

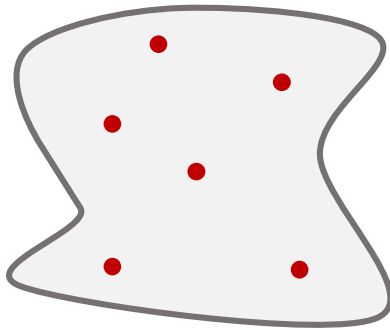
ONLINE - TSP



Online-TSP

(metric)

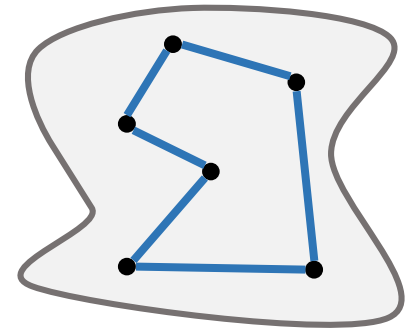
INPUT:



NP-hard!

ALG

OUTPUT:

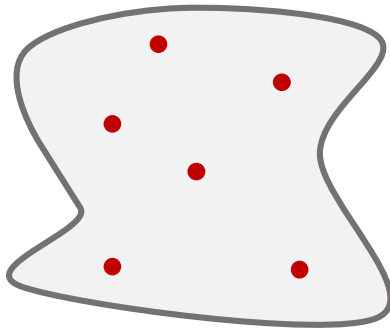


- metric space: M
(with metric d)
 - places to visit: \mathcal{S}
 - Superpolynomial Alg.
 - Approximation Alg.
e.g. *Christofides*
- Shortest tour \mathcal{T}^*
through \mathcal{S}

Online-TSP

(metric)

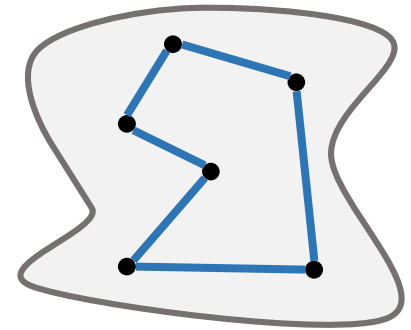
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e.g. *Christofides*

Shortest tour \mathcal{T}^*
through \mathcal{S}

Online-TSP

(metric)

Christofides Algorithm:

- (1) minimal spanning tree
- (2) minimum weighted perfect matching of odd vertices
- (3) Euler tour
- (4) Skip double visited vertices

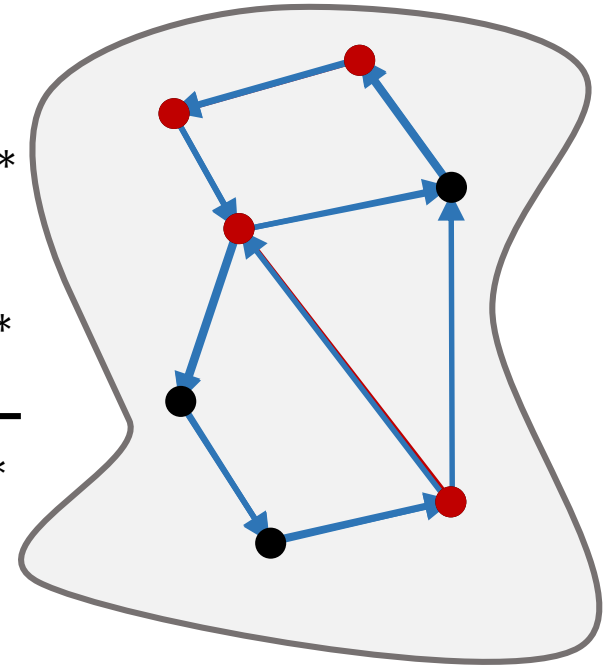
$$\leq 1 \cdot \mathcal{T}^*$$

$$\leq \frac{1}{2} \cdot \mathcal{T}^*$$

$$\leq \frac{3}{2} \cdot \mathcal{T}^*$$



1,5-approximative solution



Online-TSP

(metric)

