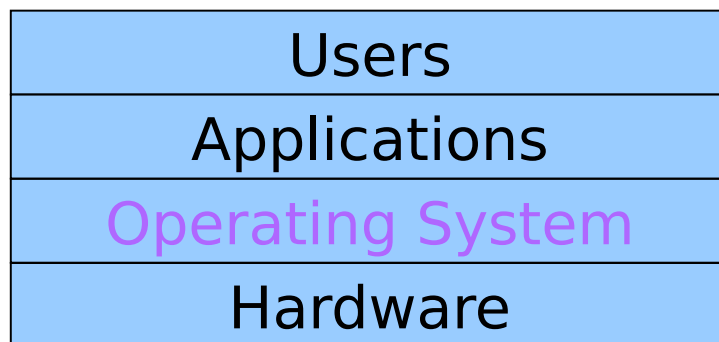
No single definition

OS = all software minus applications

Operating system goals:

Helps the applications use the computer

Helps manage the resources of the computer in an efficient manner



Operating System Evolution

Single user

Batch processing

Multiprogramming

Graphical user interface

Single User

One user at a time on console

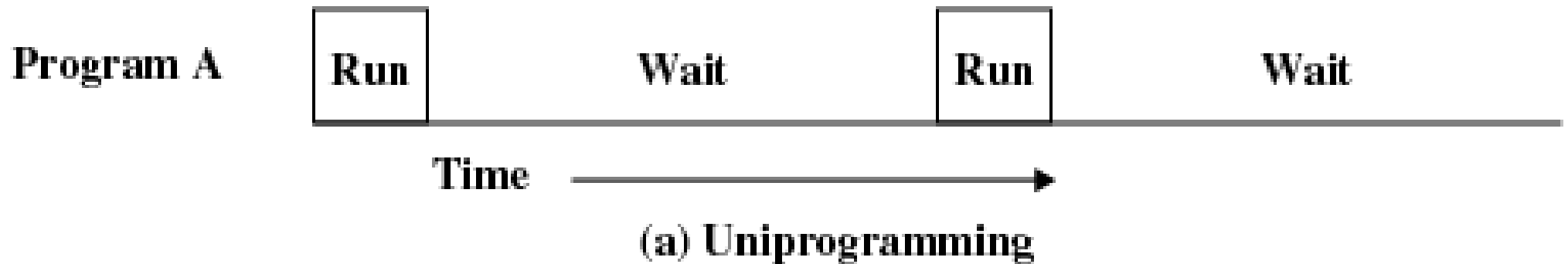
Computer executes one function at a time

No overlap: computation & I/O

User must be at console

Uniprogramming (~1960)

Processor must wait for I/O instruction to complete before proceeding



Batch Processing

Execute multiple consecutive “jobs” in batch:

Load program

Run

Print results, dump machine state

Repeat

Users submit jobs (on cards or tape)

Human schedules jobs

OS loads and runs jobs

Multiprogramming

Allow several programs to run at the same time

Run one job until I/O

Run another job, etc.

OS manages interaction between programs

Which jobs to start

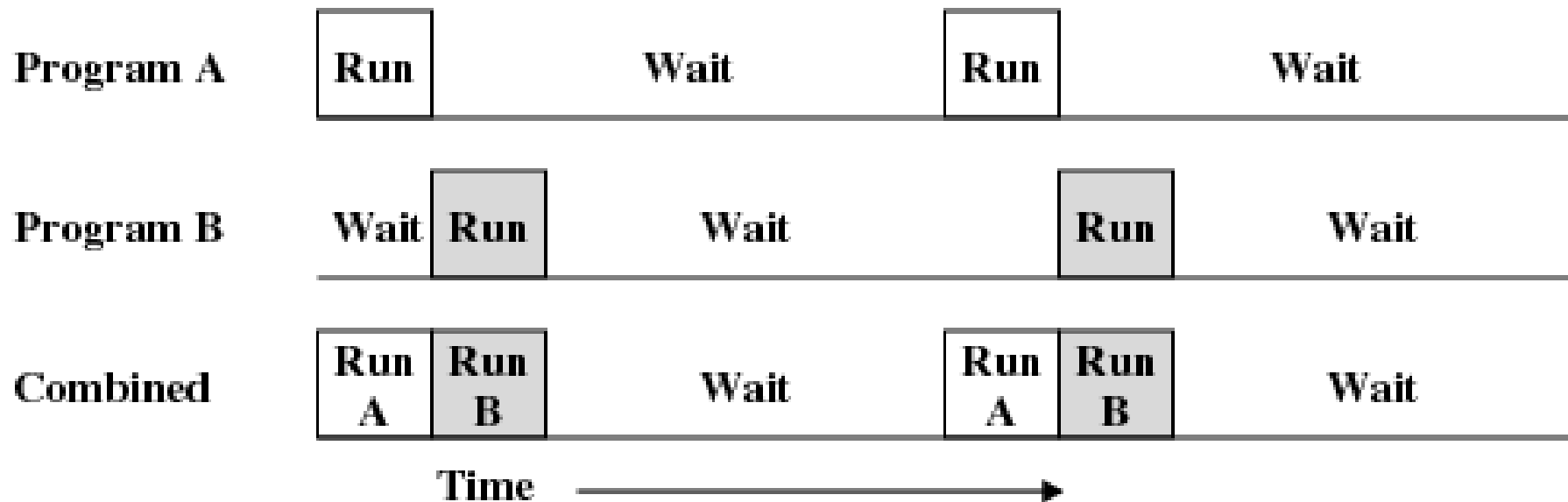
Protects program's memory from others

Decides which to resume when CPU available

TIME-SHARING

Multiprogramming (~1961)

