

The Extended Euclidean Algorithm

Definition : -----

We say that two numbers $a, b \in \mathbb{Z}$ are relatively prime if $\gcd(a, b) = 1$

Property :

if $a, b \in \mathbb{Z}$ are relatively prime then by the Euclidian Algorithm for $\exists n \in \mathbb{Z}^+, r_{n-1} = 1$

Proof :

we know that for any $a, b \in \mathbb{Z}$ Euclidian Algorithm after n iteration will find that

$d = \gcd(a, b) = \gcd(r_{n-1}, r_n)$ and $r_n = 0$ (by the well-ordering theorem). we know that all iteration share the same gcd, and because a and b are relatively prime

$\gcd(a, b) = \gcd(r_{n-1}, r_n) = \gcd(r_{n-1}, 0) = 1$, hence $r_{n-1} = 1$

* \rightarrow Algorithm for finding a^{-1}

let $a, b \in \mathbb{Z}$ be relatively prime By the above property if we apply the Euclidian Algorithm we will have that $r_{n-1} = 1$ meaning that $r_{n-3} \equiv 1 \pmod{r_{n-2}}$, because $r_{n-3} = q_{n-1} * r_{n-2} + r_{n-1}$, now i will try to just reformulate how we found r_n 's the results would be the same

algstart :

we start with normal Euclidean Algorithm and we will reformulate it,

$$\exists q_1, r_1 \in \mathbb{Z} \times \mathbb{Z}^+ \quad a = q_1 * b + r_1 \quad q_1 = \left\lfloor \frac{a}{b} \right\rfloor \text{ and } 0 \leq r_1 < b$$

we can rewrite r_1 as

$$r_1 = a - q_1 b = ax_1 + by_1$$

$$x_1 = 1$$

$$y_1 = -q_1.$$

if $r_1 = 0$ then $b = 1$ and $a = q_1$ because $\gcd(a, b) = 1$

else

we do the next euclidean division for b and r_1

$$\exists q_2, r_2 \in \mathbb{Z} \times \mathbb{Z}^+ \quad b = q_2 * r_1 + r_2 \quad q_2 = \left\lfloor \frac{b}{r_1} \right\rfloor \text{ and } 0 \leq r_2 < r_1$$

we can rewrite r_2 as $r_2 = b - q_2 * r_1 = b - q_2 * (a * x_1 + b * y_1)$

$$= b - q_2 * a * x_1 - q_2 * b * y_1$$

$$= a * (-q_2 * x_1) + b * (1 - q_2 * y_1)$$

$$= ax_2 + by_2$$

$$x_2 = -q_2 x_1$$

$$y_2 = 1 - q_2 y_1$$

if $r_2 = 0$ then $r_1 = 1$ and $b = q_2$

else

we do the next euclidean division for b and r_1

$$\exists q_3, r_3 \in \mathbb{Z} \times \mathbb{Z}^+ \quad r_1 = q_3 * r_2 + r_3 \quad q_3 = \left\lfloor \frac{r_1}{r_2} \right\rfloor \text{ and } 0 \leq r_3 < r_2$$

we can rewrite r_3 as

$$\begin{aligned} r_3 &= r_1 - q_3 * r_2 = a * x_1 + b * y_1 - q_3 (a * x_2 + b * y_2) \\ &= a * (x_1 - q_3 * x_2) + b * (y_1 - q_3 * y_2) \\ &= a x_3 + b y_3 \end{aligned}$$

$$x_3 = x_1 - q_3 x_2$$

$$y_3 = y_1 - q_3 y_2$$

if $r_3 = 0$ then $r_2 = 1$ and $r_1 = q_2$

else

...

so after $n - 1$ iteration of if else $r_n = 0$ and $r_{n-1} = 1$

$$\exists q_{n-1}, r_{n-1} \in \mathbb{Z} \times \mathbb{Z}^+ \quad r_{n-3} = q_{n-1} * r_{n-2} + r_{n-1} \quad q_{n-1} = \left\lfloor \frac{r_{n-3}}{r_{n-2}} \right\rfloor \text{ and } 0 \leq r_{n-1} < r_{n-2}$$

we can rewrite r_{n-1} as

$$\begin{aligned} r_{n-1} &= r_{n-3} - q_{n-1} * r_{n-2} \\ &= a x_{n-1} + b y_{n-1} \end{aligned}$$

$$x_{n-1} = x_{n-3} - q_{n-1} * x_{n-2}$$

$$y_{n-1} = y_{n-3} - q_{n-1} y_{n-2}$$

so we can see that $x_{n-1} = x$ and $y_{n-1} = y$ (because $r_n = 0$), so we found $r_{n-1} = a * x + b * y$ which implies $a * x = -b * y + 1$, we did found $x \in \mathbb{Z}$ for wich $a * x \equiv 1 \pmod{b}$ we give x the symbole a^{-1} and we name it the multiplicative inverse

:endalg

for a clean generalisation, first in the following line $r_1 = a - q_1 b = a x_1 + b y_1$ with $x_1 = 1$ and $y_1 = -q_1$, we see here that x_1 and y_1 doesn't have the general form neither does r_1 , so we put $r_{-1} = a$ and $r_0 = b$ wich gives $x_{-1} = 1$ and $x_0 = 0$ and finally $y_{-1} = 0$ and $y_0 = 1$ which gives the following initial and iteration condition's

$$r_1 = r_{-1} - q_1 r_0, \text{ with } n \text{ iteration } r_{n-1} = r_{n-3} - q_{n-1} r_{n-2} \text{ and } r_n = 0$$

$$x_1 = x_{-1} - q_1 x_0, \text{ with } n \text{ iteration } x_{n-1} = x_{n-3} - q_{n-1} x_{n-2} \text{ and } x = x_{n-1}$$

$$y_1 = y_{-1} - q_1 y_0, \text{ with } n \text{ iteration } y_{n-1} = y_{n-3} - q_{n-1} y_{n-2} \text{ and } y = y_{n-1}$$