

Today

537 : Part 3

=> CPU virtualization
→ mechanisms

(Limited Direct
Execution)

LDE

→ policies : scheduler
which process
should run?

last time:

→ Shortest Job First (SJF)
(assumes knowledge
of run time)

→ Round Robin :



not RR

ABAB
↓

time slice (quantum)

=> responsive (interactive)

Today: develop real sched
policy => [classic Unix
scheduler]
=> Multi-level Feedback Queue
(MLFQ)

Later today: (Virtual Memory)

Later later today : Project
Zo
(Processes, Shell)

Life ↔ Work

OS sched Policy: MLFQ

Problem: Don't know very
much about processes!

like to
learn: short jobs, longer running?

how? \Rightarrow measure : using past
to predict future)

many queues : job is on
priority: one queue @
highest Q₂ any given
middle Q₁, B, C time
low Q₀, D (might
change over
time)

each queue has priority

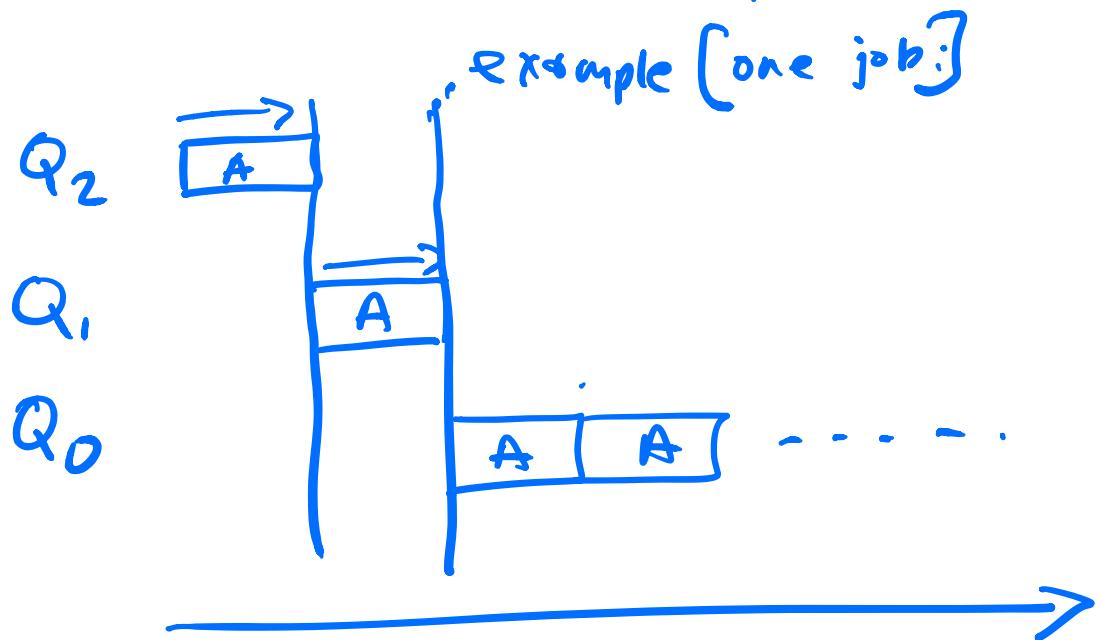
Rules:

1) if Priority(A) \geq Priority(B)
 \Rightarrow A runs (B doesn't)

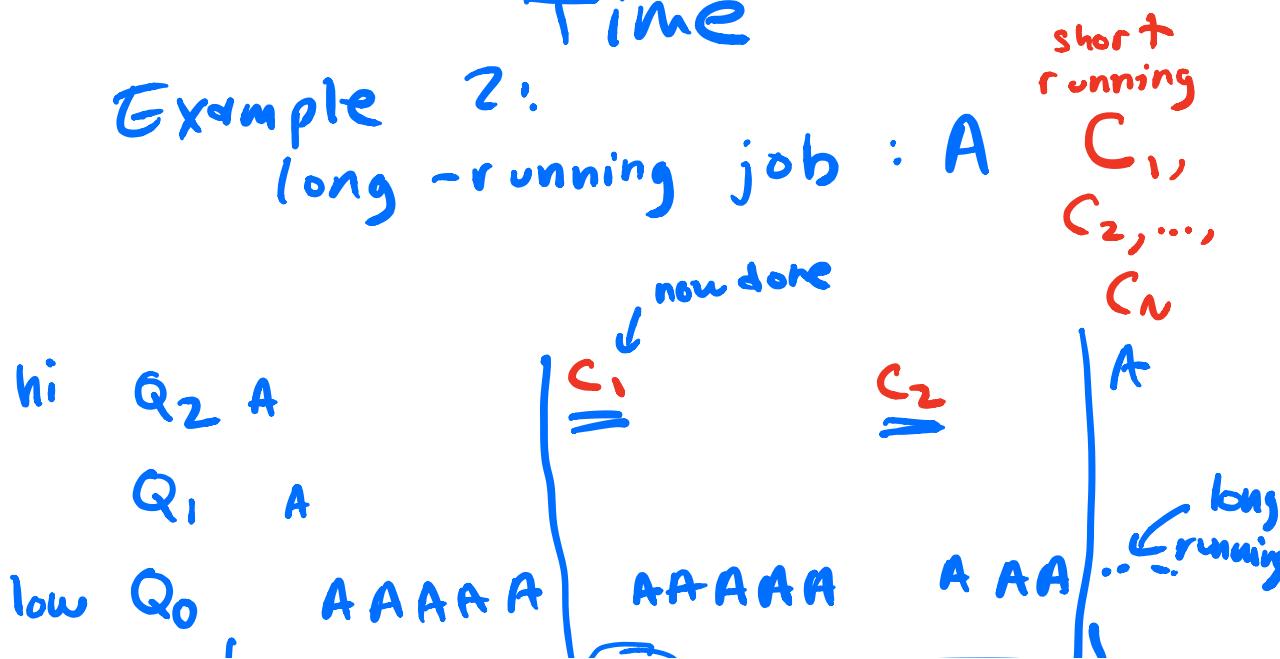
2) if Pri(A) = Pri(B)
 \rightarrow Round Robin between them

