

If we run the code the memory storage in RAM.

If we save the code the memory storage in Hddisk.

ALU → +, -, *, ÷, And, not, OR
 ADDER, MULTIPLEIR, DIV → Logic Gates

CU → Decoder Max

Number System

$(1011001011)_2 \rightarrow (\quad)_8 \quad (\quad)_{16}$

\Rightarrow

$(1313)_8$

\Rightarrow

$(2AB)_{16}$

$$\begin{array}{l} \text{Base 9} \\ 10^3 \\ \text{Base 2} \\ 10, 13 \\ \text{Base 4} \\ 10, 12, 3 \\ \text{Base 8} \\ 10^{-7} \\ \text{Base 10} \\ 10^{-9}, A \\ \text{Base 11} \\ 10^{-9}, A \\ \text{Base 16} \\ 10^{-9}, A, B, C, D, E, F \end{array}$$

$$\begin{array}{l} b_1 \rightarrow b_2 \\ \text{if } \exp_{\text{base}} \text{ in } 2 \\ b_1 \rightarrow \text{binary} \\ \text{binary} \rightarrow b_2 \\ b_1 \rightarrow 10 \\ 10 \rightarrow b_2 \\ \text{not express in 2} \\ b_1 \rightarrow b_2 \\ \text{not express in 2} \end{array}$$

$$\begin{array}{r} 64 \mid 65 \\ 1 \quad | \quad 1 \\ 0 \quad | \quad 1 \\ \hline (65)_{10} \rightarrow (101)_6 \\ (11)_2 \rightarrow (101)_2 \\ (100001)_2 \end{array}$$

$$\begin{array}{r} (1056)_{16} \rightarrow (1056)_2 \\ (1000101110)_2 \rightarrow (00010000101010)_2 \\ (1056)_2 \rightarrow (1056)_10 \\ 1056 = 1 \times 10^3 + 0 \times 10^2 + 5 \times 10^1 + 6 \times 10^0 \\ 1056 = 1000 + 0 + 50 + 6 \\ 1056 = 1000 + 50 + 6 \\ 1056 = 1056 \end{array}$$

$$b_1 \rightarrow b_2$$

$m_i < x_i$

$$b_1 \rightarrow \text{binary}$$

$$\text{binary} \rightarrow b_2$$

$$\text{Decimal to binary convert}$$

$$\begin{array}{r} 2^7 \\ 2 \mid 3 \\ 1 \quad | \quad 1 \\ 2 \quad | \quad 1 \\ 0 \quad | \quad 1 \\ \hline (3)_{10} \rightarrow (11)_2 \end{array}$$

$$\begin{array}{r} 2^1 \\ 2 \mid 1 \\ 1 \quad | \quad 1 \\ 2 \quad | \quad 1 \\ 0 \quad | \quad 1 \\ \hline (1)_{10} \rightarrow (1)_2 \end{array}$$

$$\begin{array}{r} 2^4 \\ 2 \mid 1 \\ 1 \quad | \quad 1 \\ 2 \quad | \quad 1 \\ 0 \quad | \quad 1 \\ \hline (1)_{10} \rightarrow (1)_2 \end{array}$$

$$\begin{array}{r} 2^3 \\ 2 \mid 1 \\ 1 \quad | \quad 1 \\ 2 \quad | \quad 1 \\ 0 \quad | \quad 1 \\ \hline (1)_{10} \rightarrow (1)_2 \end{array}$$

$$\begin{array}{r} 2^2 \\ 2 \mid 1 \\ 1 \quad | \quad 1 \\ 2 \quad | \quad 1 \\ 0 \quad | \quad 1 \\ \hline (1)_{10} \rightarrow (1)_2 \end{array}$$

$$\begin{array}{r} 2^1 \\ 2 \mid 1 \\ 1 \quad | \quad 1 \\ 2 \quad | \quad 1 \\ 0 \quad | \quad 1 \\ \hline (1)_{10} \rightarrow (1)_2 \end{array}$$

