

CS & IT ENGINEERING

Data Structure



Tree
Chapter- 5
Lec- 04



By- Pankaj Sharma sir

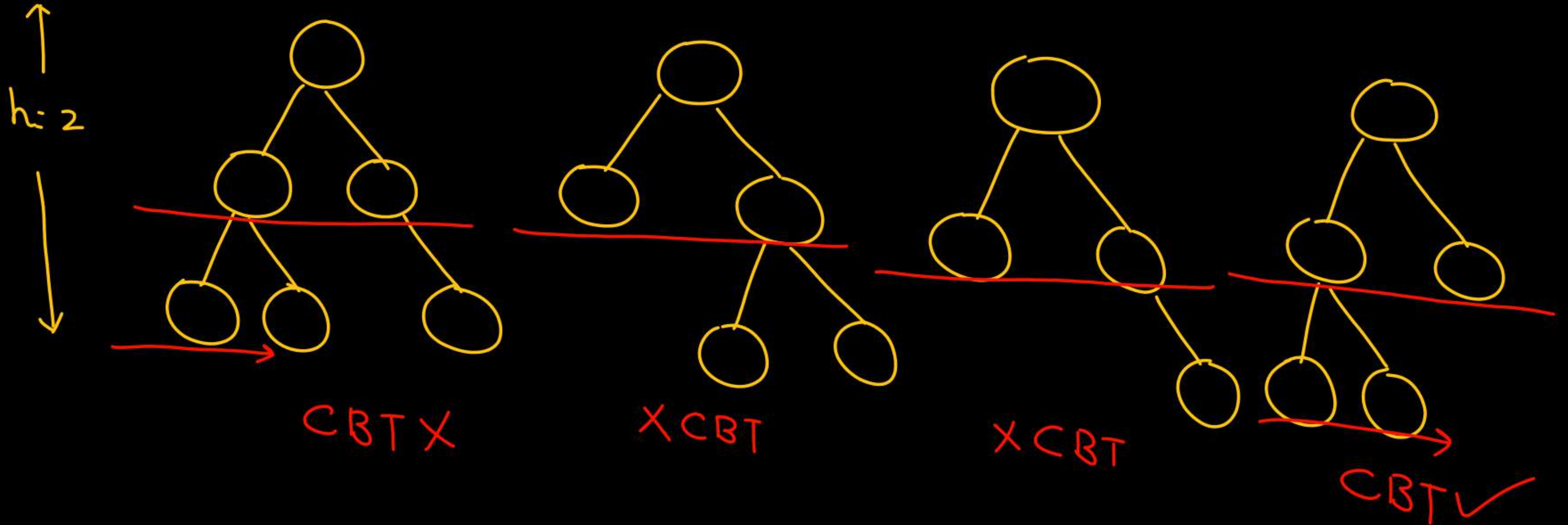
TOPICS TO BE
COVERED

Tree-IV

Complete binary tree

CBT is a binary tree which is Full till second last level.

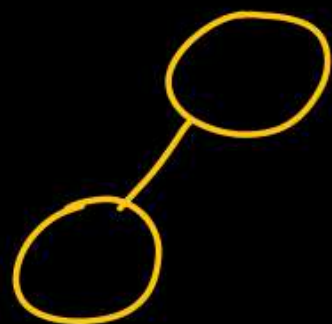
& nodes at last level are filled from left to right.



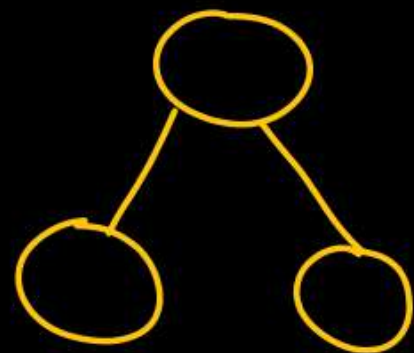
1. Structure of a CBT with 1 node



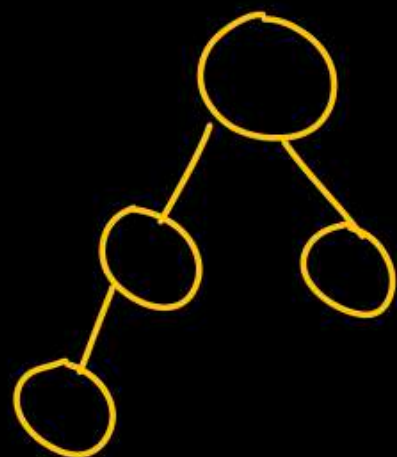
2. Structure of a CBT with 2 node



3. Structure of a CBT with 3 node



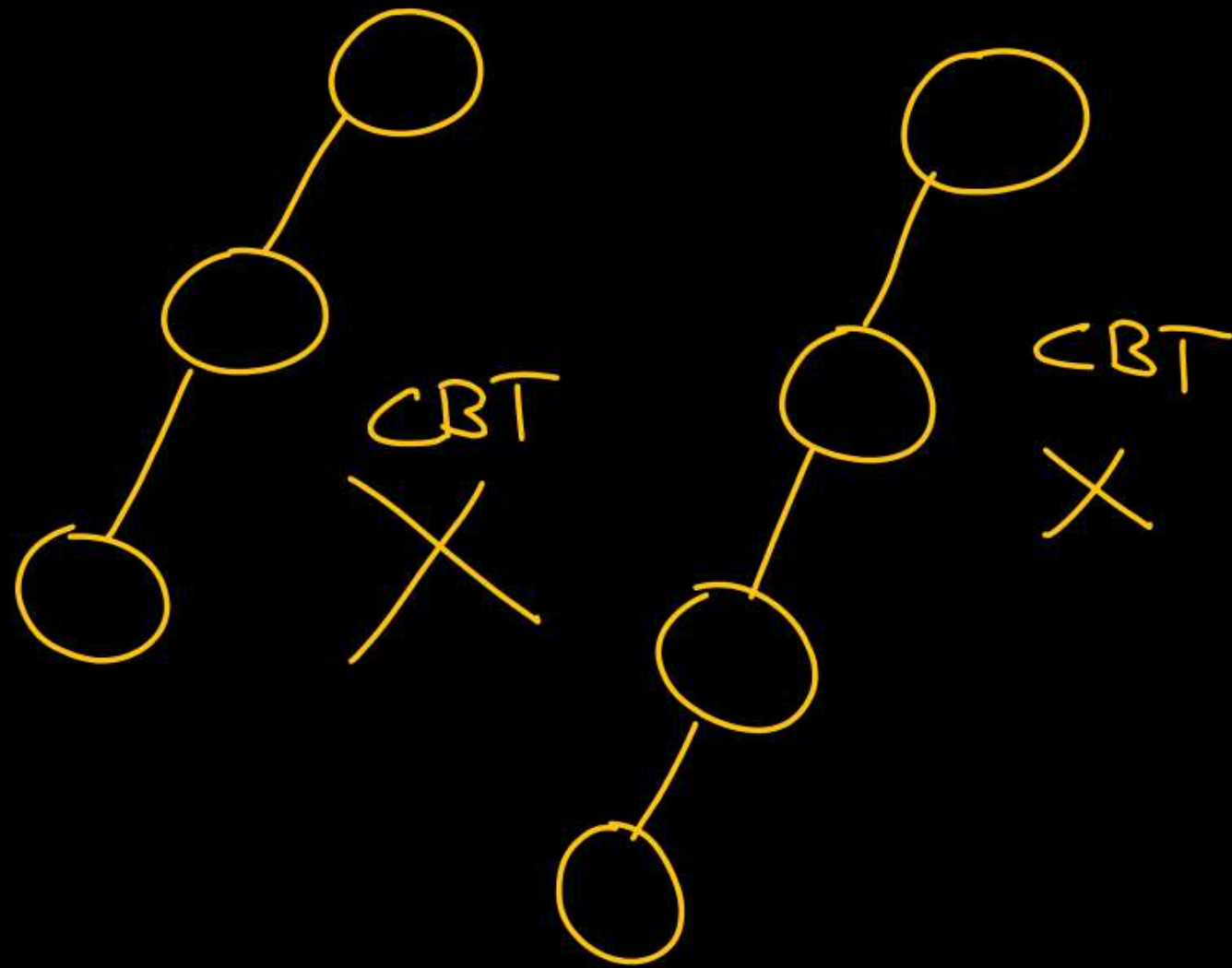
4. Structure of a CBT with 4 node



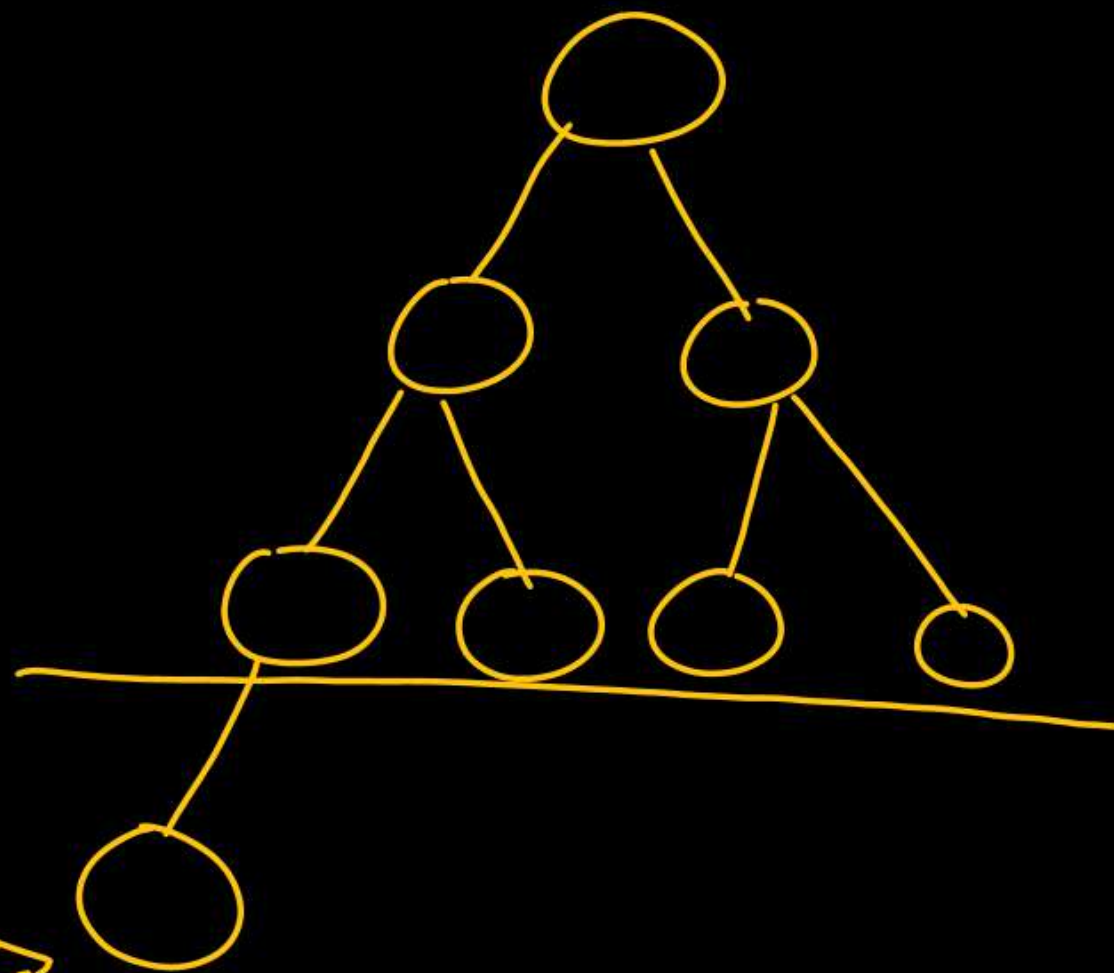
** The structure of a CBT with K nodes is always fix.

