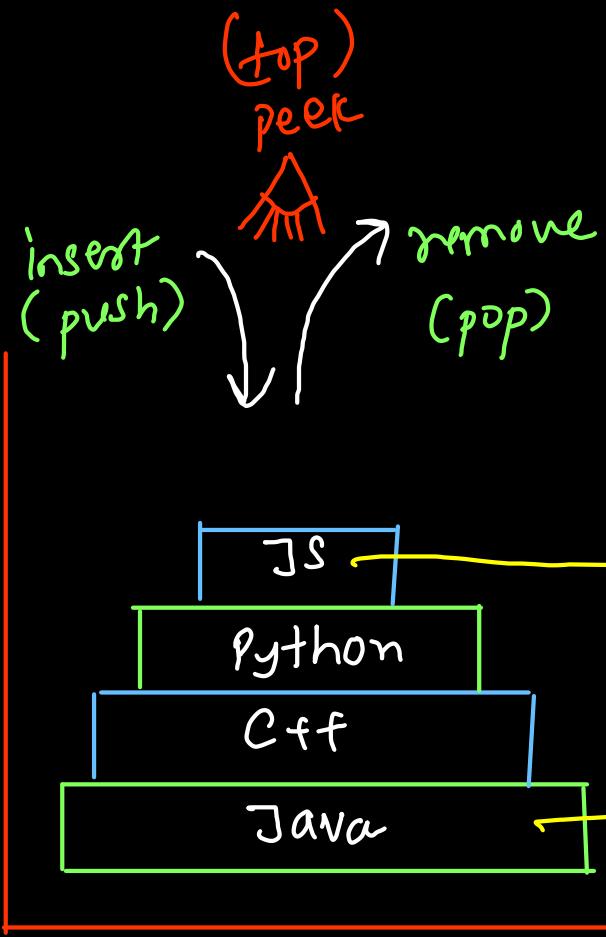


Stack



Depth First Traversal (DFS)
Programming \Rightarrow Recursion call stack
or Function call stack

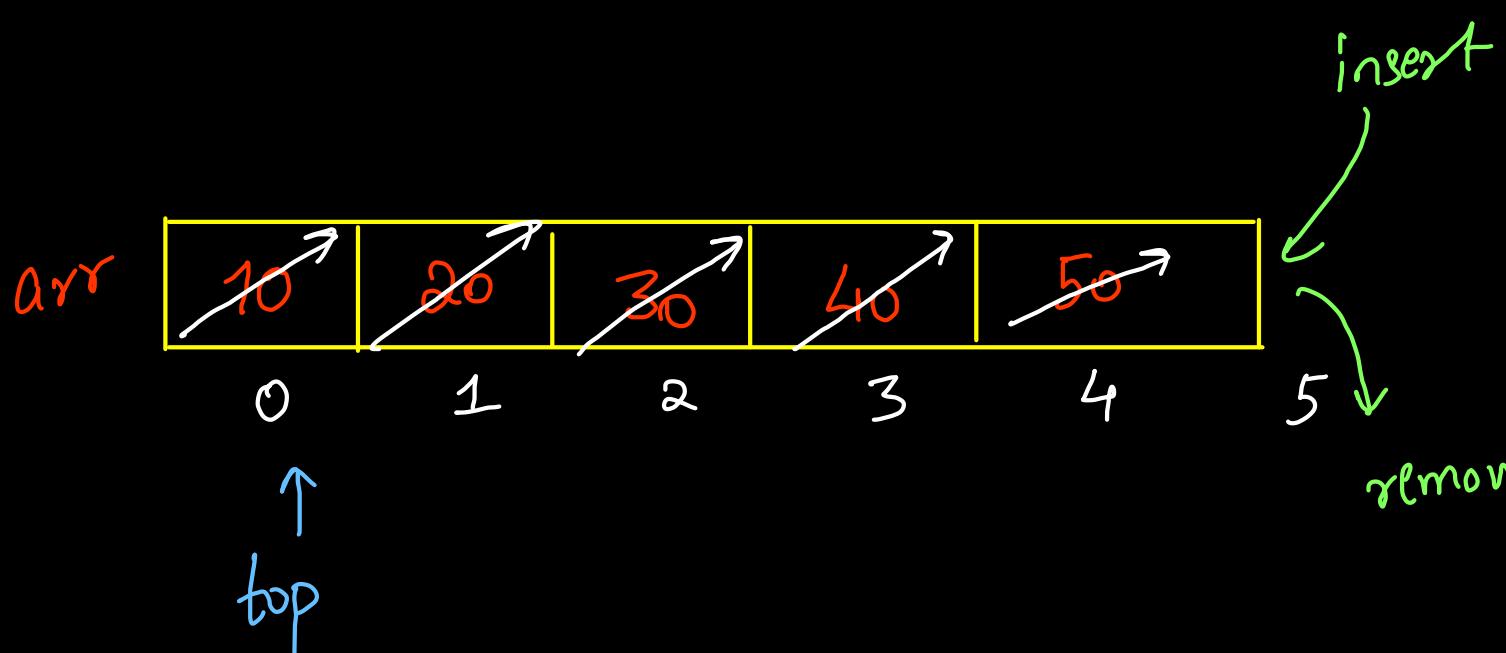
Text Editor (undo & redo)
Expressions Evaluation $\xrightarrow{\text{Operators}}$ $\xrightarrow{\text{Parens}}$

Last In/First Out { LIFO }

First In/Last out { FILO }

Java Class (Collection Framework)

```
Stack<Integer> stk = new Stack<>();  
  
stk.push(10); // O(1)  
stk.push(20);  
stk.push(30);  
stk.push(40);  
  
System.out.println(stk.peek()); // O(1)  
System.out.println(stk); // O(n)  
  
// Total Loop = O(n)  
while(stk.size() > 0){  
    System.out.println(stk.pop()); // O(1)  
    System.out.println(stk);  
}
```



push

```
arr[top] = value;
top++;

```

pop

```
top--;
return arr[top];

```

array

push → O(1) from right

pop → O(1) from right

remove
push(10)

push(20)

push(30)

push(40)

push(50)

push(60)
→ Stack Overflow

pop() = 50

pop() = 40

pop() = 30

pop() = 20

pop() = 10

pop() ⇒ Stack Underflow

Stack using Array

```
class MyStack
{
    int top = 0;
    int arr[] = new int[1000];

    void push(int value)
    {
        if(top == arr.length){
            return; // Stack Overflow
        }
        arr[top] = value;
        top++;
    }

    int pop()
    {
        if(top == 0){
            return -1; // Stack Underflow
        }
        top--;
        return arr[top];
    }
}
```

$O(1)$

Stack using ArrayList

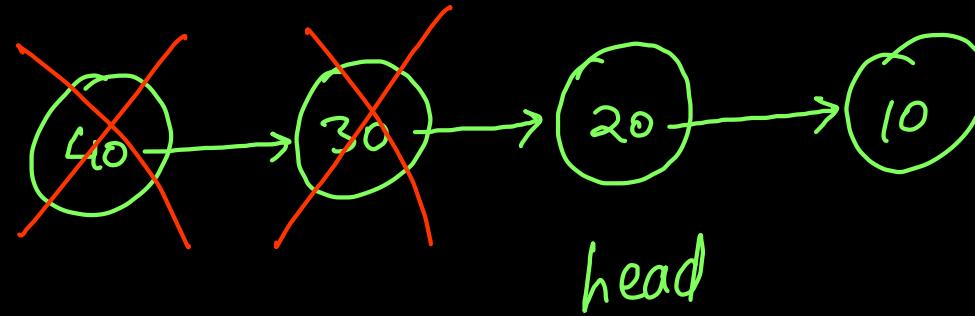
has a (composition) $O^2 - O^8$

```
class MyStack
{
    ArrayList<Integer> arr = new ArrayList<>();
    void push(int value) {
        arr.add(value); // O(1) Insertion At End
    }

    int pop()
    {
        if(arr.size() == 0){
            return -1; // Stack Underflow
        }
        return arr.remove(arr.size() - 1);
    }
}
```

$\text{add}(40)$	40 ✓
$\text{add}(30)$	30 ✓
$\text{add}(20)$	20 ✓
$\text{add}(10)$	10 ✓

Stack using hist



singly linked list with only head

① add At Head , remove At Head
 $\hookrightarrow O(1)$ $\hookrightarrow O(1)$

② add At tail , remove At tail
 $\hookrightarrow O(n)$ $\hookrightarrow O(n)$

```
class MyStack
{
    StackNode head = null;

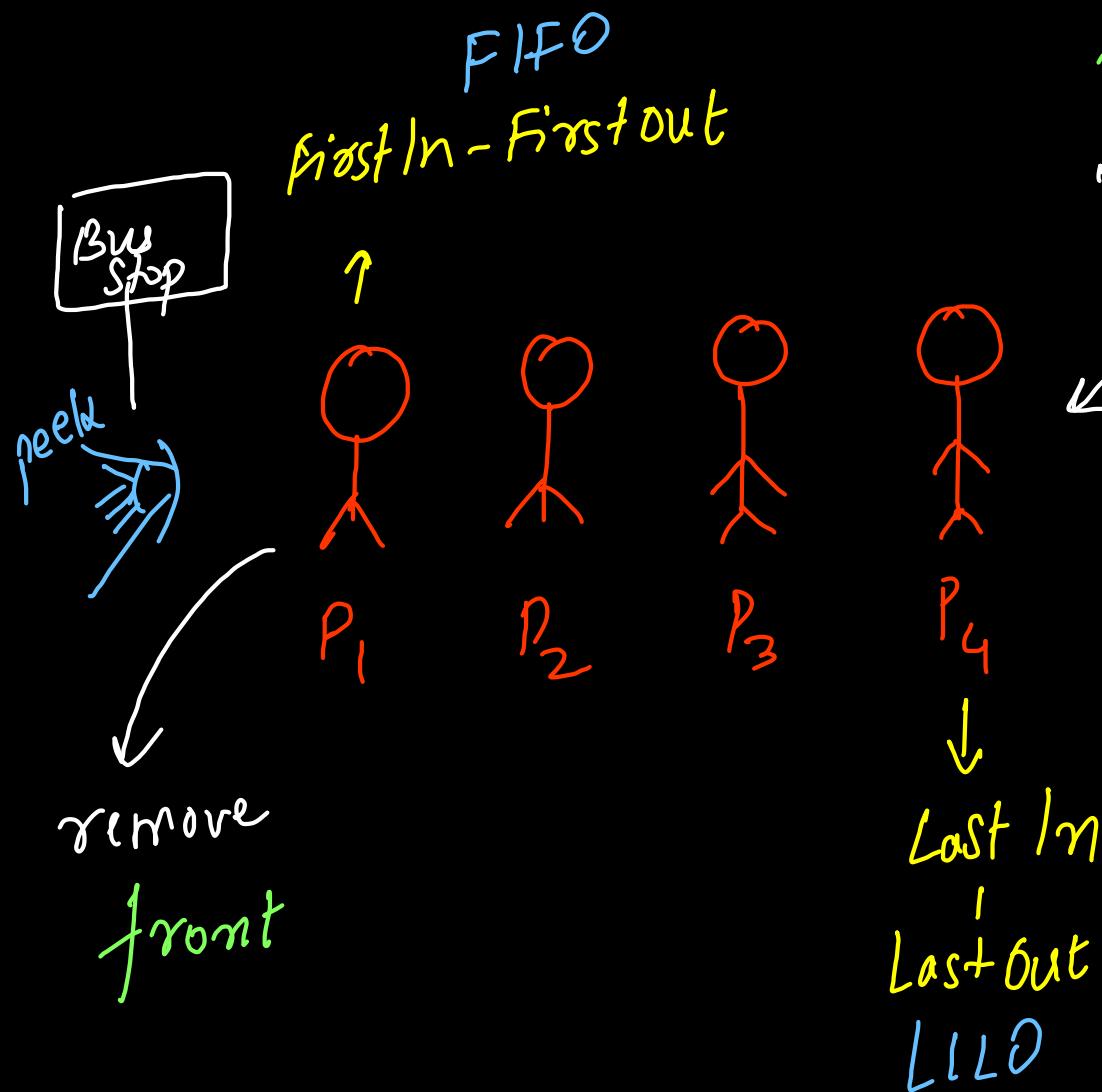
    void push(int value) //insert at head
    {
        StackNode node = new StackNode(value);
        node.next = head;
        head = node;
    }

    int pop() // remove at head
    {
        if(head == null){
            return -1; // Stack Underflow
        }
        int ans = head.data;
        head = head.next;
        return ans;
    }
}
```

$O(1)$

$O(1)$

Queue



rear
insert

Breadth First Traversal

or BFS

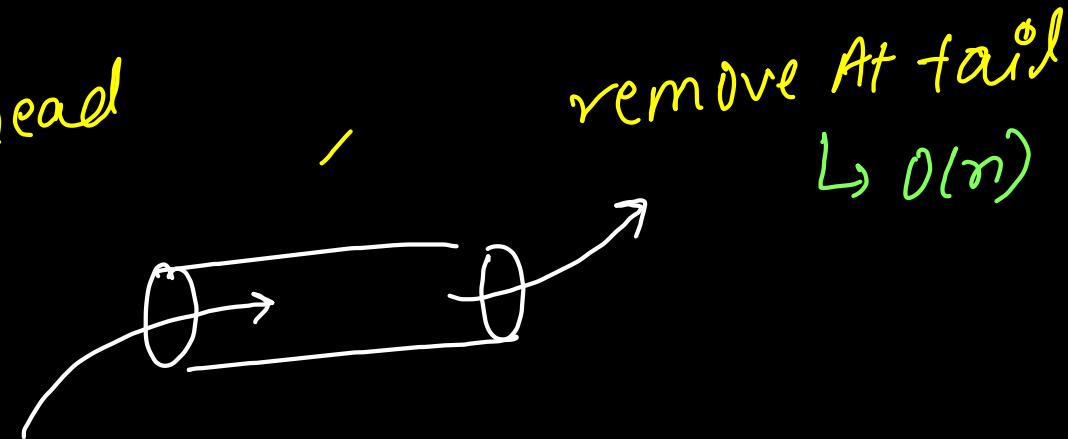
Sliding window

Binary Nos (1-N)
Generat'n

Security \Rightarrow bag {airport}

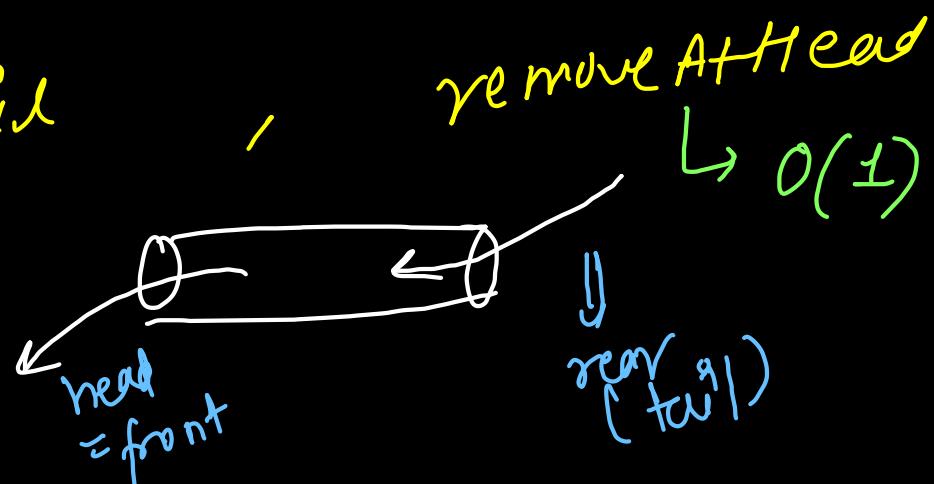
Queue using linked list
{ singly linked list with head & tail }

