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Lecture 10: Structural Testing - Paths and Data Flow

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Test Adequacy Criteria

Compromise between
the impossible and the inadequate



- Can we measure “good testing”?
- **Test adequacy criteria** “score” tests by measuring completion of **test obligations**.
 - Checklists of properties that must be met by test cases.

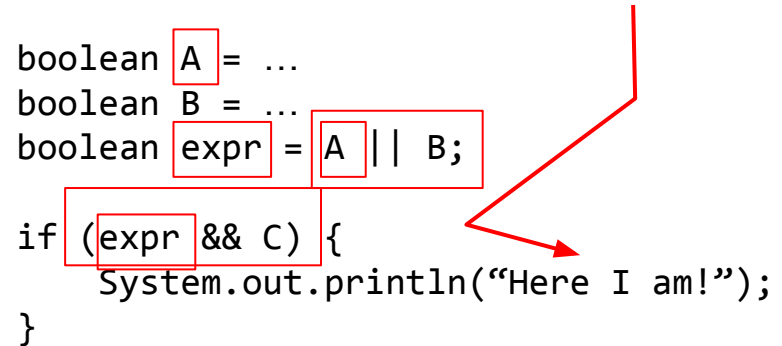
Structural Coverage Criteria

- Criteria based on exercising:
 - Statements (nodes of CFG)
 - Branches (edges of CFG)
 - Decisions and Conditions
 - Paths
 - ... and many more
- Measurements used as adequacy criteria

Elements Vs. Paths

- Statement, Branch, Condition Coverage all focus on one element at a time.
- A test executes a path, not a single element.
- Each element on that path is dependent on the others.

```
boolean A = ...  
boolean B = ...  
boolean expr = A || B;  
if (expr && C) {  
    System.out.println("Here I am!");  
}
```



Elements Vs. Paths

- There are different control paths through a program...
- ... And different ways that data passed along paths can influence execution.
- Important to examine not just elements, but paths.

```
boolean A = ... Fault in definition
boolean B = ...
boolean expr = A || B; Corrupts definition of expr if B = False

if (expr && C) {
    System.out.println("Here I am!");
}
```

expr can corrupt outcome if C = True

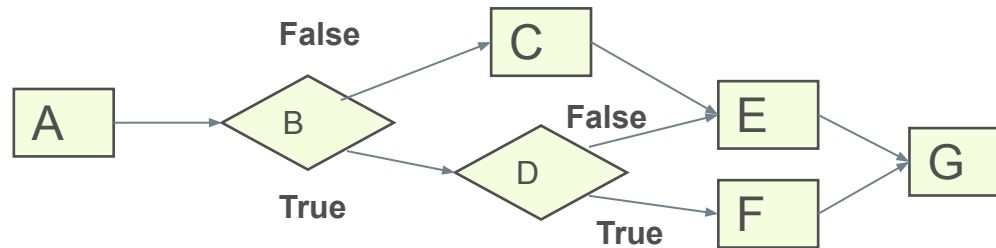
Today's Goals

- Introduce Path Coverage
- Data Flow Coverage Criteria
 - Focus on how information spreads through a program.
 - Based on Definition-Use Pairs
 - (Where is X defined? Where is each definition of X used?)

Path Coverage

Path Coverage

- Path coverage requires that all paths through the CFG are covered.



Paths:

A, B, C, E, G

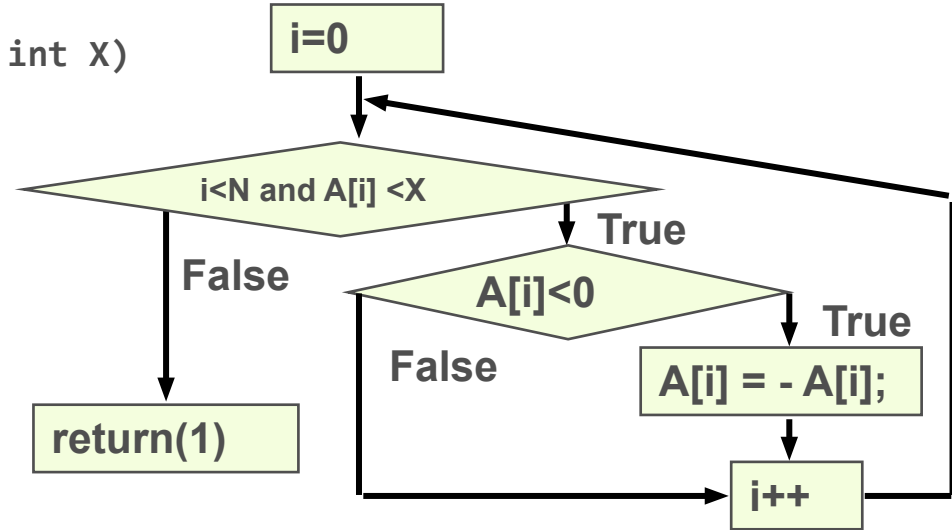
A, B, D, E, G

A, B, D, F, G

- $$\text{Coverage} = \frac{\text{Number of Paths Covered}}{\text{Number of Total Paths}}$$

Path Coverage

```
public int flipSome(int[] A, int N, int X)
{
    int i=0;
    while (i<N and A[i] <X)
    {
        if (A[i]<0)
            A[i] = - A[i];
        i++;
    }
    return A;
}
```

