

# PATH PLANNING

Undergraduate course (Spring 2020)

# MOTION PLANNING

# What is Motion Planning

A robot arm is to build an assembly from a set of parts.

Tasks for the robot:

- Grasping: position gripper on object  
design a path to this position
- Transferring: determine geometry path for arm  
avoid obstacles + clearance
- Positioning

# Information Required

- Knowledge of spatial arrangement of workspace. E.g., location of obstacles
- **Full knowledge** -> full motion planning
- **Partial knowledge** -> combine planning and execution

**motion planning = collection of problems**

# Basic Problem

A simplified version of the problem assumes

- Robot is the only moving object in the workspace
- No dynamics, no temporal issues
- Only non-contact motions

Motion Planning = pure “geometrical” problem

# World consists of

- Obstacles
  - Already occupied spaces of the world
  - In other words, robots can't go there
- Free Space
  - Unoccupied space within the world
  - Robots “might” be able to go here
  - To determine where a robot can go, we need to discuss what a *Configuration Space* is

# Notion of Configuration Space

- *Main Idea:* Represent the robot as a point, called a configuration, in a parameter space, the *configuration space* (or C-space).
- *Importance:* Reduce the problem of planning the motion of a robot in Euclidean Space to planning the motion of a point in C-space.

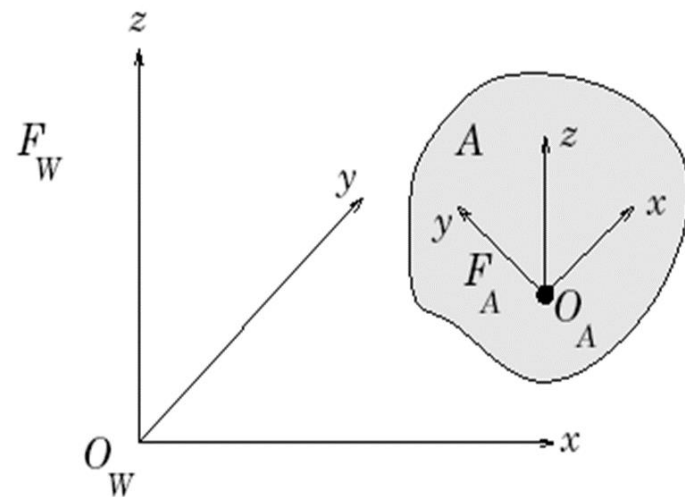
# C-space of Rigid Objects

## Workspace $\mathcal{W}$ : physical workspace

- represented as  $N$ -dimensional Euclidean Space  $\mathbb{R}^N$ , where  $N = 2, 3$
- $\mathcal{F}_\mathcal{W}$ : *fixed* Cartesian coordinate system (frame) of  $\mathcal{W}$
- $\mathcal{O}_\mathcal{W}$ : *fixed* origin of  $\mathcal{F}_\mathcal{W}$

## Robot $\mathcal{A}$ : moving rigid object/robot

- represented as compact subset of  $\mathbb{R}^N$  (at reference position and orientation)
- $\mathcal{F}_\mathcal{A}$ : frame of  $\mathcal{A}$  (aka 'local' frame of  $\mathcal{A}$ )
  - *fixed* wrt  $\mathcal{A}$  (i.e., each point in  $\mathcal{A}$  has fixed coordinates in  $\mathcal{F}_\mathcal{A}$ )
  - *moving* wrt  $\mathcal{F}_\mathcal{W}$
- $\mathcal{O}_\mathcal{A}$ : origin of  $\mathcal{F}_\mathcal{A}$  (aka the **reference point** of  $\mathcal{A}$ )





# C-space of Rigid Objects

## Definitions:

- A **configuration**  $\mathbf{q}$  of  $\mathcal{A}$  is a specification of the position and orientation of  $\mathcal{F}_{\mathcal{A}}$  wrt  $\mathcal{F}_{\mathcal{W}}$
- The **configuration space** of  $\mathcal{A}$  is the space  $\mathcal{C}$  of all the possible configurations of  $\mathcal{A}$

## Notation:

- $\mathcal{A}(\mathbf{q})$ : subset of  $\mathcal{W}$  occupied by  $\mathcal{A}$  at configuration  $\mathbf{q}$
- $a(\mathbf{q})$ : position of point  $a \in \mathcal{A}$  in  $\mathcal{W}$  when  $\mathcal{A}$  at configuration  $\mathbf{q}$

# C-space of Rigid Objects

## Robot Configurations Can be:

- 1. Free configurations: robot and obstacles do not overlap
- 2. Contact configurations: robot and obstacles touch
- 3. Blocked configurations: robot and obstacles overlap
- Configuration Space partitioned into free ( $C_{\text{free}}$ ), contact, and blocked sets.

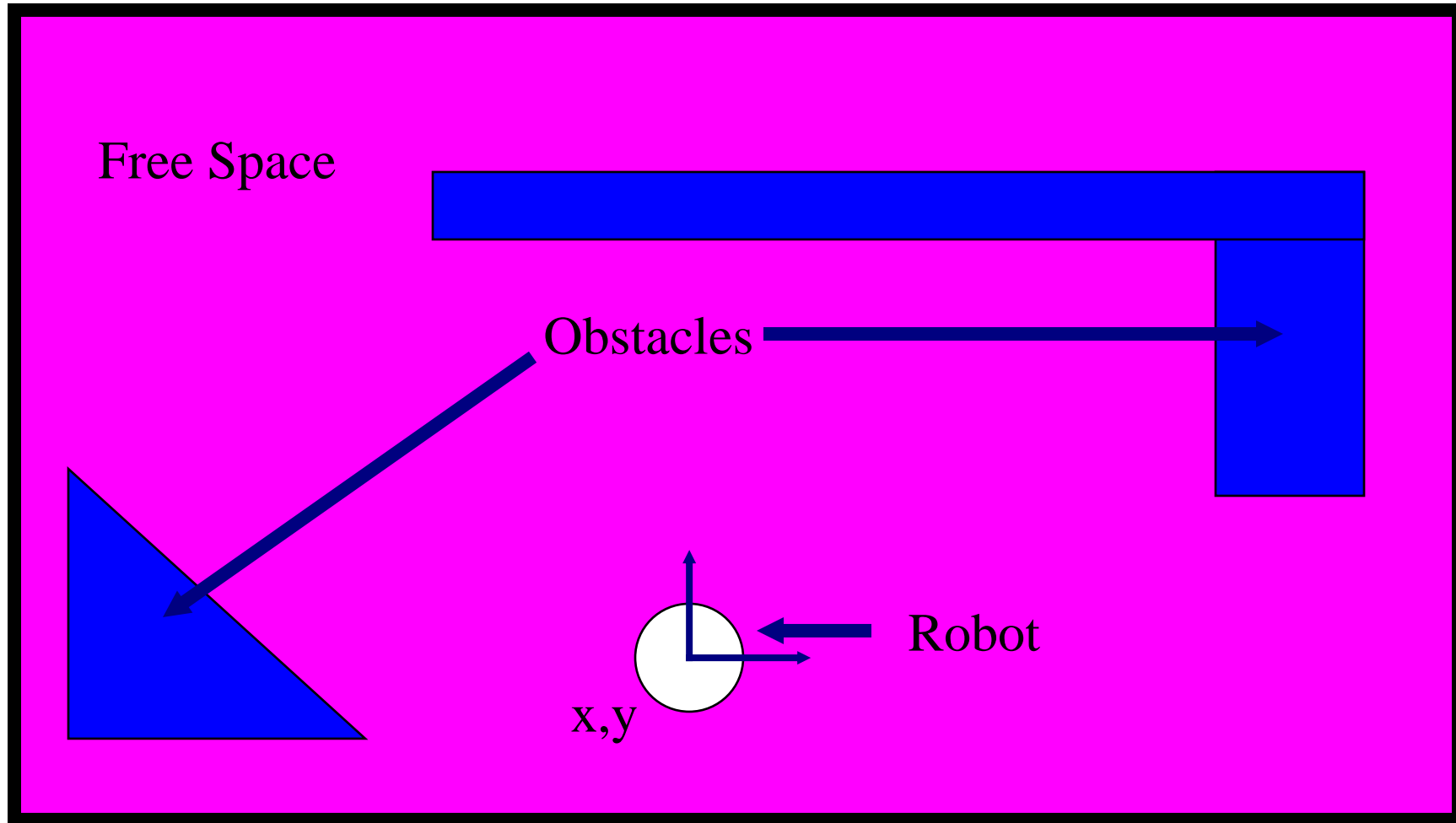
**Definition:** The obstacle  $\mathcal{B}_i$  in  $\mathcal{W}$  maps in  $\mathcal{C}$  to the region

$$\mathcal{CB}_i = \{\mathbf{q} \in \mathcal{C} \mid \mathcal{A}(\mathbf{q}) \cap \mathcal{B}_i \neq \emptyset\}$$

$\mathcal{CB}_i$  is called a *C-obstacle*. The union of all C-obstacles is called the *C-obstacle region*.

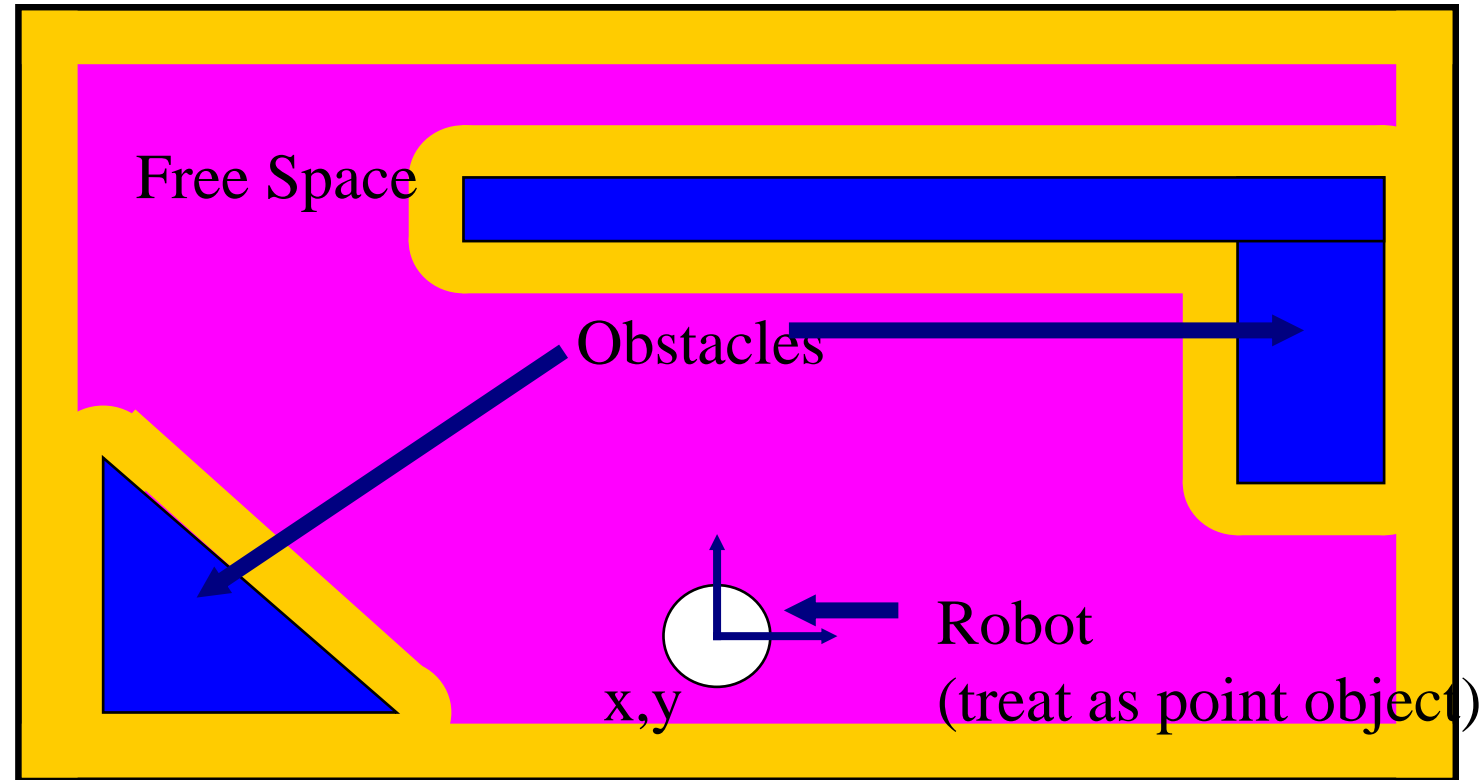
$$C_{\text{space}} = C_{\text{free}} + C_{\text{obstacle}}$$

# Example of a World (and Robot)



# The Configuration Space

- How to create it
  - First abstract the robot as a point object.
  - Then, enlarge the obstacles to account for the robot's footprint and degrees of freedom
  - In example, the robot was circular, so it enlarge obstacles by the robot's radius.



# Components

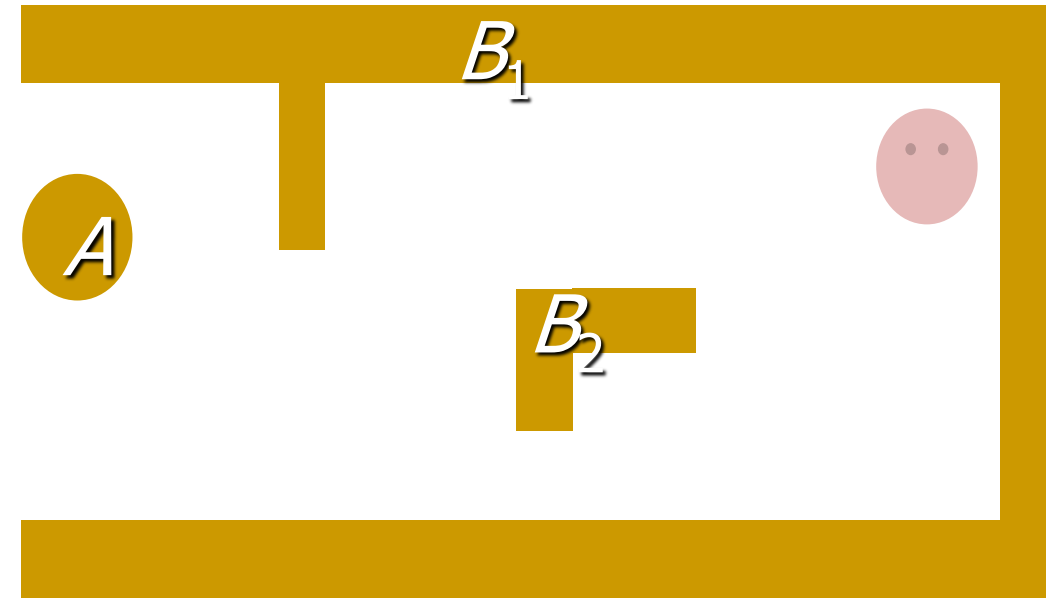
- $A$ : single rigid object - the robot - moving in Euclidean space  $W$  (the workspace).

$$W = R^N, N=2,3$$

- $B_i, i=1,\dots,q$ . Rigid objects in  $W$ . The obstacles

## Assume

- Geometry of  $A$  and  $B_i$  is perfectly known
- Location of  $B_i$  is known
- No kinematic constraints on  $A$ : a “free flying” object



# Path in C-Space

- Mathematically definition

A *path in C-space* is a continuous map:

$$\tau : s \in [0, 1] \mapsto \tau(s) \in \mathcal{C}$$

where  $\tau(0) = \mathbf{q}_{init}$  is the initial configuration and  $\tau(1) = \mathbf{q}_{goal}$  is the goal configuration of the path.

“Continuous map” means that:

$$\forall s_1, s_2 \in [0, 1] : \lim_{s_2 \rightarrow s_1} d(\tau(s_1), \tau(s_2)) = 0$$

where  $d : \mathcal{C} \times \mathcal{C} \rightarrow \mathbf{R}^+ \cup \{0\}$  is the chosen metric over  $\mathcal{C}$ .

