

CS & IT ENGINEERING

DISCRETE
MATHS
GRAPH THEORY



Lecture No. 16



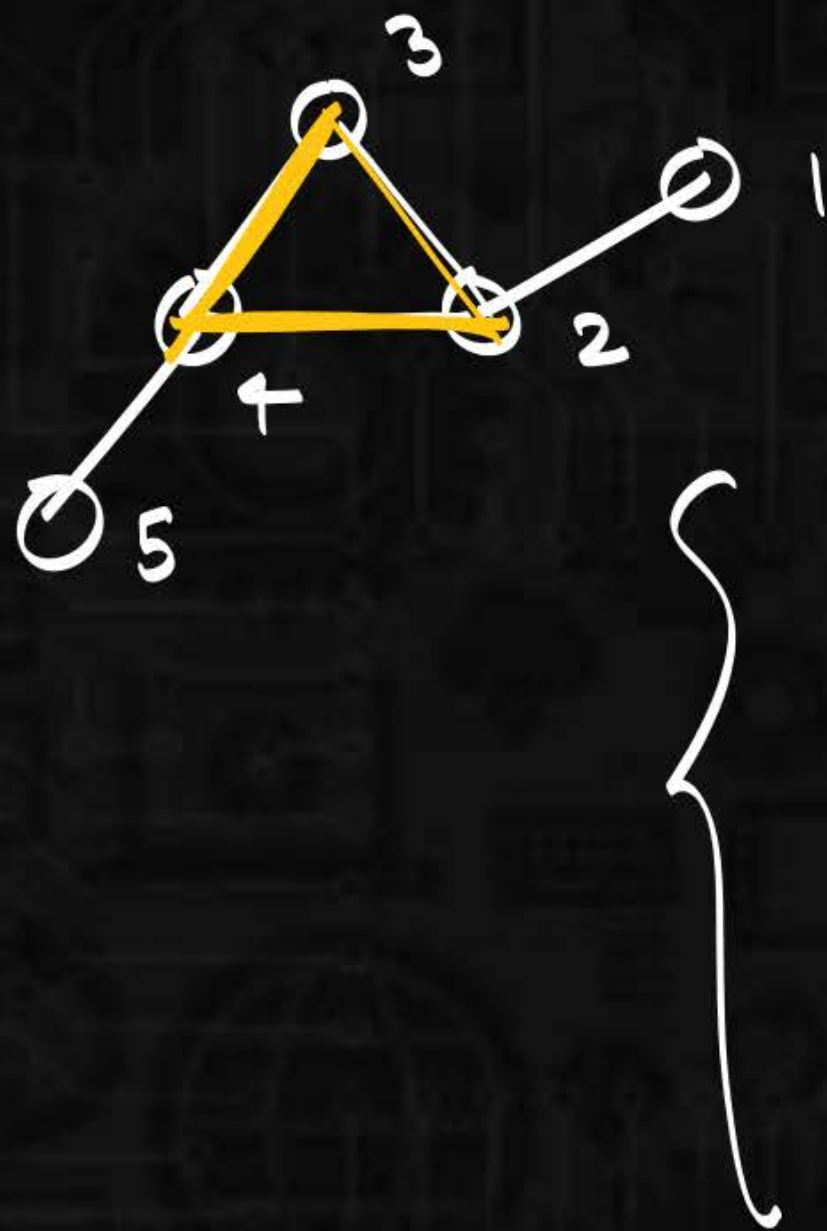
By- SATISH YADAV SIR

TOPICS

01 INDEPENDCE NO.

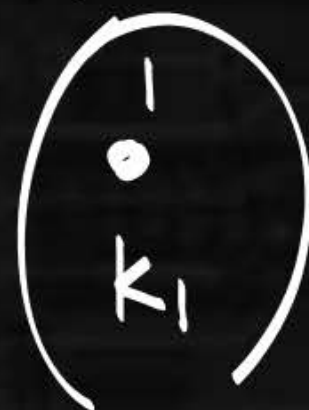
02 DOMINATION NO.

