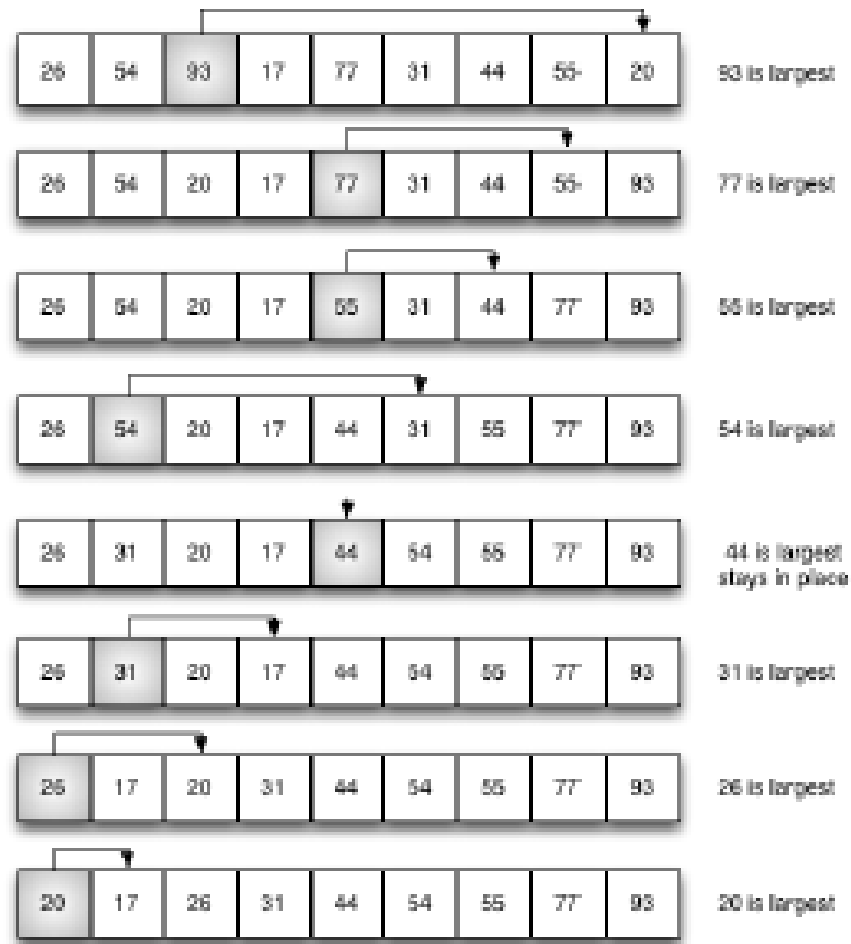




# Selection Sort

---



# Selection Sort

---

```
#ArrayLst -> ArrayLst
#sort using selection sort
def selection_sort(lst):
    lst_size = lst.size
    i = 0
    while i < lst_size-1:
        maxIndex = 0
        j = 1
        while j < lst_size-1-i:
            if lst.value[j] > lst.value[maxIndex]:
                maxIndex = j
            j += 1
        temp = lst.value[lst_size-1-i]
        lst.value[lst_size-1-i] = lst.value[maxIndex]
        lst.value[maxIndex] = temp
        i += 1
```

# Insertion Sort

---

54	26	93	17	77	31	44	55	20	Assume 54 is a sorted list of 1 item
26	54	93	17	77	31	44	55	20	inserted 26
26	54	93	17	77	31	44	55	20	inserted 93
17	26	54	93	77	31	44	55	20	inserted 17
17	26	54	77	93	31	44	55	20	inserted 77
17	26	31	54	77	93	44	55	20	inserted 31
17	26	31	44	54	77	93	55	20	inserted 44
17	26	31	44	54	55	77	93	20	inserted 55
17	20	26	31	44	54	55	77	93	inserted 20