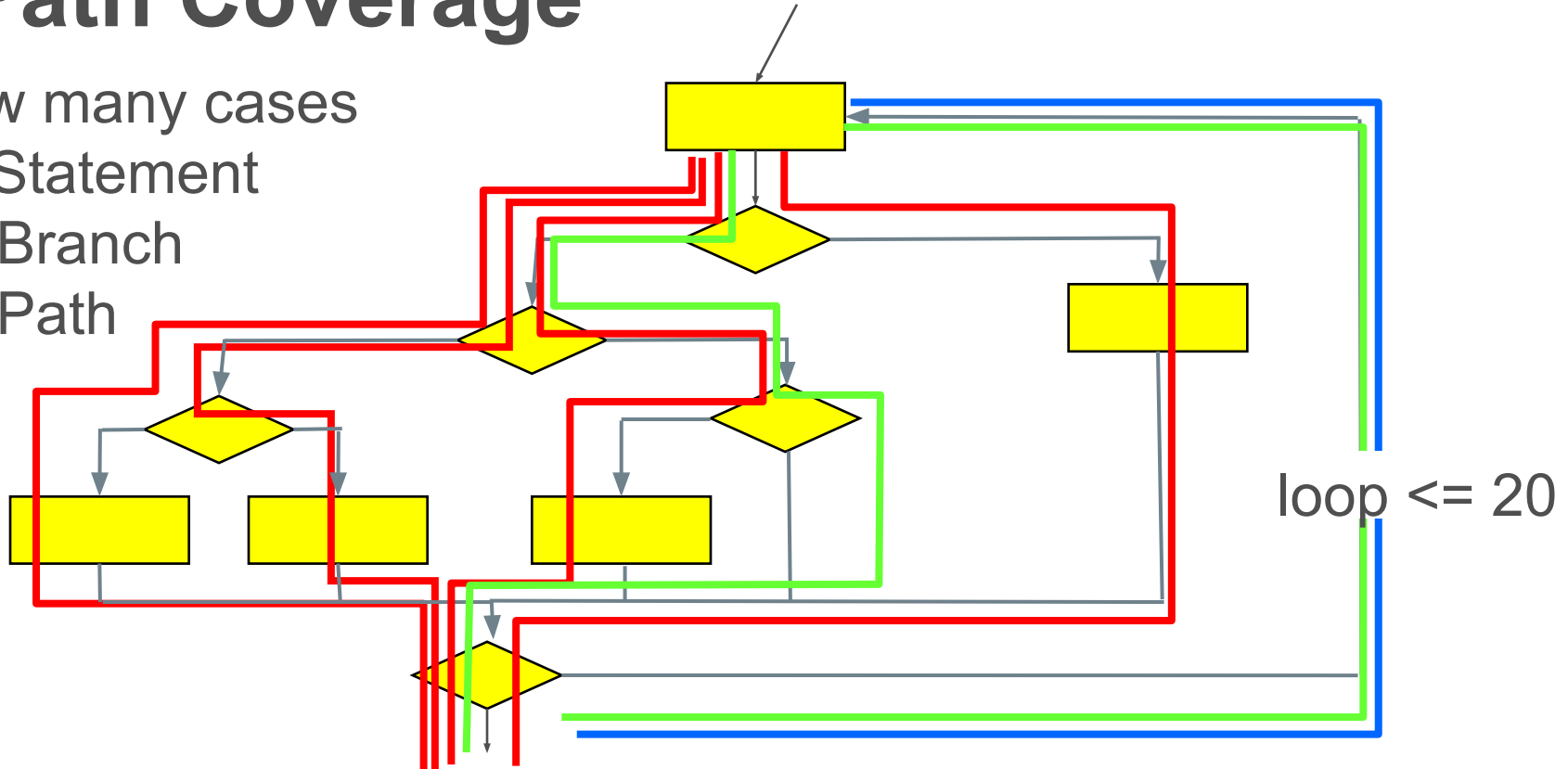


Path coverage is a powerful coverage metric, but is often impractical.

- How many paths does this have?
- Each loop cycle is a separate path!

Path Coverage

How many cases
for Statement
Branch
Path



Path coverage with (loop ≤ 20) requires:
3,656,158,440,062,976 test cases

If you run 1000 tests per second, this will
take **116,000 years**.

However, there are ways to get some of the benefits of
path coverage without the cost...

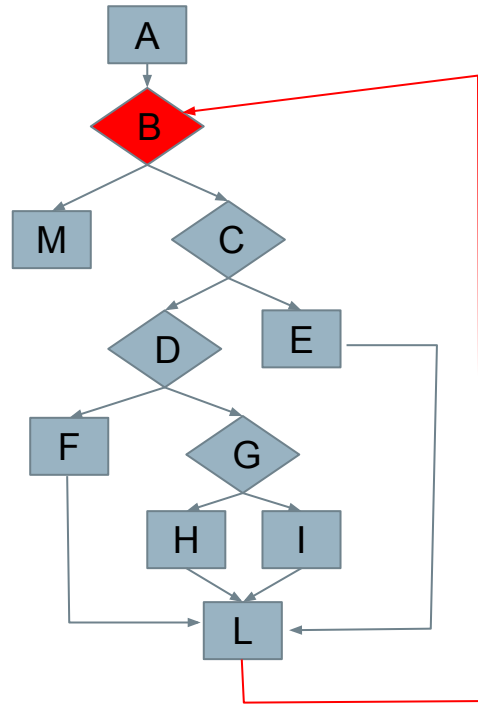
Path Coverage

- Theoretically, a very strong coverage metric.
 - Many faults emerge through sequences of interactions.
- But... Generally impossible to achieve.
 - Loops result in an infinite number of path variations.
 - Even bounding number of loop executions leaves an infeasible number of tests.