3) Start: Highest Priority

4) if process uses time slice
 @ given priority,
 \Rightarrow at end of time slice,
 move down one level



Example 2:
 long-running job : A



c_1, \dots, c_n : $\frac{1}{T}$ large # of short jobs!
=> never ending
=> what happens to A?
(long running)?
=> [Starvation]

How to ensure long-running jobs make progress?

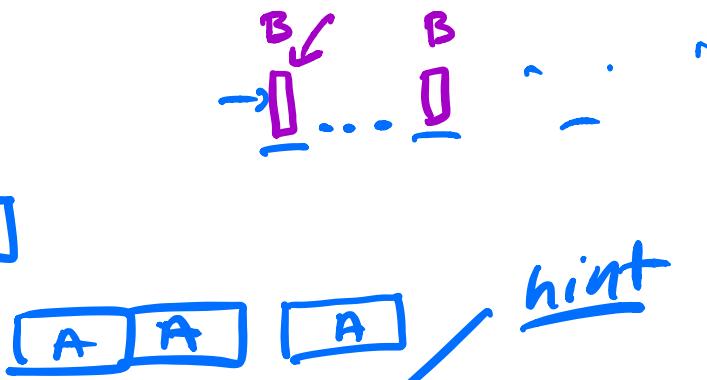
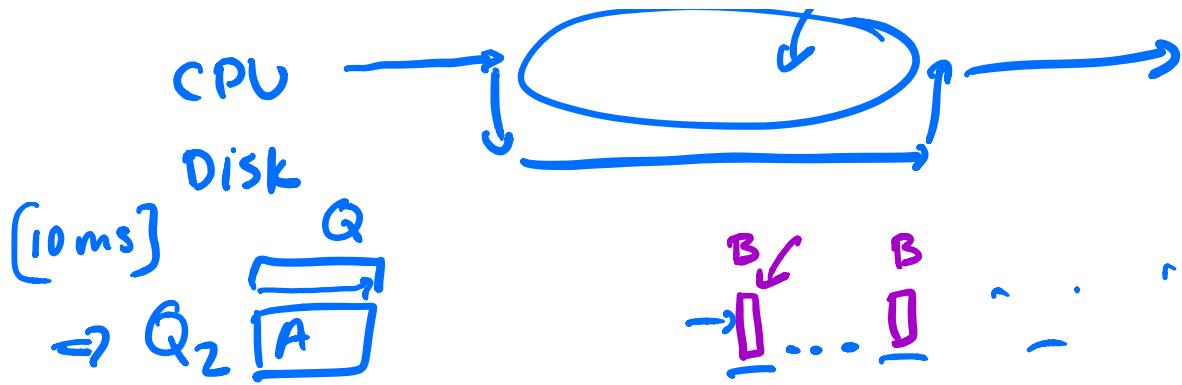
General idea: long-running need to move up

Rule: Every T seconds, move all jobs to highest priority

=> nature of job might change between interactive and batch

new worry: [I/O]

run something else here



A : long running
 B : I/O job

naive rule for I/O:
 if job runs for
 less than timeslice,
stay at same level

=> "Gaming" the scheduler



better accounting: if job uses up
quantum, moves
 down

Parameters: ML FQ

Q_{N-1}

:

Q_0



N Queues

Q_i

quantum
length
(per queue)

T

reset
period

\Rightarrow How to configure:

\Rightarrow mystery

(use defaults)

