

HW #3

Ryan English

11.2-2

11.3-3

11.4-1

11.4-3

11.3-4

11.2-2

Demonstrate what happens when we insert the keys 5, 28, 19, 15, 20, 33, 12, 17, 10 into a hash table with collision resolving chaining. Let the table have 9 slots and let the hash func be $h(k) = k \bmod 9$

$h(k) = k \bmod 9$ 0, 1, 2, 3, 4, 5, 6, 7, 8, 9

$h(k)$	Keys	
0 mod 9		$5 = 5 \bmod 9 = 5$
1 mod 9	28 → 19 → 10	$28 = 28 \bmod 9 = 1$
2 mod 9	20	$19 = 19 \bmod 9 = 1$
3 mod 9	12	$15 = 15 \bmod 9 = 6$
4 mod 9		$20 = 20 \bmod 9 = 2$
5 mod 9	5	$33 = 33 \bmod 9 = 6$
6 mod 9	15 → 33	$12 = 12 \bmod 9 = 3$
7 mod 9		$17 = 17 \bmod 9 = 8$
8 mod 9	17	$10 = 10 \bmod 9 = 1$

11.3-3

Consider a version of the division method in which $h(k) = k \bmod m$ where $m = 2^p - 1$ and k is a character string interpreted in radix 2^p . Show that if we can derive string x from string y by permuting its characters, then x and y hash to the same value. Give an example.

$$h(k) = k \bmod m$$

$$m = 2^p - 1$$

Anagrams would map to same value

Base case

$x = 'a'$ or any letter

$y = 'a'$ or any letter

$$h(a) = 0 \bmod 1 = 0$$

$$h(i) = 1 \bmod 1 = 1$$

Becomes a sum of the length of characters

11.4-1

Consider inserting the keys 10, 22, 31, 4, 15, 28, 17, 88, 59 into the hash table $m=11$ using open addressing with the aux hash $h'(k)=k$. Illustrate the results of inserting these keys using linear probing. Use quadratic $c_1=1$ $c_2=3$, and use double hashing $h_1(k)=k$ and $h_2(k)=1+(k \bmod (m-1))$

[10, 22, 31, 4, 15, 28, 17, 88, 59] $h'(k)=k$ $m=11$
 $h(k)=k \bmod 11$

Linear

$$h(k, i) = (h'(k) + i) \bmod m = (k + i) \bmod 11$$

Probe up to m $T = 8$ $i = \text{collisions}$

	T_0	T_1	T_2	T_3	T_4	T_5	T_6	T_7	T_8
0 mod 11		22							
1 mod 11								*	
2 mod 11								88	
3 mod 11									
4 mod 11				4	*				*
5 mod 11					15				*
6 mod 11						28	*		*
7 mod 11							17		*
8 mod 11									
9 mod 11									59
10 mod 11	10		31						

11.4-1 Continued

Quadratic $c_1 = 1$ $c_2 = 3$

$$A(k, i) = (k + i + 3i^2) \bmod 11$$

	T_0	T_1	T_2	T_3	T_4	T_5	T_6	T_7	T_8
0 mod 11		22						*	
1 mod 11								0+1+3	
2 mod 11									
3 mod 11									
4 mod 11					4	*		88	*
5 mod 11					4+1+3=8				4+1+3
6 mod 11						28	*		
7 mod 11							6+1+3=10		
8 mod 11					15		17		
9 mod 11		31							59
10 mod 11	10						*		
							10+2+6		

Double Hash $A(k, i) = (k + i + 5i^2) \bmod 10$

	T_0	T_1	T_2	T_3	T_4	T_5	T_6	T_7	T_8
0 mod 11		22					*	*	
1 mod 11							0+2+5 mod 10		
2 mod 11							17	*	*
3 mod 11								1+5	6
4 mod 11					4	*			
5 mod 11					4+1+5 mod 10				
6 mod 11						28	*		6+3
7 mod 11							6+1+5 mod 10		*
8 mod 11									7+5
9 mod 11		31			15				59
10 mod 11	10								

11.4-3

Consider an open-address hash table with uniform hashing. Give upper bounds on the expected number of probes in an unsuccessful search on the expected number of probes in a successful search when the load factor is $3/4$ and when it is $7/8$

$$\alpha = 1/1 - \alpha$$

$$1/\alpha \ln 1/1 - \alpha$$

$$\alpha = 3/4$$

$$\alpha = 3/4$$

$$1/1 - 3/4 = 4 \text{ probes unsuccessful}$$

$$1/3/4 \ln 1/1 - \alpha \approx 1.8 \text{ successful}$$

$$\alpha = 7/8$$

$$1/1 - 7/8 = 8 \text{ unsuccessful}$$

$$1/7/8 \ln 1/1 - 7/8 \approx 2.3 \text{ successful}$$

11.423 3-4

Consider a hash table of size $m=1000$
and a corresponding hash function $h(k) = \lfloor m(kA \bmod 1) \rfloor$ for $A = (\sqrt{5}-1)/2$. Compute
the location to which the keys $61, 62, 63, 64$,
and 65 are mapped

$$h(k) = \lfloor m(kA \bmod 1) \rfloor \quad m=1000$$

$$A = (\sqrt{5}-1)/2 \approx .61$$

$$61, 62, 63, 64, 65$$

$$h(61) = \lfloor 1000(61 \cdot .61 \bmod 1) \rfloor = 700$$

$$h(62) = \lfloor 1000(62 \cdot .61 \bmod 1) \rfloor = 318$$

$$h(63) = \lfloor 1000(63 \cdot .61 \bmod 1) \rfloor = 934$$

$$h(64) = \lfloor 1000(64 \cdot .61 \bmod 1) \rfloor = 554$$

$$h(65) = \lfloor 1000(65 \cdot .61 \bmod 1) \rfloor = 172$$