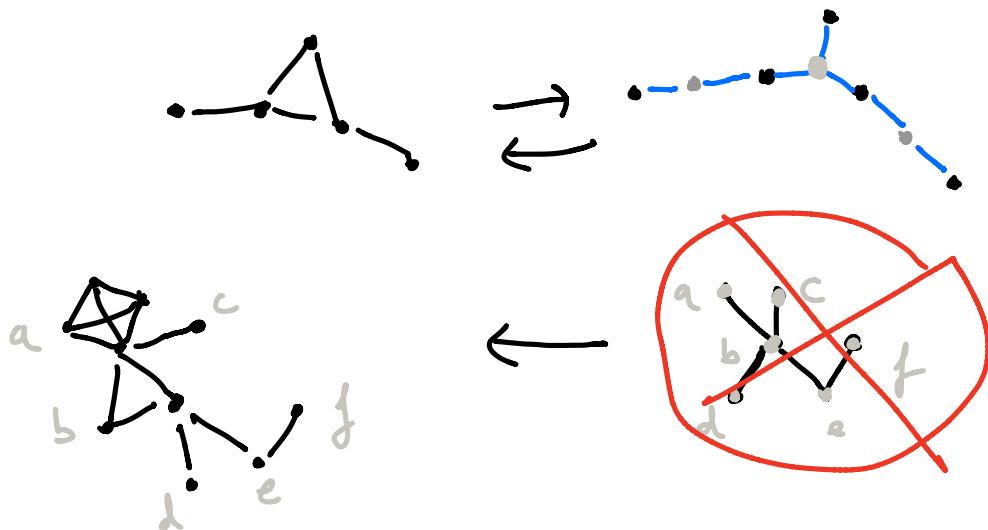


" G est un arbre de digue"

1) mettre un blanc au milieu de chaque digue



"plongement dans un arbre"
bipartite

⊕ $\delta(i,j)$ est l'^{l'unique} chemin géodésique (le + court entre)
2 points
[par les seuls graphes avec cette propriété]

- 1) S sépare $X \cup Y$ mais n: X ni Y
- 2) ~~$X \rightarrow Y$~~ (cohérence avec la suite) X, Y, Z
- 3) "does not separate" non! faut être contreposée
- 4) nif. d'épine: "pairwise separated"

Theorem 1.31. Let G be a block graph. The spine fan $\mathcal{F}(G)$ is a complete simplicial fan.

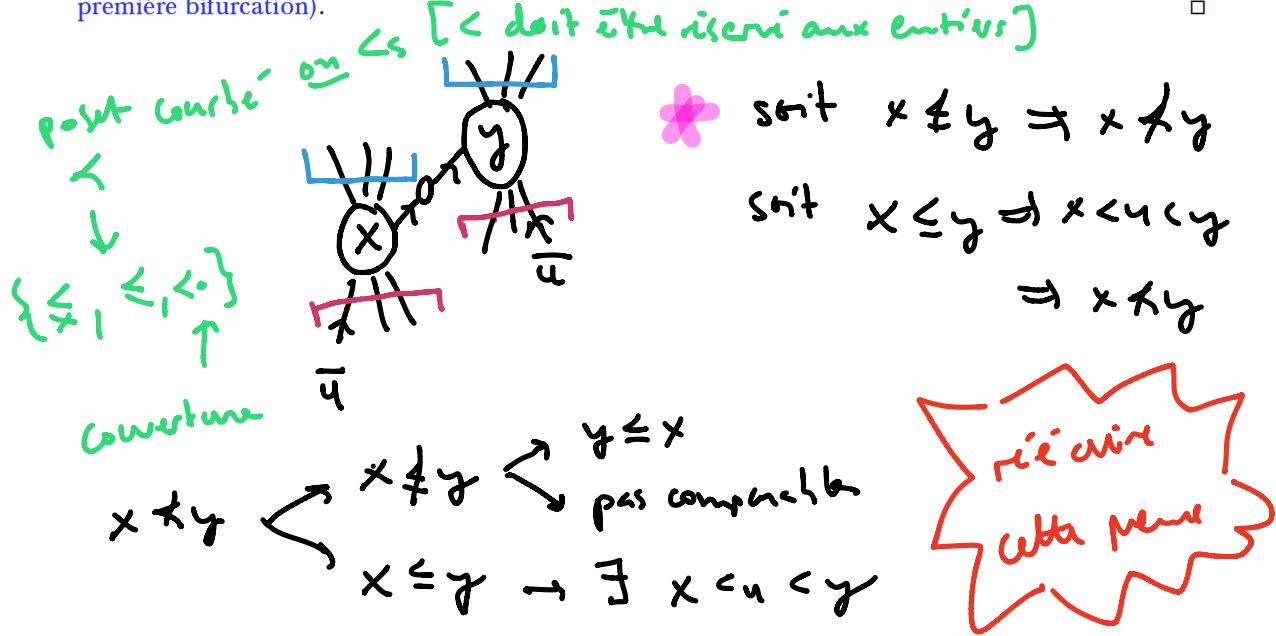
We will show that any linear order σ on $V(G)$ is a linear extension of a unique minimal spine S on G . Let $\sigma : V(G) \rightarrow \{1, \dots, n\}$ be a linear order on $V(G)$. We see it as a directed linear tree $\sigma(x_1) < \sigma(x_2) < \dots < \sigma(x_n)$.

