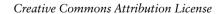
### Dissecting Partitioning

BRUCE MOMJIAN



This presentation shows the purpose, effect, and optimizations of Postgres's partitioning implementation.

https://momjian.us/presentations





Last updated: May 2023

#### Outline

- 1. Purpose and history
- 2. Partitioning in action
- 3. Partition pruning
- 4. Partition-local optimizations
- 5. Time-based partitioning
- 6. Row migration
- 7. psql support
- 8. Limitations
- 9. Complex architectures

## 1. Purpose and History



https://www.flickr.com/photos/photo\_oto/

### Purpose of Partitioning

- Query efficiency
  - faster sequential scans for partition-key-qualified queries
  - shallower indexes for faster index scans
  - improve cache efficiency for frequently accessed partitions
  - partition-local optimizations
- Data customization
  - partition-specific indexes, constraints, and triggers
- Maintenance
  - improve efficiency of modified row cleanup
  - movement of partitions to tablespaces with different storage characteristics
  - easily attach/detach partitions for maintenance
  - simpler bulk operations, including removal of old time-based partitions

https://www.postgresql.org/docs/current/ddl-partitioning.html

https://hevodata.com/learn/postgresql-partitions/

https://www.youtube.com/watch?v=edQZauVU-ws

### Partitioning Prior to Postgres 10

Prior to Postgres 10, partitioning was accomplished by combining:

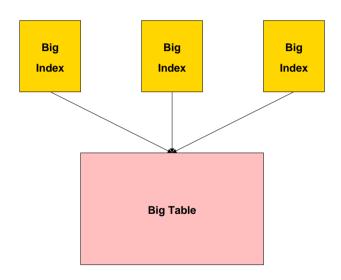
- Inheritance
- CHECK constraints and constraint\_exclusion
- Triggers for INSERT routing of new rows
- Triggers for UPDATEs that change partitioned key columns
- Use of external tools like pg\_partman (https://github.com/pgpartman/pg\_partman)

## Enhancing Partitioning Capabilities

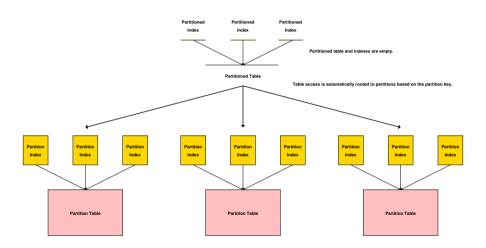
#### Postgres 10 and later added:

- Range, hash, and list partition syntax
- DEFAULT partitions\*
- PRIMARY KEYs on partitioned\* tables and FOREIGN KEYs to and from partitioned tables
- More efficient pruning
- Executor-stage pruning
- Partition-local optimizations
- Used with foreign data wrappers
- \* "Partitioned tables" are parent tables that are referenced in most queries and store no data, while "partitions" are tables where the data is stored.

## Without Partitioning



## With Partitioning



# 2. Partitioning in Action



https://www.flickr.com/photos/105105658@N03/

### Range Partitioned Table Creation

```
CREATE TABLE range partitioned (name TEXT)
    PARTITION BY RANGE (name):
    CREATE TABLE range partition less j
    PARTITION OF range partitioned
    FOR VALUES FROM (MINVALUE) TO ('i');
    CREATE TABLE range partition j to s
    PARTITION OF range partitioned
    FOR VALUES FROM ('i') TO ('s');
    CREATE TABLE range partition s greater
    PARTITION OF range partitioned
    FOR VALUES FROM ('s') TO (MAXVALUE):
Summary and causal information is in blue; significant information is in red. All the
queries used in this presentation are available at https://momjian.us/main/writings/pgsql/
partitioning.sql.
```

#### **DEFAULT Partition Table Creation**

```
-- Range partitioned tables require DEFAULT partitions for NULL storage.
-- CHECK prevents non-NULLs and avoids partition scan checking for
-- values that might be in the newly-created partition.

CREATE TABLE range_partition_nulls

PARTITION OF range_partitioned
(CHECK (name IS NULL))

DEFAULT;
```

#### The Result

#### Hash Partitioned Table Creation

```
CREATE TABLE hash partitioned (name TEXT)
PARTITION BY HASH (name);
CREATE TABLE hash partition mod 0
PARTITION OF hash partitioned
FOR VALUES WITH (MODULUS 3, REMAINDER 0);
CREATE TABLE hash partition mod 1
PARTITION OF hash partitioned
FOR VALUES WITH (MODULUS 3, REMAINDER 1):
CREATE TABLE hash partition mod 2
PARTITION OF hash partitioned
FOR VALUES WITH (MODULUS 3, REMAINDER 2);
```

#### The Result

#### List Partitioned Table Creation

```
CREATE TYPE employment status type AS ENUM ('employed', 'unemployed', 'retired');
CREATE TABLE list partitioned (
    name TEXT.
    employment status employment status type
PARTITION BY LIST (employment status);
CREATE TABLE list partition employed
PARTITION OF list partitioned
FOR VALUES IN ('employed'):
CREATE TABLE list partition unemployed
PARTITION OF list partitioned
FOR VALUES IN ('unemployed');
-- allow NULL partition key values
CREATE TABLE list partition retired and null
PARTITION OF list partitioned
FOR VALUES IN ('retired', NULL);
```

#### The Result

```
\d+ list_partitioned

Partitioned table "public.list_partitioned"

Column | Type | Collation | Nullable | Default |...

name | text | | | | |...

employment_status | employment_status_type | | |...

Partition key: LIST (employment_status)

Partitions: list_partition_employed FOR VALUES IN ('employed'),

list_partition_retired_and_null FOR VALUES IN ('retired', NULL),

list_partition_unemployed FOR VALUES IN ('unemployed')
```

### Populating the Range Partitioned Table

```
-- This method of generating random data is explained at
-- https://momjian.us/main/blogs/pgblog/2012.html#July 24 2012
INSERT INTO range partitioned
SELECT.
    SELECT initcap(string agg(x, ''))
    FROM (
        SELECT chr(ascii('a') + floor(random() * 26)::integer)
        FROM generate series(1, 2 + (random() * 8)::integer + b * 0)
    ) AS v(x)
FROM generate series(1, 100000) AS a(b);
```

