# Discrete Optimization

The Knapsack Problem: Part II

#### Goals of the Lecture

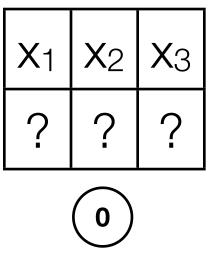
Introduce branch and bound

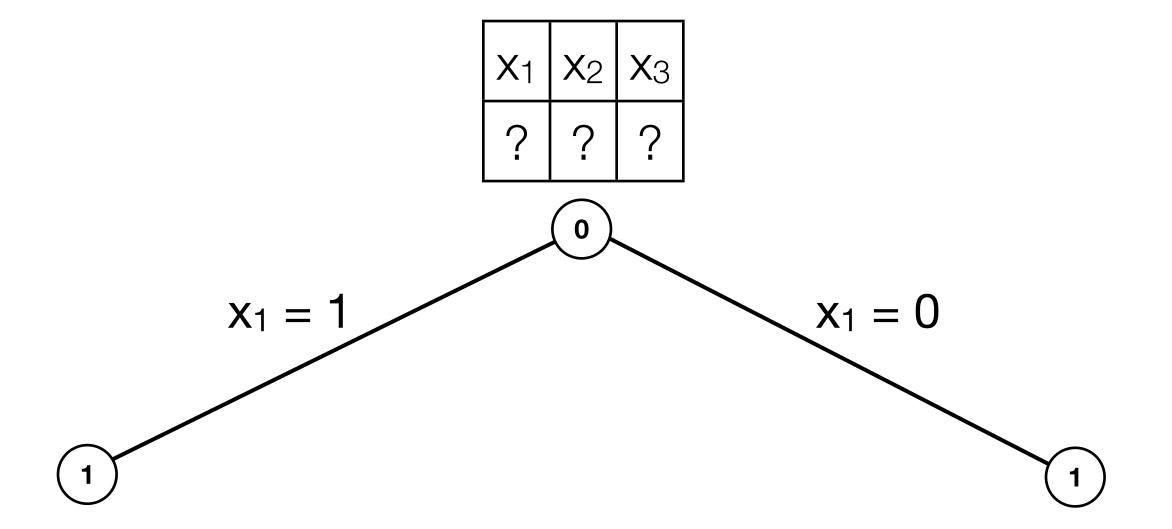
## One-Dimensional Knapsack

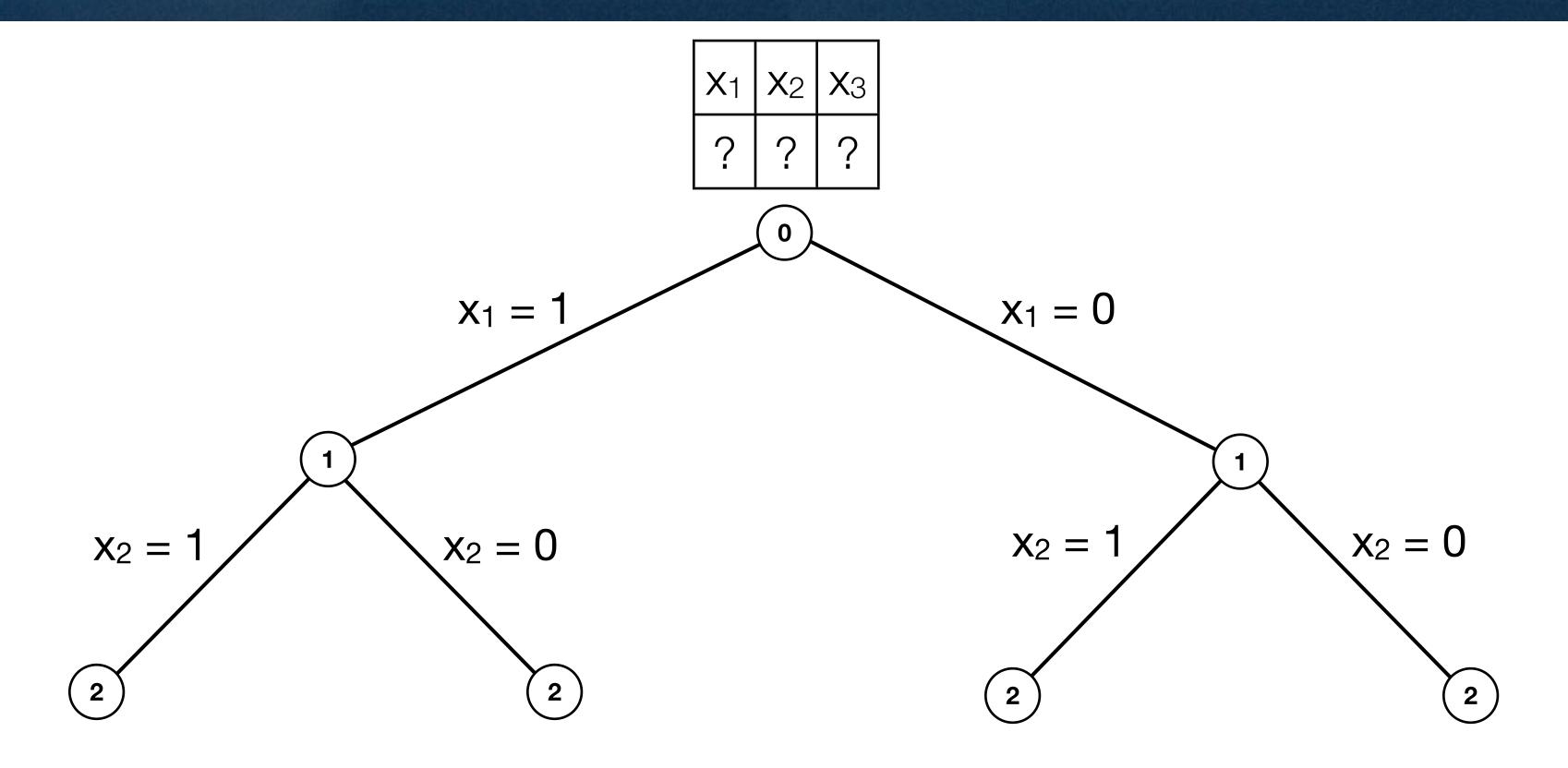
maximize subject to

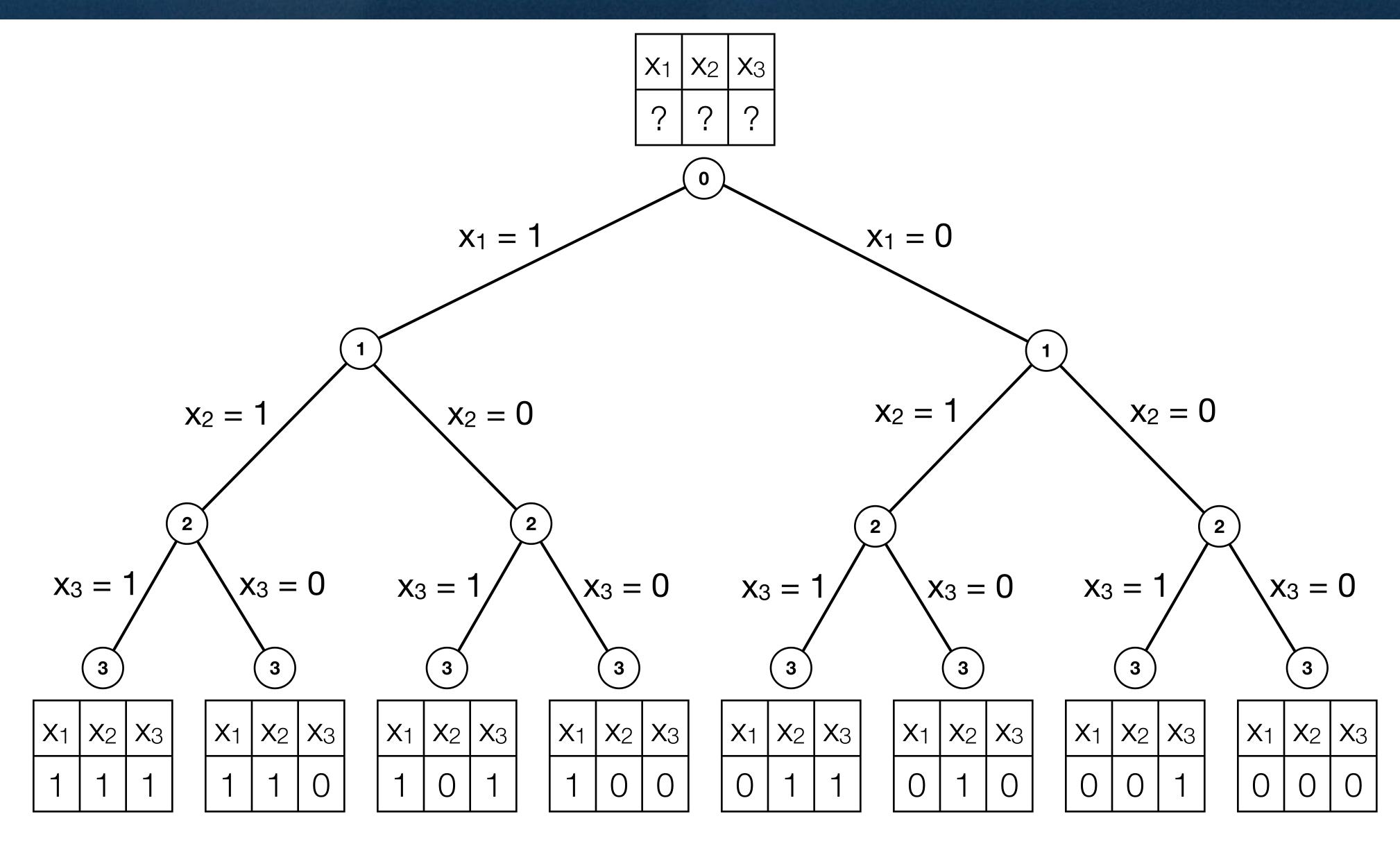
$$45x_1 + 48x_2 + 35x_3$$

$$5x_1 + 8x_2 + 3x_3 \le 10$$
$$x_i \in \{0, 1\} \quad (i \in 1..3)$$









4

- Iterative two steps
  - branching
  - bounding

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  - -split the problem into a number of subproblems
    - like in exhaustive search

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- Branching
  - -split the problem into a number of subproblems
    - like in exhaustive search
- Bounding
  - -find an *optimistic estimate* of the best solution to the subproblem
    - maximization: upper bound
    - minimization: lower bound

► How to find this optimistic estimate?

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  - -Relaxation!

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-Relaxation!

Optimization is the art of relaxation

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► What can we relax?

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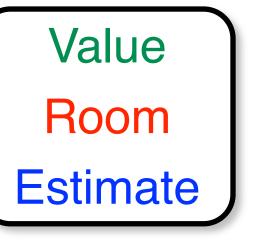
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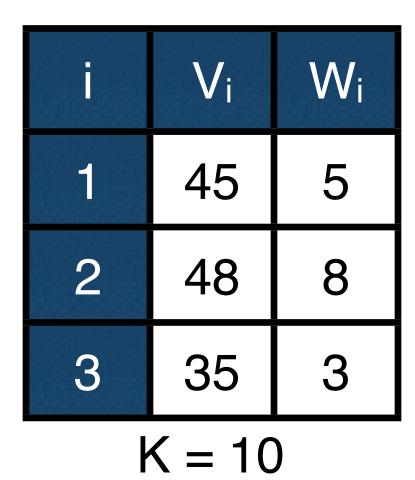
- What can we relax?
  - we can relax the capacity constraint

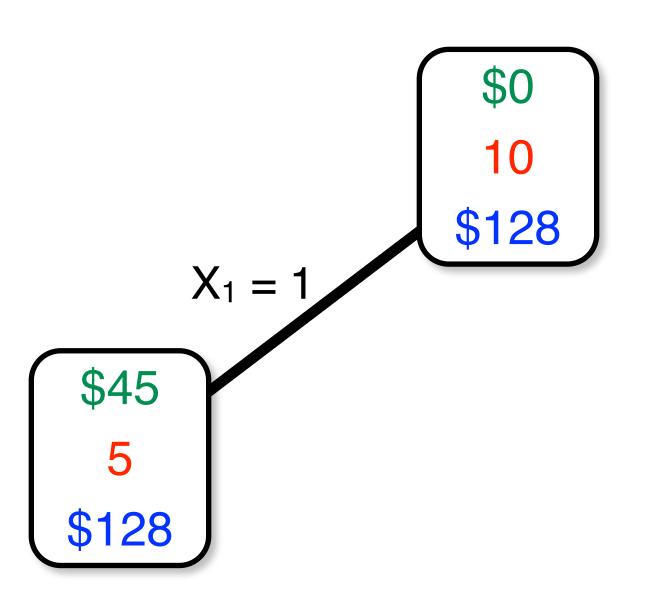
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1	45	5
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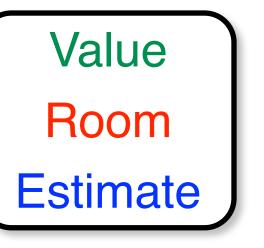
K = 10

