

(Adjoint) Algorithmic Differentiation [(A)AD]

A Hands-On Introduction

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Introduction

First-Order(A)AD

Prerequisites Tangents Adjoints

Second-(and Higher-)Order (A)AD

Tangents of Tangents
Tangents of Adjoints

"Getting Serious" with AAD

Implementation by Overloading Checkpointing Symbolic Adjoints Adjoint Code Design Patterns

Progress



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The Art of Differentiating Computer Programs Bumping?



For differentiation, is there anything else?

Perturbing the inputs – can't imagine this fails.

I pick a small Epsilon, and I wonder ...



from: "Optimality" (Lyrics: Naumann; Music: Think of Fool's Garden's "Lemon Tree") in Naumann: The Art of Differentiating Computer Programs. An Introduction to Algorithmic Differentiation. Number 24 in Software, Environments, and Tools, SIAM, 2012. Page xvii



The Art of Differentiating Computer Programs Story



- inspired by sensitivity analysis, uncertainty quantification, calibration / optimization
- finite differences (first- and second-order), symbolic, algorithmic
- ▶ implementation by overloading, source trafo, hand-coding
- real code
- sensitivity analysis as modelling and software engineering tool
- what matters
 - user expertise
 - tool quality
 - tool sustainability and support



Let
$$y = F(\mathbf{x}), F : \mathbb{R}^n \to \mathbb{R}$$
:

- 1. tangent AD: $y^{(1)} = \nabla F \cdot \mathbf{x}^{(1)} \Rightarrow \nabla F$ at $O(n) \cdot \mathsf{Cost}(F)$
- 2. adjoint AD: $\mathbf{x}_{(1)} = \nabla F^T \cdot y_{(1)} \Rightarrow \nabla F$ at $O(1) \cdot \mathsf{Cost}(F)$
- 3. 2nd-order tangent AD: $y^{(1,2)}=\mathbf{x}^{(1)}^T\cdot \nabla^2 F\cdot \mathbf{x}^{(2)}\Rightarrow \nabla^2 F$ at $O(n^2)\cdot \mathrm{Cost}(F)$
- 4. 2nd-order adjoint AD: $\mathbf{x}_{(1)}^{(2)} = y_{(1)} \cdot \nabla F^2 \cdot \mathbf{x}^{(2)} \Rightarrow \nabla^2 F$ at $O(n) \cdot \mathsf{Cost}(F)$ and $\nabla^2 F \cdot \mathbf{x}^{(2)}$ at $O(1) \cdot \mathsf{Cost}(F)$

Aims of this Course



You will learn how to

- implement tangent and adjoint versions of a Monte Carlo / Euler-Maruyama solver for parameterized scalar SDEs
- ensure feasibility of adjoint Monte Carlo simulation through pathwise adjoints
- "get serious" with AAD (tools, checkpointing, symbolic adjoints, design patterns, ...)

Euler-Maruyama



We are looking for the expected value $\mathbb{E}(x)$ of the solution $x(\mathbf{p},T),T>0$ of the scalar stochastic initial value problem

$$dx = f(x(\mathbf{p}, t), \mathbf{p}, t))dt + g(x(\mathbf{p}, t), \mathbf{p}, t)dW$$

with Brownian Motion dW and for $x(\mathbf{p},0)=x^0$.

Forward finite differences in time with time step $0<\delta t\ll 1$ yield the explicit Euler-Maruyama evolution

$$x^{i+1} := x^i + \delta t \cdot f(x^i, \mathbf{p}, i \cdot \delta t) + \sqrt{\delta t} \cdot g(x^i, \mathbf{p}, i \cdot \delta t) \cdot dW^i$$

for $i=0,\ldots,n-1,$ target time $T=n\cdot\delta t,$ parameter vector $\mathbf{p}\in\mathbb{R}^l,$ and with random numbers dW^i drawn from the standard normal distribution N(0,1).

The solution $\mathbb{E}(x)$ is approximated using Monte Carlo simulation over (a sufficiently large number of) Euler-Maruyama paths.

We are interested in sensitivities of the final state $\mathbb{E}(x)$ wrt. \mathbf{p} .

Motivation

Race (Euler-Maruyama $m=10^4, n=10^2, l=10^2$)



- primal: primal.cpp (inspect)
- bumping: fd.cpp (inspect)
- ▶ tangent: tangent.cpp (live)
- vector tangent: tangent_vector.cpp (inspect)
- adjoint: adjoint.cpp (live)
- pathwise adjoint: adjoint_pathwise.cpp (inspect)

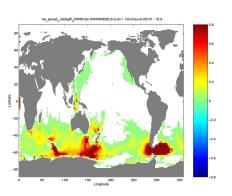
mode	run time (s)	memory size (b)	accuracy
bump	$10.9 \sim O(l)$	$\sim P$	_
tangent	$21.5 \sim O(2 \cdot l)$	$\sim 2 \cdot P$	+
vector tangent	$13.6 \sim O(2 \cdot l)$	$\sim P + P \cdot l$	+
adjoint	$0.3 \sim O(1)$	$\sim 2 \cdot P + 2 \cdot m \cdot n \cdot 8$	+
pathwise adjoint	$0.5 \sim O(1)$	$\sim 2 \cdot P + 2 \cdot (m+n) \cdot 8$	+

where P denotes the memory requirement of the primal code.

Motivation

Adjoint Nice To Have?





MITgcm, (EAPS, MIT)

in collaboration with ANL, MIT, Rice, UColorado

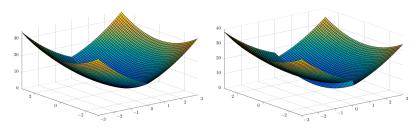
J. Utke, U.N. et al: OpenAD/F: A modular, open-source tool for automatic differentiation of Fortran codes . ACM TOMS 34(4), 2008.

Plot: A tangent computation / finite difference approximation for 64,800 grid points at 1 min each would keep us waiting for a month and a half ... :-(((We can do it in less than 10 minutes thanks to adjoints computed by a differentiated version of the MITgcm :-)

Fundamental Mathematics



- continuity
- ▶ differentiability?



- ▶ gradient, Jacobian, Hessian, higher-order derivative tensors
- ► Taylor expansion
- ► chain rule

Progress



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Prerequisites: Feasible Target Code I



1. The given implementation of $F : \mathbb{R}^n \to \mathbb{R}^m : \mathbf{y} = F(\mathbf{x})$, can be decomposed into a single assignment code (SAC)

$$v_i = \varphi_i(x_i) = x_i \qquad i = 0, \dots, n-1$$

$$v_j = \varphi_j \left((v_k)_{k \prec j} \right) \qquad j = n, \dots, n+q-1$$

$$y_k = \varphi_{n+q+k}(v_{n+p+k}) = v_{n+p+k} \quad k = 0, \dots, m-1$$

where q=p+m and $k\prec j$ denotes a direct dependence of v_j on v_k as an argument of φ_j .