When one job needs to wait for I/O, the processor can switch to the other job



(b) Multiprogramming with two programs

Graphical User Interface

Separate window for each application
Menu-driven application environment

Operating System History

Multics (1960s), Unix

**Users share system via terminals
(mainframes)**

DOS, MacOS (1980s)

**Simple OS in PCs without
multiprogramming, concurrency, memory
protection, virtual memory**

NT, MacOS X, Linux (1990s)

“Real” operating systems on PCs

Today

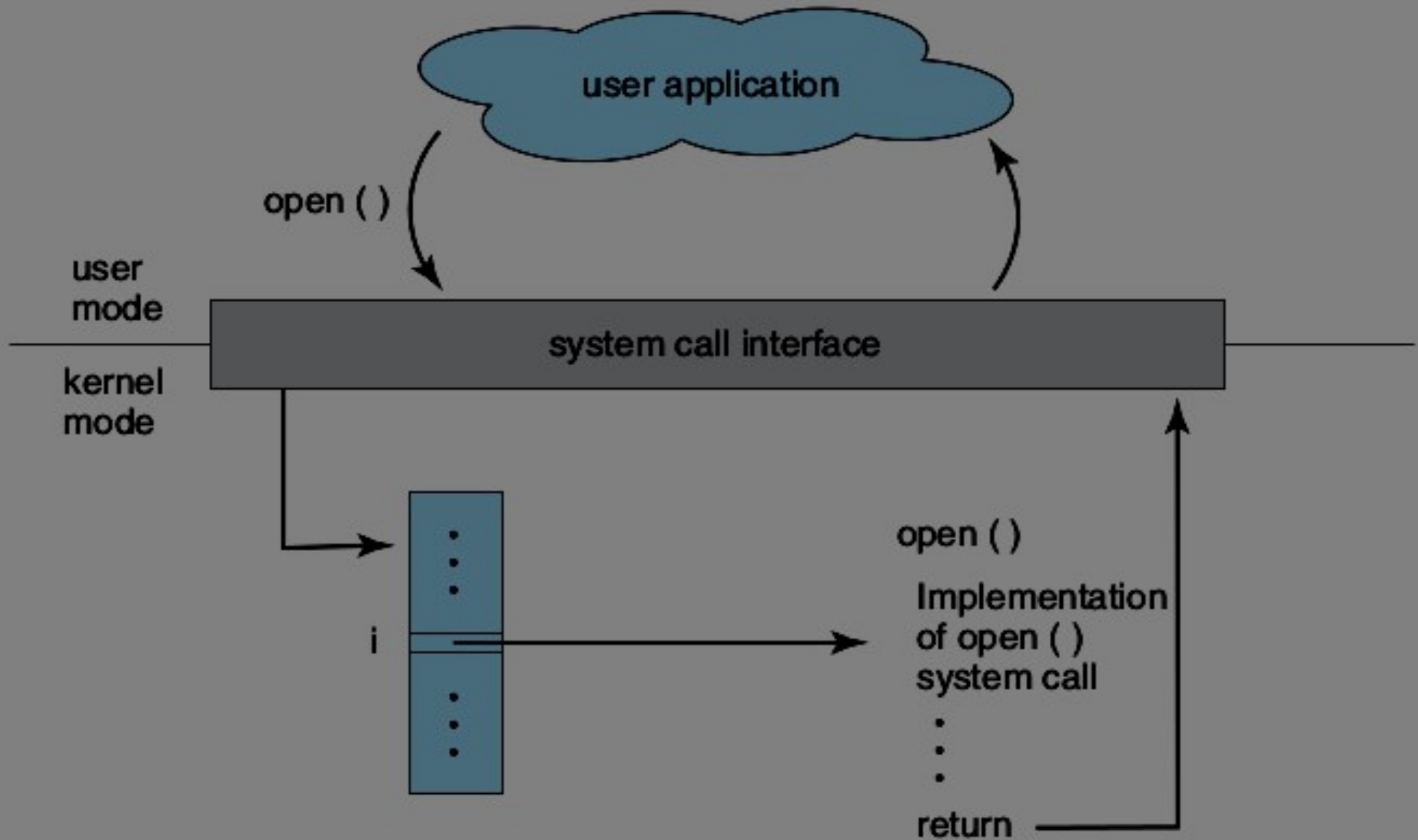
Parallel systems, mobile devices

System Calls

Programming interface to the services provided by the OS

Mostly accessed by programs via a high-level Application Programming Interface (API) rather than direct system call use

API - System Call - OS Relationship



Operating System Functionality

- **Certain services to users**
- **Load programs, run them**
- **Controls execution of programs to prevent errors and improper use of the computer**
- **Manages all resources (e.g., CPU, memory, disk)**
- **Decides between conflicting requests for efficient and fair resource use (prevent of applications from starvation and deadlock conditions)**