1.) Max. no. of nodes in a CBT of height $h = 2^{h+1} - 1$

Min no. of nodes in a CBT of height $h = ?$



$h=3$



#Nodes
 2^0

level
0

2^1

1

2^2

2

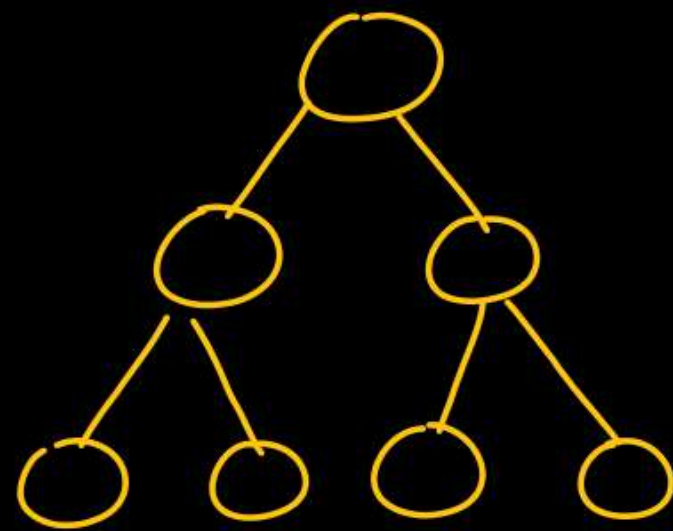
\vdots

2^{h-1}

$h-1$

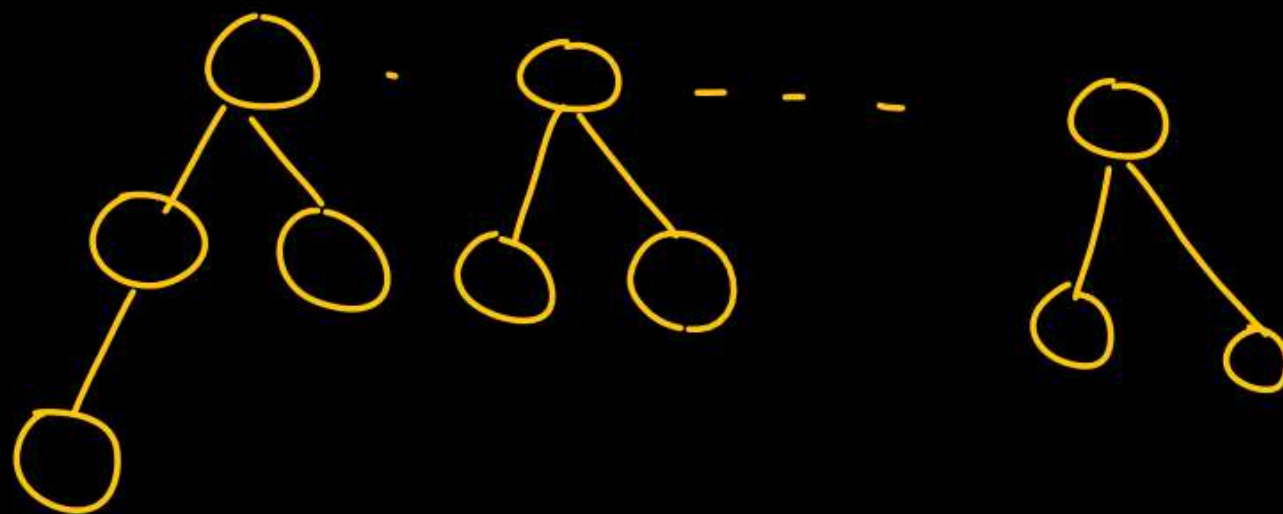
1

h



$$S = \frac{a(r^n - 1)}{r - 1}$$

$$\begin{bmatrix} a = 1 \\ r = 2 \\ n = h \end{bmatrix}$$



$$n_{\min} = (2^0 + 2^1 + 2^2 + \dots + 2^{h-1}) + 1 = \left(\frac{2^h - 1}{2 - 1} \right) + 1 = 2^h - 1 + 1 = 2^h$$

$$2^h \leq n \leq 2^{h+1} - 1$$

$$x \leq 4 \xleftarrow{\text{max}}$$

$$x \geq 2 \xleftarrow{\text{min}}$$

$$2 \leq x \leq 4$$

$$2^h \leq n$$

$$\log_2 2^h \leq \log_2 n$$

$$h \log_2 2 \leq \log_2 n$$

$$h \leq \frac{\log_2 n}{\log_2 2}$$

$$h \leq \log_2 n$$

$$\log_2(n+1) - 1 \leq h \leq \log_2(n)$$

$$n \leq 2^{h+1} - 1$$

$$(n+1) \leq 2^{h+1}$$

$$\log_2(n+1) \leq (h+1) \log_2 2$$

$$\frac{\log_2(n+1)}{\log_2 2} \leq h+1$$

$$(h+1) \geq \log_2(n+1)$$

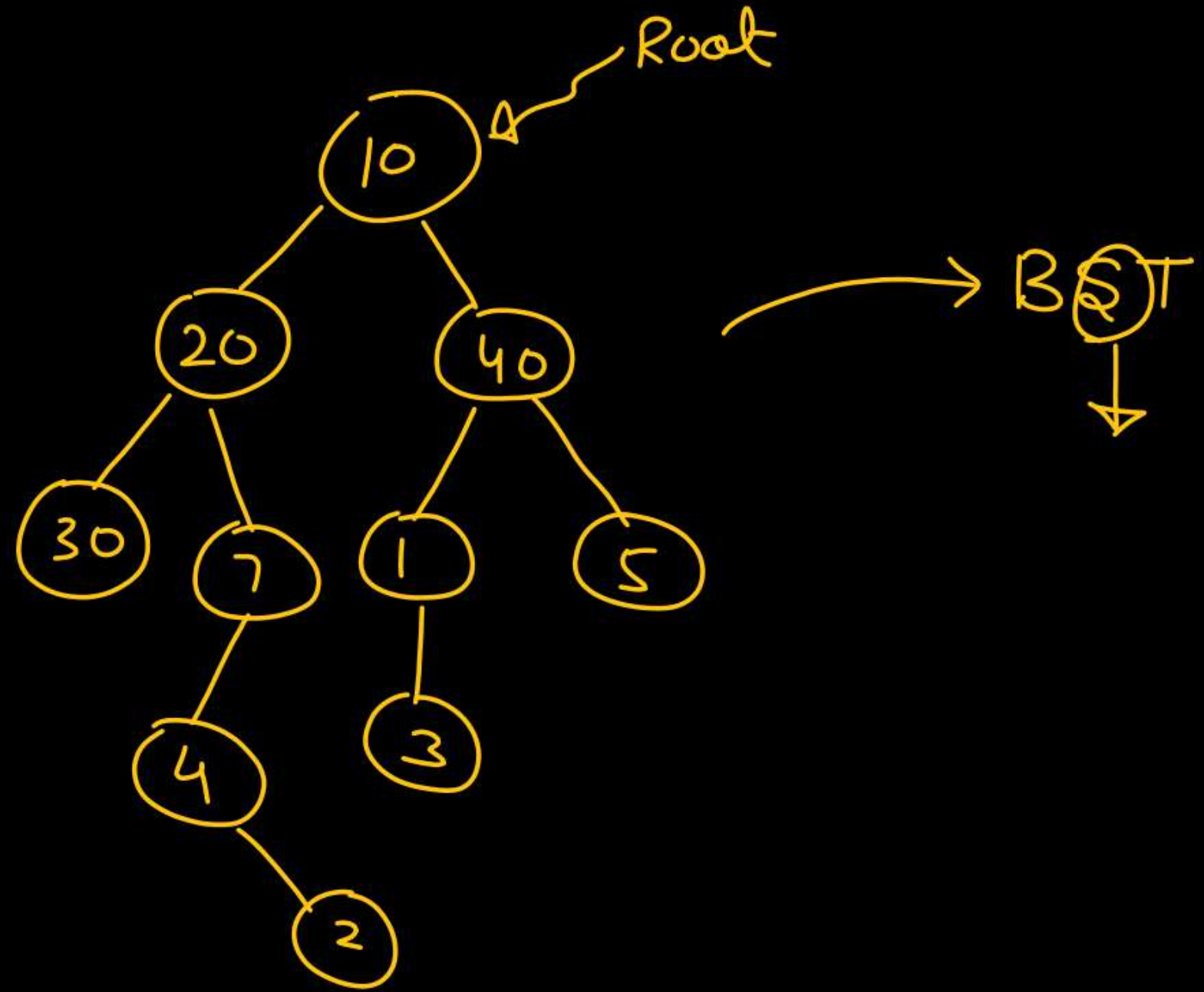
$$h \geq \log_2(n+1) - 1$$

Binary Search Tree

Why?

Given, Binary tree and a key and we need to find whether the key is present in the tree or not?

