Comparison and Contrast between the OSI and TCP/IP Model

Introduction

- This presentation would discuss some comparison and contrast between the 2 main reference models which uses the concept of protocol layering.
- Open System Interconnection Model (OSI)
- Transport Control Protocol /Internet Protocol (TCP/IP)

Introduction

The topics that we will be discussing would be based on the

OSI	TCP / IP
Application (Layer7)	Application
Presentation (Layer6)	
Session (Layer 5)	
Transport (Layer 4)	Transport
Network (Layer 3)	Internet
Data Link (Layer 2)	Subnet
Physical (Layer 1)	

Outline

- Compare the protocol layers that correspond to each other.
- General Comparison
 - Focus of Reliability Control
 - Roles of Host system
 - De-jure vs. De-facto

The Upper Layers

OSI	TCP/IP
Application (Layer7)	Application
Presentation (Layer6)	
Session (Layer 5)	

- Session
- Presentation
- Application

The Session Layer

The Session layer permits two parties to hold ongoing communications called a session across a network.

- Not found in TCP/IP model
- In TCP/IP, its characteristics are provided by the TCP protocol.
 (Transport Layer)

The Presentation Layer

The Presentation Layer handles data format information for networked communications. This is done by converting data into a generic format that could be understood by both sides.

- Not found in TCP/IP model
- In TCP/IP, this function is provided by the Application Layer.
 - **e.g.** External Data Representation Standard (XDR) Multipurpose Internet Mail Extensions (MIME)

The Application Layer

The Application Layer is the top layer of the reference model. It provides a set of interfaces for applications to obtain access to networked services as well as access to the kinds of network services that support applications directly.

- OSI- FTAM, VT, MHS, DS, CMIPTCP/IP FTP, SMTP, TELNET, DNS, SNMP
- Although the notion of an application process is common to both, their approaches to constructing application entities is different.

Approaches use in constructing application entities

The diagram below provides an overall view on the methods use by both the OSI and TCP/IP model.

Directory services Electronic mail Terminal-host File transfer File transfer (Distributed applicationconceptual) FTAM CMIP OSI DS DNS, Application SNMP SMTP Telnet protocols, services Common applicationdevelopment infrastructure Common Common transport transport services services

ISO Approach

- Sometime called Horizontal Approach
- OSI asserts that distributed applications operate over a strict hierarchy of layers and are constructed from a common tool kit of standardized application service elements.
- In OSI, each distributed application service selects functions from a large common "toolbox" of application service element (ASEs) and complements these with application service elements that perform functions specific to given end-user service.

TCP/IP Approach

- Sometime called Vertical Approach
- In TCP/IP, each application entity is composed of whatever set of function it needs beyond end to end transport to support a distributed communications service.
- Most of these application processes builds on what it needs and assumes only that an underlying transport mechanism (datagram or connection) will be provided.

Transport Layer

OSI	TCP/IP
Transport (Layer 4)	Transport (TCP/UDP)

■ The functionality of the transport layer is to provide "transparent transfer of data from a source end open system to a destination end open system" (ISO / IEC 7498: 1984).

Transport Layer

Transport is responsible for creating and maintaining the basic end-to-end connection between communicating open systems, ensuring that the bits delivered to the receiver are the same as the bits transmitted by the sender; in the same order and without modification, loss or duplication

- It takes the information to be sent and breaks it into individual packets that are sent and reassembled into a complete message by the Transport Layer at the receiving node
- Also provide a signaling service for the remote node so that the sending node is notified when its data is received successfully by the receiving node

Transport Layer protocols include the capability to acknowledge the receipt of a packet; if no acknowledgement is received, the Transport Layer protocol can retransmit the packet or time-out the connection and signal an error

- Transport protocols can also mark packets with sequencing information so that the destination system can properly order the packets if they're received out-of-sequence
- In addition, Transport protocols provide facilities for insuring the integrity of packets and requesting retransmission should the packet become garbled when routed.

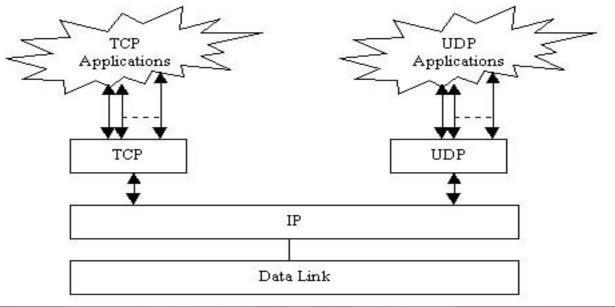
Transport protocols provide the capability for multiple application processes to access the network by using individual local addresses to determine the destination process for each data stream

- Defines two standard transport protocols: TCP and UDP
- TCP implements a reliable datastream protocol
 - connection oriented
- UDP implements an unreliable data-stream
 - connectionless

- TCP provides reliable data transmission
- UDP is useful in many applications
 - eg. Where data needs to be broadcasted or multicasted
- Primary difference is that UDP does not necessarily provide reliable data transmission

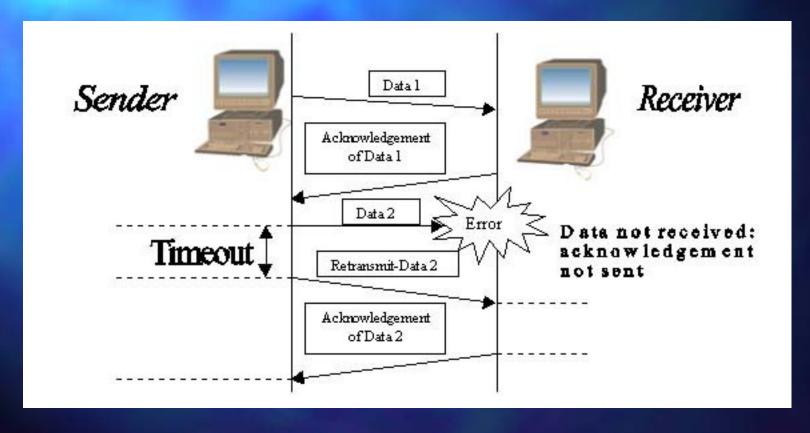
Many programs will use a separate TCP connection as well as a UDP

connection

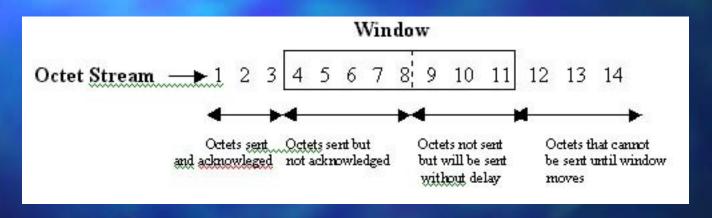


- TCP is responsible for data recovery
 - by providing a sequence number with each packet that it sends
- TCP requires ACK (ackowledgement) to ensure correct data is received
- Packet can be retransmitted if error detected

Use of ACK



- Flow control with Window
 - via specifying an acceptable range of sequence numbers



- TCP and UDP introduce the concept of ports
- Common ports and the services that run on them:

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FTP 21 and 20
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■ telnet 23

■ SMTP 25

http80

■ POP3 110

- By specifying ports and including port numbers with TCP/UDP data, multiplexing is achieved
- Multiplexing allows multiple network connections to take place simultaneously
- The port numbers, along with the source and destination addresses for the data, determine a socket

Comparing Transport for both Models

The features of UDP and TCP defined at TCP/IP Transport Layer correspond to many of the requirements of the OSI Transport Layer. There is a bit of bleed over for requirements in the session layer of OSI since sequence numbers, and port values can help to allow the Operating System to keep track of sessions, but most of the TCP and UDP functions and specifications map to the OSI Transport Layer.

Comparing Transport for both Models

■ The TCP/IP and OSI architecture models both employ all connection and connectionless models at transport layer. However, the internet architecture refers to the two models in TCP/IP as simply "connections" and datagrams. But the OSI reference model, with its penchant for "precise" terminology, uses the terms connection-mode and connection-oriented for the connection model and the term connectionless-mode for the connectionless model.

OSI	TCP/IP
Network (Layer 3)	Internet

Like all the other OSI Layers, the network layer provides both connectionless and connection-oriented services. As for the TCP/IP architecture, the internet layer is exclusively connectionless.

X.25 Packet Level Protocol – OSI's Connection-oriented Network Protocol The CCITT standard for X.25 defines the DTE/DCE interface standard to provide access to a packetswitched network. It is the network level interface, which specifies a virtual circuit (VC) service. A source host must establish a connection (a VC) with the destination host before data transfer can take place. The network attempts to deliver packets flowing over a VC in sequence.

- Connectionless Network Service
 - Both OSI and TCP/IP support a connectionless network service: OSI as an alternative to network connections and TCP/IP as the only way in use.
- Internetworking Protocols
 - OSI's CLNP (ISO/IEC 8473: 1993) is functionally identical to the Internet's IP (RPC 791). Both CLNP and IP are best-effort-delivery network protocols. Bit niggling aside, they are virtually identical. The major difference between the two is that CLNP accommodates variable-length addresses, whereas IP supports fixed, 32-bit address.

Internet (IP) Addresses

The Internet network address is more commonly called the "IP address." It consists of 32 bits, some of which are allocated to a high-order network-number part and the remainder of which are allocated to a loworder host-number part. The distribution of bits - how many form the network number, and how many are therefore left for the host number - can be done in one of three different ways, giving three different classes of IP address

OSI Network Layer Addressing

ISO/IEC and CCITT jointly administer the global network addressing domain. The initial hierarchical decomposition of the NSAP address is defined by (ISO/IEC 8348). The standard specifies the syntax and the allowable values for the high-order part of the address - the Initial Domain Part (IDP), which consists of the Authority and Format Identifier (AFI) and the Initial Domain Identifier (IDI) - but specifically eschews constraints on or recommendations concerning the syntax or semantics of the domain specific part (DSP).

- OSI Routing Architecture
 - End systems (ESs) and intermediate systems (ISs) use routing protocols to distribute ("advertise") some or all of the information stored in their locally maintained routing information base. ESs and ISs send and receive these routing updates and use the information that they contain (and information that may be available from the local environment, such as information entered manually by an operator) to modify their routing information base.

TCP/IP Routing Architecture

The TCP/IP routing architecture looks very much like the OSI routing architecture. Hosts use a discovery protocol to obtain the identification of gateways and other hosts attached to the same network (subnetwork). Gateways within autonomous systems (routing domains) operate an interior gateway protocol (intradomain IS-IS routing protocol), and between autonomous systems, they operate exterior or border gateway protocols (interdomain routing protocols). The details are different but the principles are the same.

TCP/IP
Subnet

Data link layer

The function of the *Data Link Layer* is "provides for the control of the physical layer, and detects and possibly corrects errors which may occur" (IOS/IEC 7498:1984). In another words, the Data Link Layer transforms a stream of raw bits (0s and 1s) from the physical into a data frame and provides an error-free transfer from one node to another, allowing the layers above it to assume virtually error-free transmission

- Physical layer
 - The function of the *Physical Layer* is to provide "mechanical, electrical, functional, and procedural means to activate a physical connection for bit transmission" (ISO/IEC 7498:1984). Basically, this means that the typical role of the physical layer is to transform bits in a computer system into electromagnetic (or equivalent) signals for a particular transmission medium (wire, fiber, ether, etc.)

- Comparing to TCP/IP
 - These 2 layers of the OSI correspond *directly* to the subnet layer of the TCP/IP model.
 - Majority of the time, the lower layers below the Interface or Network layer of the TCP/IP model are seldom or rarely discussed. The TCP/IP model does nothing but to high light the fact the host has to connect to the network using some protocol so it can send IP packets over it. Because the protocol used is not defines, it will vary from host to host and network to network

- Comparing to TCP/IP
 - After much deliberation by organizations, it was decided that the Network Interface Layer in the TCP/IP model corresponds to a combination of the OSI Data Link Layer and network specific functions of the OSI network layer (eg IEEE 203.3).
 - Since these two layers deal with functions that are so inherently specific to each individual networking technology, the layering principle of grouping them together related functions is largely irrelevant.

General Comparison

- Focus of Reliability Control
- Roles of Host System
- De-jure vs. De-facto

Focus of Reliability Control

- Implementation of the OSI model places emphasis on providing a reliable data transfer service, while the TCP/IP model treats reliability as an end-to-end problem.
- Each layer of the OSI model detects and handles errors, all data transmitted includes checksums. The transport layer of the OSI model checks source-to-destination reliability.
- In the TCP/IP model, reliability control is concentrated at the transport layer. The transport layer handles all error detection and recovery. The TCP/IP transport layer uses checksums, acknowledgments, and timeouts to control transmissions and provides end-to-end verification.

Roles of Host System

Hosts on OSI implementations do not handle network operations (simple terminal), but TCP/IP hosts participate in most network protocols. TCP/IP hosts carry out such functions as endto-end verification, routing, and network control. The TCP/IP internet can be viewed as a data stream delivery system involving intelligent hosts.

De-jure vs. De-facto (OSI)

OSI

- Standard legislated by official recognized body. (ISO)
- The OSI reference model was devised before the protocols were invented. This ordering means that the model was not biased toward one particular set of protocols, which made it quite general. The down side of this ordering is that the designers did not have much experience with the subject and did not have a good idea of which functionality to put in which layer.
- Being general, the protocols in the OSI model are better hidden than in the TCP/IP model and can be replaced relatively easily as the technology changes.
- Not so widespread as compared with TCP/IP. (complex , costly)
- More commonly used as teaching aids.

De-jure vs. De-facto (TCP/IP)

TCP/IP

- Standards adopted due to widespread use. (Internet)
- The protocols came first, and the model was really just a description of the existing protocols. There was no problem with the protocols fitting the model, but it is hardly possible to be use to describe other models.
- "Get the job done" orientation.
 Over the years it has handled most challenges by growing to meet the needs.
- More popular standard for internetworking for several reasons :
 - relatively simple and robust compared to alternatives such as OSI
 - available on virtually every hardware and operating system platform (often free)
 - the protocol suite on which the Internet depends.

The End



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