# Motion Planning

- **General Goal**: compute motion commands to achieve a goal arrangement of physical objects from an initial arrangement
- **Basic problem**: Collision-free path planning for one rigid or articulated object (the “robot”) among static obstacles.

## **Inputs**

- geometric descriptions of the obstacles and the robot
- kinematic and dynamic properties of the robot
- initial and goal positions (configurations) of the robot

## **Output**

- Continuous sequence of collision-free configurations connecting the initial and goal configurations

# Motion Planning

- **Path planning**

design of only geometric (**kinematic**) *specifications* of the **positions** and **orientations** of robots

- **Trajectory = Path** + assignment of **time** to points along the path

- **Trajectory planning**

**path planning** + design of **linear** and **angular velocities**

- **Motion Planning (MP)**, a general term, either:

- **Path** planning, or
- **Trajectory** planning

- **Path planning < Trajectory planning**



# Classification of MP algorithms

## Completeness

- **Exact**
  - usually computationally expensive
- **Heuristic**
  - aimed at generating a solution in a short time
  - may fail to find solution or find poor one at difficult problems
  - important in engineering applications
- **Probabilistically complete** (probabilistic completeness  $\rightarrow 1$ )

# Classification of MP algorithms

## Scope

- **Global**
  - take into account all environment information
  - plan a motion from start to goal configuration
- **Local**
  - avoid obstacles in the vicinity of the robot
  - use information about nearby obstacles only
  - used when start and goal are close together
  - used as component in global planner, or
  - used as safety feature to avoid unexpected obstacles not present in environment model, but sensed during motion

# Point-to-point Path Planning

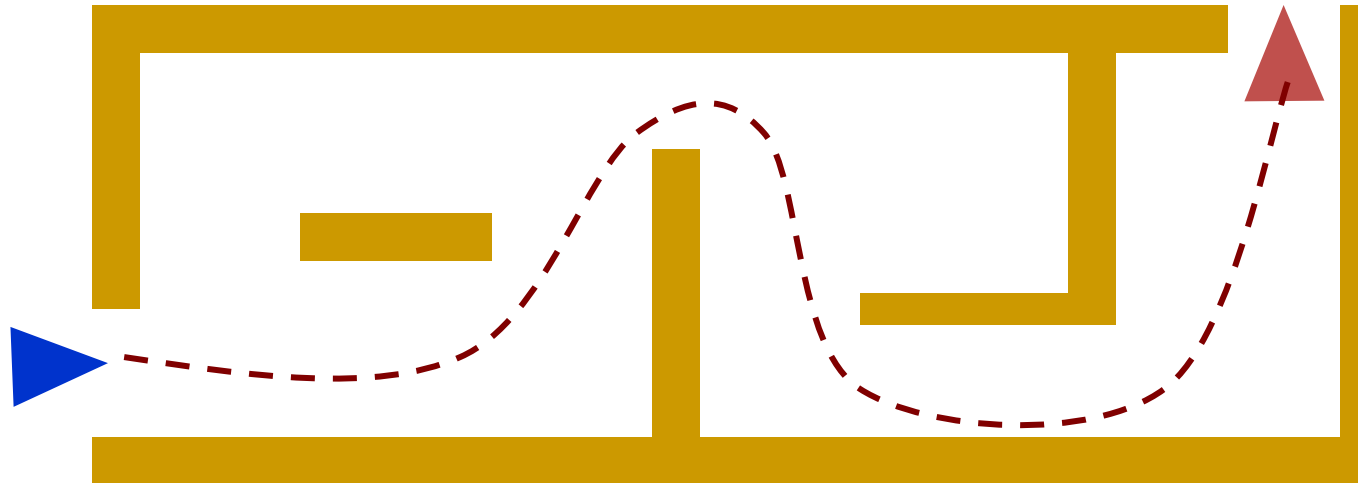
- Point-to-point path planning:
  - Looks for the best route to move an entity from point A to point B
  - Avoiding known obstacles in its path
  - Not leaving the map boundaries, and not violating the entity's mobility constraints.
- This type of path planning is used:
  - Finding routes for autonomous robots
  - Planning the manipulator's movement of a stationary robot
  - For moving entities to different locations in a map to accomplish certain goals in a gaming application.

# Region Filling Path Planning

- Tasks such as vacuuming a room, plowing a field, or mowing a lawn require region filling path planning operations that are defined as follows:
  - The mobile robot must move through an entire area, i.e., the overall travel must cover a whole region.
  - Continuous and sequential operation without any repetition of paths is required of the robot.
  - The robot must avoid all obstacles in a region.
  - An "optimal" path is desired under the available conditions.

# Components

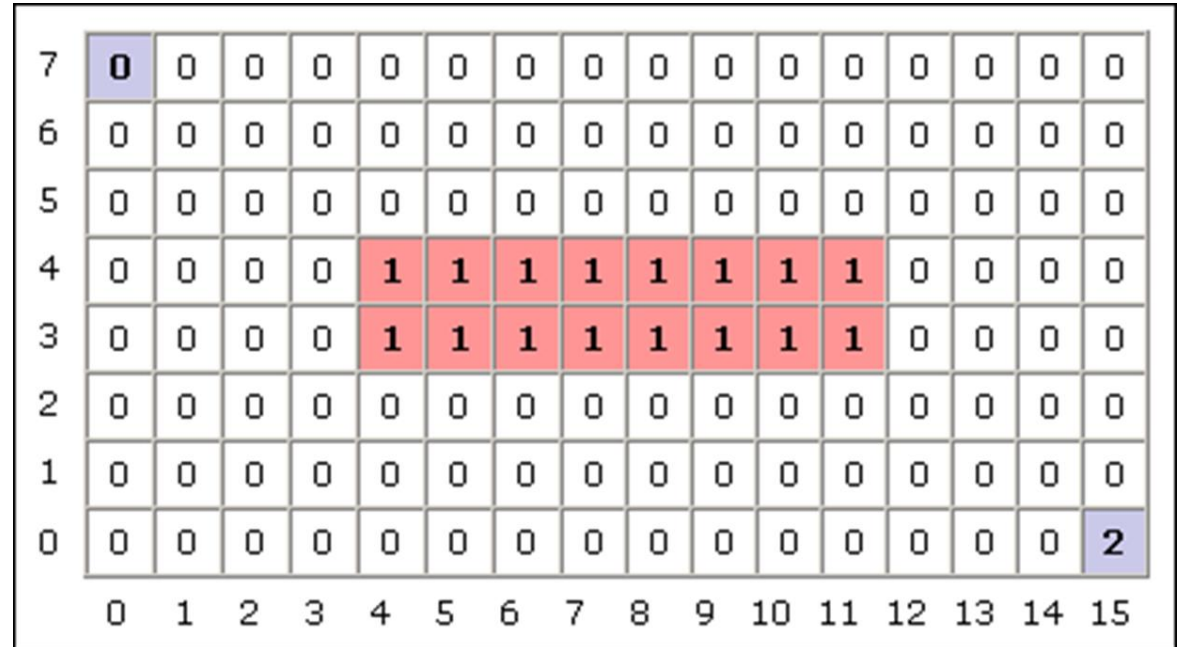
- The Problem:
  - Given an initial position and orientation  $PO_{init}$
  - Given a goal position and orientation  $PO_{goal}$
  - Generate: continuous path  $t$  from  $PO_{init}$  to  $PO_{goal}$
- $t$  is a continuous sequence of Pos'



# THE WAVEFRONT PLANNER

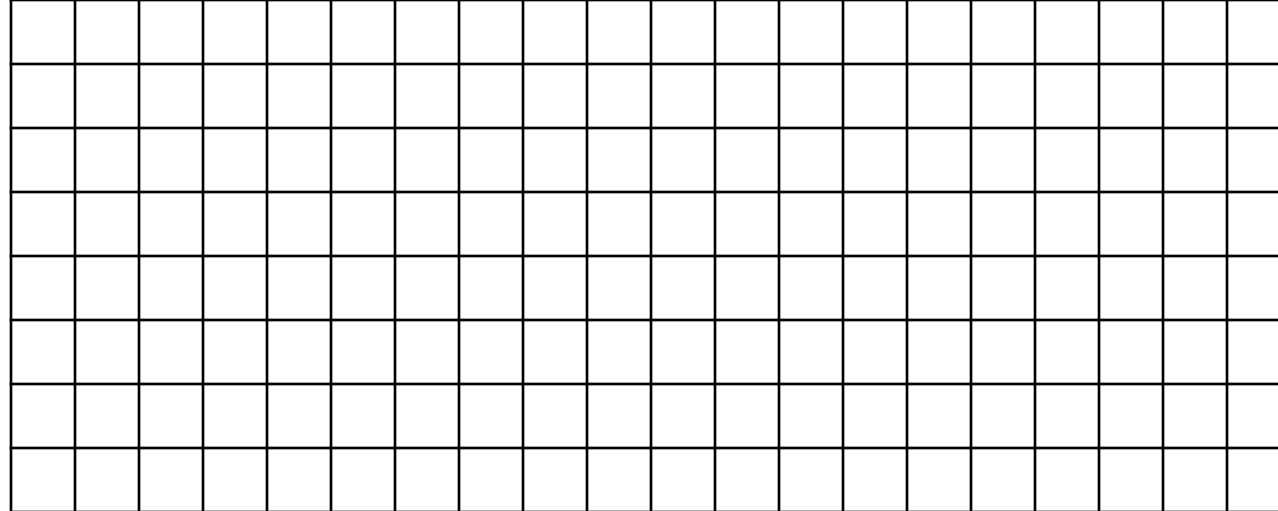
# The Wavefront Planner

- A common algorithm used to determine the shortest paths between two points
  - In essence, a breadth first search of a graph
- For simplification, we'll present the world as a two-dimensional grid
- Setup:
  - Label free space with 0
  - Label C-Obstacle as 1
  - Label the destination as 2



# Representations: A Grid

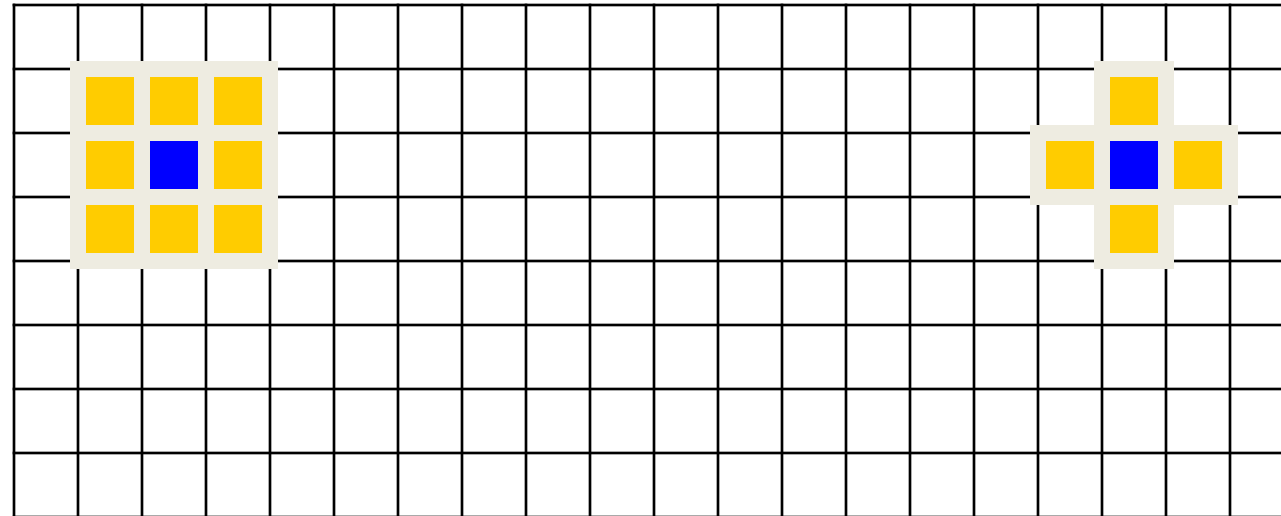
- Distance is reduced to discrete steps
  - For simplicity, we'll assume distance is uniform
- Direction is now limited from one adjacent cell to another





# Representations: Connectivity

- 8-Point Connectivity
- 4-Point Connectivity



# The Wavefront in Action (Part 1)

- Starting with the goal, set all adjacent cells with “0” to the current cell + 1
  - 4-Point Connectivity or 8-Point Connectivity?
  - Your Choice. We’ll use 8-Point Connectivity in our example

7	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
6	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
5	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
4	0	0	0	0	1	1	1	1	1	1	1	1	0	0	0	
3	0	0	0	0	1	1	1	1	1	1	1	1	0	0	0	
2	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	
1	0	0	0	0	0	0	0	0	0	0	0	0	0	3	3	
0	0	0	0	0	0	0	0	0	0	0	0	0	0	3	2	
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

## The Wavefront in Action (Part 2)

- Now repeat with the modified cells
  - This will be repeated until no 0's are adjacent to cells with values  $\geq 2$ 
    - 0's will only remain when regions are unreachable

7	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
6	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
5	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
4	0	0	0	0	1	1	1	1	1	1	1	1	0	0	0	0
3	0	0	0	0	1	1	1	1	1	1	1	1	0	0	0	0
2	0	0	0	0	0	0	0	0	0	0	0	0	0	4	4	4
1	0	0	0	0	0	0	0	0	0	0	0	0	0	4	3	3
0	0	0	0	0	0	0	0	0	0	0	0	0	0	4	3	2
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

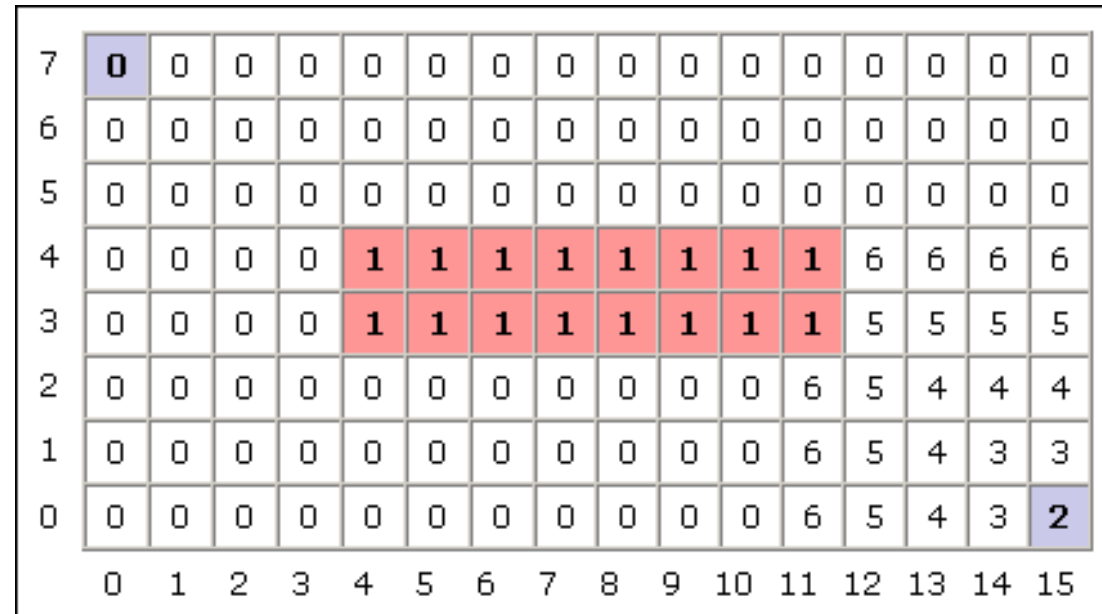
## The Wavefront in Action (Part 3)

- Repeat again...

7	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
6	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
5	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
4	0	0	0	0	1	1	1	1	1	1	1	1	0	0	0
3	0	0	0	0	1	1	1	1	1	1	1	1	5	5	5
2	0	0	0	0	0	0	0	0	0	0	0	0	5	4	4
1	0	0	0	0	0	0	0	0	0	0	0	0	5	4	3
0	0	0	0	0	0	0	0	0	0	0	0	0	5	4	3
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14

# The Wavefront in Action (Part 4)

- And again...



