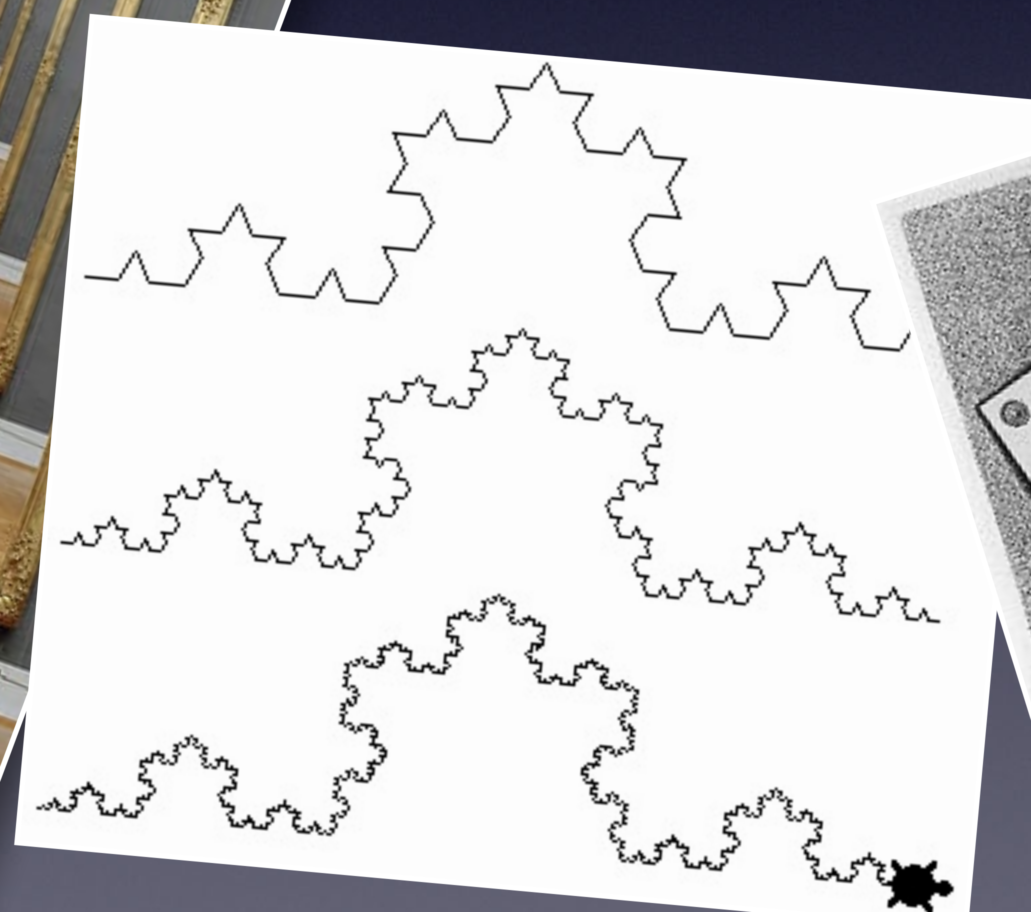
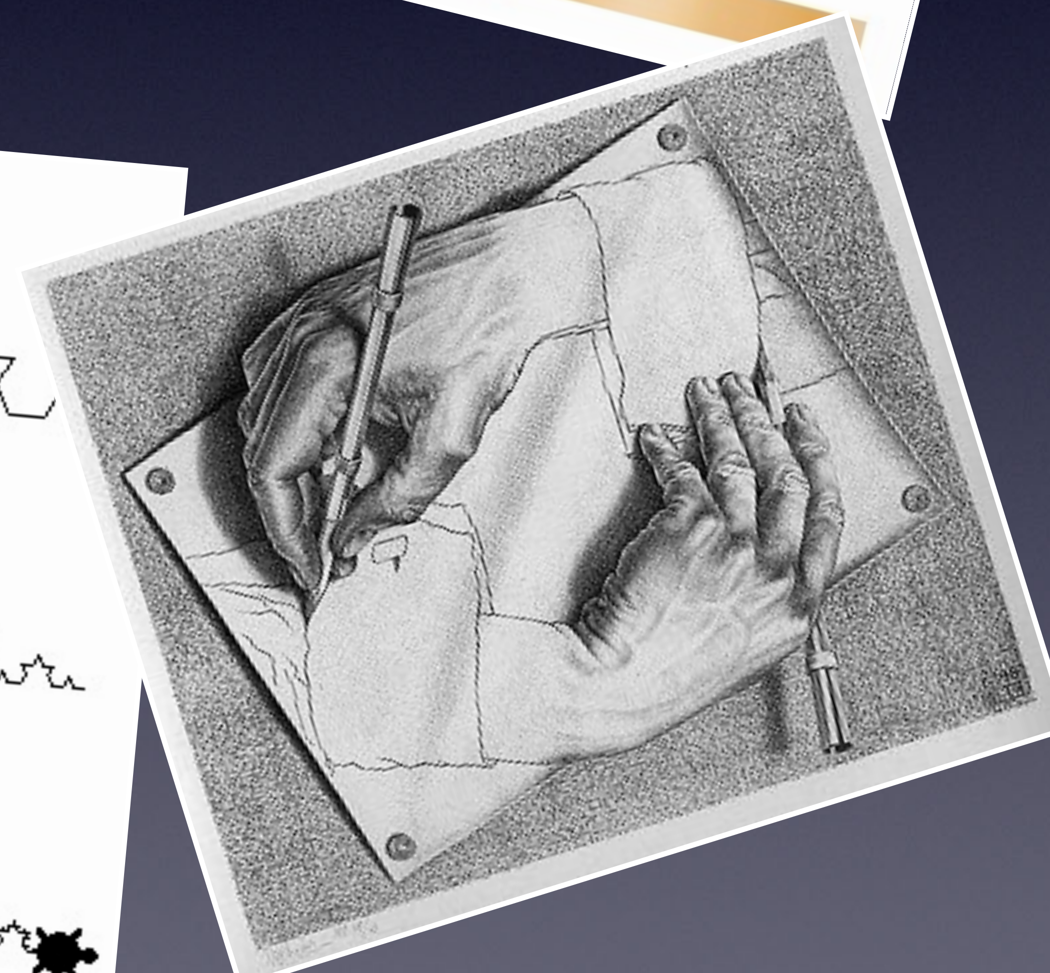
Recursion