\Rightarrow new:
ML

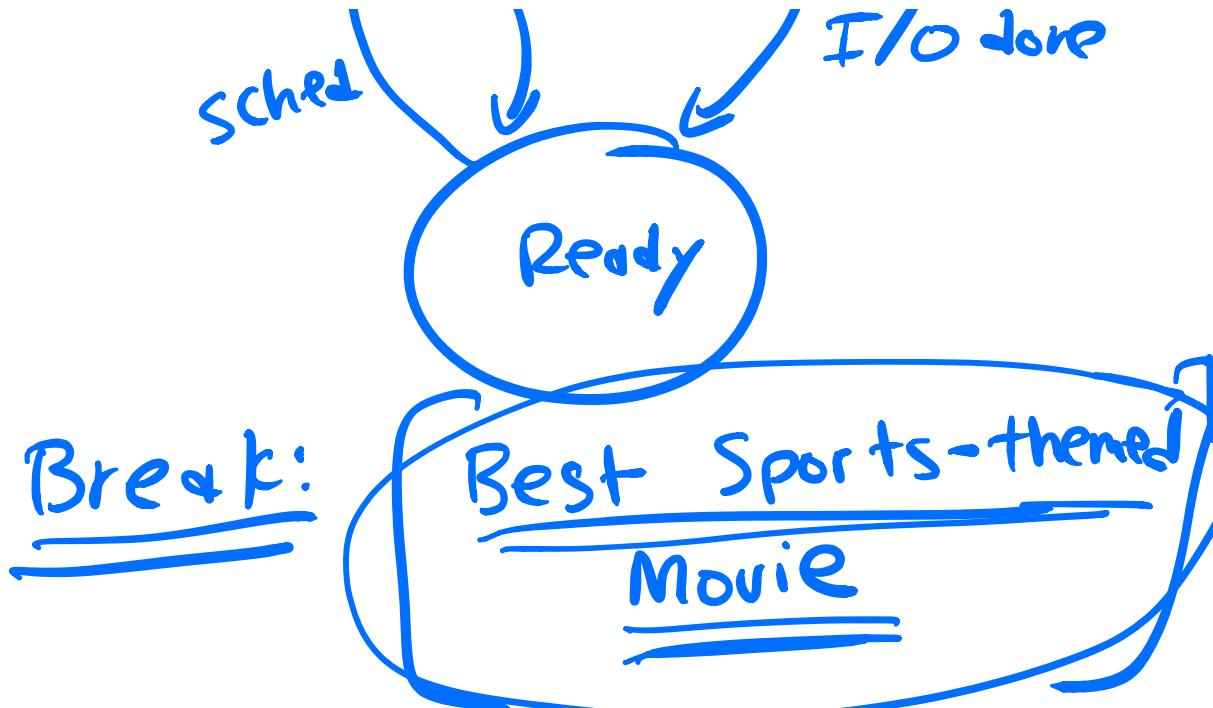
account

Run
(CPU)

issued
I/O

Wait
on I/O

desched



- $\Rightarrow \cancel{\text{Waterboy}} \Rightarrow \cancel{\text{Caddyshack}}$
- $\Rightarrow \cancel{\text{Sanctum}} \Rightarrow \cancel{\text{Dodgeball}}$
- $\Rightarrow \cancel{\text{Miracle}} \Leftarrow \Rightarrow \cancel{\text{Shaolin Soccer}}$
- $\Rightarrow \cancel{\text{Space Jam}} \Rightarrow \cancel{\text{Blades of Glory}}$
- $\Rightarrow \cancel{\text{Remember the Titans}} \Rightarrow \cancel{\text{Mafarkland}}$
- $\Rightarrow \cancel{\text{Balls of Fury}}$