03 MATCHING NO .



clique..

subgraph

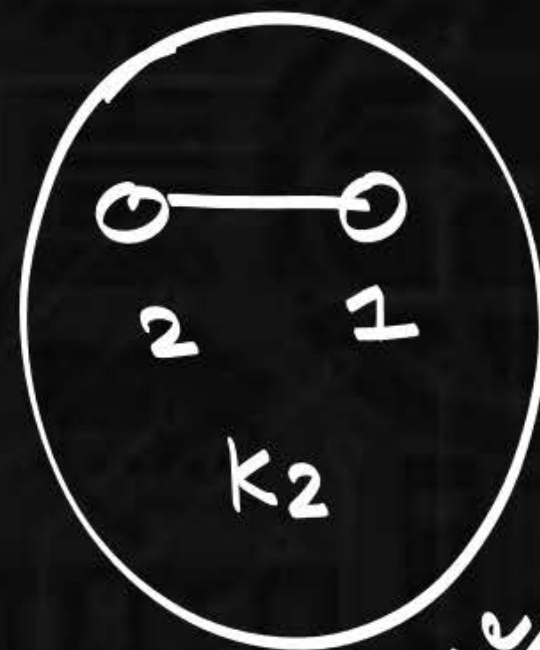


subgraph



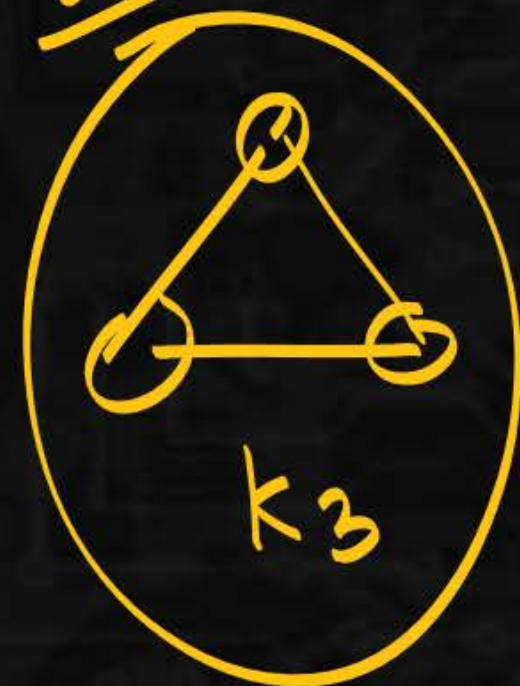
Clique no : $\boxed{w(G) = 3}$

subgraph

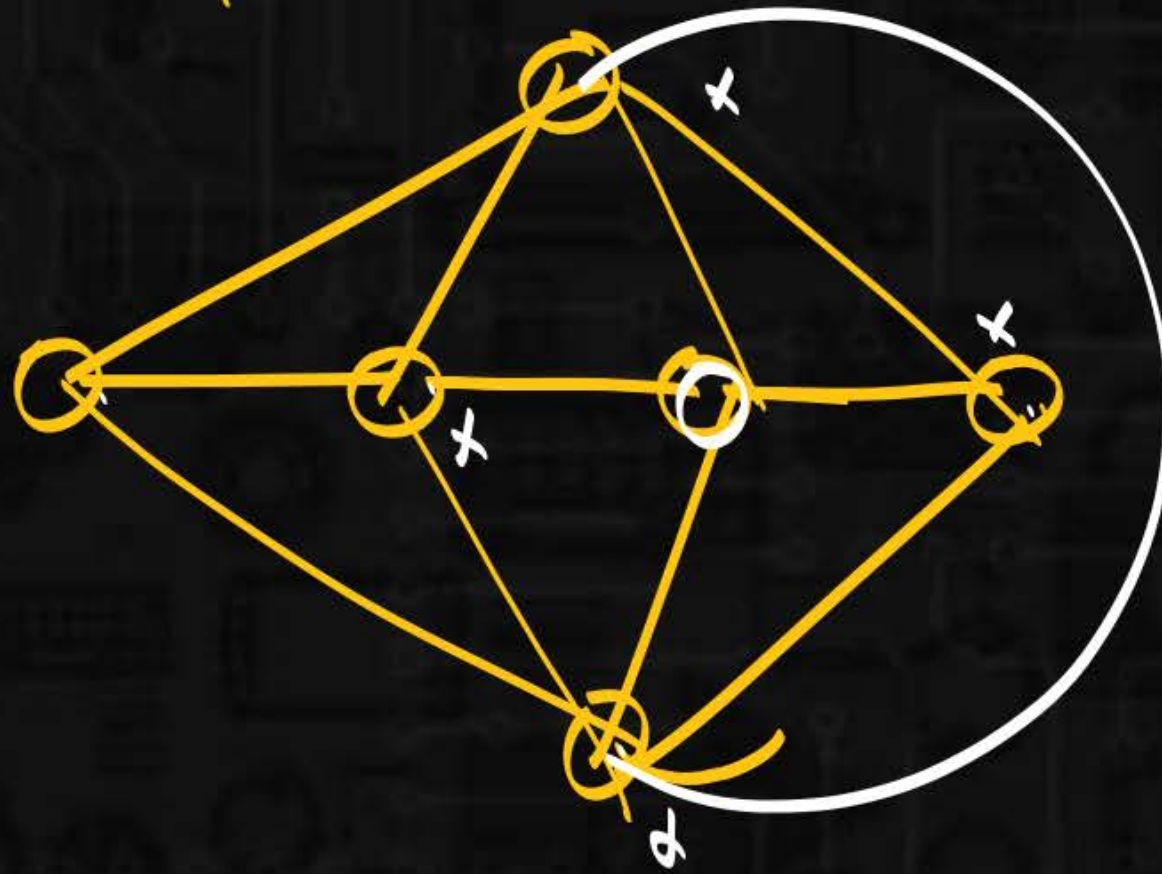


