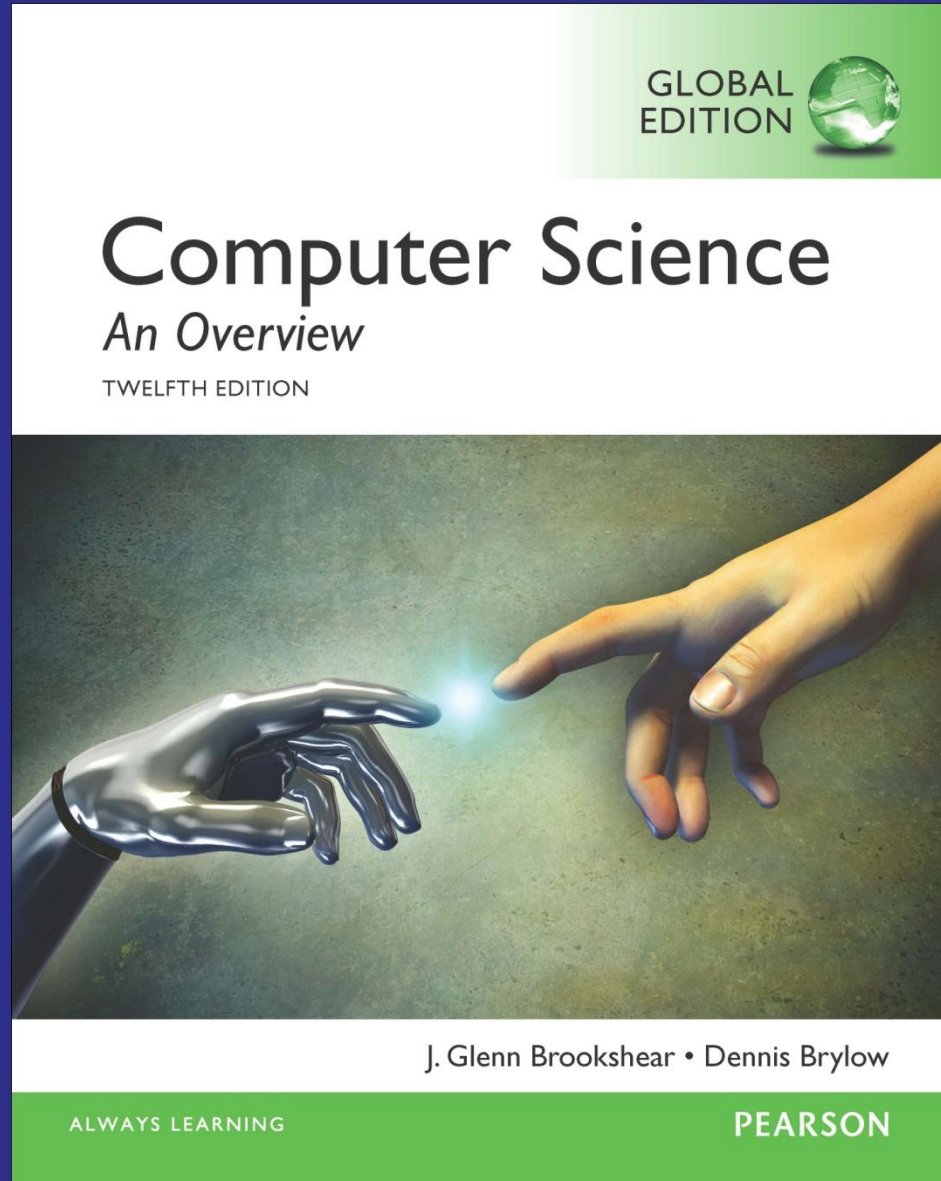


Chapter 8: Data Abstractions

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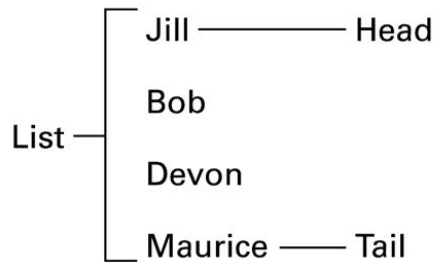
Chapter 8: Data Abstractions

- 8.1 Basic Data Structures
- 8.2 Related Concepts
- 8.3 Implementing Data Structures
- 8.4 A Short Case Study
- 8.5 Customized Data Types
- 8.6 Classes and Objects
- 8.7 Pointers in Machine Language

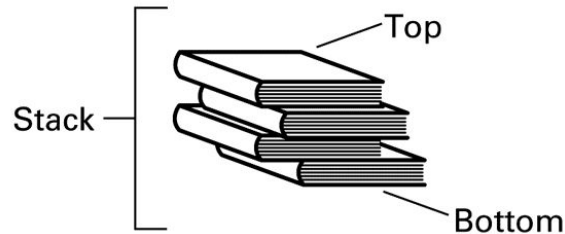
Basic Data Structures

- Arrays
- Aggregates
- List
 - Stack
 - Queue
- Tree

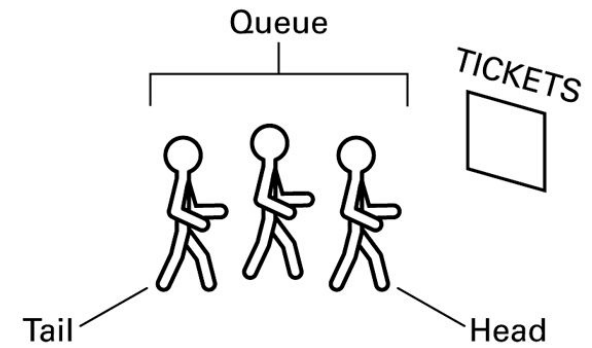
Figure 8.1 Lists, stacks, and queues



a. A list of names



b. A stack of books



c. A queue of people

Terminology for Arrays

- **Array:** A block of data whose entries are of same type
- A two **dimensional** array consists for rows and columns
- **Indices** are used to identify positions

Terminology for Aggregates

- **Aggregate:** A block of data items that might be of different type or sizes
- Each data item is called a **field**
- Fields are usually accessed by name

Terminology for Lists

- **List:** A collection of data whose entries are arranged sequentially
- **Head:** The beginning of the list
- **Tail:** The end of the list

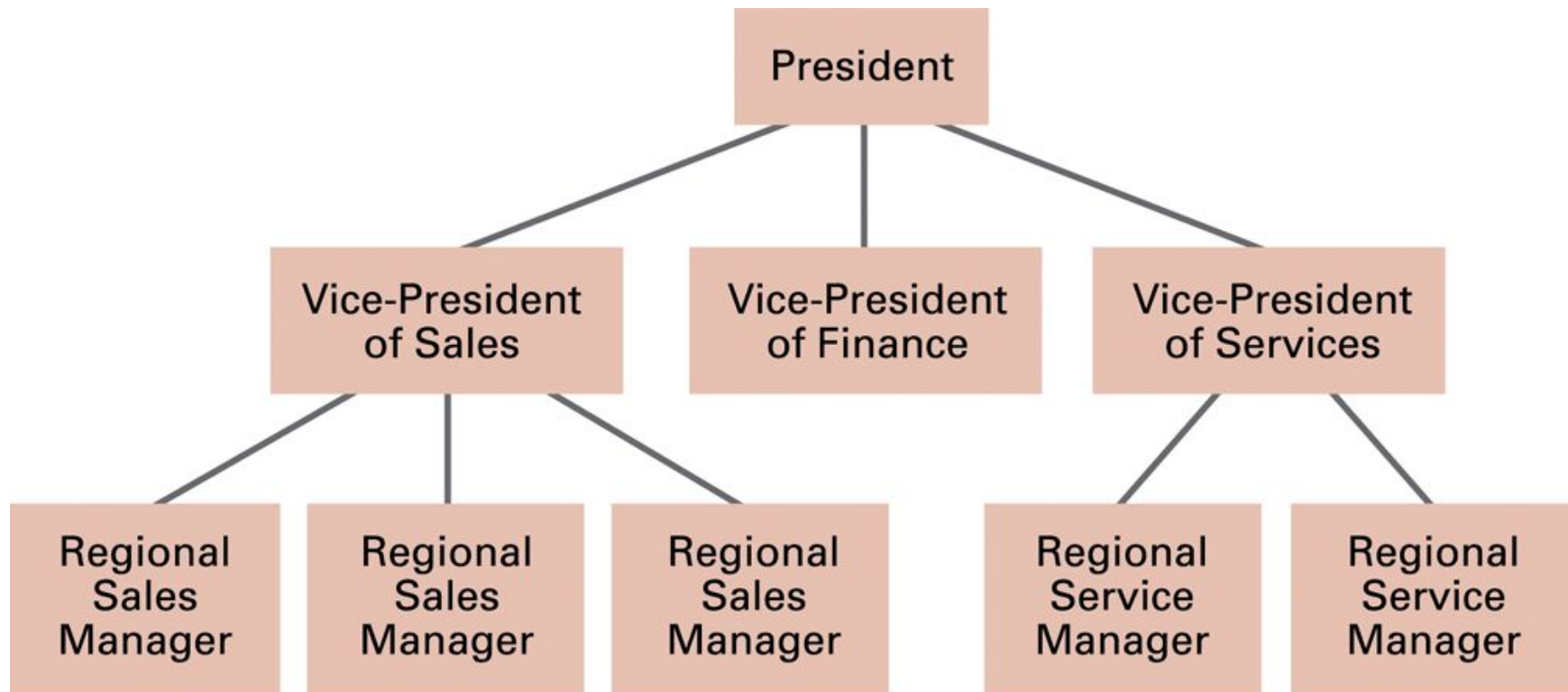
Terminology for Stacks

- **Stack:** A list in which entries are removed and inserted only at the head
- **LIFO:** Last-in-first-out
- **Top:** The head of list (stack)
- **Bottom** or **base:** The tail of list (stack)
- **Pop:** To remove the entry at the top
- **Push:** To insert an entry at the top

Terminology for Queues

- **Queue:** A list in which entries are removed at the head and are inserted at the tail
- **FIFO:** First-in-first-out

Figure 8.2 An example of an organization chart



Terminology for a Tree

- **Tree:** A collection of data whose entries have a hierarchical organization
- **Node:** An entry in a tree
- **Root node:** The node at the top
- **Terminal or leaf node:** A node at the bottom

