A Tour of PostgreSQL Data Types

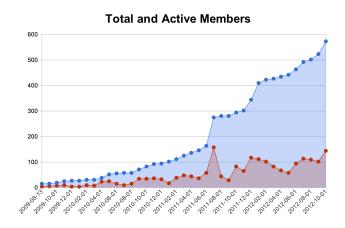
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Jim "Jimbo" Mlodgenski
PGCon 2013 – May 21, 2013

Who We Are

- Jonathan S. Katz
 - CTO, VenueBook
 - Co-Organizer NYC PostgreSQL User Group (NYCPUG)
- Jim Mlodgenski
 - CEO, StormDB
 - Co-Organizer NYCPUG

A Brief Note on NYCPUG

- Active since 2010
- 700 members
- Monthly Meetups
- PGDay NYC 2013
 - March 22
 - 100 attendees
- Part of PG.US
- PGConf NYC 2014



Why Data Types

- Fundamental
 - -0 => 1
 - -00001111
- Building Blocks
 - -0x41424344
- Accessibility
 - -1094861636
 - 'ABCD'

Why Data Types

- Primitive Data Types
 - Integers, floating points, booleans, characters
- Primitive Data Structures
 - Strings, arrays, linked lists, hash tables
- Data Structures++
 - Classes, structs, trees, matrices

Data Storage

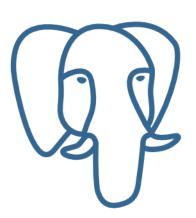
"Persistence of Memory"

Data Access and Retrieval

- Recall what we have stored
- Represent as it originally was
- Interface between disk <=> application

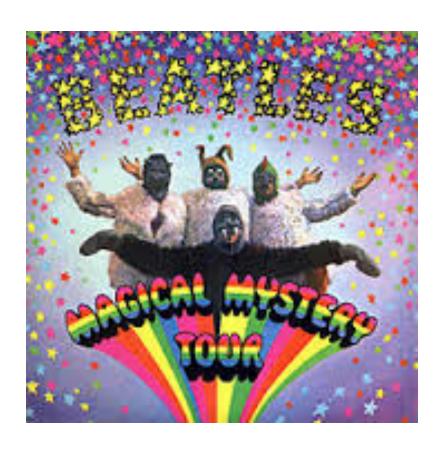
PostgreSQL

- Roots from "INGRES"
 - Image storage
- Data integrity = foremost concern
- Data representation
- Robustness
- Performance



The PostgreSQL Data Type Tour

- Data types
- Functions
- Features
- Indexes
- Use cases
- Extensions



The PostgreSQL Data Type Tour

- Assumptions
 - PostgreSQL 9.2+
 - Some looks at PostgreSQL 9.3beta1

Number Types

Name	Storage Size	Range
smallint	2 bytes	-32768 to +32767
integer	4 bytes	-2147483648 to +2147483647
bigint	4 bytes	-9223372036854775808 to 9223372036854775807
decimal	variable	up to 131072 digits before the decimal point; up to 16383 digits after the decimal point
numeric	variable	up to 131072 digits before the decimal point; up to 16383 digits after the decimal point
real	4 bytes	6 decimal digits precision
double	8 bytes	15 decimal digits precision

Integers

- smallint
 - Use only if disk space is a premium, e.g.
 embedded devices
- bigint
 - Slower than int
- int
 - For everything else...

numeric

- numeric
 - Provides scale and precision
 - *Scale* count of decimal places
 - 987.123456 has a scale of 6
 - *Precision* total count of significant digits
 - 987.123456 has a precision of 9
 - Declarations
 - numeric(precision, scale)
 - max declarable is (1000, 100)
 - numeric(precision)
 - essentially an integer
 - Numeric
 - Precision & scale up to limit (147455, 16383)
 - Contrary to SQL standard

numeric & NYC Sales Tax

```
SELECT 100 * (0.08875)::numeric;
8.875
SELECT 100 * (0.08875)::numeric(7,2);
9.0
SELECT (100 * 0.08875)::numeric(7,2);
8.88
```

numeric

- Storage
 - Determined by size of numeric type, no padding

Numbers – numeric

- 'NaN'
- decimal equivalent to numeric

Numbers – floating point

- IEEE 754
- Inexact
 - Unexpected behavior may occur
 - Overflow/underflow
 - Equality
- Constants
 - 'Nan', 'Infinity', '-Infinity'
- Types
 - real => 1E-37 <=> 1E+37
 - double precision => 1E-308 <=> 1E+308
 - float(1) <=> float(24) = real
 - float(25) <=> float(53) = double precision

numeric vs double precision

```
\timing
CREATE TABLE floats (x double precision);
CREATE TABLE numerics (x numeric(15, 15));
INSERT INTO floats
SELECT random() FROM generate series(1,1000000);
INSERT INTO numerics
SELECT * FROM floats;
CREATE INDEX floats idx ON floats (x);
CREATE INDEX numerics idx ON numerics (x);
SELECT * FROM floats WHERE x \ge 0.7;
-- avg 280ms
SELECT * FROM numerics WHERE x \ge 0.7;
-- avg 120ms
```

numeric vs floating points

- generally it is better to use numeric
- floating point usage is application specific
 - reading data from a thermometer
 - IEEE 754 specific programs
 - too many rows for larger numeric data type
 - do not require precision
- understand ramifications before making choice

Number Functions

- ceil/ceiling, floor
- exp (exponential), In, log
- greatest, least
- random, setseed
- round, truncate
- sign
- •to_number
- degrees(radians), radians(degrees)
- cos, acos, sin, asin
- cot (cotangent), tan, atan
- •atan2(x, y) = atan(x/y)

Serial Types

Name	Storage Size	Range
smallserial	2 bytes	1 to 32767
serial	4 bytes	1 to 2147483647
bigserial	8 bytes	1 to 9223372036854775807

Serial "Types"

Not truly a data type, but a convenience

```
CREATE TABLE awesome (
   id serial
);

or

CREATE SEQUENCE awesome_id_seq;
CREATE TABLE awesome (
   id integer NOT NULL DEFAULT nextval('awesome_id_seq')
);
ALTER SEQUENCE awesome_colname_seq OWNED BY awesome.id;
```

Serial Functions

- nextval
 - advances sequence and returns new value

```
SELECT nextval('sequence name');
```

- setval
 - sets the current value of the sequence

```
SELECT setval('sequence_name', 2); -- nextval returns 3
SELECT setval('sequence_name', 2, true); -- nextval returns 3
SELECT setval('sequence name', 2, false); -- nextval returns 2
```

- currval
 - returns current value of sequence if sequence has been manipulated in session
 SELECT currval('sequence name');
- lastval
 - returns current value of last sequence that has been manipulated in session SELECT lastval();

Monetary Types

Name	Storage Size	Range
money	8 bytes	-92233720368547758.08 to 92233720368547758.07

