

Seminario de teoría de categorías

Paula González, Clara Narciso, Guillermo Peñafort

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Chapter 1

Categories: objects and arrows

Definition 1.0.1 (Category). A category \mathcal{C} consists of the following data:

- A class of *objects*,

$$Ob(\mathcal{C})$$

(we usually write $X \in \mathcal{C}$ instead of $X \in Ob(\mathcal{C})$).

- For every two objects $X, Y \in \mathcal{C}$, a class of *morphisms* (or *arrows*) from X to Y ,

$$Hom_{\mathcal{C}}(X, Y)$$

(we sometimes write $Hom(X, Y)$ or $\mathcal{C}(X, Y)$ instead of $Hom_{\mathcal{C}}(X, Y)$, and $f: X \rightarrow Y$ instead of $f \in Hom(X, Y)$).

- A *composition law*, associating to morphisms $f: X \rightarrow Y$ and $g: Y \rightarrow Z$, a *composition morphism*

$$g \circ f: X \rightarrow Z.$$

Arrows and the composition law are subject to the following conditions:

- *Associativity*: given composable morphisms f, g, h ,

$$h \circ (g \circ f) = (h \circ g) \circ f.$$

- Existence of *identity arrows*: for any X , there exists a morphism $id_X: X \rightarrow X$, such that, for all $f: X \rightarrow Y$ and $g: Z \rightarrow X$,

$$f \circ id_X = f, \quad id_X \circ g = g.$$

Example. *Some basic examples:*

1. The category **Set** of all sets:

- $Ob(\mathbf{Set})$ is the class of all sets.
- $Hom(X, Y)$ is the set of all mappings $f: X \rightarrow Y$
- $g \circ f$ is the usual composition of mappings.

2. The category **Top** of topological spaces has

- topological spaces (X, \mathcal{T}) as objects (in other words, $Ob(\mathbf{Set})$ is the class of all topological spaces).
- continuous mappings $f: (X, \mathcal{T}) \rightarrow (Y, \mathcal{T})$ as morphisms, with the usual composition.

Notice that this requires two results from topology: (1) composition of continuous mappings is continuous and (2) the identity mapping is continuous.

3. The category **Group** of groups has

- groups as objects.
- group homomorphisms as arrows, with the usual composition.

This requires the result that compositions of group homomorphisms are group homomorphisms.

All of the above examples consist of certain sets (with extra structure in some cases) and mappings between them, but this is not always the case:

1. Posets give rise to categories:

For example, we may regard (\mathbb{N}, \leq) as a category having

- natural numbers as objects.
- for any $n, m \in \mathbb{N}$,

$$Hom(n, m) = \begin{cases} \{n \rightarrow m\}, & \text{if } n \leq m \\ \emptyset, & \text{if } n \not\leq m \end{cases}$$

Here $\{n \rightarrow m\}$ is just any set with a single element, which we regard as the arrow $n \rightarrow m$. These sets must be chosen so that $n \rightarrow m$ equals $n' \rightarrow m'$ if and only if $n = n'$ and $m = m'$.

Since there is at most one arrow between objects, composition is defined in the only possible way.

Observe that the reflexivity and transitivity properties of partial orders are needed in order for this to form a category.

Similarly, any set X gives rise to a category

$$(\mathcal{P}(X), \subseteq),$$

and any topological space (X, \mathcal{T}) gives rise to a category

$$(\mathcal{T}, \subseteq).$$

2. Groups as categories:

Given a group G , with binary operation “ \cdot ”, the category BG has

- a single object, written \bullet .
- elements $g \in G$ as arrows $\bullet \xrightarrow{g} \bullet$.
- The composition of two arrows $g, g' \in G$ is the arrow $g \circ g' = g \cdot g'$.

3. The category associated to a unitary ring:

Given a unitary ring R , with binary operations “ $+$ ” and “ \cdot ”, the category \mathcal{C}_R has

- a single object, written \bullet .
- elements $r \in R$ as arrows $\bullet \xrightarrow{r} \bullet$.
- The composition of two arrows $r, r' \in R$ is the arrow $r \circ r' = r \cdot r'$.

Observe that the condition that R is unitary is required by the axiom that categories must have identity arrows.

4. Given a field k , the matrix category \mathbf{Mat}_k has

- natural numbers as objects.
- For any $n, m \in \mathbb{N} \setminus \{0\}$, a morphism $n \rightarrow m$ is an $m \times n$ matrix with entries in k . Composition of such morphisms corresponds to product of matrices. There is a single arrow $0 \rightarrow m$ and $m \rightarrow 0$, for every $m \in \mathbb{N}$, and composition with these arrows is determined by lack of choice.