code.google.com/p/rzzzwilson/source/browse/talks/recursion_1

Recursion

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> To understand recursion,
> one must first understand recursion.

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> You write a recursive solution as if
> you already have a solution.

It's <u>programming by faith</u>!

Recursion

A solution is recursive if it:

- solves a problem with smaller recursive cases
- has one or more non-recursive base cases

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An example: the Factorial function:

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base
case

Recursion

The "Tower of Hanoi" puzzle was invented by Édouard Lucas in 1883:

Recursion

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- Base case (N = 1)

  We know how to move one disk.

- Recursive case

  We can move the top N-1 disks to the other position, then the single largest disk to the target position, then move the N-1 disks to the target.

Recursion

Let's write some python:

```
def hanoi(n, src, dst, tmp):
    if n == 1:
        print('move %s to %s' % (src, dst))
    else:
        hanoi(n-1, src, tmp, dst)
        hanoi(1, src, dst, tmp)
        hanoi(n-1, tmp, dst, src)
```

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        hanoi(n-1, tmp, dst, src)
```

This solution is easy to split into base and recursive cases. It fails if N is less than zero.

Recursion

A more robust and more compact solution:

```
def hanoi(n, src, dst, tmp):
    if n > 0:
        hanoi(n-1, src, tmp, dst)
        print('move %s to %s' % (src, dst))
        hanoi(n-1, tmp, dst, src)
```

More efficient than the first solution?

Recursion

```python
import sys

def hanoi(n, src, dst, tmp):
    if n > 0:
        hanoi(n-1, src, tmp, dst)
        print('move %s to %s' % (src, dst))
        hanoi(n-1, tmp, dst, src)