$$\begin{array}{r} 2^0 \\ 2 \mid 1 \\ 1 \quad | \quad 1 \\ 2 \quad | \quad 1 \\ 0 \quad | \quad 1 \\ \hline (1)_{10} \rightarrow (1)_2 \end{array}$$

$$\begin{array}{r} 2^7 \\ 2 \mid 3 \\ 1 \quad | \quad 1 \\ 2 \quad | \quad 1 \\ 0 \quad | \quad 1 \\ \hline (3)_{10} \rightarrow (11)_2 \end{array}$$

$$\begin{array}{r} 2^3 \\ 2 \mid 1 \\ 1 \quad | \quad 1 \\ 2 \quad | \quad 1 \\ 0 \quad | \quad 1 \\ \hline (1)_{10} \rightarrow (1)_2 \end{array}$$

$$\begin{array}{l} \text{What is the minimum number of bits required} \\ \text{to represent } (6728)_{10} \text{ in binary.} \\ (-\frac{\log m}{\log 2})n \\ (m-1)m^{n-1} + (m-1)^m = (m-1)m + (m-1)^m \\ (m-1)(1+m+m^2+\dots+m^n) \\ (m-1)\left(\frac{m^n-1}{m-1}\right) = m^n-1 \end{array}$$

$$\begin{array}{l} m^n - 1 \geq 6728 \\ m^n \geq 6729 \\ n = \lceil \frac{\log 6729}{\log 2} \rceil \\ n = 13 \end{array}$$

$$(ABC)_n$$

$$A \times n^2 + B \times n^1 + C \times n^0$$

$$1000 \quad 10 \quad 8$$

$$1001 \quad 11 \quad 9$$

$$1010 \quad 12 \quad A$$

$$(123)_5 = (m8)_4 \quad \text{Find all possible values of } m \text{ & } y$$

$$1 \times 5^2 + 8 \times 5^1 + 3 \times 5^0 = m8^1 + 8 \cdot 4^0$$

$$0.5 + 10 + 3 = my + 8$$

$$\Rightarrow 38 = my + 8$$

$$n < y \\ 8 < y$$

$$m > 30$$

$$(m, y) \in$$

$$(1, 30)$$

$$(2, 30)$$

$$(3, 10)$$

$$\dots$$

$$(1000001)$$

$$m > 30 \\ m < 120 \\ \therefore m < 30$$

$$(1, 30)$$

$$(2, 30)$$

$$(3, 10)$$

$$\dots$$

Q. How many value of m & y are possible for $(m, y) \in (m, y)$

$$4x9^1 + 2x9^0 = my^1 + 3x4^0$$

$$36 + 8 = my + 3$$

$$35 = my$$

$$\text{or } (88+3) \sqrt{21}$$

$$35 = ny$$

$$\text{or } (88+3) \sqrt{21}$$

$$m^2 + m - 10 = 0$$

$$\left(\frac{1-\sqrt{5}}{2}, \frac{1+\sqrt{5}}{2} \right) \cap \mathbb{Z}$$

In any number system, the range of numbers are -3 to 3 represented as $C, B, A, 0, 1, 2, 3$ express $(102)_{10}$ in the number system base $\rightarrow 7$

$$(2003)_{10} = (16342)_{7}$$

$$(11100000000000011)_{10} =$$

$$1024$$

$$4096$$

$$16384$$

$$7168$$

$$4096$$

$$2048$$

$$1024$$

$$512$$

$$256$$

$$128$$

$$64$$

$$32$$

$$16$$

$$8$$

$$4$$

$$2$$

$$1$$

$$0$$

minimum no. of bits
required

$$(516)_3 = \left\{ \begin{array}{l} \\ \\ \end{array} \right.$$

$$n^{th} > 512$$

$$n^{th} > 512$$

$$4^n > 512$$

$$4^n > 512$$

$$\log_{10} 2^n > \log_{10} 512$$

$$2^n \log_{10} 2 > \log_{10} 512$$

$$2^n > \log_{10} 512$$

Base - b

complement

1's complement

n → (b-1)^s complement of n = (b^n - 1) - n

b's complement

2's complement

"

