

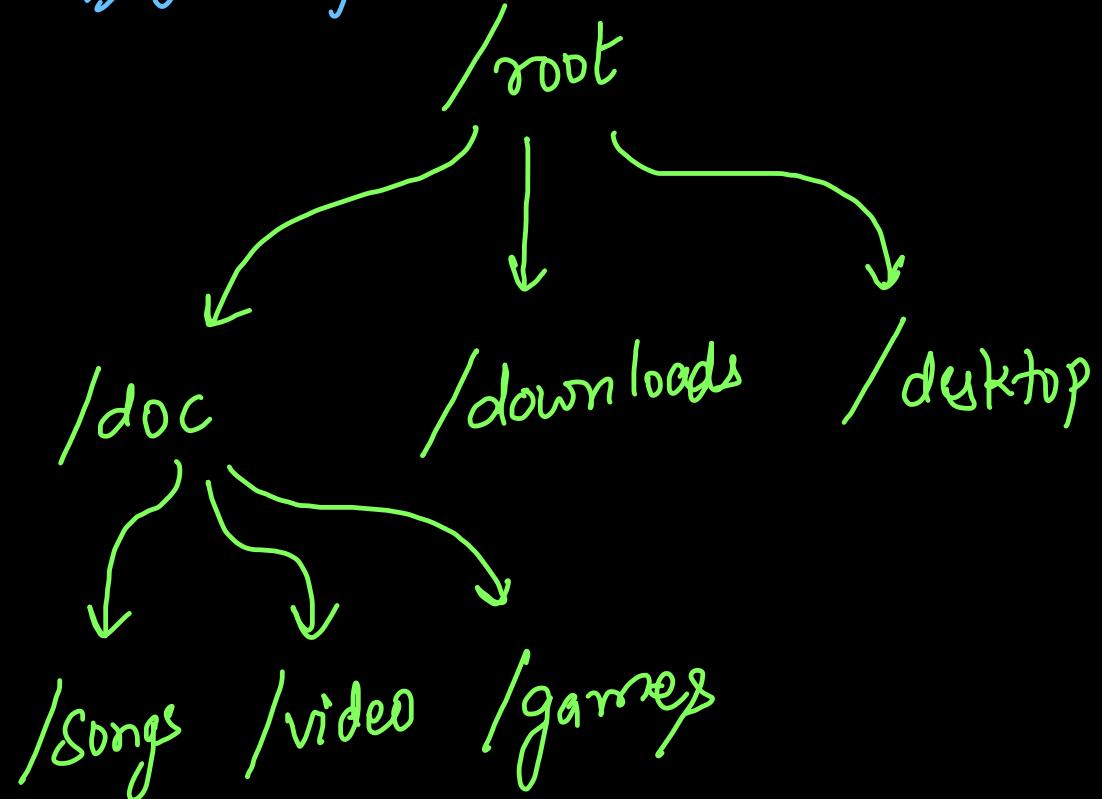
Binary Trees & Binary Search Trees

Data Structures

- Array, String → static
 - ArrayList, StringBuilder → dynamic
 - LinkedList → dynamic & non-contiguous
 - Stack, Queue, Deque
 - ↓ LIFO
 - ↓ FIFO
 - ↓ addFirst/Last
removeFirst/Last
- } linear Data Structures

Real-life examples of Heirarchical Data Structures

eg file system

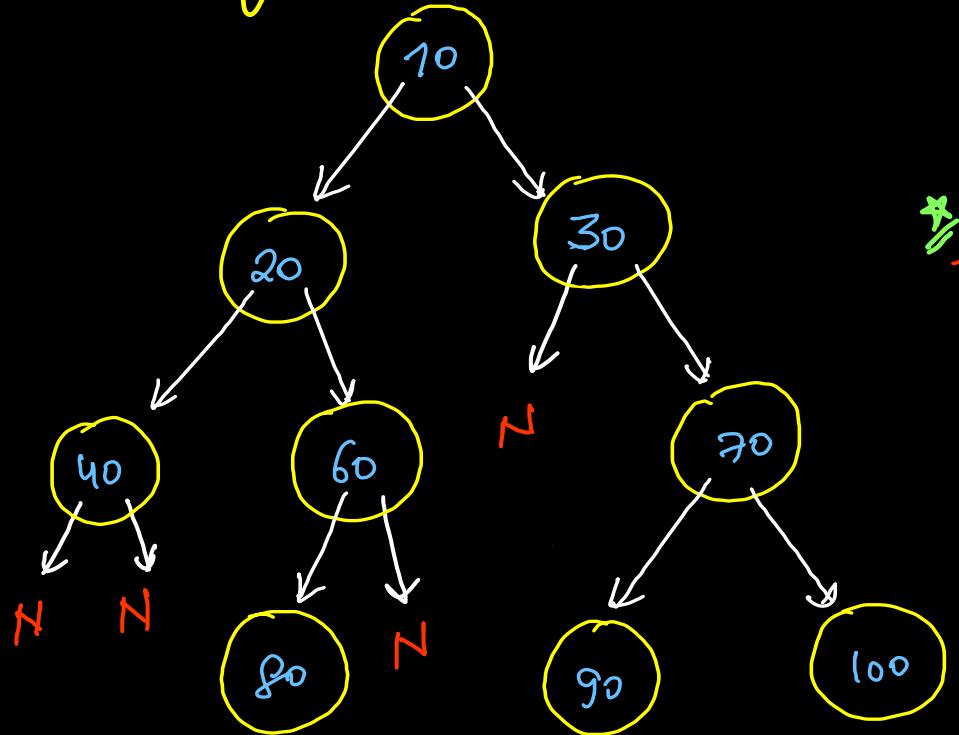


eg github repository management

eg flutter widget tree

eg organization hierarchy

Binary Tree



Single child nodes \rightarrow no siblings
(80, 20)

Single child parent \rightarrow either only left child or only right child exist (30, 70)

root node : 10

- no parent of root node
- entire tree can be traversed

* left child & right child

parent

- every node (except root) will have exactly one parent.

cousins

- nodes at same level but not siblings

descendants / subtree

\hookrightarrow all nodes that can be traversed from the current node

ancestors
 \downarrow
all nodes form current node till root node

siblings
 \hookrightarrow nodes with same parent

leaves
 \hookrightarrow nodes having 0 children
(40, 80, 90, 100)

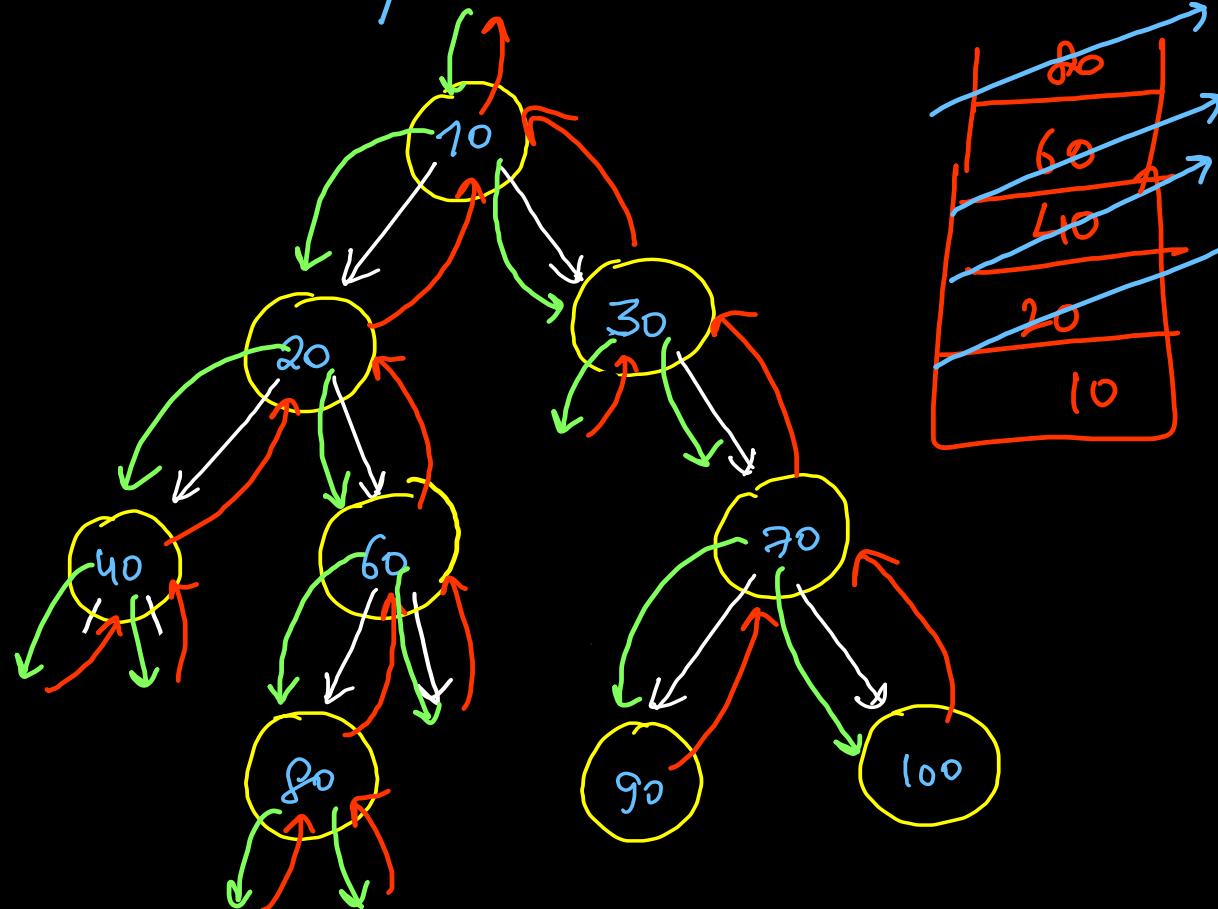
binary tree node

objects will be single node

```
class TreeNode {  
    int val;  
    TreeNode left, right;  
    TreeNode() {}  
    TreeNode(int val) { this.val = val; }  
}
```

```
// Definition for a binary tree node.  
public class TreeNode {  
    int val;  
    TreeNode left;  
    TreeNode right;  
    TreeNode() {}  
    TreeNode(int val) { this.val = val; }  
    TreeNode(int val, TreeNode left, TreeNode right) {  
        this.val = val;  
        this.left = left;  
        this.right = right;  
    }  
}
```

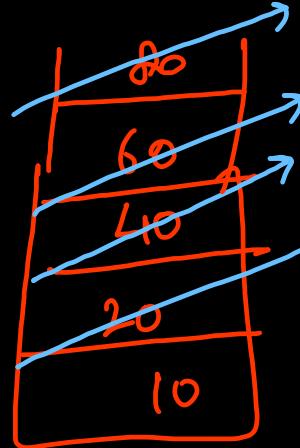
Depth First Search Traversal



preorder: 10, 20, 40, 60, 80, 30, 70, 90, 100

inorder: 40, 20, 80, 60, 10, 30, 90, 70, 100

postorder: 40, 80, 60, 20, 90, 100, 70, 30, 10,



```
public void dfs(TreeNode root)
{
    0. if (root == null) return;
    1. preorder.add(root.val);
    2. dfs(root.left);
    3. inorder.add(root.val);
    4. dfs(root.right);
    5. postorder.add(root.val);
}
```

} base case
meeting expectation (work)
expected
faithful recursive call

```
// Definition for a binary tree node.
public class TreeNode {
    int val;
    TreeNode left;
    TreeNode right;
    TreeNode() {}
    TreeNode(int val) { this.val = val; }
    TreeNode(int val, TreeNode left, TreeNode right) {
        this.val = val;
        this.left = left;
        this.right = right;
    }
}
```

hC 144

```
class Solution {
    List<Integer> preorder = new ArrayList<>();

    // Faith: Root -> Left Subtree -> Right Subtree
    public void dfs(TreeNode root){
        if(root == null) return;
        preorder.add(root.val);
        dfs(root.left);
        dfs(root.right);
    }

    public List<Integer> preorderTraversal(TreeNode root) {
        dfs(root);
        return preorder;
    }
}
```

LC 94

```
class Solution {  
    List<Integer> inorder = new ArrayList<>();  
  
    // Faith: Left Subtree -> Root -> Right Subtree  
    public void dfs(TreeNode root){  
        if(root == null) return;  
        dfs(root.left);  
        inorder.add(root.val);  
        dfs(root.right);  
    }  
  
    public List<Integer> inorderTraversal(TreeNode root) {  
        dfs(root);  
        return inorder;  
    }  
}
```

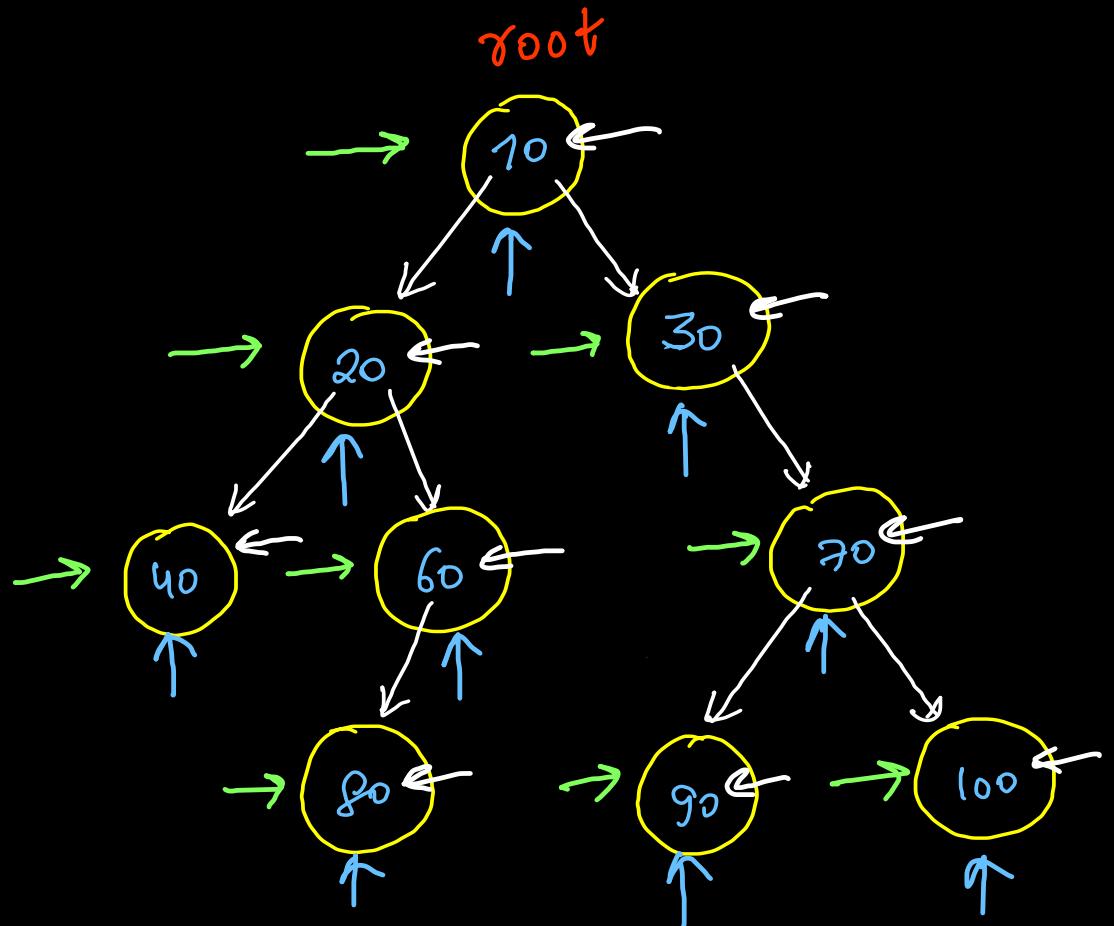
Time Complexity $\Rightarrow O(n)$
linear

Space Complexity
 $\hookrightarrow O(h)$ or $O(d)$

recursion call stack

LC 145

```
class Solution {  
    List<Integer> postorder = new ArrayList<>();  
  
    // Faith: Left Subtree -> Right Subtree -> Root  
    public void dfs(TreeNode root){  
        if(root == null) return;  
        dfs(root.left);  
        dfs(root.right);  
        postorder.add(root.val);  
    }  
  
    public List<Integer> postorderTraversal(TreeNode root) {  
        dfs(root);  
        return postorder;  
    }  
}
```



Preorder (\rightarrow)

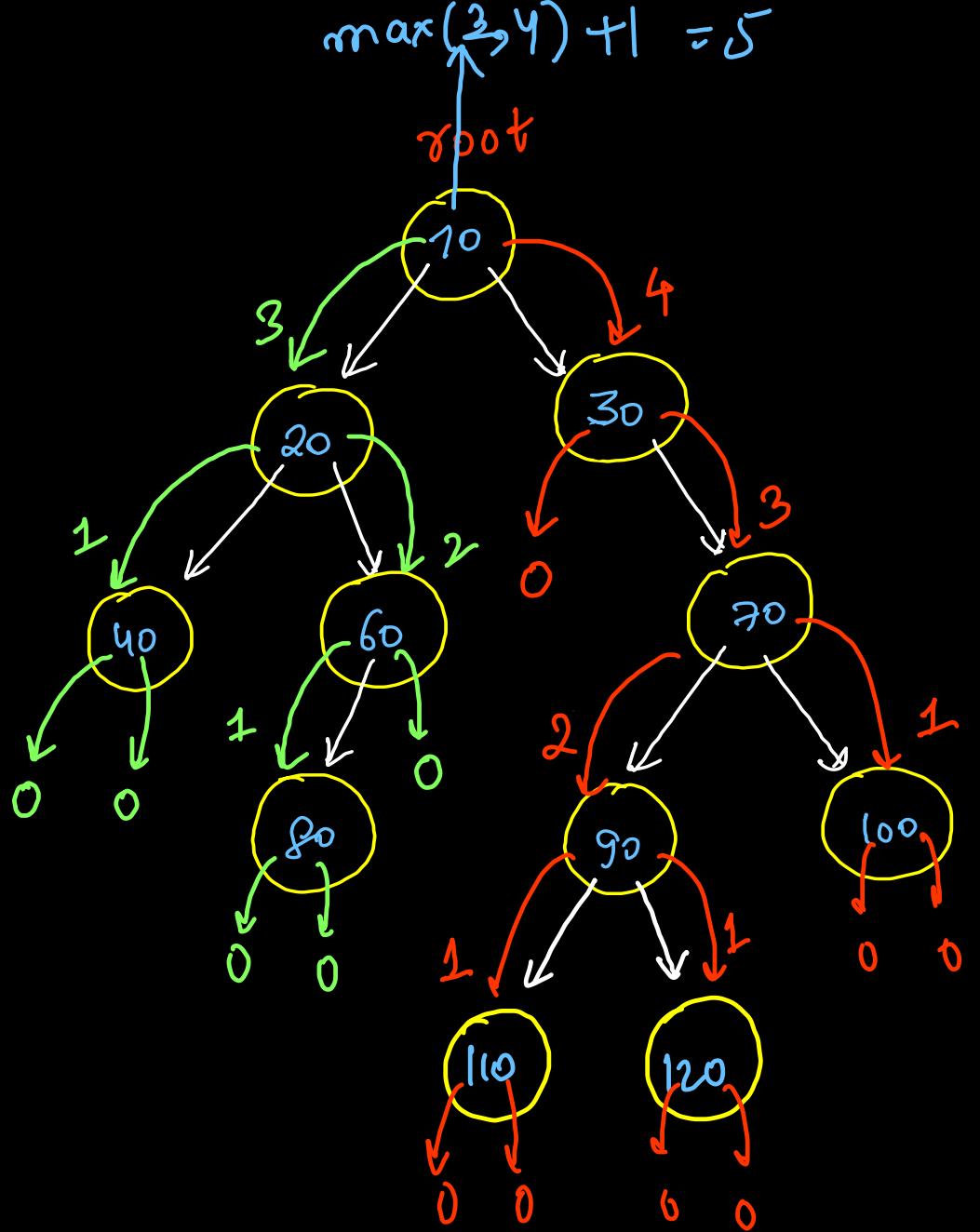
10, 20, 40, 60, 80, 30, 70, 90, 100

Inorder (\uparrow)

40, 20, 80, 60, 10, 30, 90, 70, 10

Postorder (\leftarrow)

40, 80, 60, 20, 90, 100, 70, 30, 10



Height / max^m Depth /
 Longest Path from root
 to farthest leaf node

```

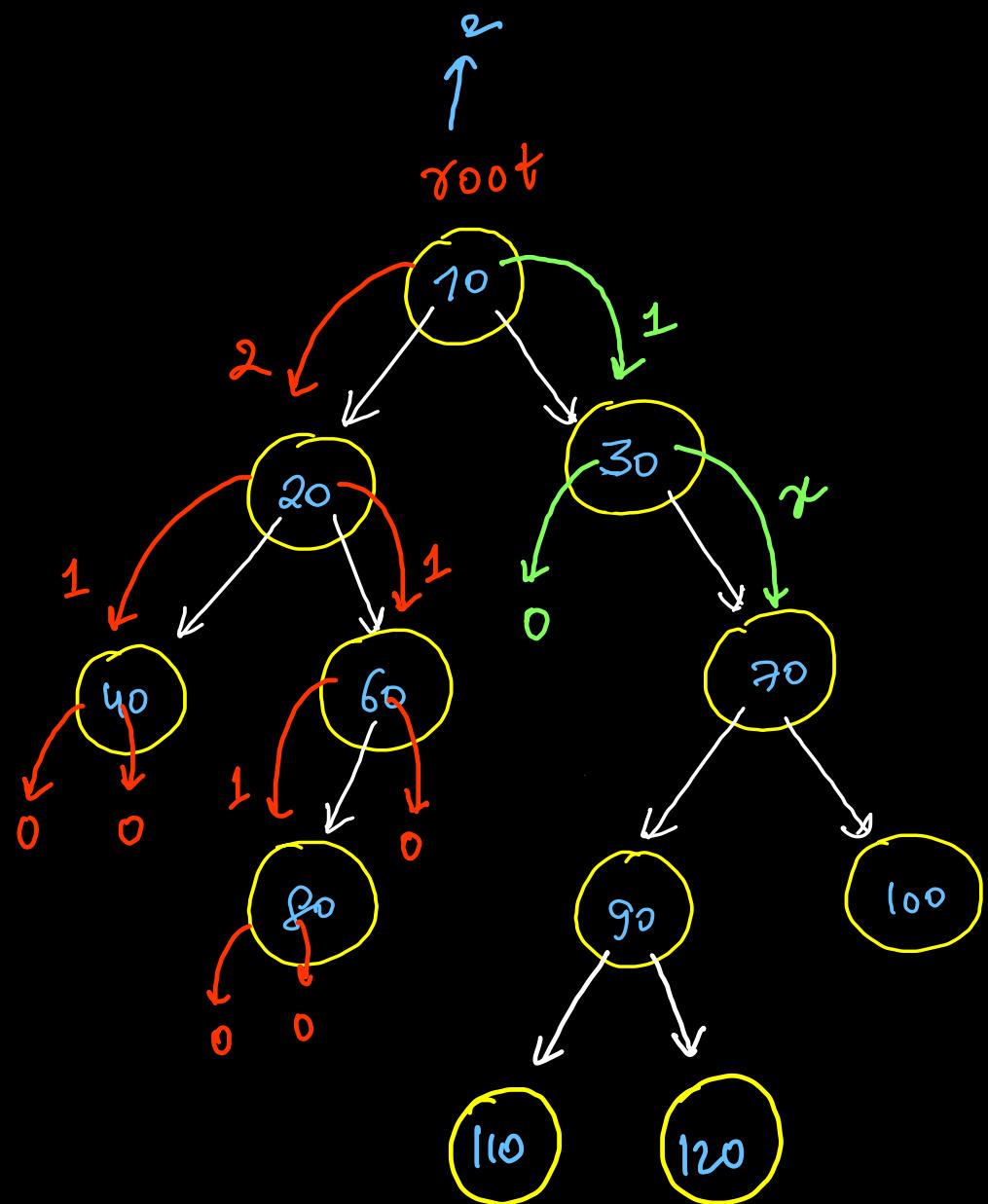
int height(TreeNode root) {
    if (root == null) return 0;
    int lh = height(root.left);
    int rh = height(root.right);
    return max(lh, rh) + 1;
}
  
```

```
class Solution {  
    public int maxDepth(TreeNode root) {  
        if(root == null) return 0;  
        int lh = maxDepth(root.left);  
        int rh = maxDepth(root.right);  
        return Math.max(lh, rh) + 1;  
    }  
}
```

Time $\rightarrow O(n)$
where $n = \text{no of nodes}$

Space $\rightarrow O(h)$
where $h = \text{avg height of tree}$

Leetcode 104

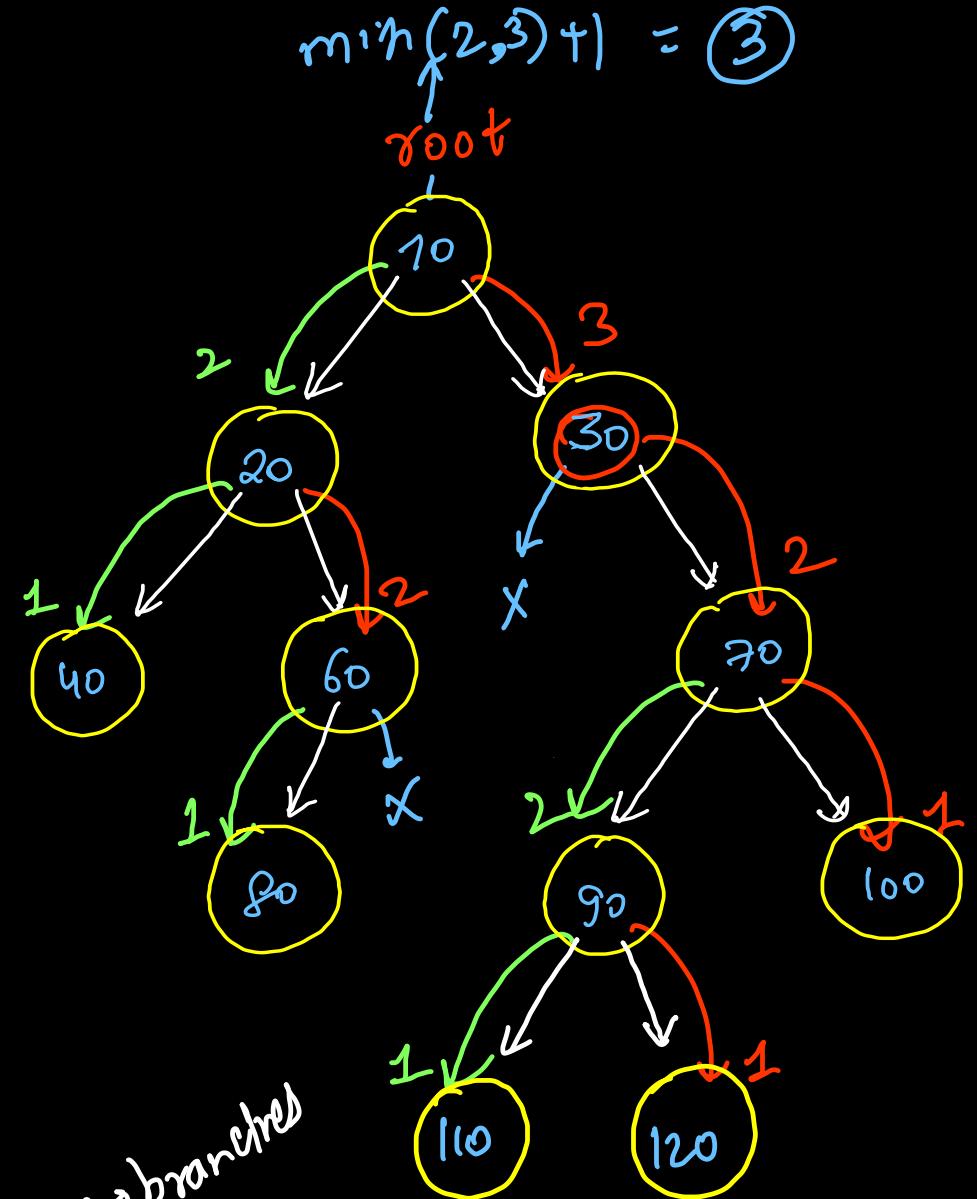


min^m Depth
 ↳ Dist of Root node to
 closest leaf node.

```

int height(TreeNode root) {
    if (root == null) return 0;
    int lh = height(root.left);
    int rh = height(root.right);
    return min(lh, rh) + 1;
  
```

↳ wrong logic



edges → branches
vertices → nodes

```
public int minDepth(TreeNode root) {
    if(root == null) return 0;
    if(root.right == null)
        return minDepth(root.left) + 1;
    if(root.left == null)
        return minDepth(root.right) + 1;
    int lh = minDepth(root.left);
    int rh = minDepth(root.right);
    return Math.min(lh, rh) + 1;
}
```

single child parent corner cases

Approach 1)

```
public int minDepth(TreeNode root) {  
    if(root == null) return 0;  
    if(root.right == null)  
        return minDepth(root.left) + 1;  
    if(root.left == null)  
        return minDepth(root.right) + 1;  
    int lh = minDepth(root.left);  
    int rh = minDepth(root.right);  
    return Math.min(lh, rh) + 1;  
}
```

preorder

Approach 2)

```
class Solution {  
    public int minDepth(TreeNode root) {  
        if(root == null) return 0;  
        int lh = minDepth(root.left);  
        int rh = minDepth(root.right);  
        if(lh == 0) return rh + 1;  
        if(rh == 0) return lh + 1;  
        return Math.min(lh, rh) + 1;  
    }  
}
```

postorder

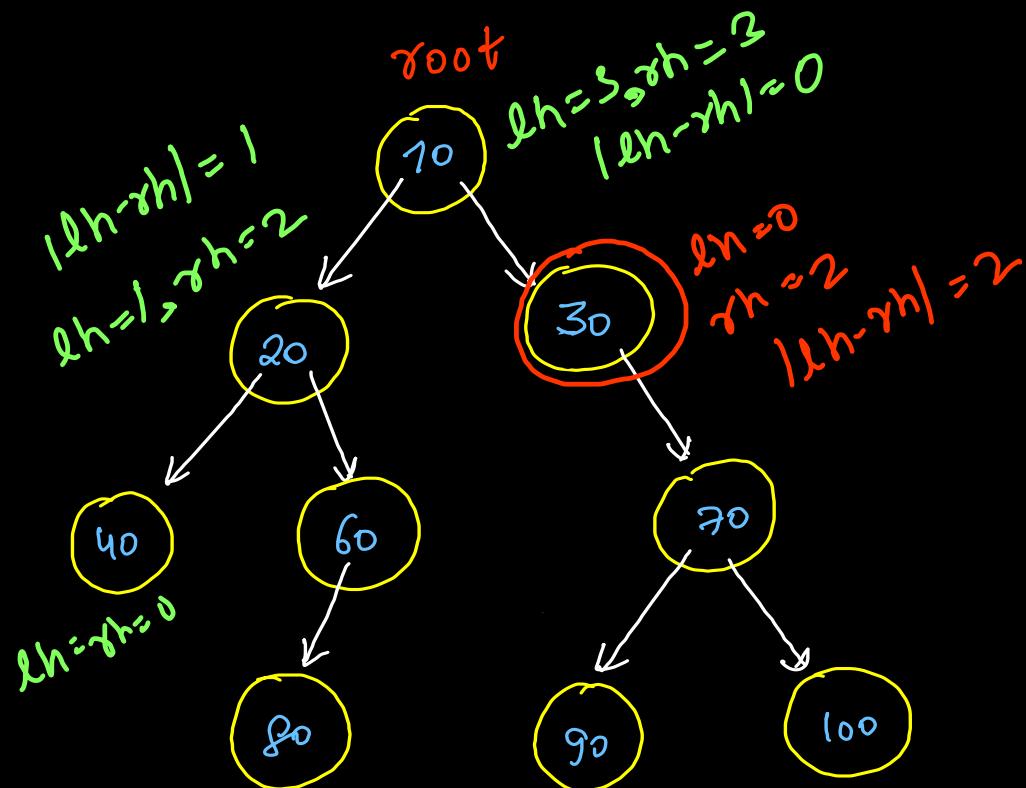
Time = $O(n)$

$O(\log n)$ avg case

Space = $O(h)$

$O(n)$ worst case

DP on Trees (Re-roooting Technique)



Leetcode 710 Balanced Binary Tree

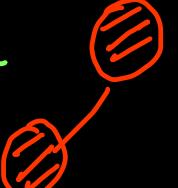
Each node

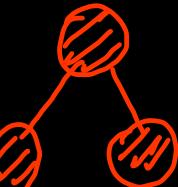
$$\Rightarrow | \text{leftheight} - \text{rightheight} | \leq 1$$

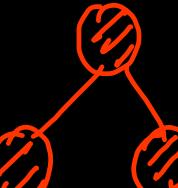
Binary Tree \rightarrow BST
(height)
{balanced}

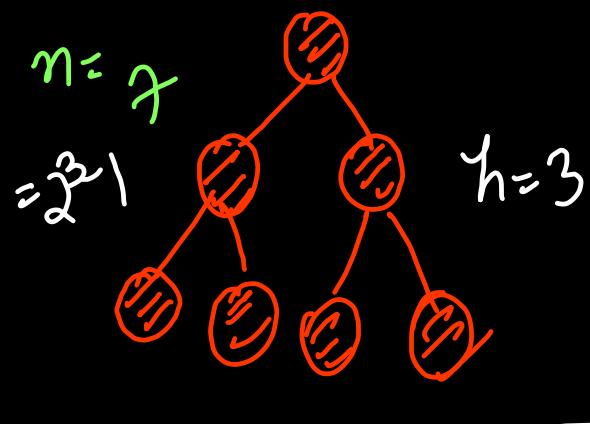
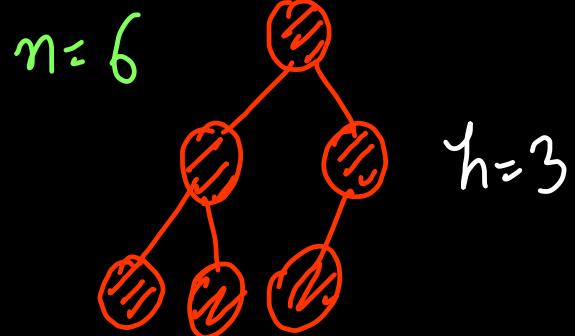
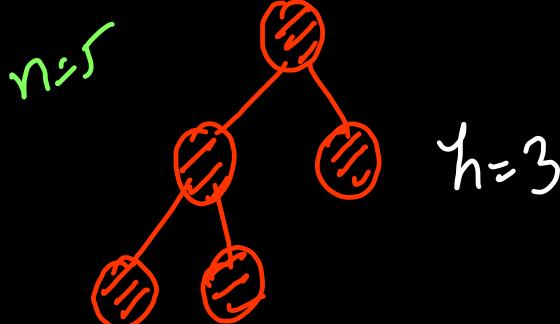
min height
balanced tree
 n nodes

$n=1$  $h=1$

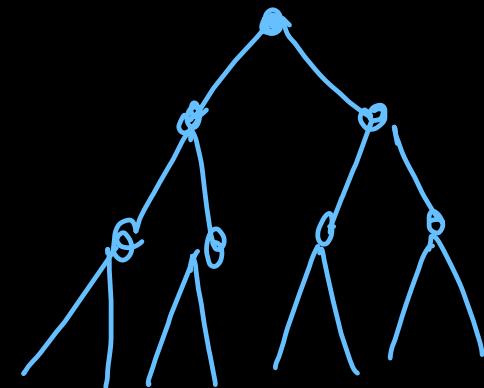
$n=2$  $h=2$

$n=3$
 $=2^2 - 1$  $h=2$

$n=4$  $h=3$



$n=8 \longrightarrow 15$ $2^4 - 1$
 $h=4$



nodes = n
min height
 $= \underline{\log_2 n}$

best case

\max^m height
Skewed tree

$$n=1 \quad h=1$$

$$n=2 \quad \begin{array}{c} \bullet \\ \swarrow \quad \searrow \end{array} \quad h=2$$

Worst case

$$\text{nodes} = n$$

$$n=3 \quad \begin{array}{c} \bullet \\ \swarrow \quad \searrow \\ \bullet \quad \bullet \end{array} \quad h=3$$

$$\text{height} = n$$

$$n=4 \quad \begin{array}{c} \bullet \\ \swarrow \quad \searrow \\ \bullet \quad \begin{array}{c} \bullet \\ \swarrow \quad \searrow \end{array} \end{array} \quad h=4$$

Approach 1) *Brute force*

```

public int height(TreeNode root){
    if(root == null) return 0;

    int lh = height(root.left);
    int rh = height(root.right);

    return Math.max(lh, rh) + 1;
}

public boolean isBalanced(TreeNode root) {
    if(root == null) return true;

    int lh = height(root.left);
    int rh = height(root.right);

    if(Math.abs(lh - rh) > 1) return false;
    return isBalanced(root.left) && isBalanced(root.right);
}

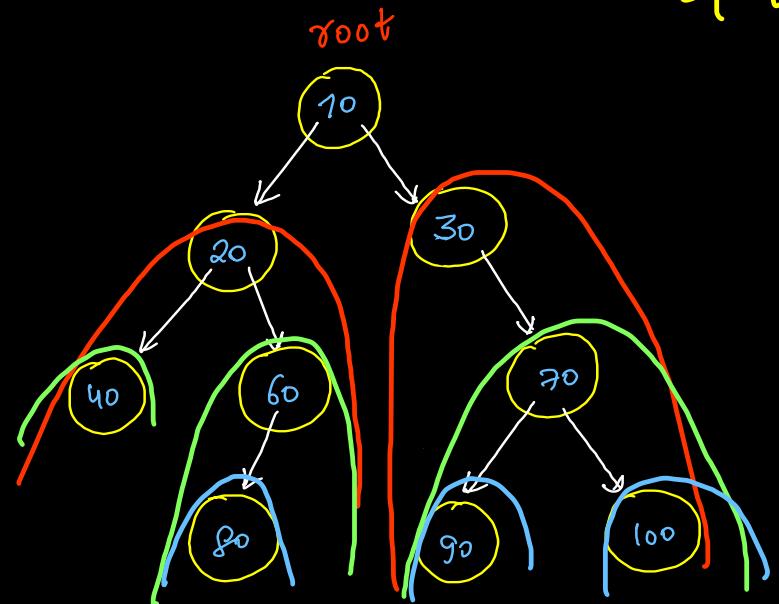
```

dfs $\approx O(n)$

$O(n^2)$

dfs $\approx O(n)$

Time = $O(n^2)$, Space = $O(h)$



$$\begin{aligned}
& 1 \cdot 1 + 2 \cdot 2 + 2^2 \cdot 3 \\
& + 2^3 \cdot 4 + \dots + 2^{\log n} \cdot \log n \\
& = \underline{\underline{n^2}}
\end{aligned}$$

```

class Solution {
    boolean balanced = true; (global)

    public int height(TreeNode root){
        if(root == null) return 0;

        int lh = height(root.left);
        int rh = height(root.right);

        if(Math.abs(lh - rh) > 1) balanced = false;
        return Math.max(lh, rh) + 1;
    }

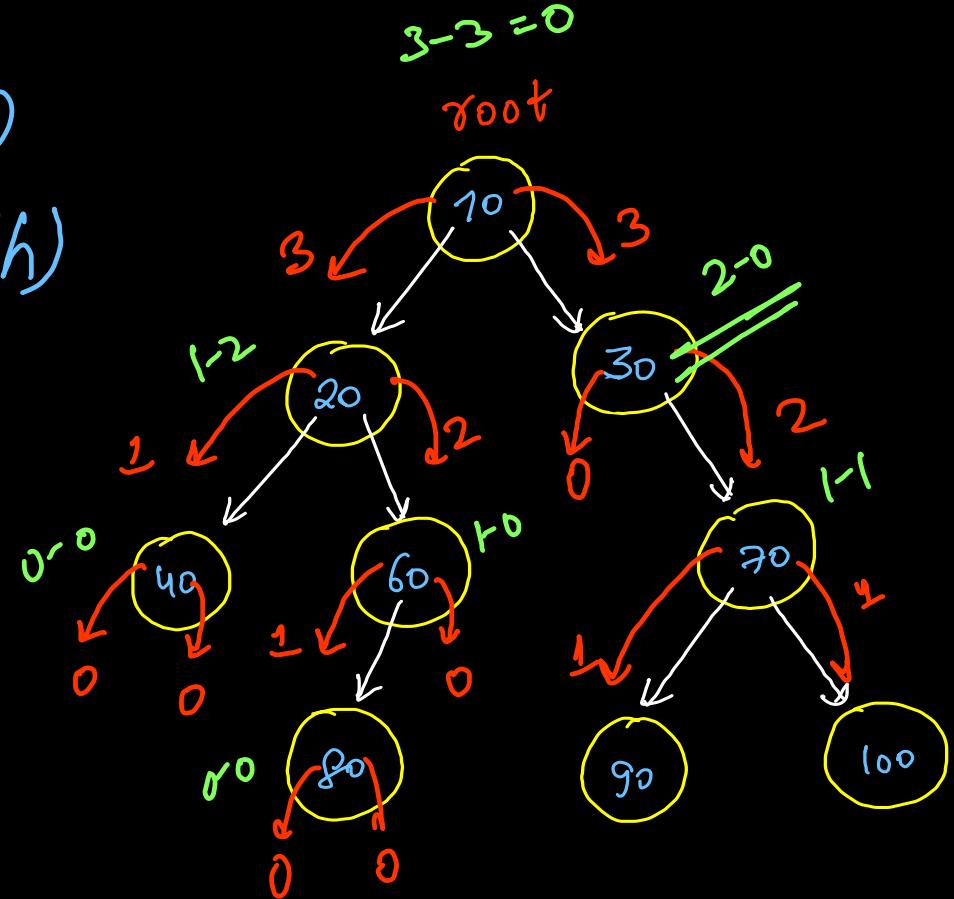
    public boolean isBalanced(TreeNode root) {
        height(root);
        return balanced;
    }
}

```

Time $\Rightarrow \Theta(n)$

Space $\Rightarrow O(h)$

dfs



Approach 2) Global variable

"global variables are considered evil/bad"

~~balanced = true~~ \Rightarrow false

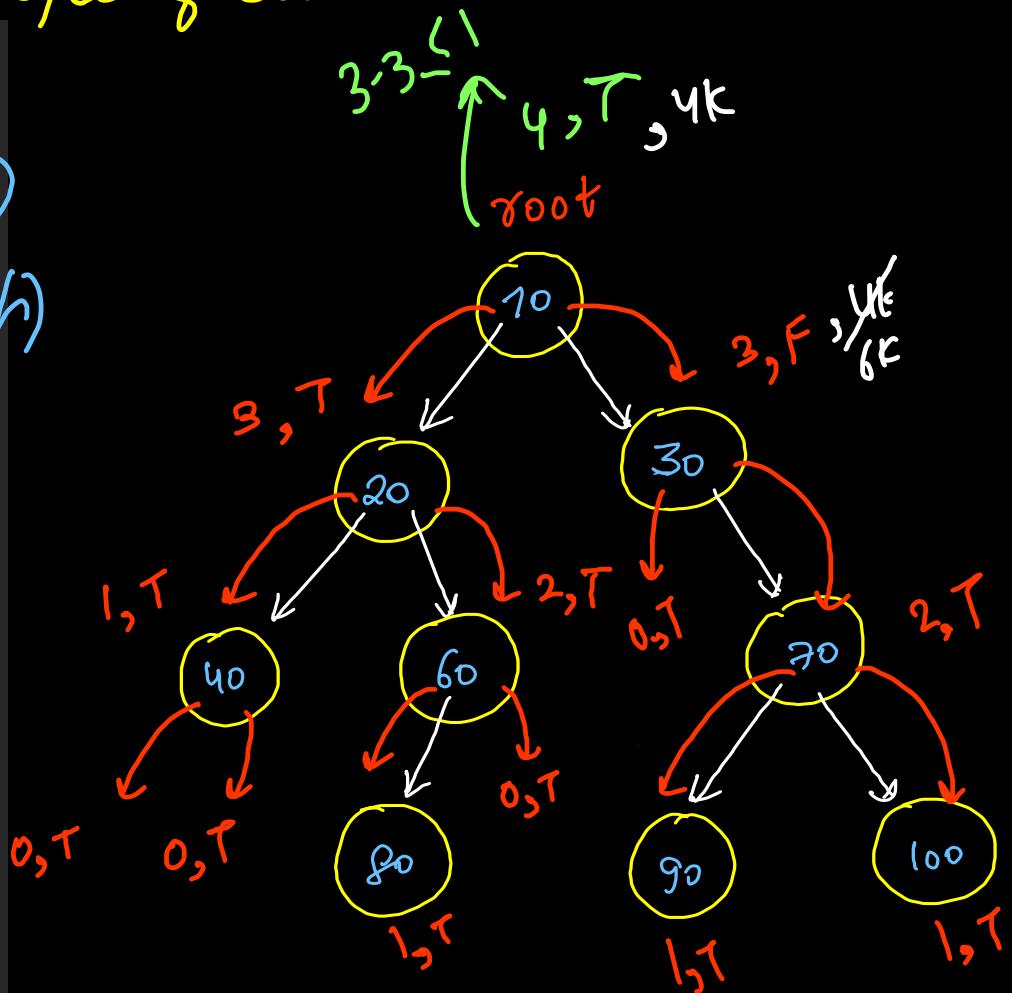
member inner class → object creation w/o object of Solution class

```
class Solution {  
    public static class Pair{  
        int height = 0;  
        boolean balanced = true;  
    }  
}
```

Approach 3)

Time = $O(n)$

Space = $O(h)$



```
public Pair dfs(TreeNode root){  
    if(root == null) return new Pair();  
  
    Pair l = dfs(root.left);  
    Pair r = dfs(root.right);  
  
    Pair curr = new Pair();  
    curr.height = Math.max(l.height, r.height) + 1;  
    curr.balanced = l.balanced && r.balanced;  
    if(Math.abs(l.height - r.height) > 1) curr.balanced = false;  
    return curr;  
}  
  
public boolean isBalanced(TreeNode root) {  
    return dfs(root).balanced;  
}
```

```
class Solution {  
    public int dfs(TreeNode root, boolean[] balanced){  
        if(root == null) return 0;  
        int lh = dfs(root.left, balanced);  
        int rh = dfs(root.right, balanced);  
  
        if(Math.abs(lh - rh) > 1) balanced[0] = false;  
        return Math.max(lh, rh) + 1;  
    }  
  
    public boolean isBalanced(TreeNode root) {  
        boolean[] balanced = {true};  
        dfs(root, balanced);  
        return balanced[0];  
    }  
}
```

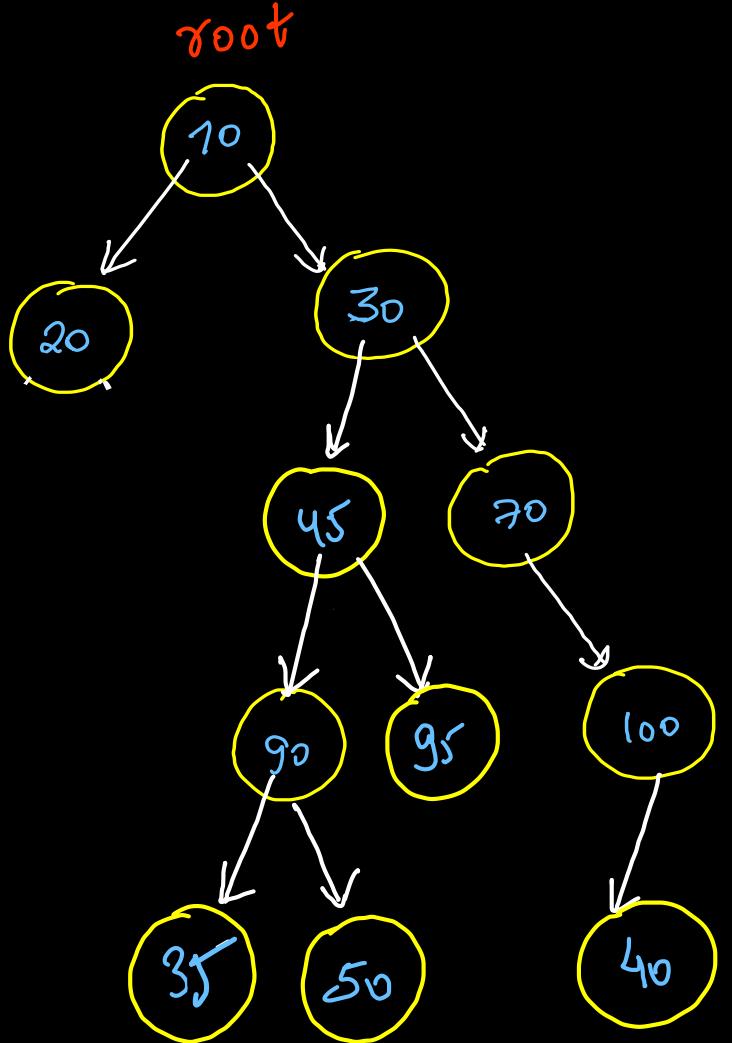
Approach 4)

Passing balanced
variable in
array parameter

LC 543

Diameter of Binary Tree

longest node to node distance



$$\text{diameter} = lh + rh + 1$$

(in terms of nodes)

```

class Solution {
    int diameter = 0;

    public int height(TreeNode root){
        if(root == null) return 0;

        int lh = height(root.left);
        int rh = height(root.right);

        diameter = Math.max(diameter, lh + rh + 1);
        return Math.max(lh, rh) + 1;
    }

    public int diameterOfBinaryTree(TreeNode root) {
        if(root == null) return 0; → empty tree
if(root.left == null & root.right == null) return 0;
        height(root);
        return diameter - 1;
    }
}

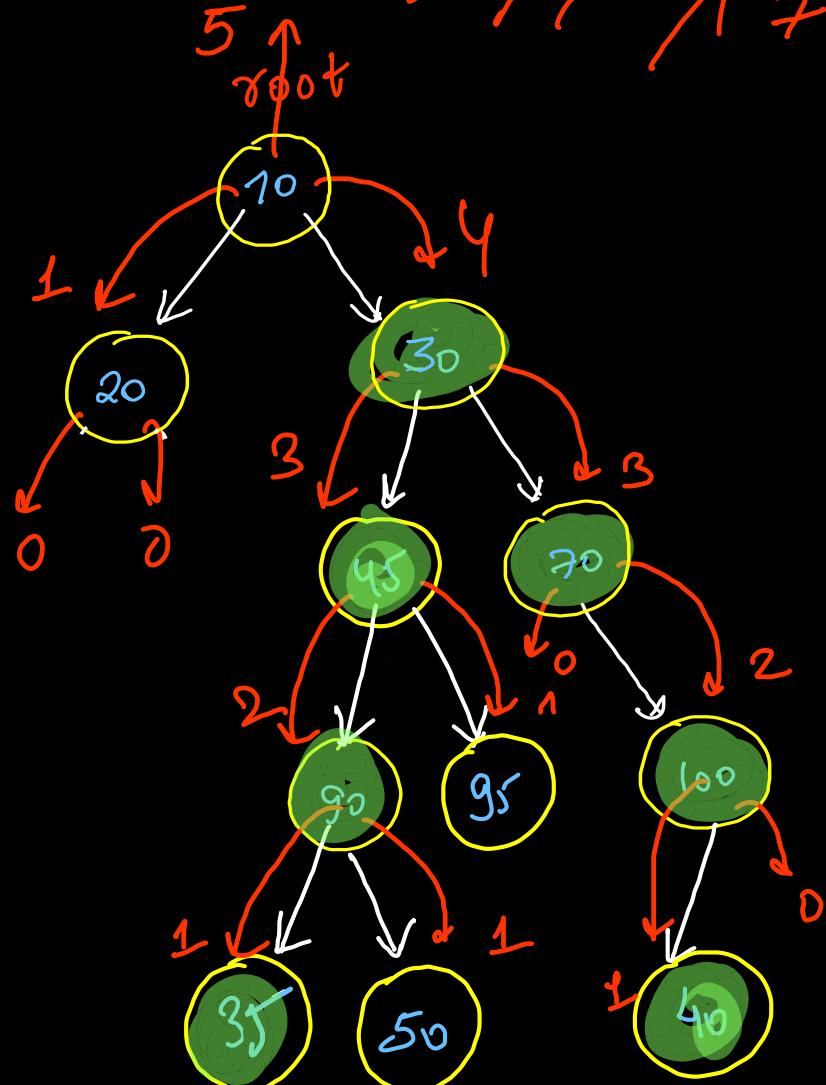
```