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- (1) minimal spanning tree
- (2) minimum weighted perfect matching of odd vertices
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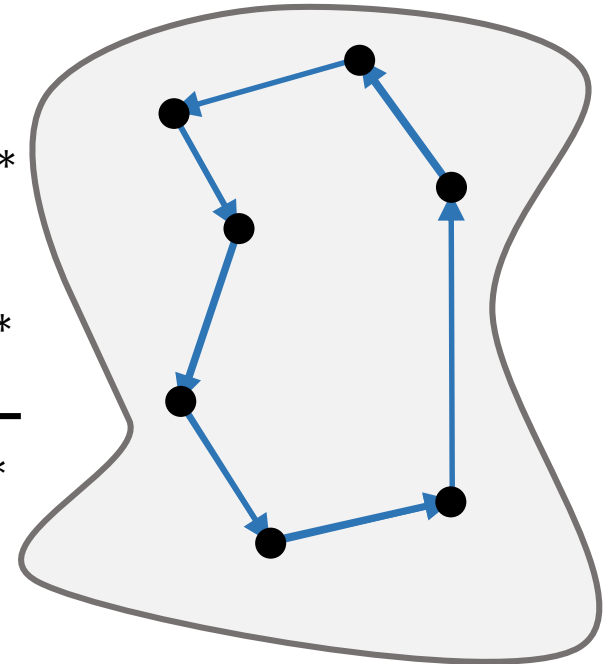
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1,5-approximative solution

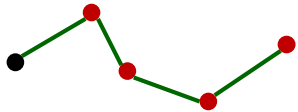


Online-TSP

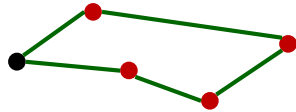
INPUT:

- metric space
- starting-point: o
- request-sequence σ :
 $, (0, x), (1, z), \dots$

N-OLTSP
“nomadic”

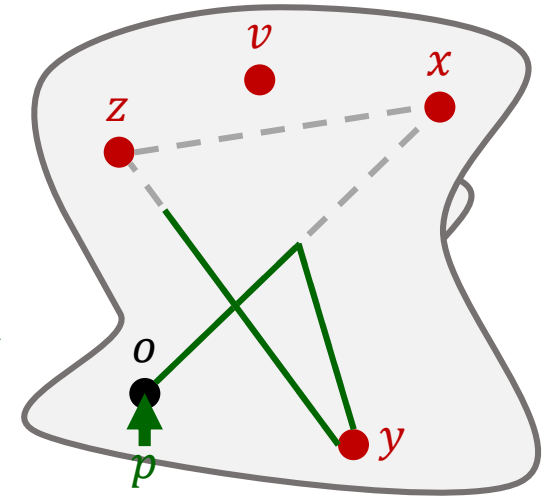


H-OLTSP
“homing”



time: 3,5

ALG



DEF: ALG is ρ -competitive

$$\Leftrightarrow |\mathcal{T}^{\text{ALG}}| \leq \rho \cdot |\mathcal{T}^{\text{OPT}}|$$

for all request- sequences

Goals

- I. Find online-algorithms (superpolynomial)
- II. Find lower bounds
- III. Find *polynomial* online-algorithms
- IV. Bonus: The real line

I. Algorithms

II. Lower Bounds

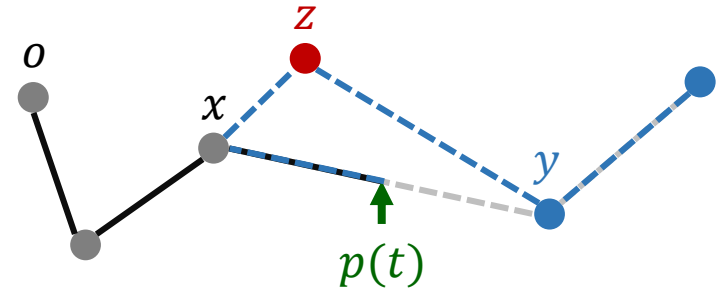
III. Polynomial Alg.

IV. Real Line

An algorithm for N-OLSTP

Invariant: always on shortest path between points in S

- (1) **New request (t, z)** at time t
and ALG between x and y
- (2) Add z to U
- (3) Follow shortest path through \mathcal{U}
beginning with x or y



