

Linear Search (Searching Algo)

Case: 1

Successful search

0	1	2	3	4	5	6
9	8	6	3	4	7	2
↑	↑	↑	↑	↑	↑	↑

target = 7

Case: 2

Unsuccessful search

0	1	2	3	4	5	6
9	8	6	3	4	7	2
↑	↑	↑	↑	↑	↑	↑

target = 10

In linear search

Avg case, Worst case = $O(n)$

Best case = $O(1)$

Binary Search

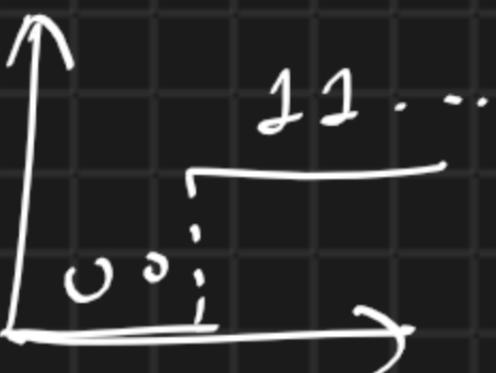
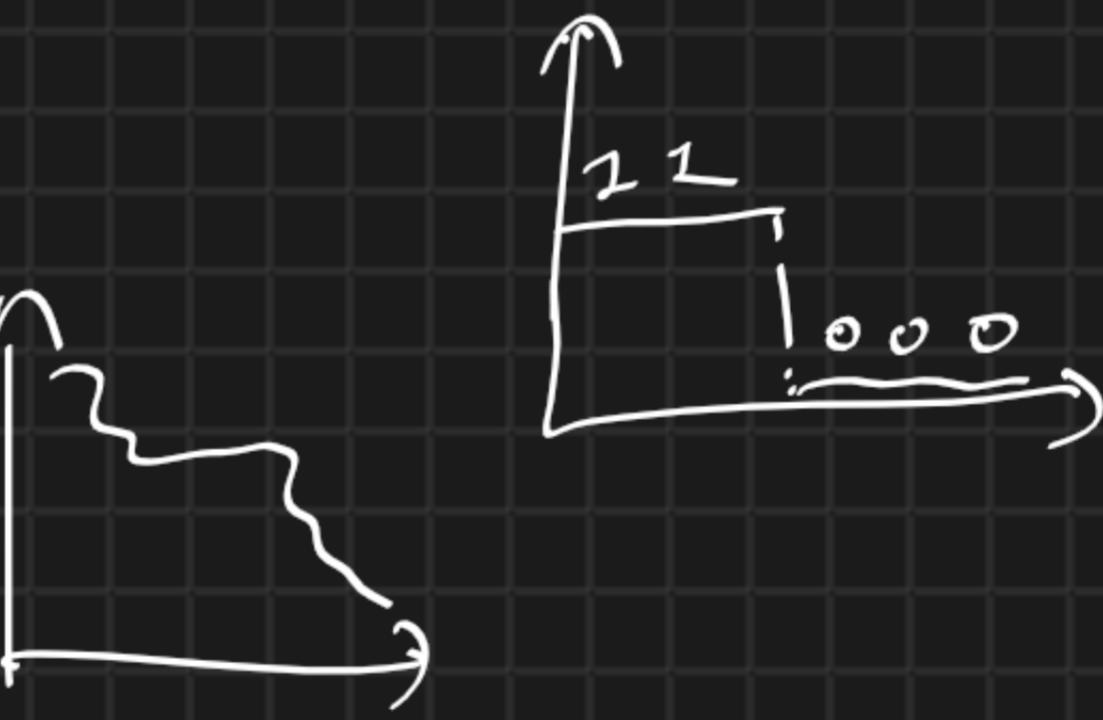
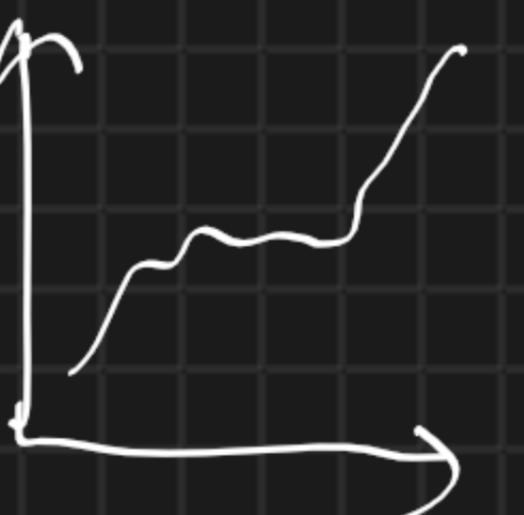
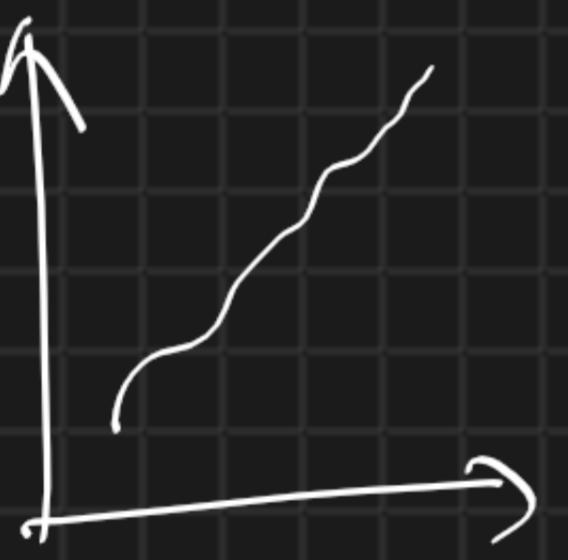
钩