"

Base - b

complement

1's complement

n → (b-1)^s complement of n = (b^n - 1) - n

b's complement

2's complement

"

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Base - b

complement

1's complement

n → (b-1)^s complement of n = (b^n - 1) - n

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Base - b

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Base - b

complement

1's complement

n → (b-1)^s complement of n = (b^n - 1) - n

b's complement

2's complement

"

"

$$\text{Base - 2} \quad 1's \rightarrow (2^{n-1}) - n$$

$$80^n = 10^n \cdot 2^n \rightarrow 10^n - 2^n = 2^n - n$$

$$2's \rightarrow 3^n - 1 - n$$

$$3's \rightarrow 3^n - n$$

$$4's \rightarrow 5^n - 1 - n$$

$$5's \rightarrow 6^n - 1 - n$$

$$6's \rightarrow 7^n - 1 - n$$

$$7's \rightarrow 8^n - 1 - n$$

$$8's \rightarrow 9^n - 1 - n$$

$$9's \rightarrow 10^n - 1 - n$$

$$10's \rightarrow 11^n - 1 - n$$

$$11's \rightarrow 12^n - 1 - n$$

$$12's \rightarrow 13^n - 1 - n$$

$$13's \rightarrow 14^n - 1 - n$$

$$14's \rightarrow 15^n - 1 - n$$

$$15's \rightarrow 16^n - 1 - n$$

$$16's \rightarrow 17^n - 1 - n$$

$$17's \rightarrow 18^n - 1 - n$$

$$18's \rightarrow 19^n - 1 - n$$

$$19's \rightarrow 20^n - 1 - n$$

$$20's \rightarrow 21^n - 1 - n$$

$$21's \rightarrow 22^n - 1 - n$$

$$22's \rightarrow 23^n - 1 - n$$

$$23's \rightarrow 24^n - 1 - n$$

$$24's \rightarrow 25^n - 1 - n$$

$$25's \rightarrow 26^n - 1 - n$$

$$26's \rightarrow 27^n - 1 - n$$

$$27's \rightarrow 28^n - 1 - n$$

$$28's \rightarrow 29^n - 1 - n$$

$$29's \rightarrow 30^n - 1 - n$$

$$30's \rightarrow 31^n - 1 - n$$

$$31's \rightarrow 32^n - 1 - n$$

$$32's \rightarrow 33^n - 1 - n$$

$$33's \rightarrow 34^n - 1 - n$$

$$34's \rightarrow 35^n - 1 - n$$

$$35's \rightarrow 36^n - 1 - n$$

$$36's \rightarrow 37^n - 1 - n$$

$$37's \rightarrow 38^n - 1 - n$$

$$38's \rightarrow 39^n - 1 - n$$

$$39's \rightarrow 40^n - 1 - n$$

$$40's \rightarrow 41^n - 1 - n$$

$$41's \rightarrow 42^n - 1 - n$$

$$(74A)_{14} \Rightarrow PPP = 74A = 9C8$$

~~9C9~~

15's
16's

complement of $a = -\frac{a}{\bar{a}}$

$$x + \bar{y} = x + (b^{n-1}) - m$$

$$= b^{n-1} \rightarrow 0$$

$\underbrace{1111 \dots}_{n \text{ times}} - \underbrace{2222 \dots}_{n \text{ times}} + \text{two representation of zero}$

