

Leetcode solution of histogram

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For any bar i the maximum rectangle is of width $r - l - 1$ where r - is the last coordinate of the bar to the **right** with height $h[r] \geq h[i]$ and l - is the last coordinate of the bar to the **left** which height $h[l] \geq h[i]$. So if for any i coordinate we know its utmost higher (or of the same height) neighbors to the right and to the left, we can easily find the largest rectangle:

From <[https://leetcode.com/problems/largest-rectangle-in-histogram/discuss/28902/5ms-O\(n\)-Java-solution-explained-\(beats-96\)](https://leetcode.com/problems/largest-rectangle-in-histogram/discuss/28902/5ms-O(n)-Java-solution-explained-(beats-96))>

```
int maxArea = 0;
for (int i = 0; i < height.length; i++) {
    maxArea = Math.max(maxArea, height[i] * (lessFromRight[i] - lessFromLeft[i] - 1));
}
```

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The main trick is how to effectively calculate `lessFromRight` and `lessFromLeft` arrays. The trivial solution is to use $O(n^2)$ solution and for each i element first find its left/right neighbour in the second inner loop just iterating back or forward:

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```
for (int i = 1; i < height.length; i++) {
    int p = i - 1;
    while (p >= 0 && height[p] >= height[i]) {
        p--;
    }
    lessFromLeft[i] = p;
}
```

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The only line change shifts this algorithm from $O(n^2)$ to $O(n)$ complexity: we don't need to rescan each item to the left - we can reuse results of previous calculations and "jump" through indices in quick manner:

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```

while(p >= 0 && height[p] >= height[i]) {
    p = lessFromLeft[p];
}

```

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Here is the whole solution:

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```

public static int largestRectangleArea(int[] height) {
    if (height == null || height.length == 0) {
        return 0;
    }
    int[] lessFromLeft = new int[height.length]; // idx of the first bar the left that is lower than current int[]
    lessFromRight = new int[height.length]; // idx of the first bar the right that is lower than
    current lessFromRight[height.length - 1] = height.length;
    lessFromLeft[0] = -1;
    for (inti = 1; i < height.length; i++) {
        int p = i - 1;
        while (p >= 0 && height[p] >= height[i]) {
            p = lessFromLeft[p];
        }
        lessFromLeft[i] = p;
    }
    for (inti = height.length - 2; i >= 0; i--) {
        int p = i + 1;
        while (p < height.length && height[p] >= height[i]) {
            p = lessFromRight[p];
        }
        lessFromRight[i] = p;
    }
    int maxArea = 0;
    for (inti = 0; i < height.length; i++) {
        maxArea = Math.max(maxArea, height[i] * (lessFromRight[i] - lessFromLeft[i] - 1));
    }
    return maxArea;
}

```

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