# The Wavefront in Action (Part 5)

- And again until...

7	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
6	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
5	0	0	0	0	0	0	0	0	0	0	0	7	7	7	7	7
4	0	0	0	0	1	1	1	1	1	1	1	1	6	6	6	6
3	0	0	0	0	1	1	1	1	1	1	1	1	5	5	5	5
2	0	0	0	0	0	0	0	0	0	0	7	6	5	4	4	4
1	0	0	0	0	0	0	0	0	0	0	7	6	5	4	3	3
0	0	0	0	0	0	0	0	0	0	0	7	6	5	4	3	2
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

# The Wavefront in Action (Done)

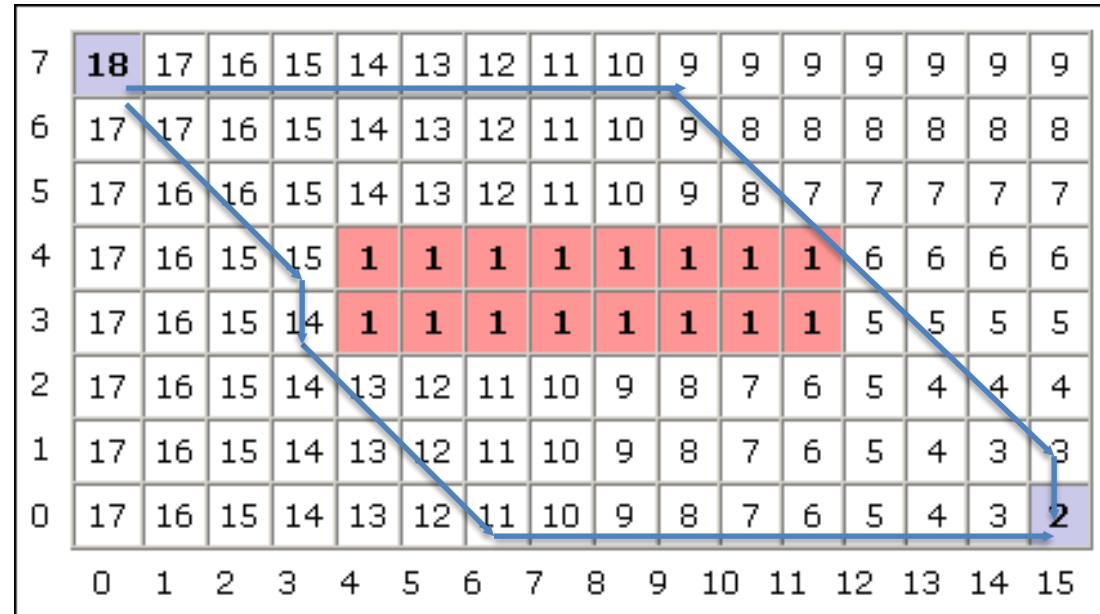
- You're done
  - Remember, 0's should only remain if unreachable regions exist

7	18	17	16	15	14	13	12	11	10	9	9	9	9	9	9
6	17	17	16	15	14	13	12	11	10	9	8	8	8	8	8
5	17	16	16	15	14	13	12	11	10	9	8	7	7	7	7
4	17	16	15	15	1	1	1	1	1	1	1	1	6	6	6
3	17	16	15	14	1	1	1	1	1	1	1	1	5	5	5
2	17	16	15	14	13	12	11	10	9	8	7	6	5	4	4
1	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3
0	17	16	15	14	13	12	11	10	9	8	7	6	5	4	2
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14

# The Wavefront, Now What?

- To find the shortest path, according to your metric, simply always move toward a cell with a lower number
  - The numbers generated by the Wavefront planner are roughly proportional to their distance from the goal

Two  
possible  
shortest  
paths  
shown

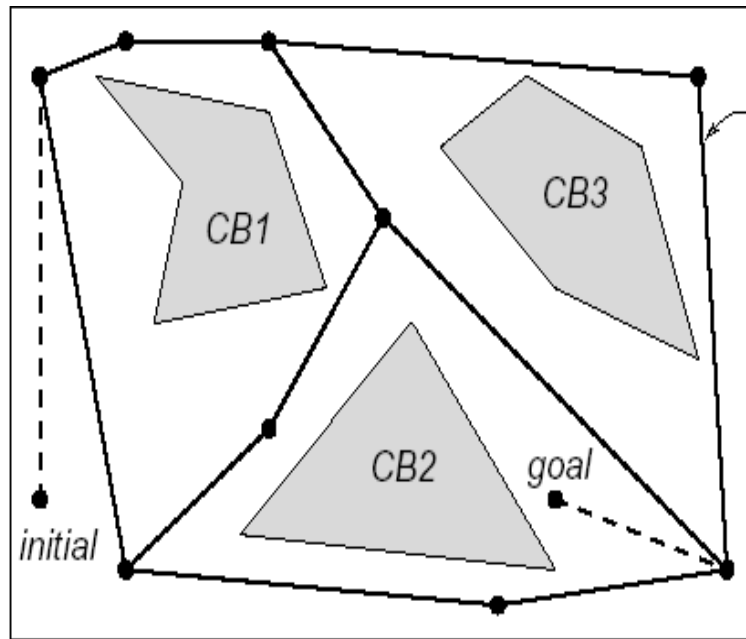




# MAP-BASED APPROACHES: VISIBILITY GRAPH

# Map-Based Approaches: Roadmap Theory

- Idea: capture the connectivity of C-free with a roadmap (graph or network) of one-dimensional curves



*Free configurations: robot and obstacles do not overlap*



# Roadmap Methods

## PATH PLANNING WITH A ROADMAP

**input:** configurations  $\mathbf{q}_{init}$  and  $\mathbf{q}_{goal}$ , and  $\mathcal{B}$

**output:** a path in  $\mathcal{C}_{free}$  connecting  $\mathbf{q}_{init}$  and  $\mathbf{q}_{goal}$

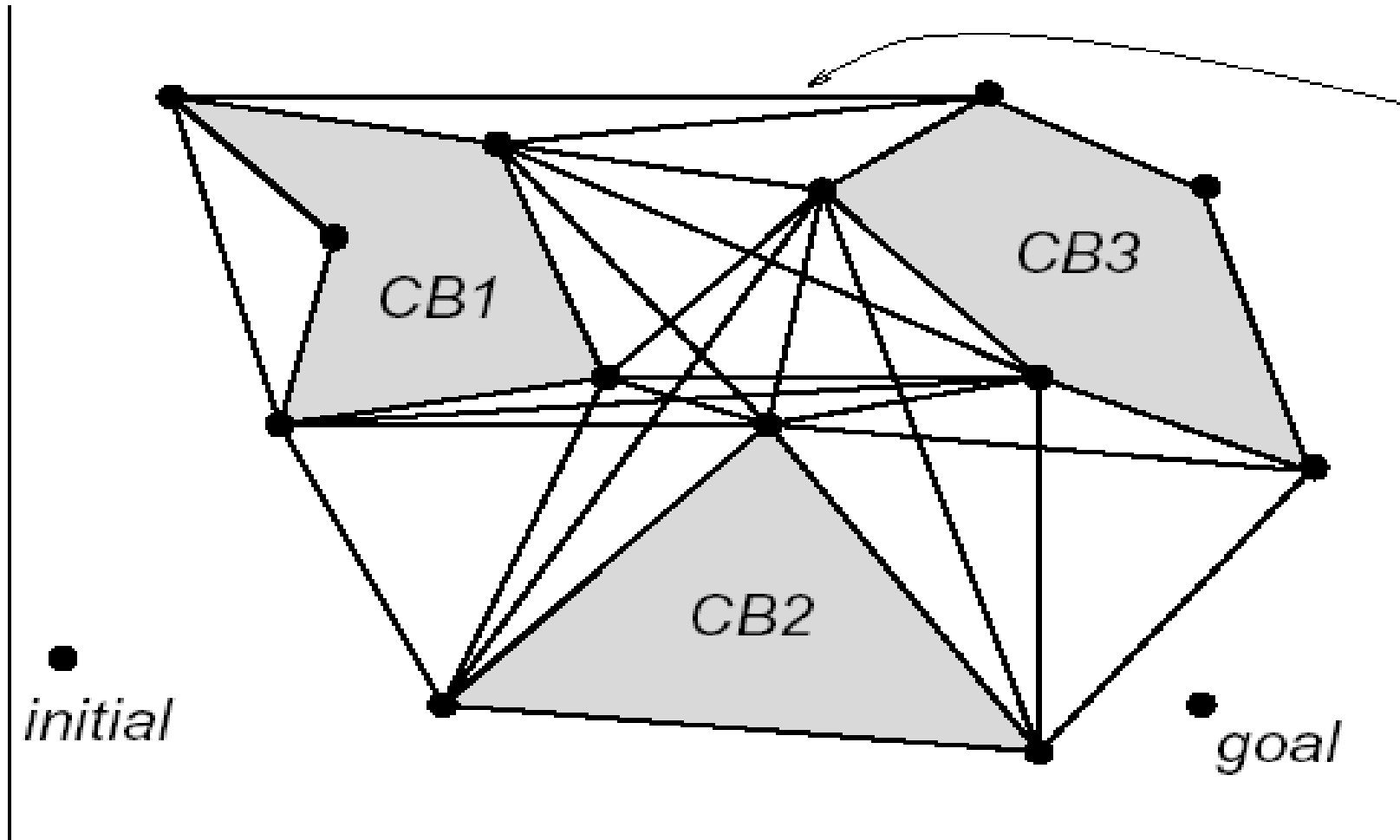
1. build a roadmap in  $\mathcal{C}_{free}$  (preprocessing)
  - roadmap nodes are free configurations (or semi-free)
  - two nodes connected by edge if can (easily) move between them
2. connect  $\mathbf{q}_{init}$  and  $\mathbf{q}_{goal}$  to roadmap nodes  $v_{init}$  and  $v_{goal}$  (in same connected component)
3. find a path in the roadmap between  $v_{init}$  and  $v_{goal}$ 
  - directly gives a path in  $\mathcal{C}_{free}$

# Roadmap Methods

- Properties of a roadmap:
  - Accessibility: there exists a collision-free path from the start to the road map
  - Departability: there exists a collision-free path from the roadmap to the goal.
  - Connectivity: there exists a collision-free path from the start to the goal (on the roadmap).
- Examples of Roadmaps
  - Visibility Graph
  - Generalized Voronoi Graph (GVG)



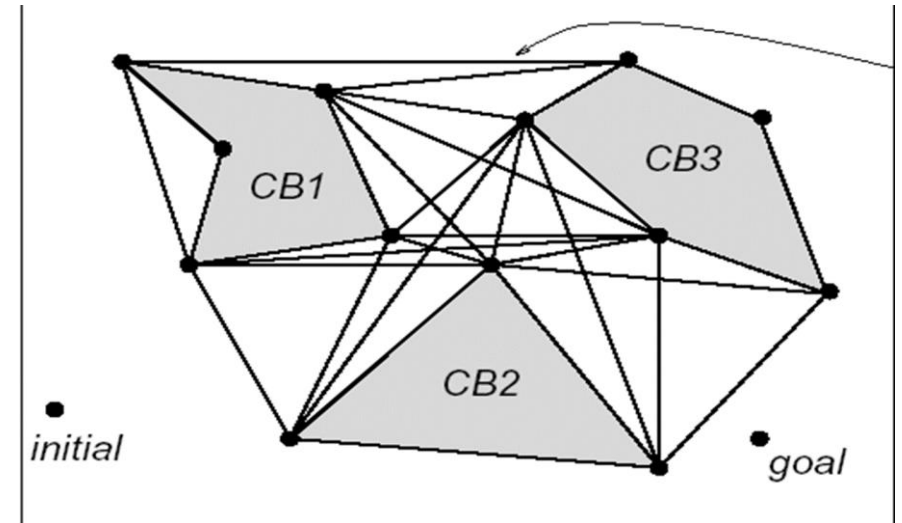
# Roadmap: Visibility Graph



# Visibility Graph of C-Space

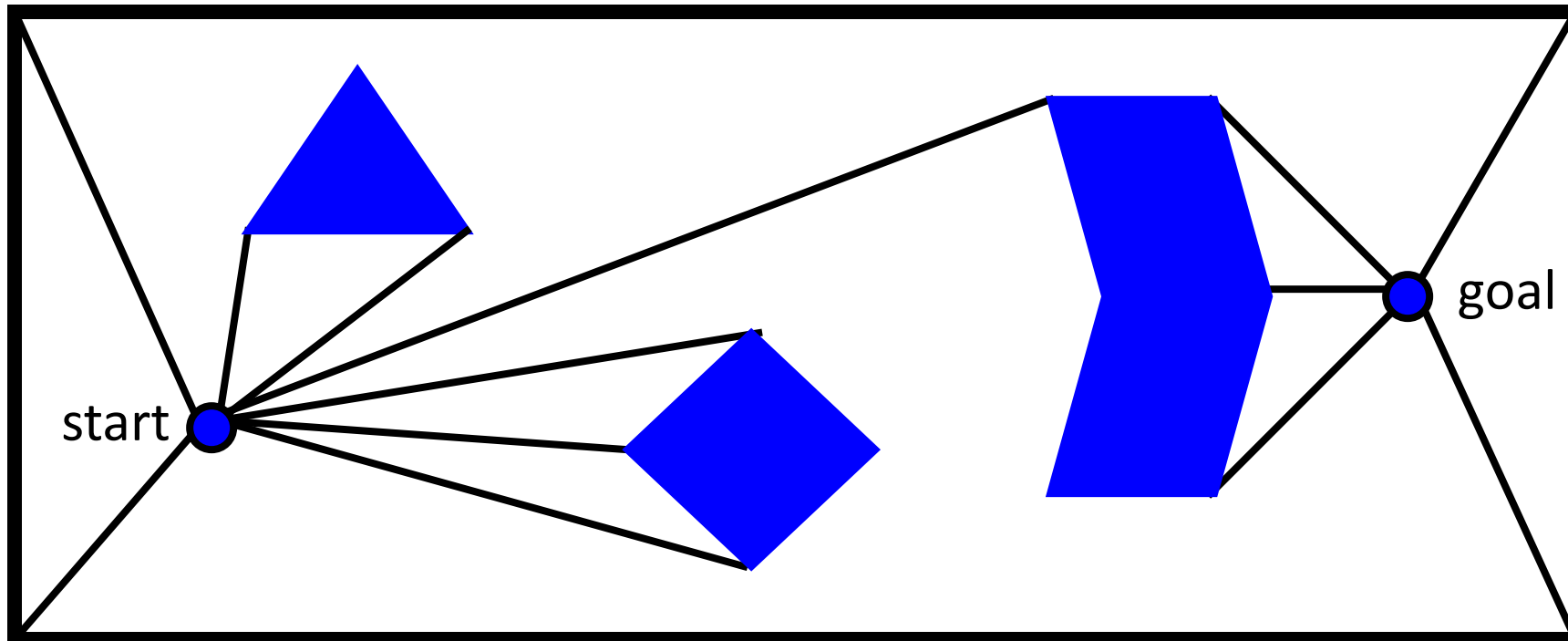
A **visibility graph** of C-space for a given  $\mathcal{CB}$  is an undirected graph  $G$  where

- nodes in  $G$  correspond to vertices of  $\mathcal{CB}$
- nodes connected by edge in  $G$  if
  - they are connected by an edge in  $\mathcal{CB}$ , or
  - the straight line segment connecting them lies entirely in  $\mathcal{C}_{free}$
- (could add  $\mathbf{q}_{init}$  and  $\mathbf{q}_{goal}$  as roadmap nodes)