Boundary Interior Coverage

- Groups paths that differ only in the subpath they follow when repeating the body of a loop.
 - Executing loop 20 times is different than executing it twice, but same *subpaths* repeat over and over.
 - Unroll loop in CFG into distinct subpaths, and cover those instead of worrying about loop cycles.

Boundary Interior Coverage



A -> B -> M

A -> B -> C -> E -> L -> B

A -> B -> C -> D -> F -> L -> B

A -> B -> C -> D -> G -> H -> L -> B

A -> B -> C -> D -> G -> I -> L -> B

Boundary Interior Coverage

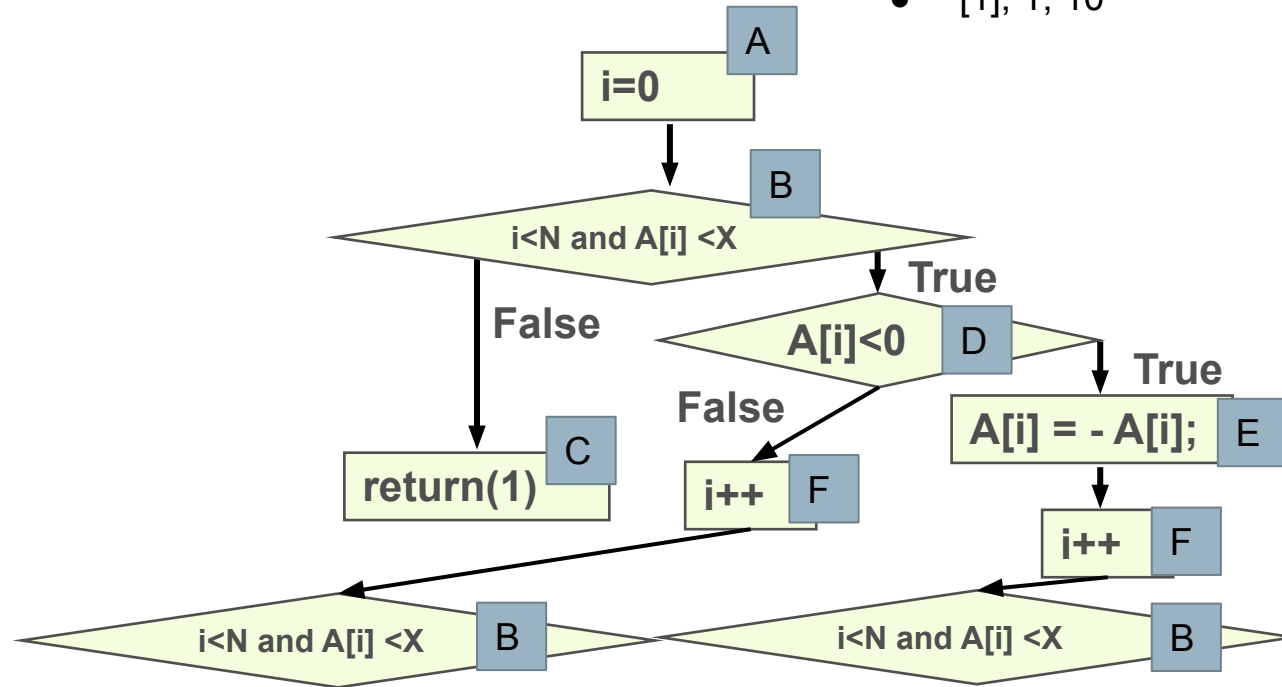
```
public int flipSome(int[] A, int N, int X)
{
    int i=0;
    while (i<N and A[i] <X)
    {
        if (A[i]<0)
            A[i] = - A[i];
        i++;
    }
    return A;
}
```

Paths:

- A, B, C
- A, B, D, F, B
- A, B, D, E, F, B

Test Input

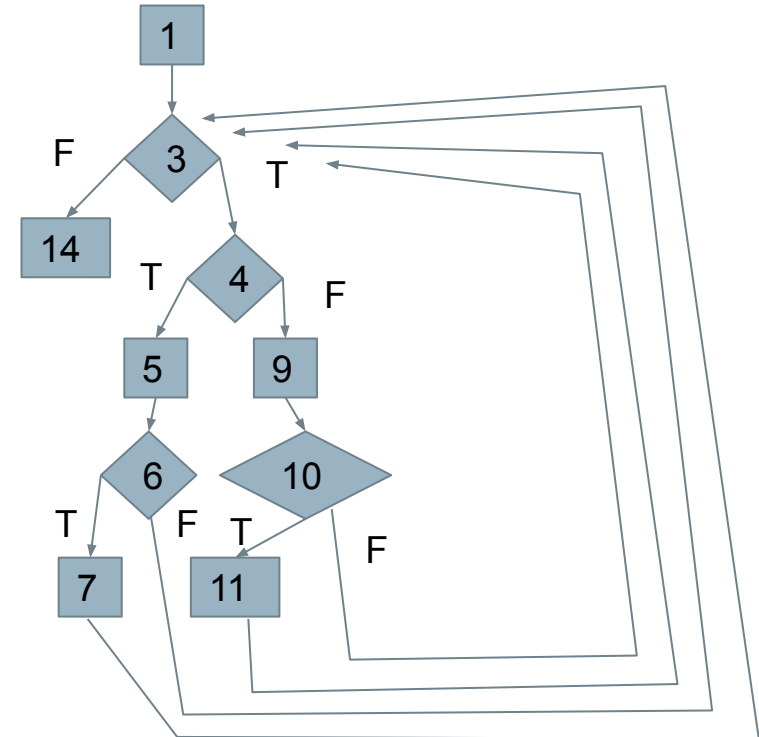
- [], 0, 10
- [-1], 1, 10
- [1], 1, 10



Boundary Interior Example

```

1. public int doSomething(int x, int y)
2. {
3.     while(y > 0) {
4.         if(x > 0) {
5.             y = y - x;
6.             if (y > 0)
7.                 System.out.println("Y: " + y);
8.         }else {
9.             x = x + 1;
10.            if (x <= 0)
11.                System.out.println(X: " + x);
12.        }
13.    }
14.    return x + y;
15. }
  
