

CS & IT ENGINEERING

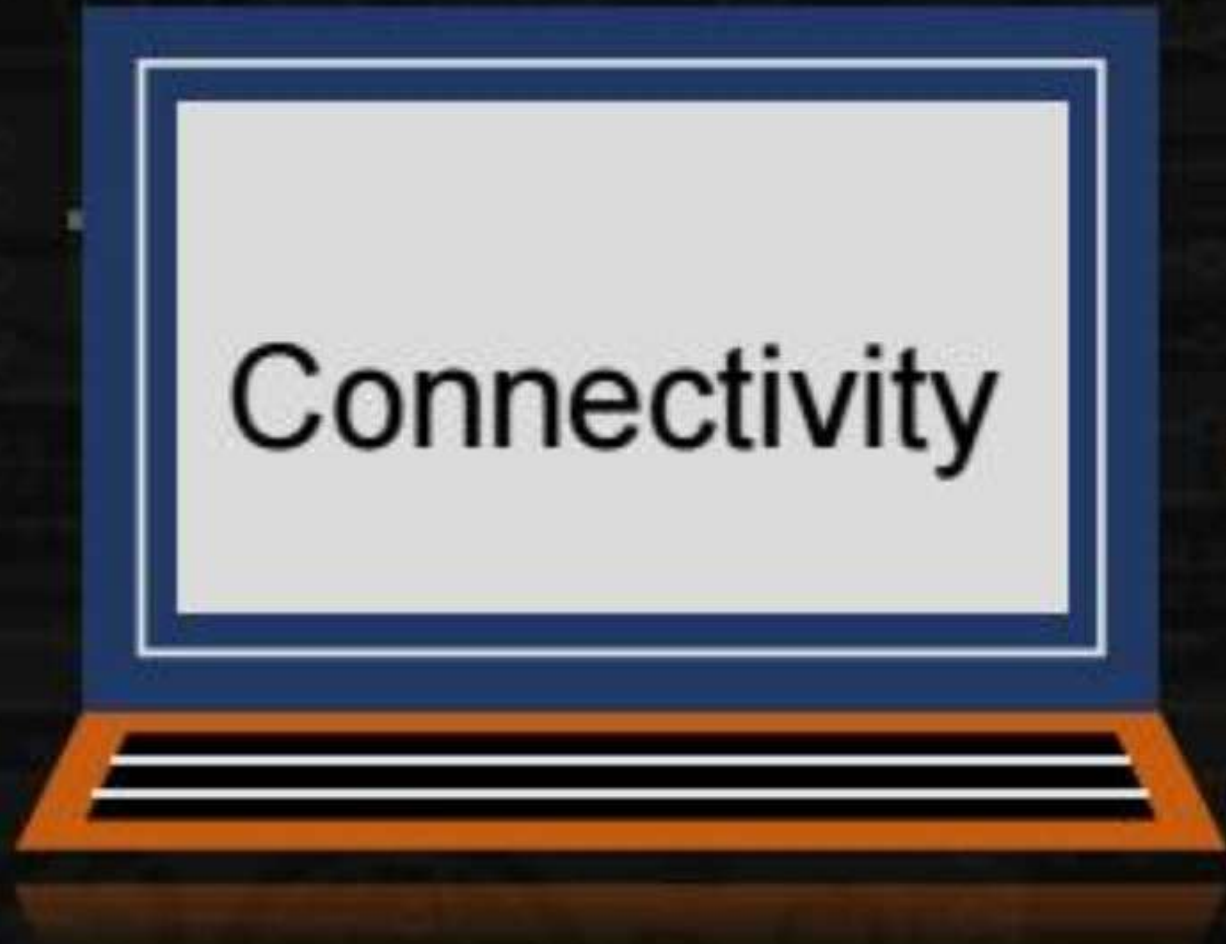
DISCRETE
MATHS
GRAPH THEORY



Lecture No. 5



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01 Types of Graphs -1

02 Types Of graphs -2

03 Practice

$$G + \bar{G} = kn$$

$$e(G) + e(\bar{G}) = \frac{n(n-1)}{2}$$

$$kn \quad n-1 \quad n-1 \quad \dots \quad n-1$$

$$G \rightarrow d_1, d_2, d_3, \dots, d_n$$

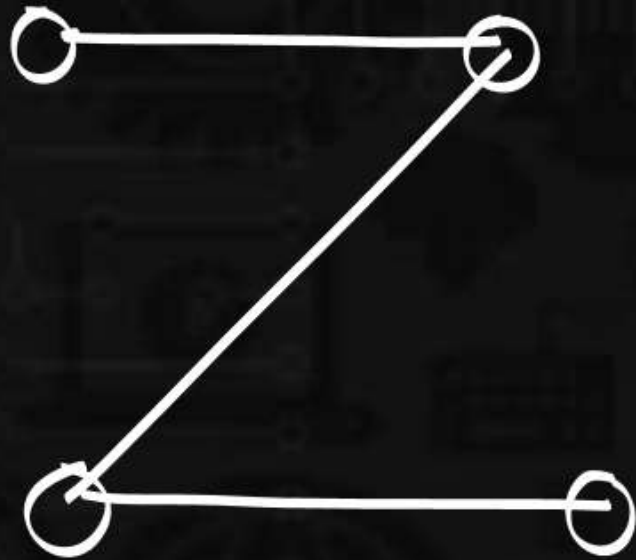


$$\bar{G} \quad n-1-d_1, n-1-d_2, n-1-d_3, \dots, n-1-d_n$$

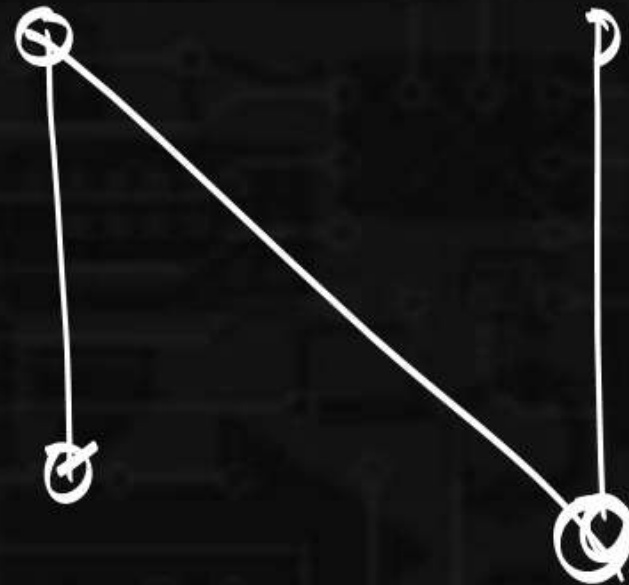
Self-complement Graph. ($G \equiv \bar{G}$)