Operating System Services - Users

- **User interface**
- **Program execution**
- **I/O operations**
- **File system manipulation**
- **Communication**
- **Error detection**

Operating System Services - Resources

- **Resource allocation for processes (system and users)**
- **Accounting**
- **Protection and security**

Operating System Design Goals

User view:

convenient to use, easy to learn and to use, reliable, safe, and fast

System view:

**easy to design, implement, and maintain;
and flexible, reliable, error free, and
efficient**

Specifying and designing an OS:

**Highly creative task of software
engineering**

Operating System Implementation

Early operating systems were written in assembly language

Now mostly C, C++

Operating System Components

- **Process scheduler - CPU**
 - **Determines when and for long each process executes**
- **Memory manager - Main memory**
 - **Determines when and how memory is allocated to processes**
 - **Decides what to do when main memory is full**
- **File system - I/O**
 - **Organizes named collections of data in persistent storage**
- **Networking - Interconnect**
 - **Enables processes to communicate with one another**
- **Protection and security**

Process Management

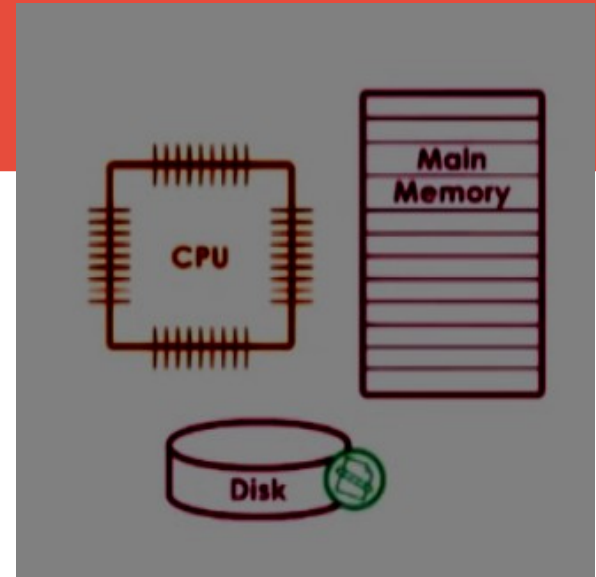
**Process needs resources to accomplish its task;
CPU, memory, I/O**

Process management issues;

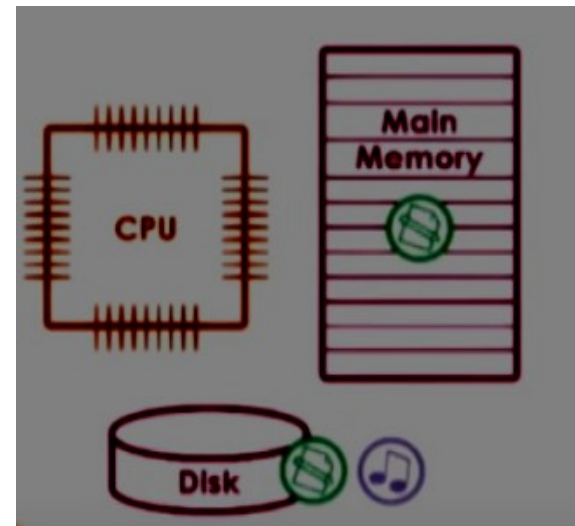
- **Scheduling processes and threads on the CPUs**
- **Creating and deleting both user and system processes**
- **Suspending and resuming processes**
- **Providing mechanisms for process synchronization and communication**