clique

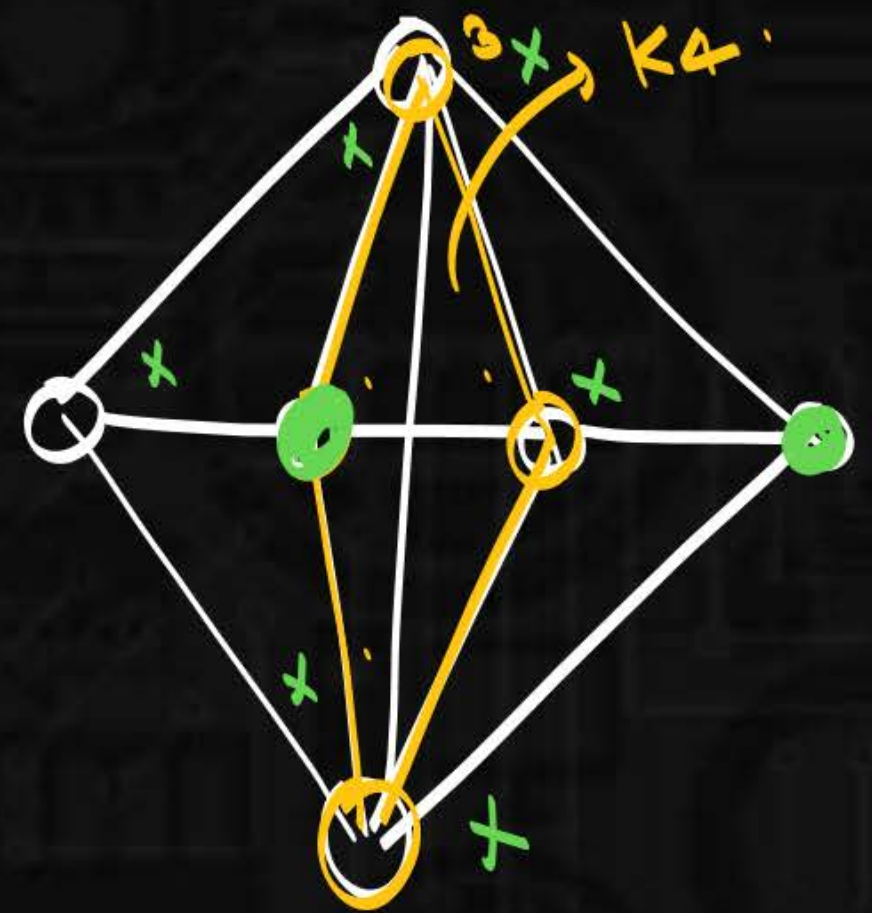
subgraph

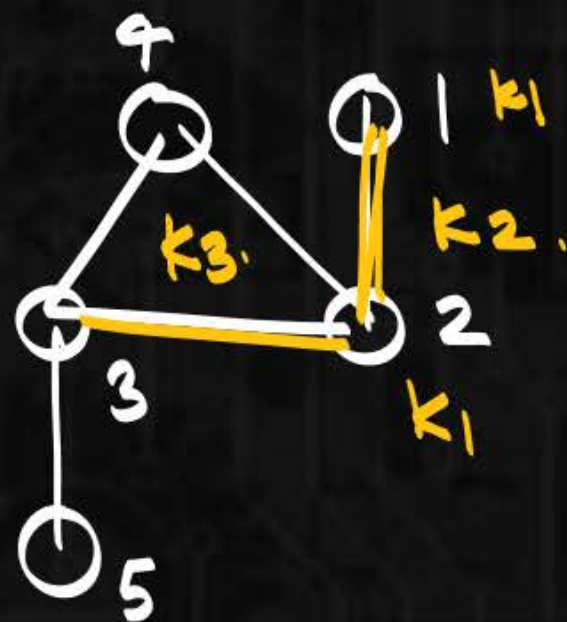


maximum size of clique. / 4



maximum size of Independent set 2





Subgraphs.