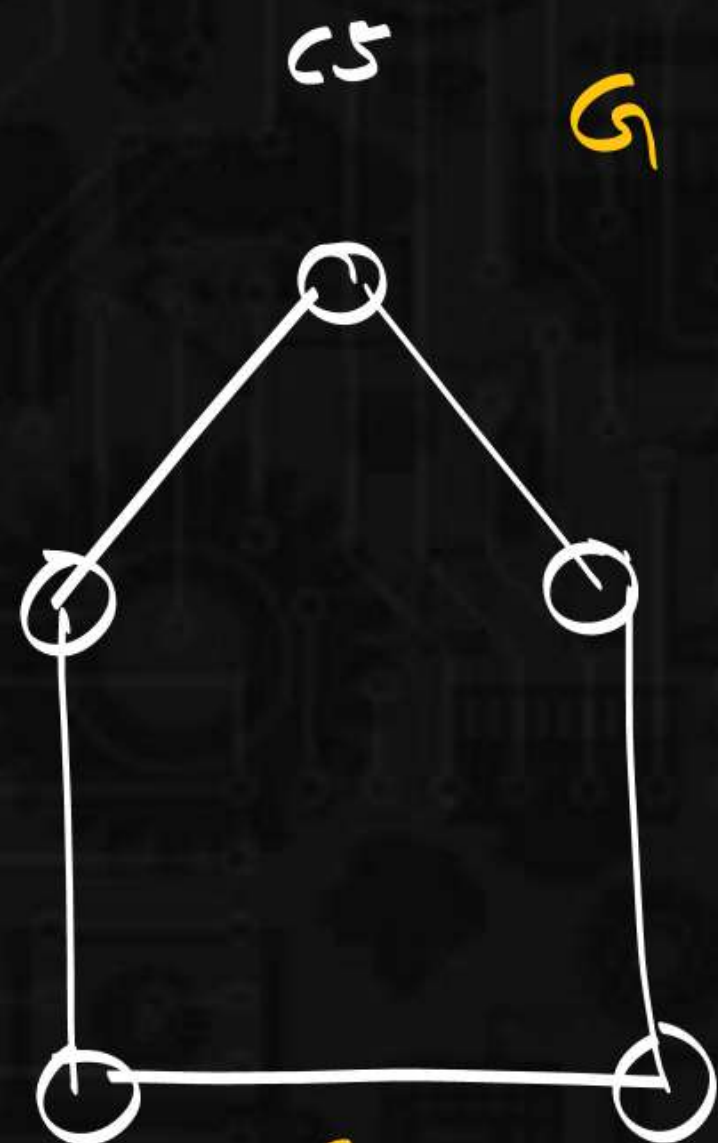
Graph is same as it's own complement.

G



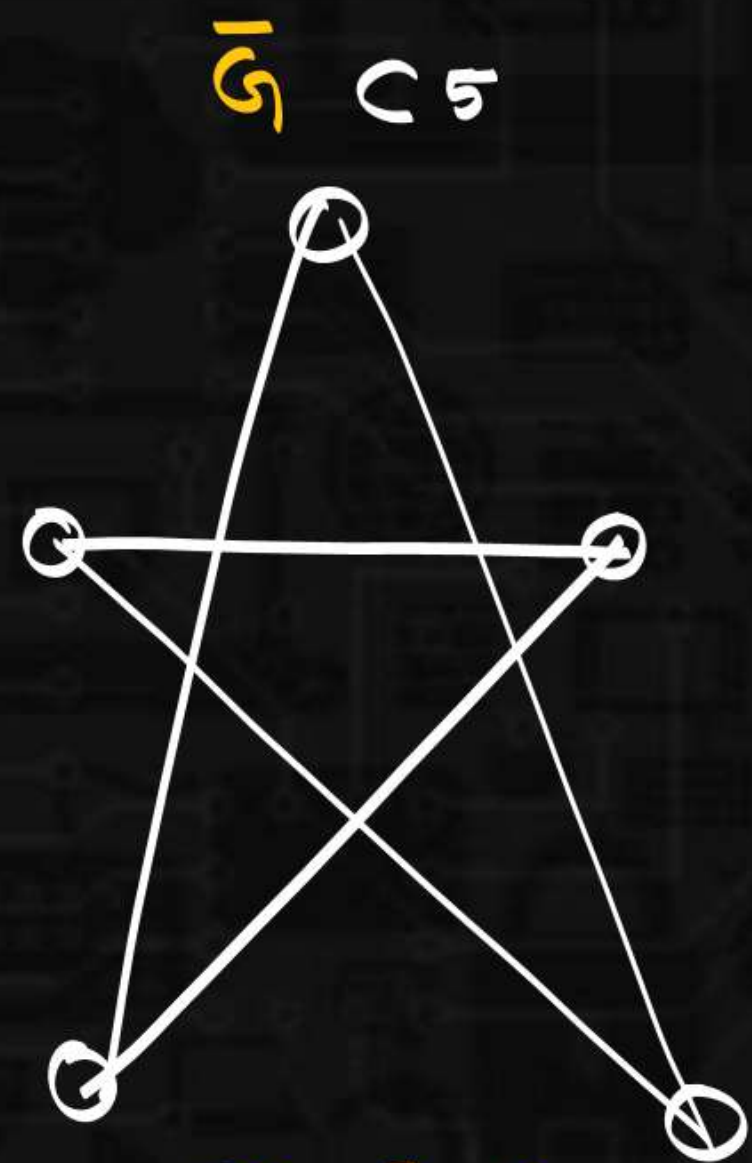
\bar{G}





C_5
 $\begin{cases} n = e = 5 \\ 2, 2, 2, 2, 2 \end{cases}$

|||



C_5
 $\begin{cases} n = 5 = e \\ 2, 2, 2, 2, 2 \end{cases}$

G is same as \bar{G}

$$G + \bar{G} = kn$$

$$\underbrace{e(G)} + \underbrace{e(\bar{G})} = \frac{n(n-1)}{2}$$

$$e + e = \frac{n(n-1)}{2}$$

$$2e = \frac{n(n-1)}{2}$$

$$e = \frac{n(n-1)}{4}$$

$$e = \frac{n(n-1)}{4}$$

$$\frac{n}{4} \text{ or } \frac{n-1}{4}$$

✓ $n=4$ $e = \frac{n(n-1)}{4} = \frac{4 \cdot 3}{4} = 3$ ✓

✓ $n=5$ $e = \frac{n(n-1)}{4} = \frac{5 \cdot 4}{4} = 5$ ✓

✗ $n=6$ $e = \frac{6 \cdot 5}{4} = \frac{15}{2} = 7.5$ ✗

✗ $n=7$ $e = \frac{7 \cdot 6}{4} = \frac{21}{2} = 10.5$ ✗

✓ $n=8$

✓ $n=9$

* $\left\{ \begin{array}{l} a \equiv b \pmod{n} \\ a \& b \end{array} \right. \text{w.r.t } n$

$\rightarrow \left\{ \begin{array}{l} 1 \& 5 \text{ are having same remainder w.r.t } 4. \\ 1 \equiv 5 \pmod{4} \end{array} \right.$

$\rightarrow \left\{ \begin{array}{l} a \equiv b \pmod{n} \\ \frac{a-b}{n} \in \mathbb{Z} \end{array} \right.$

* $\frac{a-b}{n} \in \mathbb{Z}$

$$\left\{ \frac{n}{4} \text{ or } \frac{n-1}{4} \right.$$

$$\frac{n-0}{4} \text{ or } \frac{n-1}{4}$$

$$n \equiv 0 \pmod{4} \text{ or } n \equiv 1 \pmod{4}$$

$$n \equiv 0 \text{ or } 1 \pmod{4}$$

Cycle Graph.