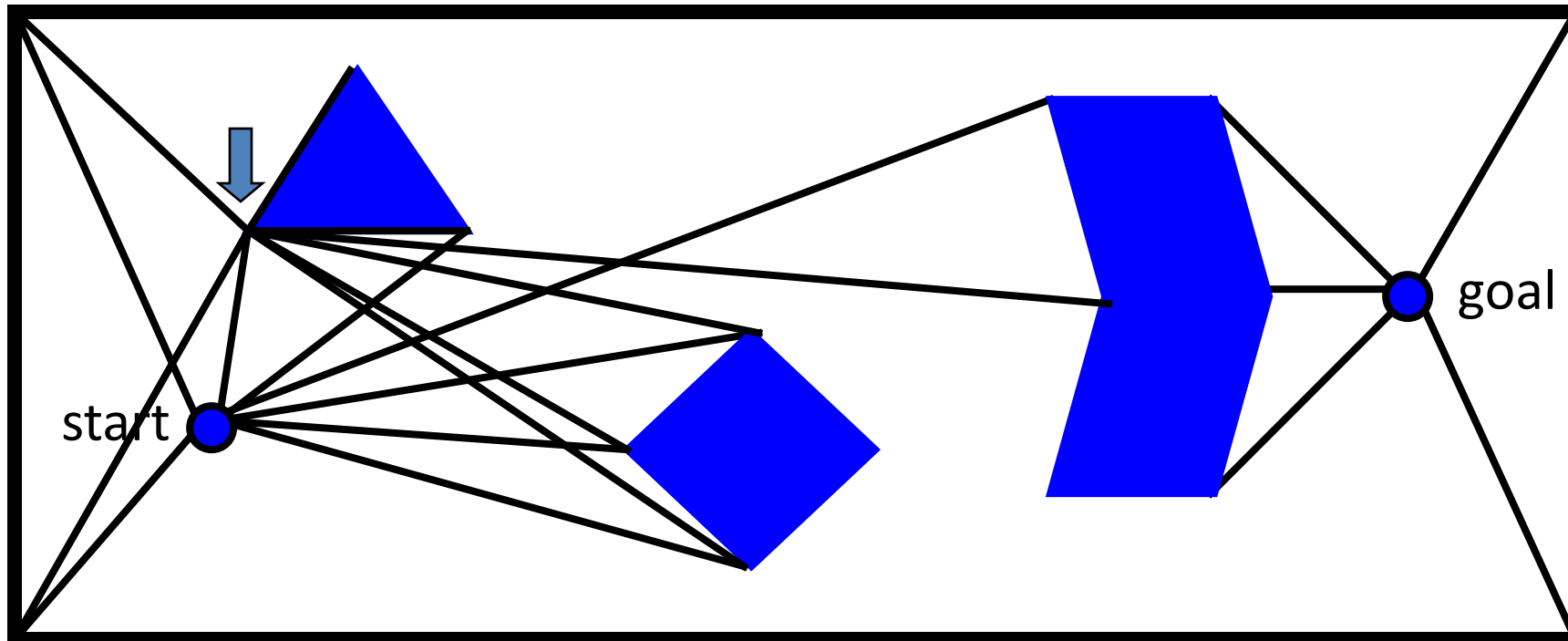
## The Visibility Graph in Action (Part 1)

- First, draw lines of sight from the start and goal to all “visible” vertices and corners of the world.



## The Visibility Graph in Action (Part 2)

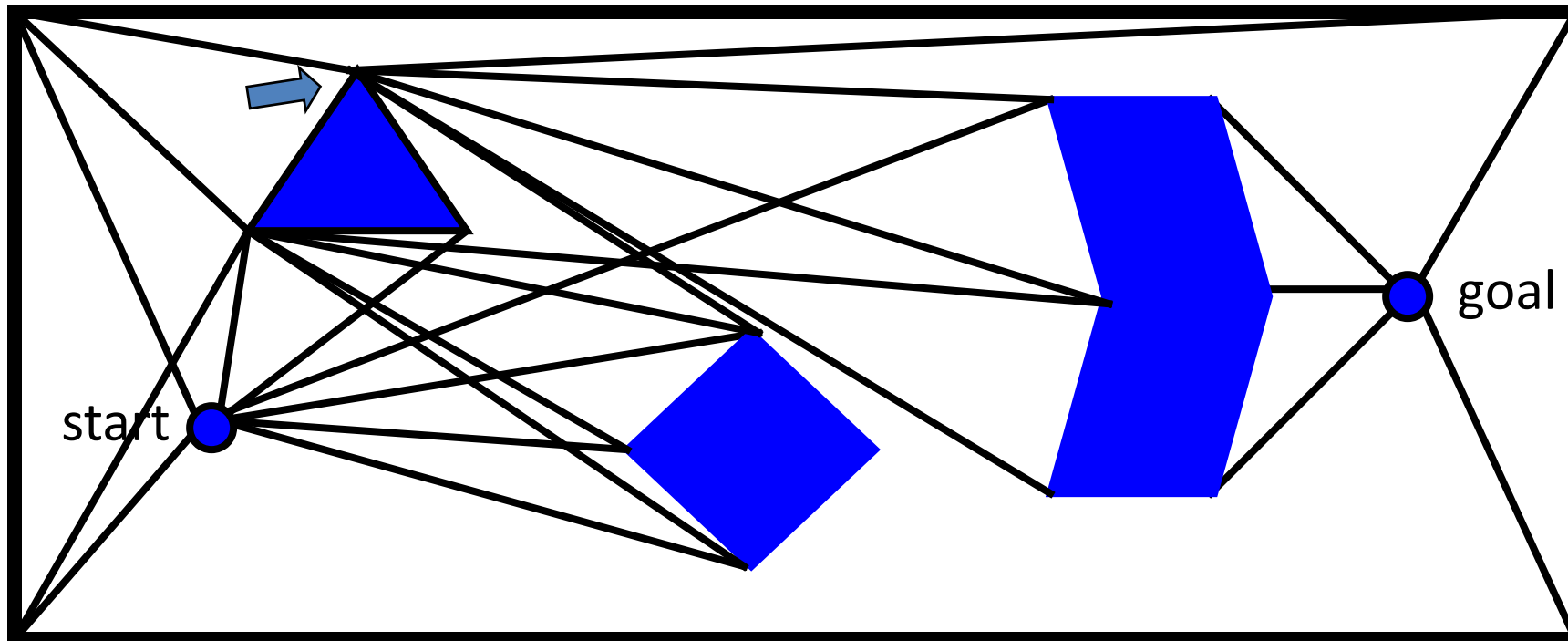
- Second, draw lines of sight from every vertex of every obstacle like before. Remember lines along edges are also lines of sight.





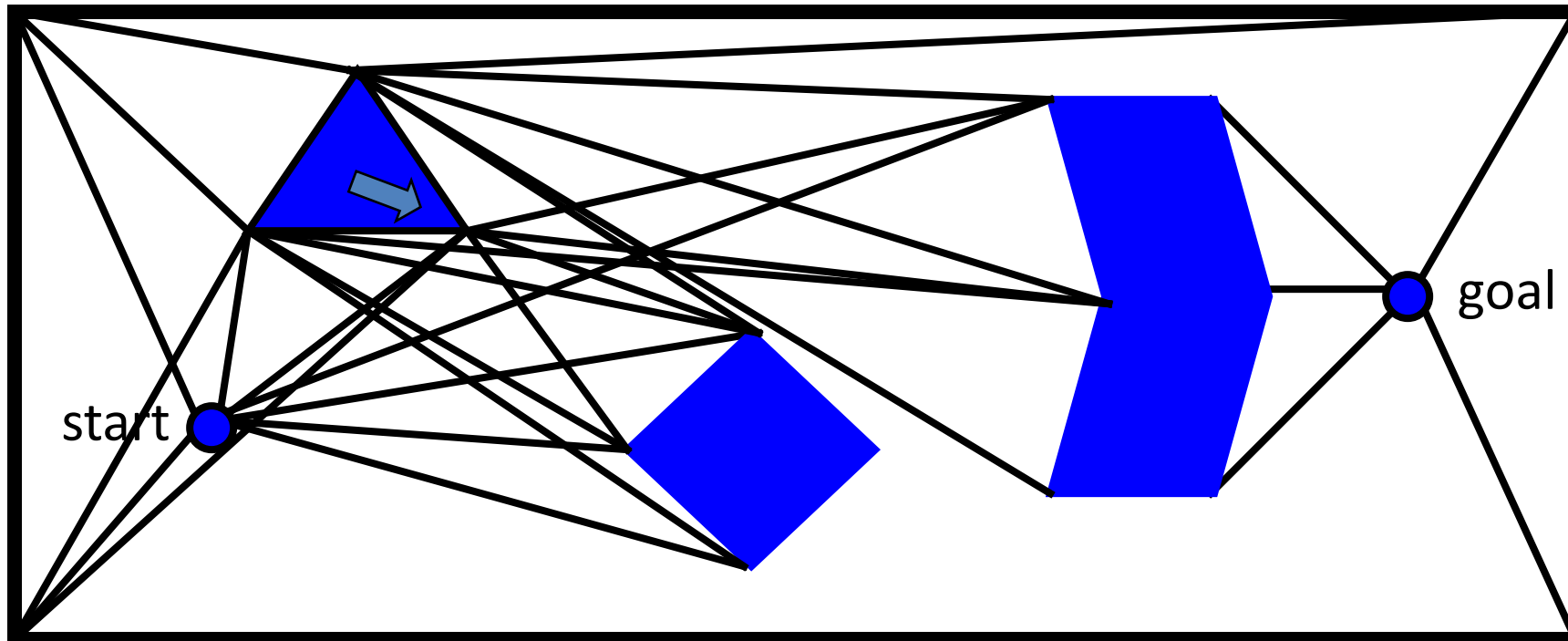
## The Visibility Graph in Action (Part 3)

- Second, draw lines of sight from every vertex of every obstacle like before. Remember lines along edges are also lines of sight.



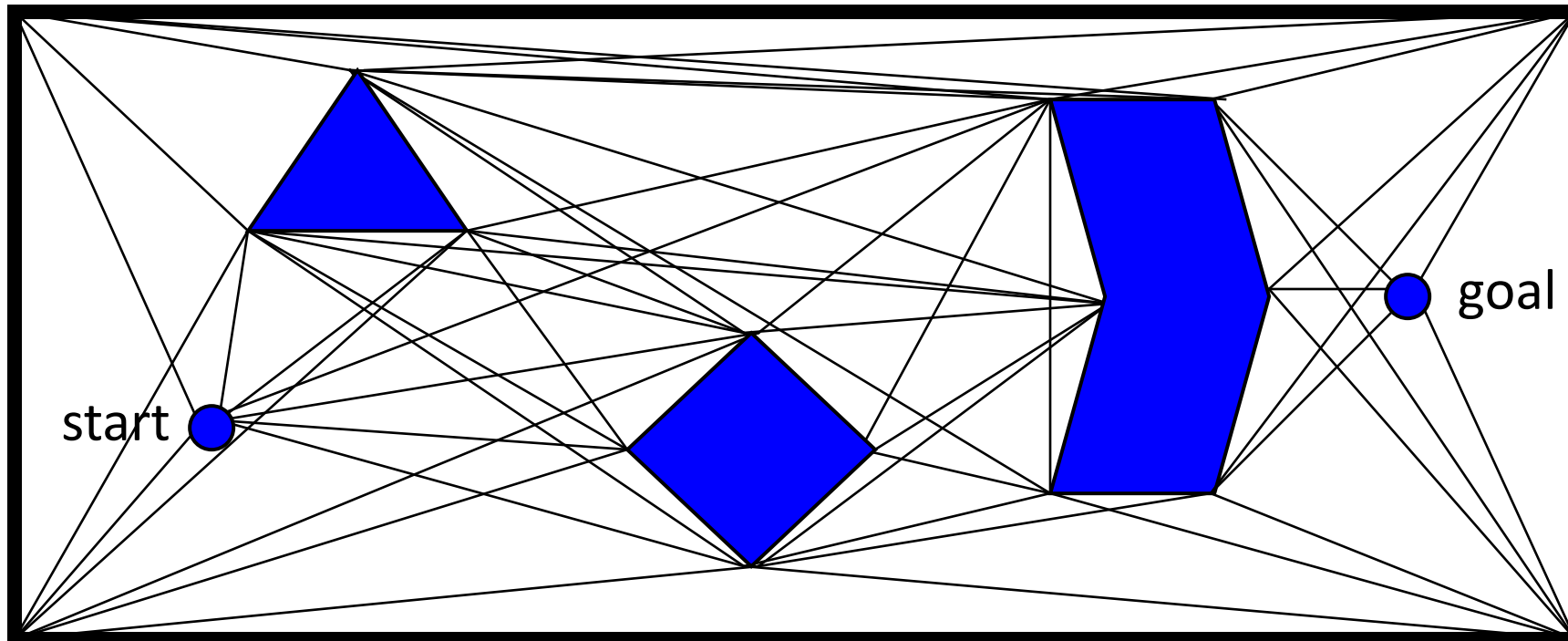
## The Visibility Graph in Action (Part 4)

- Second, draw lines of sight from every vertex of every obstacle like before. Remember lines along edges are also lines of sight.



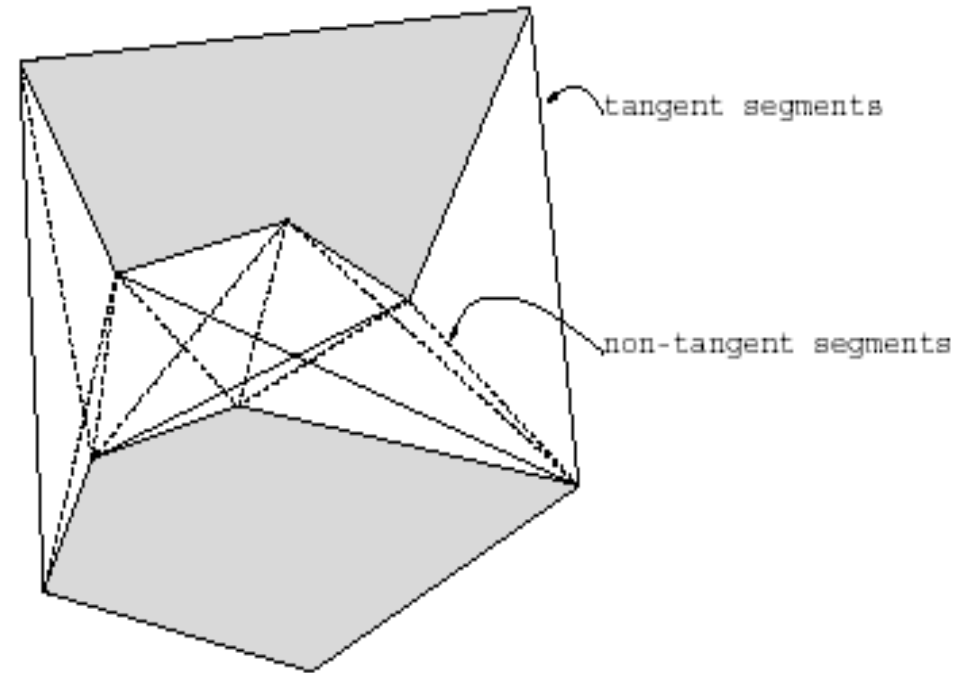
# The Visibility Graph (Done)

- Repeat until you're done.



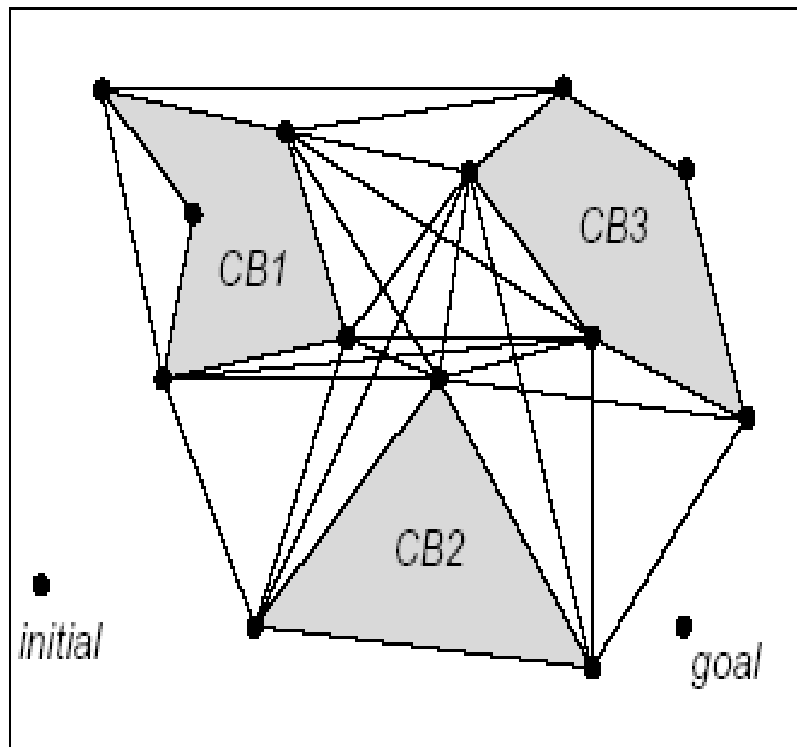
# Reduced Visibility Graphs

- *Idea:* we don't really need all the edges in the visibility graph (even if we want shortest paths)
- *Definition:* Let  $L$  be the line passing through an edge  $(x,y)$  in the visibility graph  $G$ . The segment  $(x,y)$  is a tangent segment iff  $L$  is tangent to  $CB$  at both  $x$  and  $y$ .

