

Discrete Optimization

Linear Programming: Part III

Goals of the Lecture

- ▶ Linear programming
 - the Simplex algorithm

BFS and the Naive Algorithm

1. An optimal solution is located at a vertex.

2. A vertex is a Basic Feasible Solution (BFS).

► Naive algorithm

– generate all basic feasible solutions

- select m basic variables and perform Gaussian elimination
- test whether it is feasible

– select the BFS with the best cost

► How many basic solutions?

$$\frac{n!}{m!(n-m)!}$$

Outline of the Simplex Algorithm

Goal: You want to solve a linear program

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The Simplex Algorithm

- ▶ Local search algorithm
 - move from BFS to BFS
 - guaranteed to find the global optimum
 - because of convexity

The Simplex Algorithm

- ▶ Local search algorithm
 - move from BFS to BFS
 - guaranteed to find the global optimum
 - because of convexity
- ▶ Key idea: how to move from BFS to BFS.

From BFS to BFS

$$\begin{array}{rclclclclclclclcl} 3x_1 & - & 2x_2 & + & x_3 & & & = & 1 \\ 2x_1 & & & & & + & x_4 & & + & x_6 & = & 2 \\ x_1 & & & & & & + & x_5 & + & x_6 & = & 3 \end{array}$$



$$\begin{array}{rclclclclclclclcl} x_3 & = & 1 & - & 3x_1 & - & 2x_2 & & & & & \\ x_4 & = & 2 & - & 2x_1 & & & + & x_6 & & & \\ x_5 & = & 3 & - & x_1 & & & + & x_6 & & & \end{array}$$

From BFS to BFS

$$\begin{array}{rclclcl} x_3 & = & 1 & - & 3x_1 & - & 2x_2 & & \\ x_4 & = & 2 & - & 2x_1 & & & + & x_6 \\ x_5 & = & 3 & - & x_1 & & & + & x_6 \end{array}$$

- ▶ How to move to another BFS
 - select a non-basic variable with a negative coefficient: *entering variable*
 - introduce this variable in the basis by removing a basic variable: *leaving variable*
 - perform Gaussian elimination
- ▶ Local move: swap a basic and a non-basic variables

From BFS to BFS

$$\begin{array}{rclclcl} x_3 & = & 1 & - & 3x_1 & - & 2x_2 & & \\ x_4 & = & 2 & - & 2x_1 & & & + & x_6 \\ x_5 & = & 3 & - & x_1 & & & + & x_6 \end{array}$$



$$\begin{array}{rclclclcl} x_2 & = & \frac{1}{2} & - & \frac{3}{2}x_1 & - & \frac{1}{2}x_3 & & \\ x_4 & = & 2 & - & 2x_1 & & & + & x_6 \\ x_5 & = & 3 & - & x_1 & & & + & x_6 \end{array}$$

From BFS to BFS

$$\begin{array}{rclcl} x_3 & = & 1 & - & 3x_1 & - & 2x_2 & & \\ x_4 & = & 2 & - & 2x_1 & & & + & x_6 \\ x_5 & = & 3 & - & x_1 & & & + & x_6 \end{array}$$

From BFS to BFS

x_3	$=$	1	$-$	$3x_1$	$-$	$2x_2$		
x_4	$=$	2	$-$	$2x_1$			$+$	x_6
x_5	$=$	3	$-$	x_1			$+$	x_6

