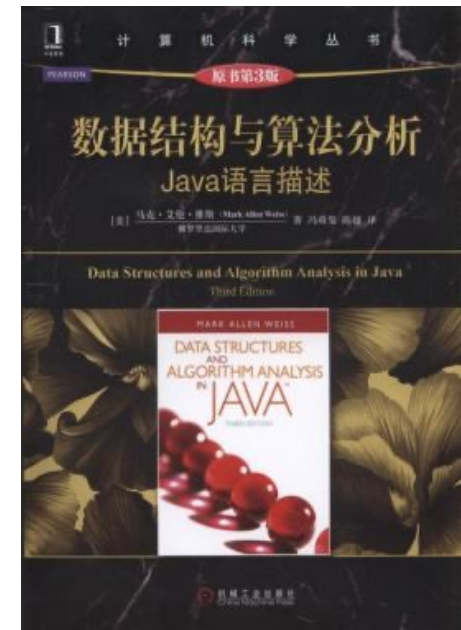
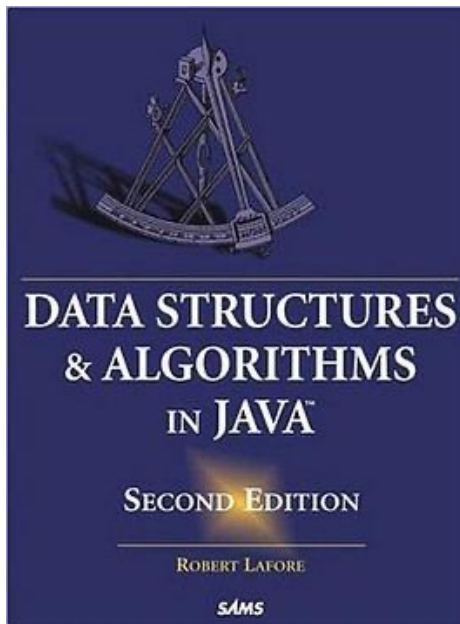
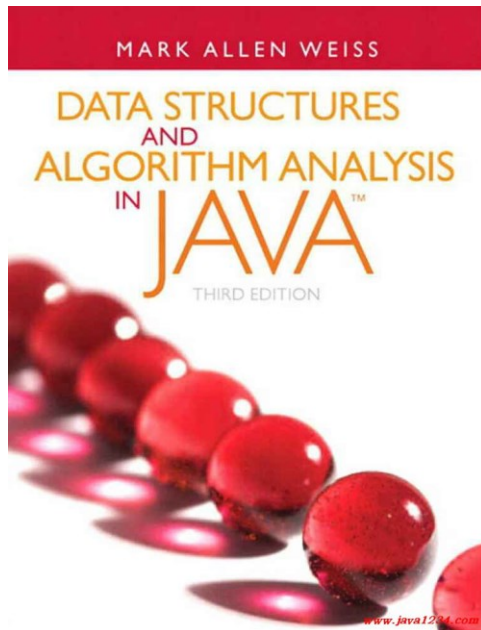


Topic 6 – Sorting algorithms



Topics

- Introduction
- Programming Revision
- Methods and Objects
- Arrays and Array Algorithms
- Big O Notation
- **Sorting Algorithms**
- Stacks and Queues
- Linked Lists
- Recursion
- Bit Manipulation

Outline

- **Introduction to Sorting**
- Bubble Sort
- Selection Sort
- Insertion Sort
- Wrap up

Sorting

- Put the data according to the rules
- Example.
 - Input:
[7 6 0 3 8 9 4 1]
 - Output:
 - from small to large
[0 1 3 4 6 7 8 9]
 - from large to small
[9 8 7 6 4 3 1 0]
 - decrease first and then increase
[9 8 7 0 1 3 4 6]

Sorting

- To be useful, data should be sorted so that it is easy to search through
- Often, we will get an array of data that is in a random order, and we will need to sort it
- We need to develop some sorting algorithms to sort arrays - the more efficient the sorting algorithm, the better



Sorting

There are many different sorting algorithms and the correct one to choose depends on the following:

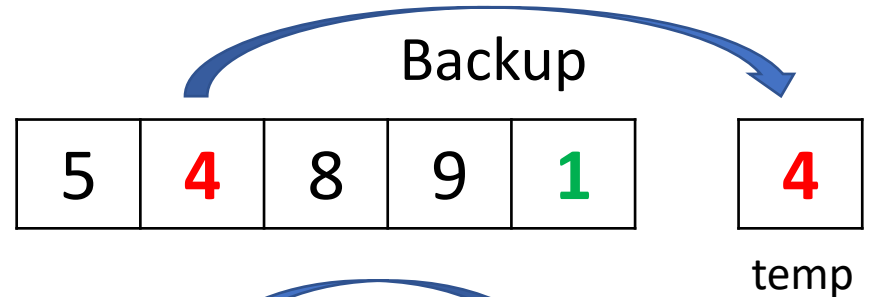
- The amount of data that needs to be sorted
- The amount of memory that is available
- The amount of time that is available
- The way the data is distributed

Sorting

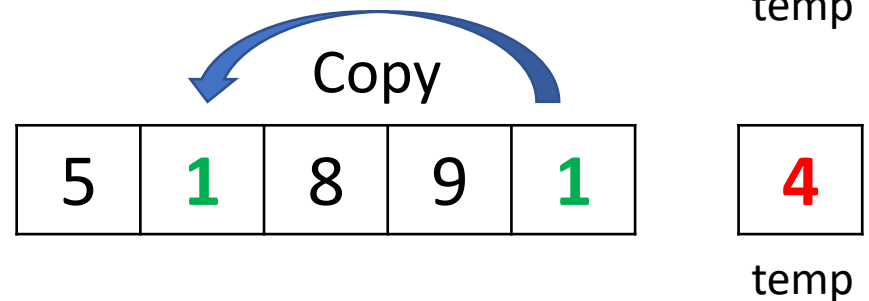
- For starters, we will consider 3 simple sorting algorithms
 - **bubble sort**
 - **selection sort**
 - **insertion sort**
- All of these algorithms have a complexity of **$O(n^2)$**
- This means that as the size of the array n increases, the time taken to sort it will increase by that amount squared
 - An array with size of $2n$ will take 4 times longer to sort
 - An array that is 3 times bigger will take 9 times longer to sort
- For that reason, these algorithms are only used to sort small amounts of data!

Swapping

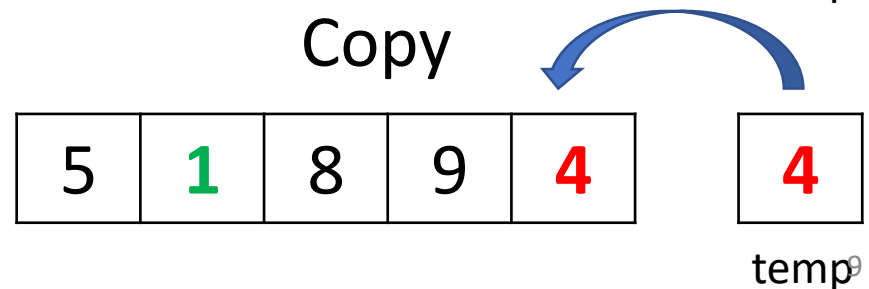
- Swapping elements in an array is crucial to be able to sort it
- Let's swap slot 2 with slot 5
 - Backup slot 2 into temp



- Copy slot 5 into slot 2



- Copy temp into slot 5



Swapping

- Swapping elements in an array is crucial to be able to sort it
- Lets swap slot 2 with slot 5