Definition 1.0.2. Given a category \mathcal{C} , the opposite category \mathcal{C}^{op} has

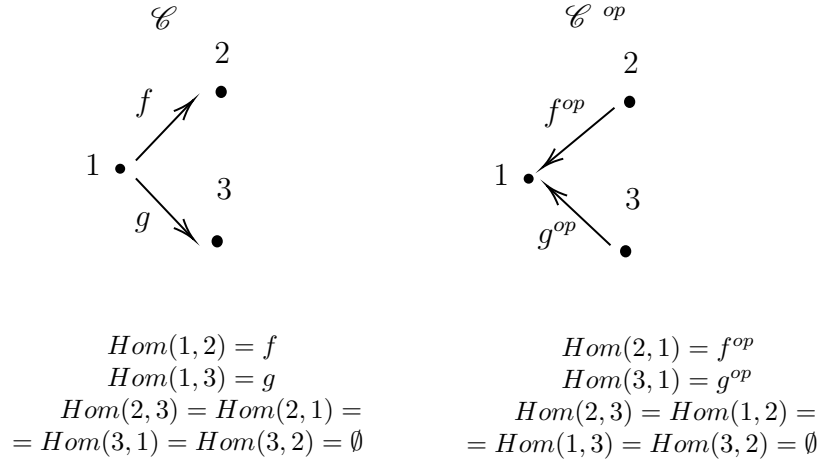
- same objects as \mathcal{C} , that is

$$Ob(\mathcal{C}^{op}) = Ob(\mathcal{C}).$$

- Given $X, Y \in \mathcal{C}^{op}$, each morphism $X \rightarrow Y$ in \mathcal{C} is an arrow $Y \rightarrow X$ in \mathcal{C}^{op} , that is,

$$\text{Hom}_{\mathcal{C}^{op}}(X, Y) = \text{Hom}_{\mathcal{C}}(Y, X).$$

Example. *This is an example of a simple category and its opposite.*



Definition 1.0.3 (Small and Locally small categories).

A category \mathcal{C} is **locally small** if, for all $X, Y \in \mathcal{C}$, $\text{Hom}(X, Y)$ is a set (as opposed to a proper class).

A category is **small** if it is locally small and $Ob(\mathcal{C})$ is also a set.

Example. *All categories introduced above are locally small. An example of non-locally small category is the one with*

- a single object \bullet .
- Any set X is an arrow $\bullet \xrightarrow{X} \bullet$. The composition of $\bullet \xrightarrow{X} \bullet$ and $\bullet \xrightarrow{Y} \bullet$ is $\bullet \xrightarrow{X \cup Y} \bullet$.

Remark. *From now on, we will always consider locally small categories.*

1.1 Subcategories

Definition 1.1.1. Given a category \mathcal{C} , we say that a subset $\mathcal{D} \subseteq \mathcal{C}$ is a **subcategory** of \mathcal{C} if it is a subcollection of objects and morphisms, satisfying that

- All domains and codomains of morphisms in \mathcal{D} are contained.
- The identity morphism of every object in \mathcal{D} is contained.
- Every possible composition between morphisms in \mathcal{D} is contained.

Example. *Some basic examples of subcategories are*

1. **Haus** is the subcategory of **Top** consisting of Hausdorff topological spaces.
2. **Ab** is the subcategory of **Group** consisting of abelian groups.

Definition 1.1.2 (Full and Wide subcategories).

Let \mathcal{D} be a subcategory of \mathcal{C} .

- \mathcal{D} is a **full subcategory** of \mathcal{C} if $\forall X, Y \in \mathcal{D}$, it is satisfied that

$$\text{Hom}_{\mathcal{D}}(X, Y) = \text{Hom}_{\mathcal{C}}(X, Y)$$

- \mathcal{D} is a **wide subcategory** of \mathcal{C} if

$$\text{Ob}(\mathcal{D}) = \text{Ob}(\mathcal{C})$$

Example. *The subcategory \mathcal{C} of **Group** consisting of*

- $\text{Ob}(\mathcal{C}) = \text{Ob}(\mathbf{Group})$.
- *The arrows in \mathcal{C} are the group isomorphisms.*

Example. *We consider the following categories:*

1. *The category **Rng** of all rings:*
 - $\text{Ob}(\mathbf{Rng})$ is conformed by all rings.
 - Ring homomorphisms are arrows, with the usual composition.
2. *The category **Ring** of all unitary rings:*

- $\text{Ob}(\mathbf{Ring})$ is conformed by all unitary rings.
- Ring homomorphisms $\mathcal{R} \longrightarrow \mathcal{S}$ satisfying $1_{\mathcal{R}} \mapsto 1_{\mathcal{S}}$ are arrows with the usual composition.