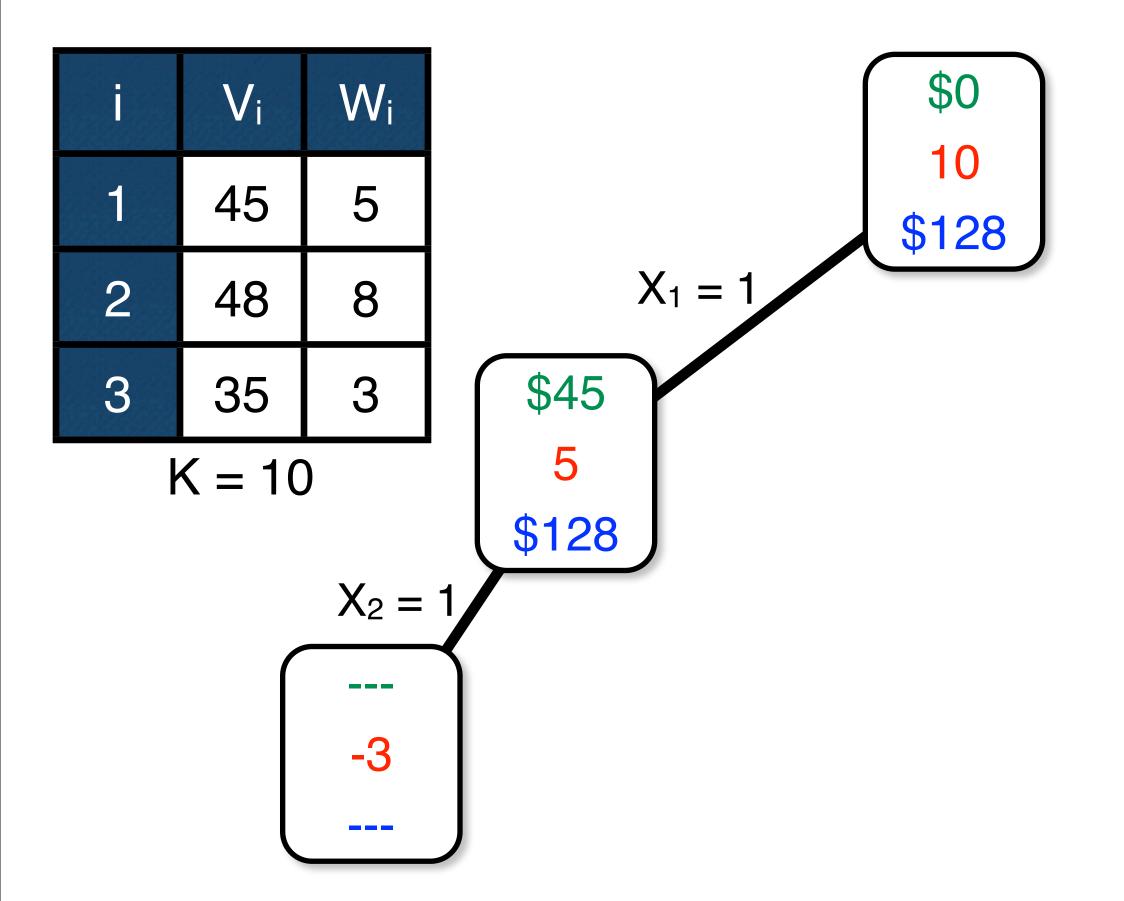


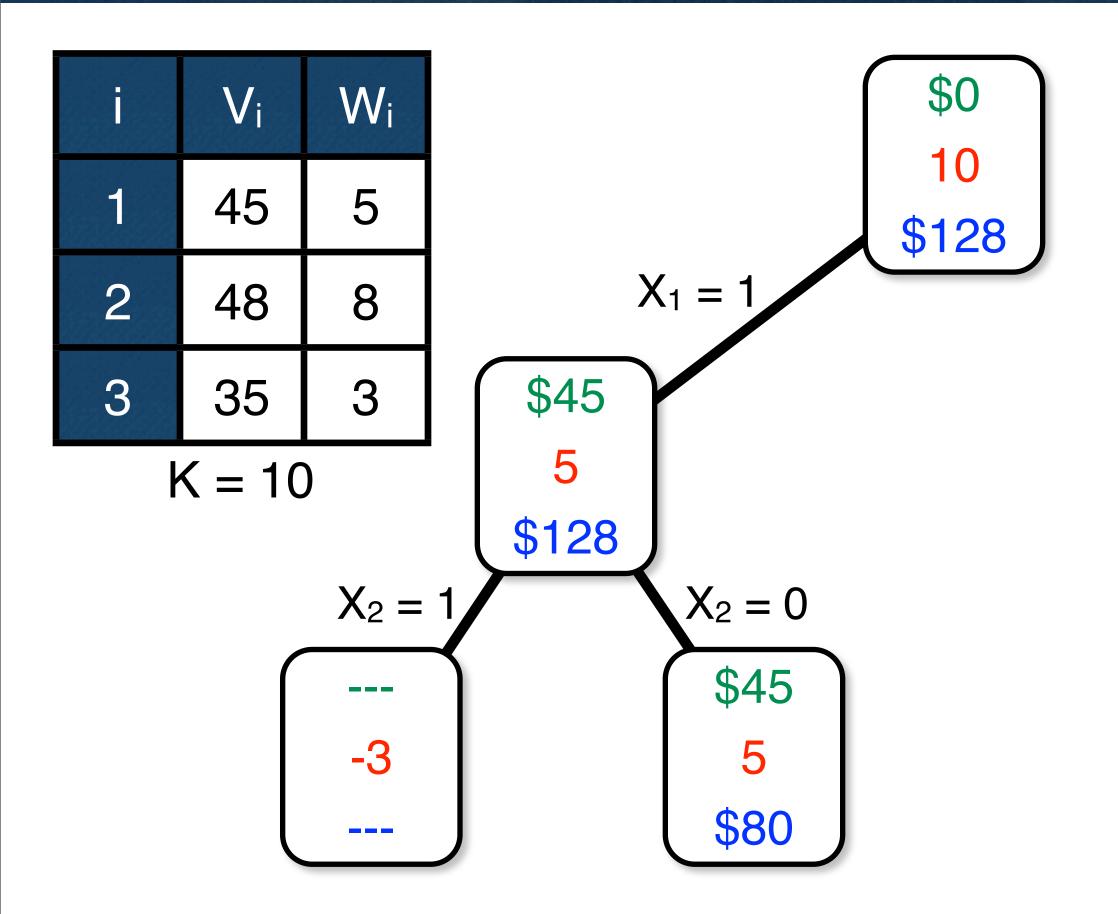


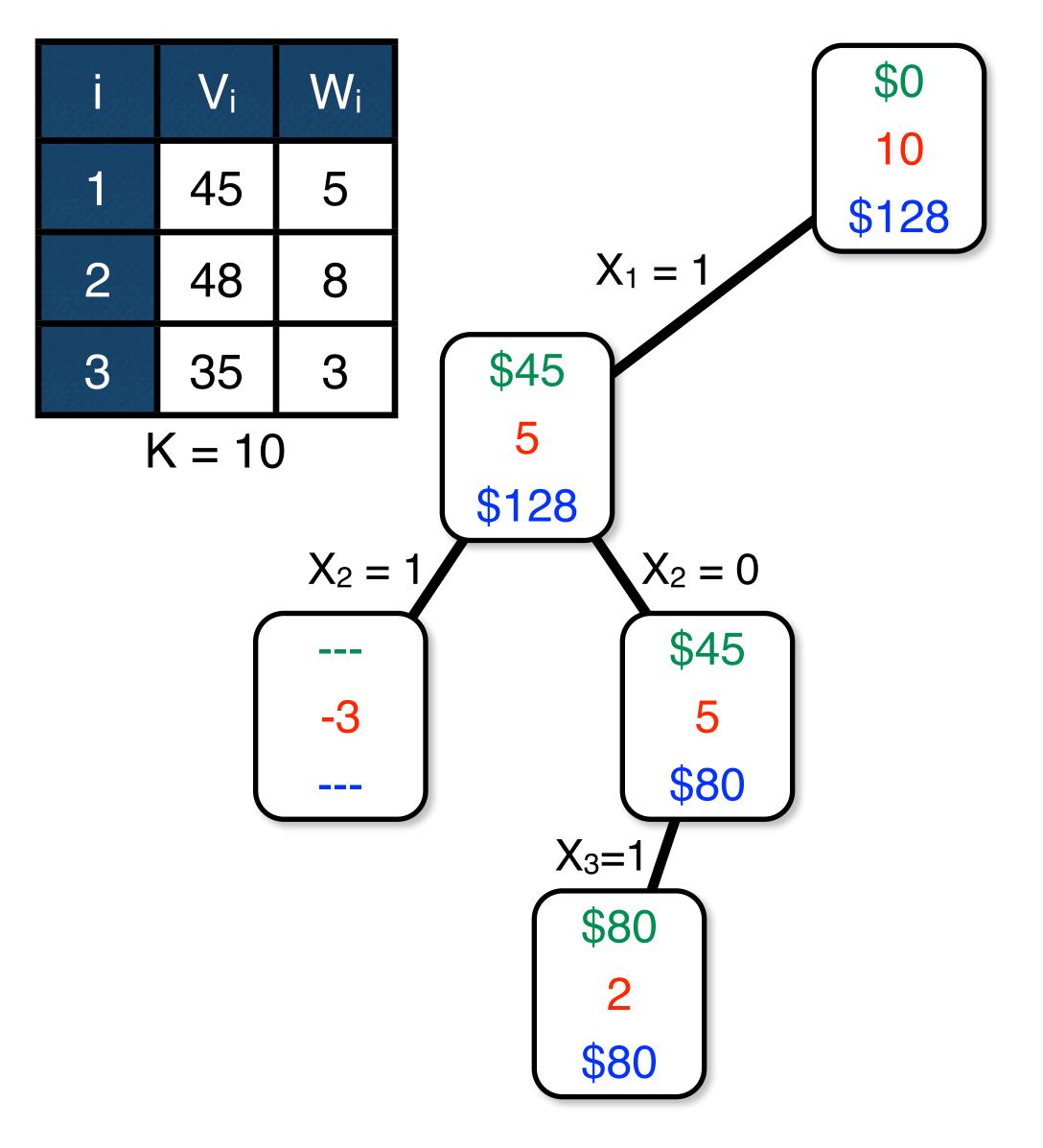


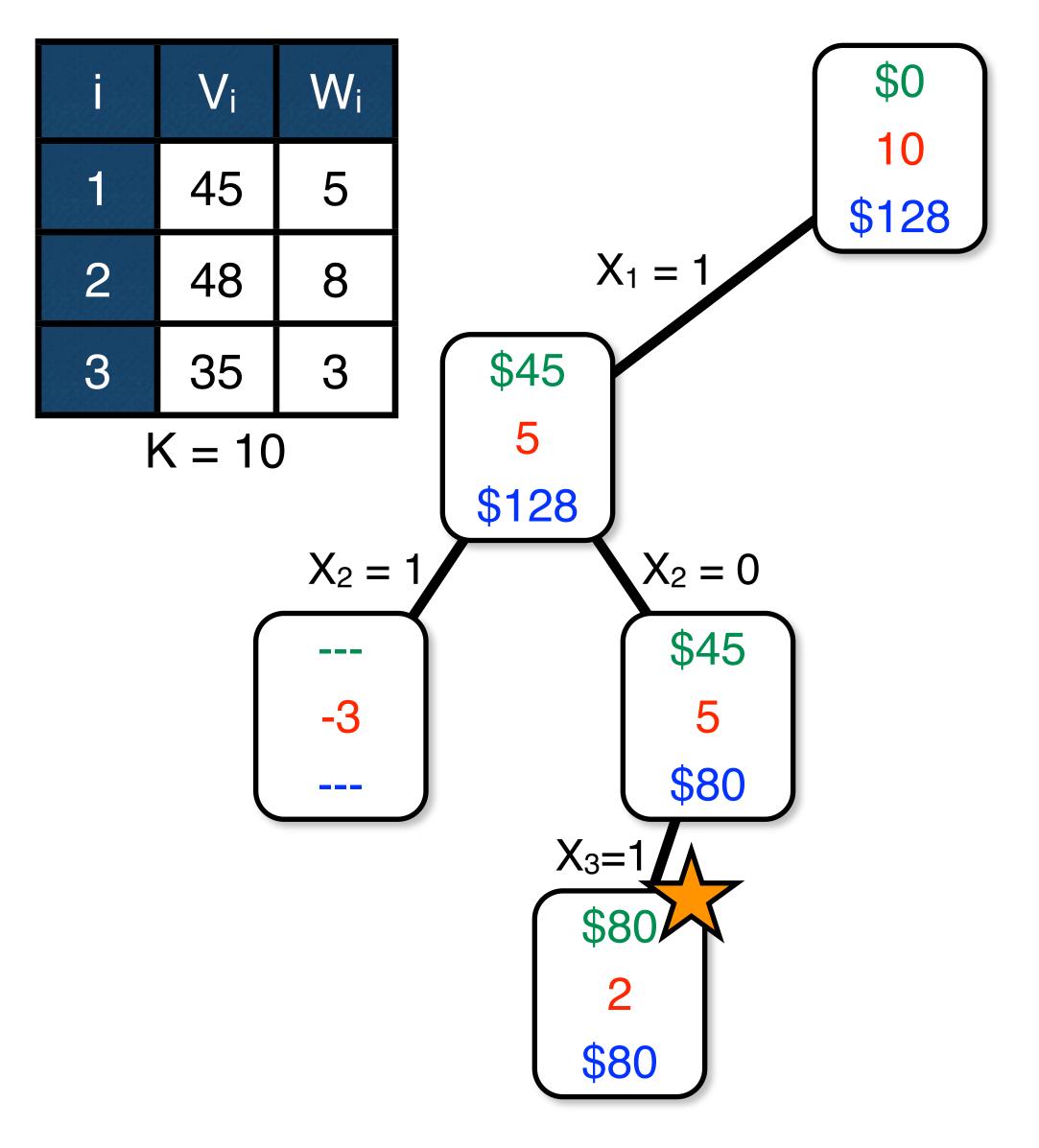


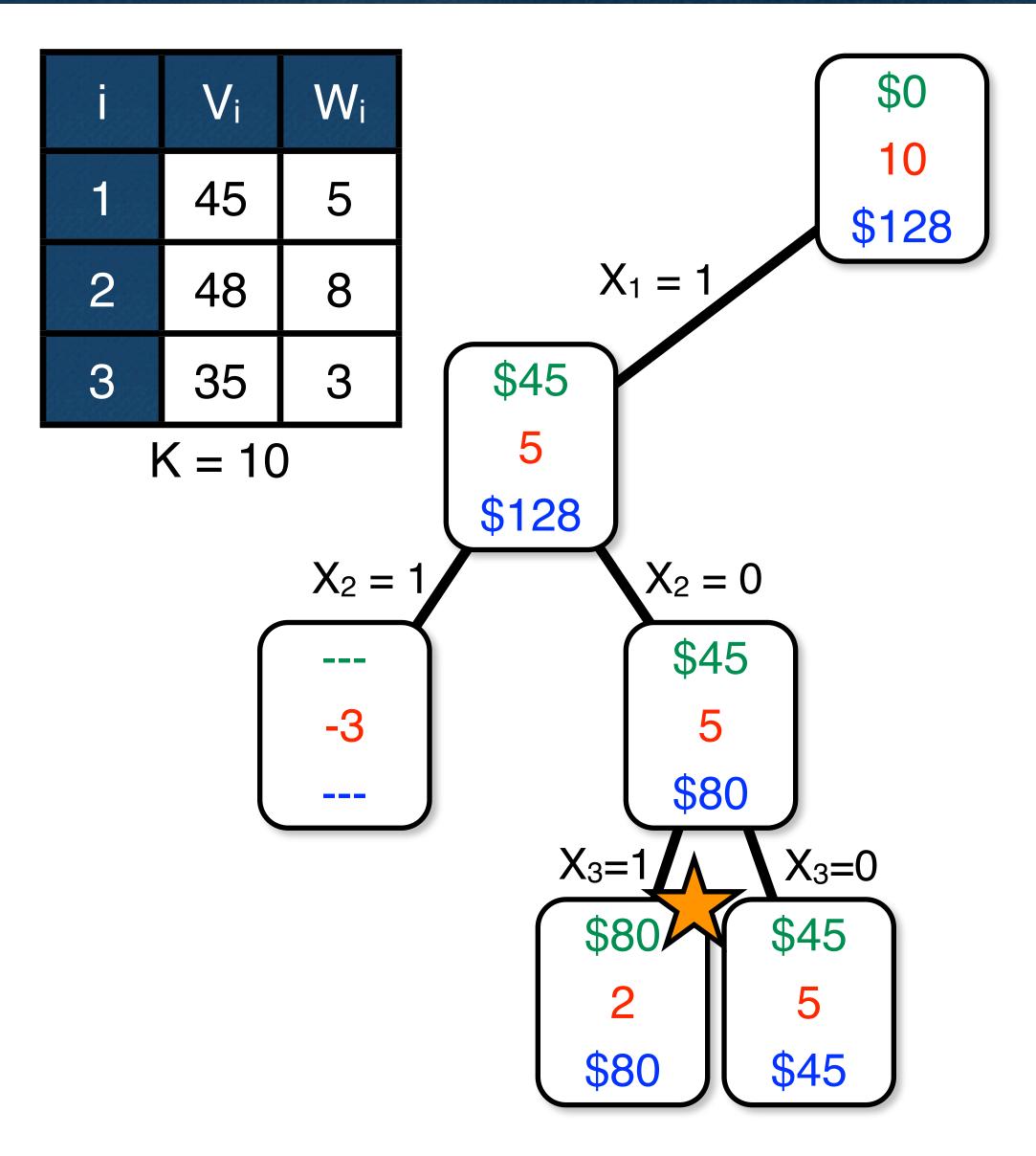


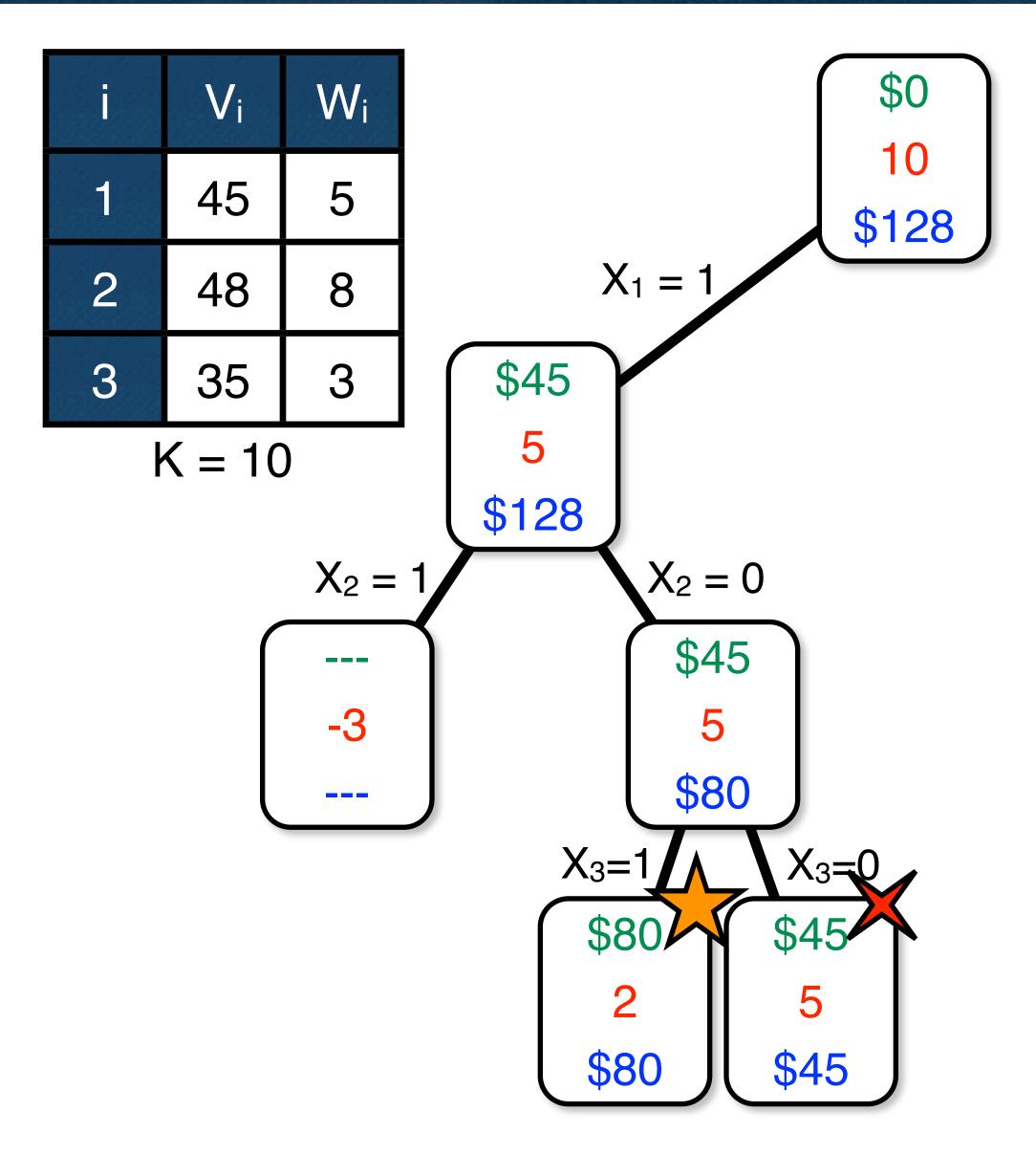


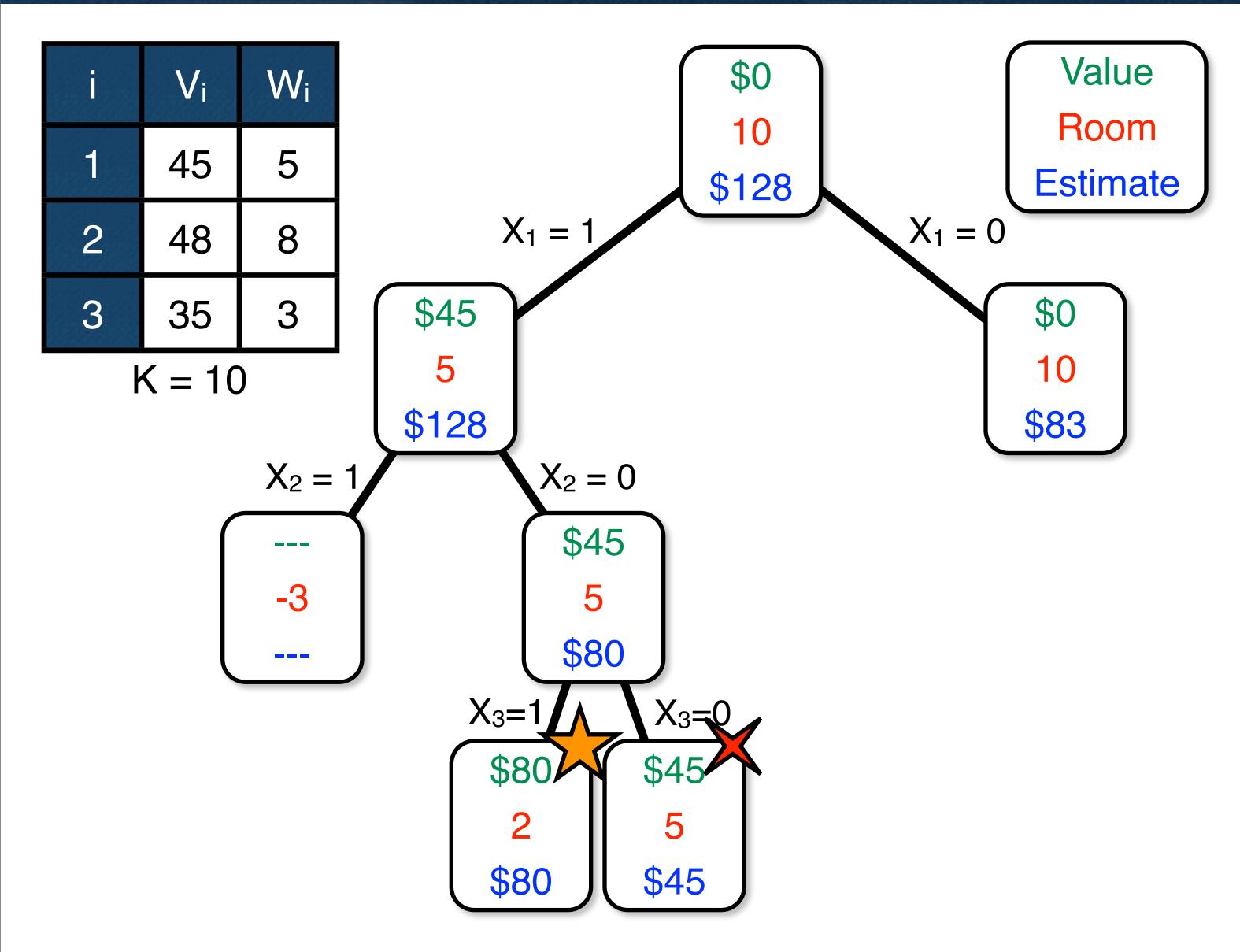


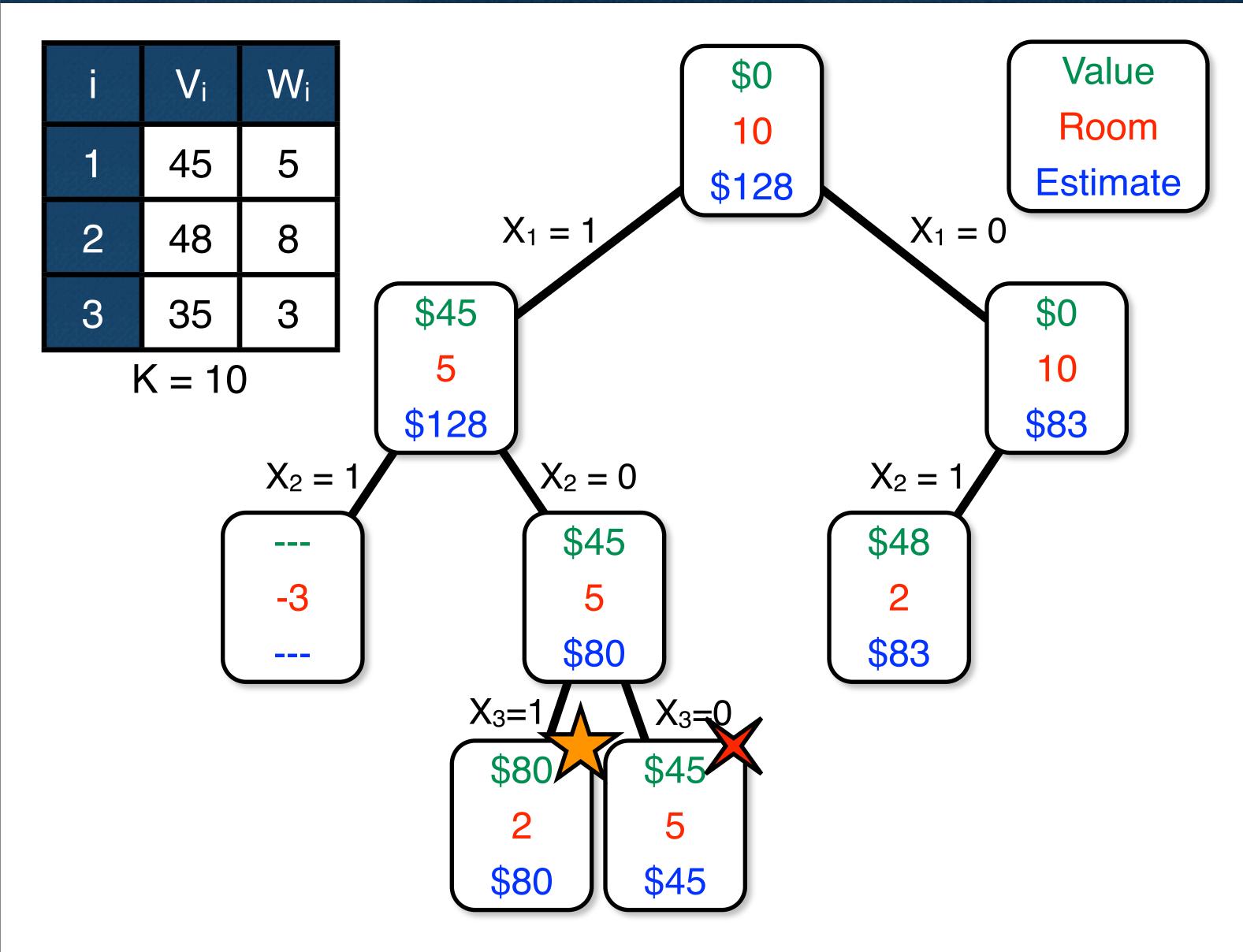


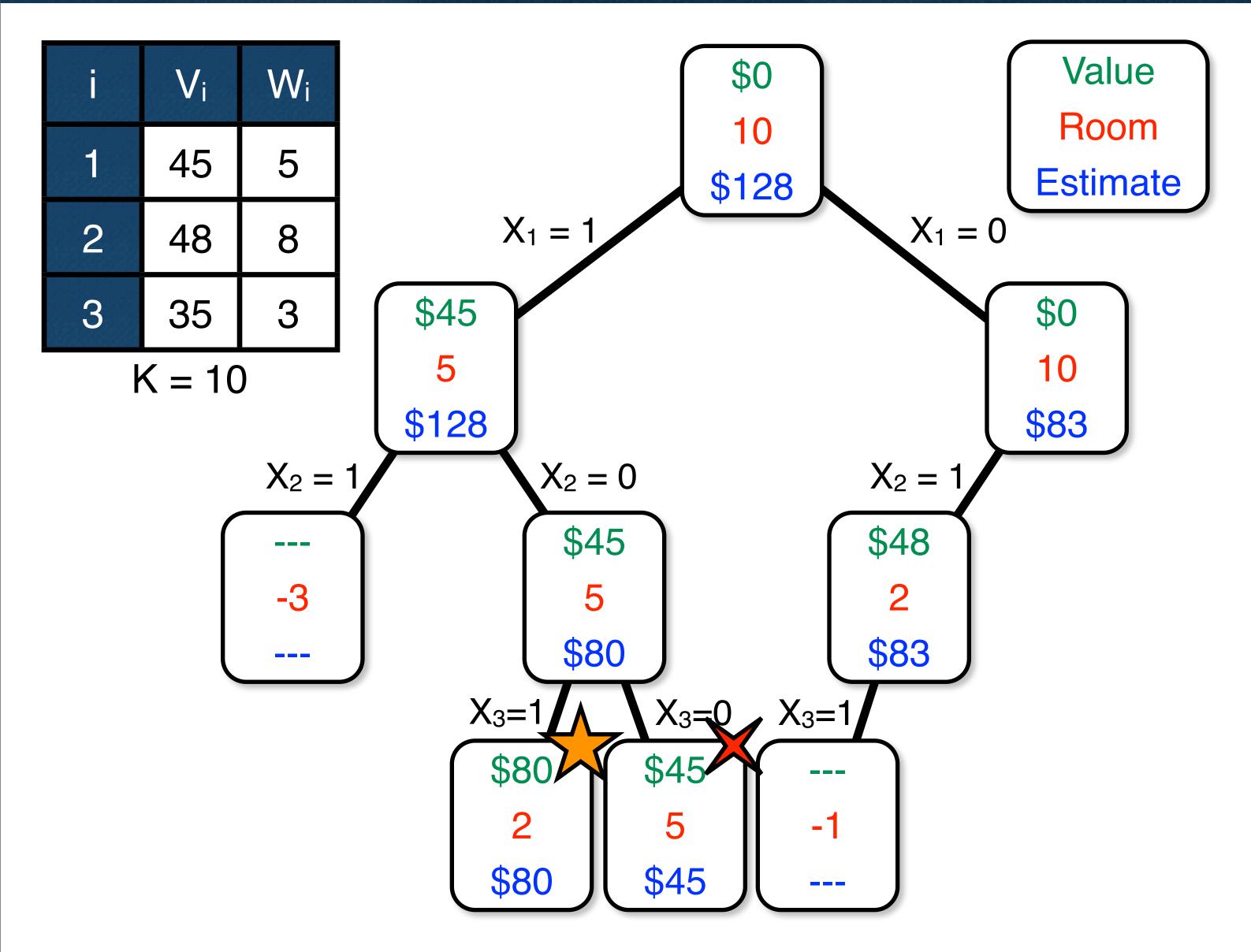


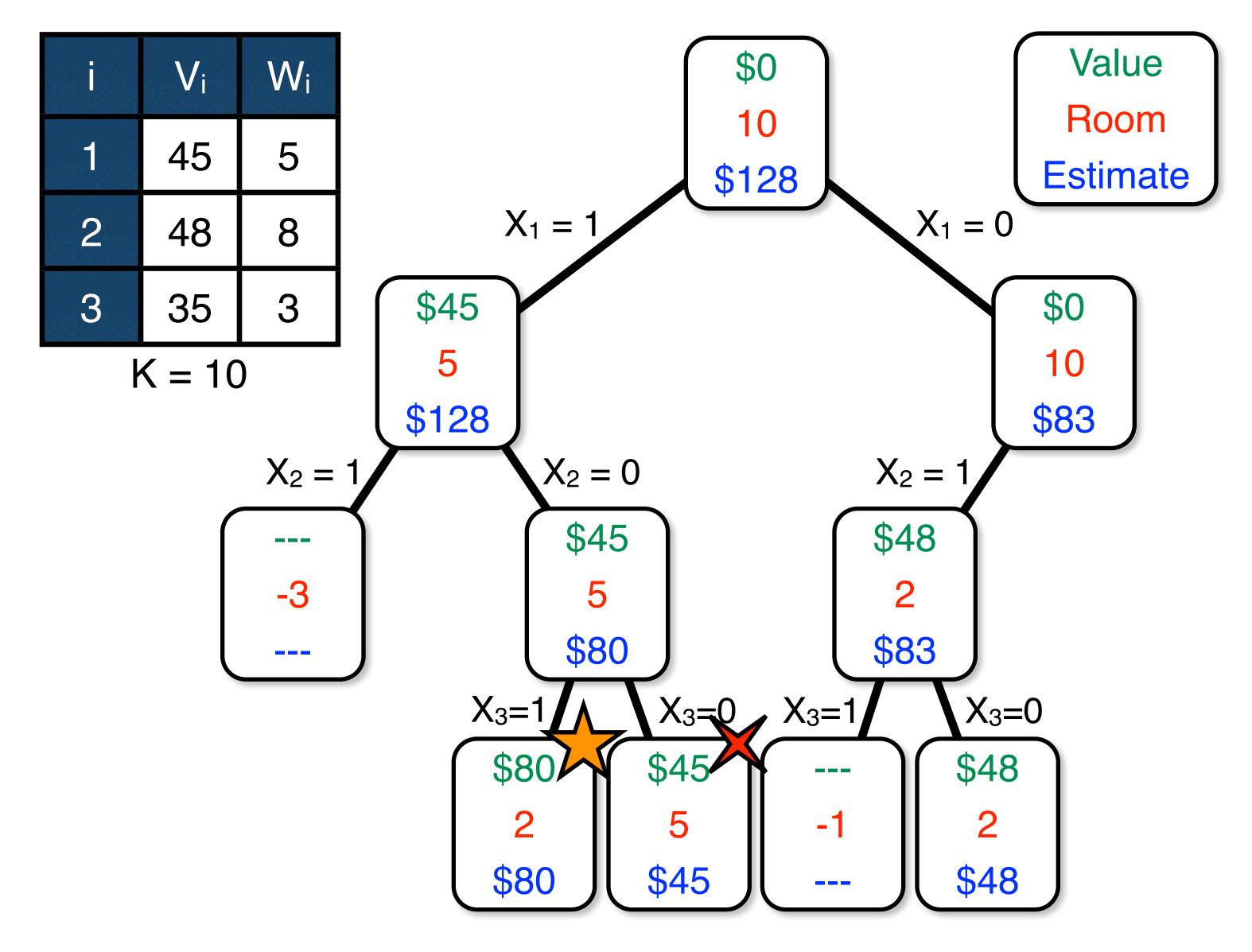


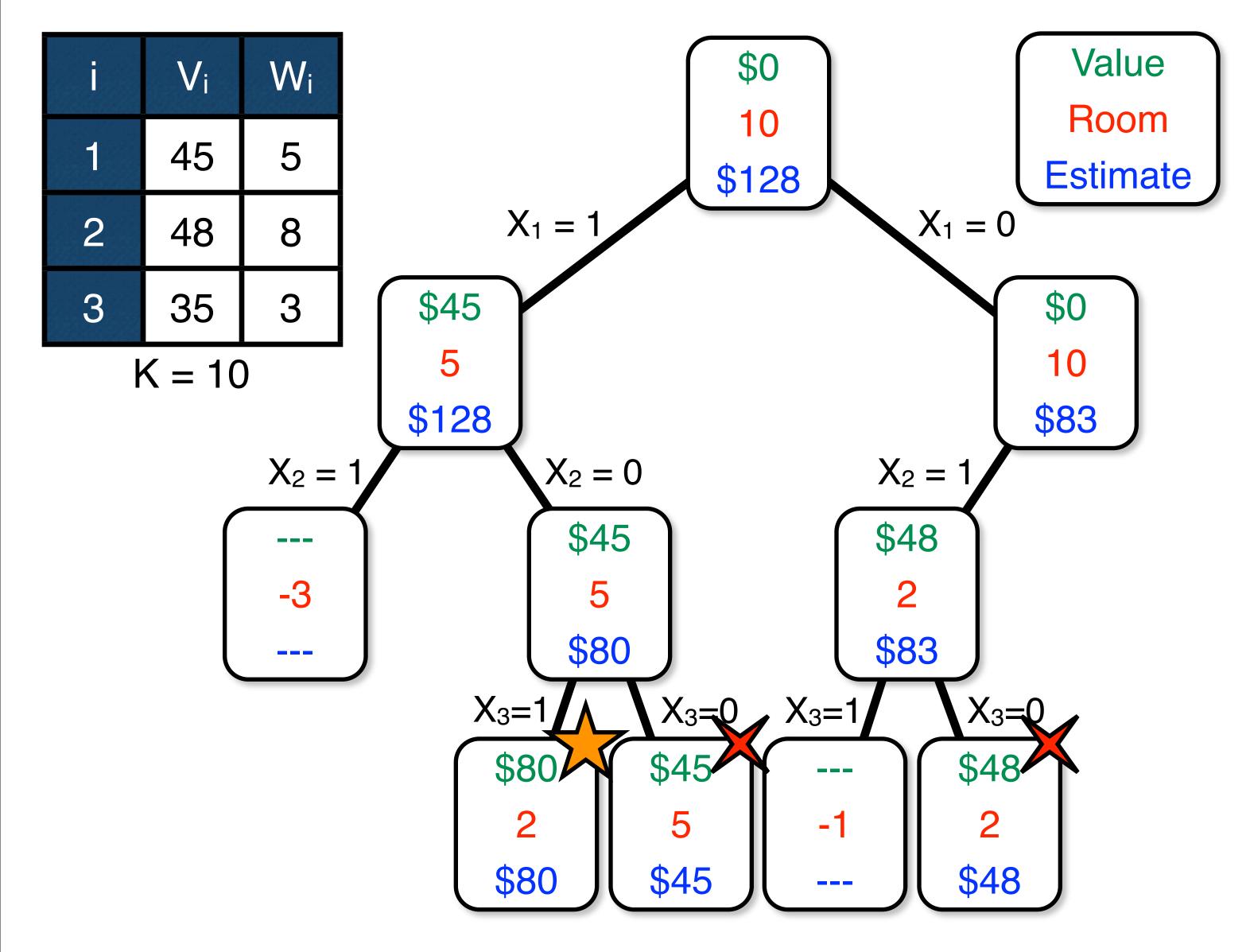


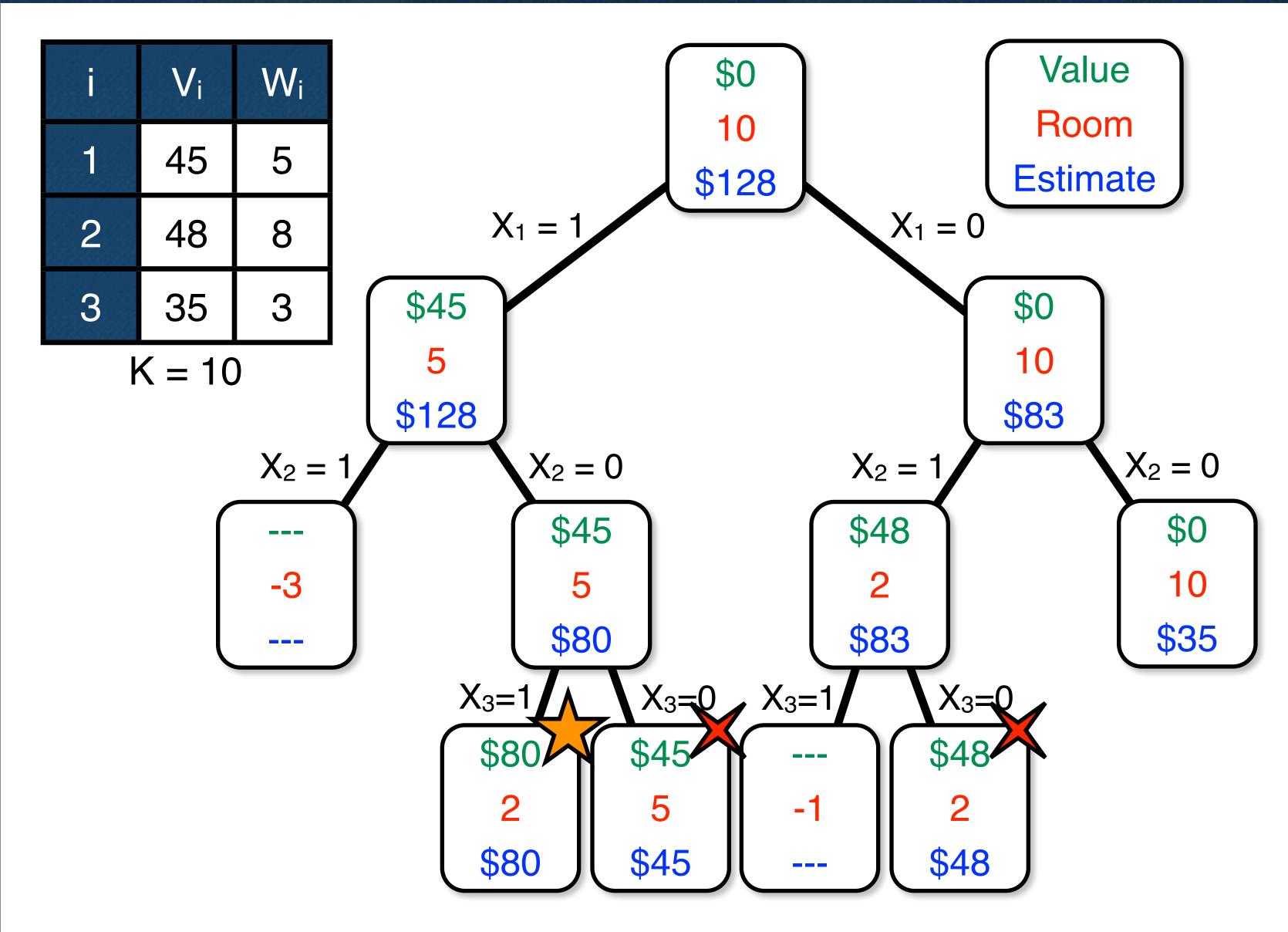


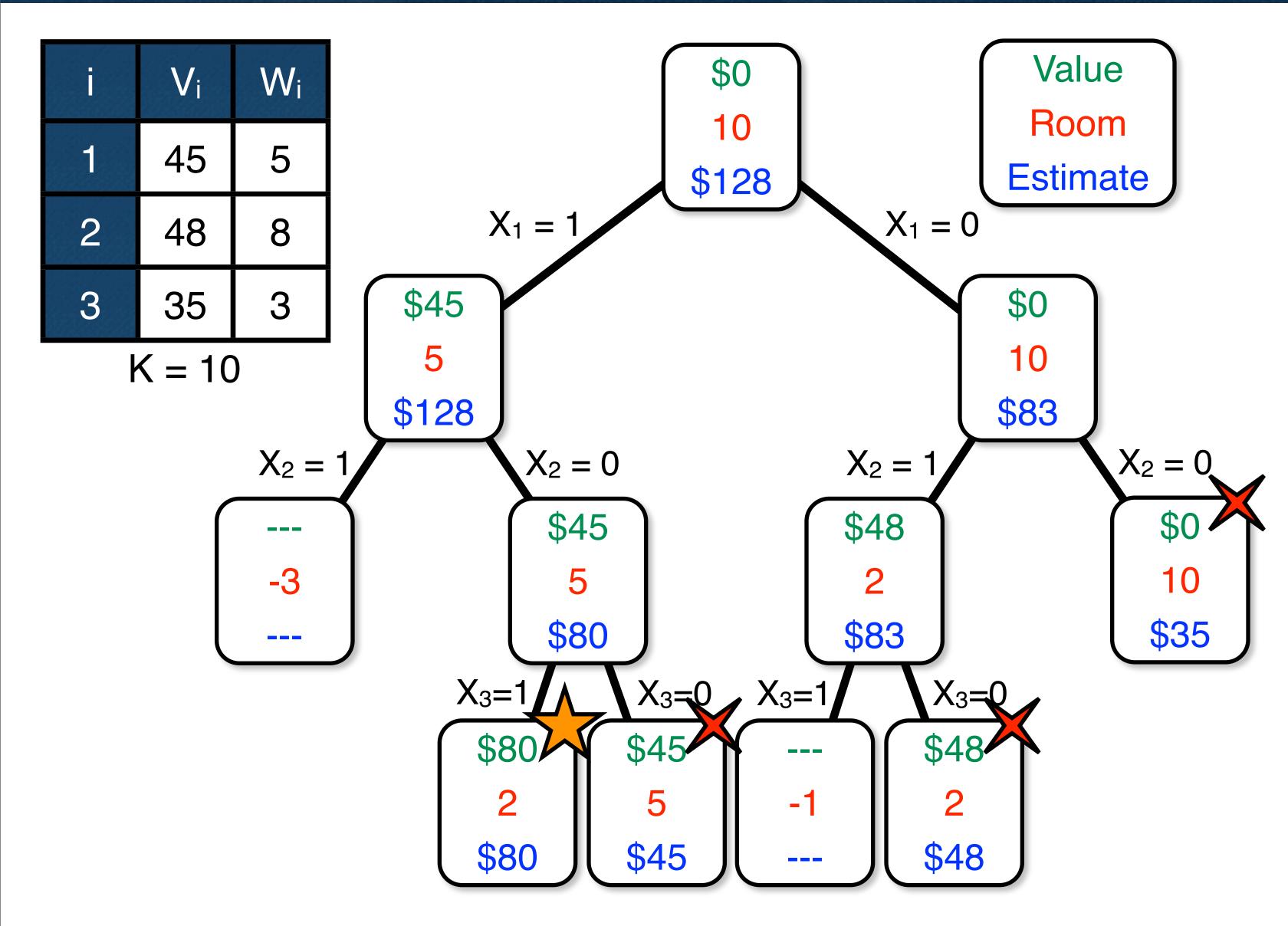












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► Can we relax something else?

What if the items are bars of Belgian chocolate?

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maximize 
$$45x_1 + 48x_2 + 35x_3$$
  
subject to  $5x_1 + 8x_2 + 3x_3 \le 10$   
 $0 \le x_i \le 1$   $(i \in 1..3)$ 

- ► This is called the linear relaxation
  - we will come back to this later in the class
  - we relax the integrality requirement

Can we solve a knapsack when we can take parts of the items?

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- Can we solve a knapsack when we can take parts of the items?
  - order the items by decreasing value of V<sub>i</sub>/W<sub>i</sub>

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- Can we solve a knapsack when we can take parts of the items?
  - order the items by decreasing value of V<sub>i</sub>/W<sub>i</sub>
  - "most value per kilo"

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- ► How to solve the relaxation now?
  - -select the items while the capacity is not exhausted
  - -select a fraction of the last item

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- ► In this example,
  - $V_1/W_1 = 9$ ,  $V_2/W_2 = 6$ ,  $V_3/W_3 = 11.7$
  - select items 3 and 1
  - select 1/4 of item 2
  - estimation: 92

► Why is correct?

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$$\det x_i = \frac{y_i}{v_i}$$

Why is correct?

maximize subject to

$$\sum_{i\in 1...j} y_i$$

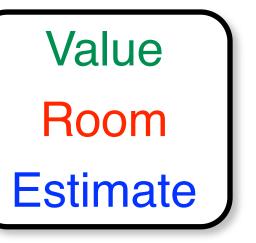
$$\sum_{i \in 1...j} \frac{w_i}{v_i} y_i \le K$$

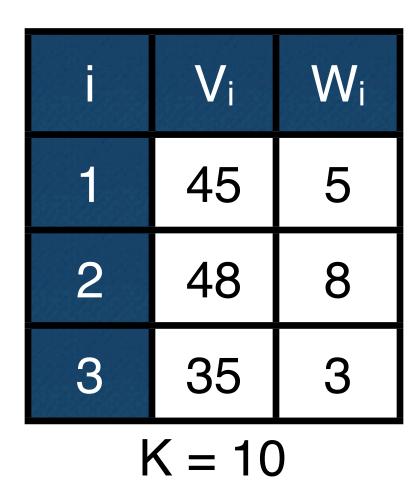
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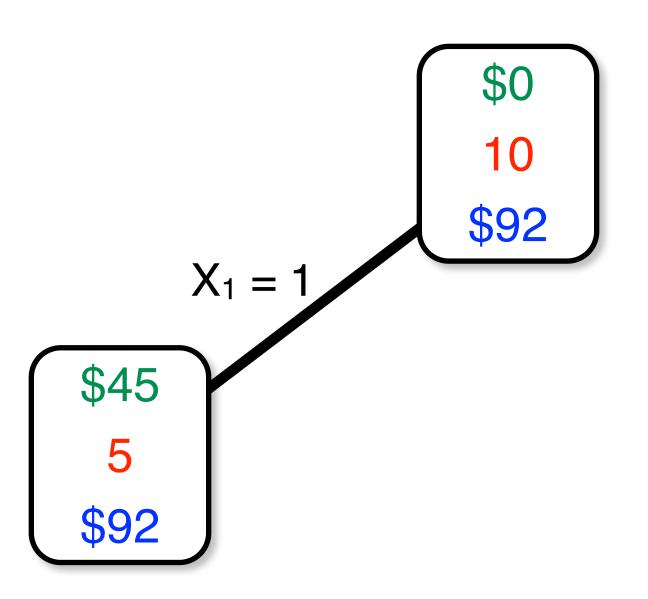
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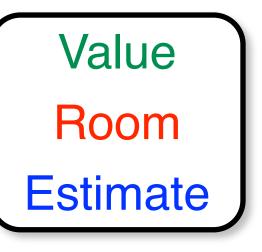
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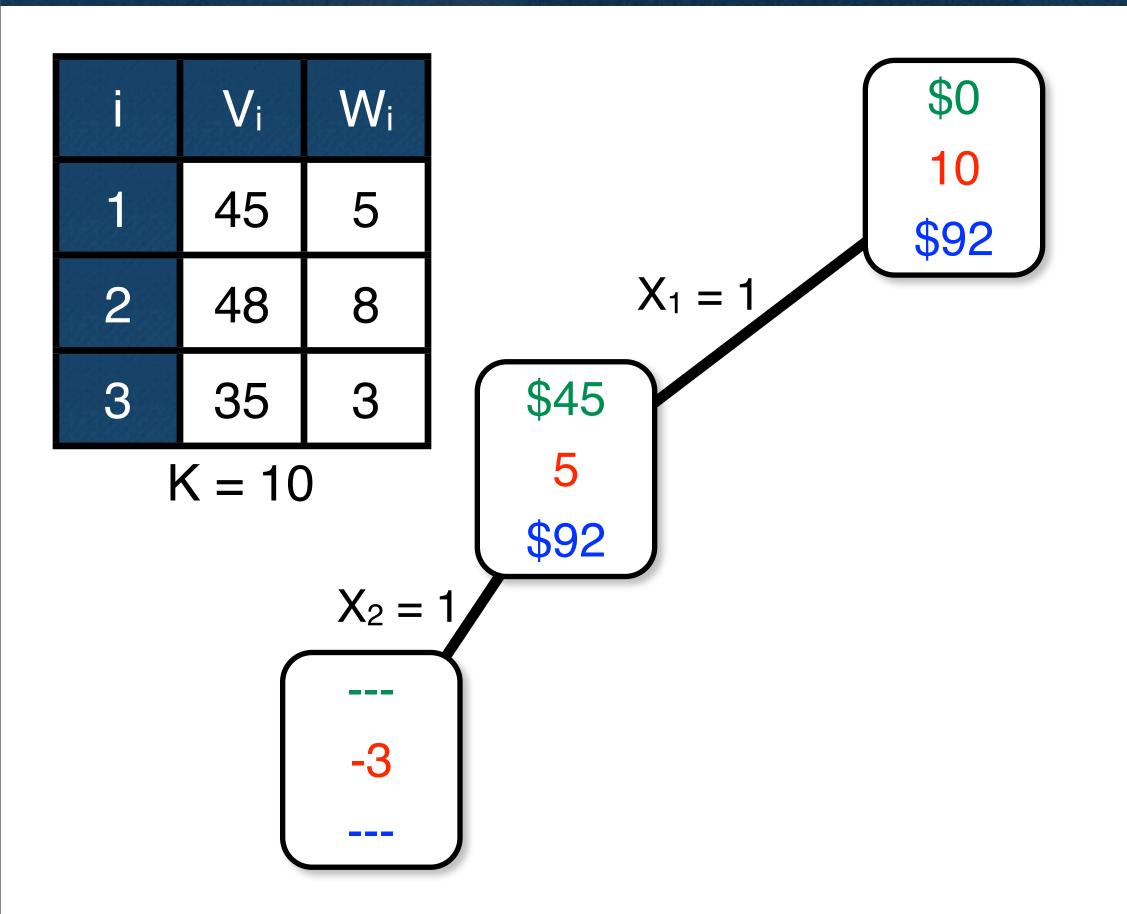