1. Backup slot 2 into temp

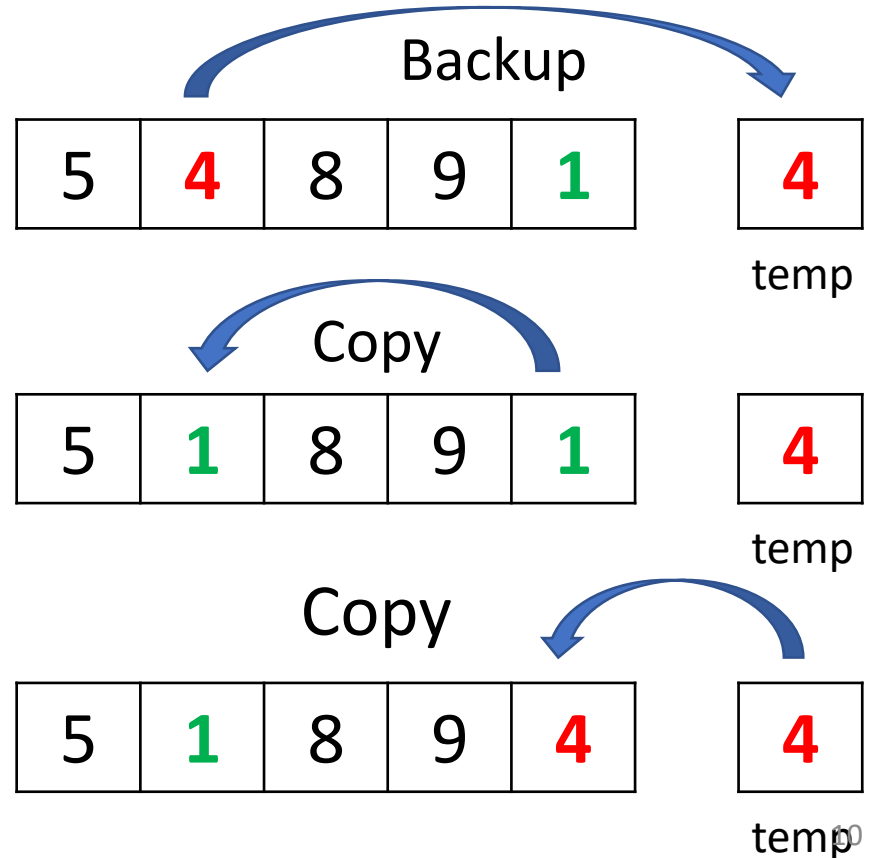
- `int temp = arr[1];`**

2. Copy slot 5 into slot 2

- `arr[1] = arr[4];`**

3. Copy temp into slot 5

- `arr[4] = temp;`**



Swapping

- In order to swap one variable with another in an array
 - Back-up slot **A** (the one that will be overwritten first) into a temporary variable
`int temp = array[A];`
 - Overwrite slot **A** with the value of slot **B**
`array[A] = array[b];`
 - Use the temporary variable to overwrite the value of slot **B** with the original value of slot **A**
`array[B] = temp;`