Base-3

$$\begin{array}{r} 222 \\ 6000 \\ \hline 0 \end{array}$$

Base-4

$$\begin{array}{r} 3333 \\ 0000 \\ \hline 0 \end{array}$$

Base-5

$$\begin{array}{r} 23-97 \\ 23+(99-97) \\ \hline 23 \end{array}$$

$$\begin{array}{r} 23-97 \\ 23+(99-97) \\ \hline 23 \end{array}$$

~~#~~

$$x + \bar{m} = 0$$

~~16~~

$$(21-4)^2 > 0$$

↳ remove carry

ADD 1

4F

$$\begin{array}{r} 101 \\ 100 \\ \hline 01 \end{array}$$

color

94 - 23

$$(41115)^{15} \times 2^3$$

$$(41115)_{15} + 23 \Rightarrow 221 - 23$$

$$\boxed{124}$$

$$- \frac{99}{73}$$

$$97 + (99-23)$$

$$(97-23) + (100-1)$$

$$74 \text{ Ans} - 0$$

$$\frac{+1}{74}$$

$$23-97$$

$$23+(99-97)$$

$$23-97$$

$$\Rightarrow 09-97$$

$$\Rightarrow 2$$

$$23$$

$$23$$

$$23$$

$$23$$

$$23$$

$$23$$

$$23$$

$$23$$

$$23$$

$$23$$

$$23$$

$$23$$

$$23$$

Q-

5-4

4-5

$$23 - 97$$

$$n - k \\ n + (b^{n-1} - 1) \rightarrow \text{it's contra}$$

$$99 \\ 97 \\ - 97 \\ \hline 02$$

$$+ 23 \\ \hline 25$$

$$15 \\ 15 \\ - 15 \\ \hline 0$$

$$- 74$$

$$\frac{1000}{100} + 1$$

$$100 \\ 100 \\ - 100 \\ \hline 0$$

$$110 \\ 110 \\ - 100 \\ \hline 10$$

$$10 \\ 10 \\ - 10 \\ \hline 0$$

$$0$$

$$99 \\ 25 \\ - 25 \\ \hline 0$$

$$15 \\ 15 \\ - 15 \\ \hline 0$$

$$0$$

$$2's$$

$$101 + 011$$

$$100 + 011$$

$$111$$

$$2's$$

$$0$$

Using j's
result
↓ - Appendix

$$97 - 23 \\ 97 + (100 - 23)$$

$$97 - 23 + 100$$

$$97 + 77$$

$$174$$

$$99 \\ 97 \\ - 97 \\ \hline 02$$

$$+ 23 \\ \hline 25$$

$$15 \\ 15 \\ - 15 \\ \hline 0$$

$$0$$

$$(97-23) \\ \frac{97}{25} \\ \frac{77}{174}$$

$$174$$

$$97 \\ 97 \\ - 97 \\ \hline 02$$

$$+ 23 \\ \hline 25$$

$$15 \\ 15 \\ - 15 \\ \hline 0$$

$$0$$

$$99 \\ 25 \\ - 25 \\ \hline 0$$

$$15 \\ 15 \\ - 15 \\ \hline 0$$

$$0$$

$$99 \\ 25 \\ - 25 \\ \hline 0$$

$$15 \\ 15 \\ - 15 \\ \hline 0$$

$$0$$

$$99 \\ 25 \\ - 25 \\ \hline 0$$

$$15 \\ 15 \\ - 15 \\ \hline 0$$

$$0$$

$$99 \\ 25 \\ - 25 \\ \hline 0$$

$$15 \\ 15 \\ - 15 \\ \hline 0$$

$$0$$

$$99 \\ 25 \\ - 25 \\ \hline 0$$

$$15 \\ 15 \\ - 15 \\ \hline 0$$

$$0$$

$$99 \\ 25 \\ - 25 \\ \hline 0$$

$$15 \\ 15 \\ - 15 \\ \hline 0$$

$$0$$

$$99 \\ 25 \\ - 25 \\ \hline 0$$

