

CS & IT ENGINEERING

Operating Systems

Memory Management



Lecture No. 4



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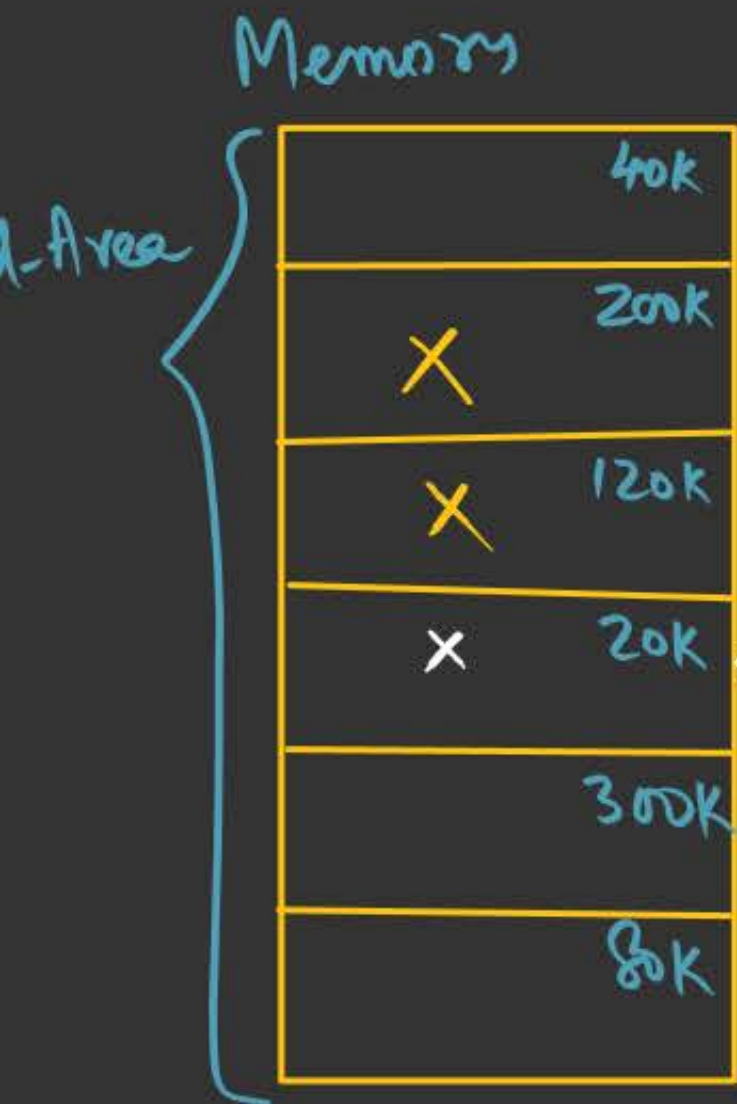
TOPICS TO BE COVERED

Fixed Partitions

Variable Partitions

Address Space

1) Fixed Partitions (MFT: Multiprogramming with Fixed No. of Tasks)



P_{process} : "15K"

→ Static

Partition Allocation Policies

15K (i) First Fit : "First free big enough"

20K (ii) Best Fit : "Smallest free big enough" (Int. Frag.)

May work faster to FF (iii) Next Fit : works like first fit, Search for free Partition begins from Last Alloc.