C_3 .

C_4 .

C_5 .

C_6 . X

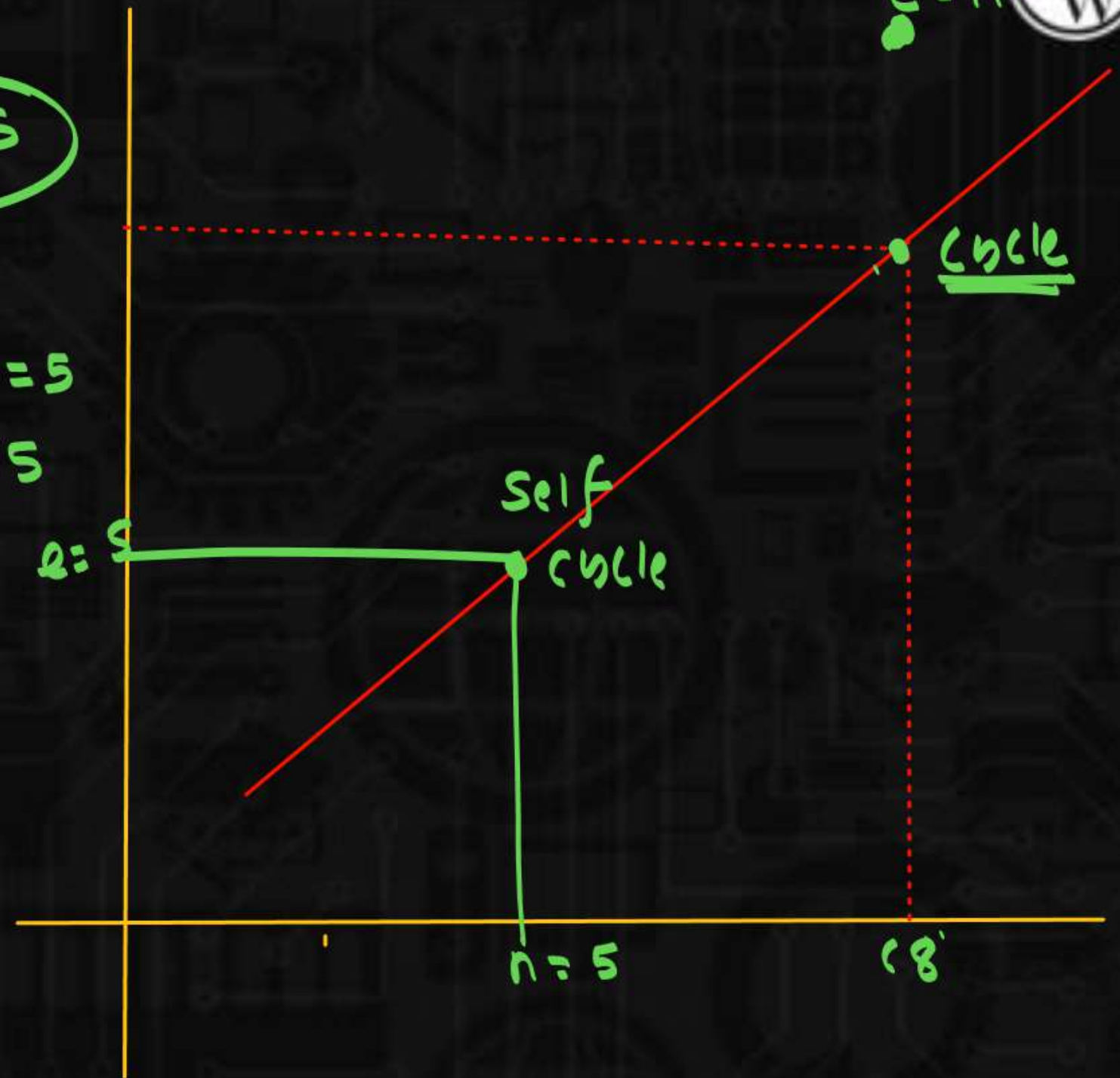
C_7 . X

C_8 . cycle graph $e=8$.

C_9 Self complement 14.

C_n
 $n=5$

Self-complement $e=5$
Cycle graph $e=5$



for self complement
Graph in C_n .

$$n=9$$

$$n=5$$

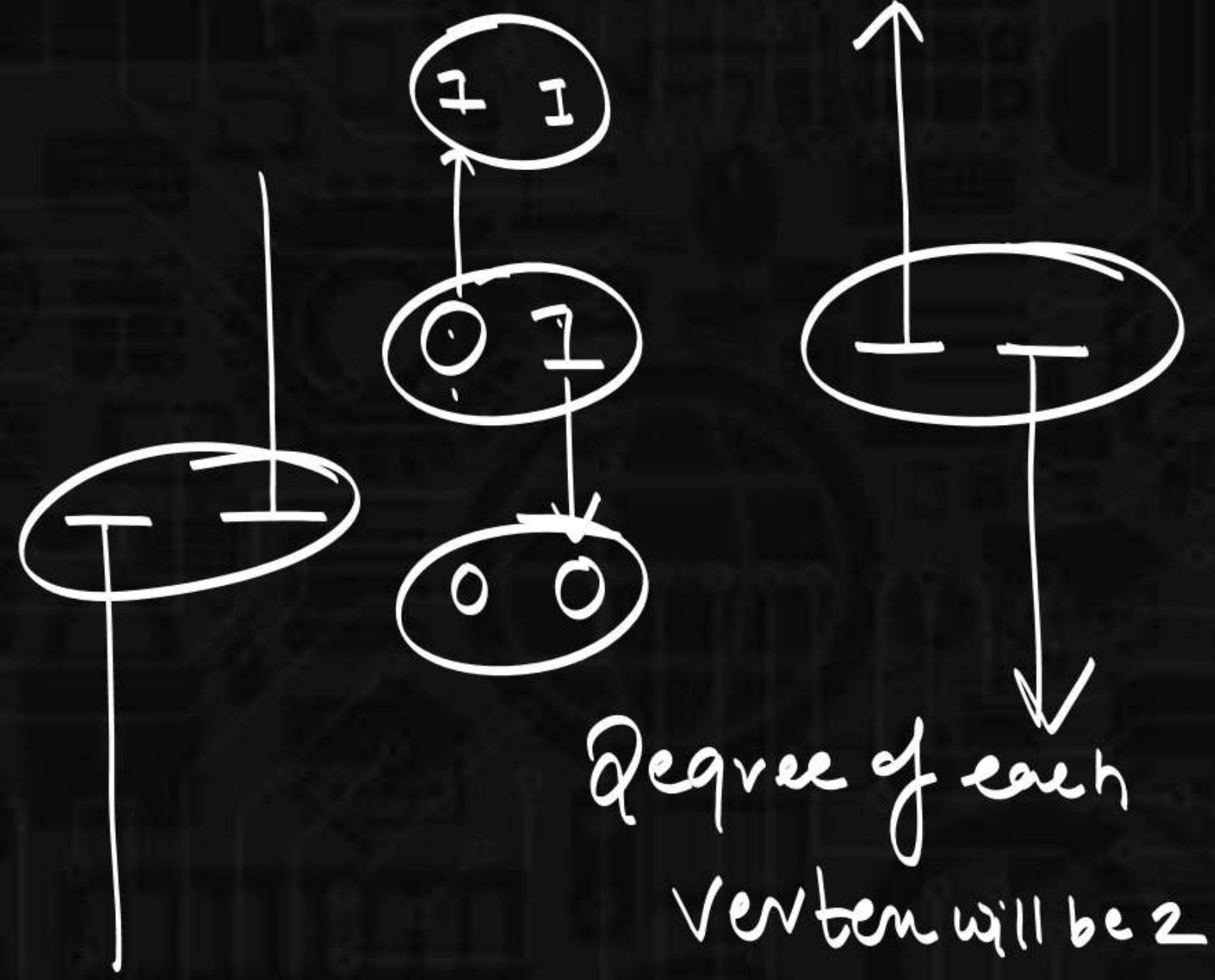
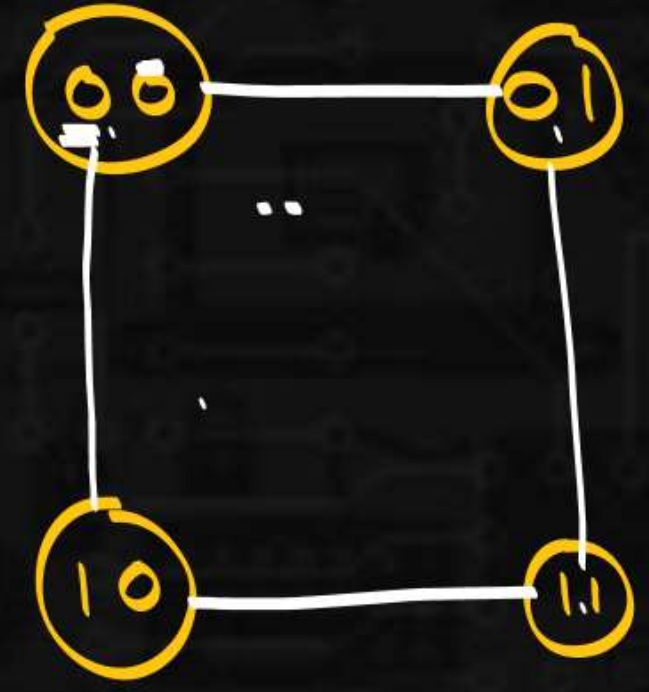
Graph vertices are represented as n-bit signal, and 2 vertices are connected if there bit position changes by 1 bit.
What will be total edges?