$n=1$

①

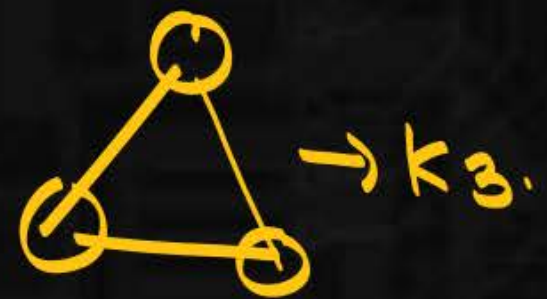
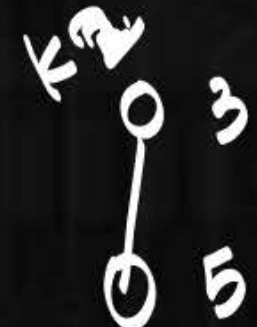
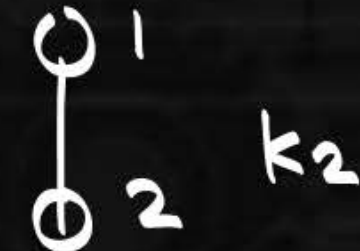
⑤

②

③

④

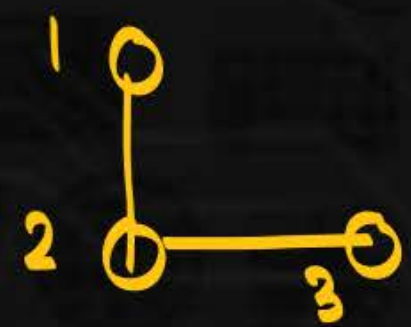
$w(h)=3$



Graph operation :

- 1) union. (\cup)
- 2) Intersection (\cap)
- 3) Ring sum (\oplus)

G_1

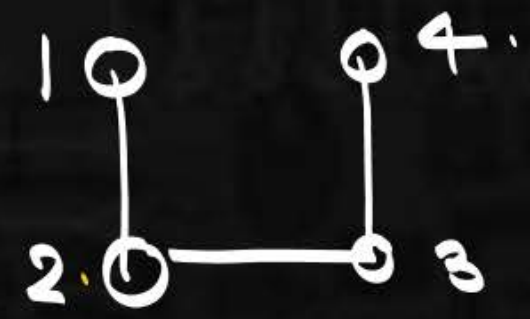


$$G_1 = (V_1, E_1)$$

$$V_1 = \{1, 2, 3\}$$

$$E_1 = \{12, 23\}$$

G_2

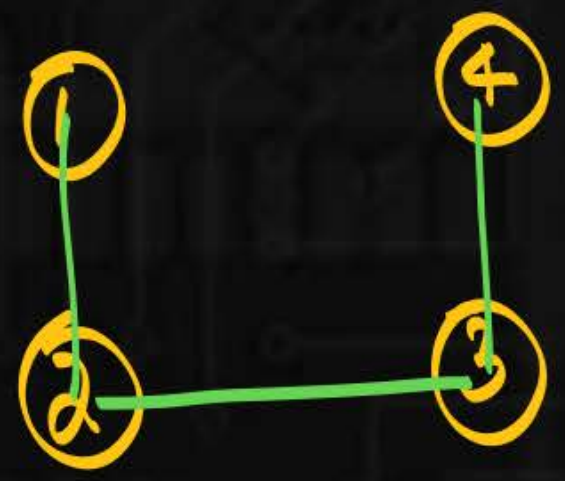


$$G_2 = (V_2, E_2)$$

$$\cup \quad V_2 = \{1, 2, 3, 4\} = V_3$$

$$\cup \quad E_2 = \{12, 23, 34\} = E_3$$

$$G_3 = (V_3, E_3)$$



$$G_3 = (V_3, E_3)$$

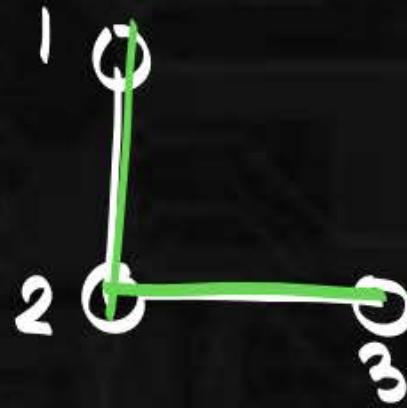
Intersection (n)