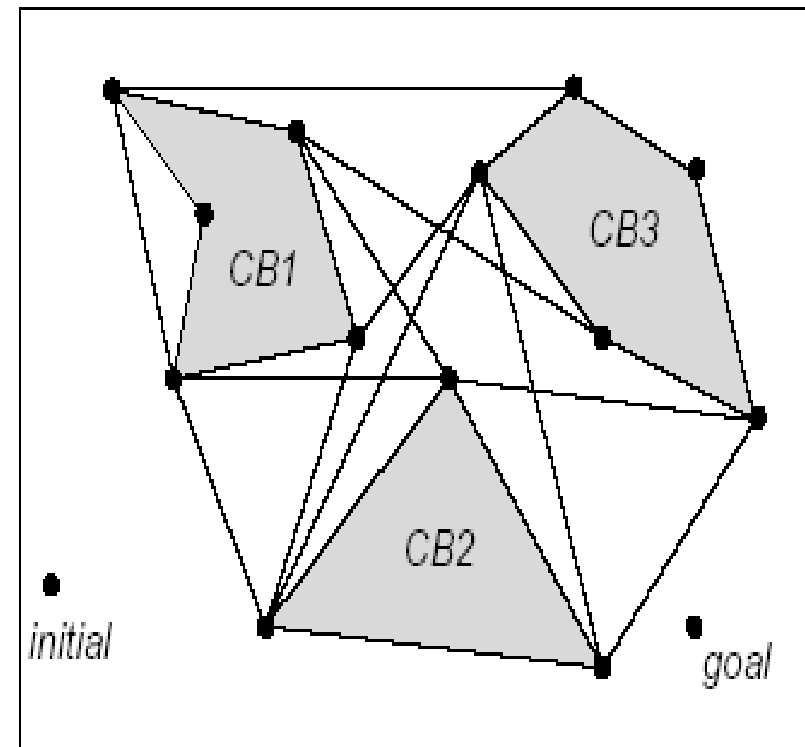


# Reduced Visibility Graphs

Visibility Graph



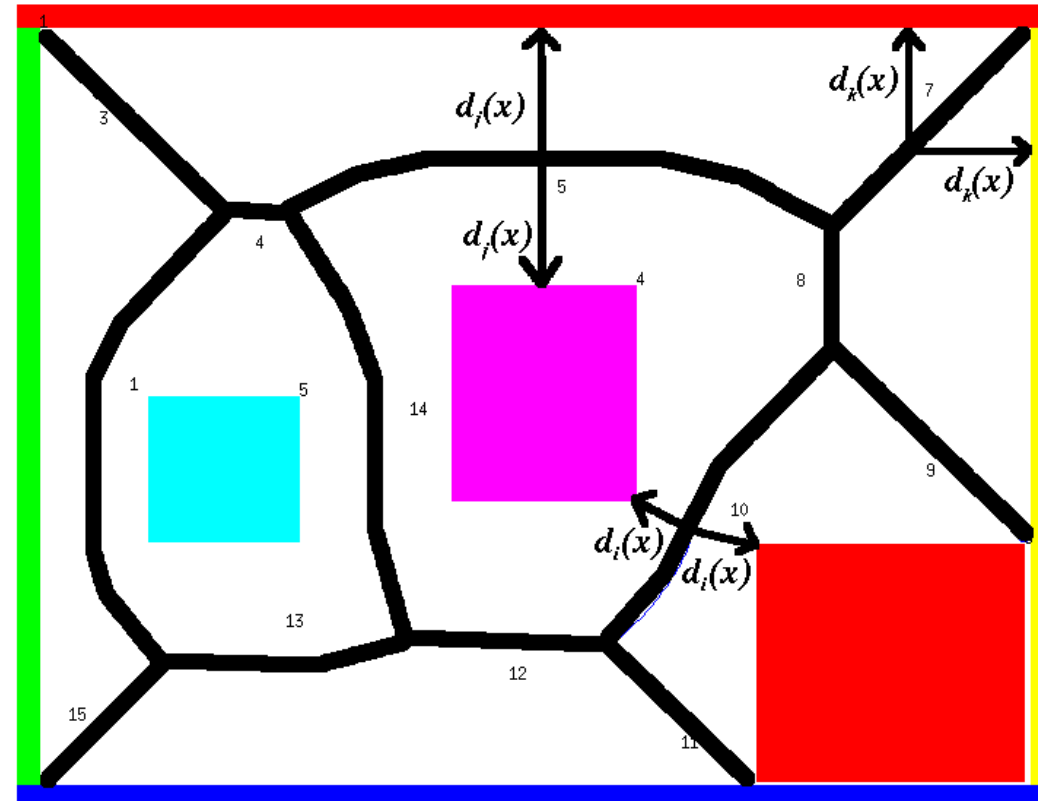
Reduced Visibility Graph



# MAP-BASED APPROACHES: THE RETRACTION APPROACH

# Retraction Example: Generalized Voronoi Diagrams

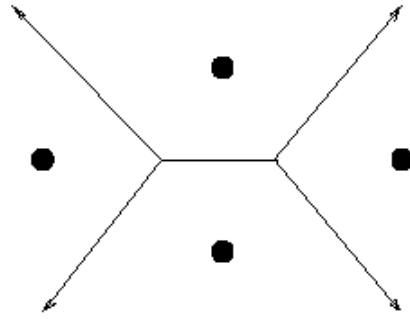
- A GVG is formed by paths equidistant from the two closest objects



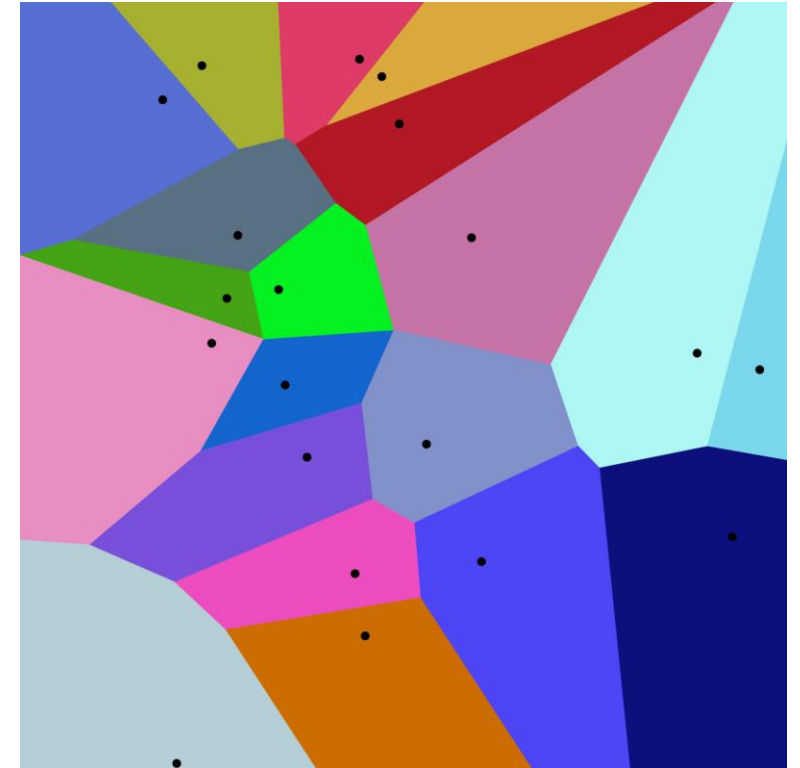
- This generates a very safe roadmap which avoids obstacles as much as possible

# Retraction Example: Generalized Voronoi Diagrams

*Example: Voronoi Diagram for point sets (original)*



- Voronoi diagram of point set  $X$  consists of straight line segments
- constructed by
  - computing lines bisecting each pair of points and their intersections
  - computing intersections of these lines
  - keeping segments with more than one nearest neighbor

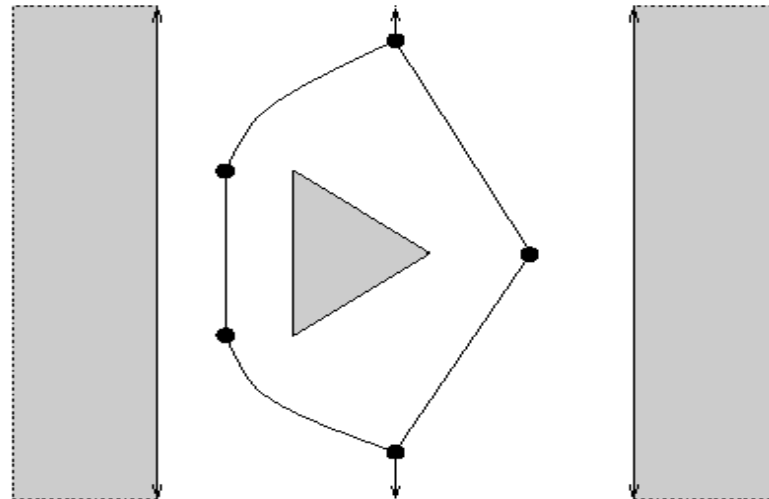




# Generalized Voronoi Diagrams

When  $\mathcal{C} = \mathbb{R}^2$  and polygonal  $\mathcal{CB}$ ,  $Vor(\mathcal{C}_{free})$  consists of a finite collection of straight line segments and *parabolic curve segments* (called **arcs**)

- **straight arcs** are defined by two vertices or two edges of  $\mathcal{CB}$ , i.e., the set of points equally close to two points (or two line segments) is a line
- **parabolic arcs** are defined by one vertex and one edge of  $\mathcal{CB}$ , i.e., the set of points equally close to a point and a line is a parabola



# Generalized Voronoi Diagrams

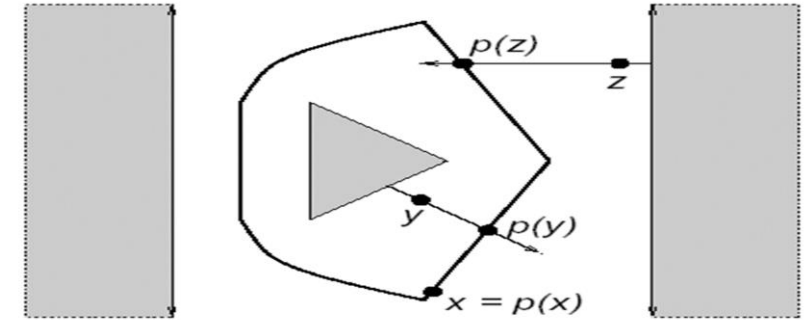
To use  $Vor(\mathcal{C}_{free})$  as our roadmap  $R$ , we need to define the retraction

$$\rho : \mathcal{C}_{free} \rightarrow Vor(\mathcal{C}_{free})$$

Case 1:  $\mathbf{q} \in Vor(\mathcal{C}_{free})$ :  $\rho(\mathbf{q}) = \mathbf{q}$

Case 2:  $\mathbf{q} \notin Vor(\mathcal{C}_{free})$

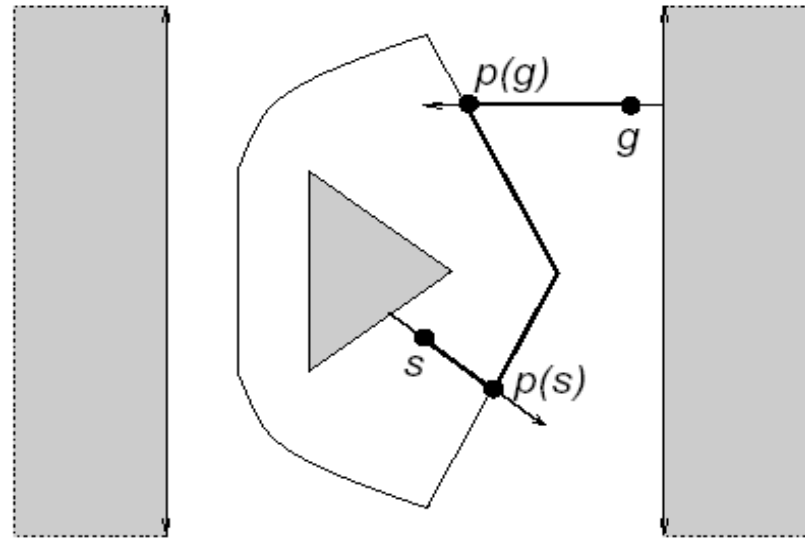
- let  $\mathbf{p}$  be the closest point of the boundary of  $\mathcal{C}_{free}$  to  $\mathbf{q}$
- let  $L$  be the ray from  $\mathbf{p}$  passing through  $\mathbf{q}$  ( $L$  follows the steepest ascent of the *clearance()* function from  $\mathbf{p}$ )
- define  $\rho(\mathbf{q})$  to be the intersection of  $L$  with  $Vor(\mathcal{C}_{free})$



# Generalized Voronoi Diagrams

To find a path:

1. compute  $Vor(\mathcal{C}_{free})$
2. find paths from  $\mathbf{q}_{init}$  and  $\mathbf{q}_{goal}$  to  $\rho(\mathbf{q}_{init})$  and  $\rho(\mathbf{q}_{goal})$ , respectively
3. search  $Vor(\mathcal{C}_{free})$  for a set of arcs connecting  $\rho(\mathbf{q}_{init})$  and  $\rho(\mathbf{q}_{goal})$



# **MAP-BASED APPROACHES: CELL DECOMPOSITION METHODS**

# Cell Decomposition Methods

## PREPROCESSING:

- represent  $\mathcal{C}_{free}$  as a collection of *cells* (connected regions of  $\mathcal{C}_{free}$ )  
     $\implies$  planning between configurations in the same cell should be 'easy'
- build *connectivity graph* representing adjacency relations between cells  
     $\implies$  cells adjacent if can move directly between them

## QUERY PROCESSING:

1. locate cells  $k_{init}$  and  $k_{goal}$  containing start and goal configurations
2. search the connectivity graph for a 'channel' or sequence of adjacent cells connecting  $k_{init}$  and  $k_{goal}$
3. find a path that is contained in the channel of cells

# Cell Decomposition Methods

Two major variants of methods:

- *exact cell decomposition:*
  - set of cells exactly covers  $\mathcal{C}_{free}$
  - complicated cells with irregular boundaries (contact constraints)
  - harder to compute
- *approximate cell decomposition:*
  - set of cells approximately covers  $\mathcal{C}_{free}$
  - simpler cells with more regular boundaries
  - easier to compute

# Exact Cell Decomposition Method

*Idea:* decompose  $\mathcal{C}_{free}$  into a collection  $\mathcal{K}$  of non-overlapping *cells* such that the union of all the cells *exactly* equals the free C-space, i.e.,  $\mathcal{C}_{free} = \cup_{k \in \mathcal{K}} k$

## Cell Characteristics:

- geometry of cells should be simple so that it is easy to compute a path between any two configurations in a cell
- it should be pretty easy to test the adjacency of two cells, i.e., whether they share a boundary

## Exact Cell Decomposition Method

### Definitions:

A **convex polygonal decomposition**  $\mathcal{K}$  of  $\mathcal{C}_{free}$  is a finite collection of convex polygons, called **cells**, such that the interiors of any two cells do not intersect and the union of all cells is  $\mathcal{C}_{free}$ .