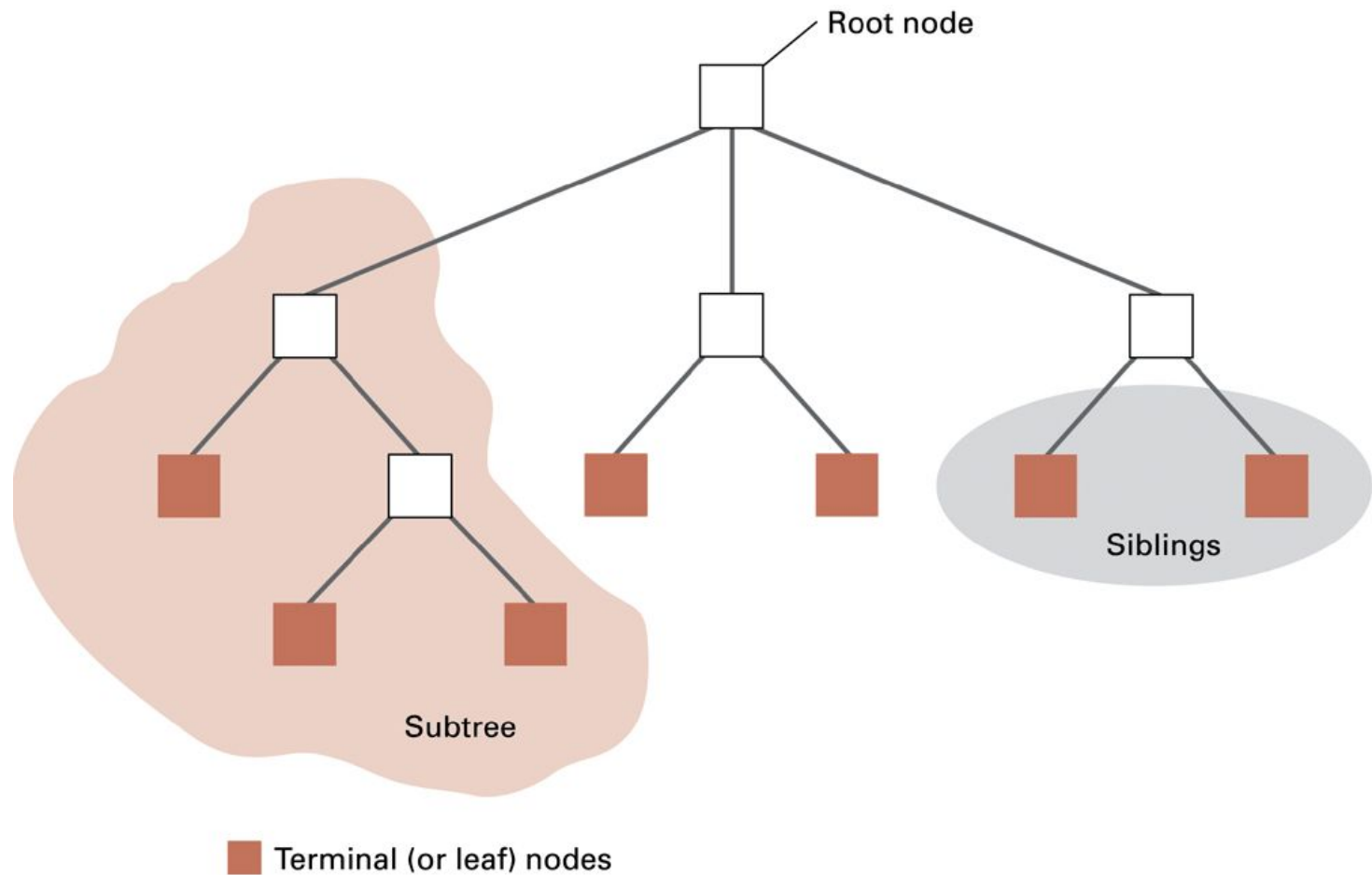
Terminology for a Tree (continued)

- **Parent:** The node immediately above a specified node
- **Child:** A node immediately below a specified node
- **Ancestor:** Parent, parent of parent, etc.
- **Descendent:** Child, child of child, etc.
- **Siblings:** Nodes sharing a common parent

Terminology for a Tree (continued)

- **Binary tree:** A tree in which every node has at most two children
- **Depth:** The number of nodes in longest path from root to leaf

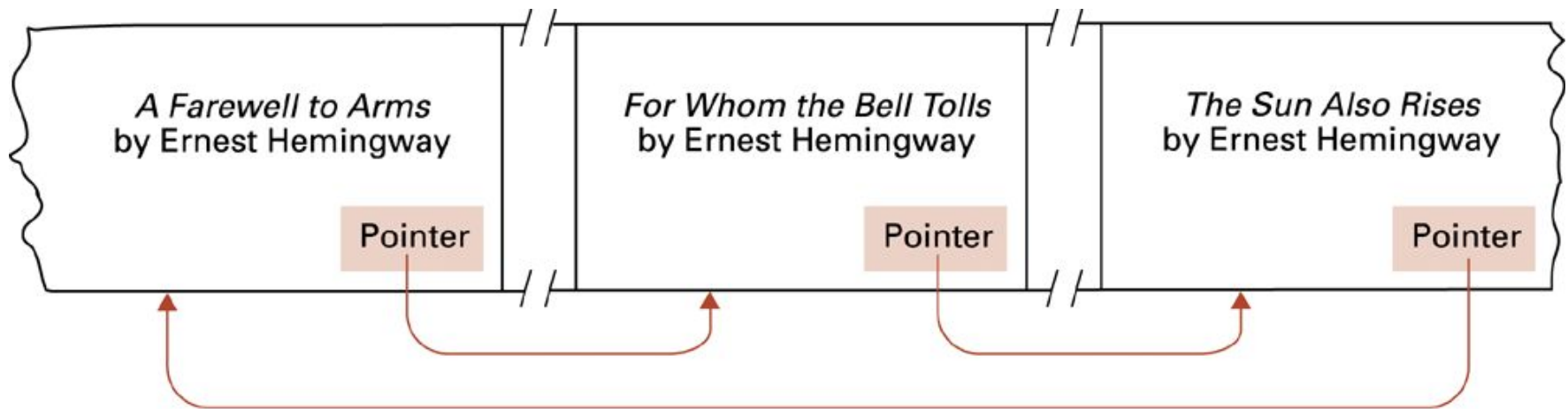
Figure 8.3 Tree terminology



Related Concepts

- Abstraction
 - Shield *users* (application software) from details of actual data storage
- Static vs. Dynamic Structure
 - Does the shape and size change over time?
- **Pointer**
 - A storage area that encodes an address where data is stored
 - Later used to access the data

Figure 8.4 Novels arranged by title but linked according to authorship



Storing Arrays

- Memory address of a particular cell can be computed
- **Row-major order** versus **column major order**
- Address polynomial

Figure 8.5 The array of temperature readings stored in memory starting at address x

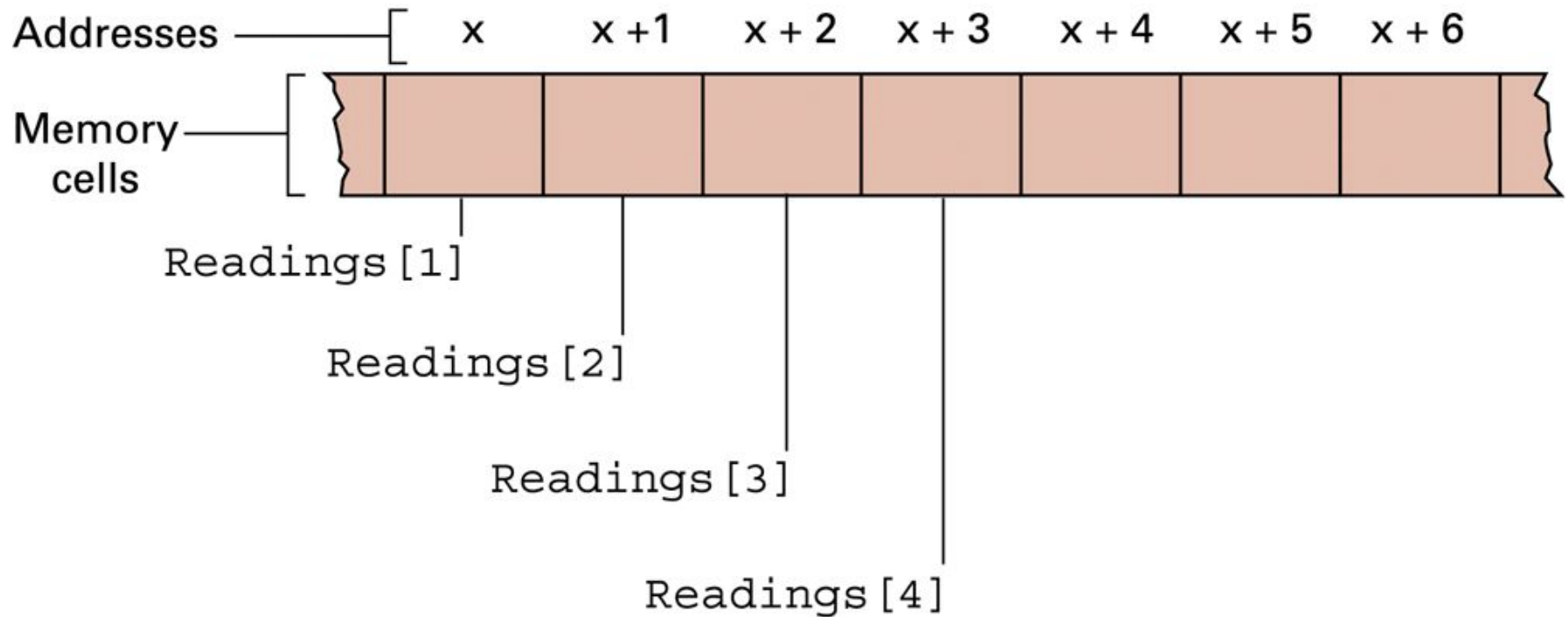
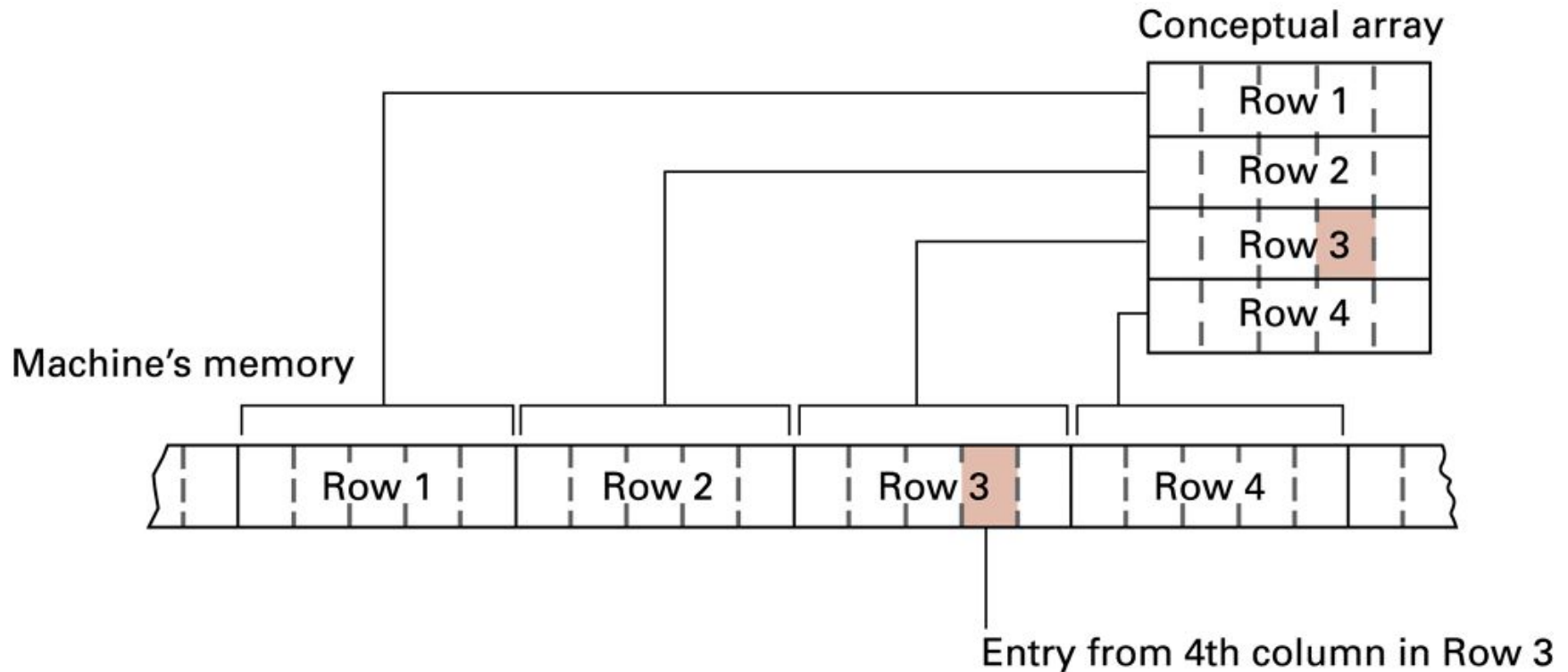