Monotonic

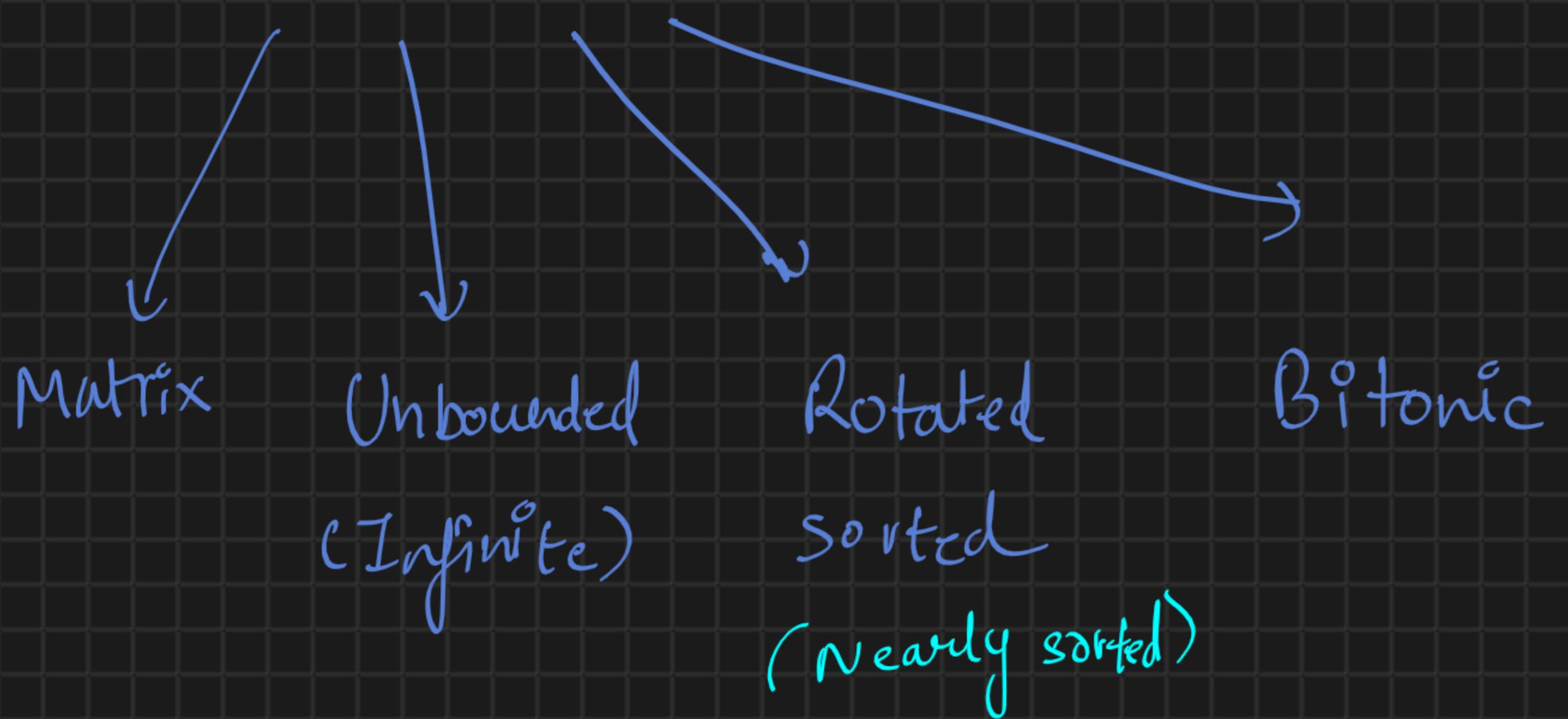
Increasing

Decreasing

Binary



Modified BS



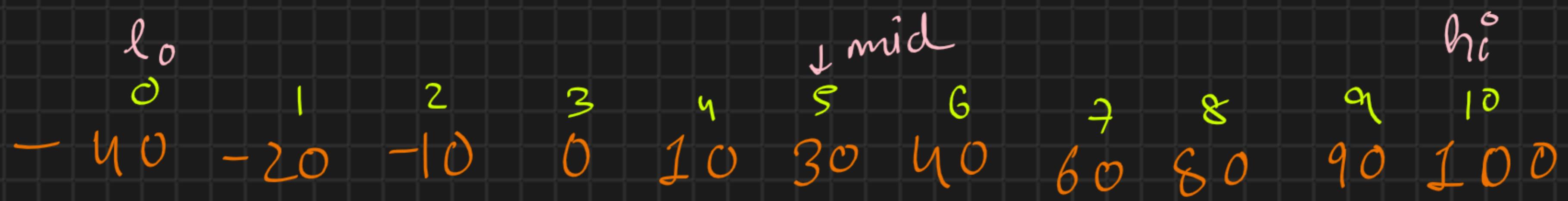
+ Binary Search on Answer

↳ Toughest questions on BS are from this category. Here, BS is not applied on the array but on the answer that we can get.

