

# 1 The algorithm

```
type Many a = [ : a : ]
type Image a = [ : [ : a : ] : ]

type Hist a = [ : a : ]

hbalanceBulk :: Many (Image Int) -> Many (Image Int)
hbalanceBulk = mapP hbalance

hbalance :: Image Int -> Image Int
hbalance img =
  let h = hist img
      a = accu h
      a0 = headP a
      agmax = lastP a
      n = normalize a0 agmax a
      s = scale gmax n
      img' = apply s img
  in img'

gmax :: Int
gmax = 255

hist :: Image Int -> Hist Int
hist =
  sparseToDenseP (gmax+1) 0
  . mapP (\g -> (headP g, lengthP g))
  . groupP
  . sortP
  . concatP

accu :: Hist Int -> Hist Int
accu = scanlP (+) 0

normalize :: Int -> Int -> Hist Int -> Hist Double
normalize a0' agmax' as =
  let a0 = P.fromIntegral a0'
      agmax = P.fromIntegral agmax'
      divisor = agmax D.- a0
  in [ : (P.fromIntegral freq' D.- a0) D./ divisor | freq' <- as : ]

scale :: Int -> Hist Double -> Hist Int
scale gmax as = [ : P.floor (a D.* P.fromIntegral gmax) | a <- as : ]

apply :: Hist Int -> Image Int -> Image Int
apply as img = mapP (mapP (as !:)) img
```

## 2 Utilised Functions

Table 1: Utilised function with their type signatures and their work and depth complexity.

Function	Work	Depth	Type
hbalanceBulkd imgs	?	?	$[[:[:[:\text{Int}]:]:]:] \rightarrow [[:[:[:\text{Int}]:]:]:]$
hbalance img	?	?	$[[:[:\text{Int}]:]:] \rightarrow [[:[:\text{Int}]:]:]$
hist img	$1 + wh(2 + \log(wh)) + gmax$	$\log gmax$	$[[:[:\text{Int}]:]:] \rightarrow [[:\text{Int}]:]$
accu h	$6n - 1$	$\log n$	$[[:\text{Int}]:] \rightarrow [[:\text{Int}]:]$
scale gmax n	$2n + 1$	1	$\text{Int} \rightarrow [[:\text{Double}]:] \rightarrow [[:\text{Int}]:]$
normalize $a_0$ $a_{gmax}$ a	$2n - 1$	1	$\text{Int} \rightarrow \text{Int} \rightarrow [[:\text{Int}]:] \rightarrow [[:\text{Double}]:]$
apply s img	$1 + w + wh$	1	$[[:\text{Int}]:] \rightarrow [[:[:\text{Int}]:]:] \rightarrow [[:[:\text{Int}]:]:]$
mapP f xs	n	1	$(a \rightarrow b) \rightarrow [[:a]:] \rightarrow [[:b]:]$
headP xs	1	1	$[[:a]:] \rightarrow a$
lastP xs	1	1	$[[:a]:] \rightarrow a$
concatP xss	1	1	$[[:a]:] \rightarrow [[:a]:]$
sortP xs	$n \log n$	$\log n$	$\text{Ord } a \Rightarrow [[:a]:] \rightarrow [[:a]:]$
groupP xs	$(n + \text{groups in xs}) \log n$	$\log n$	$\text{Eq } a \Rightarrow [[:a]:] \rightarrow [[:[:a]:]:]$
sparseToDenseP xs	$k \cdot L(\text{result})$	1	$\text{Int} \rightarrow \text{Int} \rightarrow [[:a, \text{Int}]:] \rightarrow [[:a]:]$
lengthP xs	1	1	$[[:a]:] \rightarrow \text{Int}$
scanlP f z xs	$6n - 1$	$\log n$	$(a \rightarrow a \rightarrow a) \rightarrow a \rightarrow [[:a]:] \rightarrow [[:a]:]$
indexP,!: i xs	1	1	$[[:a]:] \rightarrow \text{Int} \rightarrow a$

where  $p$  is the number of processors. Functions in *italics* denote an  $O(\cdot)$ -class rather than the actual number of steps. Note, that various array class constraints ( $\text{PAElt } a \Rightarrow \dots$ ) on array functions omitted were omitted.

