

# Beyond Joins and Indexes

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As a follow up to the presentation, *[Explaining the Postgres Query Optimizer](#)*, this talk shows the non-join and non-index operations that the optimizer can choose.

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# *Explaining the Postgres Query Optimizer*

My previous talk, *Explaining the Postgres Query Optimizer*, covered:

- Query optimization basics
- Optimizer statistics
- Join methods
- Scan methods, including indexes
- Limit

# This Presentation Covers Everything Else

- |                     |                        |                         |                  |
|---------------------|------------------------|-------------------------|------------------|
| 1. Result           | 12. Memoize            | 23. Gather Merge        | 34. Foreign Scan |
| 2. Values Scan      | 13. Group              | 24. Parallel Append     | 35. Tid Scan     |
| 3. Function Scan    | 14. Aggregate          | 25. Parallel Hash       | 36. Insert       |
| 4. Incremental Sort | 15. GroupAggregate     | 26. Parallel Hash Join  | 37. Update       |
| 5. Unique           | 16. HashAggregate      | 27. CTE Scan            | 38. Delete       |
| 6. Append           | 17. MixedAggregate     | 28. WorkTable Scan      | 39. Merge        |
| 7. Merge Append     | 18. WindowAgg          | 29. Recursive Union     | 40. Semi Join    |
| 8. Subquery Scan    | 19. Parallel Seq Scan  | 30. ProjectSet          | 41. Anti Join    |
| 9. HashSetOp        | 20. Partial Aggregate  | 31. LockRows            | 42. SubPlan      |
| 10. SetOp           | 21. Gather             | 32. Sample Scan         | 43. Others       |
| 11. Materialize     | 22. Finalize Aggregate | 33. Table Function Scan |                  |

# Controls

My previous talk covered:

- enable\_seqscan
- enable\_bitmapscan
- enable\_indexscan
- enable\_indexonlyscan
- enable\_nestloop
- enable\_hashjoin
- enable\_mergejoin
- enable\_sort

This talk will cover:

- enable\_incremental\_sort
- enable\_material

- enable\_memoize
- enable\_hashagg
- enable\_gathermerge
- enable\_parallel\_append
- enable\_parallel\_hash
- enable\_tidscan

Not covered:

- enable\_async\_append
- enable\_partition\_pruning
- enable\_partitionwise\_join
- enable\_partitionwise\_aggregate

## 1. Result

```
-- This disables EXPLAIN cost output  
\set EXPLAIN 'EXPLAIN (COSTS OFF)'
```

```
:EXPLAIN SELECT 1;  
QUERY PLAN
```

```
-----
```

Result

All the queries used in this presentation are available at <https://momjian.us/main/writings/pgsql/beyond.sql>.

## 2. Values Scan

```
:EXPLAIN VALUES (1), (2);  
      QUERY PLAN
```

```
-----  
Values Scan on "*VALUES*"
```

Optimizer choices are in red and causes are in blue.

### 3. Function Scan

```
:EXPLAIN SELECT * FROM generate_series(1,4);  
      QUERY PLAN
```

-----  
Function Scan on generate\_series

## 4. Incremental Sort

```
CREATE TABLE large (x) AS SELECT generate_series(1, 1000000);  
ANALYZE large;  
CREATE INDEX i_large ON large (x);
```

```
ALTER TABLE large ADD COLUMN y INTEGER;
```

```
:EXPLAIN SELECT * FROM large ORDER BY x,y;
```

```
QUERY PLAN
```

---

### Incremental Sort

Sort Key: x, y

Presorted Key: x

-> Index Scan using i\_large on large



# Incremental Sort

3	6
3	3
3	2
3	4
3	12
3	5
4	8
4	6
4	11
4	3

Incremental  
Sort  
→

3	2
3	3
3	4
3	5
3	6
3	12
4	3
4	6
4	8
4	11

## 5. Unique, First Example

```
:EXPLAIN SELECT DISTINCT * FROM generate_series(1, 10) ORDER BY 1;  
          QUERY PLAN
```

-----  
Unique

-> Sort

Sort Key: generate\_series

-> Function Scan on generate\_series

## Unique, Second Example

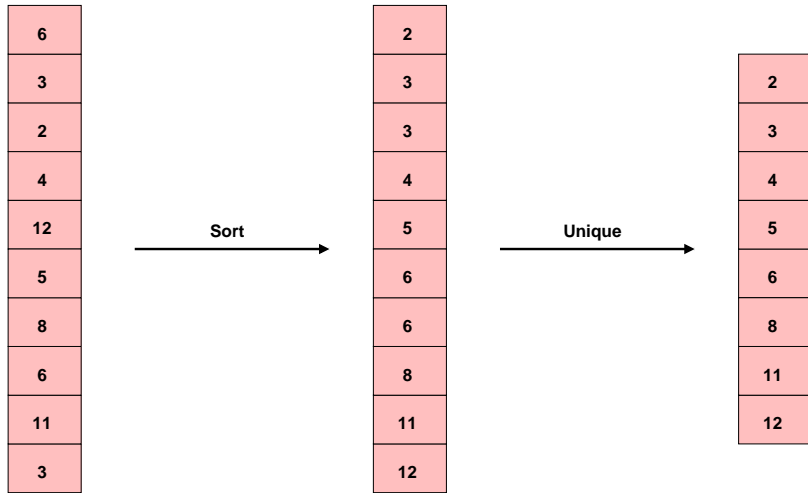
```
-- not UNION ALL  
:EXPLAIN SELECT 1 UNION SELECT 2;  
      QUERY PLAN
```