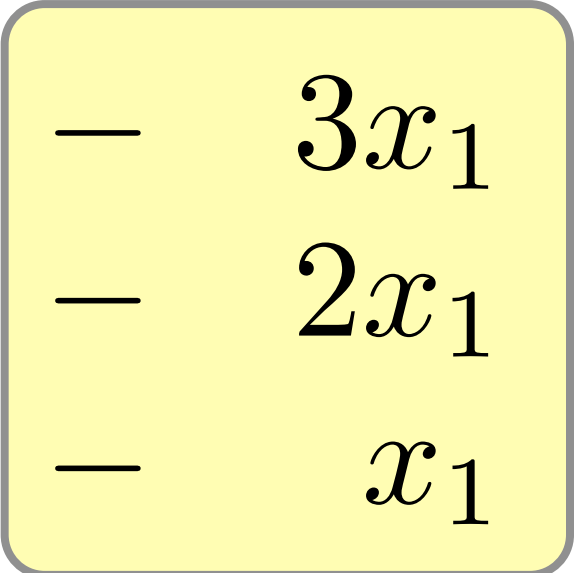
From BFS to BFS


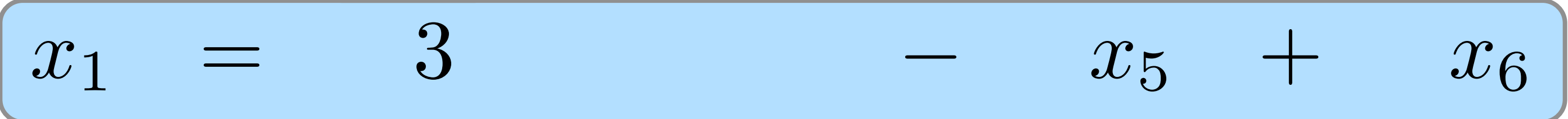
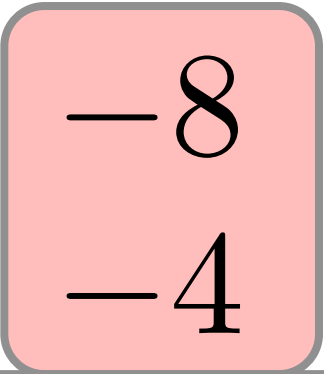
$$\begin{array}{rclclcl} x_3 & = & 1 & - & 3x_1 & - & 2x_2 \\ x_4 & = & 2 & - & 2x_1 & & + & x_6 \\ x_5 & = & 3 & - & x_1 & & + & x_6 \end{array}$$



$$\begin{array}{rclclclcl} x_3 & = & -8 & - & 2x_2 & + & 3x_5 & - & 3x_6 \\ x_4 & = & -4 & & & + & 2x_5 & - & x_6 \\ x_1 & = & 3 & & & - & x_5 & + & x_6 \end{array}$$

From BFS to BFS

$$\begin{array}{rcll} x_3 & = & 1 & - 3x_1 - 2x_2 \\ x_4 & = & 2 & - 2x_1 + x_6 \\ x_5 & = & 3 & - x_1 + x_6 \end{array}$$



$$\begin{array}{rcll} x_3 & = & -8 & - 2x_2 + 3x_5 - 3x_6 \\ x_4 & = & -4 & + 2x_5 - x_6 \\ x_1 & = & 3 & - x_5 + x_6 \end{array}$$


Not a BFS: I cannot select the leaving variable arbitrarily!

From BFS to BFS

$$\begin{array}{rclcl} x_3 & = & 1 & \boxed{\begin{array}{r} - \quad 3x_1 \\ - \quad 2x_1 \\ - \quad x_1 \end{array}} & - \quad 2x_2 \\ x_4 & = & 2 & & + \quad x_6 \\ x_5 & = & 3 & & + \quad x_6 \end{array}$$

$$\boxed{\begin{array}{c} \frac{1}{3} \\ 1 \\ 3 \end{array}}$$

► How to choose the leaving variable?

– we must maintain feasibility

$$l = \arg\min_{i: a_{ie} < 0} \frac{b_i}{-a_{ie}}$$

guarantee the feasibility because $a_{ie} < 0$

From BFS to BFS

$x_3 = 1$
 $x_4 = 2$
 $x_5 = 3$

$- 3x_1$
 $- 2x_1$
 $- x_1$

$- 2x_2$

$+ x_6$
 $+ x_6$

$\frac{1}{3}$
 1
 3

$x_1 =$

$\frac{1}{3}$

$-$

$\frac{2}{3}x_2$

$-$

$\frac{1}{3}x_3$

$x_4 =$

$\frac{4}{3}$

$+$

$\frac{4}{3}x_2$

$+$

$\frac{2}{3}x_3$

$+$

x_6

$x_5 =$

$\frac{8}{3}$

$+$

$\frac{2}{3}x_2$

$+$

$\frac{1}{3}x_3$

$+$

x_6

The Local Move

► Moving from BFS to BFS

- select the entering variable x_e

- non-basic variable with negative coefficients in the right-hand side

- select the leaving variable x_l to maintain feasibility

$$l = \arg\min_{i:a_{ie}<0} \frac{b_i}{-a_{ie}}$$

- apply Gaussian elimination

- eliminate x_e from the right-hand side

► This operation is called pivoting in linear programming

- pivot(e,l)