Monetary Types: The Story

- Stores monetary amounts with precision based on 'lc_monetary' setting
- Output based on lc_monetary
 - '\$1,000.00'

Monetary Types: The Reality

- Don't use it
- Store money as
 - integer family of types
 - numeric

Character Types (or Strings)

Name	Description
varchar(n)	variable-length with limit
char(n)	fixed-length, blank padded
text	variable unlimited length

Character Types

- char(n) and varchar(n) mostly follow the ANSI standard
 - Will throw an error if given a string longer than n characters (not bytes)
 - Trailing spaces in char(n) are ignored in char(n) comparisons,
 and stripped when converting to other string types
 - Unlike many databases, char(n) is NOT stored as a fixed-sized field in Postgres. It is treated exactly the same as varchar(n) except for being padded

Character Types

- "varlena"
 - Called internally when creating any character type
- text
 - Preferred type in practice
 - Max ~1GB
- varchar(n)
 - Use only when you have to restrict length
 - CPU overhead (marginal)
- char(n)
 - avoid
 - unexpected behavior e.g. with "LIKE" expressions

Character Types & Encoding

 What do encoding, cache management, and concurrency all have in common?

String Functions

ascii

•bit_length

btrim

•char_length

•chr

concat

convert

decode

encode

initcap

length

lower

• Ipad

•ltrim

• md5

octet_length

overlay

pg_client_encoding

position

•quote_ident

quote_literal

repeat

er rpad

·rtrim

.split_part

·strpos

·substr

 $\cdot substring \\$

∙to_ascii

∙to_hex

·translate

·trim

·upper

replace

Binary Data Types

Name	Storage Size	Description
bytea	1 to 4 bytes plus size of binary string	variable-length binary string

Binary Data Types

- Used to store "raw bytes"
- Different output formats:
 - Pre–9.0: PostgreSQL "escape"
 - -9.0+: hex
 - 'bytea_output' config parameter to choose (default: 'hex')

Binary Data Types

- Should probably not store raw binary data in PostgreSQL
- If you must, keep in its own table and JOIN when needed

Date / Time Types

- PostgreSQL second to none
- timestamp with time zone
- timestamp without time zone
- date
- time with time zone
- time without time zone
- interval

Date / Time Types

PostgreSQL – second to none

Name	Size	Range	Resolution
timestamp without timezone	8 bytes	4713 BC to 294276 AD	1 microsecond / 14 digits
timestamp with timezone	8 bytes	4713 BC to 294276 AD	1 microsecond / 14 digits
date	4 bytes	4713 BC to 5874897 AD	1 day
time without timezone	8 bytes	00:00:00 to 24:00:00	1 microsecond / 14 digits
time with timezone	12 bytes	00:00:00+1459 to 24:00:00-1459	1 microsecond / 14 digits
interval	12 bytes	-178000000 years to 178000000 years	1 microsecond / 14 digits

Date / Time General Notes

- timestamp = timestamp without time zone
- timestamptz = timestamp with time zone
- time, timestamp, and interval have optional argument "p"
 - "precision" number of fractional digits
 - p <@ [0,6], default is 6</p>
 - SELECT CURRENT_TIMESTAMP::time(4);
- interval has other storage options (more later)

Date / Time Input

- PostgreSQL is very flexible
 - ISO 8601
 - SQL
 - POSTGRES
 - and more
- Day / Month / Year ordering "datestyle" parameter
 - datestyle = 'iso, mdy'
 - DMY, MDY, YMD

Date Input

Input	Description
2013-05-03	ISO 8601, May 3 with any datestyle (recommended format)
May 3, 2013	May 3 with any datestyle
5/3/2013	May 3 with MDY, March 5 with DMY
5/21/2013	May 21 with MDY, rejected with other formats
5/2/3	May 2, 2003 with MDY, February 5 2003 with DMY, February 3, 2005 with YMD
2013-May-3 May-3-2013 3-May-2013	All equivalent with any datestyle
32-May-3	May 3 with YMD, otherwise rejected
3-May-32	rejected with YMD, otherwise May 3
May-3-32	rejected with YMD, otherwise May 3
20130503	ISO 8601
2013.123	year and day of year, in this case May 3, 2013
J2456416	Julian date, in this case May 3, 2013

Time / Time with Time Zone Input

Input	Description
22:12:34.567 22:12:34 22:12 221234	ISO 8601
10:12 AM	same as 10:12
10:12 PM	same as 22:12
22:12-5 22:12-05:00 221200-05	ISO 8601, same as 10:12 PM EST
22:12 EST	time zone specified by identifier, in this case an abbreviation; same as 10:12 PM EST

Time Zone Input

Input	Description
EST	Abbreviation (Eastern Standard Time)
America/New_York	Full name
EST5EDT	POSIX style
-5:00 -500 -5	ISO 8601 style
zulu z	Military abbreviation for UTC

http://www.postgresql.org/docs/current/static/datatype-datetime.html#DATATYPE-TIMEZONES

SELECT CURRENT_TIMESTAMP AT TIME ZONE 'CST';

Timestamp Input

Combine Date and Time and Time Zone inputs!

```
<date> <time> <timezone> (AD|BC)
```

- timestamp with time zone
 - internally stored as UTC
 - default representation is from "timezone" parameter or system default
- Caveat Emptor which of these are equivalent?

```
TIMESTAMP '2013-05-21 10:00:00'
TIMESTAMP '2013-05-21 10:00:00-05'
TIMESTAMP WITH TIME ZONE '2013-05-21 10:00:00-05'
```

Be careful with your data type declarations!

Special Inputs

Input String	Valid Types	Description
epoch	date, timestamp	1970-01-01 00:00:00+00 (Unix system time zero)
infinity	date, timestamp	later than all other time stamps
-infinity	date, timestamp	earlier than all other time stamps
now	date, time, timestamp	current transaction's start time
today	date, timestamp	midnight today
tomorrow	date, timestamp	midnight tomorrow
yesterday	date, timestamp	midnight yesterday
allballs	time	00:00:00.00 UTC

Date / Time Output

Style Specification Description		Example	
ISO	ISO 8601, SQL standard	1997-12-17 07:37:16-08	
SQL	traditional style	12/17/1997 07:37:16.00 PST	
Postgres	original style	Wed Dec 17 07:37:16 1997 PST	
German	regional style	17.12.1997 07:37:16.00 PST	

Can adjust with:

- Command: SET <datestyle>;
- postgresql.conf 'DateStyle' parameter
- environmental var: PGDATESTYLE

Intervals

- YEARMONTH
- DAY
- HOUR
- MINUTE
- SECOND
- YEAR TO MONTH
- DAY TO HOUR
- DAY TO MINUTE
- DAY TO SECOND
- HOUR TO MINUTE
- HOUR TO SECOND
- MINUTE TO SECOND

