

Project  
SNU 4190.210, Programming Principles Fall 2023  
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**due: 12/13(Wed) 23:59**

**Problem 1 (50 Points)** In Scala, implement an interpreter `interp` for the programming language E given below.

`interp` :  $E \rightarrow V$

$A$	$::=$	$x$	call by value
		$(\text{by-name } x)$	call by name
$B$	$::=$	$(\text{def } f_n (A^*) E)$	def
		$(\text{val } x E)$	val
		$(\text{lazy-val } x E)$	lazy val
$E$	$::=$	$x$	name
		$n$	integer
		$f$	float
		$s$	string
		$\text{nil}$	pair nil
		$(\text{cons } E E)$	pair constructor
		$(\text{fst } E)$	the first component of a pair
		$(\text{snd } E)$	the second component of a pair
		$(\text{int? } E)$	is int
		$(\text{float? } E)$	is float
		$(\text{string? } E)$	is string
		$(\text{nil? } E)$	is nil
		$(\text{pair? } E)$	is pair
		$(\text{substr } E E E)$	substring
		$(\text{len } E)$	length of a string or list
		$(\text{if } E E E)$	conditional
		$(\text{let } B^* E)$	name binding of def/val
		$(\text{app } E E^*)$	function call
		$(+ E E)$	addition / string concat
		$(- E E)$	subtraction
		$(* E E)$	multiplication
		$(/ E E)$	division
		$(\% E E)$	remainder
		$(= E E)$	equality
		$(< E E)$	less than
		$(> E E)$	greater than

- For ill-typed inputs, you can return arbitrary values, or raise exceptions.
- $X^*$  denotes that  $X$  can appear 0 or more times.
- **let** clauses create a new scope like a ‘block’ in Scala. Name bindings **def** and **val** work the similar way as in Scala.
  - $(\text{def } f (A^*) E)$  assigns name **f** to expression **E** with arguments  $A^*$ . Examples include  $(\text{def } f (a (\text{by-name } b)) (+ a b))$  and  $(\text{def } g () 3)$ .
  - $(\text{val } x E)$  assigns name  $x$  to the value obtained by evaluating  $E$ .
  - We do not allow the same name to be defined twice in the frame.
  - You do not have to consider forward reference in **val**. For example,  $(\text{val } x (\text{cons } 1 x))$ .
- **Environment** is collection of **Frames**. **Frame** is created when a new scope is created.

- Identifier `x` should be an alphanumeric word which does not start with a number.
- `nil` and `(cons v1 v2)` are pair type. `pair?` should return 1 for these values. Otherwise, it should return 0.
- `(int? E)` first evaluates `E` into value `v`. If `v` is `integer`, it returns 1. Otherwise, it returns 0. Also `nil?`, `float?`, `string?`, and `pair?` behave the same way.
- `(substr E1 E2 E3)` first evaluates `E1` into string `s` (If `E1` is not a string, raise any exception). `E2` and `E3` are the start and the end position of the substring of `s`. (You can simply use `String.substring` method of Scala)
- `(len E)` first evaluates `E` into value `v`. If `v` is a string or a pair (Cons or Nil), return the length of `v`. Otherwise, raise any exception.
- `len` of pair works simliar to Scala's `List[Any].length`. Since the last element of cons list from our language can be non-Nil element, `len` should caculate the number of the elements in the cons list, but must ignore the last Nil.
- e.g.) `(len (cons (cons (cons 5 4) 2) (cons 3 4))) = 3`, `(len (cons 2 (cons 3 nil))) = 2`, `(len nil) = 0`.
- For the binary operators `(+, -, *, /, %, =, <, >)`, the types of two operands must be number (except `=`). If one of the operand is float type, the result also have to be a float value. Otherwise, the result will be an integer value.
- As an exception, `+` is a string concatenation when the two operands are string values. Also you can use `=` to compare two strings.
- Comparison expressions `(=, <, >)` returns 1 if the comparison is right. Otherwise, it returns 0.
- `(if E1 E2 E3)` first evaluates `E1` into value `v`. If `v` is 0 or 0.0, it returns the result of `E3`. Otherwise, it returns the result of `E2`.
- `(lazy-val x E)` assigns name `x` to the value obtained by evaluating `E` lazily.
- Hint: Use `LazyOps`.
- For additional information, post questions on the GitHub course webpage.
- examples in `src/test/scala/TestSuite.scala`.

