

Controller Area Network

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In Vehicle Networks

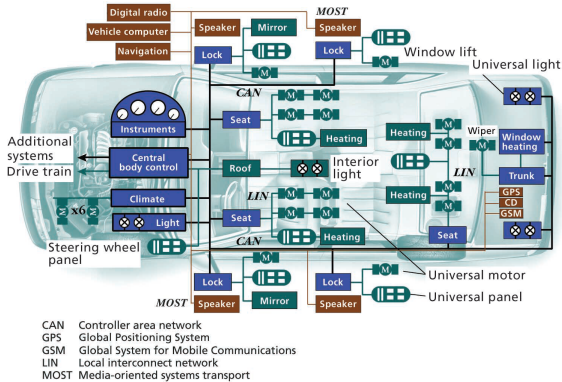


Figure: Vehicular Network Architecture

Timeline

- 1983 • Development started at Robert Bosch GmbH
- 1986 • Official release at SAE congress
- 1987 • First chips by Intel and Philips
- 1991 • Bosch publishes CAN 2.0 for standardization
- 1993 • Approved and published as ISO 11899
- 1995 • Standard split into 11899-1,2,3
- 2012 • CAN FD 1.0 released by Bosch

CAN protocol basics

- Serial Bus System
- Only physical and data link layer specified
- Physical Layer
 - Manages effective transmission of data
 - Bit timing & Synchronization
 - High & Low Speed standards exist
- Data Link Layer
 - Medium Access Control: CSMA/**CR**
 - Logical Link Control: Interface to higher layers
- Numerous application layers rely on CAN

Physical Layer

- Physical signalling specified in ISO 11898-1
- Medium access units (transceivers) described in two documents
 - ISO 11898-2: High-speed CAN
 - ISO 11898-3: Low-speed fault-tolerant CAN
- Connectors are not standardised by CAN specification

Network Topology

- CAN networks use shared bus topology
- Buses have to be terminated
- Topology should be as linear as possible
- Maximum bitrate and length directly dependent
- No effective way to share line for power and signalling

Transmission Medium

Several kinds of transmission media can be used:

- **Two-wire bus:** Enables differential signal transmission, ensures reliable communication. Requirement for high-speed CAN.
- **Single-wire bus:** Simpler/cheaper alternative, fall back in case of fault.
- **Optical transmission medium:** Ensures immunity to electromagnetic noise, used to interconnect different subnets.

ISO 11898-2: High-speed CAN

- Maximum bit rate 1 Mbps
- Linear bus end-terminated with 120Ω
- Stubs must be shorter than 30 cm

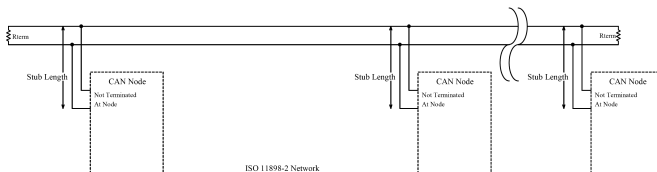


Figure: High Speed CAN Network

ISO 11898-3: Low-speed fault-tolerant CAN

- Maximum bit rate 125 kbps
- Linear or star bus terminated at node with about 100Ω

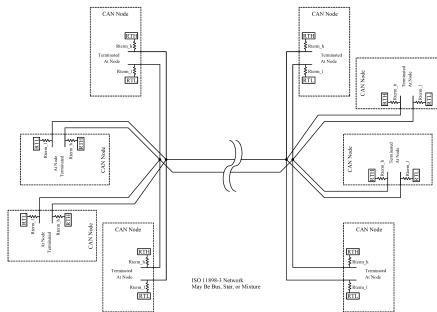


Figure: Low Speed CAN Network

Bit Encoding

- Level on bus can assume two complementary values:
 - *dominant*, usually corresponds to logical value 0
 - *recessive*, usually corresponds to logical value 1
- CAN relies on *non-return to zero* (NRZ) bit encoding

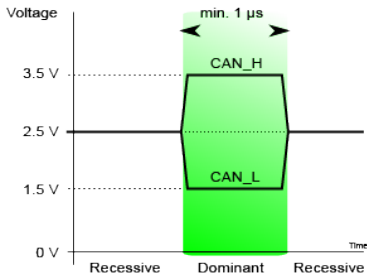


Figure: Levels on the CAN bus

Synchronization

- Timing information extracted from bit stream
- Edges of the signal are used for synchronization
- Bit stuffing to ensure sufficient number of edges

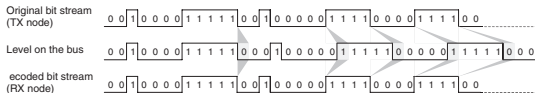


Figure: Bit stuffing technique

Frame Format

- Specifications define standard and extended frame format
 - Standard: 11 bit identifier
 - Extended: 29 bit identifier
- Standard frame format mostly used
- Identifier describes meaning of message
- Protocol foresees four kinds of frames: *data*, *remote*, *error* and *overload*