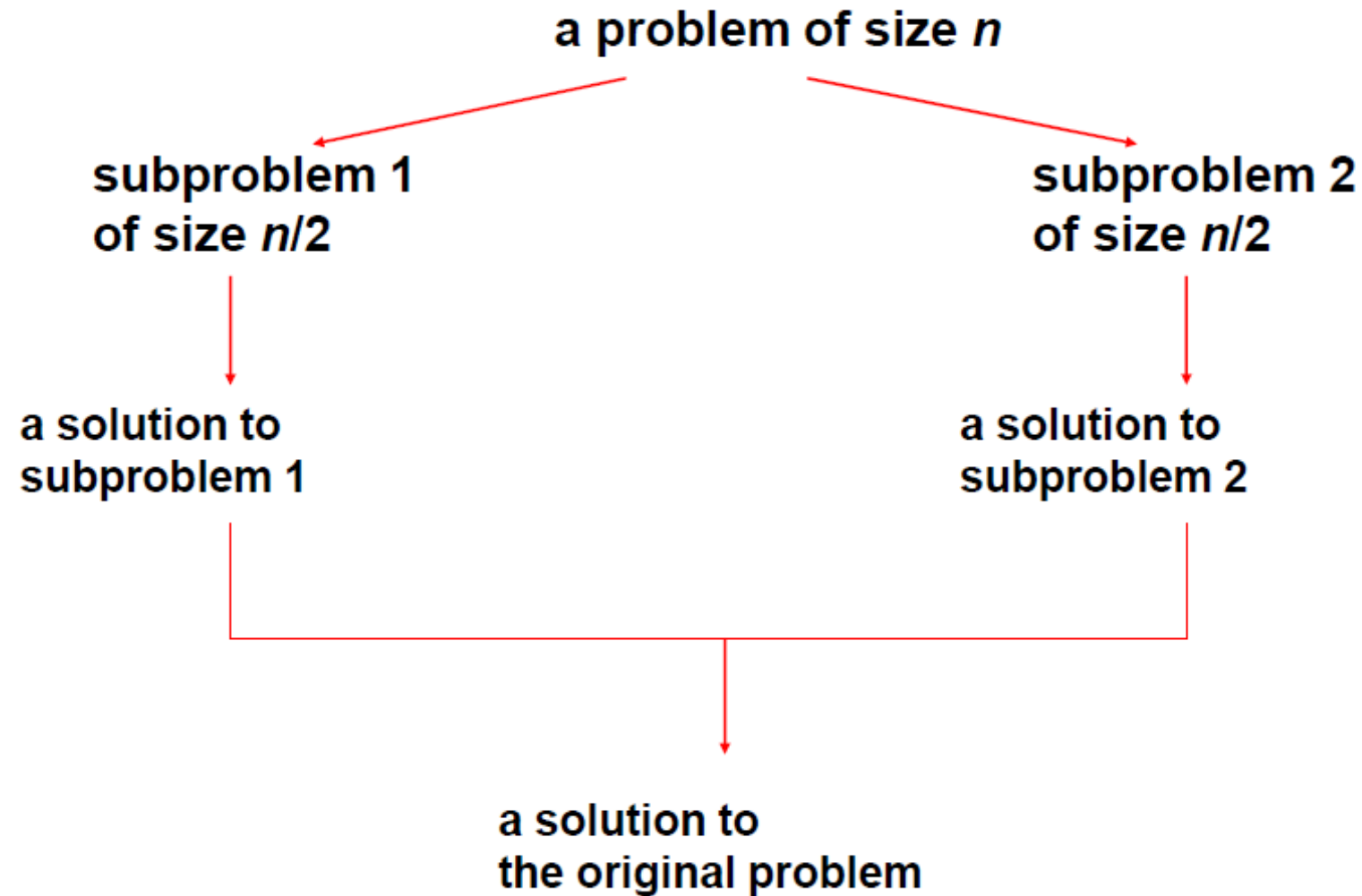
# Insertion Sort

---

```
#ArrayLst -> ArrayLst
#sort using insertion sort
def insertion_sort(lst):
    lst_size = lst.size
    swapped = 1
    i = 0
    while i < lst_size-1 and swapped == 1:
        swapped = 0
        j=0
        while j < lst_size-i-1:
            if lst.value[j] > lst.value[j+1]:
                temp = lst.value[j]
                lst.value[j] = lst.value[j+1]
                lst.value[j+1] = temp
                swapped = 1
            j += 1
        i += 1
    return lst
```