Process

application=program on disk

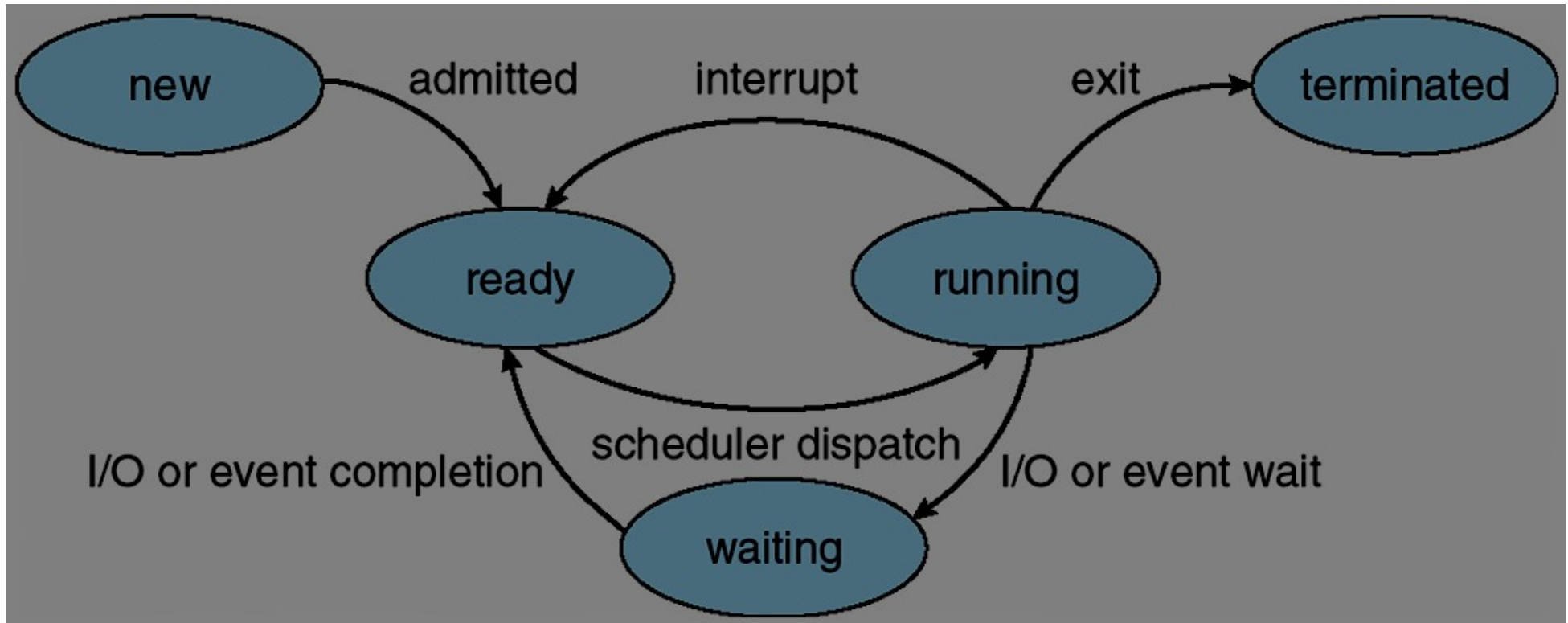


**Program becomes process when
executable file loaded
into memory**

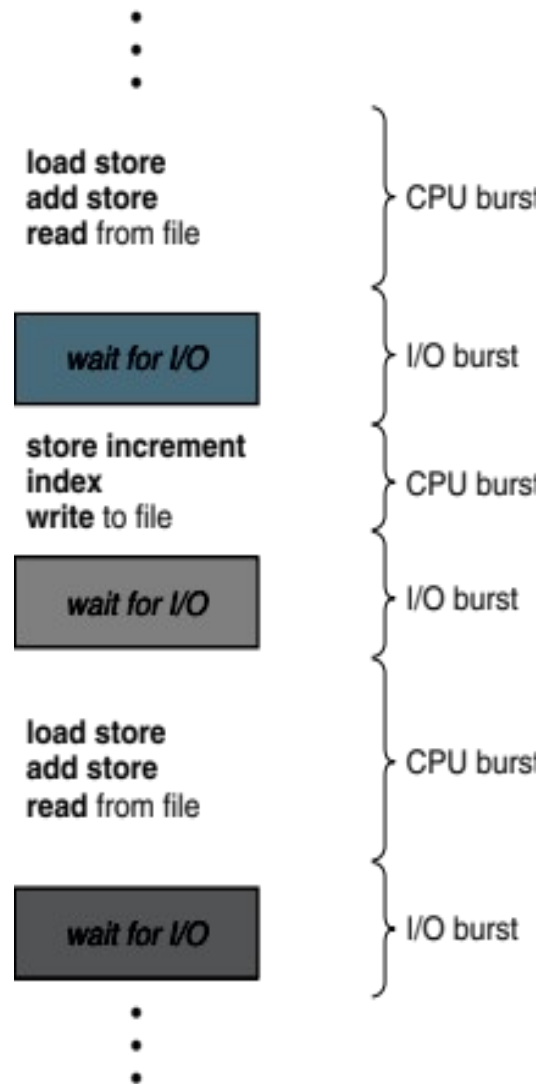


process=program in execution

Process Lifecycle



Burst Times



Process Scheduler

Determines which one of the currently ready processes will be dispatched to the CPU to start running, and how long it should run for

OS must;

- **Preempt:** interrupt and save current context
- **Schedule:** run scheduler to choose next process
- **Dispatch:** dispatch process and switch into its context

Maximize CPU use, quickly switch processes onto CPU for time sharing

Be efficient!

First- Come, First-Served (FCFS) Scheduling

| <u>Process</u> | <u>Burst Time</u> |
|----------------|-------------------|
|----------------|-------------------|

| | |
|-------|----|
| P_1 | 24 |
|-------|----|

| | |
|-------|---|
| P_2 | 3 |
|-------|---|

| | |
|-------|---|
| P_3 | 3 |
|-------|---|

Suppose that the processes arrive in the order:

P_1, P_2, P_3

The Gantt Chart for the schedule is:



Waiting time for $P_1 = 0$; $P_2 = 24$; $P_3 = 27$

Average waiting time: $(0 + 24 + 27)/3 = 17$

First- Come, First-Served (FCFS) Scheduling

Suppose that the processes arrive in the order:

P_2, P_3, P_1

The Gantt chart for the schedule is:



Waiting time for $P_1 = 6$; $P_2 = 0$; $P_3 = 3$

Average waiting time: $(6 + 0 + 3)/3 = 3$

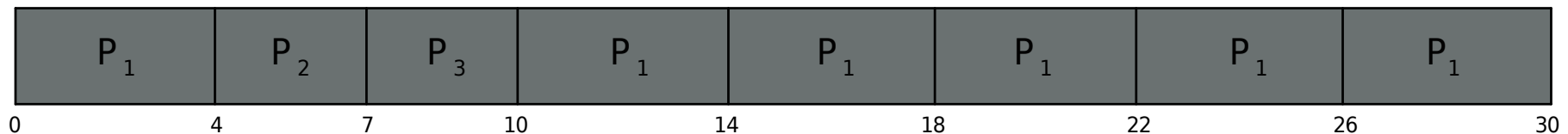
Much better than previous case

Round Robin

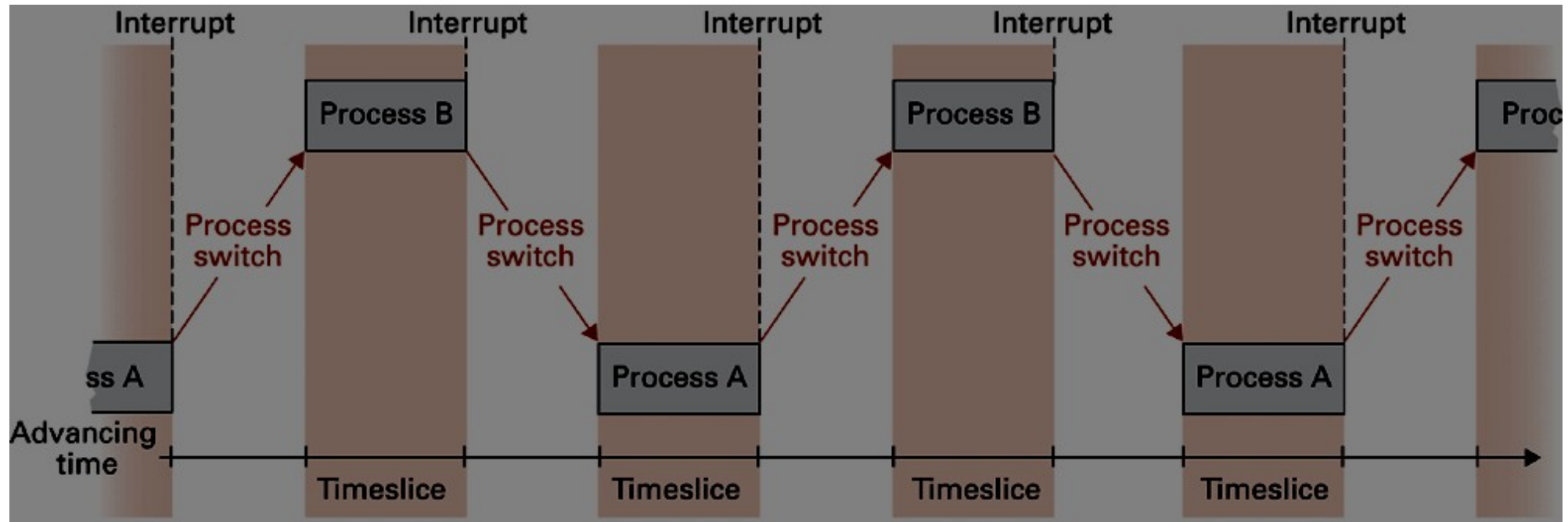
Each process gets a small unit of CPU time (**time quantum q**), usually 10-100 milliseconds. After this time has elapsed, the process is preempted and added to the end of the ready queue.

RR with $q=4$:

The Gantt chart is:

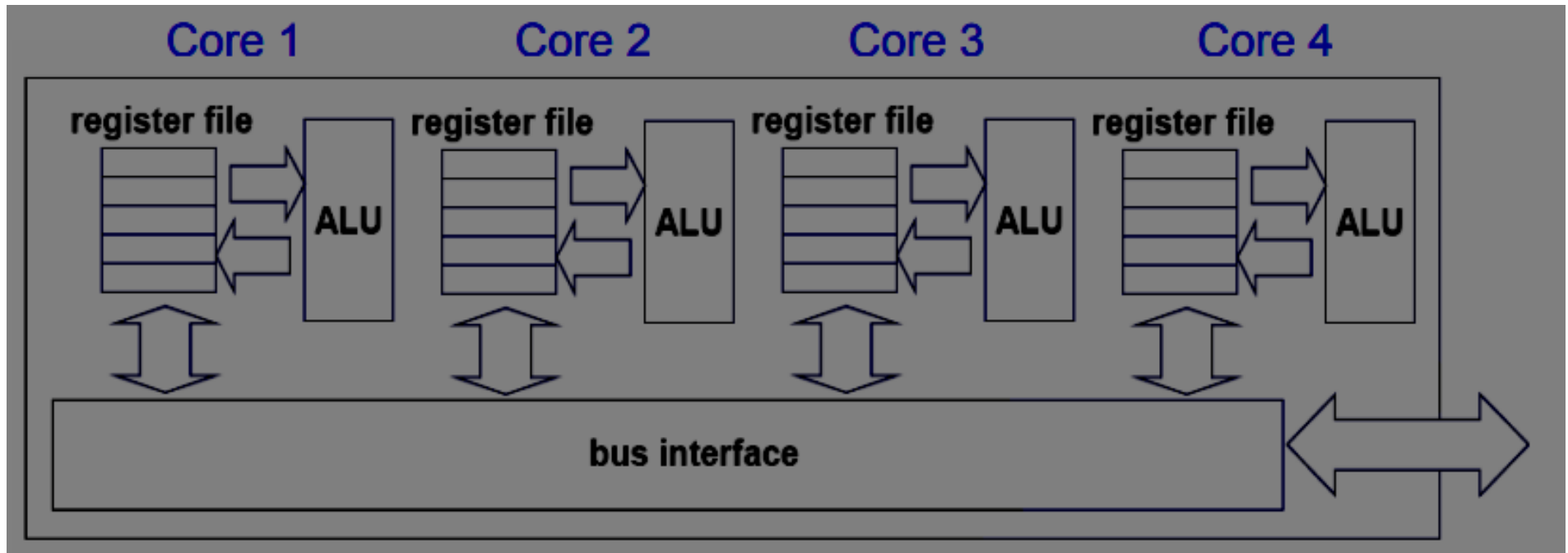


Time-Sharing Between Two Processes



Multicore Processors

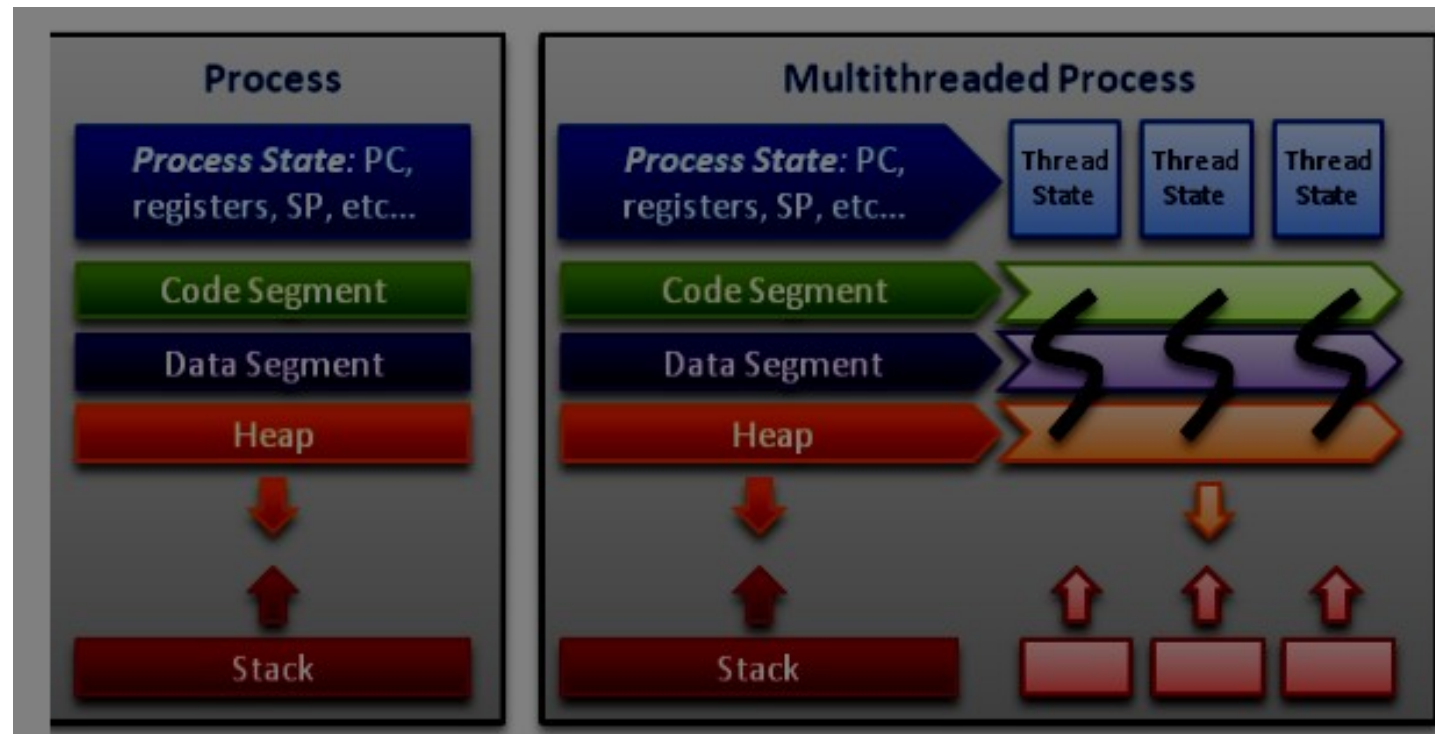
Multiple CPUs on a single chip



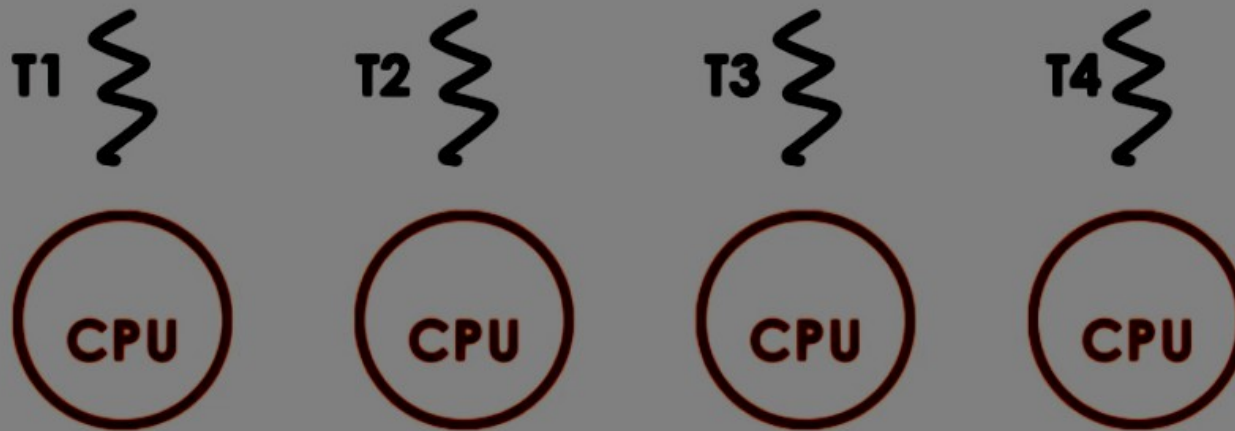
Threads

Threads use, and exist within, the process resources

Scheduled and run as independent entities



Benefits of Multithreading



Input Matrix

| |
|----|
| T1 |