Two cells  $k, k' \in \mathcal{K}$  are **adjacent** iff  $k \cap k'$  is a line segment of non-zero length (i.e., not a single point)

The **connectivity graph** associated with a convex polygonal decomposition  $\mathcal{K}$  of  $\mathcal{C}_{free}$  is an undirected graph  $G$  where

- nodes in  $G$  correspond to cells in  $\mathcal{K}$
- nodes connected by edge in  $G$  iff corresponding cells adjacent in  $\mathcal{K}$



# Path Planning with a Convex Polygonal Decomposition

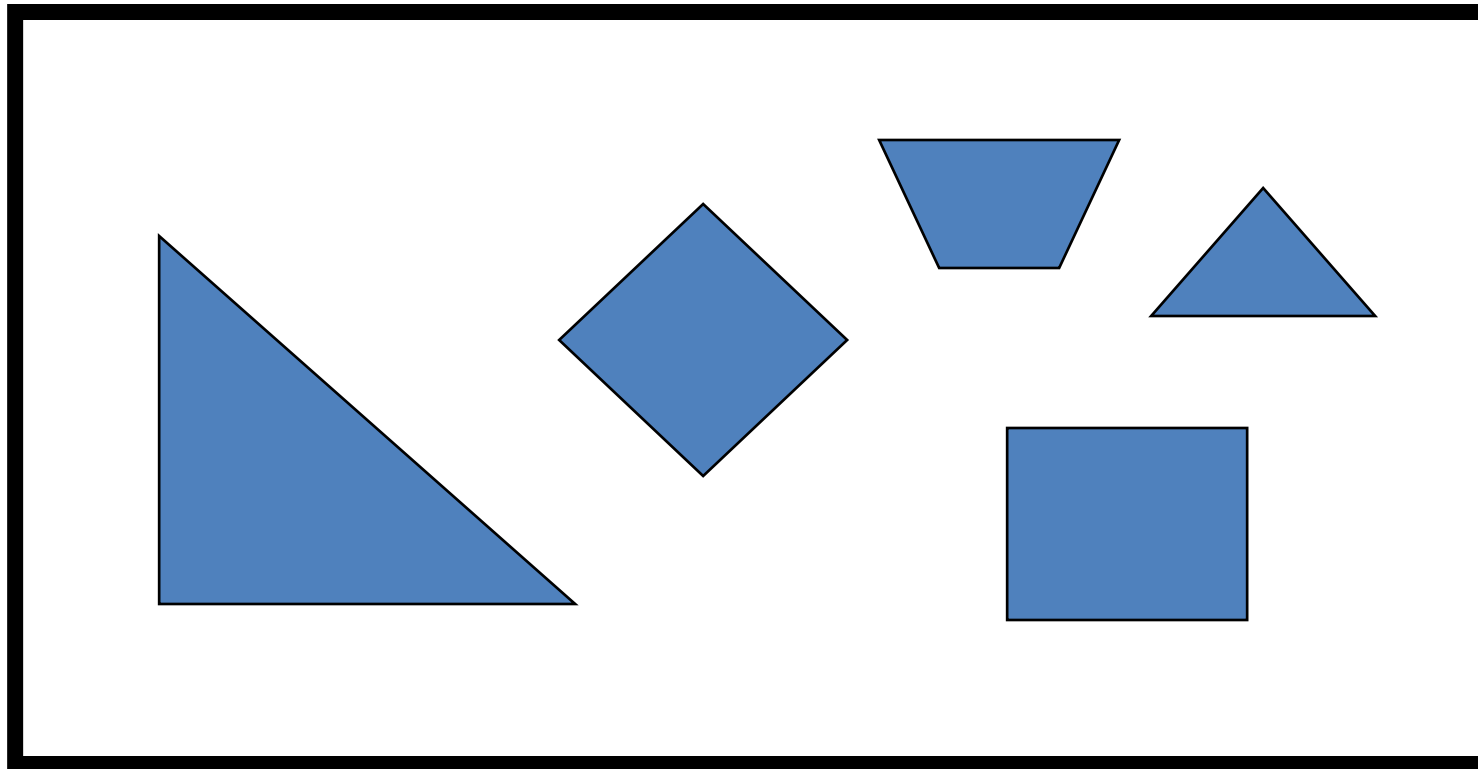
**input:** configurations  $\mathbf{q}_{init}$  and  $\mathbf{q}_{goal}$ , and  $\mathcal{CB}$  which is a polygonal region

**output:** a path in  $\mathcal{C}_{free}$  connecting  $\mathbf{q}_{init}$  and  $\mathbf{q}_{goal}$

1. Build  $\mathcal{K}$ , the convex polygonal decomposition of  $\mathcal{CB}$
2. Construct the connectivity graph  $G$  of  $\mathcal{K}$
3. locate the cells  $k_{init}$  and  $k_{goal}$  in  $\mathcal{K}$  containing  $\mathbf{q}_{init}$  and  $\mathbf{q}_{goal}$
4. find a path in  $G$  between the nodes corresponding to  $k_{init}$  and  $k_{goal}$ 
  - corresponds to a sequence of cells forming a **channel** in  $\mathcal{C}_{free}$
5. find a free path from  $\mathbf{q}_{init}$  to  $\mathbf{q}_{goal}$  in the channel

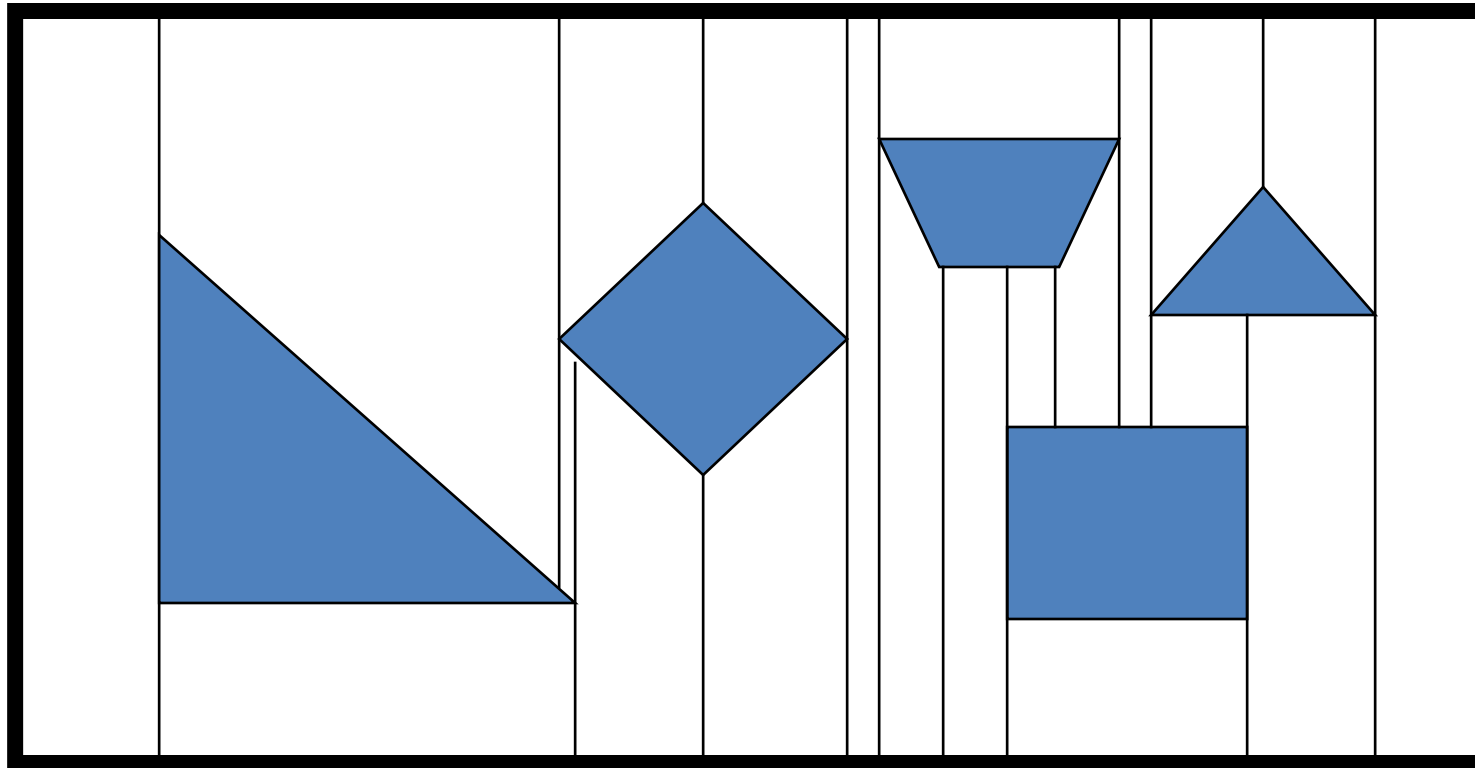
# Exact Cell Decompositions: Trapezoidal Decomposition

- A way to divide the world into smaller regions
- Assume a polygonal world



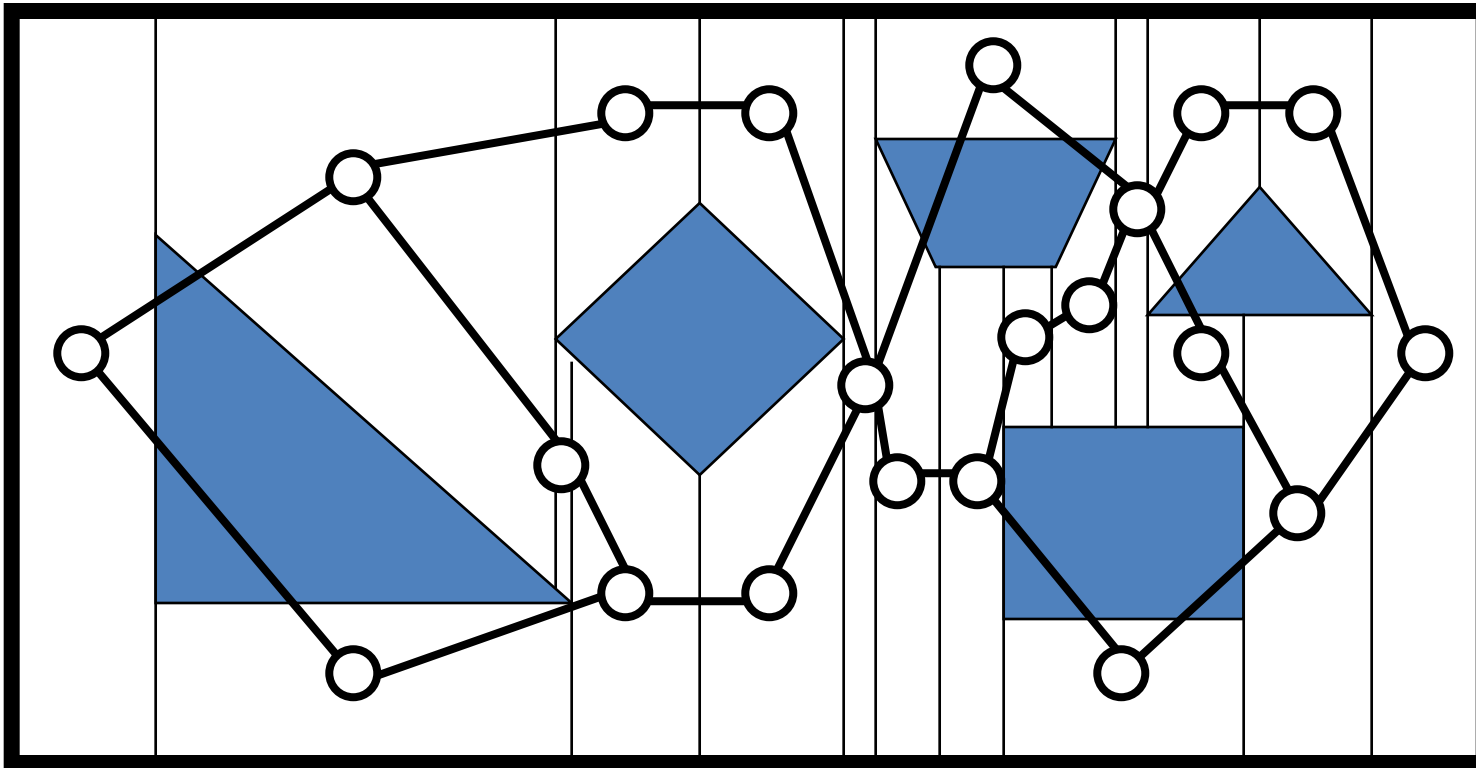
# Exact Cell Decompositions: Trapezoidal Decomposition

**Basic Idea:** at every vertex of  $\mathcal{CB}$ , extend a vertical line up and down in  $\mathcal{C}_{free}$  until it touches a C-obstacle or the boundary of  $\mathcal{C}_{free}$



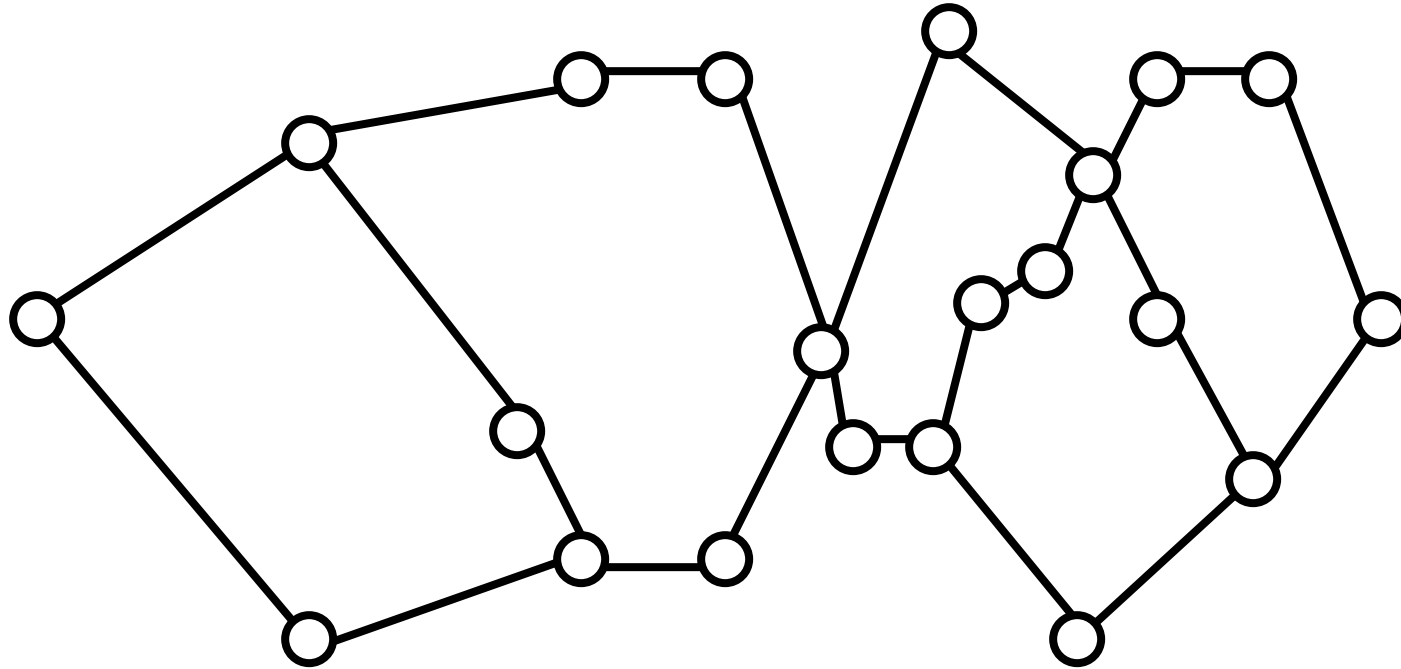
# Applications: Coverage

- By reducing the world to cells, we've essentially abstracted the world to a graph.