It is clear that **Ring** is a subcategory of **Rng** but it is not full, as morphisms in **Ring** need to take the multiplicative identity of the first ring into the multiplicative identity of the second one.

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Proposition 1.1.1. Let \mathcal{D} be a subcategory of \mathcal{C} . Then, \mathcal{D}^{op} is a subcategory of \mathcal{C}^{op}

1.2 Special Objects

Definition 1.2.1. An object $I \in \mathcal{C}$ is *initial* if, for every object $X \in \mathcal{C}$, there exists a unique morphism

$$I \rightarrow X.$$

Definition 1.2.2. An object $T \in \mathcal{C}$ is *terminal* (or *final*) if, for every object $X \in \mathcal{C}$, there exists a unique morphism

$$X \rightarrow T.$$

Terminal and initial are dual notions.

Definition 1.2.3. An object is a *zero object* if it is initial and terminal.

Example. *Some basic examples:*

1. *Special objects in category (\mathbb{N}, \leq) :*

- *The initial object in this category is 0, since $0 \leq n \forall n \in \mathbb{N}$ and so there exists a unique arrow from 0 to any other natural number. The uniqueness comes from the fact that there exists only one arrow that goes from one object to another (greater) one.*
- *There is no final object in this category, since there isn't any natural number greater than all of them.*
- *Since there's no final object, there's no zero object in this category.*

Note: if instead of \mathbb{N} we considered \mathbb{Z} , we wouldn't have an initial object either.

2. *In **Set** and **Top**:*

- *We have that \emptyset is initial since we can always consider the inclusion as the only arrow from \emptyset to any set or topological space.*
- *We have that any set with only an element is terminal, since there's only one way to define a map from a set to the singleton.*
- *There are no zero objects since \emptyset is the only initial object in these categories.*

3. *In the **Ring** category, \mathbb{Z} is an initial object since all arrows must send 0 and 1 to the new ring's neutral elements, and $\{0\}$ is terminal. However, in **Rng** we have that $\{0\}$ is a zero object.*

Exercises

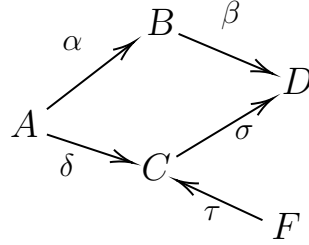
Exercise 1.2.1. *If we have two initial objects from the same category, then they're isomorphic.*

Exercise 1.2.2. *Let X_1 be an initial object, and let $X_2 \cong X_1$. Then X_2 is initial.*

Analogous results are valid for terminal objects.

Exercise 1.2.3. *If $|X| \neq 1$ then X is not terminal.*

Exercise 1.2.4. *Let's consider category \mathcal{B} that has objects A, B, C, D, F and arrows $\alpha, \beta, \gamma, \delta, \sigma, \beta\alpha = \sigma\delta$ and $\sigma\tau$. What are the initial and terminal objects?*



1.3 Monomorphisms, epimorphisms and isomorphisms

Monos & epis

Definition 1.3.1 (Monomorphism). Let \mathcal{C} be a category, and $X, Y, Z \in \mathcal{C}$.

We say that an arrow $f : X \longrightarrow Y$ is a **monomorphism** if for any pair of arrows $g, g' : Z \longrightarrow X$ holding the equality $f \circ g = f \circ g'$, they satisfy $g = g'$.

Definition 1.3.2 (Epimorphism). Let \mathcal{C} be a category, and $X, Y, Z \in \mathcal{C}$.

We say that an arrow $f : X \longrightarrow Y$ is a **epimorphism** if for any pair of arrows $h, h' : Y \longrightarrow Z$ holding the equality $h \circ f = h' \circ f$, they satisfy $h = h'$.

Example. Some examples of monomorphisms and epimorphisms are:

- The inclusion $\mathbb{Z} \hookrightarrow \mathbb{Q}$ is both a monomorphism and an epimorphism in **Ring**.
- A morphism $f : X \rightarrow Y$ in **Haus** is an epimorphism if and only if $f(X)$ is dense in Y .

Definition 1.3.3 (Split monomorphism and epimorphism).

Let \mathcal{C} be a category, $X, Y \in \mathcal{C}$ and $f \in \text{Hom}_{\mathcal{C}}(X, Y)$, $g \in \text{Hom}_{\mathcal{C}}(Y, X)$ such that

$$g \cdot f = \text{id}_X$$

Then f is called a **split monomorphism** and g is called **split epimorphism**.

Proposition 1.3.1. Split monomorphisms and epimorphisms are respectively monomorphisms and epimorphisms.

Proposition 1.3.2. Let $f : X \longrightarrow Y$ be a monomorphism (or epimorphism) in \mathcal{C} . Then, f is an epimorphism (or monomorphism) in \mathcal{C}^{op} .