Outline of the Simplex Algorithm

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The Simplex Algorithm

A BFS is optimal if its objective function, after having eliminated all the basic variables, is of the form

$$c_0 + c_1x_1 + \dots + c_nx_n$$

with

$$c_i \geq 0 \quad (1 \leq i \leq n).$$

From BFS to BFS

$$\begin{array}{ll} \min & x_1 + x_2 + x_3 + x_4 + x_5 \\ \text{subject to} & \\ & 3x_1 + 2x_2 + x_3 = 1 \\ & 5x_1 + x_2 + x_3 + x_4 = 3 \\ & 2x_1 + 5x_2 + x_3 + x_5 = 4 \end{array}$$

- First, find a BFS and eliminate the basic variable from the objective function

From BFS to BFS

$$\begin{array}{ll}
 \min & x_1 + x_2 + x_3 + x_4 + x_5 \\
 \text{subject to} & \\
 & 3x_1 + 2x_2 + x_3 = 1 \\
 & 5x_1 + x_2 + x_3 + x_4 = 3 \\
 & 2x_1 + 5x_2 + x_3 + x_5 = 4
 \end{array}$$



$$\begin{array}{ll}
 \min & \\
 \text{subject to} &
 \end{array}$$

$$6 - 3x_1 - 3x_2$$

$$\begin{array}{llll}
 x_3 & = & 1 & - 3x_1 - 2x_2 \\
 x_4 & = & 2 & - 2x_1 + x_2 \\
 x_5 & = & 3 & + x_1 - 3x_2
 \end{array}$$

From BFS to BFS

$$\begin{array}{ll} \min & 6 - 3x_1 - 3x_2 \\ \text{subject to} & \\ & x_3 = 1 - 3x_1 - 2x_2 \\ & x_4 = 2 - 2x_1 + x_2 \\ & x_5 = 3 + x_1 - 3x_2 \end{array}$$

From BFS to BFS

min

subject to

$$\begin{array}{rclclcl} & & 6 & - & 3x_1 & - & 3x_2 \\ x_3 & = & 1 & - & 3x_1 & - & 2x_2 \\ x_4 & = & 2 & - & 2x_1 & + & x_2 \\ x_5 & = & 3 & + & x_1 & - & 3x_2 \end{array}$$

From BFS to BFS

min
subject to

		6	−	$3x_1$	−	$3x_2$
x_3	=	1	−	$3x_1$	−	$2x_2$
x_4	=	2	−	$2x_1$	+	x_2
x_5	=	3	+	x_1	−	$3x_2$

From BFS to BFS

min
subject to

$$\begin{array}{rclcl}
 & 6 & - & 3x_1 & - & 3x_2 \\
 x_3 & = & 1 & - & 3x_1 & - & 2x_2 \\
 x_4 & = & 2 & - & 2x_1 & + & x_2 \\
 x_5 & = & 3 & + & x_1 & - & 3x_2
 \end{array}$$

min
subject to

$$\begin{array}{rclcl}
 & \frac{9}{2} & + & \frac{3}{2}x_1 & + & \frac{3}{2}x_3 \\
 x_2 & = & \frac{1}{2} & - & \frac{3}{2}x_1 & - & \frac{1}{2}x_3 \\
 x_4 & = & \frac{5}{2} & - & \frac{7}{2}x_1 & - & \frac{1}{2}x_3 \\
 x_5 & = & \frac{3}{2} & + & \frac{11}{2}x_1 & + & \frac{3}{2}x_3
 \end{array}$$