(1) insert At head
 $\hookrightarrow O(1)$



remove At tail
 $\hookrightarrow O(n)$

2) insert At tail
 $\hookrightarrow O(1)$



remove At Head
 $\hookrightarrow O(1)$

```

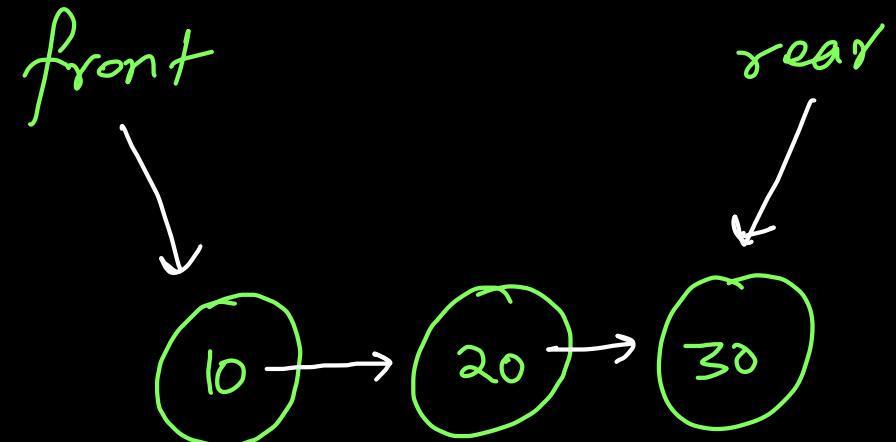
QueueNode front, rear;

//Function to push an element into the queue.
void push(int value)
{
    // Insert at Tail/Rear
    QueueNode node = new QueueNode(value);

    if(head == null){
        front = rear = node;
    } else {
        rear.next = node;
        rear = node;
    }
}

```

$O(1)$



push (10)

push (20)

push (30)

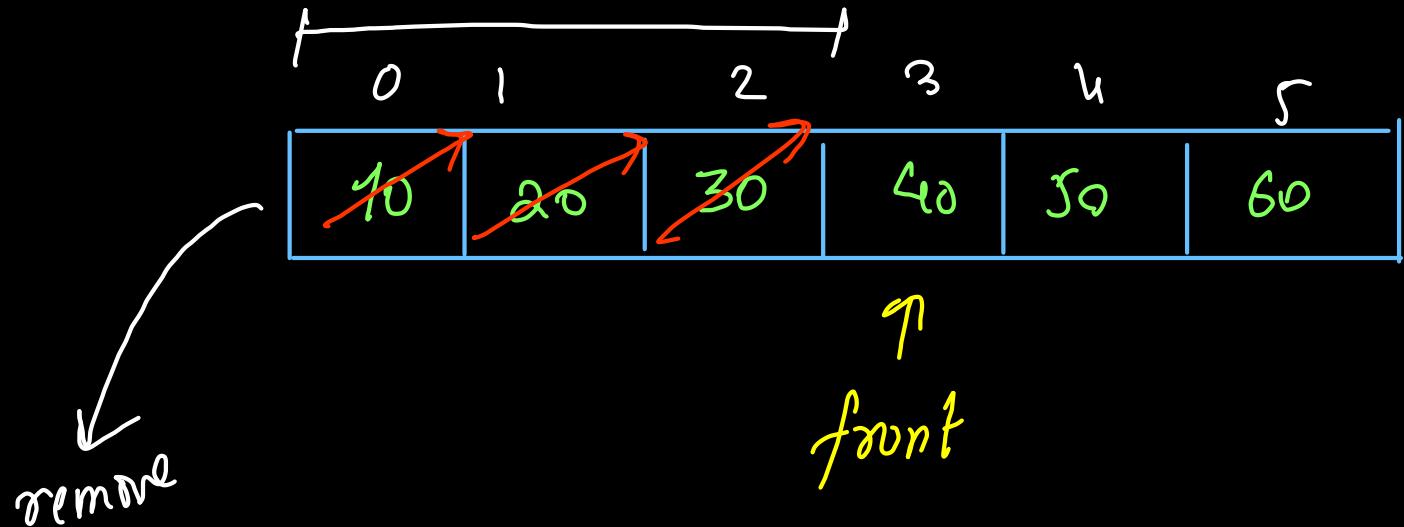
Queue using SLL

```
class MyQueue
{
    QueueNode front, rear;

    //Function to push an element into the queue.
    void push(int value) {
        // Insert at Tail/Rear
        QueueNode node = new QueueNode(value);
        if(front == null){
            front = rear = node;
        } else {
            rear.next = node;
            rear = node;
        }
    }

    int pop() {
        // Delete at Head
        if(front == null){
            return -1; // Queue Underflow
        }
        int ans = front.data;
        front = front.next;
        return ans;
    }
}
```

addr/m → not possible → wastage



push

$\text{arr}[\text{rear}] = \text{value};$

rear++

pop

$\text{ans} = \text{arr}[\text{front}]$

front

return ans;

capacity = 6
(= arr.length)

size = 3

push(10) push(50)

push(20) push(60)

push(30) push(70)

↳ Queue overflow

push(40)

remove()

remove()

remove()

Queue using normal array

```
class MyQueue {  
  
    int front = 0, rear = 0, size = 0;  
    int arr[] = new int[100005];  
  
    void push(int value)  
    {  
        if(size == arr.length) {  
            return; // Queue Overflow  
        }  
        arr[rear] = value;  
        rear++;  
        size++;  
    }  
  
    int pop()  
    {  
        if(size == 0){  
            return -1; // Queue Underflow  
        }  
        int ans = arr[front];  
        front++;  
        size--;  
        return ans;  
    }  
}
```

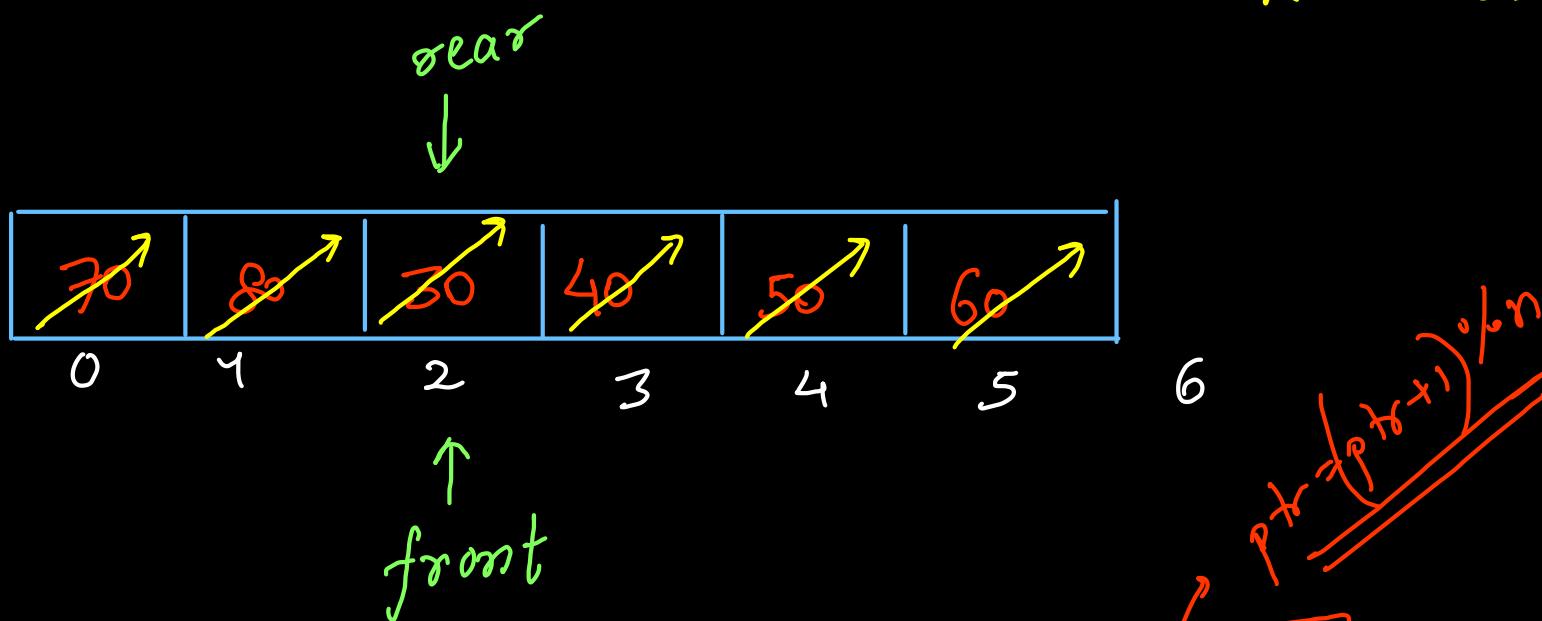
$O(1)$

$O(1)$

but wastage of space
might occur

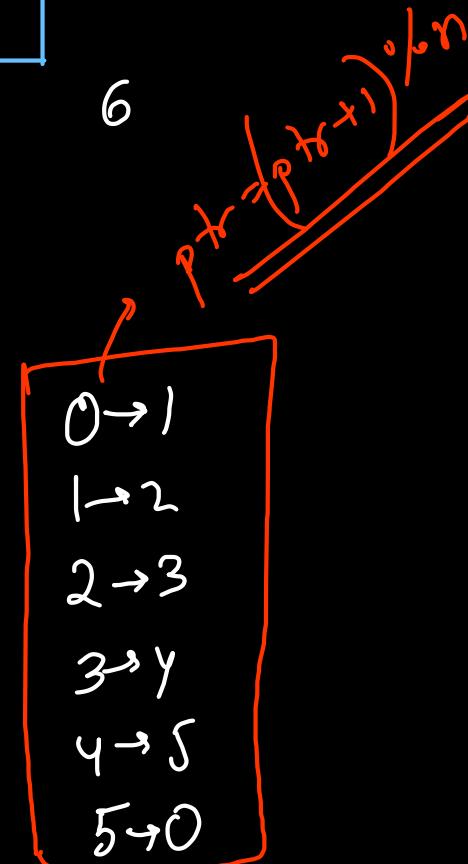
Queue using Circular Array

LC 622



$$\text{size} = 6$$

Capacity = 6



- $\text{add}(10)$ $\text{remove}()$
- $\text{add}(20)$ $\text{remove}()$
- $\text{add}(30)$ $\text{remove}()$
- $\text{add}(40)$ $\text{remove}()$
- $\text{add}(50)$ $\text{remove}()$
- $\text{add}(60)$ $\text{remove}()$
- $\text{add}(70)$ $\text{remove}()$
- $\text{add}(80)$ $\text{remove}()$
- $\text{add}(90)$ $\text{remove}()$
- Q overflow \downarrow Q underflow \downarrow

```

class MyCircularQueue {
    int[] arr;
    int front = 0, rear = 0, size = 0;

    public MyCircularQueue(int k) {
        arr = new int[k];
    }
}

```

```

public boolean enqueue(int value) {
    if(isFull()) return false; // Queue Overflow

    arr[rear] = value;
    rear = (rear + 1) % arr.length;
    size++;
    return true;
}

```

```

public boolean dequeue() {
    if(isEmpty()) return false; // Queue Underflow

    front = (front + 1) % arr.length;
    size--;
    return true;
}

```

```

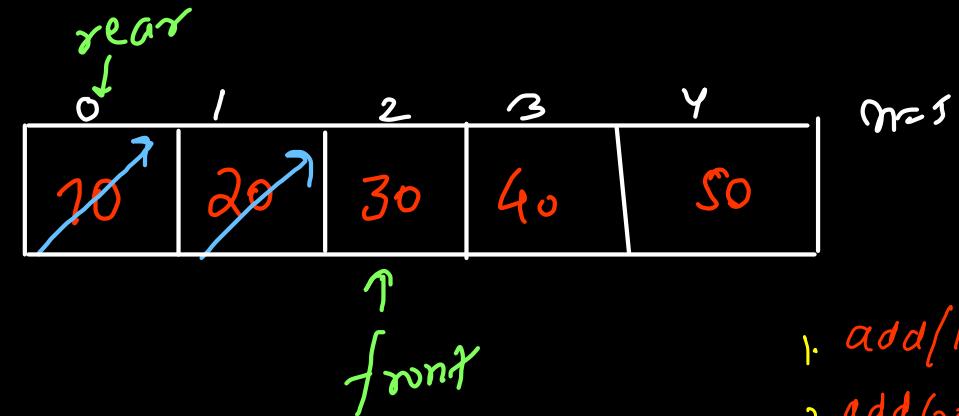
public boolean isEmpty() {
    return (size == 0); → O(1)
}

public boolean isFull() {
    return (size == arr.length); → O(1)
}

```

$4 \rightarrow 5$
 $4 \rightarrow 3$
 $3 \rightarrow 2$
 $2 \rightarrow 1$
 $1 \rightarrow 0$
 $0 \rightarrow 4$

$(\text{rear} - 1) \% n + n$



```

public int Front() {
    if(isEmpty()) return -1;
    return arr[front];
}

```

```

public int Rear() {
    if(isEmpty()) return -1;
    return arr[(rear - 1 + arr.length) % arr.length];
}

→ O(1)

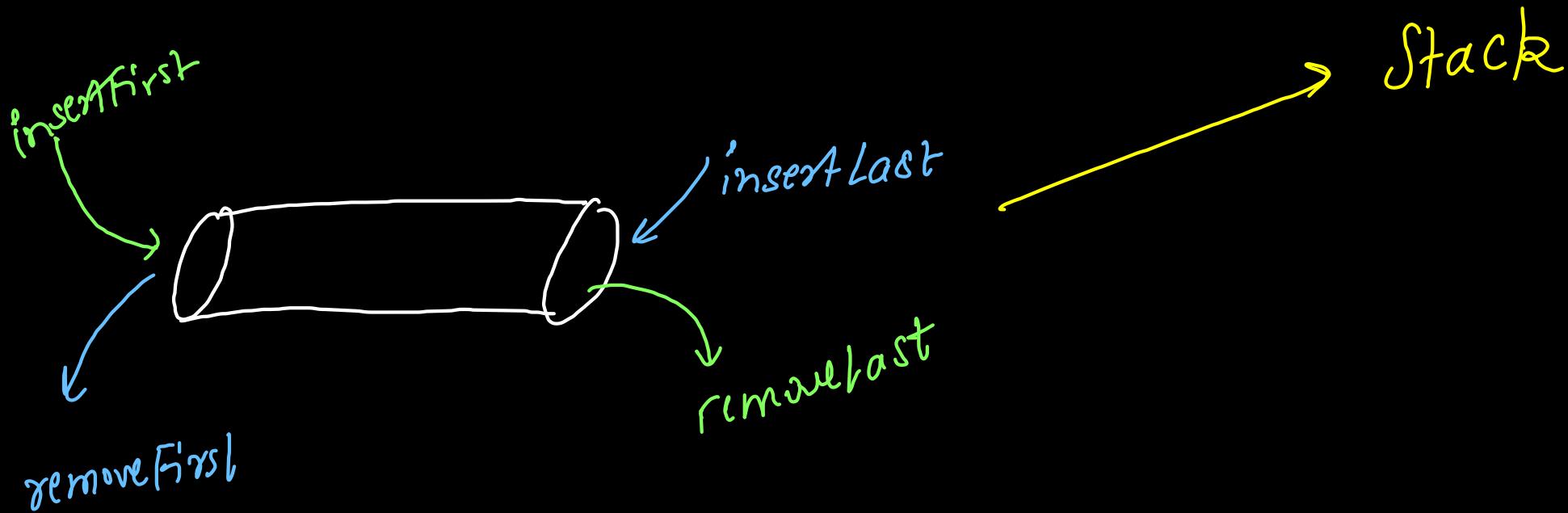
```

1. add(10)
2. add(20)
3. add(30)
4. add(40)
5. add(50)
6. remove()
7. remove()

Design Circular Deque

LeetCode 641

Doubly ended queue



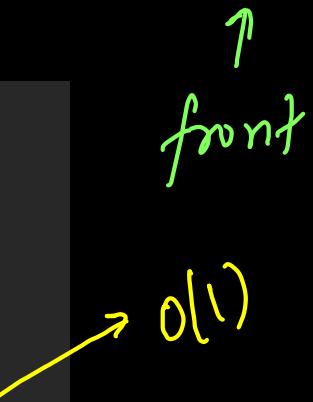


```

public boolean insertFront(int value) {
    if(isFull()) return false;

    front = (front - 1 + arr.length) % arr.length;
    arr[front] = value;
    size++;
    return true;
}