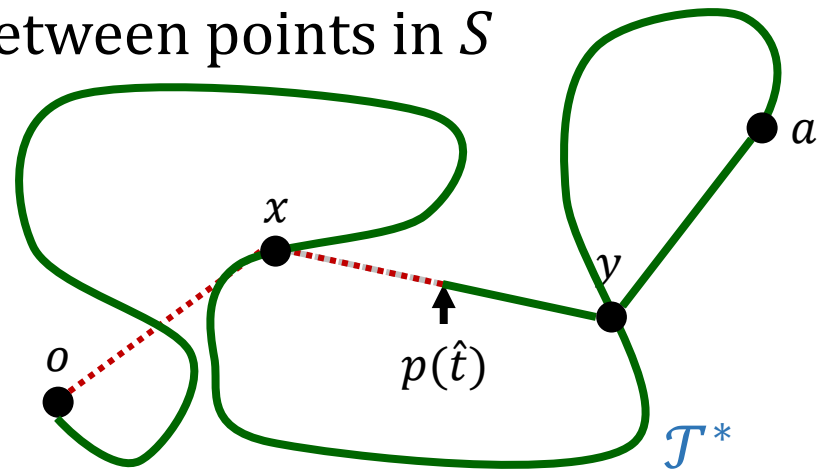
$S :=$ places requested until t
 $S \supseteq U :=$ places yet to visit at t

Greedy Travelling between Requests (GTR)

Competitiveness of GTR

Invariant: always on shortest path between points in S

- (1) Last request (\hat{t}, z) at time \hat{t} and ALG between x and y
- (3) Follow shortest path through U beginning with x or y



path found by GTR

$$|\mathcal{J}^{\text{GTR}}| \leq \hat{t} + \min\{d(o, x) + d(x, p), d(p, y) + d(y, a)\} + |\mathcal{J}^*|$$

$$\begin{aligned} &\leq \hat{t} + \frac{1}{2} \cdot |\mathcal{J}^*| + |\mathcal{J}^*| \\ \text{path found by} &\leq |\mathcal{J}^{\text{OPT}}| + \frac{3}{2} \cdot |\mathcal{J}^{\text{OPT}}| = \frac{5}{2} \cdot |\mathcal{J}^{\text{OPT}}| \\ \text{opt. offline-ALG} & \end{aligned}$$

I. Algorithms

II. Lower Bounds

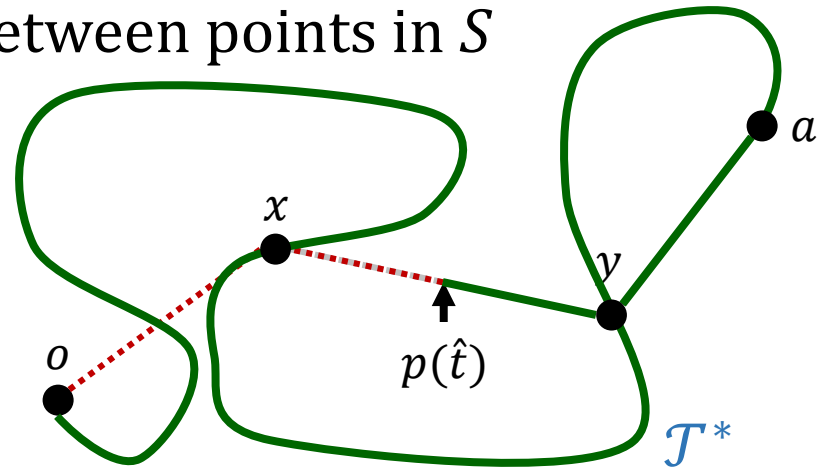
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Competitiveness of GTR

Invariant: always on shortest path between points in S

- (1) Last request (\hat{t}, z) at time \hat{t} and ALG between x and y
- (3) Follow shortest path through U beginning with x or y



THEOREM: GTR is 2,5-competitive for N-OLTSP

REMARK: tightness

REMARK: $|J^{\text{GTR}}| \leq |J^{\text{OPT}}|$ is also 2,5-competitive for H-OLTSP $|J^{\text{OPT}}|$

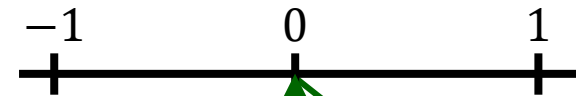
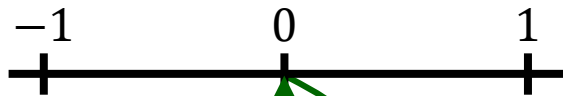
Lower Bound for N-OLTSP

time, request

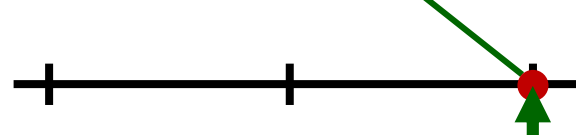
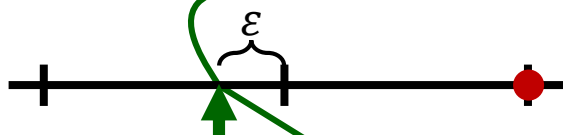
Online-ALG

Offline-ALG

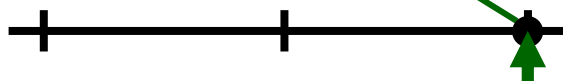
0



1, 1



$2 + \varepsilon$



$$\Rightarrow |\mathcal{T}^{\text{ON}}| = 2 + \varepsilon$$

$$\Rightarrow |\mathcal{T}^{\text{OFF}}| = 1$$

THEOREM: Any ρ -competitive ALG for N-OLTSP has $\rho \geq 2$.