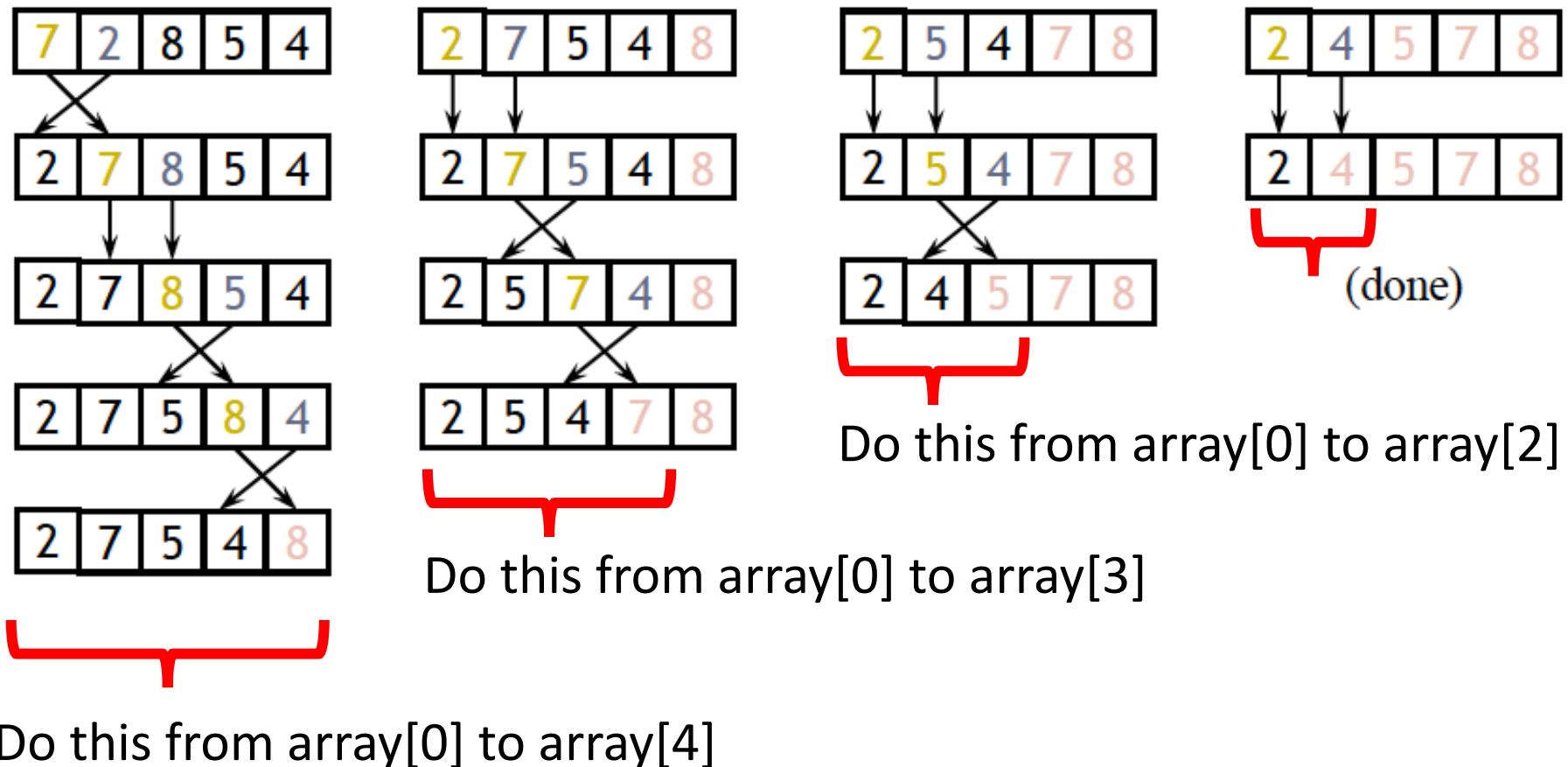
Outline

- Introduction to Sorting
- **Bubble Sort**
- Selection Sort
- Insertion Sort
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Bubble Sort

- The **bubble sort** algorithm works like this:
 - Start at the beginning of the array
 - Compare the first two numbers
 - If the one to the left is bigger, swap it with the one on the right
 - Move one position to the right and check the next two
 - Keep doing this until you reach the end of the array
 - **The biggest number** has now been '**bubbled**' to the top
 - Go back to the beginning, do the same thing and bubble the next biggest number up
 - Stop short of the top part of the array where the '**bubbles**' have arrived

Example of Bubble Sort



Example of Bubble Sort

- Red pair for current comparison
- Black squares for fixed ones

6 5 3 1 8 7 2 4



WIKIPEDIA
The Free Encyclopedia

Implementing bubble sort

- How would you write this in Java?
- Because of all the swapping, it makes sense to write a swap() method

```
public void swap(int first, int second) {  
    int temp = array[first];  
    array[first] = array[second];  
    array[second] = temp;  
}
```

//Here, we write the function for swap first.

Main bubble sort method

- Pseudo code might look something like this
 - outer 'bubbling' loop** running from end of array backwards, bubbling biggest element to the top each time it runs
 - {
 - inner 'swapping' loop** running from start of array up to last unsorted element, swapping two elements at a time
 - {
 - check if element[i] > element[i+1]**
 - if so, swap them**
 - }
 - }
 - }

Java Implementation

```
public void bubblesort() {  
    int outer, inner;  
    for(outer = nElems-1; outer > 0; outer--) {  
        for(inner = 0; inner < outer; inner++) {  
            if( array[inner] > array[inner+1]) {  
                swap(inner,inner+1);  
            }  
        }  
    }  
}
```

// inner

// outer

Exercise



- Write a function to sort an array using bubble sort
 - Input: array to be sorted