```



Boundary Interior Example

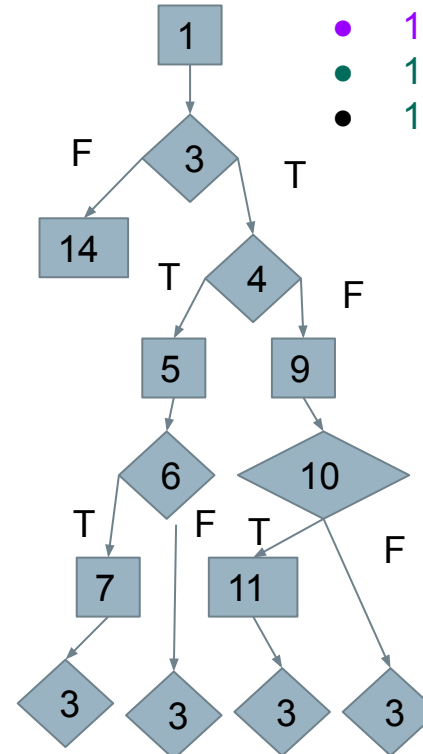
```

1. public int doSomething(int x, int y)
2. {
3.     while(y > 0) {
4.         if(x > 0) {
5.             y = y - x;
6.             if (y > 0)
7.                 System.out.println("Y: " + y);
8.         }else {
9.             x = x + 1;
10.            if (x <= 0)
11.                System.out.println(X: " + x);
12.        }
13.    }
14.    return x + y;
15. }

```

Paths:

- 1, 3-F, 14
- 1, 3-T, 4-T, 5, 6-T, 7, 3
- 1, 3-T, 4-T, 5, 6-F, 3
- 1, 3-T, 4-F, 9, 10-T, 11, 3
- 1, 3-T, 4-F, 9, 10-F, 3



Test Input:

- 10, -1
- 3, 4
- -1, 1

Number of Paths

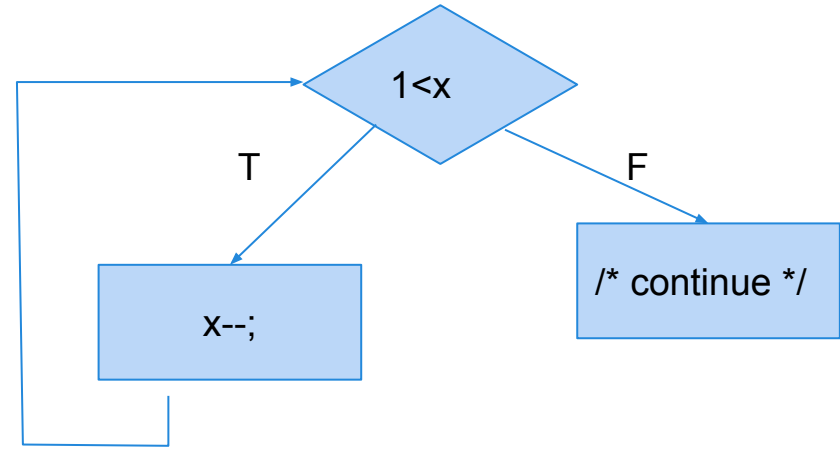
- Boundary Interior Coverage removes bounds number of loop paths.
 - However, number of paths can still be exponential.
 - N non-loop branches results in 2^N paths.
- Additional limitations may need to be imposed.

```
if (a)      S1;  
if (b)      S2;  
if (c)      S3;  
...  
if (x)      SN;
```

Data Flow

Control Flow

- Capture how execution navigates between blocks of statements.
- We care about a statement's effect **only when it affects the path**.
 - Deemphasizes information being transmitted.



Data Flow

- Program statements compute and transform data...
- Reason about data dependence
 - A variable is used here.
 - Where does its value come from?
 - Is this value ever used?
 - Is this variable properly initialized?
 - If the expression assigned to a variable is changed what else would be affected?

Data Flow

- Basis of the optimization performed by compilers.
- Used to derive test cases.
 - Have we covered the dependencies?
- Used to detect faults and other anomalies.
 - When can we cache result of a calculation instead of recalculating it?
 - Can we eliminate a variable definition?