key=100
 $O(n)$

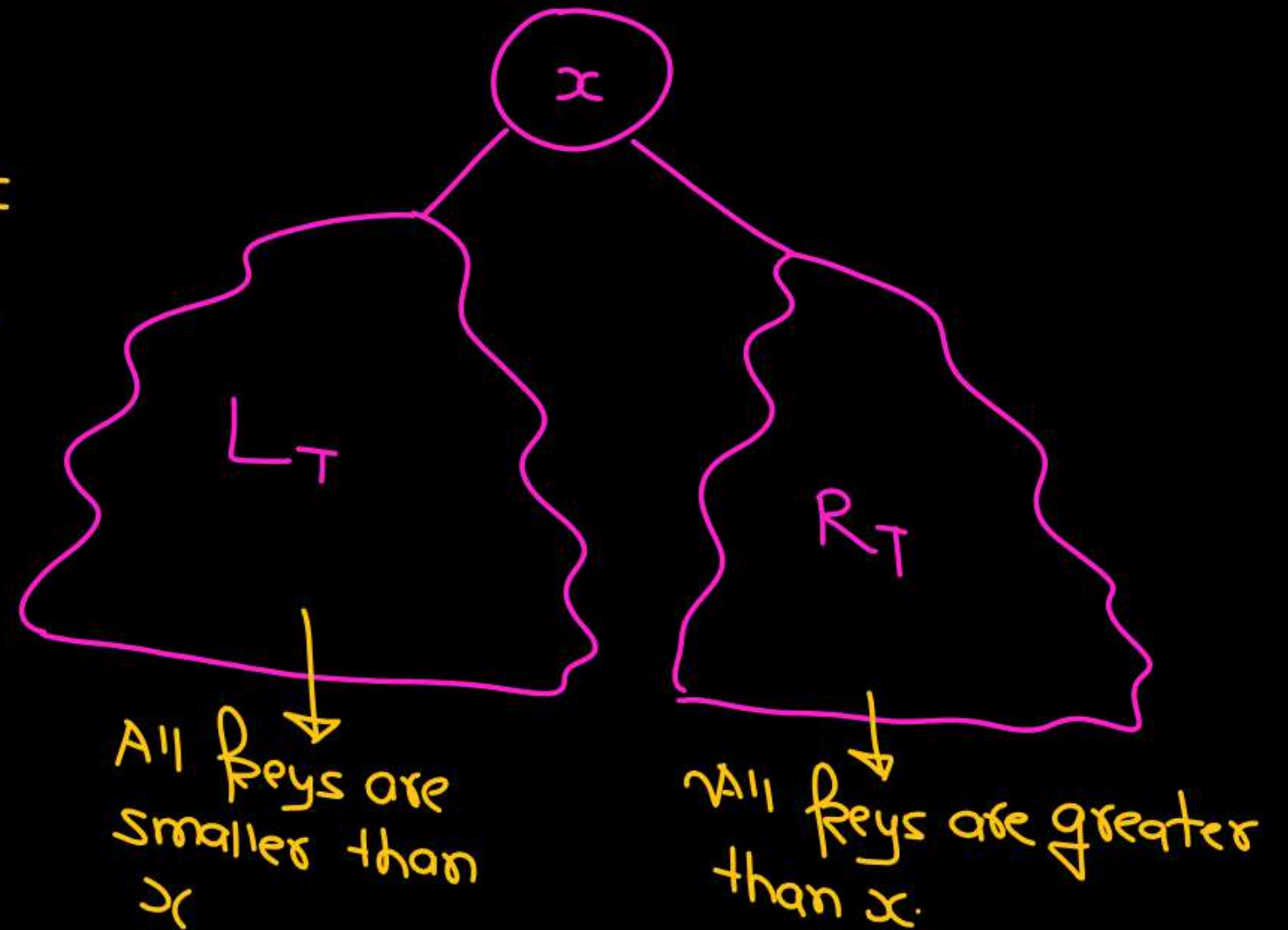


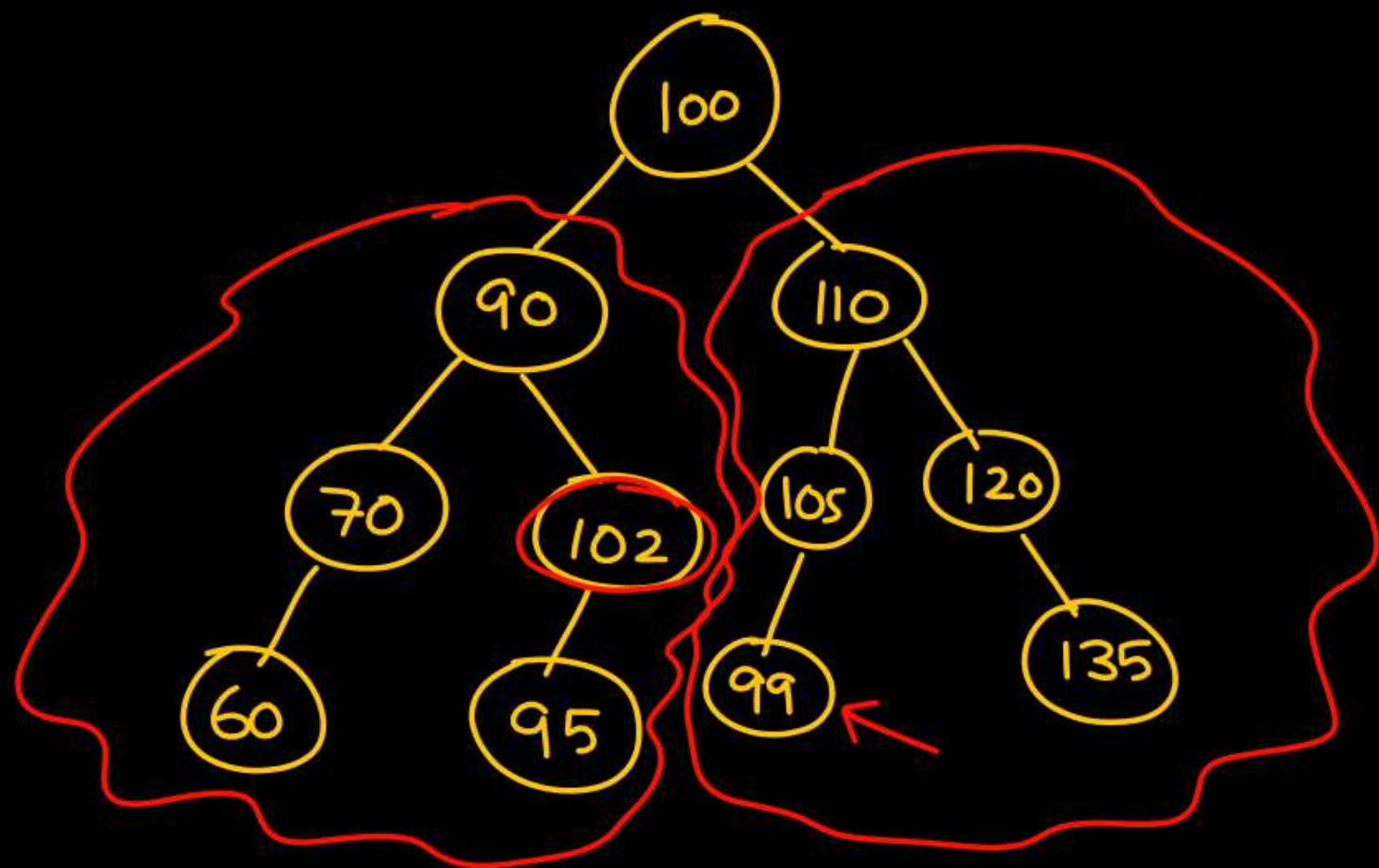
Binary Search Tree

A binary tree in which every node satisfies following property:

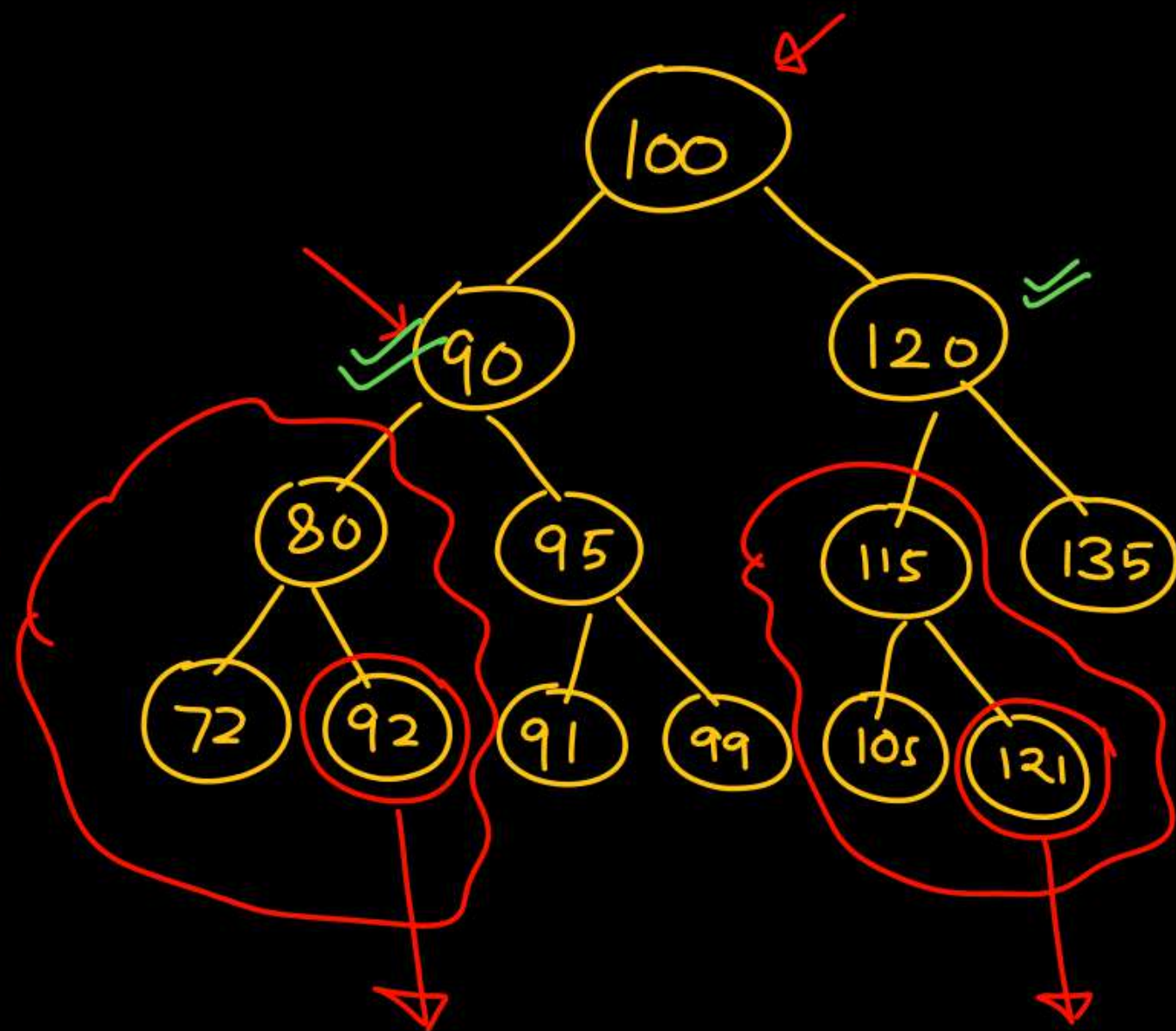
All the keys in the left subtree of a node are smaller than the node value.

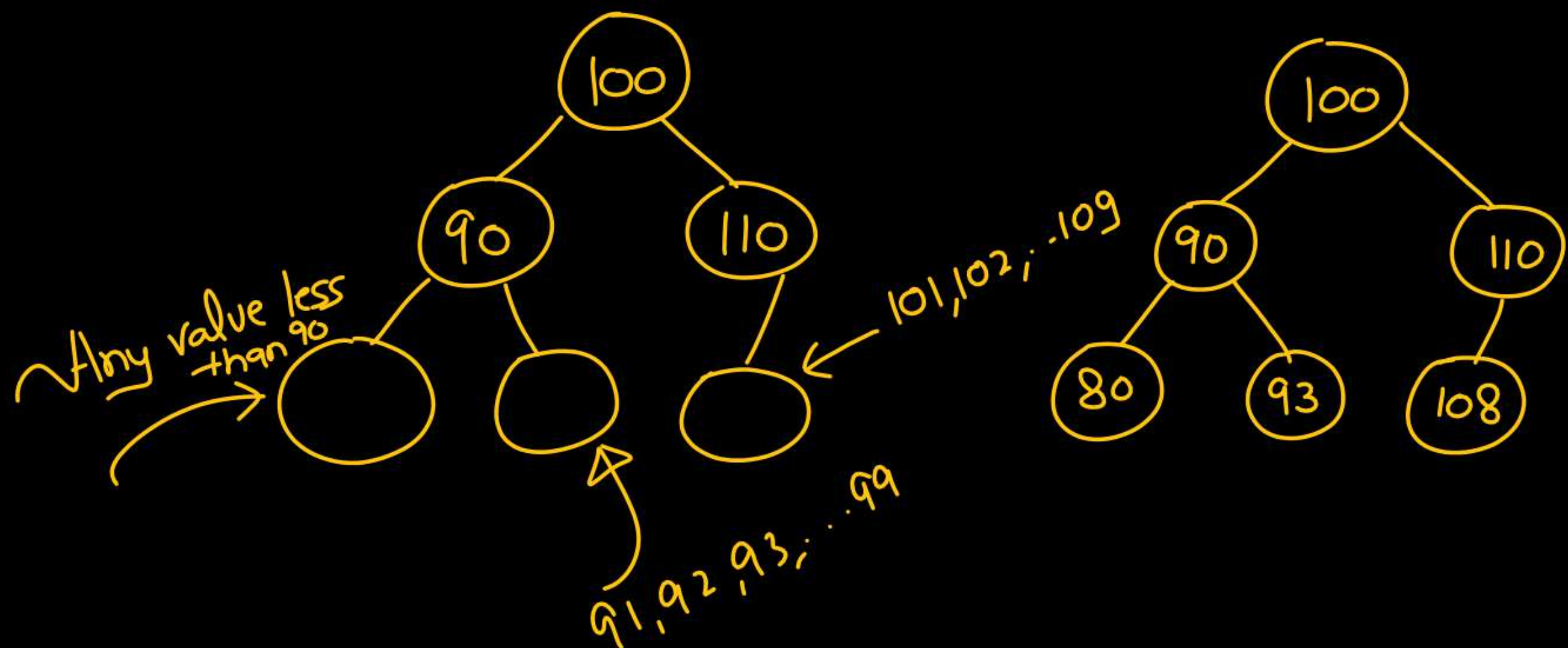
All the keys in the right subtree of a node are greater than the node value.





