# An Evaluation of the Rotor-Router Mechanism

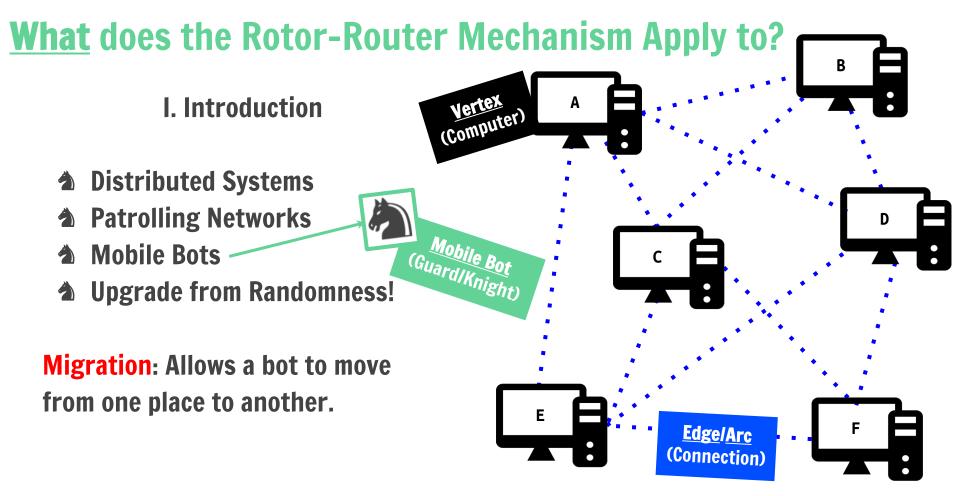
Julian M. Rice | UCLA <u>FrontierLab@OsakaU</u>: Graduate School of Information Science & Technology Masuzawa Laboratory (增澤室)

Total Slide Count: 14 | August 2018

# Research Background & Motivation

**Looking Behind the Rotor-Router Mechanism** 

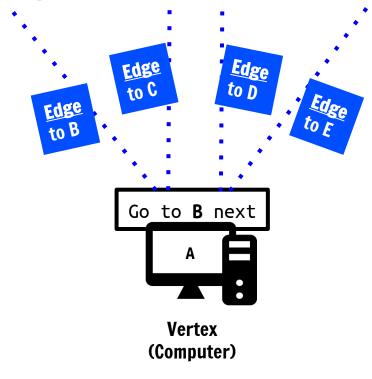
[ Research Background -> Simulation & Programming -> Results & Conclusion -> Questions ]



### **II. Rotor-Router Mechanism**

- **★** Alternative to Random Walk
- **▲** Each Vertex Points to an Edge
  - $\checkmark$  A (go to B, C, D, E, then B..)
- **△** Deterministic Algorithm
- **▲** Euler Cycle (Lock-In)

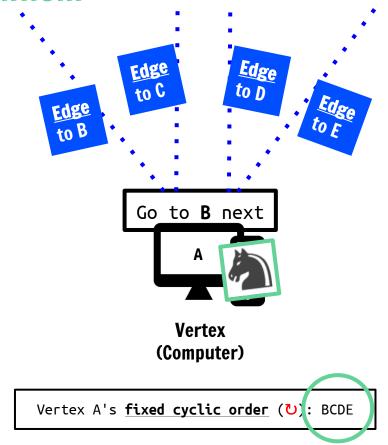
Visiting: When a bot arrives at a vertex or crosses an edge.



Vertex A's <u>fixed cyclic order</u> (♥): BCDE

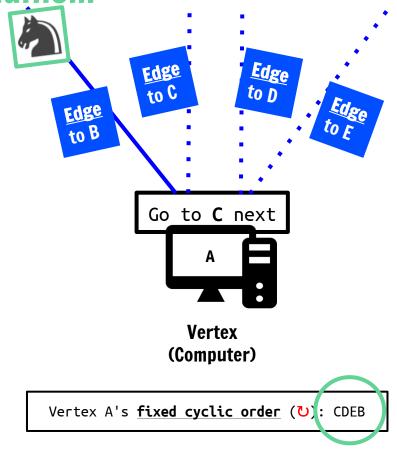
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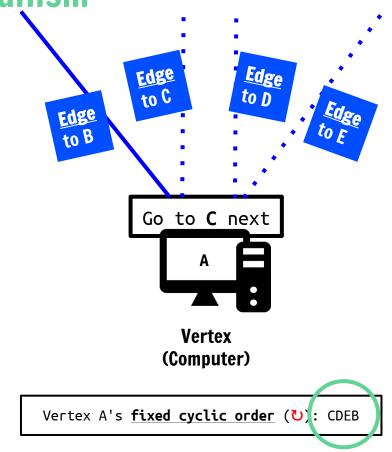
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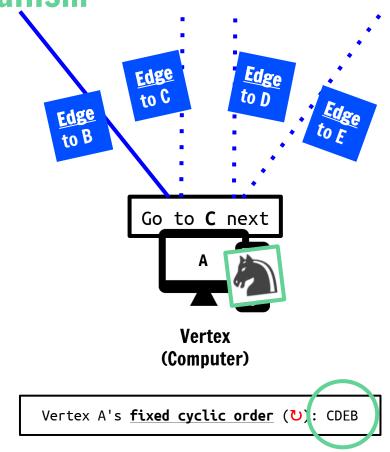
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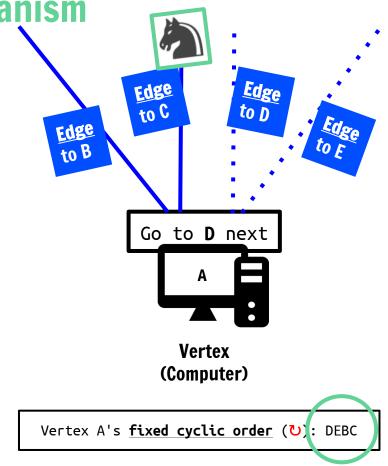
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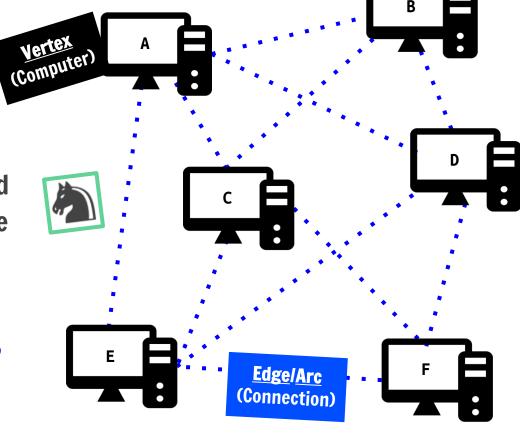


# The Rotor-Router Mechanism's **KEY Element**

**Euler Cycle:** A path where every edge is visited *exactly once*.

The bot will eventually end up entering a Euler cycle after a period of time (*exploring the network*), the stabilization period, passes.

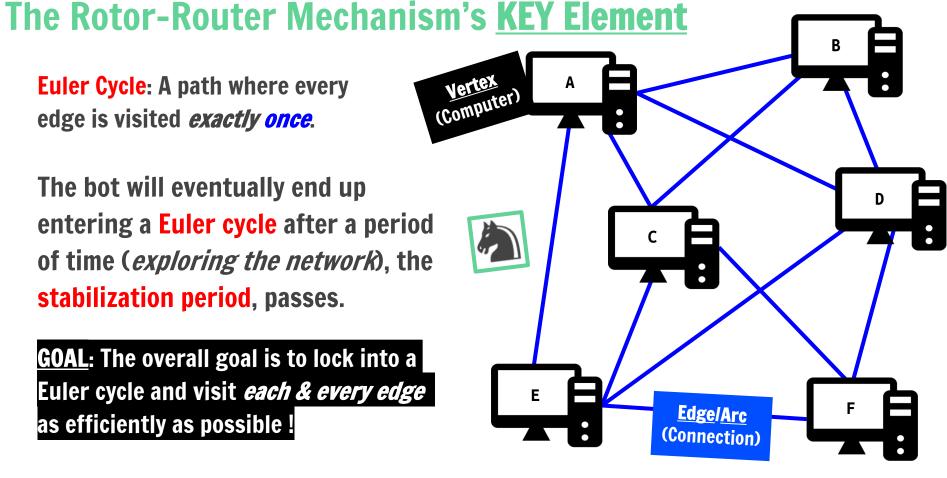
**GOAL**: The overall goal is to lock into a **Euler cycle** and visit **each & every edge** as efficiently as possible.