number = int(sys.argv[1])
hanoi(number, 'A', 'B', 'C')
```

```
% python ./hanoi.py 3
move A to B
move A to C
move B to C
move A to B
move C to A
move C to B
move A to B
```

Recursion

Why are recursive algorithms useful?

Recursion

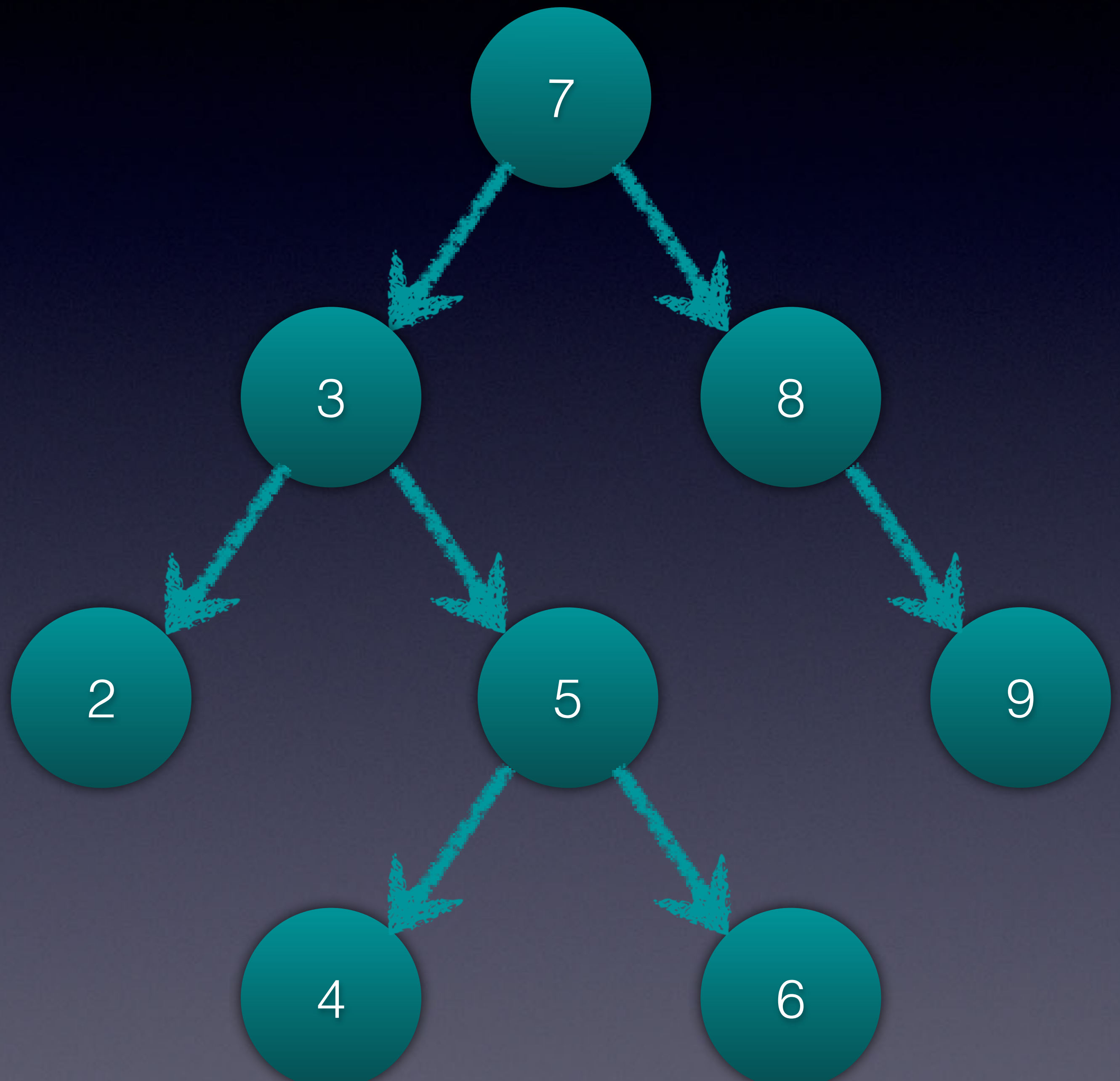
Why are recursive algorithms useful?

Some data structures are recursively defined, so recursive algorithms are a natural way to process them.

For example, binary trees.

Recursion

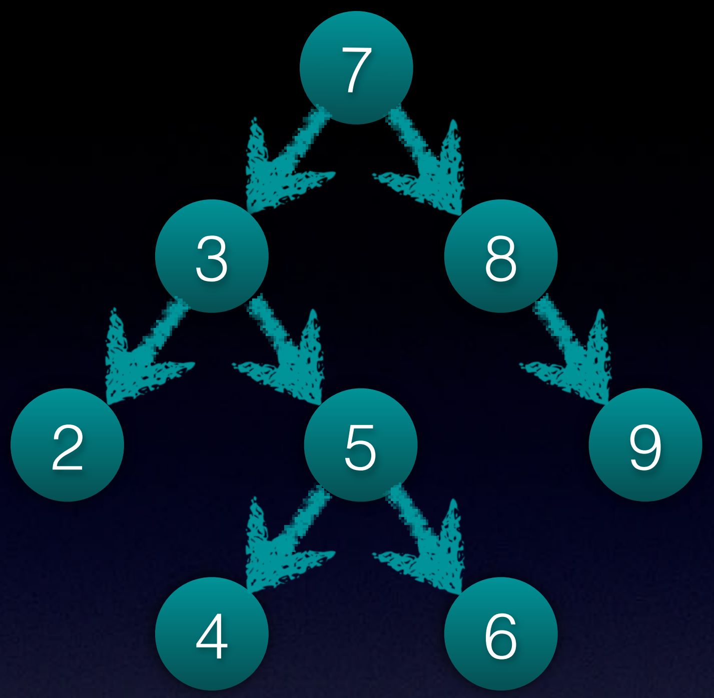
A binary tree looks like this:



A binary tree consists of a *node* which has a *value* and *left* and *right* pointers, which may be *None* or refer to another binary tree.