T_{number of edges}

$$\text{Time} = O(n)$$

$$\text{Space} = O(h)$$

diameter = ~~O(n)~~ $\beta \gamma_1 \gamma_2$



```
class Solution {
    public int height(TreeNode root, int[] diameter){
        if(root == null) return 0;

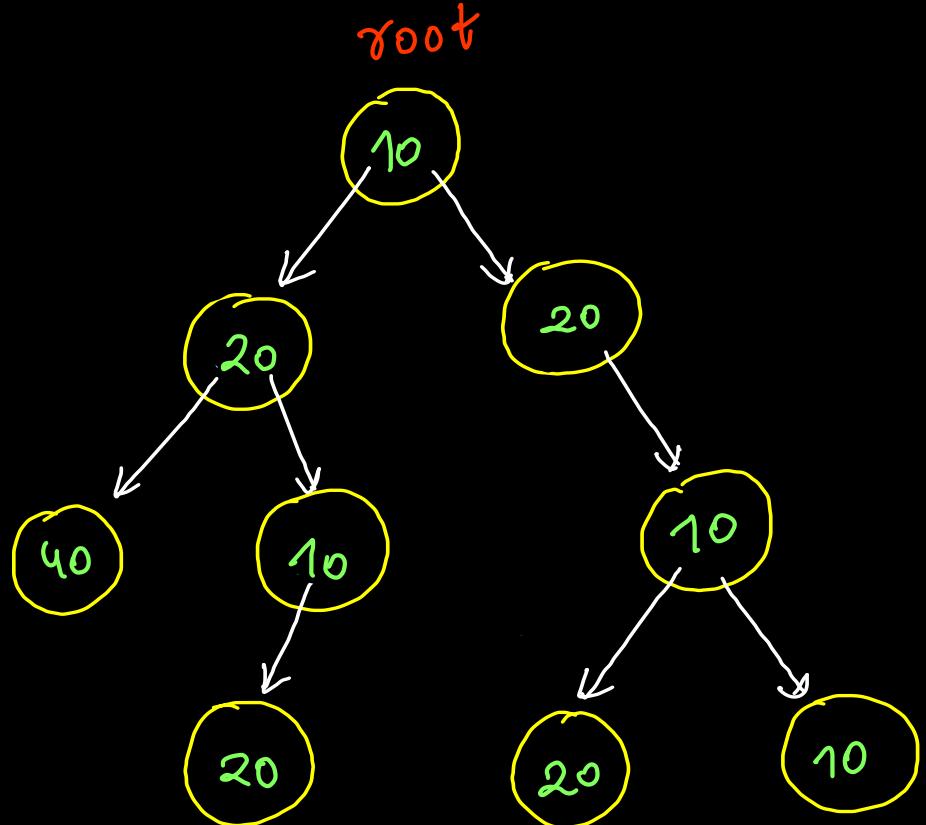
        int lh = height(root.left, diameter);
        int rh = height(root.right, diameter);

        diameter[0] = Math.max(diameter[0], lh + rh + 1);
        return Math.max(lh, rh) + 1;
    }

    public int diameterOfBinaryTree(TreeNode root) {
        if(root == null) return 0;
        int[] diameter = {0};
        height(root, diameter);
        return diameter[0] - 1;
    }
}
```

lc 1367)

linked list in Binary Tree

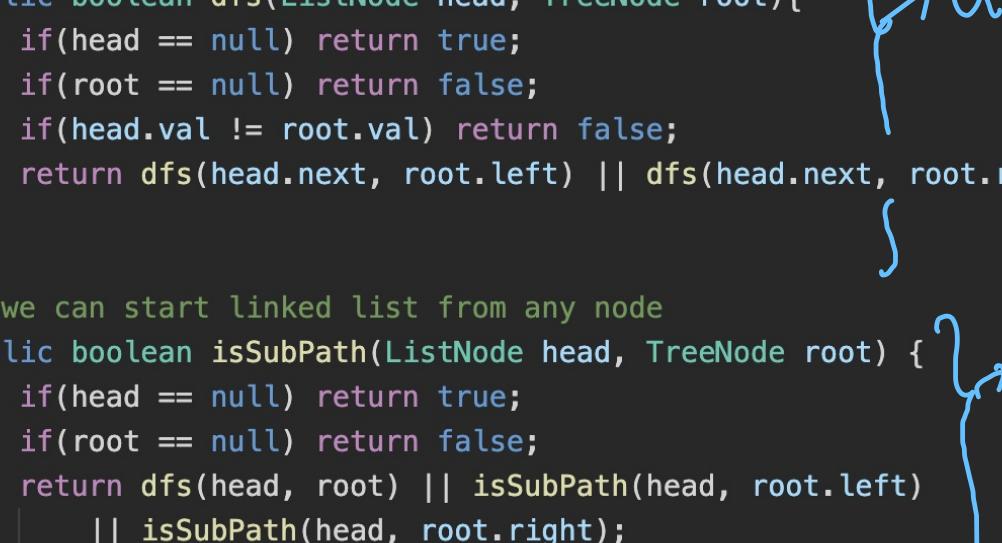


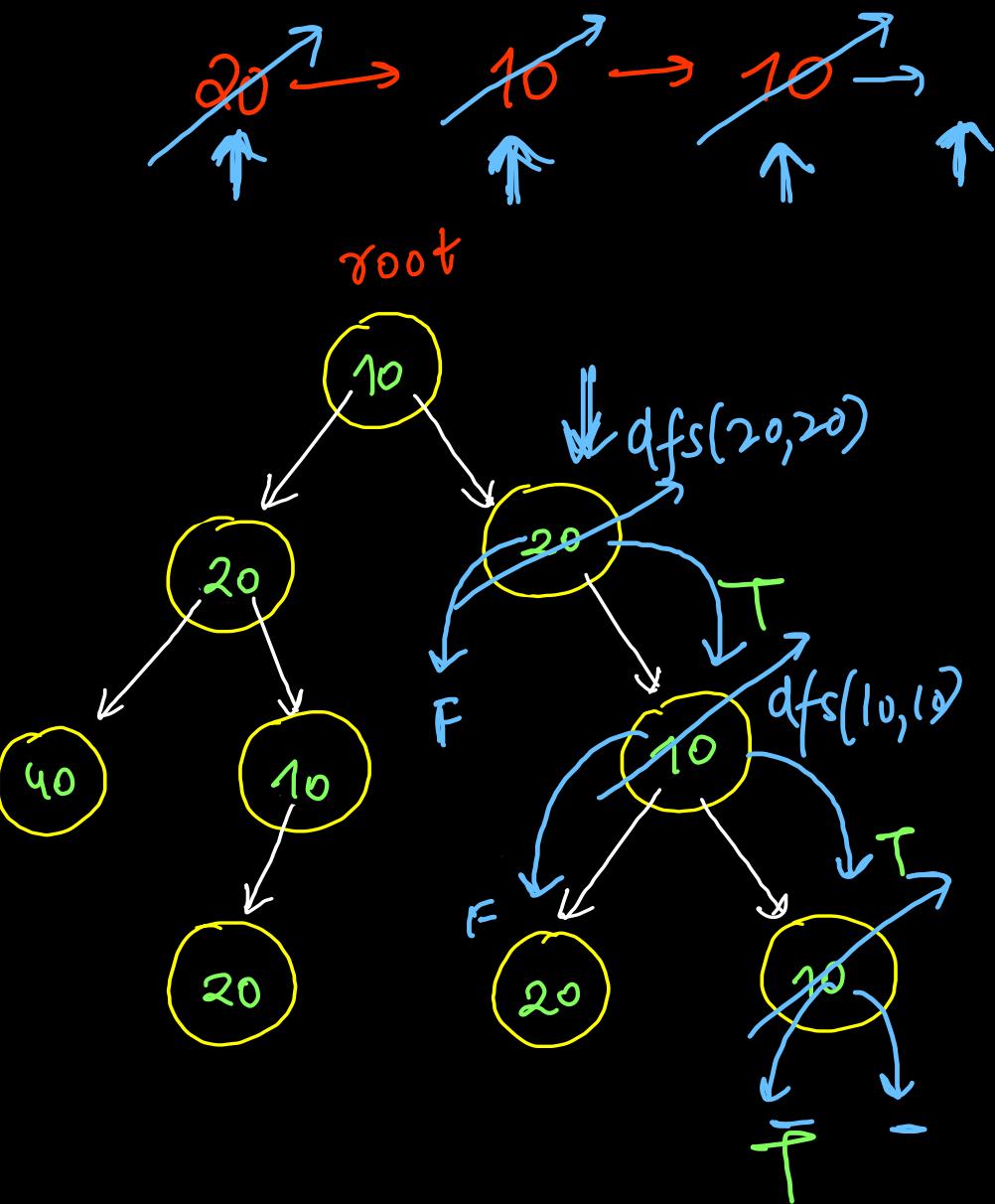
" 20 → 10 → 10 "

↑
head

```
dfs(TreeNode root, ListNode head){  
    if (head == null) return true;  
    if (root == null) return false;  
    if (root.val != head.data)  
        return false;  
    return dfs(root.left, head.next)  
        && dfs(root.right, head.next);  
}
```

```
class Solution {  
    // fixing head node as the root node  
    public boolean dfs(ListNode head, TreeNode root){  
        if(head == null) return true;  
        if(root == null) return false;  
        if(head.val != root.val) return false;  
        return dfs(head.next, root.left) || dfs(head.next, root.right);  
    }  
  
    // we can start linked list from any node  
    public boolean isSubPath(ListNode head, TreeNode root) {  
        if(head == null) return true;  
        if(root == null) return false;  
        return dfs(head, root) || isSubPath(head, root.left)  
            || isSubPath(head, root.right);  
    }  
}
```

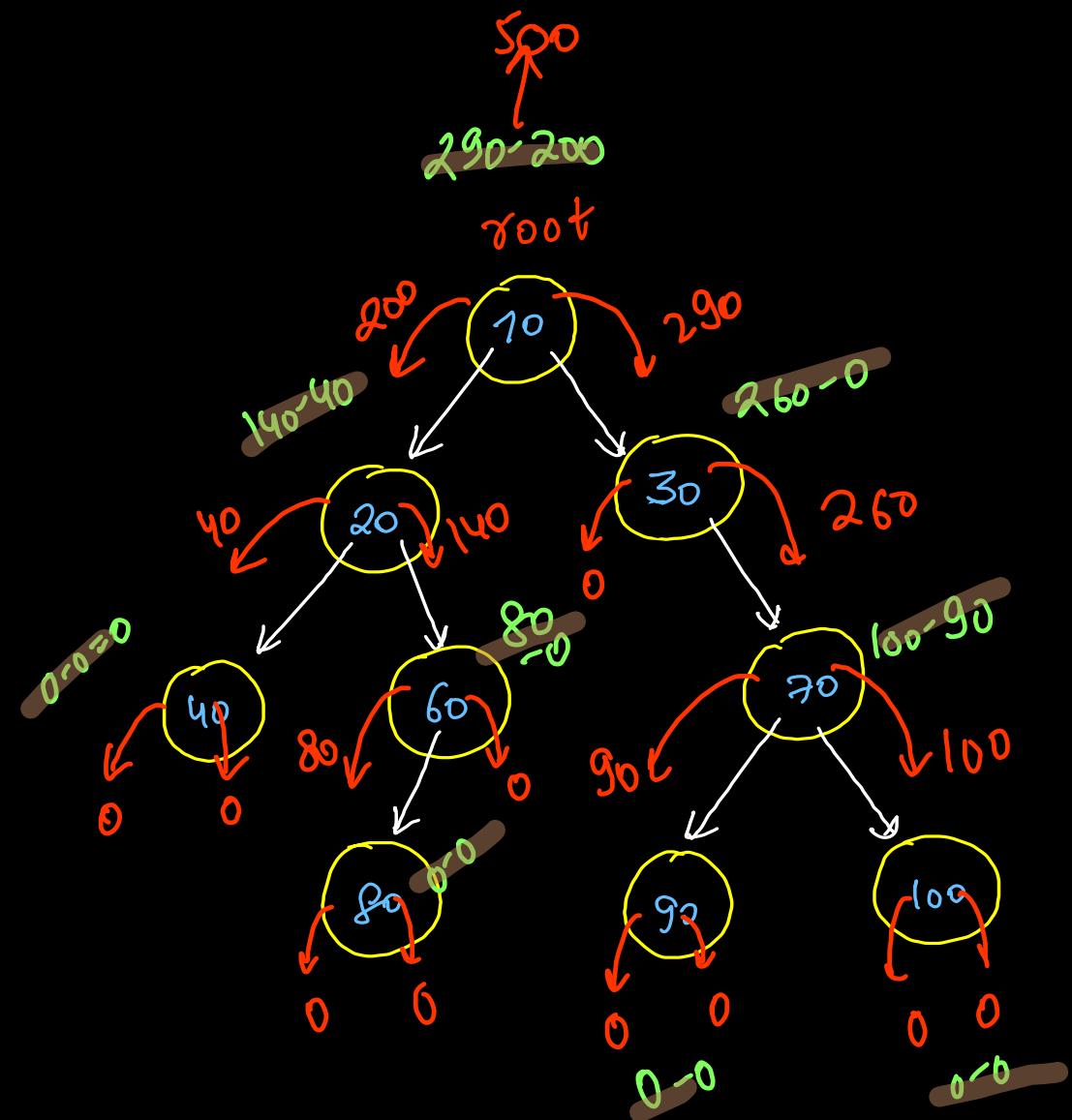


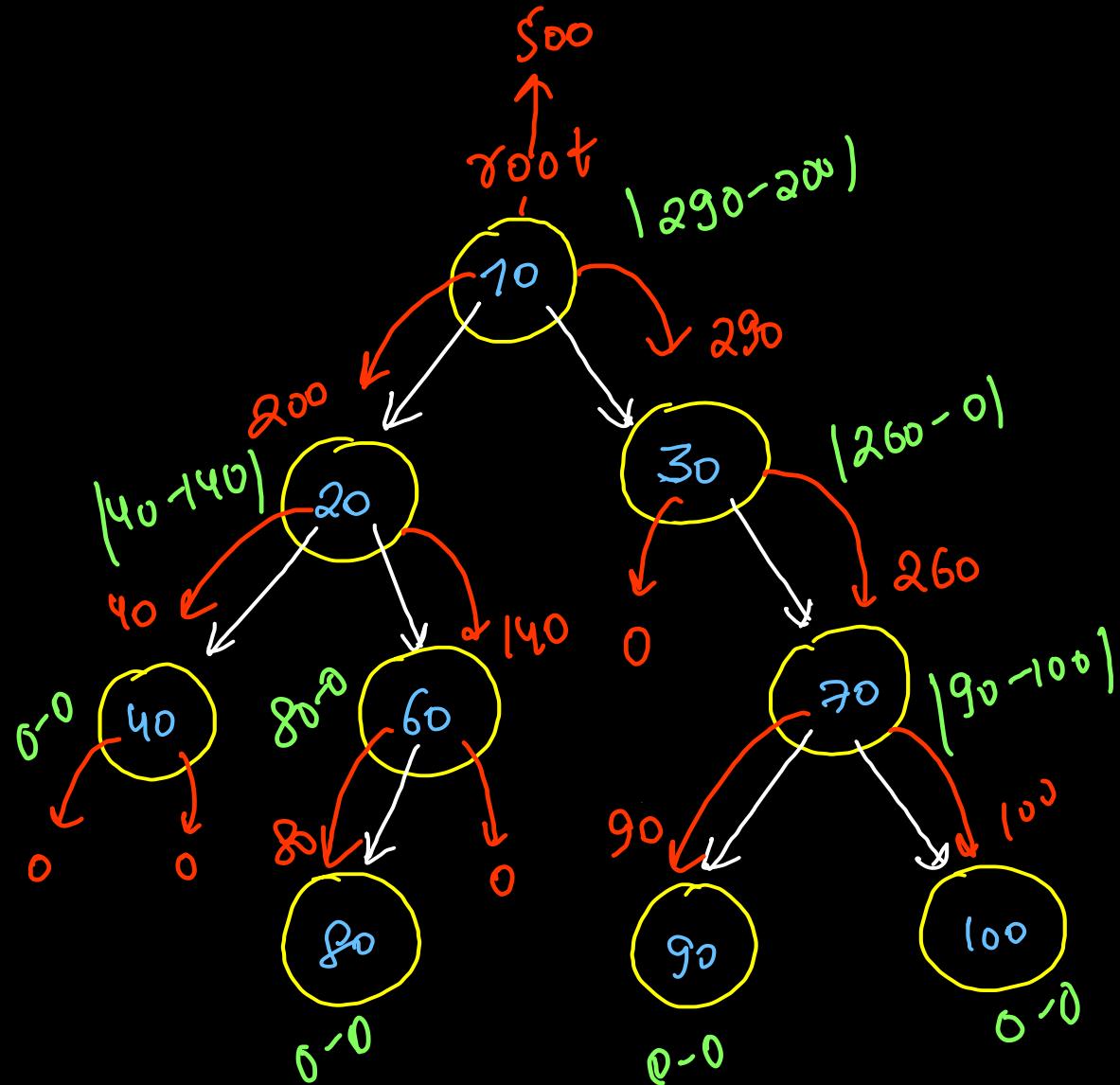


Binary Tree Tilts

Sum of all tilts

$$\Rightarrow | \text{leftSum} - \text{rightSum} |$$





Subtree sum

$\Rightarrow \text{leftsubtree sum}$

+

rightsubtree sum

+

$2001 \cdot \text{val}$

tilt

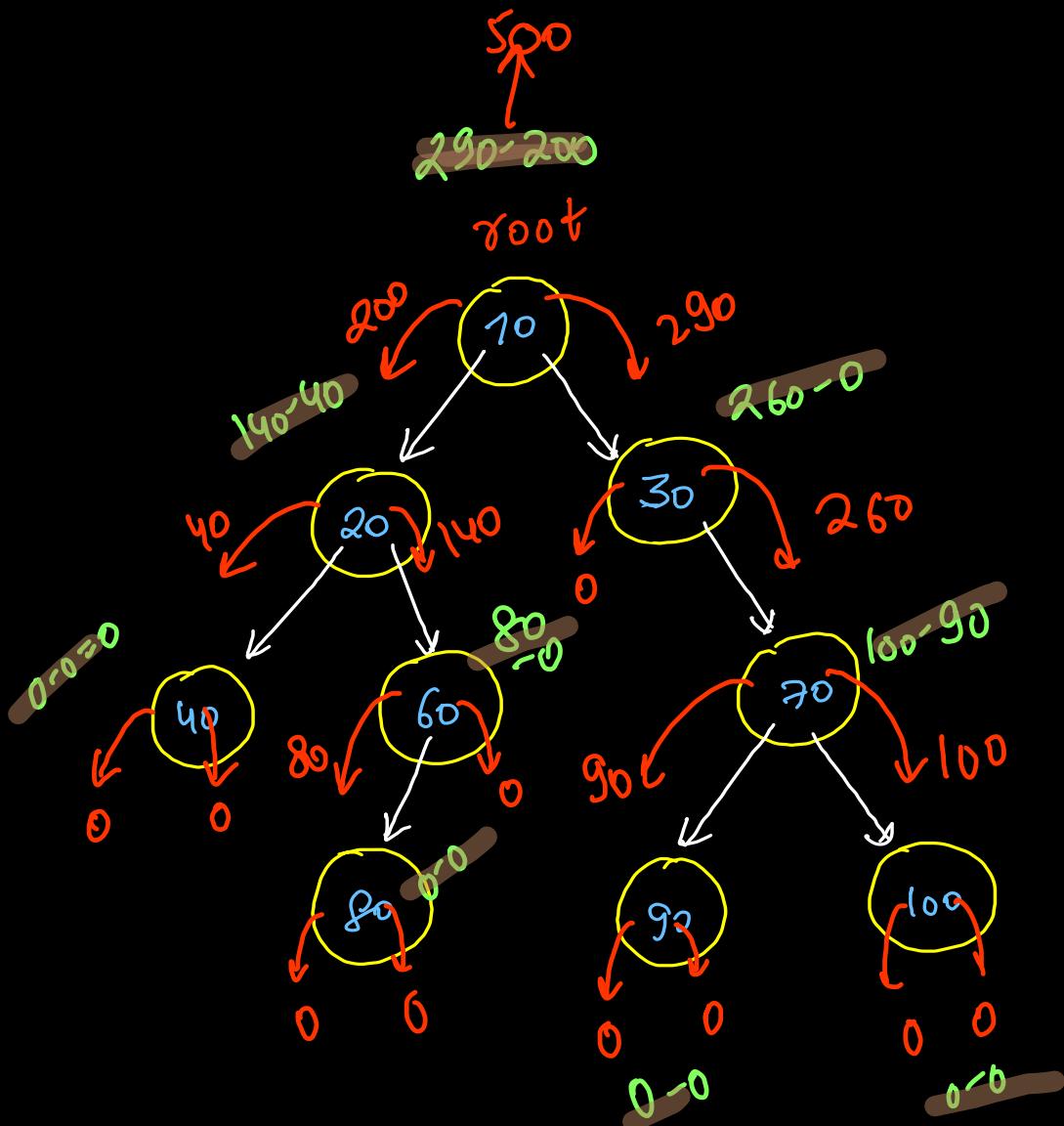
$(\text{leftsum} - \text{rightsum})$

-

$\text{rightsum})$

Sum of all tilts

tilt = 0



Time = $O(n)$, Space = $O(n)$

```
class Solution {
    int ans = 0; → green value

    public int subtreeSum(TreeNode root){
        if(root == null) return 0;

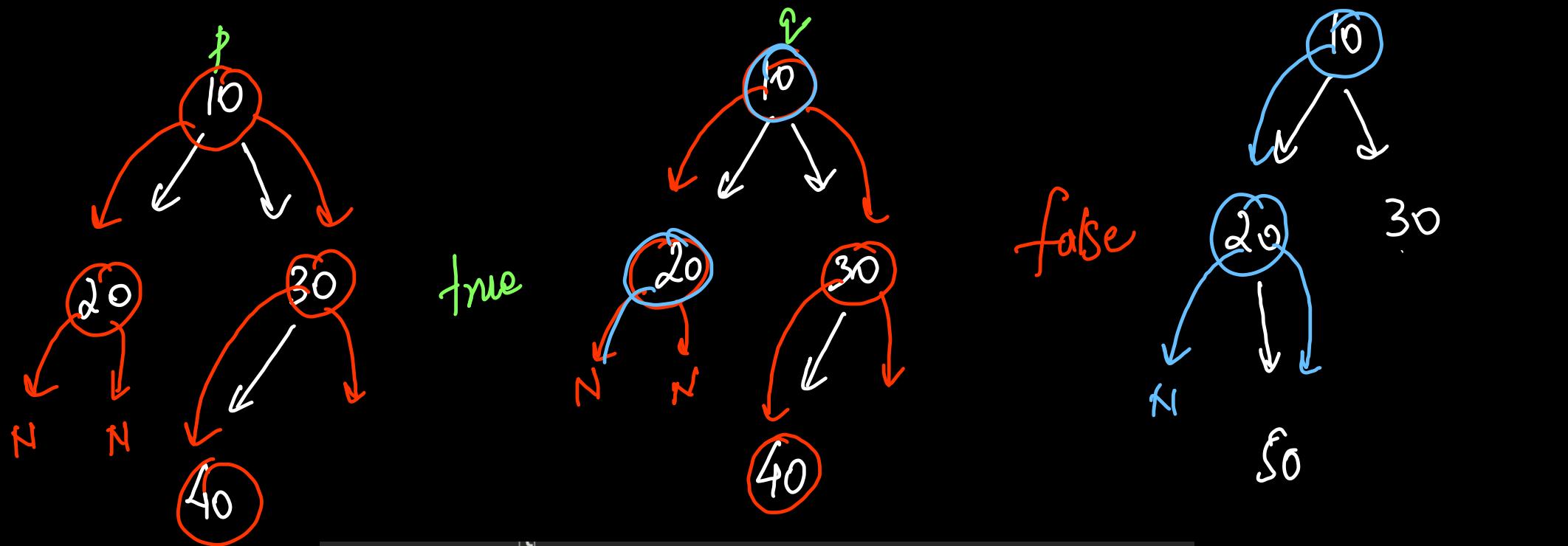
        int lsum = subtreeSum(root.left);
        int rsum = subtreeSum(root.right);

        int tilt = Math.abs(lsum - rsum); } red value
        ans += tilt;

        return lsum + rsum + root.val;
    }

    public int findTilt(TreeNode root) {
        subtreeSum(root);
        return ans;
    }
}
```

LC 100) Same Tree

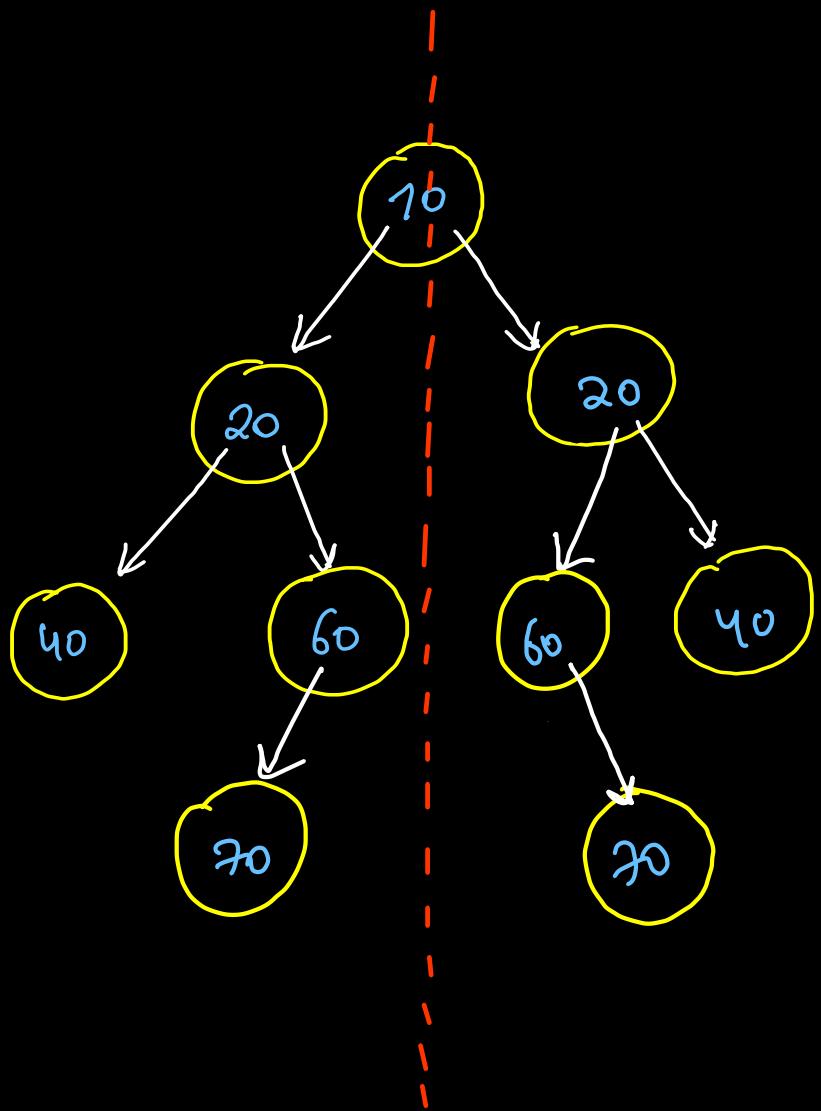


```
public boolean isSameTree(TreeNode p, TreeNode q) {  
    if(p == null && q == null) return true;  
    if(p == null || q == null) return false;  
    if(p.val != q.val) return false;  
  
    return isSameTree(p.left, q.left)  
        && isSameTree(p.right, q.right);  
}
```

Time $\Rightarrow \Theta(n)$
Space $\Rightarrow O(n)$

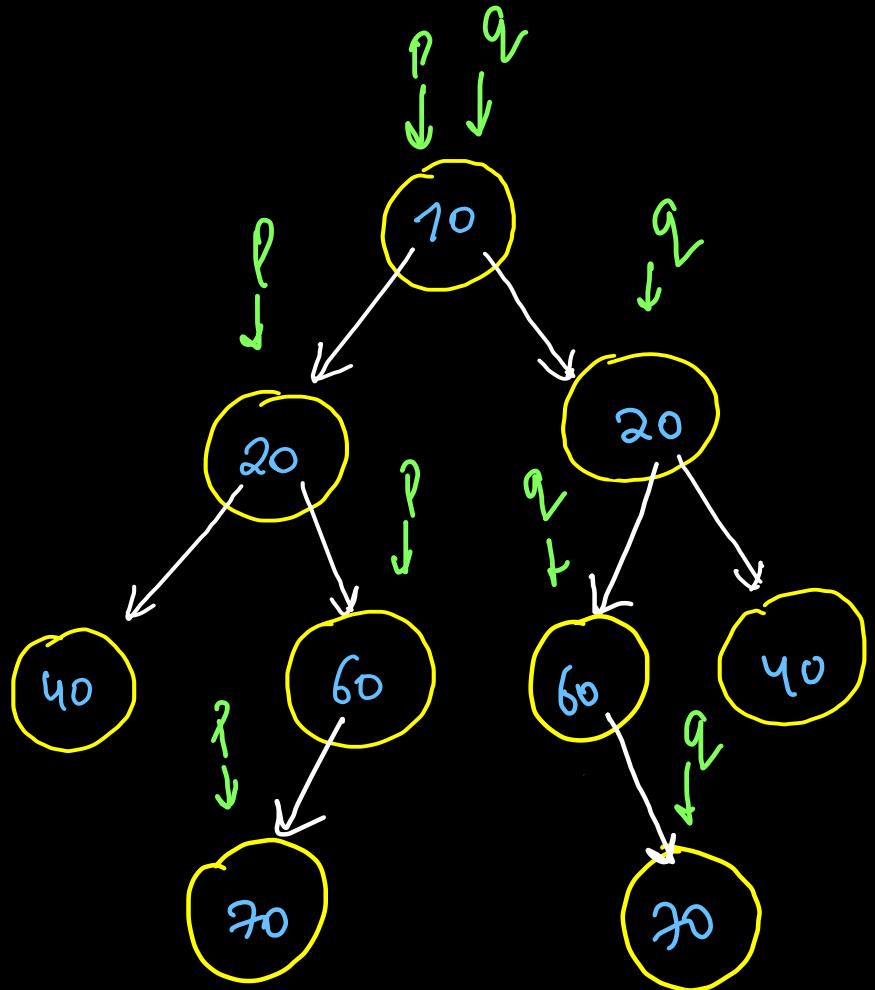
LC 101

Symmetric Tree



isMirror(TreeNode p, TreeNode q)

```
{  
    if (p == null && q == null)  
        return true;  
    if (p == null || q == null)  
        return false;  
    if (p.val != q.val) return false;  
    return isMirror(p.left, q.right)  
        && isMirror(p.right, q.left);  
}
```



```

class Solution {
    public boolean isMirror(TreeNode p, TreeNode q){
        if(p == null && q == null) return true;
        if(p == null || q == null) return false;
        if(p.val != q.val) return false;

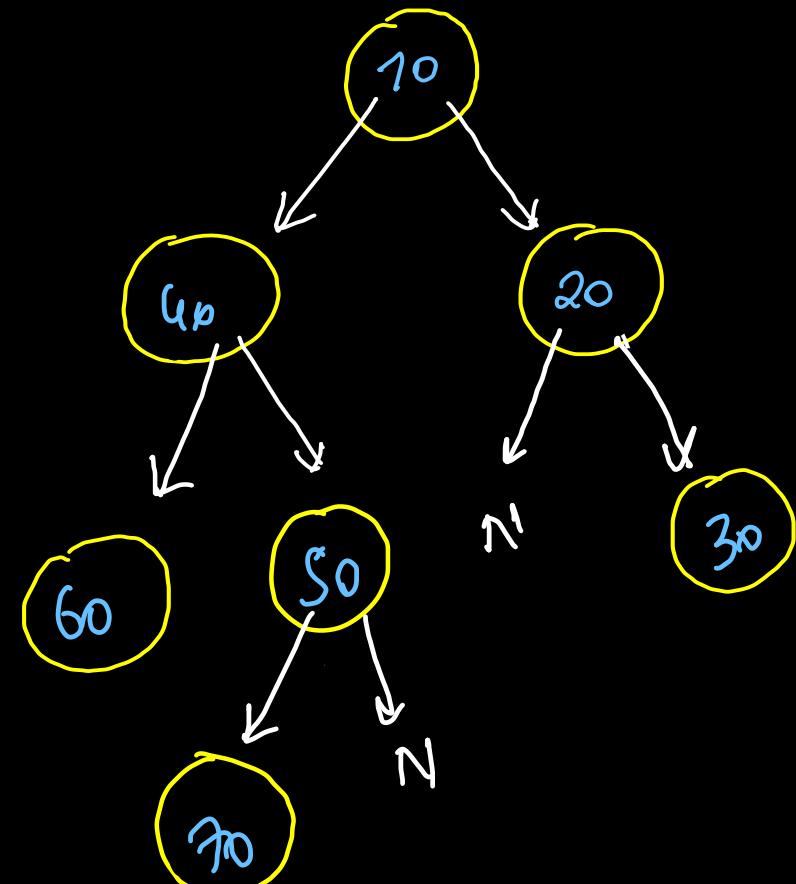
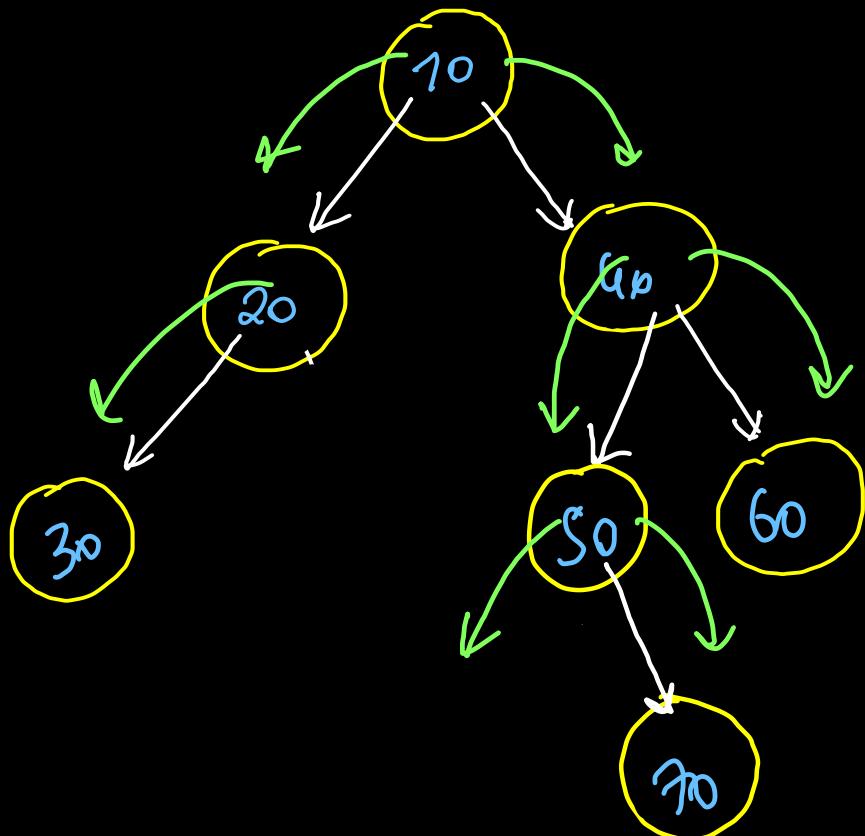
        return isMirror(p.left, q.right)
            && isMirror(p.right, q.left);
    }

    public boolean isSymmetric(TreeNode root) {
        return isMirror(root, root);
    }
}

```

↳ Time = $\Theta(n)$, Space = $\Theta(n)$

Invert Binary Tree



Postorder

```

public TreeNode invertTree(TreeNode root) {
    if(root == null) return null;

    TreeNode l = invertTree(root.left);
    TreeNode r = invertTree(root.right);

    root.left = r; root.right = l;
    return root;
}

```

Preorder

```

public TreeNode invertTree(TreeNode root) {
    if(root == null) return null;

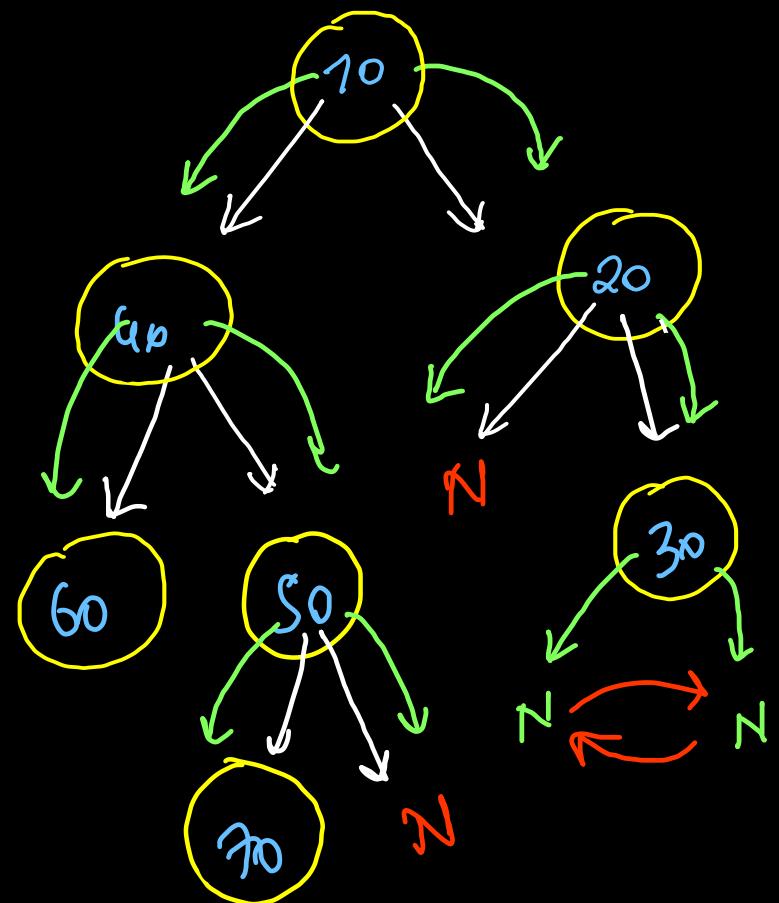
    TreeNode temp = root.left;
    root.left = root.right; root.right = temp;

    invertTree(root.left);
    invertTree(root.right);

    return root;
}

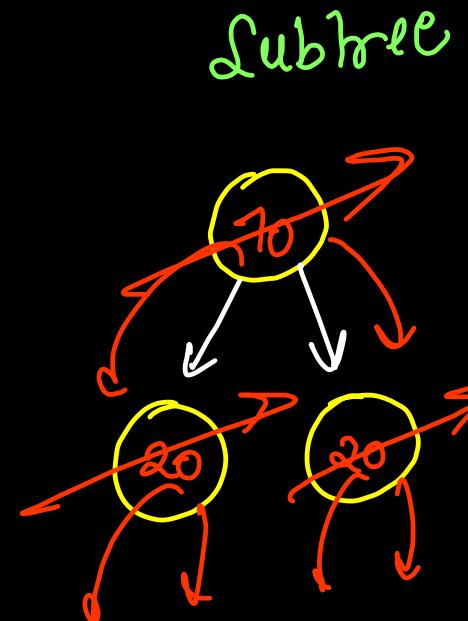
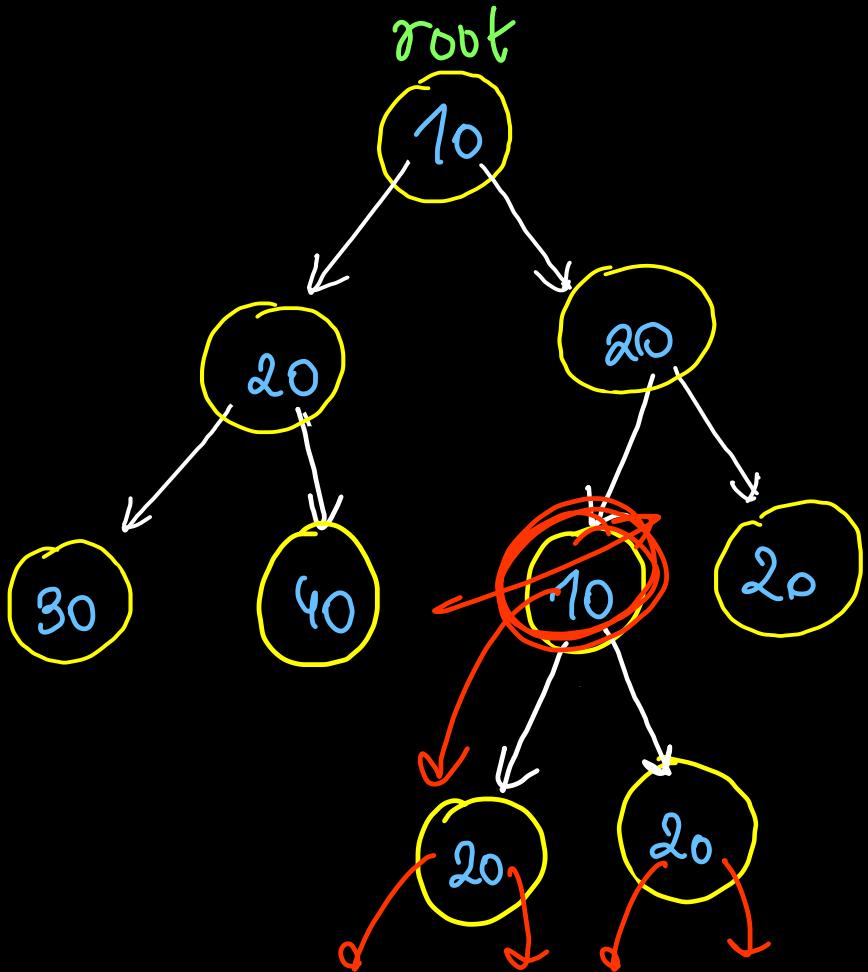
```

Time $\Rightarrow O(n)$
 Space $\Rightarrow O(h)$



LC572

Subtree of Another Tree

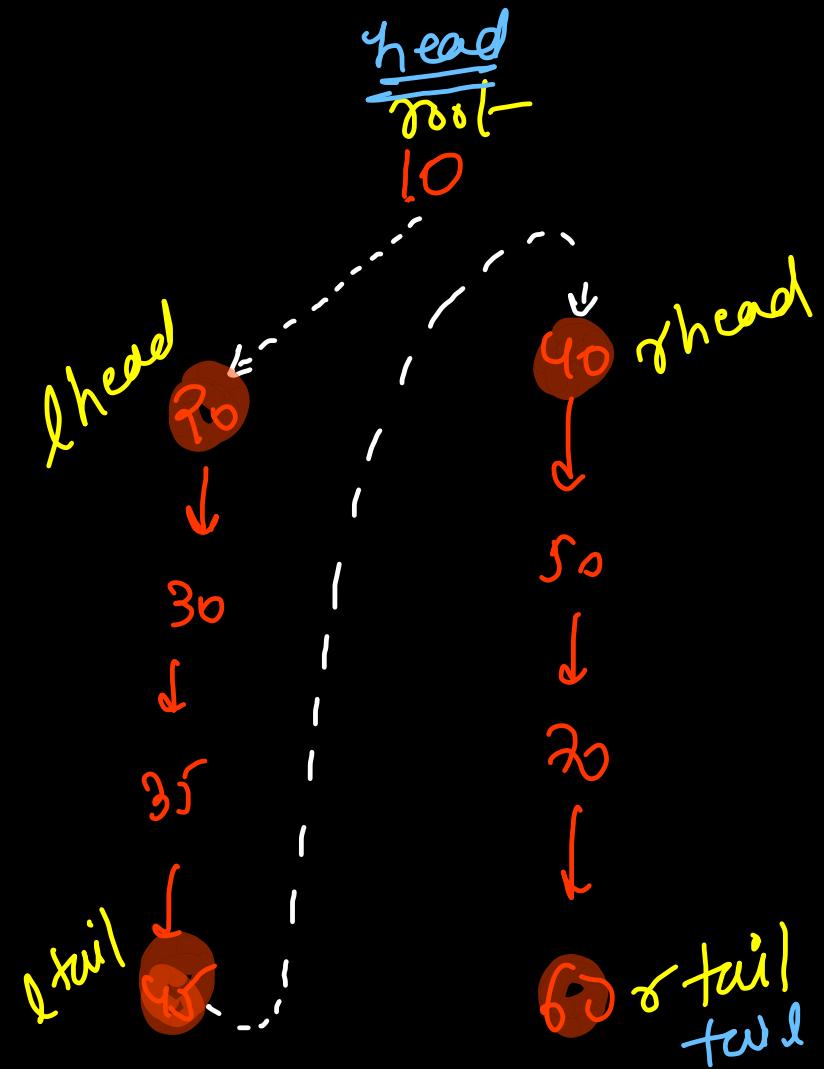
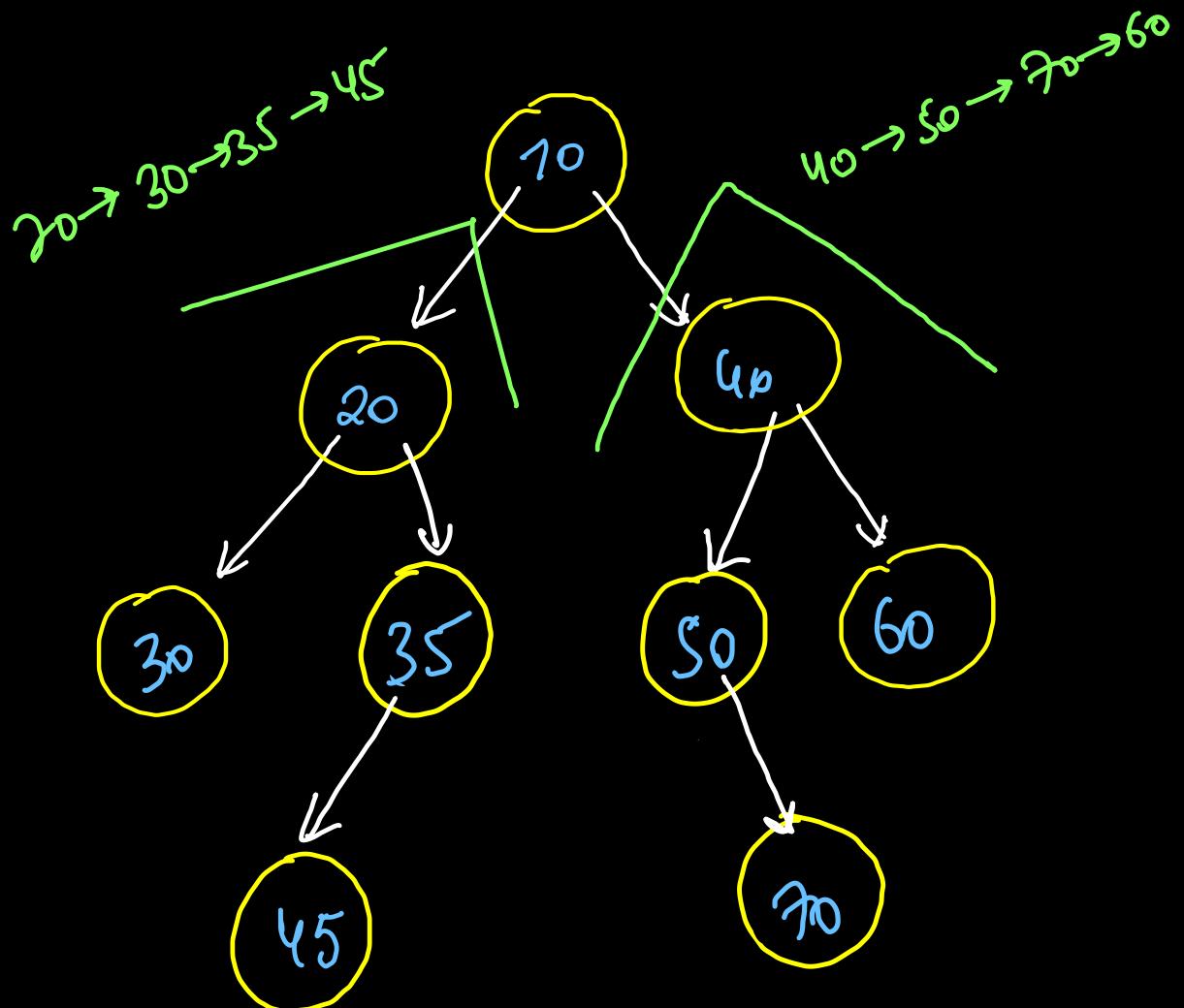


```
class Solution {  
    // fix subtree starting at the root node  
    public boolean isSame(TreeNode root, TreeNode subRoot){  
        if(root == null && subRoot == null) return true;  
        if(root == null || subRoot == null) return false;  
        if(root.val != subRoot.val) return false;  
  
        return isSame(root.left, subRoot.left)  
            && isSame(root.right, subRoot.right);  
    }  
    public boolean isSubtree(TreeNode root, TreeNode subRoot) {  
        if(root == null && subRoot == null) return true;  
        if(root == null || subRoot == null) return false;  
  
        return isSame(root, subRoot) || isSubtree(root.left, subRoot)  
            || isSubtree(root.right, subRoot);  
    }  
}
```

Time
→ $O(n^2)$

Space
→ $\Theta(n)$

Flatten/Linearize Binary Tree to Linked List



```

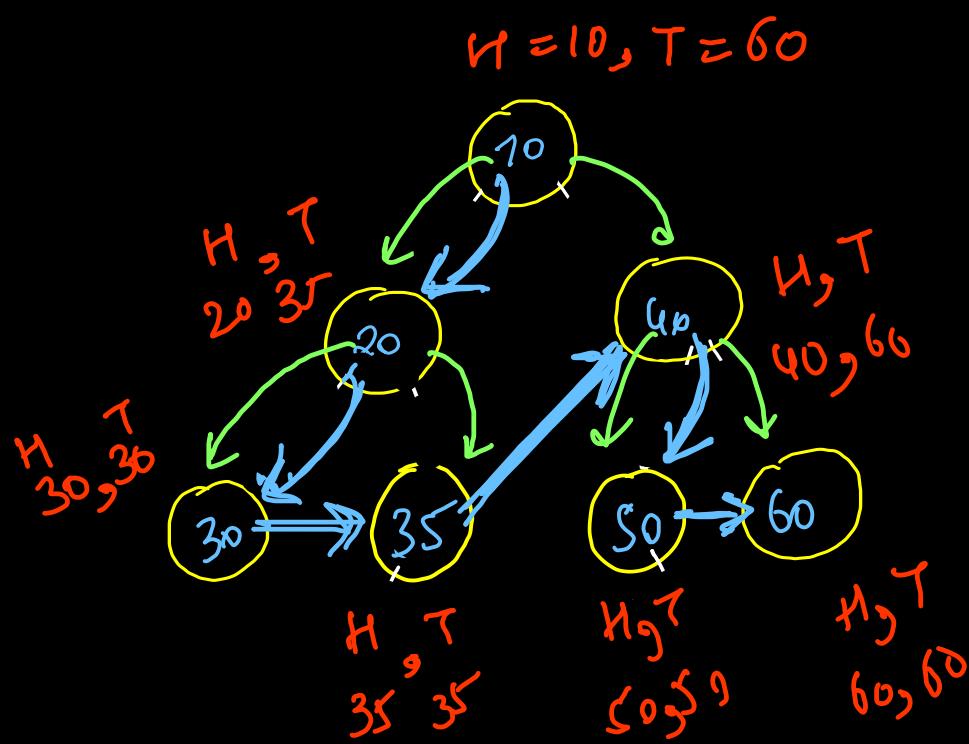
public static class Pair{
    TreeNode head, tail;
    Pair(TreeNode head, TreeNode tail){
        this.head = head;
        this.tail = tail;
    }
}

public Pair dfs(TreeNode root){
    if(root == null)
        return new Pair(null, null);
    if(leaf node  
root.left == null && root.right == null)
        return new Pair(root, root);

    Pair l = dfs(root.left);  
// linearize left subtree
    Pair r = dfs(root.right);  
& right subtree
    root.left = null;
    root.right = l.head;
    l.tail.right = r.head;

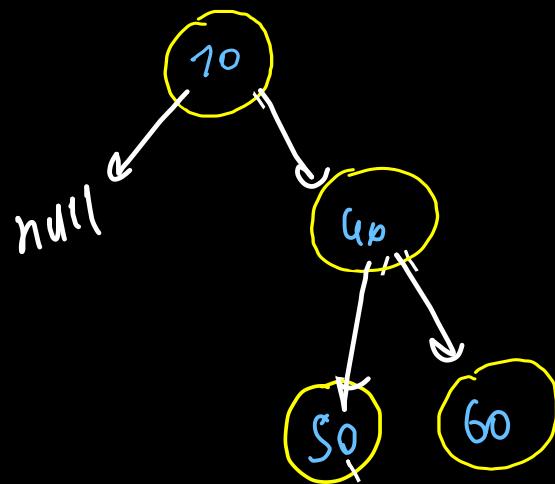
    return new Pair(root, r.tail);
}

```



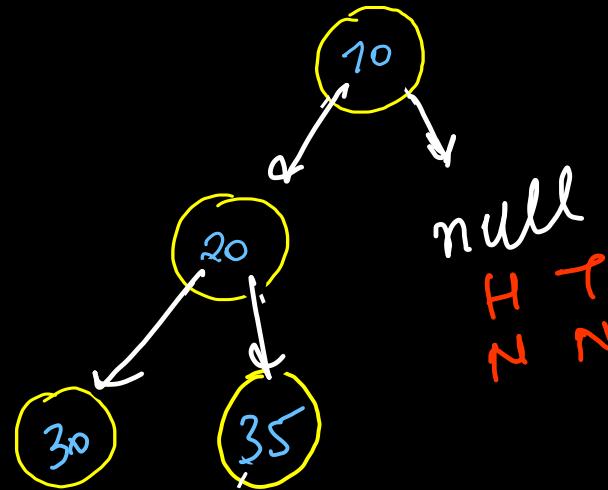
Time $\Rightarrow \mathcal{O}(n)$

Space $\Rightarrow \mathcal{O}(h)$



$10 \rightarrow 40 \rightarrow 50 \rightarrow 60$

Corner
Cases



$10 \rightarrow 20 \rightarrow 30 \rightarrow 35$

~~tail~~

null
H T
N N

```
public static class Pair{  
    TreeNode head, tail;  
    Pair(TreeNode head, TreeNode tail){  
        this.head = head;  
        this.tail = tail;  
    }  
}
```

```
public void flatten(TreeNode root) {  
    dfs(root);  
}
```

Time = $O(n)$

Space = $O(h)$

```
public Pair dfs(TreeNode root){  
    if(root == null)  
        return new Pair(null, null);  
  
    if(root.left == null && root.right == null)  
        return new Pair(root, root);  
  
    if(root.left == null){  
        Pair r = dfs(root.right);  
        return new Pair(root, r.tail);  
    }  
  
    Pair l = dfs(root.left);  
    Pair r = dfs(root.right);  
  
    root.left = null;  
    root.right = l.head;  
    if(r.head == null)  
        return new Pair(root, l.tail);  
    l.tail.right = r.head;  
    return new Pair(root, r.tail);  
}
```