$$15 \\ 15 \\ - 15 \\ \hline 0$$

$$0$$

$m = 0$

4-3

Hamming Code

Hamming Distance

position v_2
match na ho

code word

$$\begin{array}{r} 101101 \\ 100110 \\ \hline HD = 3 \end{array}$$

$\times \times \times$

Code word &
error detected

is detected

$$\begin{array}{c} 101001 \\ 100 \\ \hline 101000 \end{array}$$

$HD = 2$

$$\boxed{101010}$$

$HD = 3$

$$\boxed{101110}$$

To detect d bit error minimum Hamming distance
should be $d+1$

To correct d bit error minimum H.D. should
be $2d+1$

Hamming Code

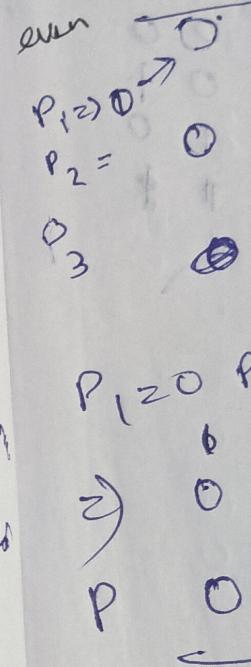
To correct 1 bit error

$$m+p = m+p+1$$

$$2^p \geq (m+p+1)$$

$$4 \text{ bit } + 3 \text{ bit } = (0, 1, 2, 3, 4, 5, 6, 7)$$

$$\Rightarrow 2^3 = 8$$



$$m = \underline{0101} \quad 4 \text{ Bit} + 3 \text{ Bit} \Rightarrow 2^3 \text{ to } 7$$

$$\begin{matrix} 2^3 \\ P_1 \swarrow P_2 \searrow P_3 \end{matrix} \geq 4+3+1$$

even

$$P_1 = 0 \quad \frac{P_1 \ P_2 \ 0 \ P_3 \ 1 \ 0 \ 1}{0 \ 0 \ 1 \ 1}$$

$$P_2 = 1$$

$$P_3 =$$

$$\begin{array}{r} 0100101 \\ \hline \text{even} \rightarrow 0100001 \end{array}$$

P ₃	P ₂	P ₁
0 0	0	0
1 0	0	1
2 0	1	0
3 0	1	1
4 1	0	0
5 1	0	1
6 1	1	0
7 1	1	1

$$\begin{aligned} P_1 &\Rightarrow \{1, 3, 5, 7\} \\ P_2 &\Rightarrow \{2, 3, 6, 7\} \\ P_3 &\Rightarrow \{4, 5, 6, 7\} \end{aligned}$$

$$P_3 \ P_2 \ P_1$$

$$101 \Rightarrow 4+0+1 \Rightarrow 5^{\text{th}}$$

Ans

$$P_1 \ P_2 \ P_3$$

$$001 \Rightarrow 2^{\text{st}} \text{ position}$$

$$Q - 1010$$

$$\underline{1, 2, 4, 8, 16} \quad 2^P \geq m + p + 1$$

$$2^3 \leq 2, 0, 1, 1, 2, 3, 4, 5$$

$$1011010$$

$$\text{even} \quad \underline{1011110}$$

$$\begin{array}{r} 1 2 3 4 5 6 7 \\ P_1 \ P_2 \ 1 \ P_3 \ 0 \ 1 \ 0 \\ \hline \textcircled{5} \ 1 \ 0 \ 0 \end{array}$$

$$P_1 = 1$$

$$P_2 = 0$$

$$P_3 = 1$$

$$\begin{array}{r} P_3 \ P_2 \ P_1 \\ 1 \ 0 \ 1 \end{array} \Rightarrow 5$$

Boolean Algebra

operator
Operand
Real world problem
ckt
Solve minies

Real world problem \rightarrow Boolean expression \rightarrow ckt

A	B	C	F
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1

1st \rightarrow A \rightarrow 110
2nd \rightarrow B \rightarrow 110
3rd \rightarrow C \rightarrow 110