(iv) Worst Fit : largest free big enough

← Last Alloc

Performance of Fixed Partitions

1. Int. Fragmentation : ✓
(I.F)

2. Ext. Fragmentation : X
(E.F)

3. Degree of M.P. : Limited

4. Max. Process Size : Ltd

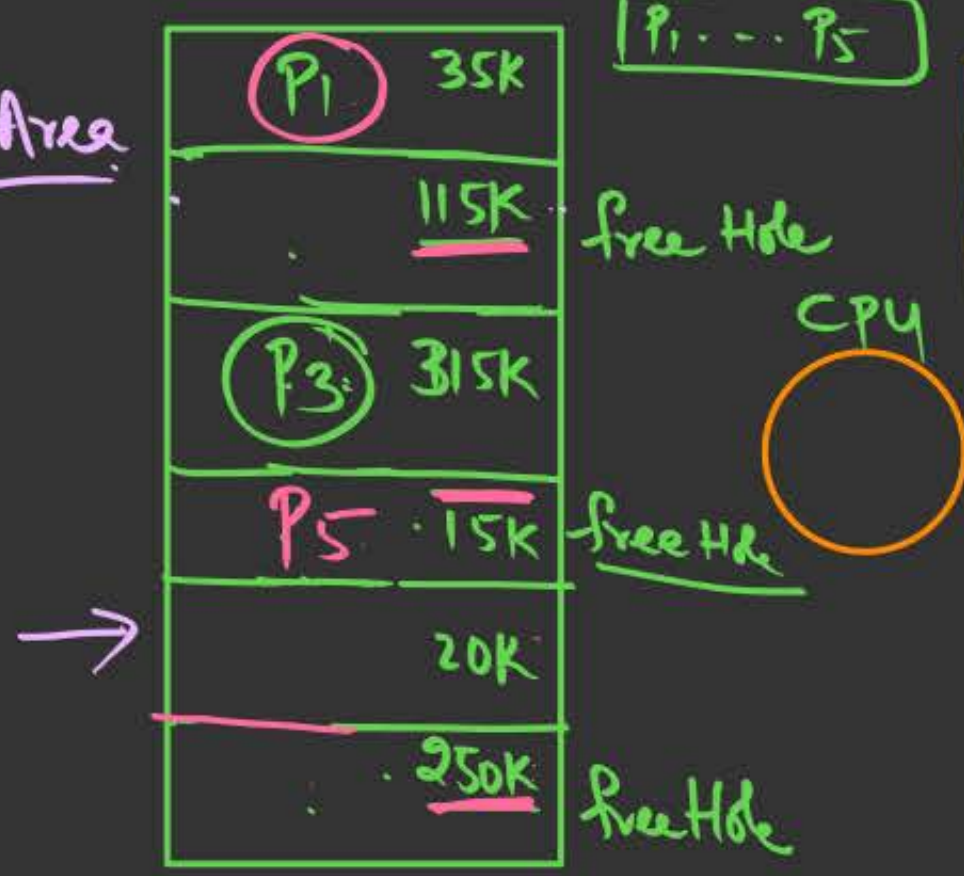
5. Part. Alloc. Policy : B.F
(less I.F)

2) Variable Partitions: (MVT: M.Pr. with variable Tasks) ^{Dynamic Partitioning}

Memory (free hole) P1...P5

Process Req's: $\langle 35K; 115K; 315K; 15K; 120K, \dots \rangle$
 t_0 P_1 P_2 P_3 P_4 P_5

U-Area



- Perf. Issues:
- 1) Int. Frag: X
 - 2) Ext. Frag: ✓
 - 3) Degree of M.Pr: FLEXI
 - 4) Max. Process Size: FLEXI
 - 5) Part. Alloc. Policy: W.F is better