```

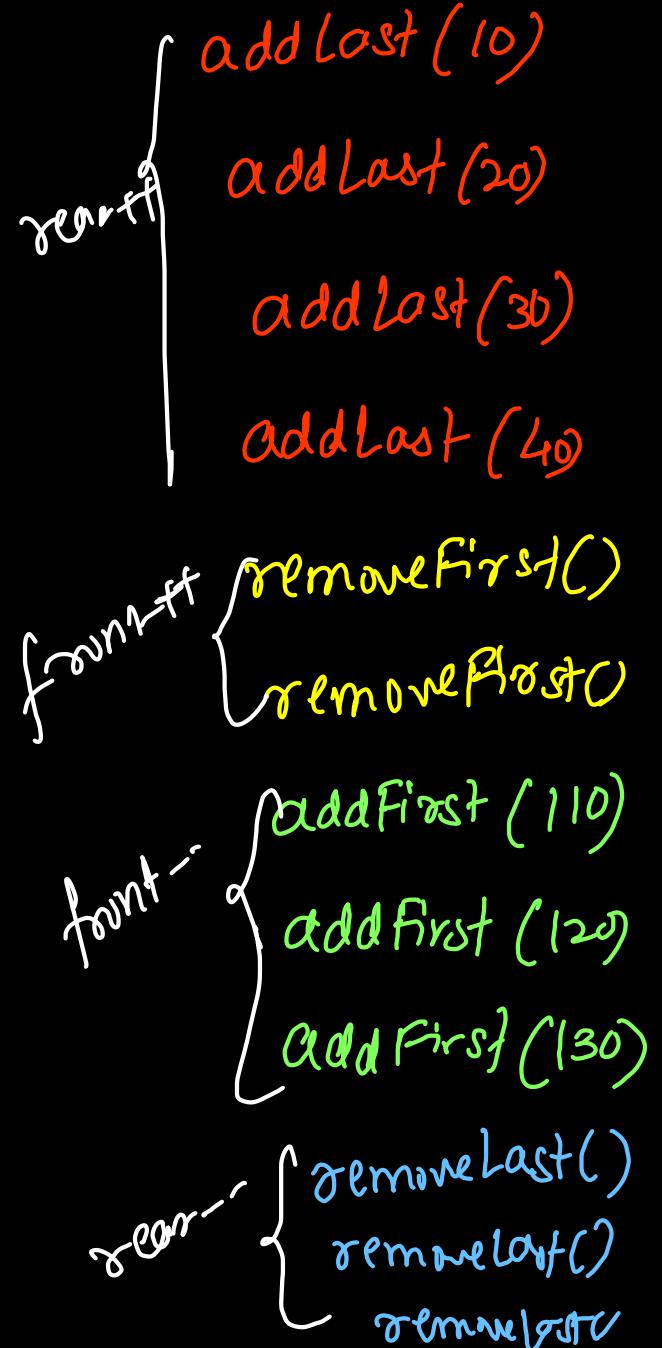


```

public boolean deleteLast() {
    if(isEmpty()) return false;

    rear = (rear - 1 + arr.length) % arr.length;
    size--;
    return true;
}

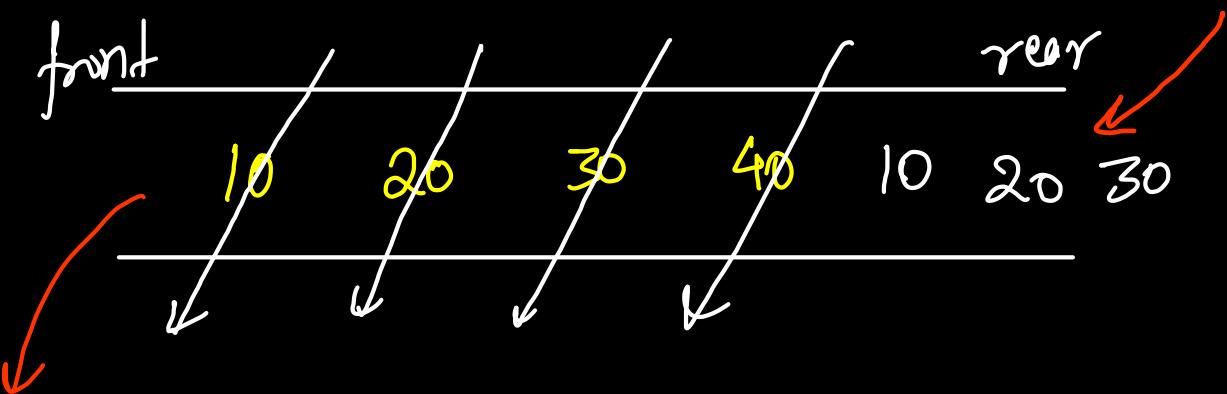
```



Implement Stack using Queue

real:

Queue
(FIFO)



logic:

Stack
(LIFO)



add(10)

add(20)

add(30)

add(40)

peek() $n \text{ pop} + n \text{ push}$

remove() $n \text{ pop}$
 $T(n-1) \text{ push}$

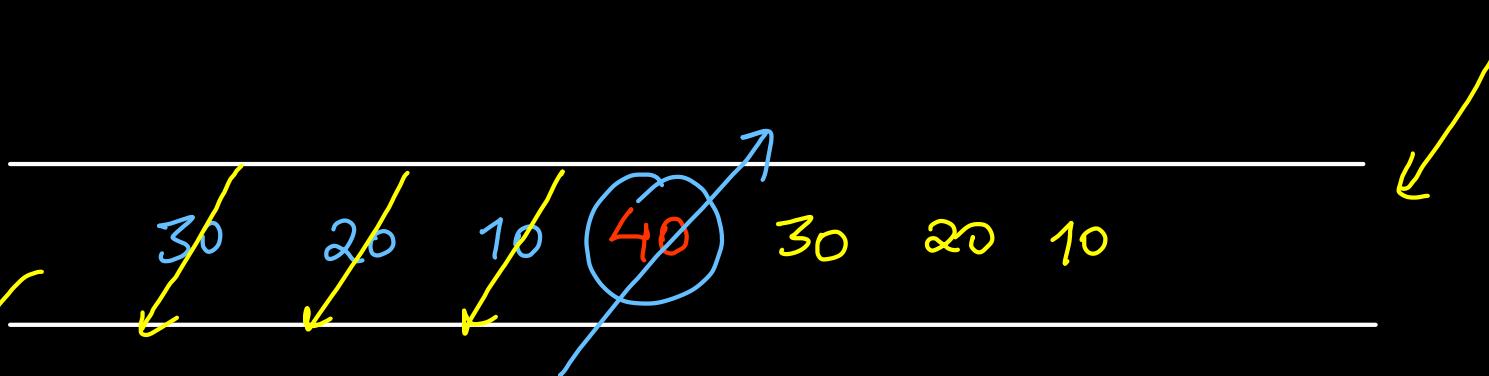
Approach 1

push Efficient $\rightarrow O(1)$

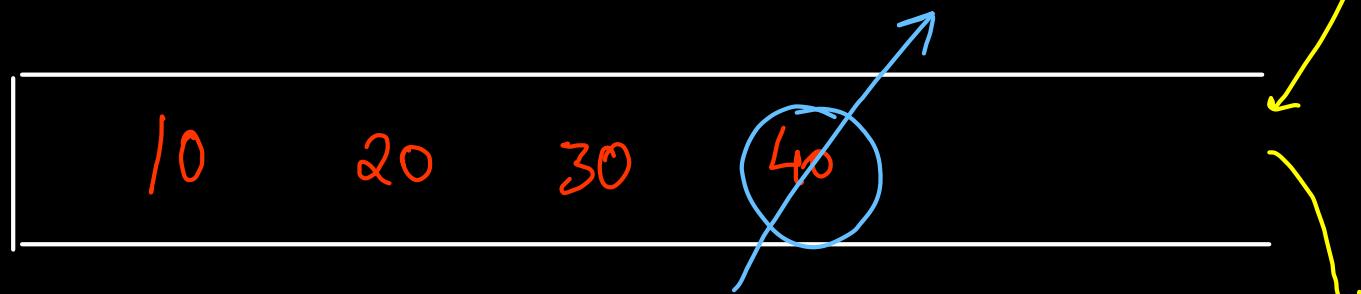
Pop Heavy $\rightarrow O(n)$

```
class MyStack {  
    Queue<Integer> q = new ArrayDeque<>();  
  
    public void push(int x) {  
        q.add(x);  
    }  
  
    public int pop() {  
        for(int idx = 1; idx < q.size(); idx++){  
            q.add(q.remove());  
        }  
        return q.remove();  
    }  
  
    public int top() {  
        for(int idx = 1; idx < q.size(); idx++){  
            q.add(q.remove());  
        }  
        int ans = q.remove();  
        q.add(ans);  
        return ans;  
    }  
  
    public boolean empty() {  
        return q.size() == 0;  
    }  
}
```

Achiral
Queue
FIFO



Logic
Stack
LIFO



$\text{add}(10)$
 $\text{add}(20)$
 $\text{add}(30)$
 $\text{add}(40)$
 peek
 remove

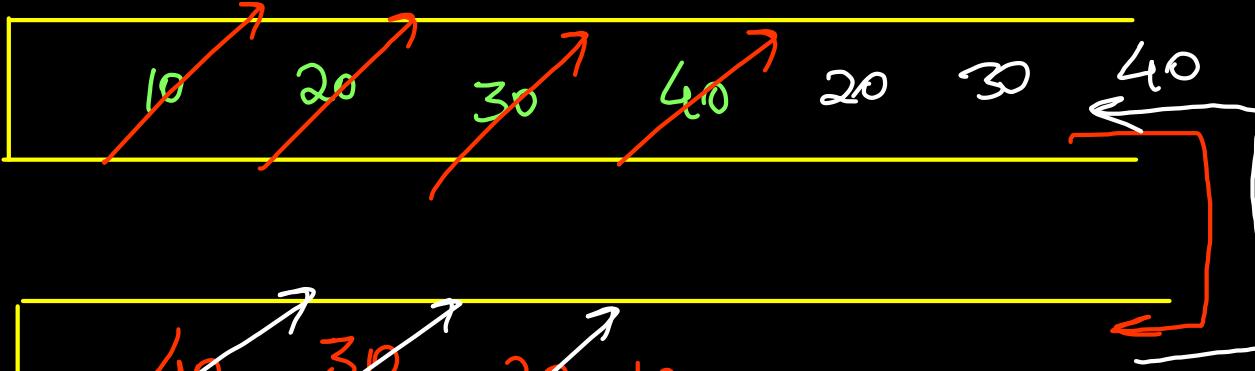
Approach 2

Push Heavy $\rightarrow O(n)$
Pop Efficient $\rightarrow O(1)$

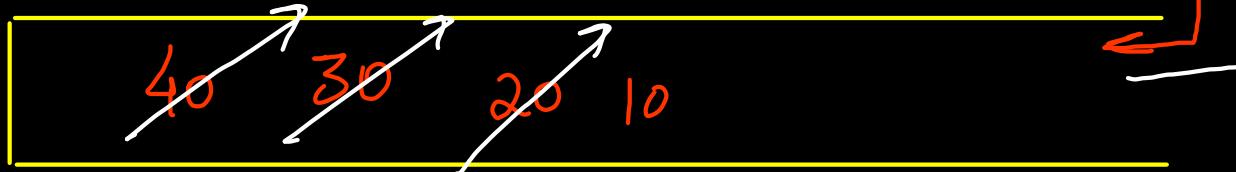
```
class MyStack {  
    Queue<Integer> q = new ArrayDeque<>();  
  
    public void push(int x) {  
        q.add(x);  
        for(int idx = 1; idx < q.size(); idx++){  
            q.add(q.remove());  
        }  
    }  
  
    public int pop() {  
        return q.remove(); → O(1)  
    }  
  
    public int top() { → O(1)  
        return q.peek();  
    }  
  
    public boolean empty() {  
        return q.size() == 0;  
    }  
}
```

Implement Queue using Stack

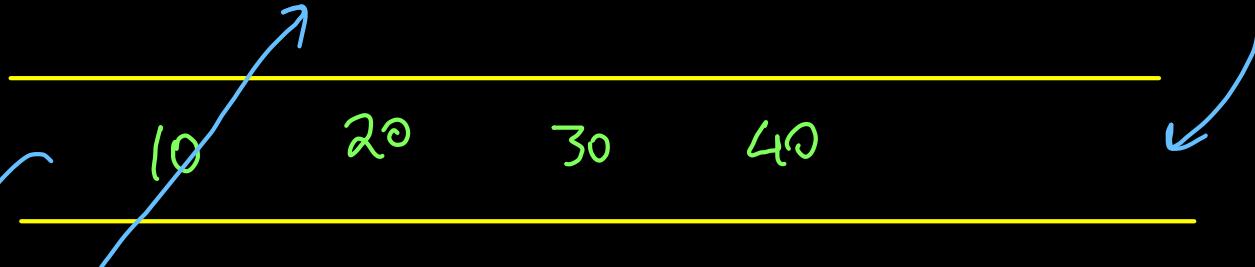
Actual
Stack
LIFO



extra
stack



Logic
Queue
FIFO



$\text{add}(10)$
 $\text{add}(20)$
 $\text{add}(30)$
 $\text{add}(40)$
 $\text{peek}()$
 $\text{remove}()$