$n = 2 \text{ bit}$

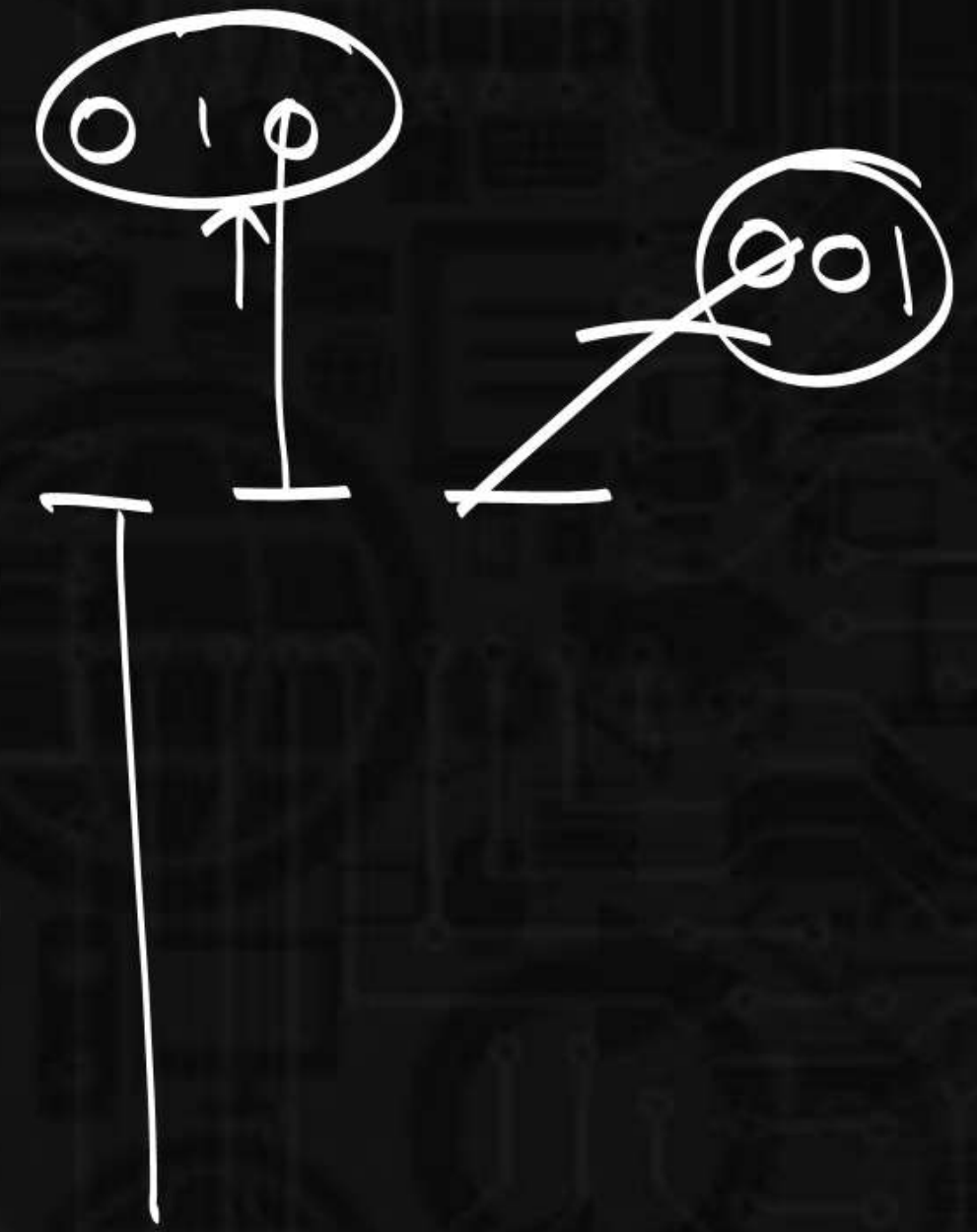
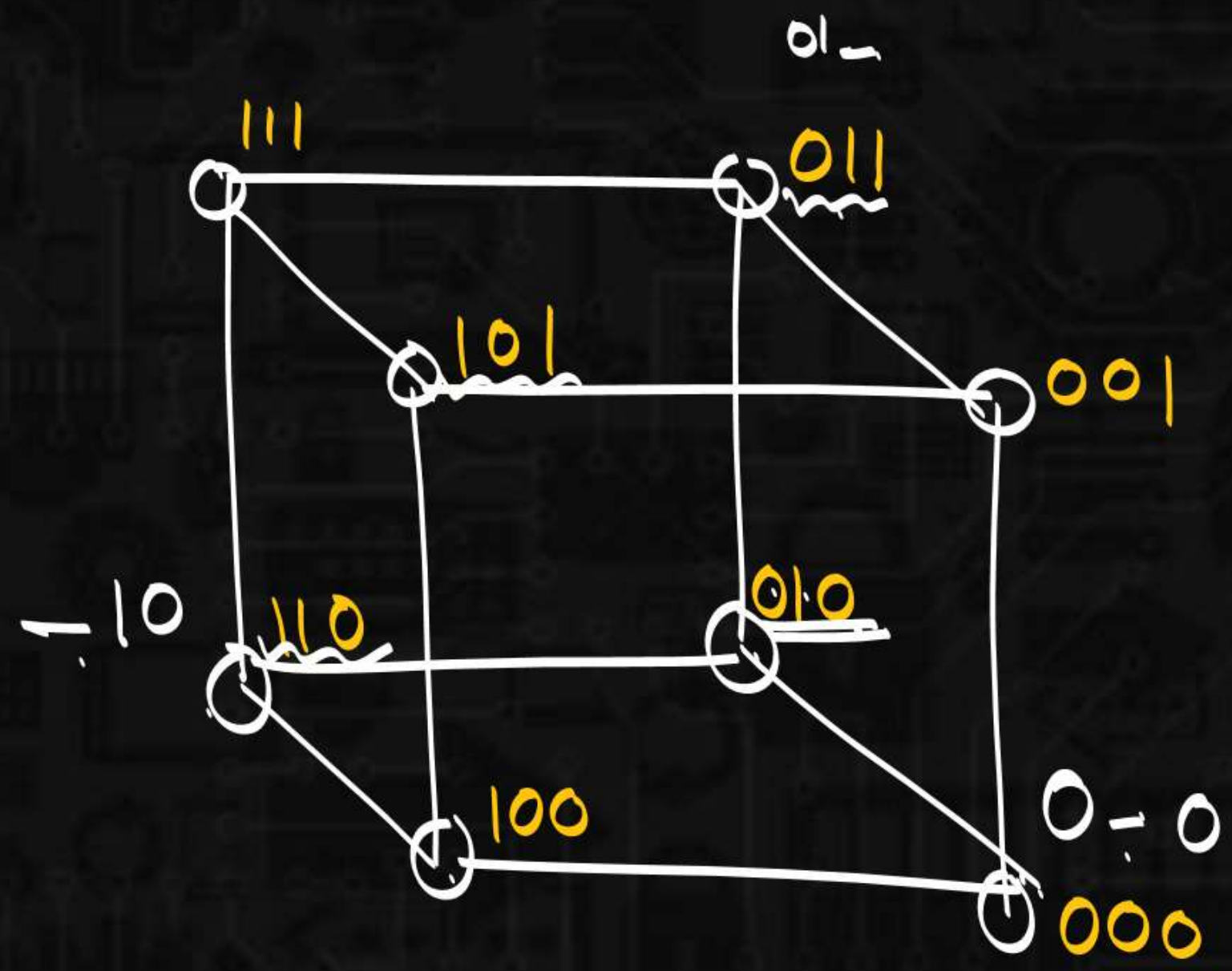
Total vertices
 $= 2^2$

