

ICS 103

Data Structures and Algorithms

International Institute of Information Technology

Hyderabad, India

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# Motivation

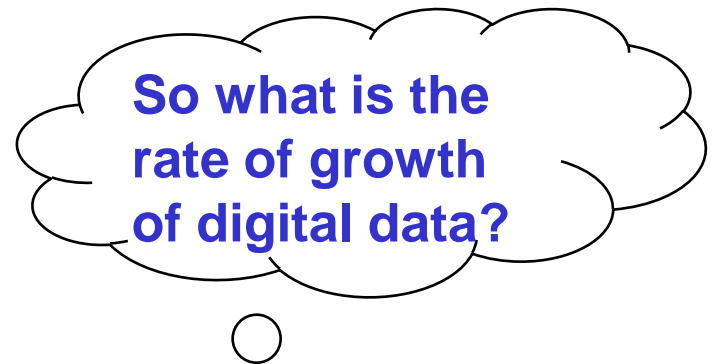
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- Amount of data that is being handled is getting huge.
- Name some examples in order of scale
  - About 1 – 5 KB
  - About 1 – 5 MB
  - About 1 – 5 GB
  - About 1 – 5 TB
  - About 1 – 5 PB
  - About 1 – 5 EB
- It is believed that every year we produce as many bits of information as is available in the past.

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- Amount of data that is being handled is getting huge.
- Name some examples in order of scale
  - About 1 – 5 KB : Compiler symbol table
  - About 1 – 5 MB : Source code of a big application
  - About 1 – 5 GB : Telephone directory
  - About 1 – 5 TB : Digital library
  - About 1 – 5 PB : Google maps
  - About 1 – 5 EB : Search engine data
- It is believed that every year we produce as many bits of information as is available in the past.



So what is the  
rate of growth  
of digital data?

# Motivation

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- What do we do with all this data?
- Think of the contacts list in your mobile phone.  
Typically,
  - look for a person by name and find his/her telephone number.
  - Once in a while, add/remove names to the list.
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  - Binary search? **Not really.**
  - We will know the answer in this course.

# Motivation

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- Consider an online application such as Google maps.
- Gives you driving directions from say IIT-H to your home.
- What is the size of this data? *468 Terabytes according to one count.*
- How can it be done?
- How to store the information? In this case, a graph.
- How to quickly find the route given two points A and B?

# Motivation

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- Need mechanisms to store the data and also to efficiently access data.
- The study of such mechanisms forms the subject matter of Data Structures.
- A fundamental part of any Computer Science curriculum.
  - several practical issues being addressed even today in important conferences.

# About this Course

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- We will cover several fundamental data structures including:
  - Arrays
  - Stacks and queues
  - Hash tables
- Other pointer based data structures such as
  - lists
  - trees, heaps
- Special data structures such as:
  - Graphs
  - Amortized data structures

# Items to Consider

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- Will introduce practical motivations to each of the considered data structures.
- Several problem solving sessions to fully understand the implications of using a data structure.
- Emphasis also on correctness and efficiency.
- Elementary analysis
- A basic introduction to **parallelism in computing** and also **parallel programming**.
  - Laboratory sessions are therefore very important.

# What is NEW this Year?

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- Class duration increased to 85 minutes from 55 minutes.
- Co-taught with Dr. Vikram Pudi.
- Term paper/Project as part of the grade.
  - Excellent ones can be extended to summer project, possibly leading to a publication.
  - So, those of you wishing to do research should aim high.
- Plan to include online surveys.

# Yet Another Look at the Syllabus

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- Syllabus by week
- Basic Data Structures
  - Processing integers (no need for data structures explicitly)
  - Analysis of algorithms
  - The need for data structures
  - The Need for Different access patterns on arrays
  - Limitations of array based data structures
- Intermediate advanced data structure
  - Hashing
  - Trees

# Yet Another Look at the Syllabus

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- Advanced data structures
  - Data structures for graphs
  - Same as week 10
  - Advanced Topics -- I
  - Advanced Topics -- II
  - Advanced Topics -- III



# Other Policies

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- Weekly three lecture hours.
- One hour of tutorial.
- Laboratory session every week
  - about 2-3 problems to be solved
  - TAs to assist.
- Several homework assignments
  - About 7, one every two weeks.
  - Each set to have about 6-7 problems
  - Late submission not allowed, unless notified earlier.
- Strictly, no plagiarism
  - Any detected case of plagiarism to be taken seriously.