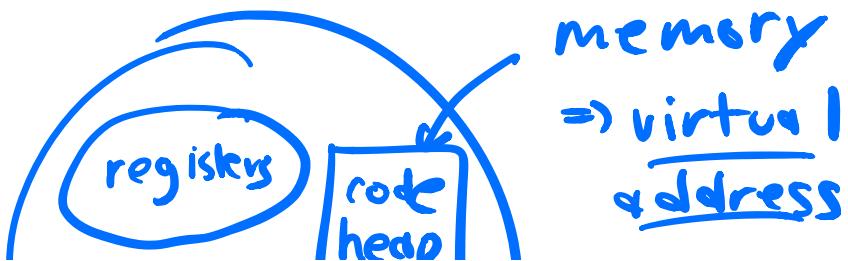
Bad news: class tomorrow

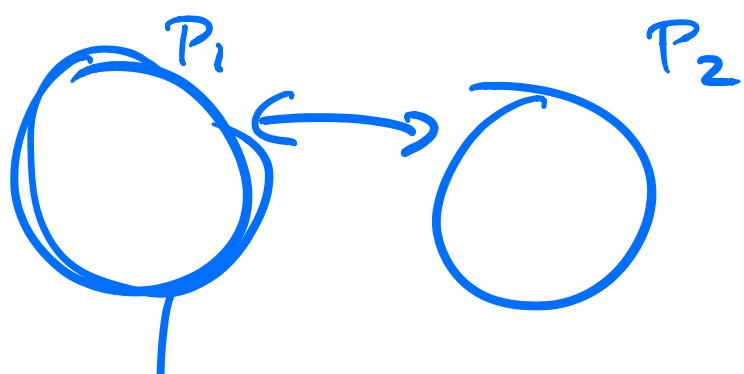
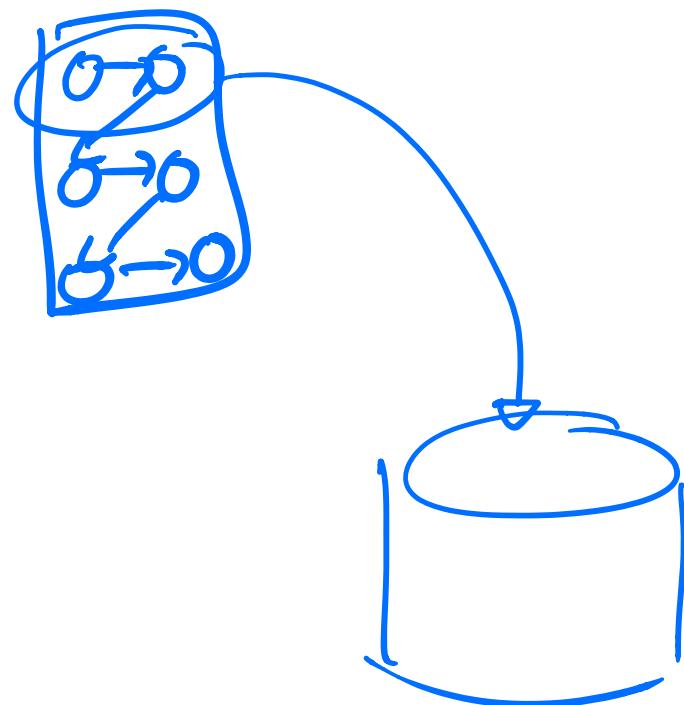
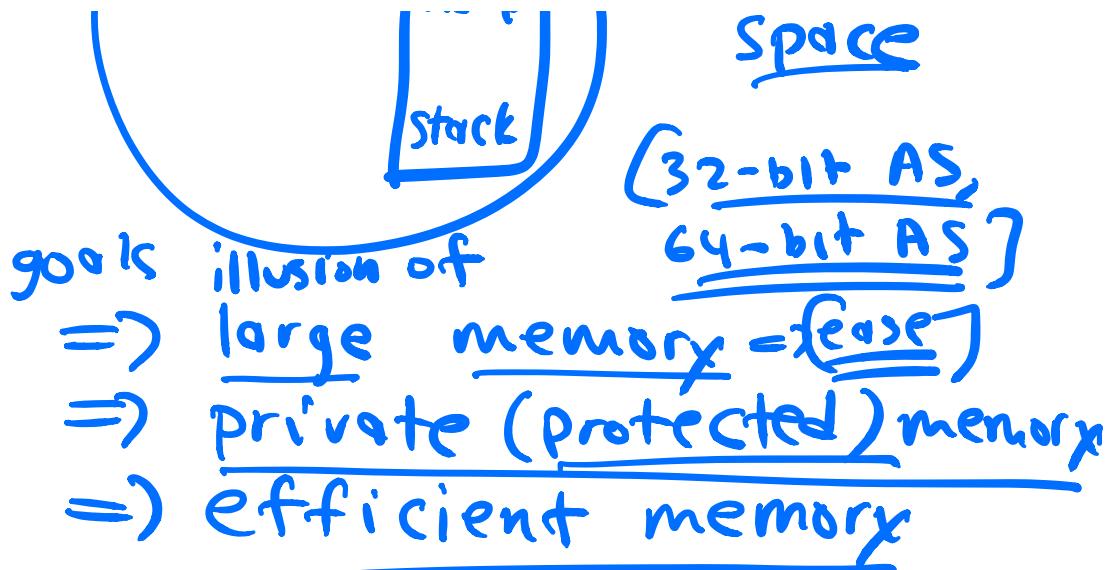
knowledge lack of knowledge

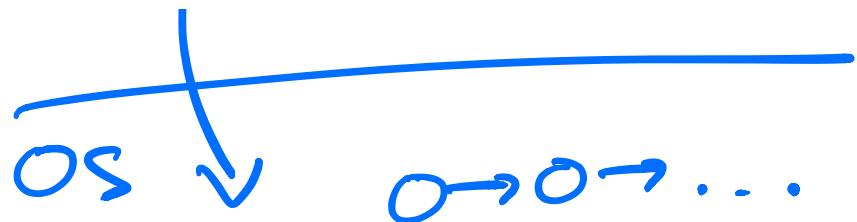
<u>X X B</u>	<u>4 5 6 7</u>	<u>knowledge</u>
<u>8 9 10 11 12 13 14</u>	<u> }</u>	

Virtual Memory: Linux CFS } later

Process: running program







efficiency : limited
Direct Execution

(loads, stores,
inst fetch)

Mechanisms : [s/w => software]
→ h/w → OS support
(hardware)

Memory Accesses:

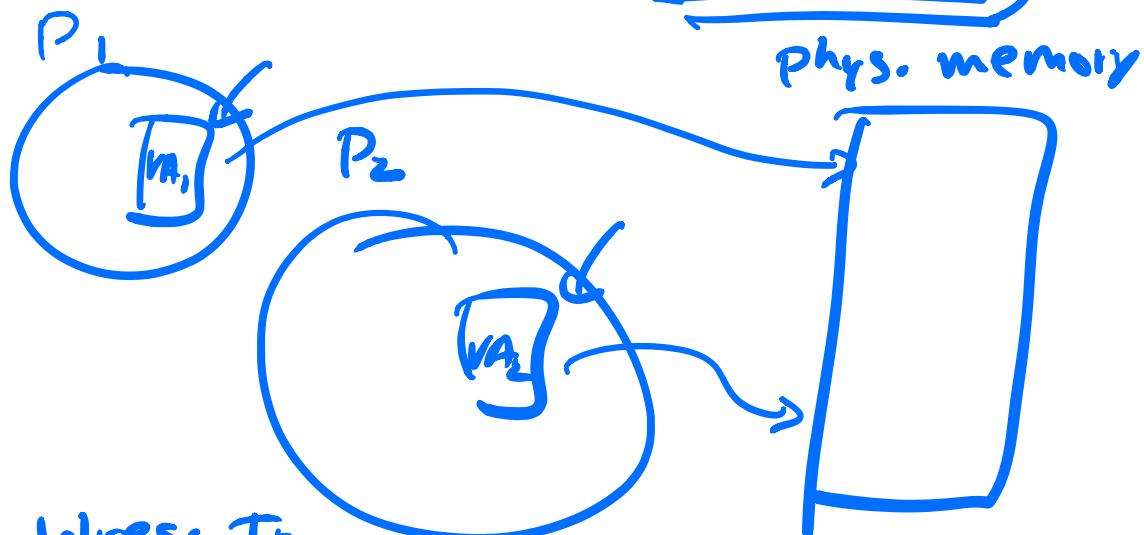
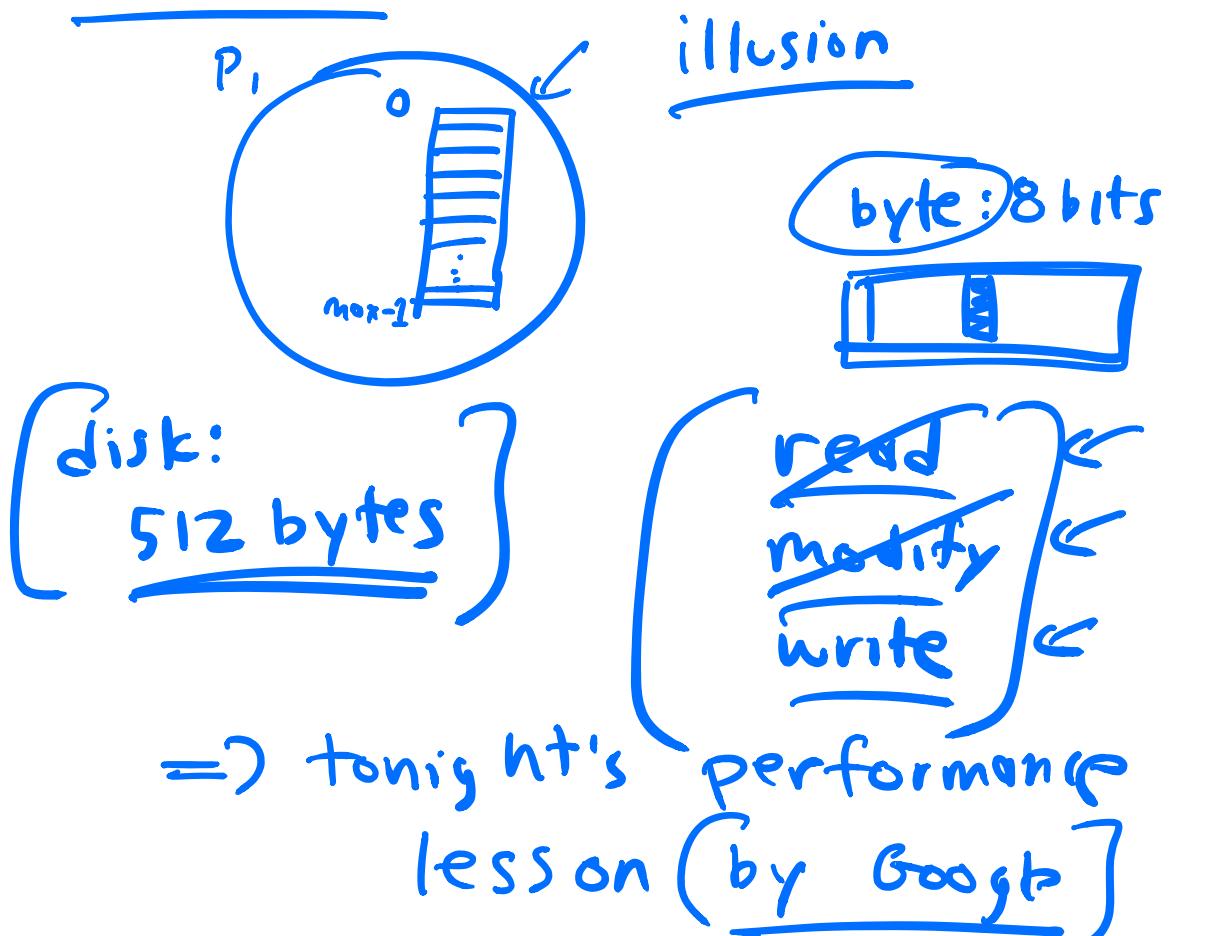
→ instruction fetch