-----

Unique

```
-> Sort  
    Sort Key: (1)  
    -> Append  
        -> Result  
        -> Result
```

# Unique



## 6. Append

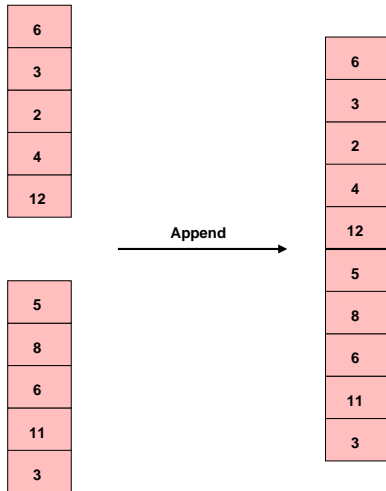
```
:EXPLAIN SELECT 1 UNION ALL SELECT 2;  
QUERY PLAN
```

-----

Append

- > Result
- > Result

# Append



## 7. Merge Append

```
:EXPLAIN (VALUES (1), (2) ORDER BY 1)
UNION ALL
      (VALUES (3), (4) ORDER BY 1)
ORDER BY 1;
```

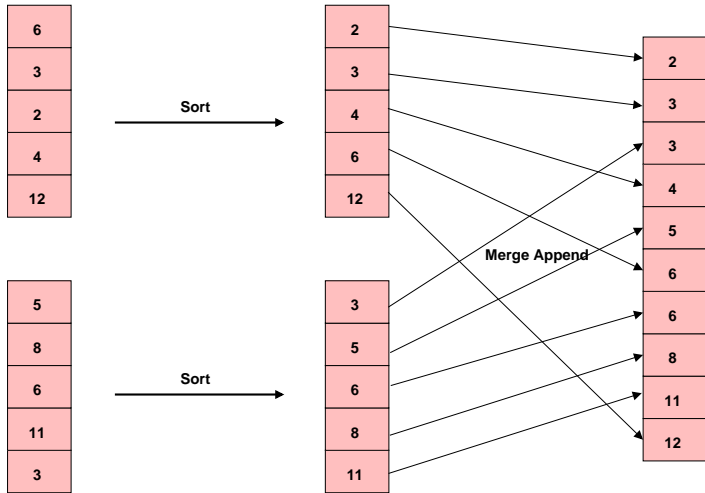
QUERY PLAN

-----

### Merge Append

```
Sort Key: "*VALUES*".column1
-> Sort
    Sort Key: "*VALUES*".column1
    -> Values Scan on "*VALUES*"
-> Sort
    Sort Key: "*VALUES*_1".column1
    -> Values Scan on "*VALUES*_1"
```

# Merge Append





## 8, 9. Subquery Scan, HashSetOp

```
CREATE TABLE small (x) AS  
SELECT generate_series(1, 1000);  
ANALYZE small;
```

```
:EXPLAIN SELECT * FROM small EXCEPT SELECT * FROM small;  
QUERY PLAN
```

-----  
HashSetOp Except

-> Append

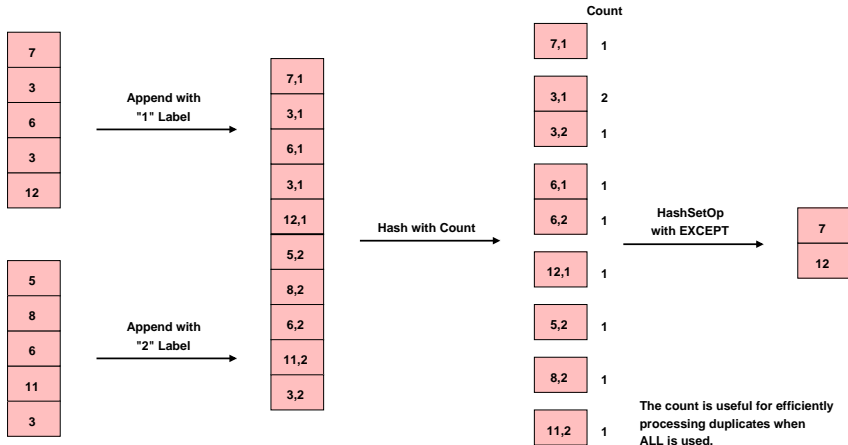
-> Subquery Scan on "\*SELECT\* 1"

-> Seq Scan on small

-> Subquery Scan on "\*SELECT\* 2"

-> Seq Scan on small small\_1

# HashSetOp



Logically UNION, EXCEPT, and INTERSECT without ALL remove duplicates in joined queries and results. For example, the query VALUES (1), (1), (2), (2) EXCEPT VALUES (1) returning 2 shows joined query removal, while UNION shows result removal.