Iterative Binary Search

while ($lo \leq hi$)

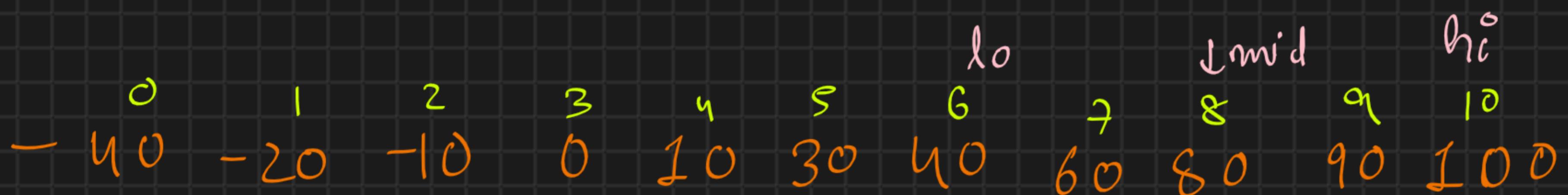
target = 60



$$mid = \frac{0+10}{2} = 5$$

if ($arr[mid] < target$)

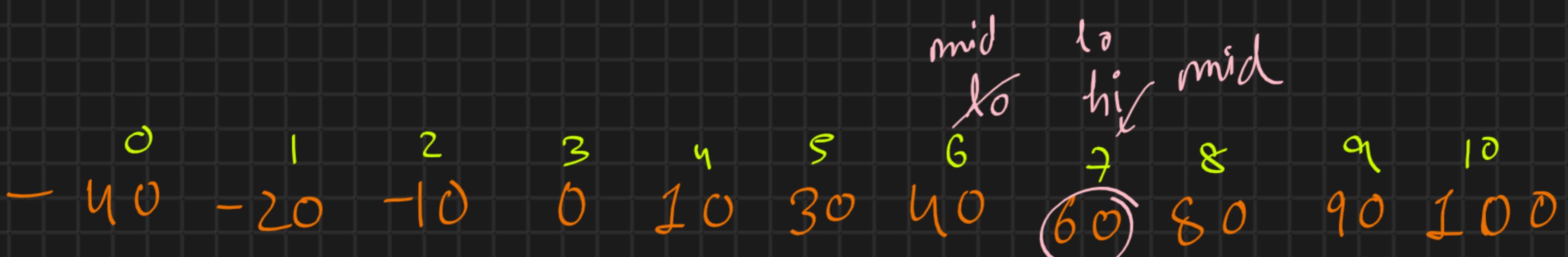
$$lo = mid + 1$$



$$mid = \frac{6+10}{2} = 8$$

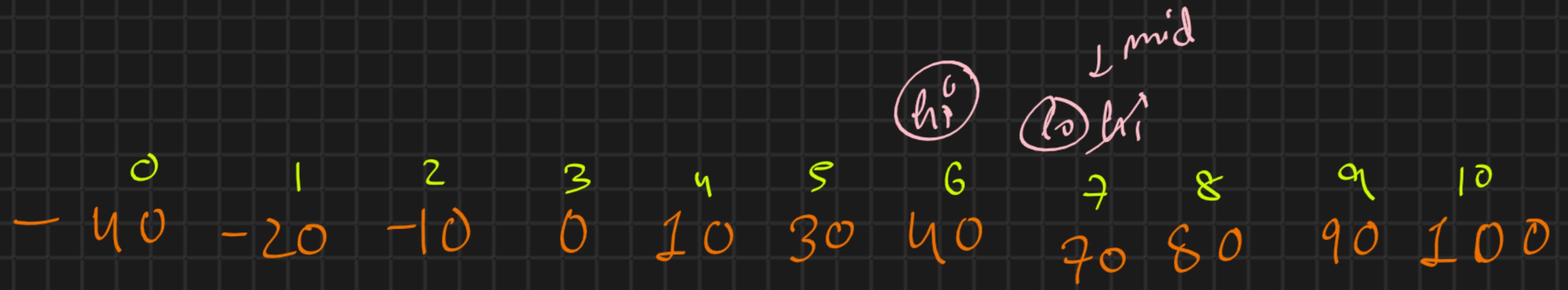
if ($arr[mid] > target$)

$$hi = mid - 1$$



$$mid = \frac{6+7}{2} = \frac{13}{2} = 6$$

Successful search



Unsuccessful Search

```

class Solution {
    public int search(int[] nums, int target) {
        int lo = 0;
        int hi = nums.length -1;

        while(lo <= hi) {
            int mid = (lo + hi) /2;

            if(nums[mid] == target) return mid;
            if(nums[mid] < target) {
                lo = mid + 1;
            } else {
                hi = mid - 1;
            }
        }

        return -1;
    }
}