Recursion

Suppose we want to print out the values in the tree nodes.



One method is, for each node:

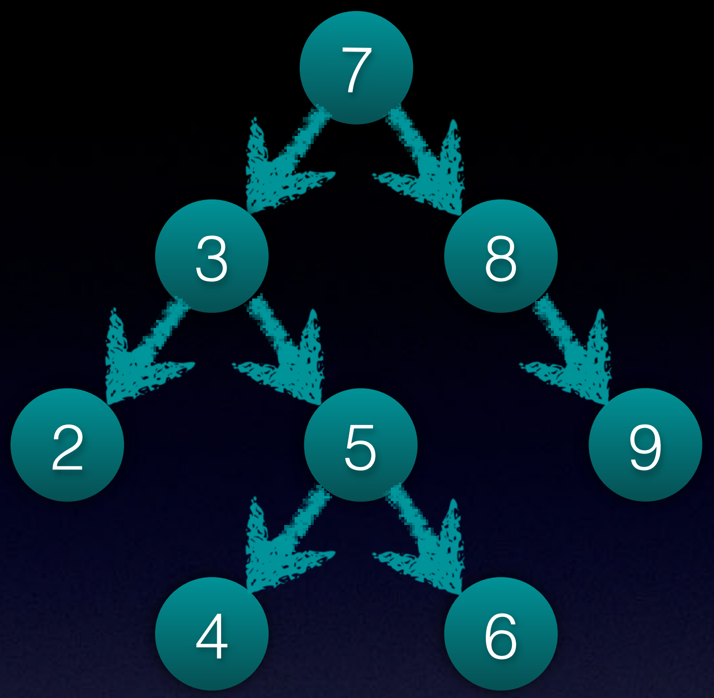
- print the left sub-tree
- print the node value
- print the right sub-tree

This is called an *in-order* walk of the tree. We print the node value in between the left and right sub-tree values.

Recursion

Let's define our tree:

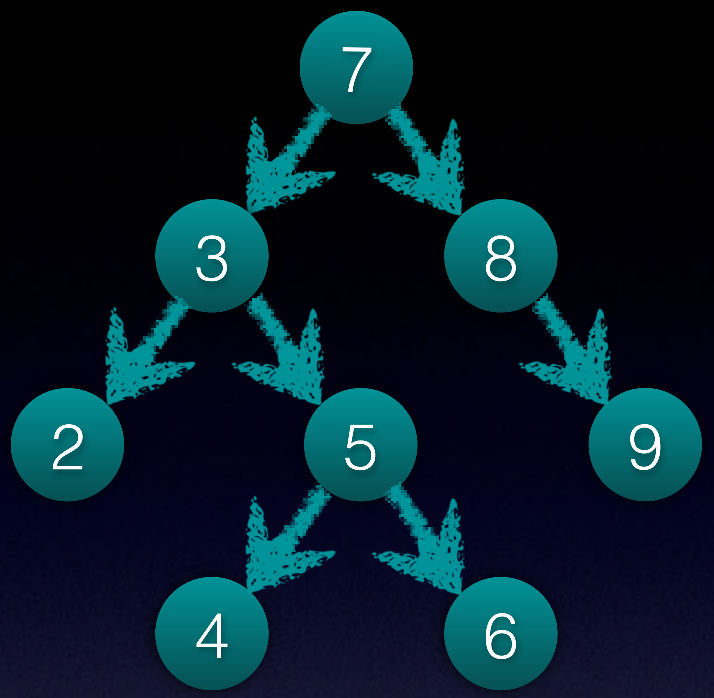
```
class Node(object):
   def __init__(self, value,
                left=None, right=None):
      self.value = value
      self.left = left
      self.right = right

