

## Exercício 29

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Introduction to the Theory of Computation

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### Passo 1

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#### Part a.

Assume that the given language  $A_1 = \{0^n 1^n 2^n \mid n \geq 0\}$  is regular. Then by pumping lemma there is pumping length  $p \in \mathbb{N}$  such that any string  $w$  whose length is at least  $p$  can be divided up into three parts as  $w = xyz$ , for which the following conditions hold:

1.  $xy^i z \in A_1$ , for any  $i \geq 0$ ,
2.  $y$  is not  $\varepsilon$ , and
3. length of initial part  $xy$  is not larger than pumping length  $p$ .

So let's take string  $w = 0^p 1^p 2^p$  from language  $A_1$ . Its length is  $3p$ , which is certainly greater than  $p$ , so it can be pumped. This means that it can be divided into three parts  $x$ ,  $y$  and  $z$ , where  $y \neq \varepsilon$ , for which  $w = xyz$  and any variant  $xy^i z$  is also a member of  $A_1$ .

We also know that  $|xy| \leq p$ , which in this case implies that  $xy = 0^k$ , for some  $k \leq p$  (because the length of first block of 0's is exactly  $p$ ). Then  $y$  consists of only 0's, which means that pumping it up produces more 0's, while number of 1's and 2's remains unchanged. This new string then can not be member of language  $A_1$ , which is in contradiction with pumping lemma.

So we must conclude that our assumption about regularity of  $A_1$  was completely wrong, i.e. language  $A_1$  is **not regular**.

**Passo 2**

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**Part b.**

Consider string  $\mathbf{a^pba^pba^pb} \in A_2$  ( $w = \mathbf{a^pb}$ , where  $p$  is the pumping length) and follow the reasoning from previous part.

**Passo 3**

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**Part c.**

Strings in language  $A_3 = \{\mathbf{a^{2^n}} \mid n \geq 0\}$  consist of only **a**'s and their length is equal to powers of 2. First we compute the difference between two consecutive powers of 2:

$$2^{n+1} - 2^n = 2 \cdot 2^n - 2^n = 2^n > n.$$

Last inequality is obvious (and can be easily proved using induction).

Let's now assume that  $A_3$  is regular. Then there is pumping length  $p$ .

$$2^p < 2^p + k = |xy^2z| \leq 2^p + p < 2^{p+1},$$

i.e. it is strictly between two consecutive powers of 2.

**Resultado**

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Try to discover which strings to pump.

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