### Populating the Hash Partitioned Table

```
INSERT INTO hash_partitioned
SELECT
(
    SELECT initcap(string_agg(x, ''))
    FROM (
        SELECT chr(ascii('a') + floor(random() * 26)::integer)
        FROM generate_series(1, 2 + (random() * 8)::integer + b * 0)
    ) AS y(x)
)
FROM generate_series(1, 100000) AS a(b);
```

### Populating the List Partitioned Table

```
INSERT INTO list partitioned
SELECT.
    SELECT initcap(string agg(x, ''))
    FROM (
        SELECT chr(ascii('a') + floor(random() * 26)::integer)
        FROM generate series(1, 2 + (random() * 8)::integer + b * 0)
    ) AS y(x)
    SELECT CASE floor(random() * 3 + b * 0)
           WHEN O THEN 'employed'::employment status type
               WHEN 1 THEN 'unemployed'::employment status type
               WHEN 2 THEN 'retired'::employment status type
       END
FROM generate series(1, 100000) AS a(b);
```

### Inserting NULL Values

```
INSERT INTO range_partitioned VALUES (NULL);
INSERT INTO hash_partitioned VALUES (NULL);
INSERT INTO list_partitioned VALUES ('test', NULL);
```

### Creating Indexes

```
CREATE INDEX i_range_partitioned ON range_partitioned (name); CREATE INDEX i_hash_partitioned ON hash_partitioned (name); CREATE INDEX i_list_partitioned ON list_partitioned (name); ANALYZE;
```

#### Where are NULLs Stored?

```
-- NULLs are stored in the DEFAULT range partition.
SELECT *, tableoid::regclass
FROM range partitioned
WHERE name IS NULL:
               tableoid
 (null) | range partition nulls
-- NULLs are always stored in the REMAINDER O hash partition:
-- see src/backend/partitioning/partbounds.c:compute partition hash value()
SELECT *. tableoid::regclass
FROM hash partitioned
WHERE name IS NULL:
               tableoid
 (null) | hash partition mod 0
-- NULLs are stored in the list partition for NULL values.
SELECT *. tableoid::regclass
FROM list partitioned
WHERE employment status IS NULL:
 name | employment status |
 test | (null) | list partition retired and null
```

## First Five Range Partitioned Rows

```
SELECT *, tableoid::regclass
FROM range partitioned
ORDER BY 2, 1
LIMIT 5:
name |
        tableoid
Aa
       range partition less j
       range partition less j
Aa
       range partition less j
Aa
       range partition less j
Aa
       range partition less j
Aa
```

## Random Range Partitioned Rows

```
WITH sample AS
    SELECT *, tableoid::regclass
    FROM range partitioned
    ORDER BY random()
    LIMIT 5
SELECT * FROM sample
ORDER BY 2. 1:
                     tableoid
   name
 Gmgarubn
             range partition less j
 Ousrtai
             range partition j to s
             range partition j to s
 Pbtmufce
             range partition j to s
 Qqt
             range partition s greater
 Ymsapxkxm
```

#### Random Hash Partitioned Rows

```
WITH sample AS
    SELECT *, tableoid::regclass
   FROM hash partitioned
    ORDER BY random()
    LIMIT 5
SELECT * FROM sample
ORDER BY 2. 1:
                  tableoid
   name
 Yxh
            hash partition mod 0
            hash partition mod 1
Asp
           hash partition mod 1
 Bgbvewd
Jemquglx | hash partition mod 2
            hash partition mod 2
 Xtlvuac
```

#### Random List Partitioned Rows

```
WITH sample AS
    SELECT *, tableoid::regclass
   FROM list partitioned
    ORDER BY random()
    LIMIT 5
SELECT * FROM sample
ORDER BY 3, 2, 1;
             employment status
                                             tableoid
    name
              employed
                                  list partition employed
 Twxcn
                                  list partition unemployed
 Bt1faascx
              unemployed
                                  list partition unemployed
 Sz
              unemployed
                                  list partition unemployed
 Xpdi
              unemployed
                                  list partition retired and null
 Uwunkpdhmv
              retired
```

#### Selecting Range Partition Boundaries

```
-- Use lower case for range boundaries if your collation sorts lower case first for case-insensitive equal string,
-- e.g., range 'j' to 's' would include 'j', but 'J' to 'S' would not. 'ja' would be in 'J' to 'S' since 'ja' is
-- not case-insensitive equal to 'J'.
SHOW 1c collate:
1c collate
en US.UTF-8
-- Case ordering ignored because of case-insensitive inequality.
-- https://www.unicode.org/reports/tr10/#Scope
SELECT 'a' < 'J' AND 'J' < 'z':
 ?column?
-----
SELECT 'ja' < 'Jb' AND 'Jc' < 'jd';
?column?
-----
-- Case ordering only honored for case-insensitive equality.
SELECT 'ia' < 'Ja':
?column?
SELECT 'island' < 'Island' AND 'islaNd' < 'iSland';</pre>
 ?column?
------
```

## 3. Partition Pruning



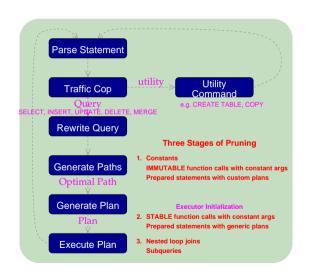
https://www.flickr.com/photos/anguskirk/

### **Pruning Stages**

Pruning eliminates access to unnecessary partitions. There are three possible stages of pruning, earlier ones being more efficient than later ones:

- 1. Optimizer
  - shown by EXPLAIN
- 2. Executor initialization
  - shown by EXPLAIN (ANALYZE) as "Subplans Removed"
  - appears under "Append" and "Merge Append"
  - see src/backend/executor/execPartition.c::ExecInitPartitionPruning()
- 3. Executor running
  - shown by EXPLAIN (ANALYZE) as "never executed"
  - happens during "Append" and "Merge Append"
  - see src/backend/executor/execPartition.c::ExecFindMatchingSubPlans()

#### Pruning Stages Diagram



### Pruning Using NULL Constants: Stage 1

```
\set EXPLAIN 'EXPLAIN (COSTS OFF)'
·FXPLAIN SFLECT *
FROM range partitioned
WHERE name IS NULL;
                     OUERY PLAN
Seq Scan on range partition nulls range partitioned
  Filter: (name IS NULL)
-- NULLs are always stored in the REMAINDER O hash partition.
· FXPLAIN SELECT *
FROM hash partitioned
WHERE name IS NULL:
                                          OUERY PLAN
 Index Only Scan using hash partition mod 0 name idx on hash partition mod 0 hash partitioned
   Index Cond: (name IS NULL)
: FXPLAIN SELECT *
FROM list partitioned
WHERE employment status IS NULL:
                          OUERY PLAN
 Seq Scan on list partition retired and null list partitioned
  Filter: (employment status IS NULL)
```

