

# Process Field Bus

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# Timeline

## 1986

- Master development plan “fieldbus” created in Germany
- 21 companies, including Siemens, involved

## 1989

- First promoted by *Bundesministerium für Bildung und Forschung (BMBF)*
- Goal to implement a bit-serial field bus for factory and process automation

## 1999

- Published openly as part of standard **IEC 61158** *Digital data communication for measurement and control - Fieldbus for use in industrial control systems*

# System Structure: Introduction

- *Profibus* is a multi-master system
- Operation of multiple systems over a single bus
- Devices are categorized in different types
  - Masters
  - Slaves

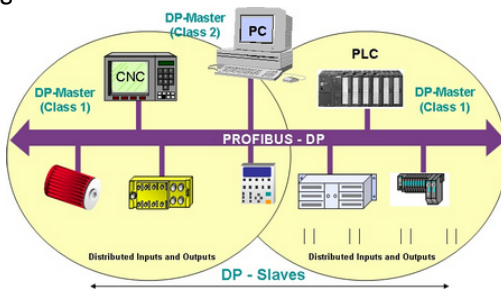


Figure: Structure of a DP system [1]

# System Structure: Master Station

## Master Station

- Active station
- Can control multiple slaves
- Requests data from slaves
- When having the *bus access token*:  
send telegrams on the bus

# System Structure: Slave Station

## Slave Station

- Passive station
- No self-initiated bus access
- Immediate response to data requested by a master
- Can only be controlled by a single master

# System Structure: Layer

## Profibus in the OSI reference model [1]

Layer	Name	Content
Layer 7	Application Layer	FMS / DP / PA protocol
Layer 2	Data Link Layer	Field Data Link (FDL) protocol
Layer 1	Physical Layer	Transmission Technology

## Protocols

- FMS (*Fieldbus Message Specification*)
- DP (*Decentralized Peripheral*)
- PA (*Process Automation*)

# Physical Layer: Transmission Technology

## Profibus FMS and Profibus DP

- Mostly using RS 485 transmission
- Optical transmission via FOC (*fibre optical cable*) possible

## Profibus PA

- Uses MBP (*Manchester bus powered*), providing power supply

Type	Transmission technology
0	copper cable with RS 485
1	synchronous MBP
2	synthetic FOC
3	glass FOC
4	HCS FOC

Transmission technology after standard **IEC 61784 [1]**

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# Physical Layer: RS 485

## RS 485

- Bus-topology
- Twisted-pair cables with  $150\Omega$
- Data rates from  $9.6\text{ kbit/s}$  to  $12\text{ Mbit/s}$
- Distance between repeaters  $100\text{ m}$  to  $1200\text{ m}$
- UART coding
  - Start = 0, Parity = EVEN, Stop = 1

■	Start	Databit 1	2	3	4	5	6	7	8	Parity	Stop
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Mainly used with *Profibus FMS* and *Profibus DP*

# Physical Layer: FOC

## Fibre Optical Cable (FOC)

- Star-, bus- or ring-topology
- Resistant against electromagnetic noise
- Distance between repeaters up to 15km

# Physical Layer: MBP

## Manchester Bus Powered (MBP)

- Bus-topology
- Stations are powered through the bus
- Safe in explosion-hazardous environments, power can be reduced to a bare minimum
- Data rate is fixed to  $31.25\text{ kbit/s}$
- Bus length up to  $1900\text{ m}$
- Allows branches up to  $60\text{ m}$  to field devices

# Physical Layer: MBP

## Manchester Bus Powered (MBP)

- Can be connected to an existing *Profibus DP* network
  - Using a DP/PA coupler
  - The faster *DP* network is used as a backbone
- Uses Manchester coding

# Physical Layer: MBP

## Manchester coding

- Every bit is coded as a change
  - Positive change: "0"
  - Negative change: "1"
  - Every bit has the same average value
  - Average used to power the peripherals
  - Time synchronization possible with every bit

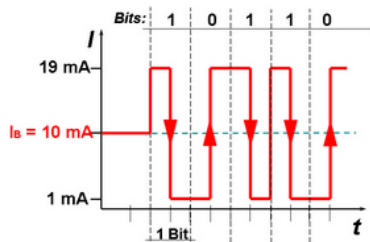


Figure: Manchester coding [1]

Mainly used with *Profibus PA*

# Data Link Layer

The **Fieldbus Data Link (FDL)** layer handles data transmission in *Profibus*

FDL consists of three functions:

- Medium Access Control (*MAC*)
- Fieldbus Link Control (*FLC*)
- Fieldbus Management (*FMA*)