```

$$TC = O(\log_2 N)$$

Int. Max-Value Int. Max-Value

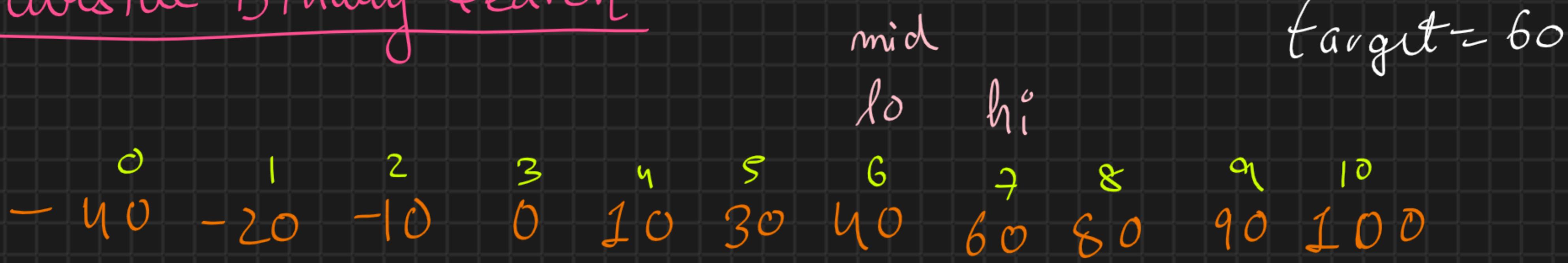
$$\text{mid} = \frac{lo + hi}{2}$$

↙ 3 out of range

$$\boxed{\text{mid} = \frac{lo + (hi - lo)}{2}}$$

↓
To avoid overflow

Recursive Binary Search



$BS(\text{arr}, 0, 10, \text{target})$

mid = 5

$$TC = (\text{calls})^h + (\text{pre} + \text{post}) * h$$

$$TC = (1)^{\log_2 N} + (K) * \log_2 N$$

$TC = O(\log_2 N)$

$BS(\text{arr}, 6, 10, \text{target})$

mid = 8

mid-1

$BS(\text{arr}, 6, 7, \text{target})$

mid = 6



$BS(\text{arr}, 7, 7, \text{target})$

mid = 7

return 7

R Relation

$$T(n) = T(n/2) + K$$

$$T(n) = \cancel{T(n/2)} + K$$

$$T(n/2) = \cancel{T(n/4)} + K$$

$$T(n/4) = \cancel{T(n/8)} + K$$

⋮
⋮

$$T(n/2^{x-1}) = \cancel{T(n/2^x)} + K$$

} x times

$$T(n) = T(n/2^x) + \underset{\text{ignore constant}}{\cancel{xK}}$$

$$\text{for } n=1 \quad T(1)=1 \rightarrow T(n=1)$$

$$T\left(\frac{n}{2^x}\right) = T(1)$$

$$\Rightarrow \frac{n}{2^x} = 1$$

$$x = \log_2 n$$

$$T(N) = T(n/2^x) + x * k$$

$$x = \log_2 N$$

$$T(N) = T(N/2^{\log_2 N}) + \log_2 N * k$$

$$T(N) = \underbrace{T(1)}_{\text{constant}} + k \log_2 N$$

$$T(N) = k \log_2 N \propto O(\log_2 N)$$

```
public int binarySearch(int[] nums, int lo, int hi, int target) {  
    if(lo > hi) return -1;  
    int mid = lo + (hi-lo)/2;  
    if(nums[mid] == target) return mid;  
    else if(nums[mid] < target) return binarySearch(nums,mid + 1,hi,target);  
    return binarySearch(nums,lo,mid -1,target);  
}  
  
public int search(int[] nums, int target) {  
    return binarySearch(nums,0,nums.length-1,target);  
}
```

Transition Point

Find Transition Point

Easy Accuracy: 53.15% Submissions: 100k+ Points: 2

Given a sorted array containing only 0s and 1s, find the transition point.

Example 1:

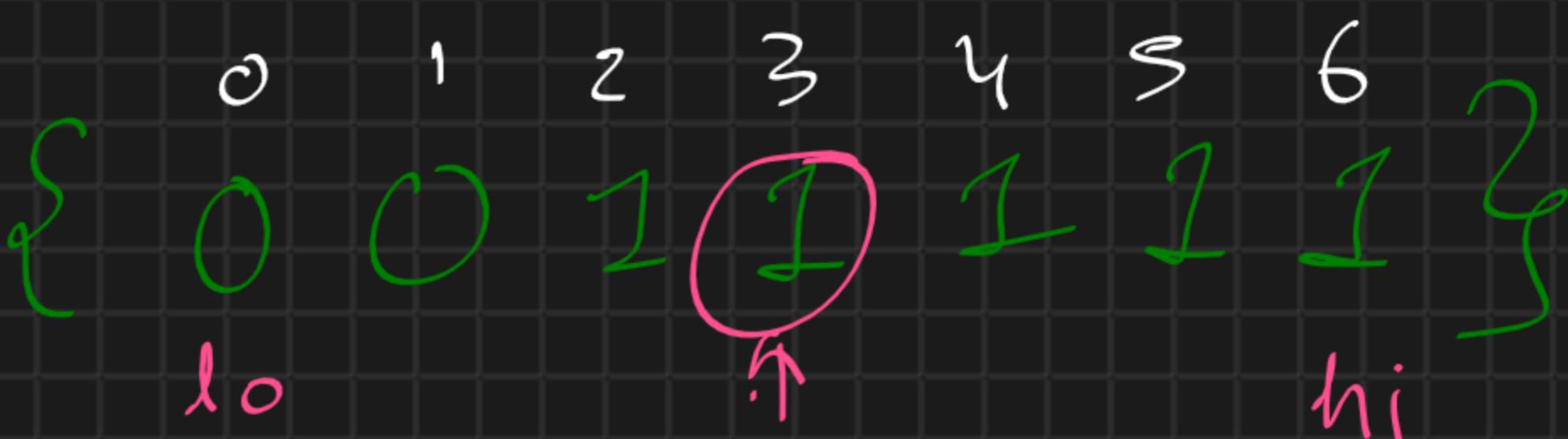
```

Input:
N = 5
arr[] = {0,0,0,1,1}
Output: 3
Explanation: index 3 is the transition
point where 1 begins.

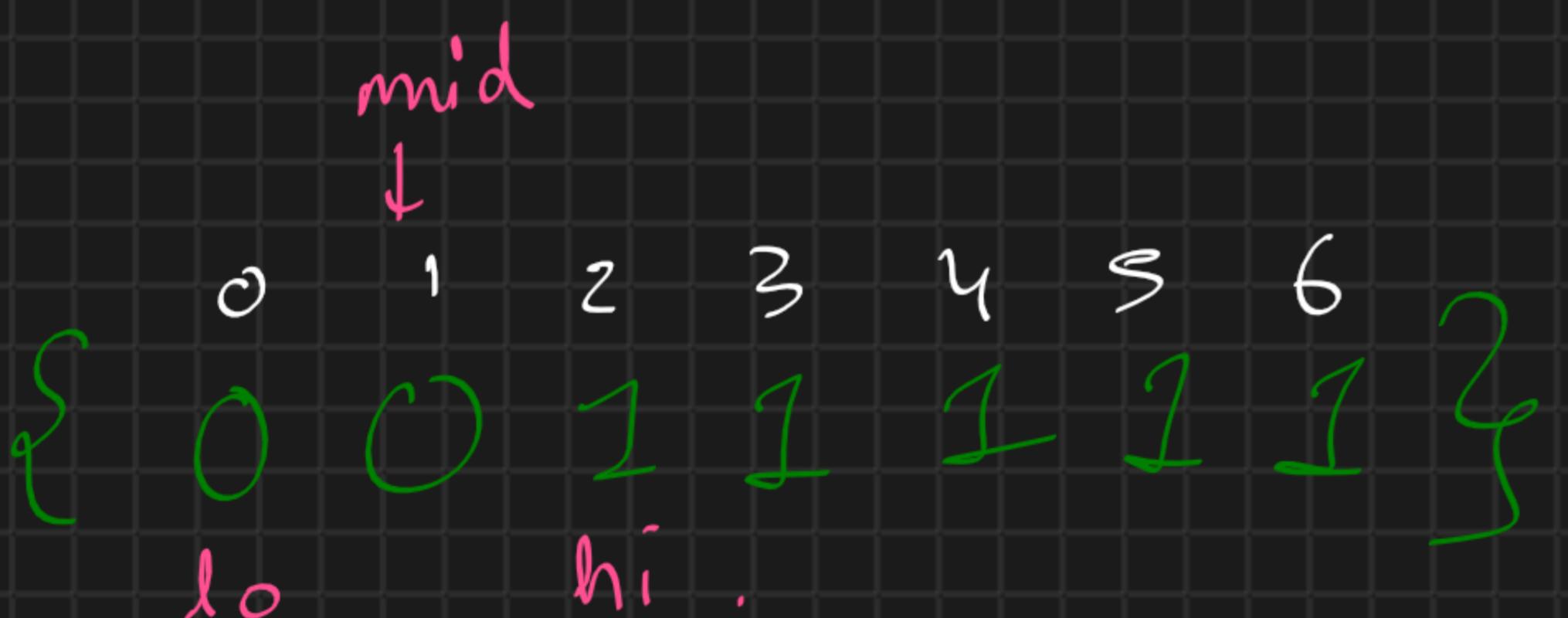
```