#### Pruning Using Non-NULL Constants: Stage 1

```
: EXPLAIN SELECT *
FROM range partitioned
WHERE name = 'Ma':
                                            OUERY PLAN
 Index Only Scan using range partition j to s name idx on range partition j to s range partitioned
   Index Cond: (name = 'Ma'::text)
· FXPLAIN SELECT *
FROM hash partitioned
WHERE name = 'Ma':
                                          OUERY PLAN
 Index Only Scan using hash partition mod 2 name idx on hash partition mod 2 hash partitioned
   Index Cond: (name = 'Ma'::text)
:EXPLAIN SELECT *
FROM list partitioned
WHERE employment status = 'retired':
                            OUERY PLAN
 Seg Scan on list partition retired and null list partitioned
  Filter: (employment status = 'retired'::employment status type)
```

### Use of Pruning and Per-Partition Index: Stage 1

```
: EXPLAIN SELECT *
FROM list partitioned
WHERE employment status = 'retired' AND
      name = 'Ma':
                                                  OUFRY PLAN
 Index Scan using list partition retired and null name idx on list partition retired and null list partitioned
   Index Cond: (name = 'Ma'::text)
   Filter: (employment status = 'retired'::employment_status_type)
\d+ list partitioned
                                            Partitioned table "public.list partitioned"
      Column.
                                              Collation | Nullable | Default | Storage | Compression | Stats target | Description
                              Type
                    text
                                                                               extended |
 name
 employment status | employment status type |
                                                                               nlain
Partition key: LIST (employment status)
Indexes:
    "i list partitioned" btree (name)
Partitions: list partition employed FOR VALUES IN ('employed'),
            list_partition_retired_and_null FOR VALUES IN ('retired', NULL),
            list partition unemployed FOR VALUES IN ('unemployed')
```

### Pruning of Custom-Plan Prepared Statements: Stage 1

```
PREPARE part test AS
SELECT *
FROM range partitioned
WHERE name = $1:
:EXPLAIN EXECUTE part test('Ba');
                                           OUERY PLAN
 Index Only Scan using range partition less i name idx on range partition less i range partitioned
   Index Cond: (name = 'Ba'::text)
:EXPLAIN EXECUTE part test('Ma');
                                           QUERY PLAN
 Index Only Scan using range partition j to s name idx on range partition j to s range partitioned
   Index Cond: (name = 'Ma'::text)
:EXPLAIN EXECUTE part test('Ta');
                                               OUFRY PLAN
 Index Only Scan using range partition s greater name idx on range partition s greater range partitioned
   Index Cond: (name = 'Ta'::text)
```

By default, custom plans are used for the first five executions, and then generic plans optionally used.

## Pruning of Generic-Plan Prepared Statements: Stage 2

```
\set EXPLAIN ANALYZE 'EXPLAIN (ANALYZE, SUMMARY OFF, TIMING OFF, COSTS OFF)'
-- force a generic plan
SET plan cache mode TO force generic plan;
-- Prunina happens during executor initialization
:EXPLAIN ANALYZE EXECUTE part test('Ba');
                                                            OHERY PLAN
 Append (actual rows=9 loops=1)
  Subplans Removed: 2
   -> Index Only Scan using range partition less j name idx on range partition less j range partitioned 1 (actual rows=9 loops=1)
         Index Cond: (name = $1)
        Heap Fetches: 0
:EXPLAIN ANALYZE EXECUTE part test('Ma');
                                                            OUERY PLAN
 Append (actual rows=6 loops=1)
   Subplans Removed: 2
   -> Index Only Scan using range partition j to s name idx on range partition j to s range partitioned 1 (actual rows=6 loops=1)
         Index Cond: (name = $1)
        Heap Fetches: 0
:EXPLAIN ANALYZE EXECUTE part test('Ta');
                                                                OUERY PLAN
 Append (actual rows=12 loops=1)
  Subplans Removed: 2
   -> Index Only Scan using range partition s greater name idx on range partition s greater range partitioned 1 (actual rows=12 loops=1)
         Index Cond: (name = $1)
                                                                                                                                     35/98
         Heap Fetches: 0
```

## Pruning of IMMUTABLE Function Calls

-- IMMUTABLE function calls are evaluated in the optimizer.  $\label{eq:calls} \begin{tabular}{ll} \begi$ 

List of operators						
Schema	Name	Left arg type	Right arg type	Result type	Function	Description
	+·	+	+	+	+	+
pg_catalog		anycompatible	anycompatiblearray	anycompatiblearray	array_prepend	prepend element onto front of array
pg_catalog		anycompatiblearray	anycompatible	anycompatiblearray	array_append	append element onto end of array
pg_catalog		anycompatiblearray	anycompatiblearray	anycompatiblearray	array_cat	concatenate
pg_catalog		anynonarray	text	text	anytextcat	concatenate
pg_catalog		bit varying	bit varying	bit varying	bitcat	concatenate
pg_catalog		bytea	bytea	bytea	byteacat	concatenate
pg_catalog		jsonb	jsonb	jsonb	jsonb_concat	concatenate
pg_catalog		text	anynonarray	text	textanycat	concatenate
pg_catalog		text	text	text	textcat	concatenate
pg_catalog		tsquery	tsquery	tsquery	tsquery_or	OR-concatenate
pg_catalog		tsvector	tsvector	tsvector	tsvector_concat	concatenate

# Pruning of IMMUTABLE Function Calls

```
\x on
\df+ textcat
-[ RECORD 1 ]-----+
Schema
                    pg_catalog
Name
                    textcat
Result data type
                    text
Argument data types
                    text, text
Type
                    func
Volatility
                    immutable
Parallel
                    safe
Owner
                    postgres
Security
                    invoker
Access privileges
                    internal
Language
Source code
                    textcat
Description
                    implementation of || operator
```

\x off

# Pruning of IMMUTABLE Function Calls: Stage 1

```
:EXPLAIN_ANALYZE SELECT *, tableoid::regclass

FROM range_partitioned

WHERE name = 'M' || 'a';

QUERY PLAN

Index Scan using range_partition_j_to_s_name_idx on range_partition_j_to_s range_partitioned (actual rows=6 loops=1)

Index Cond: (name = 'Ma'::text)
```