- 00
- 01
- 10
- 11



Degree of each vertex will be 3.

n = 3 bit
Total vertices = 2^3
000
001
111



$n \rightarrow$ bit signal

$$\text{Total vertices } (n) = 2^n$$

$$\text{Degree of each vertex} = n$$

$$\sum d(v_i) = 2e$$

$$2^n \cdot n = 2e$$

$$e = n \cdot 2^{n-1}$$

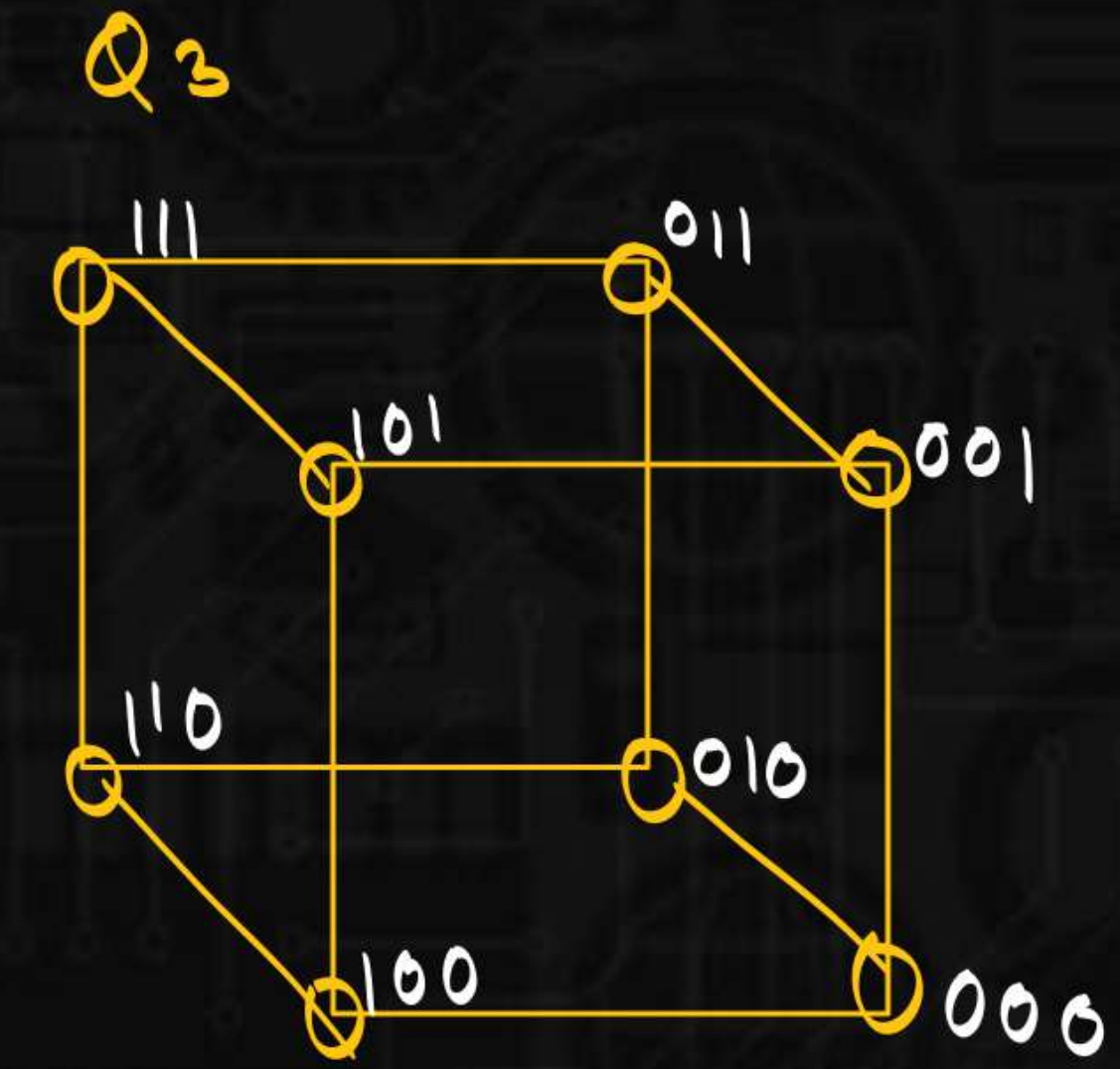
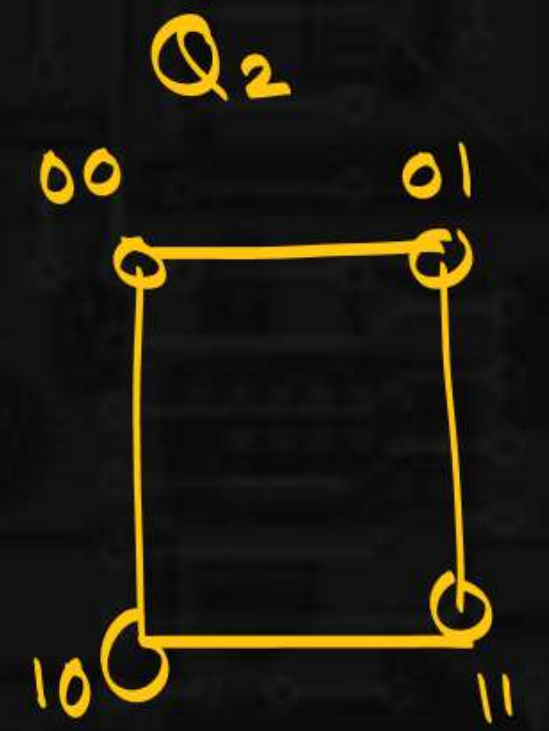
3 bit signal