Definition-Use Pairs

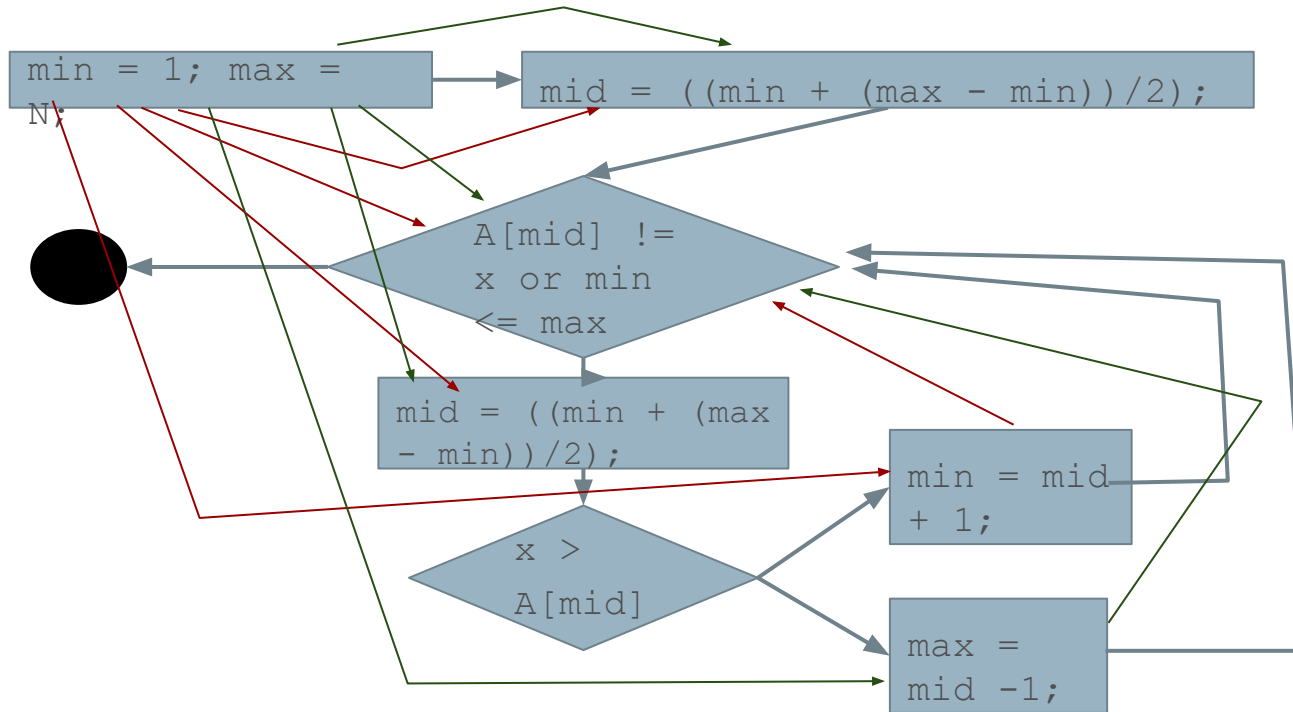
- Data is defined.
 - ... and data is used.
- Pairs of definitions and uses capture flow of information through the program.
 - Definitions occur when variables are declared, initialized, assigned values, or received as parameters.
 - Uses occur in expressions, conditional statements, parameter passing, return statements.

Definitions and Uses

```
1. min = 1;
2. max = N;
3. mid = ((min + (max - min))/2);
4. while (A[mid] != x or min <= max){
5.     mid = ((min + (max - min))/2);
6.     if (x > A[mid]){
7.         min = mid + 1
8.     } else {
9.         max = mid - 1;
10.    }
11. }
```

```
1. def - min
2. def - max, use - N
3. def - mid, use - min,
   max
4. use - A[mid], mid, x,
   min, max
5. def - mid, use - min,
   max
6. use - x, A[mid], mid
7. def - min, use - mid
8. -
9. def - max, use - mid
```


Definitions and Uses



1. **def** - `min`
2. **def** - `max`, **use** - `N`
3. **def** - `mid`, **use** - `min`, `max`
4. **use** - `A`, `mid`, `x`, `min`, `max`
5. **def** - `mid`, **use** - `min`, `max`
6. **use** - `x`, `A`, `mid`
7. **def** - `min`, **use** - `mid`
8. -
9. **def** - `max`, **use** - `mid`

Definition-Use (DU) Pairs

- We can say there is a **DU pair** when:
 - There is a **def** (definition) of variable x at location A.
 - Variable x is **used** at location B.
 - A control-flow **path** exists from A to B.
 - and the path is **definition-clear** for x from A to B.
- If variable is redefined, original def is **killed** and pair is now between new definition and its use in B.

Example - Definition-Use Pairs

```
1. min = 1;
2. max = N;
3. mid = ((min + (max - min))/2);
4. while (A[mid] != x or min <= max){
5.     mid = ((min + (max - min))/2);
6.     if (x > A[mid]){
7.         min = mid + 1
8.     } else {
9.         max = mid - 1;
10.    }
11. }
```

DU Pairs

min: (1, 3), (1, 4), (1, 5),
(7, 4), (7, 5)

max: (2, 3), (2, 4), (2, 5),
(9, 4), (9, 5)

N: (0, 2)

mid: (3, 4), (5, 6), (5, 7),
(5, 9), (5, 4)

x: (0, 4), (0, 6)

A: (0, 4), (0, 6)

Example - GCD

```
1. public int gcd(int x, int y){  
2.     int tmp;  
3.     while(y!=0){  
4.         tmp = x % y;  
5.         x = y;  
6.         y = tmp;  
7.     }  
8.     return x;  
9. }
```

```
1. def: x, y  
2. def: tmp  
3. use: y  
4. use: x, y  
   def: tmp  
5. use: y  
   def: x  
6. use: tmp  
   def: y  
7. -  
8. use: x
```

