

Process Field Bus

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Timeline

1986

- Master development plan “fieldbus” created in Germany
- 21 companies, including Siemens, involved

1989

- First promoted by *Bundesministerium für Bildung und Forschung (BMBF)*
- Goal to implement a bit-serial field bus for factory and process automation

1999

- Published openly as part of standard **IEC 61158**: *Digital data communication for measurement and control - Fieldbus for use in industrial control systems*

System Structure: Introduction

- Profibus is a multi-master system
- Operation of multiple systems over a single bus
- Three protocols available
 - FMS (*field-bus message specification*)
 - DP (*decentralized peripheral*)
 - PA (*process automation*)
- Devices are categorized in different types
 - Masters
 - Slaves

System Structure: Layer

Profibus in the OSI reference model [1]

Layer	Name	Content
Layer 8	User Layer	Profiles
Layer 7	Application Layer	FMS / DP / PA protocol
Layer 2	Data Link Layer	FDL protocol
Layer 1	Physical Layer	Transmission Technology

Device Type: Master

- Active station
- Control the data traffic on the bus
- When having the *bus access token*:
send messages without external requests

Device Type: Slave

- Passive station
- No self-initiated bus access
- Immediate response to data requested by a master
- Can only be controlled by a single master

Physical Layer

- Profibus FMS and DP
 - Mostly using RS 485 transmission
 - Optical transmission via FOC (*fibre optical cable*) possible
- Profibus PA
 - Uses MBP (*Manchester bus powered*), providing power supply

Type	Transmission technology
0	copper cable with RS 485
1	synchronous MBP
2	synthetic FOC
3	glass FOC
4	HCS FOC

Transmission technology (IRC 61784) [1]

Physical Layer: RS 485

- Bus-topology
- Twisted-pair cables with 150Ω
- Data rates from 9.6 kbit/s to 12 Mbit/s
- Distance between repeaters 100 m to 1200 m

Mainly used with *Profibus DP*

Physical Layer: FOC

- Star-, bus or ring-topology
- Fibre optical cables
- Distance between repeaters up to 15km

Physical Layer: MBP

- Bus-topology
- Stations are powered through the bus
- Safe in explosion-hazardous environments, power can be reduced to a bare minimum
- Data rate is fixed to 31.25 kbit/s
- Bus length up to 1900 m
- Allows branches up to 60 m to field devices

Mainly used with *Profibus PA*

Data Link Layer

The data transmission in Profibus is handled by the fieldbus data link (*FDL*) layer.

FDL consists of three functions:

- Medium Access Control (*MAC*)
- Fieldbus Link Control (*FLC*)
- Fieldbus Management (*FMA*)

Data Link Layer: MAC

- Make sure only one station transmits data on the bus
- When multiple masters are present
 - Masters need the access token to send data
 - Token is cyclically passed via token telegram
 - To ensure that all master stations can access the bus, token must be passed on after a certain timeout
- Slaves only respond to requests by a master

Profibus FDL combines master-slave and token passing in a hybrid access principle

Data Link Layer: FMA

Fieldbus Management provides function to manage the layer 1 and 2

- Reset the layers
- Set parameters
- Get parameters
- Inform the user about events or errors
- Activate/Deactivate service access point (*SAP*)

Data Link Layer: Error handling

Errors can be caused by

- Faulty transmitters
- Badly shielded cables
- Signal reflections
- Large divergences in time synchronization between stations

Error rate is smaller than 10^{-4} and can be reduced further by error detection and correction

Data Link Layer: Error detection and correction

Error Detection

- Hamming distance of 4 by adding a checksum to each packet
- At least 4 bits must change to result in an undetected error
- This results in *integrity class 2* after standard **IEC 870-5-1**

Send Data with No acknowledge (SDN) service

- Mainly used for synchronization and status messages
- The erroneous telegram is discarded
- Telegram from the next cycle is used instead

Data Link Layer: Error detection and correction

Send Data with Acknowledge (SDA)

- Mainly used between masters, slaves may not always send an acknowledgement
- When the sender does not get a response, the telegram is retransmitted

Send and Request Data (SRD)

- Service used between masters and slaves
- Acknowledgement is packed on top of the data telegram
- When the sender does not get a response, the telegram is retransmitted

Application Layer: Addressing

- Every station has a unique address, coded in 1 byte

Address	Use
0	reserved for tools, e.g. programming devices
1 – n	n master stations
$n - 125$	slave stations
126	reserved as <i>delivery address</i> used for changing the address of a slave during runtime
127	reserved as broadcast address

Components used for the infrastructure, e.g. repeaters transmit the data transparently and do not require an address

Application Layer: FMS

- FMS master controls the relationship with FMS slaves
- Replaced by Profibus DP

Application Layer: DP

- DP masters are separated into classes
 - Class 1: control a DP system and the slaves assigned, mostly PLC based
 - Class 2: tool for commissioning, engineering and maintenance, mostly PC based
 - Class 3: clock master, used for time synchronization

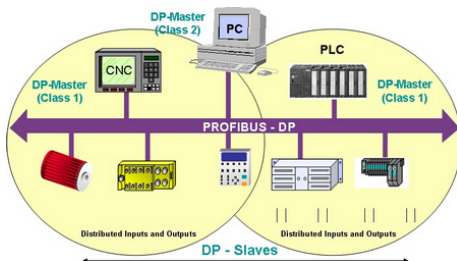


Figure: Structure of a DP system [1]

Application Layer: DP - Cyclic process data

Data exchange between masters and slaves is separated into three phases: [1]

Phase	Action
Diagnosis	Master requests diagnostic data from slaves
Initialization	Master sets parameters and checks configuration of slaves
Data Exchange	Master sends and requests data from the slaves

Application Layer: DP - Data Exchange

Class 1 master station:

- Relationship with a slave is called *MS0*
- Data exchange is cyclic
- Master sends output data to a slave
- Slave immediately responds with input data
- Master continues with the next slave or restarts the cycle

The minimum cycle time T_{BCycle} can be calculated:

$$T_{BCycle} = \frac{380 + (N_{Slaves} \cdot 300) + (N_{Bytes} \cdot 11)}{Bitrate} + 75\mu s \quad (1)$$

Application Layer: DP - Data Exchange

Class 2 master station:

- Can exist in addition class 1 masters
- Can simultaneously be a class 1 master
- Relationship with a slave is called *MS1*
- Acyclic communication with a slave in an existing MS0 relationship

Application Layer: PA

■ foo <+todo+>

<+++>

User Layer

■ foo <+todo+>

<+++>

References I



Max Felser

Profibus Manual: A collection of information explaining PROFIBUS networks

<http://www.profibus.felser.ch>