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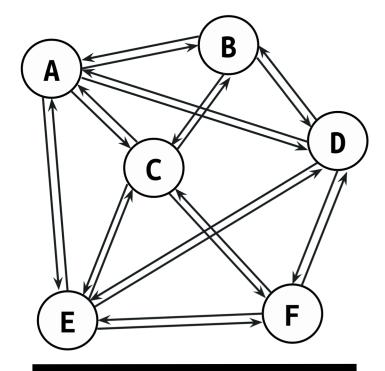


# Preparing to **Observe** the Rotor-Router Algorithm

I created a custom "network" with arrows (arcs, rather than edges) that start at one vertex and end at another.

The supposed limit is  $4 \cdot m \cdot D$ , or the number of edges (m = 11) multiplied by the diameter of the graph (D = 2)  $\Rightarrow$  44 steps

What are the **main questions** that rise from this algorithm, and **why perform a simulation**?

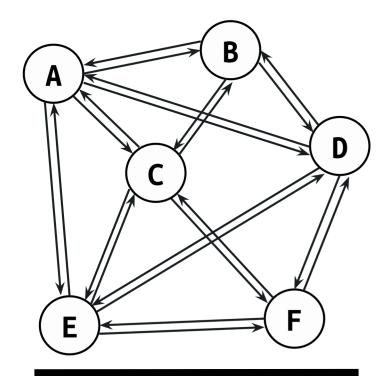


Simplified visual of the previous diagram

# Preparing to **Observe** the Rotor-Router Algorithm: **Mysteries**

The supposed limit is  $\mathbf{4} \cdot \mathbf{m} \cdot \mathbf{D}$ , or the number of edges (m = 22) multiplied by the diameter of the graph (D = 2)  $\Rightarrow$  44 steps

Mysteries (Why, Why, Why?)
[1] No matter where the bot starts, it will always end up in a forever lasting loop.
[2] The worst case is 44 steps (for this example), but what is the average case? Is it close to 44, close to 0, or in the middle?



Simplified visual of the previous diagram

# **Simulations & Programming**

What simulations did I run to get the data I needed?

[ Research Background -> Simulation & Programming -> Results & Conclusion -> Questions ]

Input:

# **Simulation Program (C++)**

Files: Edge.h, Vertex.h, Graph.h, StatList.h, Functions.h, Edge.cpp, Vertex.cpp, Graph.cpp, StatList.cpp, Functions.cpp, main.cpp

```
mobileAgent(Graph graph, Vertex* start, int iterations, string& temp) {
                    Vertex* destination = NULL:
                   for (int i = 0: i < iterations: i++) {
                       int currentIndex = start->getIndex();
                      start->edgeToCross(start, currentIndex)->incrementVi: -RESULTS-
                      -Number of TESTS: 1
                      cout << "[Step " << i+1 << "] Agent crosses the EDGE</pre>
                      << cross->getStartVertex()->getName() << cross->getEnding
                                                                  Average Visited Count for Vertices
                       << " and is now at VERTEX " << cross->getEndVertex()
                                                                    A: 6
                      graph.insertEdgeToOverall(cross); //New
                                                                    B: 3
                      destination->incrementVisited();
start->changeIndex();
                                                                    C: 5
                       start = destination:
                                                                    D: 6
[Console] Please input the number of times the mobile agent
                                                                     E: 6
[Console] Please input the number of times the test will be
                                                                     F: 4
[Step 1] Agent crosses the EDGE BA and is now at VERTEX A
                                                                  Average Visited Count for Edges
[Step 2] Agent crosses the EDGE AE and is now at VERTEX E
                                                                    AB: 3
[Step 3] Agent crosses the EDGE EC and is now at VERTEX C
                                                                    AC: 2
[Step 4] Agent crosses the EDGE CE and is now at VERTEX E
                                                                    AD: 3
[Step 5] Agent crosses the EDGE ED and is now at VERTEX D
                                                                    AE: 3
[Step 6] Agent crosses the EDGE DF and is now at VERTEX F
                                                                    BC: 2
[Step 7] Agent crosses the EDGE FD and is now at VERTEX D
                                                                    BD: 2
[Step 8] Agent crosses the EDGE DA and is now at VERTEX A
                                                                    CE: 4
[Step 9] Agent crosses the EDGE AB and is now at VERTEX B
                                                                    CF: 2
[Step 10] Agent crosses the EDGE BC and is now at VERTEX C
```

```
# of steps (for the agent)
        # of tests (for accuracy)
Output: Average visit count (edge)
```

**Average visit count (vertex)** Stabilization period (average, min, max, distr) **Euler cycle (repeated forever)** 

Stabilization Period Data Worst Case: 12 Average Case: 1.24542

**Line Count: 673 (11 files)** 

Best Case: 0

All Euler Cycles

Cycle 1: AE->EC->CE->ED->DF->FD->DA->AB->BC->CF->FE->EF->FC->CA->AC->CB->BD->DB->BA->AD->DE->EA->

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Julian M. Rice

What results did I get from running 1,000,000 tests on a graph?

