

LeetCode 238 - Product of Array Except Self (Full Explanation)

LeetCode 238 - Product of Array Except Self

Problem Understanding

Given an integer array `nums`, return an array `answer` such that `answer[i]` is equal to the product of all the elements of `nums` except `nums[i]`.

You must write an algorithm that runs in $O(n)$ time and without using the division operation.

Example:

Input: `nums = [1,2,3,4]`

Output: `[24,12,8,6]`

1. Brute Force Approach ($O(n^2)$)

Idea:

For each element `nums[i]`, multiply all other elements except `nums[i]`.

Dry Run:

`nums = [1,2,3,4]`

<code>i</code>	<code>nums[i]</code>	Product of others	Result
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0	1	$2 \times 3 \times 4 = 24$	24
---	---	----------------------------	----

1	2	$1 \times 3 \times 4 = 12$	12
---	---	----------------------------	----

2	3	$1 \times 2 \times 4 = 8$	8
---	---	---------------------------	---

3	4	$1 \times 2 \times 3 = 6$	6
---	---	---------------------------	---

Output = `[24,12,8,6]`

Python Code:

```
def productExceptSelf(nums):
```

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```
n = len(nums)
ans = []
for i in range(n):
    prod = 1
    for j in range(n):
        if i != j:
            prod *= nums[j]
    ans.append(prod)
return ans
```

C++ Code:

```
-----
vector<int> productExceptSelf(vector<int>& nums) {
    int n = nums.size();
    vector<int> ans(n);
    for(int i=0;i<n;i++){
        int prod = 1;
        for(int j=0;j<n;j++){
            if(i!=j) prod *= nums[j];
        }
        ans[i] = prod;
    }
    return ans;
}
```

Java Code:

```
-----
class Solution {
    public int[] productExceptSelf(int[] nums) {
        int n = nums.length;
        int[] ans = new int[n];
        for (int i = 0; i < n; i++) {
```

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```
int prod = 1;
for (int j = 0; j < n; j++) {
    if (i != j) prod *= nums[j];
}
ans[i] = prod;
}
return ans;
}
```

2. Optimal Approach - Prefix and Suffix Product ($O(n)$)

Idea:

We calculate $\text{prefix}[i]$ = product of elements to the left of i
and $\text{suffix}[i]$ = product of elements to the right of i .
Then $\text{answer}[i] = \text{prefix}[i] * \text{suffix}[i]$

Dry Run Example: $\text{nums} = [1, 2, 3, 4]$

Prefix pass:

$\text{prefix} = [1, 1, 2, 6]$

Suffix pass updates final result:

$\text{answer} = [24, 12, 8, 6]$

Python Code:

```
def productExceptSelf(nums):
```

```
    n = len(nums)
```

```
    ans = [1] * n
```

```
    prefix = 1
```

```
    for i in range(n):
```

```
        ans[i] = prefix
```

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```
    prefix *= nums[i]
suffix = 1
for i in range(n-1, -1, -1):
    ans[i] *= suffix
    suffix *= nums[i]
return ans
```

C++ Code:

```
-----
vector<int> productExceptSelf(vector<int>& nums) {
    int n = nums.size();
    vector<int> ans(n, 1);
    int prefix = 1;
    for (int i = 0; i < n; i++) {
        ans[i] = prefix;
        prefix *= nums[i];
    }
    int suffix = 1;
    for (int i = n - 1; i >= 0; i--) {
        ans[i] *= suffix;
        suffix *= nums[i];
    }
    return ans;
}
```

Java Code:

```
-----
class Solution {
    public int[] productExceptSelf(int[] nums) {
        int n = nums.length;
        int[] ans = new int[n];
        int prefix = 1;
```

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```
for (int i = 0; i < n; i++) {  
    ans[i] = prefix;  
    prefix *= nums[i];  
}  
  
int suffix = 1;  
for (int i = n - 1; i >= 0; i--) {  
    ans[i] *= suffix;  
    suffix *= nums[i];  
}  
  
return ans;  
}
```

Time and Space Complexity

Approach	Time	Space	Notes
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Brute Force	$O(n^2)$	$O(1)$	Simple but slow
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Prefix-Suffix	$O(n)$	$O(1)$	Efficient and elegant
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