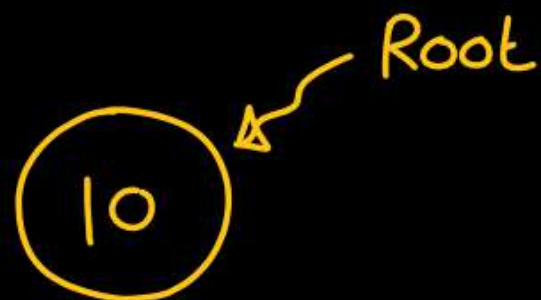
Is it a BST?





1. Construct a binary search tree by inserting keys 10, 20, 30 in order \rightarrow

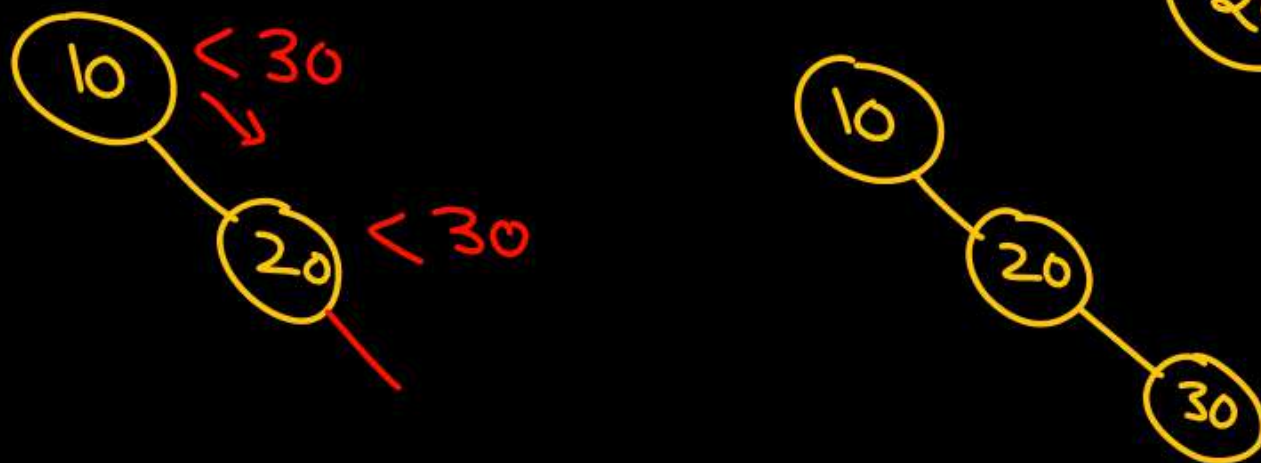
a]



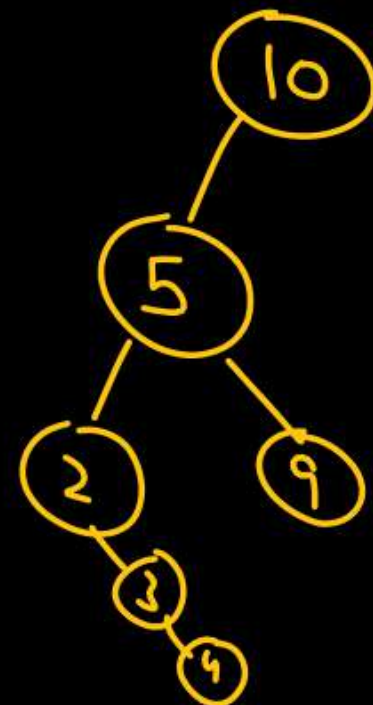
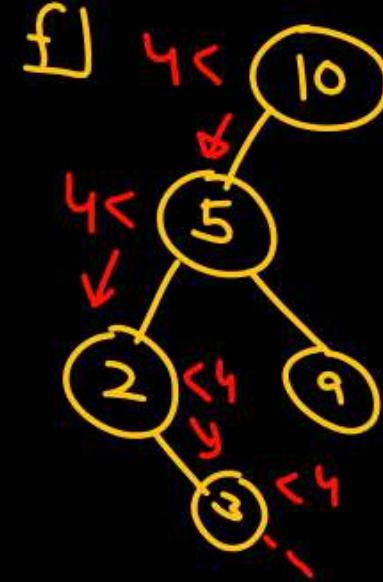
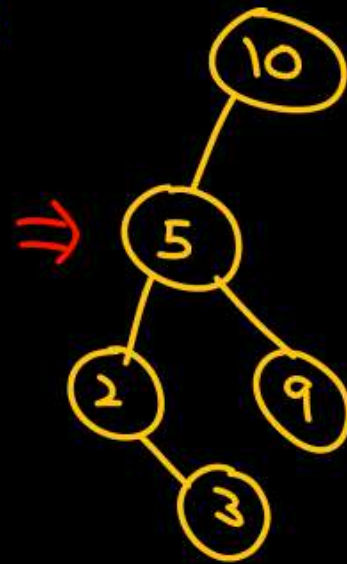
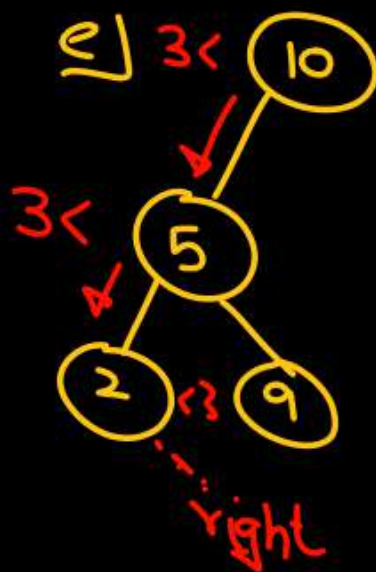
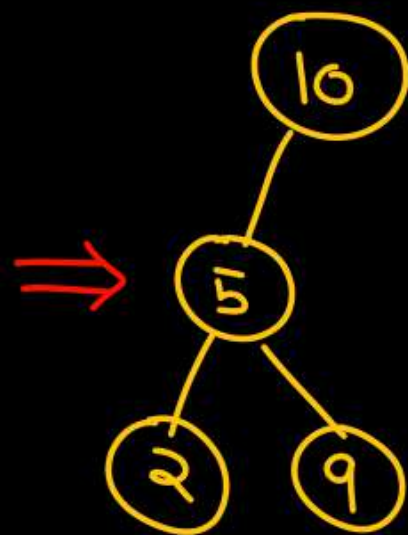
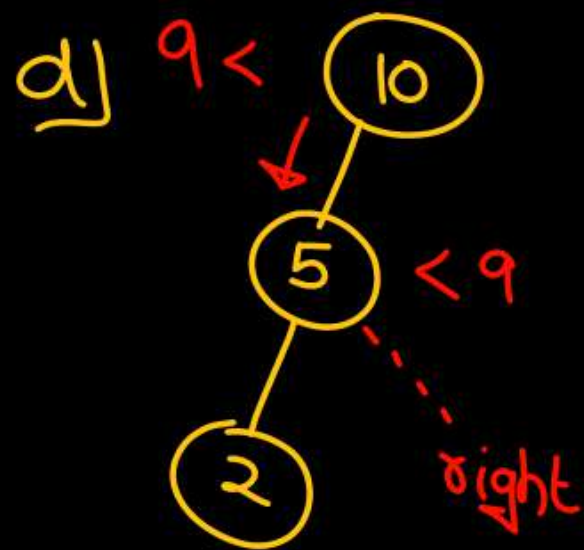
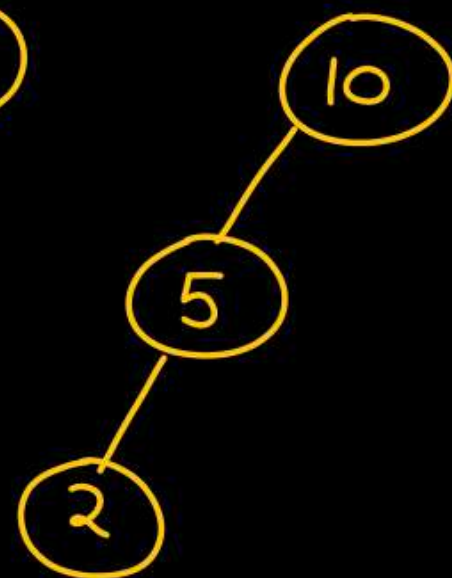
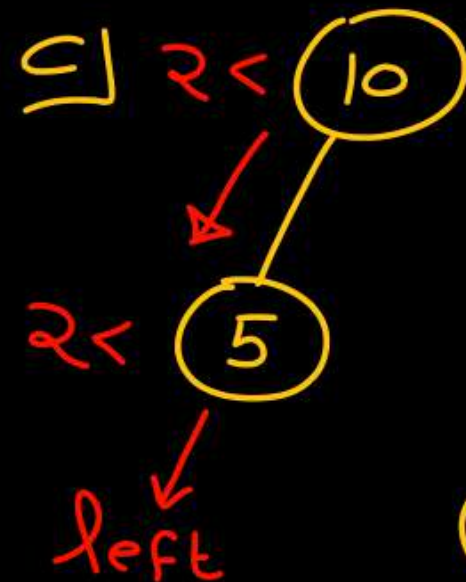
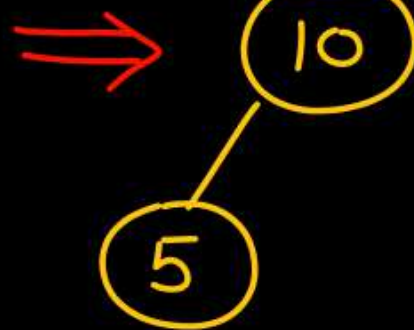
b]



c]



2. Construct binary search tree by inserting keys $\overrightarrow{10, 5, 2, 9, 3, 4}$ in order.



No. of BST when insertion order of keys are fixed

Const. BST by inserting keys 10, 20, 30 $\Rightarrow 1$

Const. BST by inserting keys 10, 5, 2, 9, 3, 4 $\Rightarrow 1$

Const. BST by inserting keys 1, 2, 3 (in any order)

