

Parte 2: Quicksort

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O quicksort

- Quicksort (ordenação rápida)
 - No pior caso é $\Theta(n^2)$
 - Na prática a eficiência na média é muito boa
 - O tempo de execução esperado $\Theta(n \log n)$
 - Os fatores ocultos na notação $\Theta(n \log n)$ são pequenos
 - Funciona bem em ambientes de memória virtual

Descrição

- Dividir
 - O arranjo $A[p..r]$ é particionado em dois subarranjos (possivelmente vazios) $A[p..q-1]$ e $A[q+1..r]$ tais que cada elemento de $A[p..q-1]$ é menor que $A[q]$ que, por sua vez, é igual ou menor a cada elemento de $A[q+1..r]$. O índice q é calculado como parte do procedimento
- Conquistar
 - Os dois subarranjos $A[p..q-1]$ e $A[q+1..r]$ são ordenados por chamadas recursivas
- Combinar
 - Os arranjos são ordenados localmente, não há a necessidade de combinação

QUICKSORT

```
QUICKSORT (A, p, r)
1: if p < r then
2:   q ← PARTITION (A, p, r);
3:   QUICKSORT (A, p, q-1);
4:   QUICKSORT (A, q+1, r);
6: end if
```

- Para ordenar um arranjo A , a chamada inicial seria
 - $\text{QUICKSORT}(A, 1, n)$
- Onde está o pulo do gato?

PARTITION

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PARTITION (A, p, r)
1:  $x \leftarrow A[r]$ ;
2:  $i \leftarrow p-1$ ;
3: for  $j \leftarrow p$  to  $r-1$  do
4:   if  $A[j] \leq x$  then
6:      $i \leftarrow i+1$ ;
7:     trocar( $A[i], A[j]$ );
8:   end if
9: end for
10: trocar( $A[i+1], A[r]$ );
11: return  $i+1$ ;
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- O PARTITION sempre seleciona um elemento $x=A[r]$ como um pivô
- O particionamento do arranjo $A[p..r]$ é feito em torno do pivô x

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A	2	8	7	1	3	5	6	4
	p							r

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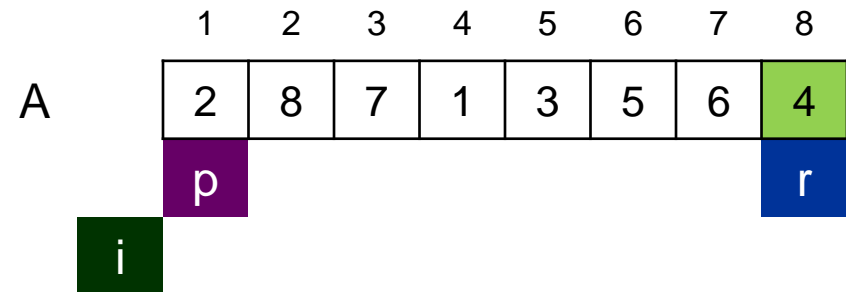
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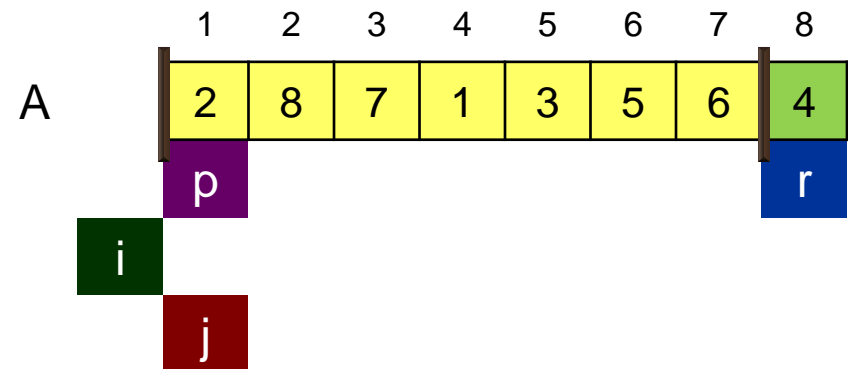
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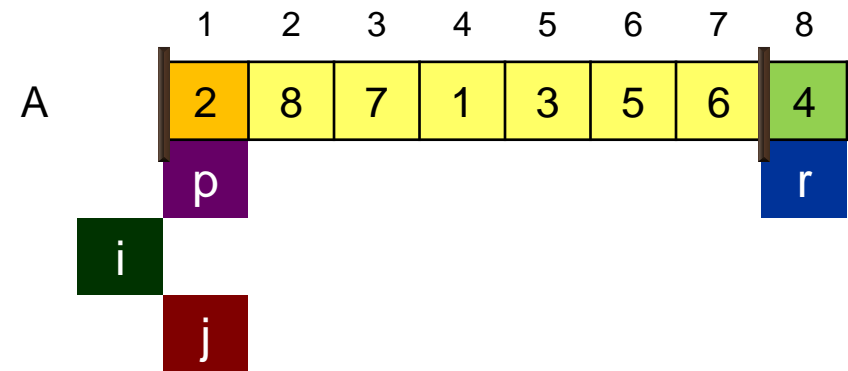


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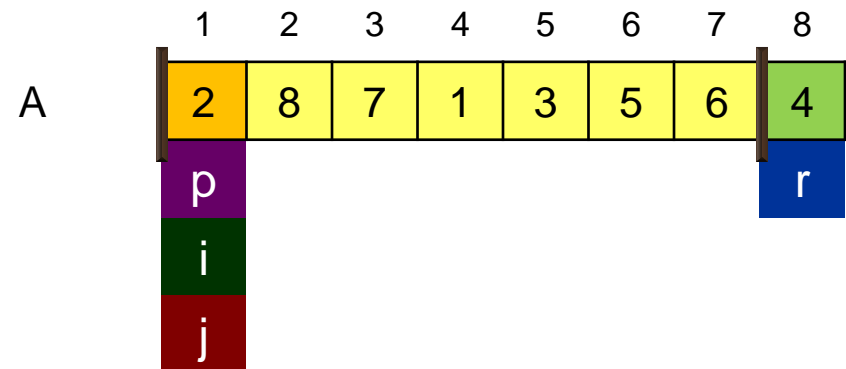


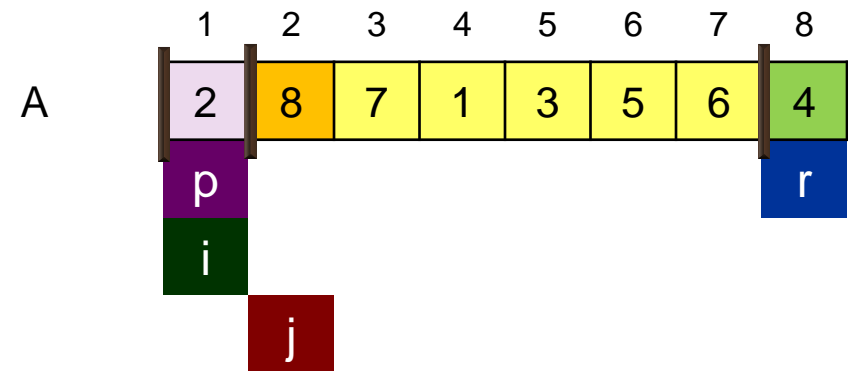
Diagram illustrating an array A with 8 slots. The slots are indexed 1 through 8. The values in the slots are: 2, 8, 7, 1, 3, 5, 6, 4. The first slot (index 1) is highlighted in light purple, and the last slot (index 8) is highlighted in light green. Below the first slot, there are two blocks labeled 'p' (purple) and 'i' (dark green). Below the last slot, there is a block labeled 'r' (dark blue). Below the third slot, there is a block labeled 'j' (dark red).