Data Frames

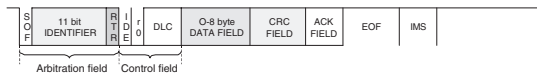


Figure: Format of standard data frames

- Dominant *start of frame* (SOF) bit
- Arbitration field: identifier and *remote transmission request* bit
- *Data length code* (DLC): Length of data field encoded in 4 bits
- *Cyclic redundancy check* (CRC) encoded in 15 bits
- ACK slot: Recessive at transmitter, dominant at receiver
- *End of frame* (EOF) slot: Seven recessive bits

Remote Frames

- Generally, source sends out data autonomously
- Protocol allows to poll for data
- Remote format similar to data format
- RTR field is recessive, data frames have higher priority
- Remote frames carry no data

Error Frames

- Notify nodes that an error has occurred
- Consist of two fields:
 - Error flag: Six dominant/recessive bits.
 - Violates bit stuffing rules, error condition is detected
 - Error delimiter: 8 recessive bits.
- Active flag: dominant, transmitted by node in state *error active*
- Passive flag: recessive, transmitted by node in state *error passive*

Fault Confinement

- Supervises correct operation of MAC sublayer
- Disconnect defective node from bus
- Uses two counters: *transmission* and *receive error count*
 - On error detect, counter is increased by a given amount
 - On success, counter is decreased by one
 - Increase amount of detecting node is higher than relying nodes
- When counter exceeds 127, node switches from error active to error passive
- When counter exceeds 255, node switches to bus off

Overload Frames

- Used to slow down operations on the bus by adding delays
- Format similar to the error frames
- Hardly ever used because today's CAN controllers are very fast

Access Technique

- CAN relies on CSMA for access control:
 - When no data is exchanged, level on the bus is recessive
 - Before transmission, nodes observe the state of the network
 - When network is idle, transmission starts immediately
- Collisions are improbable but not impossible
- CAN introduces collision resolution scheme: Bus arbitration

CSMA/CR: Bus Arbitration

Bus arbitration essentially finds out the most urgent frame.

- Level on bus is dominant if one node is sending dominant bit
- Nodes can reliably check level on bus
- On transmission, each node compares level on bus against written value
- If node transmits recessive bit but reads dominant, it backs off

CSMA/CR: Bus Arbitration

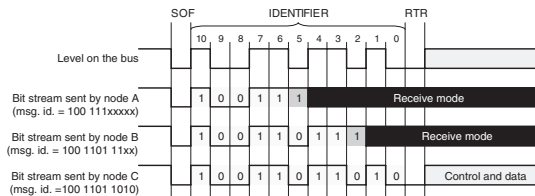


Figure: Arbitration phase in CAN

- Nodes transmit message identifier starting with MSB
- Lowest identifier corresponds to highest priority
- Message with highest priority wins contention

Error Management

- Fundamental requirement for CAN is robustness
- Specifications foresee five mechanisms for error detection:
 - 15 bit wide CRC: Discover up to five erroneous bits
 - Frame check: CRC, ACK, EOF delimiters have to be recessive
 - Acknowledgement check: Transmitter checks for set ACK bit
 - Bit monitoring: Transmitter checks level on bus against written value
 - Bit stuffing: Each node verifies if bit stuffing rules have been violated
- Residual probability for undetected corrupt message is $4.7 \cdot 10^{-11}$ times the frame error rate or less

Logical Link Layer

- Sublayer of Data Link Layer
- Provides communication services to higher layers
- Exports only two types of frames:
 - **L_DATA**: Broadcast value over the network
 - **L_REMOTE**: Ask for value over the network
- Error and overload frames invisible to higher layers
- Provides *frame acceptance filtering* function

Frame Acceptance Filtering (FAF)

- Producer transmits information on the bus
- Frame is read by every node in a receive buffer
- FAF determines if information is relevant to the node

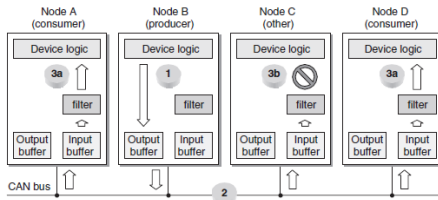


Figure: Producer/consumer model

Higher Layer Implementations

- CAN specifications do not include application layer tasks
 - Flow Control
 - Device Addressing
 - Fragmentation/Defragmentation
- Several higher layer protocols rely on CAN
- Industrial automation: CANopen, DeviceNet
- Passenger cars: Each manufacturer has its own standard

Advantages & Disadvantages of CAN

CAN implements a distributed priority-based multi-master communication system.

■ Advantages:

- Much more simple and robust than token based access schemes
- More flexible than TDMA approaches
- No message will be delayed by lower priority exchanges

■ Drawbacks:

- Relatively low maximum throughput
- Bus length limited by bandwidth, arbitration, timing
- Offers no security or authentication schemes

Thanks for your attention!

Questions? Ideas? Suggestions?

References I



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Operating Principles and Features of CAN Networks.
IEIIT-CNR, 2005.



EE JRW - Own work.

CAN ISO11898-2 Network.

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Canbus levels.

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