[ Research Background -> Simulation & Programming -> Results & Conclusion -> Questions ]

## **Results**: What did I find out?

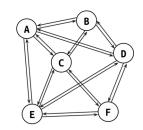
### II. Stabilization Period & Euler Cycle

\*\*Data that applies to only one graph\*\*

### (Stabilization Period)

- **★ Way lower** than **44 step** limit.
- **1** Often instantly locks in (49%)
- Random initialization (Euler Cycle)
- **△** Always different (each test)

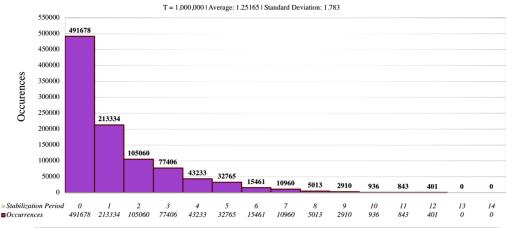
Important: There is **no way** to predict what the stabilization period average by simply looking at a network. **My simulation** brought these hidden statistics to light.



#### Stabilization Period Statistics

Locking in 2 or fewer steps => 81.01% Locking in 5 or fewer steps => 96.35% Locking in 10 or fewer steps => 99.88% Average Lock-in Time => 1.25 steps

#### Stabilization Period Distribution



#### **Sample Euler Cycles** (T = 1,000,000)

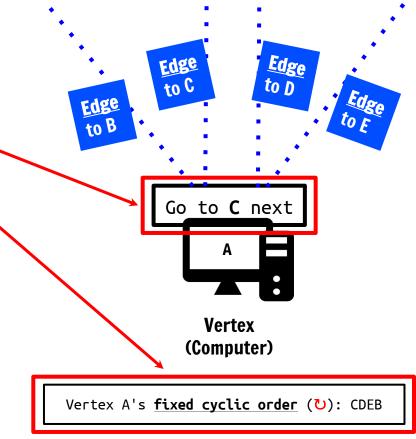
Paper 2: Ilcinkas Paper 3: Menc Paper 4: Chalopin

# **Upgrading the Rotor-Router Mechanism**

### **Enhancing the RR Mechanism**

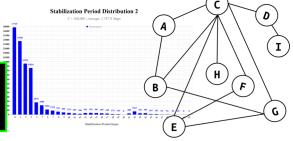
(Initialization of a Graph)

- Customizing the starting value
- Customizing the fixed order (Extra Additions to the Agent)
- **⚠** Making the agents **smarter**
- ★ Giving agents a mind (memory) (Increasing the Quantity)
- **★** Multiple Agents in a Graph



# **Conclusion**

Stabilization Period Data Worst Case: 30 Average Case: 2.78775 Best Case: 0

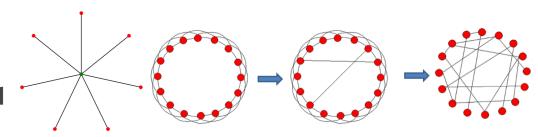


### (Thoughts & the Good Parts)

- **♦** Seeing the Rotor-Router Mechanism in action = !!
- **★** Writing the simulation from the ground up to obtain data
- **△** Creating animations, figures to understand the algorithm = !!
- **⚠** Understanding the reliability behind RR, and the details behind mobile agents, networks

### (The Further-Research Part)

- **♦** Simulations on a variety of graphs (stars, small world, etc)
- **Designing a random graph** generator that works (!!)
- **▲** Investigating stabilization period & more with *multiple agents*.
- **⚠** Creating a simulation program for newly proposed versions of the Rotor-Router mechanism (!!!!)



### (Some) References

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Total: ~140 pages

# Questions (?)

From the mechanism itself to how I ran simulations, please feel free to ask me anything!

Research Background -> Simulation & Programming -> Results & Conclusion -> Questions ]