Example - GCD

```

1. public int gcd(int x, int y){
2.     int tmp;
3.     while(y!=0){
4.         tmp = x % y;
5.         x = y;
6.         y = tmp;
7.     }
8.     return x;
9. }

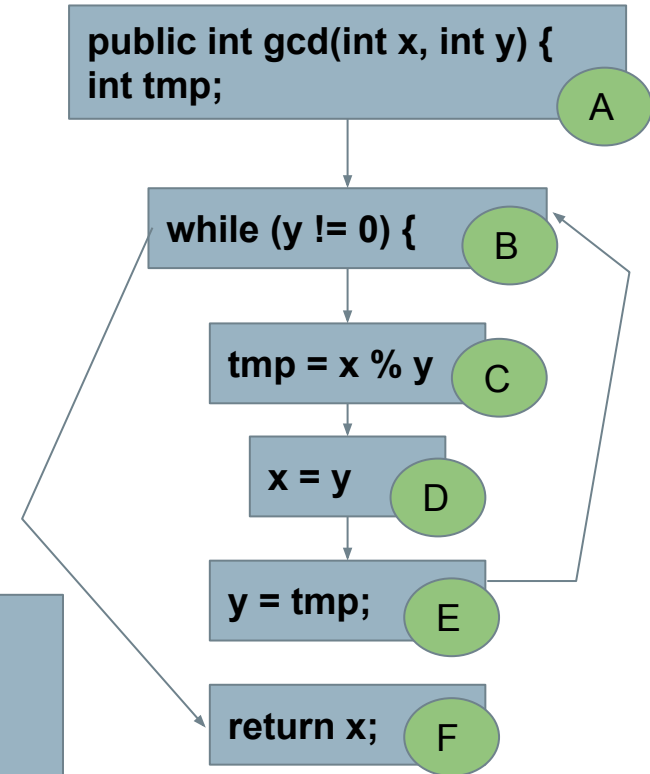
```

Def-Use Pairs

x: (1, 4), (5, 4), (5, 8), (1, 8)

y: (1, 3), (1, 4), (1, 5), (6, 3), (6, 4), (6, 5)

tmp: (4, 6)



Example - collapseNewlines

```
7. public static String collapseNewlines(String argStr)
8. {
9.     char last = argStr.charAt(0);
10.    StringBuffer argBuf = new StringBuffer();
11.
12.    for(int cldx = 0; cldx < argStr.length(); cldx++)
13.    {
14.        char ch = argStr.charAt(cldx);
15.        if(ch != '\n' || last != '\n')
16.        {
17.            argBuf.append(ch);
18.            last = ch;
19.        }
20.    }
21.
22.    return argBuf.toString();
23. }
```

| Variable | D-U Pairs |
|----------|------------------------------|
| argStr | (7, 9), (7,12), (7, 14) |
| last | (9, 15), (18, 15) |
| argBuf | (17, 22) |
| cldx | (12, 12), (12, 14) |
| ch | (14, 15), (14, 17), (14, 18) |

Let's Take a Break

Dealing With Arrays and Pointers

- Arrays and pointers (including object references and arguments) introduce issues.
 - It is not possible to determine whether two access refer to the same storage location.
 - $a[x] = 13;$
 $k = a[y];$
 - Are these a def-use pair?
 - $a[2] = 42;$
 $i = b[2];$
 - Are these a def-use pair?

Aliasing

- Two names refer to the same memory location.

- ```
int[] a = new int[3];
int[] b = a;
a[2] = 42;
i = b[2];
```

- Worse in C:

- ```
p = &b;  
*(p + i) = k;
```

Uncertainty

- Aliasing introduces uncertainty.
 - Instead of definition or use of one variable, may have a potential def or use of a set of variables.
- Safest: treat **any** use of a potential alias of V as a use of V .
 - Creates more def-use pairs (some may not be real), but avoids missed pairs.

Dealing With Uncertainty

- Treat all potential aliases as definitions and uses:

```
a[1] = 13;
```

```
k = a[2];
```

Def of **a[1]**, use of **a[2]**.

```
a[x] = 13;
```

```
k = a[y];
```

Def and use of **array a**.

- Can be very imprecise.
 - They are only the same if x and y are the same.

Dealing With Uncertainty

- Option 2: Treat uncertainty about aliases like uncertainty about control flow.

```
a[x] = 13;  
k = a[y];
```

→

```
a[x] = 13;  
if (x == y)    k = a[x];  
else          k = a[y];
```

- Rewrite code to make references explicit.
- In transformed code, all array references are distinct.

Situational Def-Use Pairs

- `++counter, counter++, counter+=1`
`counter = counter + 1`
 - Use of counter then a new definition.
- `char *ptr = *otherPtr`
 - Definition of string `*ptr`
 - Use of memory index `ptr`, string `*otherPtr`, and memory index `otherPtr`.
 - `ptr++`
 - Use of memory index `ptr`, definition of both memory index and string `*ptr` (change to index moves pointer to a new location).

Data Flow Coverage Criteria

Overcoming Limitations of Path Coverage

- We can potentially expose many faults by targeting particular paths of execution.
- What are the important paths to cover?
 - Some methods impose heuristic limitations.
 - Use data flow to select paths based on how one element can affect the computation of another.

Choosing the Paths

- Computing the wrong value leads to a failure **only when that value is used**.
 - Pair definitions with usages.
 - Ensure that definitions are actually used by covering paths from definitions to uses.
 - All DU Pair Coverage, All DU Paths Coverage, All Definitions Coverage
 - Varying power and cost.

All DU Pair Coverage

- Requires each DU pair be exercised in at least one program execution.
 - Cover **any path** between a definition and its use.
- Coverage =
$$\frac{\text{number exercised DU pairs}}{\text{number of DU pairs}}$$
- Can easily achieve structural coverage without covering all DU pairs.