I. Algorithms

II. Lower Bounds

III. Polynomial Alg.

IV. Real Line

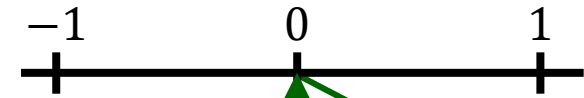
Lower Bound for H-OLTSP

time, request

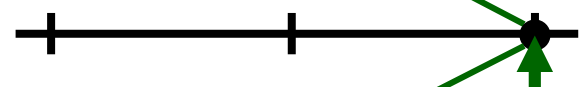
Online-ALG

Offline-ALG

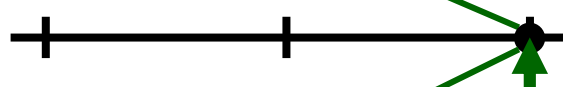
0



1, 1



$2 + \varepsilon$



$3 + \varepsilon$



$$\Rightarrow |\mathcal{T}^{\text{ON}}| = 3 + \varepsilon$$

$$\Rightarrow |\mathcal{T}^{\text{OFF}}| = 2$$

THEOREM: Any ρ -competitive ALG for H-OLTSP has $\rho \geq 1,5$.

I. Algorithms

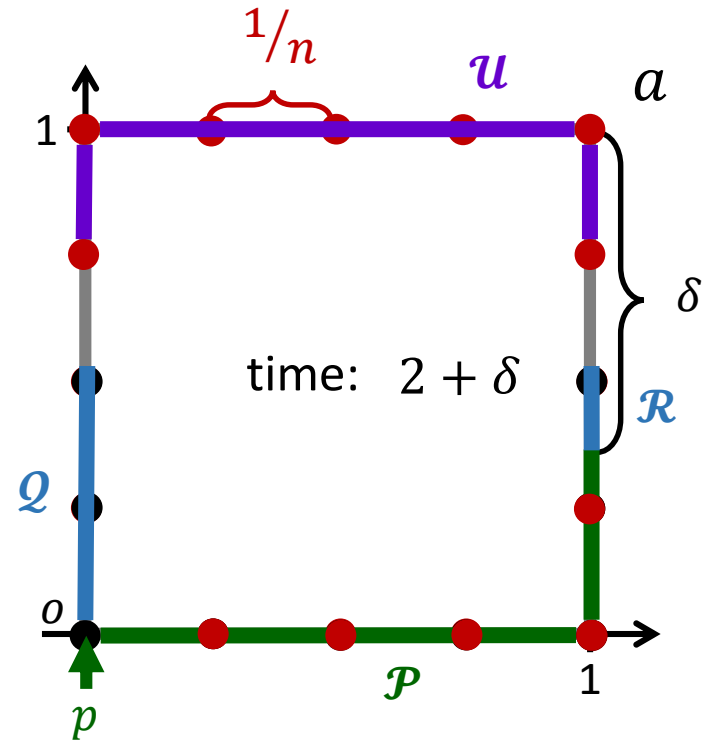
II. Lower Bounds

III. Polynomial Alg.

IV. Real Line

Lower Bound for H-OLTSP

- **requests** at time 0
- wait until time $t = 2$
- wait until $d(p, a) = t - 2$
- **new requests** on \mathcal{P}
- OPT finishes at $t = 4$
ALG finishes at $t \geq 8 - \frac{4}{n}$

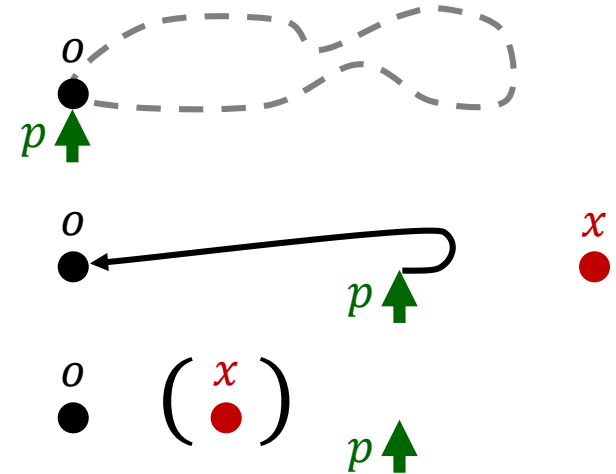


THEOREM: Any ρ -competitive ALG for H-OLTSP has $\rho \geq 2$.

