

Atividade de Grupo IV

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16. Verificar que são válidos os seguintes argumentos:

$$\begin{array}{l} \text{(a) (1)} \\ \quad x = y \vee x > y \\ \text{(2)} \\ \quad x < 4 \vee x \leq z \\ \text{(3)} \\ \quad x = y \rightarrow x < z \\ \text{(4)} \\ \quad x > y \rightarrow x < z \\ \hline \therefore x < 4 \end{array}$$

$$\begin{array}{l} \text{(b) (1)} \\ \quad 2x + y = 5 \rightarrow 2x = 2 \\ \text{(2)} \\ \quad 2x + y = 5 \vee y = 3 \\ \text{(3)} \\ \quad 2x = 2 \rightarrow x = 1 \\ \text{(4)} \\ \quad y = 3 \rightarrow 2x = 2 \\ \hline \therefore x = 1 \end{array}$$

$$\begin{array}{l} \text{(c) (1)} \\ \quad x < 3 \vee x > 4 \\ \text{(2)} \\ \quad x < 3 \rightarrow x \neq y \\ \text{(3)} \\ \quad x > 4 \rightarrow x \neq y \\ \text{(4)} \\ \quad x < y \vee x \neq y \rightarrow x \neq 4 \wedge x = 2 \\ \hline \therefore x = 2 \end{array}$$

$$\begin{array}{l} \text{(d) (1)} \\ \quad x = 3 \rightarrow 2x^2 = 18 \\ \text{(2)} \\ \quad x = 3 \vee x = -3 \\ \text{(3)} \\ \quad x = -3 \rightarrow 2x^2 = 18 \\ \text{(4)} \\ \quad 2x^2 = 18 \rightarrow x^2 = 9 \\ \hline \therefore x^2 = 9 \end{array}$$

$$\begin{array}{l} \text{(e) (1)} \\ \quad z > x \rightarrow x \neq 7 \\ \text{(2)} \\ \quad x < 6 \vee x = 3 \\ \text{(3)} \\ \quad x = 3 \rightarrow z > x \\ \text{(4)} \\ \quad x < 6 \rightarrow z > x \\ \text{(5)} \\ \quad x = 7 \vee x = 5 \\ \hline \therefore x = 5 \end{array}$$

$$\begin{array}{l} \text{(f) (1)} \\ \quad x = 3 \vee x = 4 \\ \text{(2)} \\ \quad x = 3 \rightarrow x^2 - 7x + 12 = 0 \\ \text{(3)} \\ \quad x = 4 \rightarrow x^2 - 7x + 12 = 0 \\ \text{(4)} \\ \quad x^2 - 7x + 12 = 0 \rightarrow x > 2 \\ \text{(5)} \\ \quad x^2 < 9 \rightarrow x > 2 \\ \text{(6)} \\ \quad x^2 < 9 \rightarrow x^2 = 9 \vee x^2 > 9 \\ \hline \therefore x^2 = 9 \vee x^2 > 9 \end{array}$$

$$\begin{array}{l} \text{(g) (1)} \\ \quad x > y \vee x < 4 \\ \text{(2)} \\ \quad x < 4 \rightarrow x < y \wedge y < 4 \\ \text{(3)} \\ \quad x > y \rightarrow x = 4 \\ \text{(4)} \\ \quad x \neq 4 \\ \hline \therefore x < y \end{array}$$

$$\begin{array}{l} \text{(h) (1)} \\ \quad x = \frac{5\pi}{6} \rightarrow \sin x = \frac{1}{2} \\ \text{(2)} \\ \quad x = \frac{\pi}{6} \vee x = \frac{5\pi}{6} \\ \text{(3)} \\ \quad \sin x = \frac{1}{2} \rightarrow \csc x = 2 \\ \text{(4)} \\ \quad x = \frac{\pi}{6} \rightarrow \sin x = \frac{1}{2} \\ \hline \therefore \cos x = \frac{\sqrt{3}}{2} \vee \csc x = 2 \end{array}$$

$$\begin{array}{l} \text{(i) (1)} \\ \quad x + 2y = 5 \vee 3x + 4y = 11 \\ \text{(2)} \\ \quad x > y \vee x < 2 \rightarrow y < 2 \vee y < 1 \\ \text{(3)} \\ \quad 3x + 4y = 11 \rightarrow x = 1 \\ \text{(4)} \\ \quad x > y \vee x < 2 \\ \text{(5)} \\ \quad x + 2y = 5 \rightarrow x = 1 \\ \hline \therefore x = 1 \wedge (y < 2 \vee y < 1) \end{array}$$