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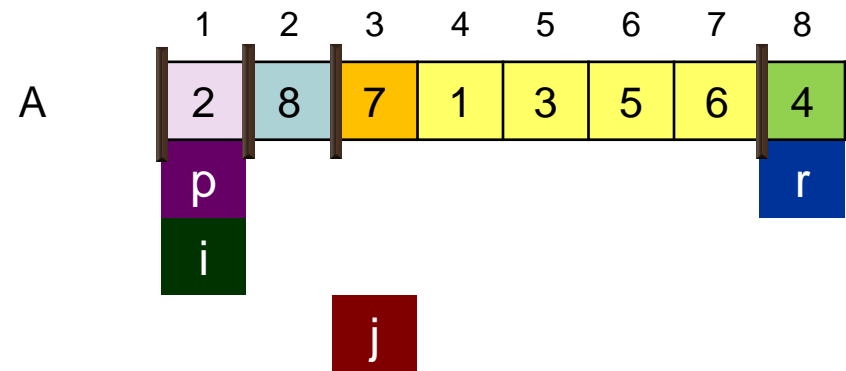


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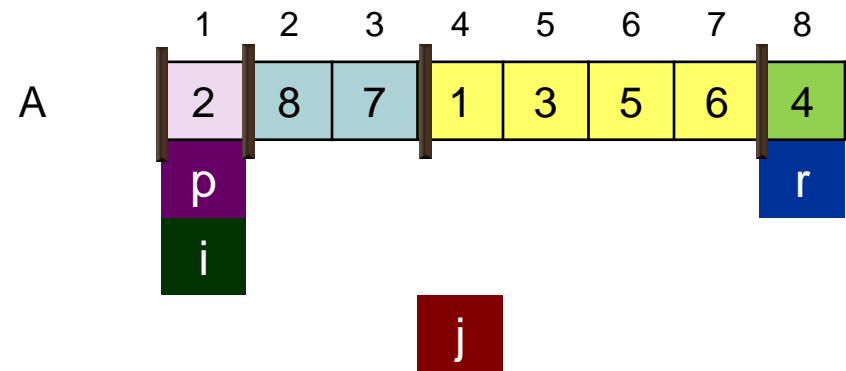
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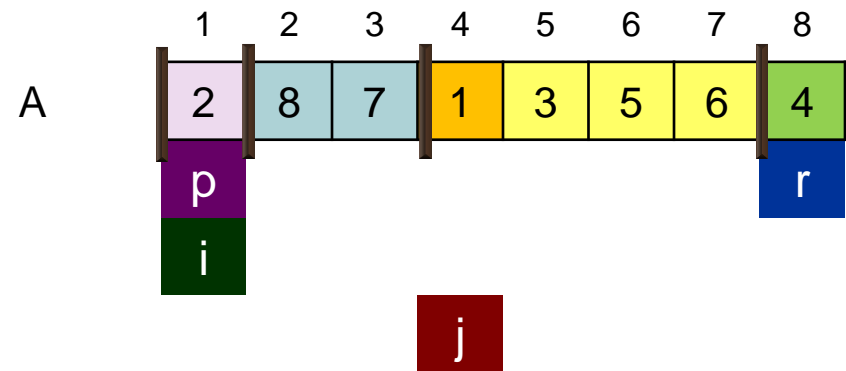


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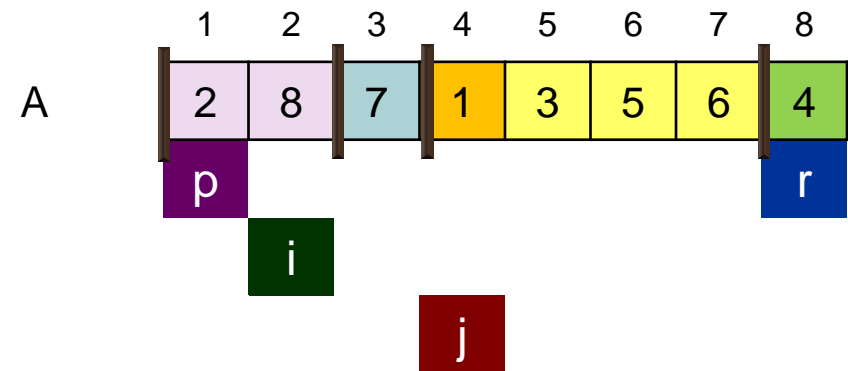


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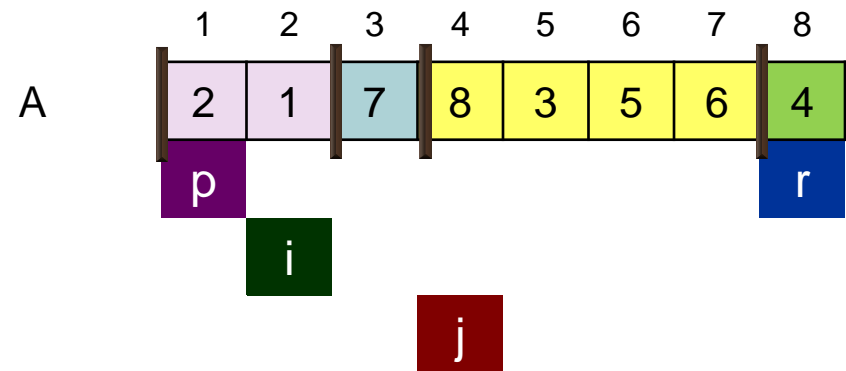


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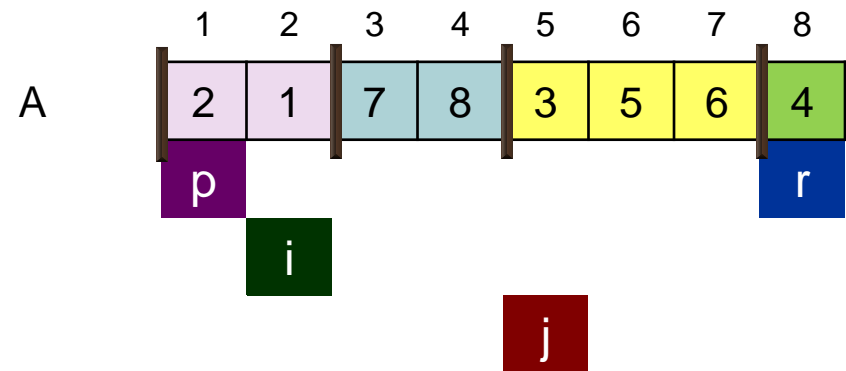
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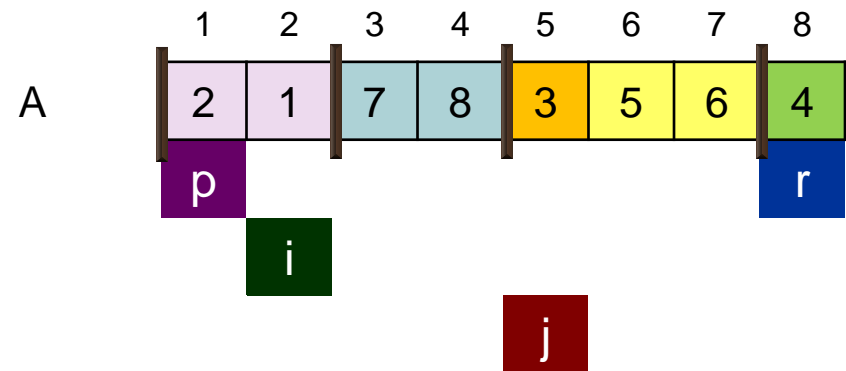