$$G_4 = (V_4, E_4)$$

$$V_4 = V_1 \cap V_2$$

$$E_4 = E_1 \cap E_2$$

$$G_1 = (V_1, E_1)$$



$$V_1 = \{1, 2, 3\}$$

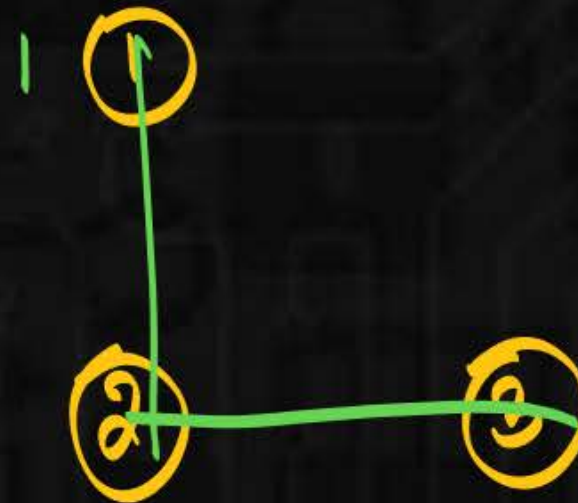
$$E_1 = \{\underline{12}, \underline{23}\}$$

$$G_2 = (V_2, E_2)$$



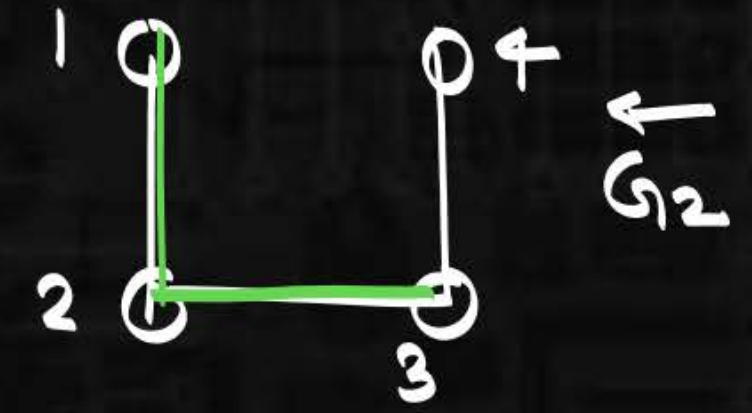
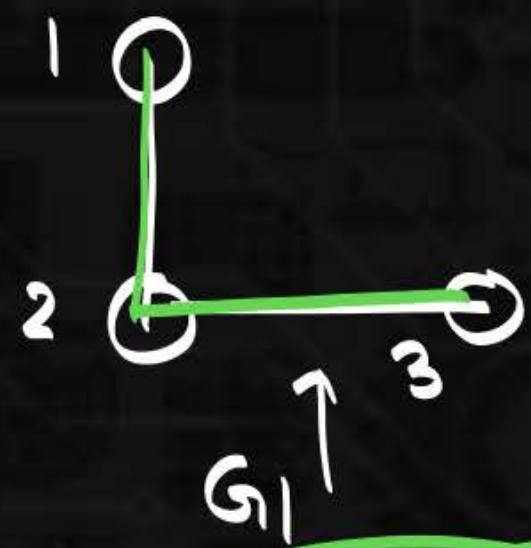
$$V_2 = \{1, 2, 3, 4\} \cap V_1 = V_4 = \{1, 2, 3\}$$

$$E_2 = \{\underline{12}, \underline{23}, 34\} \cap E_1 = \{\underline{12}, \underline{23}\}$$



Ring sum (\oplus)

