## Accumulators

- accumulate information from prior recursive calls
- three categories
  - preserve context from prior recursive calls

• result so far

lecture 18 (previous)

this lecture

worklist

# Roadmap

- The next four lectures
  - forms of data: <u>trees</u> and graphs
  - recursion: <u>structural non-tail and tail</u> and generative
  - accumulators
    - path in data: previous, upper, lower, <u>pnum</u>, <u>path</u>, <u>visited</u>...
    - rsf (result so far)
    - path in tail recursion: visited, count, ...
    - worklist
    - tandem worklist

| 118 | <u>119</u> | <u>120</u> | <u>121</u>

```
(define (sequence? lon0)
 ;; prev is Natural
 ;; invariant: the element of lon0 immediately before (first lon)
 ;; (sequence? (list 2 3 4 7 5))
 ;; (sequence? (list
                       3 4 7 5) 2)
 ;; (sequence? (list
                         4 7 5) 3)
 ;; (sequence? (list
                          7 5) 4) ==> false
 (local [(define (sequence? lon prev)
           (cond [(empty? lon) true]
                  else
                  (if (= (first lon) (+ 1 prev)) ;exploit (use)
                      (sequence? (rest lon)
                                 (first lon))
                                               ;preserve
                      false)]))]
   (if (empty? lon0)
                                    ;if original list is empty, we can't
                                    ;initialize accumulator, so special case
       true
       (sequence? (rest lon0)
                  (first lon0))))) ;initialize
```

an expression is in tail-position when no surrounding expression in that function is waiting to operate on the value of the expression

```
tail position \longrightarrow (define (mumble x)
                *(* (+ 1 x) 2))
  not in tail
              ;; for if, cond and local expressions,
   position
              ;; IF THE WHOLE EXPRESSION IS IN TAIL POSITION,
              ;; then
              ;; questions are never in tail position
              ;; answers are in tail position
              ;; body is in tail position
              (if <question>
                  <true answer>
                   <false answer>)
              (cond [<question> <answer>]
                     . . . )
              (local [<definition> ...]
                <body>)
```

### definitions

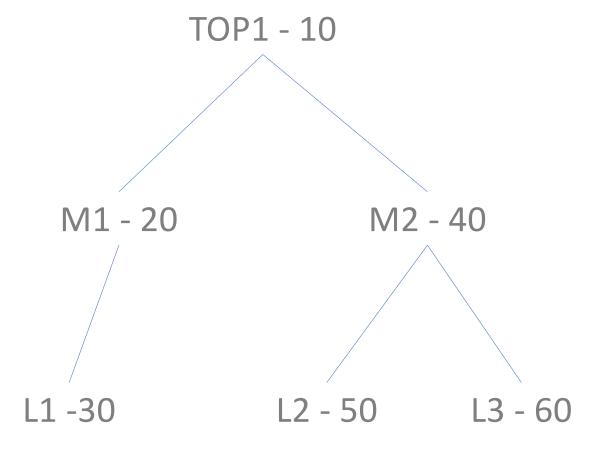
- an expression is in tail-position when no surrounding expression in that function is waiting to operate on the value of the expression
- a function is tail-recursive when ALL recursive calls are in tail position
- mutually-recursive functions are tail recursive when ALL recursive and mutually recursive calls are in tail position
- in tail recursive functions, base cases produce the final result