# Divide-and-Conquer Technique

---



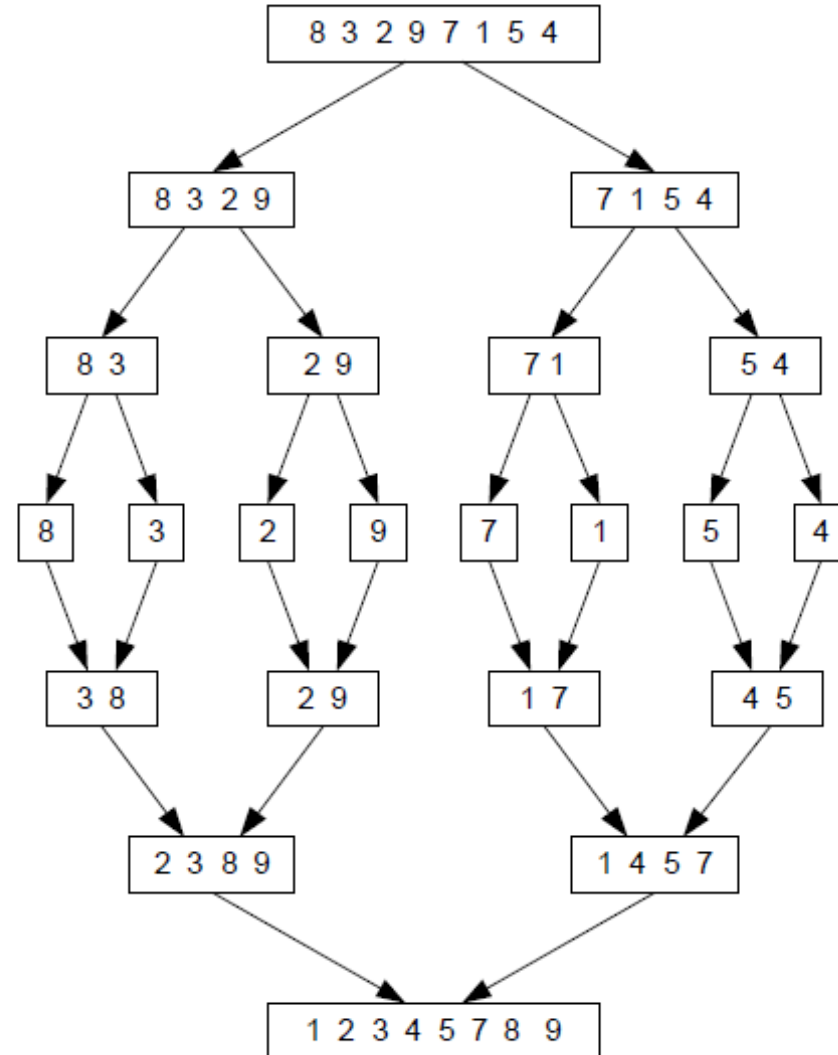
# Mergesort

---

- Split array  $A[0..n-1]$  in two about equal halves and make copies of each half in arrays B and C
- Sort arrays B and C recursively
- Merge sorted arrays B and C into array A as follows:
  - Repeat the following until no elements remain in one of the arrays:
    - » compare the first elements in the remaining unprocessed portions of the arrays
    - » copy the smaller of the two into A, while incrementing the index indicating the unprocessed portion of that array
  - Once all elements in one of the arrays are processed, copy the remaining unprocessed elements from the other array into A.

# Mergesort Example

---





# Mergesort

---

```
]def mergesort(list):  
    if (len(list) < 2):  
        return list  
  
    middle = len(list)//2  
    left = mergesort(list[:middle])  
    right = mergesort(list[middle:])  
  
    return merge(left, right)
```

```
def merge(left, right):  
    print("\nleft:", left, "\nright:", right)  
    result = []  
    i, j = 0, 0  
    while (len(result) < len (left) + len(right)):  
        if left[i] < right[j]:  
            result.append(left[i])  
            i += 1  
        else:  
            result.append(right[j])  
            j += 1  
        if i == len(left):  
            result.extend(right[j:])  
        if j == len(right):  
            result.extend(left[i:])  
    print("result", result)  
    return result
```

# Quicksort

---

- Fastest known sorting algorithm in practice
  - Caveats: not stable
- Average case complexity  $\rightarrow O(N \log N)$
- Worst-case complexity  $\rightarrow O(N^2)$ 
  - Rarely happens, if coded correctly