$$G_1 = (V_1, E_1) \quad G_2 = (V_2, E_2)$$



$$G_6 = (V_6, E_6)$$

$$V_6 = V_1 \cup V_2$$

$$E_6 = (E_1 \cup E_2) - (E_1 \cap E_2)$$

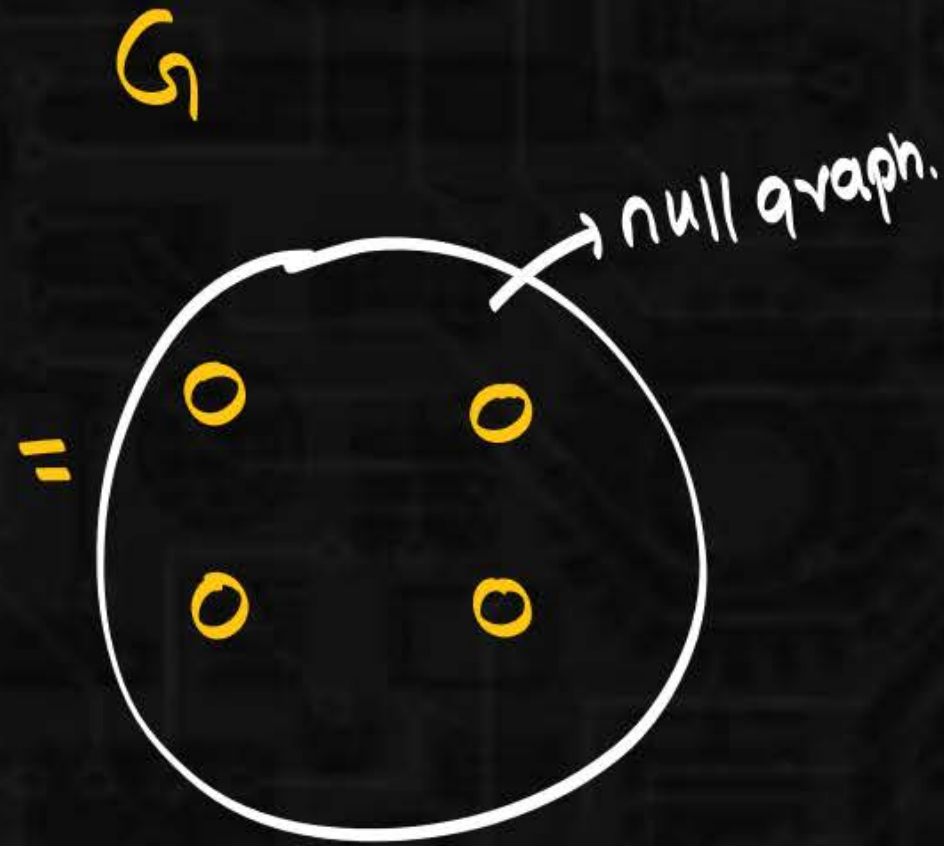
$$G_6 = (V_6, E_6)$$



→ edges in G_1 or G_2 but not in both.

$$G_1 \oplus G_1 = \text{null graph}$$

$$G \oplus G = \text{null graph.}$$

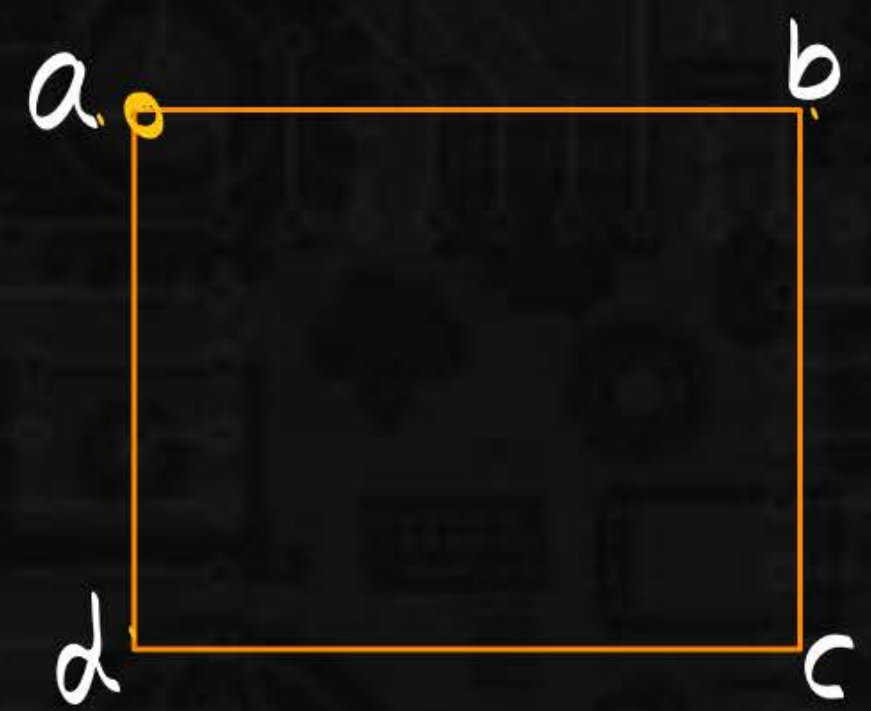


null graph.
→ set of isolated vertices.

Graph representation: $n \times n$ - symmetric binary matrix.

adjacency matrix / connection matrix.
 $x_{ij} = 1$
 $(n \times n)$ - binary matrix

connection
bet'n
vertex i
& j

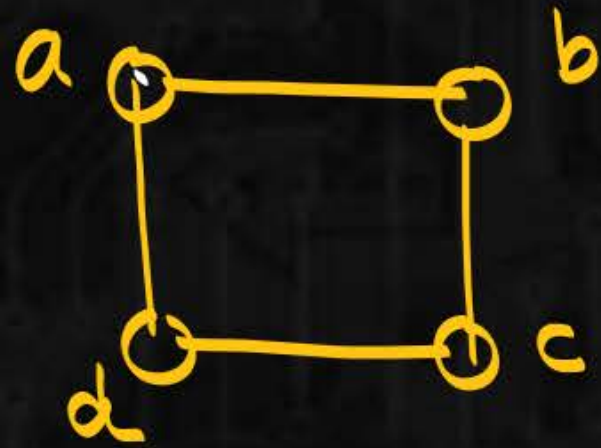


$$A(G) = \begin{matrix} & \begin{matrix} a & b & c & d \end{matrix} \\ \begin{matrix} a \\ b \\ c \\ d \end{matrix} & \begin{bmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \end{bmatrix} \end{matrix}_{4 \times 4}$$