From BFS to BFS

min
subject to

$$\begin{array}{rclclcl} & & 6 & - & 3x_1 & - & 3x_2 \\ x_3 & = & 1 & - & 3x_1 & - & 2x_2 \\ x_4 & = & 2 & - & 2x_1 & + & x_2 \\ x_5 & = & 3 & + & x_1 & - & 3x_2 \end{array}$$



min
subject to

$$\frac{9}{2} + \frac{3}{2}x_1 + \frac{3}{2}x_3$$

$$x_2 = \frac{1}{2} - \frac{3}{2}x_1 - \frac{1}{2}x_3$$

$$x_4 = \frac{5}{2} - \frac{7}{2}x_1 - \frac{1}{2}x_3$$

$$x_5 = \frac{3}{2} + \frac{11}{2}x_1 + \frac{3}{2}x_3$$

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Look Ma, My Local Move is Improving

- ▶ Assume that, in the BFS,
 - $b_1 > 0, \dots, b_m > 0$
 - there exists an entering variable e with $c_e < 0$
 - there exists a leaving variable l
- ▶ then the move $\text{pivot}(e, l)$ is improving

The Simplex Algorithm

```
while  $\exists 1 \leq i \leq n : c_i < 0$  do  
  choose  $e$  such that  $c_e < 0$ ;  
   $l = \arg\text{-min}_{i: a_{ie} < 0} \frac{b_i}{-a_{ie}}$ ;  
  pivot( $e, l$ );
```

The Simplex Algorithm

```
while  $\exists 1 \leq i \leq n : c_i < 0$  do  
    choose  $e$  such that  $c_e < 0$ ;  
     $l = \arg\text{-min}_{i: a_{ie} < 0} \frac{b_i}{-a_{ie}}$ ;  
    pivot( $e, l$ );
```

- ▶ Assume that during the execution
 - b_1, \dots, b_m are always strictly positive
 - the objective function is bounded by below
- ▶ then the simplex algorithm terminates with an optimal solution

Where is My Leaving Variable?

- Selecting the leaving variable

$$l = \arg\min_{i:a_{ie}<0} \frac{b_i}{-a_{ie}}$$

$$\begin{array}{ll} \min & 6 - 3x_1 - 3x_2 \\ \text{subject to} & \\ & x_3 = 1 + 3x_1 - 2x_2 \\ & x_4 = 2 + 2x_1 + x_2 \\ & x_5 = 3 + x_1 - 3x_2 \end{array}$$

- There are no leaving variables

Where is My Leaving Variable?

- Selecting the leaving variable

$$l = \arg\min_{i: a_{ie} < 0} \frac{b_i}{-a_{ie}}$$

min
subject to

	6	−	$3x_1$	−	$3x_2$	
x_3	=	1	+	$3x_1$	−	$2x_2$
x_4	=	2	+	$2x_1$	+	x_2
x_5	=	3	+	x_1	−	$3x_2$

- There are no leaving variables

Where is My Leaving Variable?

$$\begin{array}{ll} \min & 6 - 3x_1 - 3x_2 \\ \text{subject to} & \\ & x_3 = 1 + 3x_1 - 2x_2 \\ & x_4 = 2 + 2x_1 + x_2 \\ & x_5 = 3 + x_1 - 3x_2 \end{array}$$

► What is the basic solution?

$$\{x_1 = 0; x_2 = 0; x_3 = 1; x_4 = 2; x_5 = 3\}$$

► What happens if I increase the value of x_1

- the solution remains feasible
- the value of the objective can decrease arbitrarily

Where is My Leaving Variable?