# Quicksort Outline

---

- Divide and conquer approach
- Given array  $S$  to be sorted
  - If size of  $S < 1$  then done;
  - Pick any element  $v$  in  $S$  as the pivot
  - Partition  $S - \{v\}$  (remaining elements in  $S$ ) into two groups
    - »  $S1 = \{\text{all elements in } S - \{v\} \text{ that are } \leq \text{than } v\}$
    - »  $S2 = \{\text{all elements in } S - \{v\} \text{ that are } \geq \text{than } v\}$
  - Return {quicksort( $S1$ ) followed by  $v$  followed by quicksort( $S2$ ) }
- Trick lies in handling the partitioning (step 3).
  - Picking a good pivot
  - Efficiently partitioning in-place

# Picking the Pivot

---

Strategy 1: Pick the first element in S

- Works only if input is random
- What if input S is sorted, or even mostly sorted?
  - All the remaining elements would go into either S1 or S2!
  - Terrible performance!
- Why worry about sorted input?
  - Remember → Quicksort is recursive, so sub-problems could be sorted
  - Plus mostly sorted input is quite frequent

## Picking the Pivot (contd.)

---

Strategy 2: Pick the pivot **randomly**

- Would usually work well, even for mostly sorted input
- Unless the random number generator is not quite random!
- Plus random number generation is an expensive operation

## Picking the Pivot (contd.)

---

### Strategy 3: Median-of-three Partitioning

- *Ideally*, the pivot should be the median of input array  $S$ 
  - Median = element in the middle of the sorted sequence
- Would divide the input into two almost equal partitions
- But, its hard to calculate median quickly, without sorting first!
- So find the approximate median
  - Pivot = median of the left-most, right-most and center element of the array  $S$
  - Solves the problem of sorted input

## Picking the Pivot (contd.)

---

- Example: Median-of-three Partitioning
  - Let input  $S = \{6, 1, 4, 9, 0, 3, 5, 2, 7, 8\}$
  - $\text{left}=0$  and  $S[\text{left}] = 6$
  - $\text{right}=9$  and  $S[\text{right}] = 8$
  - $\text{center} = (\text{left}+\text{right})/2 = 4$  and  $S[\text{center}] = 0$
  - Pivot
    - » = Median of  $S[\text{left}]$ ,  $S[\text{right}]$ , and  $S[\text{center}]$
    - » = median of 6, 8, and 0
    - » =  $S[\text{left}] = 6$



# Partitioning Algorithm Idea

---

- Original input :  $S = \{6, 1, 4, 9, 0, 3, 5, 2, 7, 8\}$
- Get the pivot out of the way by swapping it with the last element

8 1 4 9 0 3 5 2 7 **6**  
pivot

- Have two 'iterators' –  $i$  and  $j$ 
  - $i$  starts at first element and moves forward
  - $j$  starts at last element and moves backwards