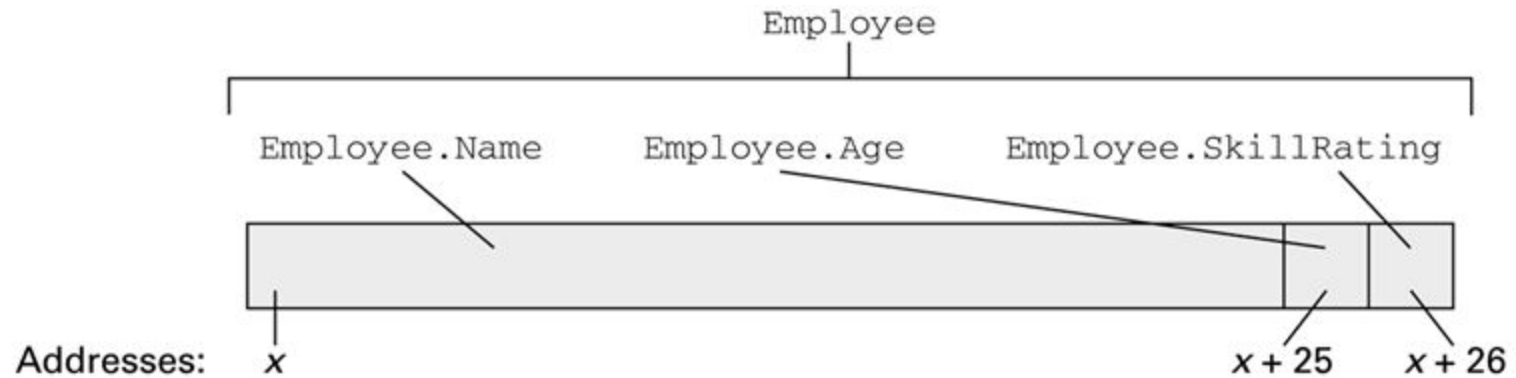
Figure 8.6 A two-dimensional array with four rows and five columns stored in row major order



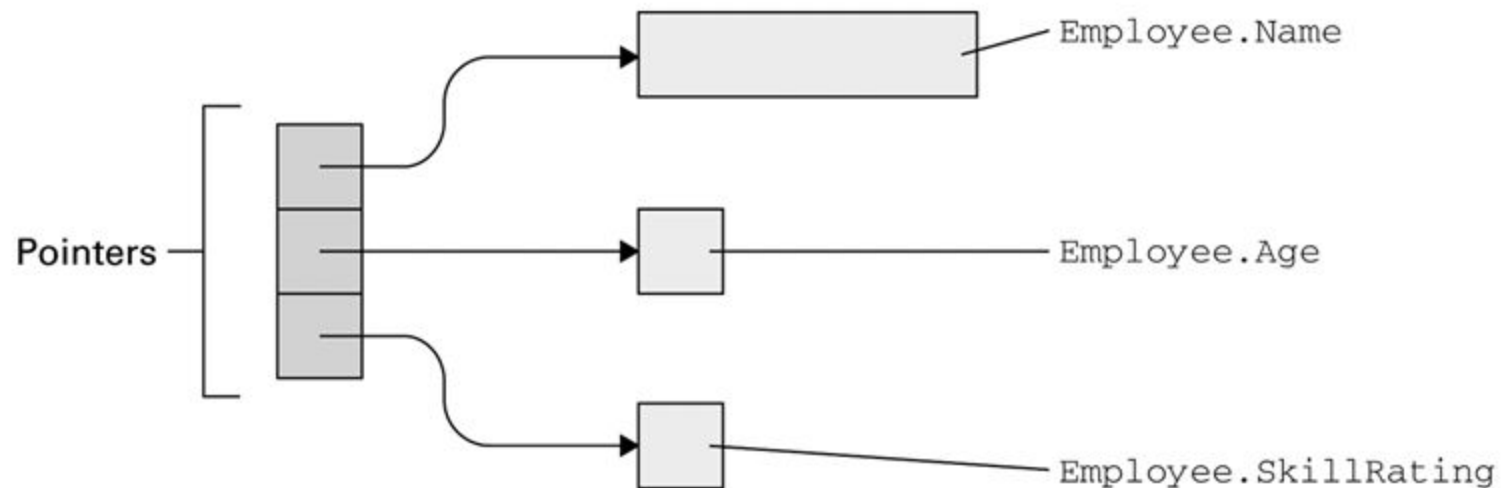
Storing Aggregates

- Fields can be stored one after the other in a contiguous block:
 - Memory cell address of each field can be computed
- Fields can be stored in separate locations identified by pointers

Figure 8.7 Storing the aggregate type Employee



a. Aggregate stored in a contiguous block



b. Aggregate fields stored in separate locations

Storing Lists

- **Contiguous list:** List in which entries are stored in an array
- **Linked list:** List in which entries are linked by pointers
 - **Head pointer:** Pointer to first entry in list
 - **null:** A “non-pointer” value used to indicate end of list