tree = Node(value=7)
tree.left = Node(value=3)
tree.left.left = Node(value=2)
tree.left.right = Node(value=5)
tree.left.right.left = Node(value=4)
tree.left.right.right = Node(value=6)
tree.right = Node(value=8)
tree.right.right = Node(value=9)
```



Recursion

The code to walk the tree is:

```
def walk_inorder(node):
    if node is not None:
        walk_inorder(node.left)
        print(node.value)
        walk_inorder(node.right)
```

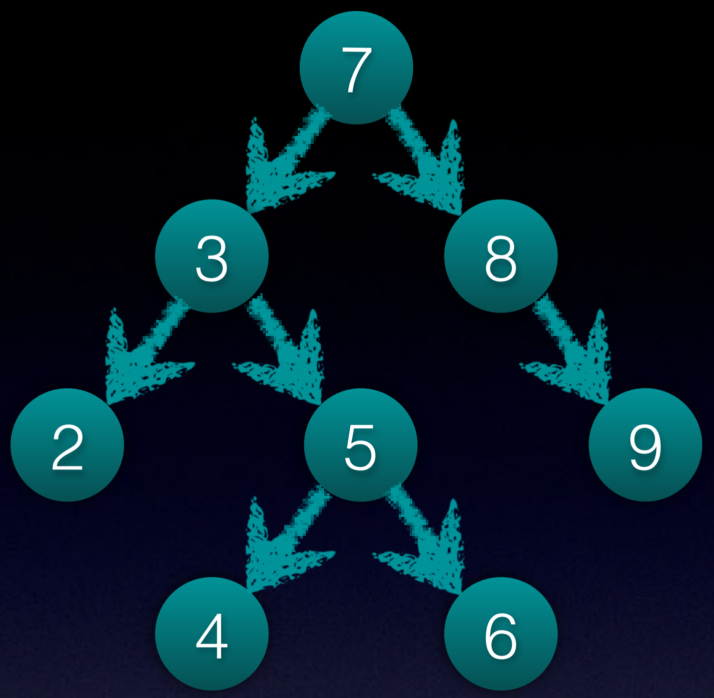


Executing this function on the example tree:

```
 2 3 4 5 6 7 8 9
```

Recursion

There are two other ways to walk through a binary tree:




```
def walk_preorder(node):
    if node is not None:
        print(node.value)
        walk_preorder(node.left)
        walk_preorder(node.right)

def walk_postorder(node):
    if node is not None:
        walk_postorder(node.left)
        walk_postorder(node.right)
        print(node.value)
```

Executing these functions on the example tree:

```
7 3 2 5 4 6 8 9    # preorder
2 4 6 5 3 9 8 7    # postorder
```

Recursion

What are the costs of a recursive solution?

- every function call costs time
- every nested function call costs memory

When the costs of using recursion are too high we might use an iterative algorithm. There are algorithms to iteratively walk through a binary tree (eg, Shorr Waite), but they are not as simple as the recursive approach.

It is possible to combine recursion with various techniques to achieve an efficient solution.

Recursion

The Fibonacci function:

```
def fibonacci(n):
    if n == 0:
        return 0
    if n == 1:
        return 1
    return fibonacci(n-1) + fibonacci(n-2)
```

Recursion

The Fibonacci function:

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def fibonacci(n):
    if n == 0:
        return 0
    if n == 1:
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    return fibonacci(n-1) + fibonacci(n-2)
```

This naive Fibonacci function is very inefficient. Try evaluating `fibonacci(40)`. It's really slow!

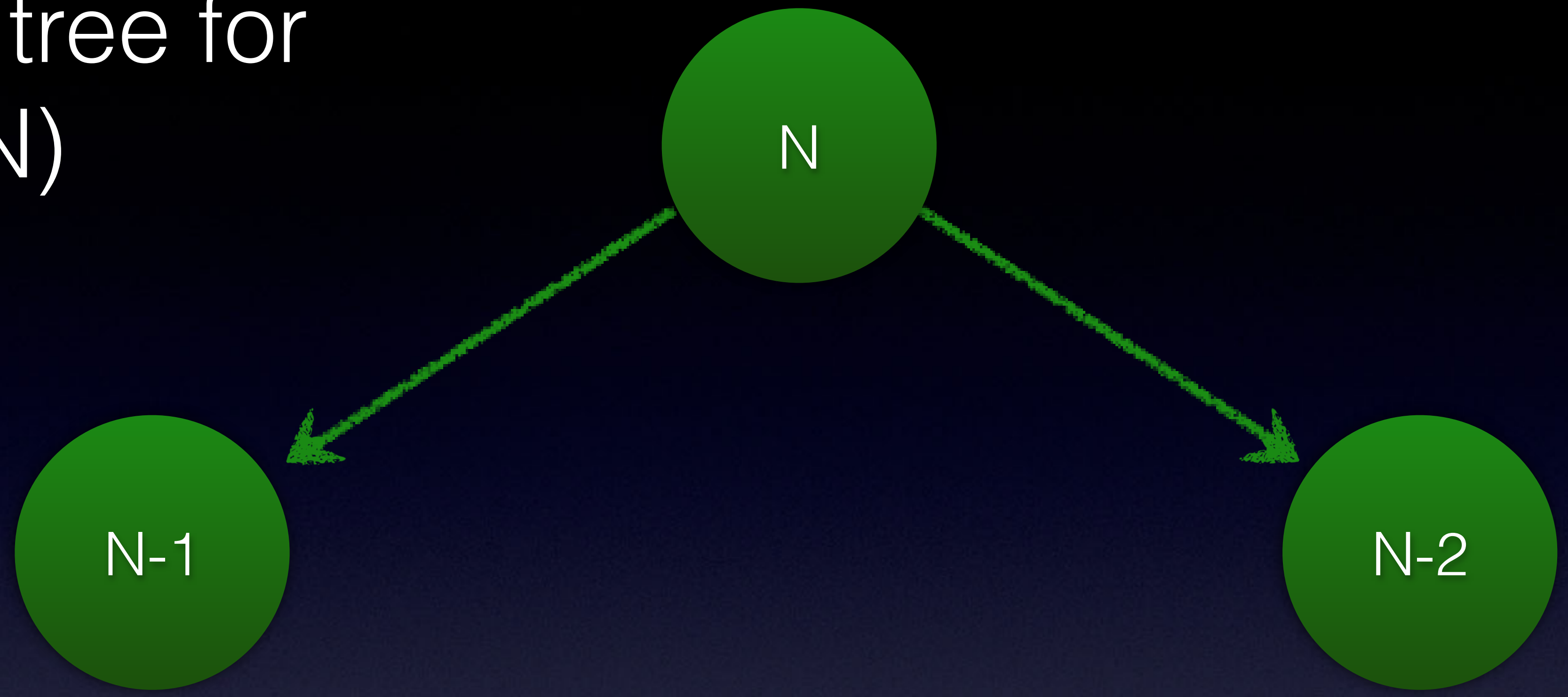
```
    fibonacci(40)=102334155  took 57.7889390s
```