# Other Policies

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- Instructor available via office hours.
- Seek an appointment for meeting outside of office hours.
- Email communication is also OK.
- Very important: **Seek help early enough.**

# Other Policies

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- Grading scheme
  - Homework – 15%
  - Mid term exam -1 – 20 %
  - Mid term exam -2 – 15 %
  - lab exam -1 – 5 %
  - lab exam -2 – 10 %
  - \*Term paper/Project – 10 % (Extra credit of 10 % for exceptional work)
  - End term exam – 25 %
- Subject to minor changes.

# A Complete Example – Number Systems

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- An example to illustrate that data structures are all pervasive.
- We will consider number systems.
- Number systems are a way to represent numbers
  - Using the representation, can do arithmetic on numbers.
  - Ability to count and do arithmetic is a fundamental civilizational trait.
  - Ancient civilizations also practised different number systems with different characteristics.

# Number Systems

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- A number system is a way to represent numbers.
- Several known number systems in practice even today.
  - Hindu/Decimal/Arabic system
  - Roman system
  - Binary, octal, hexa-decimal.
  - Unary system
  - ...
- A classification
  - positional
  - non-positional

# Number Systems

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- Hindu/Decimal system
  - Numbers represented using the digits in  $\{0, 1, \dots, 9\}$ .
  - Example: 8,732,937,309
- Roman System
  - Numbers represented using the letters I, V, X, L, C, D, and M.
  - For instance X represents 10, L represents 50.
  - LX stands for 60, VII stands for 7, **what is MMXIII?**
  - MMMDDCCCLLLXXXVVVIII – largest numbers without any overlines/subtractions. What is this number?
- Binary system
  - Numbers represented using the digits 0 and 1.
  - 10111 represents 23.

# Number Systems

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- Positional (aka value based) number systems associate a value to a digit based on the its position.
  - Example: Decimal, binary, ...
- Non-positional do not have such an association.
  - Example: Unary

# Operations on Numbers

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- Let us consider operations addition and multiplication.
- Hindu/Decimal system
  - Add digit wise
  - Carry of  $x$  from digit at position  $k$  to position  $k+1$  equivalent to a value of  $x \cdot 10^{k+1}$ ,  $k > 0$ .
  - Example: Adding 87 to 56 gives 143.
- Unary system
  - Probably, the first thing we learn.
  - To add two numbers  $x$  and  $y$ , create a number that contains the number of 1's in both  $x$  and  $y$ .
  - Example: Adding 1111 to 11111 results in 111111111.



# Operations on Numbers

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- Roman system
  - A bit complicated but possible.
  - Follow the following three steps:
    - Write the numbers side by side.
    - Arrange the letters in decreasing order of value.
    - Simplify.
  - Example: to add 32 and 67:
    - 32 = XXXII, 67 = LXVII.
    - XXXIILXVII
    - LXXXXVIII – LXLIX – XCIX
    - Simplified as: XCIX

# Operations on Numbers

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- Rules such as:
  - If there are 4I's, write it as IV.
  - If there are 4X's, write it as XL.
  - Similar rules apply.
- Careful when starting with numbers such as LXIV.
  - Can replace IV with IIII initially.

# Operations on Numbers

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- Let us now consider multiplication.
- Typically, multiplication is achieved by repeated addition.
- Decimal system
  - Known approach.
- Roman system
  - How to multiply?
  - Much complicated, but is possible.