2. All elemental functions φ_j possess continuous partial derivatives

$$d_{j,i} \equiv \frac{d\varphi_j}{dv_i}(v_k)_{k \prec j}$$

with respect to their arguments $(v_k)_{k \prec j}$ at all points of interest.

First-Order (A)AD



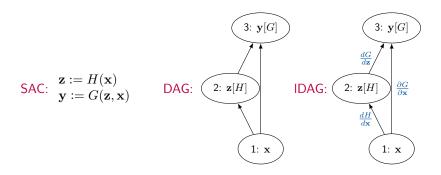


- 3. A linearized SAC (ISAC) is obtained by augmenting the elemental assignments with computations of the local partial derivatives $d_{j,i}$.
- 4. The SAC induces a directed acyclic graph (DAG) G = G(F) = (V, E) with integer vertices $V = \{0, \dots, n+q\}$ and edges $V \times V \supseteq E = \{(i,j) : i \prec j\}.$
- 5. The set of vertices representing the n inputs is denoted as $X \subseteq V$. The m outputs are collected in $Y \subseteq V$. All remaining intermediate vertices belong to $Z \subsetneq V$.
- 6. A linearized DAG (IDAG) is obtained by attaching the $d_{j,i}$ to the corresponding edges (i,j) in the DAG.

First-Order (A)AD

Prerequisites: Chain Rule on IDAG





$$\nabla F(\mathbf{x}) \equiv \frac{d\mathbf{y}}{d\mathbf{x}} = \sum_{\mathsf{path} \in \mathsf{IDAG}} \prod_{(i,j) \in \mathsf{path}} d_{j,i}$$

Combinatorics of Chain Rule



▶ U. N.: *Optimal Jacobian accumulation is NP-complete.* Math. Prog. 112(2):427–441, Springer, 2008.

Proof by reduction from Ensemble Computation

▶ U. N.: Optimal accumulation of Jacobian matrices by elimination methods on the dual computational graph. Math. Prog. 99(3):399–421, Springer, 2004.

Example: bat graph in STCE logo

▶ A. Griewank and U. N.: *Accumulating Jacobians as chained sparse matrix products.* Math. Prog. 95(3):555–571, Springer, 2003.

Example: $\mathbb{R}^4 \to \mathbb{R}^2 \to \mathbb{R}^2 \to \mathbb{R}^2 \to \mathbb{R}^4$

Mathematician's View

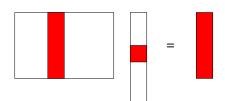


A first-order tangent model $F^{(1)}: \mathbb{R}^n \times \mathbb{R}^n \to \mathbb{R}^m \times \mathbb{R}^m$,

$$\begin{pmatrix} \mathbf{y} \\ \mathbf{y}^{(1)} \end{pmatrix} = F^{(1)}(\mathbf{x}, \mathbf{x}^{(1)}),$$

defines a directional derivative alongside with the function value:

$$\begin{aligned} \mathbf{y} &= F(\mathbf{x}) \\ \mathbf{y}^{(1)} &= \nabla F(\mathbf{x}) \cdot \mathbf{x}^{(1)} \end{aligned}$$



... definition of the whole Jacobian column-wise by input directions $\mathbf{x}^{(1)} \in \mathbb{R}^n$ equal to the Cartesian basis vectors in \mathbb{R}^n .

Computer Scientist's View



A first-order tangent code $F^{(1)}: \mathbb{R}^n \times \mathbb{R}^n \times \mathbb{R}^{\tilde{n}} \to \overline{\mathbb{R}^m \times \mathbb{R}^m \times \mathbb{R}^{\tilde{m}}}$

$$\begin{pmatrix} \mathbf{z} \\ \mathbf{z}^{(1)} \\ \tilde{\mathbf{z}} \\ \mathbf{y} \\ \mathbf{y}^{(1)} \\ \tilde{\mathbf{y}} \end{pmatrix} := F^{(1)}(\mathbf{x}, \mathbf{x}^{(1)}, \tilde{\mathbf{x}}, \mathbf{z}, \mathbf{z}^{(1)}, \tilde{\mathbf{z}}),$$

computes a Jacobian imes vector product alongside with the function value:

$$\mathbb{R}^m \times \mathbb{R}^{\tilde{m}} \ni \begin{pmatrix} \mathbf{z} \\ \tilde{\mathbf{z}} \\ \mathbf{y} \\ \tilde{\mathbf{y}} \end{pmatrix} := F(\mathbf{x}, \tilde{\mathbf{x}}, \mathbf{z}, \tilde{\mathbf{z}})$$

$$\mathbb{R}^m \ni \begin{pmatrix} \mathbf{z}^{(1)} \\ \mathbf{y}^{(1)} \end{pmatrix} := \nabla F(\mathbf{x}, \tilde{\mathbf{x}}, \mathbf{z}, \tilde{\mathbf{z}}) \cdot \begin{pmatrix} \mathbf{x}^{(1)} \\ \mathbf{z}^{(1)} \end{pmatrix}$$



Variables for which derivatives are computed are referred to as active; ${\bf x}$ and ${\bf z}$ are active inputs; ${\bf z}$ and ${\bf y}$ are active outputs.

Variables which depend on active inputs are referred to as varied.

Variables for which no derivatives are computed are referred to as passive; $\tilde{\mathbf{x}}$ and $\tilde{\mathbf{z}}$ are passive inputs; $\tilde{\mathbf{z}}$ and $\tilde{\mathbf{y}}$ are passive outputs.

Variables on which active outputs depend are referred to as useful.

Active variables are both varied and useful.

The whole (dense) Jacobian can be *harvested* column-wise from the active output directions $(\mathbf{z}^{(1)},\mathbf{y}^{(1)})^T\in\mathbb{R}^m$ by *seeding* active input directions $(\mathbf{x}^{(1)},\mathbf{z}^{(1)})^T\in\mathbb{R}^n$ with the Cartesian basis vectors in \mathbb{R}^n .

