

Today (3/13)

→ 1 Handout @ Front

→ use handout from
last time (first)

Concurrency: Primitives

→ Locks [Data Races]
→ Condition variables
[Control Races]

⇒ wait, signal

What's Left?

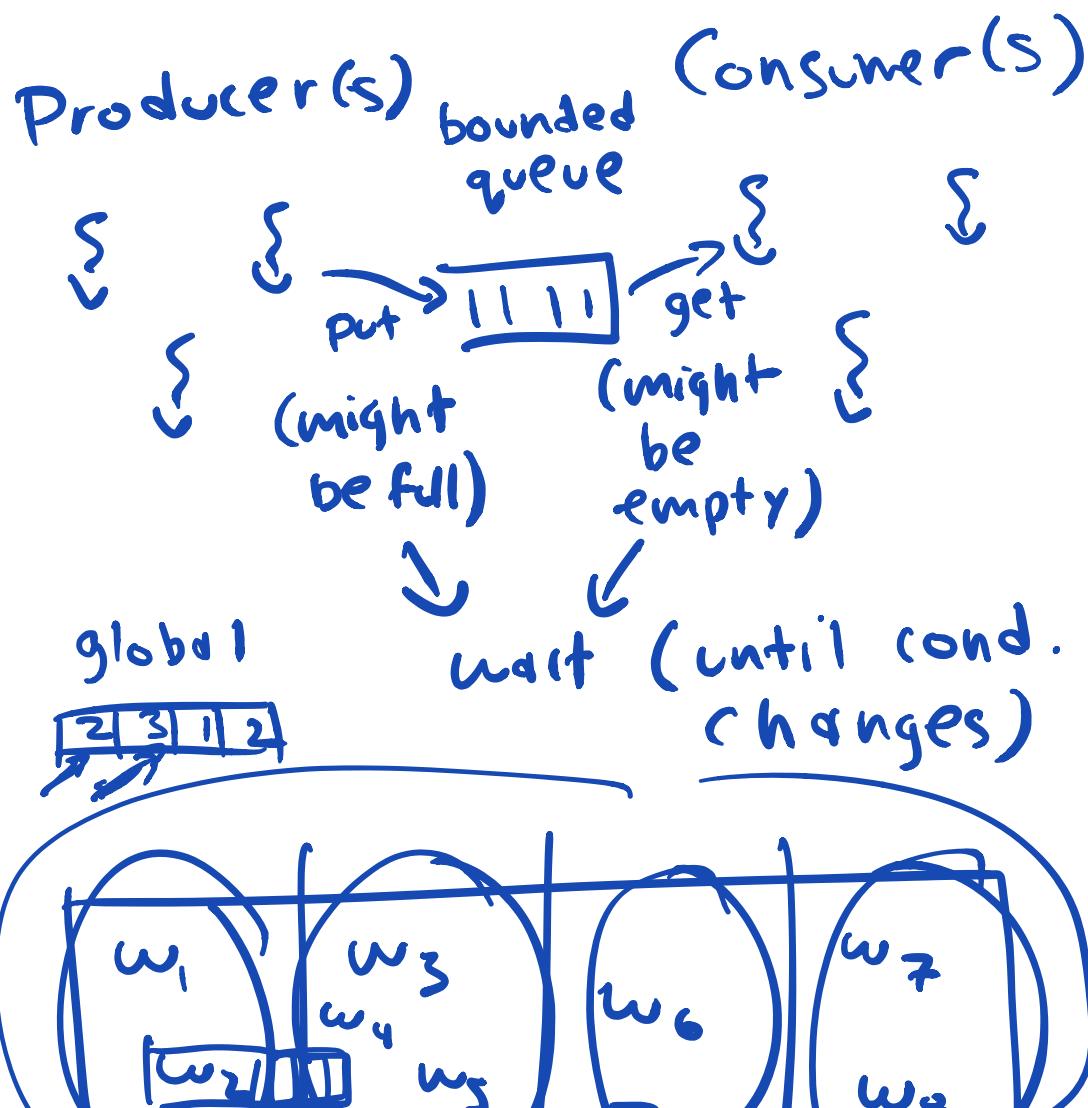
⇒ CV: Producer / Consumer
(useful in pzip)