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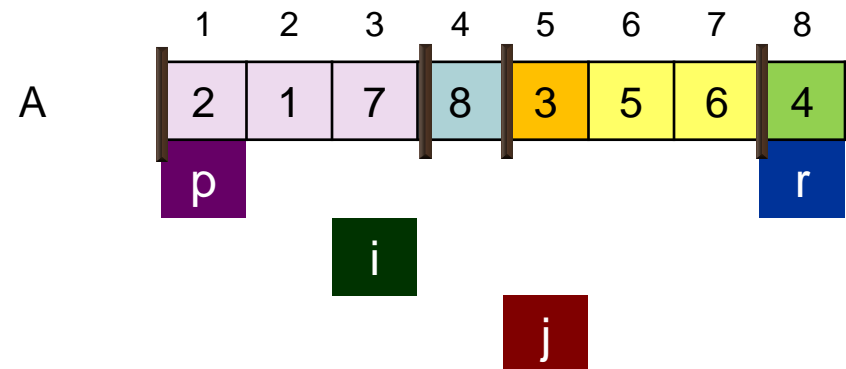


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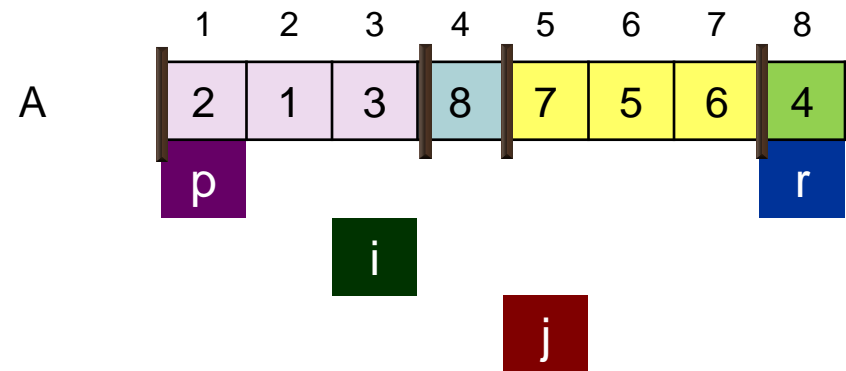


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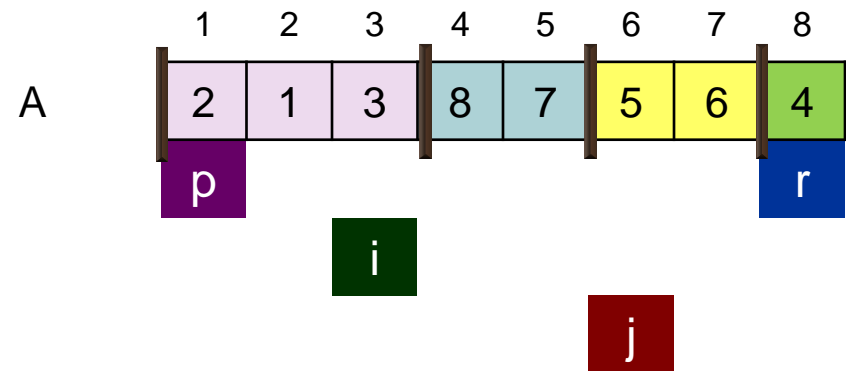
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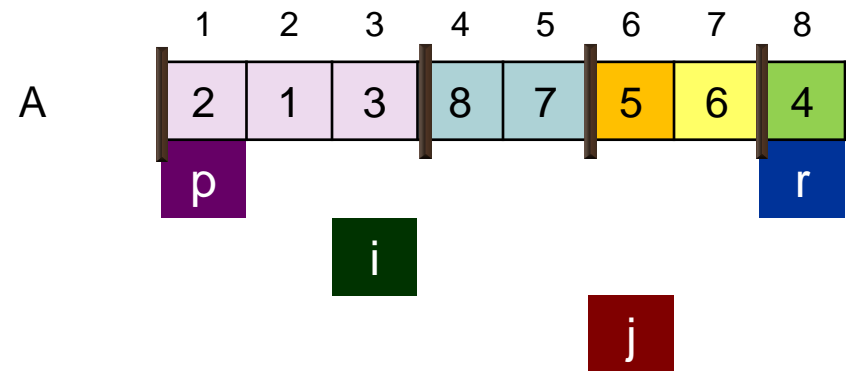


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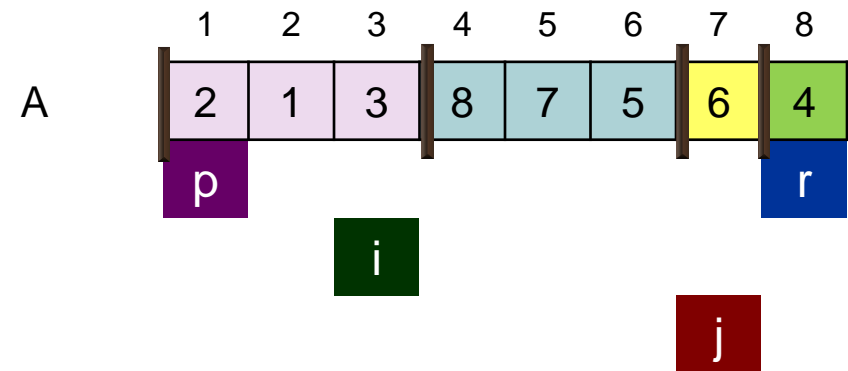
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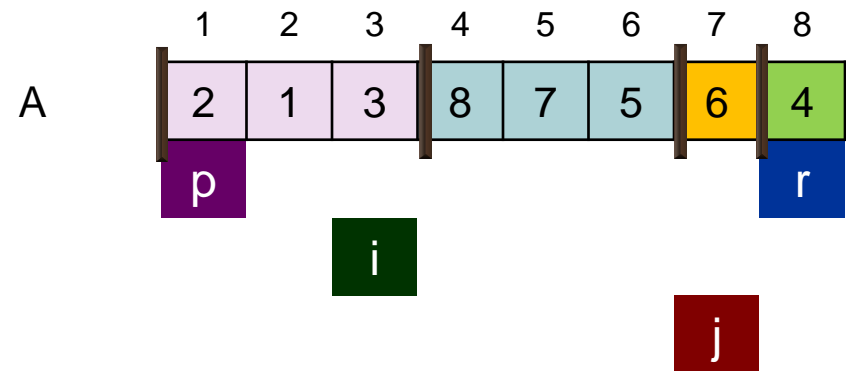