Approach 1

push Efficient $\rightarrow O(1)$

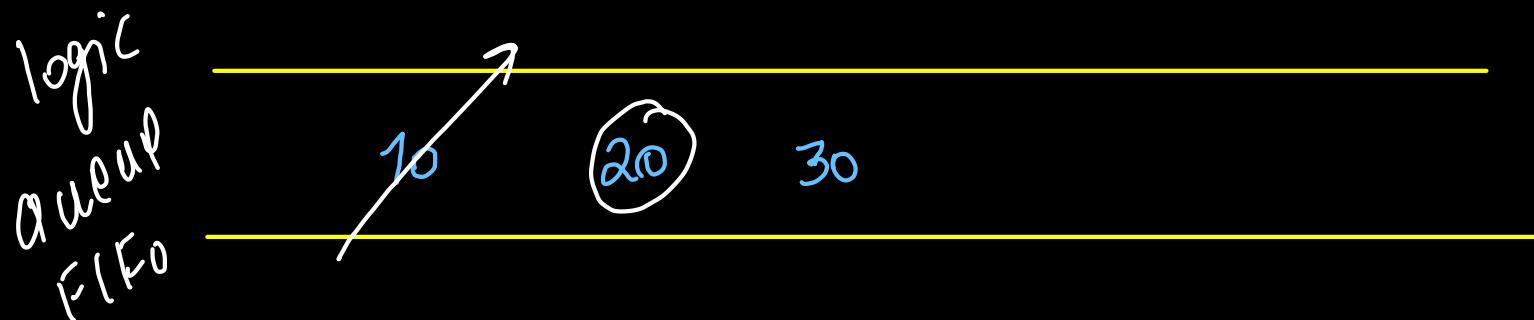
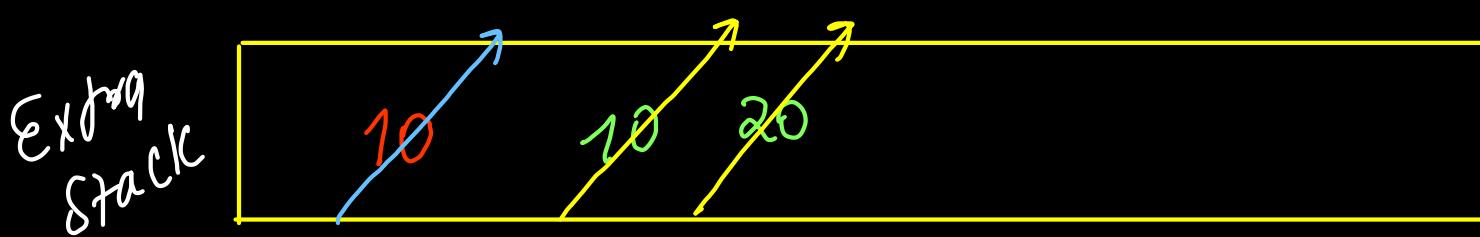
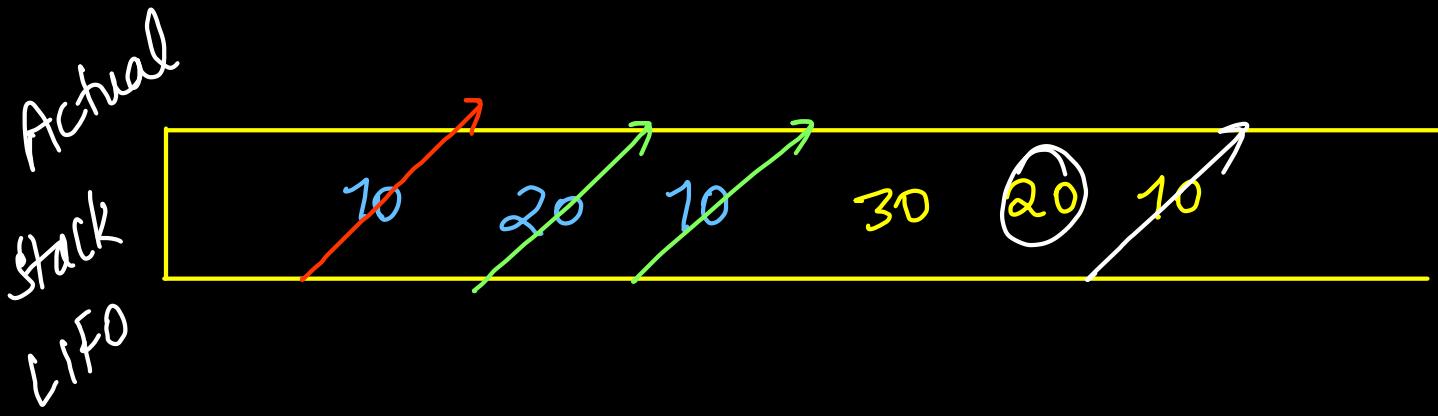
pop Heavy $\rightarrow O(n)$

```
class MyQueue {  
    Stack<Integer> actual = new Stack<>();  
    Stack<Integer> extra = new Stack<>();  
  
    public void push(int x) {  
        actual.push(x); → O(1)  
    }  
  
    public int pop() {  
        while(actual.size() > 0)  
            extra.push(actual.pop());  
        int ans = extra.pop(); → O(n)  
        while(extra.size() > 0)  
            actual.push(extra.pop());  
        return ans;  
    }  
  
    public int peek() {  
        while(actual.size() > 0)  
            extra.push(actual.pop());  
        int ans = extra.peek(); → O(n)  
        while(extra.size() > 0)  
            actual.push(extra.pop());  
        return ans;  
    }  
  
    public boolean empty() {  
        return actual.size() == 0;  
    }  
}
```

Approach 2 Push $\rightarrow O(n)$

Pop $\rightarrow O(1)$

```
class MyQueue {  
    Stack<Integer> actual = new Stack<>();  
    Stack<Integer> extra = new Stack<>();  
  
    public void push(int x) {  
        while(actual.size() > 0){  
            extra.push(actual.pop());  
        }  
        actual.push(x);  
        while(extra.size() > 0){  
            actual.push(extra.pop());  
        }  
    }  
  
    public int pop() {  
        O(1) { return actual.pop(); }  
    }  
  
    public int peek() {  
        O(1) { return actual.peek(); }  
    }  
  
    public boolean empty() {  
        O(1) { return actual.size() == 0; }  
    }  
}
```



Approach 2

Push Heavy $\rightarrow O(n)$
Pop Efficient $\rightarrow O(1)$

$\text{add}(10, 20, 30)$

Design Minimum Stack

Leetcode 155

50
30
10
20
60
40

add(40)

top() = 40 min() = 40

add(60)

top() = 60 min() = 40

add(20)

top() = 20 min() = 20

add(10)

top() = 10 min() = 10

add(30)

top() = 30 min() = 10

add(50)

top() = 50 min() = 10

```

class MinStack {
    Stack<Integer> min, stk;
    public MinStack() {
        min = new Stack<>();
        stk = new Stack<>();
        min.push(Integer.MAX_VALUE);
    }

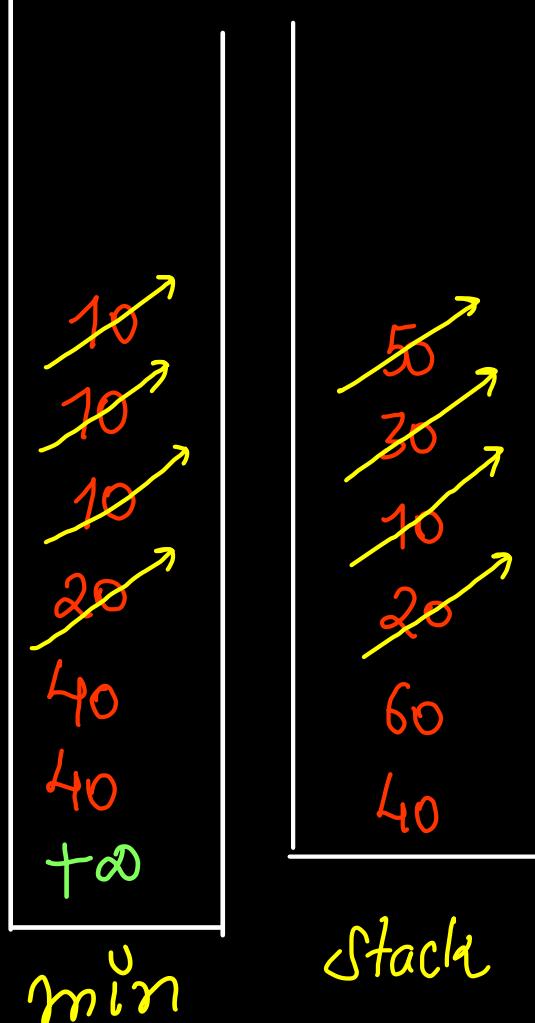
    public void push(int val) {
        stk.push(val);
        int newMin = Math.min(min.peek(), val);
        min.push(newMin);
    }

    public void pop() {
        stk.pop();
        min.pop();
    }

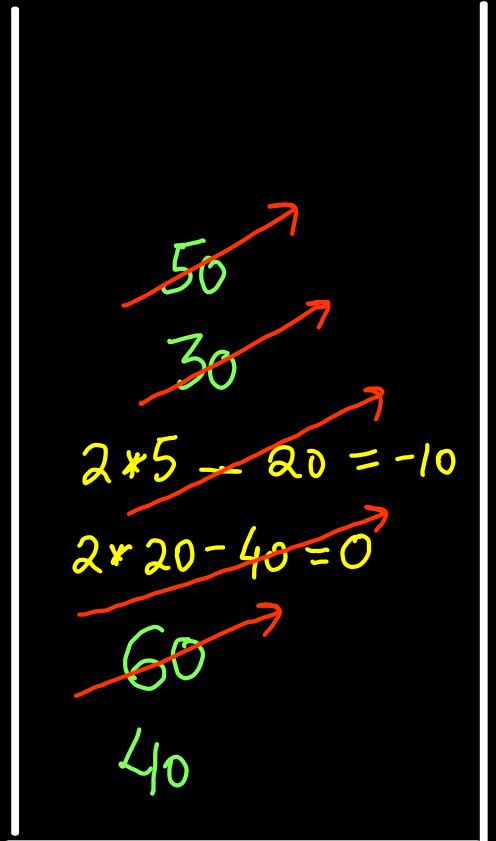
    public int top() { return stk.peek(); }
    public int getMin() { return min.peek(); }
}

```

Extra Space $\Rightarrow O(n)$
~~MinStack~~



- ✓ add(40)
 $\text{top}() = 40$ $\text{min}() = 40$
- ✓ add(60)
 $\text{top}() = 60$ $\text{min}() = 40$
- ✗ add(20)
 $\text{top}() = 20$ $\text{min}() = 20$
- ✗ add(10)
 $\text{top}() = 10$ $\text{min}() = 10$
- ✗ add(30)
 $\text{top}() = 30$ $\text{min}() = 10$
- ✗ add(50)
 $\text{top}() = 50$ $\text{min}() = 10$



Stack <Long>

Encrypted

$$\begin{aligned} \text{val} &\leftarrow \min \\ &= \text{val} + (\overbrace{\text{val} - \min}^{\text{-ve}}) \\ &= 2 * \text{val} - \min \end{aligned}$$

decryption = $2 * \min - \text{top}$

5	\downarrow	$2 * 5 - (-10)$
20	\downarrow	$2 * 20 - 0$
40	\downarrow	
\min		

- ✓ add(40)
top() = 40 min() = 40
- ✓ add(60)
top() = 60 min() = 40
- ✗ add(20)
top() = 20 min() = 20
- ✗ add(5)
top() = 5 min() = 5
- ✗ add(30)
top() = 30 min() = 10
- ✗ add(50)
top() = 50 min() = 10
- ✗ add(20)
~~add(20)~~

Extra space $\Rightarrow O(1)$

```
class MinStack {
    Stack<Long> stk = new Stack<>();
    long min = Integer.MAX_VALUE;

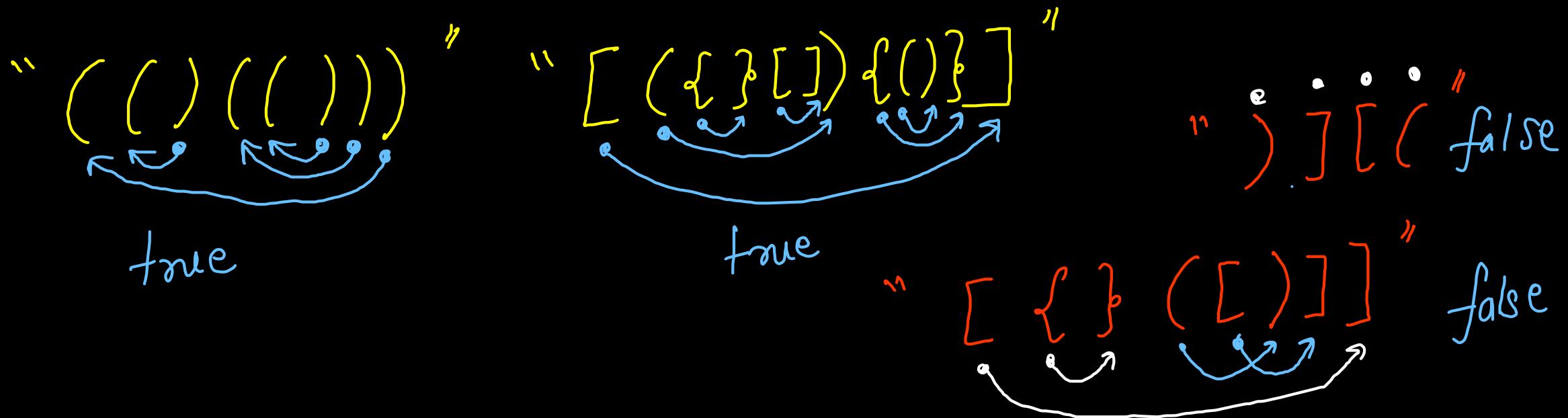
    public void push(int val) {
        if(val >= min){
            stk.push(1l * val);
        } else {
            stk.push(2l * val - min);
            min = val;
        }
    }

    public void pop() {
        long top = stk.pop();
        if(top >= min) return;
        min = 2l * min - top;
    }

    public int top() {
        long top = stk.peek();
        if(top < min)
            return (int)min;
        return (int)top;
    }

    public int getMin() {
        return (int)min;
    }
}
```

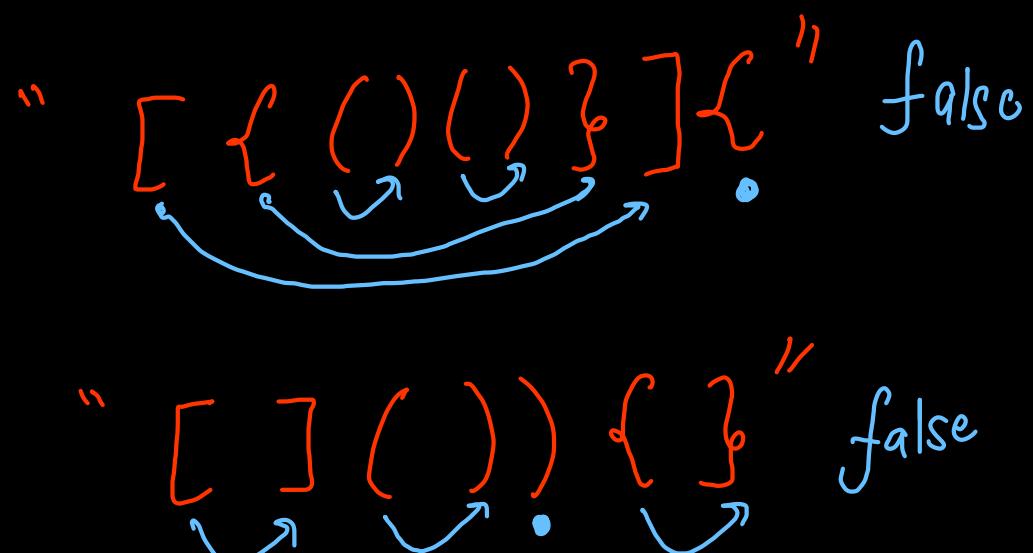
Balanced/Valid Parentheses Leetcode 20



→ open == closed

→ order matching

→ type matching



" [({ } []) { () }] " done

Diagram illustrating a balanced expression with red arrows indicating matching pairs of parentheses: [→] ← (→) ← { → } ← }

~~[{ ({ } []) { () }]]~~

Diagram illustrating an unbalanced expression where opening braces are unmatched. Red arrows point to the first three opening braces: [→ { → (→ { → [→

unbalanced opening braces

" [] ()) { } "

Diagram illustrating an unbalanced expression where closing braces are unmatched. Red arrows point to the last two closing braces:] →) →) → } → }

~~[{ ()) { }~~ stk.size() == 0 false

" [{ } ([)]] " mismatch

Diagram illustrating an unbalanced expression with a mismatched brace. A blue arrow points to a brace that does not have a corresponding closing brace.

