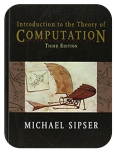
Exercício 3

Capítulo 4, Página 211



Introduction to the Theory of Computation

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Solução Certificado

Passo 1 1 de 2

It is obvious that DFA A accepts all strings, that is $L(A) = \Sigma^*$ if and only if all its states which are reachable from initial state are accepting states. Hence, the decider M for language

$$ALL_{\text{DFA}} = \{\langle A\rangle \mid A \text{ is a DFA and } L(A) = \Sigma^*\}$$

uses the marking procedure for finding reachable nodes, as in *Example 3.23*, starting from initial state. If it turns out that all reachable states are accepting, machine M *accepts* and otherwise *rejects*.

Resultado 2 de 2

This problem is the opposite of decidability of language E_{DFA} from **Theorem 4.4**.

‹ Exercício 2

Avaliar esta solução

Exercício 4 ›