Computer Scientist's View (Simplified)



A first-order tangent code $F^{(1)}: \mathbb{R}^n \times \mathbb{R}^n \to \mathbb{R}^m \times \mathbb{R}^m$,

$$\begin{pmatrix} \mathbf{y} \\ \mathbf{y}^{(1)} \end{pmatrix} := F^{(1)}(\mathbf{x}, \mathbf{x}^{(1)}),$$

computes a Jacobian \times vector product alongside with the function value:

$$\mathbf{y} := F(\mathbf{x})$$
$$\mathbf{y}^{(1)} := \nabla F(\mathbf{x}) \cdot \mathbf{x}^{(1)}$$

Conceptually



For i = 0, ..., n - 1

$$\begin{pmatrix} v_i \\ v_i^{(1)} \end{pmatrix} \coloneqq \begin{pmatrix} x_i \\ x_i^{(1)} \end{pmatrix} \qquad \text{(seed)}$$

For $i = n, \ldots, q-1$

$$\begin{pmatrix} v_i \\ v_i^{(1)} \end{pmatrix} \coloneqq \begin{pmatrix} \varphi_i(v_k)_{k \prec i} \\ \sum_{j \prec i} \frac{d\varphi_i(v_k)_{k \prec i}}{dv_j} \cdot v_j^{(1)} \end{pmatrix} \qquad \text{(propagate)}$$

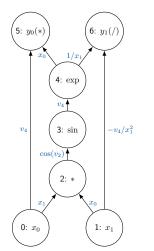
For i = 0, ..., m - 1

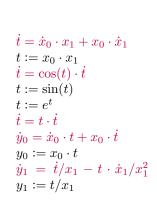
$$\begin{pmatrix} y_i \\ y_i^{(1)} \end{pmatrix} := \begin{pmatrix} v_{n+p+i} \\ v_{n+p+i}^{(1)} \end{pmatrix} \qquad \text{(harvest)}$$

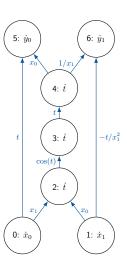
Simple Example



Tangent assignments augment primal ...







Tangents

Case Study



Euler-Maruyama live ...

Vector Tangents

Computer Scientist's View (Simplified)

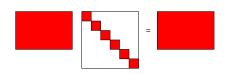


A first-order vector tangent code $F^{(1)}: \mathbb{R}^n \times \mathbb{R}^{n \times l} \to \mathbb{R}^m \times \mathbb{R}^{m \times l},$

$$\begin{pmatrix} \mathbf{y} \\ Y^{(1)} \end{pmatrix} \coloneqq F^{(1)}(\mathbf{x}, X^{(1)}),$$

computes a Jacobian \times matrix product alongside with the function value:

$$\mathbf{y} := F(\mathbf{x})$$
$$Y^{(1)} := \nabla F(\mathbf{x}) \cdot X^{(1)}$$



... harvesting of the whole Jacobian by seeding input directions $X^{(1)}[i] \in \mathbb{R}^n, \ i=0,\dots,n-1,$ with the Cartesian basis vectors in \mathbb{R}^n . Note concurrency!

Vector Tangents Case Study



Euler-Maruyama live ...

Adjoints



The Jacobian is a linear operator $\nabla F : \mathbb{R}^n \to \mathbb{R}^m$.

Its adjoint is defined as $(\nabla F)^*: {\rm I\!R}^m \to {\rm I\!R}^n$ where

$$<(\nabla F)^* \cdot \mathbf{y}_{(1)}, \mathbf{x}^{(1)}>_{\mathbf{R}^n} = <\mathbf{y}_{(1)}, \nabla F \cdot \mathbf{x}^{(1)}>_{\mathbf{R}^m}$$

and where $<.,.>_{\mathbf{R}^n}$ and $<.,.>_{\mathbf{R}^m}$ denote appropriate scalar products in \mathbb{R}^n and \mathbb{R}^m , respectively.

Theorem

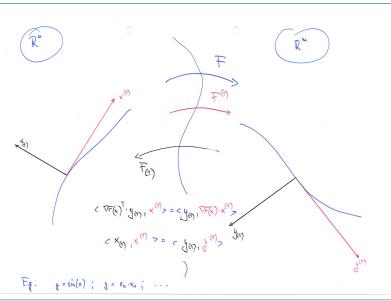
$$(\nabla F)^* = (\nabla F)^T.$$

$$<(\nabla F)^T \cdot \mathbf{y}_{(1)}, \mathbf{x}^{(1)}>_{\mathbf{R}^n} = <\mathbf{y}_{(1)}, \nabla F \cdot \mathbf{x}^{(1)}>_{\mathbf{R}^m}$$
 $[=:\mathbf{y}_{(1)}]$

Note invariant at each point in the program execution \rightarrow validation.

Adjoints





Mathematician's View

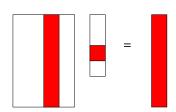


A first-order adjoint model $F_{(1)}: \mathbb{R}^n \times \mathbb{R}^m \to \mathbb{R}^m \times \mathbb{R}^n$,

$$\begin{pmatrix} \mathbf{y} \\ \mathbf{x}_{(1)} \end{pmatrix} = F_{(1)}(\mathbf{x}, \mathbf{y}_{(1)}),$$

defines an adjoint directional derivative alongside with the function value:

$$\mathbf{y} = F(\mathbf{x})$$
$$\mathbf{x}_{(1)} = \nabla F(\mathbf{x})^T \cdot \mathbf{y}_{(1)}$$



... definition of the whole Jacobian row-wise through input directions $\mathbf{y}_{(1)} \in \mathbb{R}^m$ equal to the Cartesian basis vectors in \mathbb{R}^m .

Notation



In

$$\left(\frac{dF}{d\mathbf{x}}\right)^T \cdot \mathbf{y}_{(1)}$$

the subscript on \mathbf{y} denotes the first directional differentiation of F performed in adjoint mode in direction $\mathbf{y}_{(1)} \in \mathbb{R}^m$.

Enumeration of derivatives and distinction of super- and subscripts will become relevant in the discussion of higher derivatives computed by combinations of tangent and adjoint modes.