→ explicit loads, stores

CPU ← mem CPU → mem

need: interpose

Address Translation:

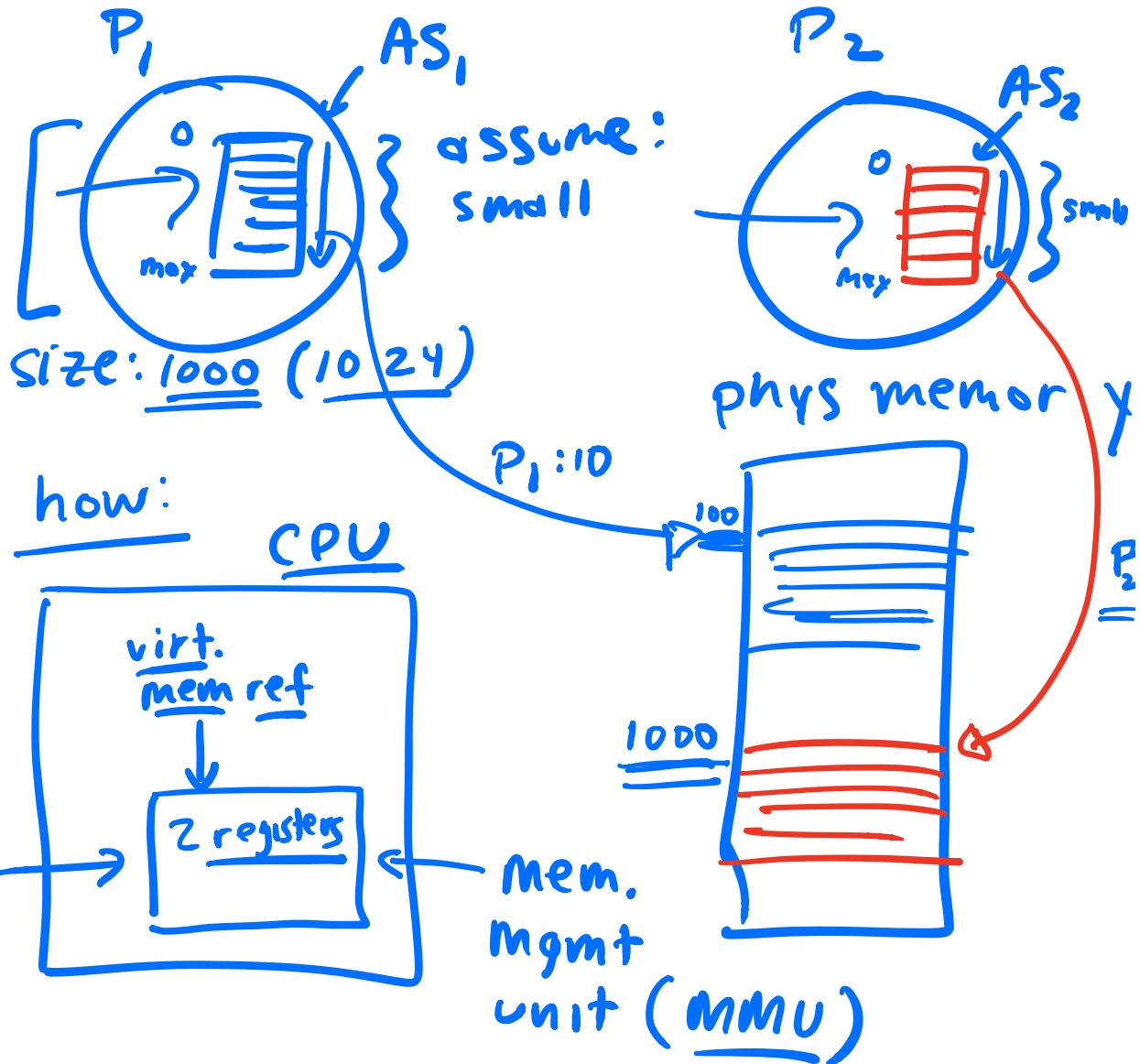


Address Translation
=> on every mem. reference,

translate virtual memory addr:
 \Rightarrow phys memory addr

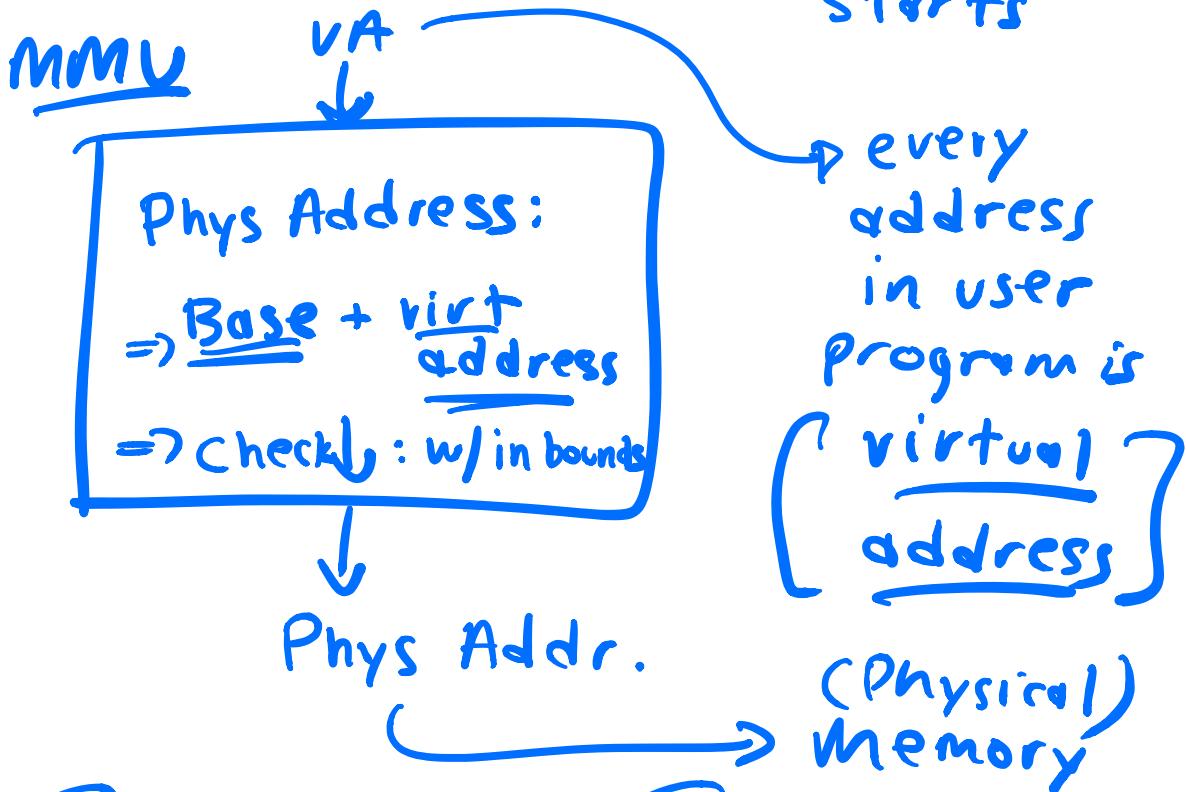
Mechanism #1:

Dynamic Relocation ("Base / Bounds")



(2 registers / CPU)

- base: address in phys mem where AS of currently running process starts
- bounds:

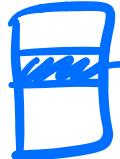


Modern OS: [LATER]

sharing of memory

P₁

P₂



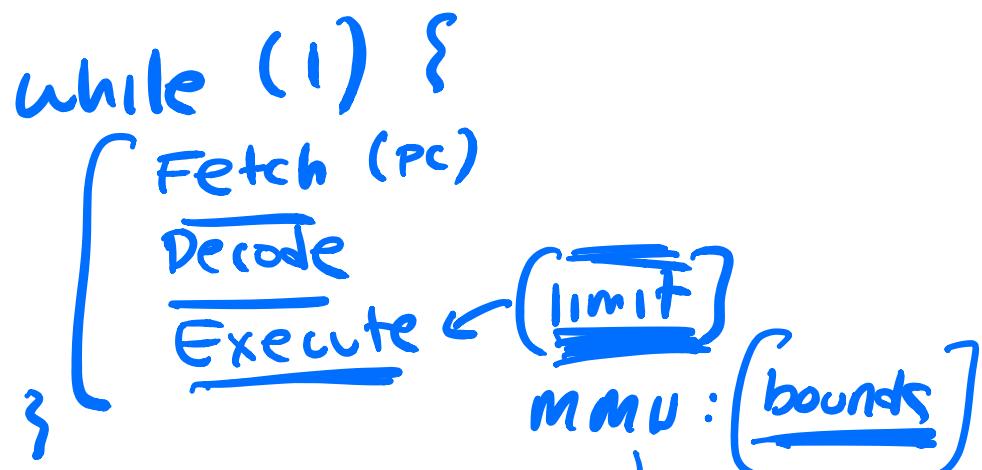


=> Lesson #2 :
(sometimes)
When ?