## 10. SetOp

*-- table has to be too large to hash*

```
:EXPLAIN SELECT * FROM large INTERSECT SELECT * FROM large;  
QUERY PLAN
```

-----  
SetOp Intersect

-> Sort

Sort Key: "\*SELECT\* 1".x, "\*SELECT\* 1".y

-> Append

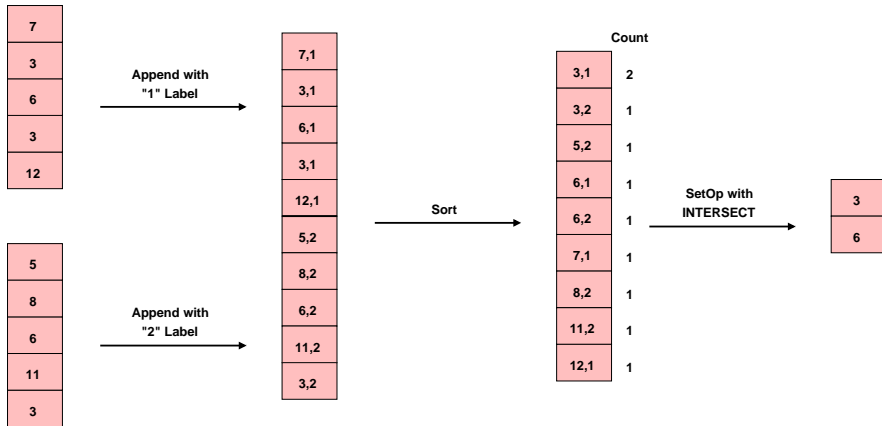
-> Subquery Scan on "\*SELECT\* 1"

-> Seq Scan on large

-> Subquery Scan on "\*SELECT\* 2"

-> Seq Scan on large large\_1

# SetOp



## 11. Materialize

```
:EXPLAIN SELECT * FROM small s1, small s2 WHERE s1.x != s2.x;  
      QUERY PLAN
```

-----

Nested Loop

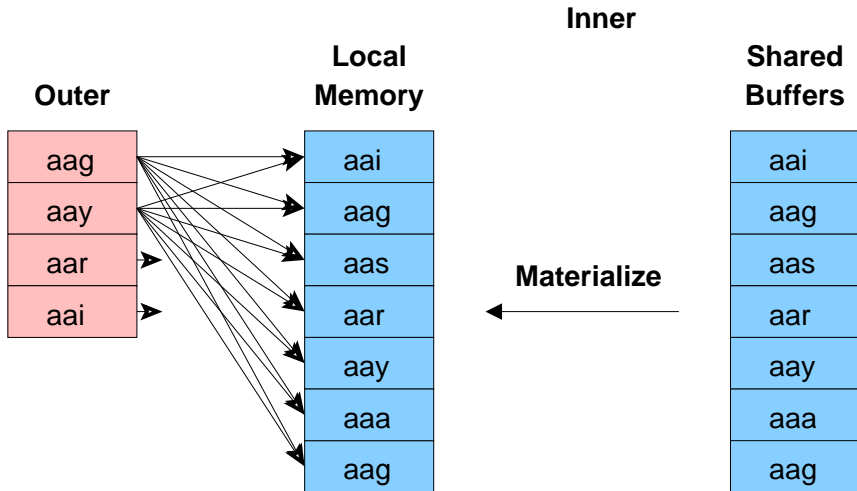
Join Filter: (s1.x <> s2.x)

-> Seq Scan on small s1

-> **Materialize**

    -> Seq Scan on small s2

# Materialize



## 12. Memoize, Setup

*-- needs duplicates and too small for a hash join*

```
CREATE TABLE small_with_dups (x) AS  
SELECT generate_series(1, 1000)  
FROM generate_series(1, 10);
```

*-- unique and too big for a hash join*

```
CREATE TABLE medium (x) AS  
SELECT generate_series(1, 100000);
```

*-- index required for this memoize example*

```
CREATE INDEX i_medium ON medium (x);  
ANALYZE;
```