# Find a path

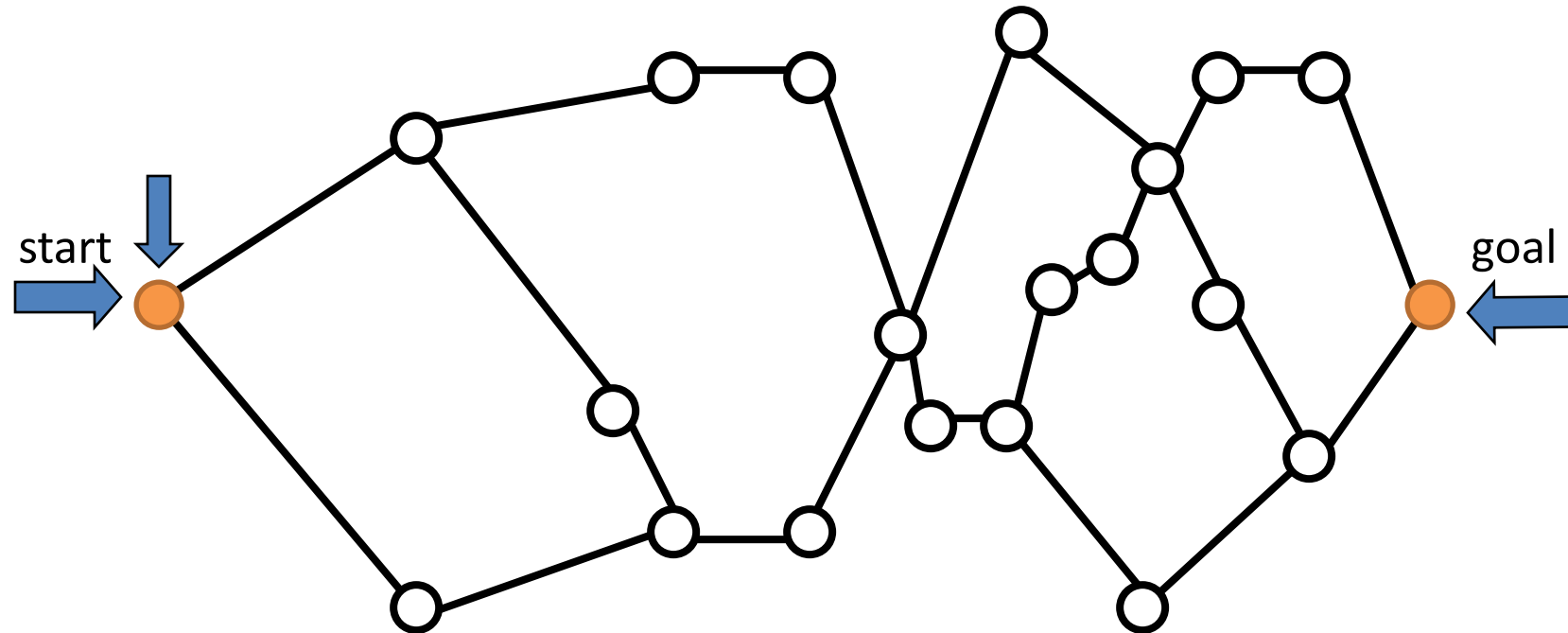
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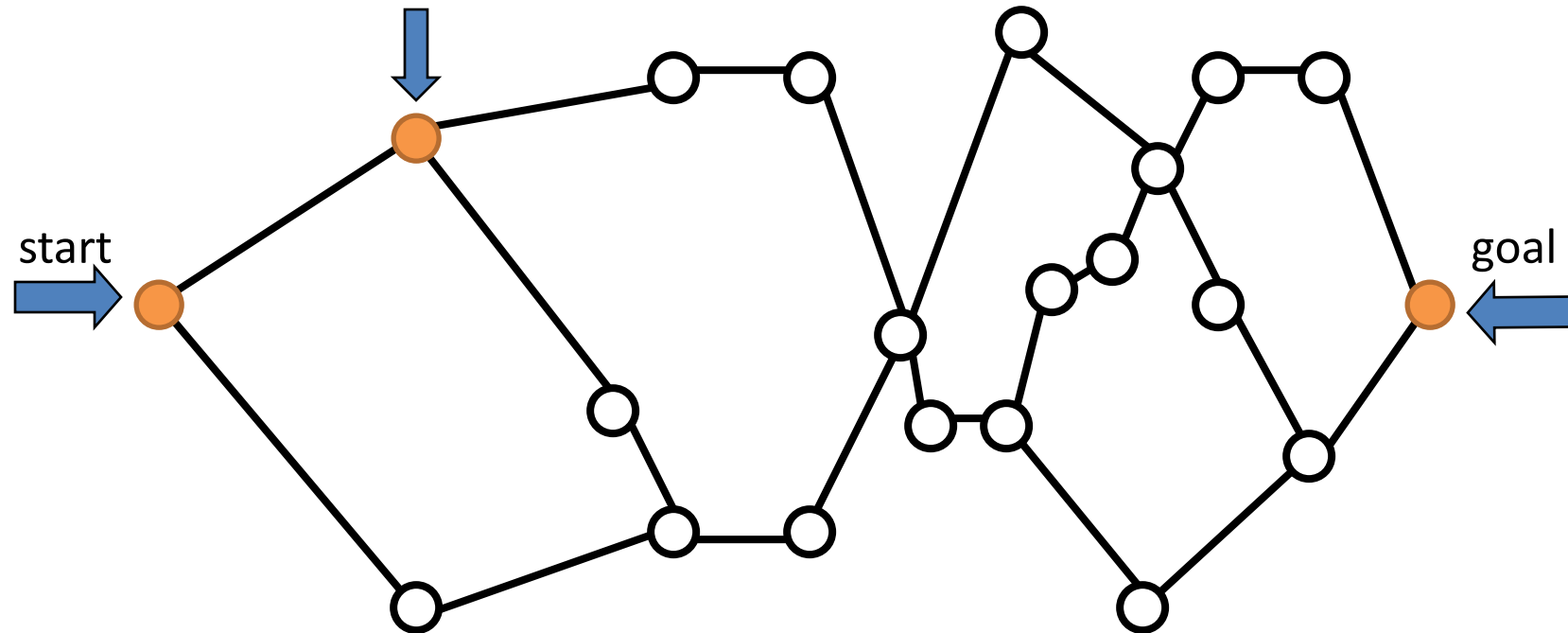
# Find a path

- With an adjacency graph, a path from start to goal can be found by simple traversal



# Find a path

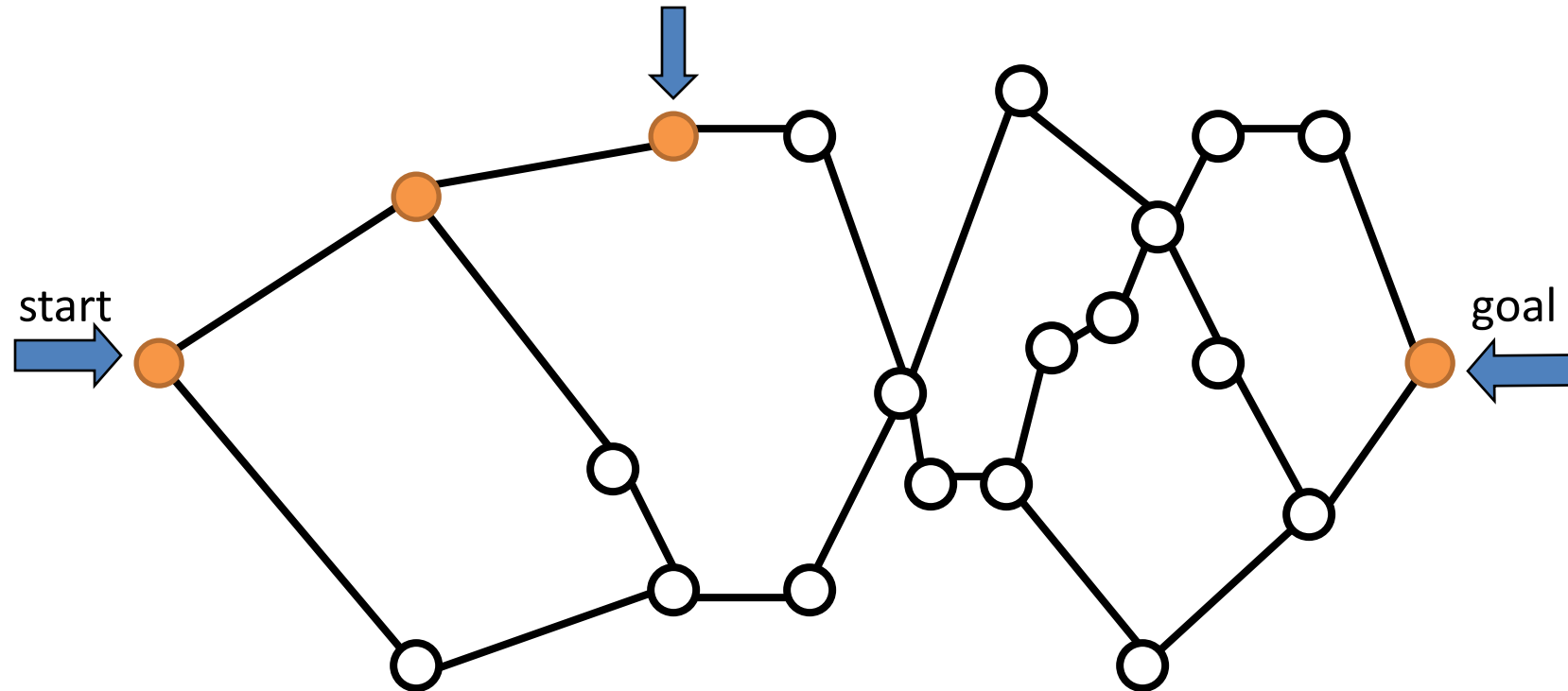
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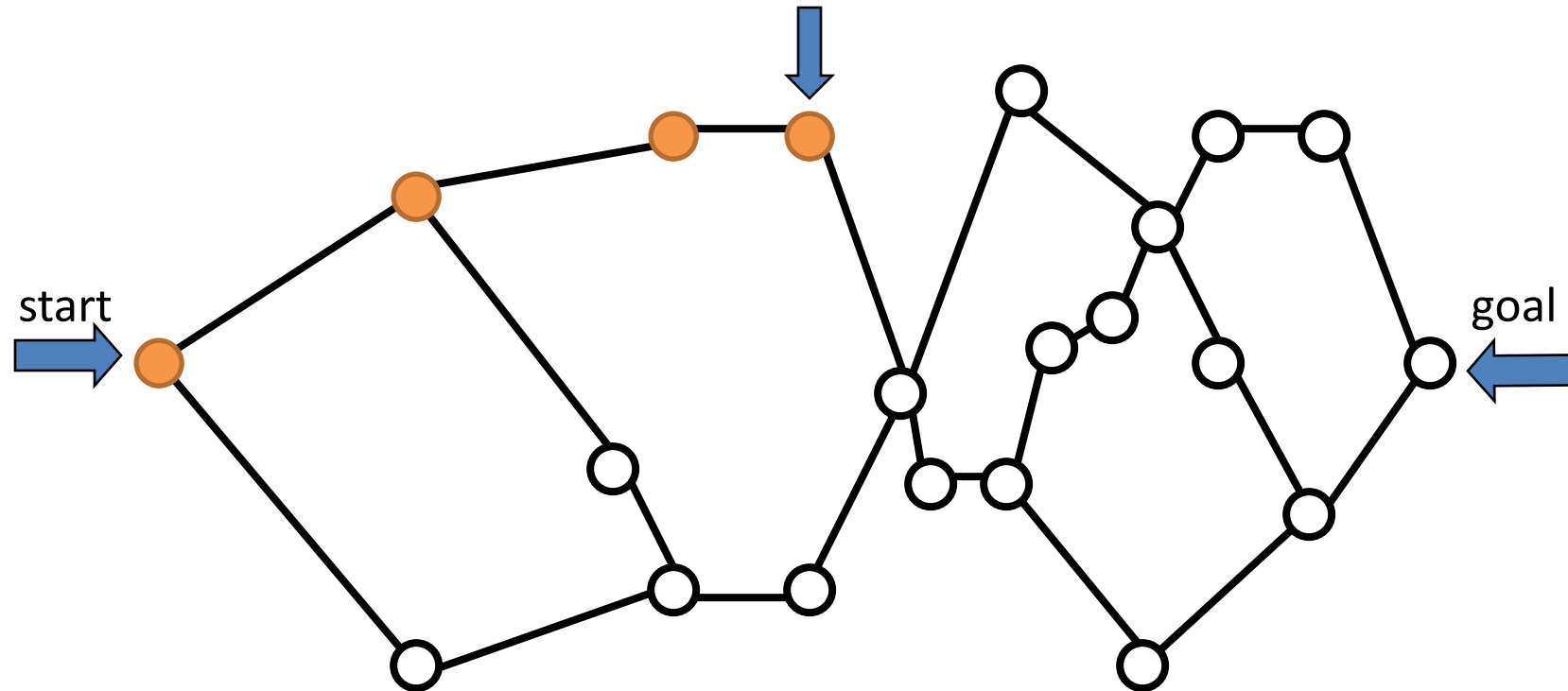
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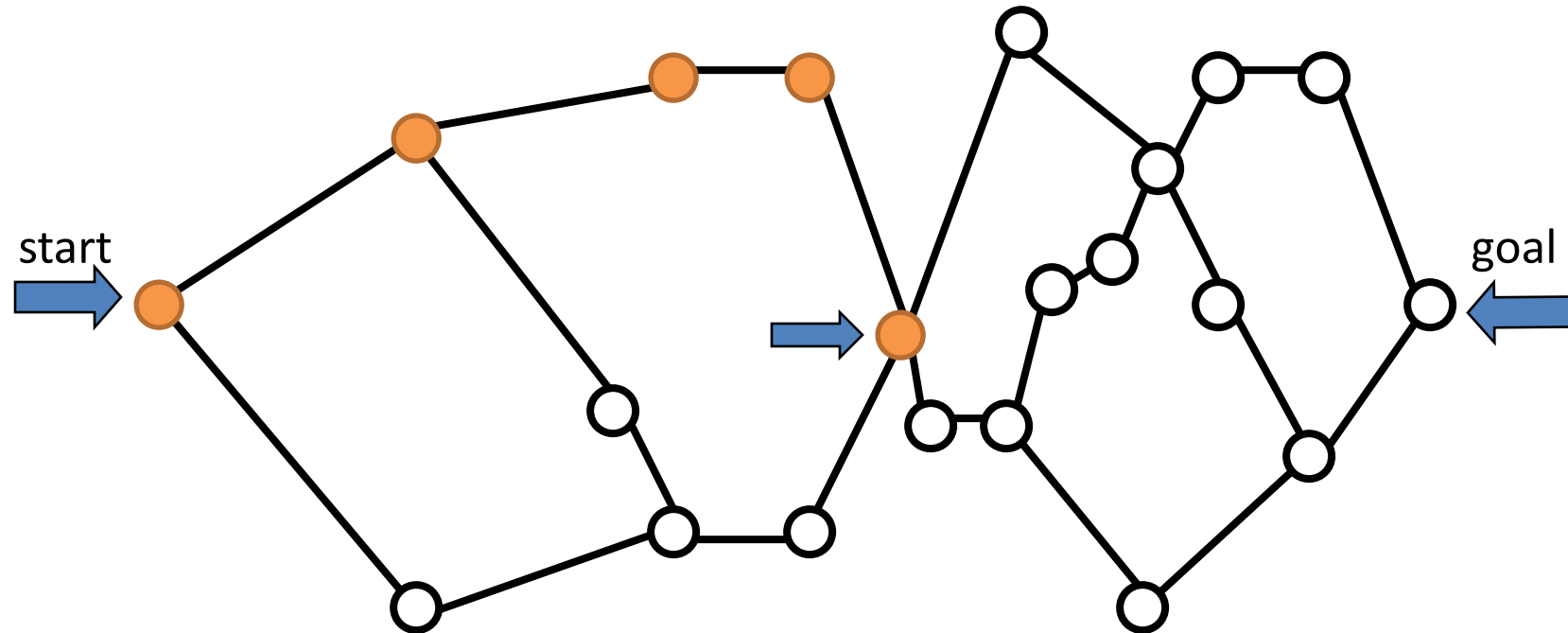
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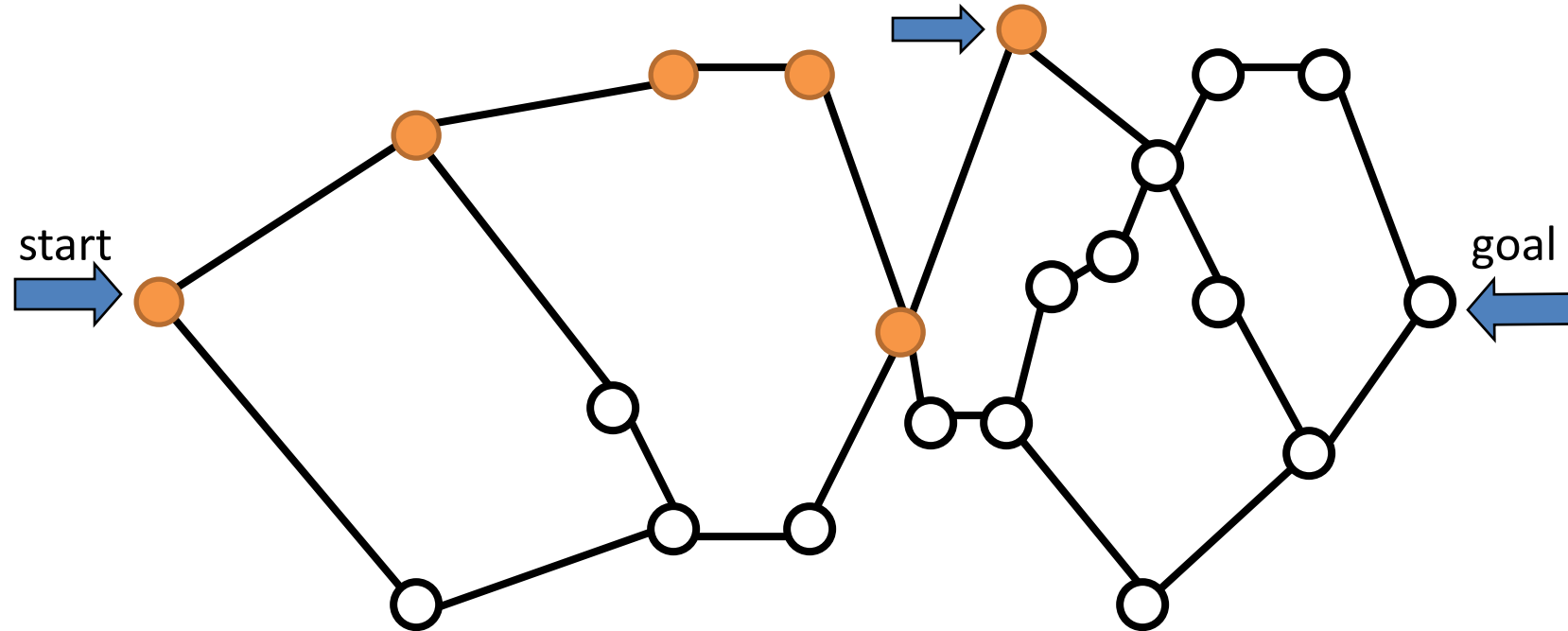
# Find a path