A better algorithm for H-OLTSP

U := places yet to visit

- (1) At o : start optimal tour through U
- (2) For new request (t, x) :
 - a) If $d(x, o) > d(p, o)$: go back to o
 - b) Else: ignore x until back at o



Plan At Home (PAH)

GOAL: PAH is 2-competitive for H-OLTSP

I. Algorithms

II. Lower Bounds

III. Polynomial Alg.

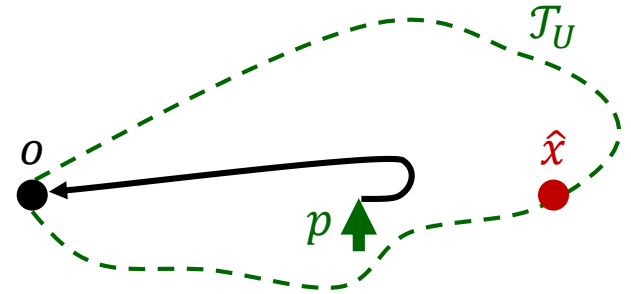
IV. Real Line

Competitiveness of PAH

$U :=$ places yet to visit, (\hat{t}, \hat{x}) last request

(2) For new request (\hat{t}, \hat{x}) :

a) If $d(\hat{x}, o) > d(p, o)$: go back to o



$$\begin{aligned} |\mathcal{J}^{\text{PAH}}| &= \underbrace{\hat{t} + d(\hat{x}, o)} + |\mathcal{J}_U| \\ &\leq |\mathcal{J}^{\text{OPT}}| + |\mathcal{J}^{\text{OPT}}| = 2 \cdot |\mathcal{J}^{\text{OPT}}| \quad \checkmark \end{aligned}$$

GOAL:

PAH is 2-competitive for H-OLTSP

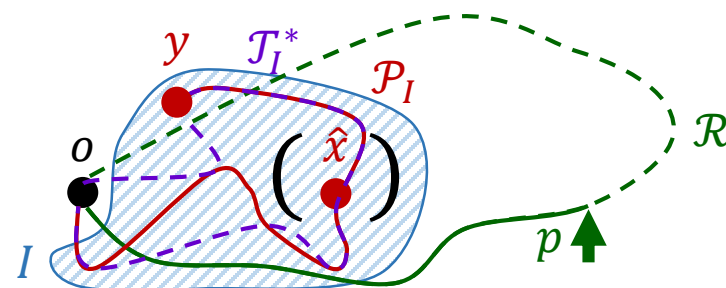
Competitiveness of PAH

U := places yet to visit, (\hat{t}, \hat{x}) last request, y 1st place in I visited by OPT

I := ignored requests

(2) For new request (t_y, y) :

b) $d(y, o) > d(p, o)$: ignore y ...



time: t_y

$$\begin{aligned}
 |\mathcal{T}^{\text{PAH}}| &\leq t_y + |\mathcal{R}| - d(o, y) + |\mathcal{J}_I^*| \\
 &\leq t_y + |\mathcal{R}| + |\mathcal{P}_I| = \underbrace{t_y + |\mathcal{P}_I|}_{|\mathcal{T}^{\text{OPT}}|} + |\mathcal{R}| \\
 &\leq |\mathcal{T}^{\text{OPT}}| + |\mathcal{T}^{\text{OPT}}| = 2 \cdot |\mathcal{T}^{\text{OPT}}|
 \end{aligned}$$