Why Intervals Are Cool

```
SELECT avg(hours)
FROM sleep
WHERE
  day BETWEEN
    CURRENT_DATE - '7 day'::interval AND
    CURRENT_DATE;
```

Basic Operators

Operator	Example	Result
+	date '2001-09-28' + integer '7'	date '2001-10-05'
+	date '2001-09-28' + interval '1 hour'	timestamp '2001-09-28 01:00:00'
+	date '2001-09-28' + time '03:00'	timestamp '2001-09-28 03:00:00'
+	interval '1 day' + interval '1 hour'	interval '1 day 01:00:00'
+	timestamp '2001-09-28 01:00' + interval '23 hours'	timestamp '2001-09-29 00:00:00'
+	time '01:00' + interval '3 hours'	time '04:00:00'
-	- interval '23 hours'	interval '-23:00:00'
-	date '2001-10-01' - date '2001-09-28'	integer '3' (days)
-	date '2001-10-01' - integer '7'	date '2001-09-24'
-	date '2001-09-28' - interval '1 hour'	timestamp '2001-09-27 23:00:00'
-	time '05:00' - time '03:00'	interval '02:00:00'
-	time '05:00' - interval '2 hours'	time '03:00:00'
-	timestamp '2001-09-28 23:00' - interval '23 hours'	timestamp '2001-09-28 00:00:00'
_	interval '1 day' - interval '1 hour'	interval '1 day -01:00:00'
-	timestamp '2001-09-29 03:00' - timestamp '2001-09-27 12:00'	interval '1 day 15:00:00'
*	900 * interval '1 second'	interval '00:15:00'
*	21 * interval '1 day'	interval '21 days'
*	double precision '3.5' * interval '1 hour'	interval '03:30:00'
/	interval '1 hour' / double precision '1.5'	interval '00:40:00'

Selected Functions

- age(timestamp, timestamp)
- age(timestamp)
- date_part(text, timestamp)
 - Same as 'EXTRACT'
- date_trunc(text, timestamp)
- justify days(interval)
- justify hours(interval)
- CURRENT_TIMESTAMP, CURRENT_DATE, CURRENT_TIME

Boolean Data Types

Postgres – second to none :-)

Name	Size
boolean	1 byte

- These are all equivalent
 - TRUE, 't', 'true', 'y', 'yes', 'on', '1'
 - FALSE, 'f', 'false', 'n', 'no', 'off', '0'
 - all case-insensitive, preferred TRUE / FALSE

Boolean Data Type Notes

• bool = boolean

NEVER CREATE AN INDEX ON A BOOLEAN TYPE

Enumerated Types

SELECT name, color_name

FROM suspect s INNER JOIN eye_color e

ON e.color_id = s.color_id

The table eye_color is (fairly) static

Declaring an Enum

CREATE TYPE enum_eye_color AS ENUM ('blue', 'brown', 'gray', 'green');

ALTER TYPE enum_eye_color ADD VALUE 'amber' BEFORE 'blue';

The declaration order is used by ORDER BY

Using an Enum

```
CREATE TABLE suspect
(name TEXT,
eye_color enum_eye_color);
```

INSERT INTO suspect VALUES
('John Doe', 'brown');

Using an Enum

Enum Alternatives

Lookup Table	
Constraint	CREATE TABLE suspect (name TEXT NOT NULL, eye_color TEXT NOT NULL, CONSTRAINT check_eye_color CHECK (eye_color IN ('blue','brown','gray','green')));
Domain	CREATE DOMAIN eye_color AS TEXT CONSTRAINT check_eye_color CHECK (VALUE IN ('blue','brown','gray','green')));

Stretch Break #1



Reading Material For Break: B-Tree Indexes

- "default" index in Postgres
- optimized for retrieving data on circular disk
 - can sometimes help with sorts
- supports <=, <, =, >, >=
 - BETWEEN, IN
 - IS NOT NULL, IS NULL
 - LIKE in specific case of 'plaintext%'
 - ~ in specific case of '^plaintext'
 - ILIKE and ~* if pattern starts with nonalpha characters
- one of many indexes in Postgres
 - some of these conditions change with other indexes...

Geometric Types

Name	Size	Representation	Format
point	16 bytes	point on a plane	(x,y)
Iseg	32 bytes	finite line segment	((x1, y1), (x2, y2))
box	32 bytes	rectangular box	((x1, y1), (x2, y2))
path	16 + 16n bytes	closed path (similar to polygon, n = total points	((x1, y1), (x2, y2),, (xn, yn))
path	16 + 16n bytes	open path, n = total points	[(x1, y1), (x2, y2),, (xn, yn)]
polygon	40 bytes + 16n	polygon	((x1, y1), (x2, y2),, (xn, yn))
circle	24 bytes	circle – center point and radius	<(x, y), r>

It Only Does Everything...

Operator	Description	Example
+	Translation	box '((0,0),(1,1))' + point '(2.0,0)'
-	Translation	box '((0,0),(1,1))' - point '(2.0,0)'
*	Scaling/rotation	box '((0,0),(1,1))' * point '(2.0,0)'
/	Scaling/rotation	box '((0,0),(2,2))' / point '(2.0,0)'
#	Point or box of intersection	'((1,-1),(-1,1))' # '((1,1),(-1,-1))'
#	Number of points in path or polygon	# '((1,0),(0,1),(-1,0))'
@ - @	Length or circumference	@-@ path '((0,0),(1,0))'
@ @	Center	@@ circle '((0,0),10)'
##	Closest point to first operand on second operand	point '(0,0)' ## lseg '((2,0),(0,2))'
<->	Distance between	circle '((0,0),1)' <-> circle '((5,0),1)'
&&	Overlaps? (One point in common makes this true.)	box '((0,0),(1,1))' && box '((0,0),(2,2))'
<<	Is strictly left of?	circle '((0,0),1)' << circle '((5,0),1)'
>>	Is strictly right of?	circle '((5,0),1)' >> circle '((0,0),1)'
&<	Does not extend to the right of?	box '((0,0),(1,1))' &< box '((0,0),(2,2))'
&>	Does not extend to the left of?	box '((0,0),(3,3))' &> box '((0,0),(2,2))'
<<	Is strictly below?	box '((0,0),(3,3))' << box '((3,4),(5,5))'
>>	Is strictly above?	box '((3,4),(5,5))' >> box '((0,0),(3,3))'

It Only Does Everything Cont'd...