- With an adjacency graph, a path from start to goal can be found by simple traversal



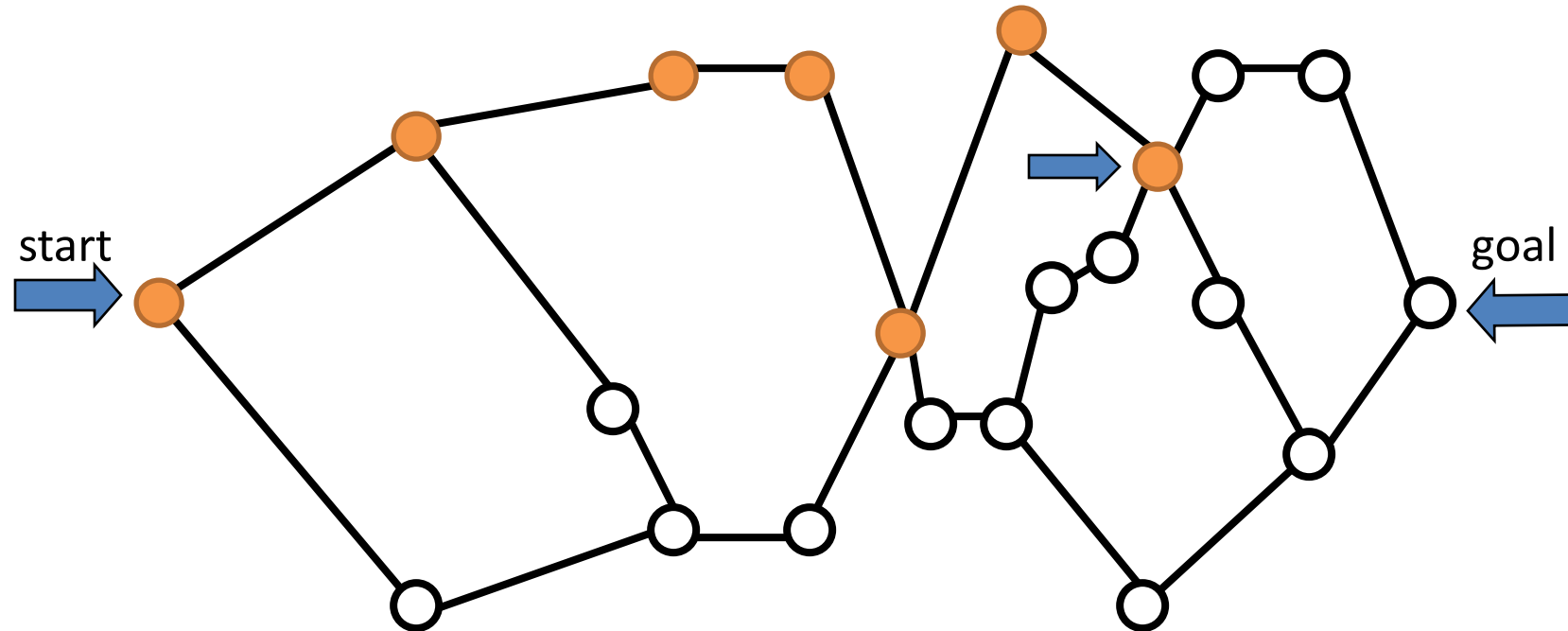
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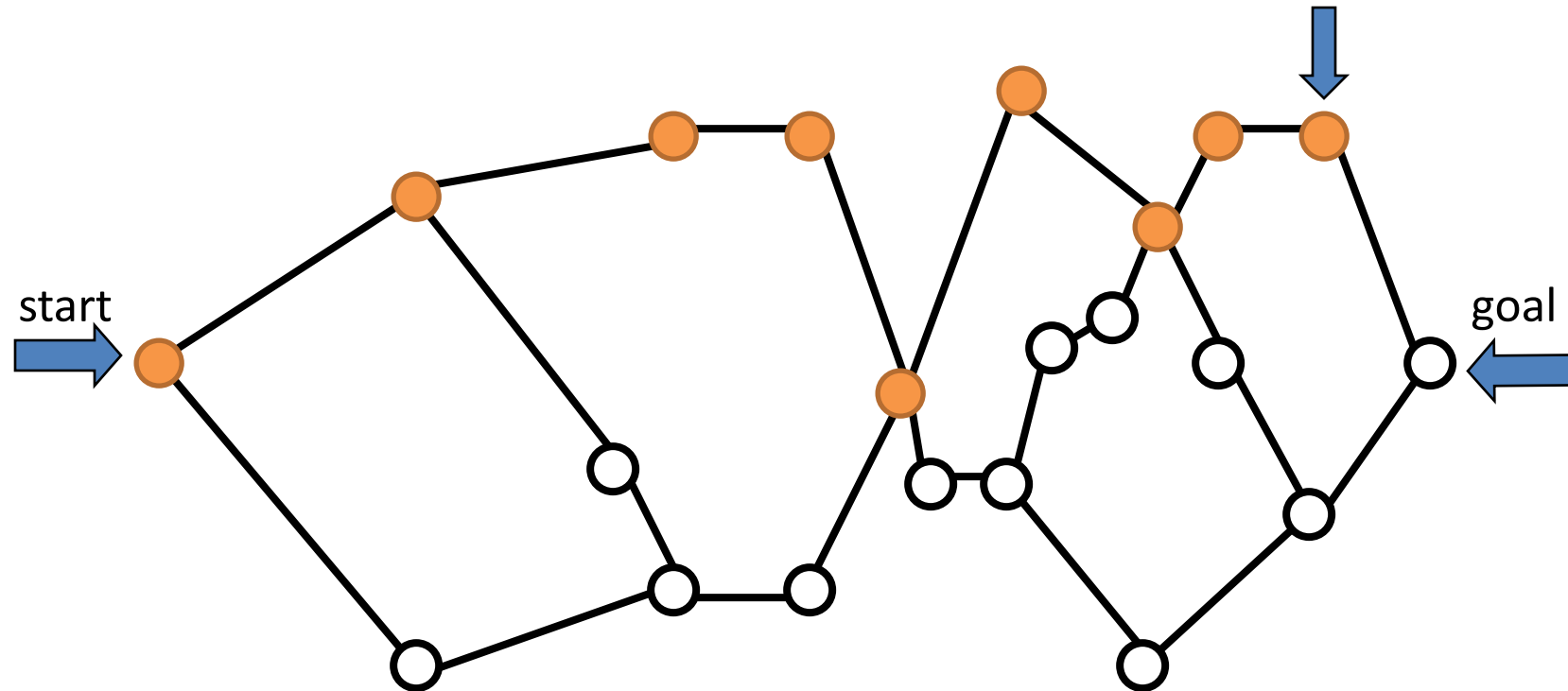
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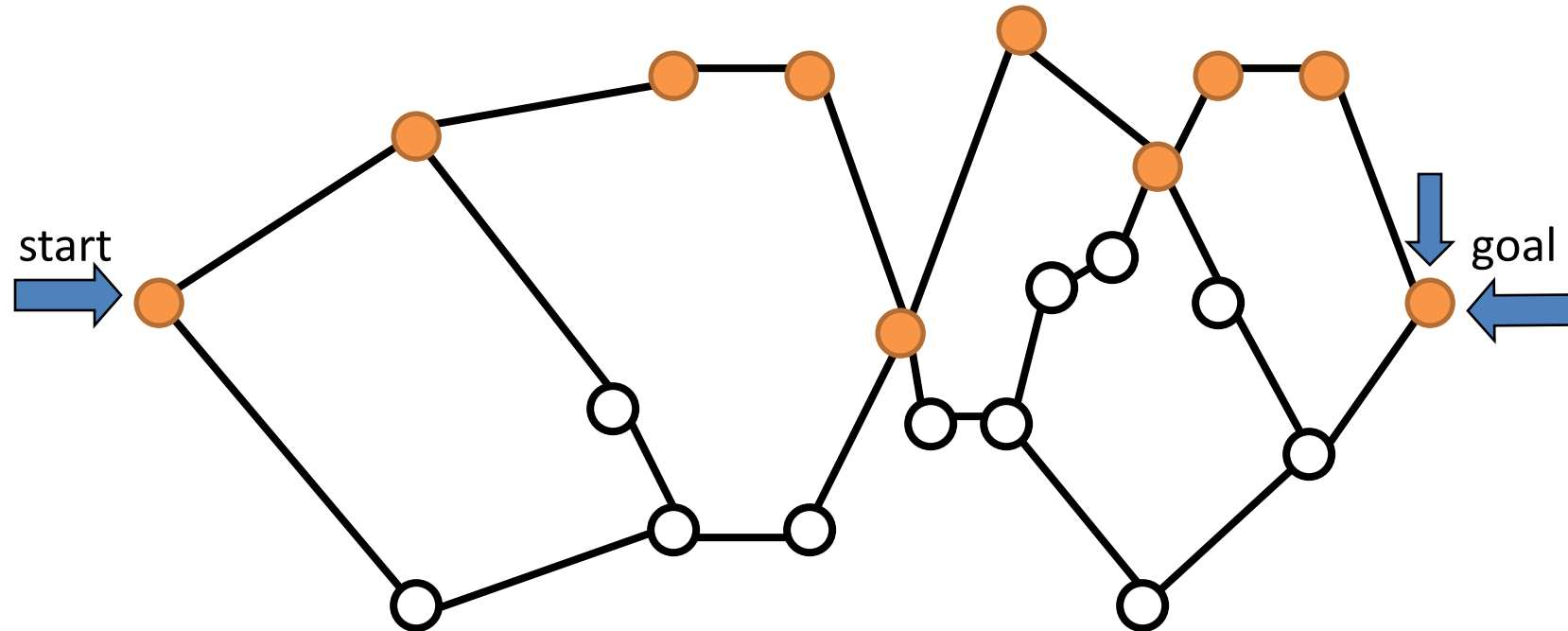
# Find a path

- With an adjacency graph, a path from start to goal can be found by simple traversal



# Find a path

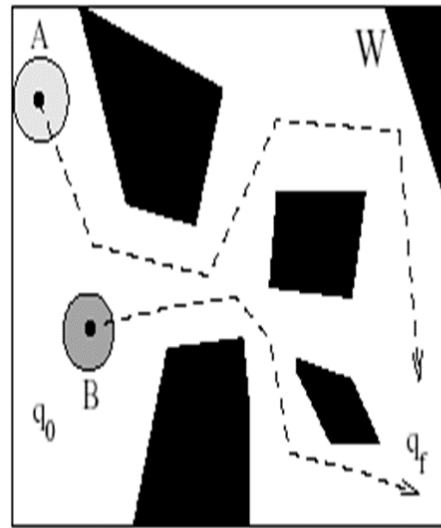
- With an adjacency graph, a path from start to goal can be found by simple traversal



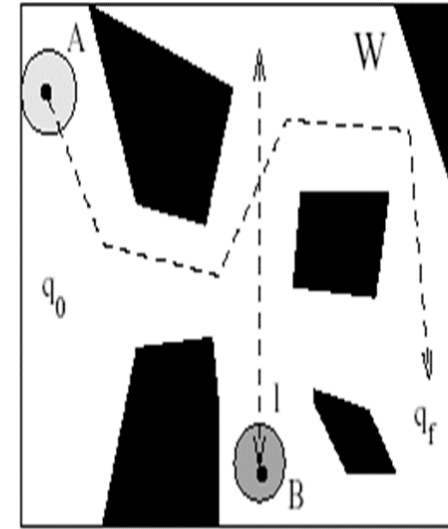


# Extensions to the Basic Problem

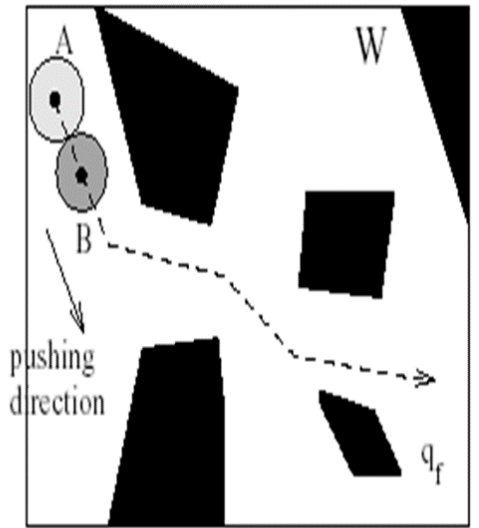
- movable obstacles
- moving obstacles
- multiple robots
- incomplete knowledge/uncertainty in geometry, sensing, etc.



(a) B is a moving robot



(b) B is an obstacle moving on a vertical line



(c) B is an object pushed by the robot