$P_{new}: 13K$
 P_6

- 1) FF: 115K
- 2) BF: 15K (2K)
- 3) NF: 250K
- 4) WF: 250K

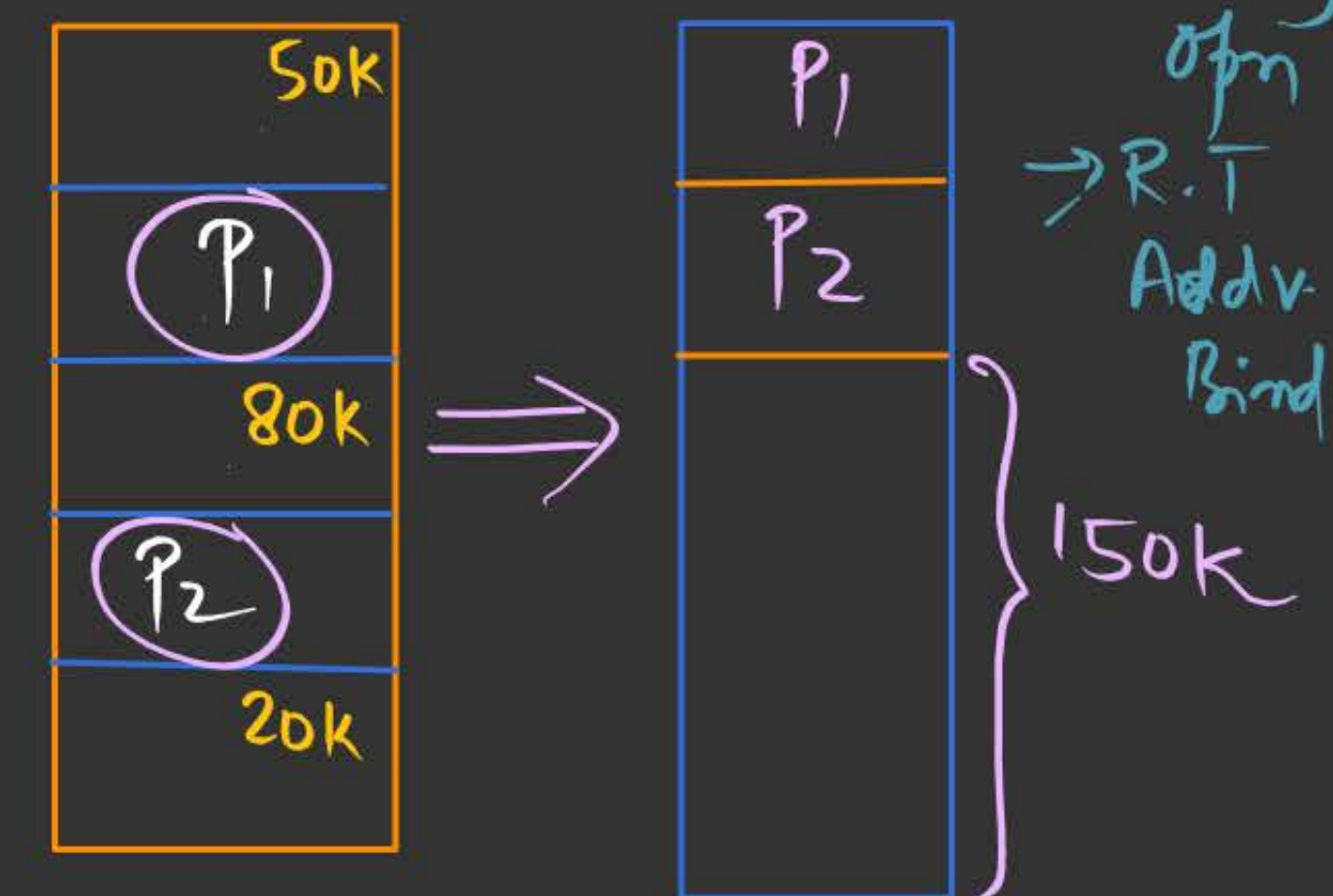
Total free Memory $P_{new}: 280K$

Adjacent free holes are automatically merged to one free hole

External Fragmentation

(i) Compaction

Relocate \rightarrow Time Consuming opn



Memory Map

(ii) N-CG Alloc

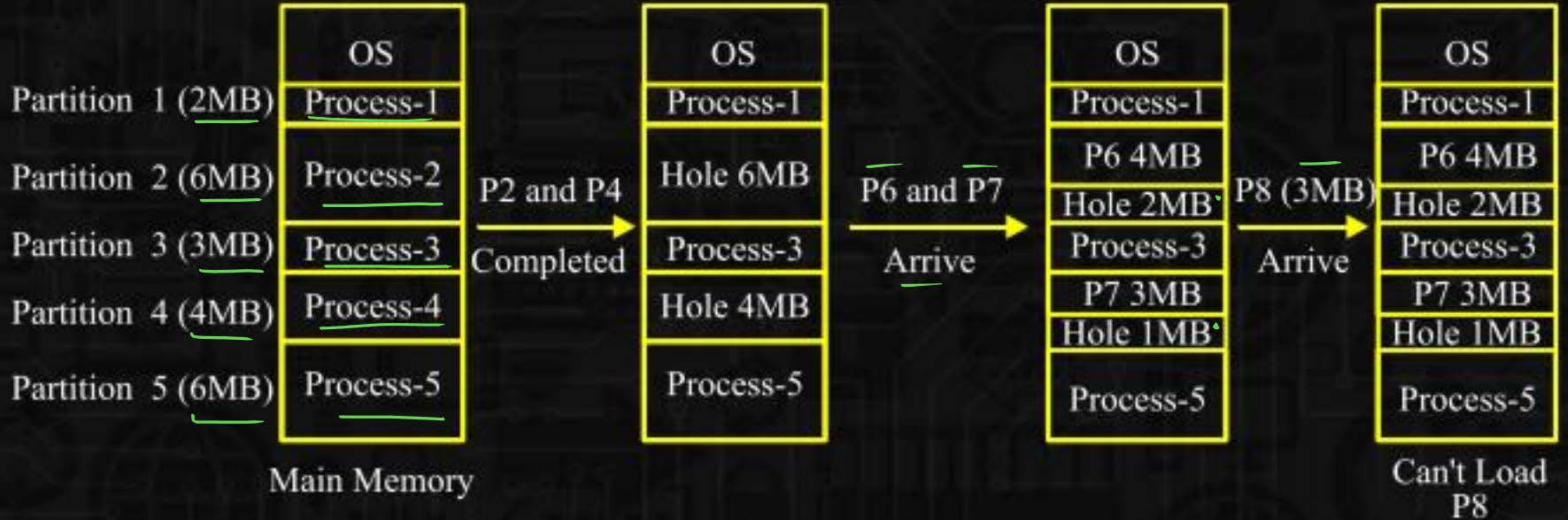
$P_{new}: 85K$ 50K
35K
 E.F: 150K

RQ → PCB₁ PCB₂ PCB₃

Variable Partitions:

E-F: 3MB

Memory



Q.1

Consider a Memory System having 6 Partitions of sizes 200K; 400K; 600K; 500K; 300K; 250K. There are 4 Processes of sizes: 357K; 210K; 468K; 49K. Using Best Fit Allocation Policy, what Partitions are not allocated/ remains Unallocated?