Computer Scientist's View



$$\begin{split} F_{(1)}: \mathbb{R}^n \times \mathbb{R}^{n_{\mathbf{x}}} \times \mathbb{R}^m \times \mathbb{R}^{\tilde{n}} &\to \mathbb{R}^m \times \mathbb{R}^n \times \mathbb{R}^{m_{\mathbf{y}}} \times \mathbb{R}^{\tilde{m}}, \\ & \left(\mathbf{z} \quad \tilde{\mathbf{z}} \quad \mathbf{y} \quad \tilde{\mathbf{y}} \quad \mathbf{x}_{(1)} \quad \mathbf{z}_{(1)} \quad \mathbf{y}_{(1)}\right)^T := F_{(1)}(\mathbf{x}, \mathbf{x}_{(1)}, \tilde{\mathbf{x}}, \mathbf{z}, \mathbf{z}_{(1)}, \tilde{\mathbf{z}}, \mathbf{y}_{(1)}), \end{split}$$

computes a shifted transposed Jacobian \times vector product alongside with the function value:

$$\mathbf{R}^{m} \times \mathbf{R}^{m} \ni \begin{pmatrix} \mathbf{z} \\ \tilde{\mathbf{z}} \\ \mathbf{y} \\ \tilde{\mathbf{y}} \end{pmatrix} := F(\mathbf{x}, \tilde{\mathbf{x}}, \mathbf{z}, \tilde{\mathbf{z}})$$
$$\begin{pmatrix} \mathbf{x}_{(1)} \\ \mathbf{z}_{(1)} \end{pmatrix} := \begin{pmatrix} \mathbf{x}_{(1)} \\ 0 \end{pmatrix} + \nabla F(\mathbf{x}, \tilde{\mathbf{x}}, \mathbf{z}, \tilde{\mathbf{z}})^{T} \cdot \begin{pmatrix} \mathbf{z}_{(1)} \\ \mathbf{y}_{(1)} \end{pmatrix}$$
$$\mathbf{y}_{(1)} := 0$$

Computer Scientist's View



The whole (dense) Jacobian can be harvested from the active input adjoints

$$\begin{pmatrix} \mathbf{x}_{(1)} \\ \mathbf{z}_{(1)} \end{pmatrix} \in \mathbb{R}^m$$

row-wise by seeding active output adjoints

$$\begin{pmatrix} \mathbf{z}_{(1)} \\ \mathbf{y}_{(1)} \end{pmatrix} \in \mathbb{R}^m$$

with the Cartesian basis vectors in \mathbb{R}^m and for $\mathbf{x}_{(1)} := 0$ on input.

Computer Scientist's View (Simplified)



A first-order adjoint code $F_{(1)}: \mathbb{R}^n \times \mathbb{R}^n \times \mathbb{R}^m \to \mathbb{R}^m \times \mathbb{R}^n$,

$$\begin{pmatrix} \mathbf{y} \\ \mathbf{x}_{(1)} \end{pmatrix} := F_{(1)}(\mathbf{x}, \mathbf{x}_{(1)}, \mathbf{y}_{(1)}),$$

computes a shifted transposed Jacobian \times vector product alongside with the function value:

$$\begin{split} \mathbf{y} &:= F(\mathbf{x}) \\ \mathbf{x}_{(1)} &:= \mathbf{x}_{(1)} + \nabla F(\mathbf{x})^T \cdot \mathbf{y}_{(1)} \end{split}$$

... harvesting of the whole Jacobian row-wise by seeding input directions $\mathbf{y}_{(1)} \in \mathbb{R}^m$ with the Cartesian basis vectors in \mathbb{R}^m and for $\mathbf{x}_{(1)} = 0$ on input.

Conceptually I



1. Augmented Primal (enable data flow reversal)

For
$$i=0,\dots,n-1$$

$$v_i \coloneqq x_i \\ \text{record } i \in V \ (v_{i_{(1)}} \coloneqq x_{i_{(1)}})$$
 For $i=n,\dots,q-1$
$$v_i \coloneqq \varphi_i(v_k)_{k \prec i} \\ \text{record } i \in V \ (v_{i_{(1)}} \coloneqq 0)$$
 For $j \prec i \colon \text{record } (i,j) \in E \ (d_{j,i} \coloneqq \frac{d\varphi_i(v_k)_{k \prec i}}{dv_j})$ For $i=0,\dots,m-1$
$$y_i \coloneqq v_{n+p+i}$$

Conceptually II



2. Adjoint

For
$$i=0,\dots,m-1$$

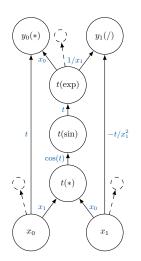
$$v_{n+p+i_{(1)}}:=y_{i_{(1)}}$$
 For $i=q-1,\dots,n$
$$\forall \ (j,i)\in E:v_{j_{(1)}}:=v_{j_{(1)}}+v_{i_{(1)}}\cdot d_{i,j}$$
 For $i=0,\dots,n-1$
$$x_{i_{(1)}}:=v_{i_{(1)}}$$

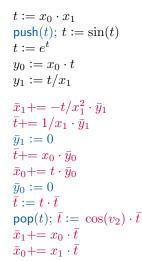
Adjoints

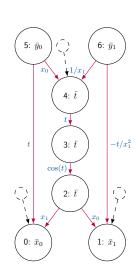
Simple Example



Mind overwrites and context ...







Adjoints

Case Study



Euler-Maruyama live ...

Progress



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Adjoints

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Adjoint Code Design Patterns

Second Derivatives

Multivariate Scalar Functions



Initially we consider multivariate scalar functions $y = F(\mathbf{x}) : D_F \subseteq \mathbb{R}^n \to I_F \subseteq \mathbb{R}$ in order to simplify the notation.

We assume F to be twice continuously differentiable over its domain D_F implying the existence of the Hessian

$$\nabla^2 F(\mathbf{x}) \equiv \frac{d^2 F}{d\mathbf{x}^2}(\mathbf{x}).$$

For multivariate vector functions the Hessian is a three-tensor complicating the notation slightly due to the need for tensor arithmetic; see later.

Numerical Approximation of Second Derivatives



Multivariate Scalar Functions

A second-order *central finite difference* quotient

$$\frac{d^2 F}{dx_i dx_j}(\mathbf{x}^0) \approx \left[F(\mathbf{x}^0 + (\mathbf{e}_j + \mathbf{e}_i) \cdot h) - F(\mathbf{x}^0 + (\mathbf{e}_j - \mathbf{e}_i) \cdot h) - F(\mathbf{x}^0 + (\mathbf{e}_i - \mathbf{e}_j) \cdot h) + F(\mathbf{x}^0 - (\mathbf{e}_j + \mathbf{e}_i) \cdot h) \right] / (4 \cdot h^2)$$
(1)

yields an approximation of the second directional derivative

$$y^{(1,2)} = \mathbf{x}^{(1)^T} \cdot \nabla^2 F(\mathbf{x}) \cdot \mathbf{x}^{(2)} \quad \text{(w.l.o.g. } m = 1)$$

as

$$\frac{d^2 F}{dx_i dx_j}(\mathbf{x}^0) \approx \frac{\frac{dF}{dx_i}(\mathbf{x}^0 + \mathbf{e}_j \cdot h) - \frac{dF}{dx_i}(\mathbf{x}^0 - \mathbf{e}_j \cdot h)}{2 \cdot h}$$

$$= \left[\frac{F(\mathbf{x}^0 + \mathbf{e}_j \cdot h + \mathbf{e}_i \cdot h) - F(\mathbf{x}^0 + \mathbf{e}_j \cdot h - \mathbf{e}_i \cdot h)}{2 \cdot h} - \frac{F(\mathbf{x}^0 - \mathbf{e}_j \cdot h + \mathbf{e}_i \cdot h) - F(\mathbf{x}^0 - \mathbf{e}_j \cdot h - \mathbf{e}_i \cdot h)}{2 \cdot h} \right] / (2 \cdot h).$$