# Memoize

```
:EXPLAIN SELECT * FROM small_with_dups JOIN medium USING (x);  
QUERY PLAN
```

---

Nested Loop

- > Seq Scan on small\_with\_dups

- > Memoize

  - Cache Key: small\_with\_dups.x

  - Cache Mode: logical

- > Index Only Scan using i\_medium on medium

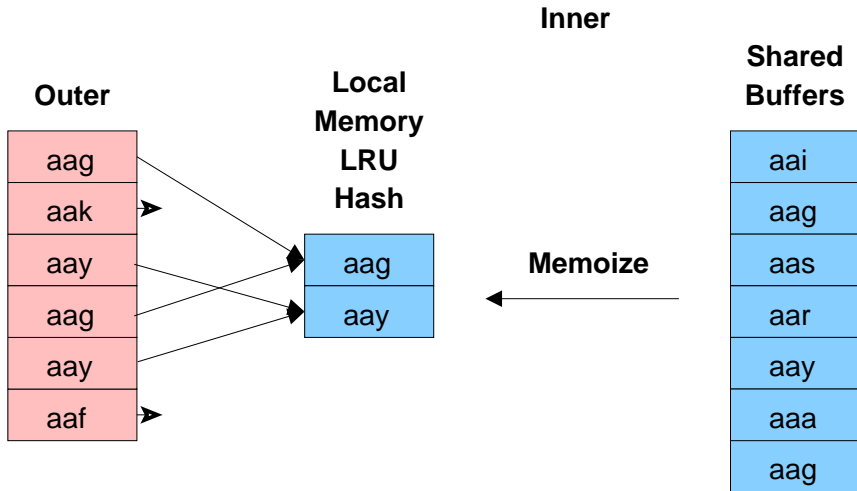
  - Index Cond: (x = small\_with\_dups.x)

Only happens in nested loops; supported in Postgres 14 and later.

[https://blog.jooq.org/postgresql-14s-enable\\_memoize-for-improved-performance-of-nested-loop-joins/](https://blog.jooq.org/postgresql-14s-enable_memoize-for-improved-performance-of-nested-loop-joins/)



# Memoize



## 13. Group, First Example

```
-- must be small enough not to trigger HashAggregate  
:EXPLAIN SELECT x FROM large WHERE x < 0 GROUP BY x;  
      QUERY PLAN
```

---

### Group

Group Key: x

-> Index Only Scan using i\_large on large  
 Index Cond: (x < 0)

## Group, Second Example

```
:EXPLAIN SELECT x FROM large GROUP BY x ORDER BY x;  
QUERY PLAN
```

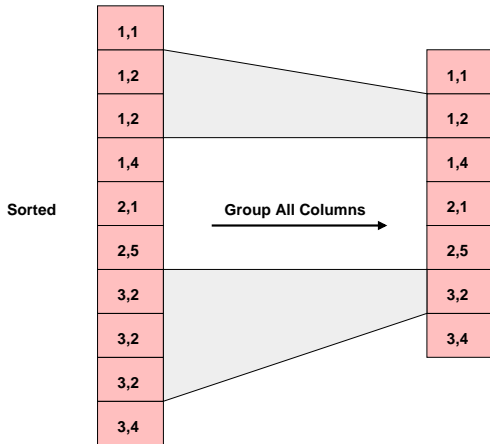
---

### Group

Group Key: x

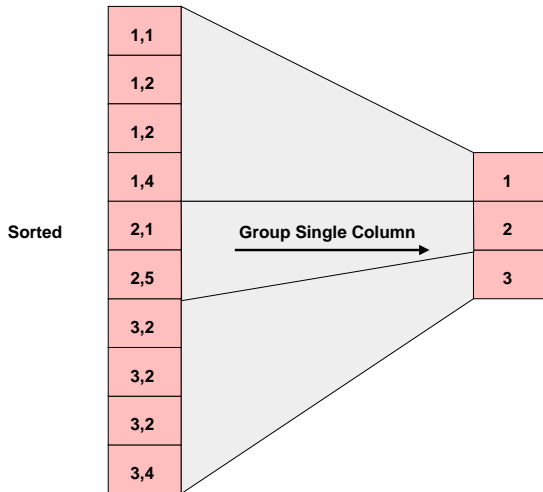
-> Index Only Scan using i\_large on large

# Group All Columns



**GROUP BY** without aggregates is similar to **SELECT DISTINCT**, except duplicate detection can consider more columns than those selected for output.