$$\text{Total vertices} = 2^3$$

$$\text{Degree of each vertex} = 3$$

Hypercube. (Q_n) ($n \geq 1$)
 \downarrow
 n -bit signal.

$$e = n \cdot 2^{n-1}$$



$$e(G) = n \cdot 2^{n-1}$$

$$\text{Total vertices} = x = 2^n$$

$$e(G) + e(\bar{G}) = kx$$

$$n \cdot 2^{n-1} + e(\bar{G}) = \frac{x(x-1)}{2}$$

$$n \cdot 2^{n-1} + e(\bar{G}) = \frac{2^n(2^n-1)}{2}$$

$$e(\bar{G}) = \frac{2^n(2^n-1)}{2} - n \cdot 2^{n-1}$$

$$N = \text{Total vertices} = 2^n.$$

$$K_N \quad N-1 \quad N-1 \quad N-1.$$

$$G \rightarrow n, n, n, n, n, \dots, n.$$



$$\bar{G} \quad N-1-n, N-1-n, N-1-n$$

$$2^n-1-n, 2^n-1-n, 2^n-1-n, \dots$$

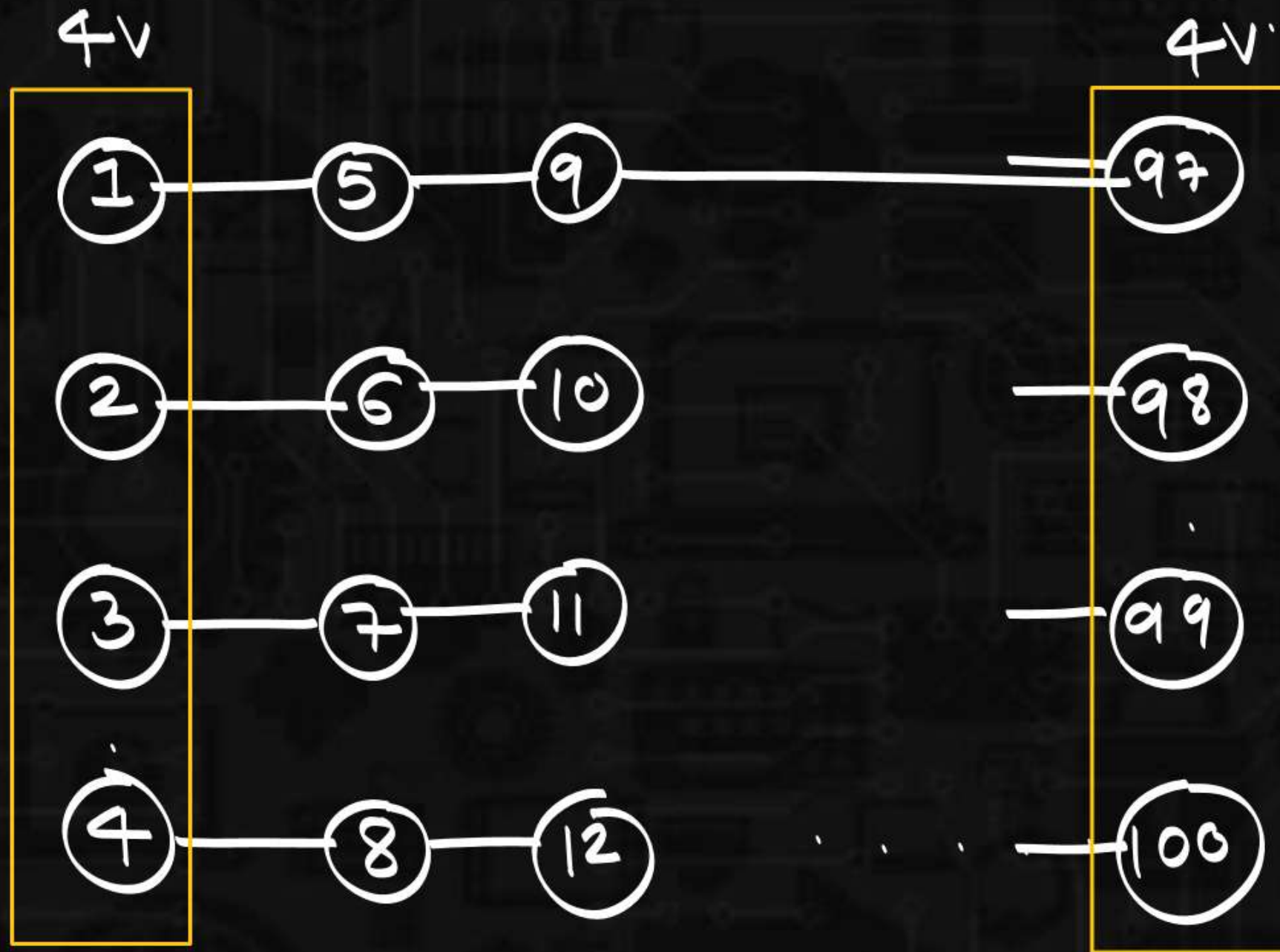
Consider a graph vertices are represented as a no. from
 $\{1 \dots 100\}$

two vertices are connected $|a-b|=4$ eg $|1-5|=4$ OR there
 what will be total edges in this ?

difference is 4.

