

**Q130.** What from the following statements concerning the intermediate representations (IR) is incorrect?

- (A) IR simplify retargeting to a new host.
- (B) IR simplify writing of a compiler for another language to the same host.
- (C) Properly fixed IR simplify semantic aspects of the source language.
- (D) IR break the compiler into manageable pieces

Answer :-D

IR break the compiler into manageable pieces .false

**Q131.** Which of the following is not a typical intermediate representation?

- (A) Three address code
- (B) Directed acyclic graphs

- (C) Abstract syntax trees  
(D) Annotated parse trees

131) Answer  
3-address code  
DAG  
Abstract syntax tree  
all are forms of intermediate representation

**Q132.** Consider the following code segment.

```
x = a + b;  
y = x * c;  
p = y + y;  
p = p + d;  
p = p - d;  
p = p + e;
```

The minimum number of total variables required to convert the above code segment to static single assignment form is \_\_\_\_\_

132) Answer 9

$$x = a + b$$

$$y = x * c$$

$$p = y + y$$

$$p = z + d$$

$$p = z - d$$

$$p = p + e$$

} These both equations ultimately results in  $p = z$

$$x = a + b$$

$$y = x * c$$

$$p = y + y$$

$$p = z$$

$$p = p + e$$

→ SSA

$$x = a + b$$

$$y = x * c$$

$$p = y + y$$

$$p1 = z$$

$$p2 = p1 + e$$

No. of temp variables

$x, a, b, c, y, p, p1, p2, e$

Q133. Consider the following code segment.

$a := b + c$

$b := c + d$

$d := b + c$

$a := a + c$

$e := a + b$

The minimum number of total variables required to convert the above code segment to static single assignment form is \_\_\_\_\_.

133) Answer 6

$$a = b + c$$

$$b = c + d$$

$$d = b + c$$

$$a_1 = a + c$$

$$e = a_1 + b$$

Thus minimum no. of total variables  
 $a, b, c, d, e, a_1$   
6

Q134. The least number of temporary variables required to create a three-address code in static single assignment form for the expression  $a = b * d - c + b * e - c$  is \_\_\_\_\_

134) Answer 4

$$a = b * d - c + b * e - c$$

$$t_1 = b * d$$

$$t_2 = b * e$$

$$t_3 = t_1 - c$$

$$t_4 = t_3 + t_2$$

$$a = t_4 - c$$

~~no temporary~~ 4 temporary variables are required

Q135. Consider the following code



```

a = 10;
b = 20;
if (a > b) {
c = 30;
e = c - b;
} else {
d = b;
e = a + d;
}

```

Determine the minimum number of registers needed to execute above code without spilling? \_\_\_\_\_.

135) Answer: 2

```

a = 10
b = 20
if (a > b) goto L1
c = 30
e = c - b
L1: d = b
    e = a + d
    t1 = φ(e1, e)

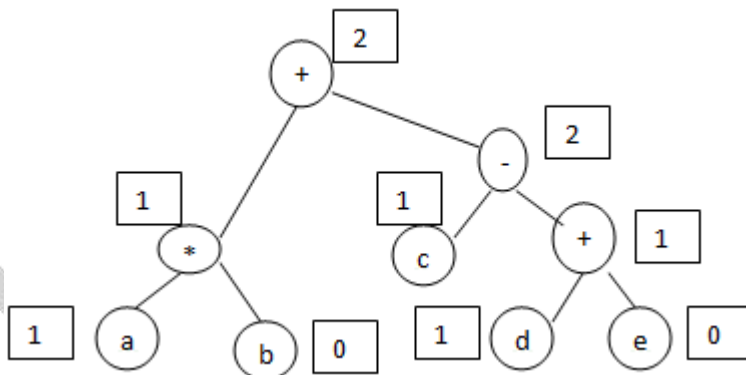
```

At one time only 2 registers are required without spilling.

**Q136.** What is the minimal number of registers necessary for the generation of code corresponding with the following expression if one operand can be a memory location.  $(a * b) + (c - (d + e))$

**Answer: 2**

**Solution:**



**Q137.** What is the minimal number of registers necessary for the generation of code corresponding with the following expression if all operands in cpu register?

$(a * b) + (c - (d + e))$

(A)2

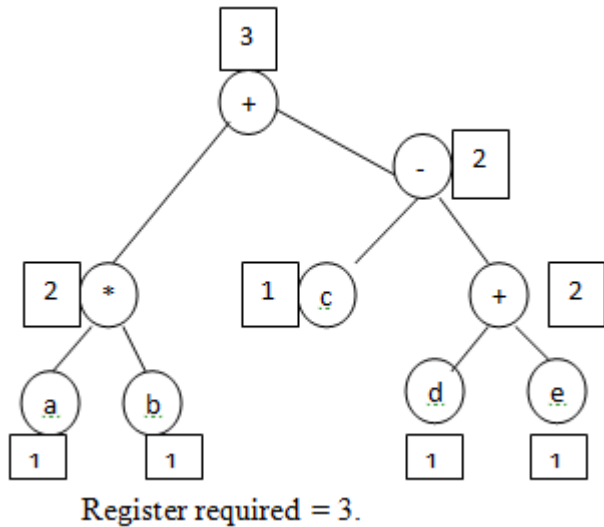
(B)3

(C)4

(D)5

Answer: B

Solution:



**Q138.** Consider the following code

```
a = 1 ;  
b = 10;  
c = 20;  
d = a + b;  
e = c + d;  
f = c + e;  
b = c + e;  
e = b + f;  
d = 5 + e;  
return d+f;
```

What is the fewest number of registers that is needed for this program, without spilling? \_\_\_\_\_

**Answer: 3**

```

a = 1 ;
b = 10;
c = 20;
d = a + b;
e = c + d;
f = c + e;
b = c + e;
e = b + f;
d = 5 + e;
return d+f;

```

```

a = 1 ;
b = 10;
c = 20;
d = a + b;
e = c + d;
f = c + e;
e = 2*f
d = 5 + e;
return d+f;

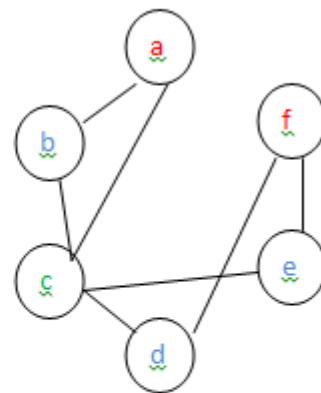
```

We can reduce the expression  $b=c+e$   
And replace  $e=b+f$  to  $e=2*f$

```

a = 1 ;
b = 10;
c = 20;
d = a + b;
e = c + d;
f = c + e;
e = 2*f
d = 5 + e;
return d+f;

```



3 colors are required thus register required will be 3

For next three questions consider the following three address code: