

# Discovering the Foundations of Expansion and Spectral Graph Theory

(Fernando Granha Jeronimo)

Under heavy construction

Last update: 09/01/24

Inspired by Babai's approach  
and the Hungarian school

## Van Gogh prize

a symbolic prize for any student  
in the course that solves an  
important open problem

(like van Gogh you are not going to receive anything)  
other than have done something amazing

**Warning:** These open problems can be challenging

## Spectral Lens

Let  $G = (V, E)$  be a  $d$ -regular graph on  $n$  vertices.

Let  $A$  be its adjacency matrix, i.e.,

$$A \in \mathbb{R}^{n \times n}, \quad A_{u,v} = \mathbf{1}_{\{u,v \in E\}}.$$

Study the spectral theorem.

Let  $\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_n$  be the eigenvalues of  $A$  with corresponding orthonormal eigenvectors  $\varphi_1, \dots, \varphi_n \in \mathbb{R}^n$ .

$$(A \varphi_i = \lambda_i \varphi_i)$$

1] Prove that  $d$  is an eigenvalue of  $A$ .

$$\text{Def } \langle x, x \rangle = \sum_{i=1}^n \bar{x}_i x_i$$

$$\text{Ex } x \in \mathbb{R}^n \text{ and } e_1, \dots, e_n \in \mathbb{R}^n$$

We can write ONB (orthonormal basis)

$$x = \sum_{i=1}^n \alpha_i e_i \quad \text{with } \alpha_i = \langle e_i, x \rangle.$$

$$\text{Ex } \langle x, x \rangle = \sum_{i=1}^n \alpha_i^2 \quad [\text{Parserval}]$$

$$\text{Ex: } \langle x, Ax \rangle = \sum_{i=1}^n \alpha_i^2 \lambda_i$$

Ex A sym  $\Rightarrow$  A has real eigenvalues

Def Rayleigh quotient  $\frac{\langle x, Ax \rangle}{\langle x, x \rangle}$  (for  $x \neq 0$ )

$$\text{Ex } \text{Tr}(A) = \sum_{i=1}^n \lambda_i$$

$$\text{Ex } \text{Tr}(A^k) = \sum_{i=1}^n \lambda_i^k$$

$$\text{Ex } \text{Tr}(A^2) = 2|E|$$

Ex If G is simple with  $\deg \geq 1$

$$\downarrow \\ \lambda_n < 0$$

$$\text{Ex Prove that } \frac{1}{n} \sum_{i=1}^n \lambda_i^2 = d$$

Def  $J_n = \begin{pmatrix} 1 & \dots & 1 \\ \vdots & \ddots & \vdots \\ 1 & \dots & 1 \end{pmatrix}$  all ones  $n \times n$  matrix

Ex Compute eigenvalues of  $J_n/n$

Def  $K_n$  complete graph on  $n$  vertices

Def  $K_{a,b}$  complete bipartite graph  
 $G = (V = L \cup R, E)$  with  $|L| = a, |R| = b$

Ex Compute the spectrum of  $K_n$

Ex Compute the spectrum of  $K_{n,n}$

Ex Compute " " " of  $K_{1,d}$

Ex " " " " of  $K_{a,b}$

Ex  $\lambda_1 \geq \max_{i \geq 2} |\lambda_i|$

Ex  $G$  bipartite  $\iff \begin{matrix} \text{Spec}(G) \\ \parallel \\ -\text{Spec}(G) \end{matrix}$

## Some Notions of Expansion

Def  $\partial(S) = E(S, \bar{S})$  "Edge boundary"

Def  $\Phi(S) = \frac{|\partial(S)|}{|S|}$  "Conductance"

Def  $\Phi(G) = \min_{\emptyset \neq S \subseteq V, |S| \leq \frac{n}{2}} \Phi(S)$  "Cheeger's constant"

Def:  $\lambda = \max \{| \lambda_2 |, | \lambda_n |\}$

two-sided spectral expansion

Def  $\lambda = \lambda_2$  one-sided spectral expansion

Def  $e(S, T) = |\{(s, t) \mid s \in S, t \in T\}|$

Ex Prove that

$$|e(S, T) - \frac{d|S||T|}{n}| \leq \lambda \sqrt{|S||T|}.$$

[Expander Mixing Lemma]

Ex Improve the error bound  $\lambda \sqrt{|S||T|}$ .