Algorithm complexity

- The complexity of a sorting algorithm is the relationship between the size of the array to be sorted and the length of time it takes to sort
- This depends on the number of comparisons and swaps
- A **comparison** involves a memory read because you are checking the difference between two numbers

if (a[inner] < a[min])

- A **swap** involves a memory write because you are changing the arrangement of data in memory

swap(inner, inner+1);

- Swaps take far longer to execute than comparisons because memory writes take longer than memory reads

Complexity of bubble sort

- Assume there are 10 numbers to be sorted (i.e. n is 10)
- Because the outer loop decreases by one each time the number of comparisons performed by the inner loop will be

$$9 + 8 + 7 + \dots + 1 = 45$$

- In general this can be expressed as

$$(n - 1) + (n - 2) + \dots + 1 = (n - 1) \times \frac{n}{2}$$

number of comparisons

Complexity of bubble sort

- On average, the algorithm will do a swap half of the time: $\frac{(n-1) \times n}{4}$ in total
 - In the **BEST CASE SCENARIO** it won't have to do any swaps
 - In the **WORST CASE SCENARIO** it will have to perform a swap after every comparison: $(n-1) \times \frac{n}{2}$ in total
 - As the size of the array n increases, the number of comparisons and swaps increases by a factor of n^2
 - Therefore, bubble sort is $O(n^2)$ → very inefficient

Outline

- Introduction to Sorting
- Bubble Sort
- **Selection Sort**
- Insertion Sort
- Wrap up

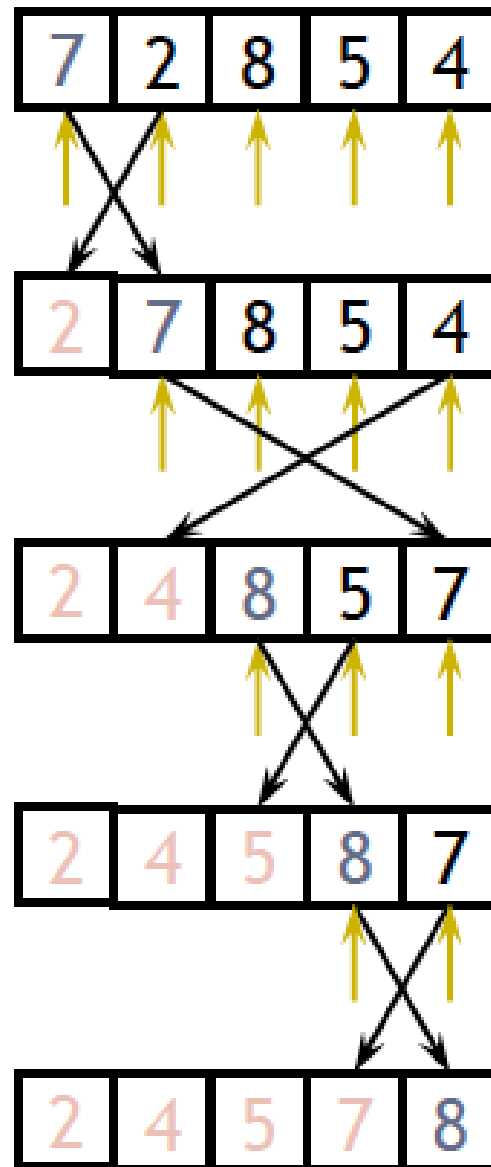
Selection Sort

- Selection sort improves on bubble sort by reducing number of swaps from $O(n^2)$ to $O(n)$
- Number of comparisons stays $O(n^2)$ but comparisons are shorter than swaps so this can be an important improvement
- Rather than continuously swapping to 'bubble' a number to the top, we find the smallest number and swap it into place

Selection Sort

- Given an array of length n
 - Search elements 0 through $n - 1$ and select the smallest
 - Swap it with the element in location 0
 - Search elements 1 through $n - 1$ and select the smallest
 - Swap it with the element in location 1
 - Search elements 2 through $n - 1$ and select the smallest
 - Swap it with the element in location 2
 - Search elements 3 through $n - 1$ and select the smallest
 - Swap it with the element in location 3
 - Continue in this fashion until there's nothing left to search