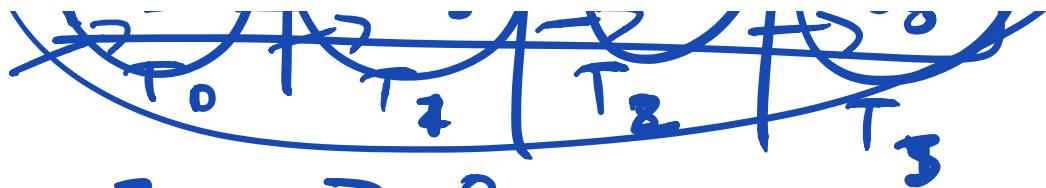
⇒ Semaphore : Historical

Next time:
Deadlock

Then: mid term

Producer / consumer:

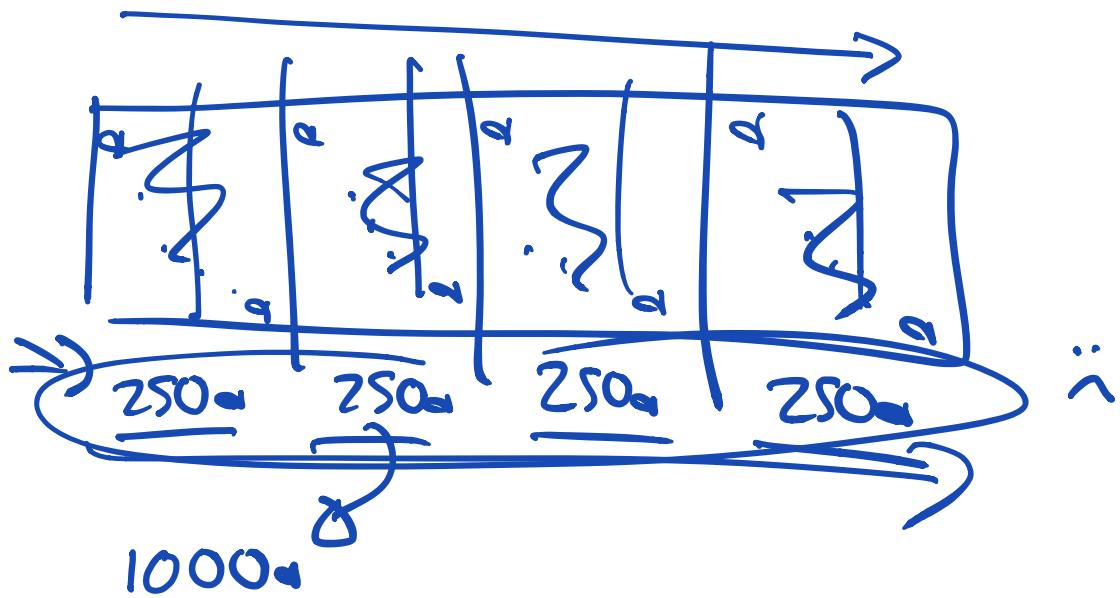




$Z \Rightarrow 8$

→ after 3 1 2
→ Wait (join)

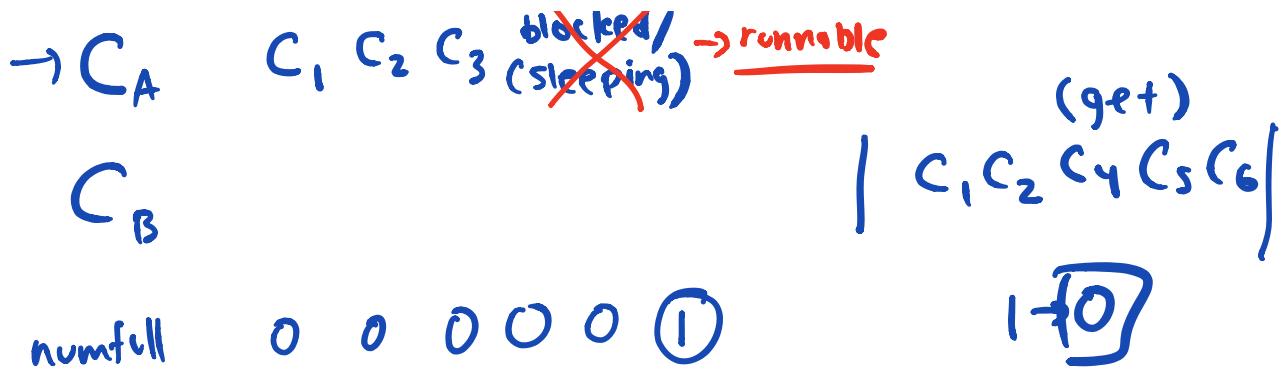
Main: add up the results



Find a problem w/
solution 2:

P

$P_1, P_2, P_4, P_5, P_6, |$
... /



when C_A runs:

calls do-get()

w/ empty buffer

when wake up:

condition might have
changed from the time
 when signal took place

\Rightarrow use while to
 recheck condition

Admin:

$\overbrace{\Rightarrow \text{contest winners}}$
 (congrats!)

=> midterm:

covers a lot

(virt, concurrency)

=> next week:
review

midterm;

march 22, => pages 370

thurs@ "light reading"

7:15 pm → done

dead, or

other skip!

(q: 15pm)

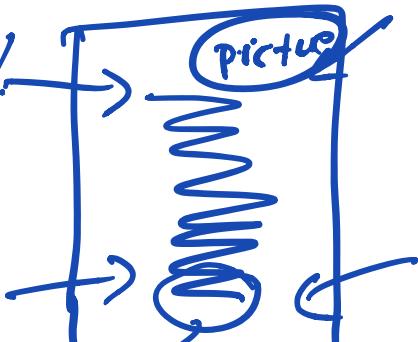
closed book,
closed notes

but: both sides

1 page

"cheat sheet"

~~reference~~



→ → → →

skip!

awesomeness
negotiations:

($8\frac{1}{2} \times 11$)
~~feet~~
~~meters~~

hand written,
printed

laser etched,

....

stickers OK!

you can
do it!

keep
going!

NO

no
computers

NO

format ? $\Rightarrow \left\{ \begin{array}{l} \text{practice} \\ \text{exams} \\ \text{yes?} \end{array} \right\}$

Semaphore : (Dijkstra
on
OS construction
single replacement
for both
"T.H.E."
locks + CVs \Rightarrow concurrency

abstract object w/ two
operations:

- wait
- post

determines
behavior

semaphore has value

could use

locks +
CUs

OR

semaphores

me

you?

use sema as
lock

=> key: init properly

T_1, T_2 call wait / post

properly

~~① sem-init(3lock);
//critical section
T1 → counter ++;
② sem-post(3lock);~~

take turns:
have to wait

→ sem-init(3lock, 1);

Fork / join : sem + cv;

```
child() {  
    // do work  
    sem-post()?  
}  
  
parent() {  
    sem-init(?cv,  
    ?D):  
    thread-create(child);  
    sem-wait()?  
}
```

⇒ yeah!

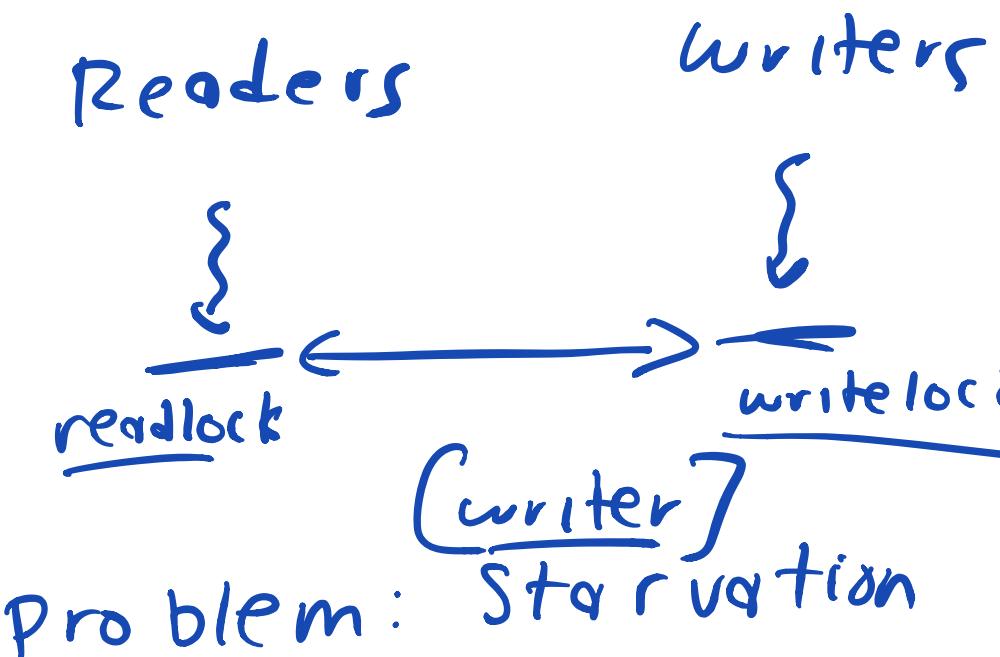
Semaphores
P(), V()