maximal

Lemma 1.32. Let G be a block graph, and let S be a minimal G -spine. Then, we have

$$x < y \text{ in } S \iff \overline{s_<(y) \setminus x} \cup \overline{t^>(x) \setminus y} \text{ does not separate } x \text{ and } y, \text{ and} \\ \exists \text{ a linear extension } \sigma \text{ such that } \sigma(x) < \sigma(y). \quad \blacksquare = x < y$$

Proof. (\implies) Suppose that $\overline{s_<(y) \setminus x} \cup \overline{t^>(x) \setminus y}$ separates x and y . Then, in virtue of Corollary 1.8, there is some $u \in \overline{s_<(y) \setminus x} \cup \overline{t^>(x) \setminus y}$ that separates x and y , and we must have $x < u < y$. The fact that $x < y$ implies the existence of σ such that $\sigma(x) < \sigma(y)$ follows from the fact that S is an oriented tree. (\impliedby) If $y \lessdot x$, there is no linear extension σ such that $\sigma(x) < \sigma(y)$. If x and y are not comparable in S , there exists u in $s_<(y) \setminus x$ or in $t^>(x) \setminus y$ that separates x and y (je prends la première bifurcation). \square



Lemma 1.32. Let G be a block graph, and let S be a minimal G -spine. Then, we have

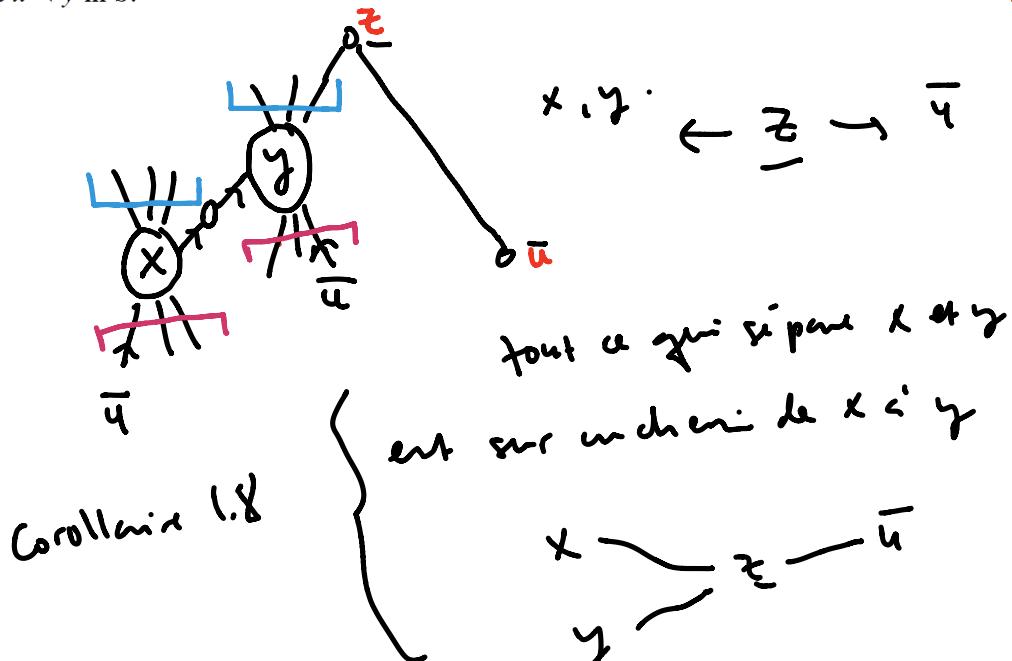
$$x < y \text{ in } S \iff \overline{s_<(y) \setminus x} \cup \overline{t^>(x) \setminus y} \text{ does not separate } x \text{ and } y, \text{ and} \\ \exists \text{ a linear extension } \sigma \text{ such that } \sigma(x) < \sigma(y).$$