Figure 8.8 Names stored in memory as a contiguous list

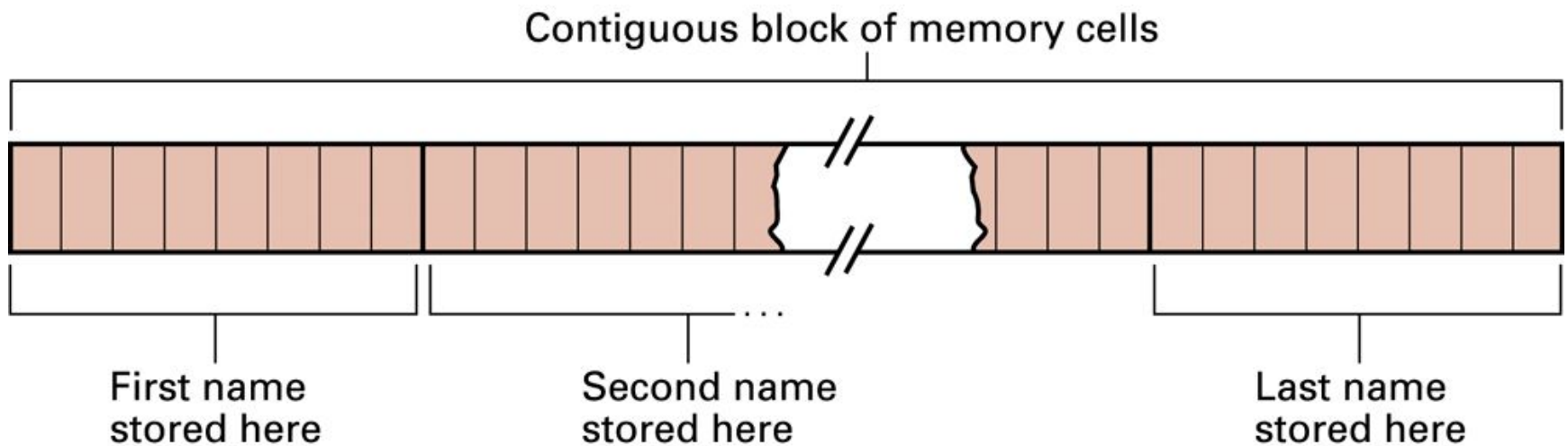


Figure 8.9 The structure of a linked list

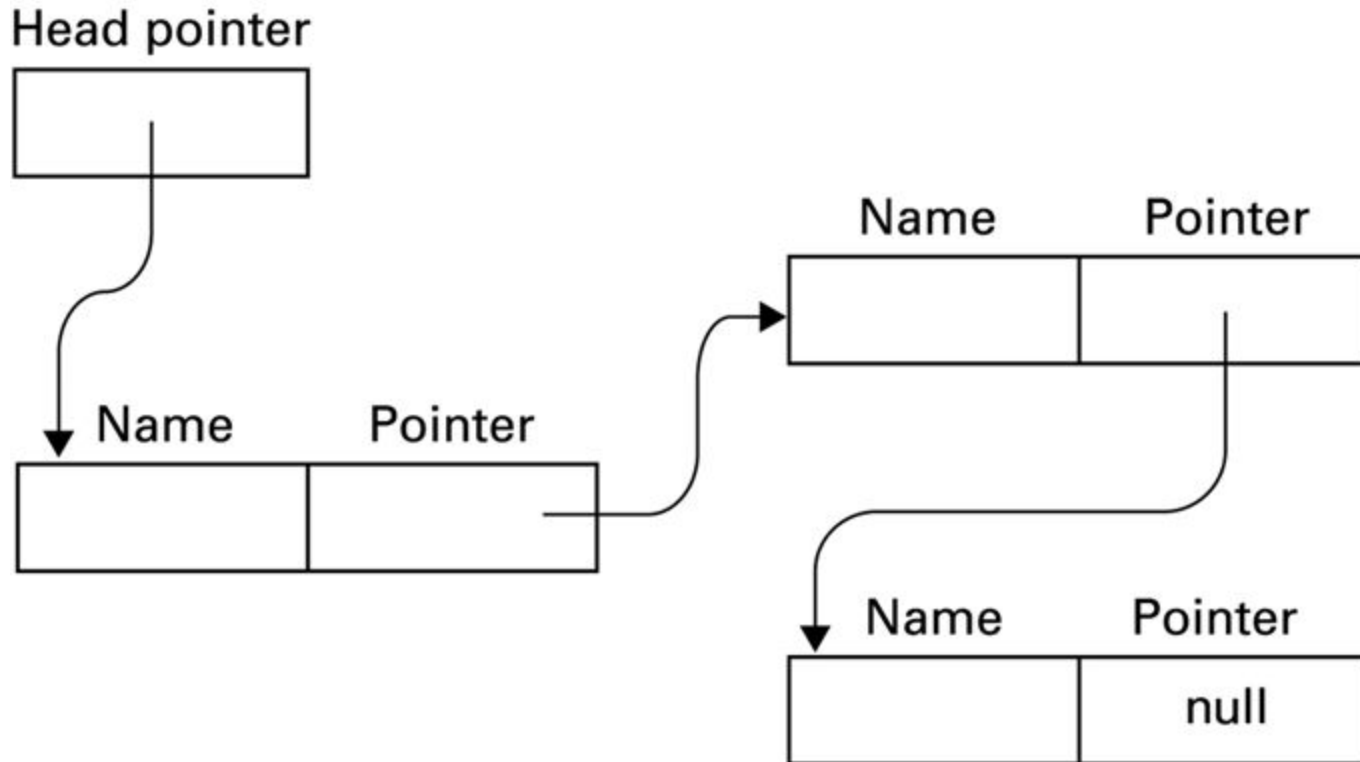


Figure 8.10 Deleting an entry from a linked list

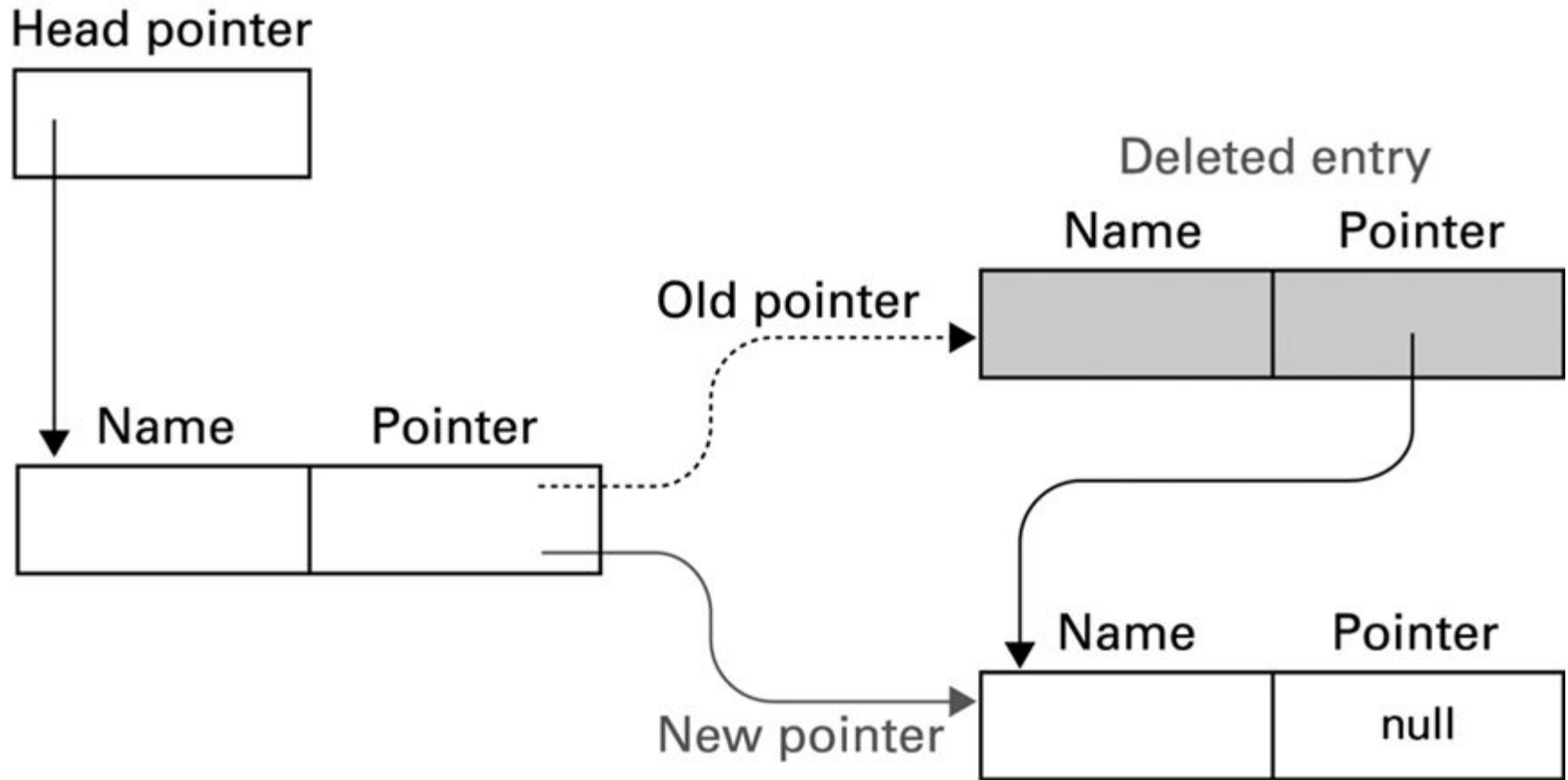
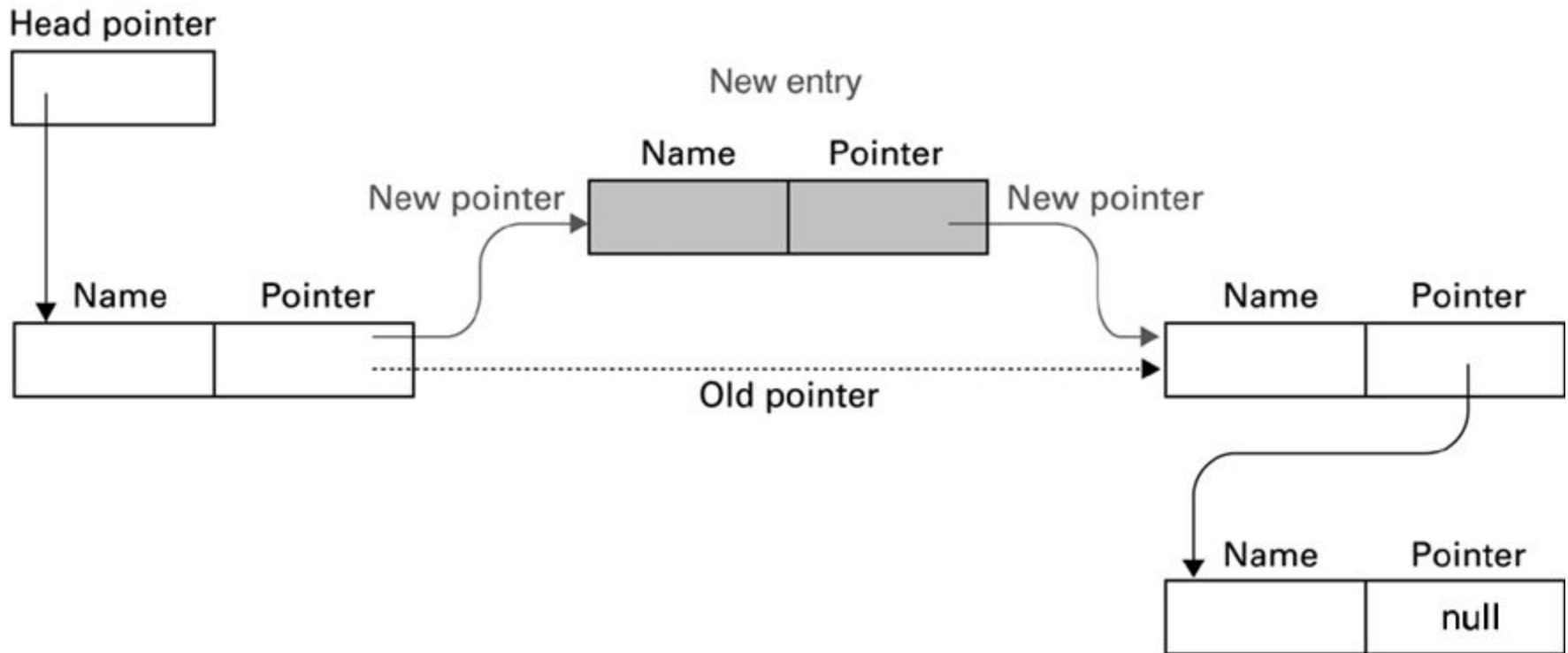


Figure 8.11 Inserting an entry into a linked list



Storing Stacks and Queues

- Stacks usually stored as contiguous lists
- Queues usually stored as **Circular Queues**
 - Stored in a contiguous block in which the first entry is considered to follow the last entry
 - Prevents a queue from crawling out of its allotted storage space