# Group Single Columns



## 14. Aggregate

```
:EXPLAIN SELECT COUNT(*) FROM medium;  
      QUERY PLAN
```

-----

Aggregate

-> Seq Scan on medium

## 15. GroupAggregate

```
:EXPLAIN SELECT x, COUNT(*) FROM medium GROUP BY x ORDER BY x;  
QUERY PLAN
```

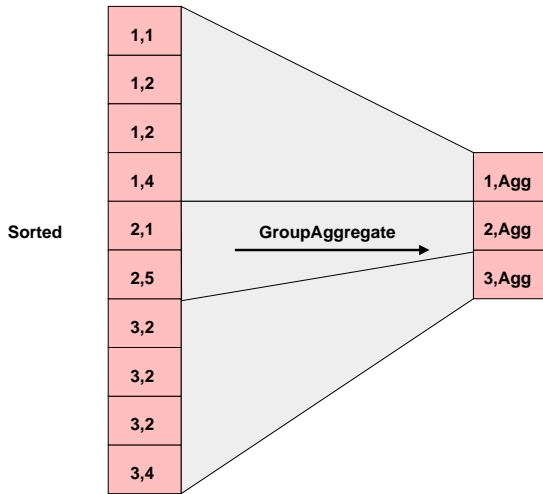
---

### GroupAggregate

Group Key: x

-> Index Only Scan using i\_medium on medium

# GroupAggregate





## 16. HashAggregate

```
:EXPLAIN SELECT DISTINCT x FROM medium;  
      QUERY PLAN
```

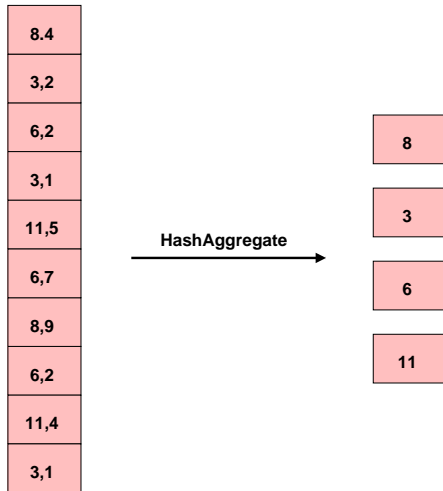
-----

HashAggregate

Group Key: x

-> Seq Scan on medium

# HashAggregate



## 17. MixedAggregate

```
:EXPLAIN SELECT x FROM medium GROUP BY ROLLUP(x);  
      QUERY PLAN
```

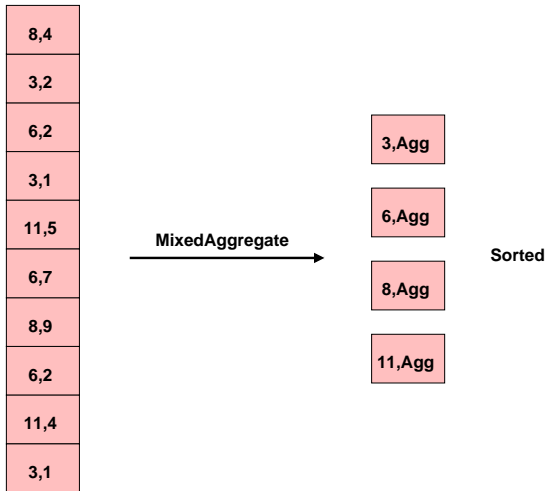
-----  
MixedAggregate

Hash Key: x

Group Key: ()

-> Seq Scan on medium

# MixedAggregate



## 18. WindowAgg

```
:EXPLAIN SELECT x, SUM(x) OVER ()  
FROM generate_series(1, 10) AS f(x);  
          QUERY PLAN
```

-----  
WindowAgg

-> Function Scan on generate\_series f

# WindowAgg

Sorted	1,1		1,1,Agg
	1,2		1,2,Agg
	1,2		1,2,Agg
	1,4		1,4,Agg
	2,1	WindowAgg →	2,1,Agg
	2,5		2,5,Agg
	3,2		3,2,Agg
	3,2		3,2,Agg
	3,2		3,2,Agg
	3,4		3,4,Agg

Window functions allow aggregates across rows while the individual rows remain.

## 19-22. Parallel Seq Scan, Partial Aggregate, Gather, Finalize Aggregate

```
:EXPLAIN SELECT SUM(x) FROM large;  
QUERY PLAN
```

---

Finalize Aggregate

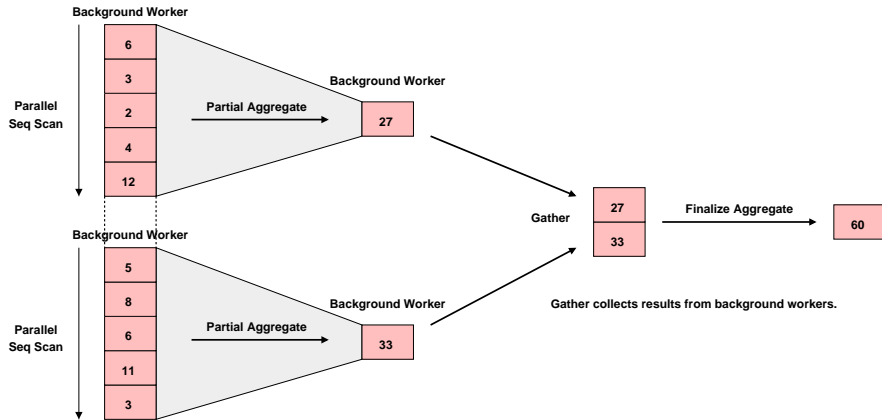
-> Gather

Workers Planned: 2

-> Partial Aggregate

-> Parallel Seq Scan on large

# Parallel Seq Scan, Partial Aggregate, Gather, Finalize Aggregate



Parallel Seq Scan uses background workers to scan different parts of a table in parallel.