Transition Point : first occ of 1

$$pa = -1$$



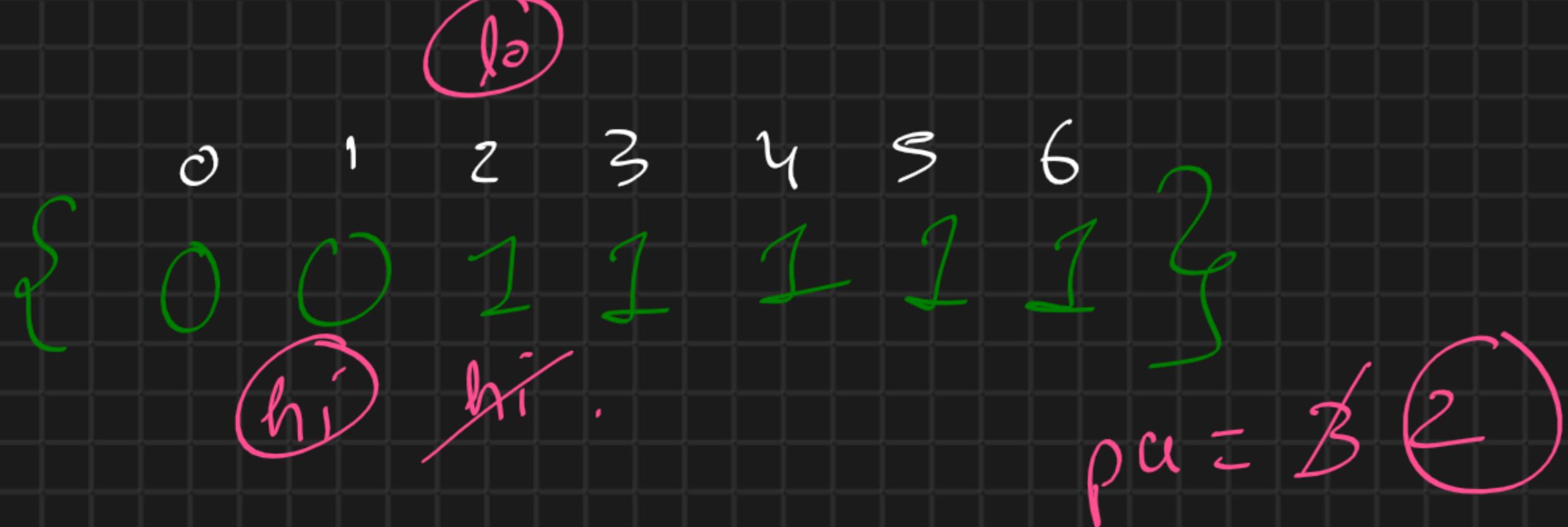
$$pa = -1$$



if ($arr[i] = -1$)
 $pa = i;$
 $hi = mid - 1;$

else

$$lo = mid + 1$$



$$pa = 3$$

```
class GfG {  
    int transitionPoint(int arr[], int n) {  
        // code here  
        int lo = 0;  
        int hi = n-1;  
  
        int pa = -1;  
        while(lo <= hi) {  
            int mid = lo + (hi-lo)/2;  
  
            if(arr[mid] == 1) {  
                pa = mid;  
                hi = mid-1;  
            } else {  
                lo = mid + 1;  
            }  
        }  
  
        return pa;  
    }  
}
```

) find Transition Point (GfG)
 $T.C = O(\log_2 n)$
 $SC = O(1)$

first Bad Version (Asked in Google)

278. First Bad Version

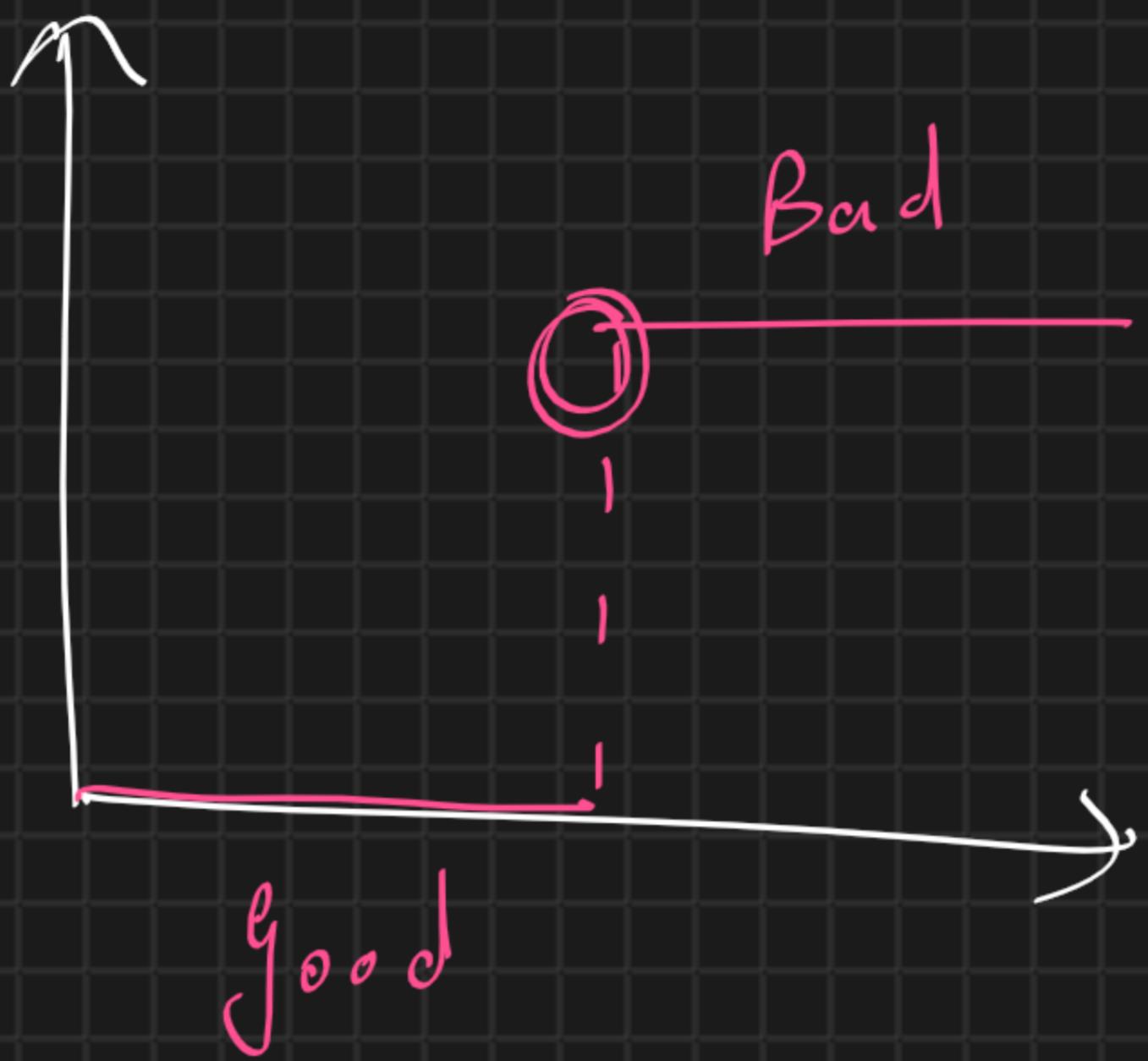
Easy 5558 2092 Add to List Share

You are a product manager and currently leading a team to develop a new product. Unfortunately, the latest version of your product fails the quality check. Since each version is developed based on the previous version, all the versions after a bad version are also bad.

