

# DISCRETE LOGARITHM

*Definition : -----*  
*for integer  $a$  relatively prime with  $n$ , and let  $M = \{k \in \mathbb{Z}, a^k \equiv 1 \pmod{n}\}$  we define the general expression  $a^m \equiv 1 \pmod{n}$  with  $m = \min M$  the existence of such  $m$  is always true, and the set  $M$  is non empty for  $\phi(n) \in M$ , and we say that  $m$  is the length of the period generated by  $a$ , and if  $m = \phi(n)$  for an  $a$  we say that  $a$  is primitive root of  $n$ .*

*Property :*  
*any integer  $z$  can be represented as  $z = q + k\phi(n)$  with  $0 \leq q < \phi(n)$*

*Proof :*  
*we divide  $z$  by  $\phi(n)$  in the euclidian sense we get the desired result.*

*Property :*  
*if  $a$  is primitive root for  $n$ , then  $a^p = a^q \pmod{n}$  if and only if  $p \equiv q \pmod{\phi(n)}$ .*

*Proof :*  
*for  $a$  is primitive root for  $n$  we take  $p = q + k\phi(n)$  we notice that if  $p \equiv q \pmod{\phi(n)}$ ,  $a^p \equiv a^q \times a^{k\phi(n)} \equiv a^q \pmod{n}$ , if  $a^p = a^q \pmod{n}$ , then  $a^{p-q} = 1 \pmod{n}$ , meaning that  $p - q = k\phi(n)$ , hence  $p \equiv q \pmod{\phi(n)}$ .*

*we introduce the discrete logarithm as  $dlog_{n,a}(b) = i$  with  $a$  is a primitive root of  $n$ , that verify the equation  $b \equiv a^i \pmod{n}$ , if  $x \equiv a^{dlog_{n,a}(x)} \pmod{n}$ ,  $y \equiv a^{dlog_{n,a}(y)} \pmod{n}$  and  $xy \equiv a^{dlog_{n,a}(xy)} \pmod{n}$  then  $xy \equiv (x \pmod{n} \times y \pmod{n}) \pmod{n}$ , hence  $a^{dlog_{n,a}(xy)} \equiv a^{dlog_{n,a}(x) + dlog_{n,a}(y)} \pmod{n}$ , and by the previous property  $dlog_{n,a}(xy) \equiv (dlog_{n,a}(x) + dlog_{n,a}(y)) \pmod{\phi(n)}$  we can generalize this to  $dlog_{n,a}(x^y) \equiv y \times dlog_{n,a}(x) \pmod{\phi(n)}$*

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