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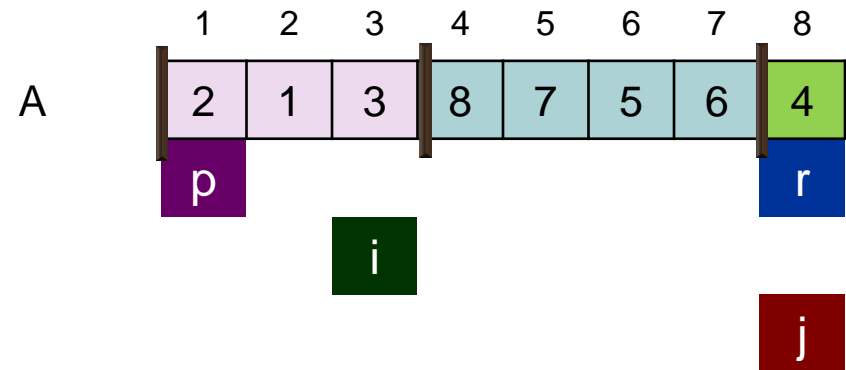
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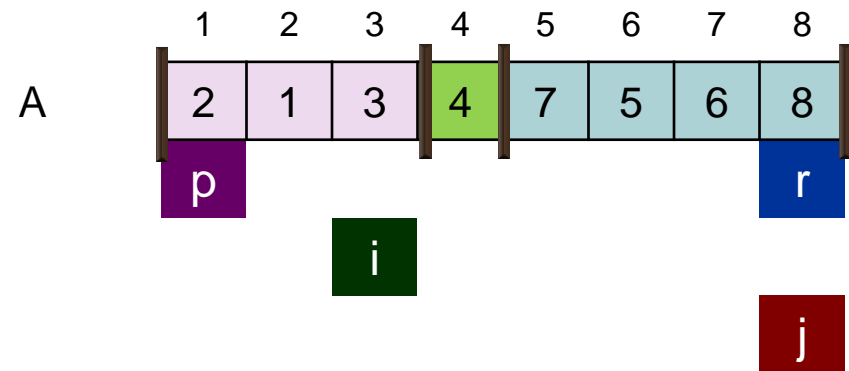


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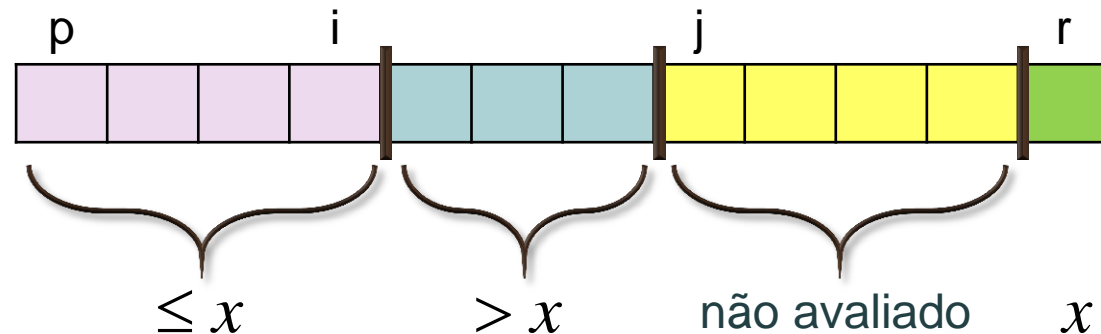


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- A cada iteração do loop (linhas 3 – 8), para qualquer k
 - Se $p \leq k \leq i$, então $A[k] \leq x$
 - Se $i < k < j$, então $A[k] > x$
 - Se $k = r$, então $A[k] = x$
 - Se $j \leq k < r$, então $A[k]$ não foi avaliado ainda



Análise do PARTITION

PARTITION(A, p, r)

$\Theta(1)$ { 1: $x \leftarrow A[r]$;
2: $i \leftarrow p-1$;

$r-p = n-1$ vezes \longrightarrow 3: **for** $j \leftarrow p$ **to** $r-1$ **do**

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$$T(n) = \Theta(1) + (n-1)\Theta(1) + \Theta(1) = \Theta(n) \quad \Rightarrow \quad T(n) = \Theta(n)$$

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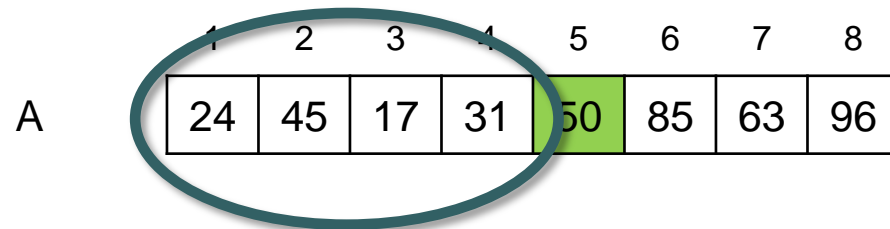
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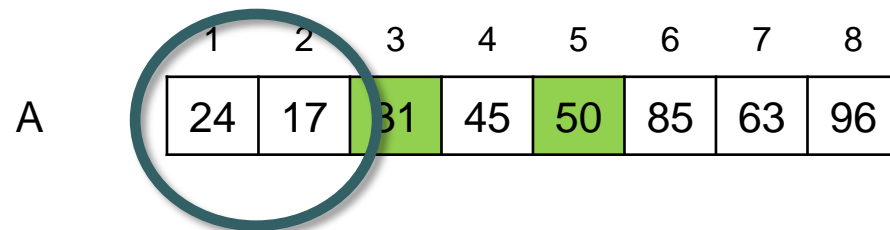
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4: QUICKSORT (A, q+1, r) ;

6: **end if**

	1	2	3	4	5	6	7	8
A	17	24	31	45	50	85	63	96

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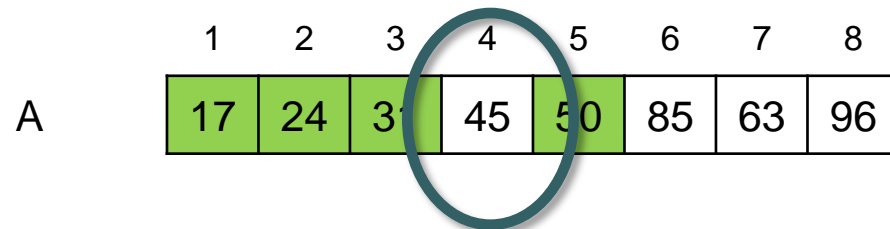
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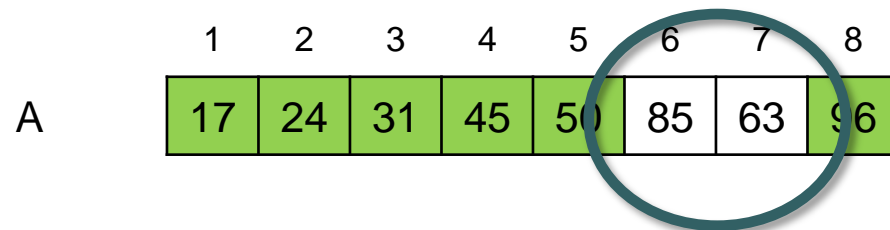
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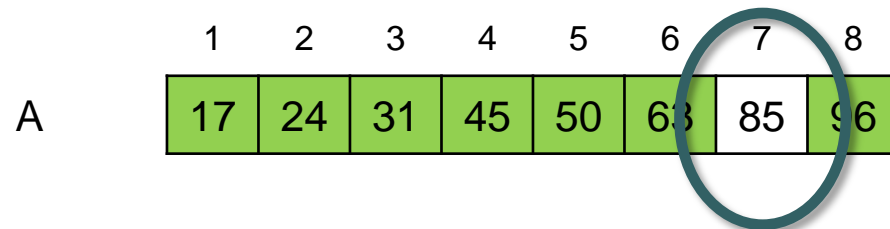
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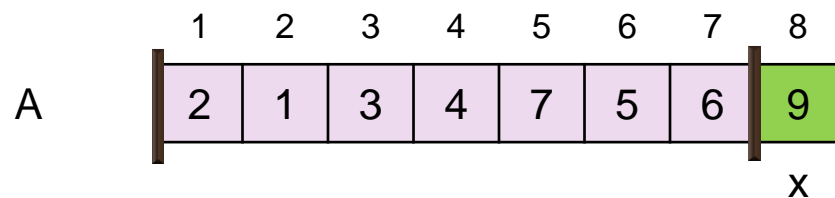
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Análise do QUICKSORT