# Pruning of STABLE Function Calls

```
-- STABLE function calls can also cause this.
\x on
\df+ concat
List of functions
-[ RECORD 1 ]-----+
Schema
                     pg catalog
Name
                     concat
Result data type
                     text
Argument data types
                     VARIADIC "any"
Type
                     func
Volatility
                     stable
Parallel
                     safe
Owner
                     postares
Security
                     invoker
Access privileges
                     internal
Language
Source code
                     text concat
                     concatenate values
Description
```

\x off

# Pruning of STABLE Function Calls: Stage 2

```
-- pruning happens during executor initialization
:EXPLAIN_ANALYZE SELECT *, tableoid::regclass
FROM range_partitioned
WHERE name = concat('M', 'a');

QUERY PLAN

Append (actual rows=6 loops=1)
Subplans Removed: 2
-> Index Scan using range_partition_j_to_s_name_idx on range_partition_j_to_s range_partitioned_1 (actual rows=6 loops=1)
Index Cond: (name = concat('M', 'a'))
```

# Pruning of Subqueries: Stage 3

```
CREATE TABLE nested outer (name) AS VALUES ('Pa'), ('Qa'), ('Ra');
ANALYZE nested outer;
-- pruning happens during executor running
:EXPLAIN ANALYZE SELECT *
FROM range partitioned
WHERE name IN (SELECT * FROM nested outer);
                                                              OUFRY PLAN
Nested Loop (actual rows=24 loops=1)
   -> HashAggregate (actual rows=3 loops=1)
         Group Key: nested outer.name
         Batches: 1 Memory Usage: 24kB
         -> Seq Scan on nested outer (actual rows=3 loops=1)
   -> Append (actual rows=8 loops=3)
         -> Index Only Scan using range partition less i name idx on range partition less i range partitioned 1 (never executed)
               Index Cond: (name = nested outer.name)
              Heap Fetches: 0
         -> Index Only Scan using range partition j to s name idx on range partition j to s range partitioned 2 (actual rows=8 loops=3)
               Index Cond: (name = nested outer.name)
              Heap Fetches: 0
         -> Index Only Scan using range partition's greater name idx on range partition s greater range partitioned 3 (never executed)
               Index Cond: (name = nested outer.name)
               Heap Fetches: 0
         -> Seq Scan on range partition nulls range partitioned 4 (never executed)
               Filter: (nested outer.name = name)
```

### Pruning of Joins: Stage 3

```
-- pruning happens during executor running
:EXPLAIN ANALYZE SELECT *
FROM nested outer JOIN range partitioned USING (name);
                                                              OUERY PLAN
Nested Loop (actual rows=24 loops=1)
   -> Seg Scan on nested outer (actual rows=3 loops=1)
   -> Append (actual rows=8 loops=3)
         -> Index Only Scan using range partition less j name idx on range partition less j range partitioned 1 (never executed)
               Index Cond: (name = nested outer.name)
              Heap Fetches: 0
         -> Index Only Scan using range partition j to s name idx on range partition j to s range partitioned 2 (actual rows=8 loops=3)
               Index Cond: (name = nested outer.name)
              Heap Fetches: 0
         -> Index Only Scan using range partition s greater name idx on range partition s greater range partitioned 3 (never executed)
               Index Cond: (name = nested outer.name)
              Heap Fetches: 0
         -> Seq Scan on range partition nulls range partitioned 4 (never executed)
               Filter: (nested outer.name = name)
```

# 4. Partition-Local Optimizations



https://www.flickr.com/photos/donaldjudge/

## Partition-Local Optimizations

Partition-local optimizations perform operations on individual partitions and combine their results, rather than operating only on partitioned tables as a whole. Postgres currently supports such optimizations for aggregates and joins.

Combining partition-local results can be expensive, and therefore these optimizations are disabled by default. Performance testing is recommended before enabling these for production queries.

# Aggregates Without partitionwise\_aggregate: Range

```
-- partitionwise aggregate is disabled by default.
:EXPLAIN ANALYZE SELECT name, COUNT(*)
FROM range partitioned
GROUP BY name:
                                           OUERY PLAN
HashAggregate (actual rows=90608 loops=1)
   Group Key: range partitioned.name
   Batches: 5 Memory Usage: 8241kB Disk Usage: 1552kB
   -> Append (actual rows=100001 loops=1)
         -> Seg Scan on range partition less j range partitioned 1 (actual rows=34626 loops=1)
         -> Seg Scan on range partition j to s range partitioned 2 (actual rows=34535 loops=1)
         -> Seq Scan on range partition s greater range partitioned 3 (actual rows=30839 loops=1)
         -> Seq Scan on range partition nulls range partitioned 4 (actual rows=1 loops=1)
```

# Aggregates Without partitionwise\_aggregate: Hash

# Aggregates Without partitionwise\_aggregate: List

:EXPLAIN ANALYZE SELECT employment status, COUNT(\*)

```
FROM list_partitioned
GROUP BY employment_status;

QUERY PLAN

HashAggregate (actual rows=4 loops=1)
Group Key: list_partitioned.employment_status
Batches: 1 Memory Usage: 24kB

-> Append (actual rows=100001 loops=1)

-> Seq Scan on list_partition_employed list_partitioned_1 (actual rows=33027 loops=1)

-> Seq Scan on list_partition_unemployed list_partitioned_2 (actual rows=33467 loops=1)

-> Seq Scan on list_partition_retired and null list_partitioned_3 (actual rows=33507 loops=1)
```

# Aggregates With partitionwise\_aggregate: Range

```
SET enable partitionwise aggregate = true;
-- needed because the cost of combining per-partition HashAggregates results
-- with many distinct values is high
SET cpu tuple cost = 0;
:EXPLAIN ANALYZE SELECT name, COUNT(*)
FROM range partitioned
GROUP BY name:
                                           OUERY PLAN
 Append (actual rows=90608 loops=1)
   -> HashAggregate (actual rows=31382 loops=1)
         Group Key: range partitioned.name
         Batches: 1 Memory Usage: 4113kB
         -> Seg Scan on range partition less i range partitioned (actual rows=34626 loops=1)
   -> HashAggregate (actual rows=31256 loops=1)
         Group Key: range partitioned 1.name
         Batches: 1 Memory Usage: 4113kB
         -> Seg Scan on range partition j to s range partitioned 1 (actual rows=34535 loops=1)
   -> HashAggregate (actual rows=27969 loops=1)
         Group Key: range partitioned 2.name
         Batches: 1 Memory Usage: 3857kB
         -> Seq Scan on range partition s greater range partitioned 2 (actual rows=30839 loops=1)
   -> HashAggregate (actual rows=1 loops=1)
         Group Key: range partitioned 3.name
         Batches: 1 Memory Usage: 24kB
         -> Seg Scan on range partition nulls range partitioned 3 (actual rows=1 loops=1)
```