~~[{ ([]]]~~ false

" [{ () () }] [] " stk.size() > 0 false

Diagram illustrating an unbalanced expression with an extra closing brace. A blue arrow points to a closing brace that has no corresponding opening brace.

```
class Solution {
    public boolean isValid(String s) {
        Stack<Character> stk = new Stack<>();
        // for(char ch: s.toCharArray())
        for(int idx = 0; idx < s.length(); idx++){
            char ch = s.charAt(idx);

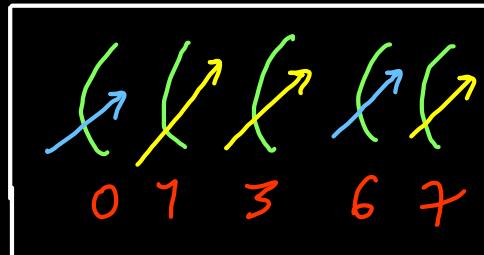
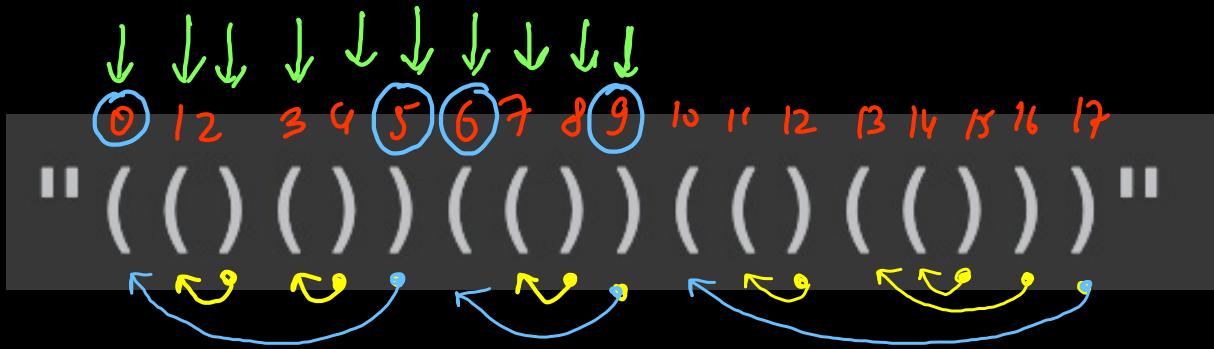
            if(ch == '(' || ch == '[' || ch == '{'){
                stk.push(ch);
            } else {
                if(stk.size() == 0) return false; // extra closing
                if(ch == ')' && stk.peek() == '(') stk.pop();
                else if(ch == '}' && stk.peek() == '{') stk.pop();
                else if(ch == ']' && stk.peek() == '[') stk.pop();
                else return false; // type or order mismatch
            }
        }

        return (stk.size() == 0); // stk.size() > 0: extra open
    }
}
```

Time $\Rightarrow O(n)$

Space $\Rightarrow O(n)$ stack

Remove Outermost Parentheses



$O(n)$

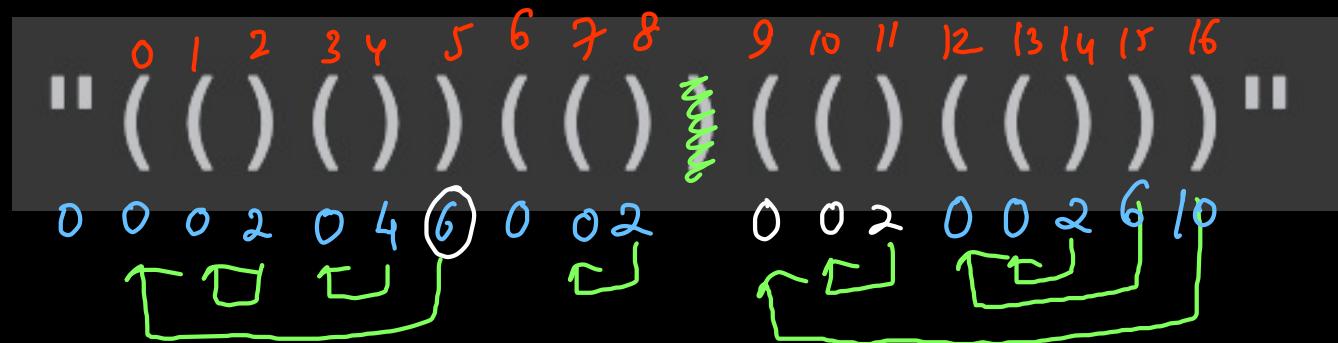
Time = $O(n)$

Space = $O(n)$

```
public String removeOuterParentheses(String s) {  
    boolean[] isOuter = new boolean[s.length()];  
  
    Stack<Integer> stk = new Stack<>();  
    for(int idx = 0; idx < s.length(); idx++){  
        if(s.charAt(idx) == '('){  
            stk.push(idx);  
        }  
        else {  
            int open = stk.pop();  
            if(stk.size() == 0){  
                // remove outermost  
                isOuter[idx] = isOuter[open] = true;  
            }  
        }  
    }  
  
    StringBuilder res = new StringBuilder();  
    for(int idx = 0; idx < s.length(); idx++){  
        if(isOuter[idx] == false){  
            res.append(s.charAt(idx));  
        }  
    }  
    return res.toString();  
}
```

$O(1)$

longest Valid Parentheses



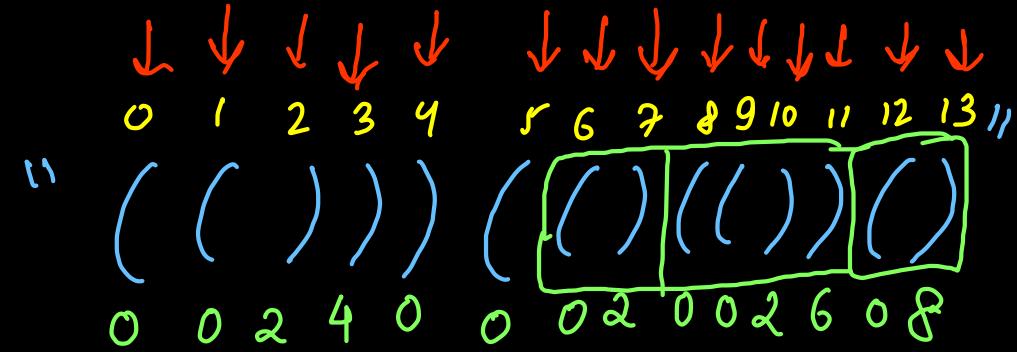
```

class Solution {
    public int longestValidParentheses(String s) {
        int[] res = new int[s.length()];
        Stack<Integer> stk = new Stack<>();

        int max = 0;
        for(int idx = 0; idx < s.length(); idx++){
            if(s.charAt(idx) == '('){
                stk.push(idx);
            } else {
                if(stk.size() == 0) continue;
                int open = stk.pop();
                res[idx] += (idx - open + 1);
                if(open > 0) res[idx] += res[open - 1];
                max = Math.max(max, res[idx]);
            }
        }
        return max;
    }
}

```

Time $\Rightarrow O(n)$
Space $\geq O(n)$



$$\max = \cancel{0} \cancel{2} \cancel{4} \cancel{8}$$

Monotonic stack

Next Greater Element

arr	200	30	10	20	70	50	60	90	10	160	130	130	150
-1	70	20	70	90	60	90	160	160	-1	150	150	150	-1

Brute force: \rightarrow Nested loops

Time $\Rightarrow O(n^2)$

Space $\Rightarrow O(1)$

1) Right side

for ($i: 0$ to $n-1$)

2) Strictly Greater

for ($j: i+1$ to $n-1$)

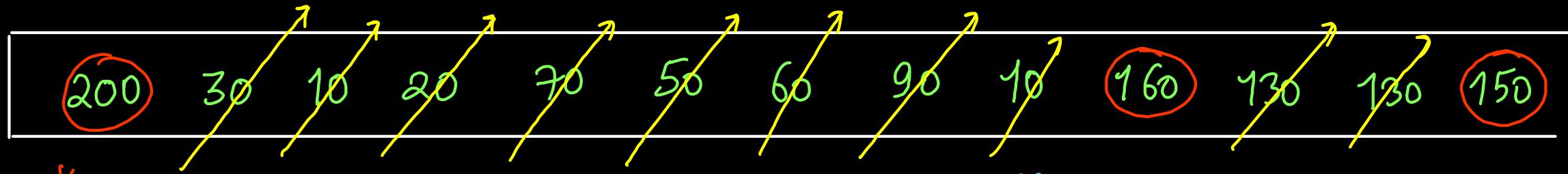
3) Nearcut

if ($arr[j] > arr[i]$)

$nge[i] = arr[j];$

break;

arr:	\downarrow												
200	30	10	20	70	50	60	90	160	10	160	130	130	150
-1	70	20	70	90	60	90	160	160	-1	150	150	150	-1



stack
node
yet to
be discovered
monotonically
decreasing

```
stack < integer > Stk;
for (int idx; 0 to N-1) {
    while (Stk.size() > 0 && arr[top] < arr[idx])
        rge.pop() = arr[idx];
    Stk.push(idx);
}
```

Approach 1)

Left to Right

Data → NGE yet to be discovered

```
public static long[] nextLargerElement(long[] arr, int n)
{
    Stack<Integer> stk = new Stack<>();
    long[] nge = new long[n];
    Arrays.fill(nge, -1);

    for(int idx = 0; idx < n; idx++){
        while(stk.size() > 0 && arr[idx] > arr[stk.peek()]){
            nge[stk.pop()] = arr[idx];
        }
        stk.push(idx);
    }

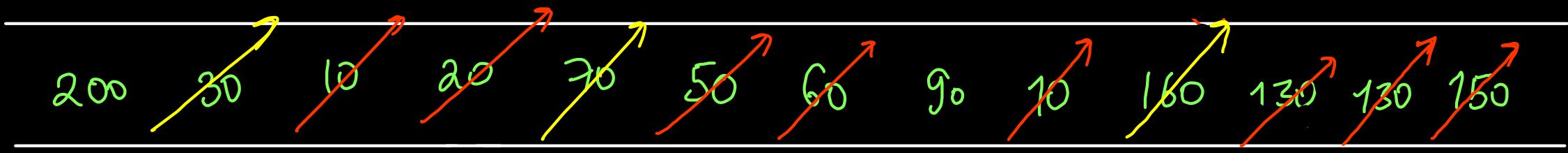
    return nge;
}
```

Time → $O(N)$

Space → $O(N)$

extra Space
or
auxiliary data
↓
Stack

\downarrow													
200	30	10	20	70	50	60	90	10	160	-1	130	130	150
-1	70	20	70	90	60	90	160	160	-1	150	150	150	-1



Stack <integer>

for (idx: n-1 to 0) {

pop smaller
or equal } while (stk.size() > 0 && arr[top] <= arr[idx])
stk.pop();

find NGE_{curr} → nge[idx] = (size > 0) ? arr[stk.peek()] : -1;

push curr → stk.push(idx);

Stack
{Potential answers}

monotonic
decreasing

Approach 2)

Right to Left

Data → Potential Nge
answers

Time $\Rightarrow O(n)$
Space $\Rightarrow O(n)$ extra
(stack)

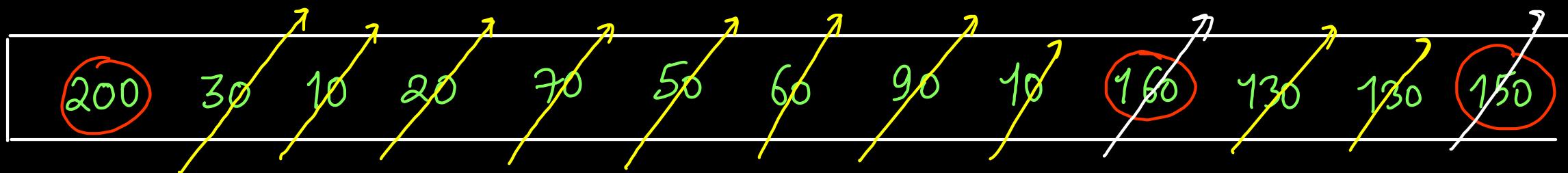
```
public static long[] nextLargerElement(long[] arr, int n)
{
    Stack<Integer> stk = new Stack<>();
    long[] res = new long[n];

    for(int idx = n - 1; idx >= 0; idx--){  $\rightarrow O(n)$  push
        while(stk.size() > 0 && arr[stk.peek()] <= arr[idx]){  $\rightarrow O(n)$  pop
            stk.pop();
        }
        res[idx] = (stk.size() > 0) ? arr[stk.peek()] : -1;
        stk.push(idx);
    }

    return res;
}
```

NGE - Circular Array

↓	200	30	10	20	70	50	60	90	10	160	130	130	150
-1	70	20	70	90	60	90	160	160	200	150	150	150	200



```
public int[] nextGreaterElements(int[] arr) {  
    int n = arr.length;  
    int[] nge = new int[n];  
    Arrays.fill(nge, -1);  
    Stack<Integer> stk = new Stack<>();  
  
    for(int idx = 0; idx < n; idx++){  
        while(stk.size() > 0 && arr[idx] > arr[stk.peek()]){  
            nge[stk.pop()] = arr[idx];  
        }  
        stk.push(idx);  
    }  
  
    for(int idx = 0; idx < n; idx++){  
        while(stk.size() > 0 && arr[idx] > arr[stk.peek()]){  
            nge[stk.pop()] = arr[idx];  
        }  
    }  
  
    return nge;  
}
```

Time $\Rightarrow O(n)$

Space $\Rightarrow O(n)$

Due to circular array

Online Algorithm

T

↓
server

↓

packets / discrete

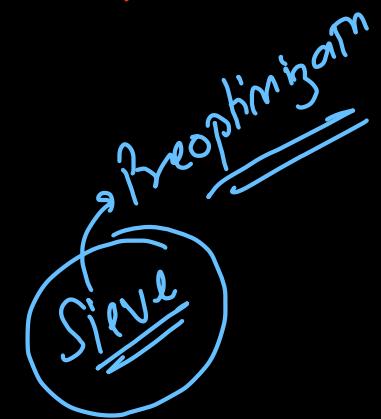
Offline Algorithm

T

↓
download

↓

complete info



Online Stock span

Online Stock Span

D_0	D_1	D_2	D_3	D_4	D_5	D_6	D_7	D_8	D_9	D_{10}	D_{11}	D_{12}
200	30	10	20	70	50	60	90	10	160	130	130	150
1	1	1	2	4	1	2	7	1	9	1	2	3

Span = Max^m consecutive days in the left
 for which the current
 price is max^m

 $= \text{current index} - \text{next greater in the left index}$

	D_3	D_2	D_1	D_0
200	30	10	20	70
	3	1	1	1

potential answers

```

Stack<Integer> stk = new Stack<>();
ArrayList<Integer> arr = new ArrayList<>();

public int next(int price) {
    while(stk.size() > 0 && arr.get(stk.peek()) <= price){
        stk.pop();
    }

    int span = (stk.size() > 0) ? arr.size() - stk.peek()
                                : arr.size() + 1;

    stk.push(arr.size());
    arr.add(price);
    return span;
}

```

Stack :



arr:

size = \emptyset
1
2

50 60 90

→ $O(N)$ total time
 $\therefore \rightarrow N / N = O(1)$ avg

739. Daily Temperatures

Hint



Medium



9.9K

227



Companies

Given an array of integers `temperatures` represents the daily temperatures, return an array `answer` such that `answer[i]` is the number of days you have to wait after the i^{th} day to get a warmer temperature. If there is no future day for which this is possible, keep `answer[i] == 0` instead.