- What if some b_i becomes zero?

$$\begin{array}{ll}
 \min & 5 - 3x_1 - 3x_2 \\
 \text{subject to} & \\
 & x_3 = 0 - 3x_1 - 2x_2 \\
 & x_4 = 2 - 2x_1 + x_2 \\
 & x_5 = 3 + x_1 - 3x_2
 \end{array}$$

- What is the leaving variable?

$$\begin{array}{ll}
 \min & 5 + \frac{3}{2}x_1 + \frac{3}{2}x_3 \\
 \text{subject to} & \\
 & x_2 = 0 - \frac{3}{2}x_1 - \frac{1}{2}x_3 \\
 & x_4 = 2 - \frac{7}{2}x_1 - \frac{1}{2}x_3 \\
 & x_5 = 3 + \frac{11}{2}x_1 + \frac{3}{2}x_3
 \end{array}$$

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We need to find a new way to guarantee termination

Termination of the Simplex Algorithm

► Many approaches

– Bland rule

- select always the first entering variable lexicographically

– Lexicographic pivoting rule

- break ties when selecting the leaving variable by using a lexicographic rule

$$l = \arg\text{-lex-min}_{i:a_{ie}<0} \frac{b_i}{-a_{ie}}$$

– Perturbation methods

Where is My BFS?

$$\begin{array}{ll} \min & x_1 + x_2 + x_3 + x_4 + x_5 \\ \text{subject to} & \\ & 3x_1 + 2x_2 + x_3 = 1 \\ & 5x_1 + x_2 + x_3 + x_4 = 3 \\ & 2x_1 + 5x_2 + x_3 + x_5 = 4 \end{array}$$

► How do I find my first BFS?

Where is My BFS?

- How do I find my first BFS?

$$\begin{array}{ll} \min & c_1 x_1 + \dots + c_n x_n \\ \text{subject to} & \\ & a_{11} x_1 + \dots + a_{1n} x_n = b_1 \\ & \dots \\ & a_{m1} x_1 + \dots + a_{mn} x_n = b_m \end{array}$$

- Introduce artificial variables

Where is My BFS?

- Introduce artificial variables

$$\begin{array}{llllllllll} \min & c_1 x_1 & + & \dots & + & c_n x_n & & & & \\ \text{subject to} & a_{11} x_1 & + & \dots & + & a_{1n} x_n & + & y_1 & & = b_1 \\ & \dots & & & & & & & & \\ & a_{i1} x_1 & + & \dots & + & a_{in} x_n & & + & y_i & = b_i \\ & \dots & & & & & & & & \\ & a_{m1} x_1 & + & \dots & + & a_{mn} x_n & & & + & y_m = b_m \end{array}$$

- I have an easy BFS

- But to another problem

Where is My BFS?

- Introduce artificial variables

$$\begin{array}{llllllllll} \min & c_1 x_1 & + & \dots & + & c_n x_n & & & & \\ \text{subject to} & a_{11} x_1 & + & \dots & + & a_{1n} x_n & + & y_1 & & = b_1 \\ & \dots & & & & & & & & \\ & a_{i1} x_1 & + & \dots & + & a_{in} x_n & & + & y_i & = b_i \\ & \dots & & & & & & & & \\ & a_{m1} x_1 & + & \dots & + & a_{mn} x_n & & & + & y_m = b_m \end{array}$$

- I have an easy BFS



- But to another problem

Where is My BFS?

- Introduce artificial variables

$$\begin{array}{llllllllll} \min & c_1 x_1 & + & \dots & + & c_n x_n & & & & \\ \text{subject to} & a_{11} x_1 & + & \dots & + & a_{1n} x_n & + & y_1 & & = b_1 \\ & \dots & & & & & & & & \\ & a_{i1} x_1 & + & \dots & + & a_{in} x_n & & + & y_i & = b_i \\ & \dots & & & & & & & & \\ & a_{m1} x_1 & + & \dots & + & a_{mn} x_n & & & + & y_m = b_m \end{array}$$

- I have an easy BFS



- But to another problem



Two-Phase Method

- ▶ First find a BFS
 - if there is one
- ▶ Then find an optimal BFS

Two-Phase Method

► First find a BFS

$$\min \quad y_1 + \dots + y_m$$

subject to

$$a_{11}x_1 + \dots + a_{1n}x_n + y_1 = b_1$$

...

$$a_{i1}x_1 + \dots + a_{in}x_n + y_i = b_i$$

...

$$a_{m1}x_1 + \dots + a_{mn}x_n + y_m = b_m$$

► Feasible if the objective value reaches 0

– all the y_i are thus zeros

- I have a BFS without the y_i
- really? always?

Until Next Time