# Multiplication in Roman Numerals

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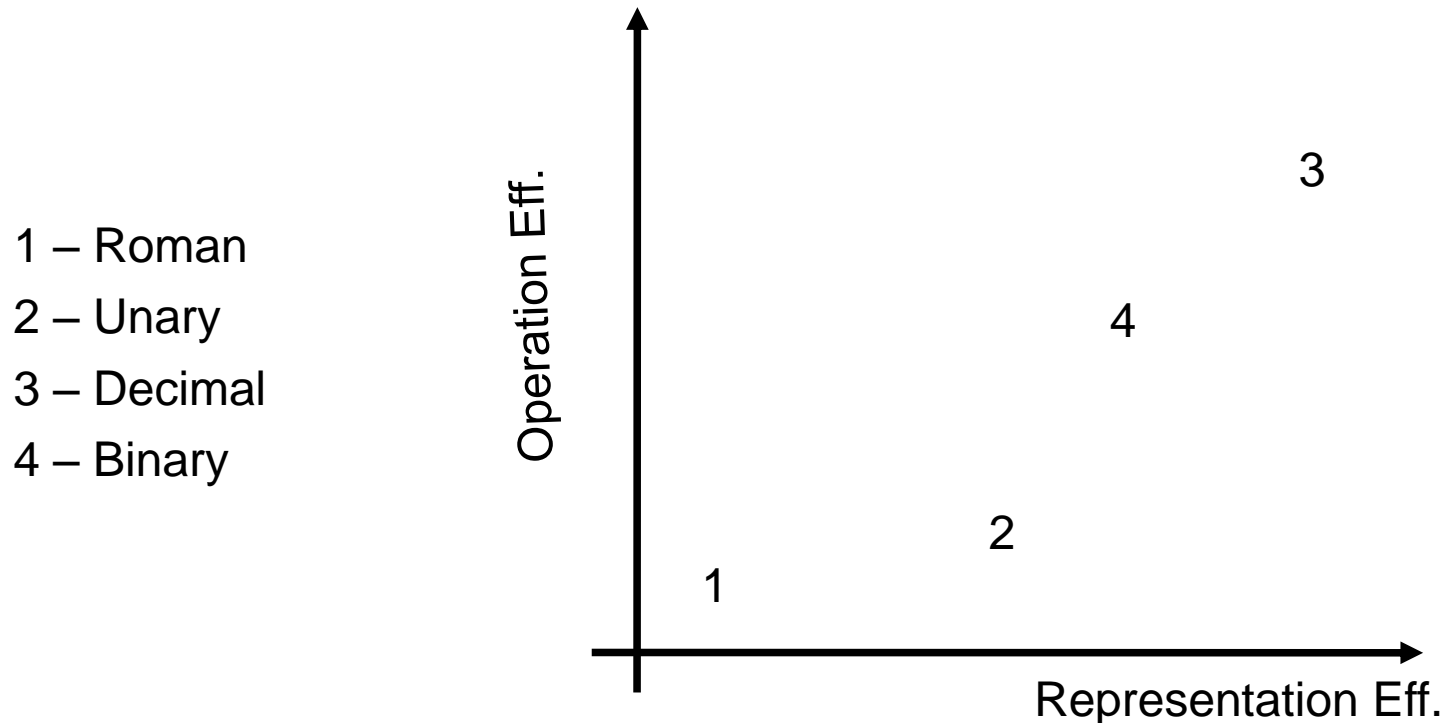
- Easy to imagine the following approach.
  - Multiplication is repeated addition
- Plus, think of a Roman number as the addition of 1000's + 100's + 50's + 10's + 5's + 1's.
- Multiply by each of these, and add as earlier.
- Example: LXII x XXXVII (62 x 37)
  - Multiply each of LXII by II. Meaning, make 2 copies of each symbol in LXII as LLXXIIII
  - Simplify using the rules of addition to CXXIV.
  - Now, multiply each of LXII by V. Start with LLLLLXXXXXIIIIIIIIII, simplify to CCCX.
  - Multiply LXII by XXX. That can be done in two ways. Either multiply by 3 followed by 10, or directly.

# Multiplication in Roman Numerals

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- Continuing the example,
  - Let us multiply by LXII by 3 as LLLXXXIIIIII and simplified as CLXXXVI.
  - Now, multiply CLXXXVI by 10 as  
 CCCCCCCCCCLLLLLLLLLLXXXXXXXXXXXXXXXXXXXX  
 XXXXXXXXXXXXXXXXXXXXVVVVVVVVVVVVIIIIIIIIII and  
 simplified as MDCCCLX.
  - Add all the constituents as CXXIV + CCCX +  
 MDCCCLX = CDXXXIV + MDCCCLX =  
 MMCCXCIV.
  - What is this number?

# Lesson Learnt



- Representation scheme for numbers influences the ease of performing operations.
- Roman system quite difficult to use.
- There are other such systems not in use today.

# Are There Other Representation Formats?

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- Yes, recall the fundamental theorem of arithmetic.
- Any number can be expressed uniquely as a product of primes.
- So, a product of primes representation is also possible.
- Not easy to add though.

# Further Operations

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- Let us now fix the decimal system as the representation scheme.
- We will now focus on the efficiency of operations.
- Let us see further operations such as finding the GCD of two numbers.



