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# 18CSC204J - Design and Analysis of Algorithms.

## Assignment - 1.

### Backward substitution.

Q1)  $T(n) = T(n-1) + 5, \quad n > 1, T(1) = 0.$

$$T(n) = T(n-1) + 5 \quad - (1)$$

$$T(n-1) = T(n-2) + 5 \quad - (2)$$

② in ①,

$$T(n) = T(n-2) + 10. \quad - (3)$$

$$T(n-2) = T(n-3) + 5 \quad - (4)$$

④ in ③,

$$T(n) = T(n-3) + 15.$$

In terms of  $k$ ,

$$T(n) = T(n-k) + k \cdot 5$$

Halting condition,  $T(1) = 0$

$$n-k = 1.$$

$$k = n-1.$$

$$T(n) = T(1) + (n-1) \cdot 5$$

$$T(n) = 5n - 5$$

$$\therefore T(n) = O(n).$$

Q2)  $T(n) = 3T(n-1), \quad n > 1, T(1) = 1.$

$$T(n) = 3T(n-1). \quad - (1)$$

$$T(n-1) = 3T(n-2) \quad - (2)$$

② in ①,

$$T(n) = 9T(n-2) \quad - (3)$$

$$T(n-2) = 3T(n-3) \quad - (4)$$

⑤ in ③,

$$x(n) = 2 + x(n-3).$$

In terms of  $k$ ,

$$x(n) = 3^{k+1} x(n-k).$$

Halting condition,  $x(1) = 4$ .

$$n - k = 1$$

$$k = n - 1.$$

$$x(n) = 3^n x(1)$$

$$x(n) = 4 \cdot 3^n.$$

$$\therefore x(n) = O(3^n).$$

Q1.  $x(n) = x(n-1) + n, \quad n > 0, \quad x(0) = 0$

$$x(n) = x(n-1) + n \quad - (1)$$

$$x(n-1) = x(n-2) + n-1 \quad - (2)$$

② in ①,

$$x(n) = x(n-2) + 2n - 1 \quad - (3)$$

$$x(n-2) = x(n-3) + n-2 \quad - (4)$$

⑤ in ③,

$$x(n) = x(n-3) + 3n - 3.$$

In terms of  $k$ ,

$$x(n) = x(n-k) + k \cdot n - \frac{k(k-1)}{2}$$

Halting condition,  $x(0) = 0$ .

$$n - k = 0.$$

$$k = n.$$

$$x(n) = x(0) + n^2 - \frac{n(n-1)}{2}$$

$$x(n) = n^2 - \frac{n^2 - n}{2}$$

$$x(n) = \frac{2n^2 - n^2 + n}{2} = \frac{n^2 + n}{2}$$

$$x(n) = O(n^2).$$

Q)  $T(n) = T(n/2) + n, n > 1, T(1) = 1.$

$T(n) = T(n/2) + n \quad - (1)$

$T(n/2) = T(n/4) + n/2 \quad - (2)$

(2) in (1),

$T(n) = T(n/4) + \frac{3n}{2} \quad - (3)$

$T(n/4) = T(n/8) + \frac{n}{4} \quad - (4)$

(3) in (4),

$T(n) = T(n/8) + \frac{7n}{4}$

In terms of  $k$ ,

$T(n) = T\left(\frac{n}{2^k}\right) + \left(\frac{2^k - 1}{2^{k-1}}\right)n.$

Halting condition,  $T(1) = 1$

$\frac{n}{2^k} = 1 \rightarrow k = \log_2 n.$

$T(n) = T(1) + \left(\frac{2^{\log_2 n} - 1}{2^{\log_2 n - 1}}\right)n = 1 + \frac{2(n-1)n}{n} = 2n - 1$

$T(n) = O(n).$

Q)  $T(n) = T(n/3) + 1, n > 1, T(1) = 1.$

$T(n) = T(n/3) + 1 \quad - (1)$

$T(n/3) = T(n/9) + 1 \quad - (2)$

(2) in (1),

$T(n) = T(n/9) + 2 \quad - (3)$

$T(n/9) = T(n/27) + 1 \quad - (4)$

(3) in (4),

$T(n) = T(n/27) + 3$

In terms of  $k$ ,

$T(n) = T\left(\frac{n}{3^k}\right) + k.$

Halting condition,  $T(1) = 1$

$\frac{n}{3^k} = 1 \rightarrow k = \log_3 n$

$T(n) = T(1) + \log_3 n$

$T(n) = \log_3 n + 1.$

$T(n) = O(\log_3 n)$