reader/writer locks:

⇒ many readers can
be in crit. section
at once

OR

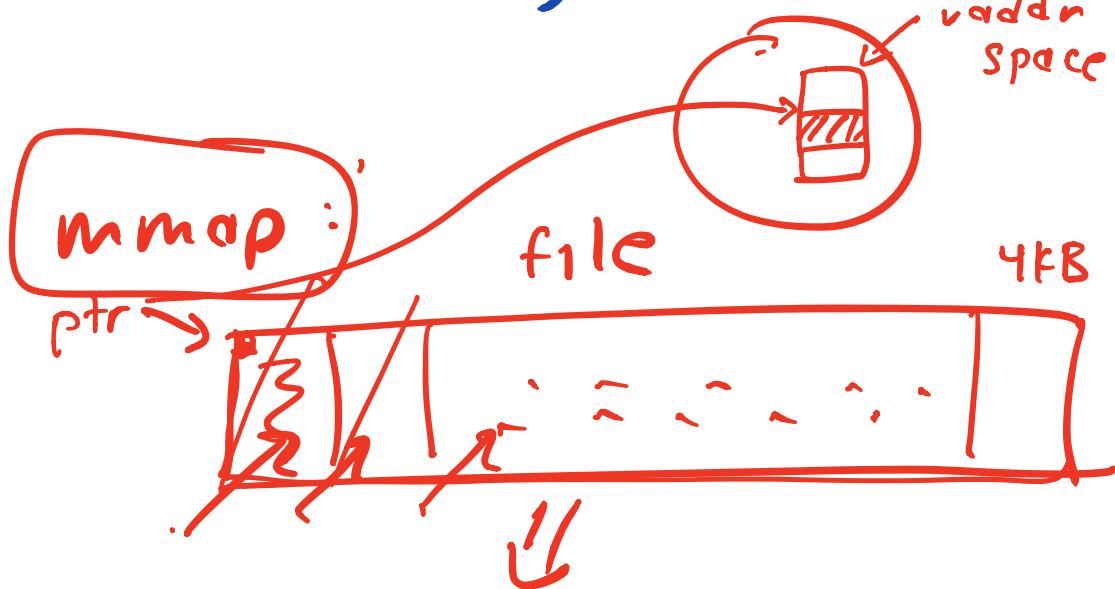
one writer



Projects : Open Questions

- Time benefits of zipping while writing?
- Pzip design pattern
- How to avoid fgetc()?

- ✓ Assume files fit into memory?
 - ✓ Zip: any characters excluded?
 - Optimal write buffer size?
- part b:
 → accessing protection bits



ptr = mmap (...);

faults in pages of
 file lazily

char *

ptr → file
single threaded: ↓
for ($i=0; i < \text{size}; i++$)
*ptr // deref

ptr++;