# GCD

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- Given two positive numbers,  $x$  and  $y$ , the largest number that divides both  $x$  and  $y$  is called the greatest common divisor of  $x$  and  $y$ . Denoted  $\text{gcd}(x,y)$ .
- Several approaches exist to find the gcd.
- Approach 1 : List all the divisors of both  $x$  and  $y$ . Find the common divisors, and the largest among the common divisors.
- Example for Approach 1:  $x = 24$ ,  $y = 42$ ,
  - divisors of 24 are  $\{1, 2, 3, 4, 6, 8, 12, 24\}$ .
  - divisors of 42 are  $\{1, 2, 3, 6, 7, 14, 21, 42\}$ .
  - Common divisors are  $\{1, 2, 3, 6\}$ . Hence,  $\text{gcd}(24, 42) = 6$ .

# GCD – Approach II

- Use the fundamental theorem of arithmetic and write  $x$  and  $y$  as:

$$- x = p_1^{a_1} \cdot p_2^{a_2} \cdots p_k^{a_k}$$

$$- y = p_1^{b_1} \cdot p_2^{b_2} \cdots p_r^{b_r}$$

$$- \text{It holds that } \gcd(x, y) = p_1^{\min\{a_1, b_1\}} \cdot p_2^{\min\{a_2, b_2\}} \cdots p_r^{\min\{a_r, b_r\}}.$$

- Example Approach II, let  $x = 24$ ,  $y = 42$ .

$$- x = 2^3 \cdot 3, y = 2 \cdot 3 \cdot 7.$$

$$- \gcd(x, y) = 2 \cdot 3 = 6.$$

# Which approach is better?

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- Both are actually bad from a computational point of view.
- Both require a number to be factorized.
  - a computationally difficult task.
- For fairly large numbers, both approaches require a lot of computation.
- Is there a better approach?
  - Indeed there is, given by the Greek mathematician Euclid.
  - Celebrated as a breakthrough.

# Euclid's algorithm for GCD

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- Based on the following lemma.
- Lemma : Let  $x, y$  be two positive integers. Let  $q$  and  $r$  be integers such that  $x = y.q + r$ . Then,  $\gcd(x, y) = \gcd(y, r)$ .
  - Argue that the common divisors of  $x$  and  $y$  are also common divisors of  $y$  and  $r$ .
  - Let  $d$  divide both  $x$  and  $y$ . Then,  $d$  divides  $x - yq = r$ .
  - The converse also applies in a similar fashion.
- The above lemma suggests the following algorithms for  $\gcd$ .
  - Apply the above lemma repeatedly till the remainder is 0.
  - Let  $r_1, r_2, \dots$ , be the remainders.

# Euclid's Algorithm for GCD

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- Let  $r_2, r_3, \dots$ , be the remainders with  $r_0 = x$  and  $r_1 = y$ .
- We have that:

$$r_0 = r_1q_1 + r_2,$$

$$r_1 = r_2q_2 + r_3$$

$$r_2 = r_3q_3 + r_4$$

and so on, till

$$r_{n-1} = r_n q_n + 0$$

- By the result of the above lemma, it also holds that:

$$\gcd(r_0, r_1) = \gcd(r_1, r_2)$$

$$= \gcd(r_2, r_3)$$

$$= \dots$$

$$= \gcd(r_{n-1}, r_n)$$

$$= \gcd(r_n, 0) = r_n$$

- Notice that  $r_n$  is the last nonzero remainder in the process.

# Euclid's Algorithm

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Algorithm GCD-Euclid(a,b)

$x := a, y := b;$

while (y  $\neq$  0)

$r := x \bmod y; x := y; y := r;$

end-while

End-Algorithm.

- Example,  $x = 42$  and  $y = 24$ .
- Iteration 1:  $r = 18, x = 24; y = 18$
- Iteration 2:  $r = 6, x = 18, y = 6$
- Iteration 3:  $r = 0$ .

# Euclid's Algorithm

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- Why is this efficient?
- It can be shown that given numbers  $x$  and  $y$ , the algorithm requires only about  $\log \min\{x, y\}$  iterations.
  - Compared to about  $\sqrt{x}$  for Approach I.
  - Why does approach 1 takes  $\sqrt{x}$  iterations?
- There is indeed a difference for large numbers.
- The example suggests that also efficient ways to perform operations are of interest.