## 23. Gather Merge

```
CREATE TABLE huge (x) AS SELECT generate_series(1, 100000000);  
ANALYZE huge;
```

```
:EXPLAIN SELECT * FROM huge ORDER BY 1;  
          QUERY PLAN
```

-----  
Gather Merge

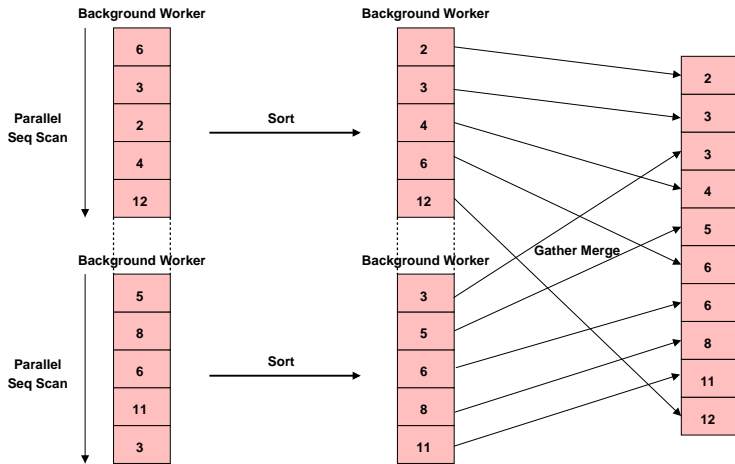
Workers Planned: 2

-> Sort

Sort Key: x

-> Parallel Seq Scan on huge

# Gather Merge



Gather Merge collects ordered results from background workers, retaining their ordering.

## 24. Parallel Append

```
:EXPLAIN SELECT * FROM huge UNION ALL SELECT * FROM huge ORDER BY 1;  
QUERY PLAN
```

-----  
Gather Merge

Workers Planned: 2

-> Sort

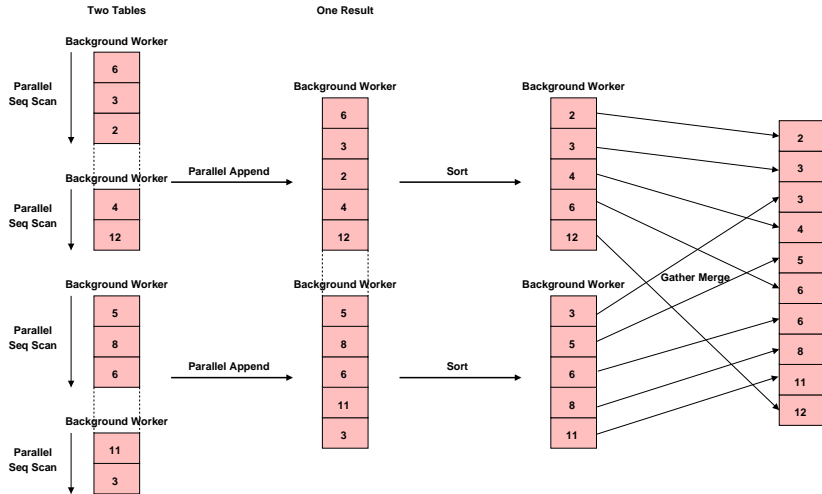
Sort Key: huge.x

-> Parallel Append

-> Parallel Seq Scan on huge

-> Parallel Seq Scan on huge huge\_1

# Parallel Append



## 25, 26. Parallel Hash, Parallel Hash Join

```
:EXPLAIN SELECT * FROM huge h1 JOIN huge h2 USING (x);  
      QUERY PLAN
```

-----  
Gather

Workers Planned: 2

-> Parallel Hash Join

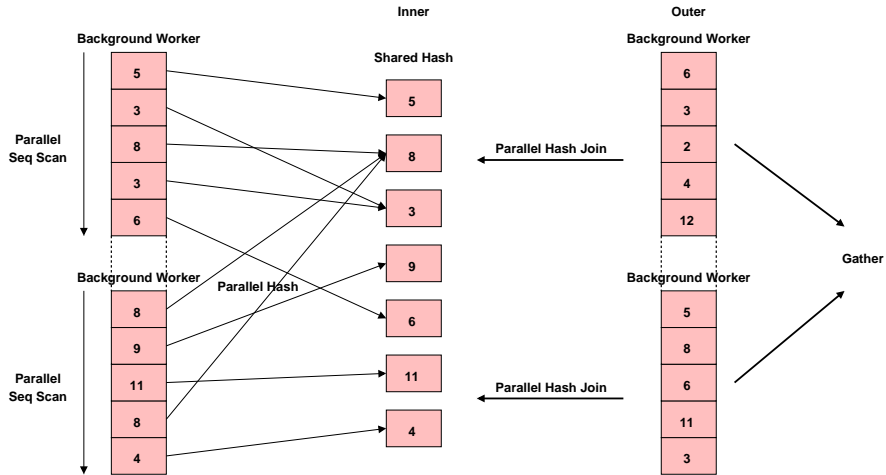
Hash Cond: (h1.x = h2.x)

