

DBMS - Transaction Concurrency Control ①

conflict equivalence:- we can't swap the conflicting pairs.

lock based concurrency control:-

★ 2PL - guarantees serializability of conflict + serializability
2 phase locking
Conflict serializability \Rightarrow Serializability

Phase 1:- Growing phase / expanding phase: acquire lock

Phase 2: Shrinking phase: release lock

Rules:

- (1) T_i can't acquire lock after a lock has been released
- (2) On commit/rollback all locks are released by default.

Note:- lock point is point in a transaction where transaction has successfully acquired all the locks and has started releasing the locks

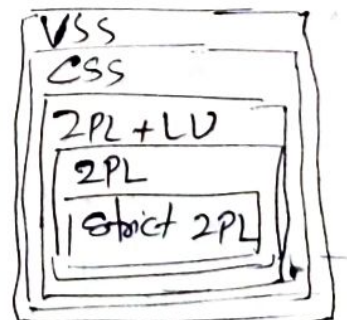
lock point used to find equivalent serializable schedule.

★ lock upgrade:- upgrade the shared lock to exclusive lock without ~~a~~ unlocking it before upgrading.
(LV)

Note:-

2PL \nRightarrow recoverability

2PL have deadlocks and starvation



★ Strict 2PL:- guarantees (CSS) + ~~strict~~ recoverable

It follows all rules of 2PL + all exclusion lock must be held till commit/rollback but there is no such restriction with shared lock.

Strict 2PL \Rightarrow deal locks & starvation

Strong Strict 2PL / Rigorous 2PL

2PL + all locks (shared or exclusive) should be held till transaction commit or rollback.

Time Stamp Ordering Protocol

(i) each transaction will given at coming time

(ii) Ordering of ~~two~~ must be $S(T_1) < S(T_2)$

and serializability must be $\boxed{T_1 \rightarrow T_2}$

(iii) If there is conflict then serializability must be $T_1 \rightarrow T_2$, if not $T_1 \rightarrow T_2$ then T_1 will be rollbacked and come as new transaction with new timestamp

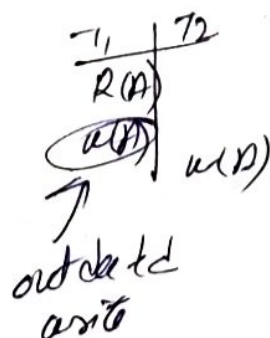
benefit

- (i) guaranty for serializability
- (ii) Deadlock free
- (iii) Starvation possible

Thomas Write Rule

Advanced version of [timestamp ordering Protocol]

→ Ignoring the out-dated writes is called Thomas write rule



this will overwrite the value