A	B	C	F
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1

10 \rightarrow complemented

$$f = \bar{A}\bar{B}C + \bar{A}B\bar{C} + A\bar{B}\bar{C}$$

$$x + \bar{x} = 1$$

$$x \cdot \bar{x} = 0$$

- Distributivity — Commutativity —
 $x + y = y + x$
 $xy = yx$
- Distributivity — Left Association
 $x + (y + z) = (x + y) + z$
Right Association
 $(xy)z = x(yz)$

Basic Properties

$$x + x = x$$

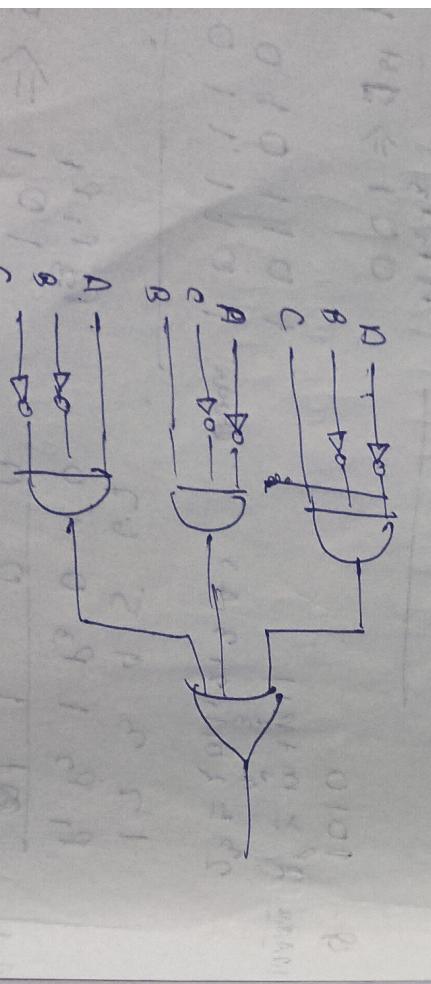
$$x \cdot x = x$$

$$x + 1 = x$$

$$x \cdot 1 = x$$

$$x + 0 = x$$

$$x \cdot 0 = 0$$



$$x + yz = (x+y)(x+z)$$

$$x^2 = x$$

$$x \cdot 1 = x$$

Every Boolean expression must satisfy -

-) complementation $\Rightarrow \bar{m} + \bar{y} = \bar{m' + y'}$
-) duality $\Rightarrow m'y \rightarrow \bar{m' + y}$ have dual and -
-) Every Boolean Exp. have dual and some property

dual -

$$f(m_1, m_2, m_3, \dots, m_n, 0, 1, +, \cdot)$$

$$\text{dual}(m_1, m_2, m_3, \dots, m_n, 1, 0, \cdot, +)$$

$$f = m + y \quad f = m + 1$$

$$f_{\text{dual}} = m(y + 1) = m \cdot 0 = 0$$

$$m'y'z + yz + m_3 z (1 + n) = (m'y')w$$

$$m'y'z + 3(y + n)$$

$$2(m'y + (y + n))$$

$$2((m' + y)(y + y') + m)$$

$$2(m' + y + n)$$

$$2(1 + y) = 2 + 2y$$

Z

Demorgan's law

$$\overline{m \cdot y} = m' + y'$$

$$\overline{m \cdot y} = m' + y'$$

$$f_{\text{compl}} = (\bar{m}_1, \bar{m}_2, \dots, \bar{m}_n, 1, 0, 1, 1, +, \cdot)$$