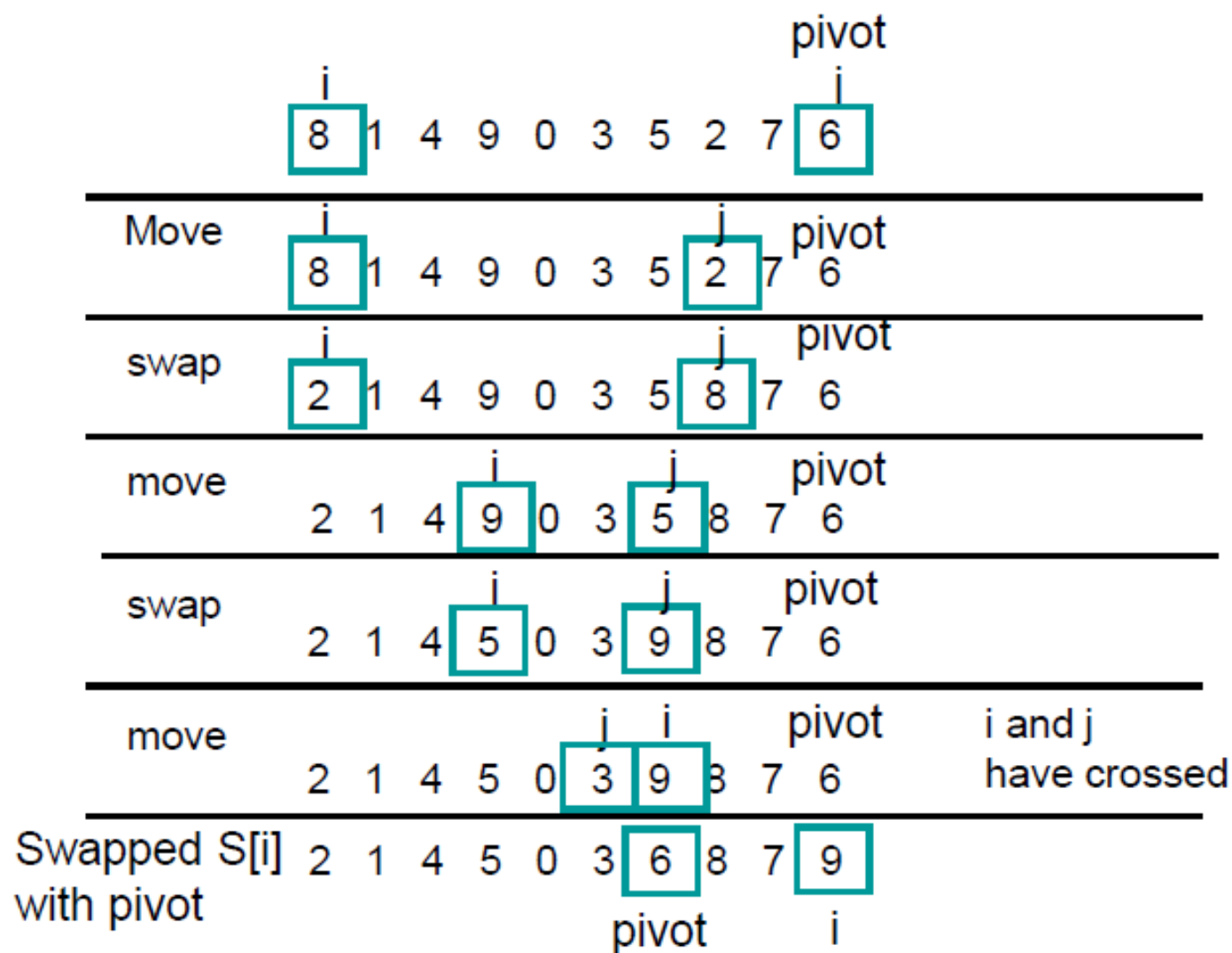
**8** 1 4 9 0 3 5 2 7 **6**  
 $i$  pivot  $j$

## Partitioning Algorithm Idea (contd.)

---

- `While (i < j)`
  - Move `i` to the right till we find a number  $\geq$  than `pivot`
  - Move `j` to the left till we find a number  $\leq$  than `pivot`
  - `If (i < j) swap(S[i], S[j])`
- Swap the `pivot` with `S[i]`

# Partitioning Algorithm Illustrated



# Quicksort

---

```
def quickSort(alist):  
    quickSortHelper(alist,0,len(alist)-1)  
  
def quickSortHelper(alist,first,last):  
    if first<last:  
        splitpoint = partition(alist,first,last)  
        quickSortHelper(alist,first,splitpoint-1)  
        quickSortHelper(alist,splitpoint+1,last)
```

```
def partition(alist, first, last):
```

```
    pivotvalue = alist[first]
```

```
    print("list start:", alist)
```

```
    print("pivot:", pivotvalue)
```

```
    leftmark = first+1
```

```
    rightmark = last
```

```
    done = False
```

```
    while not done:
```

```
        while leftmark <= rightmark and alist[leftmark] <= pivotvalue:
```

```
            leftmark = leftmark + 1
```

```
        while alist[rightmark] >= pivotvalue and rightmark >= leftmark:
```

```
            rightmark = rightmark - 1
```

```
        if rightmark < leftmark:
```

```
            done = True
```

```
        else:
```

```
            temp = alist[leftmark]
```

```
            alist[leftmark] = alist[rightmark]
```

```
            alist[rightmark] = temp
```

Just choosing first element as pivot

Could use median of three, random or other method

## Dealing with small arrays

---

- For small arrays ( $N \leq 20$ ) ,
  - Insertion sort is faster than quicksort
- Quicksort is recursive
  - So it can spend a lot of time sorting small arrays
- Hybrid algorithm:
  - Switch to using insertion sort when problem size is small (say for  $N < 20$ )