&<	Does not extend above?	box '((0,0),(1,1))' &< box '((0,0),(2,2))'
&>	Does not extend below?	box '((0,0),(3,3))' &> box '((0,0),(2,2))'
<^	Is below (allows touching)?	circle '((0,0),1)' <^ circle '((0,5),1)'
>^	Is above (allows touching)?	circle '((0,5),1)' >^ circle '((0,0),1)'
?#	Intersects?	lseg '((-1,0),(1,0))' ?# box '((-2,-2),(2,2))'
?-	Is horizontal?	?- lseg '((-1,0),(1,0))'
?-	Are horizontally aligned?	point '(1,0)' ?- point '(0,0)'
?	Is vertical?	? lseg '((-1,0),(1,0))'
?	Are vertically aligned?	point '(0,1)' ? point '(0,0)'
?-	Is perpendicular?	lseg '((0,0),(0,1))' ?- lseg '((0,0),(1,0))'
?	Are parallel?	lseg '((-1,0),(1,0))' ? lseg '((-1,2),(1,2))'
@ >	Contains?	circle '((0,0),2)' @> point '(1,1)'
<@	Contained in or on?	point '(1,1)' <@ circle '((0,0),2)'
~=	Same as?	polygon '((0,0),(1,1))' ~= polygon '((1,1),(0,0))'

It Only Does Everything Cont'd...

Function	Return Type	Description	Example
area(object)	double precision	area	area(box '((0,0),(1,1))')
center(object)	point	center	center(box '((0,0),(1,2))')
diameter(circle)	double precision	diameter of circle	diameter(circle '((0,0),2.0)')
height(box)	double precision	vertical size of box	height(box '((0,0),(1,1))')
isclosed(path)	boolean	a closed path?	isclosed(path '((0,0),(1,1),(2,0))')
isopen(path)	boolean	an open path?	isopen(path '[(0,0),(1,1),(2,0)]')
length(object)	double precision	length	length(path '((-1,0),(1,0))')
npoints(path)	int	number of points	npoints(path '[(0,0),(1,1),(2,0)]')
npoints(polygon)	int	number of points	npoints(polygon '((1,1),(0,0))')
pclose(path)	path	convert path to closed	pclose(path '[(0,0),(1,1),(2,0)]')
popen(path)	path	convert path to open	popen(path '((0,0),(1,1),(2,0))')
radius(circle)	double precision	radius of circle	radius(circle '((0,0),2.0)')
width(box)	double precision	horizontal size of box	width(box '((0,0),(1,1))')

Performance Considerations

- Size on disk
 - Consider I/O on retrievals
- Indexing
 - B-tree
 - equality operators modified for ad-hoc purposes, e.g. area
 - Are we out of luck on performance?

Index Detour #1: Expression Indexes

- allows pre-computed values to be stored in an index
- useful for "on the fly" comparisons

```
SELECT * FROM receipts WHERE (subtotal + tax) <
  numeric(1000.00);
SELECT * FROM receipts WHERE upper(name) = 'JIM';</pre>
```

- fast for searches, costly on updates
- Easy to create

```
CREATE INDEX receipts_total_idx ON
  receipts ((subtotal + tax));
CREATE INDEX receipts_upper_idx ON
  receipts((upper(name)));
```

Back to Geometric Performance: Expression Indexes

```
CREATE TABLE houses (plot box);

INSERT INTO houses

SELECT box(
  point((500 * random())::int, (500 * random())::int),
  point((750 * random() + 500)::int, (750 * random() + 500)::int)
)

FROM generate_series(1, 1000000);
```

Area without Expression Index

```
EXPLAIN SELECT * FROM houses WHERE area(plot)
BETWEEN 50000 AND 75000;

QUERY

----
Seq Scan on houses (cost=0.00..27353.00
rows=5000 width=32)
Filter: ((area(plot) >= 50000::double
precision) AND (area(plot) <= 75000::double
```

Run time average 220ms

precision)

Area with Expression Index

Average run time was 48ms

Index Detour #2: GiST

- "generalized search tree"
- balanced, tree-structured
- allows arbitrary indexing schemes
 - B-trees, R-trees
 - indexing on custom data types
- supports lots more operators
- can implement your own indexing scheme

Index Detour #2: GiST

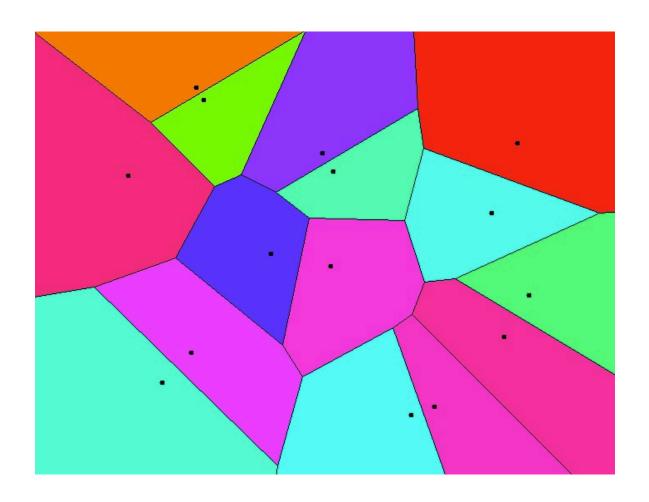
• GiST operators:

- <<
- &<
- **-** &>
- **-** >>
- <<|
- &<|
- |&>
- |>>
- **–** @>
- <@
- _ ~=
- &&

Major Detour: K-Nearest Neighbor

- PostgreSQL 9.1+
- Given a collection of n objects
- When trying to classify an unknown object
 - compute the distance between all known objects
 - find the k ($k \ge 1$) closest objects to the unknown object

K=1 Example



KNN-GiST: A Very Special GiST Index

- (almost back to geometric types!)
- Let n = size of a table
- Can index data that provides a "<->" (distance) operator
- "k" = LIMIT clause

Geometry

```
CREATE INDEX geoloc coord idx ON geoloc USING gist(coord);
EXPLAIN ANALYZE
  SELECT
    coord,
    coord <-> point (500,500)
  FROM geoloc
  ORDER BY coord <-> point(500,500)
  LIMIT 10;
```

Results

```
Limit (cost=80958.28..80958.31 rows=10
   width=20) (actual
   time=1035.313..1035.316 rows=10
   loops=1)
 -> Sort (cost=80958.28..85958.28
   rows=2000000 width=20) (actual
   time=1035.312..1035.314 rows=10
   loops=1)
    Sort Key: ((coord <->
   '(500,500)'::point))
    Sort Method: top-N heapsort
   Memory: 25kB
    -> Seq Scan on geoloc
   (cost=0.00..37739.00 rows=2000000
   width=20) (actual
   time=0.029..569.501 rows=2000000
   loops=1)
```

Total runtime: 1035.349 ms

```
Limit (cost=0.00..0.81 rows=10 width=20) (actual time=0.576..1.255 rows=10 loops=1)

-> Index Scan using geoloc_coord_idx on geoloc (cost=0.00..162068.96 rows=2000000 width=20) (actual time=0.575..1.251 rows=10 loops=1)
```

Order By: (coord <-> '(500,500)'::point)

Total runtime: 1.391 ms

Application Examples

Proximity map search – fast!