}

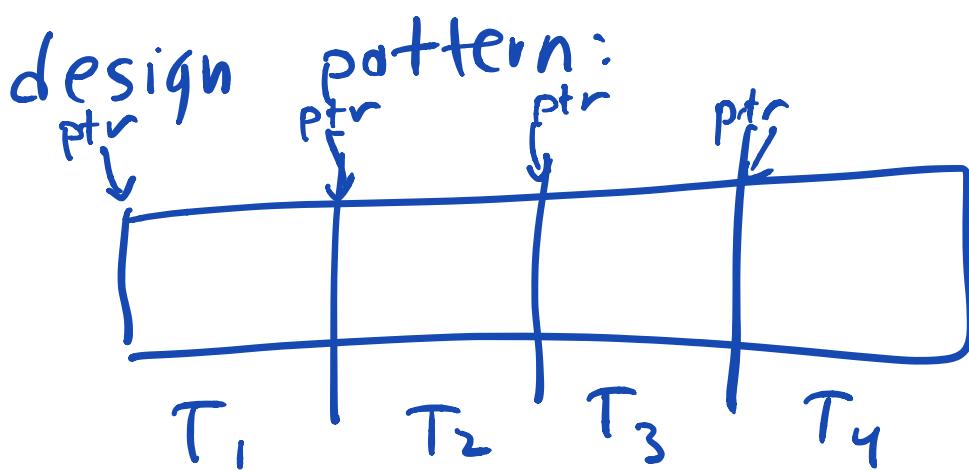
write buffer size?

write(fd, buffer, size)
system call if very
overhead small:

=> how big?

be empirical

try 4KB, 64KB, ...

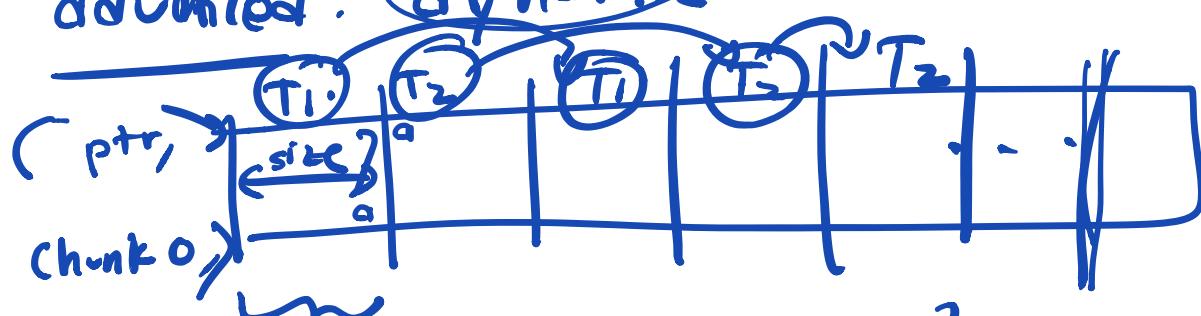


simple:

create threads

assign static work to each
wait to zip each piece
stitch together final results

advanced: dynamic



chunk size : how big?

1 MB

main:

create descriptions ↗
 for work

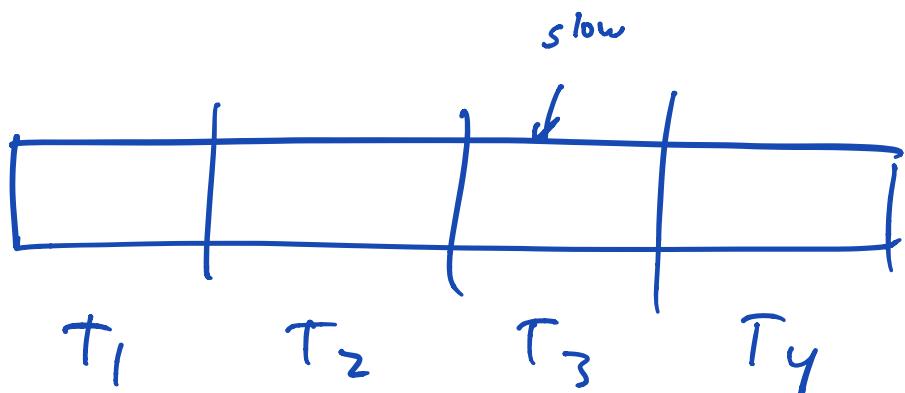
waiting for workers to
 be done
 → stitch together results
 workers:
 get a work desc,
do work: until signal:done
 put result somewhere
 return

output array: 1 entry per chunk
 $\rightarrow \boxed{[3|3] \dots |}$

list

of zipped
contents

$3_0 \rightarrow 4_b \rightarrow 16_m \rightarrow 2_r$



12

14

13
11
12
13
14
15

15