- A) 1) $x = y \vee x > y$
 2) $x < 4 \vee x \neq z$
 3) $x = y \rightarrow x < z$
 4) $x > y \rightarrow x < z$

$$\begin{array}{ll} x = y : p & \\ x > y : q & \\ x < z : t & \\ x \neq z : \sim n & \\ x < z : n & \end{array}$$

$$1) p \vee q : \sim p \rightarrow q$$

$$2) t \vee \sim n$$

$$3) p \rightarrow n$$

$$4) q \rightarrow n$$

t

$$1) \sim p \rightarrow q$$

$$4) \underline{q \rightarrow n}$$

$\sim p \rightarrow n$: Sílegismo
Hipotético.

$$1) \sim p \rightarrow n : p \vee n$$

$$2) t \vee \sim n : t \vee \sim n$$

$$3) p \rightarrow n : \sim p \vee n : n \vee \sim p$$

$$1) p \vee n$$

$$3) n \vee \sim p$$

$$\frac{(p \vee n) \wedge (n \vee \sim p)}{n \vee (p \wedge \sim p)} \Leftrightarrow \text{Regra da Conjunção}$$

$$n \vee (p \wedge \sim p) \Leftrightarrow$$

$$n \vee C \Leftrightarrow$$

$$n \Leftrightarrow \#$$

$$1) n$$

Comutativa

$$\therefore n \vee \sim n : \sim n \vee t : n \rightarrow t \mid \dots$$

$$\frac{1) n^t \quad 2) t \vee \sim n : \sim n \vee t : n \rightarrow t}{\frac{}{t}} \text{ Modus Ponens}$$

$$\begin{array}{ll}
 \text{B)} & \begin{array}{l} (b) \quad (1) \quad 2x + y = 5 \rightarrow 2x = 2 \\ \quad (2) \quad 2x + y = 5 \vee y = 3 \\ \quad (3) \quad 2x = 2 \rightarrow x = 1 \\ \quad (4) \quad y = 3 - 2x = 2 \\ \hline \therefore \quad x = 1 \end{array} \quad \begin{array}{l} 2x + y = 5 : P \\ 2x = 2 : q \\ y = 3 : r \\ x = 1 : t \end{array}
 \end{array}$$

$$\begin{array}{l}
 (1) P \rightarrow q \\
 (2) P \vee \pi \\
 (3) q \rightarrow t \\
 (4) \pi \rightarrow q \\
 \hline
 t
 \end{array}$$

$$\begin{array}{c} 1) p \rightarrow q \\ 2) q \rightarrow t \\ \hline p \rightarrow t \quad \therefore \text{Syllogism} \end{array}$$

Hypothetical Syllogism

- (1) $P \rightarrow t$
- (2) $P \vee m$
- (3) $n \rightarrow g$

$t \vee q$: Dilema
constructivo

$tvt \Rightarrow t : \text{Adig}^{\bar{AO}}!$

(c) (1) $x < 3 \vee x > 4$
 (2) $x < 3 \rightarrow x \neq y$
 (3) $x > 4 \rightarrow x \neq y$
 (4) $x < y \vee x \neq y \rightarrow x \neq 4 \wedge x = 2$

$$\therefore x = 2$$

- $x < 3$: P
- $x > 4$: Q
- $x \neq y$: R
- $x < y$: S
- $x \neq 4$: W
- $x = 2$: T

\rightarrow (Lemma)