Tangents of Tangents





A second derivative code $F^{(1,2)}: \mathbb{R}^n \times \mathbb{R}^n \times \mathbb{R}^n \times \mathbb{R}^n \to \mathbb{R} \times \mathbb{R} \times \mathbb{R} \times \mathbb{R}$, generated in Tangent-over-Tangent (ToT) mode computes

$$\begin{pmatrix} y \\ y^{(2)} \\ y^{(1)} \\ y^{(1,2)} \end{pmatrix} = F^{(1,2)}(\mathbf{x}, \mathbf{x}^{(2)}, \mathbf{x}^{(1)}, \mathbf{x}^{(1,2)}),$$

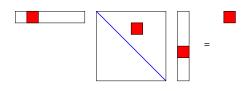
as follows:

$$\begin{pmatrix} y \\ y^{(2)} \\ y^{(1)} \\ y^{(1,2)} \end{pmatrix} := \begin{pmatrix} F(\mathbf{x}) \\ \nabla F(\mathbf{x}) \cdot \mathbf{x}^{(2)} \\ \nabla F(\mathbf{x}) \cdot \mathbf{x}^{(1)} \\ \mathbf{x}^{(1)^T} \cdot \nabla^2 F(\mathbf{x}) \cdot \mathbf{x}^{(2)} + \nabla F(\mathbf{x}) \cdot \mathbf{x}^{(1,2)} \end{pmatrix}.$$

Tangents of Tangents

Accumulation of Hessian





$$\mathbf{x}^{(1)^T} \cdot \nabla^2 F(\mathbf{x}) \cdot \mathbf{x}^{(2)}$$

... accumulation of the whole Hessian element-wise by seeding input directions $\mathbf{x}^{(1)} \in \mathbb{R}^n \ \mathbf{x}^{(2)} \in \mathbb{R}^n$ independently with the Cartesian basis vectors in \mathbb{R}^n for $\mathbf{x}^{(1,2)} = 0$; harvesting from $y^{(1,2)}$.

Note: Approximate Tangents of Tangents

Tangents of Adjoints

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Computer Scientist's View (Simplified)

A second derivative code

$$F_{(1)}^{(2)}: \mathbb{R}^n \times \mathbb{R}^n \times \mathbb{R}^n \times \mathbb{R}^n \times \mathbb{R} \times \mathbb{R} \times \mathbb{R} \to \mathbb{R} \times \mathbb{R} \times \mathbb{R}^n \times \mathbb{R}^n,$$

generated in Tangent-over-Adjoint (ToA) mode computes

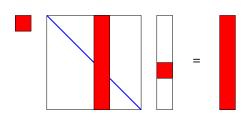
$$\begin{pmatrix} y \\ y^{(2)} \\ \mathbf{x}_{(1)} \\ \mathbf{x}_{(1)}^{(2)} \end{pmatrix} = F_{(1)}^{(2)}(\mathbf{x}, \mathbf{x}^{(2)}, \mathbf{x}_{(1)}, \mathbf{x}_{(1)}^{(2)}, \mathbf{y}_{(1)}, y_{(1)}^{(2)}),$$

as follows:

$$\begin{pmatrix} y \\ y^{(2)} \\ \mathbf{x}_{(1)} \\ \mathbf{x}_{(1)}^{(2)} \end{pmatrix} := \begin{pmatrix} F(\mathbf{x}) \\ \nabla F(\mathbf{x}) \cdot \mathbf{x}^{(2)} \\ \mathbf{x}_{(1)} + \nabla F(\mathbf{x})^T \cdot y_{(1)} \\ \mathbf{x}_{(1)}^{(2)} + y_{(1)} \cdot \nabla^2 F(\mathbf{x}) \cdot \mathbf{x}^{(2)} + \nabla F(\mathbf{x})^T \cdot y_{(1)}^{(2)} \end{pmatrix}$$

Accumulation of Hessian





$$y_{(1)} \cdot \nabla^2 F(\mathbf{x}) \cdot \mathbf{x}^{(2)}$$

... accumulation of the whole Hessian column-wise by seeding input directions $\mathbf{x}^{(2)} \in \mathbb{R}^n$ with the Cartesian basis vectors in \mathbb{R}^n for $\mathbf{x}_{(1)}^{(2)} = 0$, $y_{(1)} = 1$ and $y_{(1)}^{(2)} = 0$; harvesting from $\mathbf{x}_{(1)}^{(2)}$.

Note: Approximate Tangents of Adjoints

Progress



Introduction

First-Order(A)AD

Prerequisites

Tangents

Adjoints

Second-(and Higher-)Order (A)AD

Tangents of Tangents Tangents of Adjoints

"Getting Serious" with AAD

Implementation by Overloading Checkpointing Symbolic Adjoints Adjoint Code Design Patterns

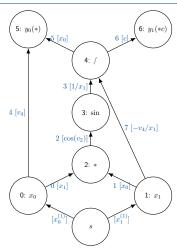


dco/c++ features

- tangents and adjoints of arbitrary order through recursive template instantiation for numerical simulation code implemented in C++
- ▶ front-ends for Fortran, C#, Matlab, Python (3x alpha)
- optimized assignment-level gradient code through expression templates
- cache-optimized internal representation in various incarnations
- vector modes / detection and exploitation of sparsity
- external adjoint / Jacobian interfaces
- user-defined intrinsics
- intrinsic NAG Library functions (e.g. Linear Algebra, Interpolation, Root Finding, Nearest Correlation Matrix)
- support for parallelism: thread-safe data structures, adjoint MPI library, GPU coupling, meta adjoint programming (map)

Tangent IDAG





Tangent IDAG

We consider

$$\begin{pmatrix} y_0 \\ y_1 \end{pmatrix} = \begin{pmatrix} x_0 * \sin(x_0 * x_1) / x_1 \\ \sin(x_0 * x_1) / x_1 * c \end{pmatrix}$$

implemented as

$$t := \sin(x_0 * x_1)/x_1$$

 $y_0 := x_0 * t; y_1 := t * c$

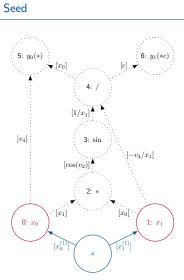
yielding SAC

$$v_2 := x_0 * x_1$$

 $v_3 := \sin(v_2)$
 $v_4 := v_3/x_1$
 $y_0 := x_0 * v_4; y_1 := v_4 * c$

for some *passive* value c, i.e, no derivatives of or with respect to required; \mathbf{x}, \mathbf{y} , and t are *active*.