Example 1:

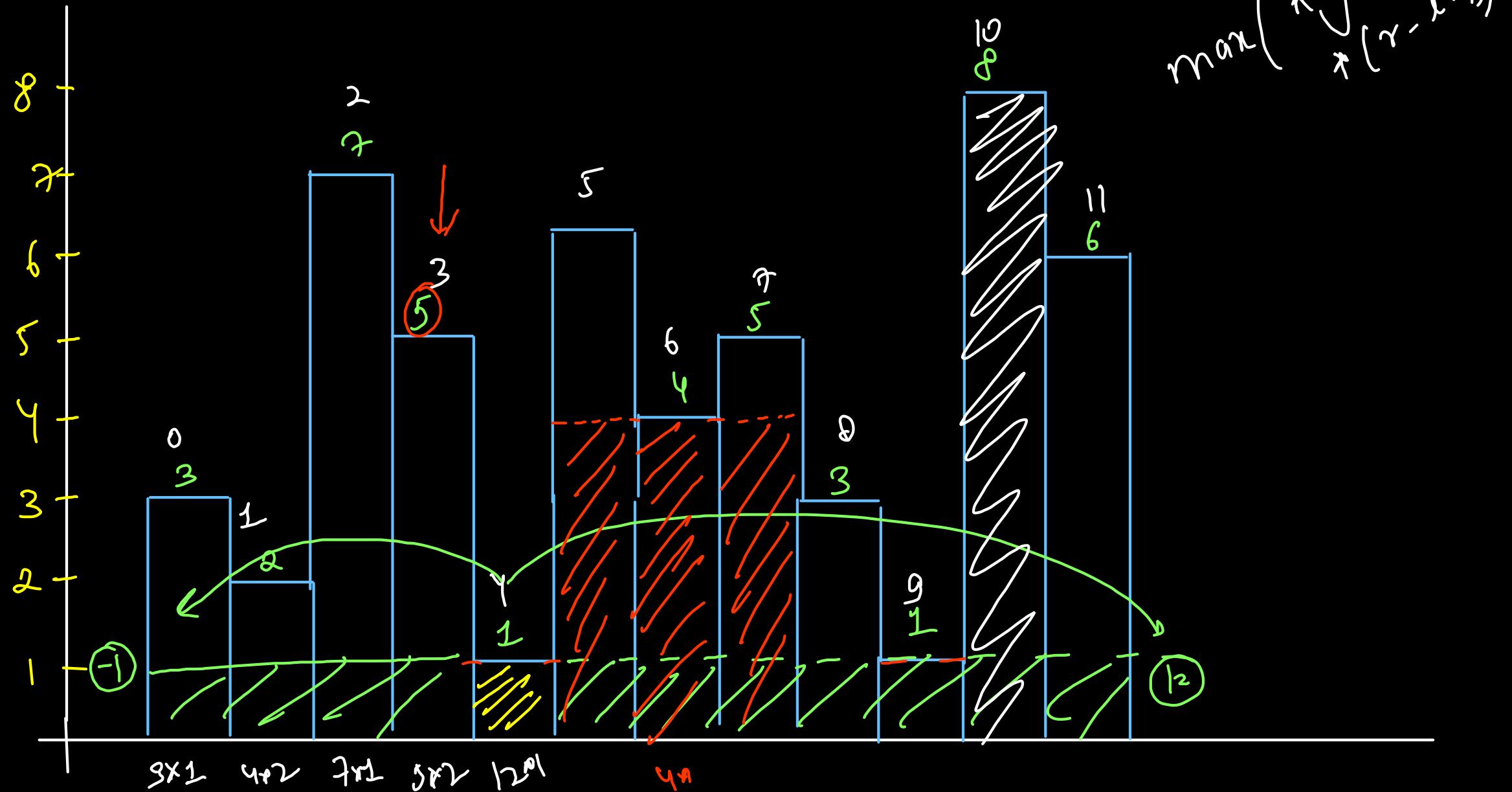
Input: `temperatures = [73, 74, 75, 71, 69, 72, 76, 73]`
Output: `[1, 1, 4, 2, 1, 1, 0, 0]`



```
public int[] dailyTemperatures(int[] arr) {  
    Stack<Integer> stk = new Stack<>();  
    int[] res = new int[arr.length];  
  
    for(int idx = arr.length - 1; idx >= 0; idx--){  
        while(stk.size() > 0 && arr[stk.peek()] <= arr[idx]){  
            stk.pop();  
        }  
  
        res[idx] = (stk.size() > 0) ? stk.peek() - idx : 0;  
        stk.push(idx);  
    }  
  
    return res;  
}
```

Time $\Rightarrow O(n)$
Space $\Rightarrow O(n)$

Largest Area Histogram



```

public int[] rightNSE(int[] arr){
    int n = arr.length;
    Stack<Integer> stk = new Stack<>();
    int[] res = new int[n];

    for(int idx = n - 1; idx >= 0; idx--){
        while(stk.size() > 0 && arr[stk.peek()] >= arr[idx]){
            stk.pop();
        }
        res[idx] = (stk.size() > 0) ? stk.peek() : n;
        stk.push(idx);
    }

    return res;
}

```

```

public int[] leftNSE(int[] arr){
    int n = arr.length;
    Stack<Integer> stk = new Stack<>();
    int[] res = new int[n];

    for(int idx = 0; idx < n; idx++){
        while(stk.size() > 0 && arr[stk.peek()] >= arr[idx]){
            stk.pop();
        }
        res[idx] = (stk.size() > 0) ? stk.peek() : -1;
        stk.push(idx);
    }

    return res;
}

```

```

public int largestRectangleArea(int[] arr) {
    int[] left = leftNSE(arr); → O(n)
    int[] right = rightNSE(arr); → O(n)
    int max = 0;

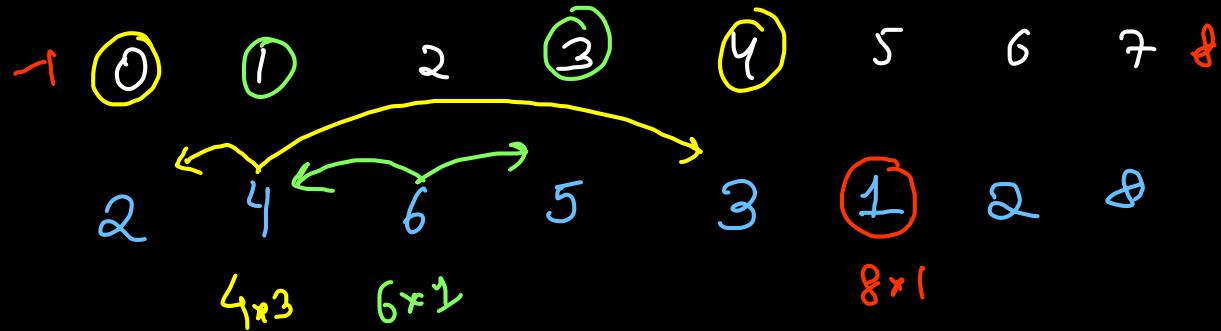
    for(int idx = 0; idx < arr.length; idx++){
        int area = arr[idx] * (right[idx] - left[idx] - 1);
        max = Math.max(max, area);
    }

    return max;
}

```

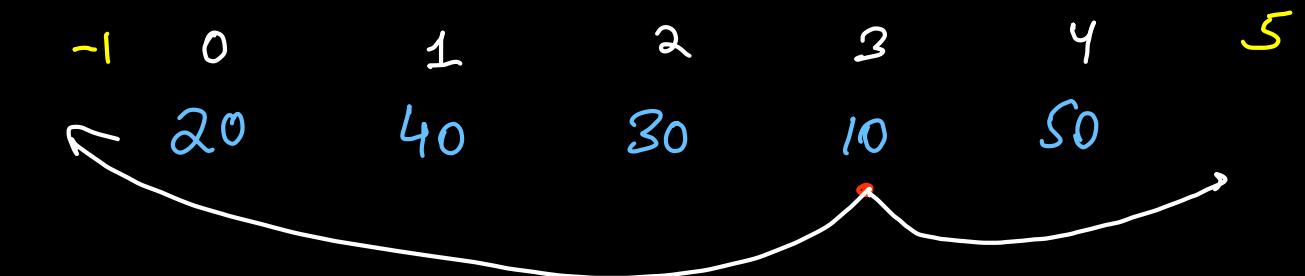
$\Theta(n)$

Time $\Rightarrow O(n)$
Space $\Rightarrow O(n)$



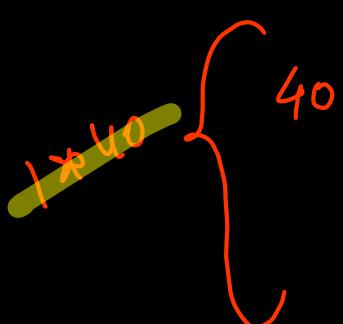
Leetcode 907

Sum of Subarray Minimums

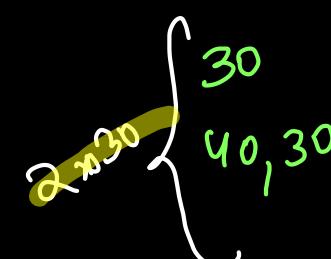


$$(r-c)(c-l)$$

$$(3-0) * (0 - (-1))$$



$$(2-1)(1-0) = 2 \cdot 1$$



20, 40, 30, 10, 50
20, 40, 30, 10

40, 30, 10

30, 10

10

10, 50

Time $O(n^2)$ Approach (1)

Find all subarrays

Space $O(1)$

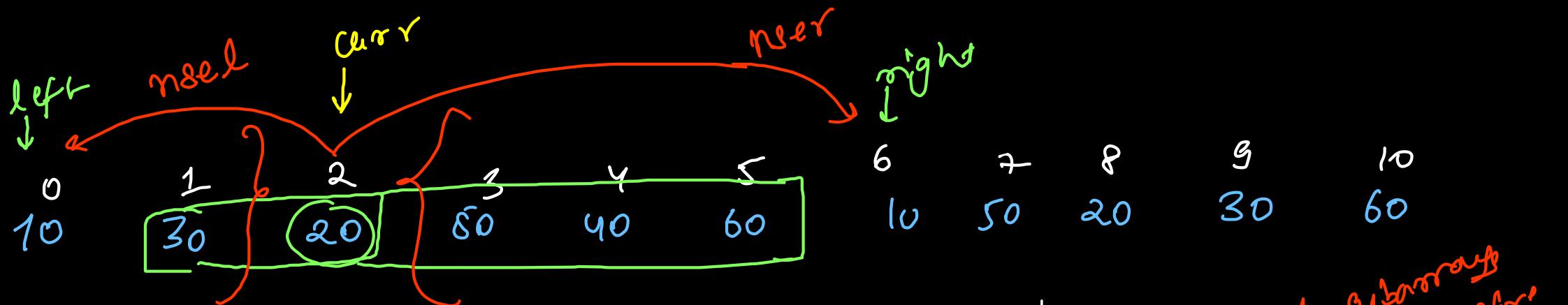
↓
Find min^m

↓
Take sum of
all min^m

Bruteforce

40, 30, 10, 50
30, 10, 50

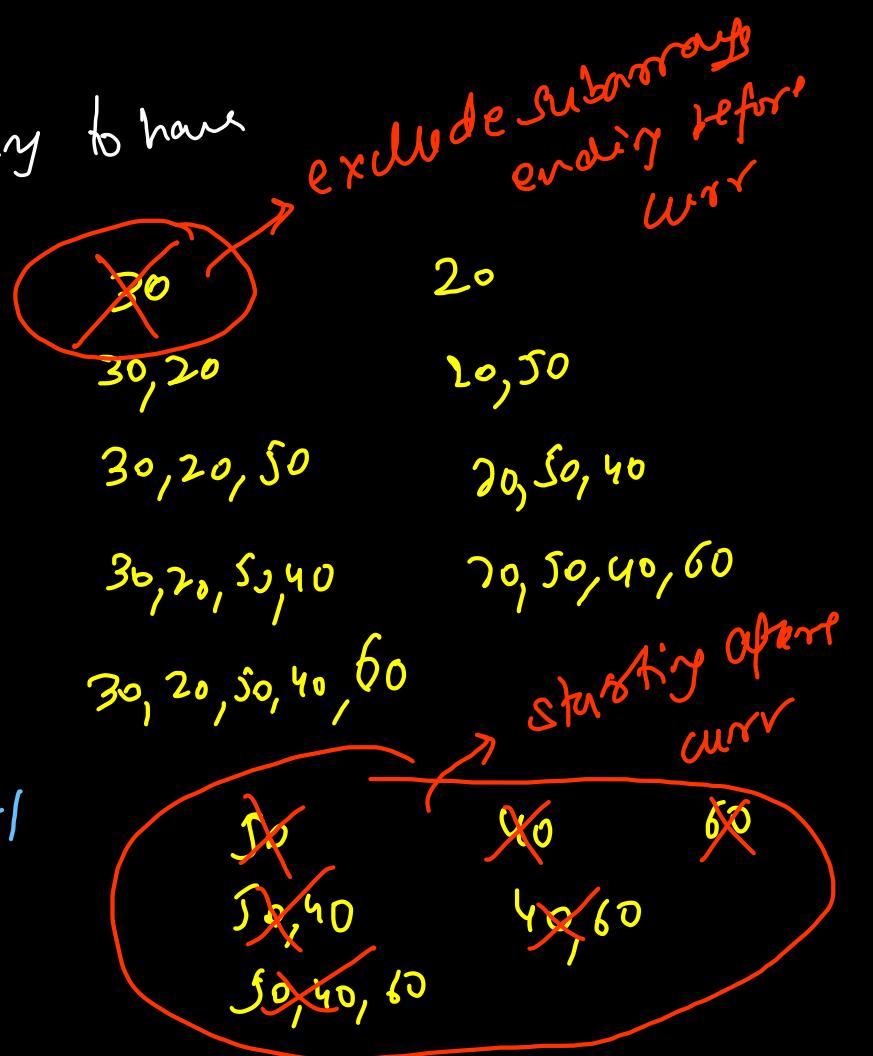
9 * 10



how many subarrays are there which are necessary to have

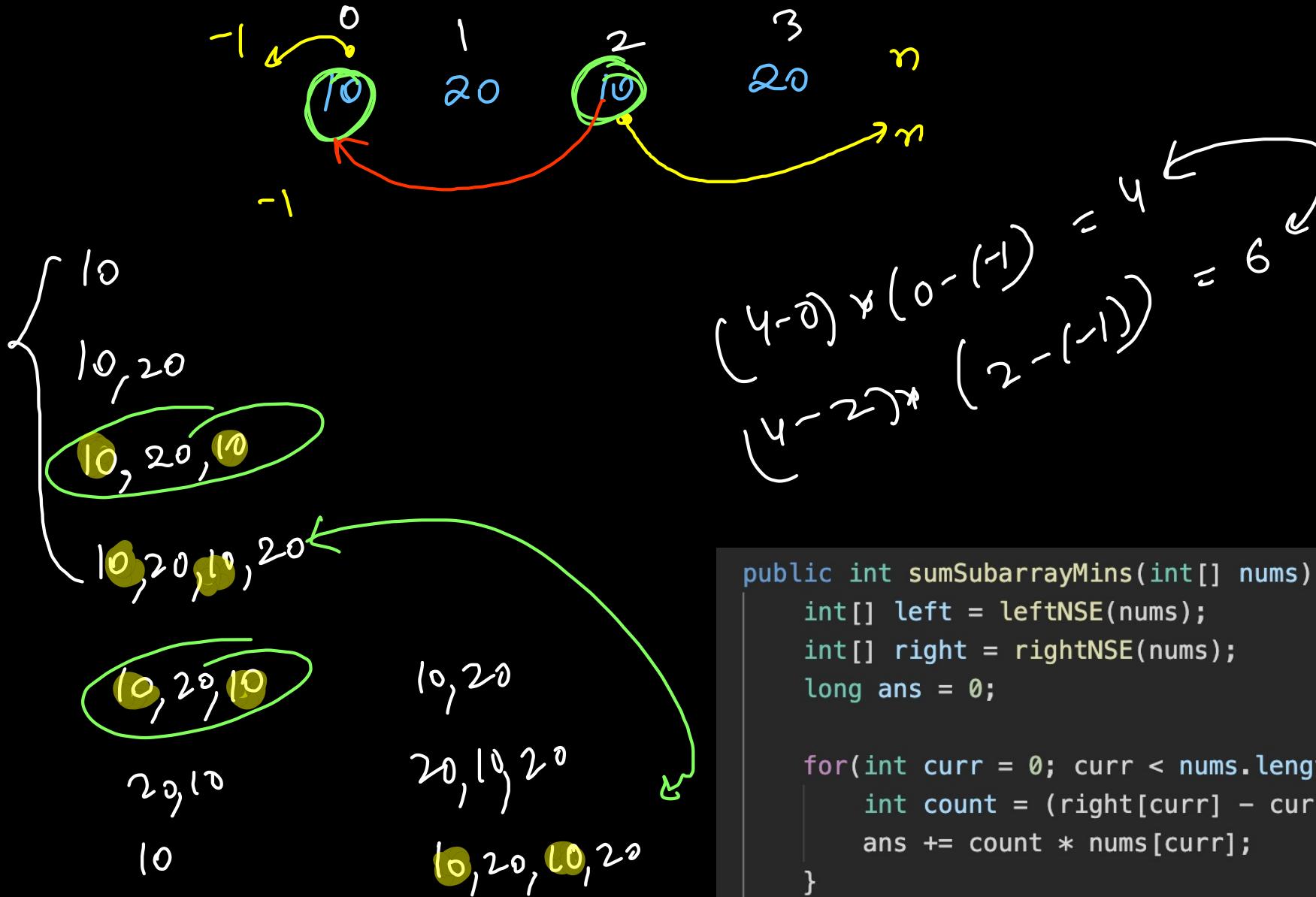
- ① curr index included
- ② start after left - (nse)
- ③ end before right (nse)

$$n = \text{right} - \text{left} - 1$$



$$\frac{(\gamma - \lambda - 1)(\gamma - \lambda)}{2} - \frac{(\underbrace{\text{curr-left}-1}_{2}) * (\text{curr-left})}{2} - \frac{(\underbrace{\text{right-curr}-1}_{2}) * (\text{right-curr})}{2}$$

$$= (\text{right-curr}) * (\text{curr-left})$$





$$20 * (1-0) * (0-(-1)) + 10 * (4-1) * (1-(-1)) + 30 * (3-2) * (2-1)$$

$\{30\}$

$\{20\}$

$\{20, 10\}$ { 20, 10, 30 } { 20, 10, 30, 10 }

$\{10\}$ { 10, 30 } { 10, 30, 10 }

Corner Case! →

① modular arithmetic
(answer can be very large)