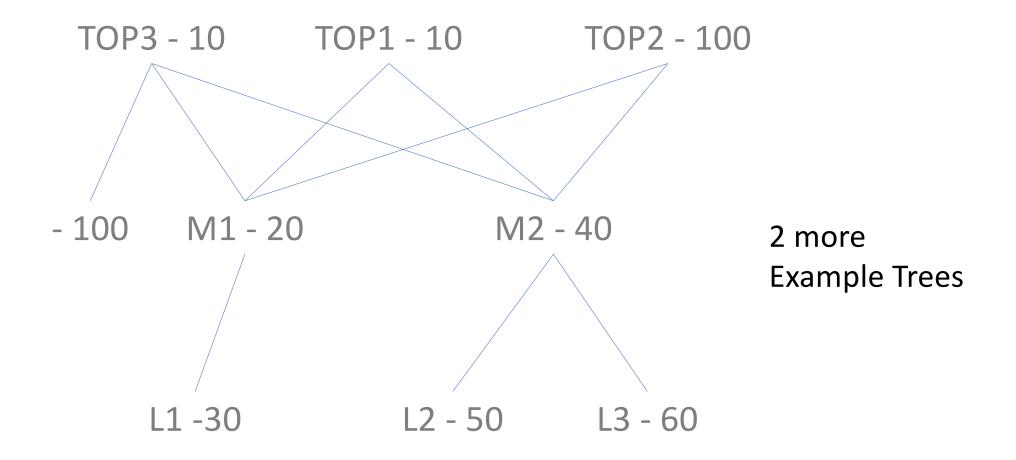
```
(define (fn-for-lox lox0)
        (cond [(empty? lox) empty]
                                                   ;base is empty
              [else___
             ;combination is cons
 not in tail
  position
      (define (rev lox0)
        (local [(define (fn-for-lox lox rsf)
                  (cond [(empty? lox) empty]
                       [else
                       ▶ (fn-for-lox ∫rest lox)
tail position
                                  ((cons (first lox)
rsf))]))l
                                                         ;combination moves
                                                         ;into rsf update
          (fn-for-lox lox0 empty)))
```

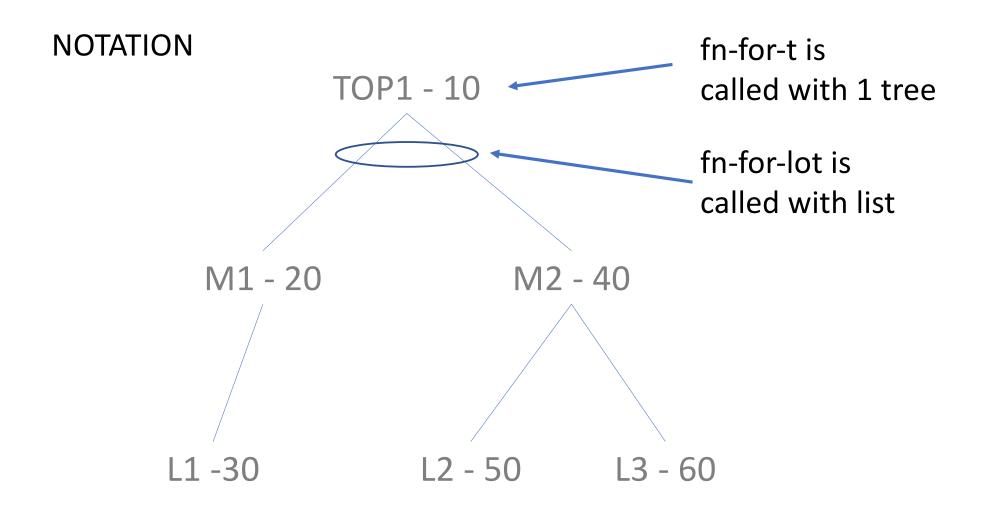
```
(define (rev lox0)
  (local [(define (fn-for-lox lox)
            (cond [(empty? lox) (... )]
                  [else
                   (... (first lox)
                        (fn-for-lox (rest lox)))]))]
    (fn-for-lox lox0)))
                                                      produce rsf at end
(define (rev lox0)
 ;; rsf is (listof X)
 ;; all elements of lox0 before (first
                                          Xox), in reverse order
  (local [(define (fn-for-lox lox rsf
            (cond [(empty? lox) rsf]
                   [else
                   (fn-for-lox (rest lox) (cons (first lox) rsf))]))]
    (fn-for-lox lox0 empty)))
                                                      combination
                                    initialize rsf (base case result)
```

```
;; QUESTION 1 [45 seconds]
;;
;; Is the call to positive? in tail position?
(define (positive-only lon)
  (cond [(empty? lon) empty]
        [else
         (if (positive? (first lon))
             (cons (first lon)
                   (positive-only (rest lon)))
             (positive-only (rest lon)))]))
;; A. Yes
;; B. No
```

