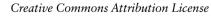
Major Features: Postgres 17

BRUCE MOMJIAN



POSTGRESQL is an open-source, full-featured relational database. This presentation gives an overview of the Postgres 17 release.

https://momjian.us/presentations





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Postgres 17 Feature Outline

- 1. Incremental backup
- 2. Improved data manipulation
- 3. Improved optimizer handling
- 4. Improved logical replicas

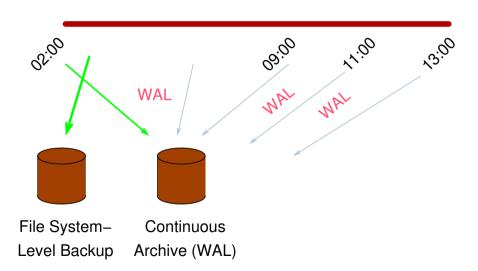
1. Incremental Backup: Backup Methods

Postgres supports three main backup methods:

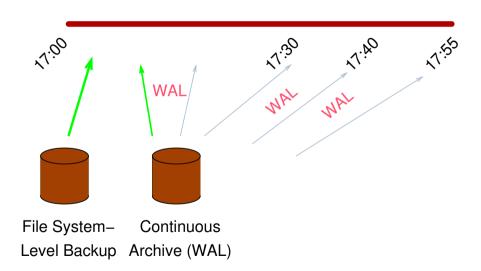
- File system backup/snapshot
- Logical backup, e.g., pg_dump
- Continuous archiving

This last method is preferred because it allows recovery to arbitrary times, including up to the most recent transactions. Replicas and delayed-replay replicas are more for failover than backup.

Continuous Archiving



Point-in-Time Recovery



Continuous Archiving Challenges

Continuous archiving requires a file system backup, even one taken while the database is active, plus write-ahead log (WAL) generated from the time of the backup to the restore time. This has some challenges:

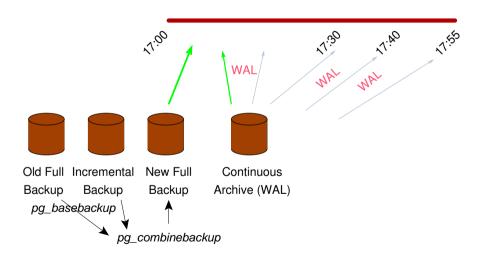
- Replay of the WAL from the time of the backup to the recovery time can be slow
- Replay time can be reduced by requiring FEWER WAL files from being processed
- This can be accomplished with more frequent file system backups
- Unfortunately file system backups are large and require a lot of I/O

Incremental Backups

Postgres's incremental backup feature solves these problems

- Can create an incremental file system backup via pg_basebackup
 - only records data blocks modified since the last full or incremental backup
- Combine a full backup with incremental backups to create a newer full backup, via pg_combinebackup
 - requires fewer WAL files to restore
 - faster restores
- Combining does not need access to the data directory so it can be done on a separate server
- Existing backup tools could already produce incremental backups, but with poorer granularity or more overhead, e.g. *pg_backrest*, *Barman*

Point-in-Time Recovery with Incremental Backup



pg combinebackup

- -- apply one incremental to create a new full backup -- full backup2 is now current as of the end of incremental1's backup \$ pg combinebackup full backup1 incremental1 -o full backup2
- -- apply three incrementals to create a new full backup \$ pg combinebackup full backup2 incremental2 incremental3 incremental4 -o full backup3

New Deployment Options

- Keep the base backup unchanged
 - create many incremental backups
 - incremental recovery is faster than WAL replay
 - discard WAL for time spans that don't need granular recovery
- Keep the base backup current
 - use *pg_combinebackup* to keep the base backup current by applying frequent incremental backups
 - this is only possible when desired recovery point is short
 - at a rapid frequency, it starts to have options like replicas, but with heavy I/O overhead
- Make a copy of the base backup
 - keep the copy current by applying incremental backups frequently
 - this can greatly reduce recovery time because there is minimal WAL to replay
 - this does not reduce the recovery time window because the old base backup, and necessary WAL, are kept

pg_combinebackup Removal

```
-- apply one incremental to create a new full backup

$ pg_combinebackup full_backup3 incremental5 -o full_backup4

-- If we don't need to recover to any time earlier than the end of incremental5,
-- we can delete the previous full backup and incremental5, and all needed WAL

$ rm -r full_backup3 incremental5

-- do it again

$ pg_combinebackup full_backup4 incremental6 -o full_backup5

$ rm -r full_backup4 incremental6
```

2. Improved Data Manipulation

- 1. JSON_TABLE()
- 2. Copy
- 3. Merge

2.1 JSON_TABLE()

```
SELECT *
FROM JSON TABLE('{"key1": "val1"}',
                '$.key1' COLUMNS (col1 text PATH '$');
 col1
val1
SELECT *
FROM JSON TABLE('{"key1": "val1", "key2": "val2"}',
                '$[*]' COLUMNS (key1 text PATH '$.key1', key2 text PATH '$.key2'));
 key1 key2
val1 | val2
```

Load JSONB Data

```
-- download sample data from https://www.mockaroo.com/
-- remove 'id' column, output as JSON, uncheck 'array'
CREATE TABLE friend (id SERIAL, data JSONB):
COPY friend (data) FROM '/tmp/MOCK DATA.json';
SELECT *
FROM friend
ORDER BY 1
LIMIT 2;
id l
                               data
 1 | {"email": "sbouzan0@wikispaces.com", "gender": "Female", ...
 2 | {"email": "ebruffell1@independent.co.uk", "gender": "Male", ...
```