Selection Sort



	8
	5
	2
	6
	9
	3
	1
	4
	0
	7

Selection Sort

- Red is current min.
- Yellow is sorted list.
- Blue is current item



Main selection sort method

- Pseudo code might look something like this

outer loop running through each place in the array looking for the correct element to swap into that place – starts at the beginning

```
{
```

inner loop which always looks for the smallest remaining unsorted item

```
{
```

```
    find minimum
```

```
    swap array[outer] with array[minimum]
```

```
}
```

```
}
```

Code for Selection Sort

```
public static void selectionSort() {  
    int min;  
    for (int outer = 0; outer < array.length-1; outer++) {  
        // outer is the point where the unsorted numbers start  
        min = outer;  
        // min's default value is the first slot to be checked  
        for (int i = outer + 1; i < array.length; i++) {  
            // inner loop checks through the unsorted numbers  
            if (array[i] < array[min]) {  
                min = i; //inner loop finds the minimum  
            }  
        } // min always refers to the min found so far  
        swap(outer, min);  
    } // all items with slot numbers less than or equal to outer are sorted  
}
```

Exercise



- Write a function to sort an array using selection sort
 - Input: array to be sorted

Analysis of selection sort

- The outer loop executes $n - 1$ times (i.e. we have to find the smallest number $n - 1$ times)
- The inner loop executes about $\frac{n}{2}$ times on average (it executes number $n - 1$ times the first time and only once for the final time)
- Number of comparisons required is roughly
$$(n - 1) \times \frac{n}{2} = \frac{1}{2}n^2 - \frac{1}{2}n$$
- The algorithm is therefore $O(n^2)$

Complexity of Selection Sort

- Selection sort performs the same number of comparisons as bubble sort :

$$(n - 1) \times \frac{n}{2}$$

- However, we have minimized the number of swaps required – we only swap something into its correct location once – a total of $n - 1$ swaps
- The algorithm is still $O(n^2)$ but it is a faster $O(n^2)$ than bubble sort since a swap takes far longer than a comparison and there are less swaps involved