# Aggregates With partitionwise\_aggregate: Hash

```
:EXPLAIN ANALYZE SELECT name, COUNT(*)
FROM hash partitioned
GROUP BY name:
                                        OUFRY PLAN
Append (actual rows=90645 loops=1)
  -> HashAggregate (actual rows=30465 loops=1)
        Group Key: hash partitioned.name
        Batches: 1 Memory Usage: 4113kB
         -> Seq Scan on hash partition mod 0 hash partitioned (actual rows=33530 loops=1)
  -> HashAggregate (actual rows=29868 loops=1)
        Group Key: hash partitioned 1.name
        Batches: 1 Memory Usage: 4113kB
        -> Seg Scan on hash partition mod 1 hash partitioned 1 (actual rows=32979 loops=1)
  -> HashAggregate (actual rows=30312 loops=1)
        Group Key: hash partitioned 2.name
        Batches: 1 Memory Usage: 4113kB
        -> Seg Scan on hash partition mod 2 hash partitioned 2 (actual rows=33492 loops=1)
```

# Aggregates With partitionwise\_aggregate: List

```
-- not needed for the next query because few distinct values
RESET cpu tuple cost;
:EXPLAIN ANALYZE SELECT employment status, COUNT(*)
FROM list partitioned
GROUP BY employment status;
                                              OUERY PLAN
 Append (actual rows=4 loops=1)
   -> HashAggregate (actual rows=1 loops=1)
         Group Key: list partitioned.employment status
         Batches: 1 Memory Usage: 24kB
         -> Seg Scan on list partition employed list partitioned (actual rows=33027 loops=1)
   -> HashAggregate (actual rows=1 loops=1)
         Group Key: list partitioned 1.employment status
         Batches: 1 Memory Usage: 24kB
         -> Seg Scan on list partition unemployed list partitioned 1 (actual rows=33467 loops=1)
   -> HashAggregate (actual rows=2 loops=1)
         Group Key: list partitioned 2.employment status
         Batches: 1 Memory Usage: 24kB
         -> Seq Scan on list partition retired and null list partitioned 2 (actual rows=33507 loops=1)
RESET enable partitionwise aggregate;
```

# Cross Partition Join: Setup

```
CREATE TABLE range partitioned2 (name TEXT)
PARTITION BY RANGE (name):
CREATE TABLE range partition less j2
PARTITION OF range partitioned2
FOR VALUES FROM (MINVALUE) TO ('.j');
CREATE TABLE range partition j to s2
PARTITION OF range partitioned2
FOR VALUES FROM ('i') TO ('s');
CREATE TABLE range partition s greater2
PARTITION OF range partitioned2
FOR VALUES FROM ('s') TO (MAXVALUE):
CREATE TABLE range partition nulls2
PARTITION OF range partitioned2
(CHECK (name IS NULL))
DEFAULT:
```

## Cross Partition Join: Populate

# Cross Partition Join Without partitionwise\_join

```
-- partitionwise aggregate is disabled by default.
:EXPLAIN ANALYZE SELECT *
FROM range partitioned JOIN range partitioned2 USING (name);
                                               OUFRY PLAN
Hash Join (actual rows=67468 loops=1)
  Hash Cond: (range partitioned.name = range partitioned2.name)
  -> Append (actual rows=100001 loops=1)
         -> Seq Scan on range partition less j range partitioned 1 (actual rows=34626 loops=1)
         -> Seg Scan on range partition j to s range partitioned 2 (actual rows=34535 loops=1)
         -> Seg Scan on range partition s greater range partitioned 3 (actual rows=30839 loops=1)
         -> Seg Scan on range partition nulls range partitioned 4 (actual rows=1 loops=1)
  -> Hash (actual rows=100000 loops=1)
        Buckets: 131072 Batches: 1 Memory Usage: 4833kB
         -> Append (actual rows=100000 loops=1)
               -> Seq Scan on range partition less j2 range partitioned2 1 (actual rows=34646 loops=1)
              -> Seg Scan on range partition j_to_s2 range_partitioned2_2 (actual rows=34606 loops=1)
              -> Seg Scan on range partition's greater2 range partitioned2 3 (actual rows=30748 loops=1)
              -> Seq Scan on range partition nulls2 range partitioned2 4 (actual rows=0 loops=1)
```

# Cross Partition Join With partitionwise\_join

```
SET enable partitionwise join = true;
: EXPLAIN ANALYZE SELECT *
FROM range partitioned JOIN range partitioned2 USING (name);
                                                          QUERY PLAN
Append (actual rows=67468 loops=1)
  -> Hash Join (actual rows=23917 loops=1)
        Hash Cond: (range partitioned 1.name = range partitioned2 1.name)
         -> Seg Scan on range partition less i range partitioned 1 (actual rows=34626 loops=1)
        -> Hash (actual rows=34646 loops=1)
              Buckets: 65536 Batches: 1 Memory Usage: 1832kB
              -> Seg Scan on range partition less i2 range partitioned2 1 (actual rows=34646 loops=1)
  -> Hash Join (actual rows=23096 loops=1)
        Hash Cond: (range partitioned 2.name = range partitioned2 2.name)
        -> Seq Scan on range partition j to s range partitioned 2 (actual rows=34535 loops=1)
        -> Hash (actual rows=34606 loops=1)
              Buckets: 65536 Batches: 1 Memory Usage: 1830kB
              -> Seg Scan on range partition i to s2 range partitioned2 2 (actual rows=34606 loops=1)
  -> Hash Join (actual rows=20455 loops=1)
        Hash Cond: (range partitioned 3.name = range partitioned2 3.name)
        -> Seg Scan on range partition s greater range partitioned 3 (actual rows=30839 loops=1)
        -> Hash (actual rows=30748 loops=1)
              Buckets: 32768 Batches: 1 Memory Usage: 1428kB
               -> Seq Scan on range partition s greater2 range partitioned2 3 (actual rows=30748 loops=1)
  -> Nested Loop (actual rows=0 loops=1)
        -> Seq Scan on range partition nulls2 range partitioned2 4 (actual rows=0 loops=1)
        -> Index Only Scan using range partition nulls name idx on range partition nulls range partitioned 4 (never executed)
               Index Cond: (name = range partitioned2 4.name)
              Heap Fetches: 0
```

# 5. Time-Based Partitioning



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# Time-Based Partitioning

Time-based partitioning stores data in partitions based on some time component. New partitions need to be created to match time-based requirements, and old partitions can be archived, deleted, or moved to tablespaces with cheaper storage. Time-based partitioning have added complexity because of time zone and daylight saving time aspects.