## 3 Variation of the number of processors

With  $P$  processors and a cross-process communication latency of  $L$ , a program with work  $W$  and depth  $D$  will run time  $T$  with such that the following holds:

$$\frac{W}{P} \leq T \leq \frac{W}{P} + L \cdot D$$

Note that,  $W$  and  $D$  have to be the exact numbers rather than an  $O(\cdot)$ -class

## 4 scanlP

$$\begin{aligned}
W(n) &= n(2 + W(f)) + W\left(\frac{n}{2}\right) \\
&= n(2 + W(f))(2 - 2^{\log_2 n}) + 1 \\
&= (2 + W(f))(2n - 1) + 1 \\
&= 6n - 1 \\
D(n) &= \log n
\end{aligned} \tag{1}$$

## 5 sparseToDenseP

```

sparseToDenseP :: Enum e => e -> a -> [:(e,a):] -> [:(a):]
sparseToDenseP      many init map      result
sparseToDenseP 0 7 9 [:(1,5),(2,4),(6,7):] == [:(0,5,4,0,0,0,7,0):]

```

sparseToDense n z map creates an array of length n where the element at the index i has x if (i,x) is in the map, or z otherwise. In effect it turns a sparse vector to a dense one

## 6 groupP

$$\begin{aligned} W(n) &= (n + \text{unique elements in xs}) \log n \\ D(n) &= \log n \end{aligned} \quad (2)$$

## 7 hbalance

Let  $|img| = w \cdot h$  and  $n = gmax$ .

$$\begin{aligned} W(w \times h) &= W(hist) + W(accum) + W(headP) + W(lastP) + W(normalize) + W(scale) + W(apply) \\ &= (1 + |img|(2 + \log |img|) + n) + (6n - 1) + 1 + 1 + (2n + 1) + (2n + 1) + (1 + w + |img|) \\ &= w + |img|(3 + \log |img|) + 11n + 5 \\ &\in O(|img| \log |img| + gmax) \\ D(w \times h) &= 1 + \max\{D(hist), D(accum), D(headP), D(lastP), D(normalize), D(scale), D(apply)\} \\ &= 1 + \max\{\log |img|, \log gmax\} \end{aligned} \quad (3)$$

With  $P$  processors we have an expected runtime of  $T$  where:

$$\begin{aligned} \frac{W}{P} &\leq T \leq \frac{W}{P} + D \\ \frac{|img| \log |img| + gmax}{P} &\leq T \leq \frac{|img| \log |img| + gmax}{P} + \max\{\log |img|, \log gmax\} \end{aligned}$$

Fall 1:  $|img| > gmax$ . Es gibt mehr Bildpixel als es zulässige Grauwerte gibt.

$$\begin{aligned} \frac{2|img| \log |img|}{P} &\leq T \leq \frac{2|img| \log |img|}{P} + \log |img| \quad | \div (2 \log |img|) \\ \frac{|img|}{P} &\leq T \leq \frac{|img|}{P} + \frac{1}{2} \end{aligned}$$

Fall 2:  $gmax > |img|$ . Der Grauwertbereich ist größer als die Anzahl der Pixel

$$\begin{aligned} \frac{2gmax \log gmax}{P} &\leq T \leq \frac{2gmax \log gmax}{P} + \log gmax \quad | \div (2 \log gmax) \\ \frac{gmax}{P} &\leq T \leq \frac{gmax}{P} + \frac{1}{2} \end{aligned} \quad (4)$$

## 8 hist

Let  $|img| = w \cdot h$  and  $n = \text{gmax}$ .

$$\begin{aligned}
W(w \times h) &= W(\text{concat}P) + W(\text{sort}P) + W(\text{group}P) + 2W(\text{map}P) + W(\text{sparseToDense}P) \\
&= 1 + |img| \log |img| + 2|img| + n \\
&= 1 + |img|(2 + \log |img|) + n \\
D(w \times h) &= \log |img|
\end{aligned} \tag{5}$$

## 9 accu

$$\begin{aligned}
W(n) &= W(\text{scanl}P) \\
&= 6n - 1 \\
D(n) &= D(\text{scanl}P) \\
&= \log n
\end{aligned} \tag{6}$$

## 10 normalize

$$\begin{aligned}
W(n) &= 2n + 1 \\
D(n) &= 1
\end{aligned} \tag{7}$$

## 11 apply

Let  $w$  and  $h$  be the width and height of the image ( $img$ ).

$$\begin{aligned}
W(w \times h) &= 1 + w \cdot W(\text{inner}) \\
&= 1 + w(1 + h) \\
&= 1 + w + wh \\
D(w \times h) &= 1
\end{aligned} \tag{8}$$