Figure 8.12 A stack in memory

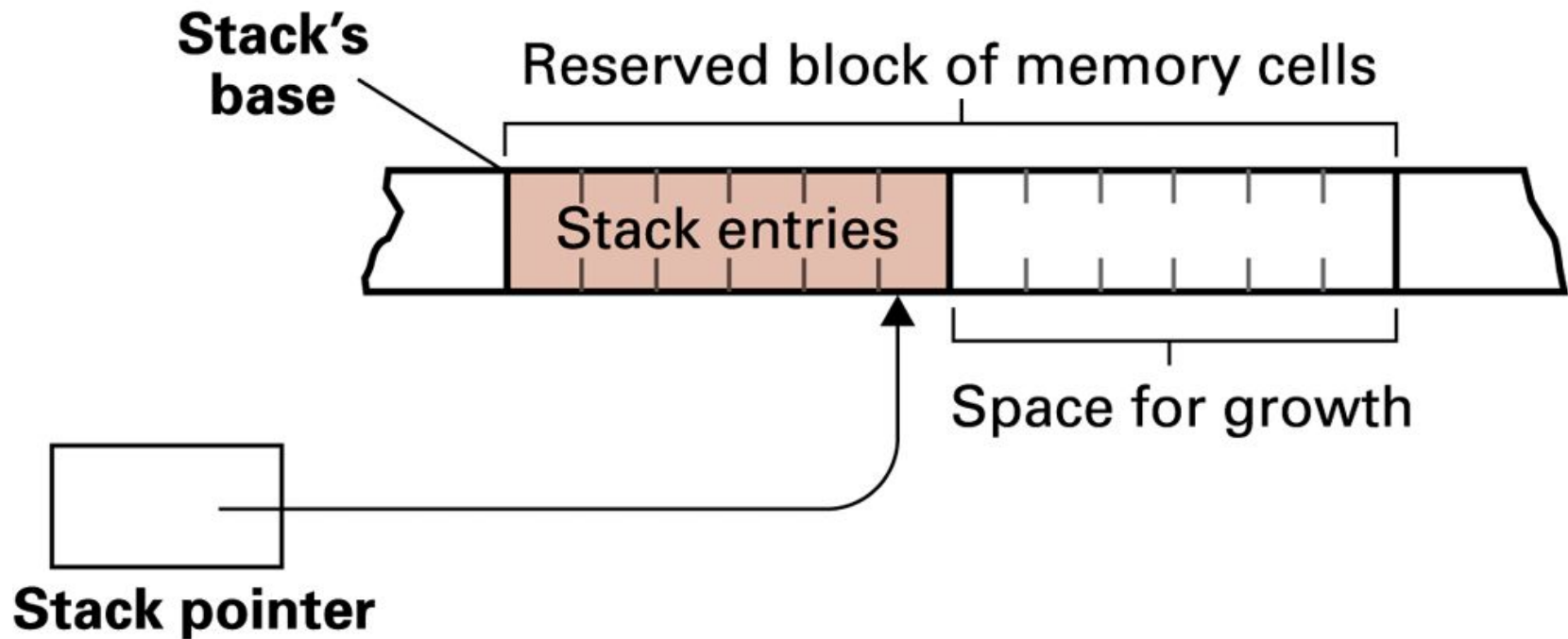
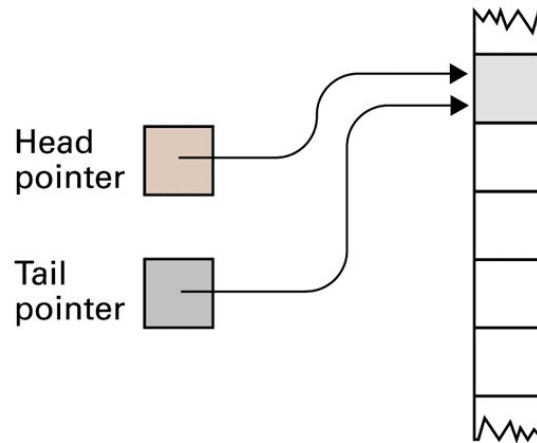
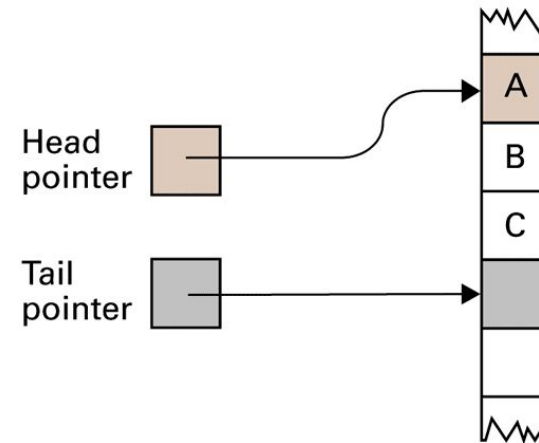


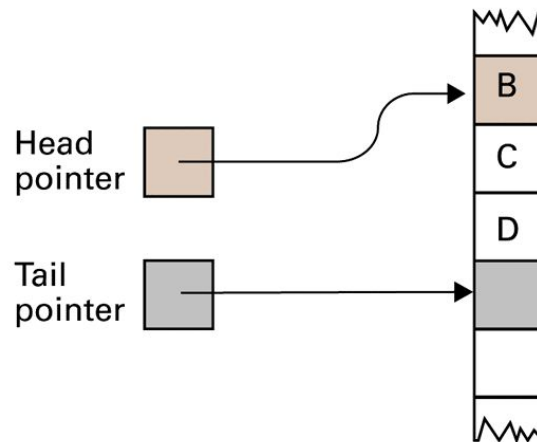
Figure 8.13 A queue implementation with head and tail pointers



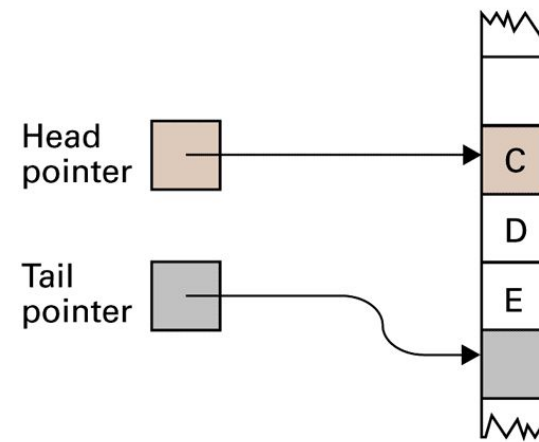
a. Empty queue



b. After inserting entries A, B, and C

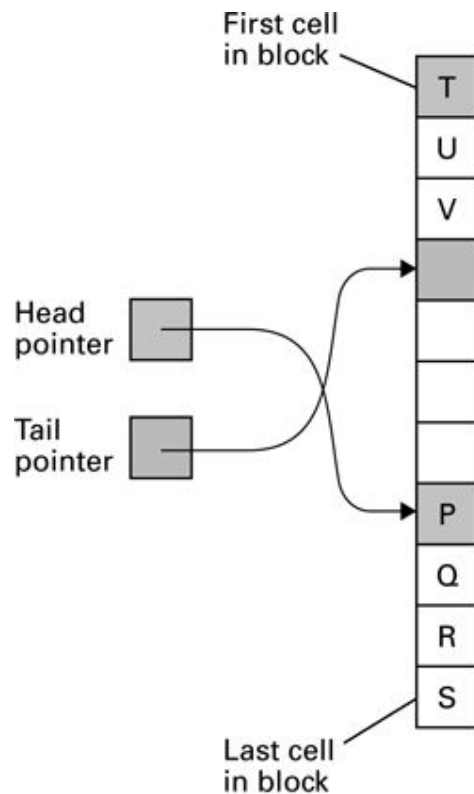


c. After removing A and inserting D

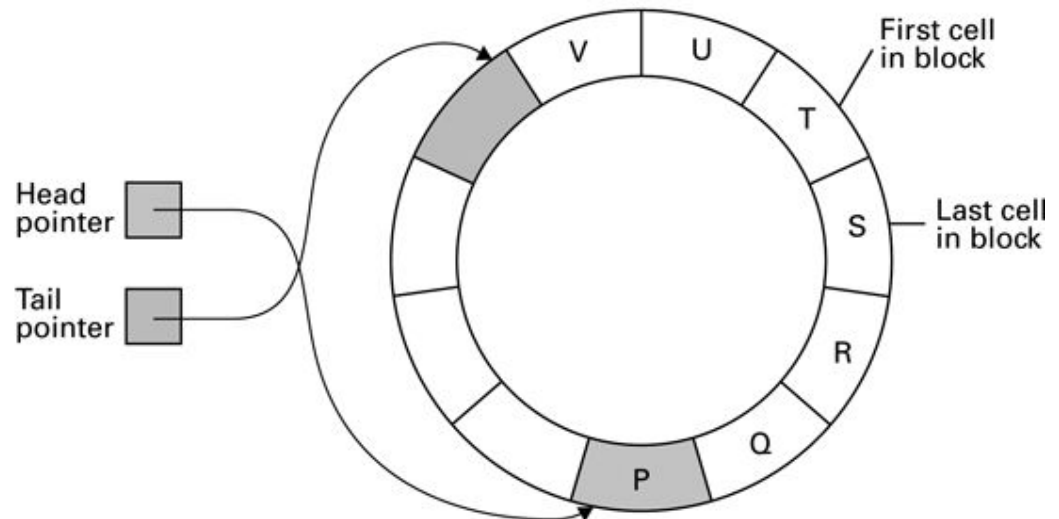


d. After removing B and inserting E

Figure 8.14 A circular queue containing the letters P through V



a. Queue as actually stored



b. Conceptual storage with last cell "adjacent" to first cell

Storing Binary Trees

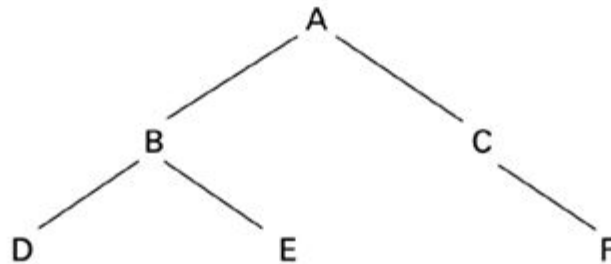
- Linked structure
 - Each node = data cells + two child pointers
 - Accessed via a pointer to root node
- Contiguous array structure
 - $A[1]$ = root node
 - $A[2], A[3]$ = children of $A[1]$
 - $A[4], A[5], A[6], A[7]$ = children of $A[2]$ and $A[3]$

Figure 8.15 The structure of a node in a binary tree



Figure 8.16 The conceptual and actual organization of a binary tree using a linked storage system

Conceptual tree



Actual storage organization

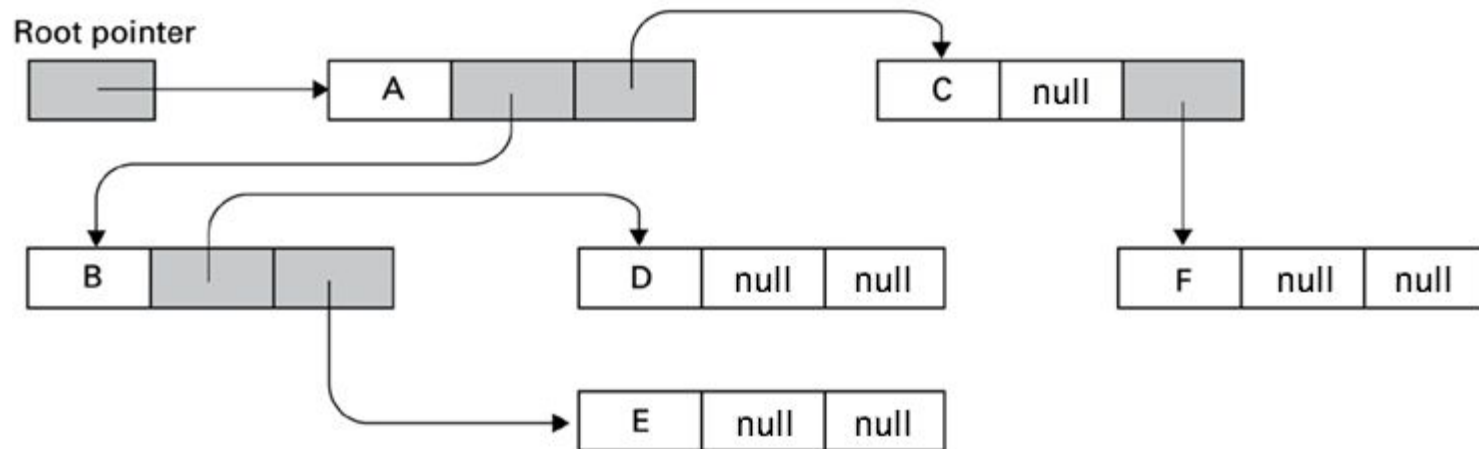
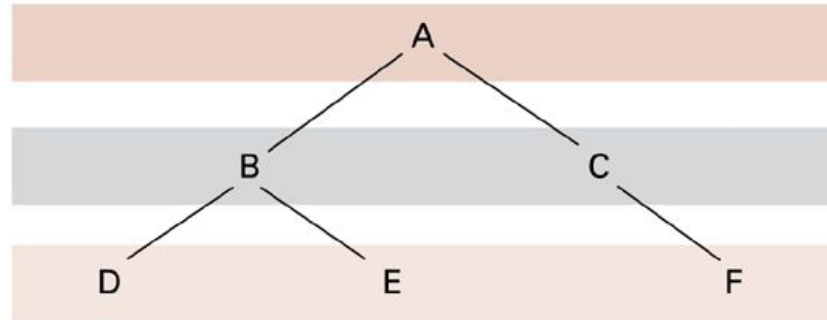