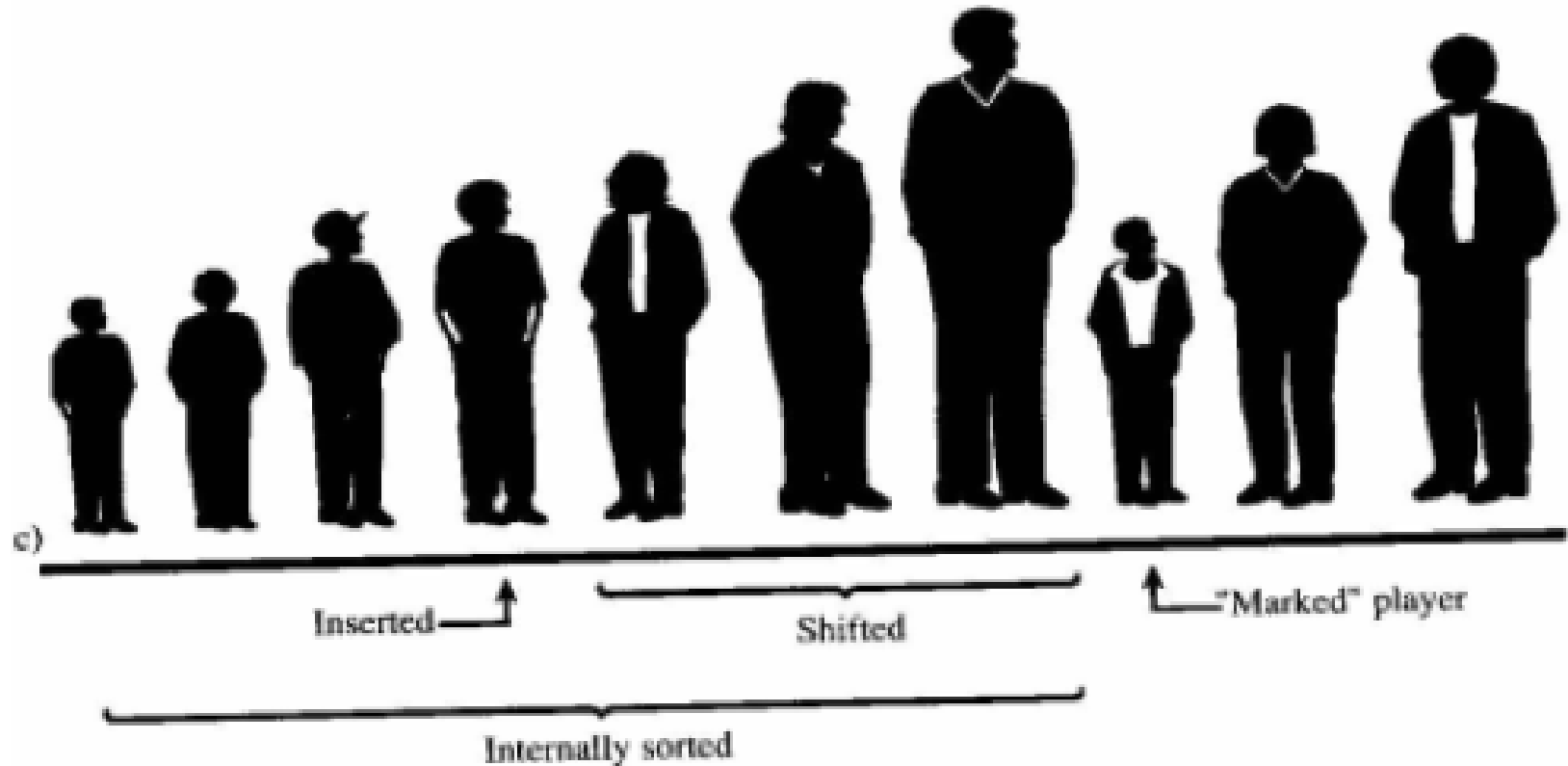
Outline

- Introduction to Sorting
- Bubble Sort
- Selection Sort
- **Insertion Sort**
- Wrap up