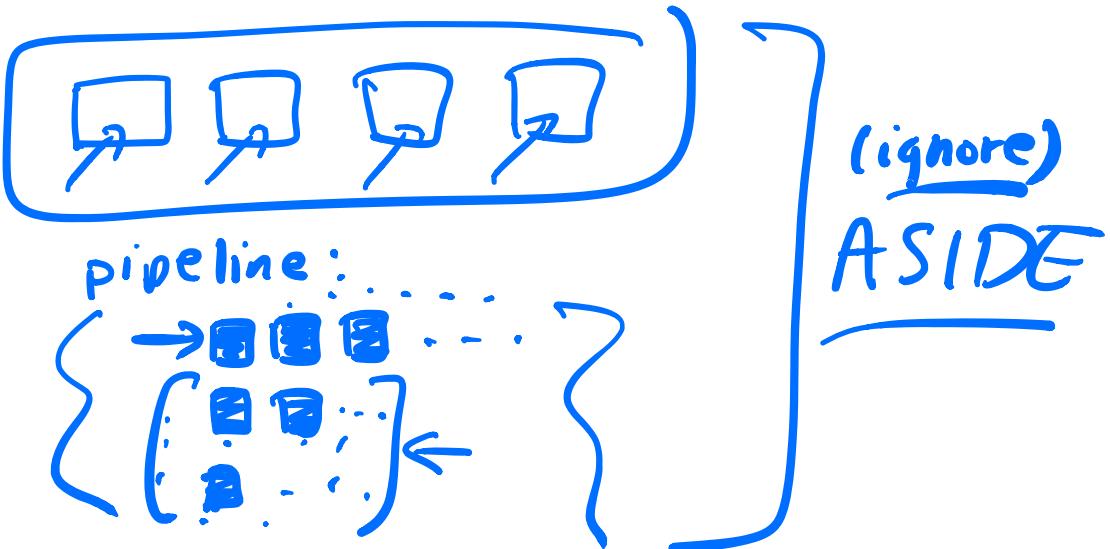
Base: Translation

Bounds: Protection
(check w/in Addr. Space)

CPU (abstractly):



if OK: go ahead
if not OK: raise exception



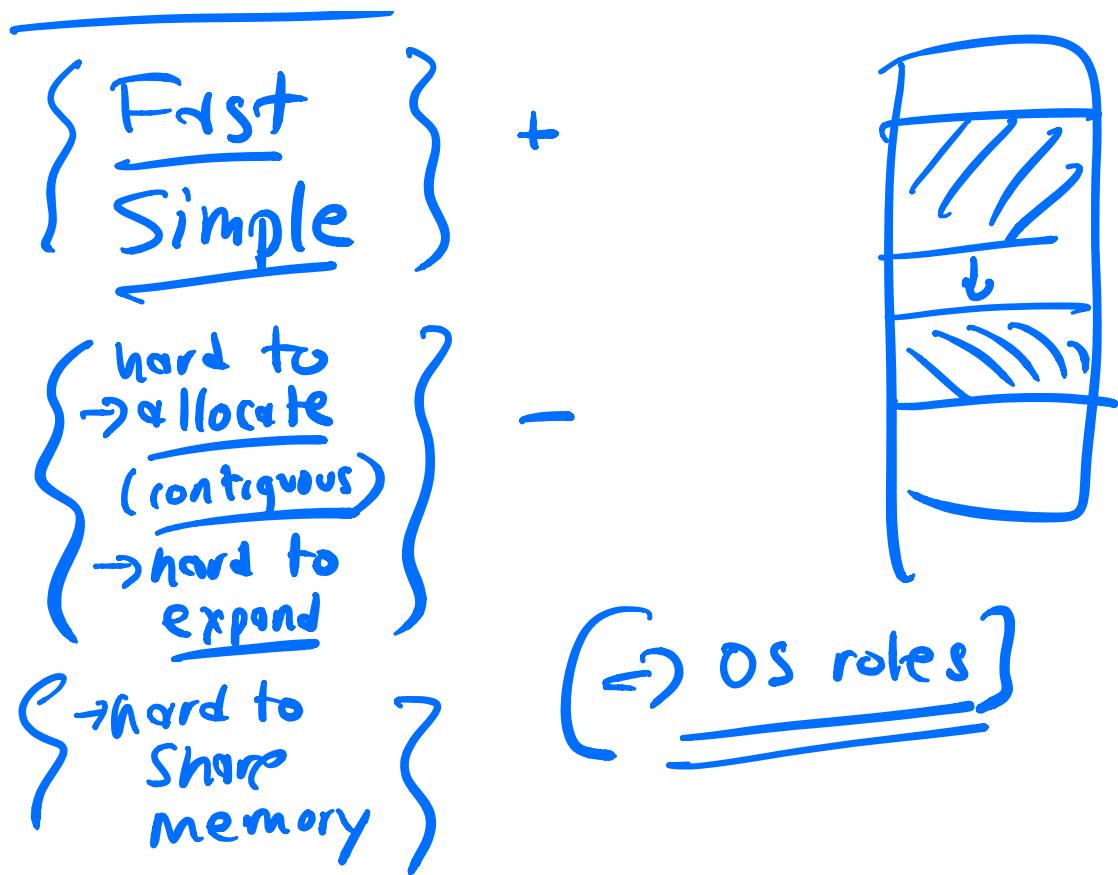
not OK: h/w raises exception
(illegal memory access)

=> H/W: OS → Process : bad
mem access

@ boot: OS set up exception
handlers

=> OS : [kill process]

Base / Bounds :



Question #2 : { Best non-sports
themed movie }