- (1) $P \vee q : q \vee P$
- (2) $P \rightarrow r$
- (3) $q \rightarrow r$
- (4) $(P \vee q) \rightarrow w \wedge t$

 t

$$\frac{\neg q \rightarrow p \\ p \rightarrow r}{\neg q \rightarrow r} \therefore \text{Silogismo Hipotético}$$

$$\Rightarrow \begin{array}{c} q \vee p \Leftrightarrow q \vee \underline{p} \\ nq \vee p \Leftrightarrow \underline{p} \vee nq \end{array}$$

$$(q \vee r) \wedge (\neg q \vee \neg r) \Leftrightarrow \therefore \text{Regeln der Konjunktion!}$$

$$n \vee (q \wedge \neg q) \Leftrightarrow$$

$n \in C \Leftrightarrow$

$\pi \Leftrightarrow \#$

$\pi \Leftrightarrow$

(1) π

(2) $(\neg v \pi) \rightarrow (\omega \wedge t)$

\Rightarrow $\frac{(1) \pi \rightarrow \text{Adicção} \quad (2) \pi \rightarrow t}{(2) \pi \rightarrow t}$ Simplificação!

$\pi \vdash \neg v \pi$: Adicção

$\Rightarrow \omega \wedge t \vdash t$: Simplificação

t

\therefore Modus Ponens

④

$$\begin{array}{l} (d) (1) x = 3 \rightarrow 2x^2 = 18 \\ (2) x = 3 \vee x = -3 \\ (3) x = -3 \rightarrow 2x^2 = 18 \\ (4) 2x^2 = 18 \rightarrow x^2 = 9 \\ \hline \therefore x^2 = 9 \end{array}$$

$x = 3 : P$

$2x^2 = 18 : q$

$x = -3 : \pi$

$x^2 = 9 : \Delta$

(1) $P \rightarrow q$

(2) $P \vee \pi$

(3) $\pi \rightarrow q$

(4) $q \rightarrow \Delta$

$(1) P \rightarrow q$

$(4) q \rightarrow \Delta$

(1) $P \rightarrow \Delta$

(2) $P \vee \pi$

(3) $\pi \rightarrow q$

$P \rightarrow \Delta : \because \text{Sílogismo Hipotético}$

$\Rightarrow (2) \neg P \rightarrow \pi$

$\Rightarrow (3) \pi \rightarrow q$

$\neg P \rightarrow q : \because \text{Sílogismo Hipotético}$

(1) $P \rightarrow \Delta \Leftrightarrow P \rightarrow \Delta$

(2) $\neg P \rightarrow q \Leftrightarrow P$

$P \vdash P \vee q : \text{Adicção}$

\Rightarrow

(3) $\pi \rightarrow q$

\Rightarrow

(3) $\pi \rightarrow q$

$\therefore \text{Modus Ponens}$

⑤

$$\begin{array}{l} (e) (1) z > x \rightarrow x \neq 7 \\ (2) x < 6 \vee x = 3 \\ (3) x = 3 \rightarrow z > x \\ (4) x < 6 \rightarrow z > x \\ (5) x = 7 \vee x = 5 \\ \hline \therefore x = 5 \end{array}$$

$z > x : P \quad x = 7 : \pi$

$x \neq 7 : q \quad x = 5 : W$

$x < 6 : \Delta$

$x = 3 : \Delta$

(1) $P \rightarrow q$

(2) $\pi \vee \Delta$

(3) $\Delta \rightarrow P$

(4) $\pi \rightarrow P$

(5) $t \vee w$

(3) $\Delta \rightarrow P$

(4) $\pi \rightarrow P$

(2) $\pi \vee \Delta$

(1) $P \rightarrow q$

(2) P

(3) $t \vee w$

$P \vee P \Leftrightarrow P : \because \text{Dilema Contrutivo}$

w

$\Rightarrow \frac{(1) P \rightarrow q}{q} \quad \frac{(2) P}{q} : \therefore \text{Modus Ponens}$

$\Rightarrow \frac{(1) q}{(2) t \vee w : \neg t \rightarrow w} : \frac{q}{w, \neg t \rightarrow w}$