Lemma 1.33. Let G be a block graph, and let S be a minimal G -spine. Then, max or ext. lin.

$$x < y \text{ in } S \iff \begin{array}{l} \text{For every linear extension } \sigma \text{ of } S, \text{ we have } \sigma(x) < \sigma(y), \\ \sigma^{-1}(s_<(\sigma(y)) \setminus \sigma(x)) \cup \sigma^{-1}(t^>(\sigma(x)) \setminus \sigma(y)) \text{ does not separate } x \text{ and } y, \text{ and} \\ \exists \text{ a linear extension } \sigma' \text{ such that } \sigma'(x) < \sigma'(y). \end{array}$$

Proof. (\implies) Suppose that $\sigma^{-1}(s_<(\sigma(y)) \setminus \sigma(x)) \cup \sigma^{-1}(t^>(\sigma(x)) \setminus \sigma(y))$ separates x and y . Then, in virtue of Corollary 1.8, there is some $u \in \sigma^{-1}(s_<(\sigma(y)) \setminus \sigma(x)) \cup \sigma^{-1}(t^>(\sigma(x)) \setminus \sigma(y))$ that separates x and y , and we must have $x < u < y$. The fact that $x < y$ implies the existence of σ' such that $\sigma'(x) < \sigma'(y)$ follows from the fact that S is an oriented tree.

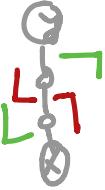
(\impliedby) There exists two linear extensions σ_1 and σ_2 , such that $\sigma_1^{-1}(s_<(\sigma_1(y)) \setminus \sigma_1(x)) = s_<(y) \setminus x$ and $\sigma_2^{-1}(t^>(\sigma_2(x)) \setminus \sigma_2(y)) = t^>(x) \setminus y$. So neither $s_<(y) \setminus x$ nor $t^>(x) \setminus y$ separate x and y , and by Proposition 1.10 we have that $s_<(y) \setminus x \cup t^>(x) \setminus y$ does not separate x and y . Now, Lemma 1.32 implies that $x < y$ in S . \square



$s_<(x) \quad t^>(y)$

Proof of the theorem. Let $\sigma : V(G) \rightarrow \{1, \dots, n\}$ be a linear order on $V(G)$. We define a poset S by the following covering relations

$$x < y \text{ in } S \iff \sigma(x) < \sigma(y) \text{ and } \overline{\sigma^{-1}(s_<(\sigma(y)) \setminus \sigma(x))} \cup \overline{\sigma^{-1}(t^>(\sigma(x)) \setminus \sigma(y))} \text{ does not separate } x \text{ and } y.$$



Using induction, we have the following characterization of S : $x' = y, y' = z$

$$x \leq y \iff \overline{\sigma(x)} \leq \overline{\sigma(y)} \text{ and } \exists x', y' \text{ with } \sigma(x) \leq \sigma(x'), \sigma(y') \leq \sigma(y) \text{ such that } \overline{\sigma^{-1}(s_<(\sigma(x')) \setminus \sigma(x))} \cup \overline{\sigma^{-1}(t^>(\sigma(y')) \setminus \sigma(y))} \text{ does not separate } x \text{ and } y.$$