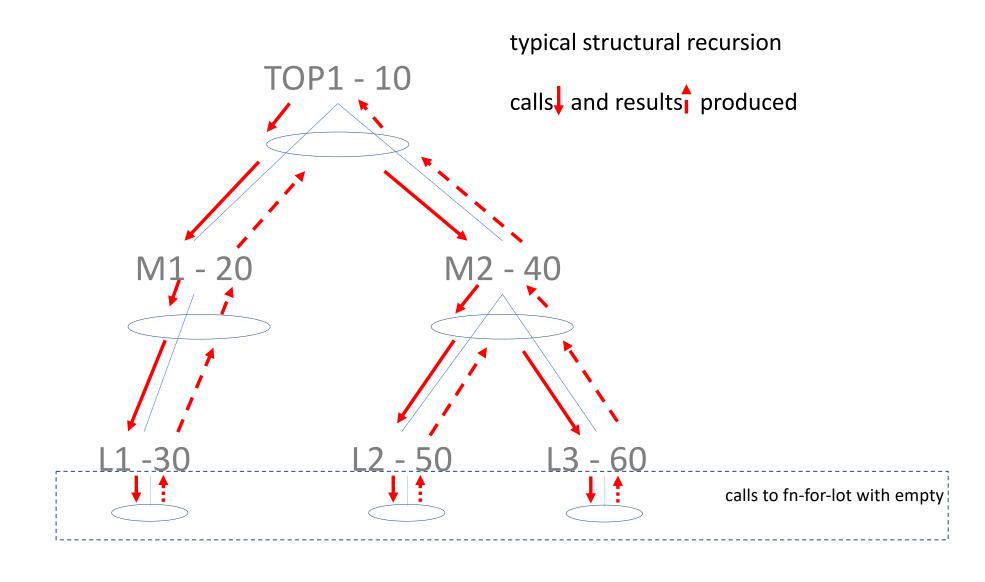
## **Example Tree**







```
(define (top->bot-sorted? t0)
  ;; pnum is Integer; immediate parent node's number
  ;; **parent means the parent in the tree **
  (local [(define (fn-for-t t pnum)
            (local [(define number (node-number t)) ;unpack the fields
                    (define subs (node-subs t))]
                                                     :for convenience
              (if (> number pnum)
                  (fn-for-lot subs number)
                  false)))
          (define (fn-for-lot lot pnum)
            (cond [(empty? lot) true]
                  [else
                   (and (fn-for-t (first lot) pnum)
                        (fn-for-lot (rest lot) pnum))]))]
    (fn-for-lot (node-subs t0) (node-number t0))))
```



# What's in the past depends on the recursion

- tail recursion means current call can have all the preceding context
  - → it can produce answer directly
- in a tree
  - ordinary recursion can carry context of what is above current call
  - but tail recursion is required to carry context of what is above and to the LEFT

```
(define (top/left->bot/right-sorted? t0)
 ;; t-wl is (listof Tree); worklist of Trees to visit
                            unvisited direct subs of visited trees
  ;; vnum is Integer; node number of most recently VISITED node
  ;; ** visited means in the dynamic flow of the tail recursion **
 ;; ** not in the static structure of the tree
 (local [(define (fn-for-t t t-wl vnum)
           (local [(define number (node-number t)) ;unpack the fields
                    (define subs (node-subs t))]
                                                     ;for convenience
              (if (> number vnum)
                  (fn-for-lot (append subs t-wl) number)
                 false)))
          (define (fn-for-lot t-wl vnum)
            (cond [(empty? t-wl) true]
                  [else
                   (fn-for-t (first t-wl) (rest t-wl) vnum)]))]
   (fn-for-t t0 empty (sub1 (node-number t0)))))
```