# Data Link Layer: MAC

## Medium Access Control (MAC)

- Make sure only one station transmits data on the bus
- When multiple masters are present
  - Masters need the access token to send data
  - Token is cyclically passed via token telegram
  - To ensure that all master stations can access the bus, token must be passed on after a certain *token rotation time*
- Slaves only respond to requests by a master

*Profibus FDL* combines master-slave and token passing in a hybrid access principle

# Data Link Layer: FMA

**Fieldbus Management (FMA)** provides functions to manage the *Physical Layer (1)* and *Data Link Layer (2)*

- Reset the layers
- Get/set parameters
- Inform the user about events or errors
- Activate/Deactivate *service access point (SAP)*



# Data Link Layer: Error handling

**Errors** can be caused by

- Faulty transmitters
- Badly shielded cables
- Signal reflections
- Large divergences in time synchronization between stations

The error rate in a typical *Profibus* network is smaller than  $10^{-4}$  and can be reduced further by error detection and correction

# Data Link Layer: Error detection

## Error Detection

- *Hamming Distance* of 4 by adding a checksum to each packet
- At least 4 bits must change to result in an undetected error
- This results in *integrity class 2* after standard **IEC 870-5-1**

# Data Link Layer: Error correction

## **Send Data with No acknowledge (SDN) service**

- Mainly used for synchronization and status messages
- The erroneous telegram is discarded
- Telegram from the next cycle is used instead

## **Send Data with Acknowledge (SDA) service**

- Mainly used between masters, slaves may not access the bus without request
- When the sender does not receive an acknowledgement, the telegram is retransmitted

# Data Link Layer: Error correction

## **Send and Request Data (SRD)** service

- Service used between masters and slaves
- Acknowledgement is piggybacked on top of the data telegram
- When the sender does not get a response, the telegram is retransmitted

# Application Layer: Addressing

Every station has a unique address, coded in 1 byte

Address	Use
0	reserved for tools, e.g. programming devices
1 – $n$	$n$ master stations
$n - 125$	slave stations
126	reserved as <i>delivery address</i> used for changing the address of a slave during runtime
127	reserved as broadcast address

Components used for the infrastructure, e.g. repeaters, couplers, transmit the data transparently and do not require an address

# Application Layer: Telegram Formats

- 1 Without data field
- 2 Variable length from 4 – 249 byte, payload 1 – 246 byte

SD2	LE	LEr	SD2	DA	SA	FC	PDU	FCS	ED
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SD2: Delimiter, LE: Length, LEr: Length repeated, DA: Destination Address, SA: Source Address, FC: Function Code, PDU: Protocol Data Unit, FCS: Frame Check Sequence, ED: End Delimiter

- 3 Fixed payload length of 8 bytes
- 4 Token telegram
- 5 Short telegram

Telegram Delimiter, featuring a *Hamming Distance* of 4

SD1	SD2	SD3	SD4	ED	SC
0x10	0x68	0xA2	0xDC	0x16	0xE5

# Application Layer: FMS Protocol

## Fieldbus Message Specification (FMS)

- FMS master controls the relationship with FMS slaves
- Replaced by *Profibus DP*

# Application Layer: DP/PA Protocol

**Profibus DP** masters are separated into classes

- Class 1: control a system and the slaves assigned, mostly *Programmable Logic Controller (PLC)* based
- Class 2: tool for commissioning, engineering and maintenance, mostly PC based
- Class 3: clock master, used for time synchronization

**Profibus PA** uses the same protocol as *Profibus DP*



# Application Layer: DP - Cyclic process data

The exchange of data between masters and slaves is separated into three phases: [1]

Phase	Action
Diagnosis	Master requests diagnostic data from slaves
Initialization	Master sets parameters and checks configuration of slaves
Data Exchange	Master cyclically sends and requests data from the slaves

# Application Layer: DP - Data Exchange

## Class 1 master station

- Relationship with a slave is called *MS0*
- Data exchange is cyclic
- Master sends output data to a slave
- Slave immediately responds with input data
- Master continues with the next slave or restarts the cycle

The minimum cycle time  $T_{BCycle}$  can be calculated:

$$T_{BCycle} = \frac{380 + (N_{Slaves} \cdot 300) + (N_{Bytes} \cdot 11)}{Bitrate} + 75\mu s$$

# Application Layer: DP - Data Exchange

## Class 2 master station

- Can exist in addition to *class 1* masters
- Can simultaneously be a *class 1* master
- Relationship with a slave is called *MS1*
- Acyclic communication with a slave in an existing MS0 relationship

# References I



## Max Felser

Profibus Manual: A collection of information explaining PROFIBUS networks

<http://www.profibus.felser.ch>



## Wikipedia

Profibus

<https://en.wikipedia.org/wiki/Profibus>



## Wikipedia

Fieldbus

<https://en.wikipedia.org/wiki/Fieldbus>