#### Create Per-Month Partitions

```
CREATE TABLE month_partitioned (day DATE, temperature NUMERIC(5,2))
PARTITION BY RANGE (day):
CREATE TABLE month partition 2023 01
PARTITION OF month partitioned
FOR VALUES FROM ('2023-01-01') TO ('2023-02-01');
CREATE TABLE month partition 2023 02
PARTITION OF month partitioned
FOR VALUES FROM ('2023-02-01') TO ('2023-03-01');
CREATE TABLE month partition 2023 03
PARTITION OF month partitioned
FOR VALUES FROM ('2023-03-01') TO ('2023-04-01'):
CREATE TABLE month partition other
PARTITION OF month partitioned
DEFAULT:
```

### Populate Per-Month Partitions

```
INSERT INTO month partitioned
SELECT
    SELECT '2023-01-01'::date +
       floor(random() * ('2023-04-01'::date - '2023-01-01'::date) + b * 0)::integer
),
   SELECT floor(random() * 10000) / 100 + b * 0
FROM generate series(1, 100000) AS a(b);
CREATE INDEX i month partitioned ON month partitioned (day);
ANALYZE:
```

#### Random Partition Rows

```
WITH sample AS
    SELECT *, tableoid::regclass
    FROM month partitioned
    ORDER BY random()
    LIMIT 5
SELECT * FROM sample
ORDER BY 3. 1:
              temperature |
                                   tableoid
    dav
 2023-01-16
                    64.45
                            month partition 2023 01
                            month partition 2023 02
 2023-02-07
                    46.03
                            month partition 2023 02
 2023-02-13
                   87.14
                            month partition 2023 03
 2023-03-07
                    97.04
                            month partition 2023 03
 2023-03-08
                    59.02
```

# Partition Pruning

```
:EXPLAIN_ANALYZE SELECT *
FROM month_partitioned
WHERE day = '2023-02-01';

QUERY PLAN

Bitmap Heap Scan on month_partition_2023_02 month_partitioned (actual rows=1169 loops=1)
Recheck Cond: (day = '2023-02-01'::date)
Heap Blocks: exact=170

-> Bitmap Index Scan on month_partition_2023_02_day_idx (actual rows=1169 loops=1)
Index Cond: (day = '2023-02-01'::date)
```

## Default Partition Usage

#### Null Uses the DEFAULT Partition

```
INSERT INTO month_partitioned VALUES (NULL, 46.24);

:EXPLAIN_ANALYZE SELECT *
FROM month_partitioned
WHERE day IS NULL;

QUERY PLAN

Seq Scan on month_partition_other month_partitioned (actual rows=1 loops=1)
Filter: (day IS NULL)
Rows Removed by Filter: 1
```

## Attaching and Detaching Partitions

```
ALTER TABLE month_partitioned DETACH PARTITION month_partition_other;

INSERT INTO month_partitioned VALUES (NULL, 46.24);

ERROR: no partition of relation "month_partitioned" found for row

DETAIL: Partition key of the failing row contains (day) = (null).

ALTER TABLE month_partitioned ATTACH PARTITION month_partition_other DEFAULT;
```

### Pruning Using a STABLE Function

```
-- Simulate CURRENT DATE by using the STABLE function concat().
\set CURRENT DATE concat('''2023-02-01''' || ''''')::date
:EXPLAIN ANALYZE SELECT *
FROM month partitioned
WHERE day = : CURRENT DATE;
                                            OUERY PLAN
Append (actual rows=1169 loops=1)
   Subplans Removed: 3
   -> Bitmap Heap Scan on month partition 2023 02 month partitioned 1 (actual rows=1169 loops=1)
         Recheck Cond: (day = (concat('2023-02-01'::text))::date)
         Heap Blocks: exact=170
         -> Bitmap Index Scan on month partition 2023 02 day idx (actual rows=1169 loops=1)
               Index Cond: (day = (concat('2023-02-01'::text))::date)
```

# Partition Expiration and Creation

```
CREATE TABLE month_partition_2023_04
PARTITION OF month_partitioned
FOR VALUES FROM ('2023-04-01') TO ('2023-05-01');
DROP TABLE month_partition_2023_01;
```

The table could also be archived before deletion or moved to a different tablespace. pg\_partman (https://github.com/pgpartman/pg\_partman) can help with auto-partition creation.

# Create Timestamp with Time Zone Partitions

```
SET timezone = 'America/New York';
CREATE TABLE month ts tz partitioned (event time TIMESTAMP WITH TIME ZONE, temperature NUMERIC(5,2))
PARTITION BY RANGE (event time);
-- date evaluated at creation time
CREATE TABLE month ts tz partition 2023 01
PARTITION OF month ts tz partitioned
FOR VALUES FROM ('2023-01-01') TO ('2023-02-01');
CREATE TABLE month ts tz partition 2023 02
PARTITION OF month ts tz partitioned
FOR VALUES FROM ('2023-02-01') TO ('2023-03-01'):
CREATE TABLE month ts tz partition 2023 03
PARTITION OF month ts tz partitioned
FOR VALUES FROM ('2023-03-01') TO ('2023-04-01'):
CREATE TABLE month ts tz partition other
PARTITION OF month ts tz partitioned
DEFAULT:
```

### DATE Data Type Has No Time Zone

```
SELECT EXTRACT(EPOCH FROM '2023-01-01'::date);
 extract
1672531200
SELECT EXTRACT(EPOCH FROM '2023-01-01 00:00:00-00'::timestamptz);
     extract
 1672531200,000000
SELECT EXTRACT(EPOCH FROM '2023-01-01'::timestamptz);
     extract
 1672549200.000000
SELECT EXTRACT(EPOCH FROM '2023-01-01 00:00:00-05'::timestamptz);
     extract
 1672549200.000000
```

# Populate Partitions

```
INSERT INTO month ts tz partitioned
SELECT
    SELECT '2023-01-01 00:00:00'::timestamptz +
       (floor(random() *
              (extract(EPOCH FROM '2023-04-01'::timestamptz) -
               extract(EPOCH FROM '2023-01-01'::timestamptz)) +
              b * 0)::integer || 'seconds')::interval
    SELECT floor(random() * 10000) / 100 + b * 0
FROM generate series(1, 100000) AS a(b);
-- add row to the DEFAULT partition
INSERT INTO month ts tz partitioned VALUES ('2023-04-05 00:00:00', 50);
CREATE INDEX i month ts tz partitioned ON month ts tz partitioned (event time);
ANALYZE:
```

We must cast to TIMESTAMPTZ, not DATE, to align with the partition boundaries.