(i) $\overrightarrow{1, 2, 3}$

(ii) 1, 3, 2

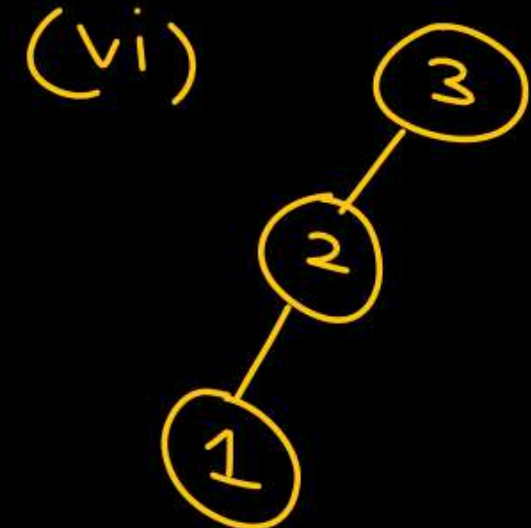
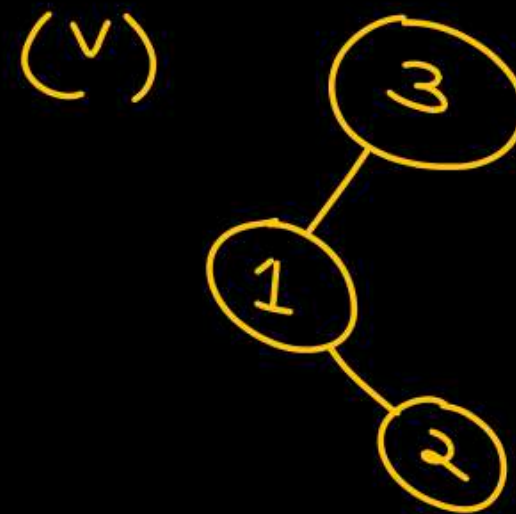
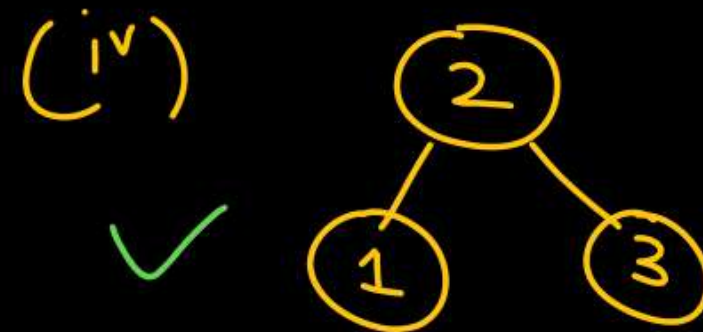
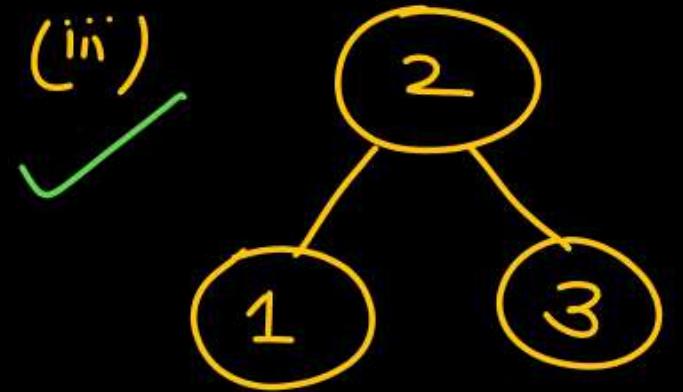
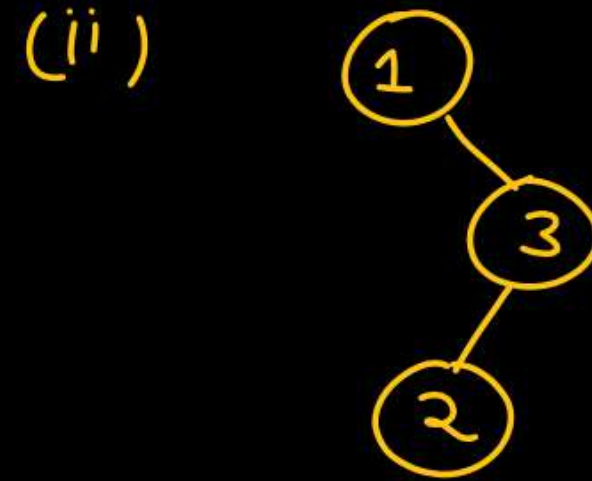
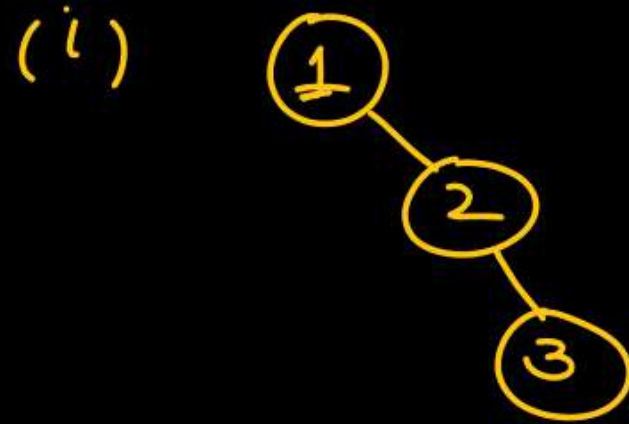
(iii) 2, 1, 3

(iv) 2, 3, 1

(v) 3, 1, 2

(vi) 3, 2, 1

6 orders possible

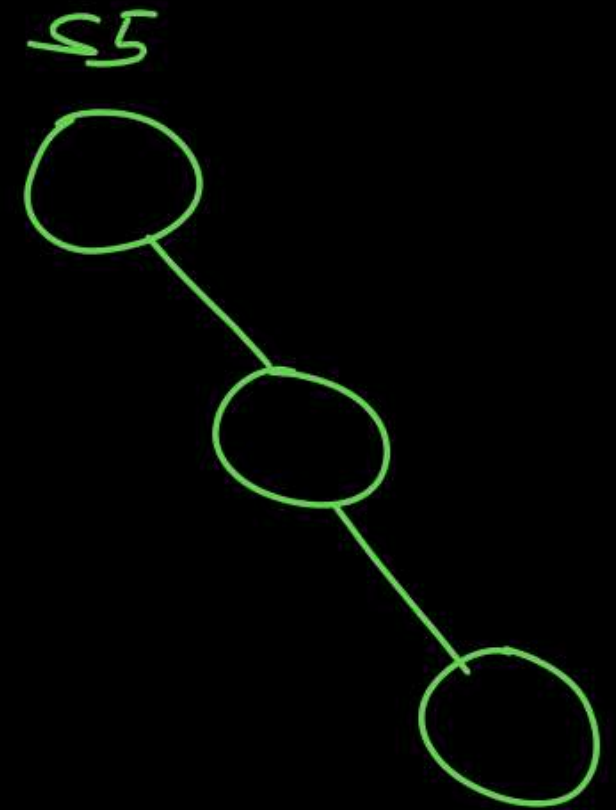
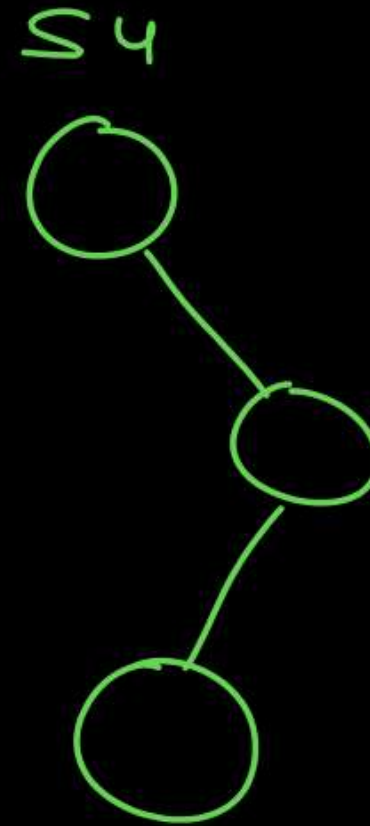
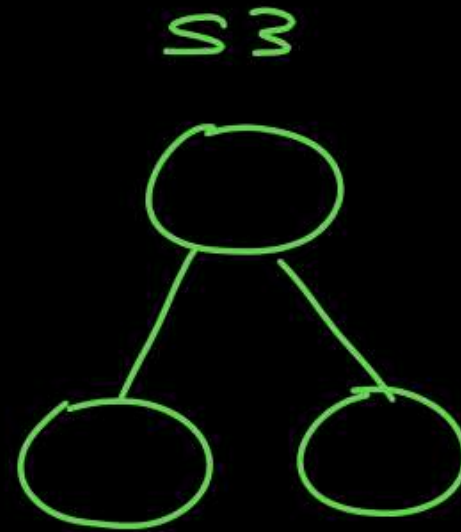
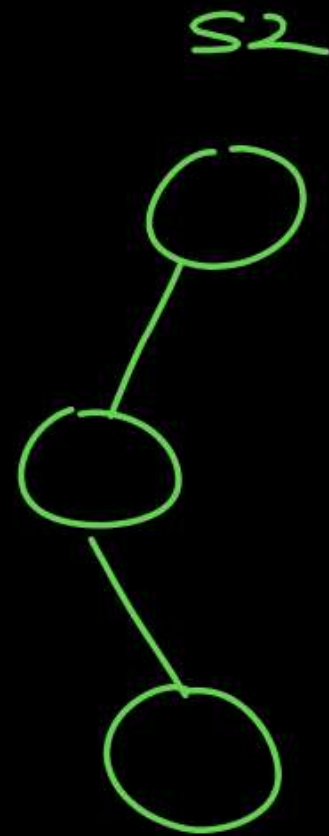
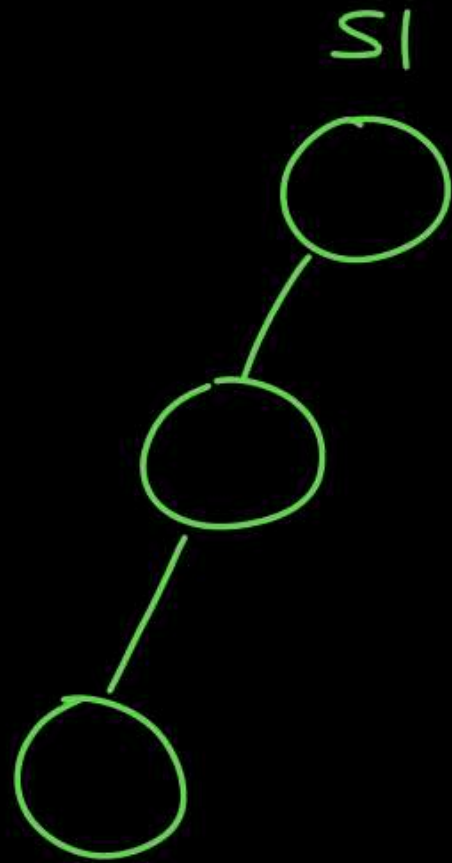


