

Solutions to Problem Set 1 (6.042J)

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Problem 1

Assume that $\log_4 6 = \frac{m}{n}$ is rational. We have:

$$4^{\frac{m}{n}} = 6 \quad (1)$$

$$(4^{\frac{m}{n}})^n = 6^n \quad (2)$$

$$4^m = 6^n \quad (3)$$

$$(2 \times 2)^m = (2 \times 3)^n \quad (4)$$

Divide both sides by 2^n :

$$\frac{2^{2m}}{2^n} = 3^n \quad (5)$$

$$2^{2m-n} = 3^n \quad (6)$$

The LHS in (6) is even while the RHS is odd. This means $2^{2m-n} \neq 3^n$. Thus, $\log_4 6$ is irrational by contradiction.

Problem 2

$P(n) ::= \{\forall n \in \mathbb{N} | n \leq 3^{\frac{n}{3}} \equiv n^3 \leq 3^n\}$. Counterexample set: $C(n) ::= \{n \in \mathbb{N} | n \neq 3^{\frac{n}{3}}\}$. Calculation of $P(n)$ with $n \leq 4$ gives:

$$P(0) ::= 0 \leq 1$$

$$P(1) ::= 1 \leq 3$$

$$P(2) ::= 8 \leq 9$$

$$P(3) ::= 27 \leq 27$$

$$P(4) ::= 64 \leq 81$$

By the well-ordering principle, C has a minimum element c . This means $c - 1$ is a non-negative integer ≥ 4 and $P(c - 1)$ is true

$$P(c - 1) ::= (c - 1)^3 \leq 3^{c-1} \equiv 3(c - 1)^3 \leq 3^c \quad (7)$$

$$P(c) ::= c^3 \leq 3^c \equiv c^3 \leq 3(c - 1)^3 \quad (8)$$

$$\equiv c \leq \sqrt[3]{3}(c - 1) \quad (9)$$

(9) is always true for $c \geq 4$ and $P(n)$ holds for $n \leq 4$. This means $C(n)$ is an empty set by contradiction and $P(n)$ holds for every non-negative integer n .

Problem 3

a)

Truth table for $(P \text{ IMPLIES } Q) \text{ OR } (Q \text{ IMPLIES } P)$:

P	Q	$P \text{ IMPLIES } Q$	$Q \text{ IMPLIES } P$	$(P \text{ IMPLIES } Q) \text{ OR } (Q \text{ IMPLIES } P)$
T	T	T	T	T
T	F	F	T	T
F	T	T	F	T
F	F	T	T	T

The OR operation always gives True value no matter what truth values its variables may have.

b)

$R ::= (P \text{ AND } Q) \text{ OR } (\text{NOT}(P) \text{ AND } \text{NOT}(Q))$

c)

Case 1: P is valid iff $\text{NOT}(P)$ is not satisfiable

P is valid which means $\text{NOT}(P)$ is always false. Thus, $\text{NOT}(P)$ is not satisfiable since it can only produce false in this case.

Case 2: $\text{NOT}(P)$ is not satisfiable iff P is valid

$\text{NOT}(P)$ is not satisfiable means it can never be true, so $\text{NOT}(P)$ is always false which means P is always true.

d)

P_1, \dots, P_k is consistent iff P_1, \dots, P_k is satisfiable. We can write S as the following:

$S ::= P_1 \text{ AND } P_2 \text{ AND } \dots \text{ AND } P_k$