



南方科技大学
SOUTHERN UNIVERSITY OF SCIENCE AND TECHNOLOGY

CS323 Lab 8

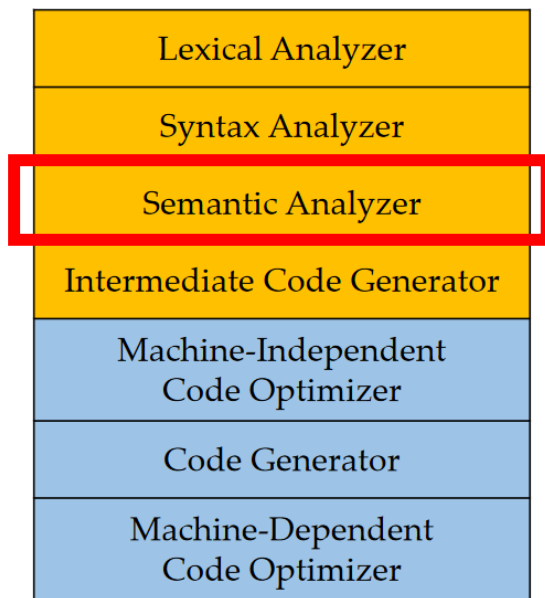
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Outline

- Overview of Project Phase 2
- Understand the parsers generated by Bison

Phase 2



- Build a semantic analyzer to check if the given SPL program satisfies semantic rules by analyzing parse tree and symbol table
 - No need to output anything if there is no rule violation
 - Otherwise, report the semantic errors (type, line number, error message, etc.)

Assumptions

Assumption 1 `char` variables only occur in assignment operations or function parameters/arguments

Assumption 2 only `int` variables can do boolean operations

Assumption 3 only `int` and `float` variables can do arithmetic operations

Assumption 4 no nested function definitions

Assumption 5 field names in `struct` definitions are unique (in any scope), i.e., the names of `struct` fields, and variables never overlap

Assumption 6 there is only the global scope, i.e., all variable names are unique

Assumption 7 using *named equivalence* to determine whether two `struct` types are equivalent

<code>struct st {int x;}</code>	The two structures are of the same type if we adopt named equivalence
<code>struct st {char y;}</code>	

Required rules

- Type 1** a variable is used without a definition
- Type 2** a function is invoked without a definition
- Type 3** a variable is redefined in the same scope
- Type 4** a function is redefined (in the global scope, since we don't have nested functions)
- Type 5** unmatching types appear at both sides of the assignment operator (=)
- Type 6** rvalue appears on the left-hand side of the assignment operator
- Type 7** unmatching operands, such as adding an integer to a structure variable
- Type 8** a function's return value type mismatches the declared type
- Type 9** a function's arguments mismatch the declared parameters (either types or numbers, or both)
- Type 10** applying indexing operator ([...]) on non-array type variables
- Type 11** applying function invocation operator (foo(...)) on non-function names
- Type 12** array indexing with a non-integer type expression
- Type 13** accessing members of a non-structure variable (i.e., misuse the dot operator)
- Type 14** accessing an undefined structure member
- Type 15** redefine the same structure type

Test cases

- We will provide test cases on GitHub (under `project/phase2/test/`)
- You are required to pass all of them

Example

```
1  struct Apple
2  {
3      int weight;
4      float round;
5  };
6  int test_2_r07()
7  {
8      struct Apple aa;
9      float weight_test = 1.0;
10     aa.weight = aa + 2;
11     return 0;
12 }
```

Error type 7 at Line 10: binary operation on non-number variables

Error type 5 at Line 10: unmatched type on both sides of assignment

What to submit?