KNN - Amsterdam Edition



Geometric Type Index Summary

- B-tree
- Expression indexes on geometric functions
- GiST indexes
 - Support special geometric operators
 - KNN-GiST

Addicted to Geometry? GIS?

PostGIS

- http://postgis.refractions.net/
- OpenGIS, WKB, WKT
- PostGIS EWKB, EWKT
- SQL-MM Part 3
- PostGIS Geography Type
- ...and takes advantage of GiST

Network Address Types

Name	Storage Size	Description
cidr	7 or 19 bytes	IPv4 or IPv6 networks
inet	7 or 19 bytes	IPv4 or IPv6 hosts or networks
macaddr	6 bytes	MAC addresses

Network Address Types

```
• inet (IPv4 & IPv6)
  - SELECT '192.168.1.1'::inet;
  - SELECT '192.168.1.1/32'::inet;
  - SELECT '192.168.1.1/24'::inet;
• cidr (IPv4 & IPv6)
  - SELECT '192.168.1.1'::cidr;
  - SELECT '192.168.1.1/32'::cidr;
  - SELECT '192.168.1.1/24'::cidr;
• macaddr
  - SELECT '08:00:2b:01:02:03'::macaddr
```

inet + cidr =

Operator	Description	Example
<	is less than	inet '192.168.1.5' < inet '192.168.1.6'
<=	is less than or equal	inet '192.168.1.5' <= inet '192.168.1.5'
=	equals	<pre>inet '192.168.1.5' = inet '192.168.1.5'</pre>
>=	is greater or equal	inet '192.168.1.5' >= inet '192.168.1.5'
>	is greater than	inet '192.168.1.5' > inet '192.168.1.4'
<>	is not equal	inet '192.168.1.5' <> inet '192.168.1.4'
<<	is contained within	inet '192.168.1.5' << inet '192.168.1/24'
<<=	is contained within or equals	inet '192.168.1/24' <<= inet '192.168.1/24'
>>	contains	inet '192.168.1/24' >> inet '192.168.1.5'
>>=	contains or equals	<pre>inet '192.168.1/24' >>= inet '192.168.1/24'</pre>
~	bitwise NOT	~ inet '192.168.1.6'
&	bitwise AND	inet '192.168.1.6' & inet '0.0.0.255'
I	bitwise OR	inet '192.168.1.6' inet '0.0.0.255'
+	addition	inet '192.168.1.6' + 25
_	subtraction	inet '192.168.1.43' - 36
_	subtraction	inet '192.168.1.43' - inet '192.168.1.19'

...even more functions

Function	Return Type	Description	Example
abbrev(inet)	text	abbreviated display format as text	abbrev(inet '10.1.0.0/16')
abbrev(cidr)	text	abbreviated display format as text	abbrev(cidr '10.1.0.0/16')
broadcast(inet)	inet	broadcast address for network	broadcast('192.168.1.5/24')
family(inet)	int	extract family of address; 4 for IPv4, 6 for IPv6	family('::1')
host(inet)	text	extract IP address as text	host('192.168.1.5/24')
hostmask(inet)	inet	construct host mask for network	hostmask('192.168.23.20/30')
masklen(inet)	int	extract netmask length	masklen('192.168.1.5/24')
netmask(inet)	inet	construct netmask for network	netmask('192.168.1.5/24')
network(inet)	cidr	extract network part of address	network('192.168.1.5/24')
set_masklen(inet, int)	inet	set netmask length for inet value	set_masklen('192.168.1.5/24', 16)
set_masklen(cidr, int)	cidr	set netmask length for cidr value	set_masklen('192.168.1.0/24'::cidr, 16)
text(inet)	text	extract IP address and netmask length as text	text(inet '192.168.1.5')

Bit Strings

Name	Storage Size	Description
bit(n)	y + ceil(n / 8) bytes	stores exactly n 0s and 1s y = 5 or 8
bit varying(n)	y + ceil(n / 8) bytes	stores up to n 0s and 1s y = 5 or 8
bit varying	variable	stores unlimited number of 0s and 1s

Bit Strings

```
SELECT B'10010010101000';
SELECT '1'::bit(3); -- '100';
CREATE TABLE bits (
 a bit (3),
 b bit varying(5),
 c bit varying
```

Bit Strings

```
SELECT B'101' | B'010'; -- 101010

SELECT B'1011' & B'0101'; -- 0001

SELECT B'1011' | B'0101'; -- 1111

SELECT B'1011' | B'0101'; -- 1110

SELECT B'1011'; -- 0100

SELECT B'1011' | 2; -- 1100

SELECT B'1011' >> 2; -- 0010
```

- built-in to PostgreSQL
- uses "tsearch2" algorithm
- appropriate data types + indexes for retrieval

tsvector

a sorted list of normalized lexemes

```
SELECT 'PGDay NYC 2013 is a conference run by the local NYC PostgreSQL User Group'::tsvector;
```

```
tsvector
-----
'2013' 'Group' 'NYC' 'PGDay' 'PostgreSQL' 'User' 'a' 'by'
   'conference' 'is' 'local' 'run' 'the
```

```
SELECT tsvector('now:1 i:2 have:3
 learned: 4 my: 5 data: 6 types: 7
 and:8 i:9 can:10 teach:11 the:12
 world:13 now:14');
 tsvector
 'and':8 'can':10 'data':6 'have':3
 'i':2,9 'learned':4 'my':5 'now':
 1,14 'teach':11 'the':12 'types':
 7 'world':13
```

```
SELECT tsvector('now:1C i:2 have:
 3 learned: 4B my: 5 data: 6A
 types:7B');
 tsvector
 'data':6A 'have':3 'i':2
 'learned':4B 'my':5 'now':1C
 'types':7B
```

- use 'to_tsvector' to normalize text
- used for indexing on actual full text search applications

- tsquery
 - lexemes that are to be searched for
 - operators: &, |, !

```
SELECT 'PostgreSQL & conference'::tsquery;
tsquery
-----
'PostgreSQL' & 'conference'
```

```
SELECT 'PostgreSQL & (conference | 2013)'::tsquery;
tsquery
 'PostgreSQL' & ( 'conference' | '2013' )
SELECT 'PostgreSQL & !conference | 2013'::tsquery;
tsquery
 'PostgreSQL' & !'conference' | '2013'
```

prefix matching

```
SELECT 'pg:* & conference'::tsquery;

tsquery
-----
'pg:*' & 'conference'
```

be aware of stemming

```
S
ELECT to_tsvector('postgraduate') @@ to_tsquery('postgres:*');
```

Full Text Search Functions

- •get_current_ts_config
- length
- numnode
- •plainto_tsquery
- querytree
- •setweight
- strip
- •to_tsquery
- •to_tsvector
- •ts_headline
- •ts_rank
- •ts_rank_cd
- •ts_rewrite
- •tsvector_update_trigger
- •tsvector_update_trigger_column