② duplicate values

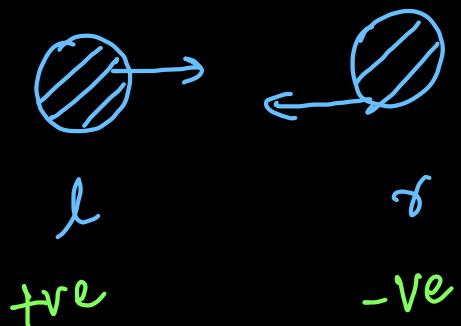
```
public int sumSubarrayMins(int[] nums) {
    int[] left = leftNSE(nums);
    int[] right = rightNSE(nums);
    long ans = 0;

    for (int curr = 0; curr < nums.length; curr++) {
        long count = (right[curr] - curr) * (curr - left[curr]);
        ans = (ans + (count * nums[curr]) % 1000000007) % 1000000007;
    }

    return (int) ans;
}
```



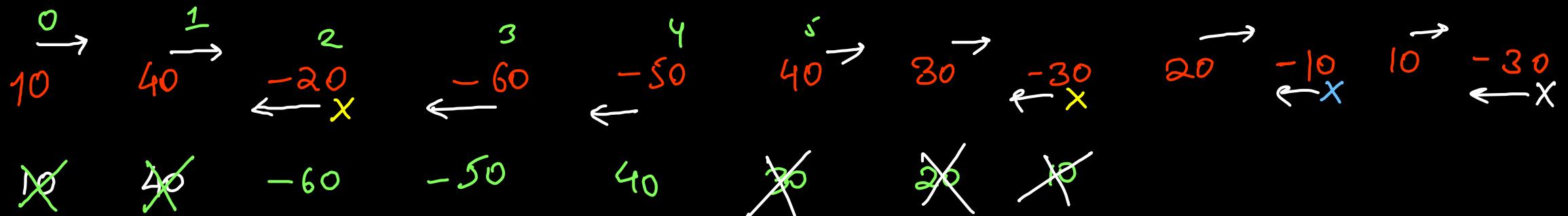
Asteroid collision



$|l| > |r| \Rightarrow$ destroy $r \Rightarrow$ don't push(r)

$|l| = |r| \Rightarrow$ destroy $l \& r \Rightarrow$ pop left $\&$ don't push(r)

$|l| < |r| \Rightarrow$ destroy $l \Rightarrow$ pop(while) $\&$ push(r)



```

public int[] asteroidCollision(int[] asteroids) {
    ArrayList<Integer> stk = new ArrayList<>();

    for(int idx = 0; idx < asteroids.length; idx++){
        int right = asteroids[idx];

        if(right > 0) stk.add(right);
        else if(stk.size() == 0 || stk.get(stk.size() - 1) < 0){
            stk.add(right);
        }
        else if(stk.get(stk.size() - 1) < -right){
            stk.remove(stk.size() - 1);
            idx--; // to come back to same stone
        }
        else if(stk.get(stk.size() - 1) == -right){
            stk.remove(stk.size() - 1);
        }
    }

    int[] res = new int[stk.size()];
    for(int i = 0; i < res.length; i++){
        res[i] = stk.get(i);
    }
    return res;
}

```

Time $\Rightarrow \Theta(n)$

Space $\Rightarrow \Theta(n)$ {stack} \equiv

Celebrity Problem

$m[\text{row}][\text{col}]$

	P_0	P_1	P_2	P_3	P_4	
P_0	-	1	1	1	0	
P_1	0	-	0	1	1	
P_2	0	0	-	1	0	
P_3	0	0	0	-	0	
P_4	1	1	1	1	-	

$1 \Rightarrow$ now knows col

$0 \Rightarrow$ now doesn't col

Celebrity

① Everybody should him

② He should know no one

Q1) Can be there more than
1 celebrity?

Q2) Can be there 0 celebrities?
↳ Yes

Assume 2 celebrities are there

p_i^o & p_j^o are celebs

	p_i^o		p_j^o
p_i^o	1 1 1 1 1 0 0 0	0 0 0 0 0 0 0	1 1 1 1 1 0 0 0 0 0
p_j^o	1 1 1 1 1 0 0 0 0 0	0 0 0 0 0 0 0	1 1 1 1 1 0 0 0
	1 1 1 1 1 0 0 0 0 0 0 0	0 0 0 0 0 0 0 0	1 1 1 1 1 0 0 0 0 0 0

Contradiction

P_0	P_1	P_2	P_3	P_4
1	1	1	0	
0	-	0	1	1
0	0	-	1	0
0	0	-	0	
1	1	1	1	-

Banbe force

Check row == 0 & col == -1
for every person

\Rightarrow Time $\Rightarrow O(n^2)$

Space $\Rightarrow O(1)$

Elimination Answer

$$P_0 \xrightarrow{\checkmark} P_1$$

$$m[P_0][P_1] = 1$$

P_0 knows P_1

Eliminate P_0

$$P_1 \xrightarrow{\alpha} P_2$$

$$m[P_1][P_2] = 0$$

P_1 doesn't know P_2

Eliminate P_2

$$P_1 \xrightarrow{\checkmark} P_3$$

$$m[P_1][P_3] = 1$$

P_1 knows P_3

Eliminate P_1

$$P_3 \xrightarrow{\alpha} P_4$$

$$m[P_3][P_4] = 0$$

P_3 doesn't know P_4

Eliminate P_3

celebrity = \nearrow ~~(3)~~

```
class Solution
{
    int celebrity(int M[][] , int n) {
        int ans = 0;
        for(int idx = 1; idx < n; idx++){
            if(M[idx][ans] == 0){
                ans = idx;
            }
        }

        for(int i = 0; i < n; i++){
            if(i == ans) continue;
            if(M[ans][i] != 0) return -1;
            if(M[i][ans] != 1) return -1;
        }

        return ans;
    }
}
```

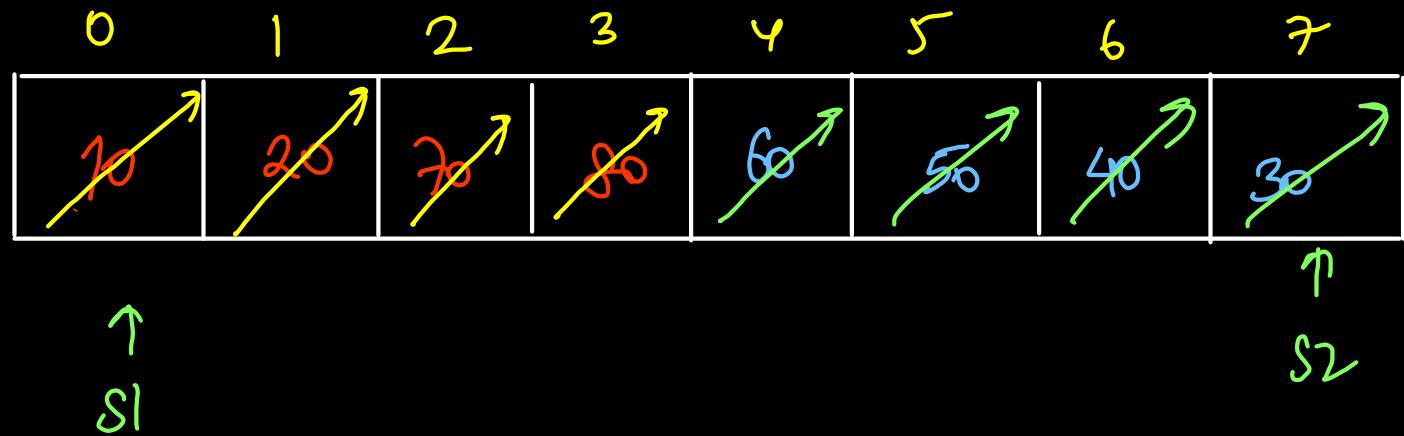
finding
the potential
celebrity

No celebrity

Time $\Rightarrow O(n)$

Space $\Rightarrow O(1)$

Implement Two Stacks in Array



$S_1 > S_2$: Stack overflow

$S_1 == 0$: pop : S_1 underflow

$S_2 == n-1$: pop : S_2 underflow

$S_1.add(10)$

$S_1.add(20)$

$S_2.add(30)$

$S_2.add(40)$

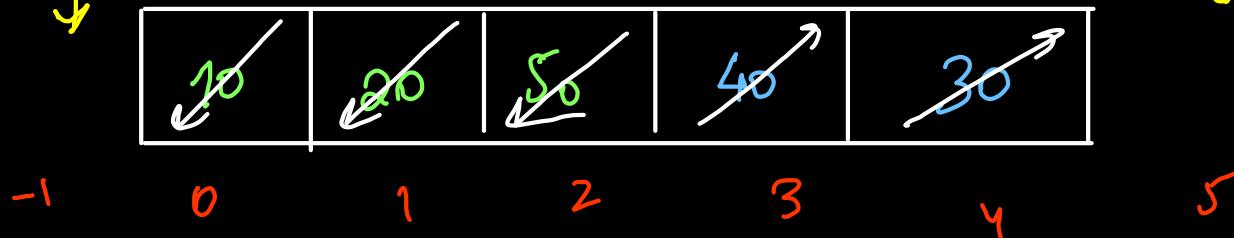
$S_2.add(50)$

$S_2.add(60)$

$S_1.add(70)$

$S_1.add(80)$

top1



-1

0

1

2

3

4

5

/* Structure of the class is

class TwoStack

{

int size;

int top1,top2;

int arr[] = new int[100];

TwoStack()

{

size = 100;

top1 = -1;

top2 = size;

}

}*/

top2

$\text{top2} = \text{top1} + 1$

class Stacks

{

void push1(int x, TwoStack sq)

$\rightarrow O(1)$

{

if(sq.top2 == sq.top1 + 1) return; // Overflow
sq.top1++;
sq.arr[sq.top1] = x;

}

void push2(int x, TwoStack sq)

$\rightarrow O(1)$

{

if(sq.top2 == sq.top1 + 1) return; // Overflow
sq.top2--;
sq.arr[sq.top2] = x;

}

int pop1(TwoStack sq)

$\rightarrow O(1)$

{

if(sq.top1 == -1) return -1;
return sq.arr[sq.top1--];

}

//Function to remove an element from top of the stack2.

int pop2(TwoStack sq)

$\rightarrow O(1)$

{

if(sq.top2 == sq.size) return -1;
return sq.arr[sq.top2++];

}

Remove K Digits

LeetCode 402

"~~8~~356~~8~~9789"

$k = 5$

answer : \rightarrow "~~8~~356~~8~~9~~7~~8~~9~~"

① place value
higher ↑ deletion
smaller ↓

✓② digit ↑
deletion ↓ small

“ 1 0 0 0 5 7 8 ” $k = 3$

“ 0 0 0 5 7 8 ”
leading zeros

Time $\Rightarrow O(n)$

Space $\Rightarrow O(1)$

```
class Solution {
    public String removeKdigits(String num, int k) {
        StringBuilder res = new StringBuilder();
        char[] chs = num.toCharArray();
        for(char ch: chs){
            while(k > 0 && res.length() > 0 && res.charAt(res.length() - 1) > ch){
                res.deleteCharAt(res.length() - 1);
                k--;
            }
            res.append(ch);
        }
        while(k-- > 0){
            res.deleteCharAt(res.length() - 1);
        }
        int idx = 0;
        while(idx < res.length() && res.charAt(idx) == '0'){
            idx++;
        }
        if(idx == res.length()) return "0";
        return res.toString().substring(idx);
    }
}
```

monotonic stack

delete last digits

remove leading zeros

```
class Solution {
    public String removeKdigits(String num, int k) {

        StringBuilder res = new StringBuilder();

        char[] chs = num.toCharArray();
        for(char ch: chs){
            while(k > 0 && res.length() > 0 && res.charAt(res.length() - 1) > ch){ \n                res.deleteCharAt(res.length() - 1);
                k--;
            }
            res.append(ch);
        }

        while(k-- > 0){
            res.deleteCharAt(res.length() - 1);
        }

        int idx = 0;
        while(idx < res.length() && res.charAt(idx) == '0')
            idx++;

        if(idx == res.length()) return "0";
        return res.toString().substring(idx);
    }
}
```

$k = 7 - 3 - 2$

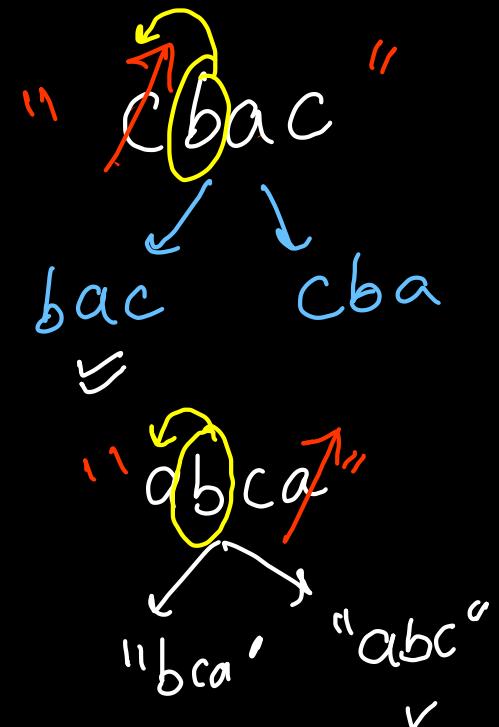
2 5 6 0 0 7 9 4 8 9 "

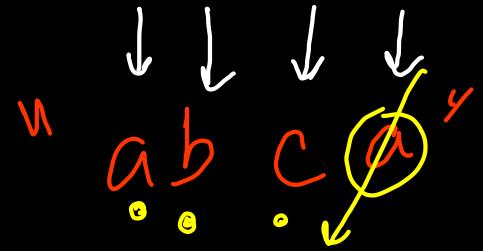
" 2 5 6 0 0 7 9 4 8 9 "

Remove Duplicate letters

lexicographical order / dictionary order

" 7 2 4 6 9 " > " 7 2 4 6 8 9 9 9 9 "





a b c
↑ ↑ ↑

a: 2

b: 1

c: 1

① monotonic stack
(String Builder)

② frequency array
(duplicates)

③ visited array

" a . d b c d a e f a d e "

" a ~~d~~ b c d e f "

a : β_2^1 ✗ ✓

b : f_0 ✗ ✓

c : f_0 ✗ ✓

d : β_2 ✓

e : f_1 ✗ ✓

f : f_0 ✗ ✓

- ① duplicate
- ② lexicographical

```

public String removeDuplicateLetters(String s) {
    int[] freq = new int[26];
    boolean[] vis = new boolean[26];

    char[] chs = s.toCharArray();
    for(char ch: chs) freq[ch - 'a']++;

    StringBuilder res = new StringBuilder();

    for(char ch: chs){
        freq[ch - 'a']--;
        if(vis[ch - 'a'] == true) continue;
        vis[ch - 'a'] = true;

        while(res.length() > 0){
            char top = res.charAt(res.length() - 1);
            if(top > ch && freq[top - 'a'] > 0){
                res.deleteCharAt(res.length() - 1);
                vis[top - 'a'] = false;
            } else break;
        }

        res.append(ch);
    }

    return res.toString();
}

```

“a c f d b a f c”

“a c f[↗]d b f ”
 . . X . . .