- Your semantic analyzer source code + a brief report
- Five test cases each team
 - At least four violations of our semantic rules
 - No lexical or syntax errors
- If you implement additional rule checkers:
 - Document the features in the project report
 - Provide test cases
- Deadline: 10:00 PM, December 3, 2023

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Example Grammar G

- $S \rightarrow CC$
- $C \rightarrow cC \mid d$

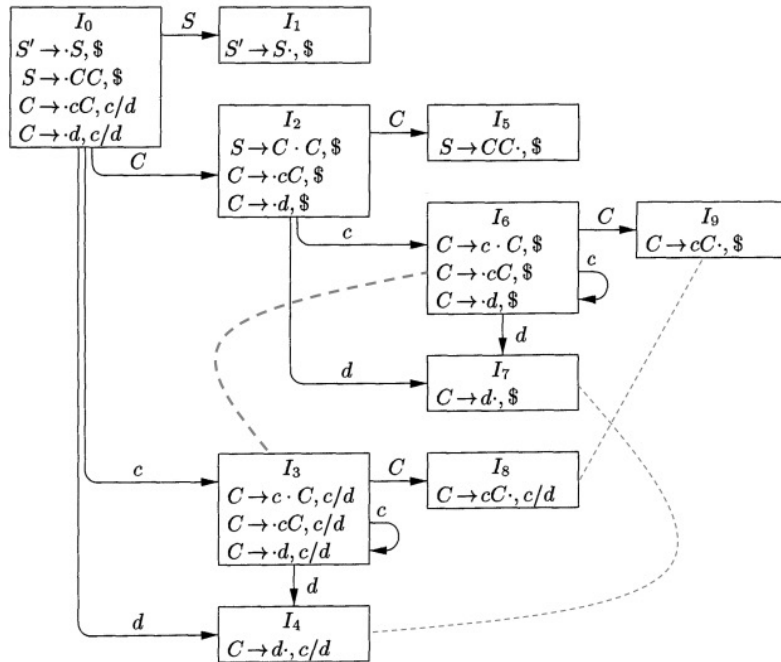
Step 1: Constructing LALR(1) Automaton

- Manually construct the LR(1) item sets following the algorithm below

```
void items( $G'$ ) {  
    initialize  $C$  to  $\{\text{CLOSURE}(\{[S' \rightarrow \cdot S, \$]\})\}$ ;  
    repeat  
        for ( each set of items  $I$  in  $C$  )  
            for ( each grammar symbol  $X$  )  
                if (  $\text{GOTO}(I, X)$  is not empty and not in  $C$  )  
                    add  $\text{GOTO}(I, X)$  to  $C$ ;  
    until no new sets of items are added to  $C$ ;  
}
```

Step 1: Constructing LALR(1) Automaton

- Merge the item sets with the same core



Merge the following item sets:

- I_3 and I_6
- I_4 and I_7
- I_8 and I_9

Step 2: Generate the LALR(1) Automaton Using Bison

- Write a Flex program to recognize two patterns “c” and “d”
- Write a Bison program to recognize the language $L(G)$
- We provide sample programs on GitHub if you cannot finish the above two programs by yourself
 - Under the lab8 directory
- Build the Bison program using the commands below:
 - `bison -d syntax.y --report all` (will generate a file “syntax.output”)
 - `flex lex.l`
 - `gcc syntax.tab.c -lfl -ly -o test.out`
 - The test.out executable can recognize strings such as “dd” (can be tested with `echo “dd” | ./test.out`)

Step 3: Understand the Generated LALR(1) Automaton

- Check the “syntax.output” file to understand the automaton

Grammar

0 \$accept: S \$end

1 S: C C

2 C: c C

3 | d

Terminals, with rules where they appear

\$end (0) 0

error (256)

c (258) 2

d (259) 3

Nonterminals, with rules where they appear

\$accept (5)

on left: 0

S (6)

on left: 1, on right: 0

C (7)

on left: 2 3, on right: 1 2

State 0

0 \$accept: . S \$end

1 S: . C C

2 C: . c C

3 | . d

c shift, and go to state 1

d shift, and go to state 2

S go to state 3

C go to state 4

Questions

- The automaton constructed by Bison has one more state than the one constructed manually by you.
 - Which state in the automaton generated by Bison does not appear in the automaton manually constructed by you?
 - Why is there such a difference?
- Is the automaton generated by Bison equivalent to the one manually constructed by you in terms of language recognition ability?