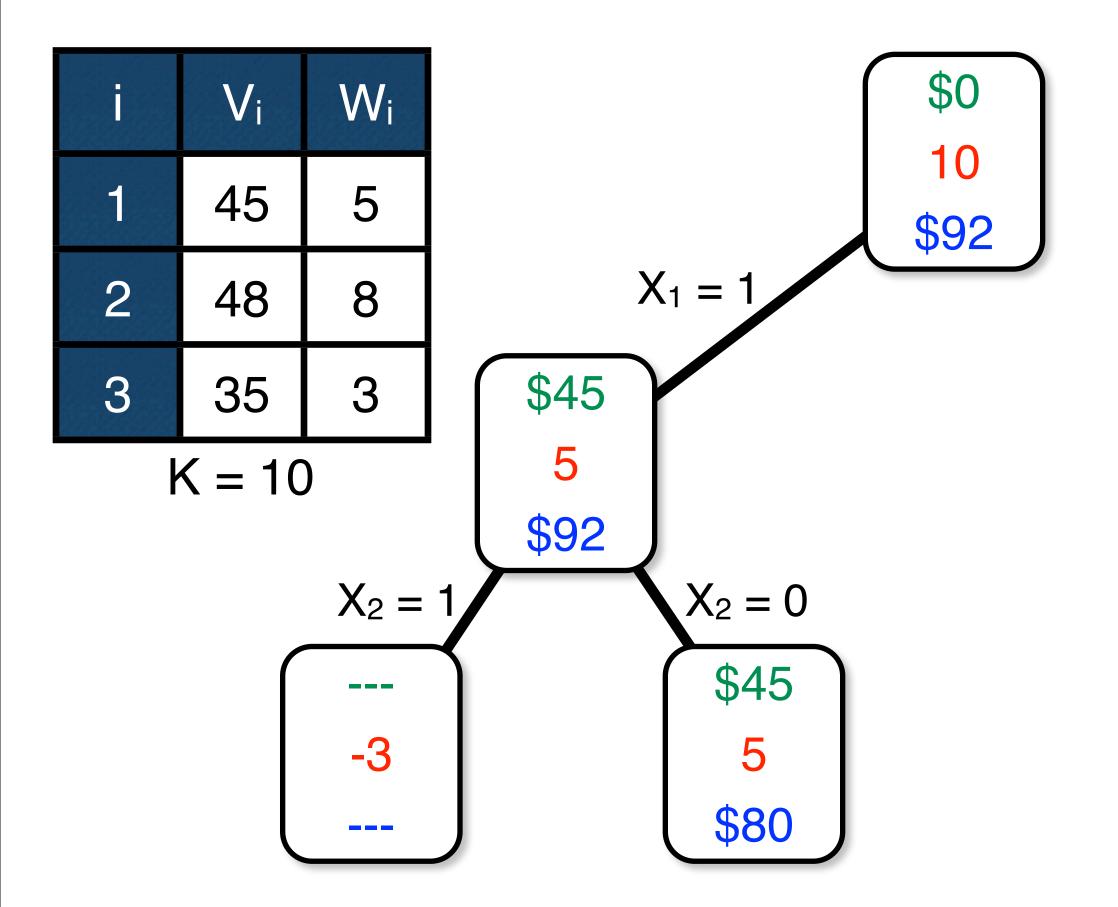


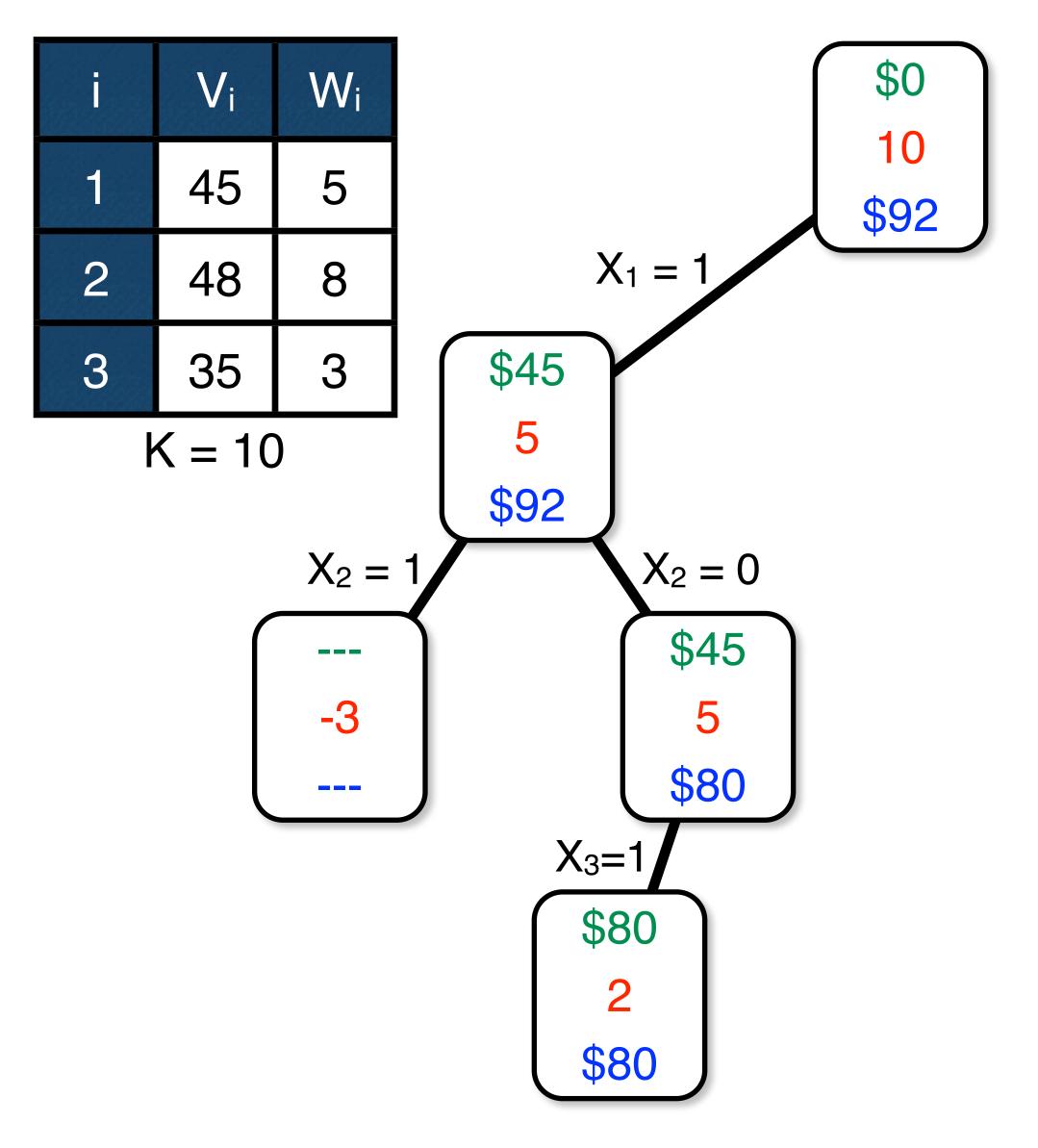


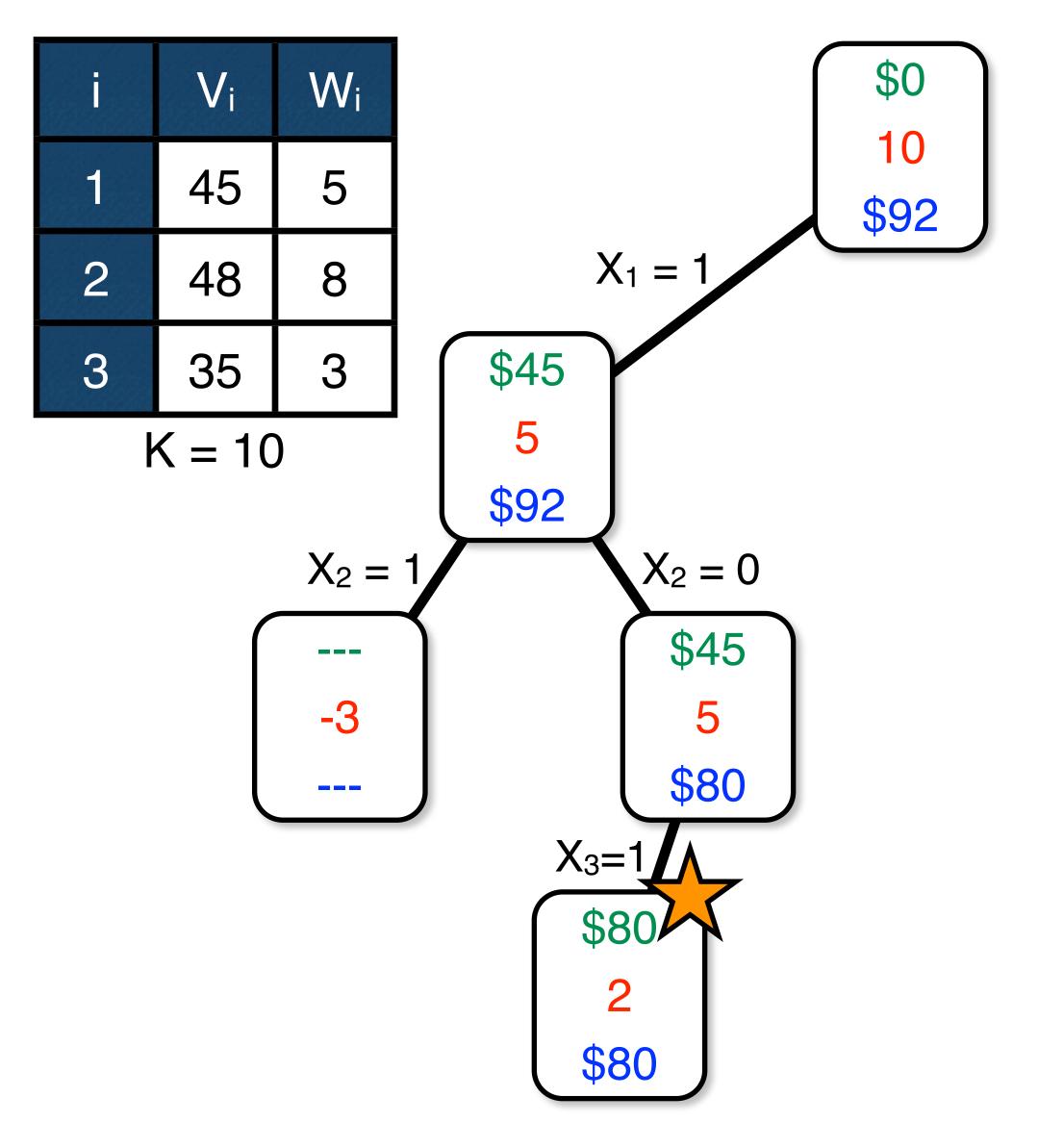


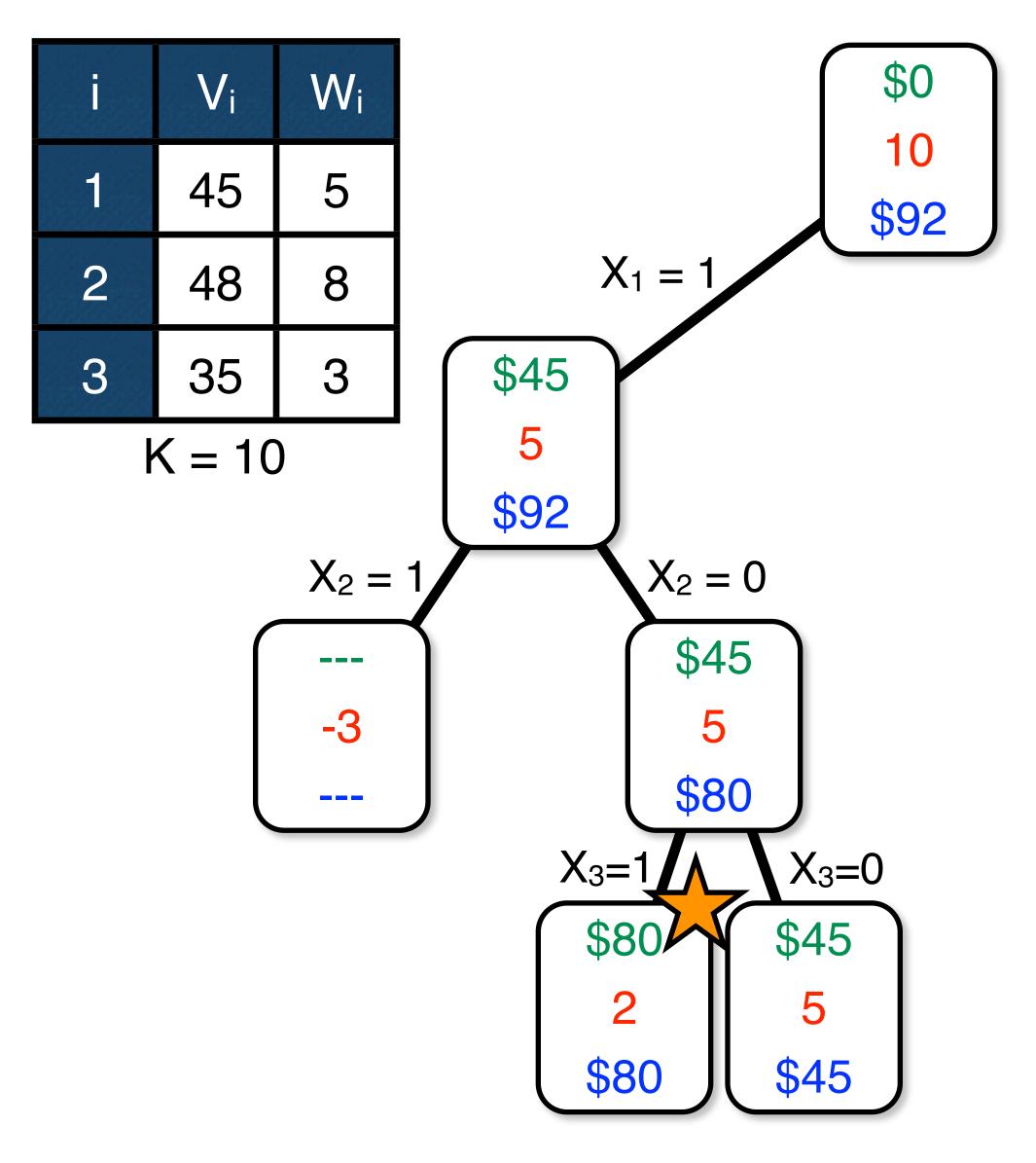


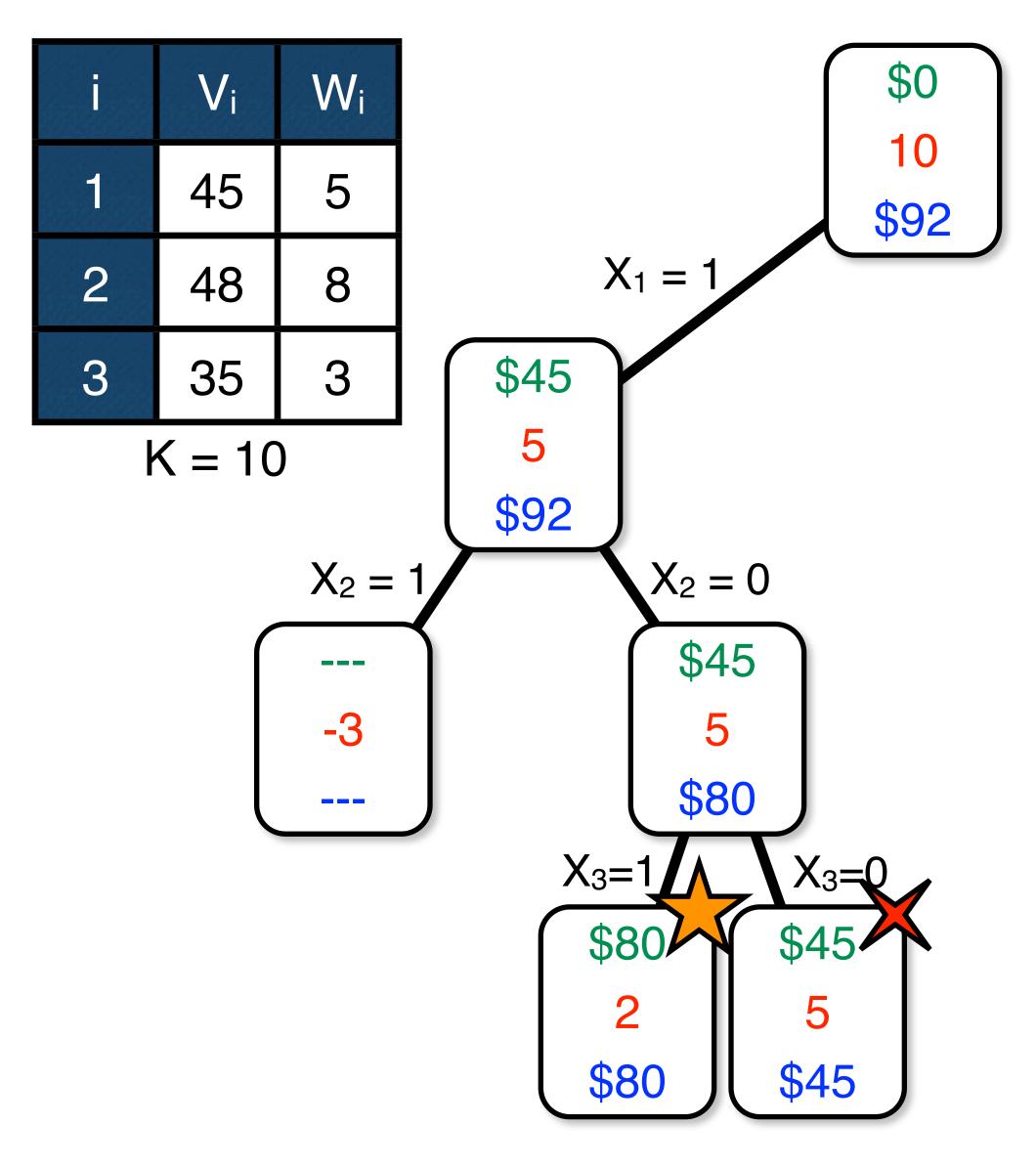


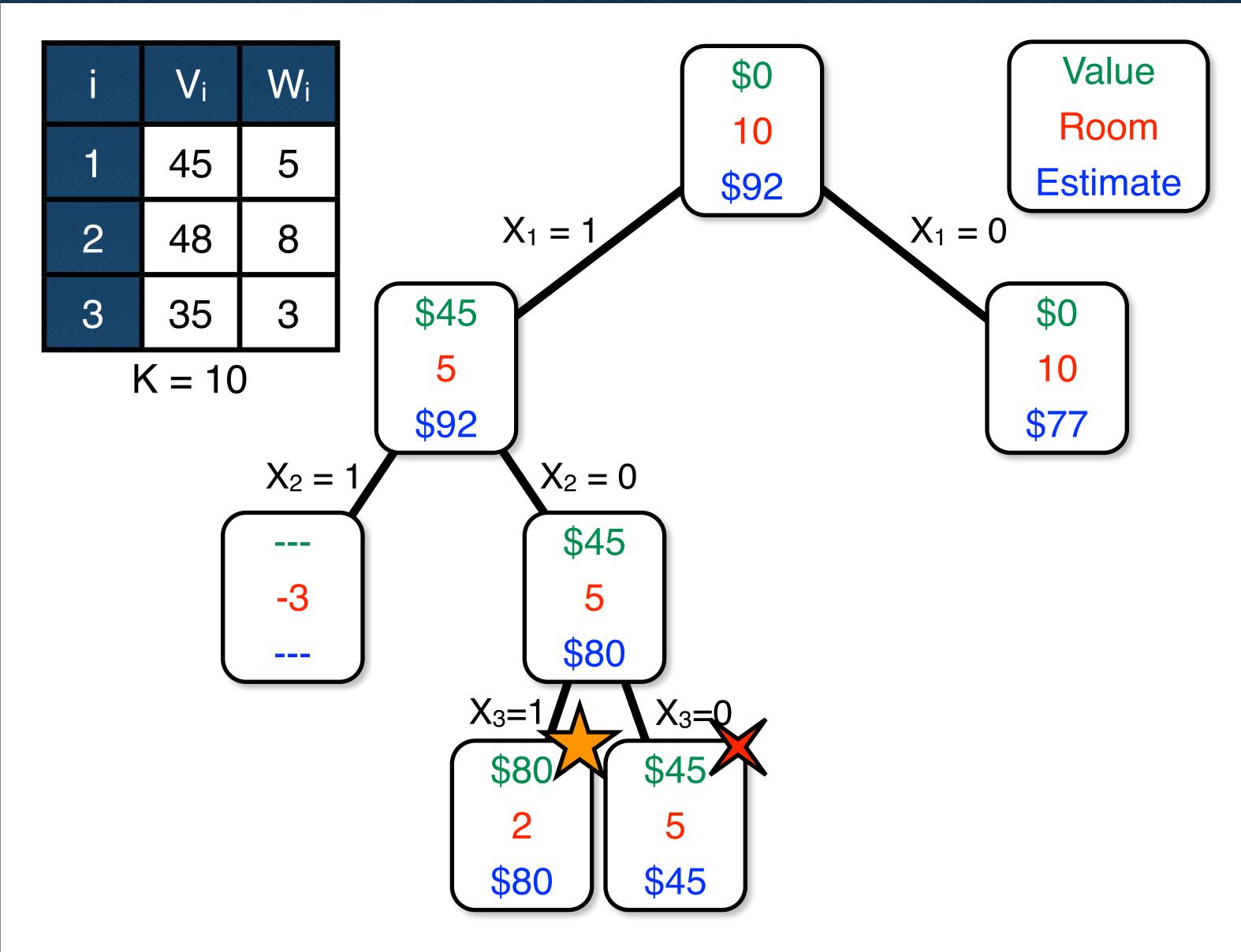


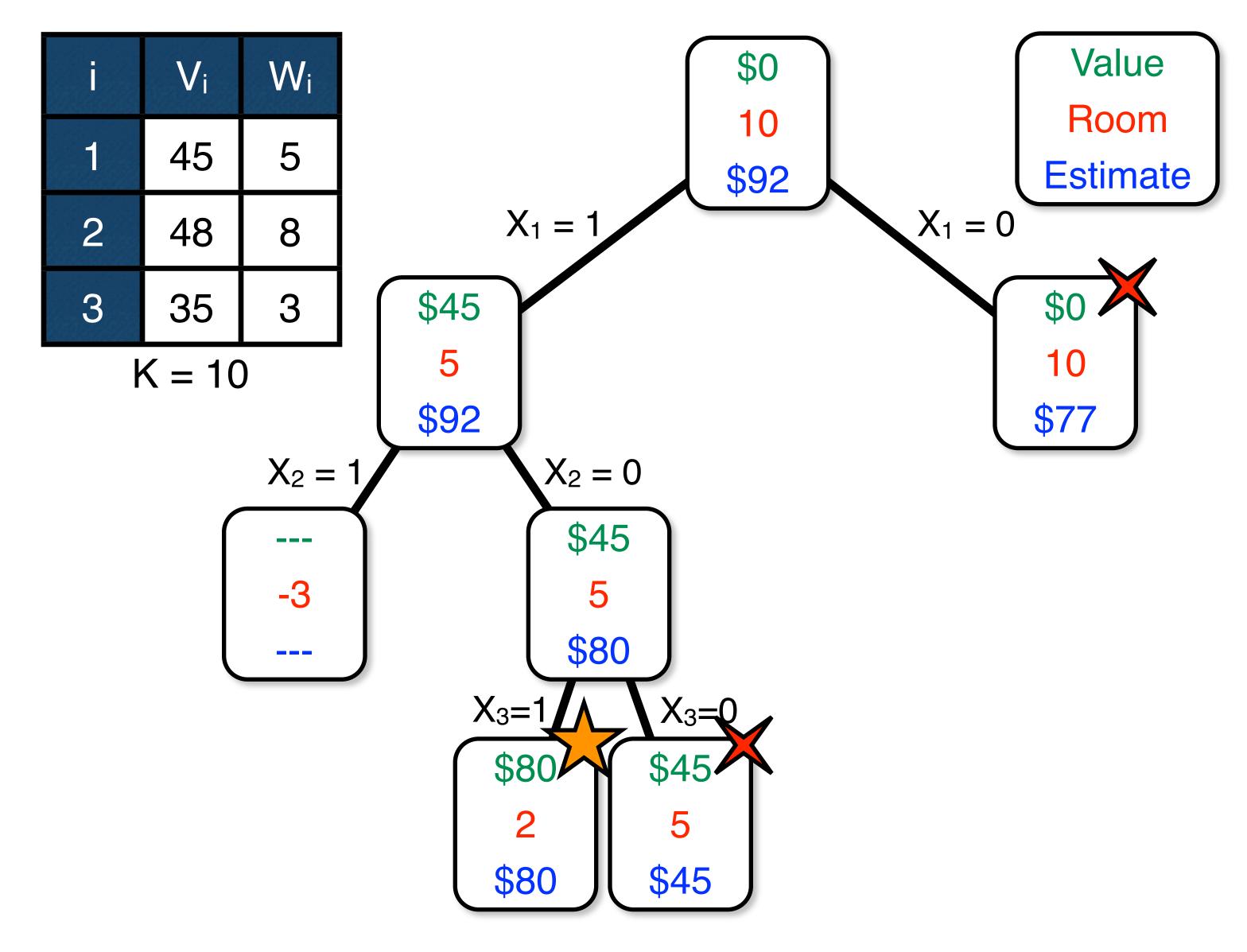












- Search Strategies
  - depth-first
  - best-first
  - -many others

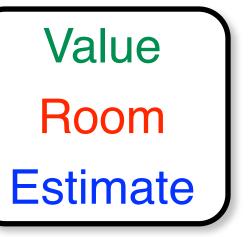
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  - prunes when a node estimation is worse than the best found solution
  - -memory efficient

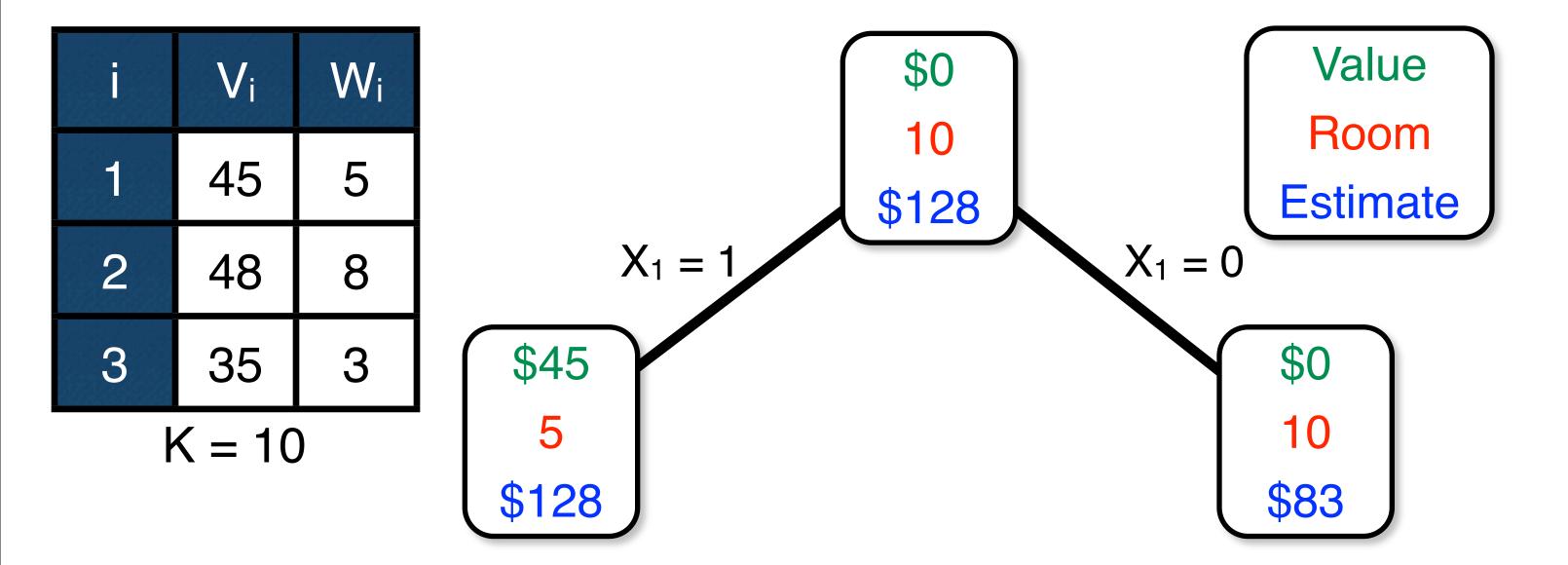
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- ► Best-First
  - select the node with the best estimation