# Modular Arithmetic

- Again, integer operations with several applications.
  - E.g., cryptography
- Let  $m$  be a positive integer. Then, for two positive integers  $a$  and  $b$ , we say that  $a = b \pmod{m}$  iff there exists an integer  $k$  such that  $a = b + km$ .
  - Example:  $12 = 2 \pmod{5}$ ,  $37 = 1 \pmod{4}$ .
- With the above definition, we can write
  - If  $a = b \pmod{m}$ , and  $c = d \pmod{m}$ , then  $a+c = (b+d) \pmod{m}$ , and  $ac = bd \pmod{m}$ .
  - Modular addition:  $(a+b) \pmod{m} = ((a \pmod{m}) + (b \pmod{m})) \pmod{m}$ .
    - Example:  $a = 2 \pmod{5}$ ,  $b = 4 \pmod{5}$ , then  $a+b = 1 \pmod{5}$ .
  - Similarly, modular multiplication:  $ab \pmod{m} = ((a \pmod{m}) \cdot (b \pmod{m})) \pmod{m}$ .
    - Example: With the above  $a$ ,  $b$ ,  $ab = 3 \pmod{5}$ .



# Modular Multiplication

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- An important operation in cryptography.
- Given integers  $b$ ,  $n$ , and  $m$ , compute  $b^n \pmod{m}$ .
- One can simply compute  $b^n$  and then take the modulo wrt  $m$ .
- But becomes impractical for even moderate values of  $b$  and  $n$ .
- One improvement : Recall the definition of modular multiplication.
  - Can take  $b^k \pmod{m}$  whenever  $b^k$  exceeds  $m$ .
  - Continue with  $b^k \pmod{m}$ .
  - Example:
    - $b = 5$ ,  $n = 4$ ,  $m = 6$ .  $b^2 = 25 = 1 \pmod{6}$ .
    - Now,  $b^3 = 5 \pmod{6}$ , and  $b^4 = 1 \pmod{6}$ .

# Modular Multiplication

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- A further improvement is possible with the following approach.
- Let  $n$  be written in binary as  $n_{r-1}n_{r-2}\dots n_1n_0$ .
- Then, the exponent of  $b$  can be rewritten as
  - $n_{r-1}2^{r-1} + n_{r-2}2^{r-2} + \dots + 2^1n_1 + 2^0n_0$
  - Now, if for any  $i$  between 0 and  $r-1$ , if any  $n_i = 0$ , then  $n_i2^i = 0$ , and  $b^0 = 1 \pmod{m}$ .
  - Such indices have no effect on  $b^n \pmod{m}$ .
  - Thus, we need to evaluate only indices where  $n_i = 1$ .

# Modular Multiplication

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- Example, let  $b = 4$ ,  $n = 9$ , and  $m = 11$ .
- $n = (1001)_2$
- So, we need to evaluate  $4^8 \pmod{11}$  and  $4^1 \pmod{11}$ .
- How to compute the former?
- Realise that for  $k = 2^r$ ,  $b^k \pmod{m}$  can be computed by computing  $b^2 \pmod{m}$ ,  $b^4 = b^2 \cdot b^2 \pmod{m}$ , and so on for  $r$  iterations.
- Putting together everything, the improved modular exponentiation algorithm follows.

# Algorithm for Modular Exponentiation

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Algorithm ModExp(b, n, m)

Begin

Let  $n = (n_{r-1}n_{r-2}\dots n_1n_0)_2$

result = 1;

power = a mod n;

for i=1 to r-1 do

if  $n_i = 1$  then

    result = (result . Power) mod m

    power = power . Power

End

End

# About the Algorithm

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- Not very difficult to execute the if condition. Why?
- Try a few examples offline.
- There are  $\log n$  iterations of the for loop.
- Each iteration requires at most two multiplications and two modulo operations.
- Each multiplication, and also modulus, can be computed in about  $\log m$  bit operations.

# More on Integers

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- Presently, some computations require us to work with numbers that are more than 200 digits.
  - Example: RSA cryptography.
- How to process such large numbers with a computer?
  - A problem of huge practical interest.
  - Few solutions, but scope for improvement is still there.
  - A current research area for some.

# Laboratory Session

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- Problem 1: Implement routines to add and multiply two Roman numbers.
- Problem 2: Implement Euclid's Gcd algorithm.
- Problem 3: Implement the routines to do modular exponentiation.

# Acknowledgements

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- To several online sources about the Roman number system.
- To Pranav for initiating a discussion on number systems in one meeting.