Def  $\alpha(G)$  = independence number

Ex Prove that  $\frac{\alpha(G)}{n} \leq \frac{-\lambda_n}{d - \lambda_n}$

[Hoffmann's bound]

Ex Prove that  $d_{avg} \leq \lambda_1 \leq \Delta(G)$

Def  $\chi(G)$  is the chromatic number

Ex Prove that  $\chi(G) \leq \lambda_1 + 1$

[Witj's bound]

## Mixing Bounds

Def  $\vec{1}$  is the all one vector

Def  $R = \frac{1}{d} A$  is the random walk matrix

$$E_x R \vec{1}_n = \vec{1}_n$$

$E_x$  Prove  $\|R^p - \vec{1}_n\|_1 \leq \left(\frac{1}{d}\right)^p \sqrt{n}$  for any distribution  $p$ .

[Mixing bound]  
in  $l_1$

Study the Perron-Frobenius theorem (useful for understanding more general Markov chains)

Eigenvalues are an optimization problem

Let  $V_k = \text{span}\{\varphi_1, \dots, \varphi_k\}$

$W_k = \text{span}\{\varphi_{k+1}, \dots, \varphi_n\}$

Ex  $\lambda_k = \min_{0 \neq x \in V_k} \frac{\langle x, Ax \rangle}{\langle x, x \rangle} = \max_{0 \neq x \in W_k} \frac{\langle x, Ax \rangle}{\langle x, x \rangle}$

Ex Prove the min-max variational theorem

$$\lambda_k = \max_{V \subseteq \mathbb{R}^n} \min_{0 \neq x \in V} \frac{\langle x, Ax \rangle}{\langle x, x \rangle} = \min_{\dim(V)=k} \max_{V \subseteq \mathbb{R}^n} \frac{\langle x, Ax \rangle}{\langle x, x \rangle}$$

$$\dim(V) = k$$

$$\dim(V) = n-k+1$$

[Courant-Fischer-Weyl]

## The Magic of Interlacing

### Ex Eigenvalue Interlacing

Let  $A \in \mathbb{R}^{n \times n}$  be real symmetric matrix

and  $B$  be a  $(n-1) \times (n-1)$  principal submatrix

$$\text{eig}(A) = \{\lambda_1 \geq \dots \geq \lambda_n\}$$

$$\text{eig}(B) = \{\tilde{\lambda}_1 \geq \dots \geq \tilde{\lambda}_{n-1}\}$$

Prove  $\lambda_1 \geq \tilde{\lambda}_1 \geq \lambda_2 \geq \dots \geq \tilde{\lambda}_{n-1} \geq \lambda_n$

*Hint: use min-max theorem for eigenvalues*

Extend to  $r \times r$  principal submatrix  $B$  with  $1 \leq r < n$

Ex:  $\lambda_j \geq \tilde{\lambda}_j \geq \lambda_{j+n-r}$  for  $j \in \{1, \dots, r\}$   
[Cauchy Interlacing Thm]

## Rephrasing on PSDness

Def A real sym matrix  $M$  is positive semi-definite (PSD) if  
 $\forall x \in \mathbb{R}^n, x^T M x \geq 0.$

Ex Prove: The following are equivalent

- 1)  $M$  is PSD
- 2)  $M$  has non-negative eigenvalues
- 3)  $\exists$  a matrix  $W$  s.t.  $M = W^T W$

Notation We write  $M \succ 0$  if  $M$  is PSD

We write  $M_1 \succ M_2$  if  $M_1 - M_2 \succ 0$ .

This gives a partial order (Loewner order)

## Laplacian Matrix

Def  $L = dI - A$  [Laplacian Matrix]

Let  $\mu_1 \leq \mu_2 \leq \dots \leq \mu_n$  be the eigenvalues of  $L$

Ex For  $d$ -regular  $G$ , we have

$$\mu_1 = d - \lambda_1, \dots, \mu_n = d - \lambda_n.$$

Ex Prove that  $\langle x, Lx \rangle = \sum_{j \sim i} (x_j - \bar{x})^2$

Ex Conclude that  $L \geq 0$  (PSD)