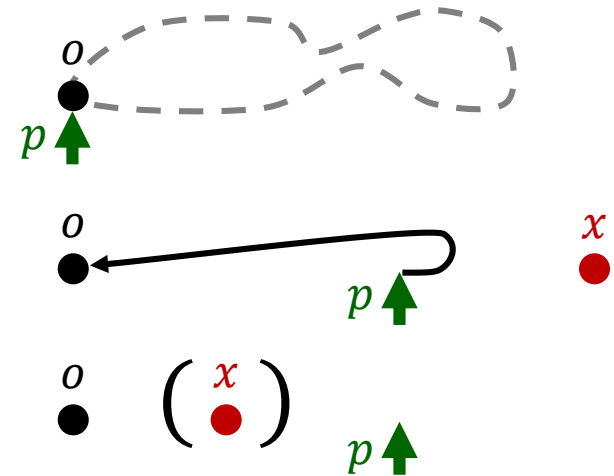
GOAL: PAH is 2-competitive for H-OLTSP



Competitiveness of PAH

U := places yet to visit

- (1) At o : start optimal tour through U
- (2) For new request (t, x) :
 - a) If $d(x, o) > d(p, o)$: go back to o
 - b) Else: ignore x until back at o



GOAL: PAH is 2-competitive for H-OLTSP r H-OLTSP

I. Algorithms

II. Lower Bounds

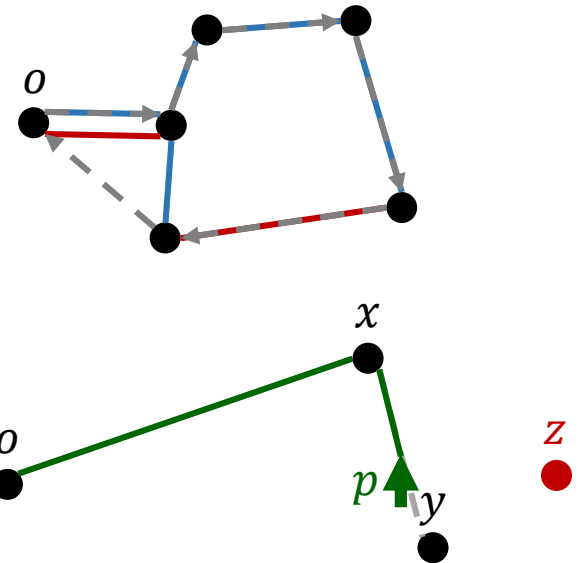
III. Polynomial Alg.

IV. Real Line

Polynomial Algorithm for H-OLTSP

Invariant: always on shortest path between points in S

- (1) At o : Find tour through $U \cup \{o\}$
with Christofides-Heuristic
- (2) For new request (t, z) at time t
and ALG between x and y :
 - a) Add z to U
 - b) go back to o via x or y
(take shortest path)

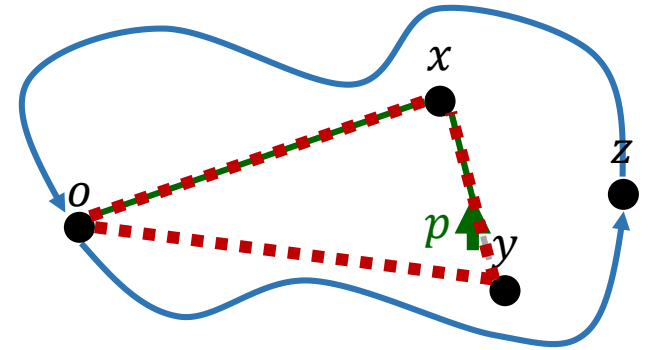


THEOREM: CHR is a polynomial (and correct).

Competitiveness of CHR

Invariant: always on shortest path between points in S

- (1) At o : Find tour through $U \cup \{o\}$
with Christofides-Heuristic



$$\begin{aligned} |\mathcal{T}^{\text{CHR}}| &= \hat{t} + \underbrace{\min\{d(o, x) + d(x, p), d(p, y) + d(y, o)\}}_{\text{CHR}(U)} + \text{CHR}(U) \\ &\leq |\mathcal{T}^{\text{OPT}}| + \frac{1}{2} \cdot |\mathcal{T}^{\text{OPT}}| + \frac{3}{2} \cdot |\mathcal{T}^{\text{OPT}}| = 3 \cdot |\mathcal{T}^{\text{OPT}}| \end{aligned}$$

REMARK: There is a 3-competitive algorithm for N-OLTSP.



I. Algorithms

II. Lower Bounds

III. Polynomial Alg.

IV. Real Line

Credits & References

- Paper: <http://citeseerx.ist.psu.edu/viewdoc/summary?doi=10.1.1.8.5620>
- Map: <http://awoiaf.westeros.org/index.php/File:WorldofIceandFire.png>
- Font: <http://www.fonts4free.net/game-of-thrones-font.html>