$= 0$ (else)

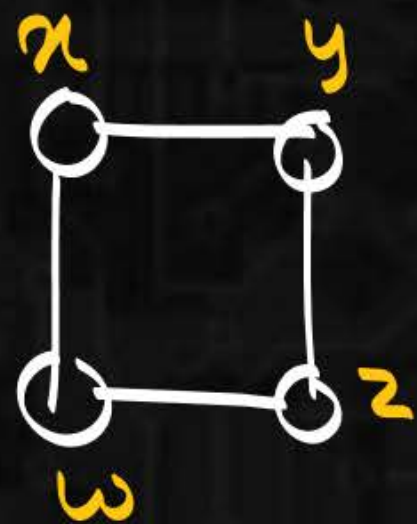
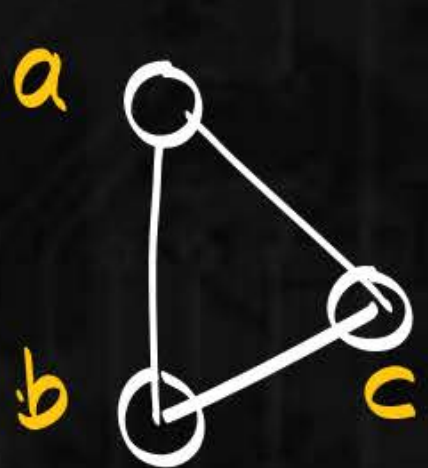
$$x_{ij} = x_{ji}$$

$$x_{ab} = x_{ba}$$



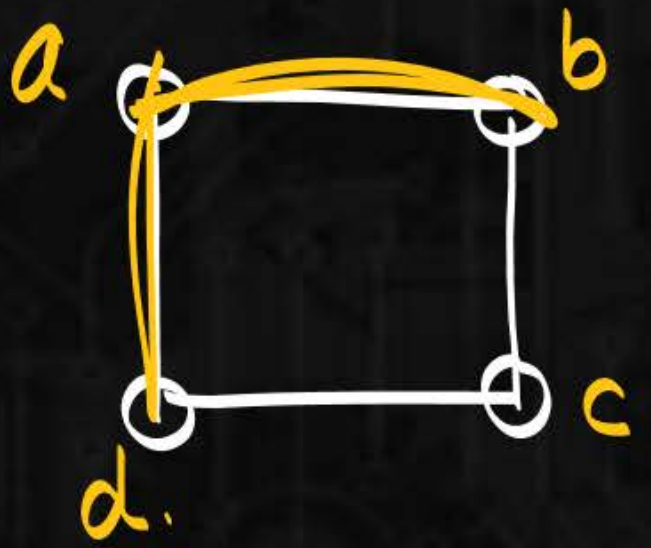
$$\rightarrow \begin{matrix} & \begin{matrix} a & b & c & d \end{matrix} \\ \begin{matrix} a \\ b \\ c \\ d \end{matrix} & \begin{bmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \end{bmatrix} \end{matrix} \quad \begin{matrix} 1. \\ 4 \times 4 \end{matrix}$$

- no self loop.
- Diagonal elements will be zero.
- no. of 1's are present in row/column will give degree.



$$A(G) = \left[\begin{array}{c|c} A(G_1) & \\ \hline & A(G_2) \end{array} \right]$$

	a	b	c	x	y	z	w
a	0	1	1	0	0	0	0
b	1	0	1	0	0	0	0
c	1	1	0	0	0	0	0
x	0	0	0	0	1	0	1
y	0	0	0				
z	0	0	0				
w	0	0	0				