Recursion

Evaluation tree for fibonacci(N)

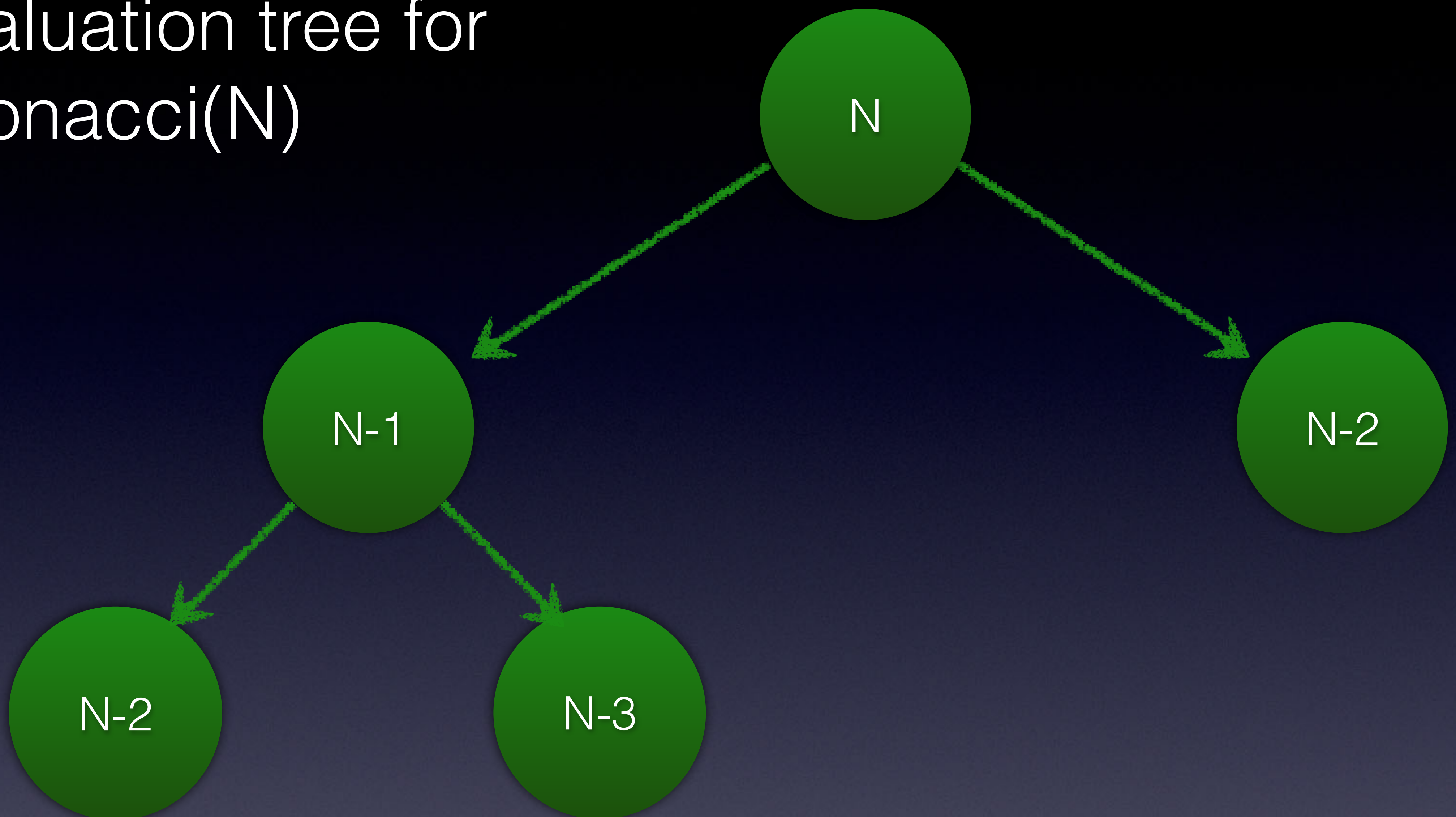
Recursion

Evaluation tree for fibonacci(N)