All DU Pairs Coverage Example

```
1. public int doSomething(int x, int y)
2. {
3.     while(y > 0) {
4.         if(x > 0) {
5.             y = y - x;
6.             if (y > 0)
7.                 System.out.println("Y: " + y);
8.         }else {
9.             x = x + 1;
10.            if (x <= 0)
11.                System.out.println(X: " + x);
12.        }
13.    }
14.    return x + y;
15. }
```

X:

(1, 4), (1, 5), (1, 9), (1, 14)
(9, 10), (9, 11), (9, 5), (9, 9), (9, 14)

Y:

(1, 3), (1, 5), (1, 14)
(5, 6), (5, 7), (5, 3), (5, 5), (5, 14)

X: (1, 4), (1, 5), (1, 9), (1, 14), (9, 10), (9, 11), (9, 5), (9, 9), (9, 14)
Y: (1, 3), (1, 5), (1, 14), (5, 6), (5, 7), (5, 3), (5, 5), (5, 14)

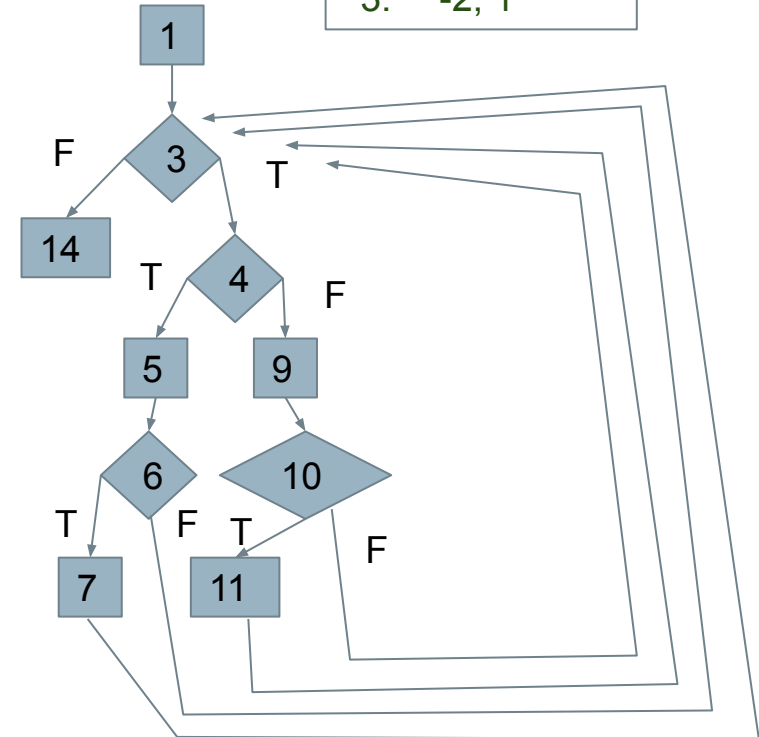
```

1. public int doSomething(int x, int y)
2. {
3.     while(y > 0) {
4.         if(x > 0) {
5.             y = y - x;
6.             if (y > 0)
7.                 System.out.println("Y: " + y);
8.         }else {
9.             x = x + 1;
10.            if (x <= 0)
11.                System.out.println(X: " + x);
12.        }
13.    }
14.    return x + y;
15. }

```

Test Input:

1. -1, 1
2. 3, 7
3. -2, 1



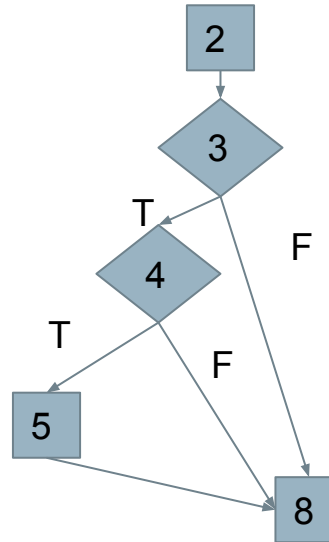
All DU Paths Coverage

- A use may be reachable along several paths from the definition.
- Cover all simple (non-looping) paths at least once.
 - Can reveal faults where a path is exercised that should use a certain definition but doesn't.

$$\text{Coverage} = \frac{\text{number of exercised DU paths}}{\text{number of DU paths}}$$

All DU Paths Example

```
1. ...
2. int x = 1;
3. if (y > 7) {
4.     if (z > 5) {
5.         z = x + 5;
6.     }
7. }
8. y = x + 7;
9. ...
```



DU Pair (2, 8) for X can be reached along multiple paths.

- 2, 3T, 4T, 5, 8
- 2, 3T, 4F, 8
- 2, 3F, 8

Test Input:

- $y = 10, z = 6$
- $y = 10, z = 3$
- $y = 2, z = (\text{anything})$

Path Explosion Problem

- Even without looping paths, number of DU paths can be exponential.
 - Code between definition and use can be irrelevant to that variable, but contains many paths.

```
public void countBits(char ch){  
    int count = 0;  
    if (ch & 1)    ++count;  
    if (ch & 2)    ++count;  
    if (ch & 4)    ++count;  
    if (ch & 8)    ++count;  
    if (ch & 16)   ++count;  
    if (ch & 32)   ++count;  
    if (ch & 64)   ++count;  
    if (ch & 128)  ++count;  
    System.out.println(ch + " has " +  
count + "bits set to 1");  
}
```

All Definitions Coverage

- All DU Pairs/All DU Paths may be too expensive in some situations.
- Pair each definition with at least one use.
 - Skips many DU pairs, but ensures each definition tried.

$$\text{Coverage} = \frac{\text{number of covered definitions}}{\text{number of definitions}}$$