A lot of functions and "weird" operators involved

```
SELECT title
FROM conferences
WHERE to_tsvector(title) @@
to_tsquery('postgres:*');
```

Index Detour #3: GIN

- Generalized Inverted Index
 - search for composite values in composite items (huh?)
 - provides general access methods for implementor to provide logic
 - stores data by "keys"
 - rows referenced by multiple keys
 - exact vs partial match
 - fast on reads, slow on writes
- Supported on
 - full text search
 - btree_gin
 - hstore
 - pg_trgm
 - one-dimensional arrays on built-in types

- Can use GiST or GIN
 - Size: GIN 2-3x larger
 - Read performance: GIN 2-3x faster
 - Index creation: GiST 2-3x faster
 - Update: GiST moderately to 10x faster
 - (FASTUPDATE on GIN)
 - "100,000 lexemes"

```
CREATE INDEX full_text_search_idx ON conferences
USING gin(to tsvector('title'));
```

More on Full Text Search

- Lecture in itself
- http://www.postgresql.org/docs/current/ static/textsearch.html

UUID

- Universally Unique Identifiers
- 16 bytes on disk
- Acceptable Formats
 - A0EEBC99-9C0B-4EF8-BB6D-6BB9BD380A11
 - {a0eebc99-9c0b-4ef8-bb6d-6bb9bd380a11}
 - a0eebc999c0b4ef8bb6d6bb9bd380a11
 - a0ee-bc99-9c0b-4ef8-bb6d-6bb9-bd38-0a11
 - {a0eebc99-9c0b4ef8-bb6d6bb9-bd380a11}

UUID Functions

CREATE EXTENSION "uuid-ossp";

- uuid_generate_v1
- uuid_generate_v1mc
- uuid_generate_v3
- uuid_generate_v4
- uuid_generate_v5

XML

- ensures that XML is valid
- no comparison methods
- caveat emptor: encoding
 - e.g. 'xpath'

XML

Ensures the value is well formed XML

```
postgres=# SELECT xml '<PUG>NYC</PUG>';
     xml
------
<PUG>NYC</PUG>

postgres=# SELECT xml '<PUG>NYC';
ERROR: invalid XML content at character 12
DETAIL: line 1: Premature end of data in tag
   PUG line 1
<PUG>NYC
```

XML Functions

- xml_is_well_formed
- xpath_string
- xpath_number
- xpath_bool
- xpath_nodeset
- xpath_nodeset
- xpath_nodeset
- xpath_list
- xpath_list

Embedded XML Fragments

```
CREATE TABLE Journey (
    JourneyId INTEGER,
    T<sub>1</sub>X XMT<sub>1</sub>
);
INSERT INTO Journey
VALUES (1, '<LX>
              <LEG LAT="52" LONG="0">
              <LEG LAT="44" LONG="5" >
             </LX>');
UPDATE Journey
   SET LegX = ' < LX > ' | |
                  xpath string(LegX, ''/lx/leg'') ||
                 '<LEG LAT="56" LONG="10" >
                 </T<sub>1</sub>X>
WHERE JourneyId = 1;
SELECT xpath number(LegX, 'fn:count(/lx/leg)') as num legs FROM Journey WHERE
    JournevId = 1;
```

Arrays

```
CREATE TABLE person (
    full name text,
    sports text[],
   cars text[][],
   numbers int[3],
    incomes int ARRAY[4],
   phrases text ARRAY

    PostgreSQL does not enforce size restrictions

  - 9.3 and below
```

Arrays

```
SELECT ARRAY[1,2,3];
SELECT ARRAY[ARRAY[1,2], ARRAY[3,4]];
SELECT '{1,2,3}';
SELECT '{{1,2},{3,4}}';
```

Arrays

arrays are 1-indexed

```
SELECT (ARRAY[1,2,3])[1]; -- returns 1

SELECT (ARRAY[1,2,3])[0]; -- returns NULL

SELECT (ARRAY[1,2,3])[1:2]; -- returns {1,2}

SELECT (ARRAY[1,2,3])[2:3]; -- returns {2,3}

SELECT (ARRAY[1,2,3])[2:3][2]; -- returns {2,3}

SELECT (ARRAY[1,2,3])[2:3])[2]; -- returns 3
```

INSERT with ARRAY

```
INSERT INTO person
 VALUES ('Rocky Bama',
  '{"baseball", "basketball"}',
  '{{"Toyota", "Prius"}, {"Chevy", "Tahoe"}}');
INSERT INTO person
 VALUES ('Rocky Bama',
  ARRAY['baseball','basketball'],
  ARRAY [['Toyota', 'Prius'],
  ['Chevy', 'Tahoe']]);
```

SELECT and ARRAY

SELECT and ARRAY

```
SELECT *
FROM person
WHERE sports[1] = 'baseball';
```

SELECT and ARRAY

```
SELECT full_name
FROM person
WHERE 'baseball' = ANY (sports);
```

"true" if any entry in sports for a person is 'baseball'

SELECT and ARRAY (2)

```
SELECT full_name
FROM person
WHERE 'baseball' = ALL
  (sports);
```

 "true" only if every entry in sports for a tuple in person is 'baseball'

UPDATE and **ARRAY**

```
UPDATE person SET sports[2] = 'tennis';

UPDATE person SET sports[2:3] = '{"hockey",
    "soccer"}';

UPDATE person SET sports = ARRAY['foozball',
    'billiards'];
```

Array Operators

```
• <, <=, =, >= >, <>

    compares each array elements

   – B-tree index = yes!
• @>,<@
   SELECT ARRAY [1, 2, 3] @> ARRAY [1, 2];
   SELECT ARRAY [1,2] < \emptyset ARRAY [1,2,3];
&&
   SELECT ARRAY[1,2,3] && ARRAY[3,4,5];
   SELECT ARRAY[1,2,3] | ARRAY[3,4,5];
   SELECT ARRAY [ARRAY [1,2], ARRAY [3,4]] | ARRAY [5,6];
   SELECT ARRAY[1,2,3] || 4;
```

can use GIN index on one dimensional arrays

Array Functions

modification

```
SELECT array_append(ARRAY[1,2,3], 4);

SELECT array_prepend(1, ARRAY[2,3,4]);

SELECT array_cat(ARRAY[1,2], ARRAY[3,4]);

SELECT array_remove(ARRAY[1,2,1,3], 1);

SELECT array replace(ARRAY[1,2,1,3], 1, -4)
```

size

```
SELECT array_length(ARRAY[1,2,3,4], 1); -- 4

SELECT array_ndims(ARRAY[ARRAY[1,2], ARRAY[3,4]]);
-- 2

SELECT array_dims(ARRAY[ARRAY[1,2], ARRAY[3,4]]);
-- [1:2][1:2]
```