Recursion

Evaluation tree for fibonacci(N)



Recursion

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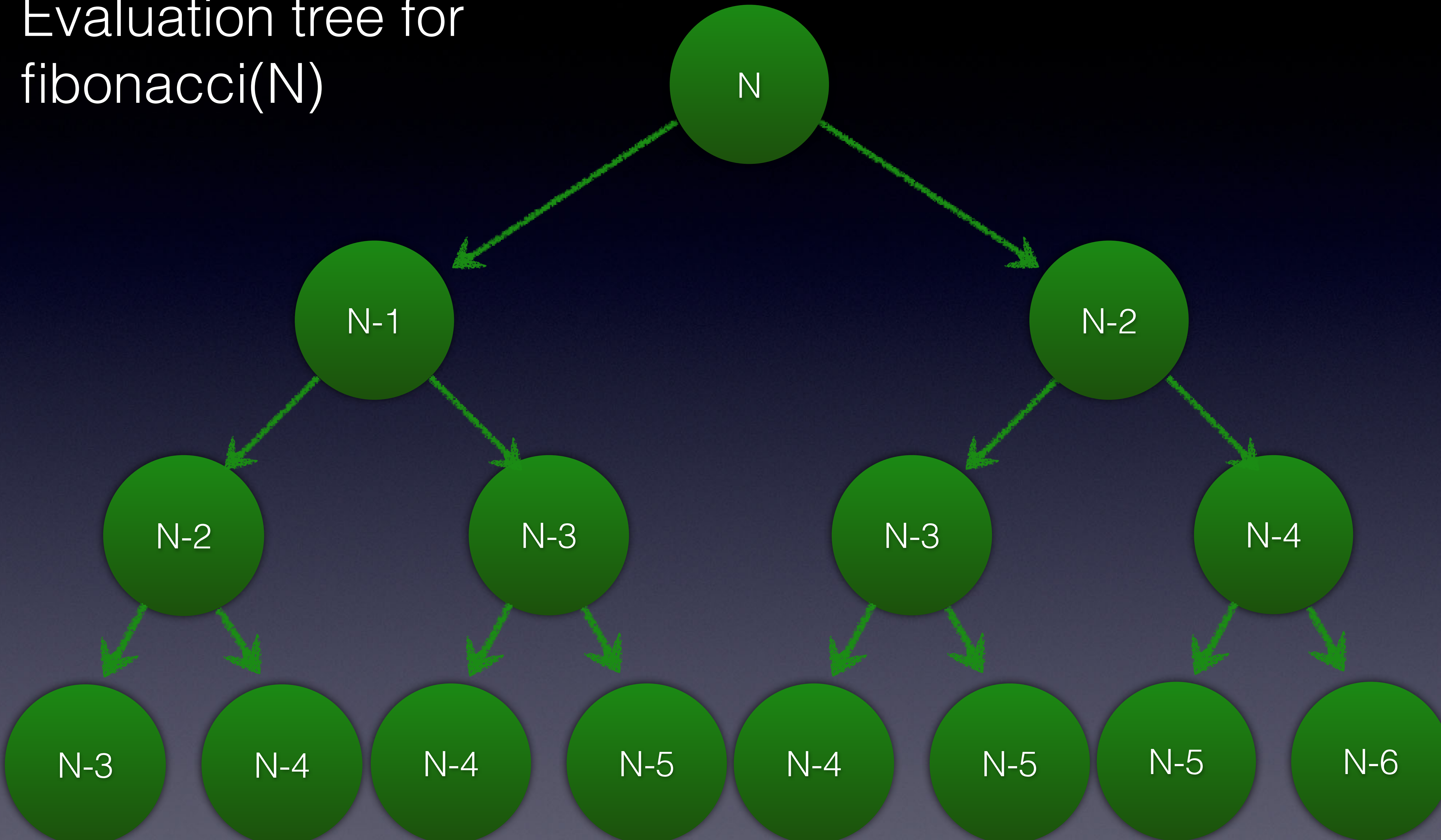