- O desempenho do Quicksort depende da qualidade do particionamento
 - Pior caso: totalmente desbalanceado
 - Melhor caso: totalmente balanceado
- Particionamento balanceado
 - Custo de mesma ordem que o Mergesort
- Particionamento desbalanceado
 - Custo de mesma ordem que o Insertionsort

Pior caso

- O pior balanceamento é quando o particionamento produz
 - Subproblema de $n-1$ elementos
 - Outro com 0 elementos
 - Ocorre quando o pivô é o maior (ou o menor) elemento do arranjo



Pior caso

- Quando isso ocorre sempre (recursivamente), temos

```
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1: if p < r then
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```

$$\begin{aligned}T(n) &= T(n-1) + T(0) + \Theta(n) \\&= T(n-1) + \Theta(1) + \Theta(n) \\&= T(n-1) + \Theta(n)\end{aligned}$$

$$T(n) = \Theta(n^2)$$

Melhor caso

- Na divisão mais balanceada possível, o PARTITION produz dois subproblemas de tamanho igual

```
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4:   QUICKSORT (A, q+1, r) ;
6: end if
```

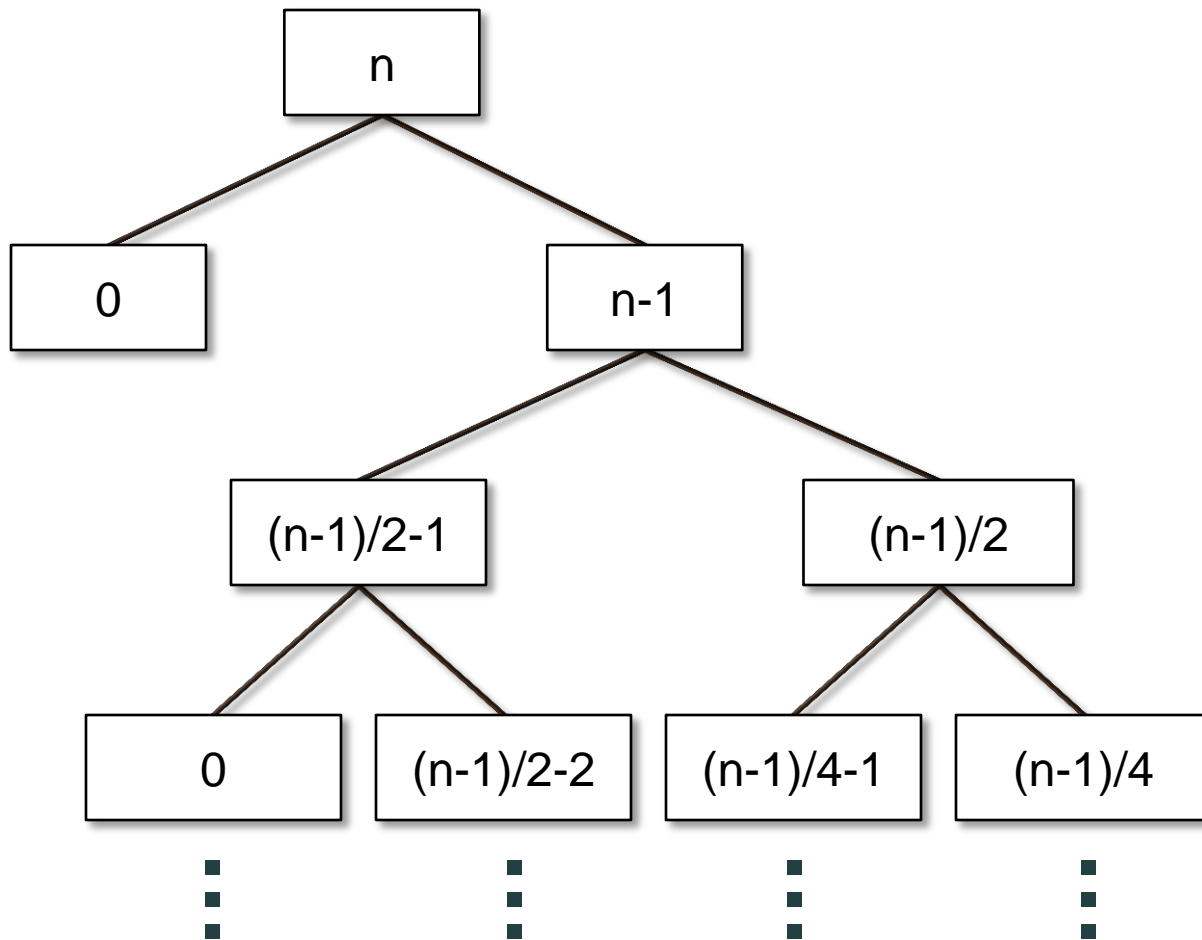
$$T(n) \leq 2T(n/2) + \Theta(n)$$

$$T(n) = O(n \log n)$$

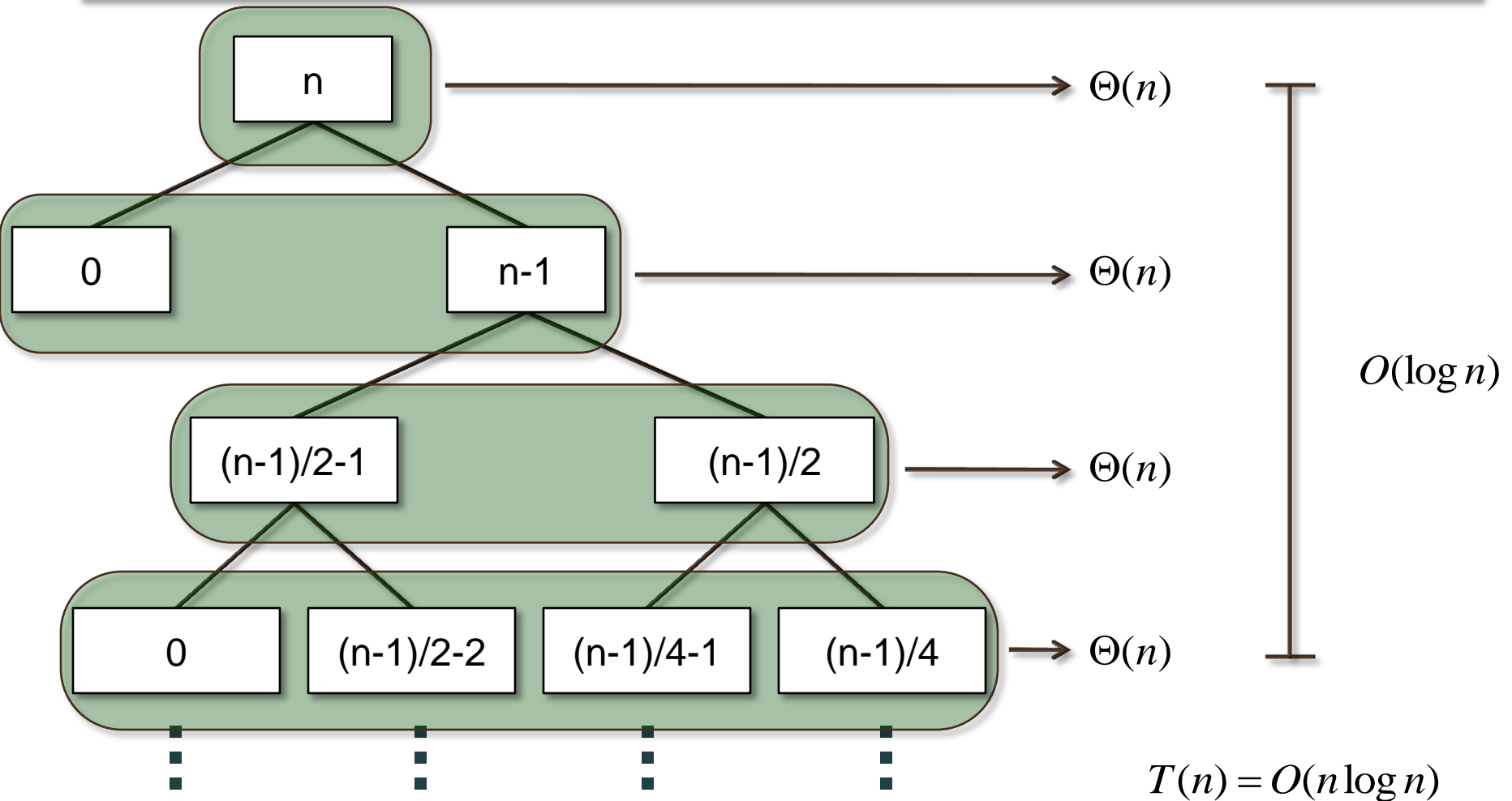
Intuição para o caso médio

- No caso médio, o PARTITION produz uma mistura de divisões BOAS e RUINS
 - As divisões boas e ruins são divididas aleatoriamente
- Para efeitos de intuição
 - Suponha que as divisões boas sejam as do melhor caso e as divisões ruins sejam as do pior caso

Intuição para o caso médio



Intuição para o caso médio



Ordenação por comparação

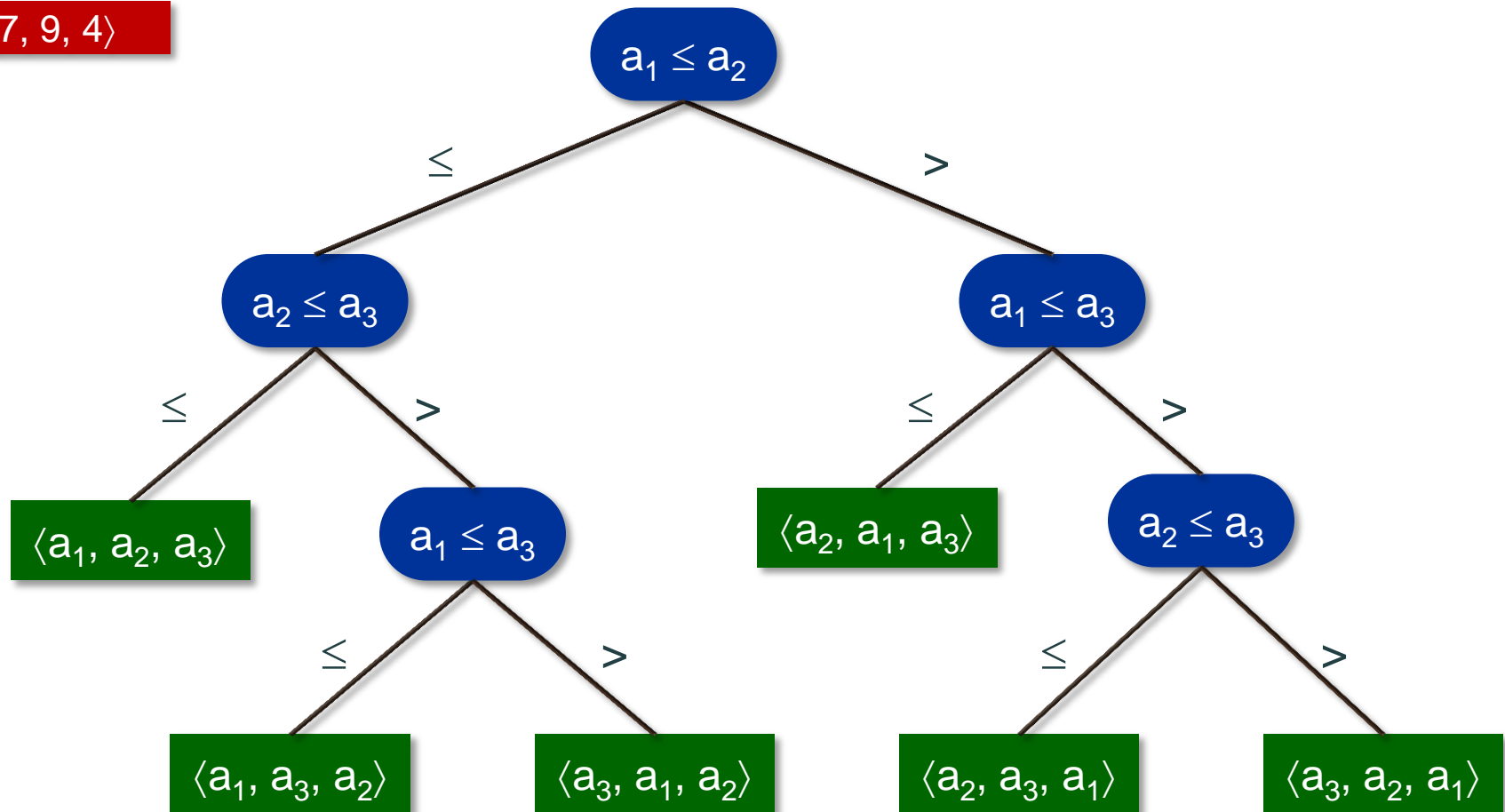