- (1) S is a poset. Reflexivity and antisymmetry are clear. For transitivity, suppose that we have $x \leq y$ and $y \leq z$. That means there exists x_1, y_1, y_2, z_2 with $\sigma(x) \leq \sigma(x_1), \sigma(y_1) \leq \sigma(y) \leq \sigma(y_2), \sigma(z_2) \leq \sigma(z)$ such that

$$\begin{aligned} \overline{\sigma^{-1}(s_<(\sigma(x_1)) \setminus \sigma(x))} \cup \overline{\sigma^{-1}(t^>(\sigma(y_1)) \setminus \sigma(y))} &\text{ does not separate } x \text{ and } y, \text{ and} \\ \overline{\sigma^{-1}(s_<(\sigma(y_2)) \setminus \sigma(y))} \cup \overline{\sigma^{-1}(t^>(\sigma(z_2)) \setminus \sigma(z))} &\text{ does not separate } y \text{ and } z. \end{aligned}$$

Since $\sigma(x_1) \leq \sigma(y_2)$ and $\sigma(y_1) \leq \sigma(z_2)$, we have that

$$\overline{\sigma^{-1}(s_<(\sigma(x_1)) \setminus \sigma(x))} \cup \overline{\sigma^{-1}(t^>(\sigma(z_2)) \setminus \sigma(z))} \text{ does not separate } x \text{ and } z.$$

- (2) S is a tree. Vincent? ~~# rel. coverage = # sources - 1~~

- (3) S is a minimal spine. Let z be a vertex of S . It is enough to prove that the source set of each incoming edge of z lies in a distinct connected component of $G \setminus \{z\}$; the proof for the outgoing edges is symmetric. Let $x, y \in V(G)$ be such that $x < z$ and $y < z$. Without loss of generality, we can suppose that $\sigma(x) < \sigma(y) < \sigma(z)$.

If $x \leq y$, we have by definition that $\overline{\sigma^{-1}(t^>(\sigma(y)))}$ does not separate x and y . Since $\sigma(z) \in t^>(\sigma(y))$, we have in particular that $\{z\}$ does not separate x and y .

Now we suppose that $x < z$ and $y < z$. If $\{z\}$ does not separate x and y , using Proposition 1.10 we have that $\overline{\sigma^{-1}(s_<(\sigma(z)) \setminus \{\sigma(x), \sigma(y)\})} \cup \overline{\sigma^{-1}(t^>(\sigma(y)))}$ does not separate x and y either. Now let w be such that $\sigma(w) < \sigma(y)$. Then,

$$\overline{\sigma^{-1}(s_<(\sigma(y)) \setminus \sigma(x))} \cup \overline{\sigma^{-1}(t^>(\sigma(w)) \setminus \sigma(y))}$$

does not separate x and y , so we have $x \leq y$, a contradiction.

- (4) S is unique. Let S' be a minimal G -spine for which σ is a linear extension. From Lemma 1.33 we have that $x < y$ in S' implies $x < y$ in S . Since S and S' have the exact same number of covering relations (they are both oriented trees with $n-1$ oriented edges), we also have that $x < y$ in S implies $x < y$ in S' . (Toute injection entre deux ensembles de même cardinalité est une bijection)

□

(2) montrer que les couples x, y qui vérifient [...] sont au plus au nombre de $n-1$

(1,3) → on peut les remplacer par l'arg. géométrique

1324

$$i \geq \bar{i} \leq j$$

$$\sigma : \begin{matrix} 1 & \rightarrow & 1 \\ 2 & \rightarrow & 3 \\ \bar{i} & \rightarrow & 2 \\ j & \rightarrow & 4 \end{matrix}$$

$\Rightarrow \sigma(i) \leq \sigma(j)$ et i, j

about $d\sigma$ \approx $c c^* d\epsilon$

$$G \setminus (\overline{s_{\sigma(i)}} \cup t_{\sigma(j)})$$

1, M ???