P4Q

P ₄	200K
P ₁	400K
	600K
P ₃	500K
	300K
P ₂	250K

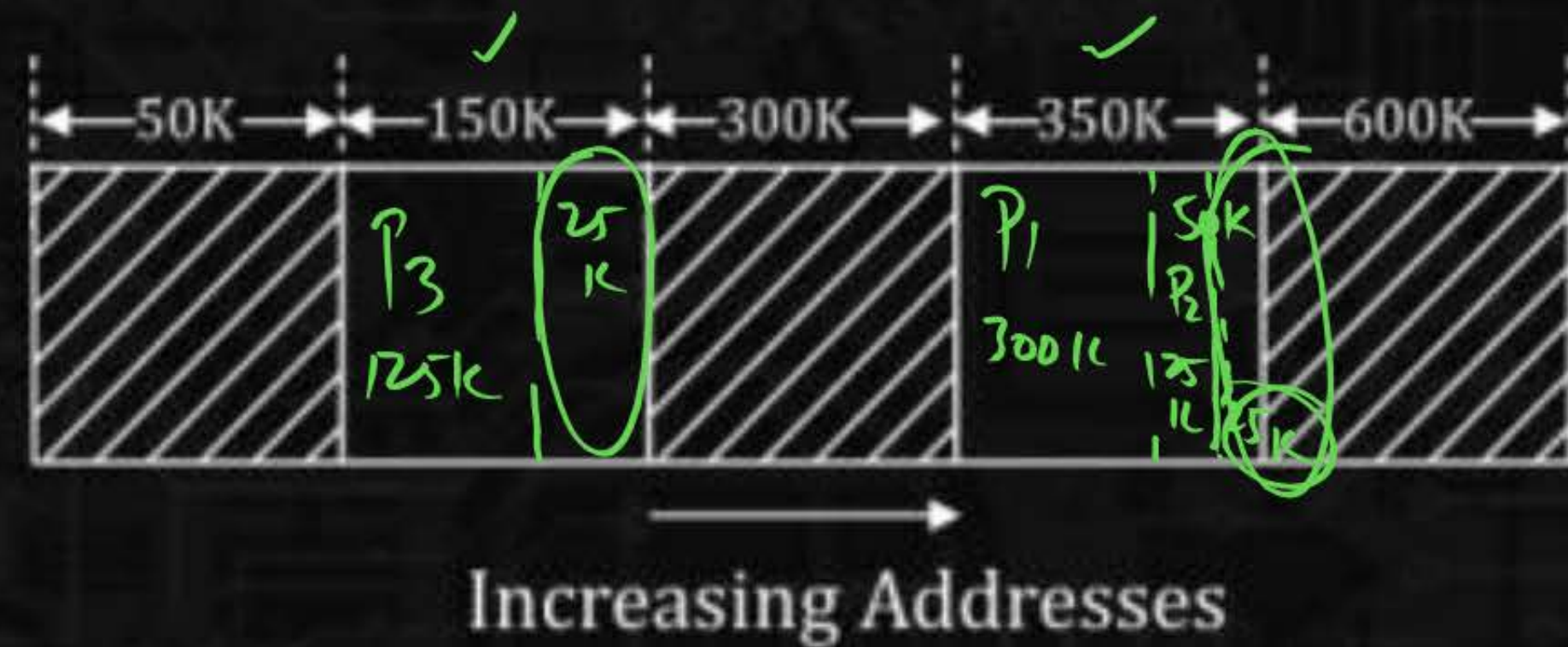
600K, 300K

Q. 2



Consider the following Memory Map in which blank regions are not in use and hatched regions are in use. Using Variable Partitions with no Compaction:

The sequence of requests for blocks of sizes 300K, 25K, 125K, 50K can be satisfied if we use:



- A. Either first fit or best fit policy (any one)
- B. First fit but not best fit policy ✓
- C. Best fit but not first fit policy
- D. None of the above.

F.F : ✓
B.F : X

Q. 3

Consider a System with Memory of size 1000KBytes. It uses Variable Partitions with no Compaction. Presently there are 2 partitions of sizes 200K & 260K respectively.

(i) What is the allocation request of the Process which would always be denied?

A. 131 K

B. 151 K

C. 181 K

D. 541 K

(ii) The smallest Allocation Request which could be denied is:

A. 131 K

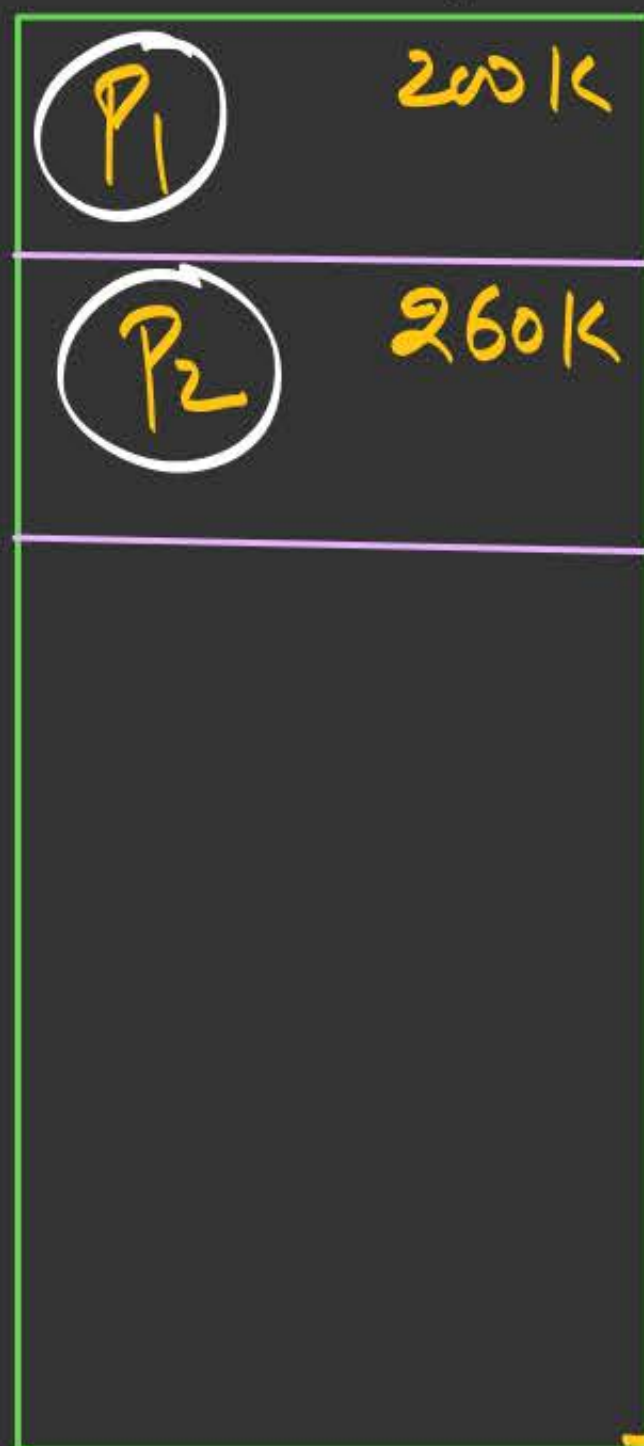
B. 151 K

C. 181 K

D. 541 K

Case I

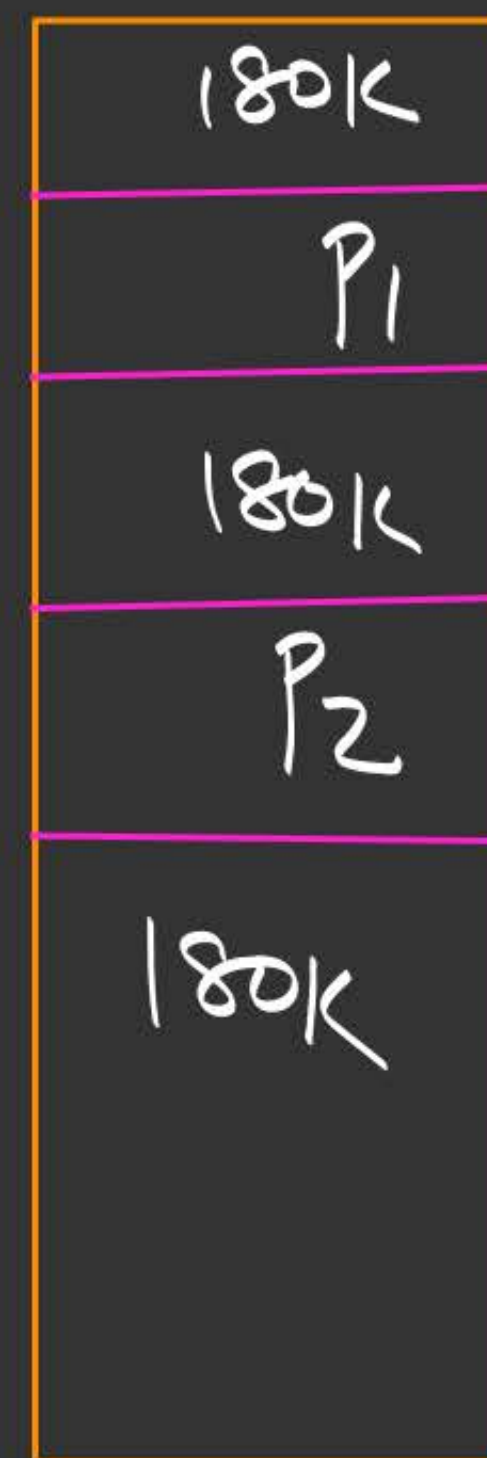
1-free Hole



541K



541



181K ; 541K

$$\frac{540}{3} = \underline{180K}$$

$$> \underline{\underline{131K}}$$

Q. 4



Consider a System having Memory of size 2^{46} Bytes, uses **Fixed Partitioning**. It is divided into fixed size Partitions each of size 2^{24} Bytes. The OS maintains a (Process Table) with one entry per Process. Each entry has, two fields: First, is a pointer pointing to Partition in which the Process is loaded and Second, Field is Process ID(PID). The Size of PID is 4 Bytes.

Calculate

- (a) The Size of Pointer to the nearest Byte. $3B$
- (b) Size of Process Table in Bytes if the System has 500 Processes.