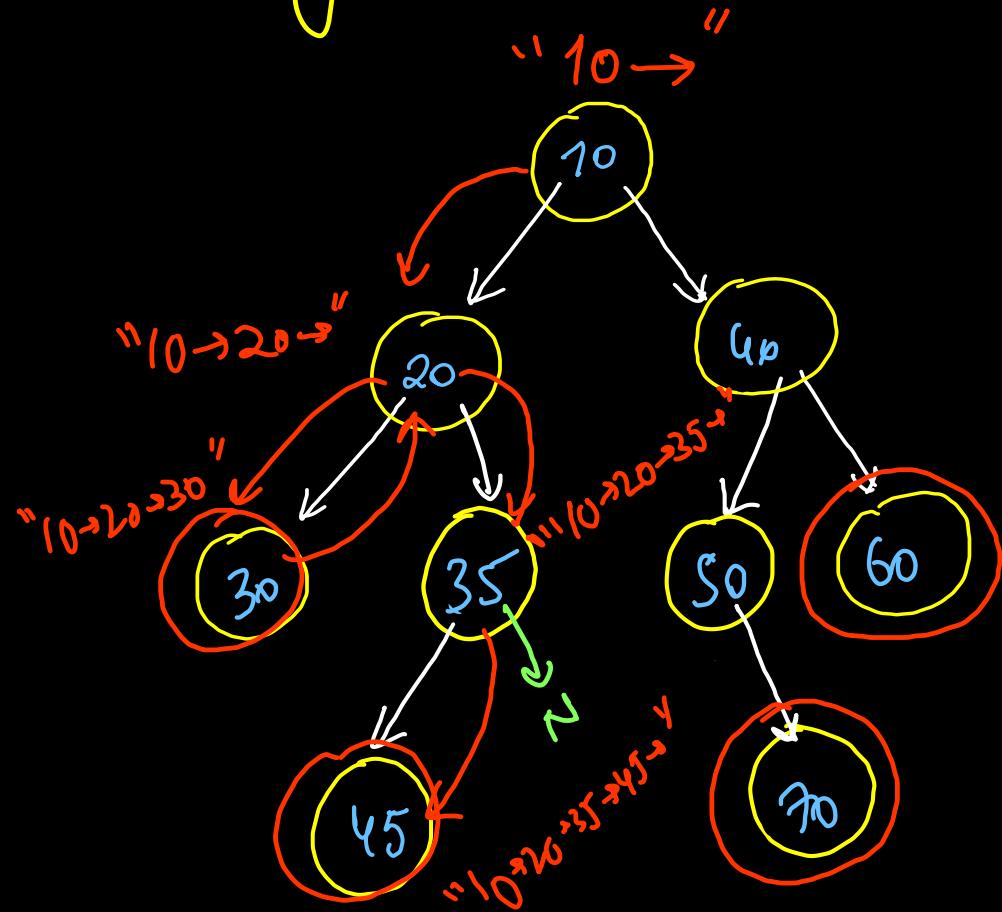
leaf node

} only right child

} only left child

Path sum Problems

Binary Tree Paths LC Q57 : all root to leaf paths



{ " 10 → 20 → 30 "
" 10 → 20 → 35 → 45 "
" 10 → 40 → 50 → 70 "
" 10 → 40 → 60 " }

```
List<String> paths = new ArrayList<>();
```

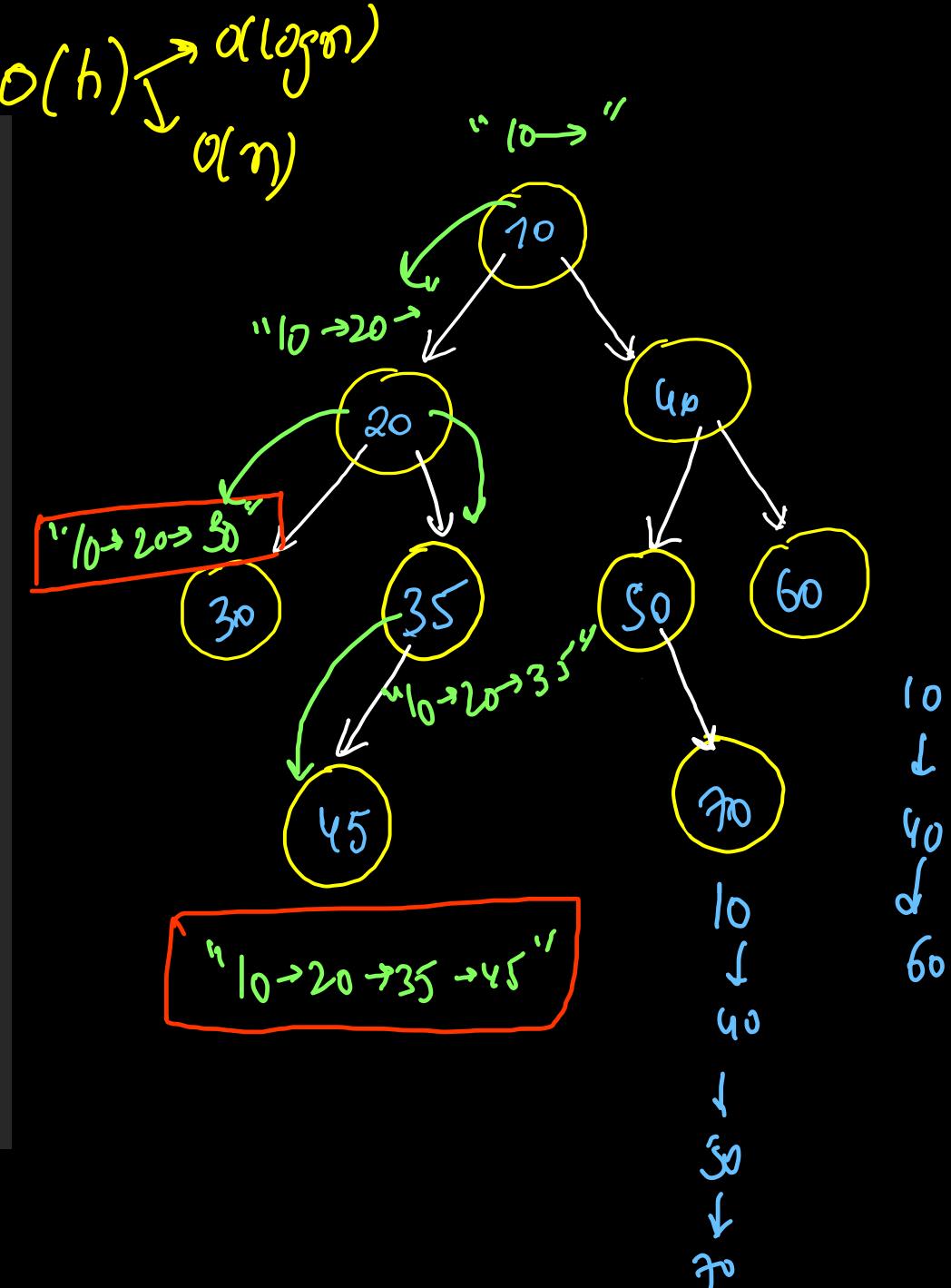
```
public void dfs(TreeNode root, String path){  
    if(root == null) return;
```

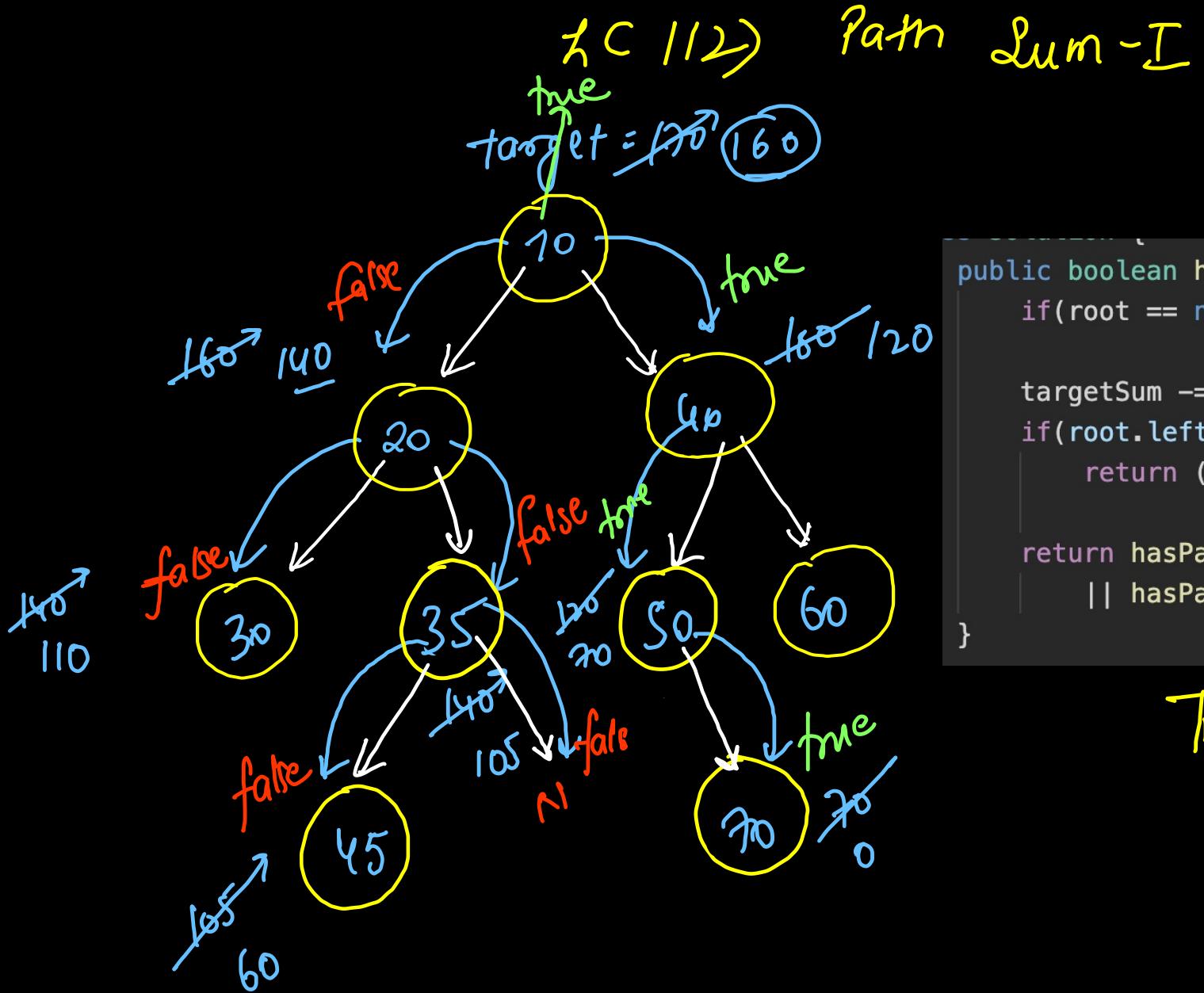
```
    if(root.left == null && root.right == null){  
        path += root.val;  
        paths.add(path);  
        return;  
    }
```

```
    dfs(root.left, path + root.val + "->");  
    dfs(root.right, path + root.val + "->");  
}
```

```
public List<String> binaryTreePaths(TreeNode root) {  
    dfs(root, "");  
    return paths;  
}
```

Time = $O(n)$, Space = $O(h) \xrightarrow{O(\log n)} O(n)$





target = 170

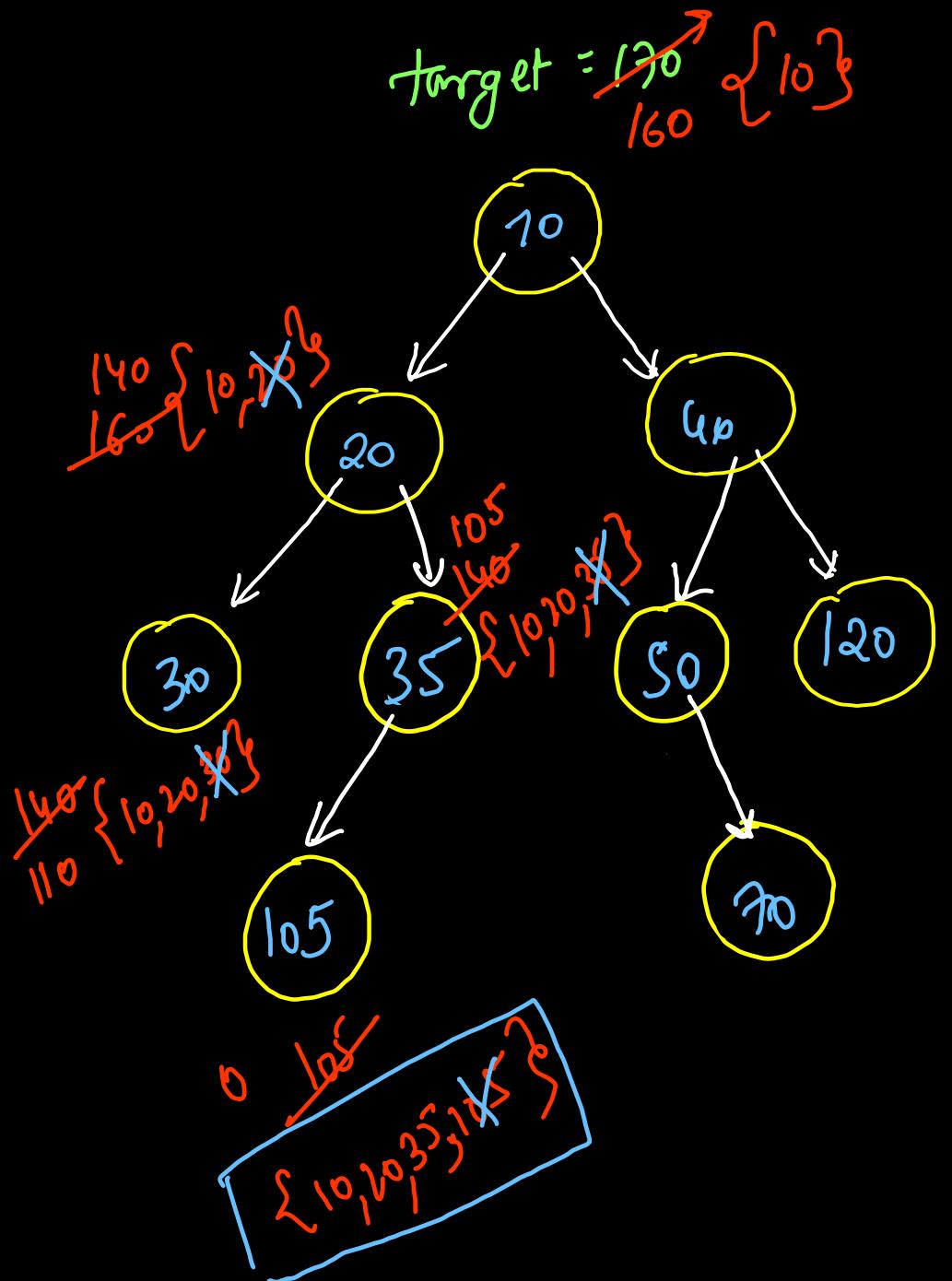
```

public boolean hasPathSum(TreeNode root, int targetSum) {
    if(root == null) return false;

    targetSum -= root.val;
    if(root.left == null && root.right == null)
        return (targetSum == 0);

    return hasPathSum(root.left, targetSum)
        || hasPathSum(root.right, targetSum);
}
  
```

Time $\Rightarrow O(n)$, Space $\Rightarrow O(h)$



LC 113) Path Sum -II
All targetsum root to leaf paths

- { {10, 20, 35, 105},
- { 10, 40, 50, 70},
- { 10, 40, 120 } }

```
class Solution {
    List<List<Integer>> paths = new ArrayList<>();

    public void dfs(TreeNode root, int targetSum, List<Integer> path) {
        if(root == null) return;

        targetSum -= root.val;
        path.add(root.val);

        if(root.left == null && root.right == null && targetSum == 0){
            paths.add(new ArrayList<>(path)); // deep copy
        }

        dfs(root.left, targetSum, path);
        dfs(root.right, targetSum, path);
        path.remove(path.size() - 1); // backtracking
    }

    public List<List<Integer>> pathSum(TreeNode root, int targetSum) {
        dfs(root, targetSum, new ArrayList<>());
        return paths;
    }
}
```

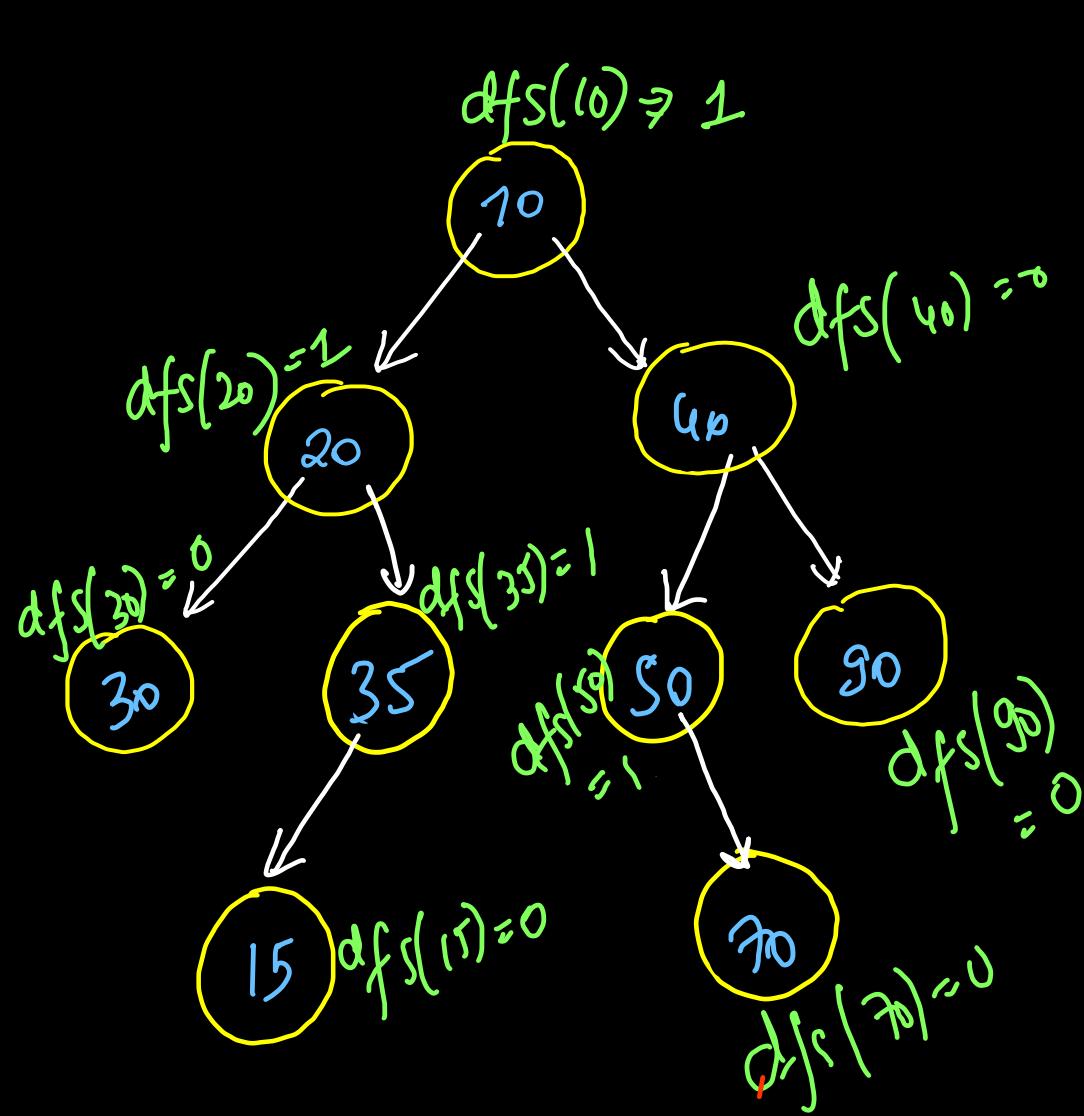
Time $\Rightarrow O(n)$

dfs

Space $\Rightarrow O(h)$

2C 437

Path Sum - III



targetsum $\Rightarrow 50$

$$\left\{ \begin{array}{l} 10 \rightarrow 40, \\ 20 \rightarrow 30, \\ 35 \rightarrow 15, \\ 50 \end{array} \right\}$$

- need not start from root
- need not end at leaf
- must go downward.

```

// downward path must start from root
public int dfs(TreeNode root, long targetSum){
    if(root == null) return 0;

    targetSum -= root.val;
    int path = (targetSum == 0) ? 1 : 0;

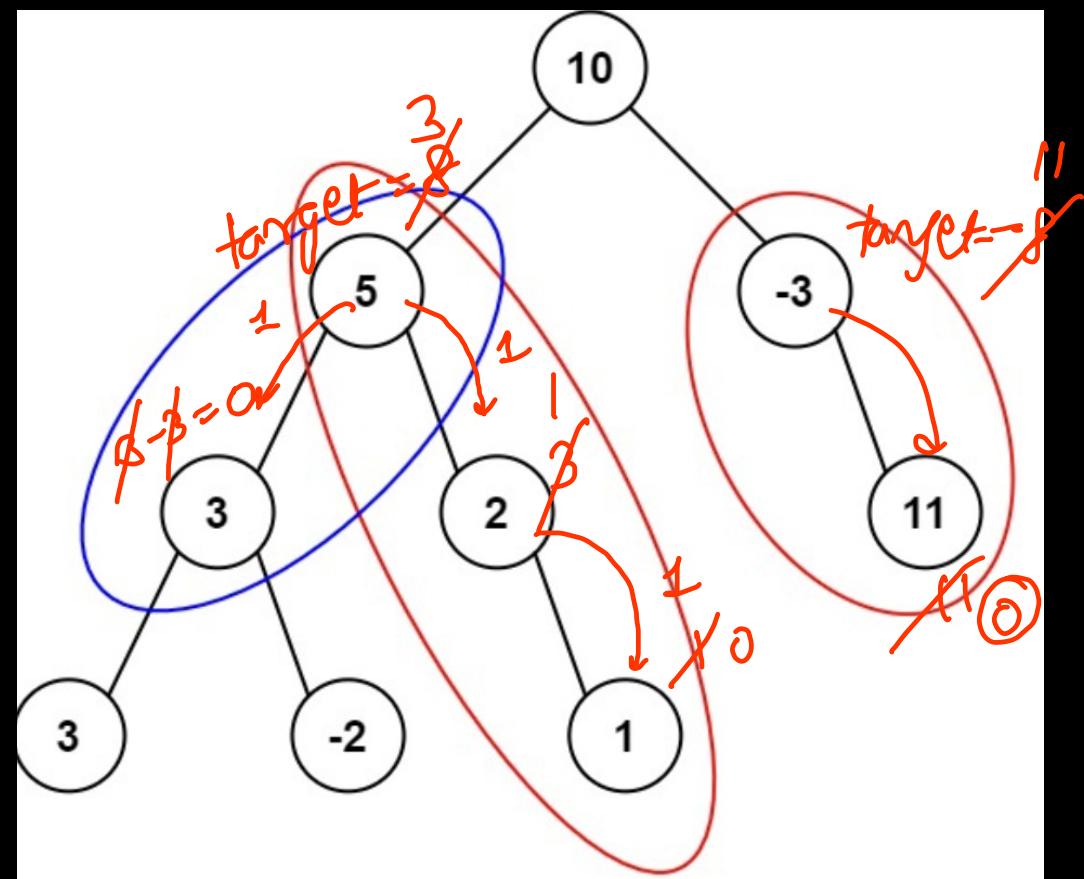
    path += dfs(root.left, targetSum);
    path += dfs(root.right, targetSum);
    return path;
}

public int pathSum(TreeNode root, int targetSum) {
    if(root == null) return 0;
    return dfs(root, targetSum)
        + pathSum(root.left, targetSum)
        + pathSum(root.right, targetSum);
}

```

targetSum=8

Time $\Rightarrow \Theta(n^2)$, Space $\Rightarrow \Theta(h)$

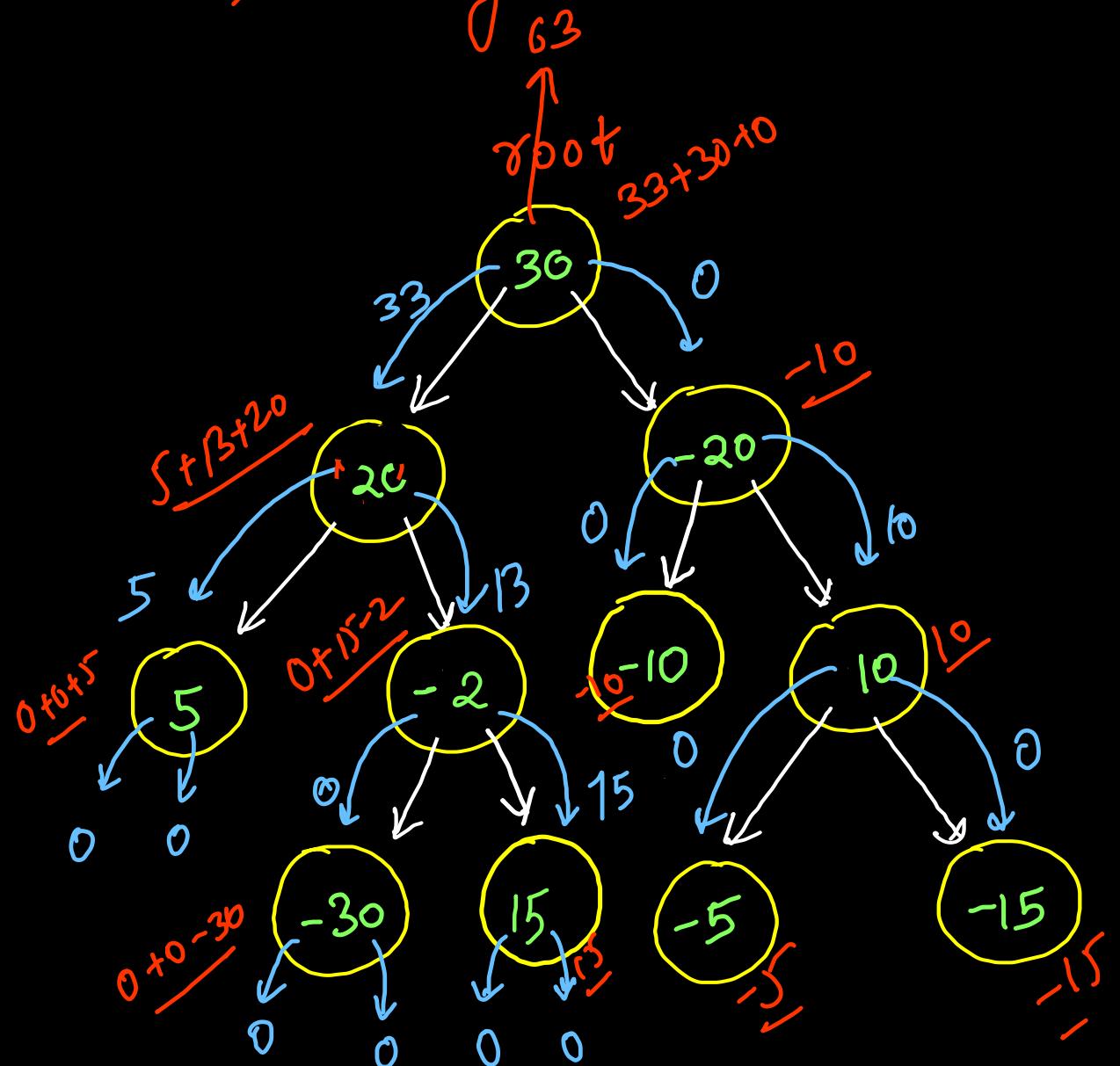


$$\text{dfs}(10) = 0$$

$$\text{dfs}(5) = 2$$

$$\text{dfs}(-3) = 1$$

LC 124) Binary Tree Max Path Sum



Path max path sum

- ① need not start from root
- ② need not end at leaf
- ③ need not be downward
- ④ must be non-empty!

diameter = (global)

$\Rightarrow \text{leftsum} + \text{rightsum} + \text{root val}$

height = $\max(\text{leftsum}, \text{rightsum}) + \underline{\text{root val}}$

```

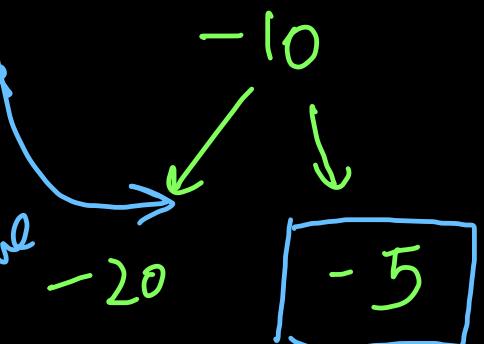
long diameter = Long.MIN_VALUE;
public long height(TreeNode root){
    if(root == null) return 0;
    long lsum = Math.max(height(root.left), 0); → to discard negative sum (path need not end at leaf)
    long rsum = Math.max(height(root.right), 0);
    diameter = Math.max(diameter, lsum + rsum + root.val);
    return Math.max(lsum, rsum) + root.val;
}

public int maxPathSum(TreeNode root) {
    height(root);
    return (int)diameter;
}

```

non-empty path

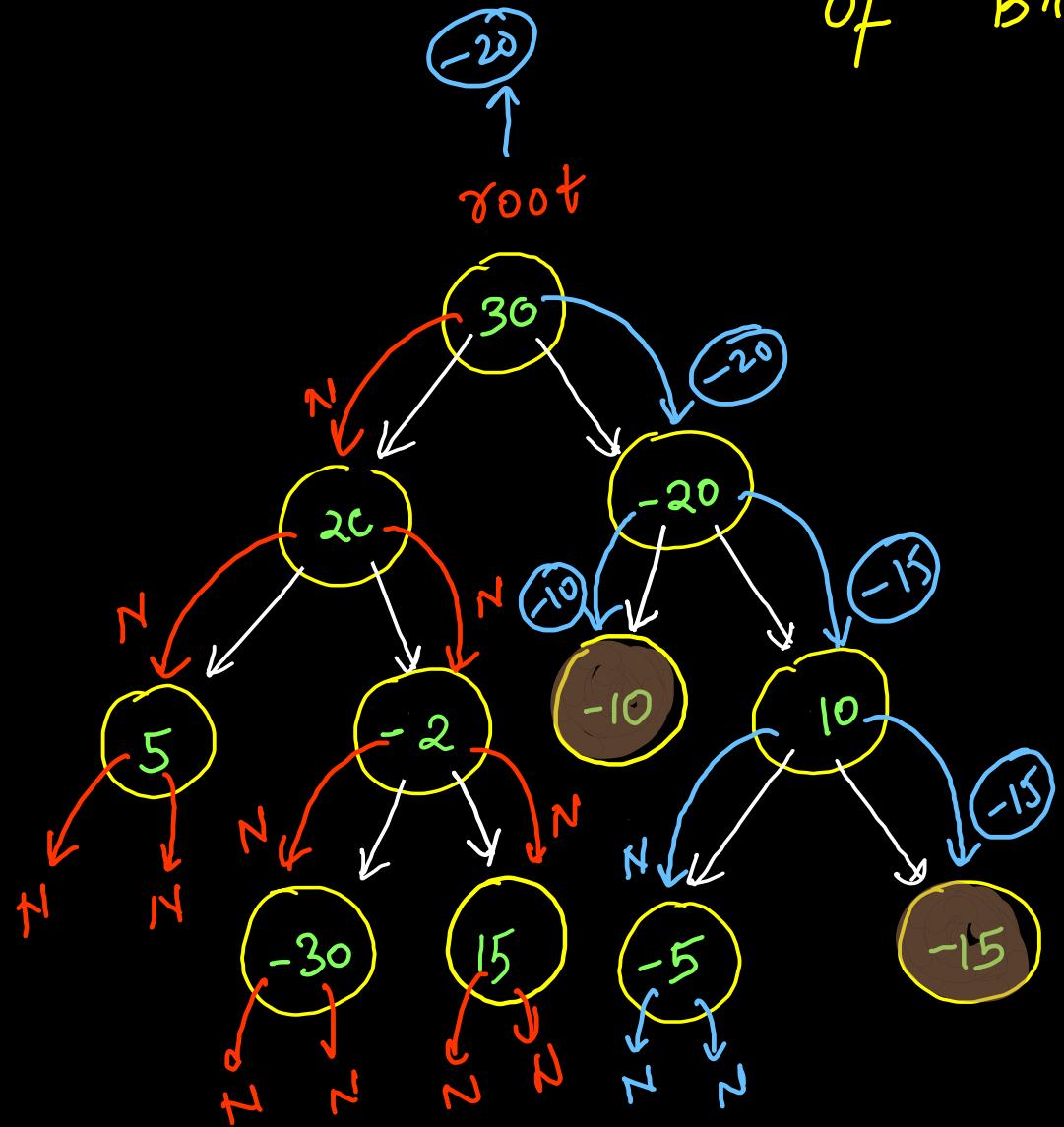
Common
ancestor
All nodes
are negative



Time = $O(n)$

Space = $O(h)$

Lowest Common Ancestor of Binary tree



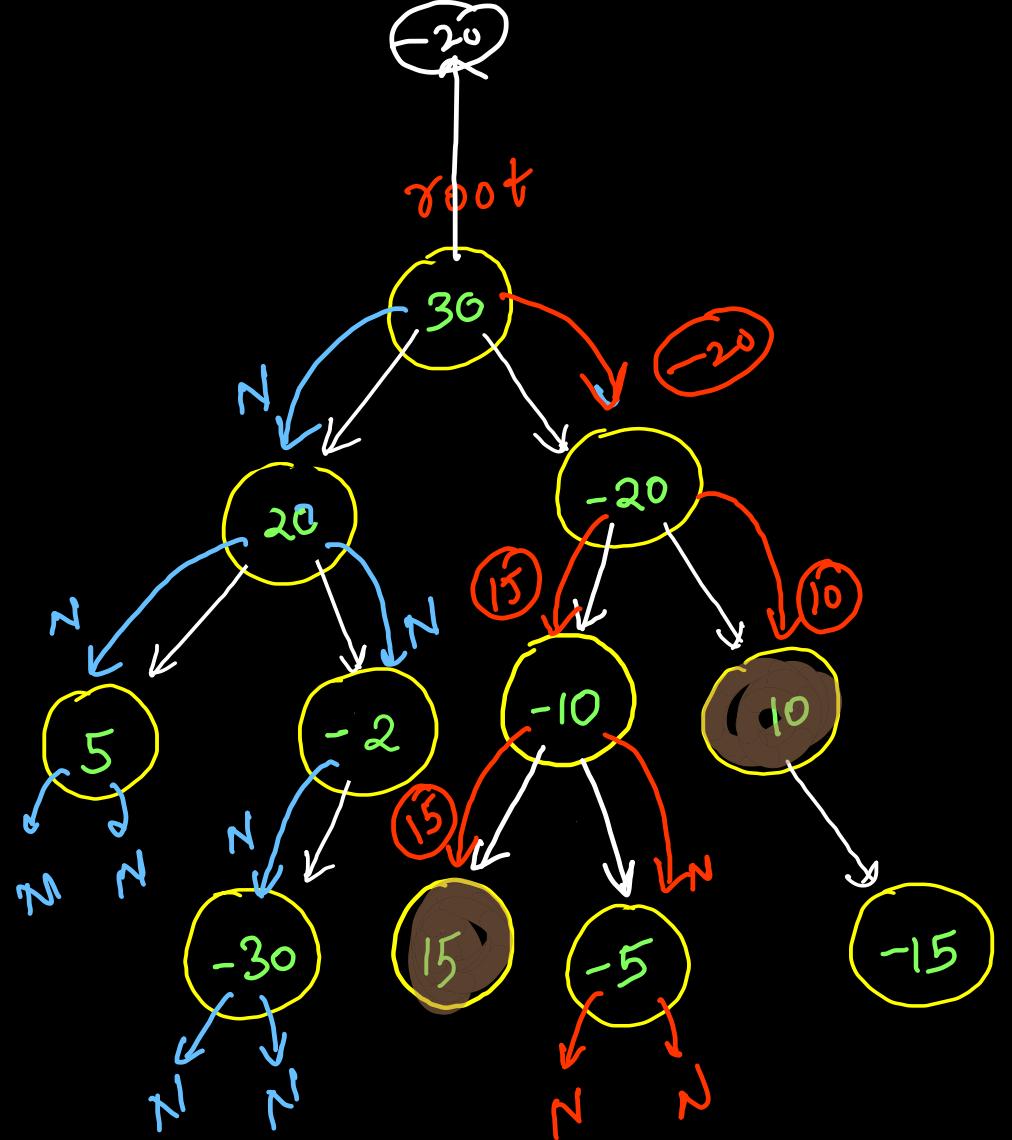
$$\text{LCA}(-2, 5) = 20$$

$$\begin{aligned} -2 &\Rightarrow \boxed{-2, 20, 30} \\ +5 &\Rightarrow \boxed{5, 20, 30} \end{aligned}$$

$$\bullet \text{LCA}(-10, -15) \Rightarrow -20$$

$$\begin{aligned} -10 &\Rightarrow \cancel{-10, -20, 30} \\ -15 &\Rightarrow \cancel{-15, 10, -20, 30} \end{aligned}$$

$$\text{LCA}(-20, -30) \Rightarrow 20$$



$$p = 15, q = 10$$

```
if (root == p || root == q || root == N)
    return root;
```

$\text{left} \rightarrow N, \text{right} \rightarrow N$
 $\Rightarrow \text{return null};$

$\text{left} != N, \text{right} == N$
 $\Rightarrow \text{return left}$

$\text{left} == N, \text{right} != N$
 $\Rightarrow \text{return right}$

$\text{left} != N, \text{right} != N$
 $\Rightarrow \text{return root}.$

```
public TreeNode lowestCommonAncestor(TreeNode root, TreeNode p, TreeNode q) {  
    if(root == null || root == p || root == q) return root;  
  
    TreeNode lans = lowestCommonAncestor(root.left, p, q);  
    TreeNode rans = lowestCommonAncestor(root.right, p, q);  
  
    if(lans != null && rans != null) return root;  
    return (lans != null) ? lans : rans;  
}
```

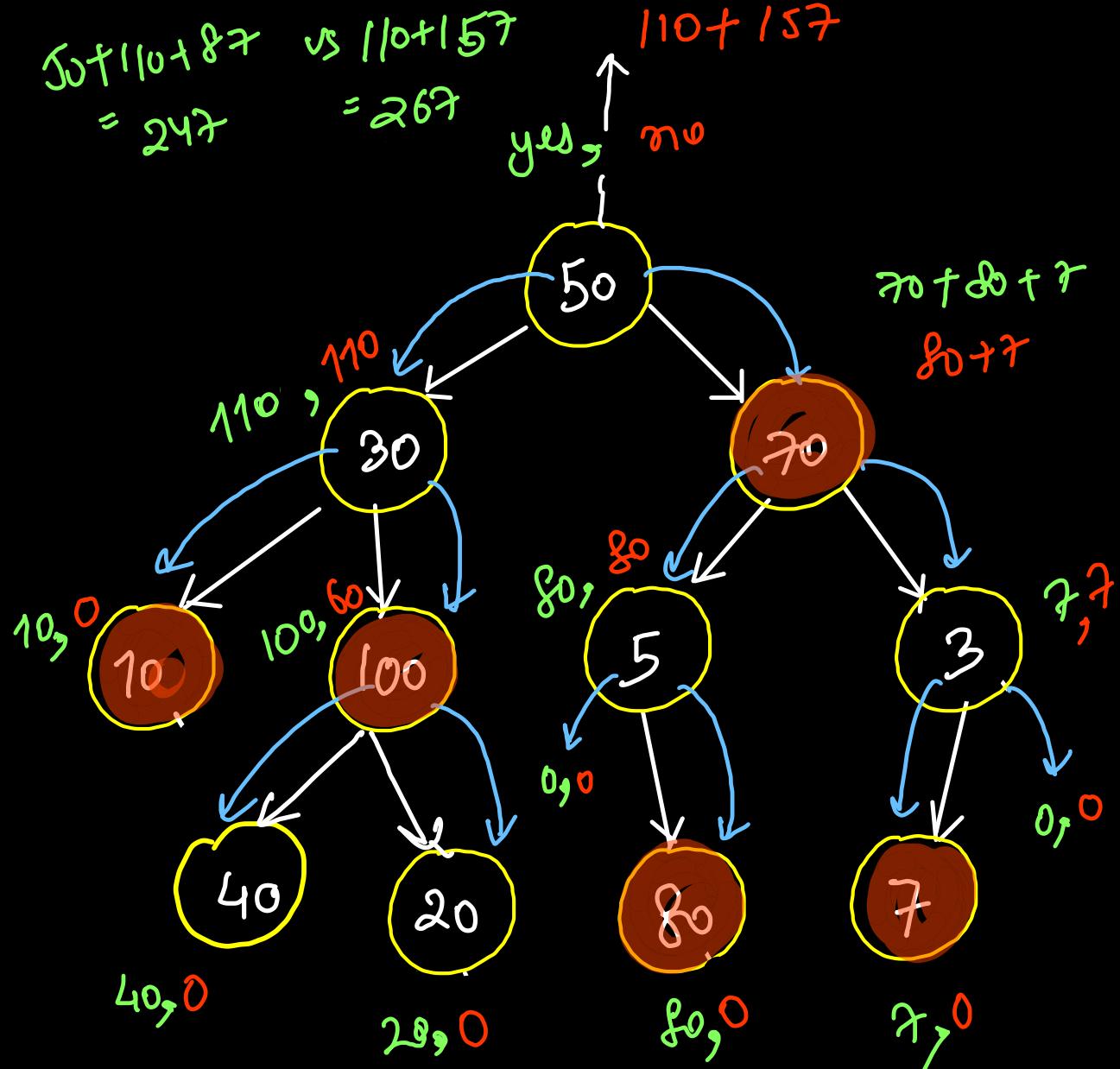
Time
 $\hookrightarrow O(n)$

Space
 $\hookrightarrow O(h)$

→ This code will only work if both nodes exist
or both nodes does not exist.

but it will fail if only one node
exists of p & q exist.

$$\begin{aligned}
 & 50 + 110 + 87 \quad \text{vs} \quad 110 + 157 \\
 & = 247 \quad \quad \quad = 267
 \end{aligned}$$



(DP on tree)
& C 327) house robber - III

max^m amount of
money w/o breaking
into 2 consecutive
houses

$$\text{root.no} = \text{left.yes} + \text{right.yes}$$

$$\text{root.yes} = \text{root.no} \rightarrow \frac{\text{root val}}{\text{f}}$$

$\xrightarrow{\text{we can skip 2 contiguous houses.}}$

$\frac{\text{left.no}}{\text{f}} + \text{right.no}$

Approach 1)

```
public static class Pair{ int yes, no; }

public Pair dfs(TreeNode root){
    if(root == null) return new Pair();

    Pair l = dfs(root.left);
    Pair r = dfs(root.right);

    Pair curr = new Pair();
    curr.no = l.yes + r.yes;
    curr.yes = Math.max(curr.no, root.val + l.no + r.no);
    return curr;
}

public int rob(TreeNode root) {
    return dfs(root).yes;
}
```

Approach 2)

```
public static class Pair{ int yes, no; }

public Pair dfs(TreeNode root){
    if(root == null) return new Pair();

    Pair l = dfs(root.left);
    Pair r = dfs(root.right);

    Pair curr = new Pair();
    curr.yes = root.val + l.no + r.no;
    curr.no = Math.max(l.yes, l.no) + Math.max(r.yes, r.no);
    return curr;
}

public int rob(TreeNode root) {
    Pair p = dfs(root);
    return Math.max(p.yes, p.no);
}
```

⇒ yes is not compulsory

{ root ko le bhi sakte hai
ya chod bhi sakte
hai }

⇒ yes is compulsory

{ root ko lena hi
lena haai }

```

public static class Pair{ int yes, no; }

public Pair dfs(TreeNode root){
    if(root == null) return new Pair();

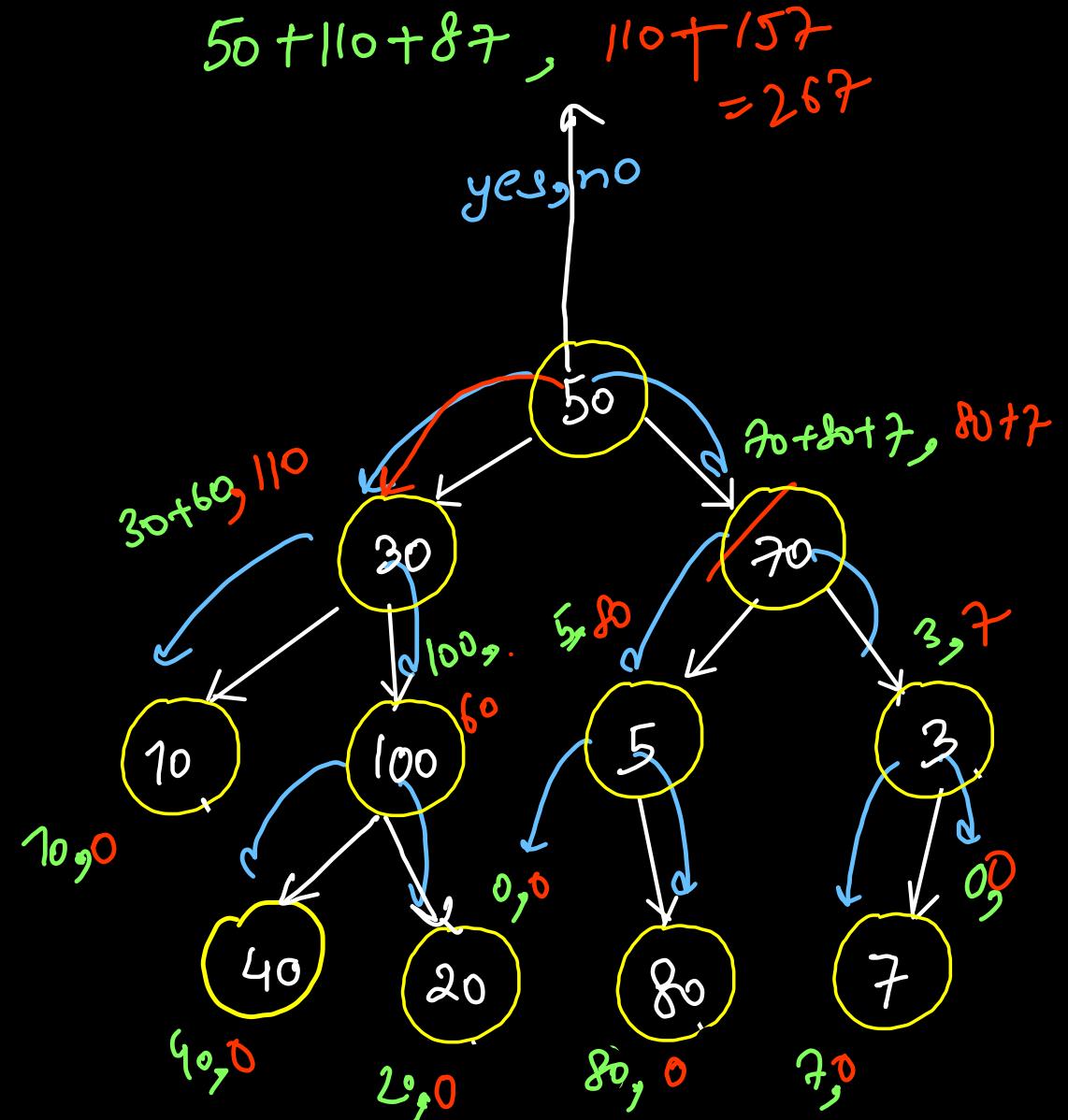
    Pair l = dfs(root.left);
    Pair r = dfs(root.right);

    Pair curr = new Pair();
    curr.yes = root.val + l.no + r.no;
    curr.no = Math.max(l.yes, l.no) + Math.max(r.yes, r.no);
    return curr;
}

public int rob(TreeNode root) {
    Pair p = dfs(root);
    return Math.max(p.yes, p.no);
}

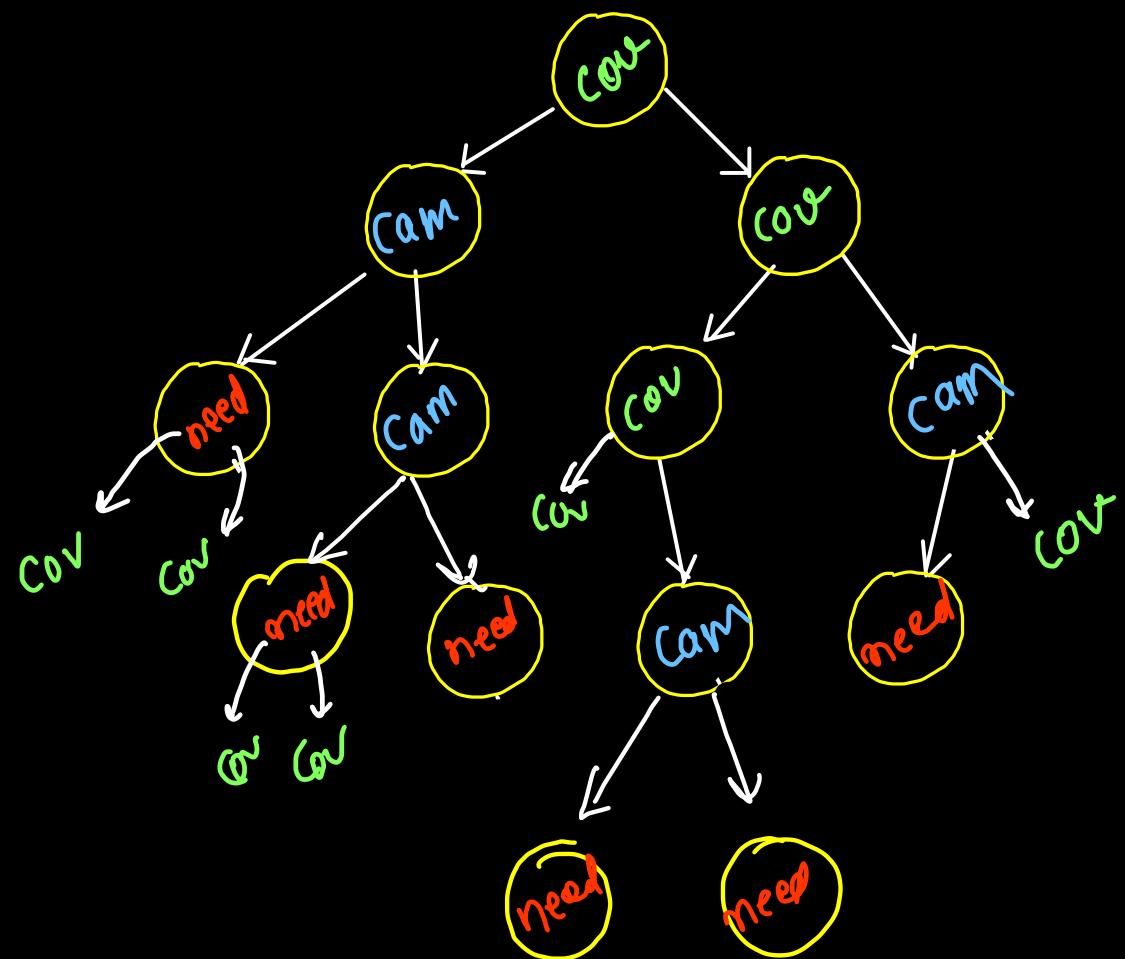
```

Time $\Rightarrow \mathcal{O}(n)$
 Space $\Rightarrow \mathcal{O}(h)$



Lc 968) Binary Tree

Cameras



greedy solⁿ

CAMERA

COVERED

→ NEED { parent }

Camera = f f f 4

```

public static enum State{
    COVERED, CAMERA, NEED;
}

int cameras = 0;

public State dfs(TreeNode root){
    if(root == null) return State.COVERED;

    State ls = dfs(root.left);
    State rs = dfs(root.right);

    if(ls == State.NEED || rs == State.NEED){
        cameras++;
        return State.CAMERA;
    }

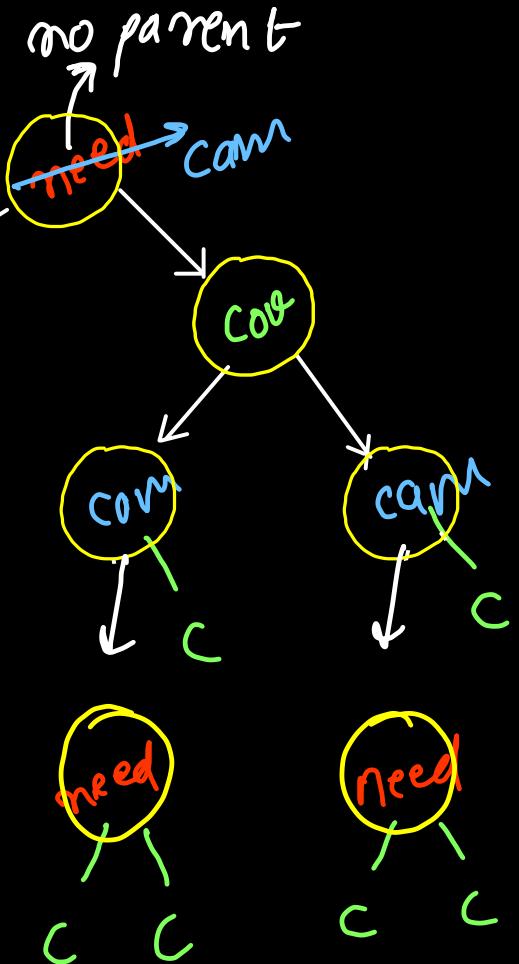
    if(ls == State.COVERED && rs == State.COVERED){
        return State.NEED;
    }

    return State.COVERED;
}

```

any one child have need \Rightarrow parent will fullfill it

greedily (parent will install camera)



```

public int minCameraCover(TreeNode root) {
    if(dfs(root) == State.NEED) {
        cameras++;
    }
    return cameras;
}