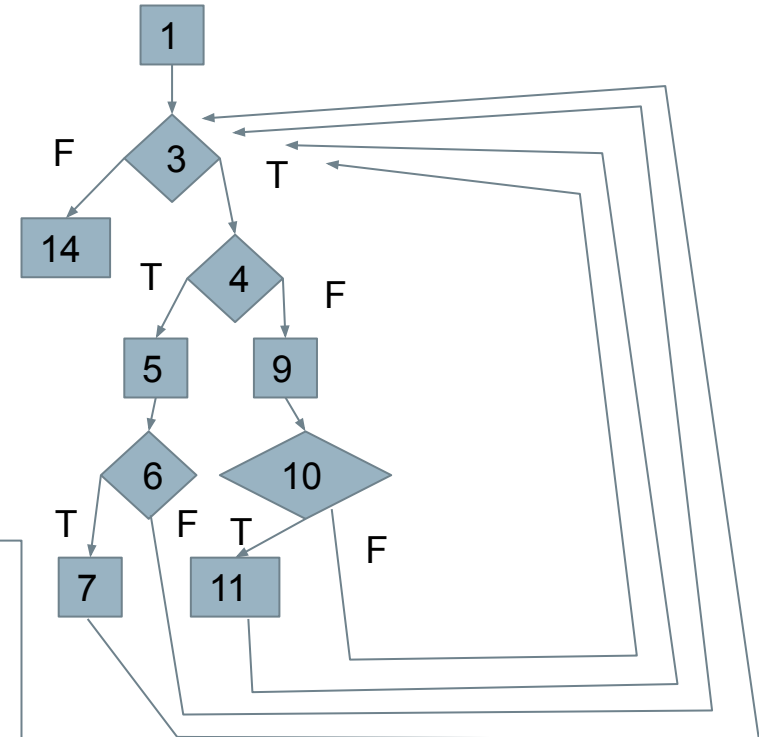
X: (1, 4), (1, 5), (1, 9), (1, 14), (9, 10), (9, 11), (9, 5), (9, 9), (9, 14)
Y: (1, 3), (1, 5), (1, 14), (5, 6), (5, 7), (5, 3), (5, 5), (5, 14)

X: Definitions on lines 1, 9
Y: Definitions on lines 1, 5

```

1. public int doSomething(int x, int y)
2. {
3.     while(y > 0) {
4.         if(x > 0) {
5.             y = y - x;
6.             if (y > 0)
7.                 System.out.println("Y: " + y);
8.         }else {
9.             x = x + 1;
10.            if (x <= 0)
11.                System.out.println(X: " + x);
12.        }
13.    }
14.    return x + y;
15. }
    
```

- Any input covers (1, -) pairs.
- Reaching lines 5, 9 covers (5,14) and (9,14) pairs.



Infeasibility Problem

- Metrics may ask for impossible test cases.
- Path-based metrics may require infeasible combinations of feasible elements.
 - Alias analysis may add additional infeasible paths.
- All Definitions, All DU-Pairs Coverage reasonable.
 - All DU-Paths is much harder!

Activity - DU Pair Coverage

- Identify all DU pair
- Write **your own** test input to achieve All DU Pair Coverage.
 - e.g., Input (1, 1)
For x, covers pairs:
(1,4), ...

```
1. int doSomething(int x, int y)
2. {
3.     while(y > 0) {
4.         if(x > 0) {
5.             y = y - x;
6.         }else {
7.             x = x + 1;
8.         }
9.     }
10.    return x + y;
11. }
```

Activity - DU Pairs

```
1. int doSomething(int x, int y)
2. {
3.     while(y > 0) {
4.         if(x > 0) {
5.             y = y - x;
6.         }else {
7.             x = x + 1;
8.         }
9.     }
10.    return x + y;
11. }
```

| Variable | Defs | Uses |
|----------|------|-------------|
| x | 1, 7 | 4, 5, 7, 10 |
| y | 1, 5 | 3, 5, 10 |

| Variable | D-U Pairs |
|----------|---|
| x | (1, 4), (1, 5), (1, 7), (1, 10), (7, 4), (7, 5), (7, 7), (7, 10) |
| y | (1, 3), (1, 5), (1, 10), (5, 3), (5, 5), (5, 10) |

Activity - DU Pairs

```

1. int doSomething(int x, int y)
2. {
3.     while(y > 0) {
4.         if(x > 0) {
5.             y = y - x;
6.         }else {
7.             x = x + 1;
8.         }
9.     }
10.    return x + y;
11. }
  
```

| Variable | Defs | Uses |
|----------|------|-------------|
| x | 1, 7 | 4, 5, 7, 10 |
| y | 1, 5 | 3, 5, 10 |

| Variable | D-U Pairs |
|----------|---|
| x | (1, 4), (1, 5), (1, 7), (1, 10), (7, 4), (7, 5), (7, 7), (7, 10) |
| y | (1, 3), (1, 5), (1, 10), (5, 3), (5, 5), (5, 10) |

Test Input 1: (x = 1, y = 2)

Covers lines 1, 3, 4, 5, 3, 4, 5, 3, 10

Test Input 2: (x = -1, y = 1)

Covers lines 1, 3, 4, 6, 7, 3, 4, 6, 7, 3, 4, 5, 3, 10

Test Input 3: (x = 1, y = 0)

Covers lines 1, 3, 8

We Have Learned

- Control-flow and data-flow capture important paths in program execution.
- Analysis of how variables are defined and then used and the dependencies between definitions and usages can help us reveal important faults.
- Many forms of analysis can be performed using data flow information.

We Have Learned

- If there is a fault in a computation, we can observe it by looking at where the computation is used.
- By identifying DU pairs and paths, we can create tests that trigger faults along those paths.
 - All DU Pairs coverage
 - All DU Paths coverage
 - All Definitions coverage

Next Time

- Exercise Session - Structural Testing
 - Using Meeting Planner code.
- Next Tuesday - Fault-Based Testing
 - Pezze & Young - Ch 16
- Assignment 2
 - Due February 26! We have covered everything on it.



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