Figure 8.17 A tree stored without pointers

Conceptual tree



Actual storage organization

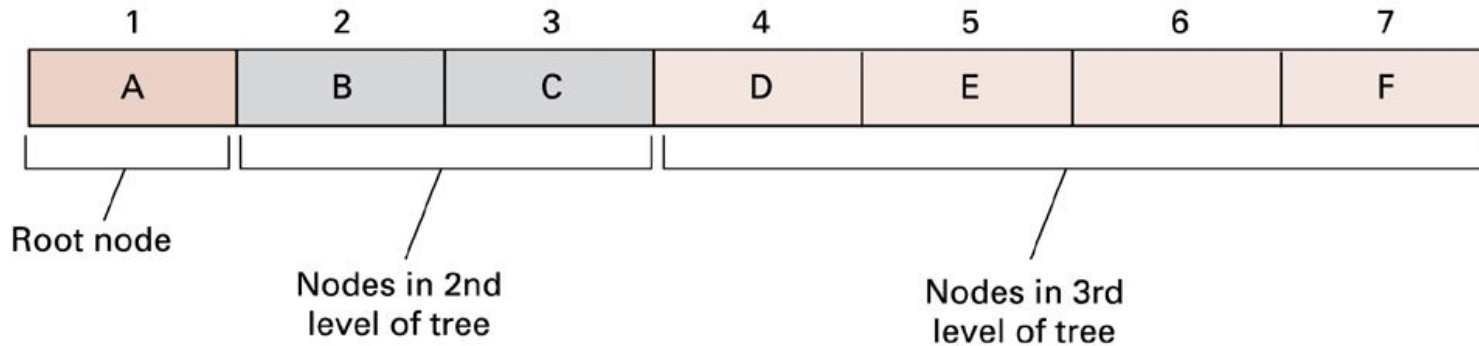
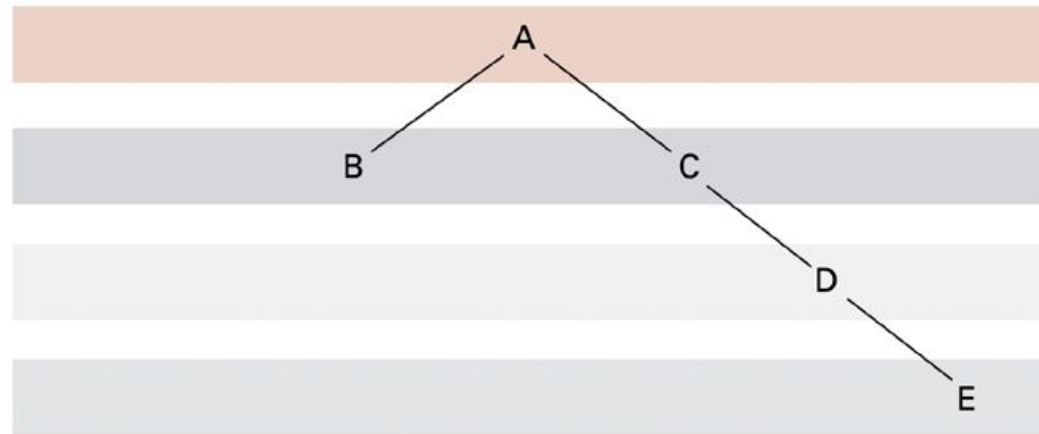
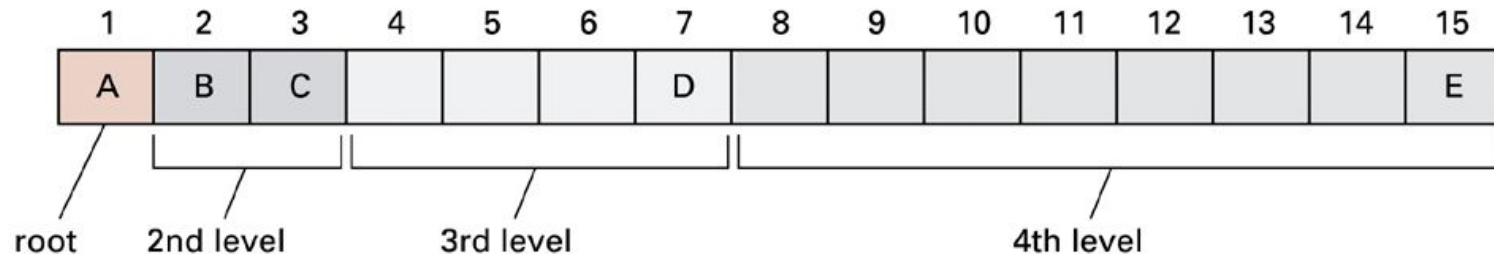


Figure 8.18 A sparse, unbalanced tree shown in its conceptual form and as it would be stored without pointers

Conceptual tree



Actual storage organization



Manipulating Data Structures

- Ideally, a data structure should be manipulated solely by pre-defined functions
 - Example: A list typically has a function `insert` for inserting new entries
 - The data structure along with these functions constitutes a complete abstract tool

Figure 8.19 A function for printing a linked list

```
def PrintList (List):  
    CurrentPointer = List.Head  
    while (CurrentPointer is not None):  
        print(CurrentPointer.Value)  
        CurrentPointer = CurrentPointer.Next
```

Case Study

Problem:

Construct an abstract tool consisting of a list of names in alphabetical order along with the operations: search, print, and insert.

Figure 8.20 The letters A through M arranged in an ordered tree

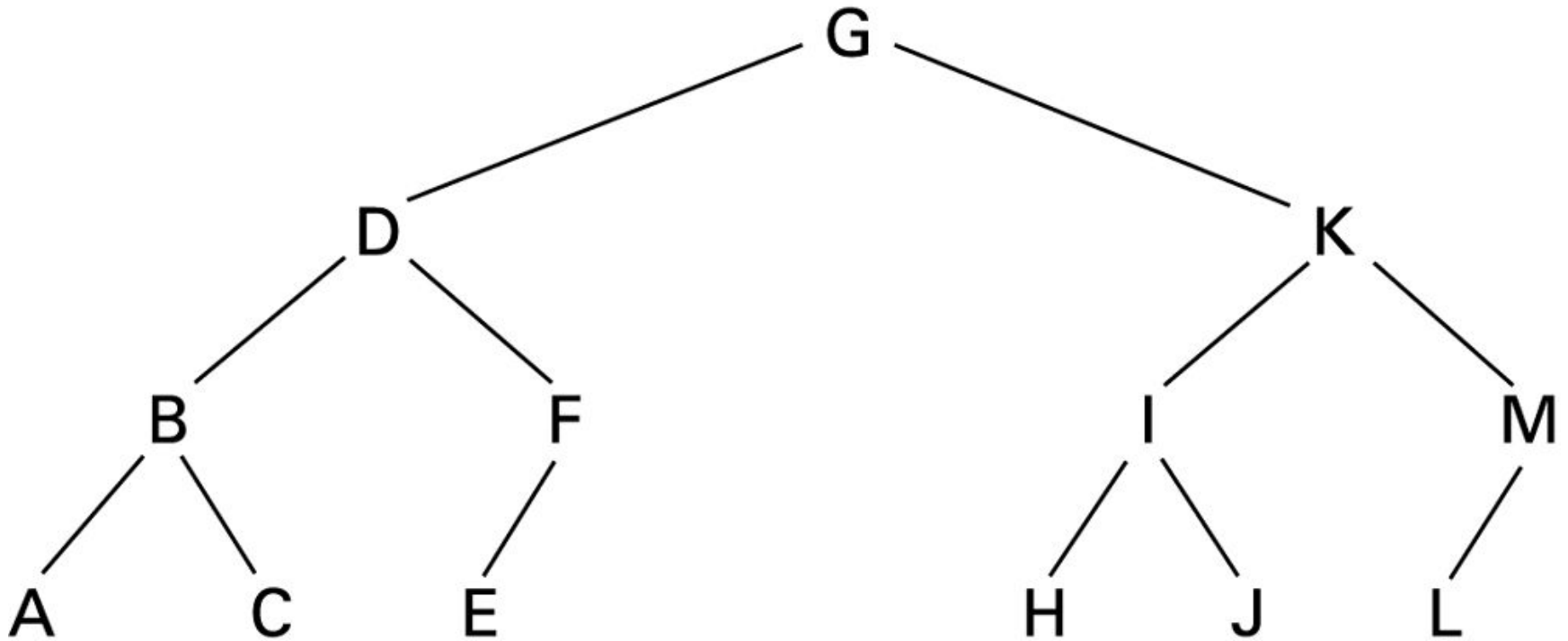


Figure 8.21 The binary search as it would appear if the list were implemented as a linked binary tree

```
def Search (Tree, TargetValue):  
    if (Tree is None):  
        return None          # Search failed  
    elif (TargetValue == Tree.Value):  
        return Tree          # Search succeeded  
    elif (TargetValue < Tree.Value):  
        # Continue search in left subtree  
        return Search(Tree.Left, TargetValue)  
    elif (TargetValue > Tree.Value):  
        # Continue search in right subtree  
        return Search(Tree.Right, TargetValue)
```


Figure 8.22 The successively smaller trees considered by the function in Figure 8.21 when searching for the letter J