$q = \neg t : \therefore \text{Modus Ponens}$

(F)

$$\begin{aligned}
 (1) \quad & x = 3 \vee x = 4 \\
 (2) \quad & x = 3 \rightarrow x^2 - 7x + 12 = 0 \\
 (3) \quad & x = 4 \rightarrow x^2 - 7x + 12 = 0 \\
 (4) \quad & x^2 - 7x + 12 = 0 \rightarrow x > 2 \\
 (5) \quad & x^2 < 9 \rightarrow x > 2 \\
 (6) \quad & x^2 < 9 \rightarrow x^2 = 9 \vee x^2 > 9 \\
 \therefore \quad & x^2 = 9 \vee x^2 > 9
 \end{aligned}$$

$$\begin{aligned}
 x = 3 : p \\
 x = 4 : q \\
 x^2 - 7x + 12 : r \\
 x > 2 : s
 \end{aligned}$$

$$\begin{aligned}
 x^2 < 9 : t \\
 x \neq 2 : \sim s \\
 x^2 < 9 : \sim t
 \end{aligned}$$

$$x^2 > 9 : u$$

(1) π (2) $\pi \rightarrow s$ (3) $t \rightarrow \sim s$ (4) $\sim t \rightarrow u \vee v$

$$(1) p \vee q$$

 \Rightarrow

$$(1) p \vee q$$

$$(2) p \rightarrow r$$

$$(3) q \rightarrow \pi$$

$$r \vee \pi : n \therefore \text{Dilema Constructivo}$$

$$(2) p \rightarrow \pi$$

$$(3) q \rightarrow \pi$$

$$(4) r \rightarrow s$$

$$(5) t \rightarrow \sim s$$

$$(6) \sim t \rightarrow u \vee v$$

 \Rightarrow

$$(1) \pi$$

$$(2) r \rightarrow s$$

 $\therefore \text{modus ponens}$

$$(1) s$$

$$(2) t \rightarrow \sim s$$

$$(3) \sim t \rightarrow u \vee v$$

$$\begin{aligned}
 (1) \quad & s \\
 \Rightarrow \quad & \underline{(2) s \rightarrow \sim t} \\
 & \sim t : \text{modus ponens}
 \end{aligned}$$

$$\begin{aligned}
 (1) \quad & \sim t \\
 \Rightarrow \quad & \underline{(2) \sim t \rightarrow u \vee v} \\
 & u \vee v : \text{modus ponens}
 \end{aligned}$$

 \swarrow

(G)

$$\begin{aligned}
 (1) \quad & x > y \vee x < 4 \\
 (2) \quad & x < 4 \rightarrow x < y \wedge y < 4 \\
 (3) \quad & x > y \rightarrow x = 4 \\
 (4) \quad & x \neq 4 \\
 \therefore \quad & x < y
 \end{aligned}$$

$$\begin{aligned}
 x > y : p \\
 x < 4 : q \\
 x < y : r \\
 y \neq 4 : s
 \end{aligned}$$

$$\begin{aligned}
 x = 4 : t \\
 x \neq 4 : w
 \end{aligned}$$

$$(1) \sim p \rightarrow r \wedge s$$

$$(2) p \rightarrow t$$

$$(3) \sim t$$

$$(1) p \vee q$$

$$(2) q \rightarrow r \wedge s$$

$$(3) p \rightarrow t$$

$$(4) w = \sim t$$

$$(1) \sim p \rightarrow q$$

$$(2) q \rightarrow r \wedge s$$

$$\sim p \rightarrow r \wedge s \therefore \text{Silegismo Hipotético}$$

$$\begin{aligned}
 \Rightarrow \quad & \underline{(2) p \rightarrow t} \\
 (3) \quad & \sim t \\
 & \sim p : \text{modus tollens}
 \end{aligned}$$

$$(1) \sim p \rightarrow r \wedge s$$

$$\sim p : \text{modus tollens}$$

Simplificado

$$r \wedge s \vdash r$$

 \swarrow

(H)

$$\begin{aligned}
 (1) \quad & p : p \\
 (2) \quad & \frac{n}{6} : q \\
 (3) \quad & \frac{\sin x}{2} = \frac{1}{2} \rightarrow \csc x = 2 \\
 (4) \quad & \frac{x}{6} = \frac{\pi}{6} \rightarrow \sin x = \frac{1}{2} \\
 \therefore \quad & \cos x = \frac{\sqrt{3}}{2} \vee \csc x = 2
 \end{aligned}$$

$$(1) p \rightarrow q$$

$$(2) r \vee p$$

$$(3) q \rightarrow s$$

$$(4) r \rightarrow q$$

$$(1) p \rightarrow q$$

$$(3) q \rightarrow s$$

$$(p \rightarrow s \therefore \text{Silegismo Hipotético})$$

$$\begin{aligned}
 \frac{}{} \quad & p \rightarrow s \therefore \text{Silegismo} \\
 & \text{Hipotético}
 \end{aligned}$$