Ex Prove that  $\langle \mathbb{1}_S, L \mathbb{1}_S \rangle = |E(S, \bar{S})|$

Ex  $G$  connected  $\Leftrightarrow \mu_2 > 0$

Ex If  $G$  is connected, then  $\mu_2 \geq \frac{1}{n} \text{diam}(G)$

$E_{\lambda} \{ K_k | \mu_k = 0 \} = \# \text{ of connected components}$

Ex G bipartite iff  $\mu_n = 2d$

Ex Prove that  $\frac{\mu_2}{2d} \leq \underline{\Phi}(\epsilon)$

[CH\*] Prove that  $\underline{\Phi}(\epsilon) \leq O\left(\sqrt{\frac{\mu_2}{d}}\right)$

Hint: use eigenvector to  $\mu_2$  to find a cut ("rounding")

[Cheeger's Inequality]

$$\frac{\mu_2}{2d} \leq \underline{\Phi}(G) \leq \sqrt{\frac{2\mu_2}{d}}$$

## Characteristic Polynomial

Def Characteristic polynomial  $\det(\lambda I - A) =: ch(\lambda)$

The roots of  $ch(\lambda)$  are the eigenvalues of  $A$

Cayley-Hamilton Theorem:  $ch(A) = 0$

Ex 6 has diam =  $K \Rightarrow A$  has at least  
connected  $K+1$  distinct eigenvalues

Hint: [Cayley-Hamilton] minimal polynomial

Ex Let  $A, B$  be two real sym  
matrices with eig ( $\lambda_1 \geq \dots \geq \lambda_n$ )  
 $\tilde{\lambda}_1 \geq \dots \geq \tilde{\lambda}_n$ )

Compute the eigenvalues of  $A \otimes B$ .

Def  $S: E \rightarrow \{ \pm 1 \}$  is an edge  
signing

Def  $(A_S)_{u,v} = \begin{cases} S(\{u,v\}) & \text{if } \{u,v\} \in E \\ 0 & \text{otherwise} \end{cases}$

Ex Prove  $\lambda_1(A_S) \leq \Delta(G)$  for any  
signing  $S$ .

## Limitations on Spectral Expansion

Ex  $G$  d-regular  $\Rightarrow \lambda \geq \sqrt{d} (1 - o_n(1))$

Ex If  $G$  has diam  $\geq 4 \Rightarrow \lambda_2 \geq \sqrt{d}$   
[Hint: look for the stars]

Ex  $\lambda_2 < 0 \iff G = K_n$   
[Hint: interlacing]

Ex Suppose  $G$  is connected.  $G$  has a unique positive eigenvalue iff  $G$  is a complete  $k$ -partite graph  
[Hint: interlacing]

$$Ch^* \quad \lambda_2 \geq 2\sqrt{d-1} \left( 1 - O\left(\frac{1}{\text{diam}}\right) \right)$$

[Alon-Boppana bound]

(con) If  $d = O(1)$

$$\lambda_2 \geq 2\sqrt{d-1} \left( 1 - O\left(\frac{1}{\log n}\right) \right)$$

Def 6 is Ramanujan if  $\lambda \leq 2\sqrt{d-1}$

a.k.a. "Optimal" Spectral Expansion

OP [van Gogh Prize]

Construct infinite families

of Ramanujan graphs for every

$$d \geq 3$$

## Vertex Expansion

Def  $N(S) = \{u \mid \exists s \in S, \{u, s\} \in E\}$

Def Vertex (or Losen) Expansion

$$\Phi^V(S) = \frac{|N(S)|}{d|S|}$$

$$\underline{\Phi}_\epsilon^V(\epsilon) = \min_{\substack{\emptyset \neq S \subseteq V \\ |S| \leq \epsilon n}} \Phi^V(S).$$

## OP [van Gogh prize]

Construct explicit family with

$$\underline{\Phi}_\epsilon^V(\epsilon) > \frac{1}{2} \quad \text{for } \epsilon = \Omega(1)$$

(on two-sided bipartite lossen)

## On the Complexity of Expansion

(Hypothesis)  $\forall \eta \in (0, 1) \exists \delta \in (0, 1)$

s.t. it is NP-hard to distinguish  
given input graph  $G = (V, E)$

(Yes)  $\exists S \subseteq V$  with  $|S| \leq \delta n$  and  
 $\bar{\Phi}(S) \leq \eta$ .

(No)  $\nexists S \subseteq V$  with  $|S| \leq \delta n$ , we

have  $\bar{\Phi}(S) \geq 1 - \eta$

---

## OP van Gogh prize

Prove or refute the above hypothesis

---