#### Partition Details

#### First Partition Row

```
SELECT CURRENT TIMESTAMP;
      current timestamp
2023-03-20 09:35:05.375191-04
SELECT *, tableoid::regclass
FROM month ts tz partitioned
ORDER BY 1
LIMIT 1:
      event time | temperature | tableoid
2023-01-01 00:01:26-05 | 8.14 | month ts tz partition 2023 01
```

#### First Partition Row in a Different Time Zone

```
SET timezone = 'Asia/Tokyo';
SELECT CURRENT TIMESTAMP;
      current timestamp
 2023-03-20 22:35:05.376393+09
SELECT *. tableoid::regclass
FROM month ts tz partitioned
ORDER BY 1
LIMIT 1:
      event time
                       | temperature |
                                               tableoid
2023-01-01 14:01:26+09 | 8.14 | month ts tz partition 2023 01
```

## Partition Bounds Adjusted

### The Same in UTC

#### The Same in UTC

# No Pruning of a Function Call on a Column

```
SET timezone = 'America/New York';
:EXPLAIN ANALYZE SELECT *
FROM month ts tz partitioned
WHERE date(event time) = '2023-02-05';
                                                 QUERY PLAN
Append (actual rows=1093 loops=1)
   -> Seq Scan on month ts tz partition 2023 01 month ts tz partitioned 1 (actual rows=0 loops=1)
         Filter: (date(event time) = '2023-02-05'::date)
         Rows Removed by Filter: 34240
   -> Seq Scan on month ts tz partition 2023 02 month ts tz partitioned 2 (actual rows=1093 loops=1)
         Filter: (date(event time) = '2023-02-05'::date)
         Rows Removed by Filter: 30075
   -> Seq Scan on month ts tz partition 2023 03 month ts tz partitioned 3 (actual rows=0 loops=1)
         Filter: (date(event time) = '2023-02-05'::date)
         Rows Removed by Filter: 34547
   -> Seq Scan on month ts tz partition other month ts tz partitioned 4 (actual rows=0 loops=1)
         Filter: (date(event time) = '2023-02-05'::date)
         Rows Removed by Filter: 45
```

# Date Range Can Be Pruned

### Date Calculation Can Be Pruned

```
-- simulate CURRENT TIMESTAMP by using the STABLE function concat().
\set CURRENT TIMESTAMP concat('''2023-02-05 23:43:51''' || ''''')::timestamptz
-- pruning happening during executor initialization
:EXPLAIN ANALYZE SELECT *
FROM month ts tz partitioned
WHERE event time > :CURRENT TIMESTAMP - '24 hours'::interval AND
     event time <= :CURRENT TIMESTAMP;</pre>
                                                              OUERY PLAN
 Append (actual rows=1091 loops=1)
   Subplans Removed: 3
   -> Bitmap Heap Scan on month ts tz partition 2023 02 month ts tz partitioned 1 (actual rows=1091 loops=1)
         Recheck Cond: ((event time > (concat('2023-02-05 23:43:51'::text))::timestamp with time zone - '24:00:00'::interval)) AND
                        (event_time <= (concat('2023-02-05 23:43:51'::text))::timestamp with time zone))
         Heap Blocks: exact=169
         -> Bitmap Index Scan on month ts tz partition 2023 02 event time idx (actual rows=1091 loops=1)
               Index Cond: ((event time > ((concat('2023-02-05 23:43:51'::text))::timestamp with time zone - '24:00:00'::interval)) AND
                            (event time <= (concat('2023-02-05 23:43:51'::text))::timestamp with time zone))
```

# Timestamp Range Can Be Pruned

# Where Are Per-Day Rows?

# Where Are Per-Day Rows?

```
SELECT *, tableoid::regclass
FROM month ts tz partitioned
WHERE event time >= '2023-03-01 00:00:00' AND
    event time < '2023-03-02 00:00:00'
ORDER BY 1
LIMIT 1;
2023-03-01 00:00:20-05 | 49.85 | month ts tz partition 2023 03
SELECT *, tableoid::regclass
FROM month ts tz partitioned
WHERE event time >= '2023-03-01 00:00:00' AND
    event time < '2023-03-02 00:00:00'
ORDER BY 1 DESC
LIMIT 1:
     event_time | temperature | tableoid
2023-03-01 23:59:35-05 | 27.94 | month ts tz partition 2023 03
```

# Query in a Different Time Zone

```
SET timezone = 'Asia/Tokyo':
: FXPLAIN ANALYZE SELECT *
FROM month ts tz partitioned
WHERE event time >= '2023-03-01 00:00:00' AND
      event time < '2023-03-02 00:00:00';
                                                    OUERY PLAN
 Append (actual rows=1109 loops=1)
   -> Bitmap Heap Scan on month ts tz partition 2023 02 month ts tz partitioned 1 (actual rows=622 loops=1)
         Recheck Cond: ((event time >= '2023-03-01 00:00:00+09'::timestamp with time zone) AND
                        (event time < '2023-03-02 00:00:00+09'::timestamp with time zone))
         Heap Blocks: exact=166
         -> Bitmap Index Scan on month ts tz partition 2023 02 event time idx (actual rows=622 loops=1)
               Index Cond: ((event time >= '2023-03-01 00:00:00+09'::timestamp with time zone) AND
                            (event_time < '2023-03-02 00:00:00+09'::timestamp with time zone))
   -> Bitmap Heap Scan on month ts tz partition 2023 03 month ts tz partitioned 2 (actual rows=487 loops=1)
         Recheck Cond: ((event time >= '2023-03-01 00:00:00+09'::timestamp with time zone) AND
                        (event time < '2023-03-02 00:00:00+09'::timestamp with time zone))
         Heap Blocks: exact=169
         -> Bitmap Index Scan on month ts tz partition 2023 03 event time idx (actual rows=487 loops=1)
               Index Cond: ((event time >= '2023-03-01 00:00:00+09'::timestamp with time zone) AND
                            (event time < '2023-03-02 00:00:00+09'::timestamp with time zone))
```