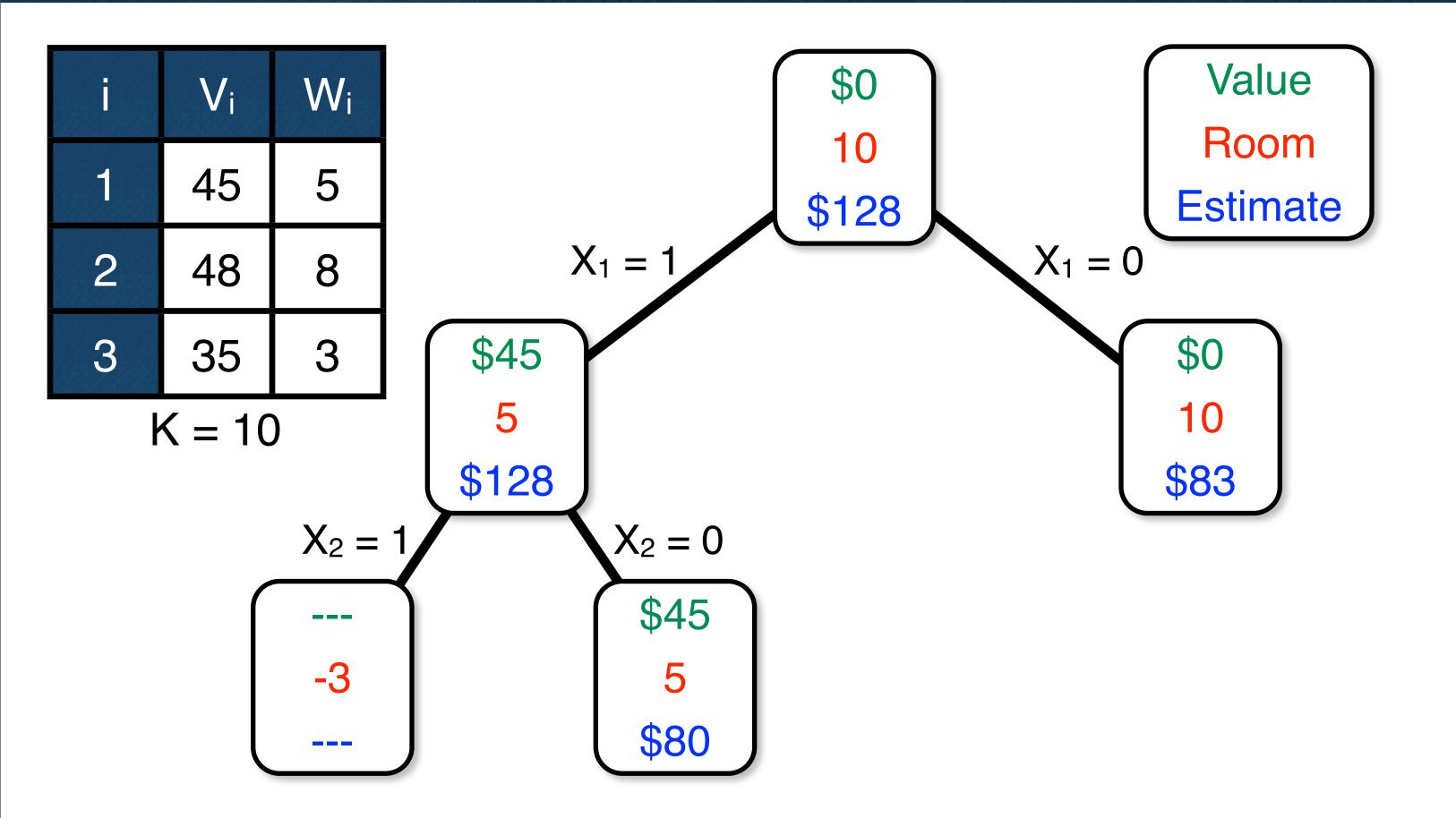
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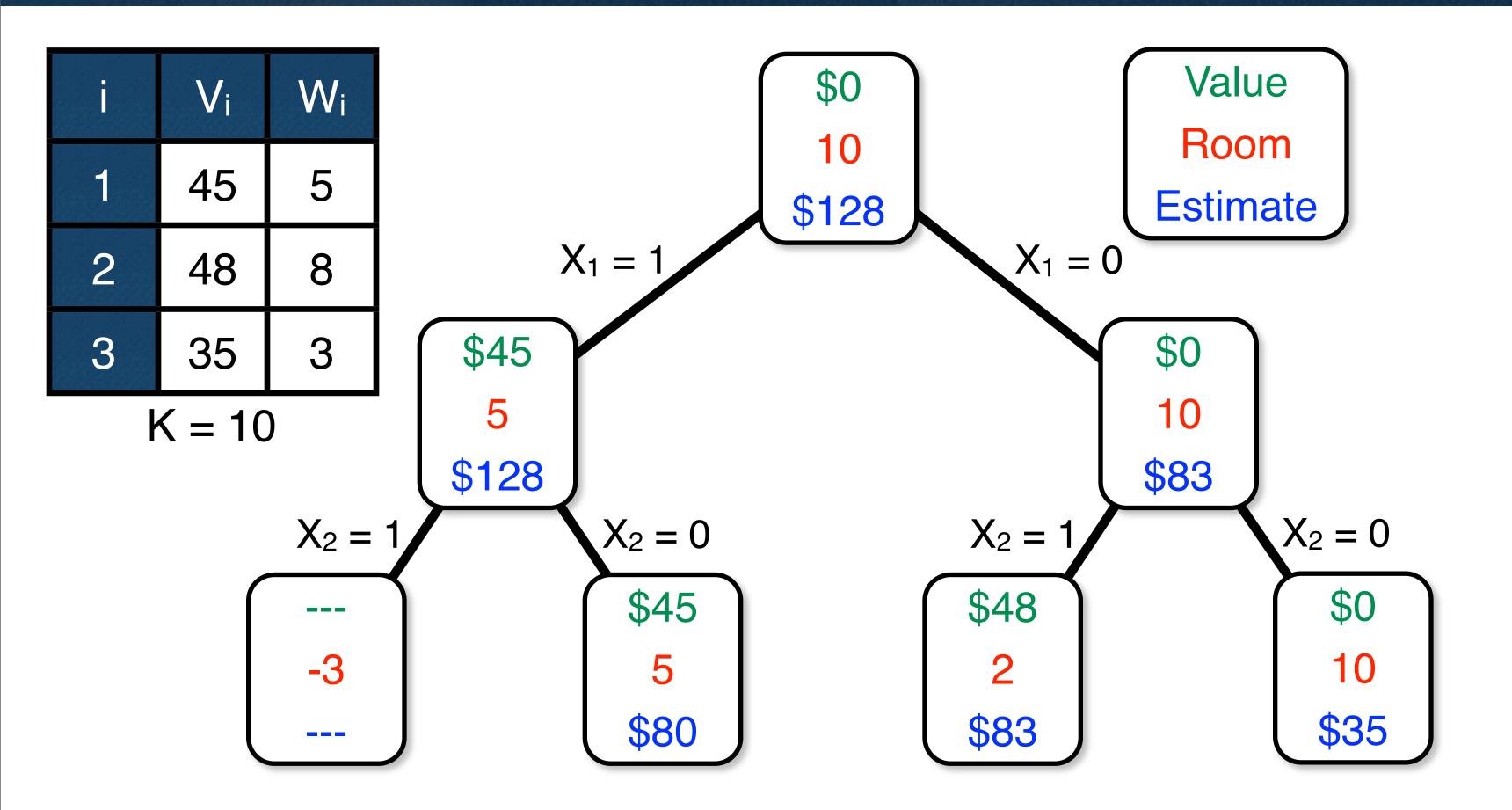
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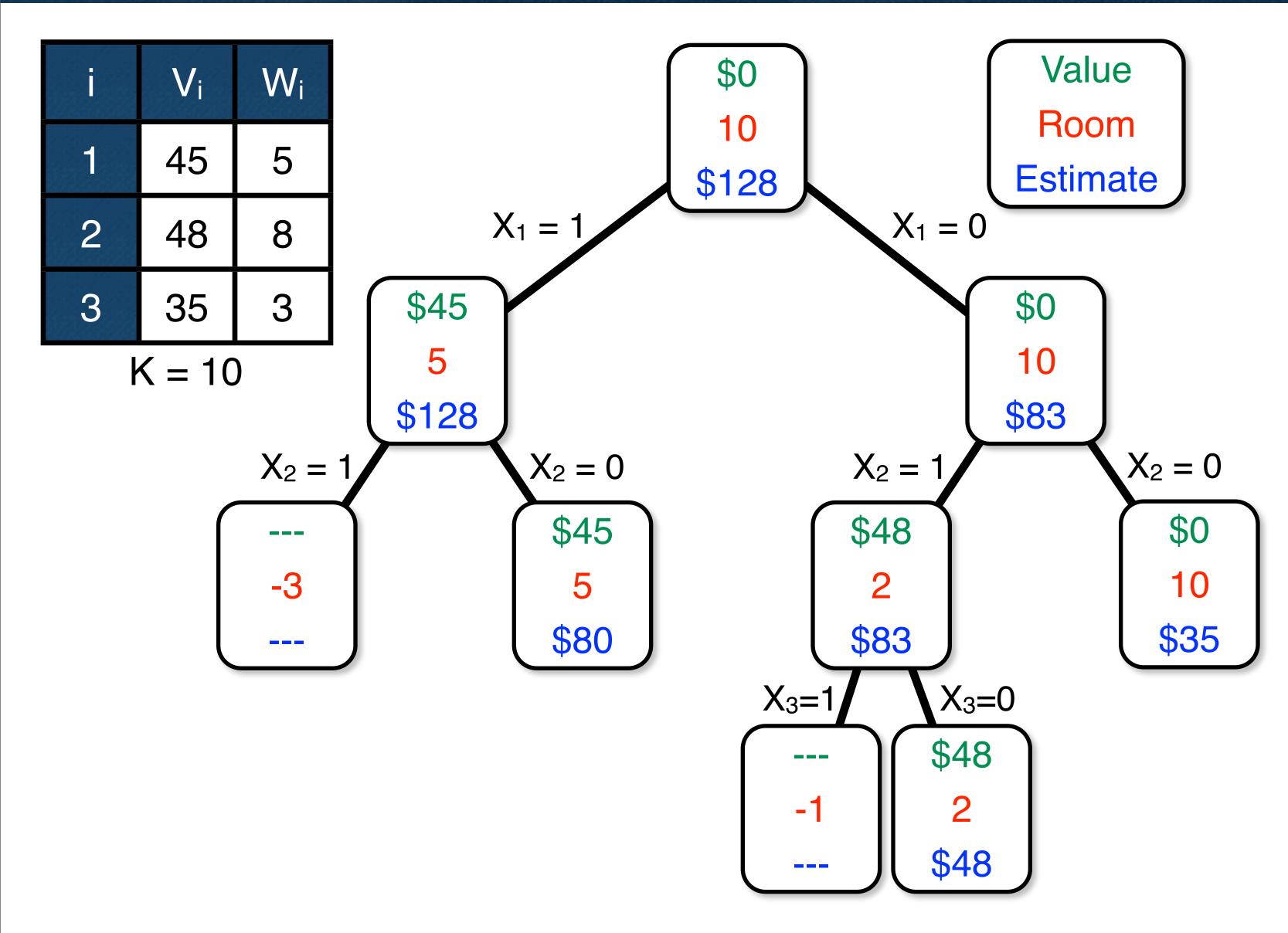