3 keys \Rightarrow 5 BST

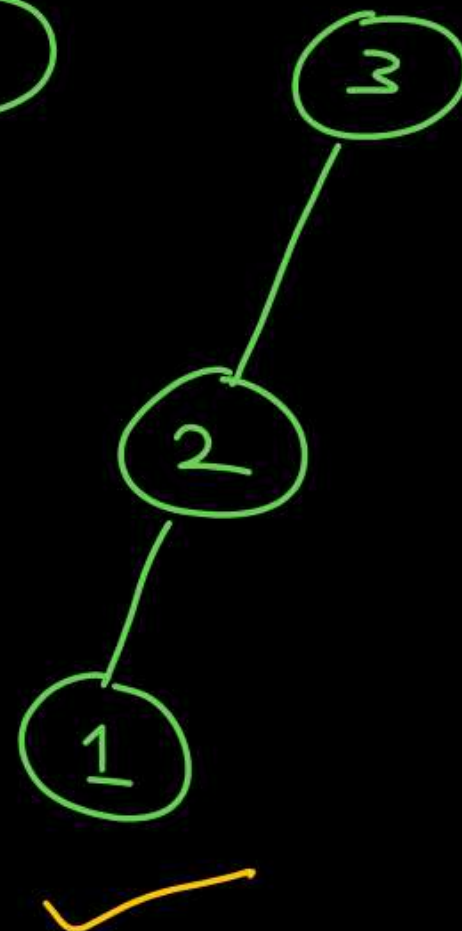
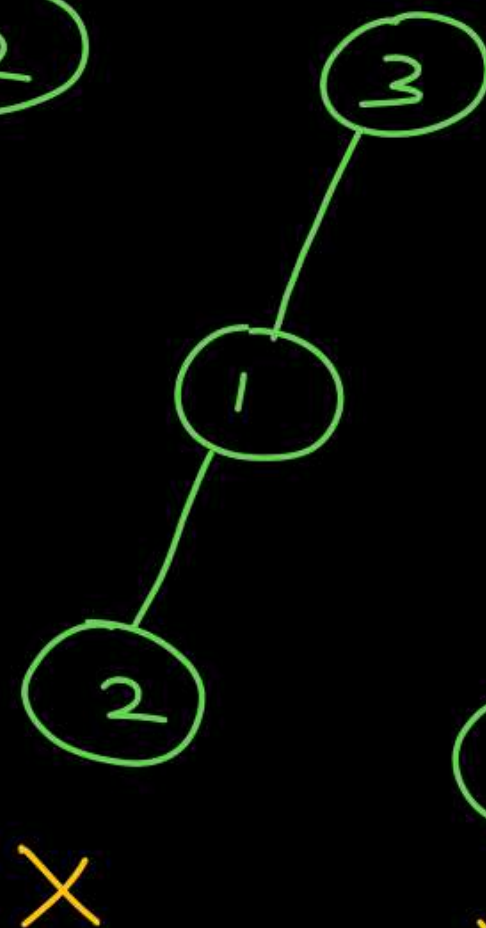
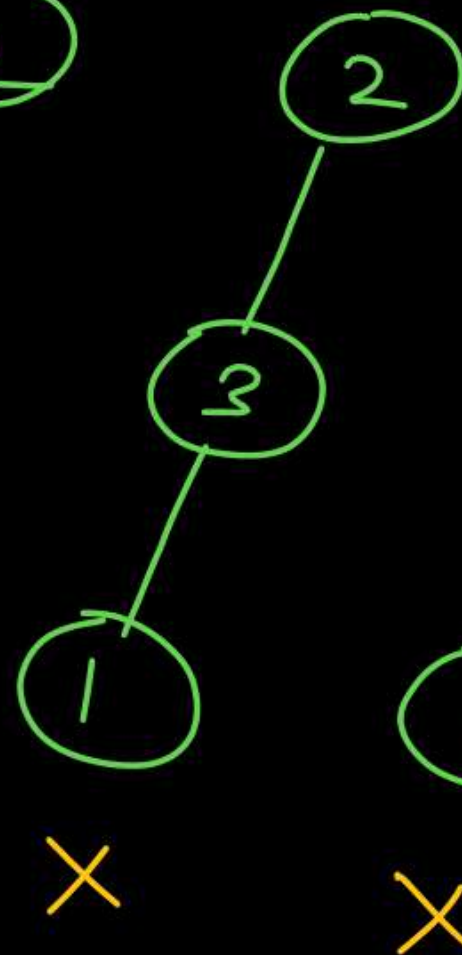
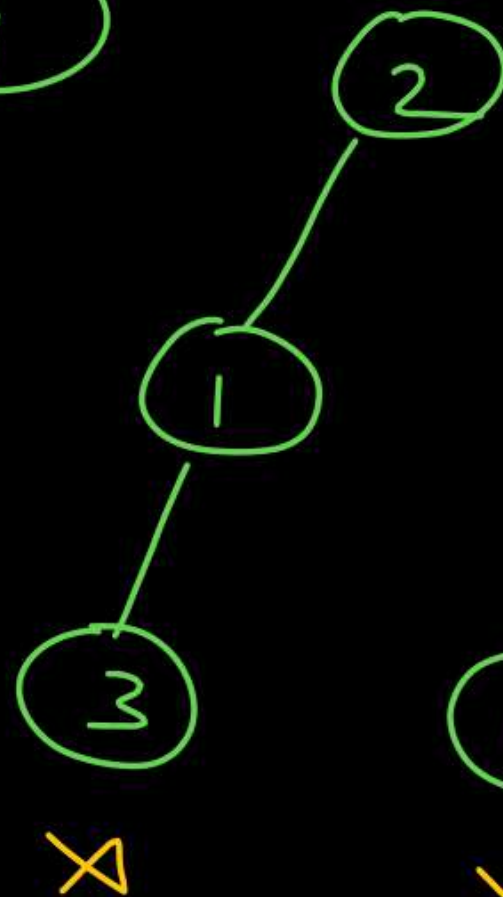
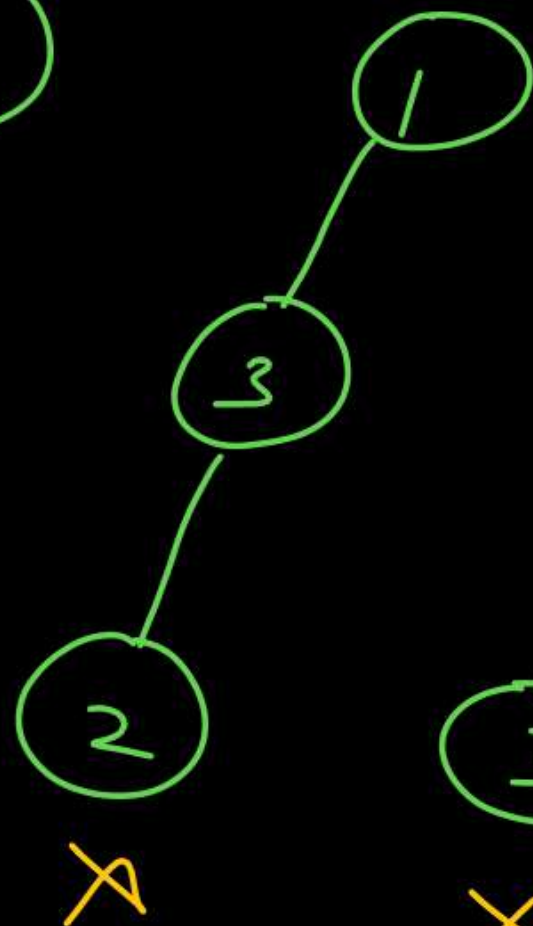
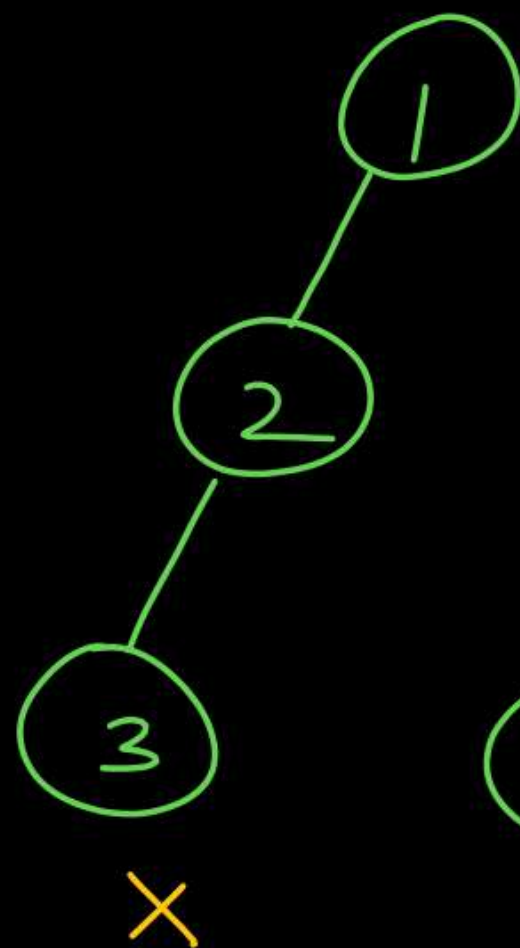
unlabelled binary tree with 3 nodes = 5

structure with n nodes = $\frac{2^n C_n}{n+1}$

3 node



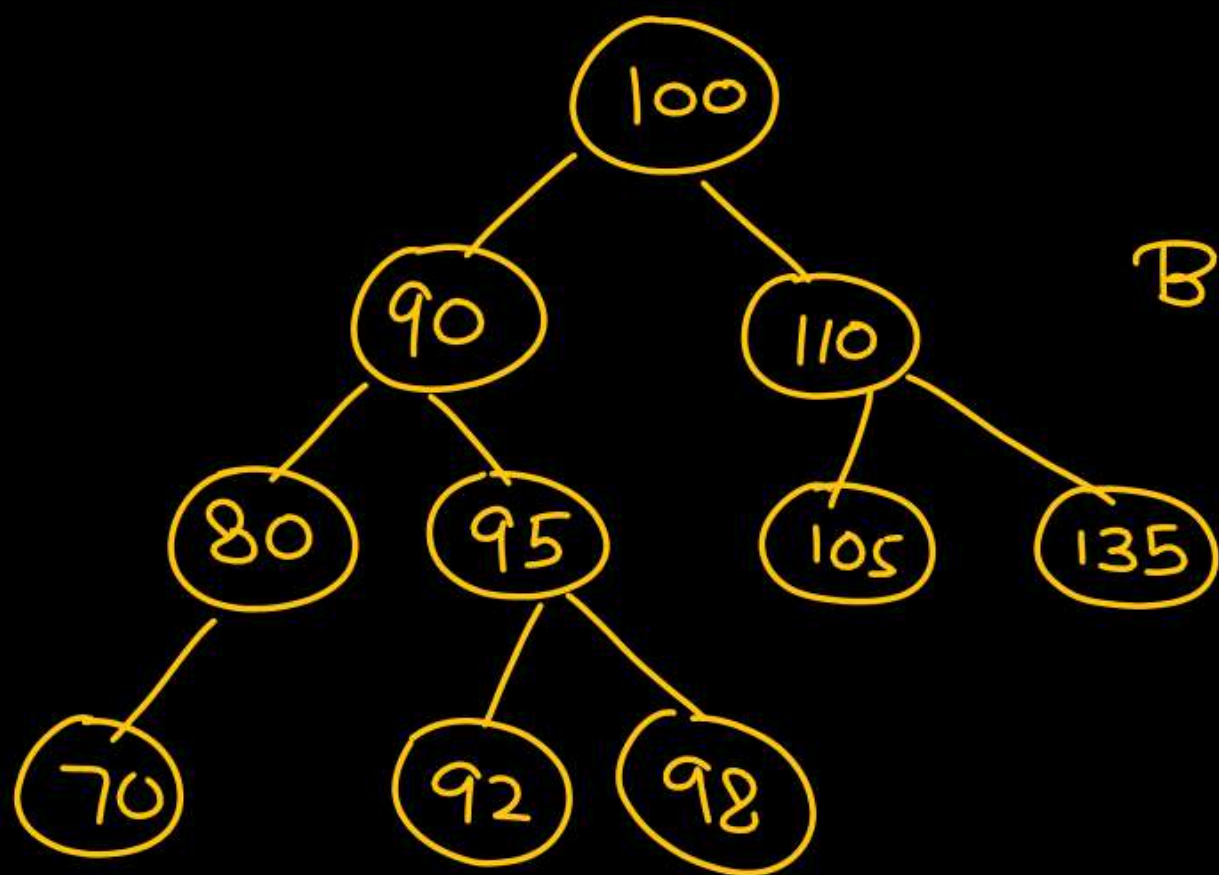
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$$\# \text{ BST with } n \text{ keys} = \frac{2^n C_n}{n+1}$$

2 ★★★★★

Inorder traversal of a BST is always increasing
order of keys.