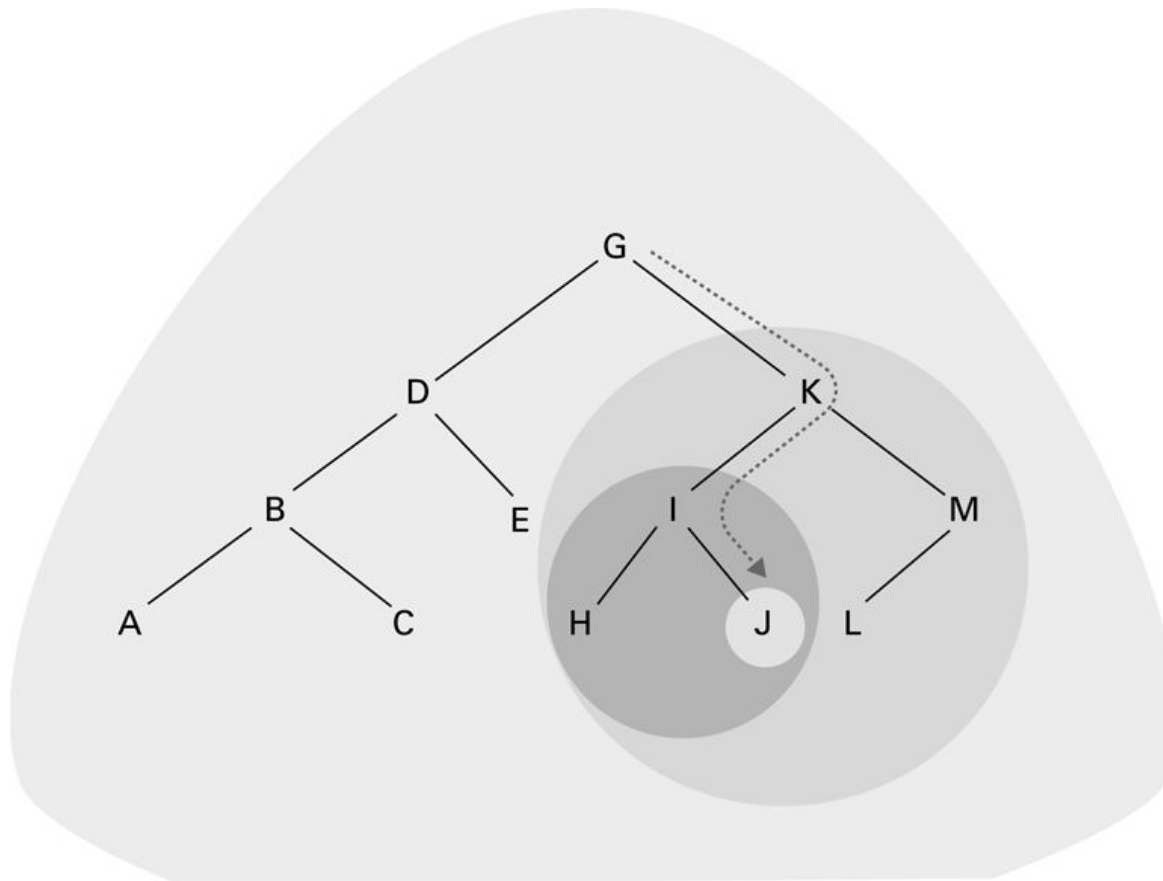


Figure 8.23 Printing a search tree in alphabetical order

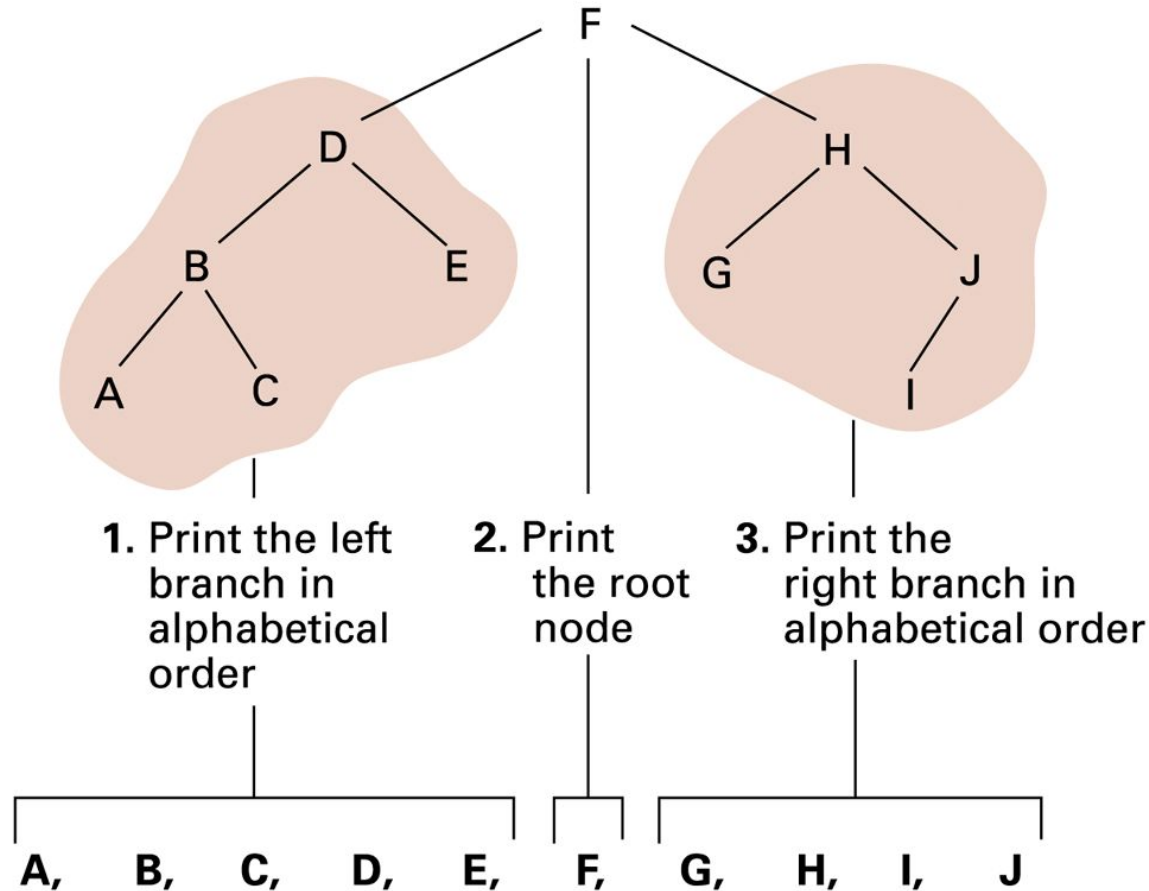
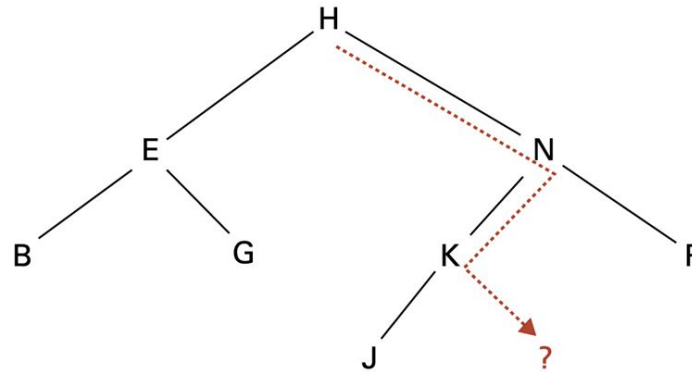


Figure 8.24 A function for printing the data in a binary tree

```
def PrintTree (Tree):  
    if (Tree is not None):  
        PrintTree(Tree.Left)  
        print(Tree.Value)  
        PrintTree(Tree.Right)
```

Figure 8.25 Inserting the entry M into the list B, E, G, H, J, K, N, P stored as a tree

a. Search for the new entry until its absence is detected



b. This is the position in which the new entry should be attached

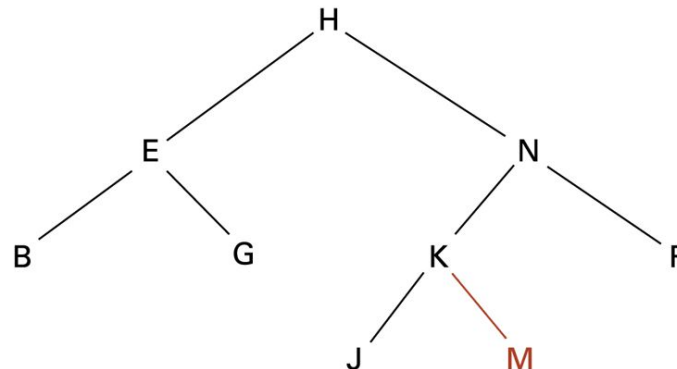


Figure 8.26 A function for inserting a new entry in a list stored as a binary tree

```
def Insert(Tree, NewValue):  
    if (Tree is None):  
        # Create a new leaf with NewValue  
        Tree = TreeNode()  
        Tree.Value = NewValue  
    elif (NewValue < Tree.Value):  
        # Insert NewValue into the left subtree  
        Tree.Left = Insert(Tree.Left, NewValue)  
    elif (NewValue > Tree.Value):  
        # Insert NewValue into the right subtree  
        Tree.Right = Insert(Tree.Right, NewValue)  
    else:  
        # Make no change  
    return Tree
```

User-defined Data Type

- Use an aggregate structure to define new type, in C:

```
struct EmployeeType
{
    char    Name[25];
    int     Age;
    real    SkillRating;
}
```

- Use the new type to define variables:

```
struct EmployeeType DistManager, SalesRep1;
```

Abstract Data Type

- A user-defined data type that can include both data (representation) and functions (behavior)
- Example:

```
interface StackType
{
    public int pop();
    public int push(int item);
    public boolean isEmpty();
    public boolean isFull();
}
```

Class

- An abstract data type with extra features
 - Properties can be inherited
 - Constructor methods to initialize new objects
 - Contents can be encapsulated

Figure 8.27 A stack of integers implemented in Java and C#

```
class StackOfIntegers implements StackType
{
    private int[] StackEntries = new int[20];
    private int StackPointer = 0;
    public void push(int NewEntry)
    {
        if (StackPointer < 20)
            StackEntries[StackPointer++] = NewEntry;
    }
    public int pop()
    {
        if (StackPointer > 0) return StackEntries[--StackPointer];
        else return 0;
    }
    public boolean isEmpty()
    {
        return (StackPointer == 0);
    }
    public boolean isFull()
    {
        return (StackPointer >= MAX);
    }
}
```

Pointers in Machine Language

- **Immediate addressing:** Instruction contains the data to be accessed
- **Direct addressing:** Instruction contains the address of the data to be accessed
- **Indirect addressing:** Instruction contains the location of the address of the data to be accessed

Figure 8.28 Our first attempt at expanding the machine language in Appendix C to take advantage of pointers