Suppose you have n versions $[1, 2, \dots, n]$ and you want to find out the first bad one, which causes all the following ones to be bad.

You are given an API `bool isBadVersion(version)` which returns whether `version` is bad. Implement a function to find the first bad version. You should minimize the number of calls to the API.

```
public class Solution extends VersionControl {  
    public int firstBadVersion(int n) {  
        int lo = 1;  
        int hi = n;  
  
        int pa = -1;  
        while(lo <= hi) {  
            int mid = lo + (hi-lo)/2;  
  
            if(isBadVersion(mid) == true) {  
                pa = mid;  
                hi = mid -1;  
            } else {  
                lo = mid + 1;  
            }  
        }  
  
        return pa;  
    }  
}
```



Same as
Transition
Point

$$TC = O(\log_2 N)$$

$$SC = O(1)$$

374. Guess Number Higher or Lower

Easy 1329 212 Add to List Share

We are playing the Guess Game. The game is as follows:

I pick a number from `1` to `n`. You have to guess which number I picked.

Every time you guess wrong, I will tell you whether the number I picked is higher or lower than your guess.

You call a pre-defined API `int guess(int num)`, which returns three possible results:

- `-1`: Your guess is higher than the number I picked (i.e. `num > pick`).
- `1`: Your guess is lower than the number I picked (i.e. `num < pick`).
- `0`: your guess is equal to the number I picked (i.e. `num == pick`).

Return *the number that I picked*.

```
public class Solution extends GuessGame {  
    public int guessNumber(int n) {  
        int lo = 1;  
        int hi = n;  
  
        while(lo <= hi) {  
            int mid = lo +(hi-lo)/2;  
  
            if(guess(mid) == 0) {  
                return mid;  
            } else if(guess(mid) == -1) {  
                hi = mid -1;  
            } else {  
                lo = mid + 1;  
            }  
  
        }  
  
        return -1; //this statement will never execute  
    }  
}
```

target

$mid \leftarrow \frac{lo+hi}{2}$

$guess(mid) = -1$

$\hookrightarrow mid > target$

$hi = mid - 1$

$guess(mid) = 1$

$\hookrightarrow mid < target$

$lo = mid + 1$

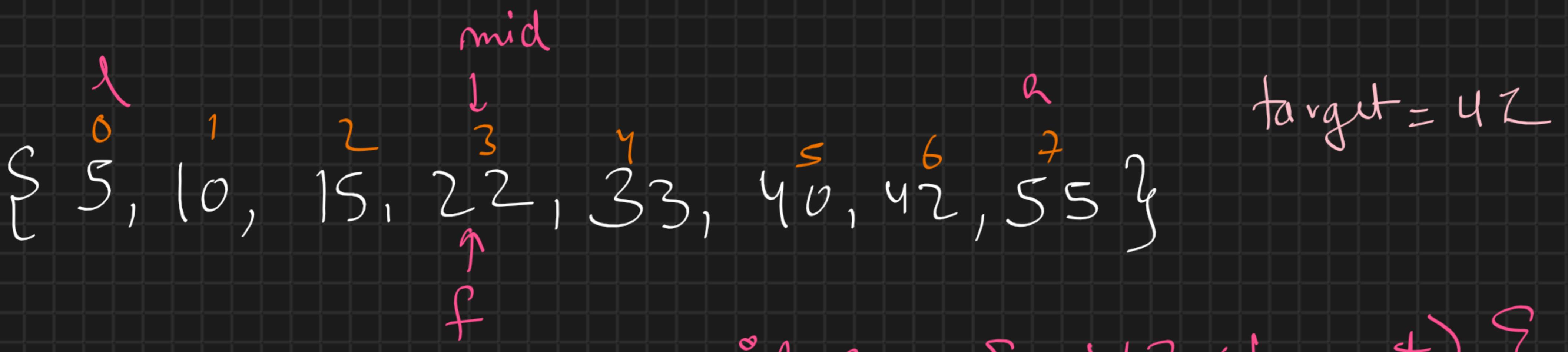
$Tc = O(\log_2 N)$

$Sc = O(1)$

Broken Economy {ceil and floor}

ceil → Bado me Sabse Chota (Min) (Initial value = ∞)