Recursion

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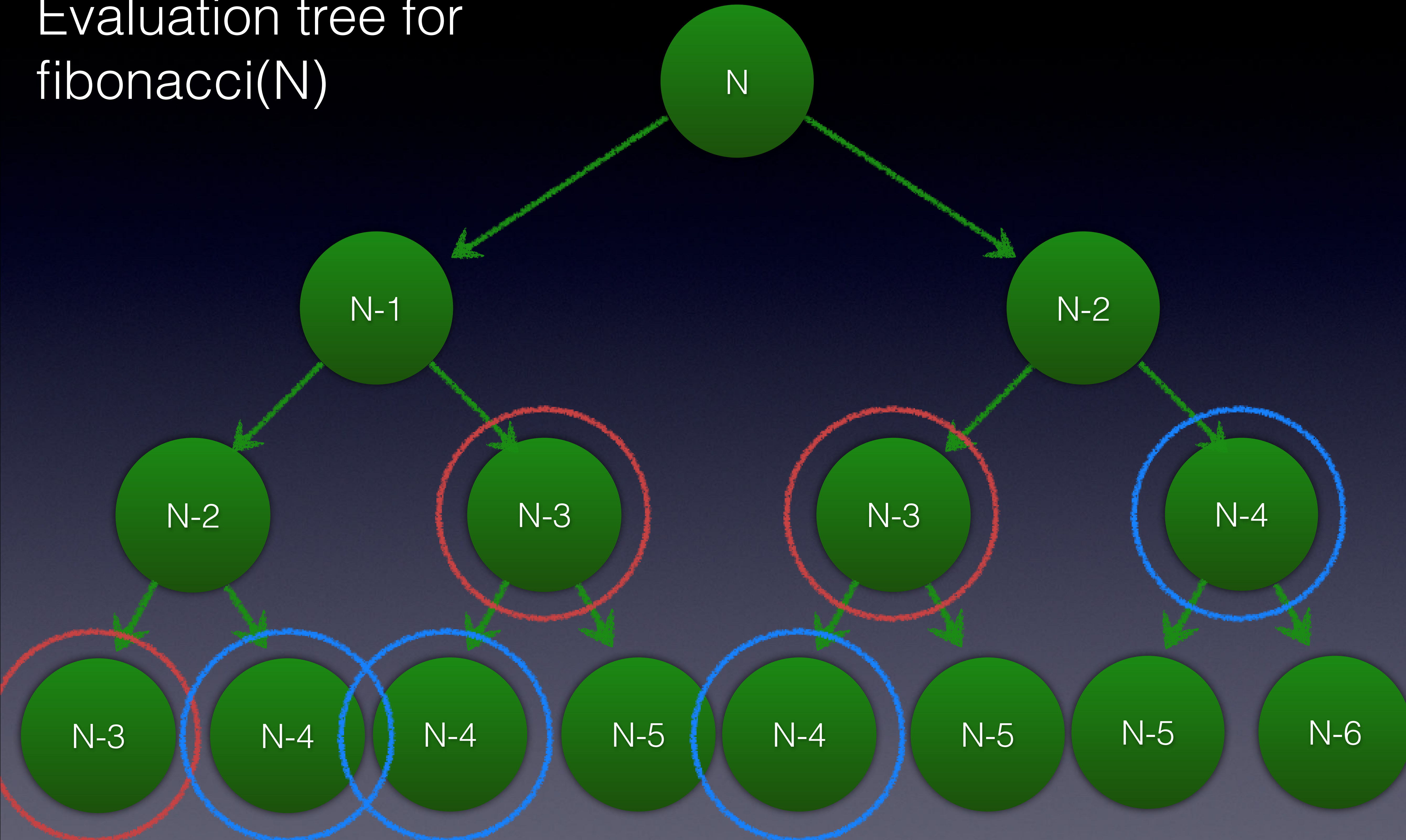

Recursion

If we could evaluate `fibonacci(m)` once and use the remembered result, we would save time.

```
fib_memo = {0: 0, 1: 1}    # memo + base cases

def fibonacci_memo(n):
   if n not in fib_memo:
      fib_memo[n] = (fibonacci_memo(n-1) +
                     fibonacci_memo(n-2))
   return fib_memo[n]
```

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fib_memo = {0: 0, 1: 1}    # memo + base cases

def fibonacci_memo(n):
   if n not in fib_memo:
      fib_memo[n] = (fibonacci_memo(n-1) +
                     fibonacci_memo(n-2))
   return fib_memo[n]
```

```
     fibonacci(40)=102334155  took 57.788939s
fibonacci_memo(40)=102334155  took  0.000025s
```

A useful speedup - more than 2,000,000 times faster!

Recursion

This approach is called memoisation. It's part of what is called dynamic programming.

There are other ways to memoise using decorators, but that's another talk.

Recursion

?

Code and PDF at:

code.google.com/p/rzzzwilson/source/browse/talks/recursion_1/