Proposition 1.3.3 (Operations with monomorphism). (*epimorphism by duality*):

- Composition of monomorphisms is monomorphisms.
- If $g \circ f$ is a monomorphism, then f is a monomorphism.

Proposition 1.3.4 (Subcategories reflect monomorphism).

Let $\mathcal{A} \hookrightarrow \mathcal{B}$ be a subcategory, and let f be an arrow in \mathcal{A} . If f is a monomorphism in \mathcal{B} , then f is a monomorphism in \mathcal{A} .

Proposition 1.3.5 (Special arrows from special objects).

If I is initial, any arrow $X \rightarrow I$ is an epimorphism

Remark. It's important to separate the usual definitions for monomorphism and epimorphisms used in representable categories from the more general, categorical one:

- Being mono doesn't always imply being injective, just like being epi doesn't always imply being surjective. The first case is more difficult to see since in general in categories made from algebraic structures we do have that mono implies injectivity, but a counterexample for epimorphisms is the inclusion $\mathbb{Z} \hookrightarrow \mathbb{Q}$ in **Ring**.
- Being a split monomorphism implies being an injection in **Sets** but the other way around is not true.

Exercises

Exercise 1.3.1. Prove proposition [1.3.1](#)

Exercise 1.3.2. The inclusion $\mathbb{Z} \hookrightarrow \mathbb{Q}$ is an epimorphism in **Ring**.

Exercise 1.3.3. Prove proposition [1.3.5](#)

Exercise 1.3.4. Split monomorphism implies injectivity in **Sets**

Isomorphisms

Definition 1.3.4 (Isomorphism). Let \mathcal{C} be a category, and $X, Y \in \mathcal{C}$.

An arrow $f : X \rightarrow Y$ is said to be an **isomorphism** if there exists another arrow $f^{-1} : Y \rightarrow X$ satisfying that $f \circ f^{-1} = id_Y$ and $f^{-1} \circ f = id_X$.

Example. *Some examples of isomorphisms are:*

1. *A mapping $f: X \rightarrow Y$ is an isomorphism in **Set** if and only if it is a bijection.*
2. *A group homomorphism ϕ is an isomorphism in **Group** if and only if it is bijective.*
3. *The isomorphisms in **Top** are the homeomorphisms.*
4. *The only isomorphisms in (\mathbb{N}, \leq) are the identities.*
5. *A morphism $X \rightarrow Y$ in \mathcal{C} is the same as an isomorphism $Y \rightarrow X$ in \mathcal{C}^{op} .*
6. *Given an object $X \in \mathcal{C}$, an **automorphism** of X is just an isomorphism $X \rightarrow X$. This gives rise to an obvious automorphism subcategory*

$$\text{Aut}(X) \hookrightarrow \text{End}(X) \hookrightarrow \mathcal{C}.$$

Proposition 1.3.6 (Basic properties of isomorphisms). *1. Inverses of isomorphisms are unique.*

2. *Any composition of isomorphisms is an isomorphism.*
3. *If the composition of two arrows is an isomorphism and one of them is an isomorphism, then so is the other.*

Proposition 1.3.7 (Isomorphisms and subcategories).

Let $\mathcal{A} \hookrightarrow \mathcal{B}$ be a subcategory, and let f be an arrow in \mathcal{A} .

1. *If f is an isomorphism in \mathcal{A} , then f is an isomorphism in \mathcal{B}*
2. *Full subcategories reflect isomorphisms:
If \mathcal{A} is a full subcategory and f is an isomorphism in \mathcal{B} , then f is an isomorphism in \mathcal{A} .*

(Obs: the fullness condition above cannot be dropped).

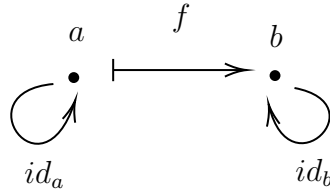
Proposition 1.3.8 (Characterization of isomorphisms).

A morphism is an isomorphism if and only if it is a split monomorphism and a split epimorphism.

Remark. *The condition of being split is necessary. We'll give some examples.*

1. Considerate a category \mathcal{C} given by

- $Ob(C) = \{a, b\}$
- Appart from the identitty arrows, there only exists an arrow $f : a \longrightarrow b$.



It is clear that f is a monomorphism and an epimorphism, but can't be an isomorphism because it doesn't exist an arrow $g : a \longrightarrow b$.

2. In **Top** category, we have that

- An arrow is a monomorphism if and only if it is injective.
- An arrow is an epimorphism if and only if it is surjective.

In conclusion, every bijective arrow in **Top** is a mono and an epimorphism but it is only an isomorphisms if its inverse is continuous.