$A =$

	a	b	c	d
a	0	1	0	1
b	1	0	1	0
c	0	1	0	1
d	1	0	1	0

$A^2(i,i) = \text{degree of vertex } v_i$

$0 \times 0 + 1 \times 1 + 0 \times 0 + 1 \times 1 = 2.$

$A^2 = A \times A =$

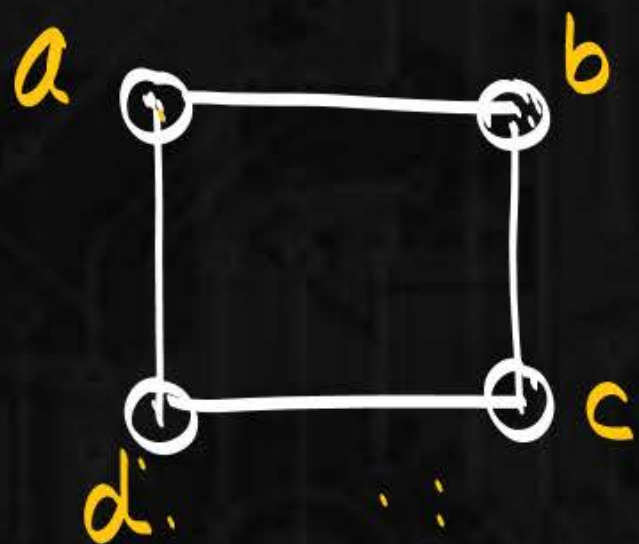
0	1	0	1
1	0	1	0
0	1	0	1
1	0	1	0

0	1	0	1
1	0	1	0
0	1	0	1
1	0	1	0

$\Rightarrow A^2$

	a	b	c	d
a	2	0	2	0
b	0	2	0	2
c	2	0	2	0
d	0	2	0	2

$a \rightarrow b$ (length 3) walk



$A^2 =$

2	0	2	0
0	2	0	2
2	0	2	0
0	2	0	2

$A^3(a,b) = \{4 \text{ walks of length 3 from } a, b\}$

$A^n(i,j) = \text{no. of walks from } i, j \text{ of length } n$

$A^3(a,b): a \rightarrow b$
4 walks of length 3:

$\begin{cases} a-d-a-b \\ a-d-c-b \\ a-b-a-b \\ a-b-c-b \end{cases}$

$A^3 =$

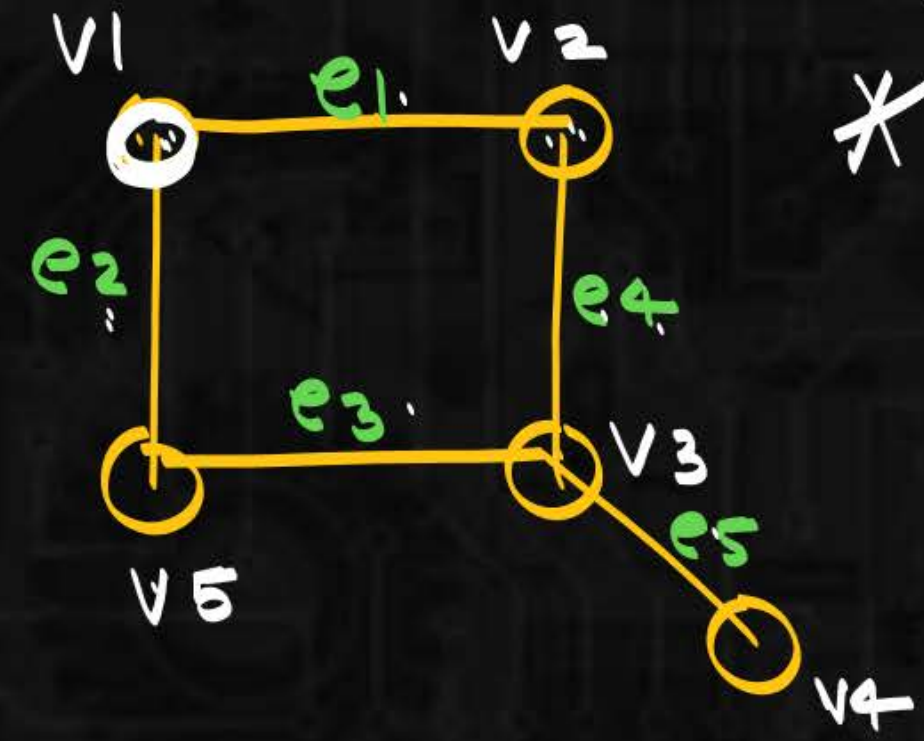
2	0	2	0
0	2	0	2
2	0	2	0
0	2	0	2

0	1	0	1
1	0	1	0
0	1	0	1
1	0	1	0

$A^3:$

	a	b	c	d
a	0	4	0	4
b	4	0	4	0
c	0	4	0	4
d	4	0	4	0

Incident matrix (m(G)) :



* $\rightarrow \underline{v1}$

$\rightarrow v2$

$v3$

$v4$

$v5$

	e_1	e_2	e_3	e_4	e_5
$\underline{v1}$	1	1	0	0	0
$v2$	1	0	0	1	0
$v3$	0	0	1	1	1
$v4$	0	0	0	0	1
$v5$	0	1	1	0	0

* $\left\{ \begin{array}{l} \text{rank}(m(G)) \\ = \underline{n-1} \end{array} \right.$

- \rightarrow Every column will have only 2 ones.
- \rightarrow every row will give degree of a vertex.