$$\text{eg } |1-5|=4$$

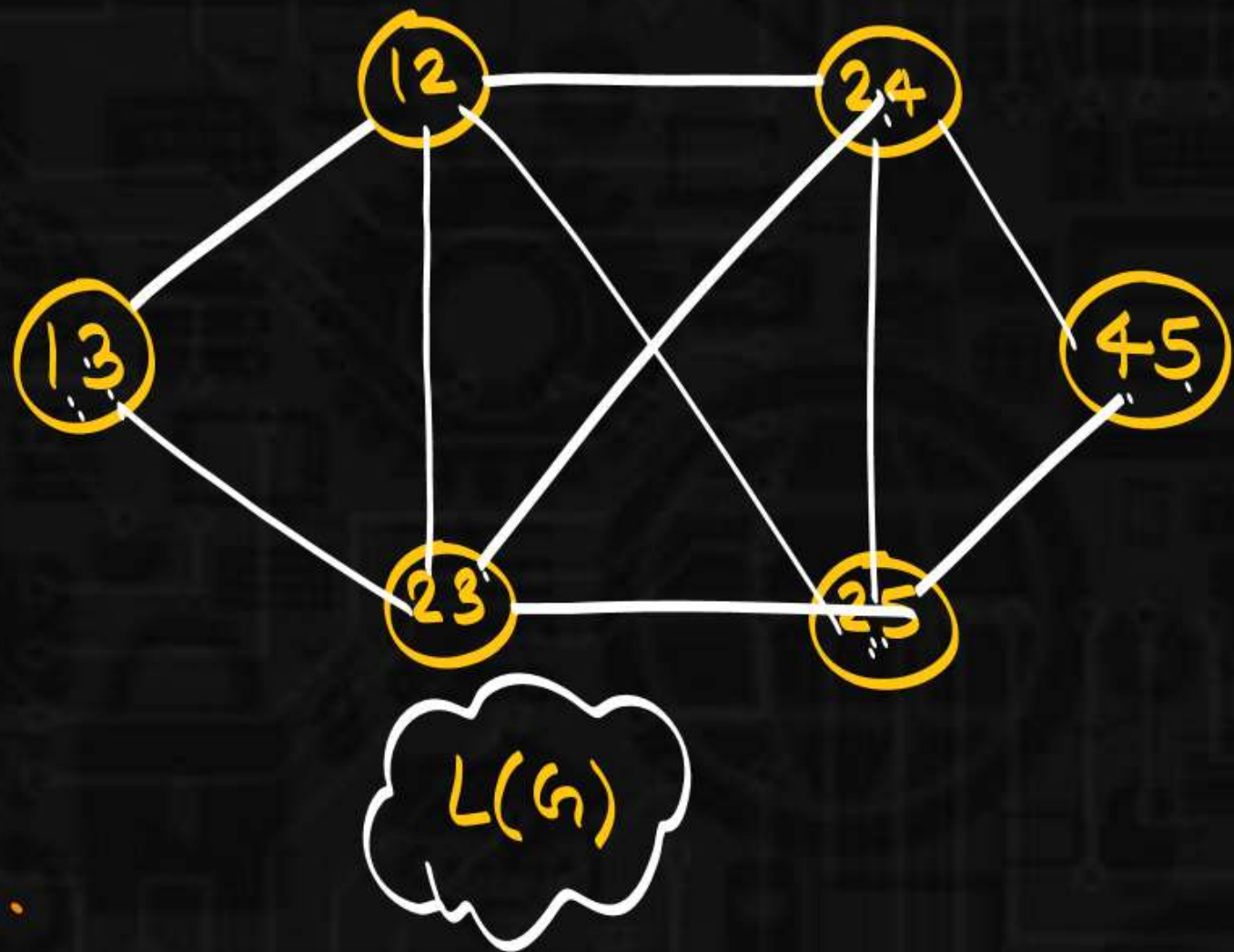
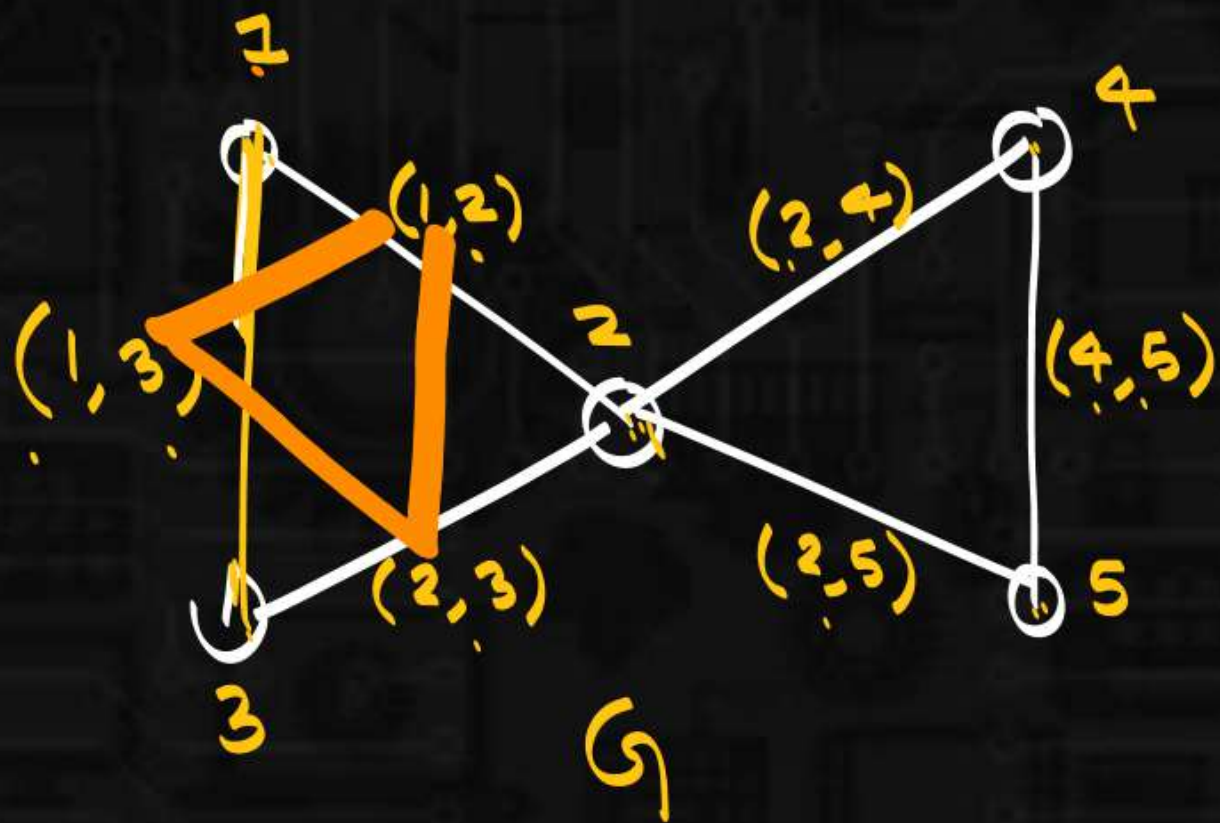
$$|5-9|=4.$$