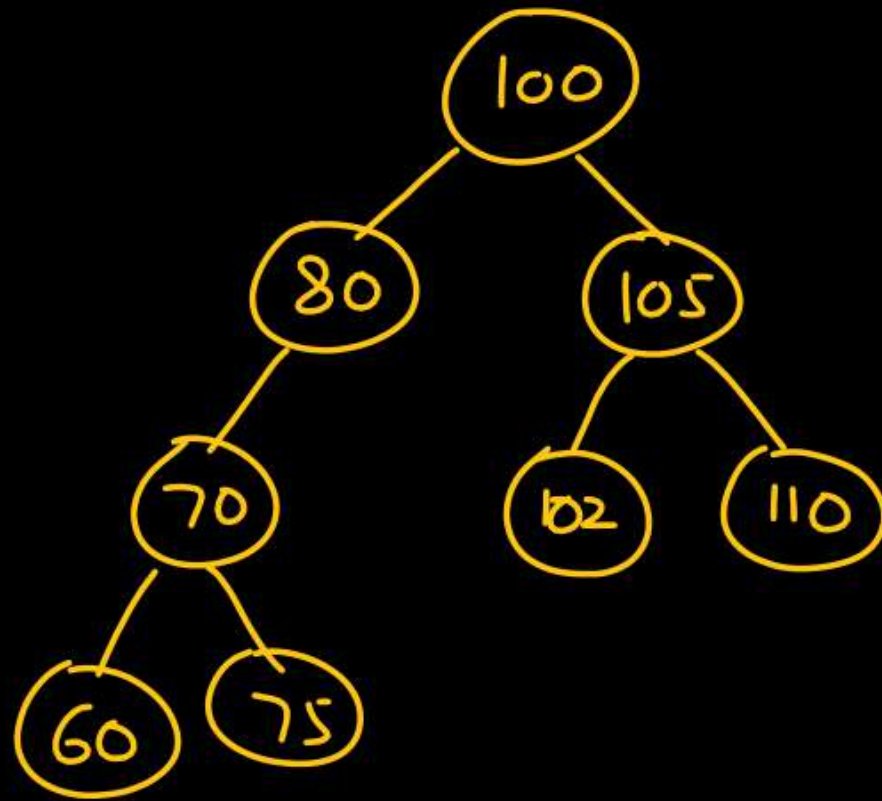


BST

Inorder : 70 80 90 92 95 98

100 105 110 135

Pre : 100 80 70 60 75 105 102 110



Pre : 100 80 70 60 75 105 102 110

Given the preorder of a BST as : 100, 80, 70, 60, 75, 105, 102, 110

the postorder of BST is —

a)

b)

c)

d)

Pre: 100 80 70 60 75 105 102 110

Given the preorder of a BST as : 100, 80, 70, 60, 75, 105, 102, 110
the postorder of BST is —

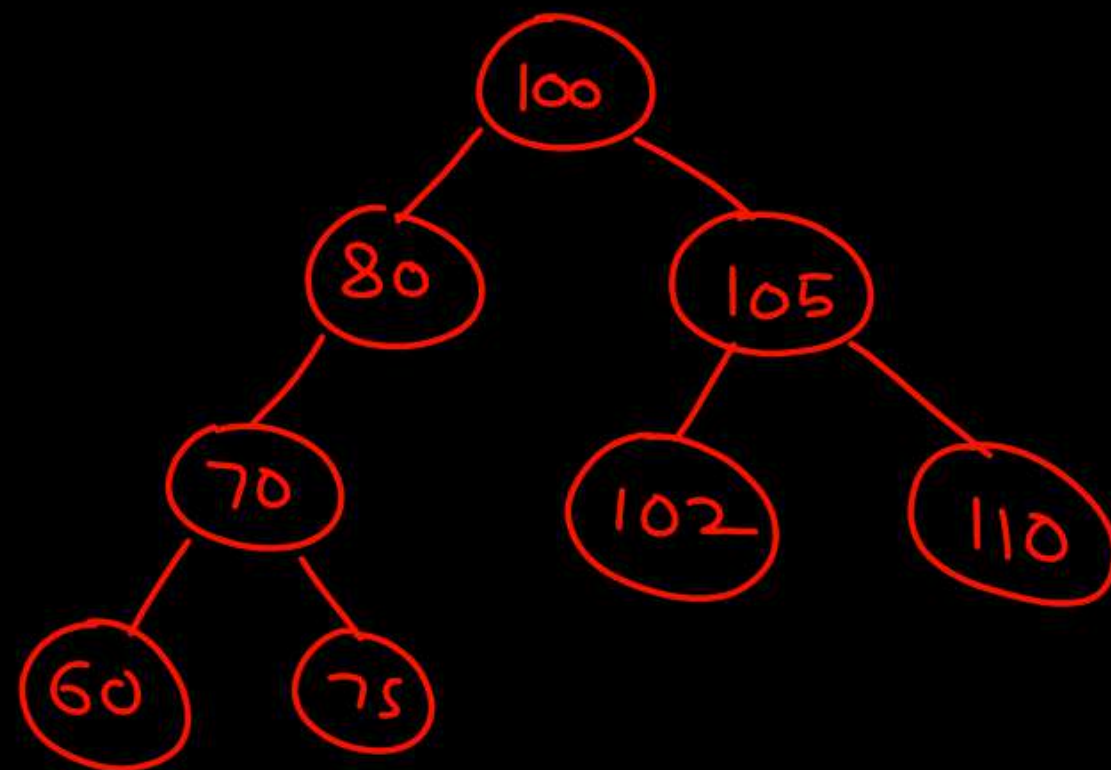
In: 60 70 75 80 100 102 105 110

Pre: 100 80 70 60 75 105 102 110

In: 60 70 75 80 100 102 105 110

Pre: 100 80 70 60 75 105 102 110

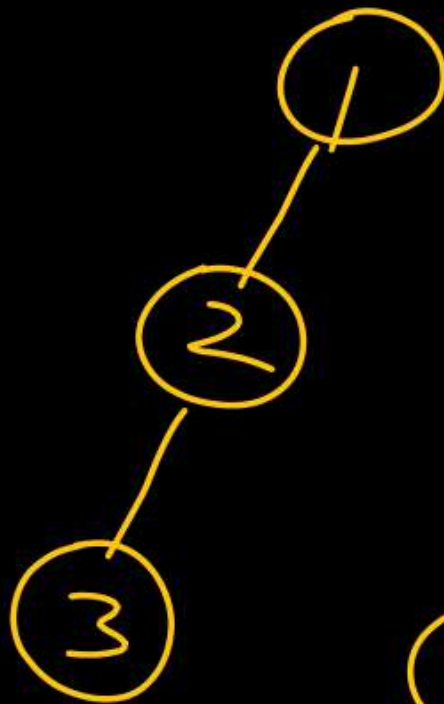
→



Q BST

Inorder:

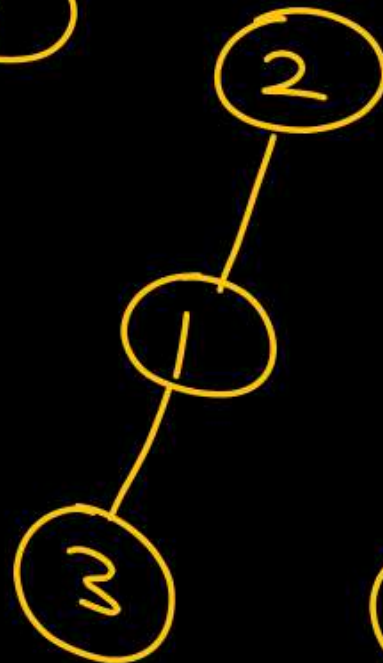
1, 2, 3 Keys



Inorder: 321



231



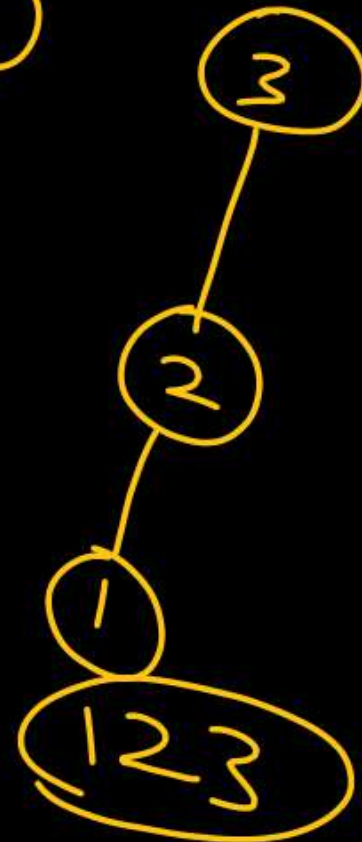
312



132

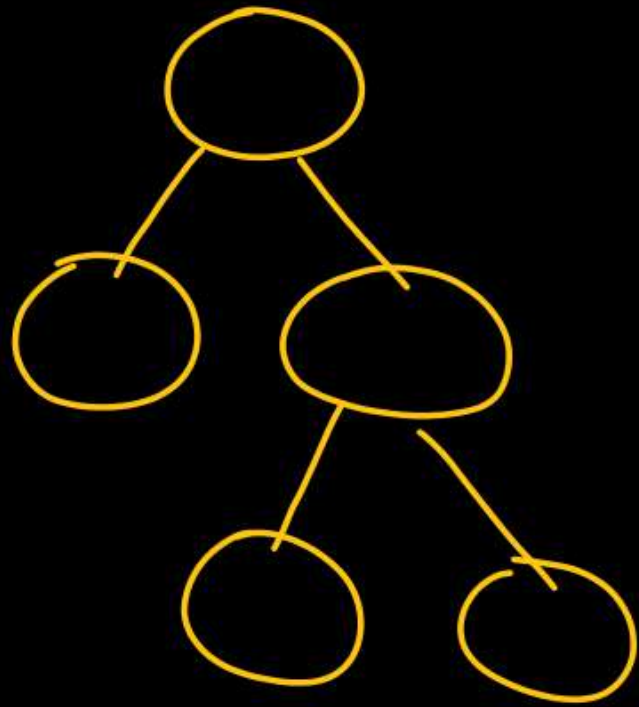


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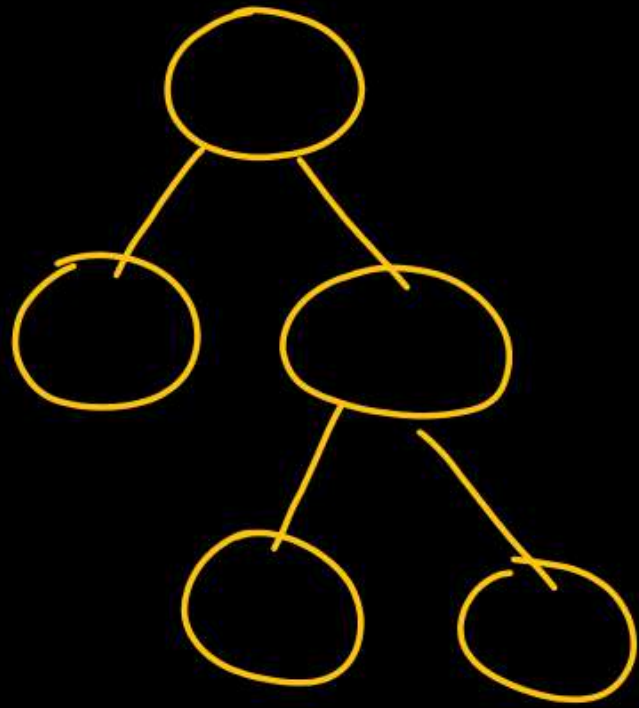