- Os algoritmos vistos compartilham uma propriedade
 - A seqüência ordenada é obtida apenas usando comparações entre os elementos de entrada
 - Esta classe de algoritmos chama-se **ordenação por comparação**
- Pertencem a esta classe
 - Bubblesort
 - Insertionsort
 - Mergesort
 - Heapsort
 - Quicksort
 - ...

Limite inferior

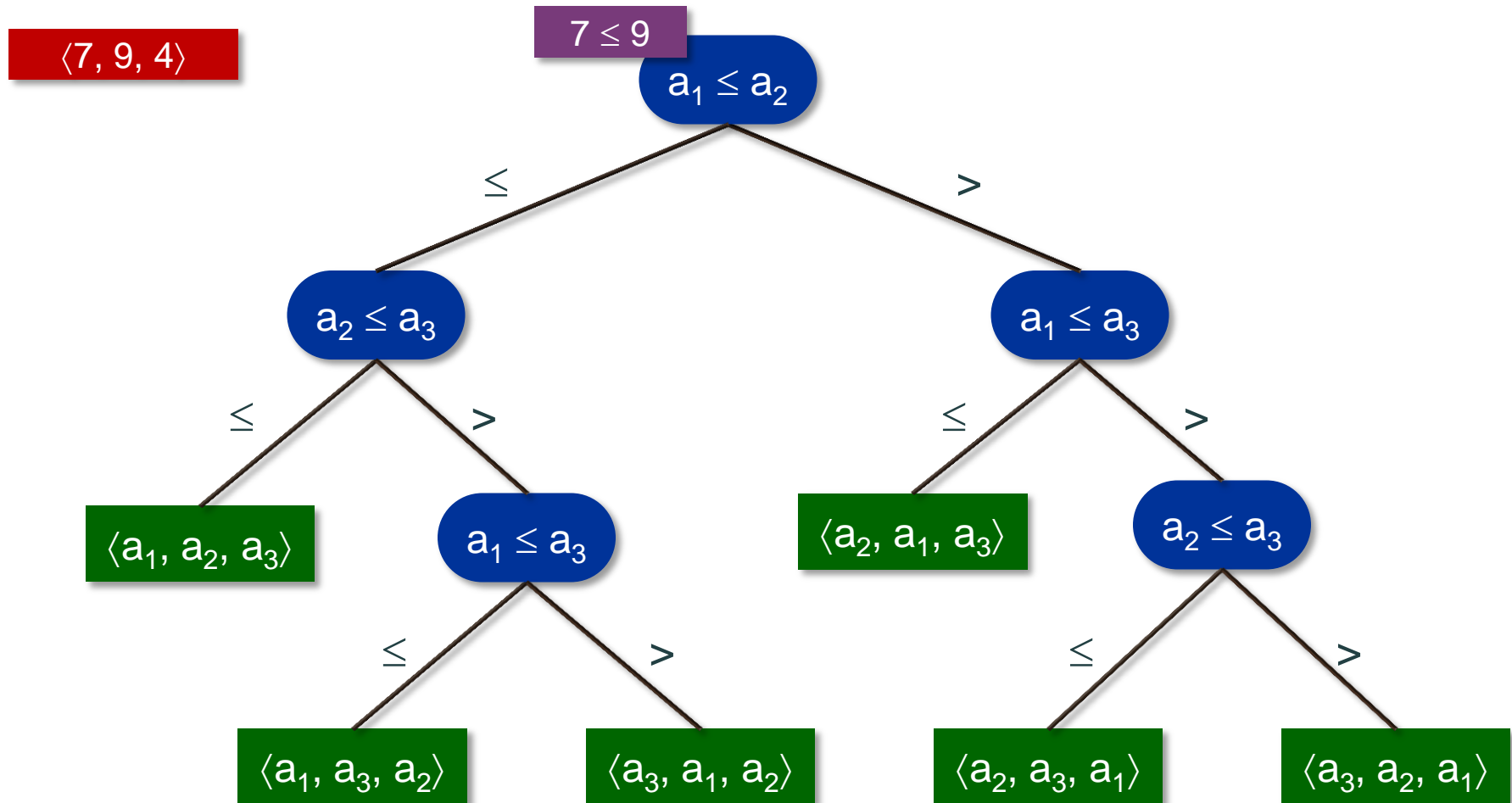
- Nenhum algoritmo de ordenação por comparação pode ser melhor do que $\Omega(n \log n)$
 - Para provar isso é necessário montar uma **árvore de decisão**
- A cada dois elementos a_i, a_j ,
 - Precisamos definir quem é o menor
 - Fazemos a comparação $a_i \leq a_j$

Árvore de decisão para o Insertionsort

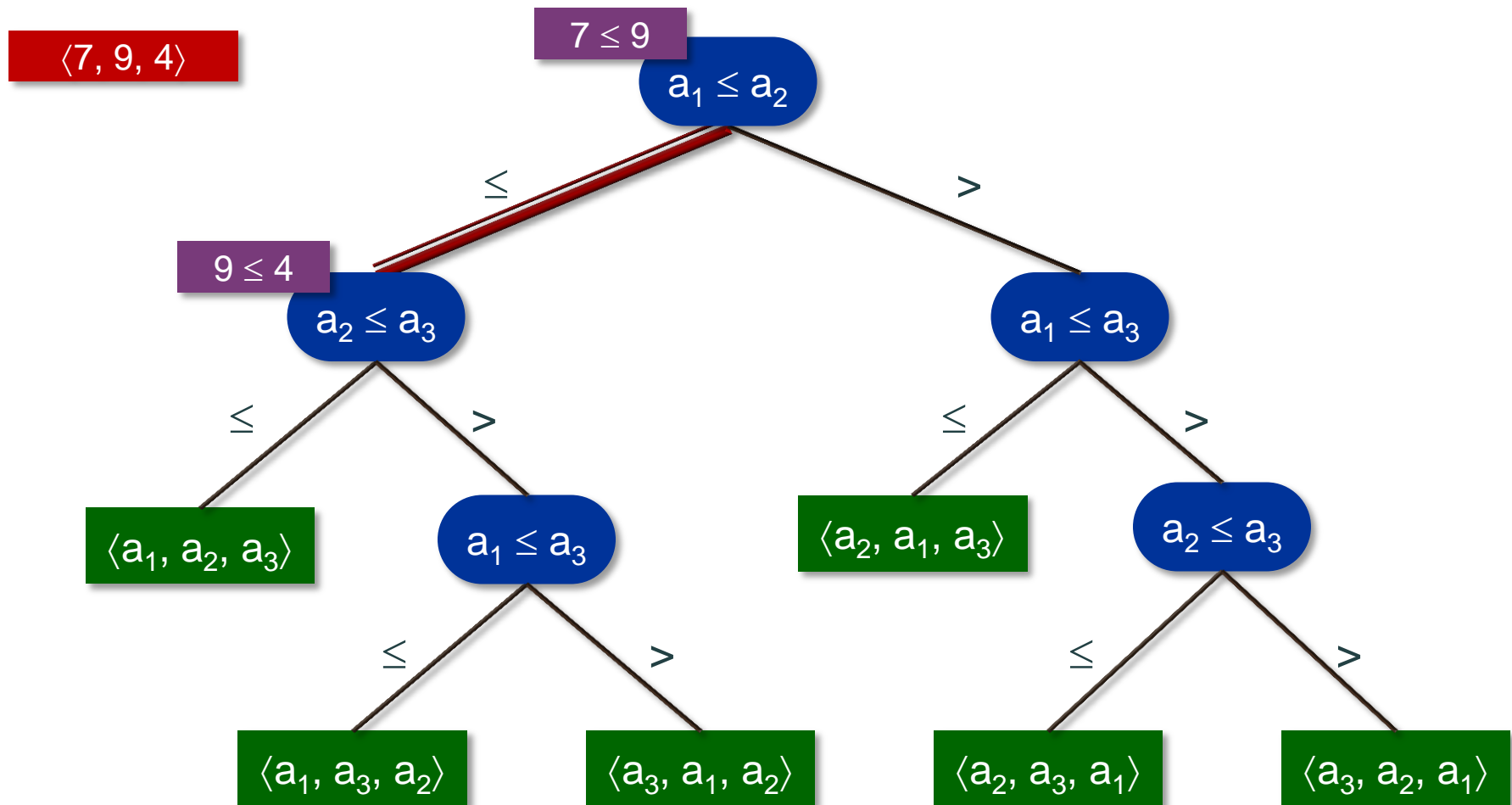
$\langle 7, 9, 4 \rangle$



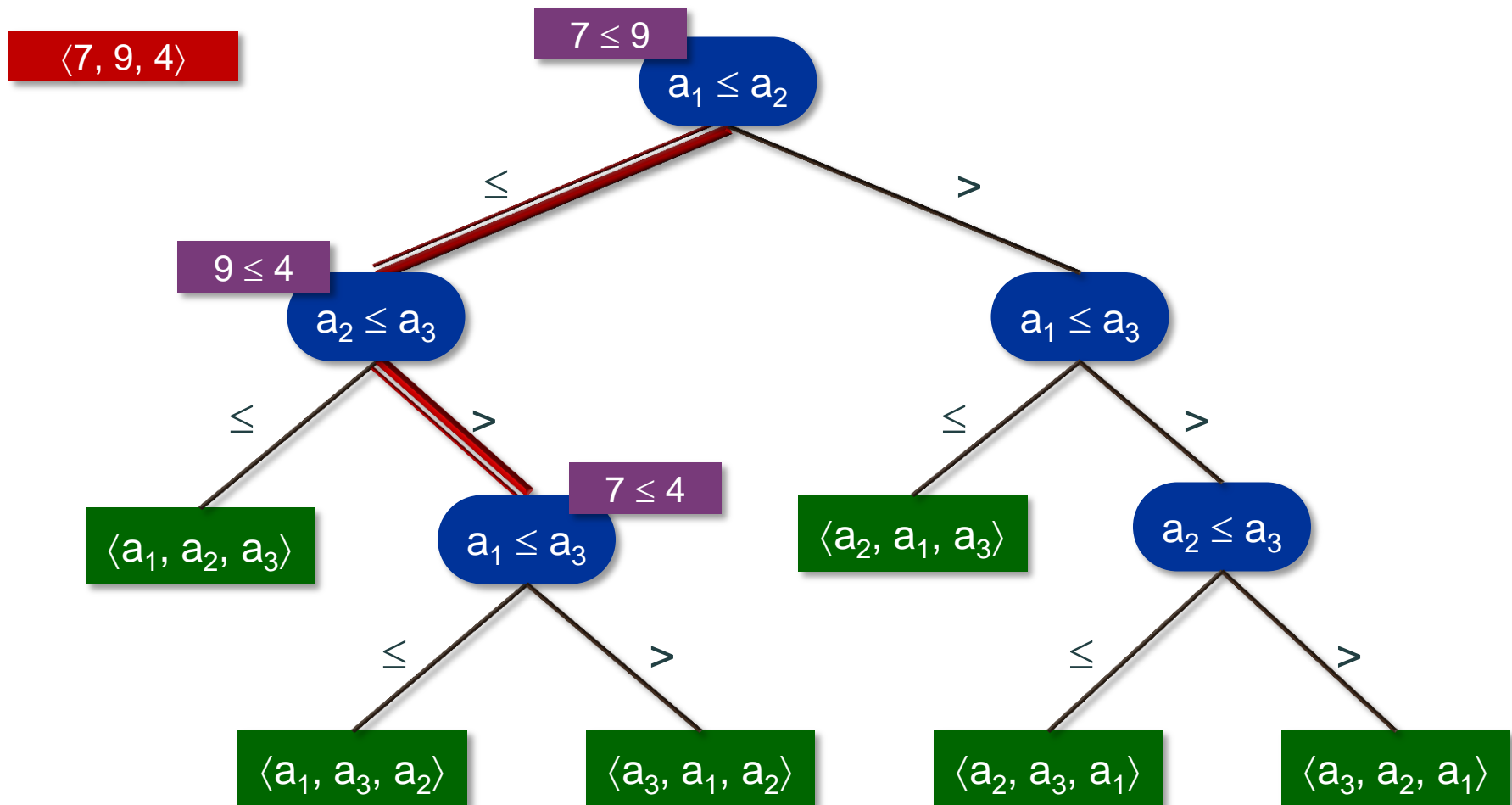