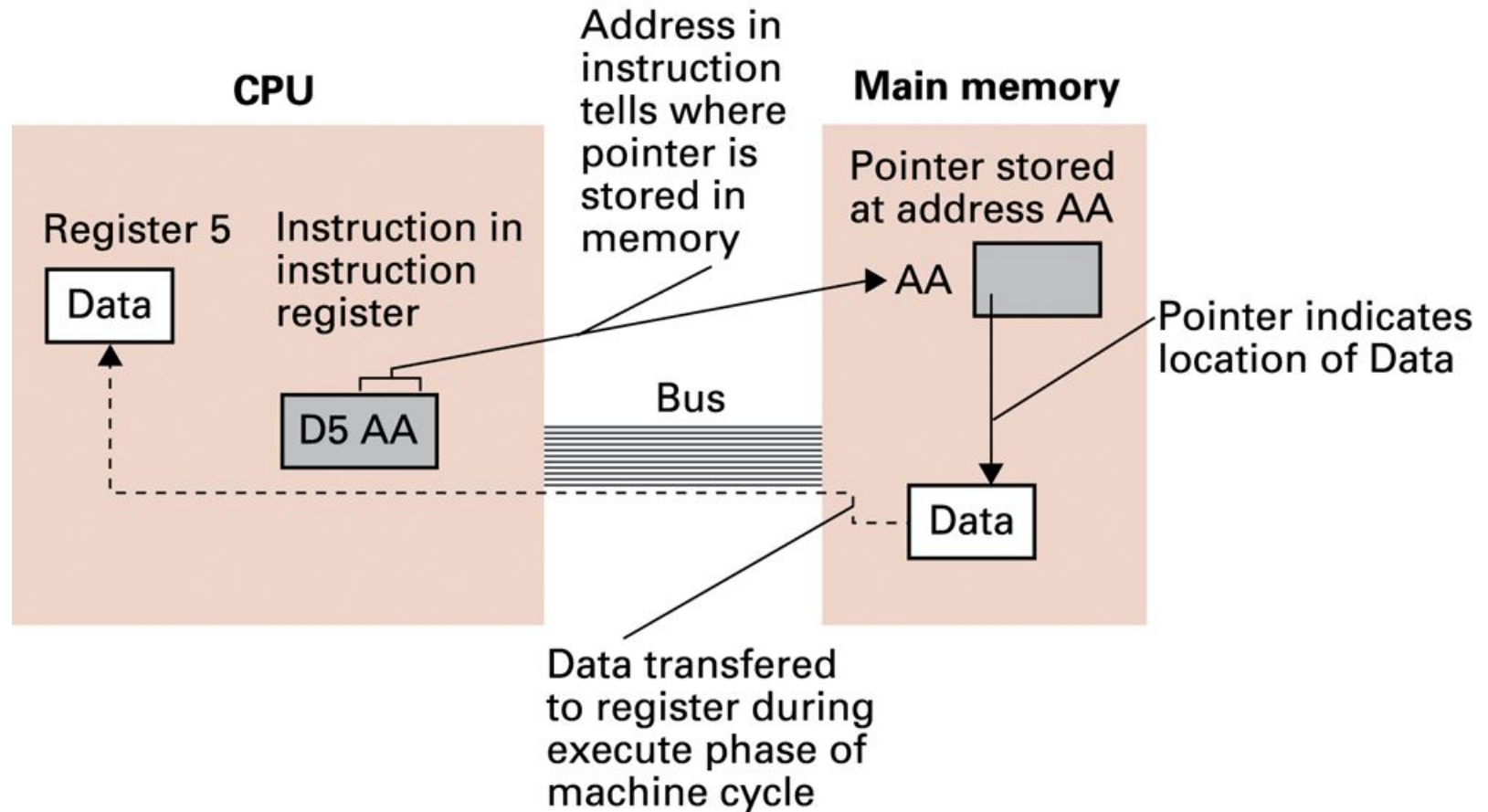
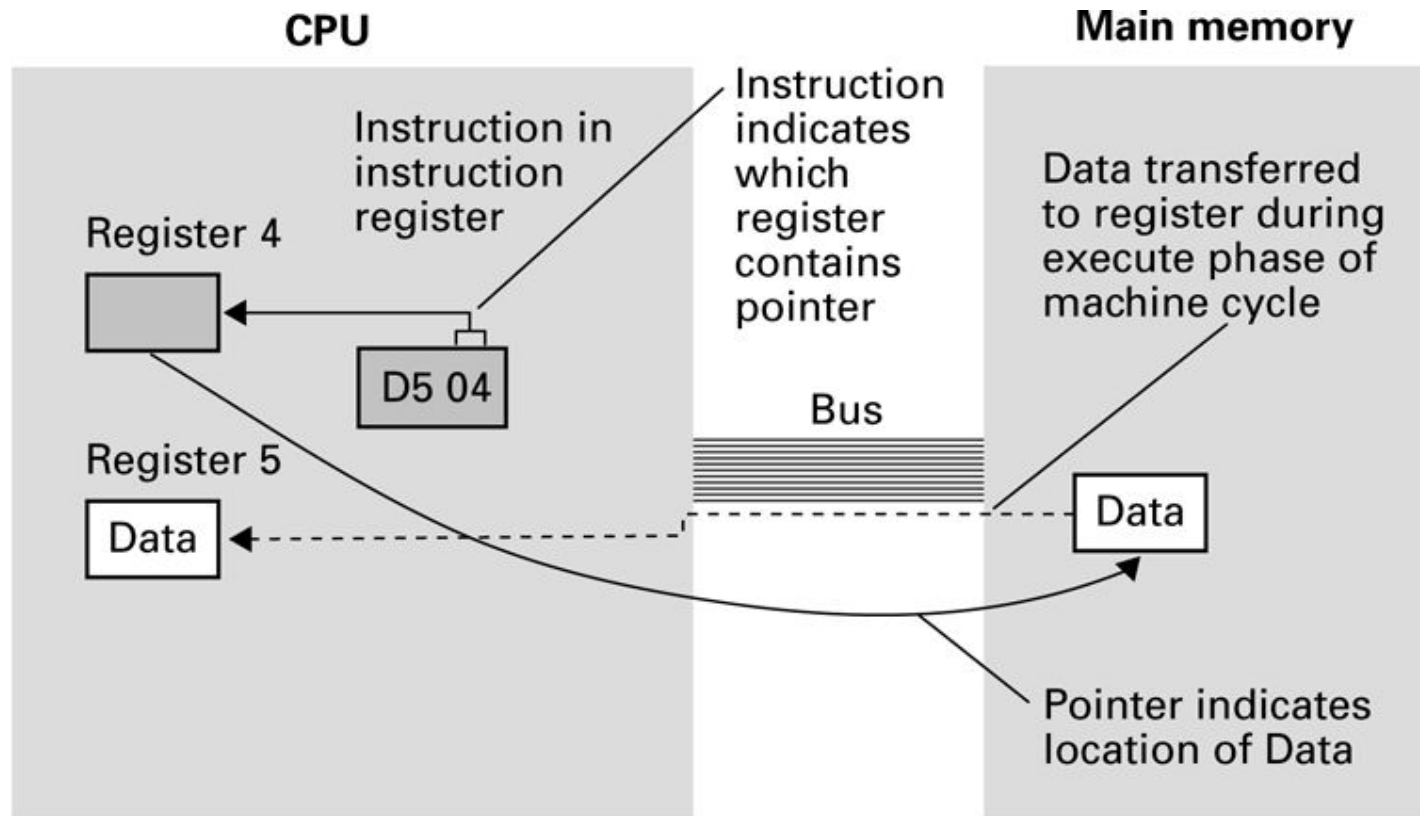
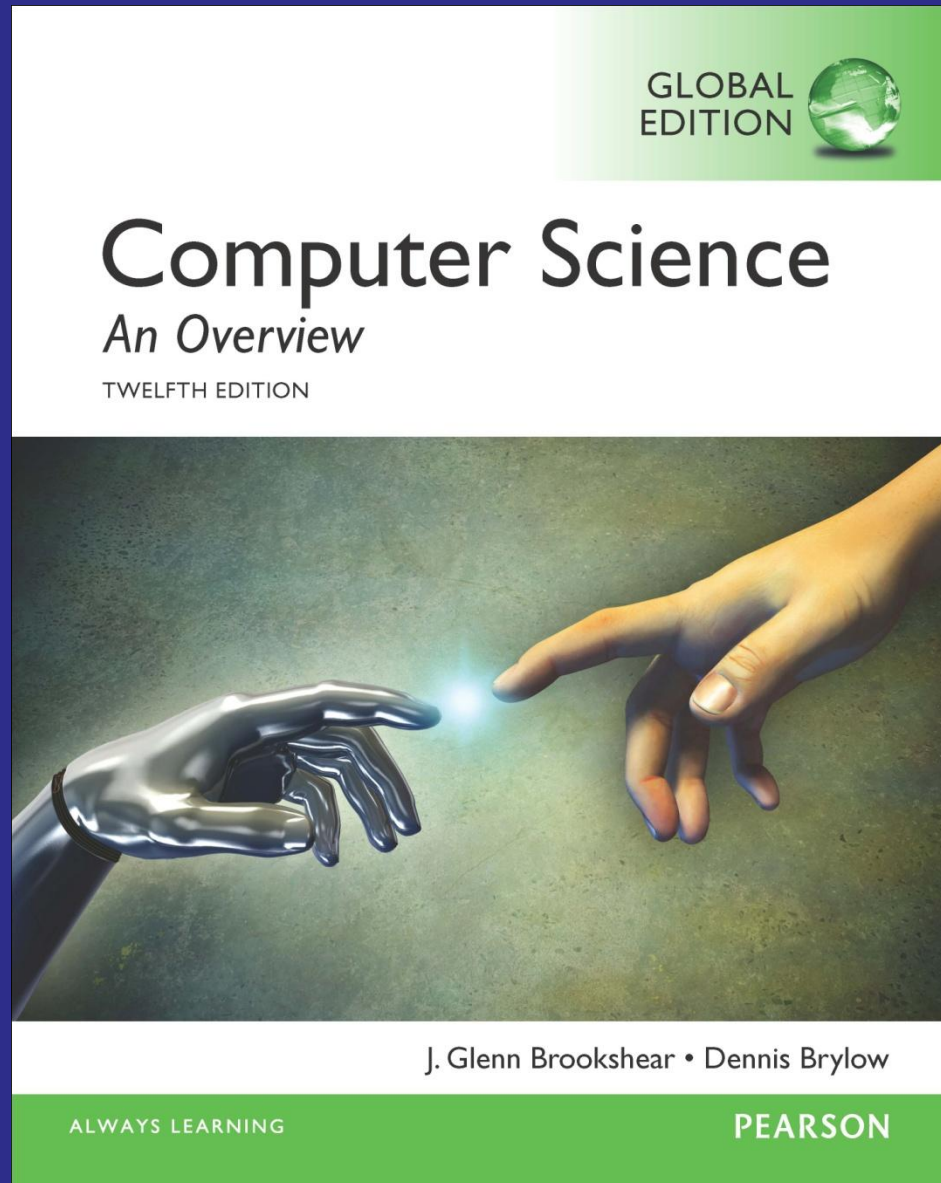


Figure 8.29 Loading a register from a memory cell that is located by means of a pointer stored in a register



End of Chapter

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