Pretty Print JSON

```
SELECT id, jsonb pretty(data)
FROM friend
ORDER BY 1
LIMIT 1;
 id l
                   jsonb pretty
          "email": "sbouzan0@wikispaces.com",+
          "gender": "Female".
          "last name": "Bouzan",
          "first name": "Sher",
          "ip address": "89.153.16.253"
```

JSON_TABLE()

first_name	last_name	email
Abbey	Terbeck	aterbeckf7@latimes.com
Giustino	Weeke	gweekerm@fastcompany.com
Bailey	Romi	bromibc@desdev.cn
Guillemette	Hastwell	ghastwell61@microsoft.com
Cherise	Biermatowicz	cbiermatowicz3c@google.com.hk

2.2 COPY Error Handling

```
CREATE TABLE copy test (int col INTEGER, date col DATE, jsonb col JSONB);
COPY copy test (int col) FROM STDIN;
test> 1
test> 2
test> x
test> \.
ERROR: invalid input syntax for type integer: "x"
CONTEXT: COPY copy test, line 3, column int col: "x"
SELECT *
FROM copy test;
 int col \overline{|} date col | jsonb col
```

COPY Ignore Errors

```
COPY copy test (int col) FROM STDIN WITH (ON ERROR ignore);
test> 3
test> 4
test> a
test> \.
NOTICE: 1 row was skipped due to data type incompatibility
SELECT * FROM copy test;
 int col | date col | jsonb col
          (null) | (null)
           (null) | (null)
```

COPY Report Error Rows

```
COPY copy test (int col) FROM STDIN WITH (ON ERROR ignore, LOG VERBOSITY verbose);
test> 5
test> 6
test> b
test> \.
NOTICE: skipping row due to data type incompatibility at line 3 for column "int col": "b"
NOTICE:
        1 row was skipped due to data type incompatibility
SELECT * FROM copy test;
 int col | date col | jsonb col
           (null)
                     | (null)
           (null)
                    | (null)
           (null) | (null)
           (null)
                      (null)
```

2.3 MERGE Improvements

- Allow MERGE to modify updatable views
- Add when not matched by source
- Allow MERGE to use the RETURNING clause

3. Improved Optimizer Handling

- 1. CTE passdown
- 2. NULL optimizations
- 3. Correlated subquery optimization
- 4. Other

3.1 CTE Passdown

```
WITH RECURSIVE dep (classid, obj) AS (
        SELECT (SELECT oid FROM pg class WHERE relname = 'pg class'),
                oid
        FROM pg class
        WHERE relname = 'deptest'
        UNION ALL
        SELECT pg depend.classid, objid
        FROM pg depend JOIN dep ON (refobjid = dep.obj)
) -- statistics and sort order are passed down here
SELECT (SELECT relname FROM pg class WHERE oid = classid) AS class,
        (SELECT typname FROM pg type WHERE oid = obj) AS type,
        (SELECT relname FROM pg class WHERE oid = obj) AS class,
        (SELECT relkind FROM pg class where oid = obj::regclass) AS kind,
        (SELECT pg get expr(adbin, classid) FROM pg attrdef WHERE oid = obj) AS attrdef,
        (SELECT conname FROM pg constraint WHERE oid = obj) AS constraint
FROM dep
ORDER BY obj;
```

3.2 NULL Optimizations

```
CREATE TABLE null test (not null col INTEGER NOT NULL,
                       opt null col INTEGER);
-- This disables EXPLAIN cost output
\set EXPLAIN 'EXPLAIN (COSTS OFF)'
:EXPLAIN SELECT *
FROM null_test;
     OUERY PLAN
Seq Scan on null test
-- no sequential scan
:EXPLAIN SELECT *
FROM null test
WHERE not null col IS NULL;
        QUERY PLAN
Result.
  One-Time Filter: false
```

NULL Optimizations

```
:EXPLAIN SELECT *
FROM null test
WHERE opt null col IS NOT NULL;
             QUERY PLAN
Seq Scan on null test
  Filter: (opt null col IS NOT NULL)
-- no 'Filter' clause
:EXPLAIN SELECT *
FROM null test
WHERE not null col IS NOT NULL;
     QUERY PLAN
Seq Scan on null_test
```

3.3 Correlated Subqueries in Postgres 16

```
:EXPLAIN SELECT COUNT(*)
FROM pg class
WHERE oid IN (
        SELECT attrelid
        FROM pg attribute
        WHERE attrelid = pg class.oid
);
                                     QUERY PLAN
Aggregate
   -> Seg Scan on pg class
         Filter: (SubPlan 1)
         SubPlan 1
           -> Index Only Scan using pg attribute relid attnum index on pg attribute
                 Index Cond: (attrelid = pg class.oid)
```

Correlated Subqueries Now as Joins

```
:EXPLAIN SELECT COUNT(*)
FROM pg class
WHERE oid IN (
        SELECT attrelid
        FROM pg attribute
        WHERE attrelid = pg class.oid
);
                                 QUERY PLAN
Aggregate
   -> Hash Join
        Hash Cond: (pg class.oid = pg attribute.attrelid)
         -> Seg Scan on pg class
         -> Hash
               -> HashAggregate
                     Group Key: pg attribute.attrelid, pg attribute.attrelid
                     -> Seg Scan on pg attribute
```

3.4 Other Optimizer Improvements

- Partition pruning for boolean columns
- Range value containment
- Optimize LIMIT on partitions
- Allow GROUP BY to be reordered to match ORDER BY
- Improve Merge Append *
- More parallelism

4. Improved Logical Replication

- Add tool pg_createsubscriber to allow creation of logical replicas from physical ones
- Enable failover of logical replication slots
- Enable pg_upgrade to preserve logical replication slots in future major upgrades

Conclusion

