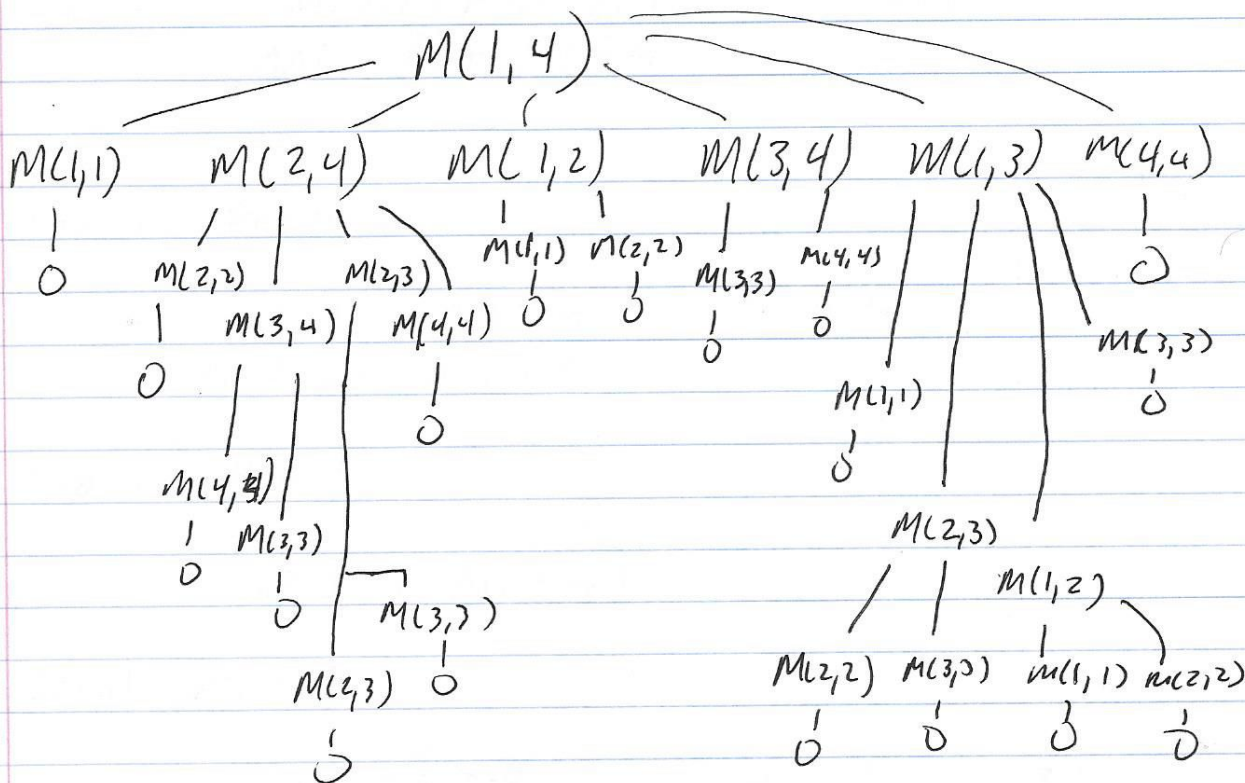


Project 2 Ryan English 03/06/21

Draw Recursive Tree at  $M(1,4)$



Depth = 4 = j

i = j return 0 → length of 1

for i to j      n  
2 calls      2<sup>n</sup>

$$T(n) = \begin{cases} 1 & \text{if } n=1 \\ 1 + \sum_{k=1}^{n-1} T(k) + T(k+1) & \text{if } n \geq 2 \end{cases}$$

$$T(n) = \Omega(2^n)$$

Recursive  
 $T(n) = \Omega(2^n)$

Iteration  
 $T(n) = O(n^3)$

Recursive is much slower