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Given a binary tree structure with n nodes (unlabelled)

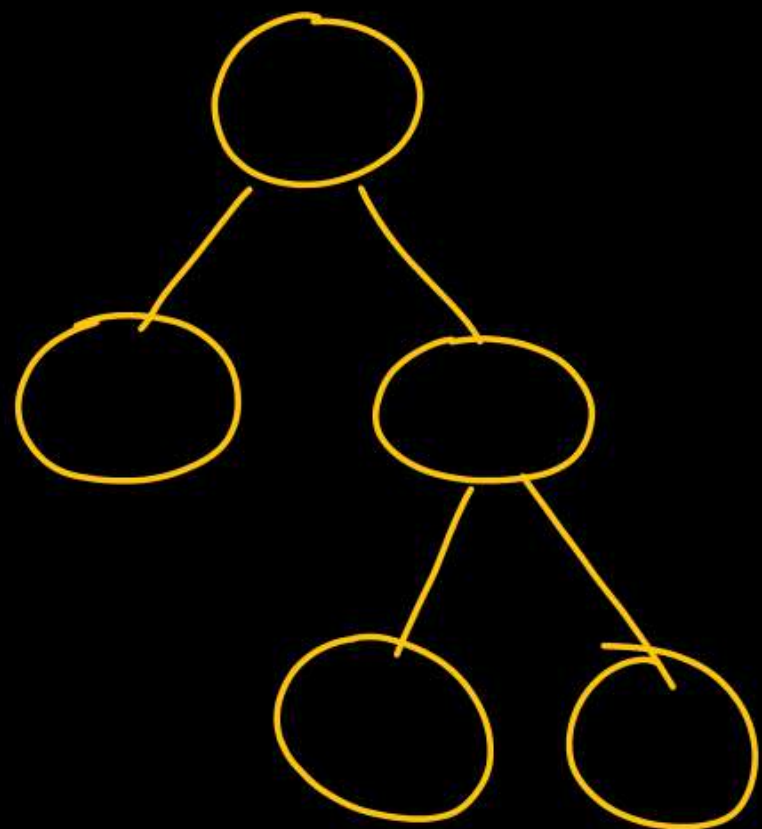


Given a binary tree structure with n nodes (unlabelled)



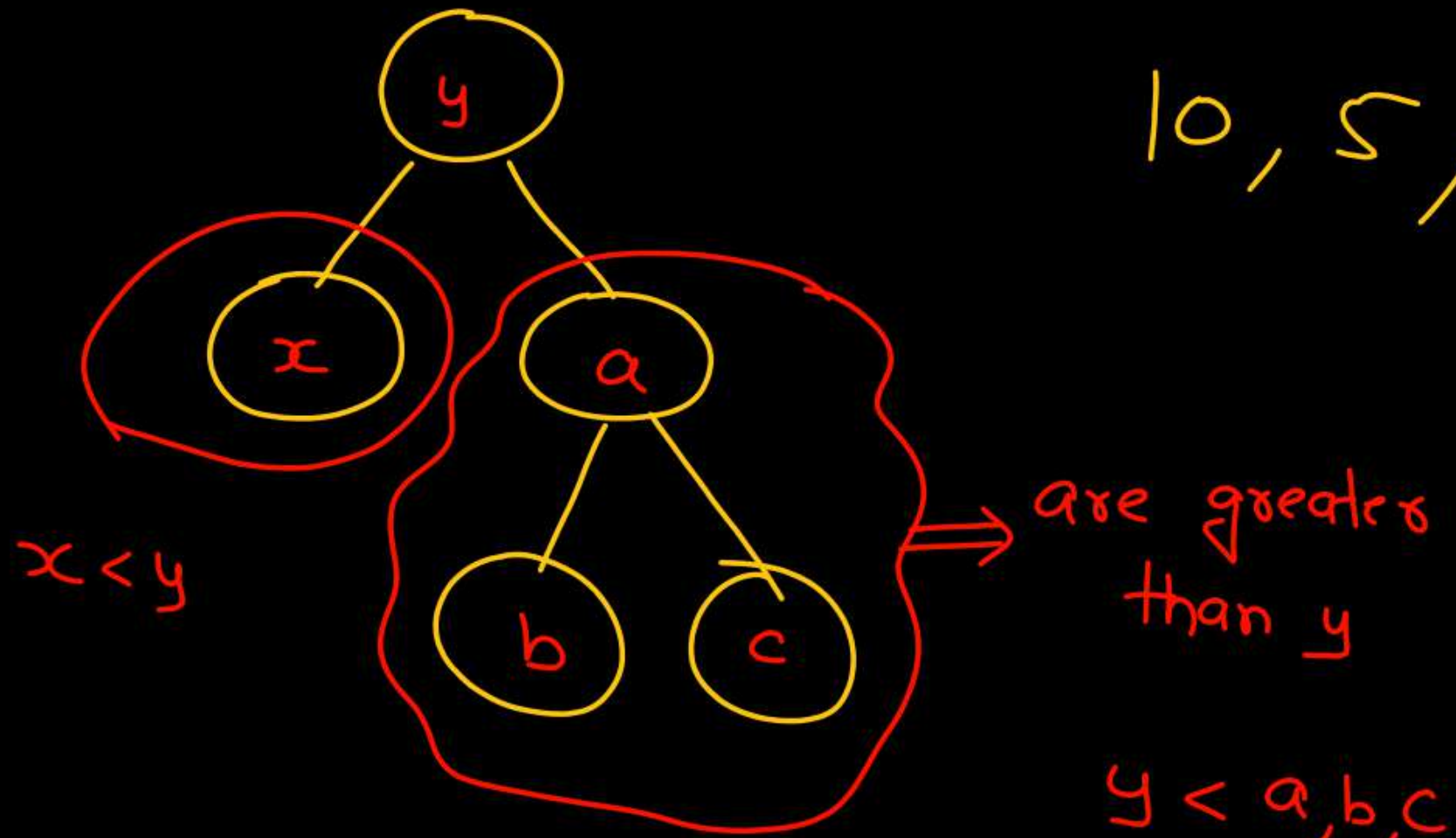
and n keys are also given

how many BSTs are
possible

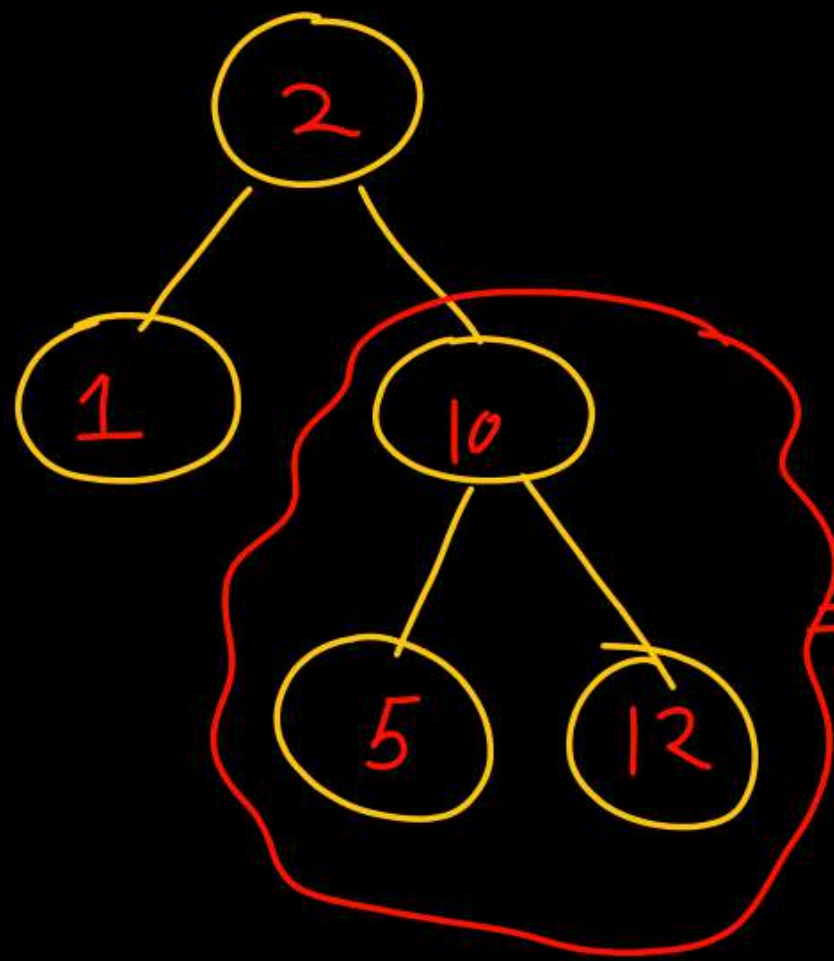


10, 5, 12, 2, 1

10, 5, 12, 2, 1



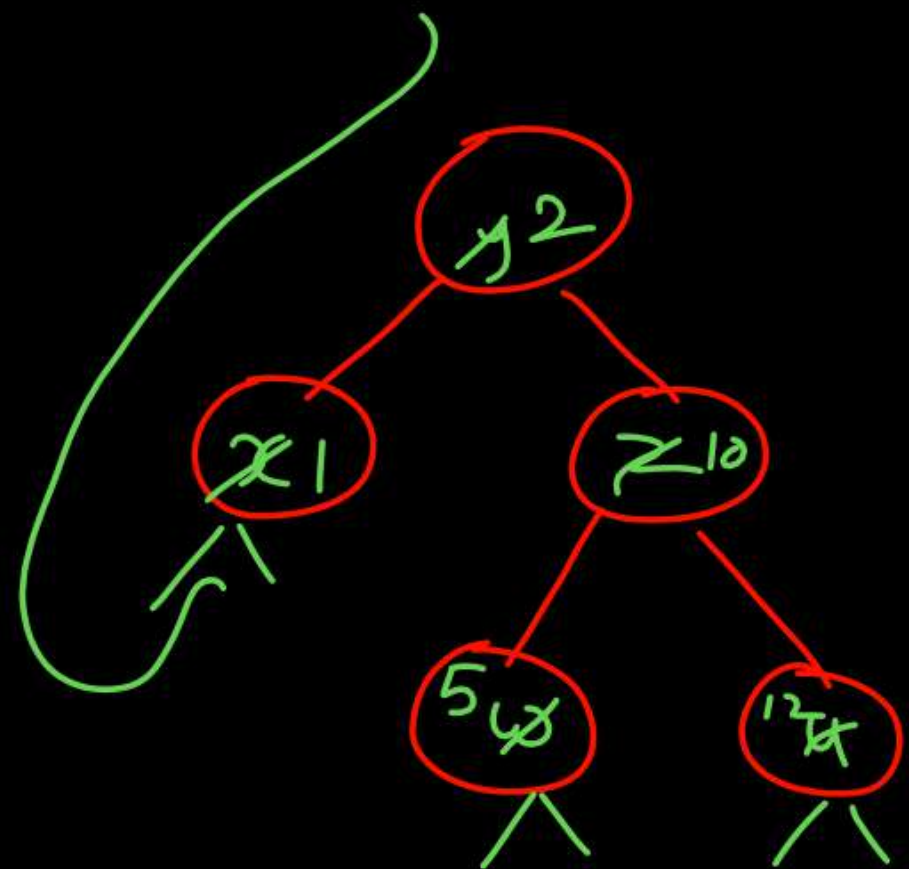
$x < y < a, b, c$
↓
Min



10, 5, 12, 2, 1

$$y < a, b, c$$

$$\text{Min} \rightarrow x < y < a, b, c$$



x, y, w, z, u

10, 5, 12, 2, 1
→
1 2 5 10 12