a: ✓/o

b: ✓/o

c: ✓/i

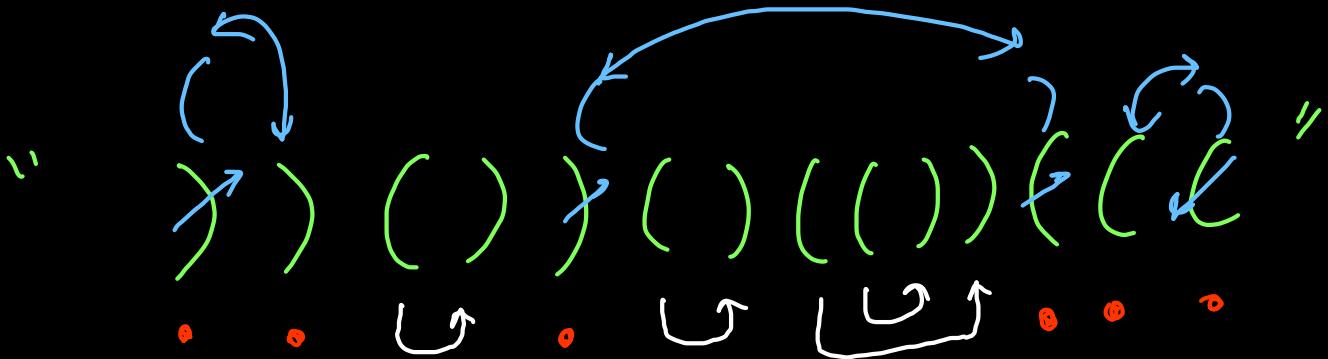
d: ✓/o

f: ✓/o

Time $\Rightarrow O(n)$

Space $\Rightarrow O(2^6) \approx O(1)$

Min^m Additions/Removals/Swaps
to make it Balanced



⑥ Additions \Rightarrow number of unbalanced characters

⑥ Removals \Rightarrow number of unbalanced characters

Swaps \Rightarrow no of open = no of closing ; no of open/close

③

fm monir {
 → Recursion & Backtracking
 → Hashmaps & Heaps

Bit manipulation
↓
CP batch

fm monir {
 → Trees
 → Graphs

fm monir {
 → Dynamic Programming
 → Tries

```

public int minAddToMakeValid(String s) {
    int unbalanced_open = 0, unbalanced_close = 0;

    for(char ch: s.toCharArray()){
        if(ch == '(') unbalanced_open++;
        else {
            if(unbalanced_open > 0) unbalanced_open--;
            else unbalanced_close++;
        }
    }

    return unbalanced_open + unbalanced_close;
}

```

$\text{Open} = 0 \nearrow \neq 8 \neq 1 \phi$
 $\text{Close} = \phi \nearrow \neq 3 \quad \nearrow 3$

Time $\Rightarrow O(n)$

Space $\Rightarrow O(1)$

" l) e (e) t) (c (o) d) (e "



Time $\Rightarrow \mathcal{O}(n)$

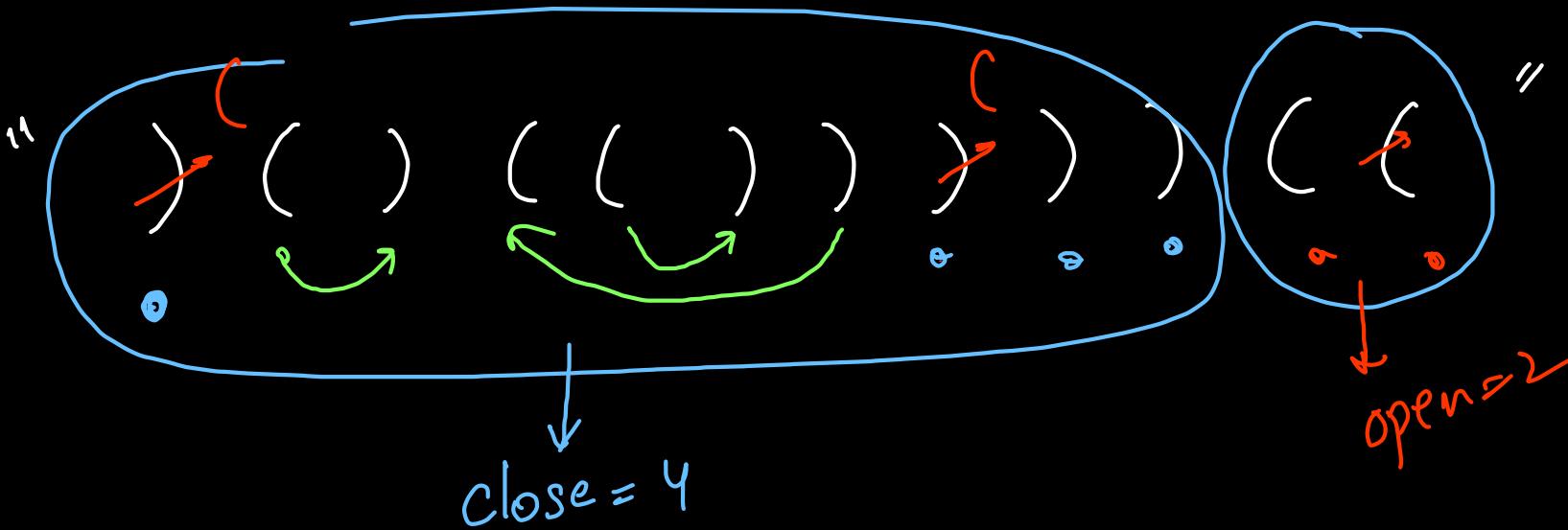
Space $\Rightarrow \mathcal{O}(n)$
 \downarrow
 boolean stack

```

public String minRemoveToMakeValid(String s) {
    boolean[] unbalanced = new boolean[s.length()];

    Stack<Integer> stk = new Stack<>();
    for(int idx = 0; idx < s.length(); idx++){
        if(s.charAt(idx) == '('){
            stk.push(idx);
            unbalanced[idx] = true;
        } else if(s.charAt(idx) == ')'){
            if(stk.size() > 0){
                unbalanced[stk.pop()] = false;
            } else {
                unbalanced[idx] = true;
            }
        }
    }

    StringBuilder res = new StringBuilder();
    for(int idx = 0; idx < s.length(); idx++){
        if(unbalanced[idx] == false){
            res.append(s.charAt(idx));
        }
    }
    return res.toString();
}
  
```



() ()) ())

close = 1

open = 3

```

int countRev (String s)
{
    int open = 0, close = 0;

    for(char ch: s.toCharArray()){
        if(ch == '{') open++;
        else {
            if(open > 0) open--;
            else close++;
        }
    }

    int sum = open + close;
    if(sum % 2 == 1) return -1; //odd length
    if(open % 2 == 1) return sum / 2 + 1;
    return sum / 2;
}

```

")) ((("

$$\text{open} = 2 \quad \text{close} = 3$$

")) ((() "

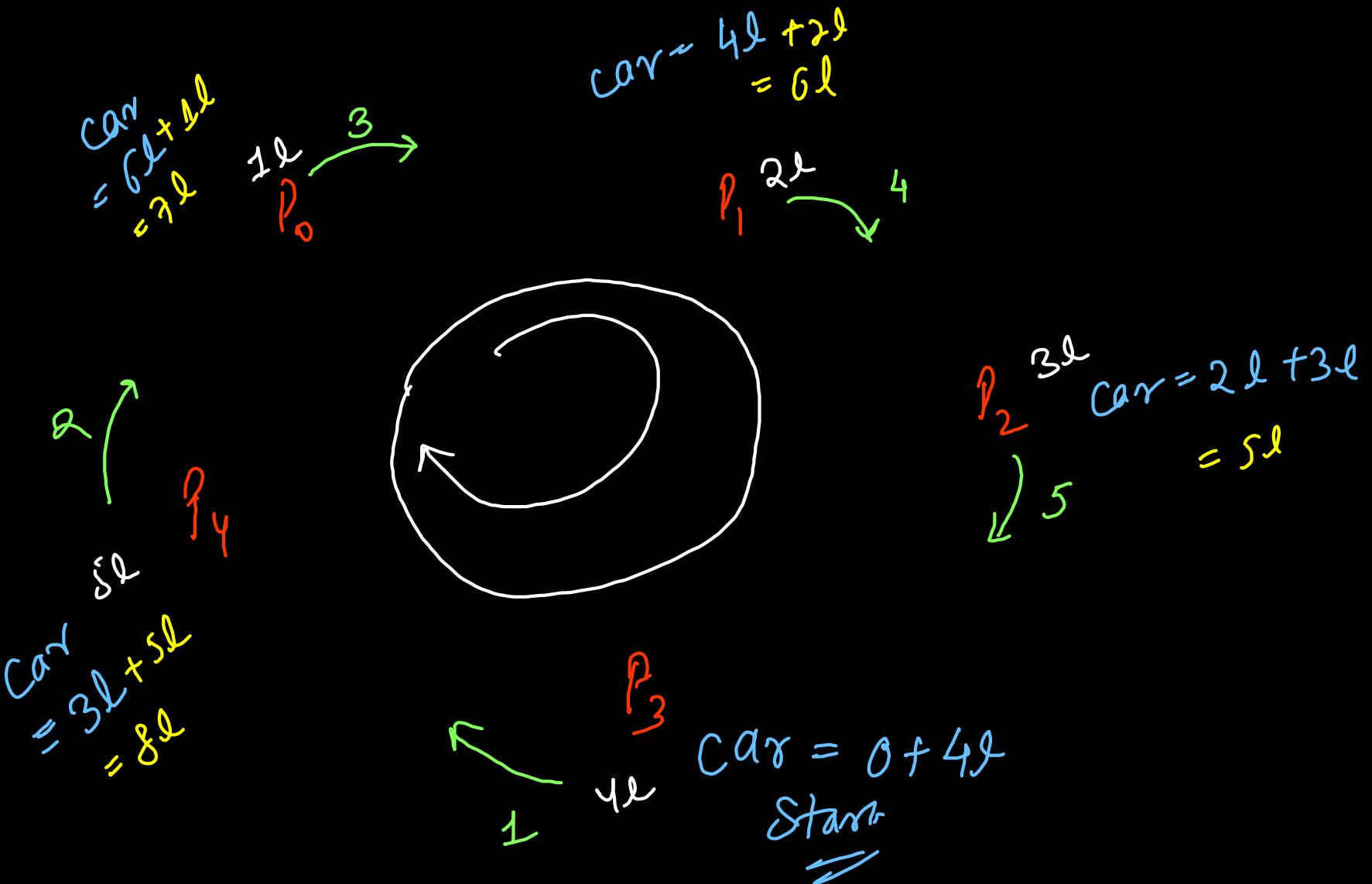
$$\text{open} = 1 \quad \text{close} = 3 \\ (3+1)/2 + 1 = 3$$

")) (() "

$$\text{open} = 2 \quad \text{close} = 2 \\ (2+2)/2 = 2$$

Circular Tour / Gas Station

LC 134



```

public int canCompleteCircuit(int[] gas, int[] cost) {
    int n = gas.length;
    for(int st = 0; st < n; st++){
        int fuel = 0;
        boolean ans = true;

        for(int jump = 0; jump < n; jump++){
            int curr = (st + jump) % n;
            fuel += gas[curr];

            if(fuel < cost[curr]){
                ans = false;
                break;
            }
            fuel -= cost[curr];
        }

        if(ans == true) return st;
    }
    return -1;
}

```

no starting pt available

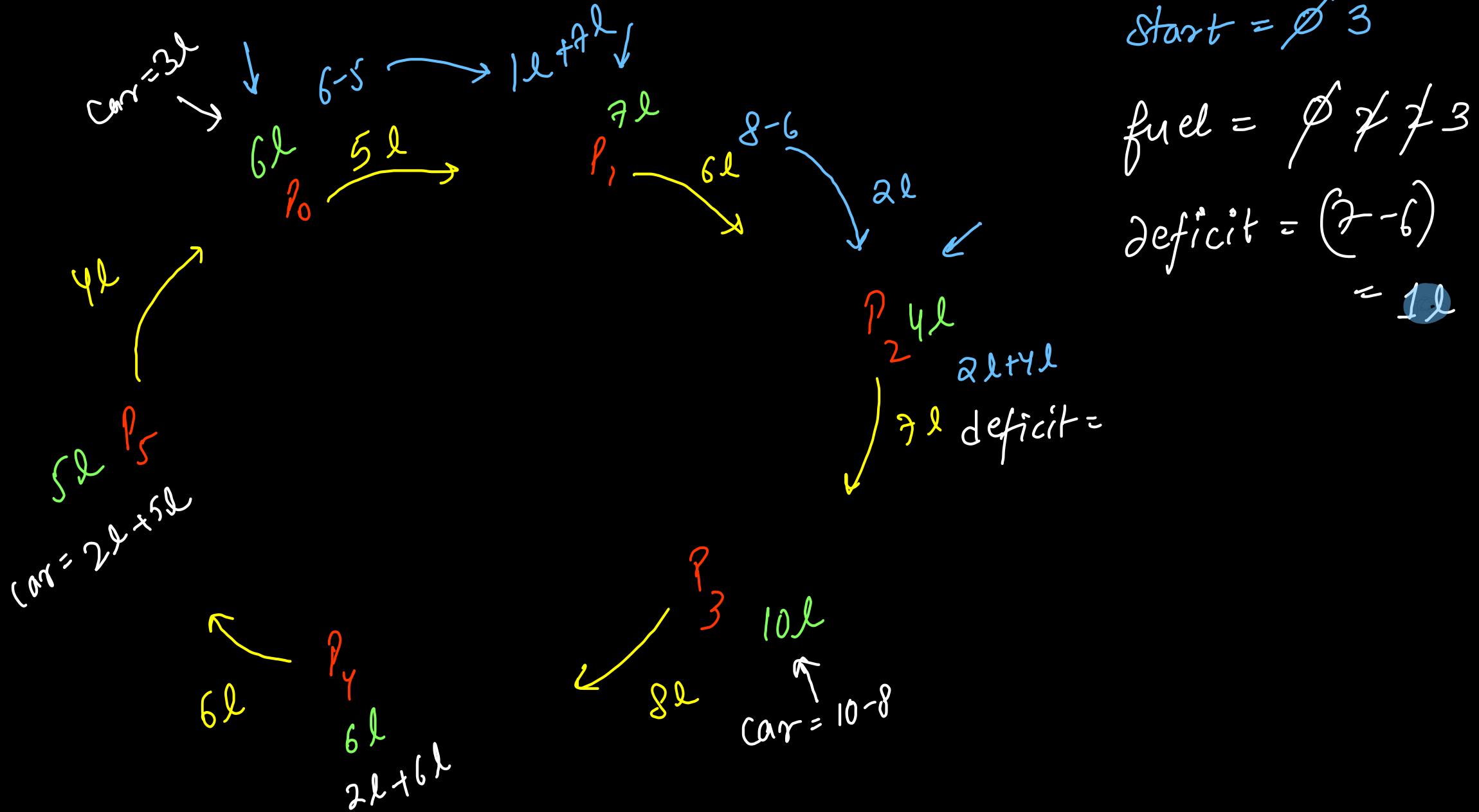
Brute force

↳ consider all pts as
Starting index

Time $\Rightarrow O(n^2)$ ~~TLE~~

worst case

Space $\Rightarrow O(1)$
extra



```

class Solution {
    public int canCompleteCircuit(int[] gas, int[] cost) {
        int deficit = 0, fuel = 0, start = 0;

        for(int idx = 0; idx < gas.length; idx++){
            fuel += gas[idx] - cost[idx];

            if(fuel < 0){      ↗ skip starting pt as
                start = idx + 1;   p1, p2
                deficit += -fuel;
                fuel = 0;
            }
        }

        if(fuel >= deficit) return start;
        return -1;
    }
}

```

Time $\Rightarrow \Theta(n)$

Space $\Rightarrow \Theta(1)$

BODMAS

⇒ Bracket

⇒ Division = multiplication
[same priority
left to right
associative]

Expressions

$$a + b$$

$$a - b$$

$$a / b$$

$$a * b$$

↓

Infix

$\sqrt{1} \text{ op } \sqrt{2}$

$5 + (2 * (6 / 3))$

$$5 * 3 / 2$$

Maths

$$(5 * 3) / 2$$

7

5

1

$5 * (3 / 2)$