```

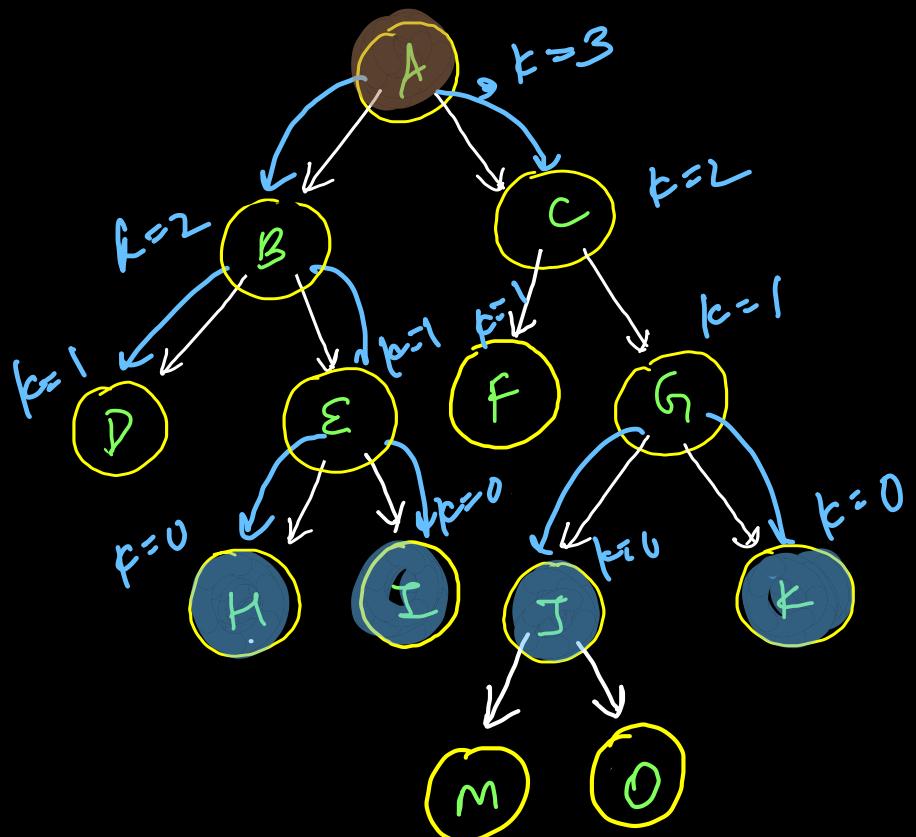
if root has need \Rightarrow it can be fulfilled by root only

All nodes at K distance

$v_{target} = G$

target = G , $k=2$

root



$k=0 \rightarrow G$

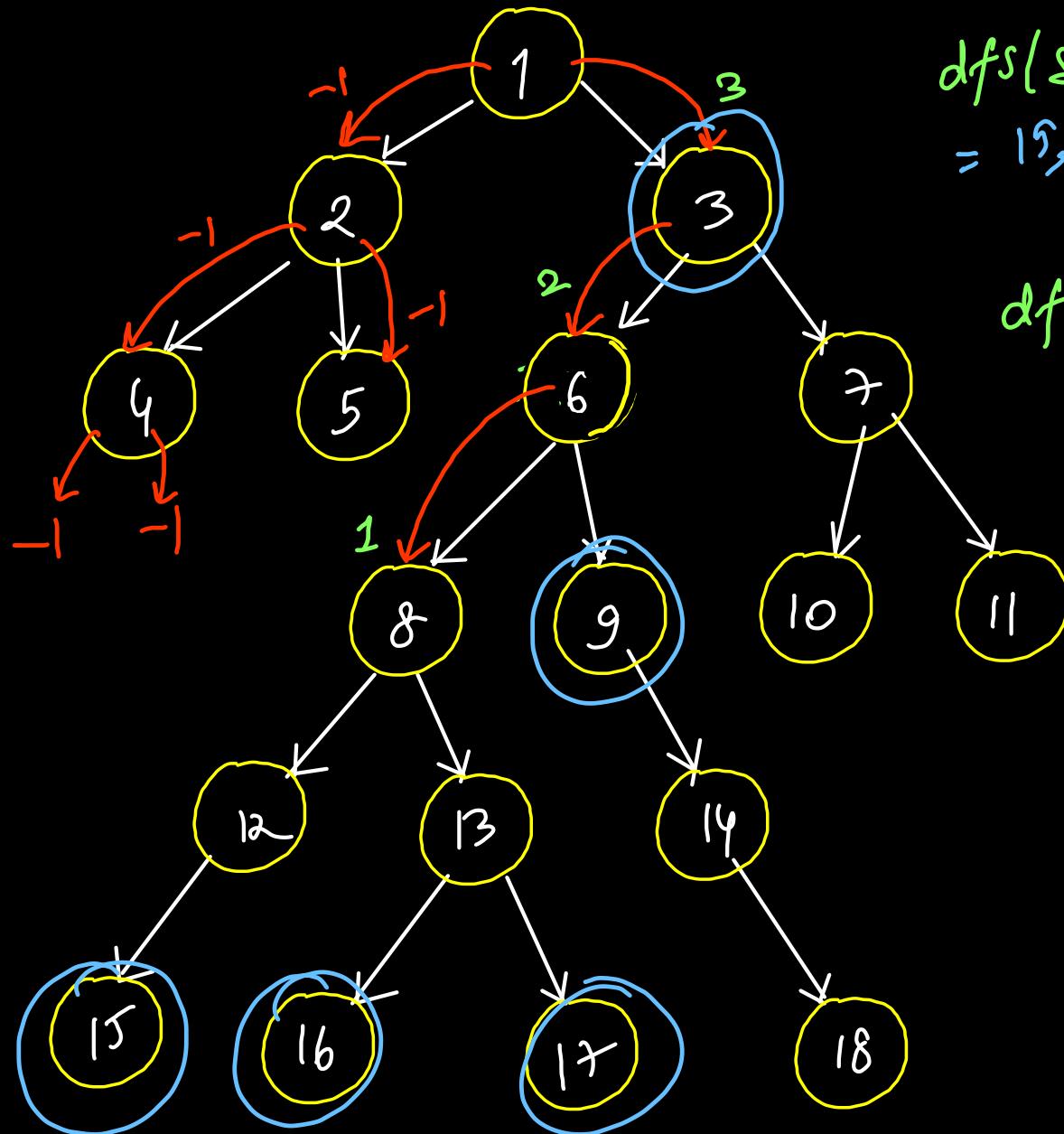
$k=1 \rightarrow C, J, K$

$\vee k=2 \rightarrow A, F, M, D$

$k=3 \rightarrow B$

$k=4 \rightarrow P, S$

$k=5 \rightarrow H, I$



$\text{dfs}(8, 2)$
 $= 15, 16, 17$

$\text{dfs}(9, 0)$
 $= 9$

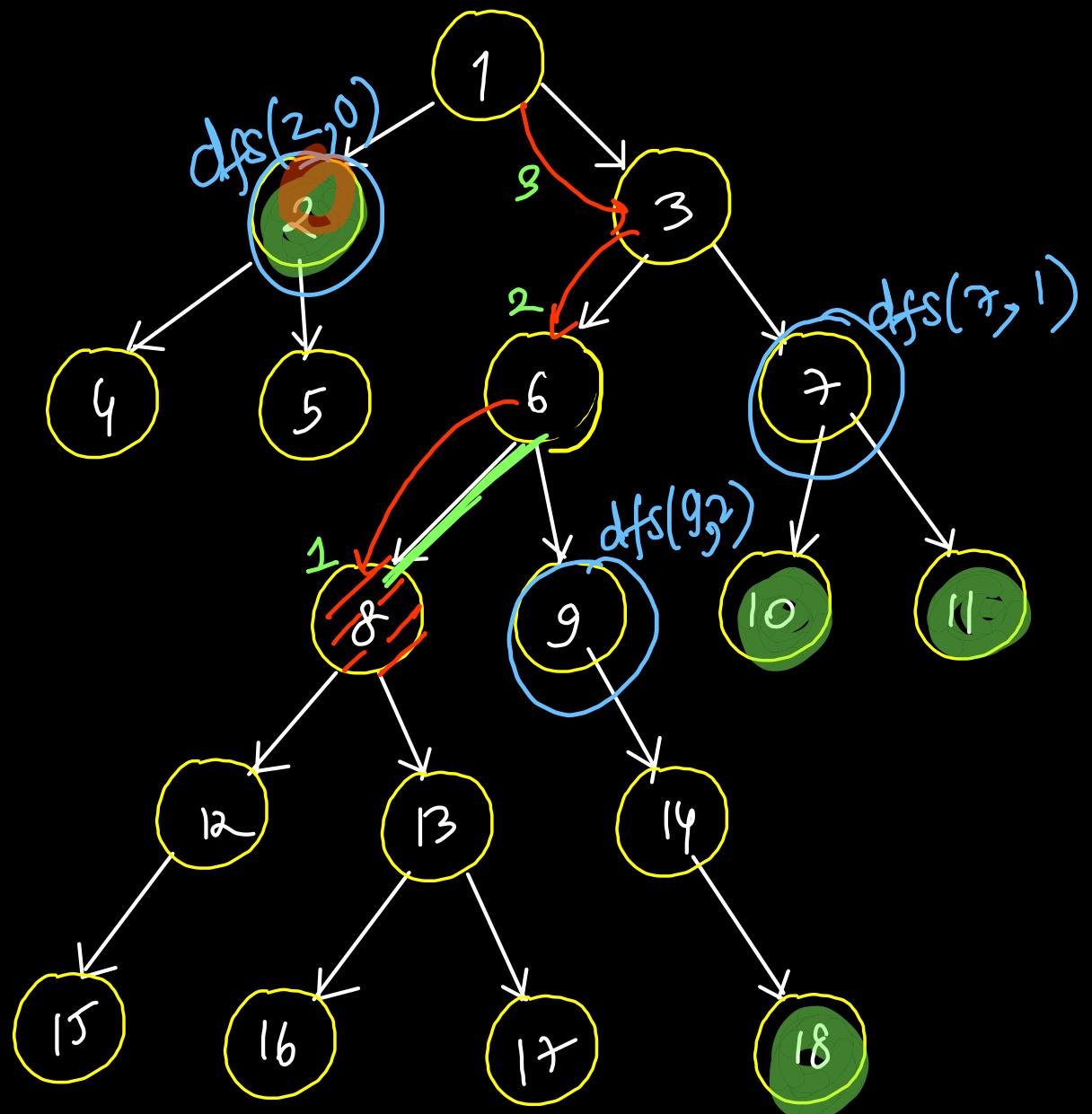
$\text{dfs}(3, 0)$
 $=$

target = 8, k = 2

```
// if root is the target node
public void dfs(TreeNode root, int k){
    if(k < 0 || root == null) return;
    if(k == 0) ans.add(root.val);
    dfs(root.left, k - 1);
    dfs(root.right, k - 1);
}
```

```
public int find(TreeNode root, TreeNode target, int k){
    if(root == null) return -1;
    if(root == target){
        dfs(root, k);
        return 1;
    }
}
```

```
int left = find(root, target, k);
int right = find(root, target, k);
```



target = 8, k = 4

```

public int find(TreeNode root, TreeNode target, int k){
    if(root == null) return -1;
    if(root == target){
        dfs(root, k);
        return 1; //dist of target to parent
    }

    int left = find(root.left, target, k);
    if(left > 0) {
        dfs(root.right, k - left - 1);
        if(left == k) ans.add(root.val);
        return 1 + left;
    }

    int right = find(root.right, target, k);
    if(right > 0) {
        dfs(root.left, k - right - 1);
        if(right == k) ans.add(root.val);
        return 1 + right;
    }

    return -1;
}

```

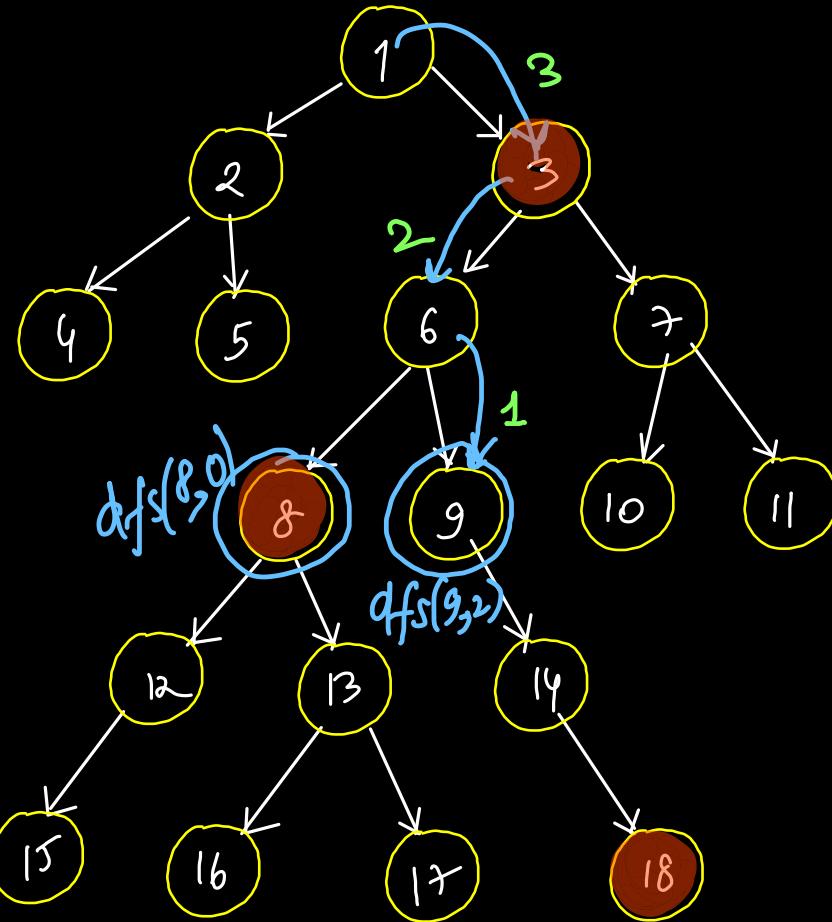
*Time \Rightarrow linear
Space \Rightarrow linear*

```

public List<Integer> distanceK(TreeNode root, TreeNode target, int k) {
    find(root, target, k);
    return ans;
}

```

⑨ $k=2$

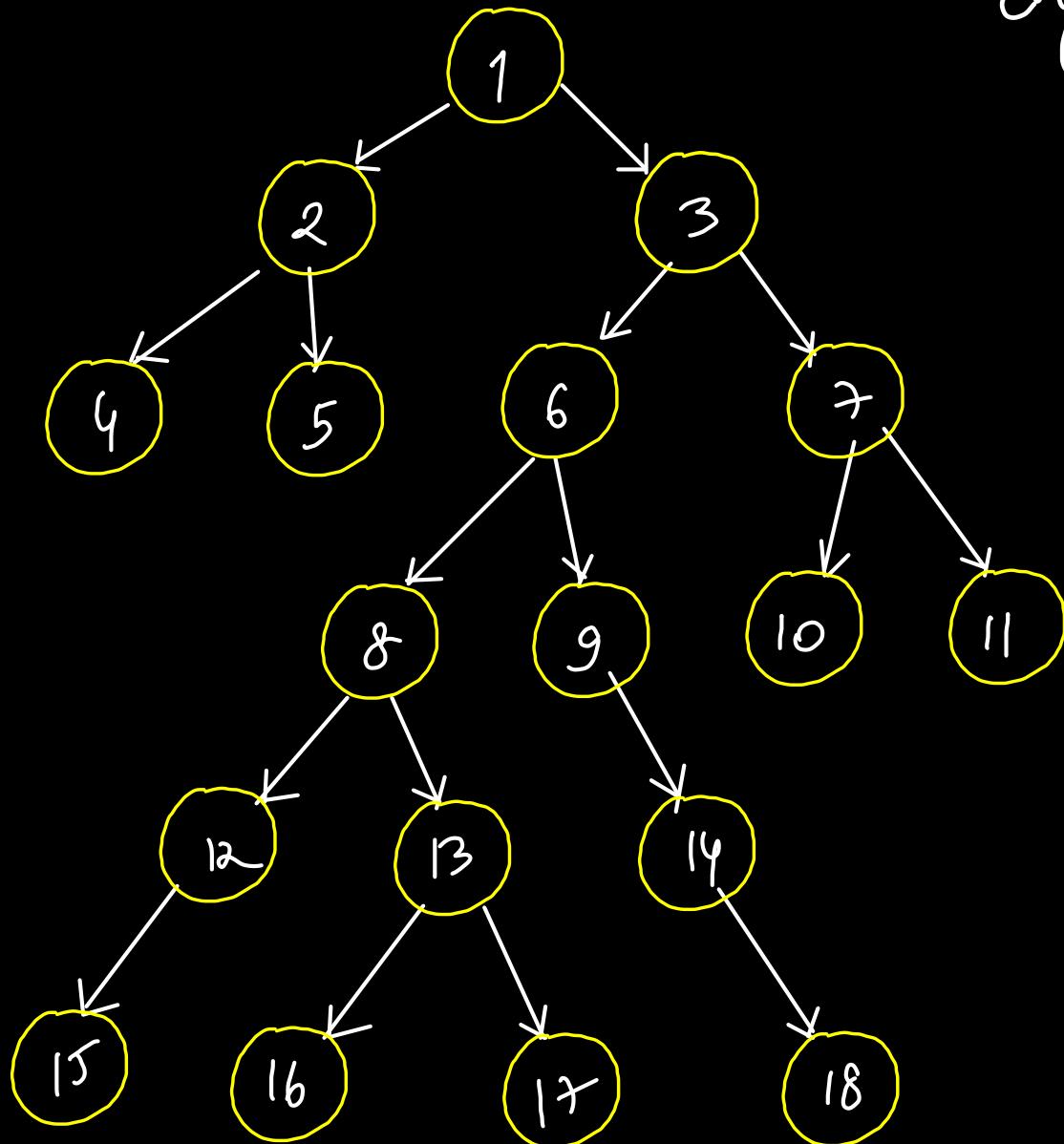


```

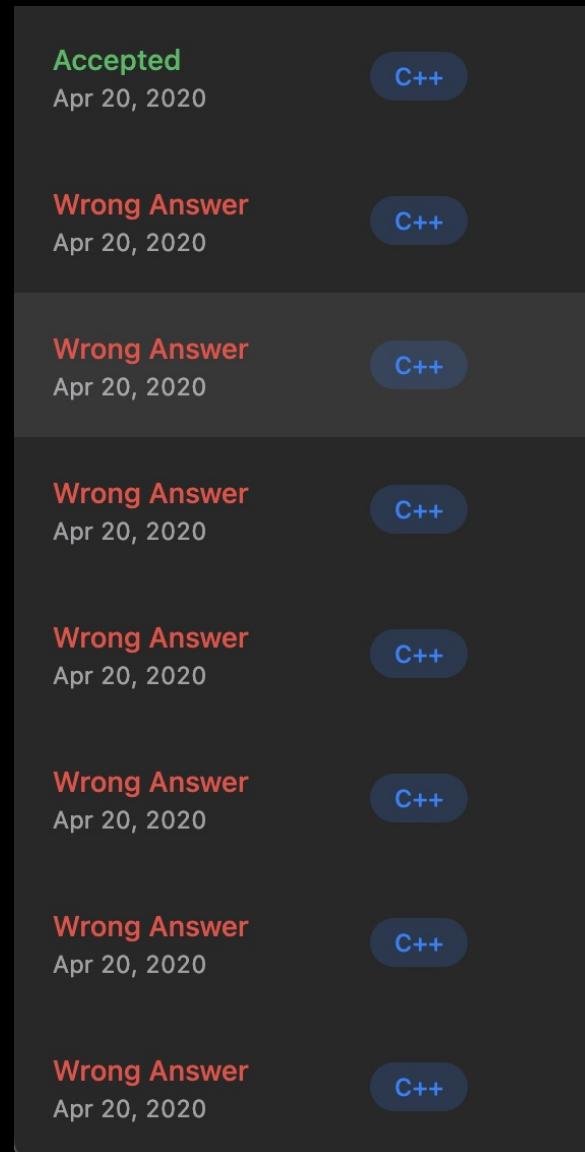
List<Integer> ans = new ArrayList<>();

// if root is the target node
public void dfs(TreeNode root, int k){
    if(k < 0 || root == null) return;
    if(k == 0) ans.add(root.val);
    dfs(root.left, k - 1);
    dfs(root.right, k - 1);
}

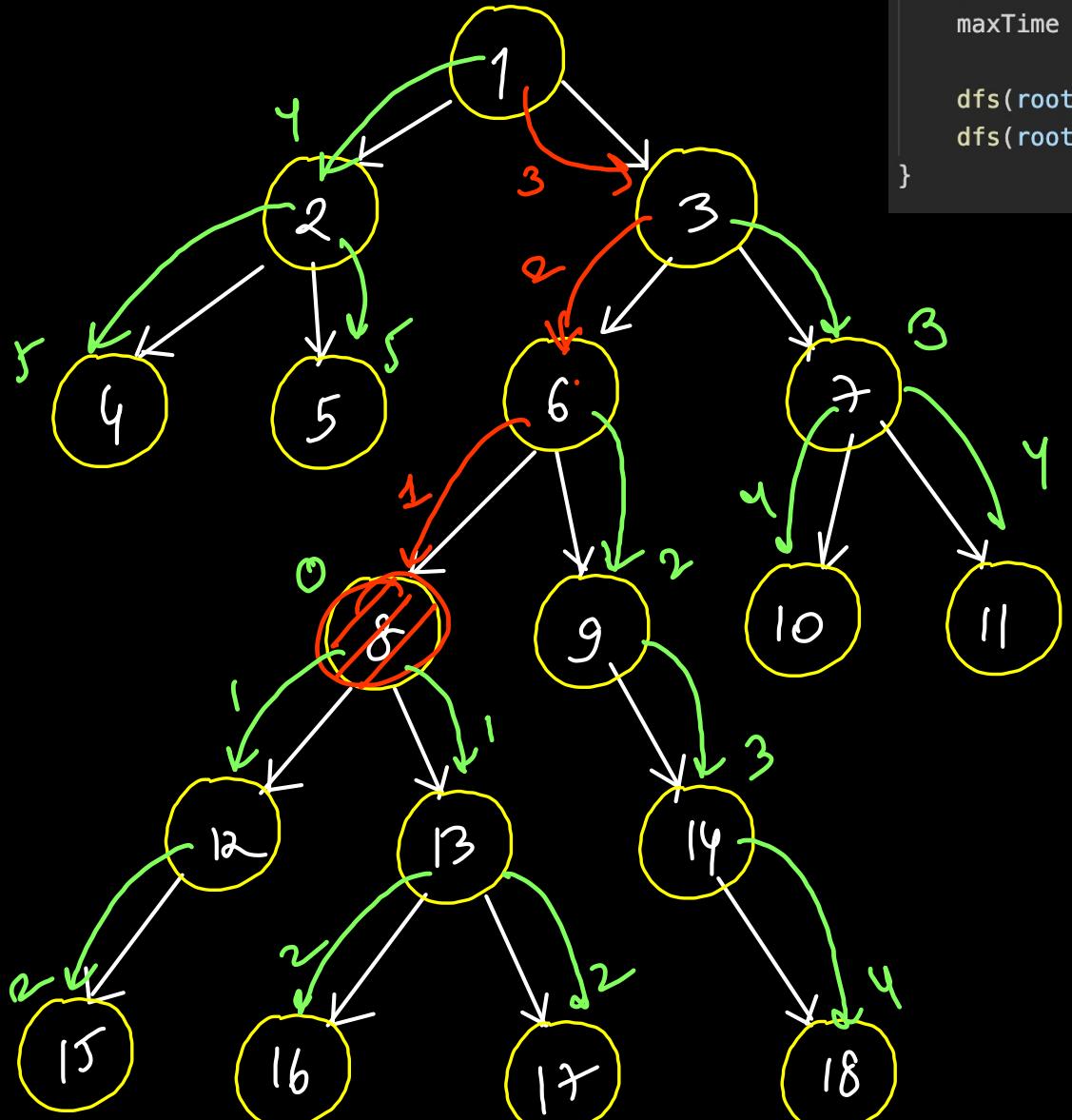
```



Engifest → 17 → Sunidhi Chaudhary
18 → Amit Trivedi
19 → AP Dhillon



~~maxtime = 6 / 1 / 2 / 3 / 4 / 5~~



```
int maxTime = 0;  
// subtract to get infected  
public void dfs(TreeNode root, int time){  
    if(root == null) return;  
    maxTime = Math.max(maxTime, time);  
  
    dfs(root.left, time + 1);  
    dfs(root.right, time + 1);  
}
```

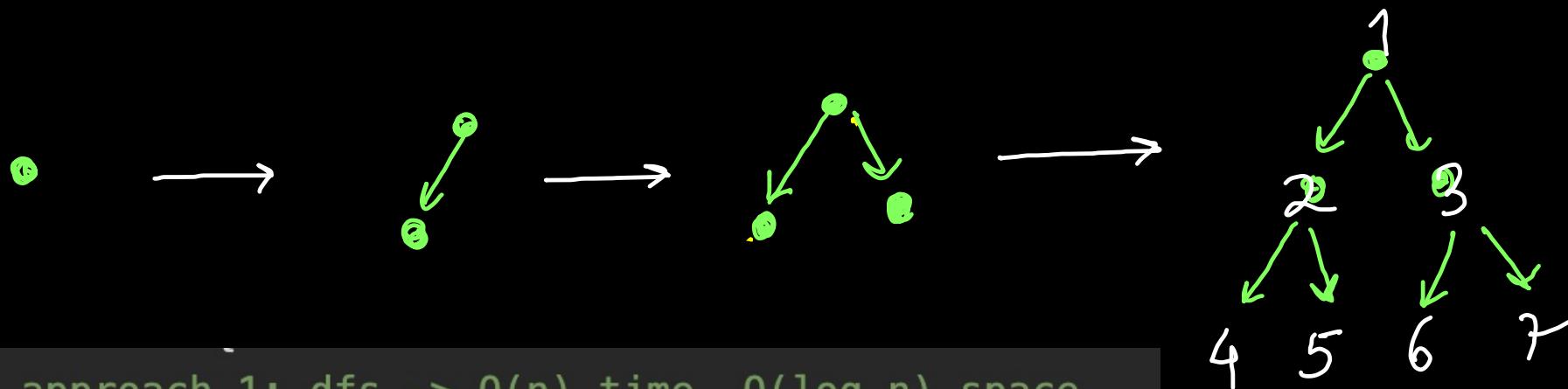
Burning Tree

```
public int find(TreeNode root, int target){  
    if(root == null) return -1;  
    if(root.val == target){  
        dfs(root, 0);  
        return 1;  
    }  
  
    int left = find(root.left, target);  
    if(left > 0) {  
        maxTime = Math.max(maxTime, left);  
        dfs(root.right, left + 1);  
        return 1 + left;  
    }  
  
    int right = find(root.right, target);  
    if(right > 0) {  
        maxTime = Math.max(maxTime, right);  
        dfs(root.left, right + 1);  
        return 1 + right;  
    }  
  
    return -1;  
}  
  
public int amountOfTime(TreeNode root, int target) {  
    find(root, target);  
    return maxTime;  
}
```

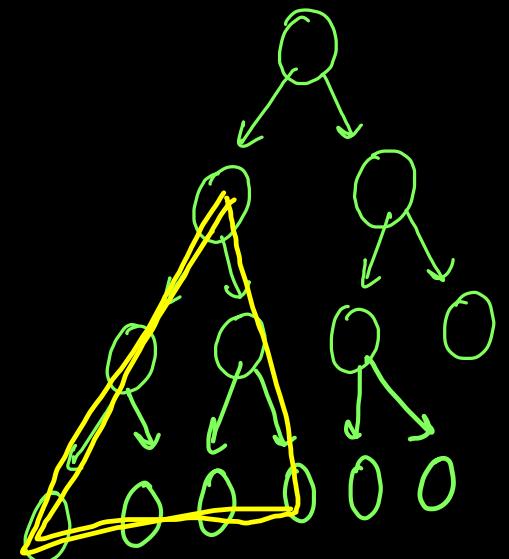
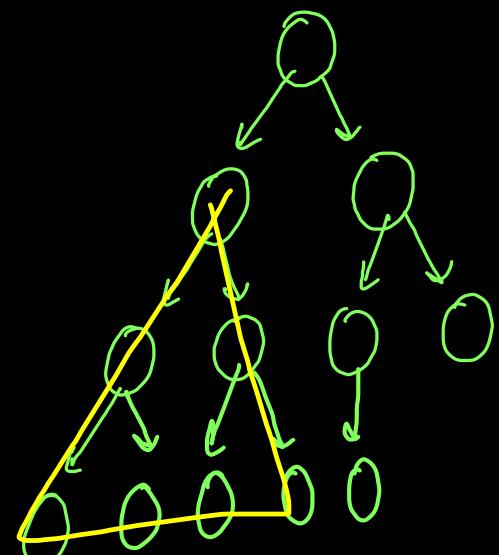
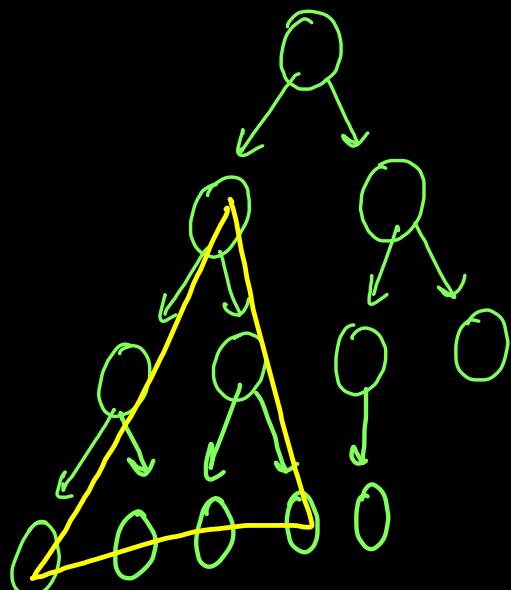
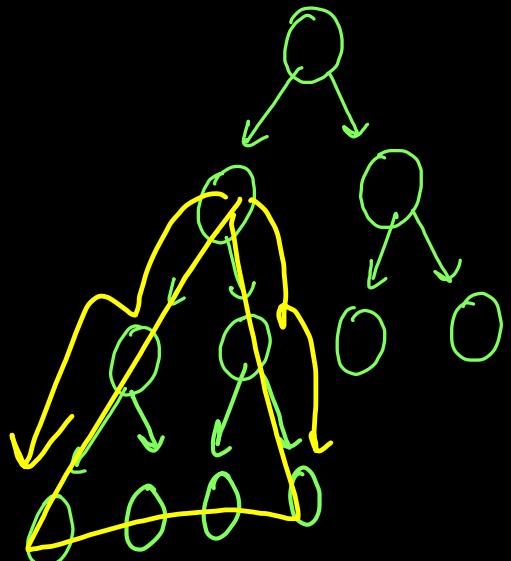
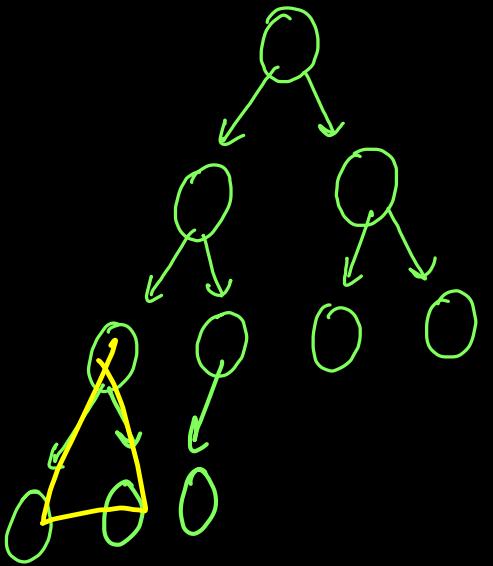
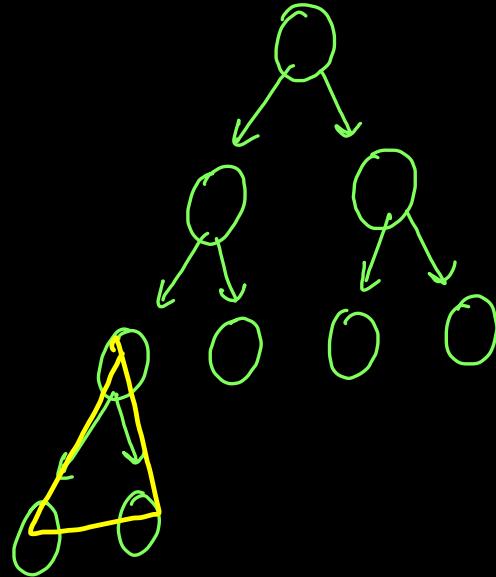
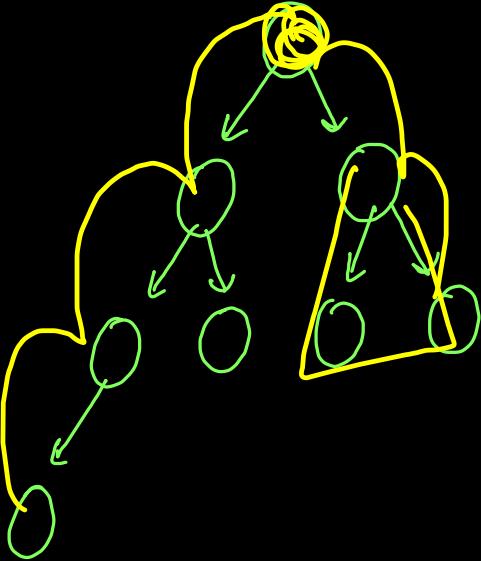
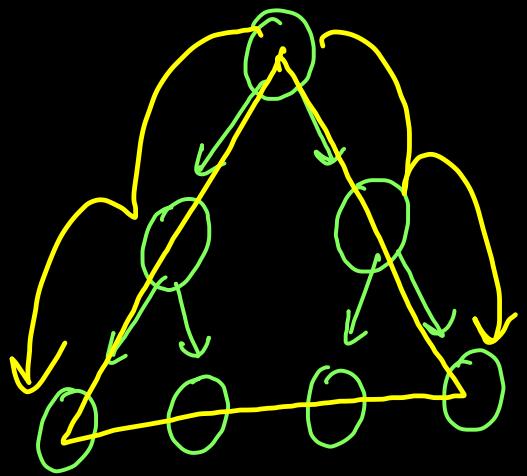
Time → linear
Space → linear

Complete Binary Tree → Order of insertion will be
top to bottom left to right

↳ all levels (except the last) are complete
& last level nodes are as left as possible



```
// approach 1: dfs -> O(n) time, O(log n) space
public int size(TreeNode root){
    if(root == null) return 0;
    return 1 + size(root.left) + size(root.right);
}
```



$$0 \rightarrow 1$$

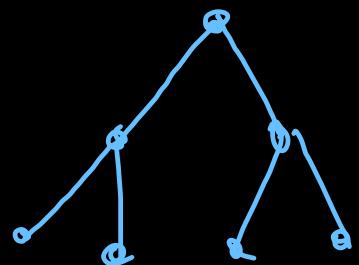


$$\rightarrow 1 + 2 = 3$$

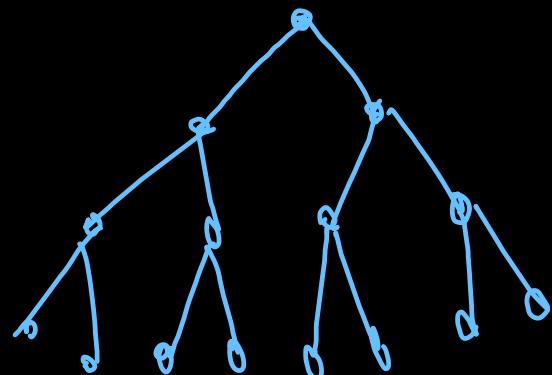
height = h { Perfect binary tree }

$$= 2^0 + 2^1 + 2^2 + \dots + 2^{h-1}$$

$$= 2^h - 1 \{ \text{nodes} \}$$



$$\rightarrow 1 + 2 + 2^2 = 7$$



$$\rightarrow 1 + 2 + 2^2 + 2^3 \\ = 15$$

complete binary tree is always balanced $\Rightarrow h = \log_2 n$

```

public int leftMost(TreeNode root){
    if(root == null) return 0;
    return 1 + leftMost(root.left);
}

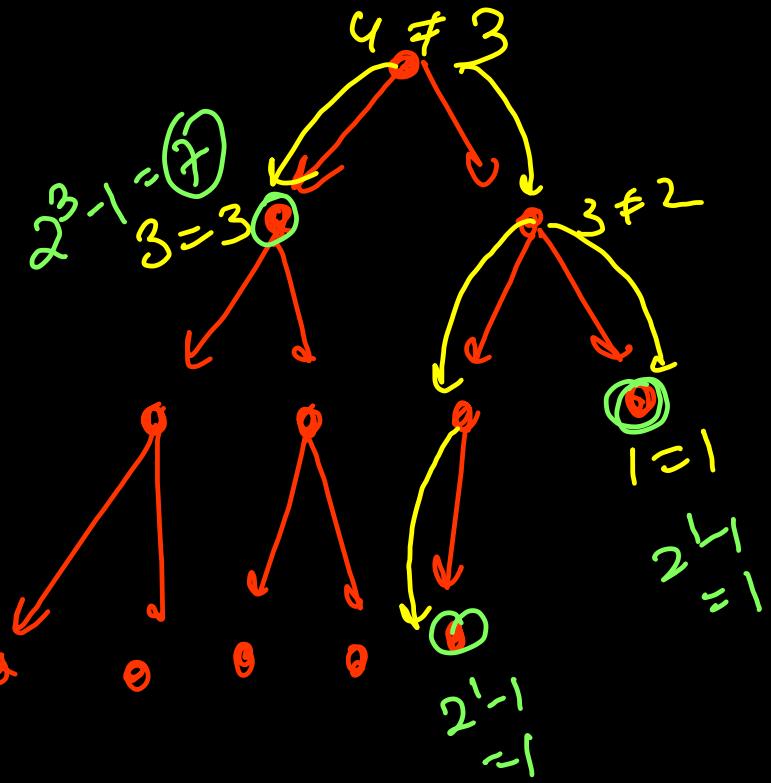
public int rightMost(TreeNode root){
    if(root == null) return 0;
    return 1 + rightMost(root.right);
}

public int countNodes(TreeNode root) {
    if(root == null) return 0;
    int lh = leftMost(root);
    int rh = rightMost(root); } O(h)

    if(lh == rh) return (1 << lh) - 1;
    // perfect binary tree

    return countNodes(root.left) + countNodes(root.right) + 1;
}

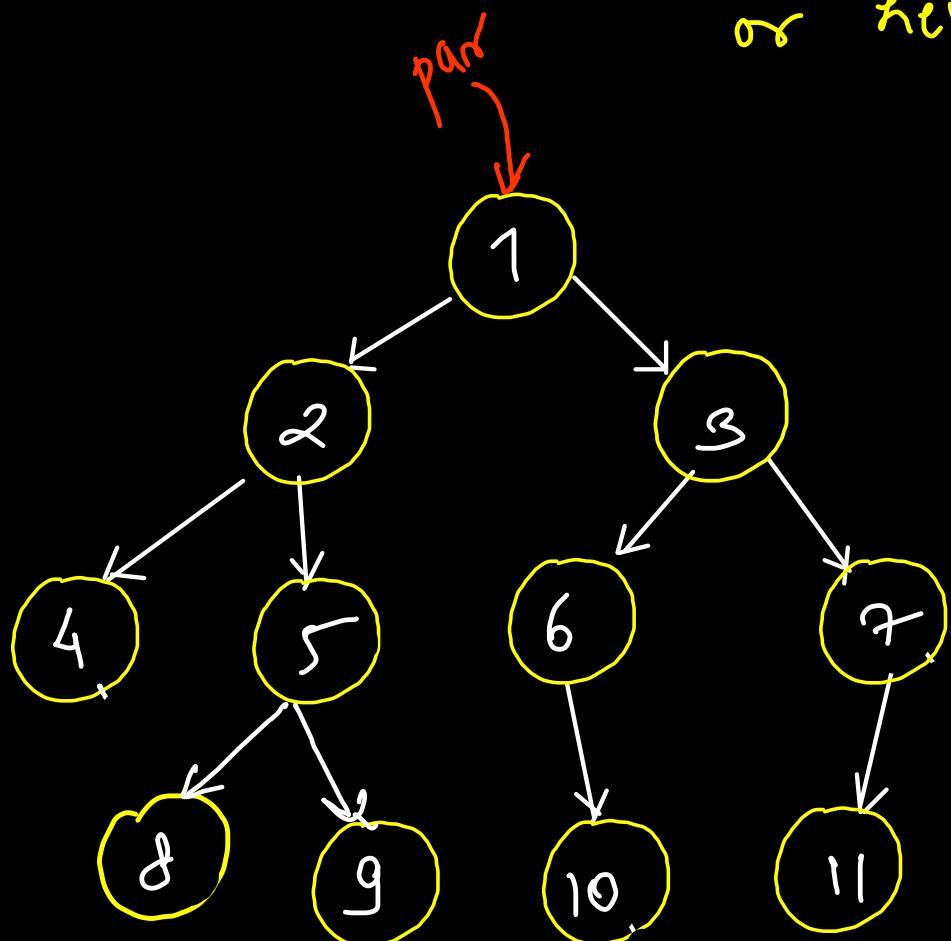
```



Total time $\Rightarrow O(h^2)$
 $\Rightarrow O((\log n)^2)$
 \Rightarrow

Breadth First Traversal (BFS)

or level order traversal

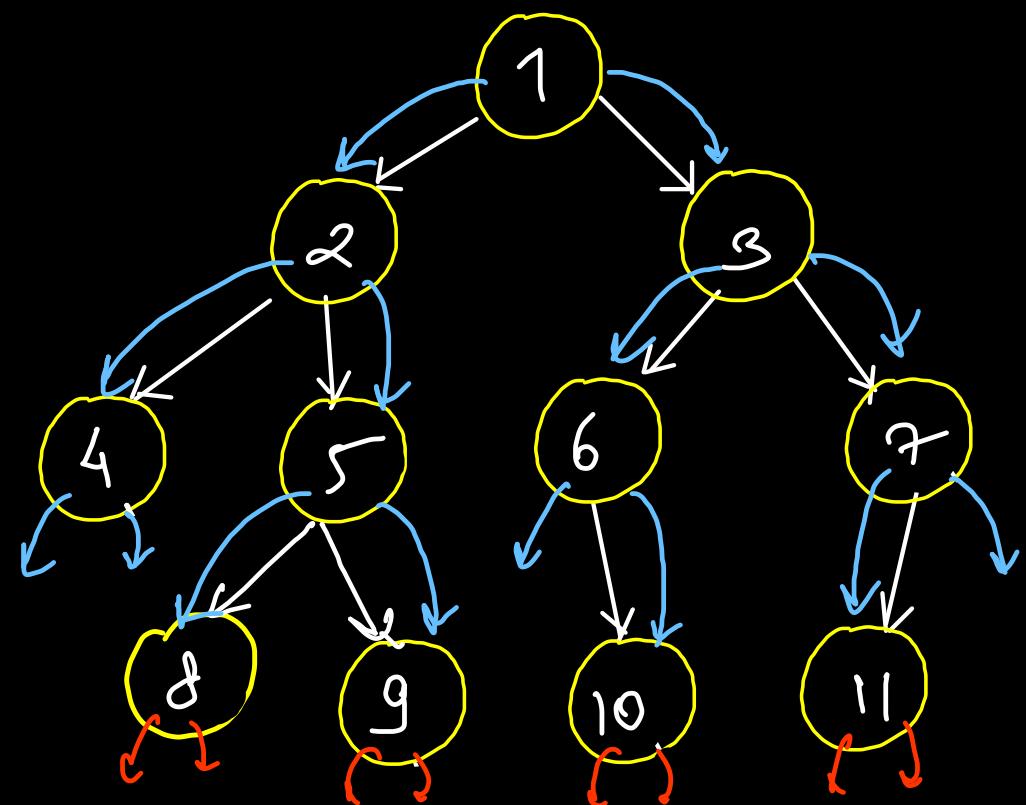
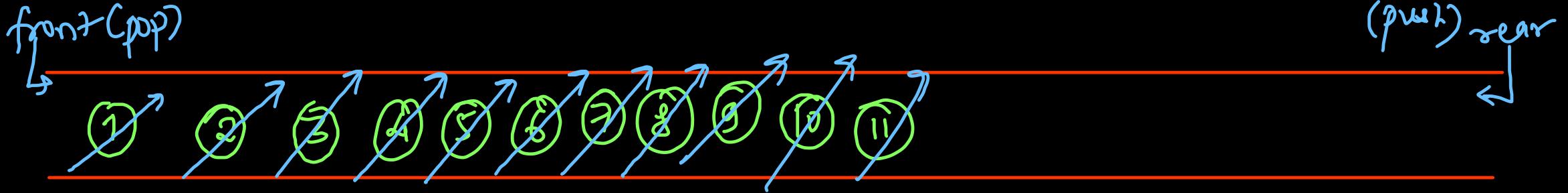


GFG → top to bottom
left to right

{ 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11 }

dfs → recursion → stack

bfs → fifo → queue



1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11

0-1 bfs → Queue will atmost
contain 2 levels

```
static ArrayList <Integer> levelOrder(Node root)
{
    ArrayList<Integer> bfs = new ArrayList<>();
    if(root == null) return bfs;

    Queue<Node> q = new ArrayDeque<>();
    q.add(root);

    while(q.size() > 0) {
        root = q.remove();
        bfs.add(root.data);

        if(root.left != null) q.add(root.left);
        if(root.right != null) q.add(root.right);
    }

    return bfs;
}
```

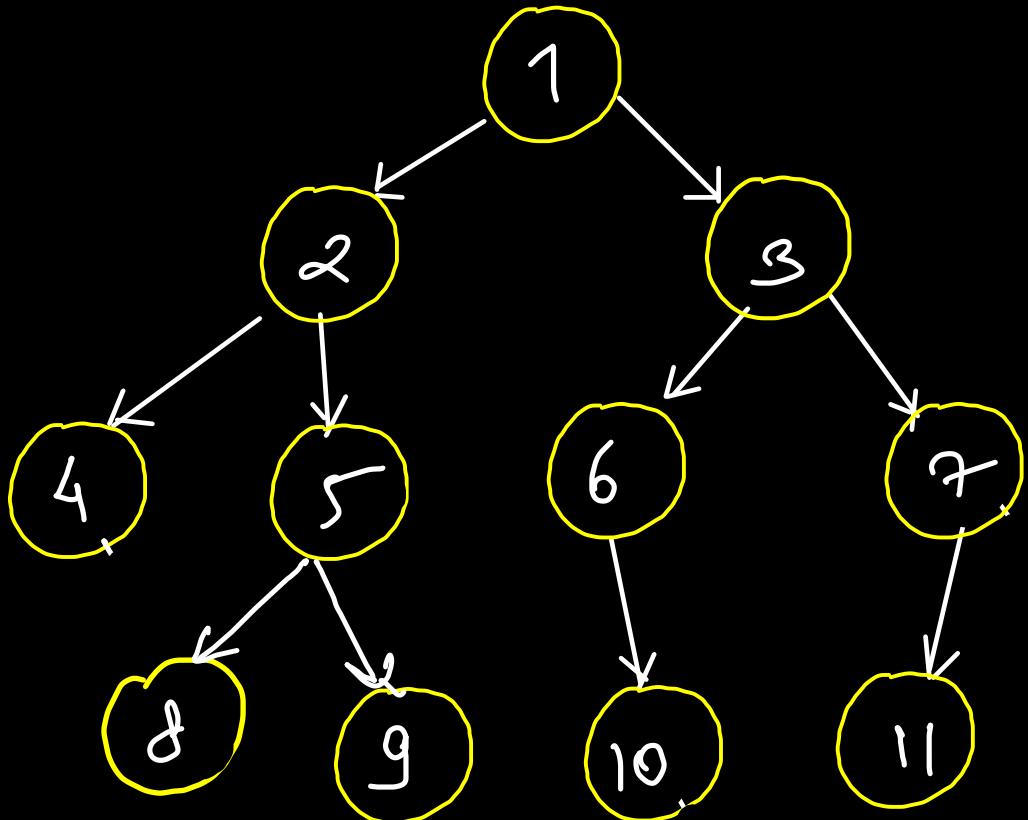
Time Complexity
↳ $O(n)$

Space Complexity
→ Queue Auxiliary Data Structure
→ $O(n)$

Leetcode 102

Level order level by level

top to down, left to right, level
by level



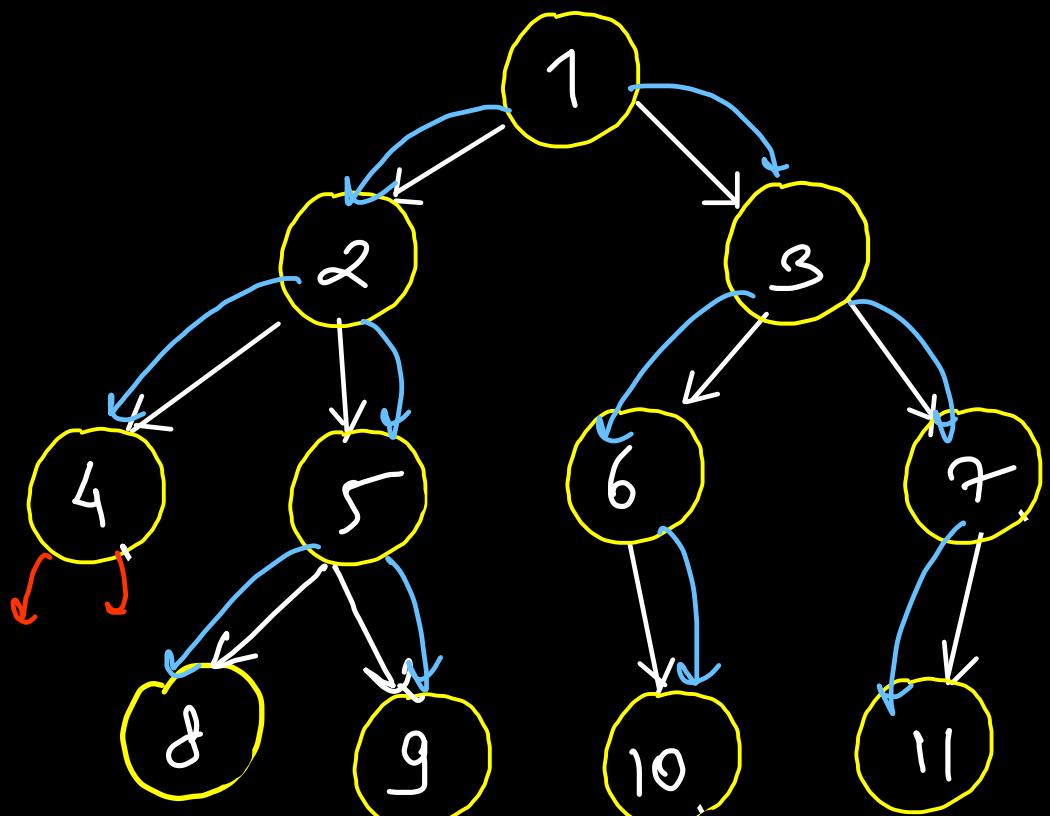
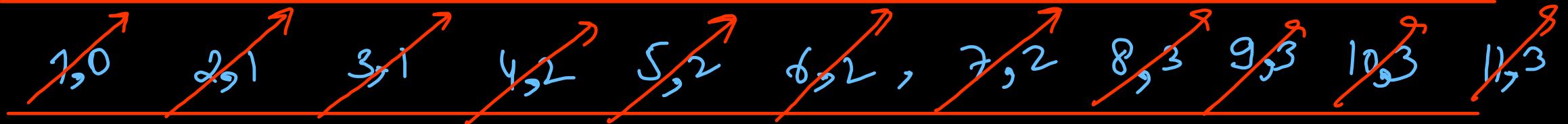
{1}

{2,3}

{4,5,6,7}

{8,9,10,11}

Approach 1) Queue → node
 height/level/depth → Pair



$\{ \{1\}, \{2, 3\}, \{4, 5, 6, 7\}, \{8, 9, 10, 11\} \}$
 $4k \quad 6k \quad 8k \quad 10k$
 $\{ 4k, 6k, 8k, 10k \}$

```
public static class Pair{
    TreeNode root;
    int level;
    Pair(TreeNode root, int level){
        this.root = root;
        this.level = level;
    }
}
```