$$500 \times 7B = 3500B$$

P.T

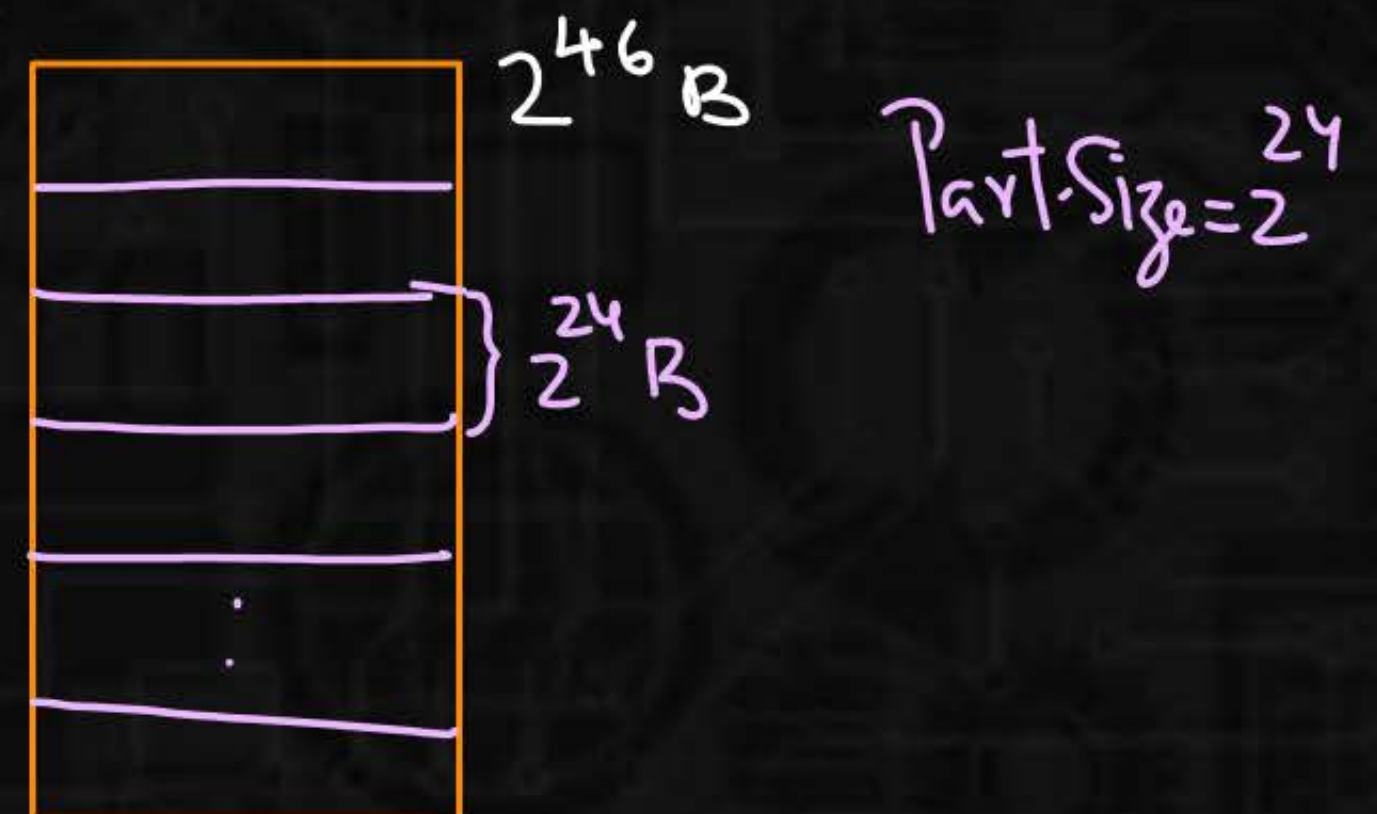
	Ptr (3B)	Pid (4B)
1		
2		
3		
⋮		
5n		

No. of Partitions: $\frac{2^{46}}{2^{24}} = 2^{22}$

(N)

Partition Addr: $22 \text{ bits} = 3B$

(Ptr)



Q. 5

Consider a System Using Variable Partition with no Compaction



Free holes	4K; 8K; 20K; 2K
Program size	2K; 14K; 3K; 6K; 10K; 20K; 2K
Time for Execution	4; 10; 2; 1; 4; 1; 8

Using **Best Fit Allocation Policy** and **FCFS CPU Scheduling Technique**, Find the Time of Loading & Time of Completion of each program. The Burst Times are in Seconds.

Q. 6



Consider allocation of memory to a new process. Assume that none of the existing holes in the memory will exactly fit the process's memory requirement. Hence, a new hole of smaller size will be created if allocation is made in any of the existing holes. Which one of the following statements is TRUE?

- A. The hole created by next fit is never larger than the hole created by best fit
- B. The hole created by worst fit is always larger than the hole created by first fit
- C. The hole created by first fit is always larger than the hole created by next fit
- D. The hole created by best fit is never larger than the hole created by first fit