-> Parallel Seq Scan on huge h1

-> Parallel Hash

-> Parallel Seq Scan on huge h2

# Parallel Hash, Parallel Hash Join

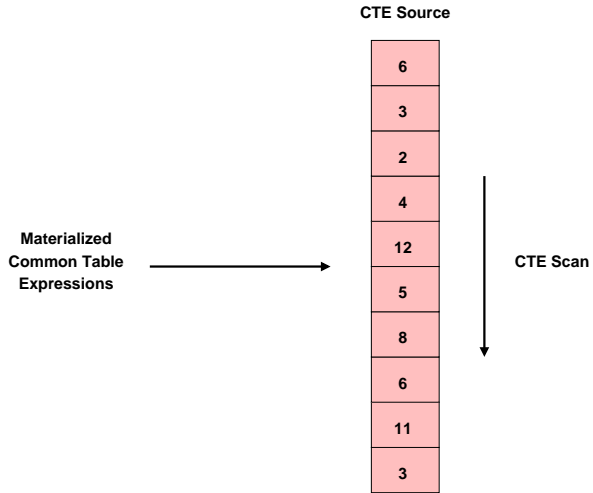


## 27. CTE Scan

```
:EXPLAIN WITH source AS MATERIALIZED (  
    SELECT 1  
)  
SELECT * FROM source;  
    QUERY PLAN
```

```
-----  
CTE Scan on source  
  CTE source  
    -> Result
```

# CTE Scan





## 28, 29. WorkTable Scan, Recursive Union

```
:EXPLAIN WITH RECURSIVE source (counter) AS (  
    SELECT 1  
    UNION ALL  
    SELECT counter + 1  
    FROM source  
    WHERE counter < 10  
)  
SELECT * FROM source;
```

QUERY PLAN

-----  
CTE Scan on source

CTE source

-> Recursive Union

-> Result

-> WorkTable Scan on source source\_1

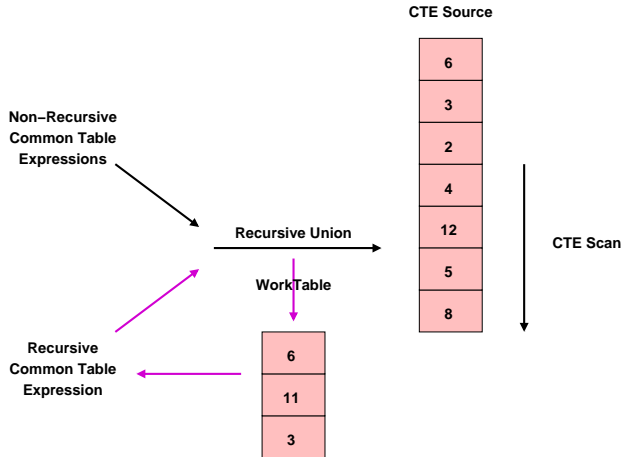
Filter: (counter < 10)

## CTE Query Flow

WITH RECURSIVE **source** AS (  
 SELECT 1  
 UNION ALL  
 SELECT 1 FROM **source**  
)  
SELECT \* FROM **source**;

```
graph TD; 1((1)) --> source1[source]; 2((2)) --> close_paren[)]; 3((3)) --> source2[source];
```

# WorkTable Scan, Recursive Union



WorkTable is cleared before every iteration. Recursion stops when the recursive CTE returns no rows.

## 30. ProjectSet

```
:EXPLAIN SELECT generate_series(1,4);
```

```
QUERY PLAN
```

```
-----
```

```
ProjectSet
```

```
-> Result
```

## 31. LockRows

```
:EXPLAIN SELECT * FROM small FOR UPDATE;
```

```
QUERY PLAN
```

```
-----
```

```
LockRows
```

```
-> Seq Scan on small
```

## 32. Sample Scan

```
:EXPLAIN SELECT * FROM small TABLESAMPLE SYSTEM(50);  
      QUERY PLAN
```

-----

```
Sample Scan on small  
  Sampling: system ('50'::real)
```

### 33. Table Function Scan

```
:EXPLAIN SELECT *  
FROM XMLTABLE('/ROWS/ROW'  
PASSING  
$$  
  <ROWS>  
    <ROW id="1">  
      <COUNTRY_ID>US</COUNTRY_ID>  
    </ROW>  
  </ROWS>  
$$  
COLUMNS id int PATH '@id',  
_id FOR ORDINALITY);  
      QUERY PLAN
```

-----  
Table Function Scan on "xmltable"

## 34. Foreign Scan

```
CREATE EXTENSION postgres_fdw;
```

```
CREATE SERVER postgres_fdw_test  
FOREIGN DATA WRAPPER postgres_fdw  
OPTIONS (host 'localhost', dbname 'fdw_test');
```

```
CREATE USER MAPPING FOR PUBLIC  
SERVER postgres_fdw_test  
OPTIONS (password '');
```

```
CREATE FOREIGN TABLE other_world (greeting TEXT)  
SERVER postgres_fdw_test  
OPTIONS (table_name 'world');
```

```
:EXPLAIN SELECT * FROM other_world;  
      QUERY PLAN
```

```
-----  
Foreign Scan on other_world
```

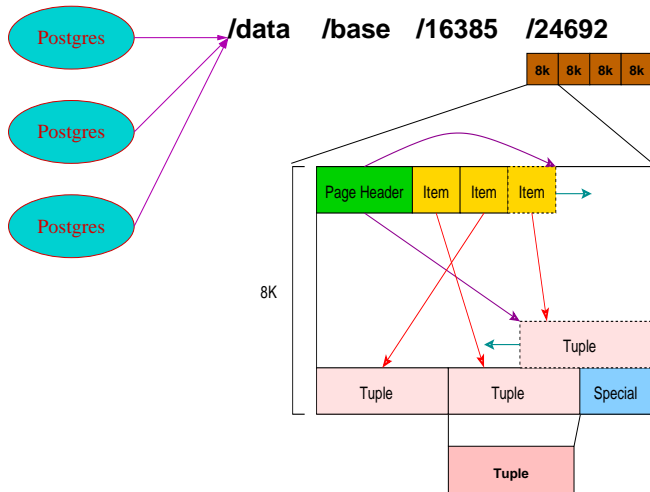


## 35. Tid Scan

```
:EXPLAIN SELECT * FROM small WHERE ctid = '(0,1)';  
      QUERY PLAN
```

```
-----  
Tid Scan on small  
  TID Cond: (ctid = '(0,1)::tid)
```

# Tid Scan



## 36. Insert

```
:EXPLAIN INSERT INTO small VALUES (0);
```

```
QUERY PLAN
```

```
-----
```

```
Insert on small
```

```
-> Result
```

## 37. Update

```
:EXPLAIN UPDATE small SET x = 1 WHERE x = 0;  
      QUERY PLAN
```

-----

Update on small

-> Seq Scan on small  
 Filter: (x = 0)

## 38. Delete

```
:EXPLAIN DELETE FROM small;  
      QUERY PLAN
```

```
-----
```

```
Delete on small
```

```
-> Seq Scan on small
```

*-- You cannot run EXPLAIN on utility commands like TRUNCATE.*

```
:EXPLAIN TRUNCATE small;
```

```
ERROR:  syntax error at or near "TRUNCATE"
```

```
LINE 1: EXPLAIN (COSTS OFF) TRUNCATE small;
```

```
      ^
```

## 39. Merge

```
CREATE TABLE mergetest (x, y) AS VALUES (1, NULL), (3, NULL), (5, NULL);
```

```
:EXPLAIN MERGE INTO mergetest  
USING (VALUES (1), (2), (3), (4), (5), (6)) m (x)  
ON mergetest.x = m.x  
WHEN NOT MATCHED THEN  
    INSERT (x) VALUES (m.x)  
WHEN MATCHED THEN  
    UPDATE SET y = TRUE;
```

QUERY PLAN

-----  
Merge on mergetest

-> Hash Right Join

Hash Cond: (mergetest.x = "\*VALUES\*".column1)

-> Seq Scan on mergetest

-> Hash

-> Values Scan on "\*VALUES\*"

## 40. Semi Join, First Example

```
:EXPLAIN SELECT *  
FROM small  
WHERE EXISTS (SELECT * FROM medium WHERE medium.x = small.x);  
QUERY PLAN
```

-----  
Hash Semi Join

Hash Cond: (small.x = medium.x)

-> Seq Scan on small

-> Hash

-> Seq Scan on medium

Stop scan after first inner match.

## Semi Join, Second Example

```
:EXPLAIN SELECT *  
FROM small  
WHERE small.x IN (SELECT medium.x FROM medium);  
QUERY PLAN
```

-----  
Hash Semi Join

Hash Cond: (small.x = medium.x)

-> Seq Scan on small

-> Hash

-> Seq Scan on medium

EXISTS and IN are equivalent in handling of NULLs because EXISTS only checks for row existence while IN logically does OR comparisons that can ignore non-true results from NULL comparisons.



## 41. Anti Join

```
:EXPLAIN SELECT *  
FROM medium  
WHERE NOT EXISTS (SELECT * FROM small WHERE small.x = medium.x);  
      QUERY PLAN
```

-----  
Hash Anti Join

Hash Cond: (medium.x = small.x)

-> Seq Scan on medium

-> Hash

-> Seq Scan on small

Stop scan after first inner match; negate result.

## 42. SubPlan

```
:EXPLAIN SELECT *  
FROM small  
WHERE small.x NOT IN (SELECT medium.x FROM medium);  
QUERY PLAN
```

```
-----  
Seq Scan on small  
  Filter: (NOT (hashed SubPlan 1))  
    SubPlan 1  
      -> Seq Scan on medium
```

NOT IN and NOT EXISTS are not equivalent for NULLs because NOT IN logically does repeated not-equal AND comparisons which must all be true to return true; NULL affects this.

## 43. Others: Outer Join Removal

*-- UNIQUE index guarantees at most one right row match*

```
CREATE UNIQUE INDEX i_small ON small (x);
```

*-- LEFT JOIN guarantees every left row is returned*

```
:EXPLAIN SELECT medium.x FROM medium LEFT JOIN small USING (x);
```

```
QUERY PLAN
```

```
-----
```

```
Seq Scan on medium
```

# Not Covered

- Named Tuplestore Scan: after triggers
- Custom Scan: custom scan providers

# Conclusion



<https://momjian.us/presentations>

<https://www.flickr.com/photos/glassholic/>