## Different Count

# Partition Range

```
SELECT *, tableoid::regclass
FROM month ts tz partitioned
WHERE event time >= '2023-03-01 00:00:00' AND
    event time < '2023-03-02 00:00:00'
ORDER BY 1
LIMIT 1;
2023-03-01 00:00:32+09 | 16.32 | month ts tz partition 2023 02
SELECT *, tableoid::regclass
FROM month ts tz partitioned
WHERE event time >= '2023-03-01 00:00:00' AND
     event time < '2023-03-02 00:00:00'
ORDER BY 1 DESC
LIMIT 1:
     event_time | temperature | tableoid
2023-03-01 23:56:22+09 | 19.78 | month ts tz partition 2023 03
```

# Range Boundaries Are Set at Creation

```
-- caused by mismatch with America/New_York time zone boundaries
CREATE TABLE month_ts_tz_partition_2023_04
PARTITION OF month_ts_tz_partitioned
FOR VALUES FROM ('2023-04-01') TO ('2023-05-01');
ERROR: partition "month_ts_tz_partition_2023_04" would
    overlap partition "month_ts_tz_partition_2023_03"
LINE 3: FOR VALUES FROM ('2023-04-01') TO ('2023-05-01');
```

# Matching Rows in the DEFAULT Partition

```
SET timezone = 'America/New_York';

-- caused by daylight saving time change

CREATE TABLE month_ts_tz_partition_2023_04

PARTITION OF month_ts_tz_partitioned

FOR VALUES FROM ('2023-04-01') TO ('2023-05-01');

ERROR: updated partition constraint for default partition

"month_ts_tz_partition_other" would be violated by some row
```

### Move DEFAULT Rows to a New Partition

```
BEGIN WORK:
-- lock table and/or detach DEFAULT partition?
CREATE TEMP TABLE tmp default AS
SELECT *
FROM month ts tz partition other
WHERE event time >= '2023-04-01 00:00:00' AND
      event_time < '2023-05-01 00:00:00':
DELETE FROM month ts tz partition other
WHERE event time >= '2023-04-01 00:00:00' AND
      event_time < '2023-05-01 00:00:00':
CREATE TABLE month ts tz partition 2023 04
PARTITION OF month_ts_tz_partitioned
FOR VALUES FROM ('2023-04-01') TO ('2023-05-01'):
INSERT INTO month ts tz partitioned
SELECT * FROM tmp default;
SELECT * FROM month ts tz partition other;
 event time | temperature
COMMIT:
```

### **New Partition Contents**

```
SELECT * FROM month_ts_tz_partition_2023_04;
event_time | temperature

2023-04-05 00:00:00-04 | 50.00

SELECT * FROM month_ts_tz_partition_other;
event_time | temperature
```

# 6. Row Migration



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# Rows in 'j' to 's' Partition

```
SELECT *, tableoid::regclass
FROM range partitioned
WHERE name = 'Ma'
ORDER BY 2, 1;
 name |
               tableoid
Ma
      range partition j to s
 Ma
        range partition j to s
Ma
        range partition j to s
 Ma
        range partition j to s
Ma
        range partition j to s
 Ma
        range partition j to s
```

# Migration to Greater than 's' Partition

```
UPDATE range partitioned
SET name = 'zz ' || name
WHERE name = 'Ma';
SELECT *, tableoid::regclass
FROM range partitioned
WHERE name = 'zz Ma'
ORDER BY 2, 1;
 name
          tableoid
 zz Ma | range partition s greater
 zz Ma | range partition s greater
zz Ma | range partition s greater
zz Ma | range partition s greater
 zz_Ma | range_partition_s_greater
 zz Ma
         range partition s greater
```

# 8. psql Support



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## psql support

```
COMMENT ON TABLE range_partitioned IS 'Section 2';
COMMENT ON TABLE hash_partitioned IS 'Section 2';
COMMENT ON TABLE list_partitioned IS 'Section 2';
COMMENT ON TABLE range_partitioned2 IS 'Section 4';
COMMENT ON TABLE month_partitioned IS 'Section 5';
COMMENT ON TABLE month_ts_tz_partitioned IS 'Section 5';
```

#### \dPt+

#### List of partitioned tables

Schema			Total size	•
public   public   public		+   postgres   postgres   postgres   postgres	3888 kB   4280 kB   2896 kB	Section 2 Section 2 Section 5 Section 5
public	range_partitioned range_partitioned	postgres   postgres   postgres	3904 kB	Section 2 Section 4

## 8. Limitations



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# Partitioning Limitations

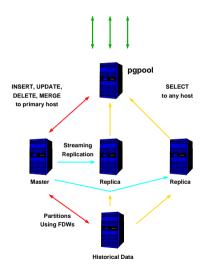
- Adding partitions
  - requires locking
    - partitions can be created, populated, and then attached to reduce locking
  - attached partition rows are verified unless CHECK constraints match partition bounds
  - requires DEFAULT partition scans unless CHECK constraints make it unnecessary
  - foreign table rows are not verified
- Removing partitions
  - requires locking
  - partitions can be detached CONCURRENTLY to reduce locking
- Function calls on columns cannot be pruned unless specified in the range
  - function calls on constants often can
- Partitioned table indexes must start with the partition columns
  - no global indexes, see https://momjian.us/main/blogs/pgblog/2020.html# July 1 2020
- CONCURRENT index creation on partitioned tables is not supported
  - CONCURRENT index creation on partitions is supported, and can then be attached

# 9. Complex Architectures

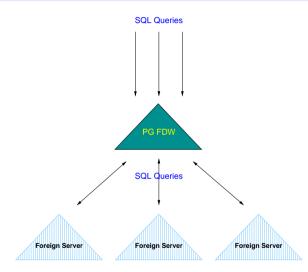


https://www.flickr.com/photos/alexdrop/

# Foreign Servers for Archive Data



# Sharding



https://momjian.us/main/presentations/performance.html#sharding

## Conclusion