**Problem 2 (15 Points)** Optimize `interp` to handle tail recursive input programs, such as the example code shown below.

```
(let (def f (x sum) (if (> x 0) (app f (- x 1) (+ x sum)) sum))
(app f 10 0))
```

Hint: You don't need to reuse `Frame`. Just make `app` handler tail recursive, then you will get what you want.

**Problem 3 (15 Points)** Add IO action to `interp` by implementing `defIO`, `runIO`, `readline`, and `print` following:

$B$	$::=$	$\dots$	
		$(\text{defIO } fn(x^*) IO^* E)$	define an IO action
$IO$	$::=$	$(\text{runIO } x E)$	run an IO action and bind to $x$
		$(\text{readline } x)$	read a line text and bind to $x$
		$(\text{print } E)$	print

- $(\text{defIO } fn(A^*) (IO^*) E)$  defines a function that returns **IO action**.  $fn$  takes a list of parameters  $A^*$  like a normal function. Then, it executes  $IO^*$  sequentially, and finally  $E$  becomes the result of the IO action.
- The result of  $(\text{app } fn E^*)$  is **NOT** an evaluation of  $E$ . Instead, it returns an runnable *thunk* (pending action) which has a type of `IO[Value]` where the last  $E$  is `Value` type. From this stage, it does not actually runs  $(IO^*)$  inside  $fn$ . Those will be run when `runIO` calls the thunk which will be explained below.
- $(\text{runIO } x E)$  actually **runs** an IO action when  $E$  returns an IO action. (i.e. `IO[Value]`). The result of the IO action is bound to  $x$  (i.e. The type of  $x$  is `Value`). You can also bind any value to  $x$  other than IO action.  $x$  can be used in the following  $IO$  actions and the result  $E$ , but not in the previous  $IO$  actions. (i.e. does not allow forward reference).
- $(\text{readline } x)$  reads a line from the input (e.g. standard input) and binds it to  $x$ . The read value does not contain newline characters (`\n`, `\r`, `\r\n`, ...). You can assume that there will be only ASCII characters in the input.
- $(\text{print } E)$  prints the result of  $E$  to the output (e.g. standard output). It does not automatically append newline characters.
- If the top-level (i.e. the most outer) expression is an IO action, that IO action will be automatically executed if you run a main function.

You must use `Reader` and `Writer` trait to actually execute `readline` and `print`. Also, you have to use `Show` trait to convert a value to string for `print`.

For example,

```
(let ((defIO f1 (x) (runIO a readline) (+ a x)) (defIO f2 (b) (runIO
ax (f1 b)) (runIO b readline) (+ ax b))) (app f2 10))
```

reads two numbers `a`, `b`, and prints the value of `a + b + 10`.

**Problem 4 (20 Points)** Implement a calculator REPL (read-eval-print-loop) with the language E.

$$\begin{array}{lll} \textit{Exp} & ::= & n \quad \text{positive integer} \\ & | & \textit{Exp} + \textit{Exp} \quad \text{add} \\ & | & \textit{Exp} - \textit{Exp} \quad \text{sub} \\ & | & \textit{Exp} * \textit{Exp} \quad \text{mul} \\ & | & \textit{Exp} / \textit{Exp} \quad \text{div} \end{array}$$

- For each text taken from `readline`, if the text is a valid expression, evaluate the expression and print the result.
- If the text is not a valid expression, print `parse error`.
- If the text is `exit`, terminate the REPL.
- There will be no whitespaces in the input.