An Abstract Application Layer Interface to Transport Services draft-trammell-taps-interface-01

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Interface Design Principles (§3) (a review)

We set out to define a *single interface* to a variety of transport protocols to be used in a variety of application design patterns, to enable applications written to a single API to make use of multiple transport protocols in terms of the features they provide, providing:

- explicit support for security properties as first-order transport features;
- asynchronous connection, transmission, and reception;
- support for multistreaming and multipath transport protocols; and
- atomic transmission of data, using application-assisted framing and deframing where necessary.

Interface Diagram (as of -01)

```
Properties (related to Send() properties)
Require() Prefer() Ignore()
                                            Avoid()
                                                           Prohibit()
Security parameters
                                        Connection
Preconnection
              Initiate() → Ready<>
                                        Clone()
                                                   → Connection Group
              Listen() → CReceived<>
Clone()
              Rendezvous() → RDone<>
                                        Send(MCtx, EOM) →
              Stop() → Stopped<>
                                                Sent<>, Expired<>
                                        Receive() →
                                                Received<Data/Metadata>
   Endpoints
                                        Close() → Closed<>
     Local
                                        Abort() → Aborted<>
   Remote
```

(non-editorial) changes since -00

- #201 Restructure Transport Properties
- #200 Rework Partial Sends and Receives
- #198 Message Receive Metadata
- #195 Ordering of API Events
- #181 Rework Interface Types
- #171 Batching Sends

#201 Transport Parameters Rework

- All of the various ways to configure stacks (pre-selection, post-selection, and per-send) are related, but were spread throughout the document
- New approach: group all (non-security) parameters into into Properties (new §12), attempt to reclassify them.
 - Note: the authors do not think we have this right yet, but we do think it's less intentionally confusing than it was.
 - Definitely needs reordering (order is kind of random)
 - May need new / renamed axes / classifications.
- Preferences still expressed using Require(), Prefer(), Avoid(), Prohibit(); send properties are bound to MessageContext passed on Send().

#201: current property "axes"

Data type

Boolean / Enumeration / Integer / Preference

Scope

Preconnection / Connection / Message

Classification

Affected Aspects		Path & Protocol Selection	Protocol Operation	Control Flow	
je	Immediate	Selection Property	Protocol Property	Control Property	
Level of Abstraction	Interpreted				

Properties (1/3)

protocol/control prop.

	Type	Dep.	Select	post- Select	per- Send
12.3.1. Final	bool				√/?
12.3.2. Reliable Message Transfer	pref		✓		
12.3.3. Configure Reliability Per Message	pref		\checkmark		
12.3.4. Reliable Transfer (Message)	bool	†			✓
12.3.5. Preservation of Data Ordering	pref		\checkmark		
12.3.6. Ordered	bool	†			✓
12.3.7. Direction of communication	enum		?	?	
12.3.8. 0-RTT Establishment w/ldem.	pref		✓		
12.3.9. Idempotent	bool	†			\checkmark
12.3.10. Multistream in Group	pref		✓		
12.3.11. Excessive RTX Notification	pref		\checkmark		
12.3.12. Exc. RTX Notification Threshold	int	†	√	✓	

Properties (2/3)

protocol/control prop.

	Туре	Dep.	Select	post- Select	per- Send
12.3.13. Soft Error Notification	pref		✓		
12.3.14. Checksum Coverage Control	pref		✓		
12.3.15. Checksum Coverage Length	int	†			✓
12.3.16. Recv Checksum Requirement	int		√	✓	
12.3.17. Interface Instance / Type	(enum,pref)		✓		
12.3.18. PvD Instance / Type	(enum,pref)		✓		
12.3.19. Capacity Profile (intent)	enum		✓	\checkmark	✓
12.3.20. Congestion Control	pref		✓		
12.3.21. Niceness	int			\checkmark	✓
12.3.22. Abort Timeout	int		✓		
12.3.23. Connection Group TX Scheduler	enum		\checkmark	✓	

Properties (3/3)

protocol/control prop.

	Туре	Dep.	Select	post- Select	per- Send
12.3.24. Max Idempotent Send Size	int			r/o	
12.3.25. Max No-Frag Send Size	int			r/o	
12.3.26. Max (non-partial?) Send Size	int			r/o	
12.3.27. Max (non-partial?) Recv Size	int			r/o	
12.3.28. PR Send Lifetime	int	12.3.3.			✓

Some Observations from the Editor (+discussion)

- Calling these axes is a little misleading: they're not orthogonal
- We have only six distinct kinds of thing:
 - Preference used for selection, scoped to preconnection, read-only after connection.
 - Property used to control how messages are sent, scoped to message (boolean or integer, usually linked to selection preference).
 - Property used to control protocol operation, scoped to preconnection + connection (usually integer, e.g. sizes/timeouts), possibly also usable for selection.
 - Property used to inspect protocol operation, scoped to connection, read-only (usually integer, e.g. buffer size).
 - Enumeration/preference tuples for selecting interface/PvD.
 - Intents, which can influence selection, configuration, scheduling, etc. at a higher level.

#200 Partial Send and Receive

- API is organized around atomic write/read of messages
 - (using application-supplied deframing when the underlying transport doesn't do framing, see §8.4)
- But sometimes you have a message (or a real stream) that won't fit into a buffer.
- Solution: partial read/write
 - Introduce optional EOM parameter to Send(); calls with EOM =
 false → still writing to a partial message identified by a given
 MessageContext.
 - ReceivedPartial<> event fires when a partial message is received.
- Partial read/write boundaries are not preserved.

Open issue: API for idempotent Send on establishment (#112 / #124)

- How does the application tell the stack that it wants to send some 0RTT data?
 - Some tradeoffs here, but mainly a bikeshed.
- Option 1: as in #124, hold any data sent until an explicit Connection. Start() call.
 - Send() before Start() is ORTT if idempotent.
 - Start() is always required, even if you don't know what ORTT is.
- Option 3: 0RTT behavior is implied by 0RTT selection properties.
 - When Initiate() is called and selects a ORTT-capable stack, the actual initiation is delayed slightly to wait for the first Send(), which is ORTT if idempotent.
 - Note this makes racing 0RTT-capable and 0RTT-incapable stacks impossible.
- Option 3.5: as 3, but with a Preconnection.InitiateNow() to override the wait-for-Send() behavior (e.g. for application protocols where the server sends first)
- Option 5: Add Preconnection.Send(), which initiates with ORTT data.

Next steps

There are still some open issues: github.com/taps-api/drafts/issues