Time $\Rightarrow O(n)$ } bfs
Space $\Rightarrow O(n)$

Approach 1) Using pair

```
public List<List<Integer>> levelOrder(TreeNode root) {
    List<List<Integer>> levels = new ArrayList<>();
    if(root == null) return levels;

    Queue<Pair> q = new ArrayDeque<>();
    q.add(new Pair(root, 0));
    int level = -1;

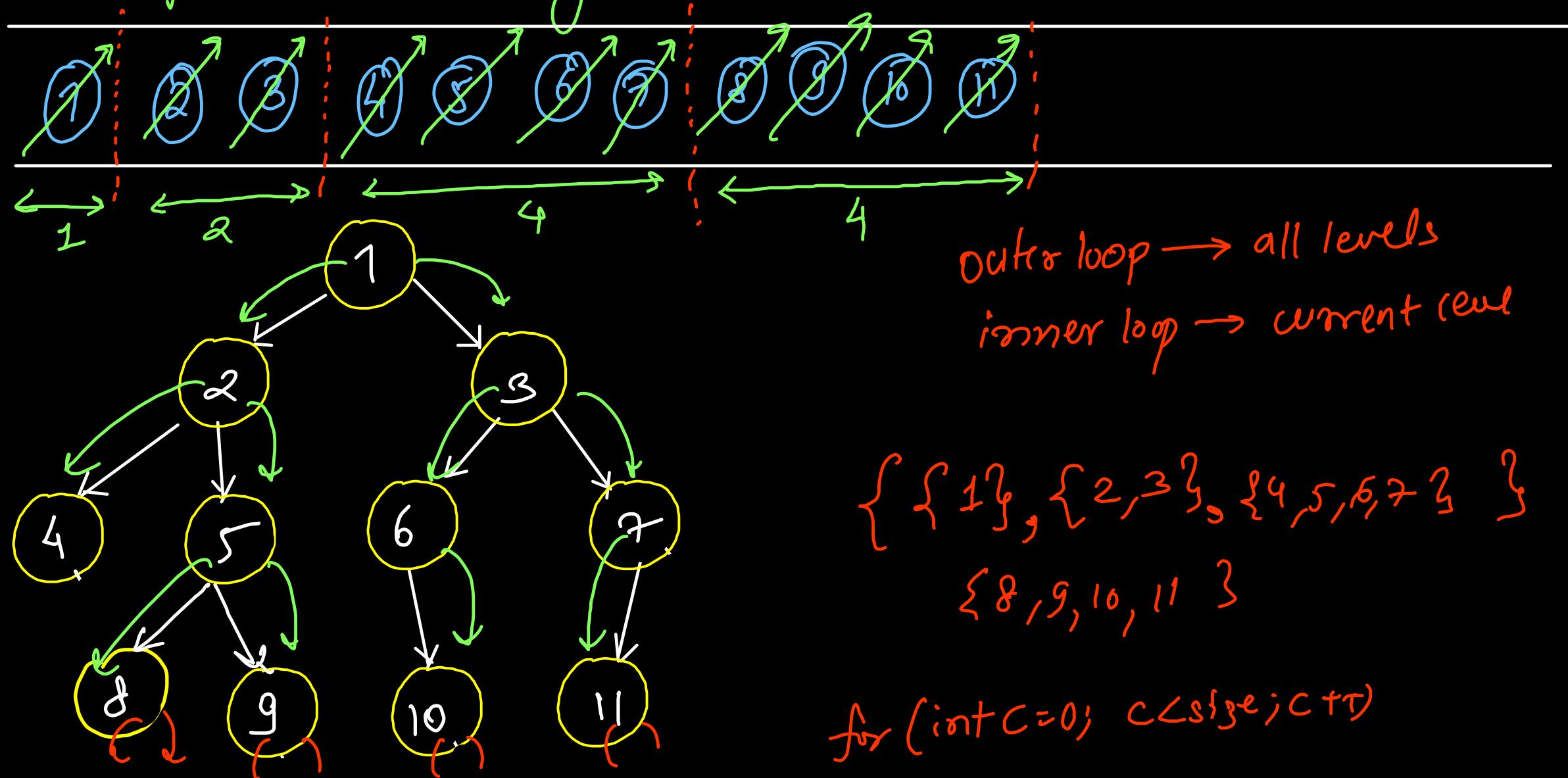
    while(q.size() > 0){
        Pair p = q.remove();

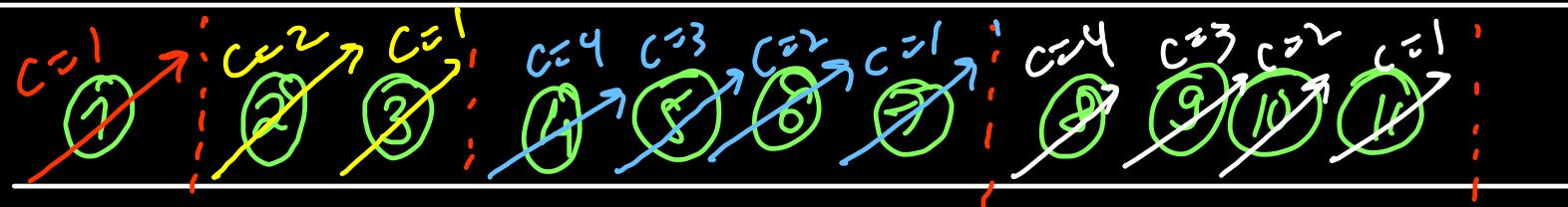
        if(p.level > level){
            levels.add(new ArrayList<>());
            level = p.level;
        }
        levels.get(levels.size() - 1).add(p.root.val);

        if(p.root.left != null)
            q.add(new Pair(p.root.left, p.level + 1));
        if(p.root.right != null)
            q.add(new Pair(p.root.right, p.level + 1));
    }

    return levels;
}
```

Approach 2) Using count/nested loops





```

public List<List<Integer>> levelOrder(TreeNode root) {
    List<List<Integer>> levels = new ArrayList<>();
    if(root == null) return levels;

    Queue<TreeNode> q = new ArrayDeque<>();
    q.add(root);
    level by level
    while(q.size() > 0) {
        List<Integer> level = new ArrayList<>();
        for(int c = q.size(); c > 0; c--) {
            root = q.remove();
            level.add(root.val);

            if(root.left != null) q.add(root.left);
            if(root.right != null) q.add(root.right);
        }
        levels.add(level);
    }

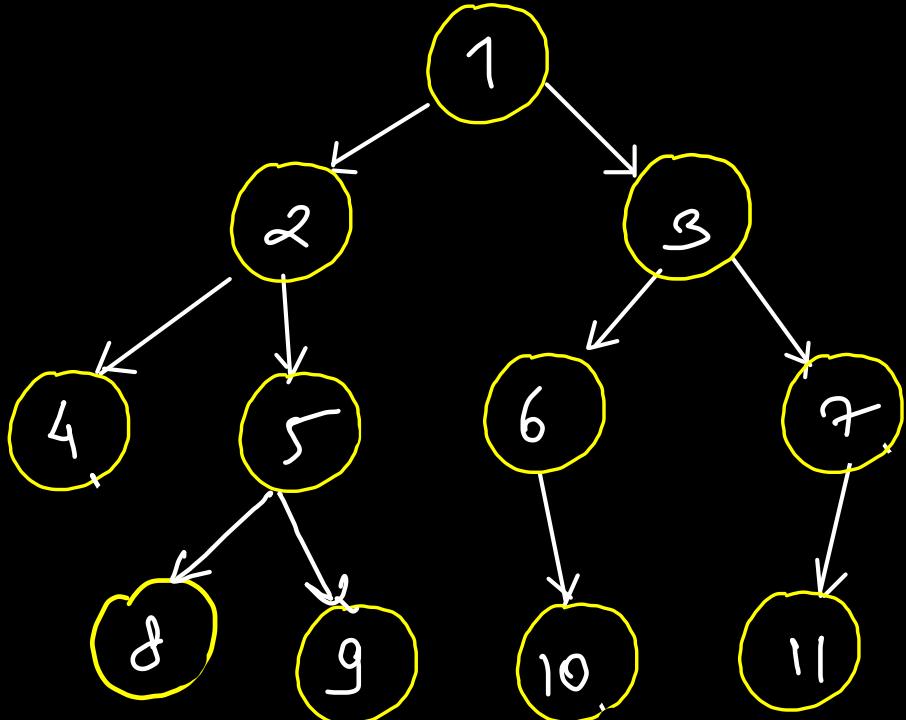
    return levels;
}

```

current level

Time $\Rightarrow O(n)$
linear

Space $\Rightarrow O(n)$
queue



{ {1} } {2, 3} {4, 5, 6, 7} {8, 9, 11} }

```
class Solution {
    public List<List<Integer>> levelOrder(TreeNode root) {
        List<List<Integer>> levels = new ArrayList<>();
        if(root == null) return levels;

        TreeNode marker = new TreeNode(-1);
        Queue<TreeNode> q = new ArrayDeque<>();
        q.add(root); q.add(marker);

        List<Integer> level = new ArrayList<>();
        while(q.size() > 1 || q.peek() != marker){
            root = q.remove();

            if(root != marker){
                level.add(root.val);
                if(root.left != null) q.add(root.left);
                if(root.right != null) q.add(root.right);
            }
            else {
                q.add(marker);
                levels.add(level);
                level = new ArrayList<>();
            }
        }

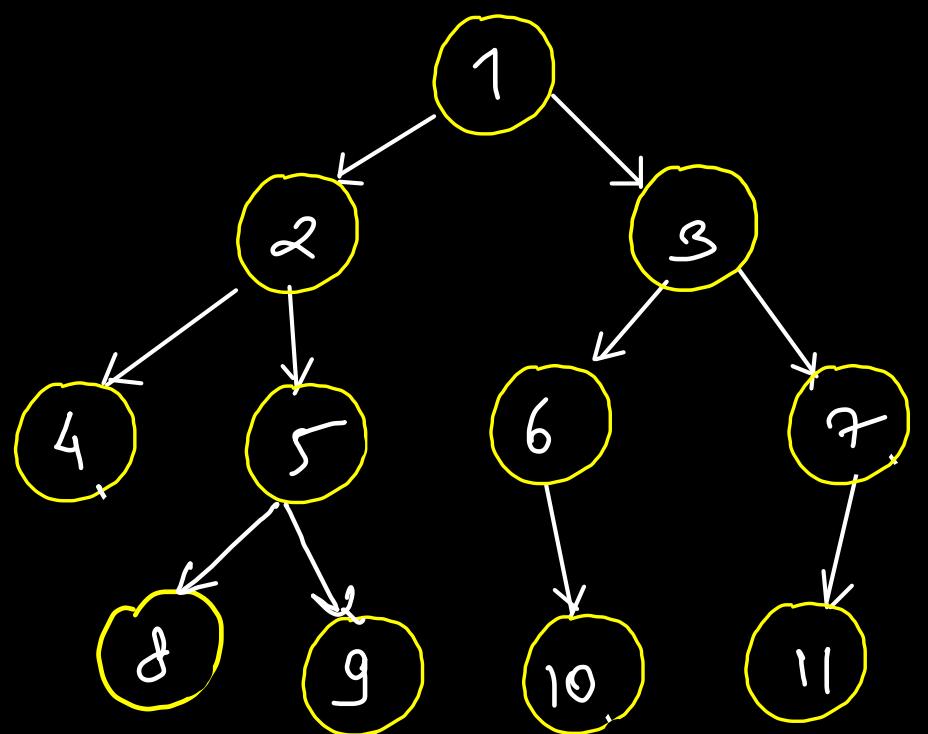
        levels.add(level);
        return levels;
    }
}
```

Approach3) using
marker/dummy/delimiter

Time $\Rightarrow \mathcal{O}(n)$

Space $\Rightarrow \mathcal{O}(n)$

Ex (0-7) Level order (Bottom to top)



$\{ \{1\}, \{2, 3\}, \{4, 5, 6, 7\}, \{8, 9, 10, 11\} \}$

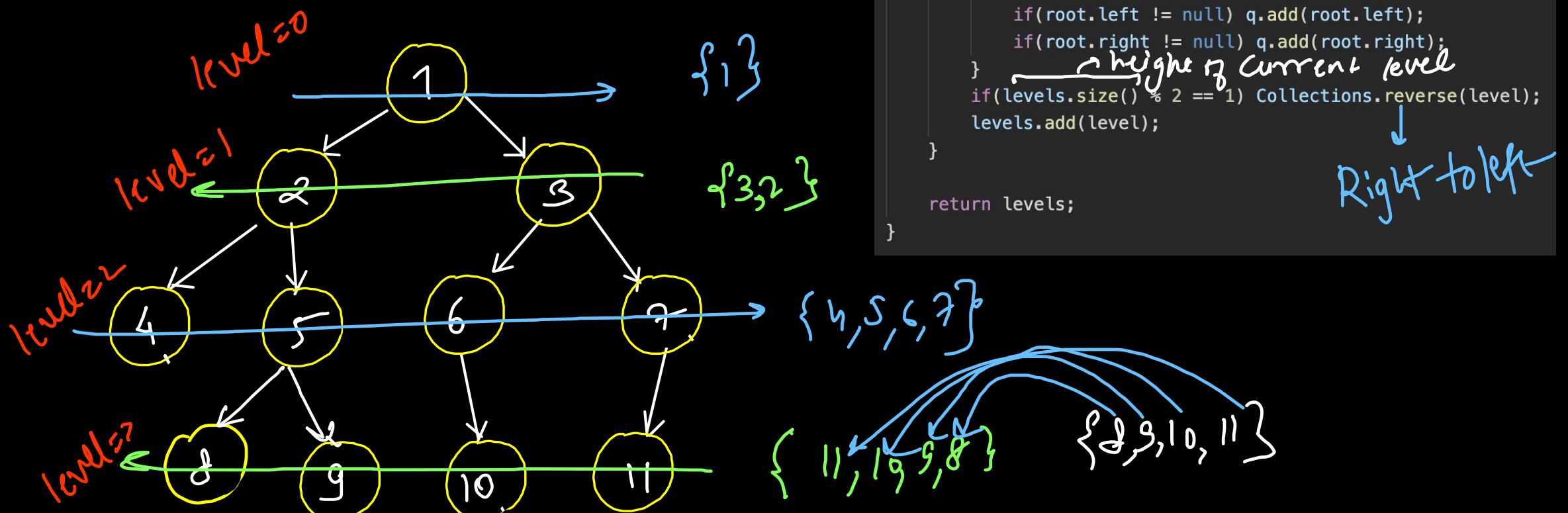
$\underline{O(h)}$ ↓ Collections.reverse(levels)

$\{ \{8, 9, 10, 11\}, \{4, 5, 6, 7\}, \{2, 3\}, \{1\} \}$

Time $\Rightarrow O(n)$

Space $\Rightarrow O(n)$

LC 103
Zigzag
Level order



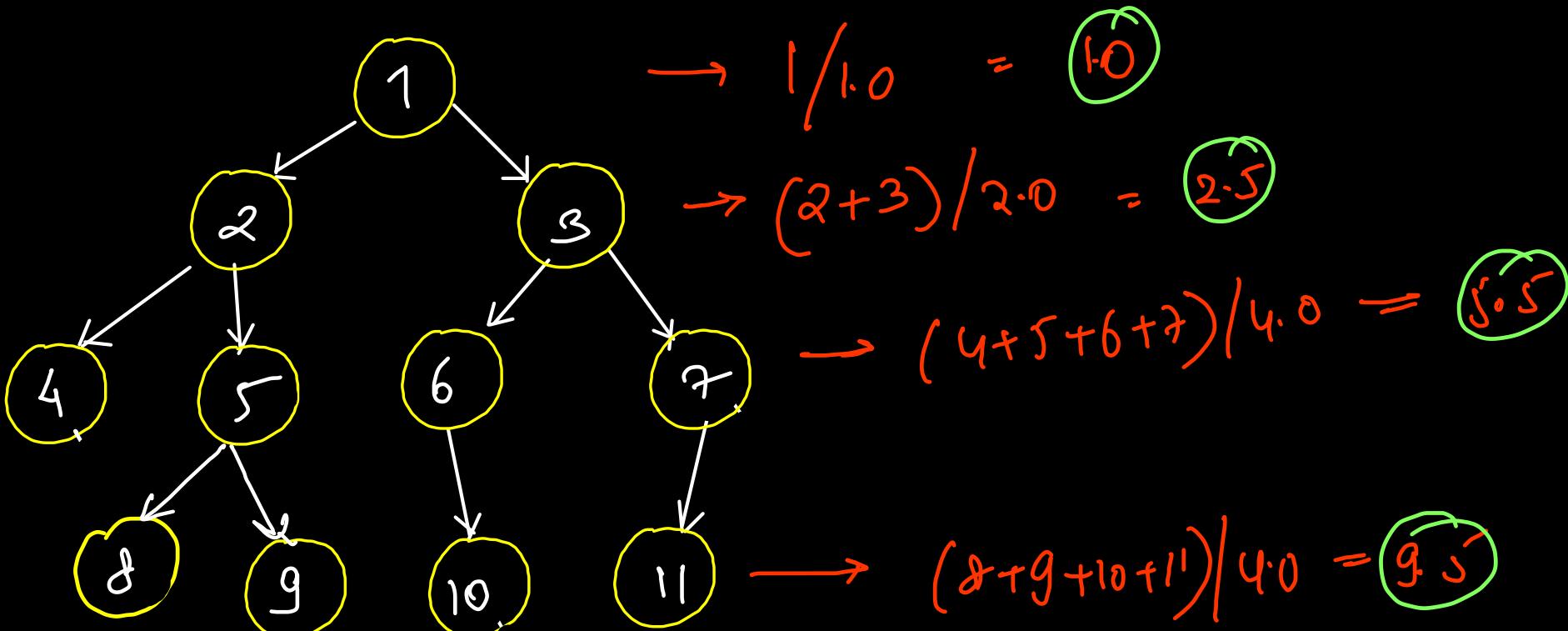
```
public List<List<Integer>> zigzagLevelOrder(TreeNode root) {
    List<List<Integer>> levels = new ArrayList<>();
    if(root == null) return levels;

    Queue<TreeNode> q = new ArrayDeque<>();
    q.add(root);
    while(q.size() > 0)
    {
        List<Integer> level = new ArrayList<>();
        for(int c = q.size(); c > 0; c--)
        {
            root = q.remove();
            level.add(root.val);

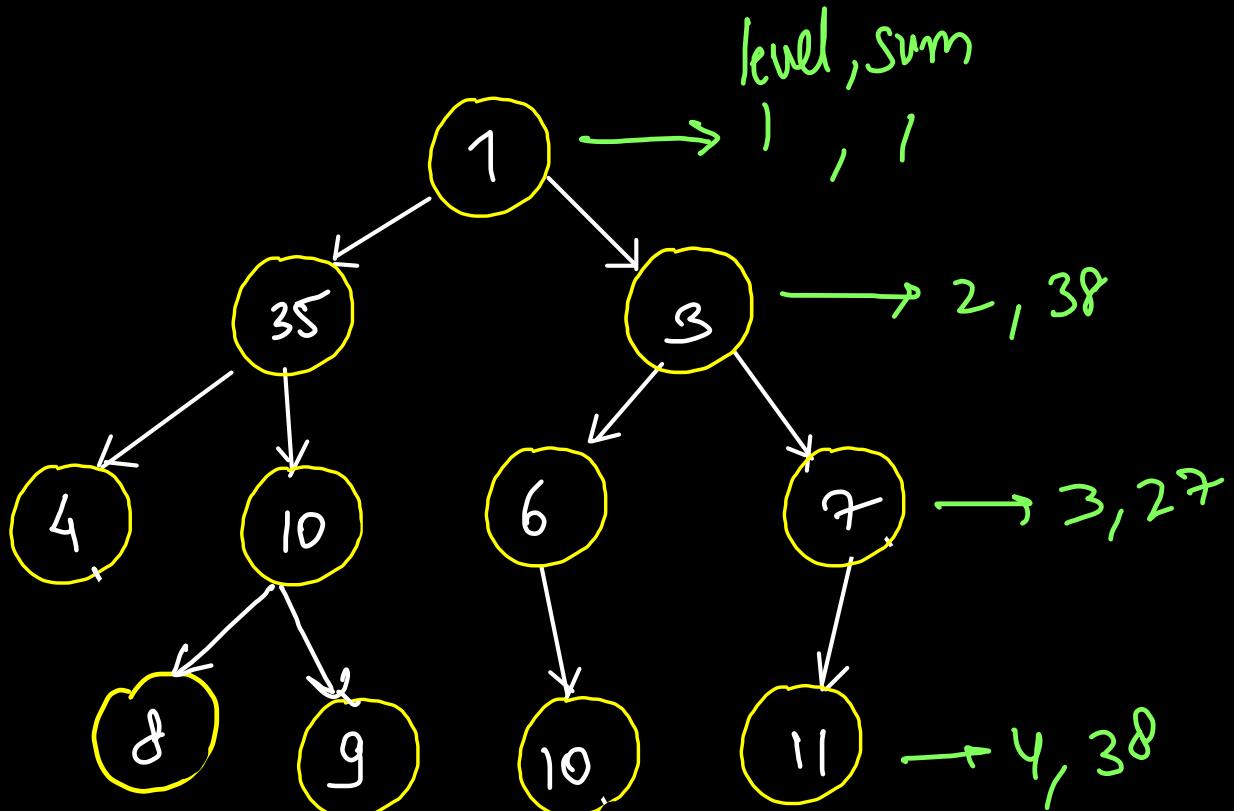
            if(root.left != null) q.add(root.left);
            if(root.right != null) q.add(root.right);
        }
        if(levels.size() % 2 == 1) Collections.reverse(level);
        levels.add(level);
    }
    return levels;
}
```

height of current level
Right to left

LC 637) Average of levels



LC 1161) Max^m level sum



$$\text{maxSum} = \cancel{0}/\cancel{1} 38$$

$$\text{maxLevel} = \cancel{0}/\cancel{1} 2$$

time = $O(n)$, space = $O(n)$

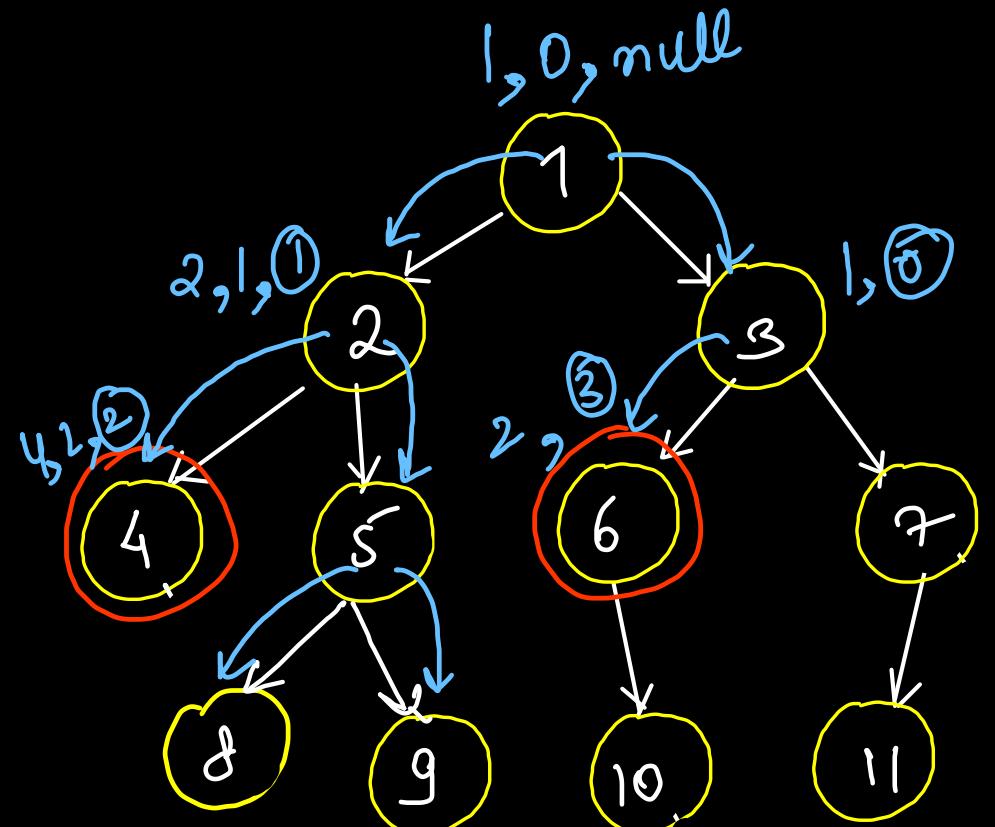
```
public List<Double> averageOfLevels(TreeNode root) {  
    List<Double> levels = new ArrayList<>();  
    if(root == null) return levels;  
  
    Queue<TreeNode> q = new ArrayDeque<>();  
    q.add(root);  
  
    while(q.size() > 0) {  
        double sum = 0, count = q.size();  
        for(int c = q.size(); c > 0; c--) {  
            root = q.remove();  
            sum += root.val;  
            if(root.left != null) q.add(root.left);  
            if(root.right != null) q.add(root.right);  
        }  
        levels.add(sum / count);  
    }  
  
    return levels;  
}
```

LC 637)

time = $O(n)$, space = $O(n)$

```
public int maxLevelSum(TreeNode root) {  
    Queue<TreeNode> q = new ArrayDeque<>();  
    q.add(root);  
  
    int maxsum = Integer.MIN_VALUE, maxlevel = 0;  
    int currsum = 0, currlevel = 0;  
    while(q.size() > 0) {  
        currsum = 0; currlevel++;  
  
        for(int c = q.size(); c > 0; c--) {  
            root = q.remove();  
            currsum += root.val;  
            if(root.left != null) q.add(root.left);  
            if(root.right != null) q.add(root.right);  
        }  
  
        if(currsum > maxsum){  
            maxsum = currsum;  
            maxlevel = currlevel;  
        }  
    }  
  
    return maxlevel;
```

LC 1161)



$x.parent = 2$	$y.parent = 3$
$x.level = 2$	$y.level = 2$

2,3 → sibling false

4,6 → cousins true

5,7 → cousins true

7,10 → neither cousin
nor siblings. false

$x.parent \neq y.parent$ &
 $x.level = y.level$

```
class Solution {
    int xparent = -1, yparent = -1, xlevel = -1, ylevel = -1;

    public void dfs(TreeNode root, int level, int parent, int x, int y){
        if(root == null) return;

        if(root.val == x) {
            xparent = parent;
            xlevel = level;
        }

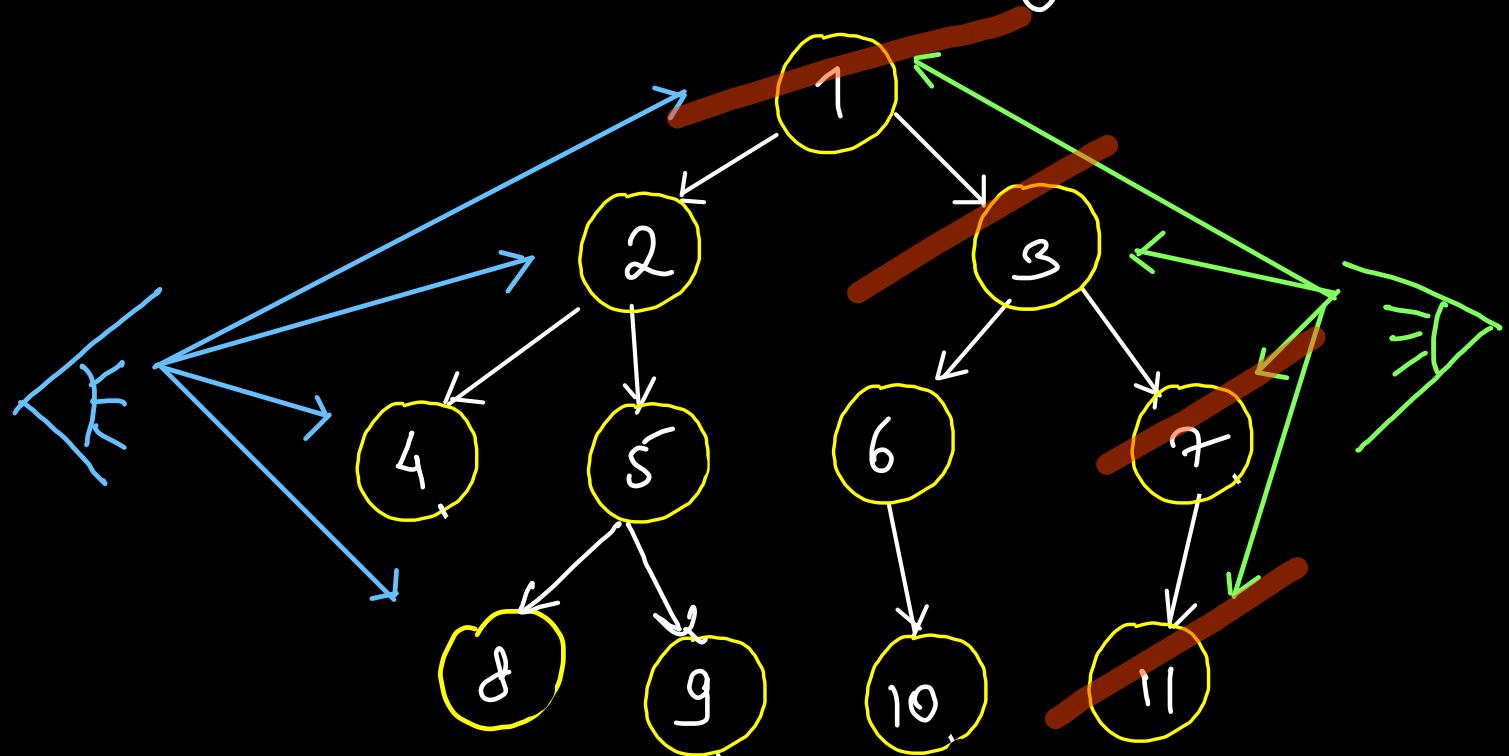
        if(root.val == y){
            yparent = parent;
            ylevel = level;
        }

        dfs(root.left, level + 1, root.val, x, y);
        dfs(root.right, level + 1, root.val, x, y);
    }

    public boolean isCousins(TreeNode root, int x, int y) {
        dfs(root, 0, -1, x, y);
        return (xparent != yparent) && (xlevel == ylevel);
    }
}
```

Time = $O(n)$
Space = $O(h)$

Left View & Right view



left view $\rightarrow 1, 2, 4, 8$

(leftmost/first node) of each level

right

view $\rightarrow 1, 3, 7, 11$

(rightmost/last node) of each level

Left View

```
ArrayList<Integer> leftView(Node root){  
    ArrayList<Integer> left = new ArrayList<>();  
    if(root == null) return left;  
  
    Queue<Node> q = new ArrayDeque<>();  
    q.add(root);  
  
    while(q.size() > 0){  
        left.add(q.peek().data);  
        for(int c = q.size(); c > 0; c--) {  
            root = q.remove();  
            if(root.left != null) q.add(root.left);  
            if(root.right != null) q.add(root.right);  
        }  
    }  
  
    return left;  
}
```

Right View

```
public List<Integer> rightSideView(TreeNode root) {  
    List<Integer> right = new ArrayList<>();  
    if(root == null) return right;  
  
    Queue<TreeNode> q = new ArrayDeque<>();  
    q.add(root);  
  
    while(q.size() > 0){  
        for(int c = q.size(); c > 0; c--) {  
            root = q.remove();  
            if(root.left != null) q.add(root.left);  
            if(root.right != null) q.add(root.right);  
        }  
        right.add(root.val);  
    }  
  
    return right;  
}
```

Time = $O(n)$

Space = $O(n)$

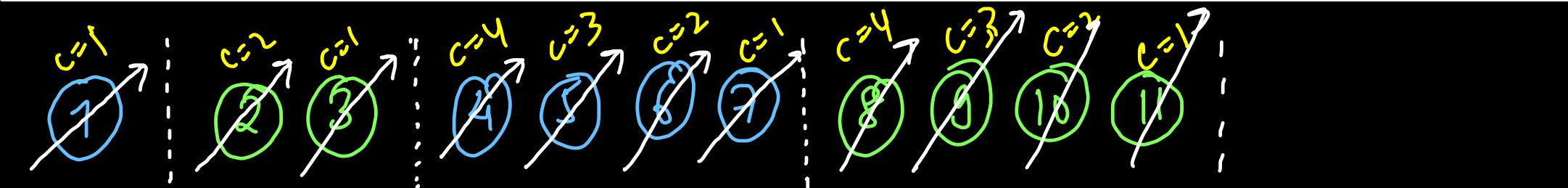
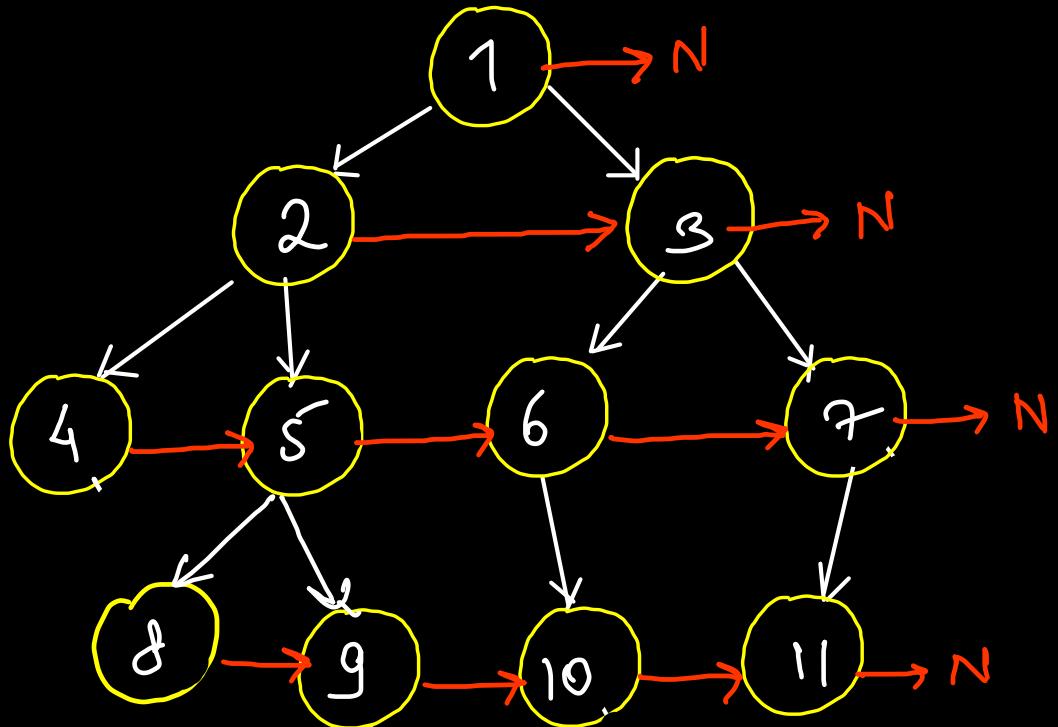
Populating next right pointer

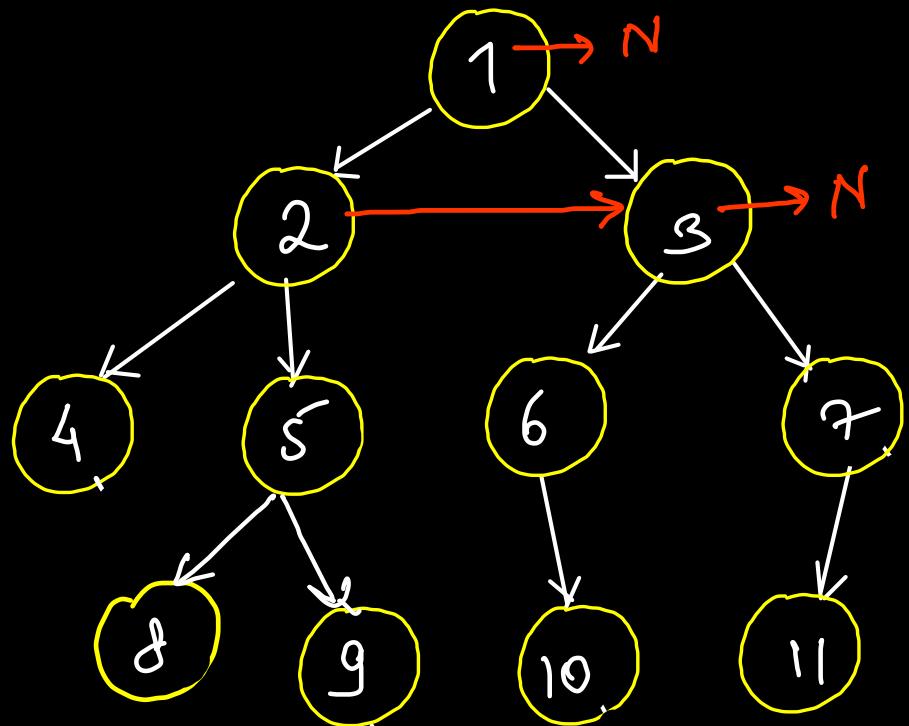
fc 116 & 117

Approach 1)

Queue \rightarrow Time = $O(n)$
Space = $O(n)$

bfs level order





```

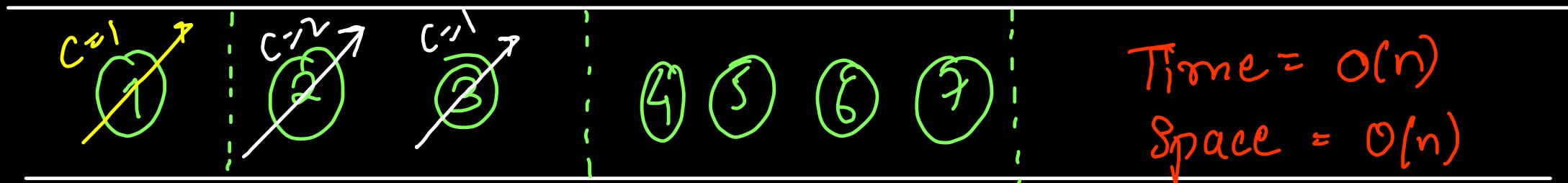
public Node connect(Node root) {
    if(root == null) return null;

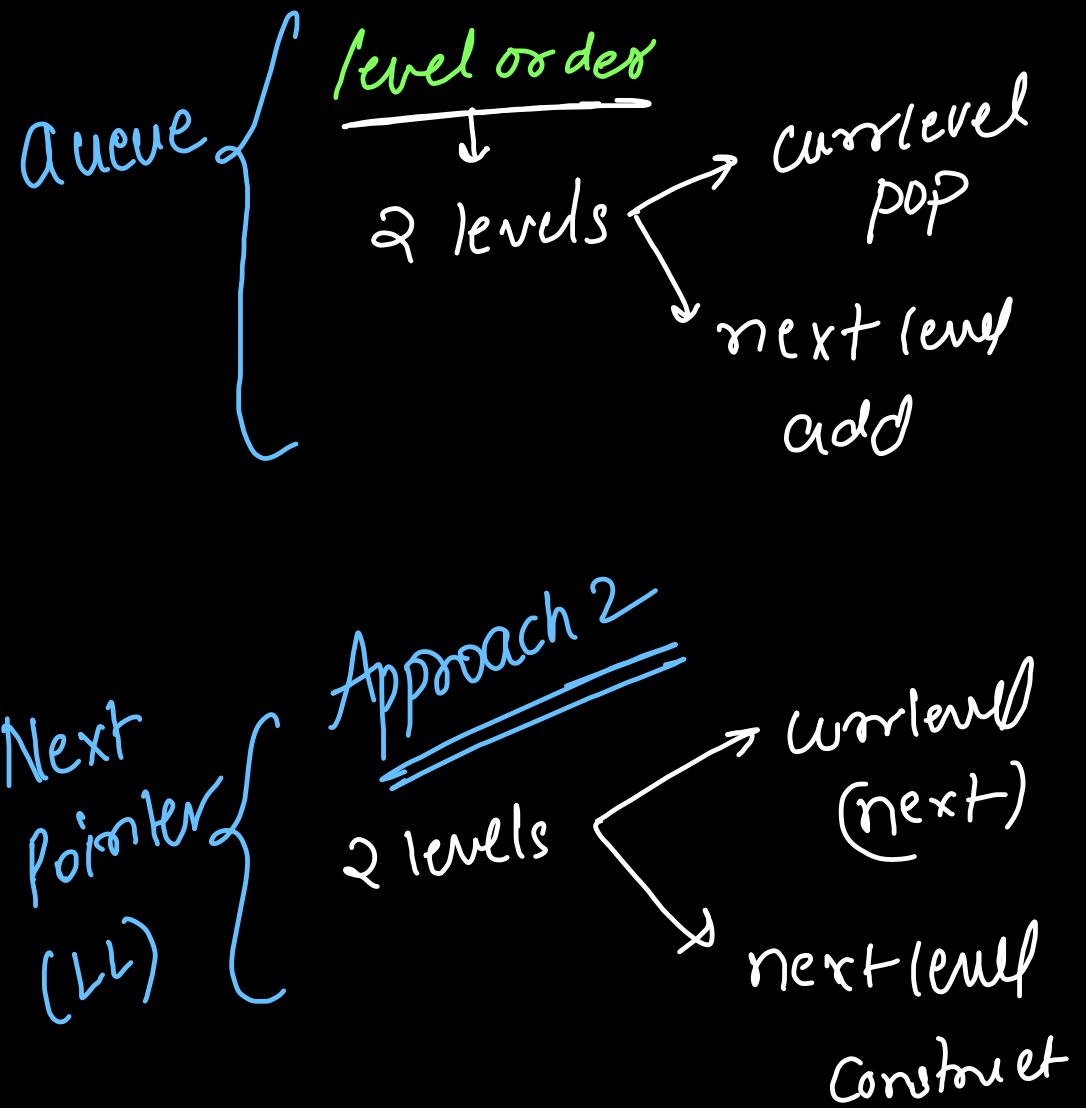
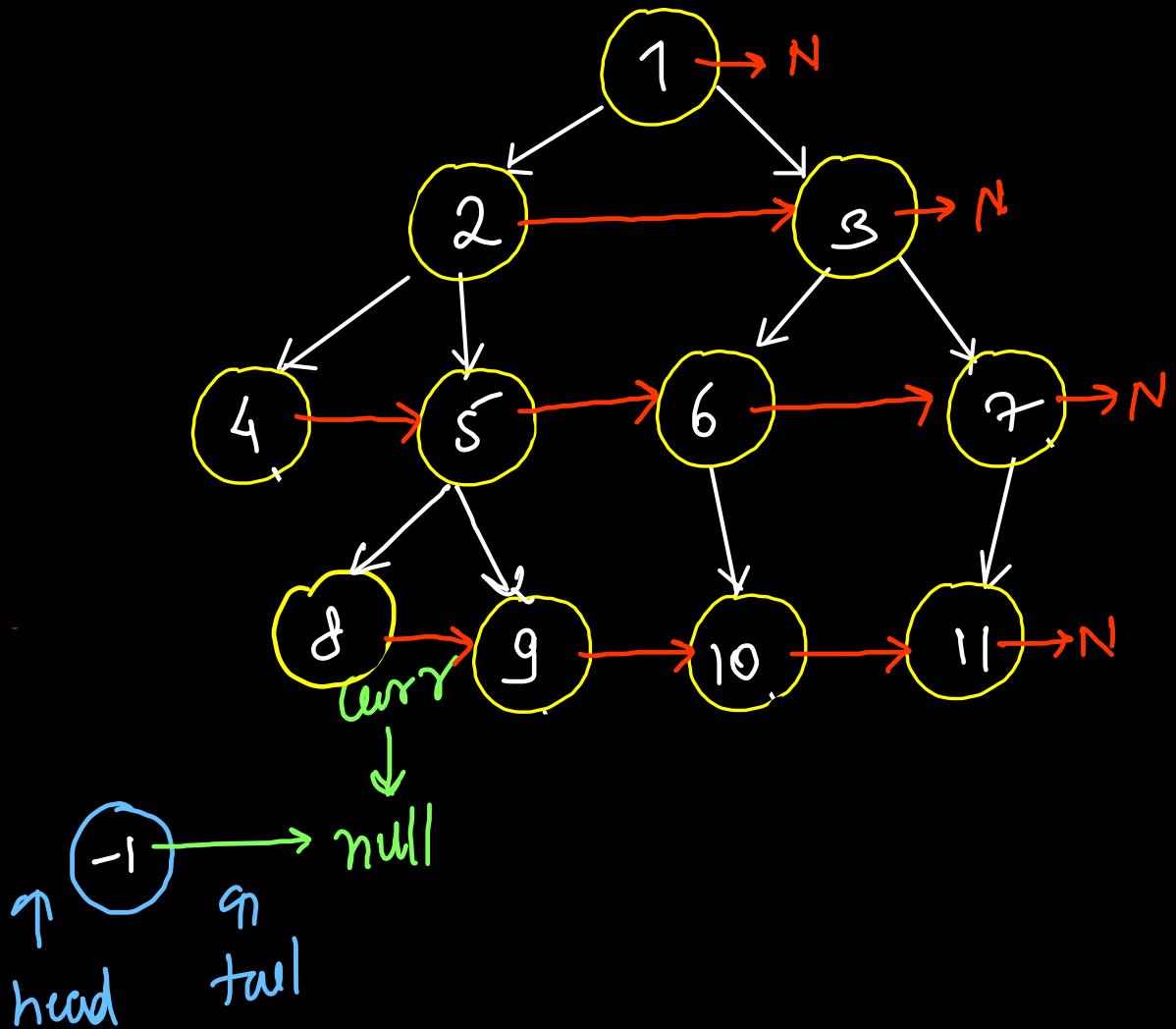
    Queue<Node> q = new ArrayDeque<>();
    q.add(root);

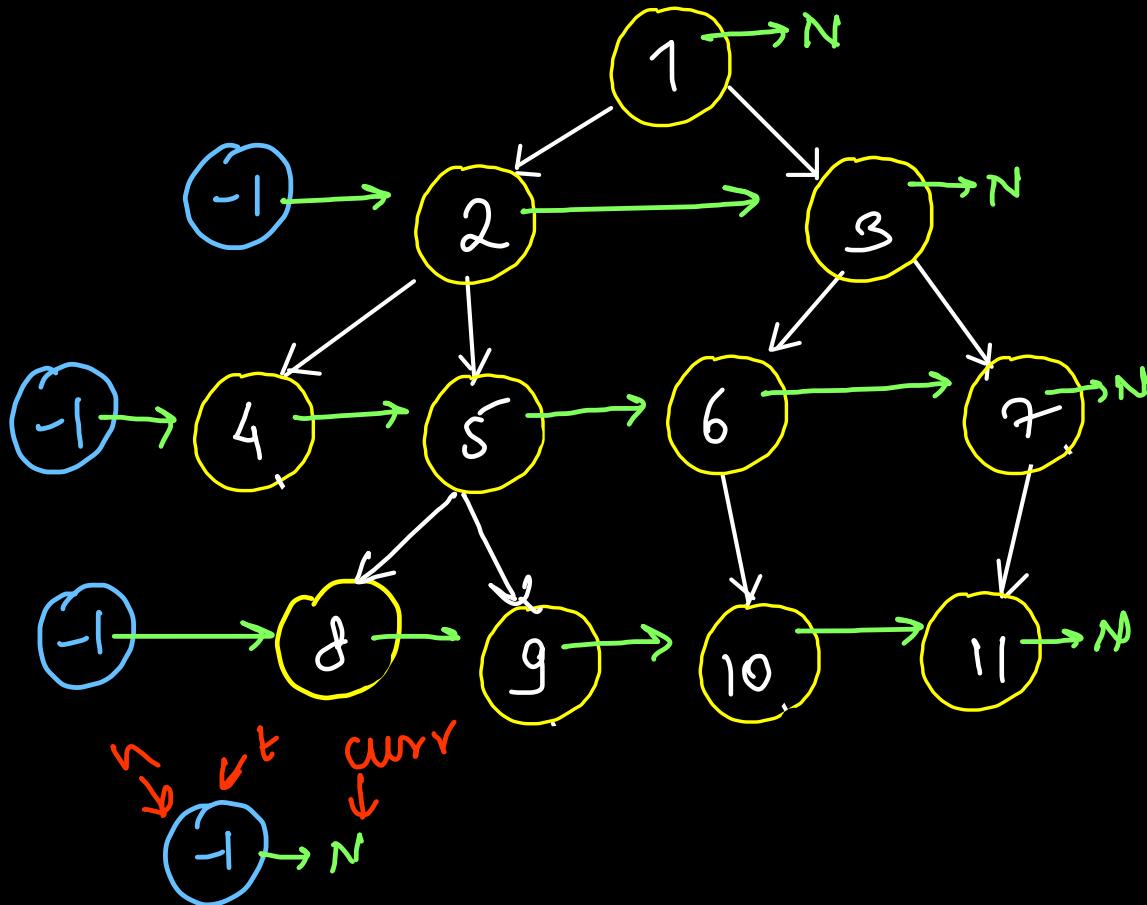
    while(q.size() > 0){
        for(int c = q.size(); c > 0; c--){
            Node curr = q.remove();
            if(c > 1) curr.next = q.peek();

            if(curr.left != null) q.add(curr.left);
            if(curr.right != null) q.add(curr.right);
        }
    }

    return root;
}
  
```







```

public Node connect(Node root) {
    Node curr = root;
    while(curr != null)
    {
        Node head = new Node(-1);
        Node tail = head;
        for( ; curr != null; curr = curr.next){
            if(curr.left != null){
                tail.next = curr.left;
                tail = tail.next;
            }
            if(curr.right != null) {
                tail.next = curr.right;
                tail = tail.next;
            }
        }
        curr = head.next;
    }
    return root;
}

```

Time = O(n)
Space = O(1)

```
public Node connect(Node root) {  
    Node curr = root;  
    while(curr != null)  
    {  
        Node head = new Node(-1);  
        Node tail = head;  
        for( ; curr != null; curr = curr.next){  
            if(curr.left != null)  
                tail = tail.next = curr.left;  
            if(curr.right != null)  
                tail = tail.next = curr.right;  
        }  
  
        curr = head.next;  
    }  
    return root;  
}
```

Smaller Version

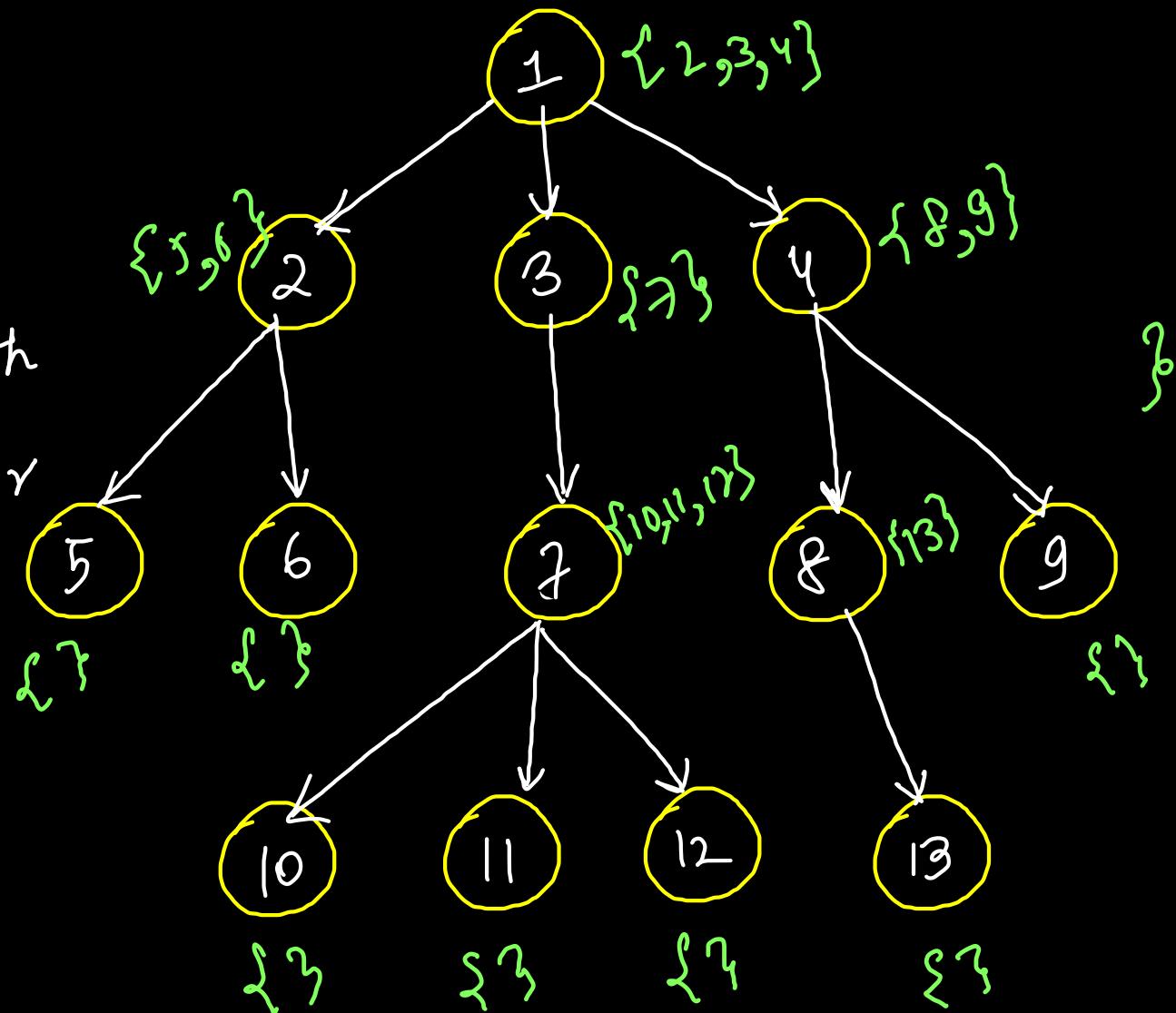
Nary or Generic Tree

LC 589) Preorder

LC 590) Postorder

LC 559) Maxⁿ depth

LC 429) Level order



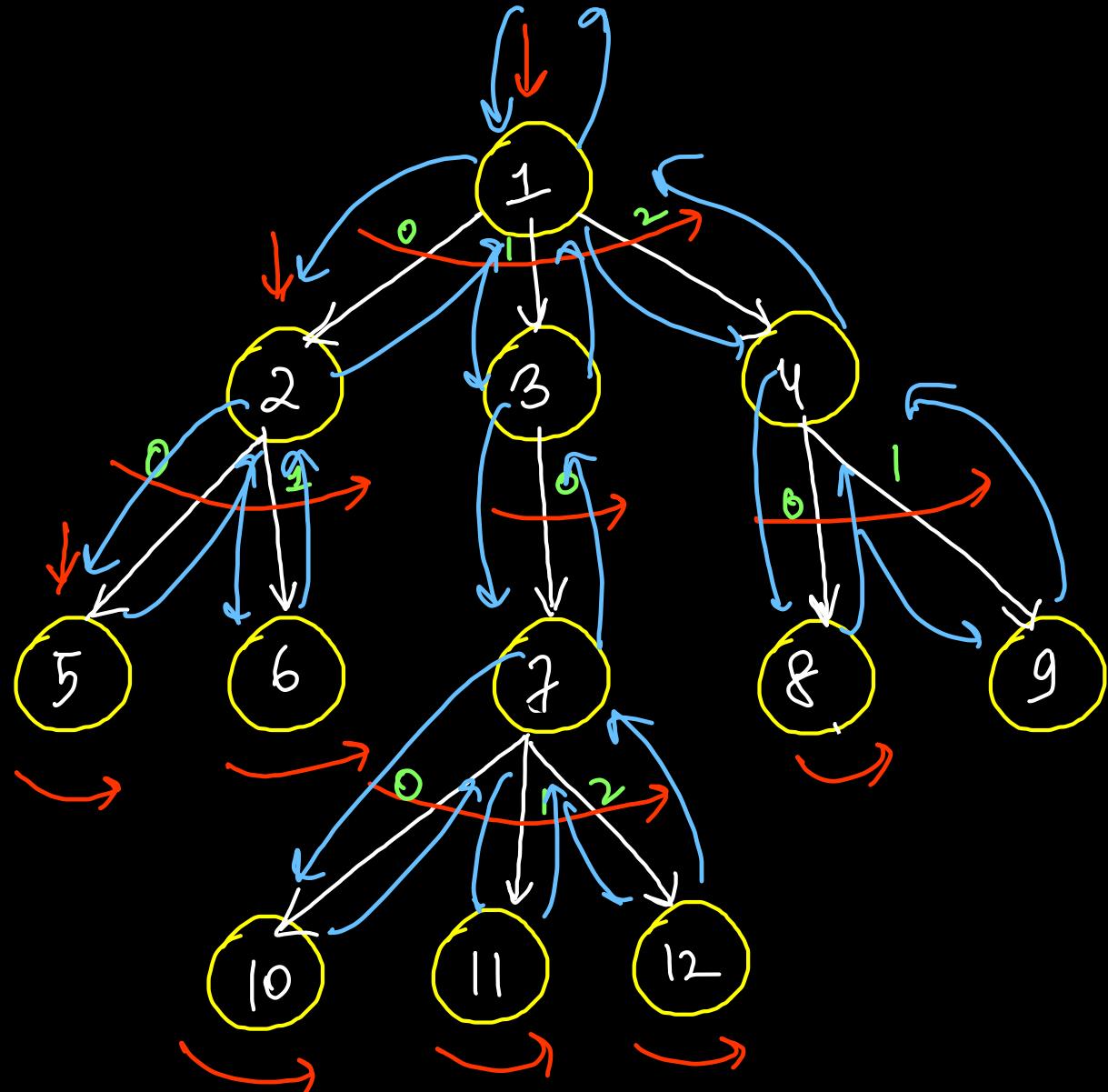
```
// Definition for a Node.
class Node {
    public int val;
    public List<Node> children;

    public Node() {}

    public Node(int _val) {
        val = _val;
    }

    public Node(int _val, List<Node> _children) {
        val = _val;
        children = _children;
    }
};
```

Input Space Complexity
: $O(\text{nodes} + \text{edges})$
 $= O(N + N^{-1})$
 $\approx O(2N - 1)$
Linear



dfs

- Preorder
 - root → children
- Postorder
 - children → root
 - (ltr)

Preorder: →
 1, 2, 5, 6, 3, 7, 10, 11, 12, 4, 8
 Postorder: →
 5, 6, 2, 10, 11, 12, 7, 3, 8, 9, 4, 1

Preorder LC 589)

```
class Solution {  
    List<Integer> ans = new ArrayList<>();  
    public void dfs(Node root){  
        ans.add(root.val);  
        for(Node child: root.children)  
            dfs(child);  
    }  
  
    public List<Integer> preorder(Node root) {  
        if(root != null) dfs(root);  
        return ans;  
    }  
}
```

Time \Rightarrow Recursion + Loop
 \downarrow \downarrow
 $O(n)$ + $O(e) = O(n-1)$

$O(2n-1)$ linear total

Postorder LC 590)

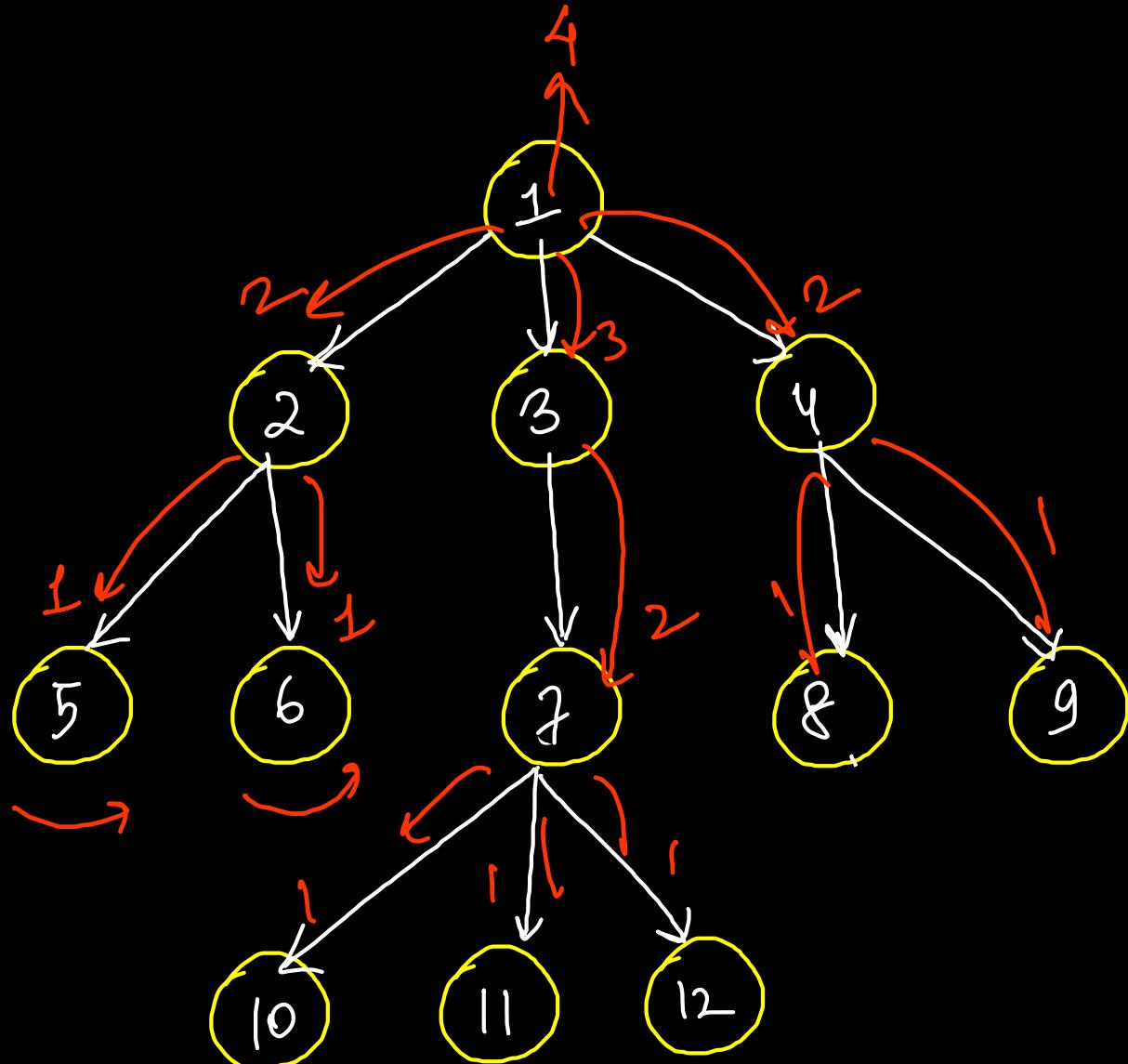
```
class Solution {  
    List<Integer> ans = new ArrayList<>();  
    public void dfs(Node root){  
        for(Node child: root.children)  
            dfs(child);  
        ans.add(root.val);  
    }  
  
    public List<Integer> postorder(Node root) {  
        if(root != null) dfs(root);  
        return ans;  
    }  
}
```

Space \Rightarrow recursion call stack

$O(\text{height})$
avg $O(\log H)$ worst $O(n)$

Max Depth of Nairy Tree

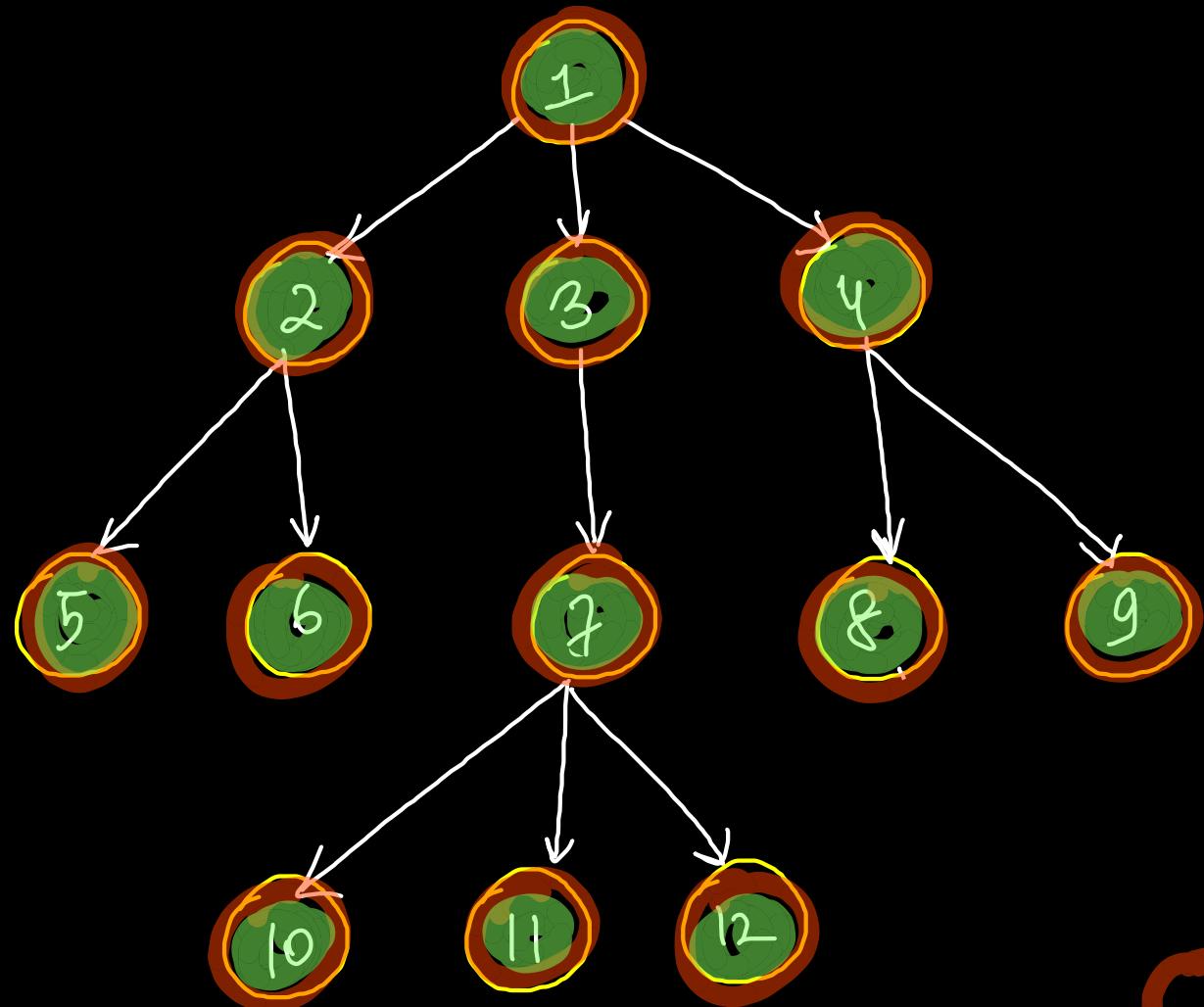
LC 559



```
public int maxDepth(Node root) {  
    if(root == null) return 0;  
  
    int height = 0;  
    for(Node child: root.children){  
        height = Math.max(height, maxDepth(child));  
    }  
    return 1 + height;  
}
```

Time $\Rightarrow \Theta(n)$ linear

Space $\Rightarrow \Theta(h)$ dfs



~~h(429)~~
Level order

{ { 1 } }

{ 2, 3, 4 }

{ 5, 6, 7, 8, 9 }

{ 10, 11, 12 }

○ → push

● → pop

```
public List<List<Integer>> levelOrder(Node root) {  
    List<List<Integer>> levels = new ArrayList<>();  
    if(root == null) return levels;  
    Queue<Node> q = new ArrayDeque<>();  
    q.add(root);  
  
    while(q.size() > 0){  
        List<Integer> level = new ArrayList<>();  
        for(int c = q.size(); c > 0; c--){  
            root = q.remove();  
            level.add(root.val);  
            for(Node child: root.children)  
                q.add(child);  
        }  
        levels.add(level);  
    }  
    return levels;  
}
```

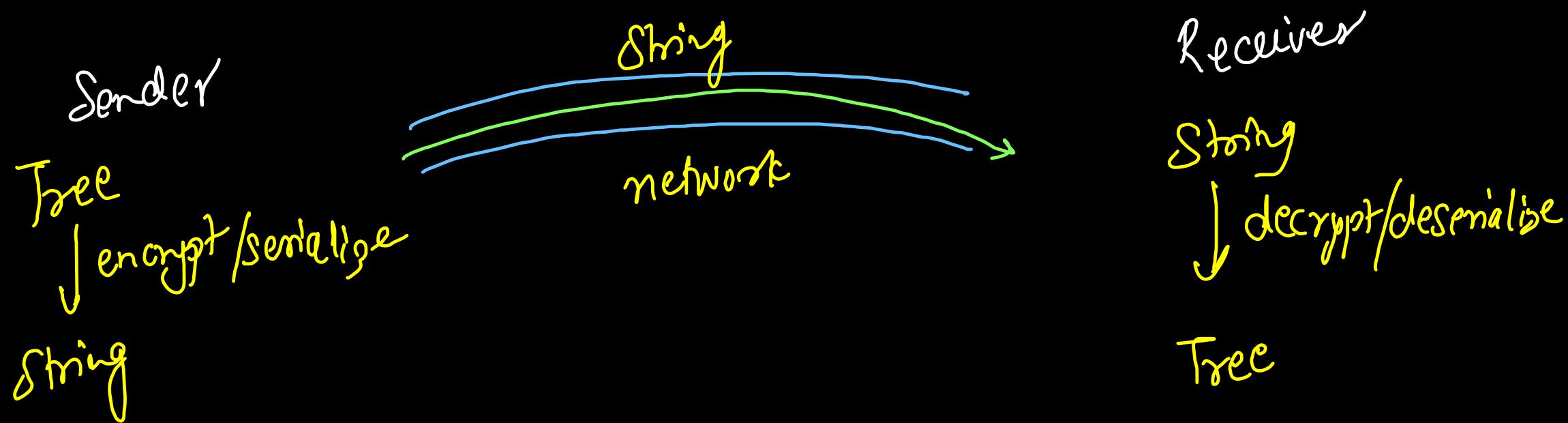
LC 429

Time = $O(n \times e) = O(2^n)$
linear ↴

Space = $O(n)$

Serialization is the process of converting a data structure or object into a sequence of bits so that it can be stored in a file or memory buffer, or transmitted across a network connection link to be reconstructed later in the same or another computer environment.

Design an algorithm to serialize and deserialize a binary tree. There is no restriction on how your serialization/deserialization algorithm should work. You just need to ensure that a binary tree can be serialized to a string and this string can be deserialized to the original tree structure.

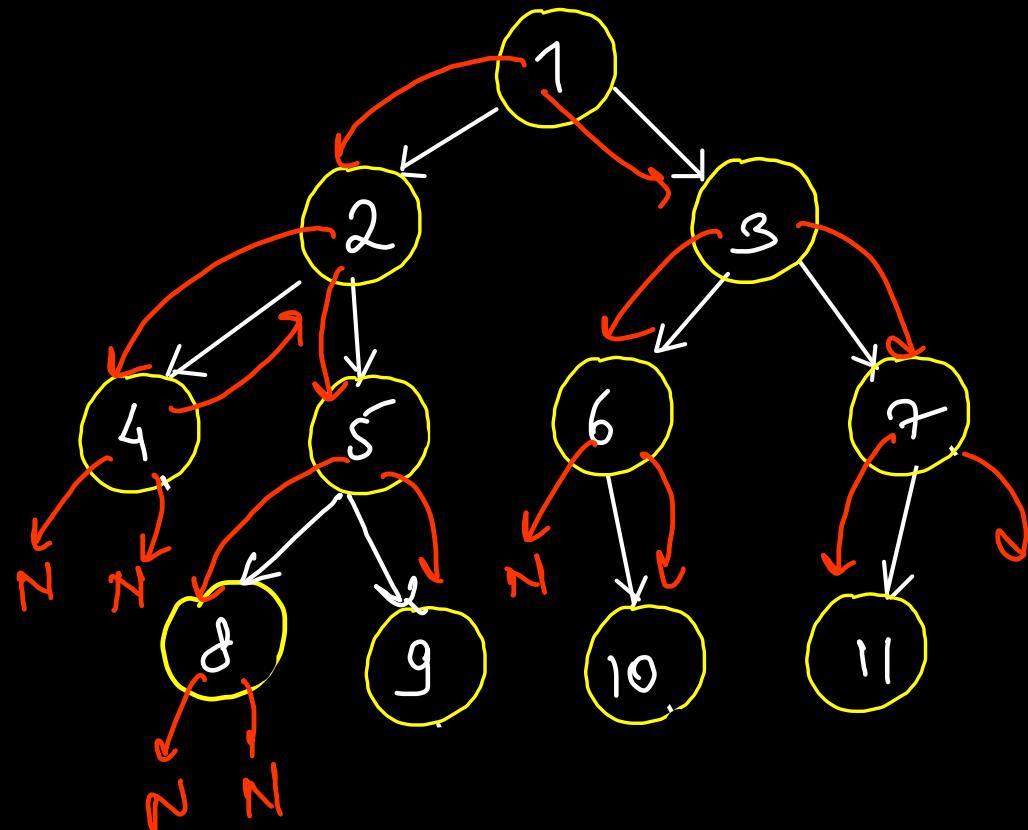


Approach 1) dfs

Preorder Serializatn

String:

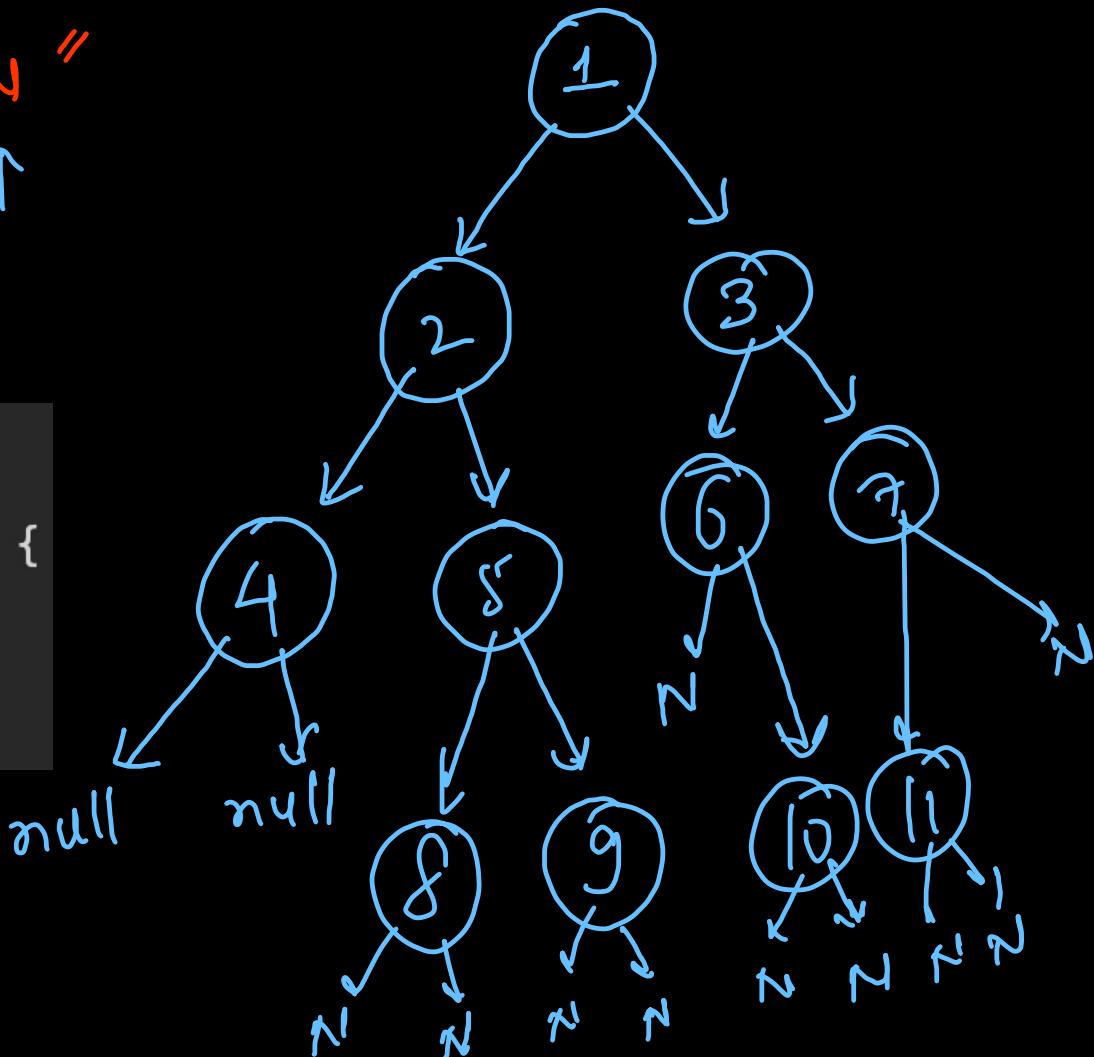
"1, 2, 4, N, N, 5, 8, N, N, 9, N, N,
3, 6, N, 10, N, N, 7, 11, N, N, N"



"↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓ ↓
 "1, 2, 4, N, N, 5, 8, N, N, 9, N, N,
 "3, 6, N, 10, N, N, 7, 11, N, N, N "

root → root.left → root.right

```
// Decodes your encoded data to tree.
public TreeNode deserialize(String data) {
  String[] nodes = data.split(",");
  return dfs(nodes);
}
```

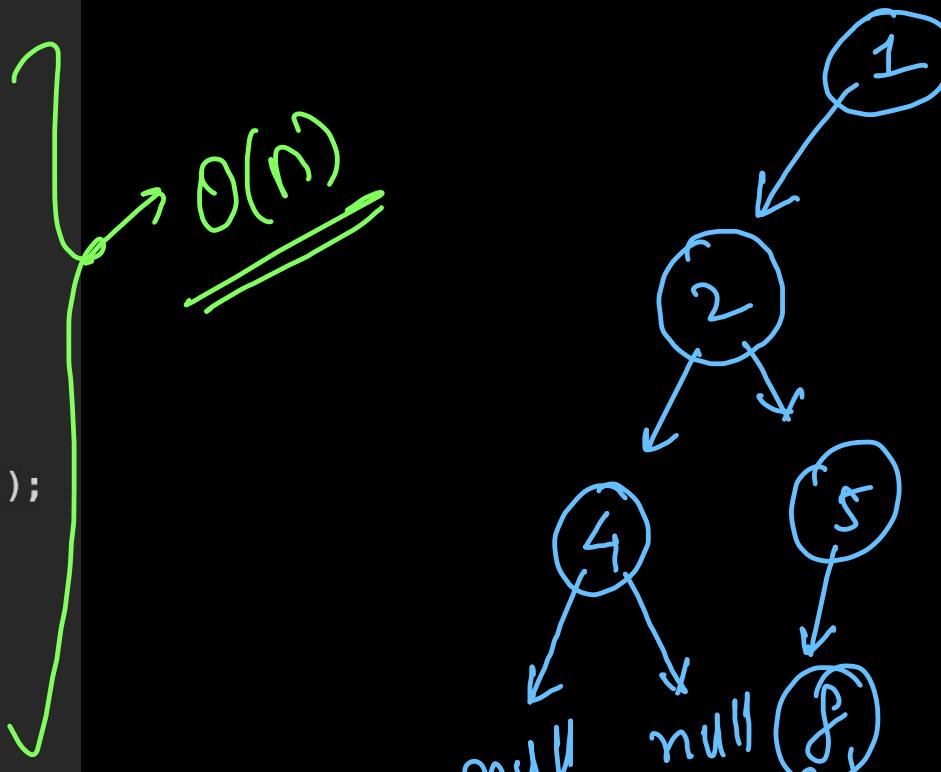


```

int idx = 0;
public TreeNode dfs(String[] nodes){
    if(idx == nodes.length) return null;
    if(nodes[idx].equals("N")) {
        idx++;
        return null;
    }

    int val = Integer.parseInt(nodes[idx++]);
    TreeNode root = new TreeNode(val);
    root.left = dfs(nodes);
    root.right = dfs(nodes);
    return root;
}

```



"~~1, 2, 4, N, N, 5, 8, N, N, 9, N, N,~~
 1, 2, 4, N, N, 5, 8, N, N, 9, N, N,
 3, 6, N, 10, N, N, 7, 11, N, N, N "

Approach 1

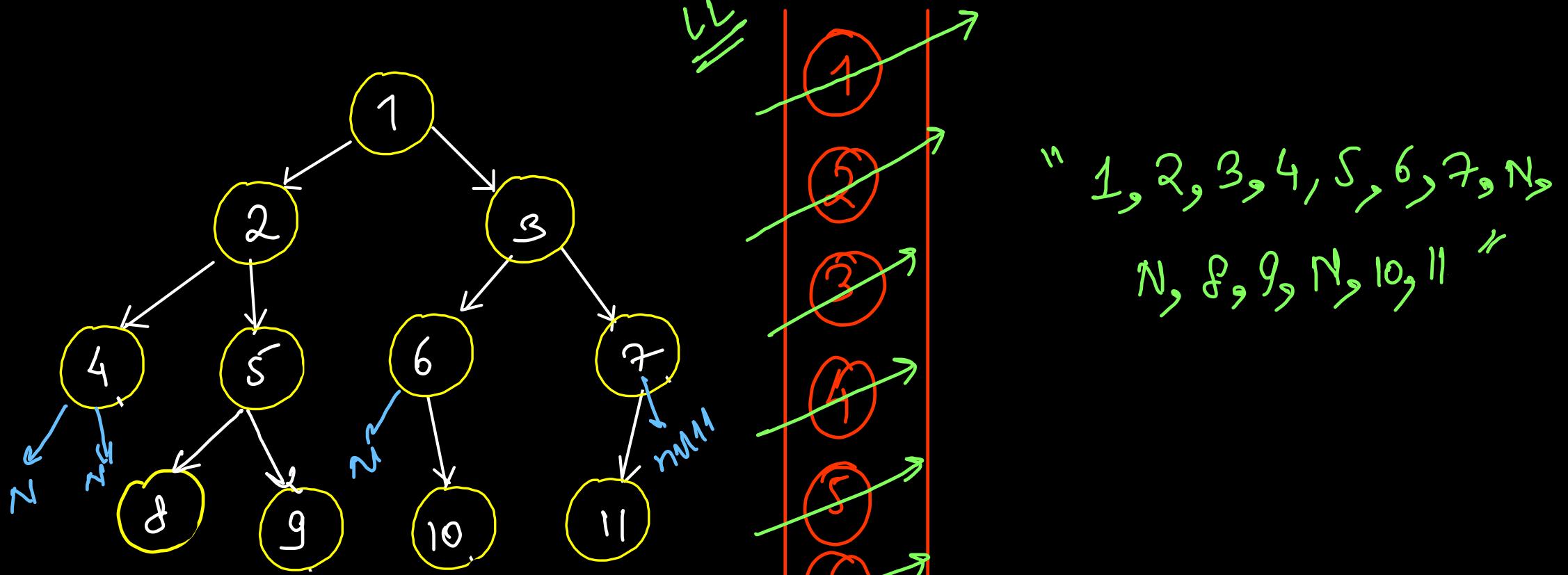
```
// Encodes a tree to a single string.  
public String serialize(TreeNode root) {  
    if(root == null) return "N,";  
    String ans = root.val + ",";  
    ans += serialize(root.left);  
    ans += serialize(root.right);  
    return ans;  
}
```

$O(n)$

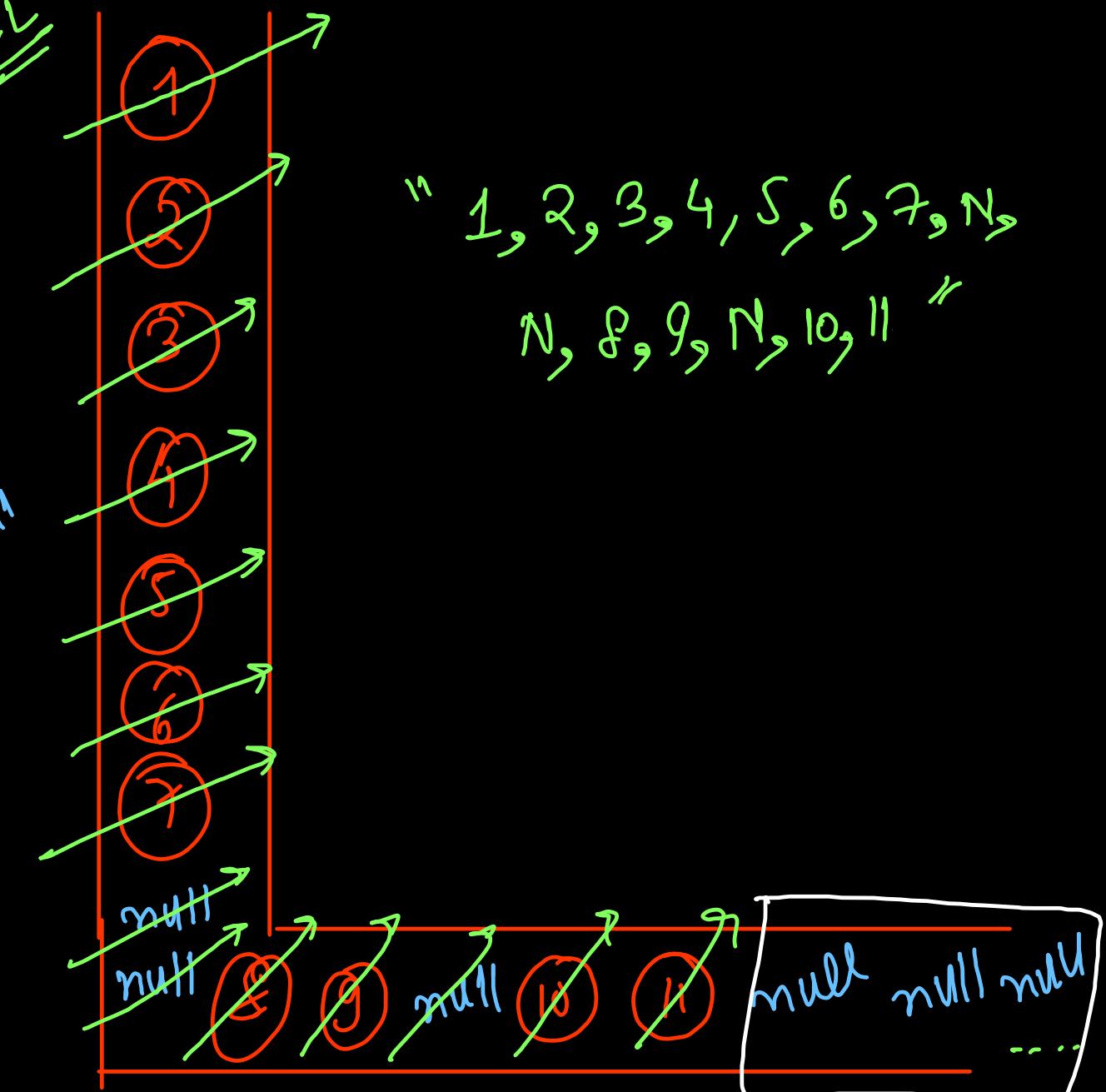
```
int idx = 0;  
public TreeNode dfs(String[] nodes){  
    if(idx >= nodes.length) return null;  
    if(nodes[idx].equals("N")) {  
        idx++;  
        return null;  
    }  
  
    int val = Integer.parseInt(nodes[idx++]);  
    TreeNode root = new TreeNode(val);  
    root.left = dfs(nodes);  
    root.right = dfs(nodes);  
    return root;  
}
```

$O(n)$

```
// Decodes your encoded data to tree.  
public TreeNode deserialize(String data) {  
    String[] nodes = data.split(",");  
    return dfs(nodes);  
}
```



" 1, 2, 3, 4, 5, 6, 7, N,
N, 8, 9, N, 10, 11 "

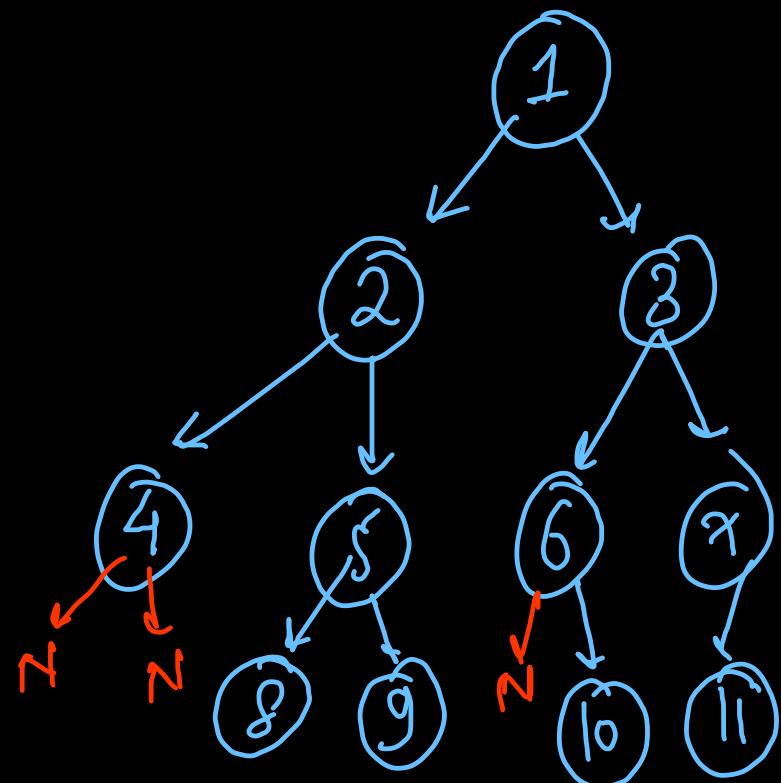
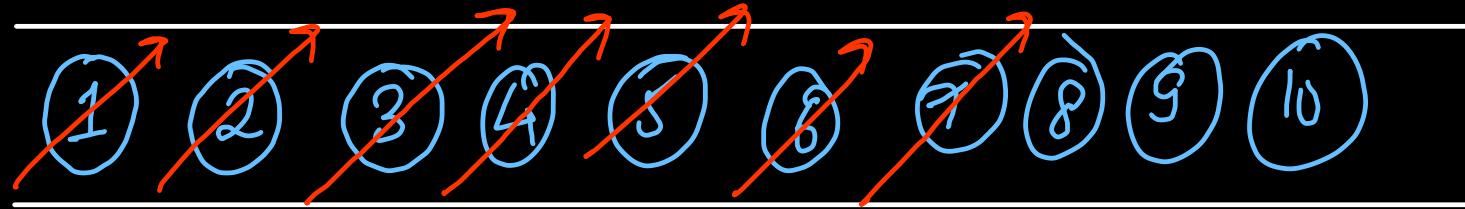


Approach 2) Level order

```
public String serialize(TreeNode root) {  
    LinkedList<TreeNode> q = new LinkedList<>();  
    q.addLast(root);  
  
    StringBuilder res = new StringBuilder();  
    while(q.size() > 0){  
        root = q.removeFirst();  
  
        if(root == null) {  
            res.append("N,");  
            continue;  
        }  
  
        res.append(root.val + ",");  
        q.addLast(root.left);  
        q.addLast(root.right);  
    }  
  
    return res.toString();  
}
```

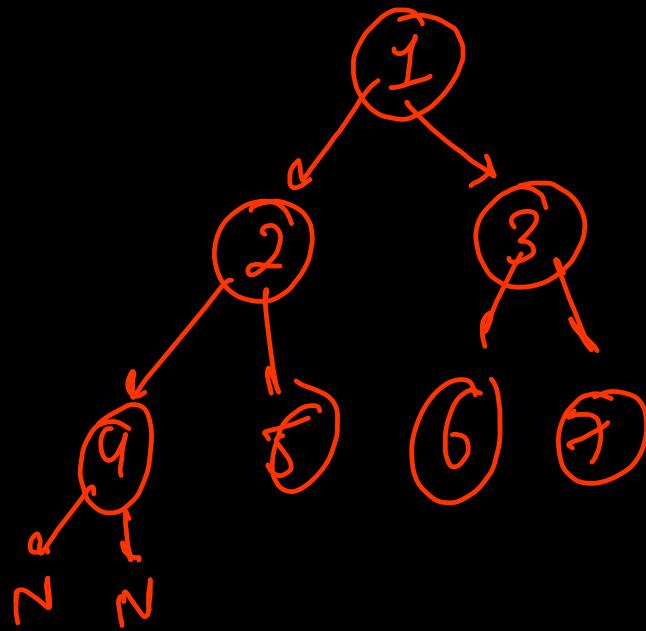
$O(n)$ time

$\underline{\underline{O(n)}}$ space



$\Theta(n)$

```
public TreeNode deserialize(String data) {  
    String[] nodes = data.split(",");  
    if(nodes[0].equals("N") == true) return null;  
  
    TreeNode root = new TreeNode(Integer.parseInt(nodes[0]));  
    LinkedList<TreeNode> q = new LinkedList<>();  
    q.addLast(root);  
    int idx = 1;  
  
    while(q.size() > 0){  
        TreeNode curr = q.removeFirst();  
  
        if(nodes[idx].equals("N") == false) {  
            curr.left = new TreeNode(Integer.parseInt(nodes[idx]));  
            q.addLast(curr.left);  
        }  
        idx++;  
  
        if(nodes[idx].equals("N") == false) {  
            curr.right = new TreeNode(Integer.parseInt(nodes[idx]));  
            q.addLast(curr.right);  
        }  
        idx++;  
    }  
  
    return root;  
}
```

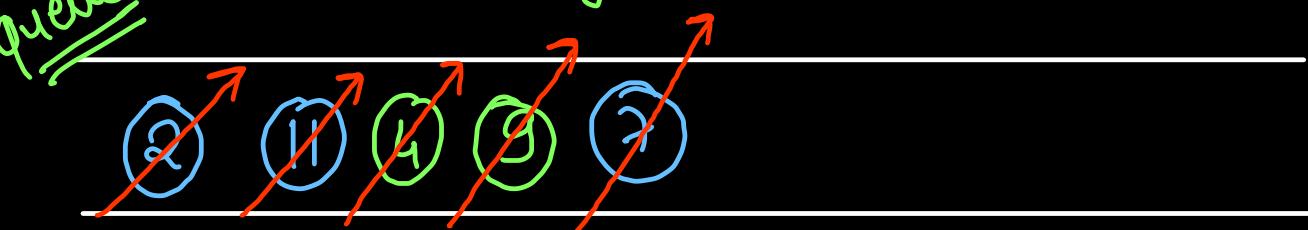


Diagonal Traversal of Binary Tree

Same diagonal → right child

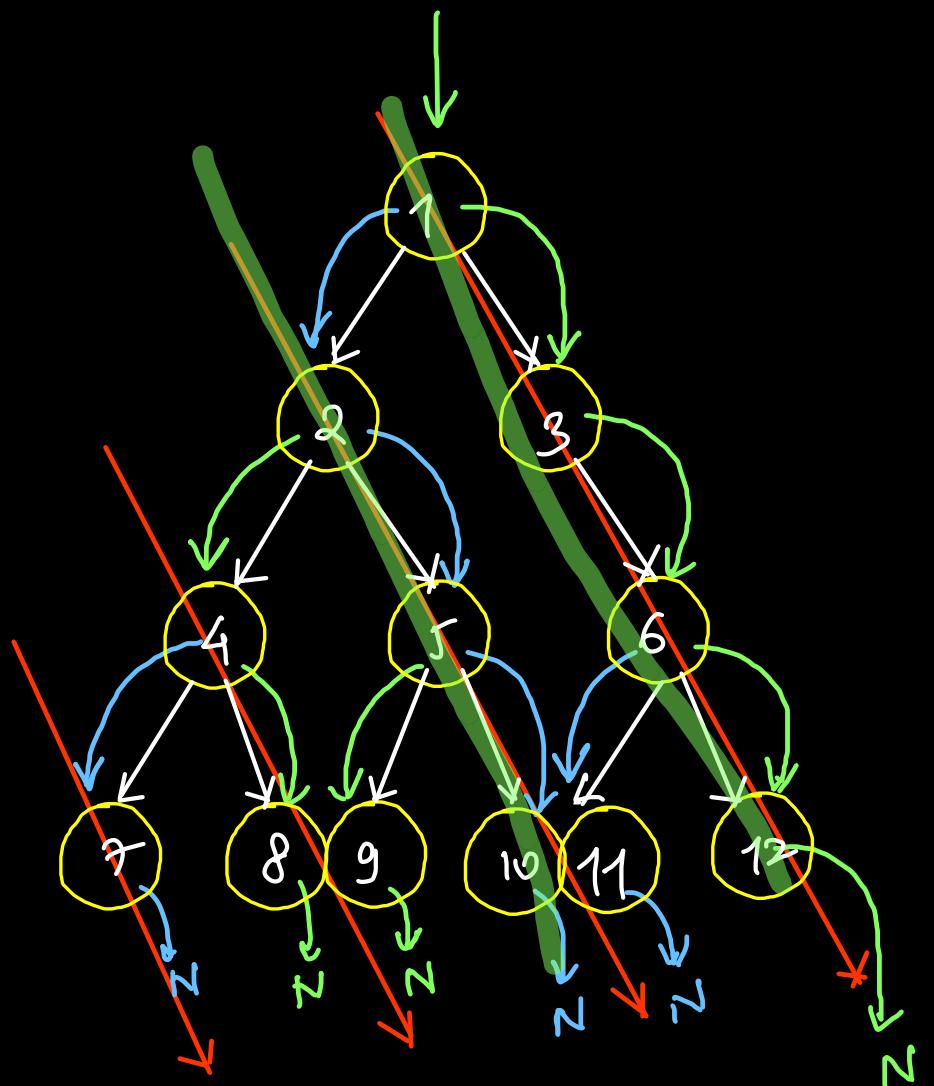
next diagonal → left child

queue



rec

{ 1, 3, 6, 12, 2, 5, 10, 11, 4, 8, 9, 7 }



```
public ArrayList<Integer> diagonal(Node root)
{
    ArrayList<Integer> res = new ArrayList<>();
    Queue<Node> q = new ArrayDeque<>();

    while(root != null || q.size() > 0){
        if(root == null){
            root = q.remove();
        } else {
            res.add(root.data);
            if(root.left != null) q.add(root.left);
            root = root.right;
        }
    }

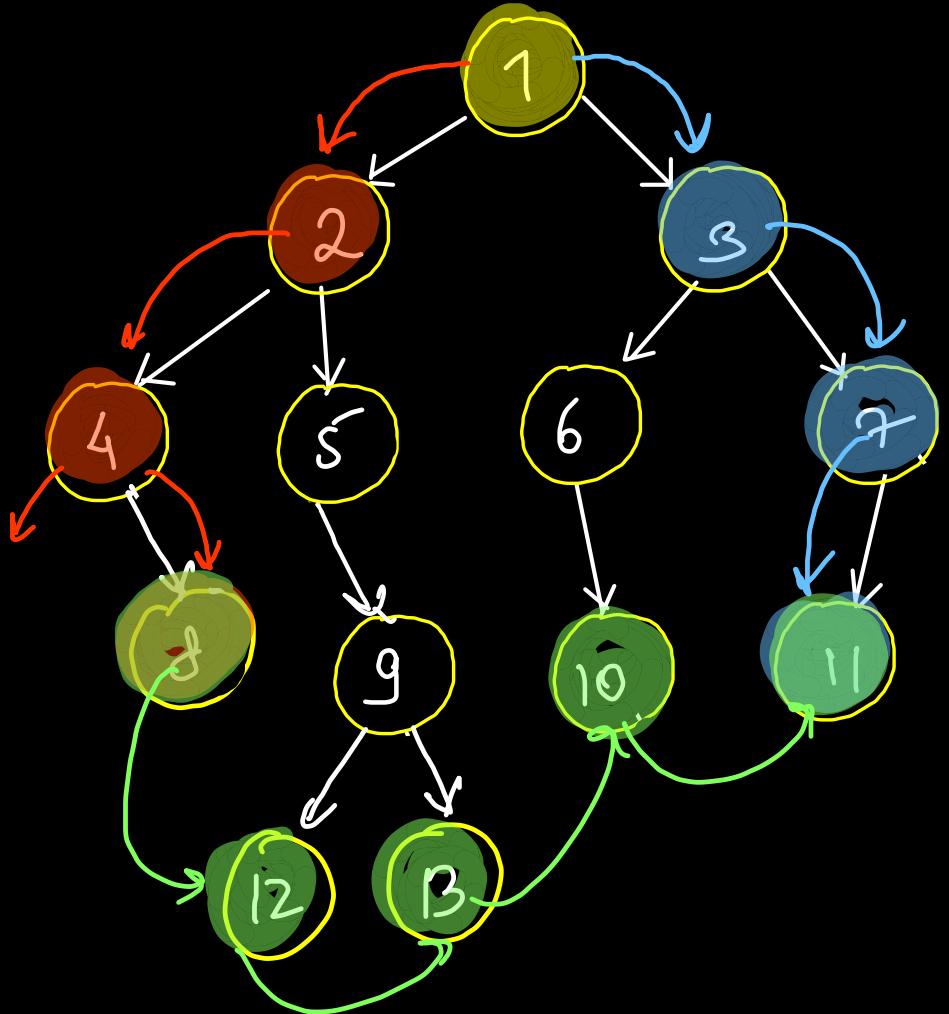
    return res;
}
```



```
while(!root == null && q.size() == 0)
```

Space $\Rightarrow O(\text{diagonal})$
(quadratic) or $O(n)$

Time \Rightarrow $O(n)$



preorder Left Boundary (Top to Bottom)

1, 2, 4, 8
≠ Left view

Leaf nodes (left to right)

8, 12, 1, 10, 11

postorder Right Boundary (Bottom to Top)

11, 7, 3, 1
≠ Right view

```

ArrayList<Integer> res;
void leftBoundary(Node root){ → O(h)
    if(root == null) return;
    if(root.left == null && root.right == null)
        return; // leaf node is ignored

    res.add(root.data); // preorder
    if(root.left != null) leftBoundary(root.left);
    else leftBoundary(root.right);
}

void rightBoundary(Node root){ → O(h)
    if(root == null) return;
    if(root.left == null && root.right == null)
        return; // leaf node is ignored

    if(root.right != null) rightBoundary(root.right);
    else rightBoundary(root.left);
    res.add(root.data); // postorder
}

```

```

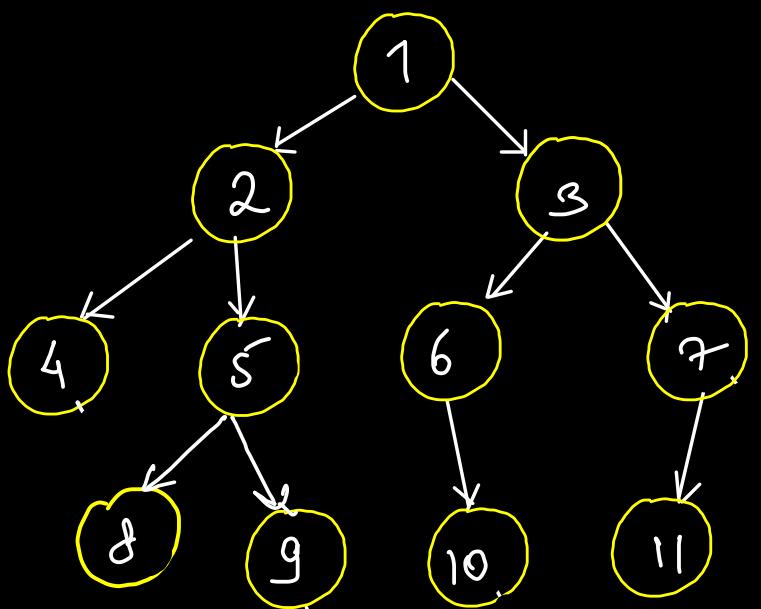
void dfs(Node root){ → O(n)
    if(root == null) return;
    if(root.left == null && root.right == null)
        res.add(root.data);
    dfs(root.left);
    dfs(root.right);
}

ArrayList <Integer> boundary(Node root)
{
    res = new ArrayList<>();
    res.add(root.data);
    leftBoundary(root.left); → O(h)
    dfs(root.left); O(n)
    dfs(root.right); → O(h)
    rightBoundary(root.right); → O(h)
    return res;
}

```

$$\text{time} = O(H + N + H)$$

Space = $O(H)$ recursion.



root → rootleft → rootright
 rootidx
 preorder: { 1, 2, 4, 5, 8, 9, 3, 6, 10, 7, 11 }

inorder: { 4, 2, 8, 5, 9, 1, 6, 10, 3, 11, 7 }
 low ↑ mid ↑ high

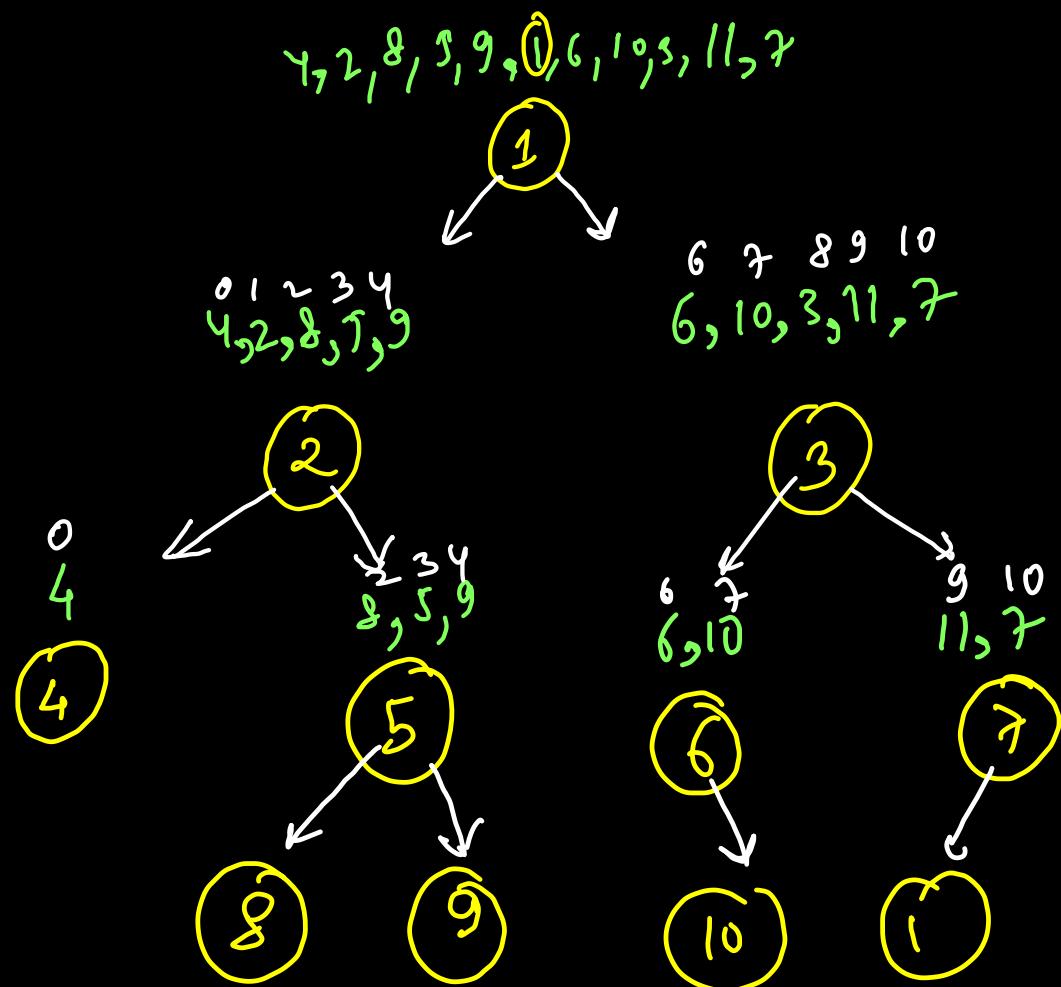
left → root → right

$\gamma_{rootidx, low, high}$

- $\xrightarrow{\text{left subtree}}$ $\gamma_{rootidx+1, low, mid-1}$
- $\xrightarrow{\text{right subtree}}$ $\gamma_{rootidx+1, mid+1, high}$
 $f(mid-low)$

preorder: $\{ \overset{0}{1}, \overset{1}{2}, \overset{2}{4}, \overset{3}{5}, \overset{4}{8}, \overset{5}{9}, \overset{6}{\text{6}}, \overset{7}{3}, \overset{8}{6}, \overset{9}{10}, \overset{10}{7}, \overset{11}{11} \}$

inorder: $\{ \overset{4}{4}, \overset{2}{2}, \overset{8}{8}, \overset{5}{5}, \overset{9}{9}, \overset{1}{1}, \overset{6}{6}, \overset{10}{10}, \overset{3}{3}, \overset{11}{11}, \overset{7}{7} \}$



```

class Solution {
    int[] preorder, inorder;

    public int search(int val){
        for(int i = 0; i < inorder.length; i++){
            if(inorder[i] == val) return i;
        }
        return -1;
    }

    public TreeNode construct(int rootidx, int low, int high){
        if(low > high) return null;
        if(low == high) return new TreeNode(preorder[rootidx]);

        TreeNode root = new TreeNode(preorder[rootidx]);
        int mid = search(preorder[rootidx]);
        root.left = construct(rootidx + 1, low, mid - 1);
        root.right = construct(rootidx + 1 + (mid - low), mid + 1, high);
        return root;
    }

    public TreeNode buildTree(int[] preorder, int[] inorder) {
        this.preorder = preorder;
        this.inorder = inorder;
        return construct(0, 0, inorder.length - 1);
    }
}

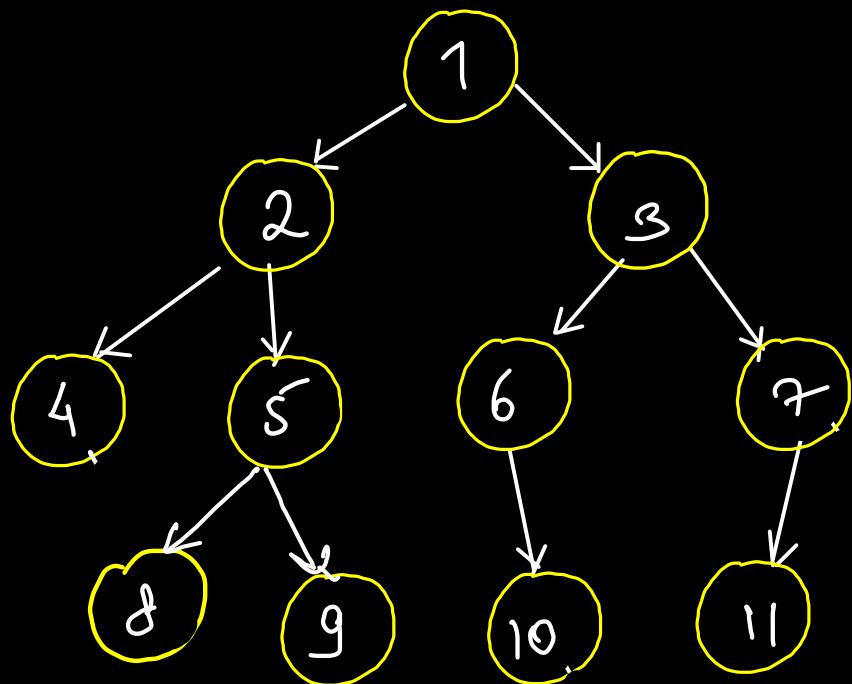
```

→ assuming search is constant

if balanced binary tree
 $O(n \log n)$ and
else : $O(n^2)$
worst case

Similar to Divide & Conquer !

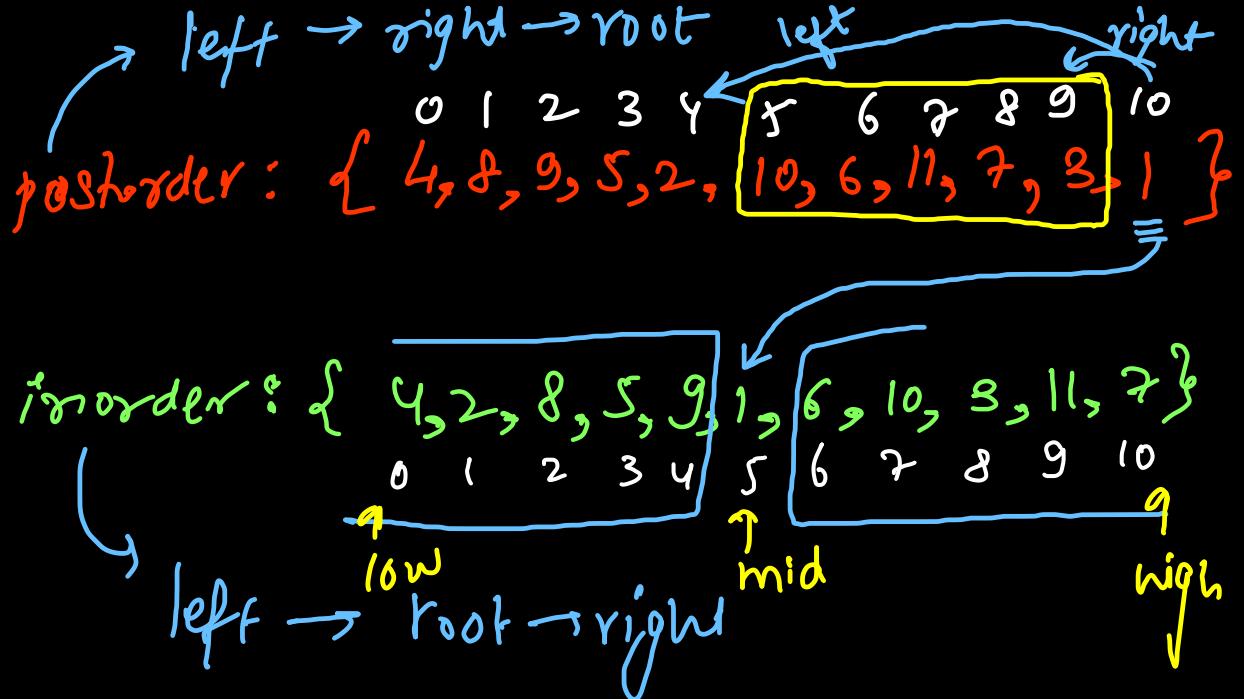
LC 106)



postidx, low, high

left \rightarrow postidx-1 - (high-mid), low, mid-1

right \rightarrow postidx-1 , mid+1, high



```
class Solution {
    HashMap<Integer, Integer> in = new HashMap<>();

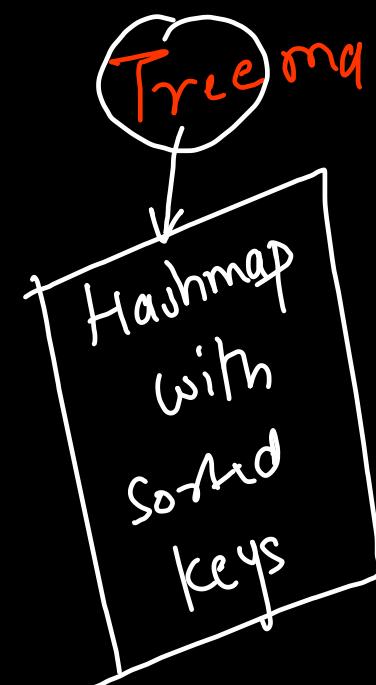
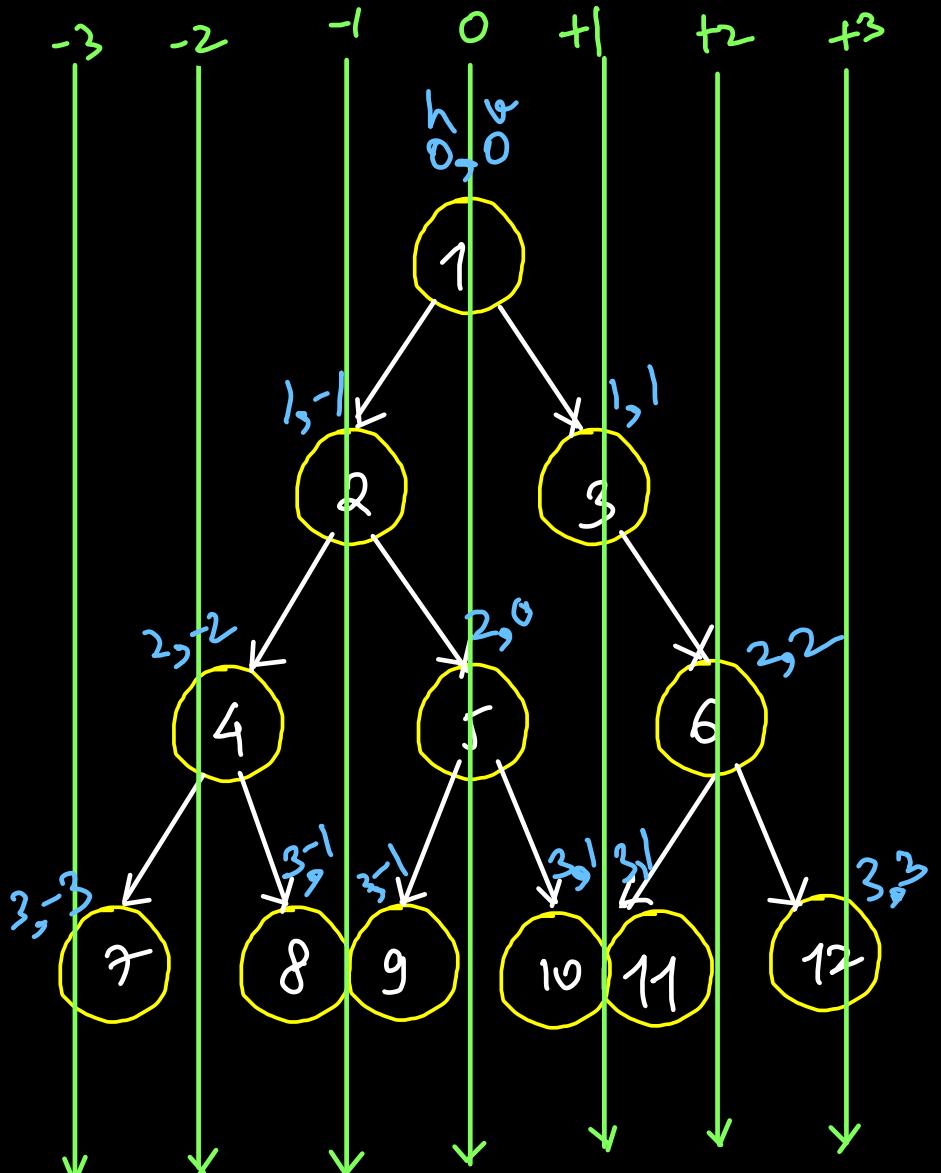
    public TreeNode construct(int[] postorder, int i, int l, int r){
        if(l > r) return null;
        TreeNode root = new TreeNode(postorder[i]);

        int j = in.get(postorder[i]);
        root.left = construct(postorder, i - (r - j) - 1, l, j - 1);
        root.right = construct(postorder, i - 1, j + 1, r);

        return root;
    }

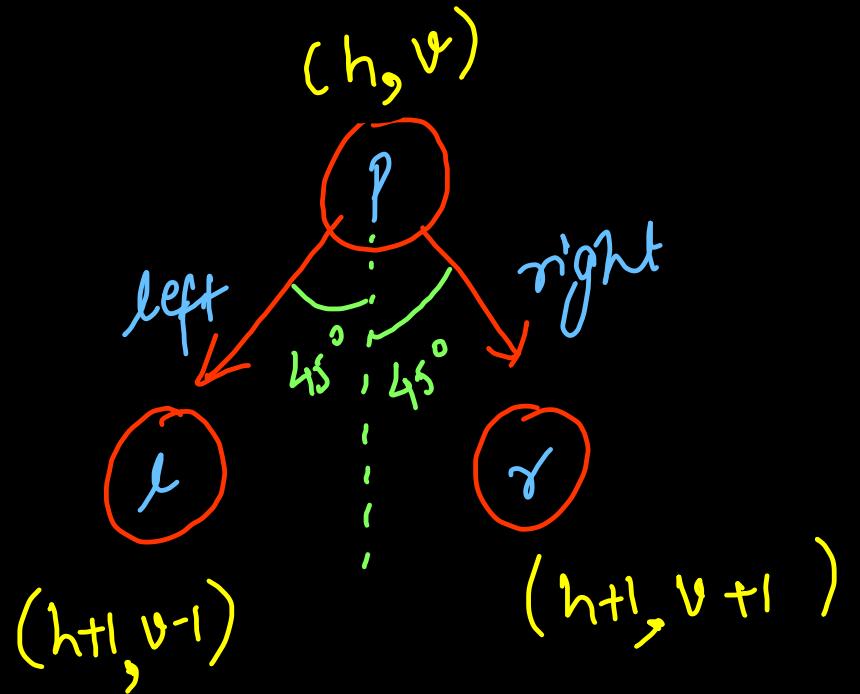
    public TreeNode buildTree(int[] inorder, int[] postorder) {
        for(int i = 0; i < inorder.length; i++)
            in.put(inorder[i], i);
        return construct(postorder, postorder.length - 1, 0, inorder.length - 1);
    }
}
```

Vertical Traversal of Binary Tree

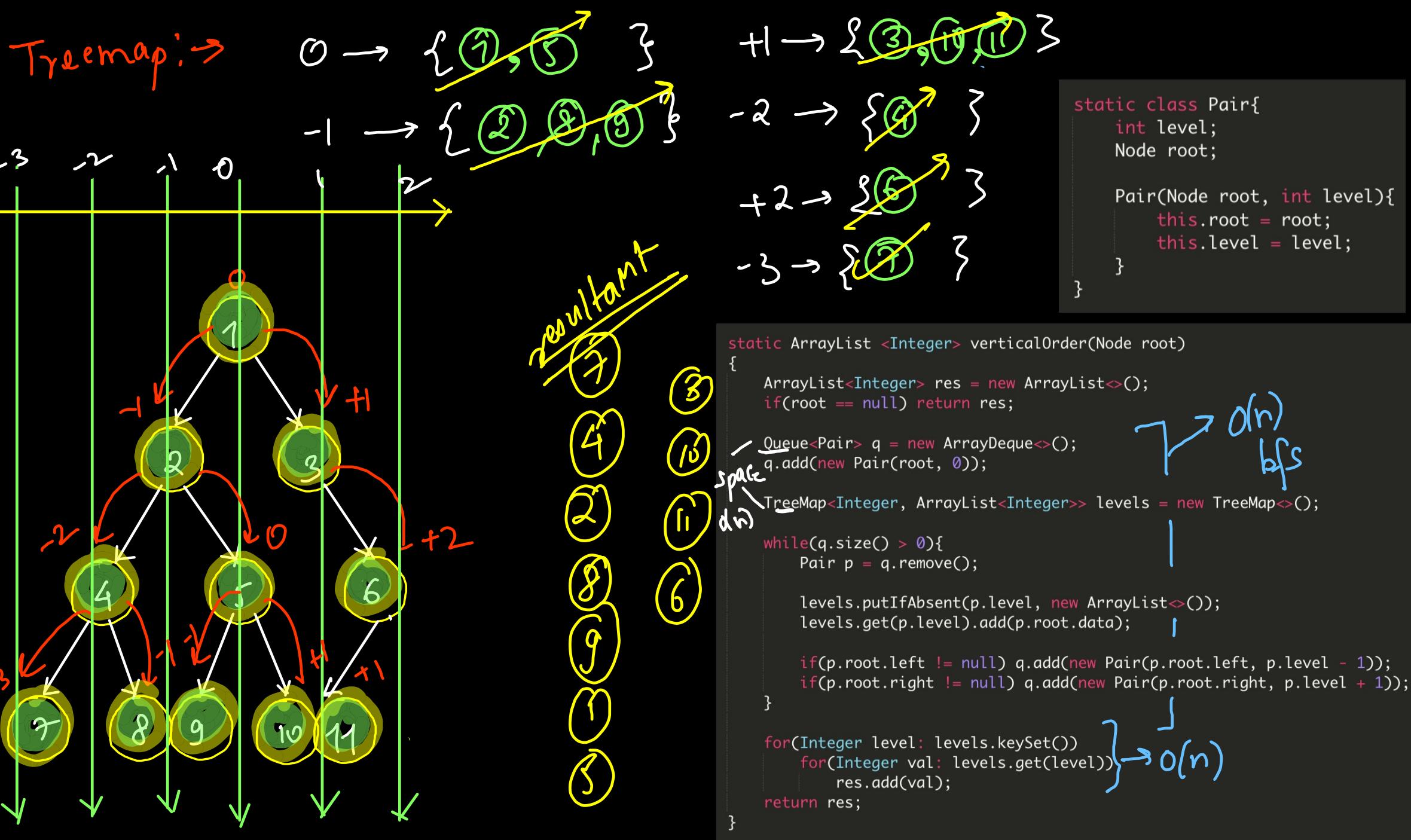


ArrayList<Integer>

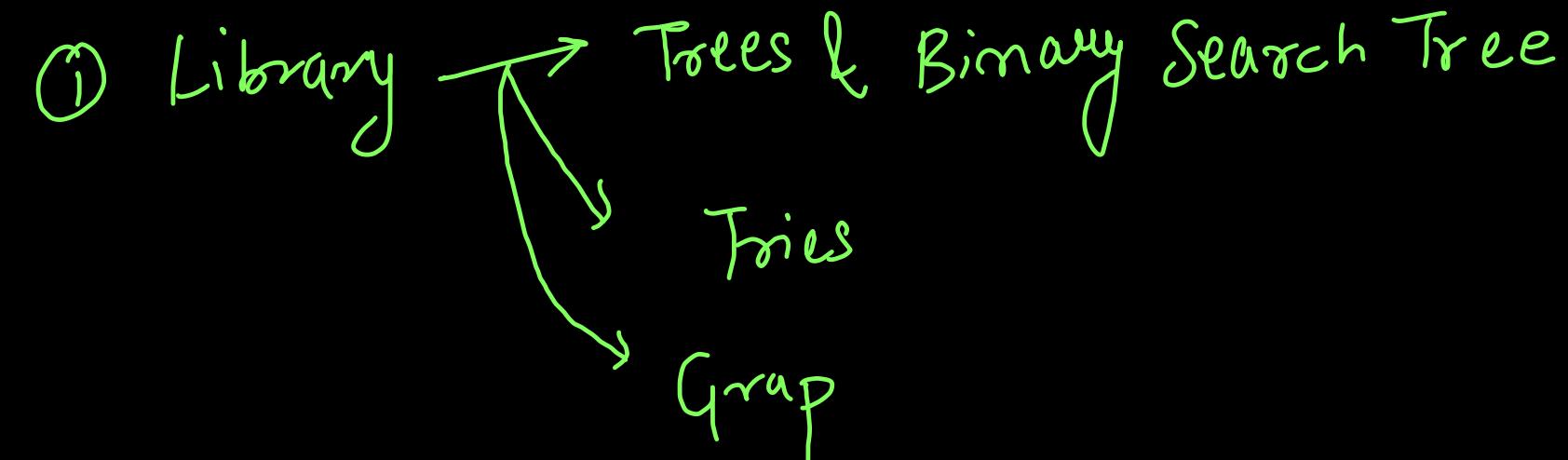
0 →	{①, ⑤}
-1 →	{②, ⑧, ⑨}
-2 →	{④}
-3 →	{⑦}
1 →	{③, ⑩, ⑪}
2 →	{⑥}
3 →	{⑫}



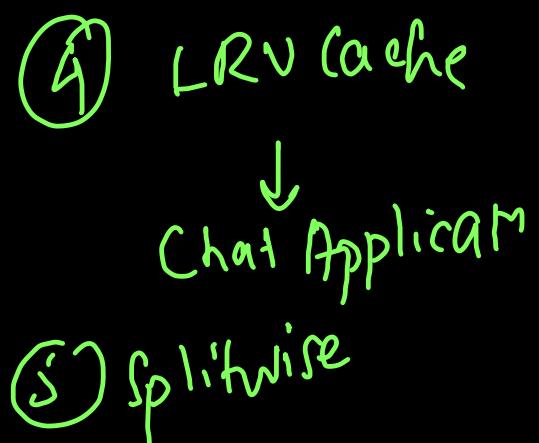
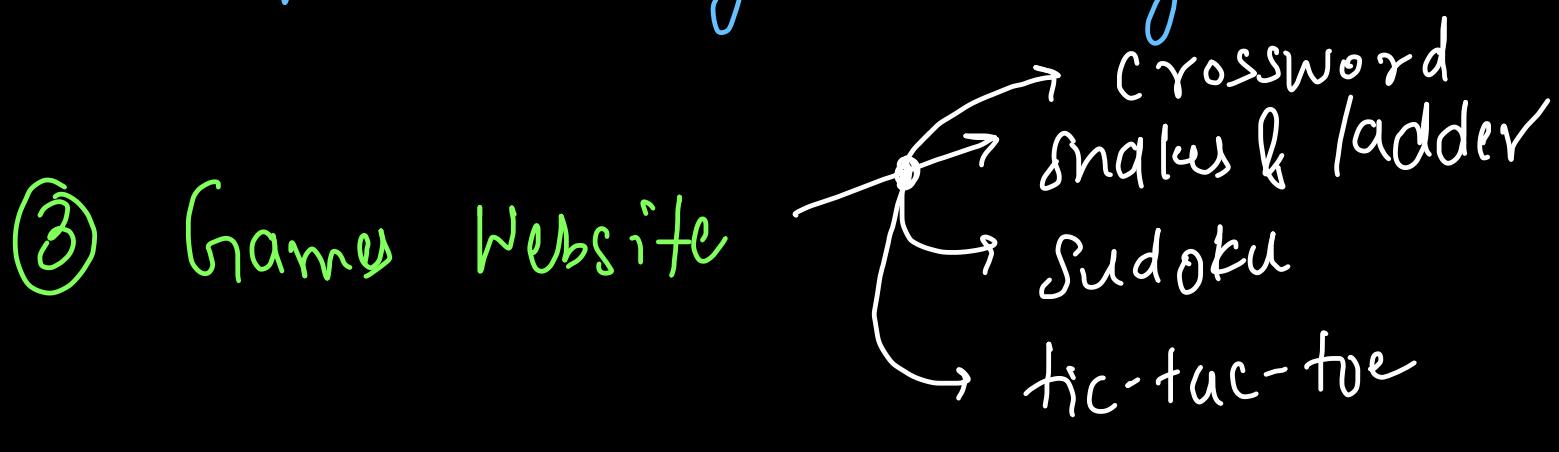
$p \circ \text{horizontal} = h$
 $p \circ \text{level} = v$

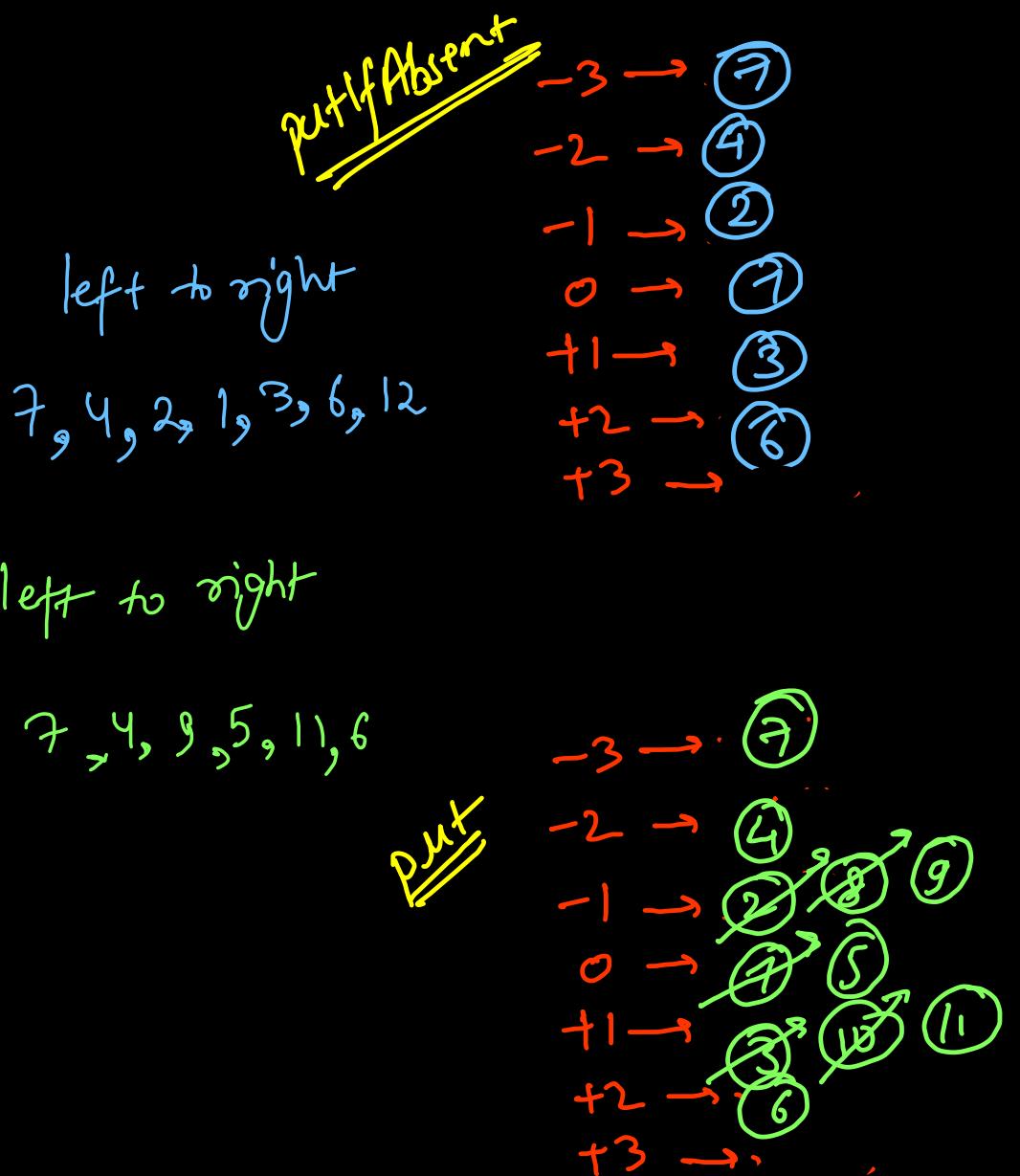
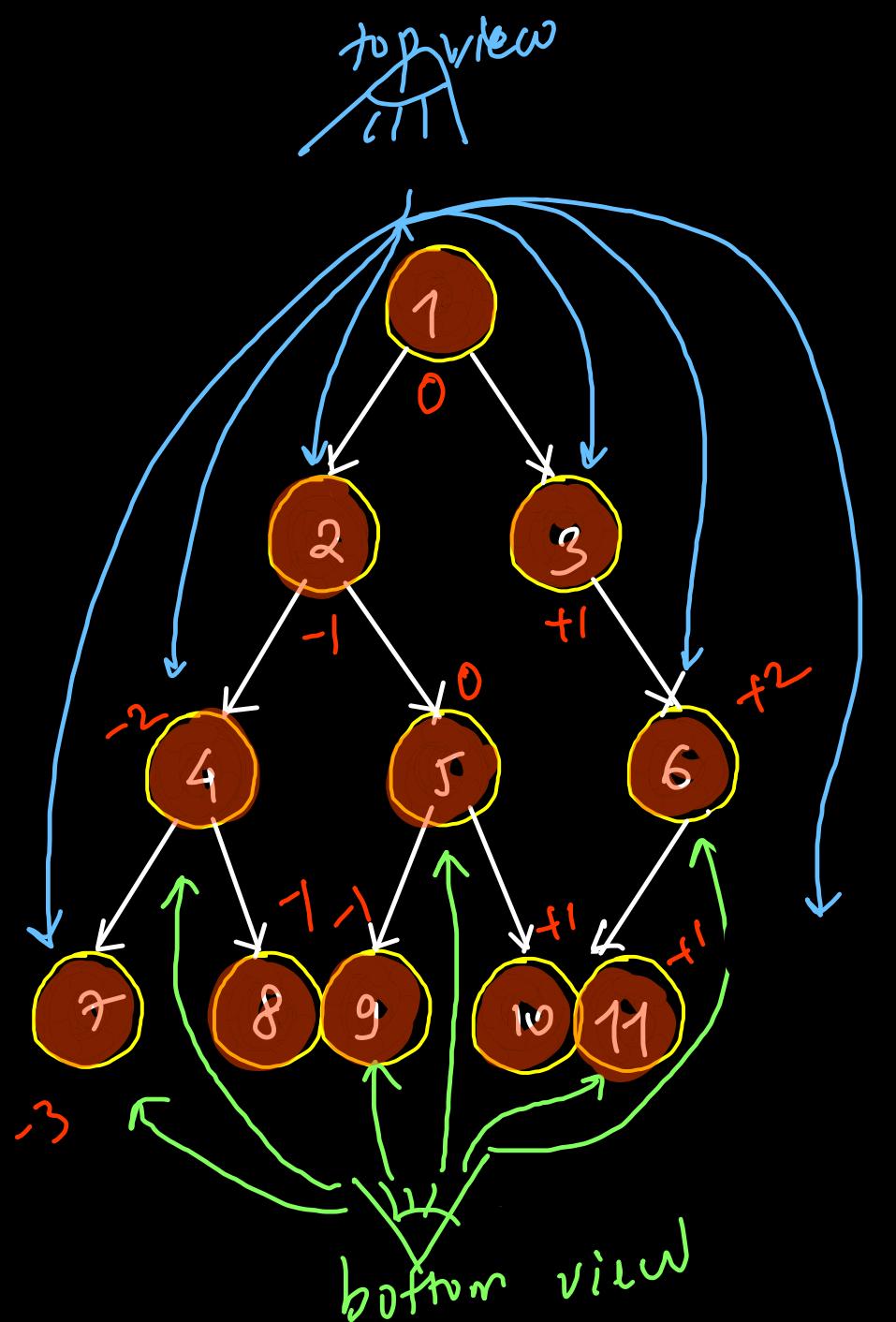


DSA based Projects



② Huffman Encoding & Decoding → zipper/compress





```

static class Pair{
    int level;
    Node root;

    Pair(Node root, int level){
        this.root = root;
        this.level = level;
    }
}

static ArrayList<Integer> topView(Node root)
{
    ArrayList<Integer> res = new ArrayList<>();
    if(root == null) return res;

    Queue<Pair> q = new ArrayDeque<>();
    q.add(new Pair(root, 0));

    TreeMap<Integer, Integer> levels = new TreeMap<>();

    while(q.size() > 0){
        Pair p = q.remove();

        levels.putIfAbsent(p.level, p.root.data);
        if(p.root.left != null) q.add(new Pair(p.root.left, p.level - 1));
        if(p.root.right != null) q.add(new Pair(p.root.right, p.level + 1));
    }

    for(Integer key: levels.keySet())
        res.add(levels.get(key));
    return res;
}

```

Top view

```

static class Pair{
    int level;
    Node root;

    Pair(Node root, int level){
        this.root = root;
        this.level = level;
    }
}

static ArrayList<Integer> bottomView(Node root)
{
    ArrayList<Integer> res = new ArrayList<>();
    if(root == null) return res;

    Queue<Pair> q = new ArrayDeque<>();
    q.add(new Pair(root, 0));

    TreeMap<Integer, Integer> levels = new TreeMap<>();

    while(q.size() > 0){
        Pair p = q.remove();

        levels.put(p.level, p.root.data);
        if(p.root.left != null) q.add(new Pair(p.root.left, p.level - 1));
        if(p.root.right != null) q.add(new Pair(p.root.right, p.level + 1));
    }

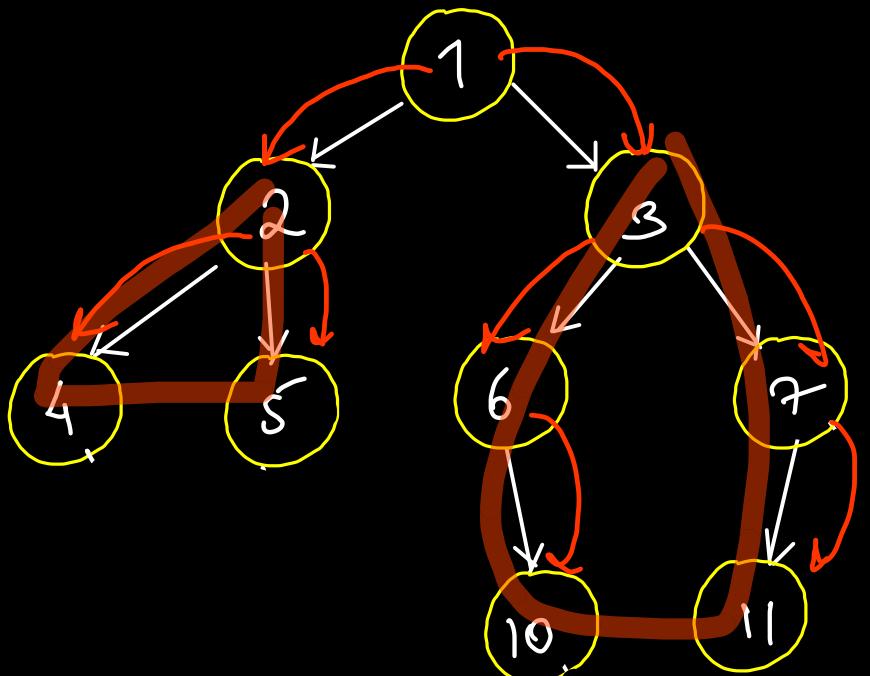
    for(Integer key: levels.keySet())
        res.add(levels.get(key));
    return res;
}

```

Bottom view

$O(n)$ time, $O(n)$ space

LC 606 Construct String from Tree



1 (2 (4) (5)) (3 (6 ()(10)) (7(11)))

leaf node

only right child

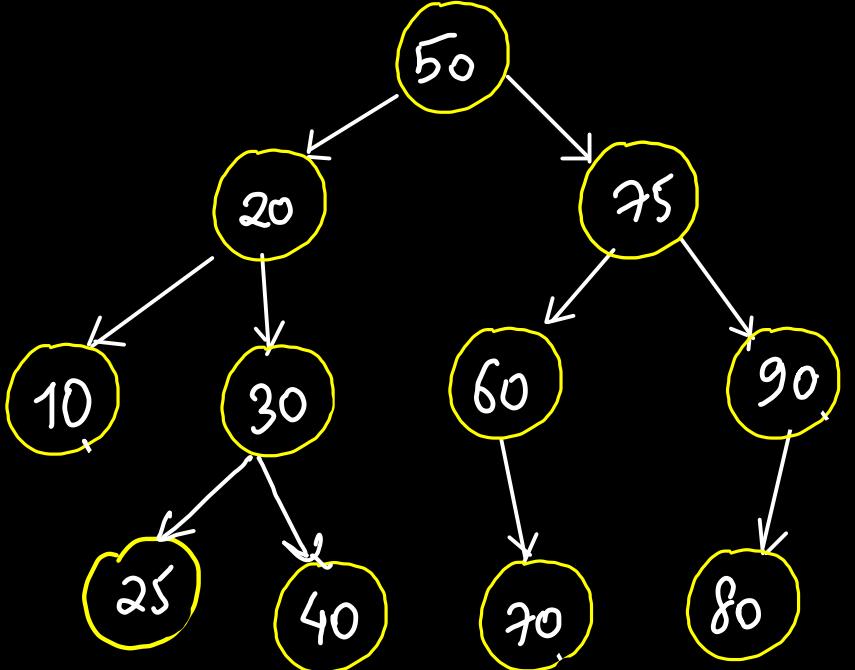
only left child

```
public String tree2str(TreeNode root) {  
    if(root == null) return "";  
    if(root.left == null && root.right == null)  
        return Integer.toString(root.val); // leaf  
    if(root.right == null){  
        String left = tree2str(root.left);  
        return root.val + "(" + left + ")";  
    }  
    if(root.left == null){  
        String right = tree2str(root.right);  
        return root.val + "()" + "(" + right + ")";  
    }  
  
    String left = tree2str(root.left);  
    String right = tree2str(root.right);  
    return root.val + "(" + left + ")" + "(" + right + ")";  
}
```

→ assuming string operations are $O(1)$

Time = $O(n)$
space = $O(h)$

Binary Search Tree



⇒ Binary Tree + Searching

⇒ BinarySearch + Tree
(Divide & Conquer)

① For every node

left subtree $<$ root $<$ right node's

② Inorder Traversal \Rightarrow increasing order

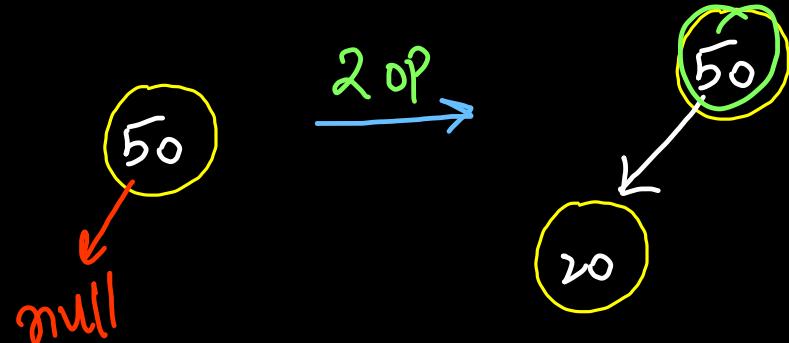
10, 20, 25, 30, 40, 50, 60, 70, 75, 80, 90

Insert into BST { Will be always at leaf }

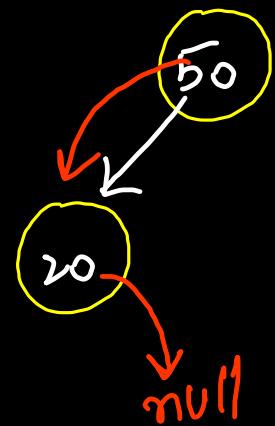
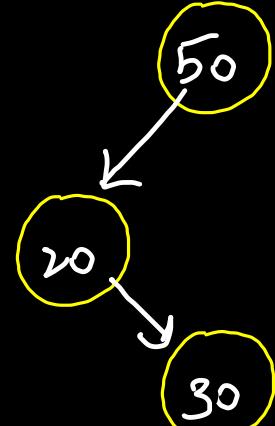
1) insert 50



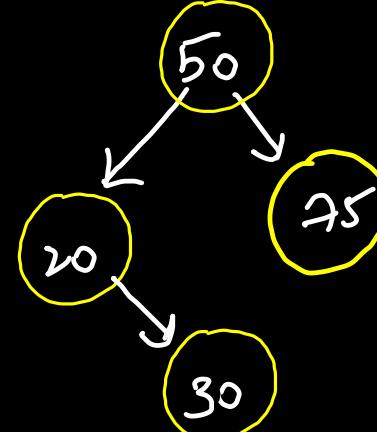
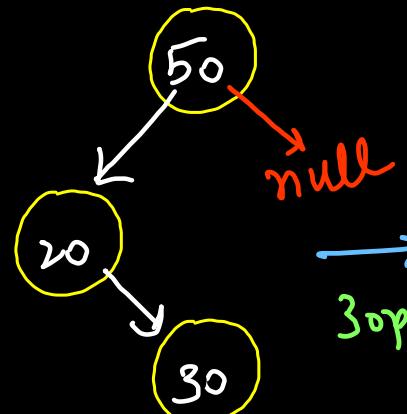
2) insert 20



3) insert 30

 $\xrightarrow{2\text{op}}$ 

4) insert 75



5) insert 70



Total insertn

$$\Rightarrow 1 + 2 \times 2 + 3 \times 4 + 4 \times 8 \Rightarrow \sum h \times 2^{h-1}$$

$n \log n$ for n nodes

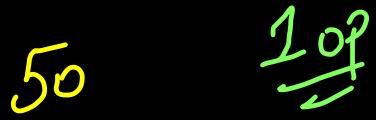
Avg case $\Rightarrow O(\log n)$
per insertn

```
public TreeNode insertIntoBST(TreeNode root, int val) {  
    if(root == null) return new TreeNode(val);  
    if(val < root.val) root.left = insertIntoBST(root.left, val);  
    else root.right = insertIntoBST(root.right, val);  
    return root;  
}
```

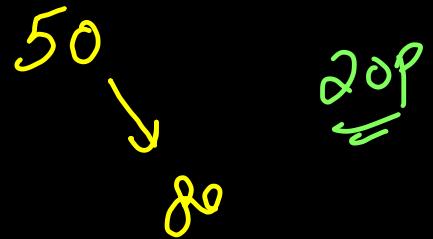
{ balanced }
bst

Worst case analysis (skewed binary tree)

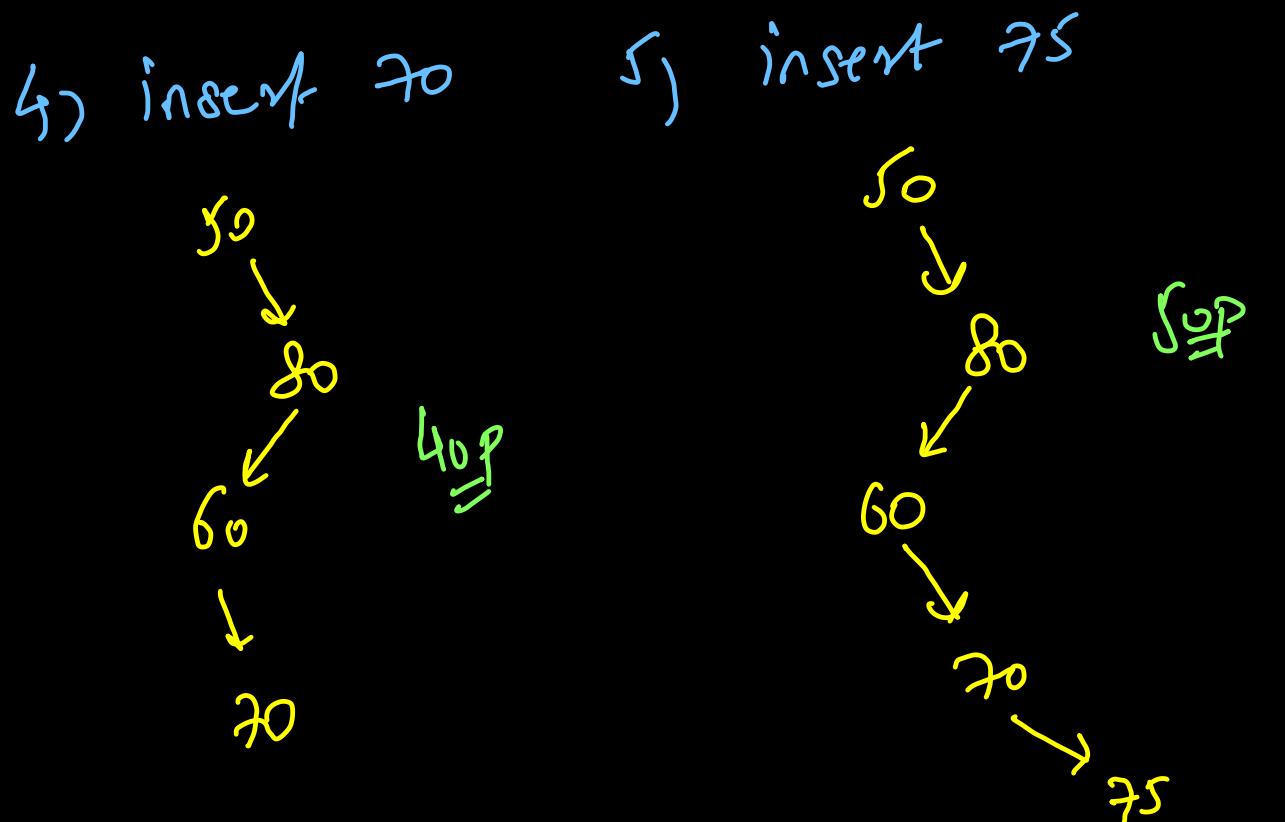
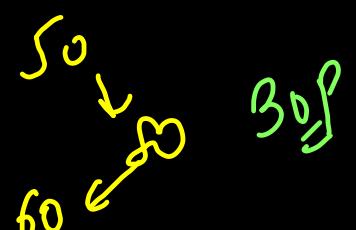
1) insert 50



2) insert 80



3) insert 60

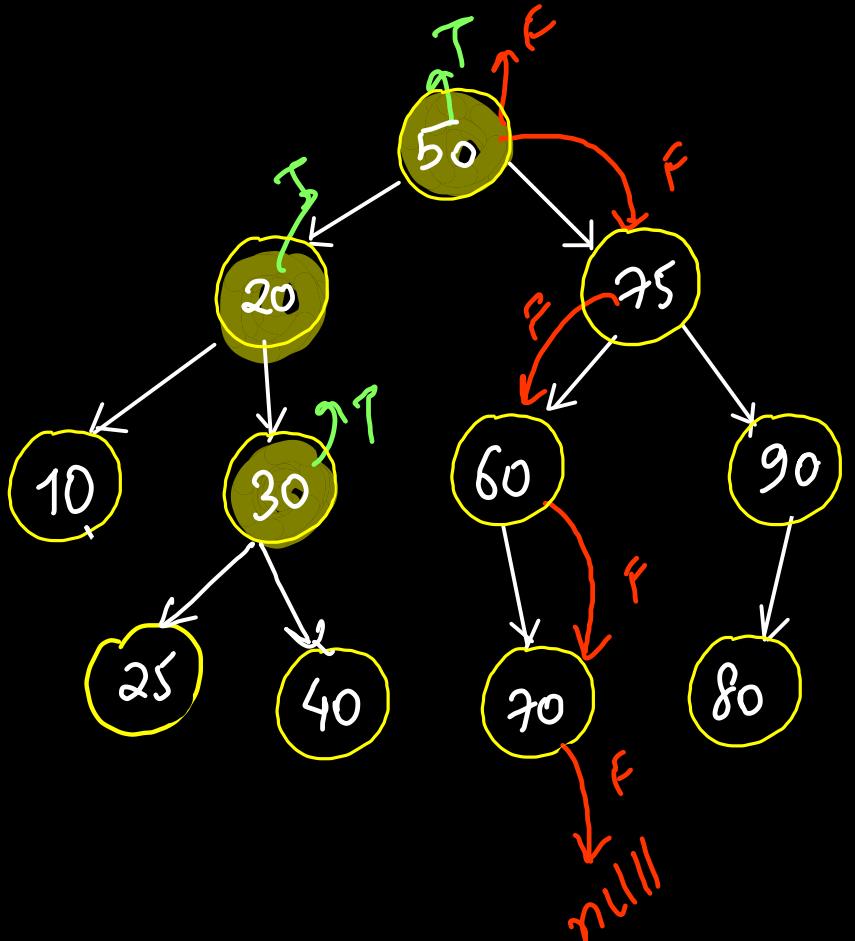


$$\begin{aligned}
 \text{Worst case total} &= 1 + 2 + 3 + 4 + \dots + n \\
 &= \underline{n(n+1)} = O(n^2)
 \end{aligned}$$

Single insert $\Rightarrow O(n^2/n) = O(n)$

Search in BST

LC 700



Search(30) \Rightarrow true

search(72) \Rightarrow false
(null)

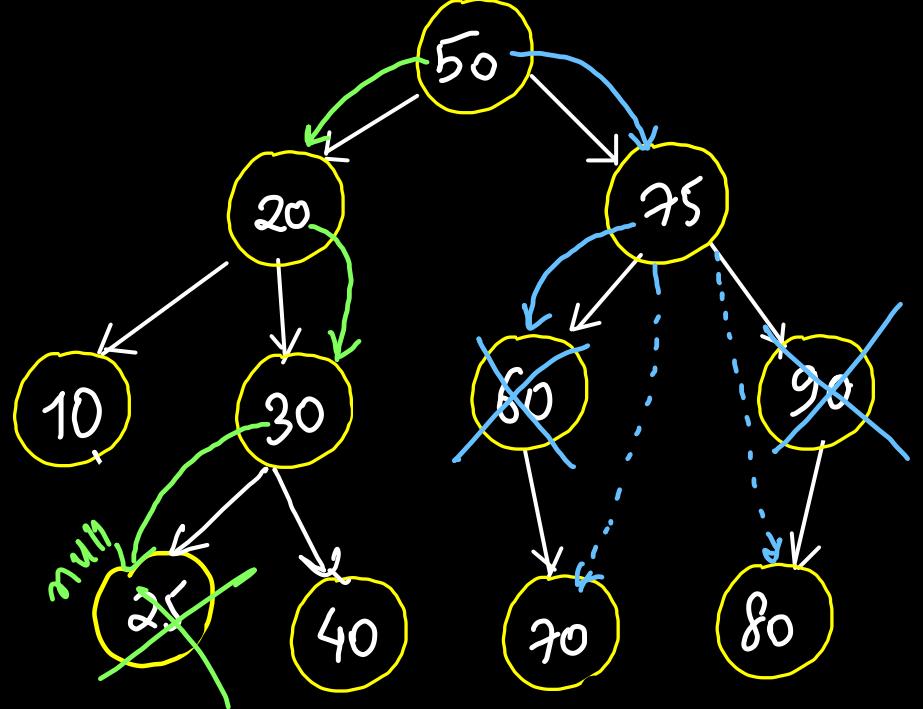
```
public TreeNode searchBST(TreeNode root, int target) {  
    if(root == null || root.val == target) return root;  
    if(target < root.val) return searchBST(root.left, target);  
    return searchBST(root.right, target);  
}
```

time $O(h)$

space

$O(\log n)$ balanced

$O(n)$ skewed

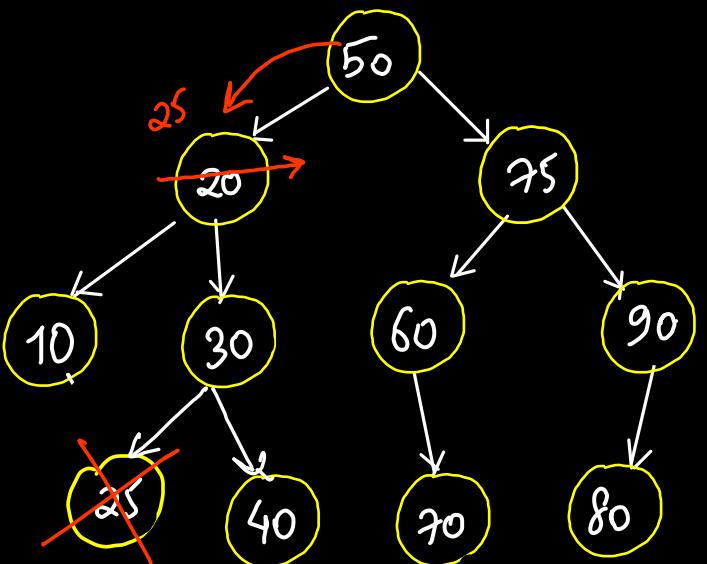


LC 450) Delet'n im BST

Case 1 ~~Node not present~~ delete 100
 \Rightarrow nothing happens
 (node not found)

leaf node ~~Case 2~~ delete 25
 \Rightarrow directly delete it
 (replace with null)

only one child present ~~case 3 present~~ delete 60 \Rightarrow delete node
 delete 90 (replace with the child)



born child
present ~~#~~

delete 20

$$= \text{replace}(x) + \text{delete}(x)$$

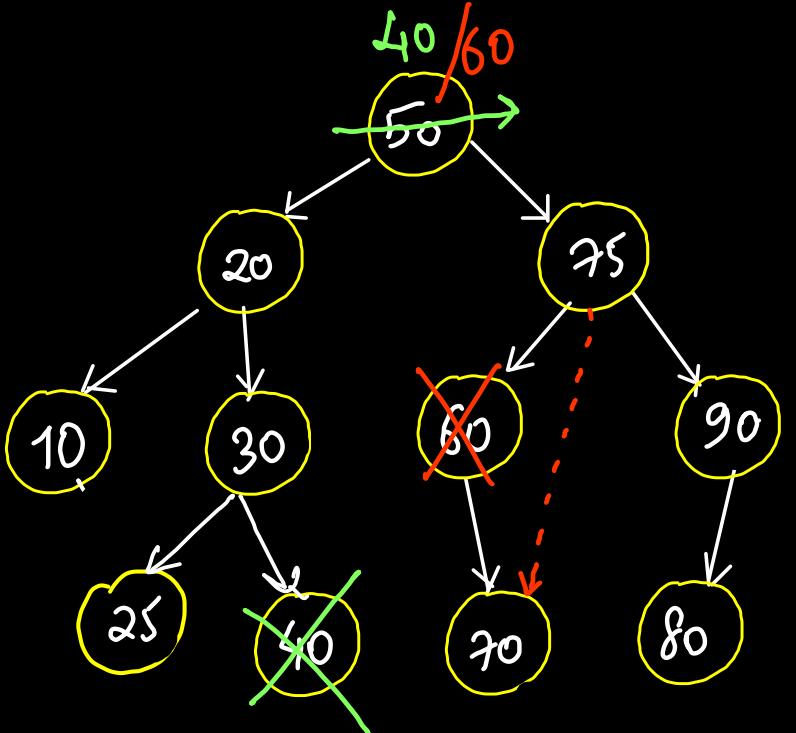
$x = \text{just smaller}(\text{floor}) = 10$