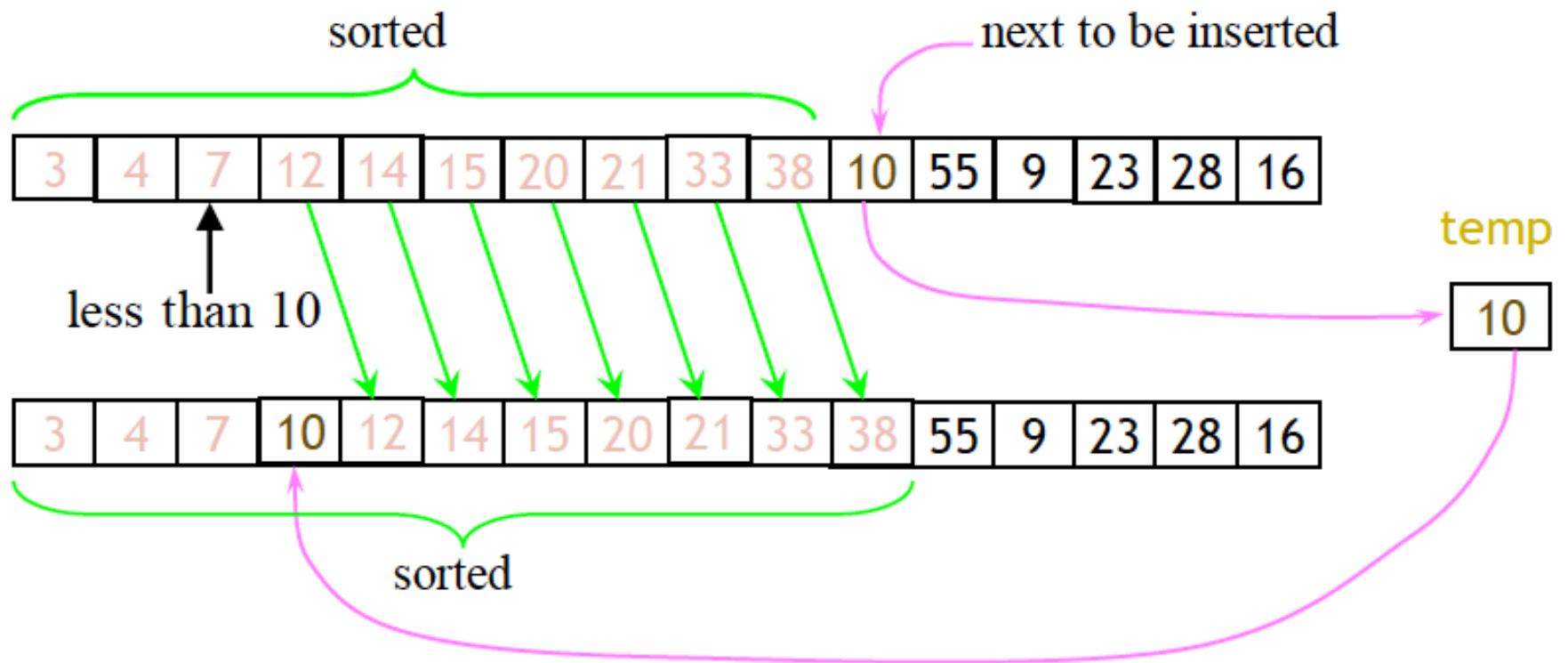
Insertion Sort

- Still $O(n^2)$ but can be faster than bubble and selection sort
- Usually used as the final stage of more sophisticated sorts, like quicksort
- Doesn't use a swap – instead shifts elements up to make space
- The idea is that the elements on the left of a marker are already sorted
- You make space to slot in the new unsorted element by moving all the other elements up one
- Exactly like insertion into an ordered array

Insertion Sort



Insertion Sort



Insertion Sort

6 5 3 1 8 7 2 4

- The partial sorted list (**black**) initially contains only the first element in the list.
- With each iteration one element (**red**) is removed from the "not yet checked for order" input data and inserted in-place into the sorted list.



Code for Insertion Sort

```
public static void insertionSort() {  
    for (int outer = 1; outer < array.length; outer++) {  
        // outer is the next element to be sorted  
        int temp = array[outer]; //back it up  
        int inner = outer; // inner used to track shifts  
        while (inner > 0 && array[inner - 1] >= temp) {  
            array[inner] = array[inner - 1]; // swap  
            inner--;  
        } //shift them all right until one is smaller  
        array[inner] = temp;  
    }  
}
```

Exercise



- Write a function to sort an array using insertion sort
 - Input: array to be sorted

Analysis of insertion sort

- We run once through the outer loop, inserting each of n elements
- On average, there are $\frac{n}{2}$ elements already sorted
 - The inner loop checks and moves half of these
 - This gives a second factor of $\frac{n}{4}$
- Hence, the time required for an insertion sort of an array of n elements is proportional to $\frac{n^2}{4}$
- Discarding constants, we find that insertion sort is $O(n^2)$
- For already sorted data, runs in $O(n)$ time
- For data arranged in inverse order, runs no faster than bubble sort

Outline

- Introduction to Sorting
- Bubble Sort
- Selection Sort
- Insertion Sort
- **Wrap up**

Remembering

- Bubble Sort
 - Bubbles the biggest to the end by swapping every time
- Selection Sort
 - Selects the min and swaps that element to the beginning
- Insertion Sort
 - Finds where the element should go in the sorted part and moves all the elements up one to make space
- All three algorithms are considered $O(n^2)$

Comparison

- **Bubble sort** is useful for **small** amounts of data because it is easy to code
 - Comparisons: Always $O(n^2)$
 - Swaps: Depends on how sorted the list is → varies from 0 to $O(n^2)$
- **Selection sort** can be used when amount of data is small and memory writing is far more time consuming than memory reading
 - Comparisons: Always $O(n^2)$
 - Swaps: Always $n - 1$
- **Insertion sort** performs well when the list is already partially sorted
 - Comparisons: Depends on how sorted the list is → varies from $O(n)$ to $O(n^2)$
 - Swaps: Depends on how sorted the list is → varies from $O(n)$ to $O(n^2)$