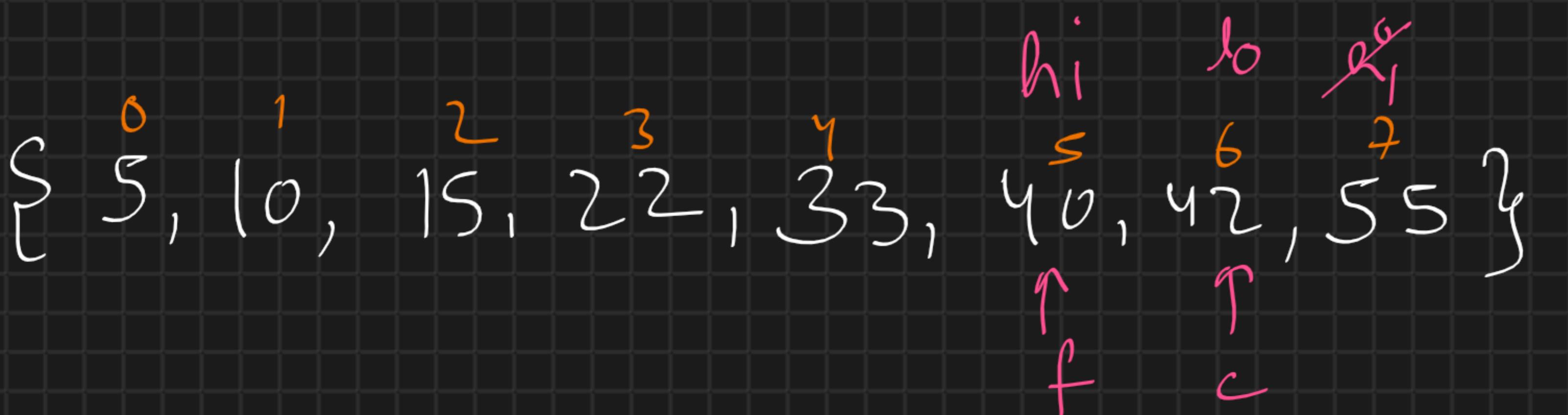
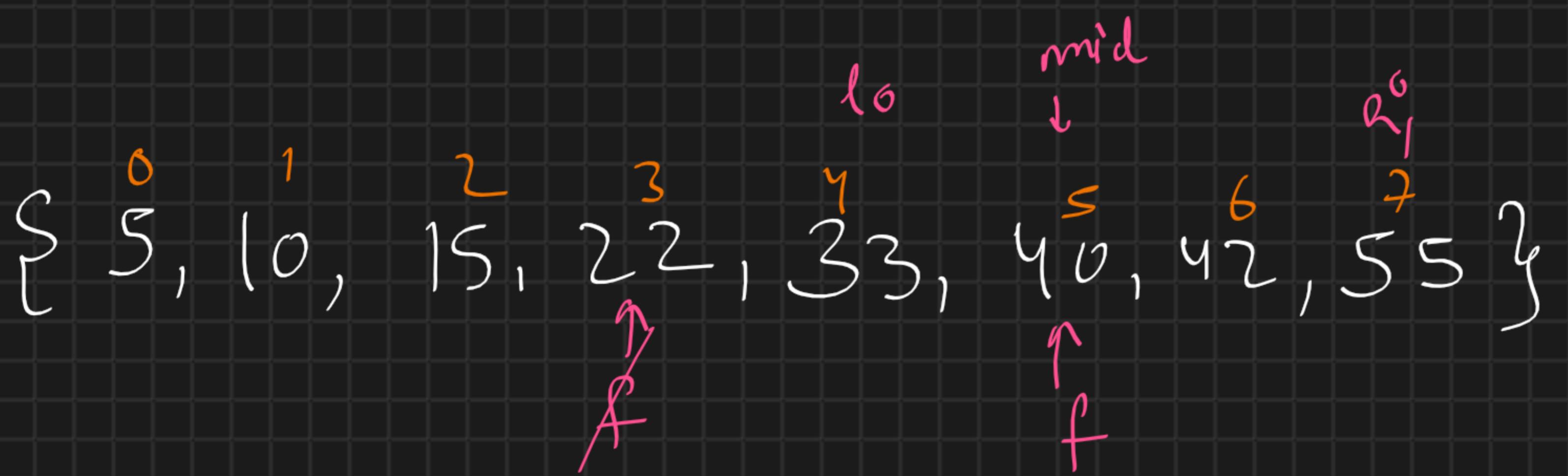
floor → Choto me Sabse Bada (Max) (Initial value = $-\infty$)



```

if (arr[mid] < target) {
    floor = arr[mid];
    lo = mid + 1;
}

```



```

if (arr[mid] > target) {
    ceil = arr[mid];
    hi = mid - 1;
}

```

$\{5, 10, 15, 22, 33, 40, 42, 55\}$

hi lo
↑ ↑
f c

floor = curv [hi]

ceil = curv [lo]

```
public static void solve(int[] arr,int target) {  
    int ceil = Integer.MAX_VALUE;  
    int floor = Integer.MIN_VALUE;  
  
    int lo = 0;  
    int hi = arr.length-1;  
  
    while(lo <= hi) {  
        int mid = lo + (hi-lo)/2;  
  
        if(arr[mid] == target) {  
            System.out.println(arr[mid]);  
            return;  
        } else if(arr[mid] < target) {  
            if(arr[mid] > floor) {  
                floor = mid;  
            }  
            lo = mid + 1;  
        } else {  
            if(arr[mid] < ceil) {  
                ceil = mid;  
            }  
            hi = mid -1;  
        }  
    }  
  
    //ceil  
    System.out.println(arr[lo]);  
    //floor  
    System.out.println(arr[hi]);  
}
```

$$TC \in O(\log_2 N)$$

$$SC \in O(1)$$