Array Functions

bounds SELECT array lower (ARRAY[2,3,4], 1); SELECT array upper (ARRAY[2,3,4], 1); join SELECT array to string(ARRAY[1,2,NULL,4], ',', '*'); **--** 1,2,*,4 expand SELECT unnest (ARRAY[1,2,3]);unnest

array_agg

- useful for variable-length lists or "unknown # of columns"
 - e.g. "find all speakers for a talk"

- Added in 9.2
- Ensures the value is valid JSON

• Enhanced functionality added in 9.3

Operator	Description	Example
->	return JSON array element OR JSON object field	'[1,2,3]'::json -> 0; '{"a": 1, "b": 2, "c": 3}'::json -> 'b';
->>	return JSON array element OR JSON object field AS text	['1,2,3]'::json ->> 0; '{"a": 1, "b": 2, "c": 3}'::json ->> 'b';
#>	return JSON object using path	'{"a": 1, "b": 2, "c": [1,2,3]}'::json #> '{c, 0}';
#>>	return JSON object using path AS text	'{"a": 1, "b": 2, "c": [1,2,3]}'::json #> '{c, 0}';

array_to_json

```
SELECT array to json(ARRAY[ARRAY[1,2], ARRAY[3,4]])
[[1,2],[3,4]]
row_to_json
SELECT row to json(ROW(1,2,3));
{"f1":1, "f2":2, "f3":3}
SELECT row_to_json(x) FROM x LIMIT 1;
{"a":56, "b":42, "c":63}
```

- json_extract_path, json_extract_path_text
 - LIKE (#>, #>>) but with list of args

```
SELECT json_extract_path(
   '{"a": 1, "b": 2, "c": [1,2,3]}'::json,
   'c', '0');
```

- to_json
- json_each, json_each_text

json_object_keys

```
SELECT * FROM json_object_keys('{"a": 1,
    "b": [2,3,4], "c": { "e":
    "wow" }}'::json);
-----
a
b
c
```

json_populate_record

json_populate_recordset

json_agg

Stretch Break #2



Stretch Break #2

• If you don't want to stretch, try this puzzle: With this table:

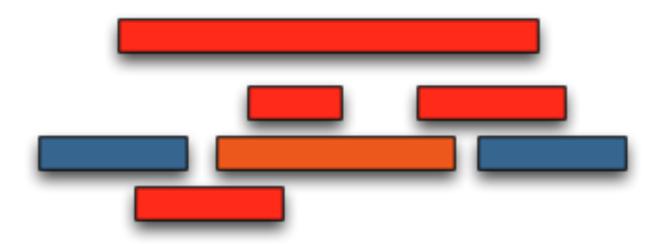
name (text)	low (int)	high (int)
а	20	30
b	25	35
С	10	15
d	17	24
е	40	50
f	26	36

Write a query that finds all the names whose range of values are between 18 and 26

Ranges

- Scheduling
- Probability
- Measurements
- Financial applications
- Clinical trial data
- Intersections of ordered data

Why Range Overlaps Are Difficult



Before Postgres 9.2

OVERLAPS

```
SELECT
('2013-01-08`::date, '2013-01-10'::date) OVERLAPS
('2013-01-09'::date, '2013-01-12'::date);
```

- Limitations:
 - Only date/time
 - Start <= x <= End</p>

Postgres 9.2+

- INT4RANGE (integer)
- INT8RANGE (bigint)
- NUMRANGE (numeric)
- TSRANGE (timestamp without time zone)
- TSTZRANGE (timestamp with time zone)
- DATERANGE (date)

Range Type Size

Size on disk = 2 * (data type) + 1
 –sometimes magic if bounds are equal

Range Bounds

Ranges can be inclusive, exclusive or both

$$-[2,4] \Rightarrow 2 \leq x \leq 4$$
 $-[2,4] \Rightarrow 2 \leq x \leq 4$
 $-(2,4] \Rightarrow 2 \leq x \leq 4$
 $-(2,4] \Rightarrow 2 \leq x \leq 4$

Can also be empty

Infinite Ranges

Ranges can be infinite

$$-[2,) => 2 \le x < \infty$$

 $-(2) => -\infty < x \le 2$

- CAVEAT EMPTOR
 - "infinity" has special meaning with timestamp ranges

```
- [CURRENT_TIMESTAMP,) = [CURRENT_TIMESTAMP,]
```

```
- [CURRENT_TIMESTAMP, 'infinity') <>
  [CURRENT_TIMEAMP, 'infinity']
```

Constructing Ranges

```
SELECT '[1,10]'::int4range;
int4range
-----
[1,11)
(1 row)
```

Constructing Ranges

- Constructor functions too
 - Defaults to '[)'

```
test=# SELECT numrange(9.0, 9.5);
  numrange
-----
[9.0,9.5)
(1 row)
```

Bonus

Can have arrays of ranges

Simple Overlaps

Range Indexes

 Creating a GiST index on ranges speeds up queries with these operators:

```
= & & 
 < @ 
 @ > 
 < < 
 > > 
 - | - 
 & < 
 & >
```

Range Indexes

Performance

```
test=# EXPLAIN ANALYZE SELECT * FROM ranges WHERE
int4range(500,1000) && bounds;

QUERY PLAN
------
Bitmap Heap Scan on ranges
(actual time=0.283..0.370 rows=653 loops=1)
   Recheck Cond: ('[500,1000)'::int4range && bounds)
   -> Bitmap Index Scan on ranges_bounds_gist_idx
        (actual time=0.275..0.275 rows=653 loops=1)
        Index Cond: ('[500,1000)'::int4range &&
bounds)
Total runtime: 0.435 ms
```

What If the Range is Much Larger?

Another Index Detour: SP-GiST

- "space-partitioned generalized search tree"
- designed for handling unbalanced data structures
 - quadtrees
 - k-d trees
 - radix trees
- searches are fast if match partitioning rules

```
CREATE INDEX ranges_bounds_spgist_idx ON ranges
spgist(bounds);
```

SP-GiST and Ranges (9.3+)

SP-GiST indexes support ranges

```
QUERY PLAN

Bitmap Heap Scan on ranges (cost=20.41..1748.32 rows=516 width=17) (actual time=0.558..1.463 rows=1502 loops=1)
Recheck Cond: (500 <@ bounds)

-> Bitmap Index Scan on ranges_bounds_spgist_idx (cost=0.00..20.28 rows=516 width=0) (actual time=0.413..0.413 rows=1502 loops=1)

Index Cond: (500 <@ bounds)

Total runtime: 1.585 ms
```

Scheduling

```
CREATE TABLE travel log (
  id serial PRIMARY KEY,
 name varchar (255),
  travel range daterange,
 EXCLUDE USING gist (travel range WITH &&)
);
INSERT INTO travel log (name, trip range) VALUES ('Chicago',
  daterange('2012-03-12', '2012-03-17'));
INSERT INTO travel log (name, trip range) VALUES ('Austin',
  daterange('2012-03-16', '2012-03-18'));
ERROR: conflicting key value violates exclusion constraint
  "travel log trip range excl"
DETAIL: Key (trip range) = ([2012-03-16,2012-03-18)) conflicts
  with existing key (trip range) = ([2012-03-12,2012-03-17)).
```

Extending Ranges

```
CREATE TYPE inetrange AS RANGE (
  SUBTYPE = inet
);
SELECT '192.168.1.8'::inet <@ inetrange('192.168.1.1',
  '192.168.1.10');
?column?
t.
SELECT '192.168.1.20'::inet <@ inetrange('192.168.1.1',
  '192.168.1.10');
 ?column?
```

...back to the original problem

```
SELECT name FROM ranges WHERE range && int4range(18,26,'[]');
```

Composite Types

```
CREATE TYPE address AS (
    street TEXT,
    city TEXT,
    state TEXT,
    zip CHAR(10)
);
```

Composite Types

```
CREATE TABLE customer (
   full_name TEXT,
   mail_address address
);
```

Composite Types

```
INSERT INTO customer VALUES
('Joe Lee',
ROW('100 Broad Street', 'Red
 Bank', 'NJ', '07701'));
INSERT INTO customer VALUES
('Joe Lee',
 ('100 Broad Street', 'Red Bank',
 'NJ', '07701'));
```

Composite Types with SELECT

```
SELECT (mail_address).city
FROM customer
WHERE (mail_address).state = 'NJ';
SELECT (customer.mail_address).city
FROM customer
WHERE (customer.mail_address).state = 'NJ';
```

Composite Types and JSON

```
SELECT row to json(customer)
FROM customer;
  "full name": "Joe Lee",
  "mail address": {
      "street": "100 Broad Street",
      "city": "Red Bank",
      "state": "NJ",
      "zip":"07701"
```

Composite Type Operators

- Create a new function using CREATE FUNCTION that accepts one or two arguments using this type
- Use CREATE OPERATOR to choose what operator should be used to invoke this function

```
CREATE OPERATOR = (
    PROCEDURE = addr_eq,
    LEFTARG=address,
    RIGHTARG=address
);
```

- Still not enough choices? Create your own.
 - PostGIS geometry
 - Hstore
 - BioPostgres

Needs an input function

```
Datum
tinyint_in(PG_FUNCTION_ARGS)
{
  char *num = PG_GETARG_CSTRING(0);
  PG_RETURN_TINYINT(pg_atoi(num, sizeof(tinyint), '\0'));
}
```

Needs an output function

```
Datum
tinyint_out(PG_FUNCTION_ARGS)
{
  tinyint arg1 = PG_GETARG_TINYINT(0);
  /* sign, 3 digits, '\0' */
  char *result = (char *) palloc(5);

  pg_itoa(arg1, result);
  PG_RETURN_CSTRING(result);
}
```

And a type defintion

And that's it...

Should add operators

```
Datum
tinyint_eq(PG_FUNCTION_ARGS)
{
    PG_RETURN_BOOL(PG_GETARG_TINYINT(0) == PG_GETARG_TINYINT(1));
}
```

Extensions

- "pg_contrib"
 - additional supplied modules
 - some provide additional data types outside of core
- Postgres 9.1+
 - CREATE EXTENSION "extension-name";
- Postgres <=9.0
 - psql -f path/to/contrib/install.sql yourdb
- must be database owner or superuser for both methods

cube

- data type for n-dimensional cubes
- stored as 64-bit floats
- CREATE EXTENSION cube;

```
SELECT '1'::cube;
SELECT '(1,2,3)'::cube;
SELECT '(1,2,3), (4,5,6)'::cube;
```

cube

- supports <, <=, =, >=, >, <>
- &&
 - cube overlap
- <@
 - a <@ b cube a is contained by cube b</p>
- @>
 - a @> b cube a contains cube b

cube

- Indexing
 - B-tree
 - GiST

- key-value store in PostgreSQL
- stores keys and values as strings
- installation
 - "CREATE EXTENSION hstore"

```
SELECT 'jk=>1, jm=>2'::hstore;
------
"jk"=>"1", "jm"=>"2"
```

```
SELECT hstore (ARRAY['jk', 'jm'], ARRAY['1',
  '2']);
 "jk"=>"1", "jm"=>"2"
SELECT hstore(ARRAY['jk', '1', 'jm', '2']);
 "jk"=>"1", "jm"=>"2"
SELECT hstore(ROW('jk', 'jm'));
 "f1"=>"jk", "f2"=>"jm"
```

```
SELECT ('jk=>1, jm=>2'::hstore) \rightarrow 'jk';
SELECT ('jk=>1, jm=>2'::hstore) \rightarrow ARRAY['jk','jm'];
 {1,2}
SELECT delete('jk=>1, jm=>2'::hstore, 'jm');
 "jk"=>"1"
```

```
SELECT ('jk=>1, jm=>2'::hstore) @> 'jk=>1'::hstore;
t
SELECT ('jk=>1, jm=>2'::hstore) ? 'sf';
f
SELECT ('jk=>1, jm=>2'::hstore) ?& ARRAY['jk', 'sf'];
f
SELECT ('jk=>1, jm=>2'::hstore) ?| ARRAY['jk', 'sf'];
t
```

```
SELECT hstore_to_array('jk=>1, jm=>2'::hstore);
 {jk,1,jm,2}
SELECT hstore to matrix('jk=>1, jm=>2'::hstore);
 \{\{jk,1\},\{jm,2\}\}
SELECT hstore_to_json('jk=>1, jm=>2'::hstore);
 {"jk": "1", "jm": "2"}
SELECT hstore_to_json_loose('jk=>1, jm=>2'::hstore);
 {"jk": 1, "jm": 2}
```

- akeys, avals
 - array
- skeys, svals
 - set
- each
 - set of all keys + valls
- slice
 - similar to "hstore -> ARRAY[]"
- delete

```
SELECT delete('jk=>1, jm=>2'::hstore, 'jm');
-----
"jk"=>"1"
```

- supports GiST and GIN indexes
 - @>, ?, ?&, ?|
- supports B-tree and hash indexes
 - "=" comparisons
 - enables UNIQUE hstore columns
 - DISTINCT, GROUP BY, ORDER BY

Stretch Break #3

- Exercise:
 - come up with ideas for new data type extensions

Just Kidding...Conclusion

- There are a *lot* of data types in PostgreSQL
- ...and if there are not enough, you can create more

References

- PostgreSQL 9.2
 - http://www.postgresql.org/docs/current/static/ index.html
- PostgreSQL 9.3beta1
 - http://www.postgresql.org/docs/devel/static/ index.html
- Other talks
 - https://wiki.postgresql.org/images/4/46/Knn.pdf
 - https://wiki.postgresql.org/images/f/f0/Rangetypes.pdf

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- Feedback please!
 - https://papers.pgcon.org/feedback/PGCon2013/ event/633.en.html