$$(1) P \rightarrow \Delta$$

$$(2) \neg \vee P$$

$$(3) \neg \rightarrow \neg$$

$$\underline{\Delta \vee \neg} \therefore \text{Dilema Constitutivo}$$

$$\begin{array}{l} (1) x+2y=5 \vee 3x+4y=11 \\ (2) x>y \vee x < 2 \rightarrow y < 2 \vee y < 1 \\ (3) 3x+4y=11 \rightarrow x=1 \\ (4) x>y \vee x < 2 \\ (5) x+2y=5 \rightarrow x=1 \\ \hline \therefore x=1 \wedge (y < 2 \vee y < 1) \end{array}$$

$$\begin{array}{ll} x+2y=5 : P & y < 2 : t \\ 3x+4y=11 : q & y < 1 : u \\ x>y : n & x=1 : w \\ x \neq 2 : \Delta & \end{array}$$

$$\begin{array}{l} (1) P \vee q \\ (2) \neg \vee \Delta \rightarrow t \vee u \\ (3) q \rightarrow w \\ (4) \neg \vee \Delta \\ (5) P \rightarrow w \\ \hline w \wedge (t \vee u) \end{array}$$

$$\begin{array}{l} (1) P \vee q \\ (2) q \rightarrow w \\ (3) P \rightarrow w \\ \hline w \vee w \therefore \text{Dilema Constitutivo} \\ w \end{array}$$

$$\begin{array}{l} (4) w \\ (2) \neg \vee \Delta \rightarrow t \vee u \\ (3) \neg \vee \Delta \end{array}$$

$$\begin{array}{l} \Rightarrow (2) \neg \vee \Delta \rightarrow t \vee u \\ \underline{(3) \neg \vee \Delta} \\ t \vee u \therefore \text{Modus Ponens} \end{array}$$

$$\begin{array}{l} (1) t \vee u \\ \Rightarrow \underline{(2) w} \\ w \wedge (t \vee u) \end{array}$$

17. Usar as Regras de Inferência para mostrar que são válidos os seguintes argumentos:

- (a) $p \vee q \rightarrow \neg r, p, s \rightarrow r \vdash \neg s$
- (b) $p \wedge (q \vee r), q \vee r \rightarrow \neg s, s \vee t \vdash t$
- (c) $p \vee q \rightarrow \neg r, q, s \wedge t \rightarrow r \vdash \neg(s \wedge t)$
- (d) $p \rightarrow q, \neg q, \neg p \vee \neg r \rightarrow s \vdash s$
- (e) $p \vee (q \wedge r), q \rightarrow s, r \rightarrow t, s \wedge t \rightarrow p \vee r, \neg p \vdash r$
- (f) $q \vee (r \rightarrow t), q \rightarrow s, \neg s \rightarrow (t \rightarrow p), \neg s \vdash r \rightarrow p$
- (g) $p \vee q \rightarrow (p \rightarrow s \wedge t), p \wedge r \vdash t \vee u$

$$\textcircled{A} \quad \begin{array}{l} (1) P \vee q \rightarrow \neg r \\ (2) P \\ (3) \Delta \rightarrow r \\ \hline \neg b \end{array}$$

$$P \vdash P \vee q \therefore \text{Adic}\overline{\text{e}} \quad \Rightarrow$$

$$(1) P \vee q \rightarrow \neg r$$

17. Dado

$$\begin{array}{l} (1) P \vee q \rightarrow \neg r \\ \hline P \vee q \end{array} \quad \Rightarrow$$

$$\begin{array}{c}
 (1) P \vee q \rightarrow N\pi \\
 (2) P \vee q \\
 (3) \Delta \rightarrow \pi
 \end{array} \Rightarrow \frac{\begin{array}{c} (1) P \vee q \\ (2) P \vee q \\ (3) \Delta \rightarrow \pi \end{array}}{N\pi \therefore \text{Modus Ponens}} \Rightarrow$$

$$\begin{array}{c}
 (1) \sim \pi \\
 (2) \Delta \rightarrow \pi \\
 \hline
 \sim \Delta \therefore \text{Modus Tolleris}
 \end{array}$$