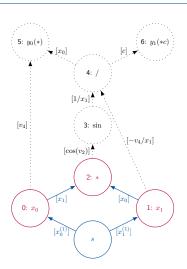




$$x_0 := ?$$
 $x_1 := ?$
 $x_0^{(1)} := ?$
 $x_1^{(1)} := ?$

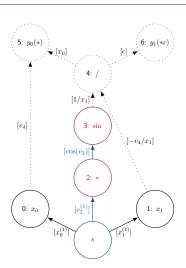
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Propagate (Local Directional Derivatives)



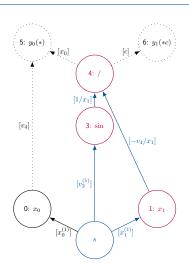
$$\begin{array}{l} v_2 := x_0 * x_1 \\ v_2^{(1)} := x_1 * x_0^{(1)} + x_0 * x_1^{(1)} \end{array}$$





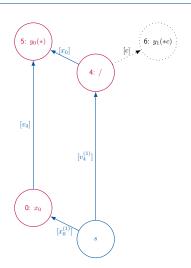
$$v_2 := x_0 * x_1 v_2^{(1)} := x_1 * x_0^{(1)} + x_0 * x_1^{(1)} v_3 := \sin(v_2) v_3^{(1)} := \cos(v_2) * v_2^{(1)}$$





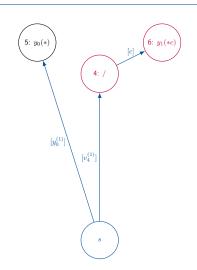
$$\begin{array}{l} v_2 := x_0 * x_1 \\ v_2^{(1)} := x_1 * x_0^{(1)} + x_0 * x_1^{(1)} \\ v_3 := \sin(v_2) \\ v_3^{(1)} := \cos(v_2) * v_2^{(1)} \\ v_4 := v_3/x_1 \\ v_4^{(1)} := (v_3^{(1)} - v_4 * x_1^{(1)})/x_1 \end{array}$$





$$\begin{array}{l} v_2 \coloneqq x_0 * x_1 \\ v_2^{(1)} \coloneqq x_1 * x_0^{(1)} + x_0 * x_1^{(1)} \\ v_3 \coloneqq \sin(v_2) \\ v_3^{(1)} \coloneqq \cos(v_2) * v_2^{(1)} \\ v_4 \coloneqq v_3/x_1 \\ v_4^{(1)} \coloneqq (v_3^{(1)} - v_4 * x_1^{(1)})/x_1 \\ y_0 \coloneqq x_0 * v_4 \\ y_0^{(1)} \coloneqq v_4 * x_0^{(1)} + x_0 * v_4^{(1)} \end{array}$$

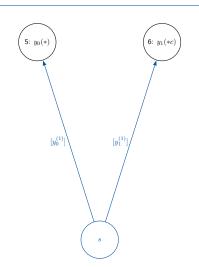




$$\begin{aligned} v_2 &:= x_0 * x_1 \\ v_2^{(1)} &:= x_1 * x_0^{(1)} + x_0 * x_1^{(1)} \\ v_3 &:= \sin(v_2) \\ v_3^{(1)} &:= \cos(v_2) * v_2^{(1)} \\ v_4 &:= v_3/x_1 \\ v_4^{(1)} &:= (v_3^{(1)} - v_4 * x_1^{(1)})/x_1 \\ y_0 &:= x_0 * v_4 \\ y_0^{(1)} &:= v_4 * x_0^{(1)} + x_0 * v_4^{(1)} \\ y_1 &:= v_4 * c \\ y_1^{(1)} &:= c * v_4^{(1)} \end{aligned}$$

Harvest





$$\begin{aligned} v_2 &\coloneqq x_0 * x_1 \\ v_2^{(1)} &\coloneqq x_1 * x_0^{(1)} + x_0 * x_1^{(1)} \\ v_3 &\coloneqq \sin(v_2) \\ v_3^{(1)} &\coloneqq \cos(v_2) * v_2^{(1)} \\ v_4 &\coloneqq v_3/x_1 \\ v_4^{(1)} &\coloneqq (v_3^{(1)} - v_4 * x_1^{(1)})/x_1 \\ y_0 &\coloneqq x_0 * v_4 \\ y_0^{(1)} &\coloneqq v_4 * x_0^{(1)} + x_0 * v_4^{(1)} \\ y_1 &\coloneqq v_4 * c \\ y_1^{(1)} &\coloneqq c * v_4^{(1)} \end{aligned}$$

Tangents with dco/c++

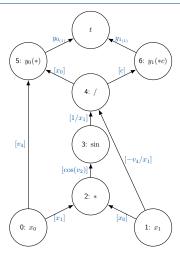


```
User Guide: \mathbf{y}^{(1)} := \nabla F(\mathbf{x}) \cdot \mathbf{x}^{(1)}
```

```
void driver(const vector<double>& xv, double & yv, vector<double> & g) {
    typedef dco::gt1s<double>::type DCO_T;
    size_t n=xv.size();
    DCO_T y=0;
    for (size_t i=0;i<n;i++) {
        vector<DCO_T> x(n,0);
        for (size_t j=0;j<n;j++) x[j]=xv[j];
        dco::derivative(x[i])=1; // seed directions
        f(x,y); // overloaded primal
        g[i]=dco::derivative(y); // harvest directional derivatives
    }
    yv=dco::value(y); // extract function value
}</pre>
```

Adjoint IDAG (Tape)





Adjoint IDAG

We consider

$$\begin{pmatrix} y_0 \\ y_1 \end{pmatrix} = \begin{pmatrix} x_0 * \sin(x_0 * x_1) / x_1 \\ \sin(x_0 * x_1) / x_1 * c \end{pmatrix}$$

implemented as

$$t := \sin(x_0 * x_1)/x_1$$

 $y_0 := x_0 * t$
 $y_1 := t * c$

yielding SAC

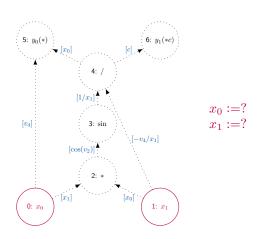
$$v_2 := x_0 * x_1$$

 $v_3 := \sin(v_2)$
 $v_4 := v_3/x_1$
 $y_0 := x_0 * v_4$
 $y_1 := v_4 * c$

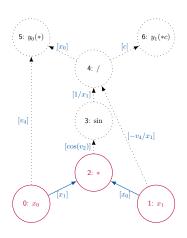
for some passive value c.