$$4 \times 1 + 4 \times 1 + 92 \times 2 = 2e$$

$$e = 96.$$

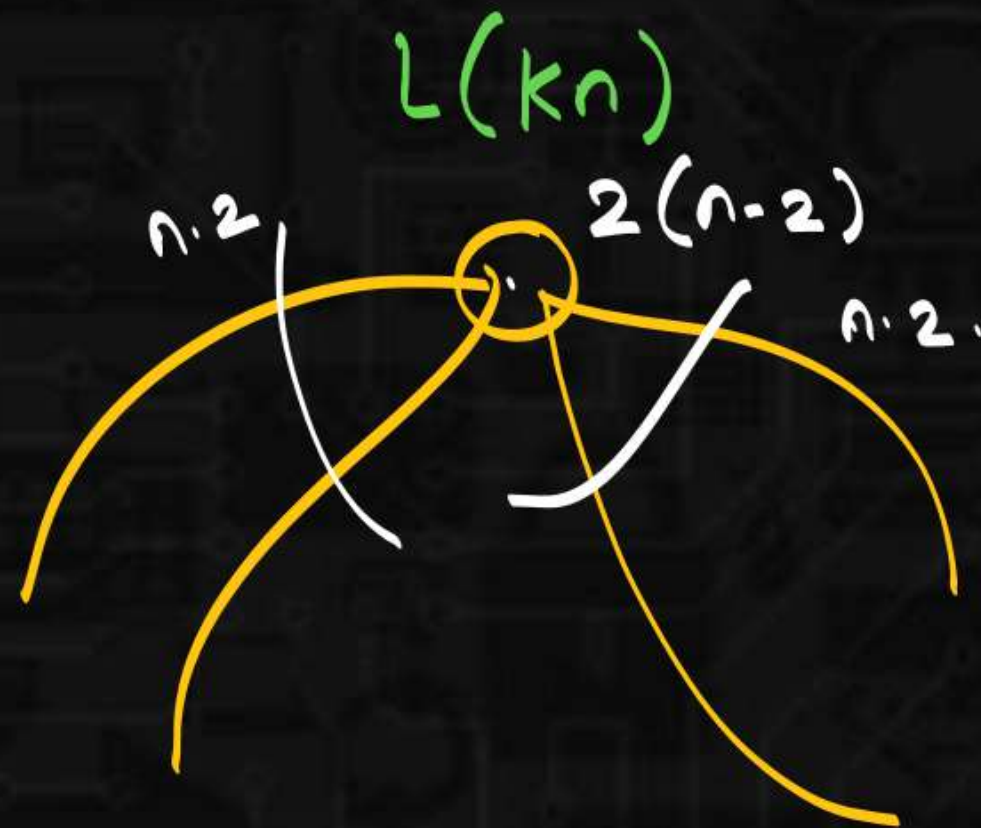
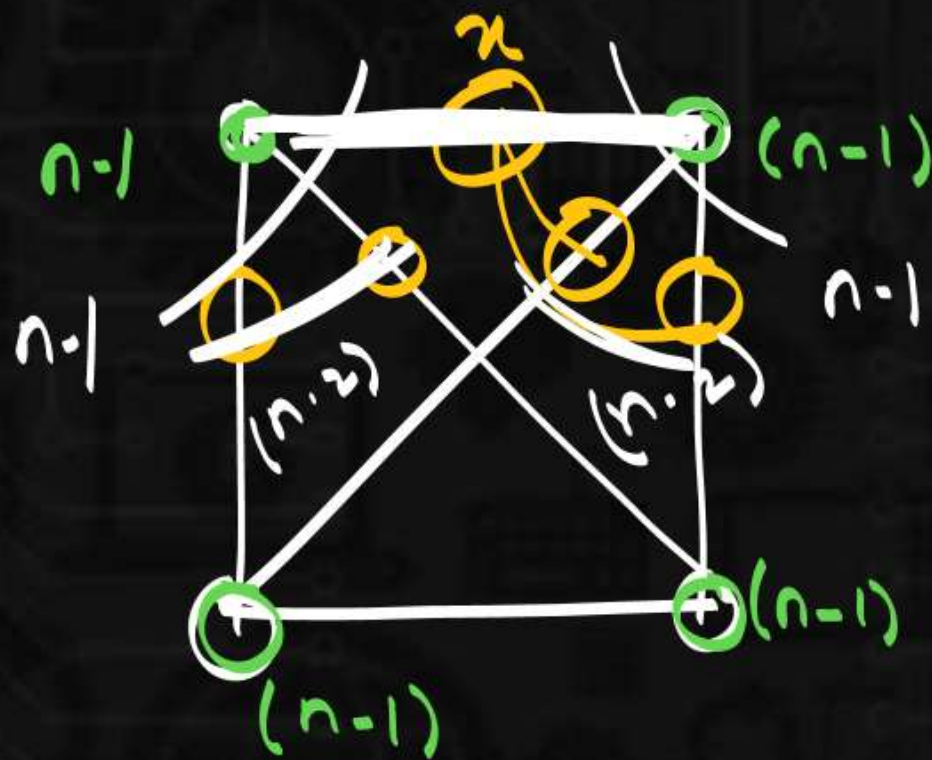
Line Graph. $L(G)$

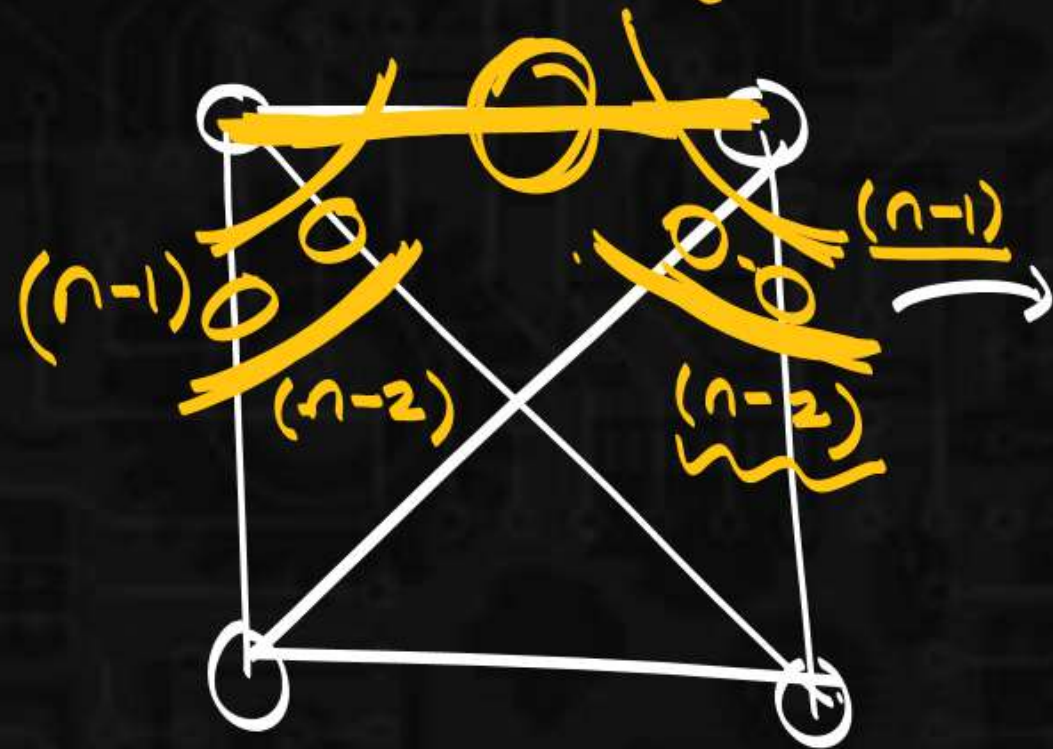


Line graph
↓
edge graph

What will be degree sequence of $L(K_n)$?

K_n/K_4





Degree of each vertex in $L(kn)$ is $2(n-2)$.