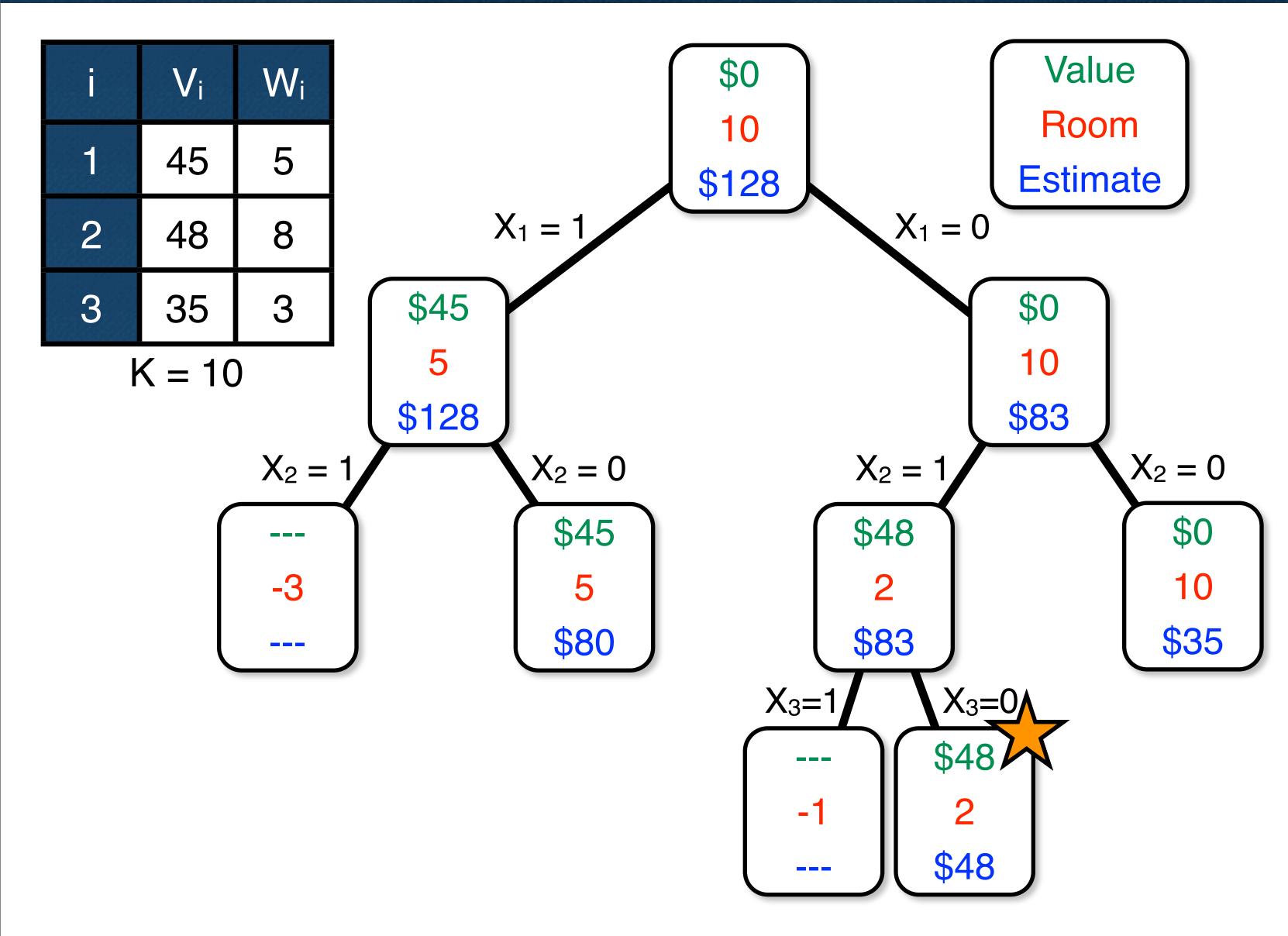


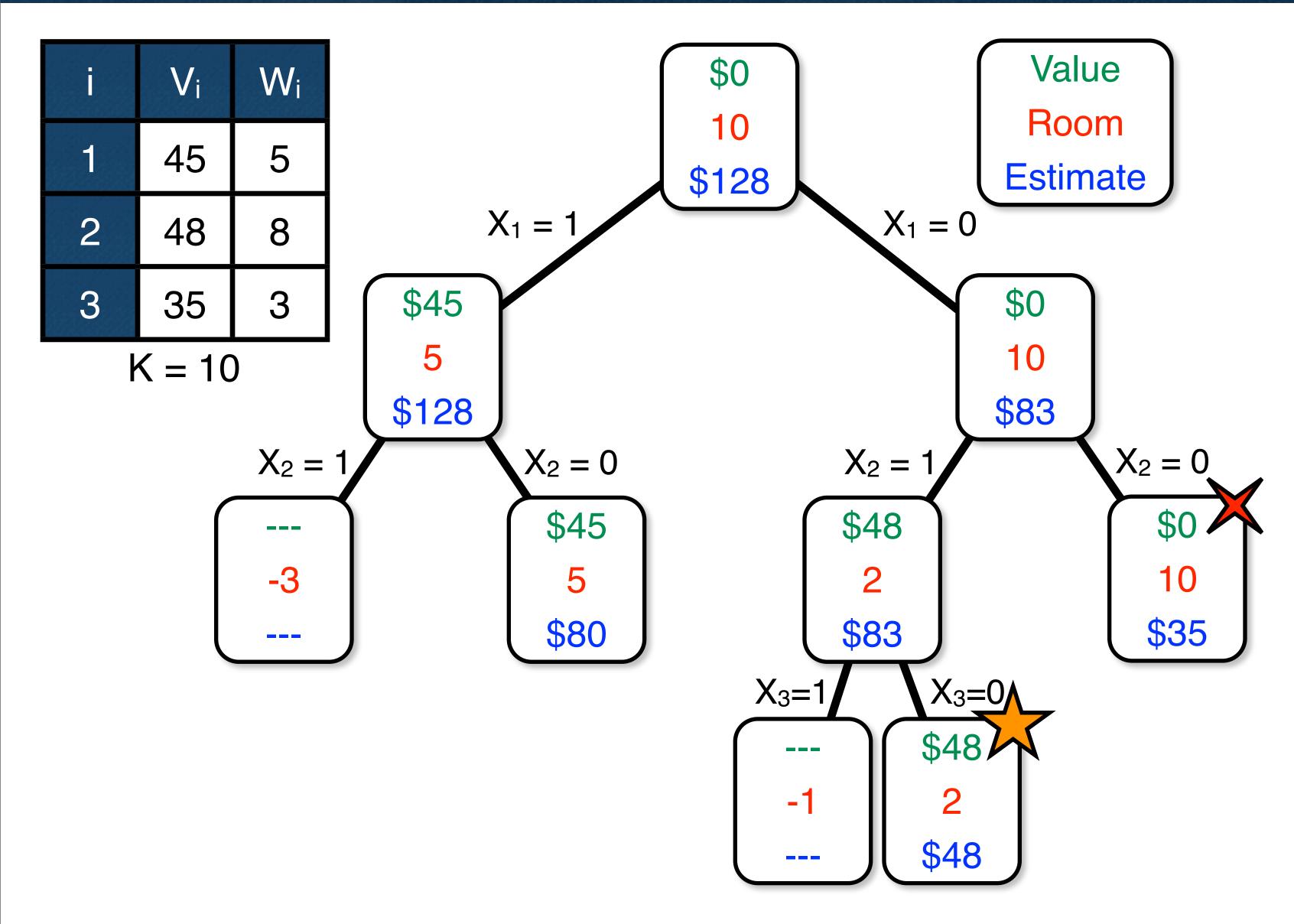


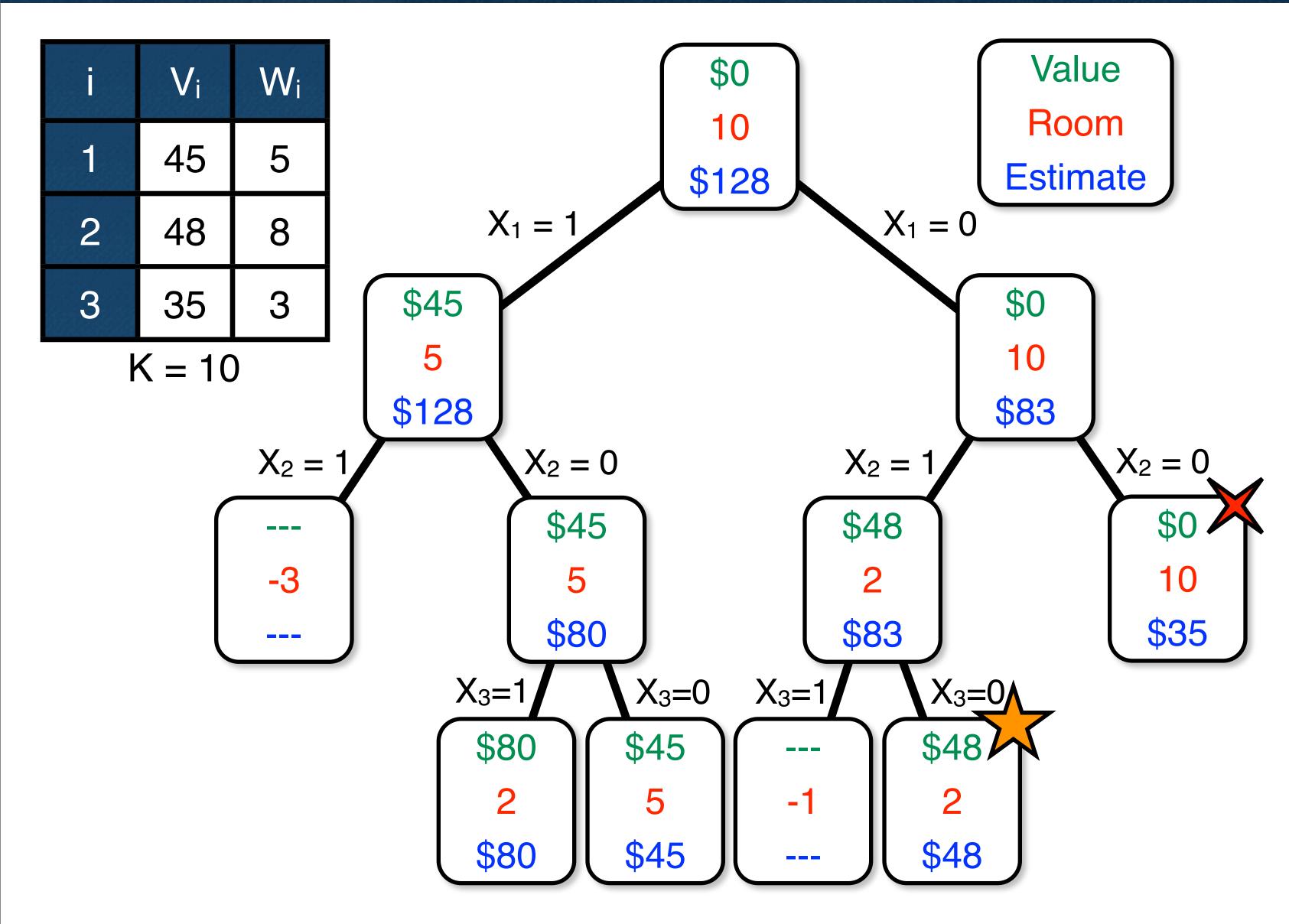


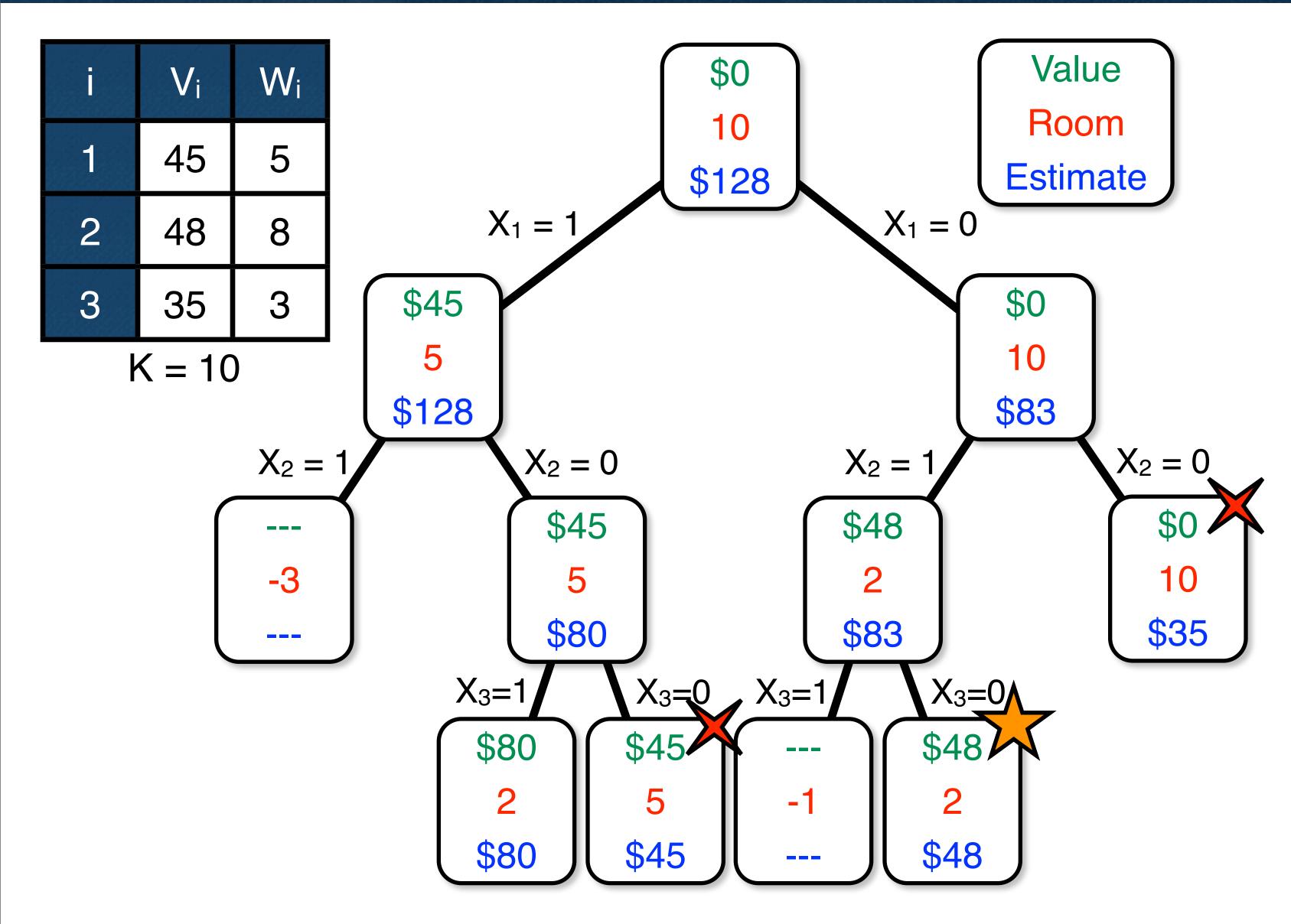


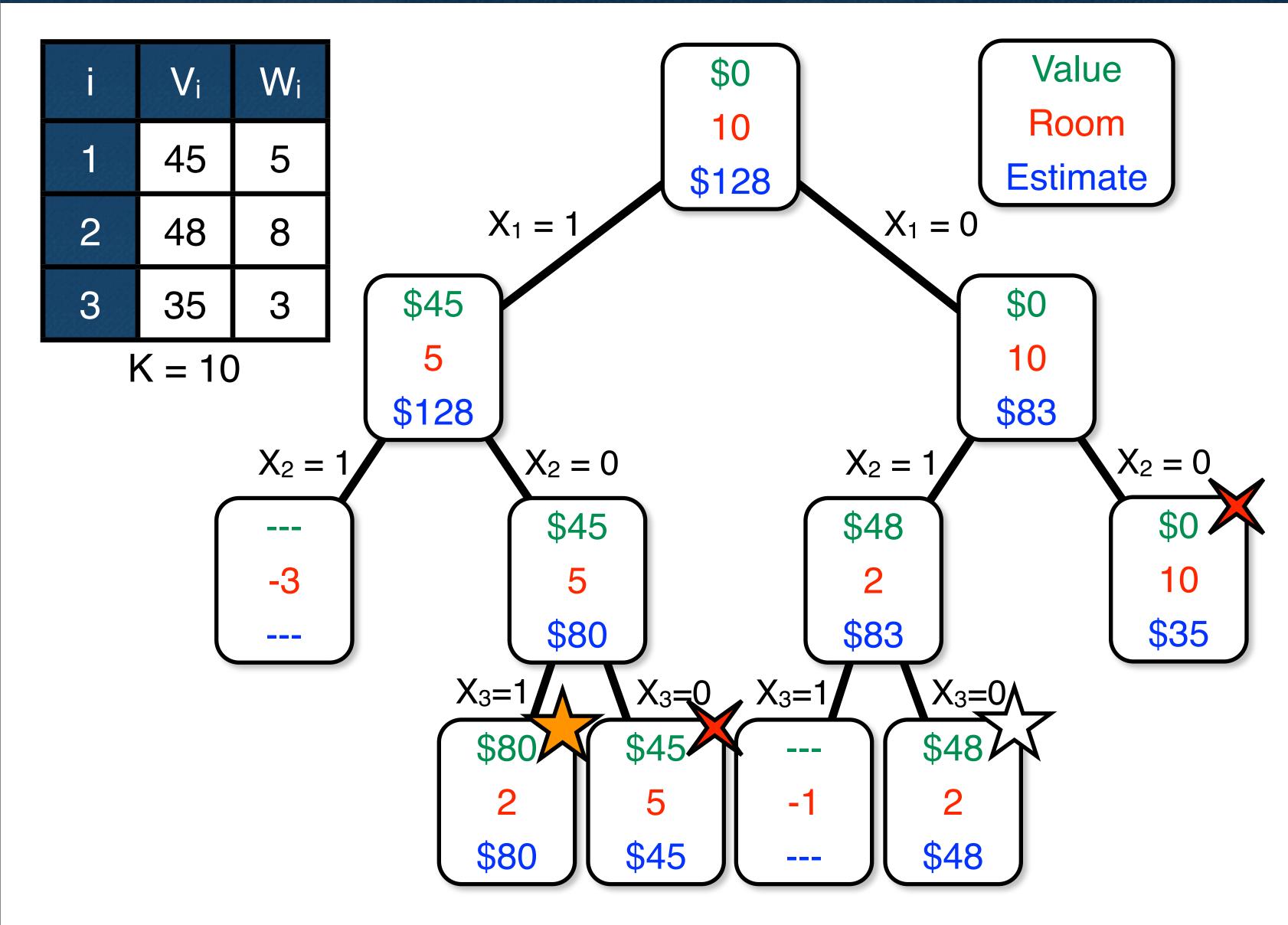












- Best-First
  - select the node with the best estimation
  - -when does it prune?
    - when all the nodes are worse than a found solution
  - is it memory efficient?
    - exaggerate!

# Dynamic Programming and Branch and Bound

Which one is better?

# Dynamic Programming and Branch and Bound

Which one is better?

► Can you combine the two?

## Until Next Time