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Register (Independent Inputs with Tape)

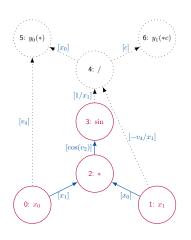






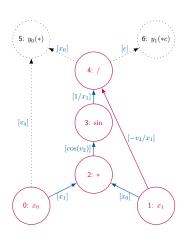
$$v_2 := x_0 * x_1$$





$$v_2 := x_0 * x_1$$
$$v_3 := \sin(v_2)$$



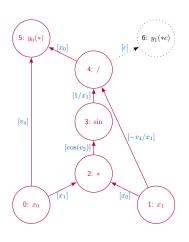


$$v_2 := x_0 * x_1$$

$$v_3 := \sin(v_2)$$

$$v_4 := v_3/x_1$$

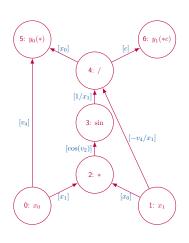




$$v_2 := x_0 * x_1$$

 $v_3 := \sin(v_2)$
 $v_4 := v_3/x_1$
 $y_0 := x_0 * v_4$





$$v_2 := x_0 * x_1$$

$$v_3 := \sin(v_2)$$

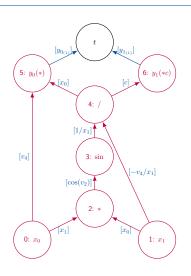
$$v_4 := v_3/x_1$$

$$y_0 := x_0 * v_4$$

$$y_1 := v_4 * c$$

Seed

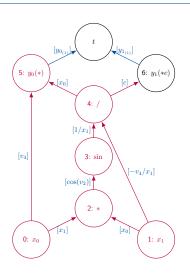




$$\begin{array}{l} v_2 \coloneqq x_0 * x_1 \\ v_3 \coloneqq \sin(v_2) \\ v_4 \coloneqq v_3/x_1 \\ y_0 \coloneqq x_0 * v_4 \\ y_1 \coloneqq v_4 * c \\ y_{0_{(1)}} \coloneqq ? \\ y_{1_{(1)}} \coloneqq ? \\ x_{0_{(1)}} \coloneqq ? \\ x_{1_{(1)}} \coloneqq ? \\ v_{2_{(1)}} \coloneqq 0 \\ v_{3_{(1)}} \coloneqq 0 \\ v_{4_{(1)}} \coloneqq 0 \end{array}$$

Interpret (Tape)



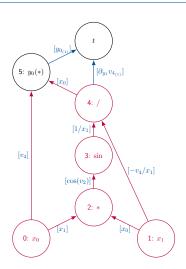


$$\begin{array}{l} v_2 := x_0 * x_1 \\ v_3 := \sin(v_2) \\ v_4 := v_3/x_1 \\ y_0 := x_0 * v_4 \\ y_1 := v_4 * c \\ v_{4_{(1)}} + = c * y_{1_{(1)}} \end{array}$$

Note C++ Syntax:

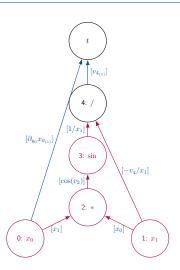
$$\begin{aligned} v_{4_{(1)}} + &= c * y_{1_{(1)}} \\ \Leftrightarrow \\ v_{4_{(1)}} &:= v_{4_{(1)}} + c * y_{1_{(1)}}. \end{aligned}$$





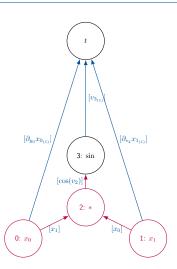
$$\begin{split} v_2 &\coloneqq x_0 * x_1 \\ v_3 &\coloneqq \sin(v_2) \\ v_4 &\coloneqq v_3 / x_1 \\ y_0 &\coloneqq x_0 * v_4 \\ y_1 &\coloneqq v_4 * c \\ v_{4_{(1)}} + &= c * y_{1_{(1)}} \\ v_{4_{(1)}} + &= x_0 * y_{0_{(1)}} \\ x_{0_{(1)}} + &= v_4 * y_{0_{(1)}} \end{split}$$

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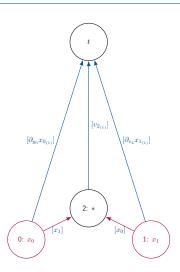
$$\begin{aligned} v_2 &:= x_0 * x_1 \\ v_3 &:= \sin(v_2) \\ v_4 &:= v_3/x_1 \\ y_0 &:= x_0 * v_4 \\ y_1 &:= v_4 * c \\ v_{4(1)} + &= c * y_{1(1)} \\ v_{4(1)} + &= x_0 * y_{0(1)} \\ x_{0(1)} + &= v_4 * y_{0(1)} \\ u &:= 1/x_1 \\ v_{3(1)} + &= u * v_{4(1)} \\ x_{1(1)} - &= v_4 * u * v_{4(1)} \end{aligned}$$





$$\begin{aligned} v_2 &:= x_0 * x_1 \\ v_3 &:= \sin(v_2) \\ v_4 &:= v_3/x_1 \\ y_0 &:= x_0 * v_4 \\ y_1 &:= v_4 * c \\ v_{4(1)} + &= c * y_{1(1)} \\ v_{4(1)} + &= x_0 * y_{0(1)} \\ x_{0(1)} + &= v_4 * y_{0(1)} \\ u &:= 1/x_1 \\ v_{3(1)} + &= u * v_{4(1)} \\ x_{1(1)} - &= v_4 * u * v_{4(1)} \\ v_{2(1)} + &= \cos(x_2) * v_{3(1)} \end{aligned}$$