Árvore de decisão para o Insertionsort



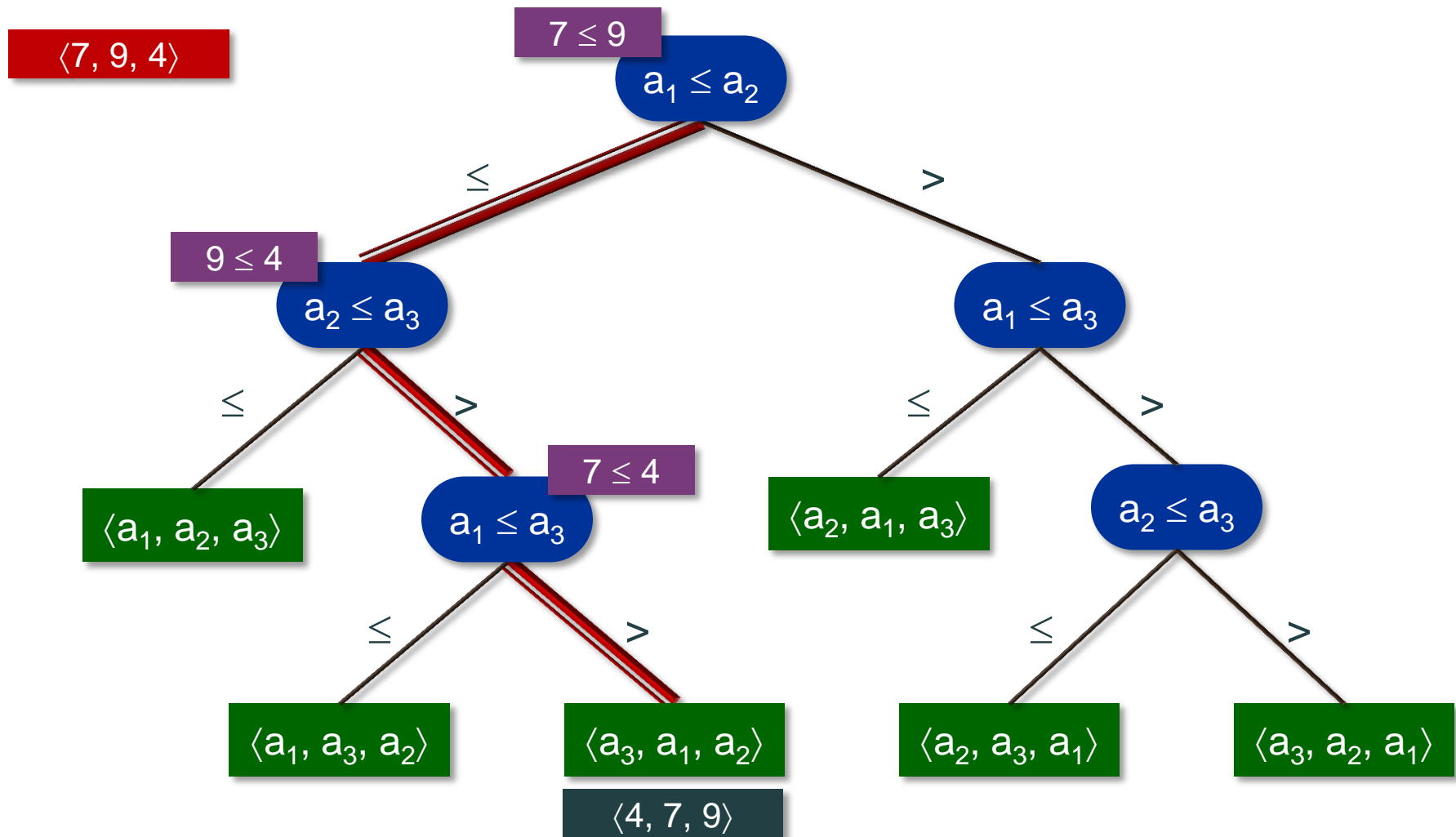
Árvore de decisão para o Insertionsort



Árvore de decisão para o Insertionsort



Árvore de decisão para o Insertionsort



Limite inferior

- O número de folhas (f) da árvore corresponde a todas as possíveis permutações do arranjo
 - $f = n!$
- O número de folhas (f) de uma árvore binária de altura h é no máximo 2^h
 - $f \leq 2^h$
- Logo
 - $n! \leq 2^h \Rightarrow h \geq \log(n!) \Rightarrow h \geq \Theta(n \log n) \Rightarrow h = \Omega(n \log n)$
- No pior caso o custo será altura da árvore!

Exercício

- Qual valor de q o PARTITION retorna quando todos os elementos do arranjo $A[p..r]$ são iguais?
- Modifique o PARTITION para que $q = (p+r) / 2$ quando todos os elementos do arranjo $A[p..r]$ são iguais