| T2 |
| T3 |
| T4 |

Parallelization --> speed up

Specialization --> hot cache!

**Efficiency --> lower memory requirement,
cheaper communication overhead**

Memory Management

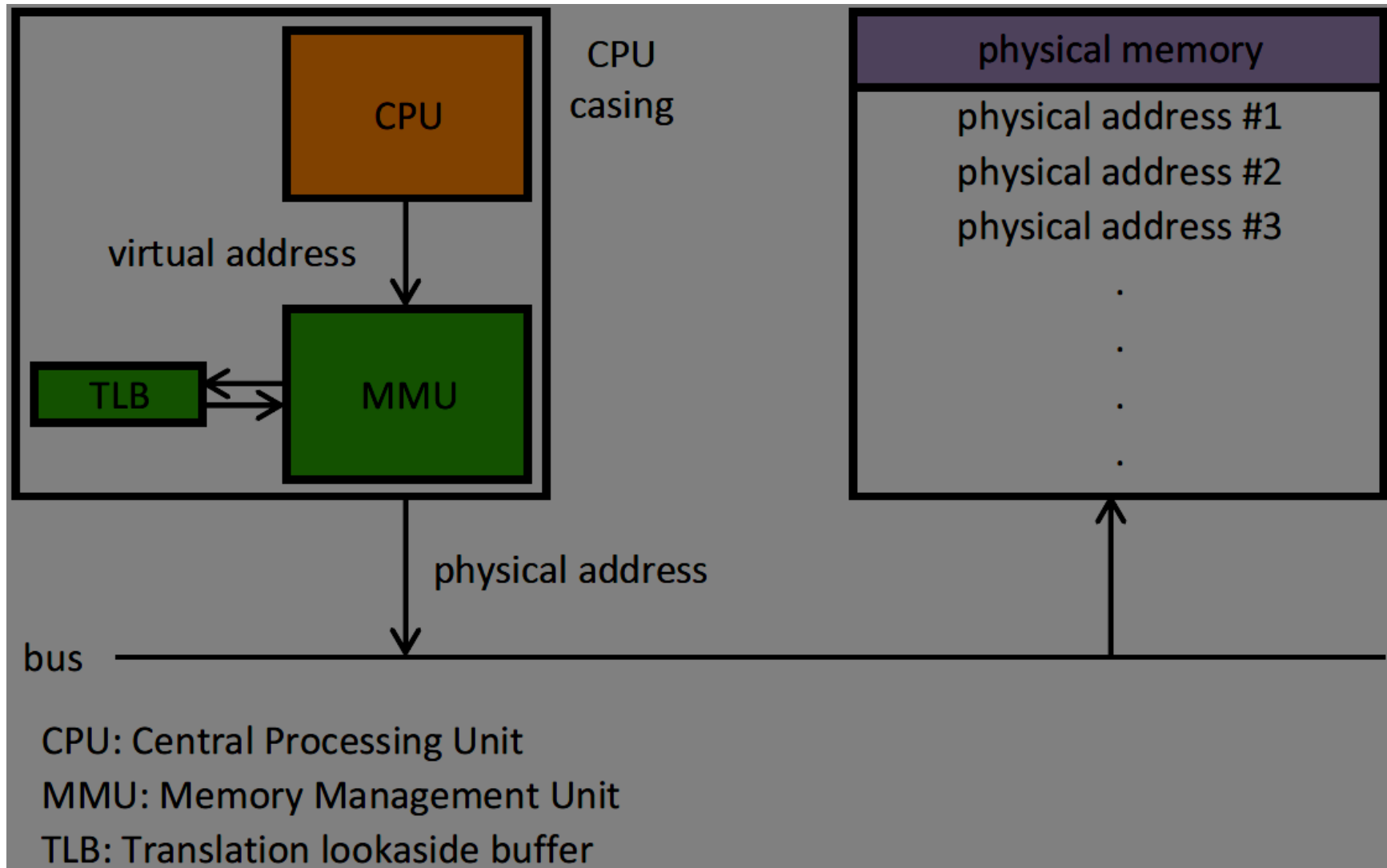
To execute a program all (or part) of the instructions must be in memory

All (or part) of the data that is needed by the program must be in memory

Memory management issues:

- **Keeping track of which parts of memory are currently being used and by whom**
- **Deciding which processes and data to move into and out of memory**
- **Allocating and deallocating memory space as needed**

Memory Management Unit



Storage Management

File system management:

- Creating and deleting files
- Supporting primitives for manipulating files and directories
- Mapping files onto secondary storage
- Backing up files on stable (nonvolatile) storage media

Mass storage management:

- Usually disks used to store data that does not fit in main memory or data that must be kept for a “long” period of time

I/O subsystem:

- To hide features of hardware devices from the user

- **49 Drivers for hardware devices**

Networking

Let another process know that some event has occurred

Transfer data from one process to another

Protection and Security

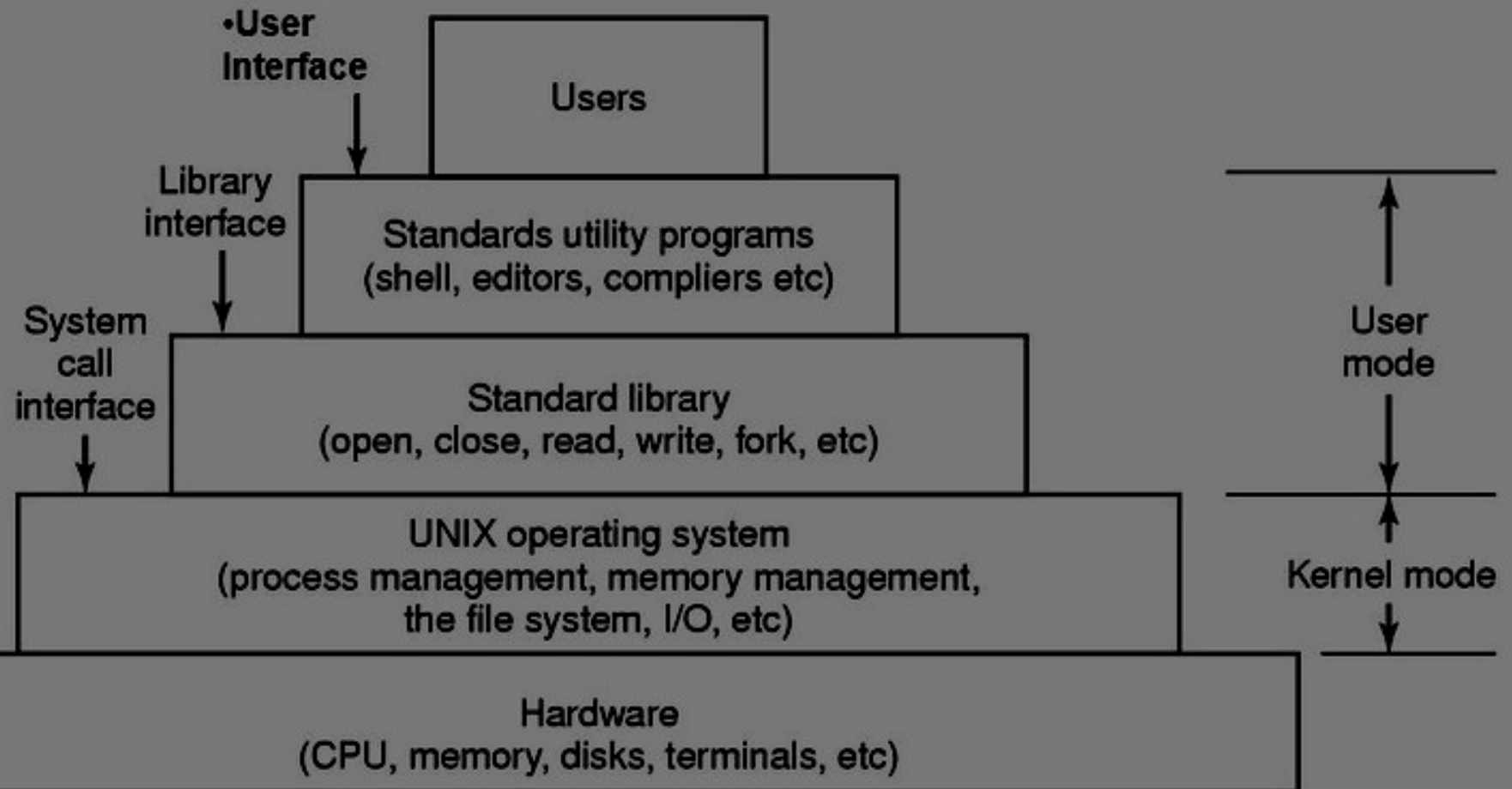
Protection

- **any mechanism for controlling access of processes or users to resources defined by the OS**

Security

- **defense of the system against internal and external attacks**
- **Insecure passwords, auditing software**
- **Unruly processes**

Linux Architecture



When You Press the Power Button?

| | |
|----------|--|
| BIOS | Basic Input/Output System executes MBR |
| MBR | Master Boot Record executes GRUB |
| GRUB | Grand Unified Bootloader executes Kernel |
| Kernel | Kernel executes <code>/sbin/init</code> |
| Init | Init executes runlevel programs |
| Runlevel | Runlevel programs are executed from <code>/etc/rc.d/rc*.d/</code> |