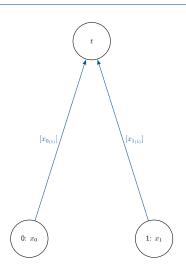




$$\begin{aligned} v_2 &:= x_0 * x_1 \\ v_3 &:= \sin(v_2) \\ v_4 &:= v_3/x_1 \\ y_0 &:= x_0 * v_4 \\ y_1 &:= v_4 * c \\ v_{4(1)} + &= c * y_{1(1)} \\ v_{4(1)} + &= x_0 * y_{0(1)} \\ x_{0(1)} + &= v_4 * y_{0(1)} \\ u &:= 1/x_1 \\ v_{3(1)} + &= u * v_{4(1)} \\ x_{1(1)} - &= v_4 * u * v_{4(1)} \\ v_{2(1)} + &= \cos(x_2) * v_{3(1)} \\ x_{0(1)} + &= x_1 * v_{2(1)} \\ x_{1(1)} + &= x_0 * v_{2(1)} \end{aligned}$$

Harvest





$$\begin{split} v_2 &\coloneqq x_0 * x_1 \\ v_3 &\coloneqq \sin(v_2) \\ v_4 &\coloneqq v_3/x_1 \\ y_0 &\coloneqq x_0 * v_4 \\ y_1 &\coloneqq v_4 * c \\ v_{4(1)} + &= c * y_{1(1)} \\ v_{4(1)} + &= x_0 * y_{0(1)} \\ x_{0(1)} + &= v_4 * y_{0(1)} \\ u &\coloneqq 1/x_1 \\ v_{3(1)} + &= u * v_{4(1)} \\ x_{1(1)} - &= v_4 * u * v_{4(1)} \\ v_{2(1)} + &= \cos(x_2) * v_{3(1)} \\ x_{0(1)} + &= x_1 * v_{2(1)} \\ x_{1(1)} + &= x_0 * v_{2(1)} \end{split}$$

Adjoints with dco/c++



Driver

5

6

10

11

12

13

14

15 16

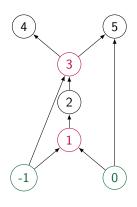
17

```
User Guide: \mathbf{x}_{(1)} := \nabla F(\mathbf{x})^T \cdot \mathbf{y}_{(1)}
void driver(const vector<double> &xv, double &yv, vector<double> &g) {
  typedef dco::ga1s<double> DCO_M; // dco mode
  typedef DCO_M::type DCO_T; // dco type
  typedef DCO_M::tape_t DCO_TAPE_T; /dco tape type
  size_t n=xv.size():
  vector<DCO_T> x(n); DCO_T y;
  DCO_M::global_tape=DCO_TAPE_T::create(); // tape creation
  for (size_t i=0; i< n; i++)  { // independent tape entries
    x[i]=xv[i]; DCO_M::global_tape—>register_variable(x[i]);
  f(x,y); // overloaded primal
  DCO_M::global_tape—>register_output_variable(y); // dependent tape entry
  yv=dco::value(y); dco::derivative(y)=1; // seed
  DCO_M::global_tape—>interpret_adjoint(); // tape interpretation
```

for $(size_t i=0; i< n; i++) \{ g[i]=dco::derivative(x[i]); \} // harvest$

DCO_TAPE_T::remove(DCO_M::global_tape); // release tape





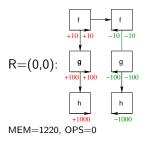
 v_5, v_4, \ldots, v_{-1}

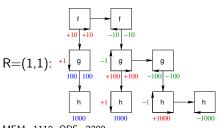
- ► U.N.: DAG REVERSAL is NP-Complete, J. Disc. Alg. 7(4), 402-410 (2009).
- ► U.N.: CALL TREE REVERSAL is NP-Complete, LNCSE 64, 13-22 (2008).
- ▶ J. Lotz et al.: Mixed Integer Programming for Call Tree Reversal, SIAM CSC (2016).

CALL TREE REVERSAL



Example: Let $\overline{\text{MEM}} = 1110 \dots$



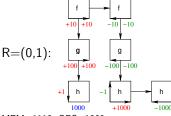


CALL TREE REVERSAL

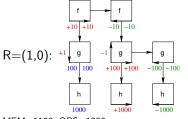


Example: Let $\overline{\text{MEM}} = 1110$... Greedy Heuristics

Smallest Memory Increase starts from R=1 and yields ... Largest Memory Decrease (LMD) starts from R=0 and yields ...



MEM=1110, OPS=1000



Largest Memory Increase (LMI) remains at R = 1 as R = (1,0) infeasible

(Embedded) Symbolic Adjoints Story



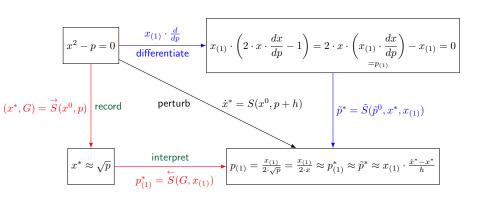
Algorithmic Differentiation (AD) is based on Symbolic Differentiation (SD). AD approaches vary in terms of choice of SD level.

- ▶ U.N. et al.: Algorithmic differentiation of numerical methods: Tangent and adjoint solvers for parameterized systems of linear equations, RWTH Technical Report AIB-2012-10 (2012).
- U.N. et al.: Algorithmic differentiation of numerical methods: Tangent and adjoint solvers for parameterized systems of nonlinear equations, ACM TOMS 41 (4), 26 (2015).
- ▶ N. Safiran et al.: Algorithmic Differentiation of Numerical Methods: Second-Order Adjoint Solvers for Parameterized Systems of Nonlinear Equations, Procedia Computer Science 80, 2231-2235 (2016).
- ▶ J. Lotz et al.: A Case Study in Adjoint Sensitivity Analysis of Parameter Calibration, Procedia Computer Science 80, 201-211 (2016).

(Embedded) Symbolic Adjoints

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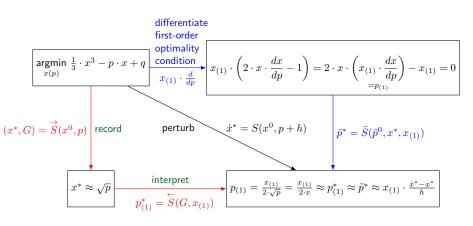
Case Study: Root Finding



(Embedded) Symbolic Adjoints

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Case Study: Optimization



Adjoint Code Design Patterns

Case Study: Ensemble of Evolutions