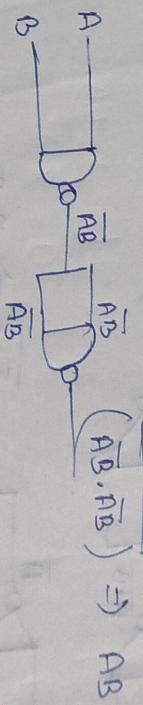
$$f_{\text{compl}} = \bar{m}' \cdot \bar{y}' \cdot \bar{o}' \cdot \bar{a}' \quad \text{and } \bar{m} \cdot \bar{y}$$

$$f_{\text{dual}} = my$$

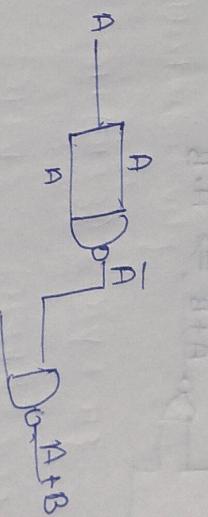
Universal Gate -

Using NAND

AND



OR

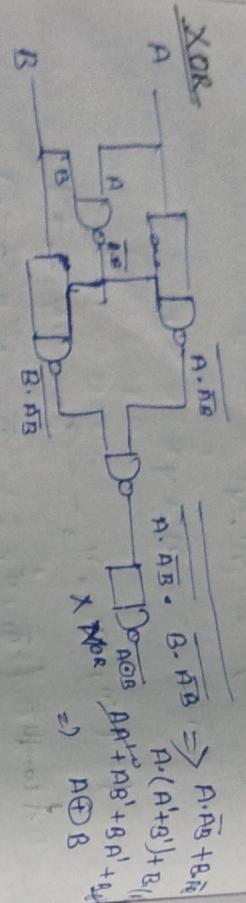


$$\overline{A} \cdot \overline{B} = \overline{A+B}$$

\overline{A}

NOT $A \rightarrow D \rightarrow \overline{A}$

XOR



NAND

$\overline{A \cdot B}$

$\Rightarrow A' + B'$

dual

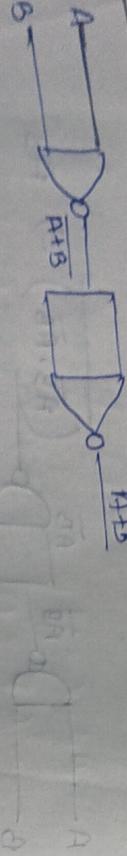
$A'B'$

using NOR
dual

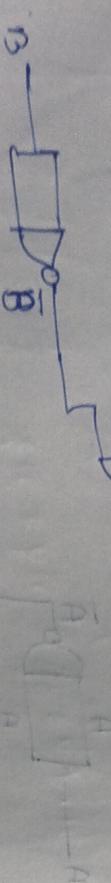
using NOR

$A + B = \overline{\overline{A} \cdot \overline{B}}$

OR $A \rightarrow D \rightarrow \overline{A+B}$



AND $A \rightarrow D \rightarrow \overline{A \cdot B}$



NOT

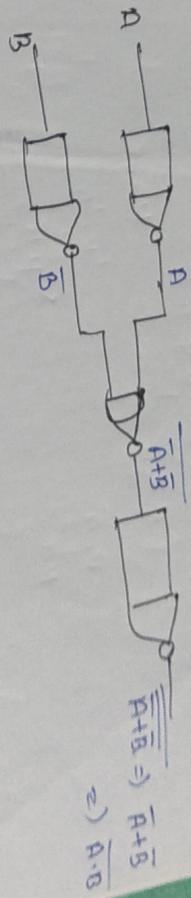
$A \rightarrow D \rightarrow \overline{A}$

XNOR
 $A \cdot \overline{B} + \overline{A} \cdot B$

NOT



NAND using NOR



$\overline{A} \cdot \overline{B} = \overline{A+B}$

\overline{A}

$\overline{A+B}$

$\Rightarrow \overline{A \cdot B}$

$\Rightarrow \overline{A+B}$

$\Rightarrow \overline{A+B}$