$$\textcircled{B} \quad \begin{array}{c}
 (1) P \wedge (\neg q \vee \pi) \\
 (2) \neg q \vee \pi \rightarrow \sim \Delta \\
 (3) \Delta \vee t \\
 \hline
 t
 \end{array} \Rightarrow P \wedge (\neg q \vee \pi) \vdash \neg q \vee \pi \therefore \text{Simplificatio}$$

$$\Rightarrow \begin{array}{c}
 (1) \neg q \vee \pi \\
 (2) \neg q \vee \pi \rightarrow N\Delta \\
 \hline
 N\Delta \therefore \text{Modus Ponens}
 \end{array} \Rightarrow \frac{\begin{array}{c} (1) \sim \Delta \\ (2) \sim \Delta \rightarrow t \end{array}}{t \therefore \text{Modus Ponens}}$$

$$\textcircled{C} \quad \begin{array}{c}
 (1) P \vee q \rightarrow \sim \pi \\
 (2) \neg q \\
 (3) \Delta \wedge t \rightarrow \pi \\
 \hline
 \sim(\Delta \wedge t)
 \end{array} \Rightarrow \neg q \vdash P \vee q \therefore \text{Adicatio}$$

$$\Rightarrow \begin{array}{c}
 (1) P \vee q \rightarrow \sim \pi \\
 (2) P \vee q \\
 \hline
 N\pi \therefore \text{Modus Ponens}
 \end{array} \Rightarrow \frac{\begin{array}{c} (1) \sim \pi \\ (2) \Delta \wedge t \rightarrow \pi \end{array}}{(1) \sim \pi \Rightarrow (2) \Delta \wedge t \rightarrow \pi} \Rightarrow \frac{\begin{array}{c} (1) \sim \pi \\ (2) \sim \pi \rightarrow N(\Delta \wedge t) \end{array}}{N(\Delta \wedge t) \therefore \text{Modus Ponens}}$$

$$\textcircled{D} \quad \begin{array}{c}
 (1) P \rightarrow q \\
 (2) \sim q \\
 (3) \sim P \vee N\pi \rightarrow \Delta \\
 \hline
 \Delta
 \end{array} \Rightarrow \frac{\begin{array}{c} (1) \sim q \rightarrow \sim P \\ (2) \sim q \end{array}}{\sim P \therefore \text{Modus Ponens}} \Rightarrow$$

$$\Rightarrow \begin{array}{l} (1) \sim P \\ (2) \sim P \vee \sim M \rightarrow \Delta \end{array} \Rightarrow (1) \sim P \vdash \sim P \vee \sim M \therefore \text{Adic} \bar{\exists}$$

$$\Rightarrow \frac{\begin{array}{l} (1) \sim P \vee \sim M \\ (2) \sim P \vee \sim M \end{array}}{\Delta : \text{Modus ponens}}$$

(E) (1) $P \vee (q \wedge M)$

(6) $(q \wedge M)$ 1,5: SD

(2) $q \rightarrow \Delta$

\Rightarrow (7) M : SIMP

(3) $n \rightarrow t$

(4) $\Delta \wedge t \rightarrow P \vee M$

(5) $\sim P$

n

(F) (1) $q \vee (n \rightarrow t)$

(5) $t \rightarrow P$: 3,4: MP

(2) $q \rightarrow \Delta$

(6) $\sim \Delta \rightarrow \sim q$: 2: CP

(3) $\sim \Delta \rightarrow (t \rightarrow P)$

(7) $\sim q$: 4,6: MP

(4) $\sim \Delta$

(8) $n \rightarrow t$: 1,7: SD

$n \rightarrow P$

(9) $n \rightarrow P$: 5,8: SH

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(G) (1) $P \vee q \rightarrow (P \rightarrow \Delta \wedge t)$

(2) $P \wedge M$

$t \vee u$

(3) P : 2: SIMP

(4) $P \vee q$: 3: AD

(5) $P \rightarrow \Delta \wedge t$: 1,4: MP

(6) $\wedge \neg t : 3, 5 : MP$

(7) $t : 6 : SIMP$

(8) $\neg t \vee u : 7 : AD$