

Lecture 2

Wednesday, January 13, 2021 11:21 AM

③ Time Complexity

	Cost	Count
INSERTION_SORT(A)		
n = length(A);	C1	1
for j = 2 : n	C2	n
temp = A(j);	C3	n-1
i = j - 1;	C4	n-1
while ((i > 0) && (x(i) > temp))	C5	t_j
A(i+1) = A(i);	C6	$t_j - 1$
i = i - 1;	C7	$t_j - 1$
A(i+1) = temp;	C8	n-1

; t_j : random Var.

$$\min(t_j) = 1 \quad ; \quad \max(t_j) = j$$

We can write $T(n)$ as

$$T(n) = C' + B'n + D' \sum_{j=2}^n t_j$$

C', B', D' are constants

Two Cases: [1] $A(1..n)$ is already sorted
 $t_j = 1$ for every j

$$T(n) = C' + B'n + D' [1 + 1 + \dots + 1]$$

$$T(n) = \alpha + \beta n \quad \leftarrow \text{Linear} \quad - (A)$$

[1] Worst Case "A is sorted in the reverse order"

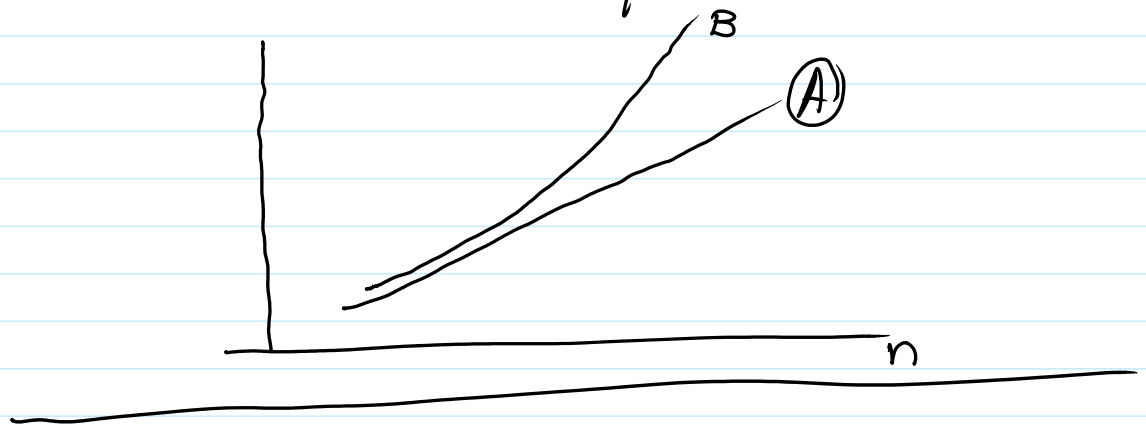
$$T(n) = C' + B'n + D' [2 + 3 + 4 + \dots + n]$$

$$T(n) = C + B'n + D' \left[\underbrace{2+3+4+\dots+n}_{\text{arithmetic seq.}} \right]$$

$$= C' + B'n + D' \left(\frac{2+n}{2} \right) (n-1)$$

$$= \alpha + \beta n + \gamma n^2 \quad \text{--- (B)}$$

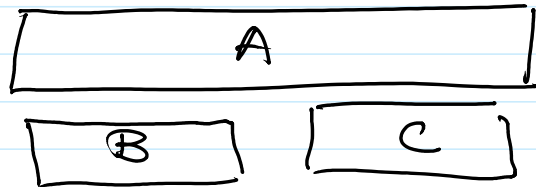
In this case running time
is Quadratic



Merge Sort

Basic Idea

1. Divide Array A into two sub-array



2. Conquer - which means to sort B and C
3. Combine or Merge the Conquered B and C

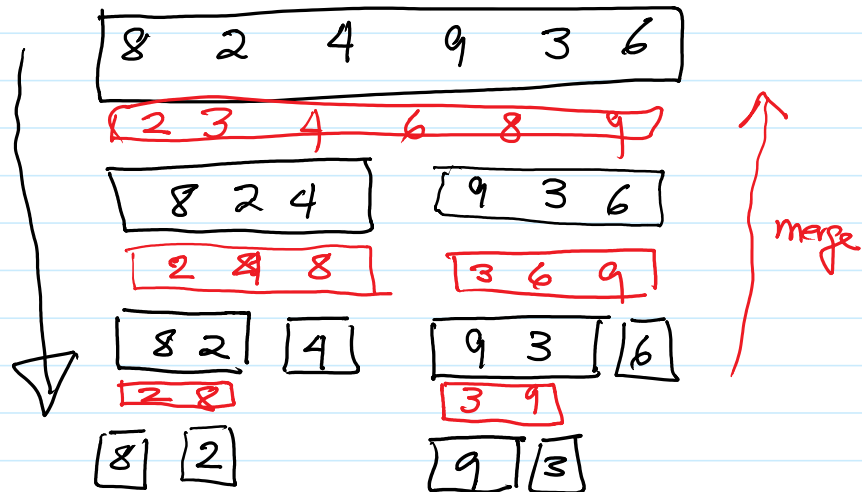
Numerical Example



Numerical Example

$n = 6$

Divide



Alg.

Merge ($A[1..n]$)

If $n == 1 \Rightarrow$ Done

else

merge ($A[1.. \lfloor \frac{n}{2} \rfloor]$)

merge ($A[\lfloor \frac{n}{2} \rfloor + 1, n]$)

Combine the two Sublists

MATLAB DEMO

```
function y = INSERTION_SORT(x)
n = length(x);
for j = 2 : n
    temp=x(j);
    i = j - 1;
    while ((i > 0) && (x(i) > temp))
        x(i+1) = x(i);
        i = i - 1;
    end
    x(i + 1) = temp;
end
y = x;
end
```

```
>> for j=1 : 3000 x(j)=floor(rand()*100);end;
>> tic;y=INSERTION_SORT(x);t=toc;
>> t
```

t =

0.1198

```
>> plot(x)
>> figure;plot(y)
>>
```