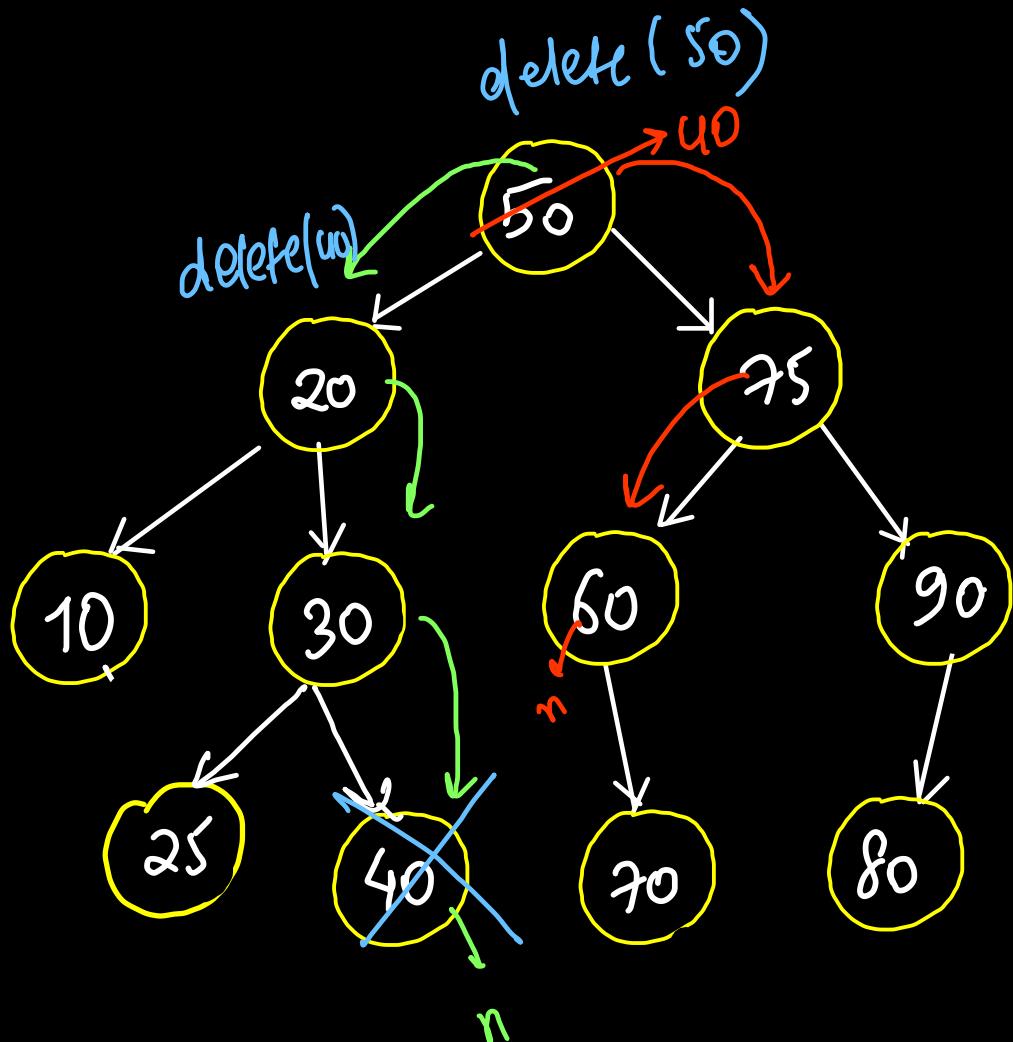
or $\text{just larger}(\text{ceil}) = 25$

delete 50

$$= \text{replace}(x) + \text{delete}(x)$$

$x \rightarrow \text{floor} = 40$
or
 $\text{ceil} = 60$





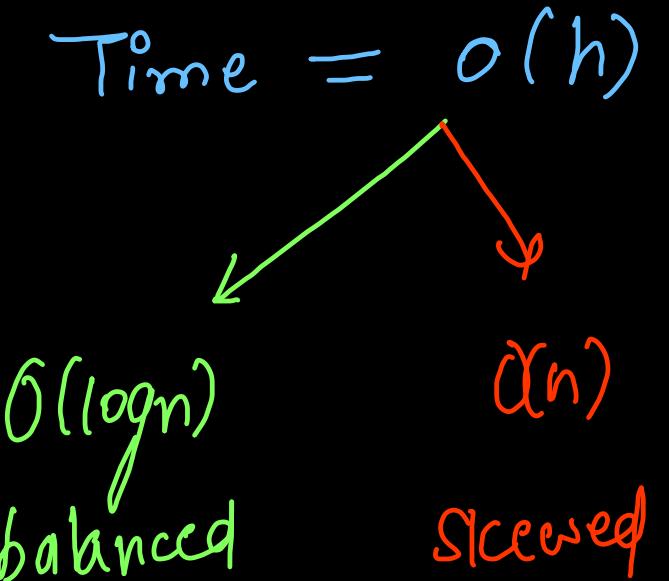
Sort_{inorder} predecessor
 ↑
 floor = max^m node
 in root.left
 or
 ceiling = min^m node
 in root.right
 Sort_{inorder} successor

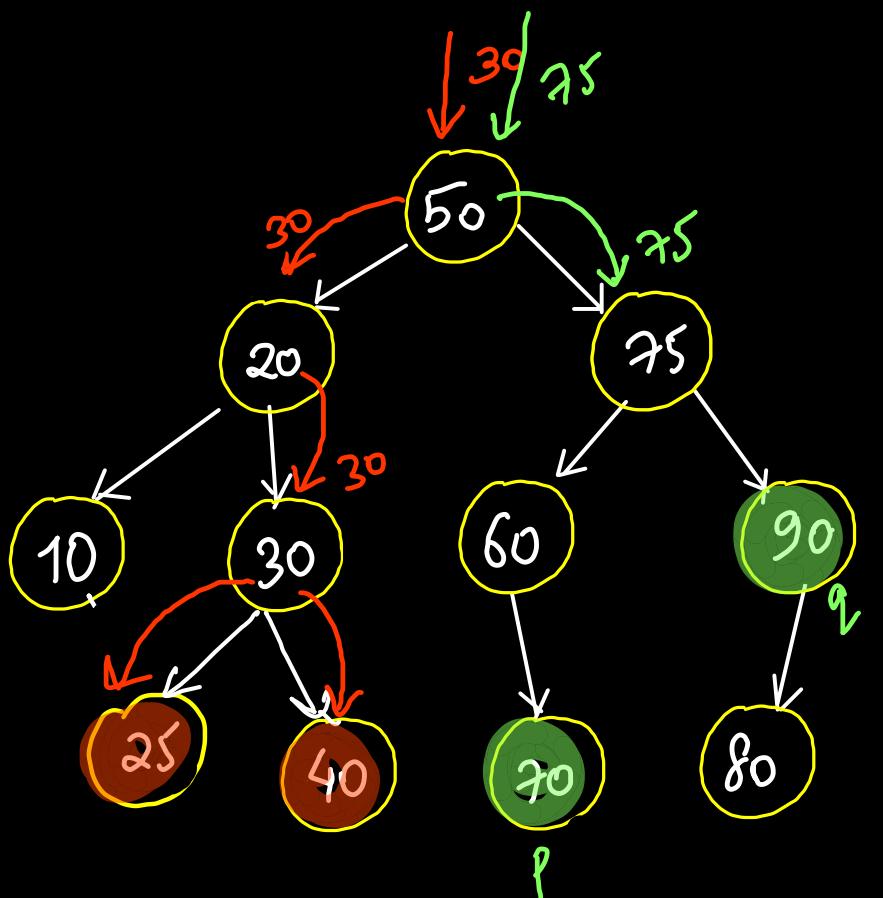
```

class Solution {
    public int floor(TreeNode root){
        root = root.left;
        while(root.right != null){
            root = root.right;
        }
        return root.val;
    }

    public TreeNode deleteNode(TreeNode root, int key) {
        if(root == null) return null;
        if(key < root.val) root.left = deleteNode(root.left, key);
        if(key > root.val) root.right = deleteNode(root.right, key);
        if(key == root.val){
            if(root.left == null) return root.right;
            if(root.right == null) return root.left;
            else {
                root.val = floor(root);
                root.left = deleteNode(root.left, root.val);
            }
        }
        return root;
    }
}

```





LCA^{235} Lowest Common Ancestor

$$\text{LCA}(30, 10) = 20$$

$$\text{LCA}(70, 80) = 75$$

$$\text{LCA}(25, 60) = 50$$

$$\text{LCA}(30, 40) = 30$$

App1) Brute force $\Rightarrow \text{dfs } \underline{\underline{O(n)}}$

App2)

```
public TreeNode lowestCommonAncestor(TreeNode root, TreeNode p, TreeNode q) {  
    if(p.val < root.val && q.val < root.val)  
        return lowestCommonAncestor(root.left, p, q);  
    if(p.val > root.val && q.val > root.val)  
        return lowestCommonAncestor(root.right, p, q);  
    return root;  
}
```

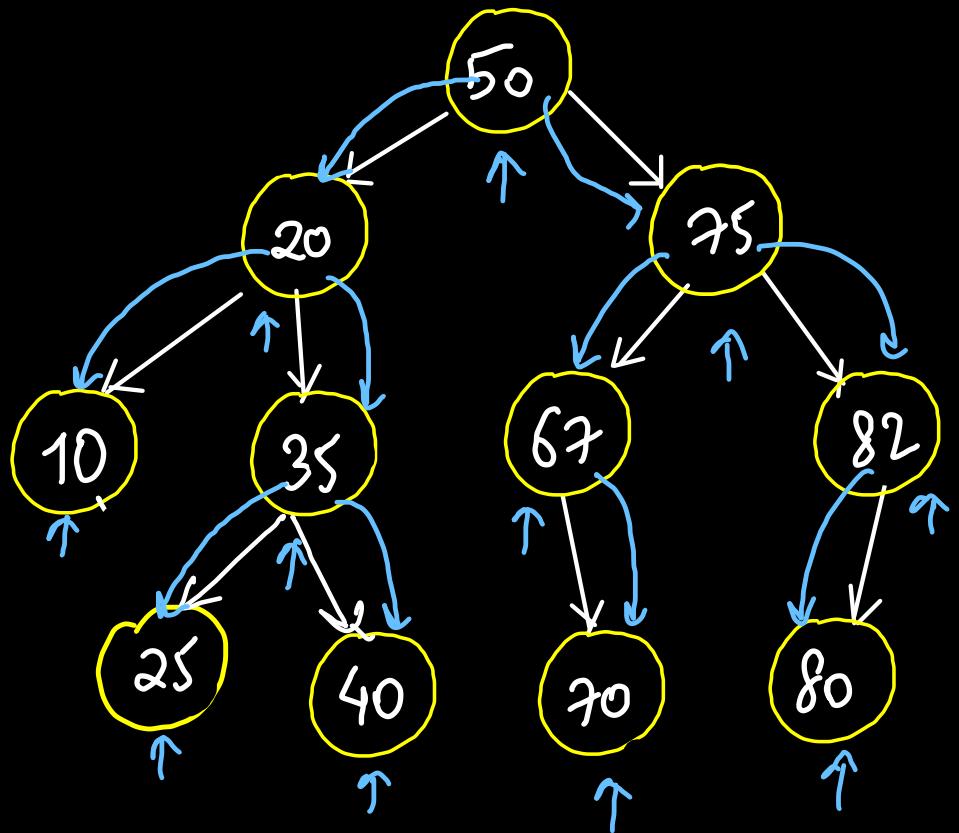
$\mathcal{O}(\log n)$ avg (balanced)

$\mathcal{O}(h)$ time/space

$\mathcal{O}(n)$ worst (skewed)

~~LC 530~~ min Absolute Difference in BST

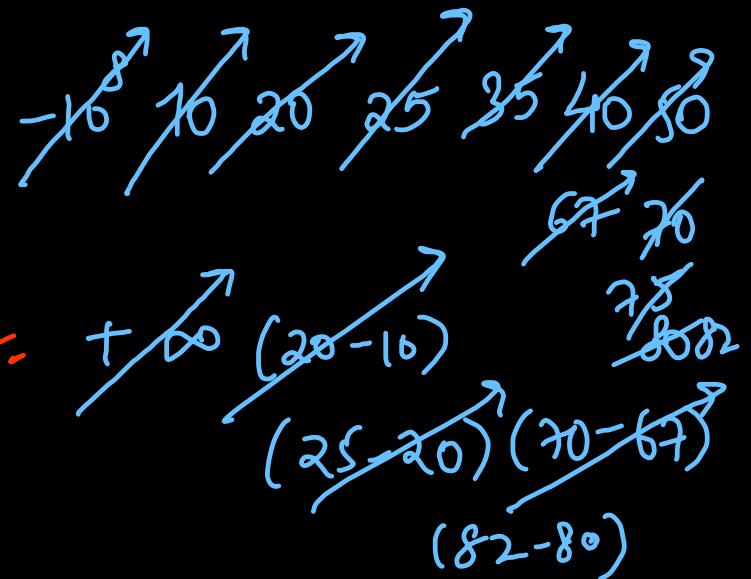
~~O(m)~~ inorder \rightarrow sorted increasing



prev =

.

mindiff =



```
int prev = (int)-1e8, diff = (int)1e8;

public void inorder(TreeNode root){
    if(root == null) return;
    inorder(root.left);
    diff = Math.min(diff, root.val - prev);
    prev = root.val;
    inorder(root.right);
}

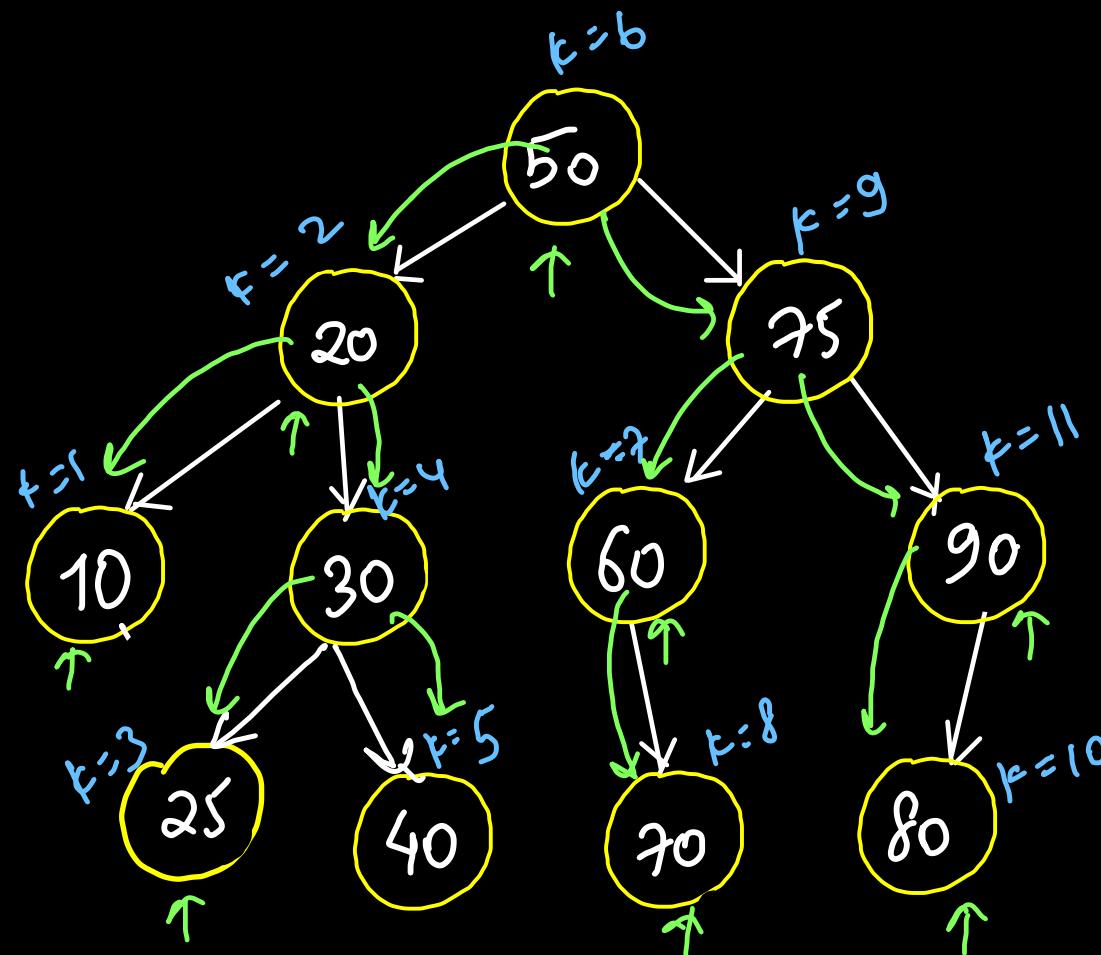
public int getMinimumDifference(TreeNode root) {
    inorder(root);
    return diff;
}
```

Time = $O(n)$
dfs

Space = $O(h)$
recursion call

\downarrow
 $O(\log n)$
avg
 $O(n)$
worst

k th smallest Element



$k=8$?

count = $\emptyset / \{5\} / \{4, 5\} / \{4, 7\} / \{4, 8\}$
ans = $\nearrow 20$
 g
 10
 11

```
class Solution {
    int count, ans;

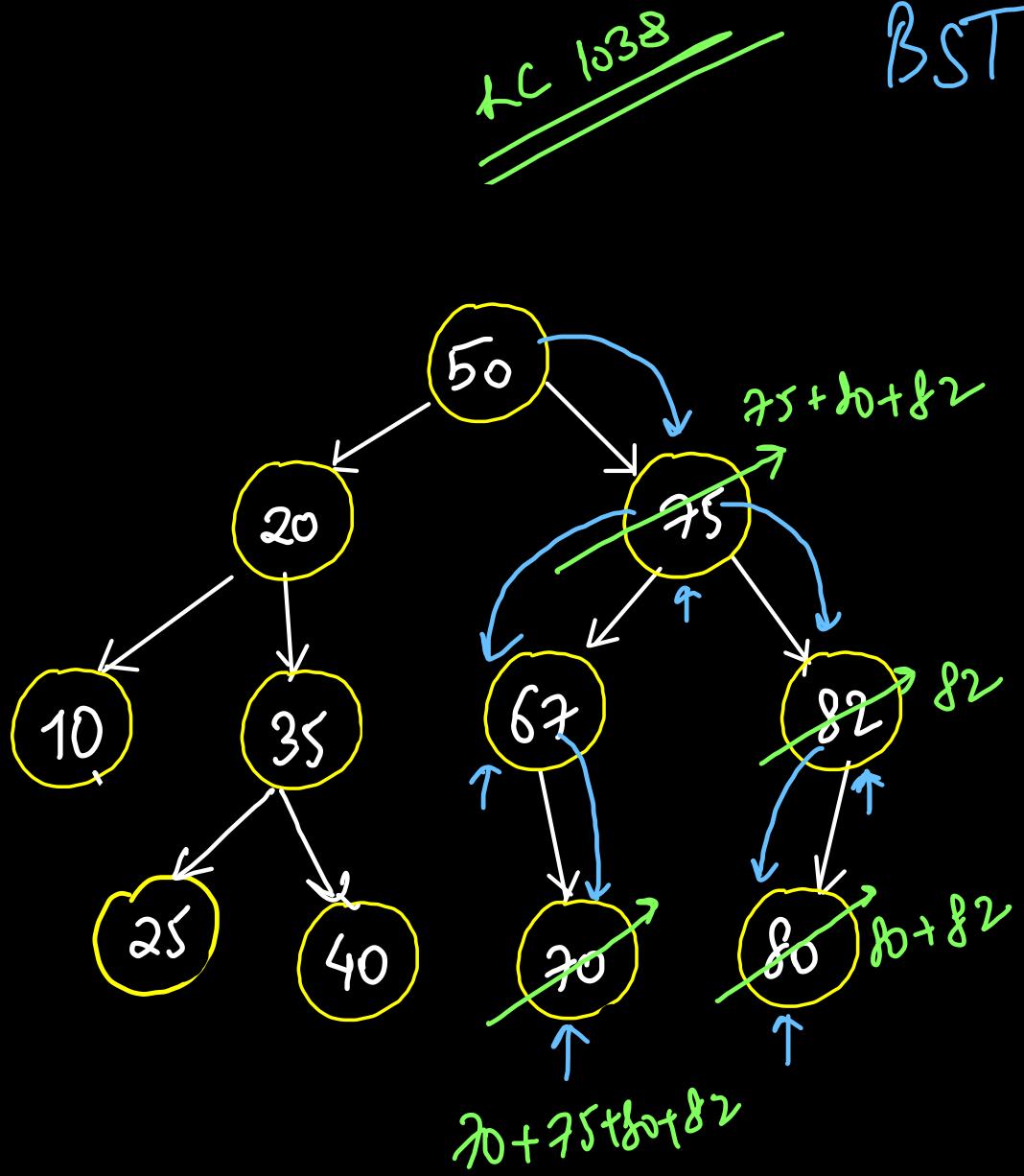
    public void inorder(TreeNode root){
        if(root == null) return;
        inorder(root.left);
        count--;
        if(count == 0) ans = root.val;
        inorder(root.right);
    }

    public int kthSmallest(TreeNode root, int k) {
        this.count = k;
        inorder(root);
        return ans;
    }
}
```

Time = $O(n)$
dfs

Space = $O(h)$
recursion call

$O(\log n)$ ↘ $O(n)$ ↓
avg worst



BST +> GST

reverse/inorder
(decreasing order)



`dfs(root.right)`

`sum += root.val`

`root.val = sum`

`dfs(root.left)`

$$\text{Sum} = 0 + 82 + 80 + 75 + 70 + 67$$

```
class Solution {
    int sum = 0;

    public void dfs(TreeNode root){
        if(root == null) return;
        dfs(root.right);
        sum += root.val;
        root.val = sum;
        dfs(root.left);
    }

    public TreeNode bstToGst(TreeNode root) {
        dfs(root);
        return root;
    }
}
```

Time = $O(n)$

dfs

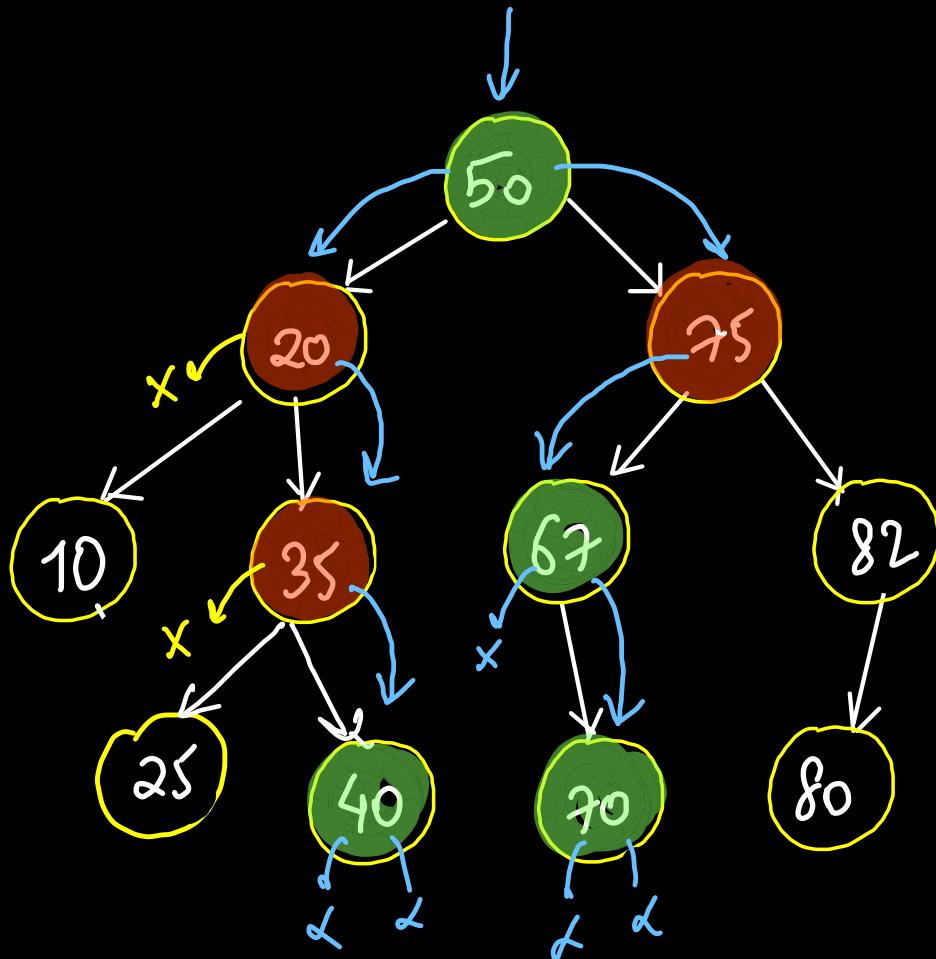
Space = $O(h)$

recursion call

$O(\log n)$ ↘ $O(n)$ ↓
avg worst

~~LC 938~~

Range Sum of BST



Approach 1) Normal dfs
 $O(n)$ time

Approach 2) dfs + prunings

[37, 72]

$$\text{sum} = 0 + 50 + 40 + 67 + 70$$

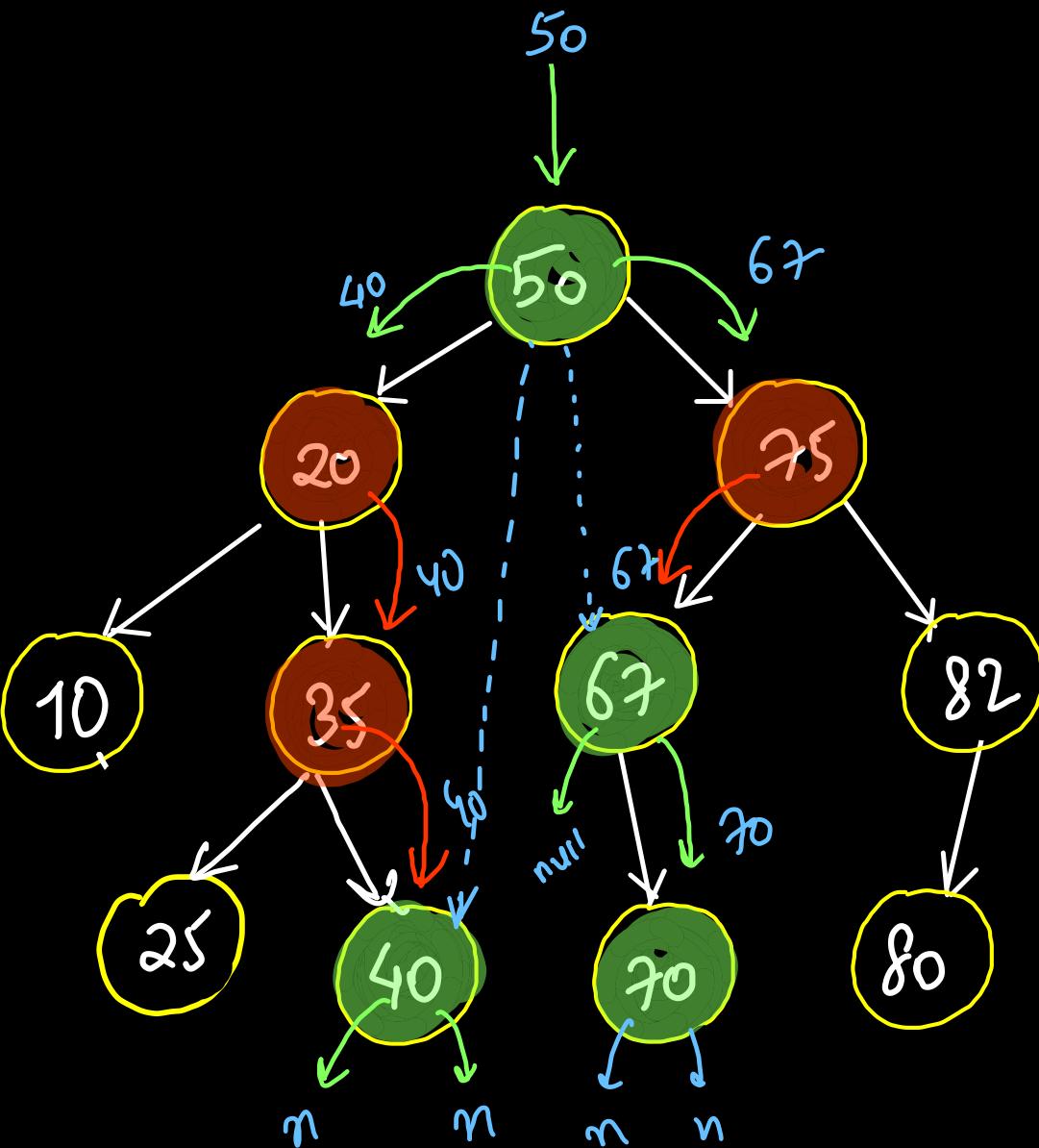
```
public int rangeSumBST(TreeNode root, int low, int high) {  
    if(root == null) return 0;  
    if(root.val < low) return rangeSumBST(root.right, low, high);  
    if(root.val > high) return rangeSumBST(root.left, low, high);  
  
    int sum = root.val;  
    sum += rangeSumBST(root.left, low, high);  
    sum += rangeSumBST(root.right, low, high);  
    return sum;  
}
```

$O(n)$ worst case

$O(\text{range})$ avg

Trim a BST

LC 669



[37, 72]

```
public TreeNode trimBST(TreeNode root, int low, int high) {
    if(root == null) return null;
    c1 if(root.val < low) return trimBST(root.right, low, high);
    c2 if(root.val > high) return trimBST(root.left, low, high);

    c3 { root.left = trimBST(root.left, low, high);
          root.right = trimBST(root.right, low, high); }
    return root;
}
```

root is in range if it should not be deleted

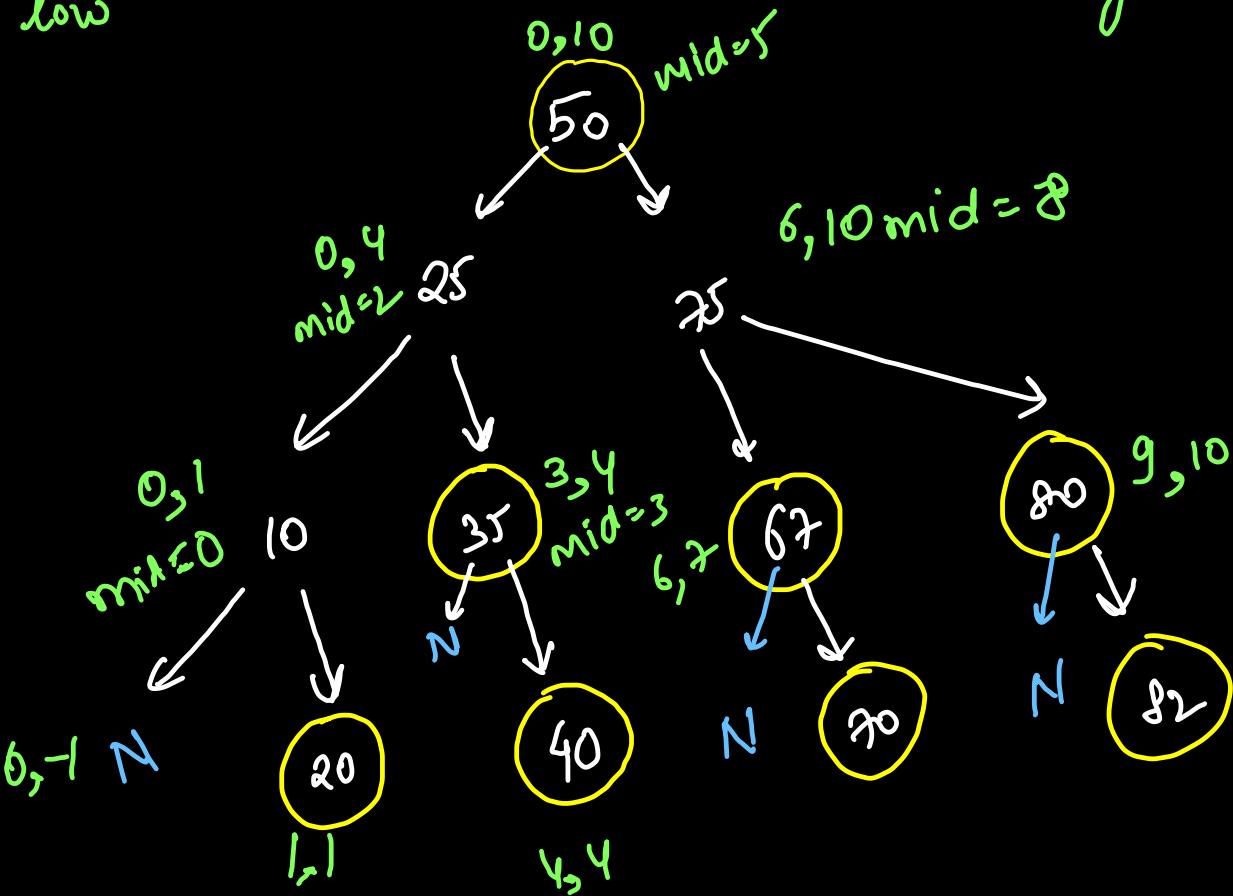
case1: delete root + left subtree
eg 20

case2: delete root + right subtree
eg 25

~~xc108~~

Sorted Array to BST $\xrightarrow{\text{height-balanced}} \text{height} = \lceil h \rceil = \lceil \log n \rceil$

{ 10, 20, 25, 35, 40, 50, 67, 70, 75, 80, 82 }
low ↑ mid ↑ high ↑



$$T(n) = 2T(n/2) + O(1)$$

$$\begin{aligned} T(n) &= 2^h \cdot h + 1 \\ &= 2^{\lceil \log n \rceil} \cdot \lceil \log n \rceil \\ &= n \log n \end{aligned}$$

```

public TreeNode dfs(int[] nums, int low, int high){
    if(low > high) return null;

    int mid = low + (high - low) / 2;
    TreeNode root = new TreeNode(nums[mid]);

    root.left = dfs(nums, low, mid - 1);
    root.right = dfs(nums, mid + 1, high);
    return root;
}

public TreeNode sortedArrayToBST(int[] nums) {
    return dfs(nums, 0, nums.length - 1);
}

```

Time $\Rightarrow O(n)$

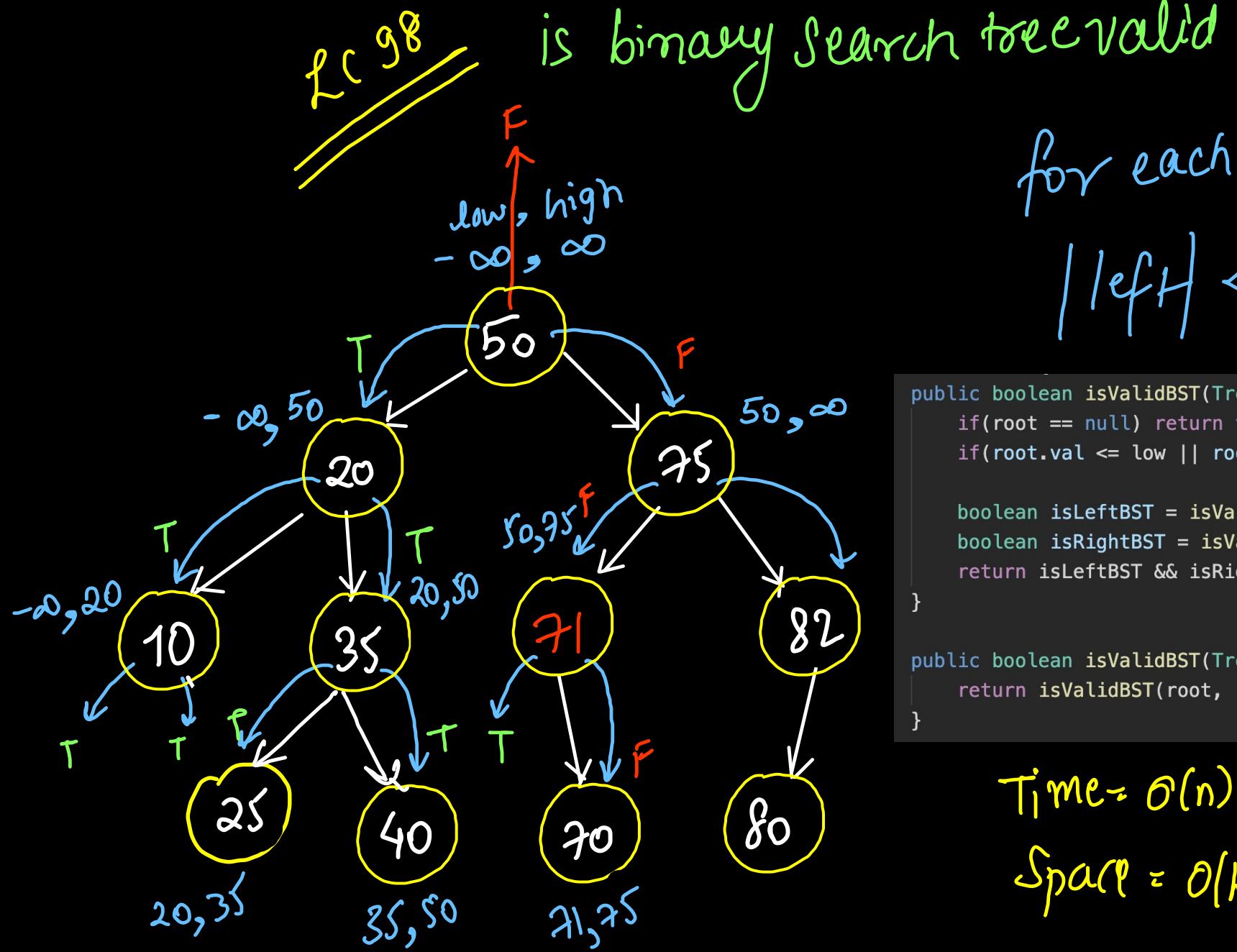
Space $\Rightarrow O(\log n)$

recursion

Recursion
calls height + workheight

$2^{\log_2 n} + O(1) * \log_2 n$

$n + \log_2 n$



for each node

$$|\text{left}| < \text{root} < |\text{right}|$$

```
public boolean isValidBST(TreeNode root, long low, long high){
    if(root == null) return true;
    if(root.val <= low || root.val >= high) return false;

    boolean isLeftBST = isValidBST(root.left, low, root.val);
    boolean isRightBST = isValidBST(root.right, root.val, high);
    return isLeftBST && isRightBST;
}

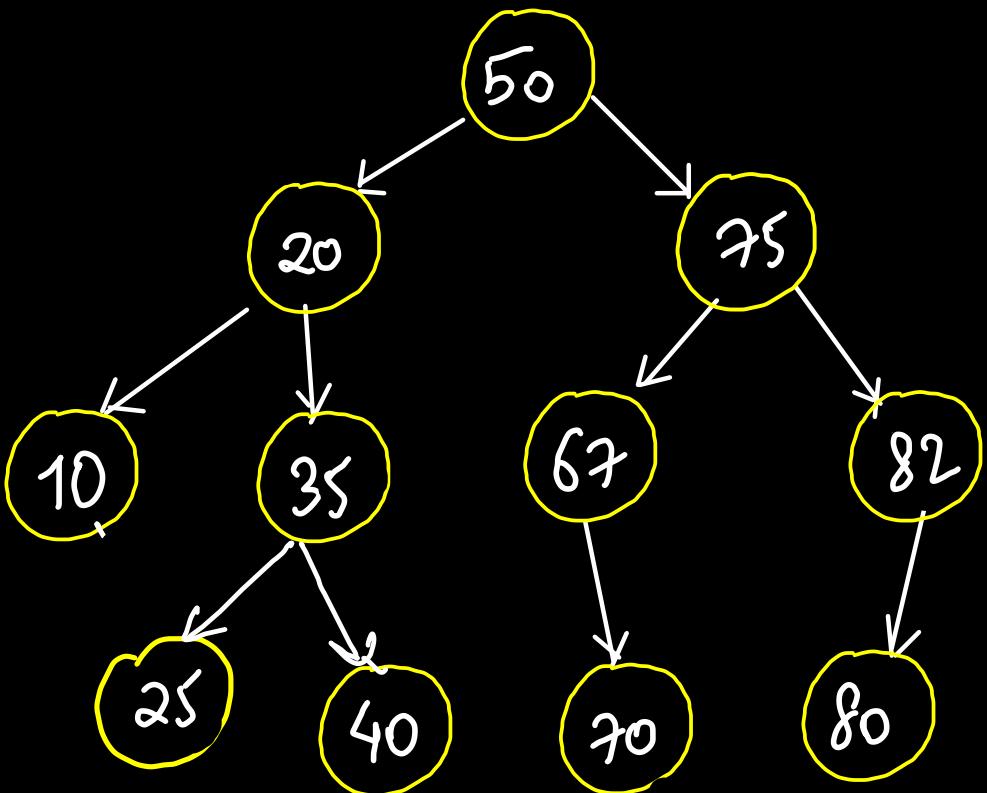
public boolean isValidBST(TreeNode root) {
    return isValidBST(root, (long)-1e15, (long)1e15);
}
```

Time = $O(n)$

Space = $O(h)$

Iterative DFS Traversal of Binary Tree

→ Preorder/Inorder/Postorder



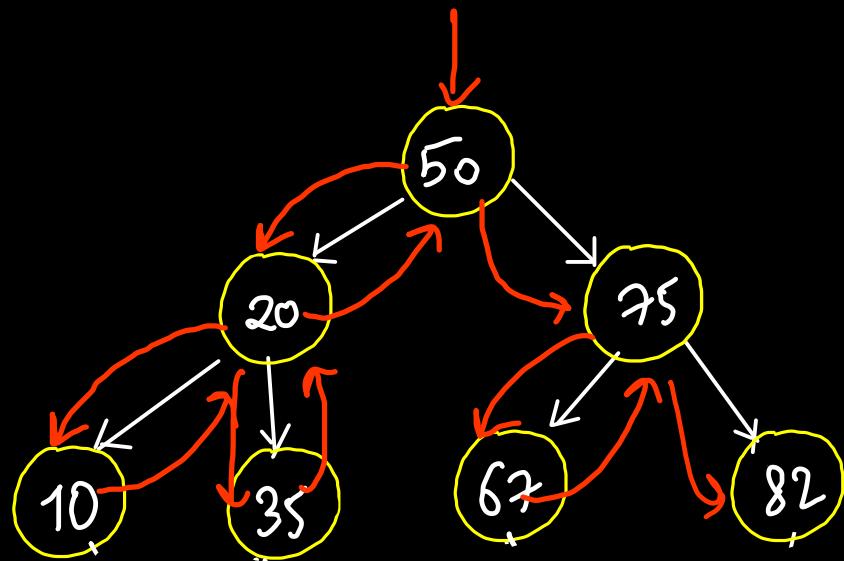
Stack <Pair>

node

Integer

1, 2, 3

↙ ↓ ↓
Pre In Postorder



Preorder : 50 20 10 35 75 67 82

Inorder: 10 20 35 50 67 75 82

Postorder: 10 35 20 67 82 75 50

Stack (DFS)

```
static class Pair{  
    TreeNode root;  
    int state = 1; // 1 -> preorder, 2 -> inorder, 3 -> postorder  
    Pair(TreeNode root) { this.root = root; }  
}
```

Time = $O(n)$ dfs

Space = $O(h)$ stack space

$\delta(\log n)$
avg

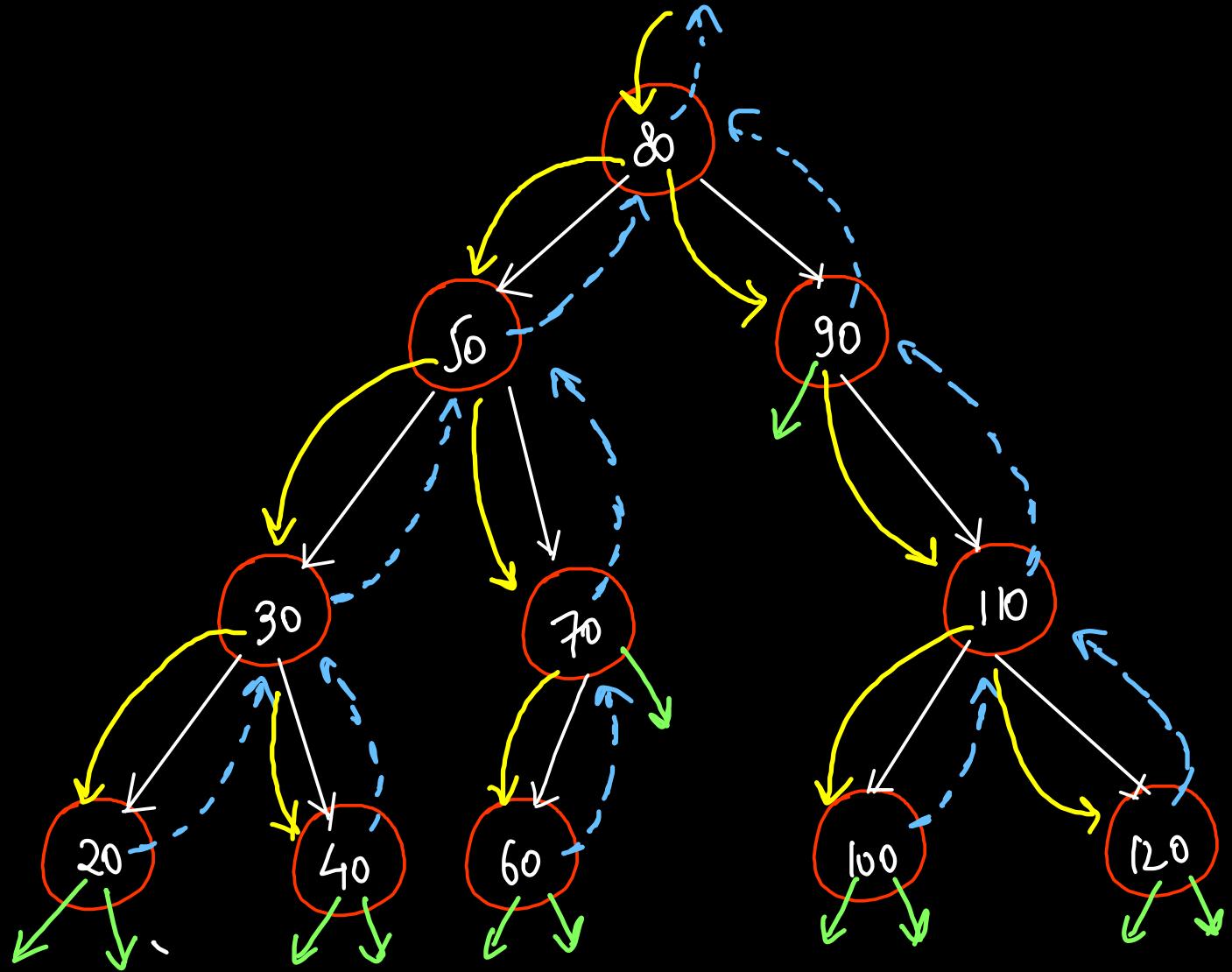
$O(n)$
worst-

No
Stack
Overflow
Exceptions

```
List<Integer> preorder = new ArrayList<>();  
List<Integer> inorder = new ArrayList<>();  
List<Integer> postorder = new ArrayList<>();  
  
Stack<Pair> stk = new Stack<>();  
if(root != null) stk.push(new Pair(root));  
  
while(stk.size() > 0){  
    root = stk.peek().root;  
    int state = stk.peek().state++;  
  
    if(state == 1){  
        preorder.add(root.val);  
        if(root.left != null)  
            stk.push(new Pair(root.left));  
    }  
    else if(state == 2){  
        inorder.add(root.val);  
        if(root.right != null)  
            stk.push(new Pair(root.right));  
    }  
    else {  
        postorder.add(root.val);  
        stk.pop();  
    }  
}
```

Tree Traversals

Recursive DFS	Iterative DFS	BFS or Level order	Morris Traversal DFS (Pre/In/Postorder)
DS: Recursion call stack (at max 20000)	DS: Stack(Pair) (inside heap) (at max $10^5 - 10^6$)	DS: Queue or Linked List (at max $10^5 - 10^6$)	DS: Threaded Binary Tree (same input tree)
Time = $O(n)$	Time = $O(n)$	Time = $O(n)$	Time = $O(n)$ \propto linear
Space = $O(h)$ $O(\log n)$ $O(n)$	Space = $O(h)$ $O(\log n)$ $O(n)$	Space = $O(2 \times \text{breadth})$ = $O(n)$ <u>balanced</u>	Space = $O(1)$ no extra/constant Space



1. $\text{dfs}(\text{root-left})$

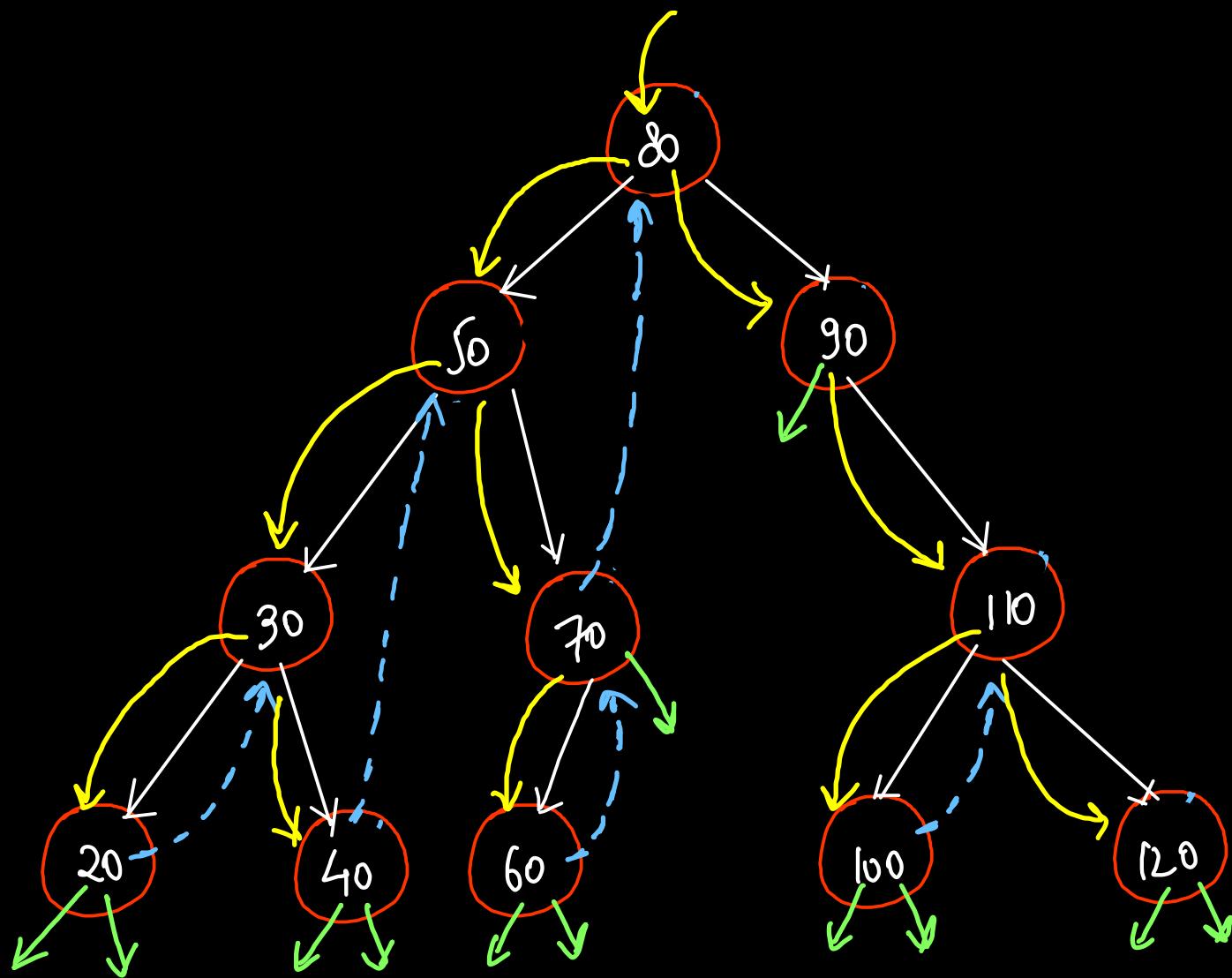
2. $\text{dfs}(\text{root-right})$

white \rightarrow tree edges
 yellow \rightarrow forward edges
 blue \rightarrow backtracking /
 backward edge

(1) Threads → create

⇒ ~~loops/cycle~~

if threads
are created
initially

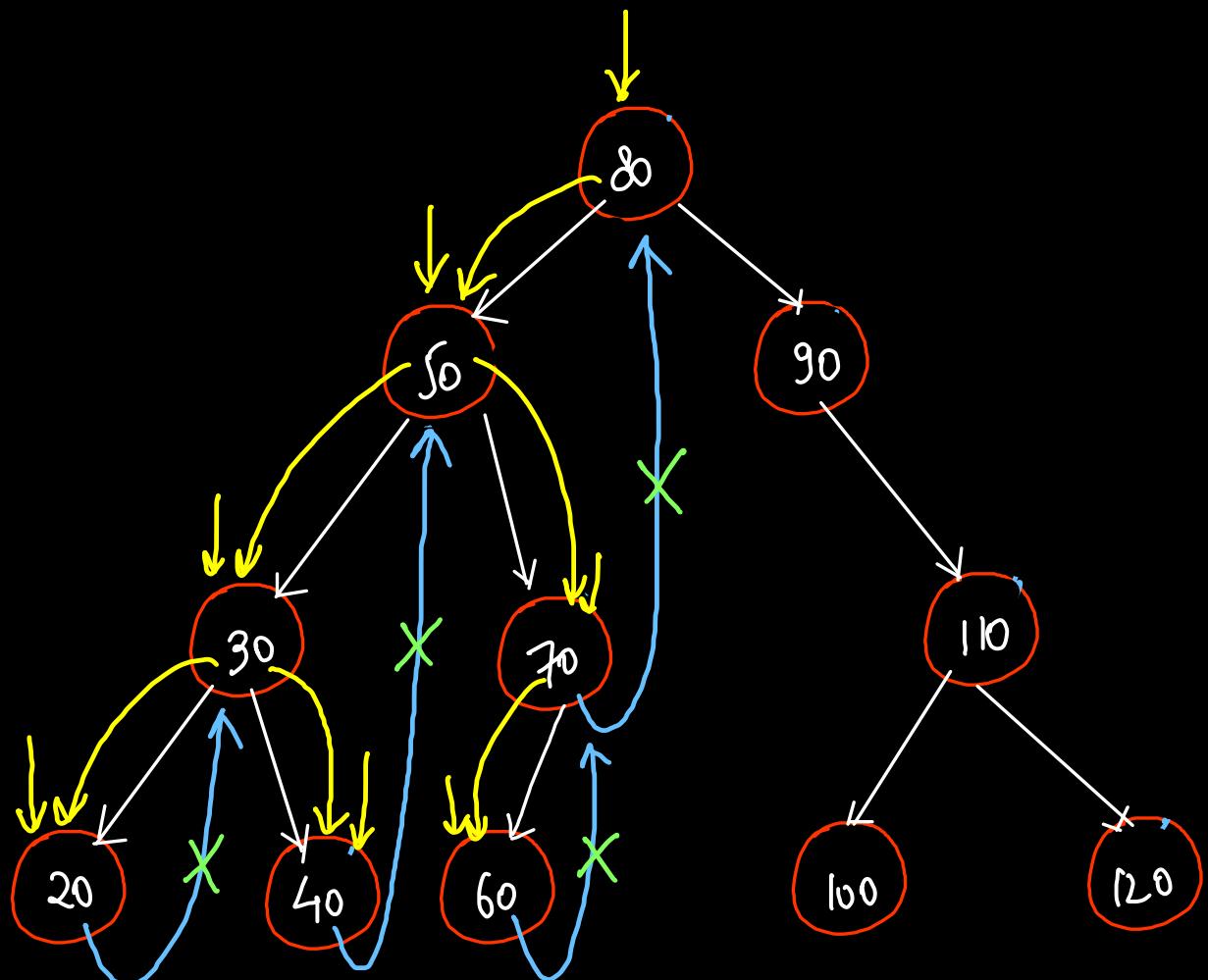


20 → 30 40 → 50 60 → 70 70 → 80 100 → 110

 110 → 120

 { threads
 ① floor
 ② in-order
 predecessor}

Intuition



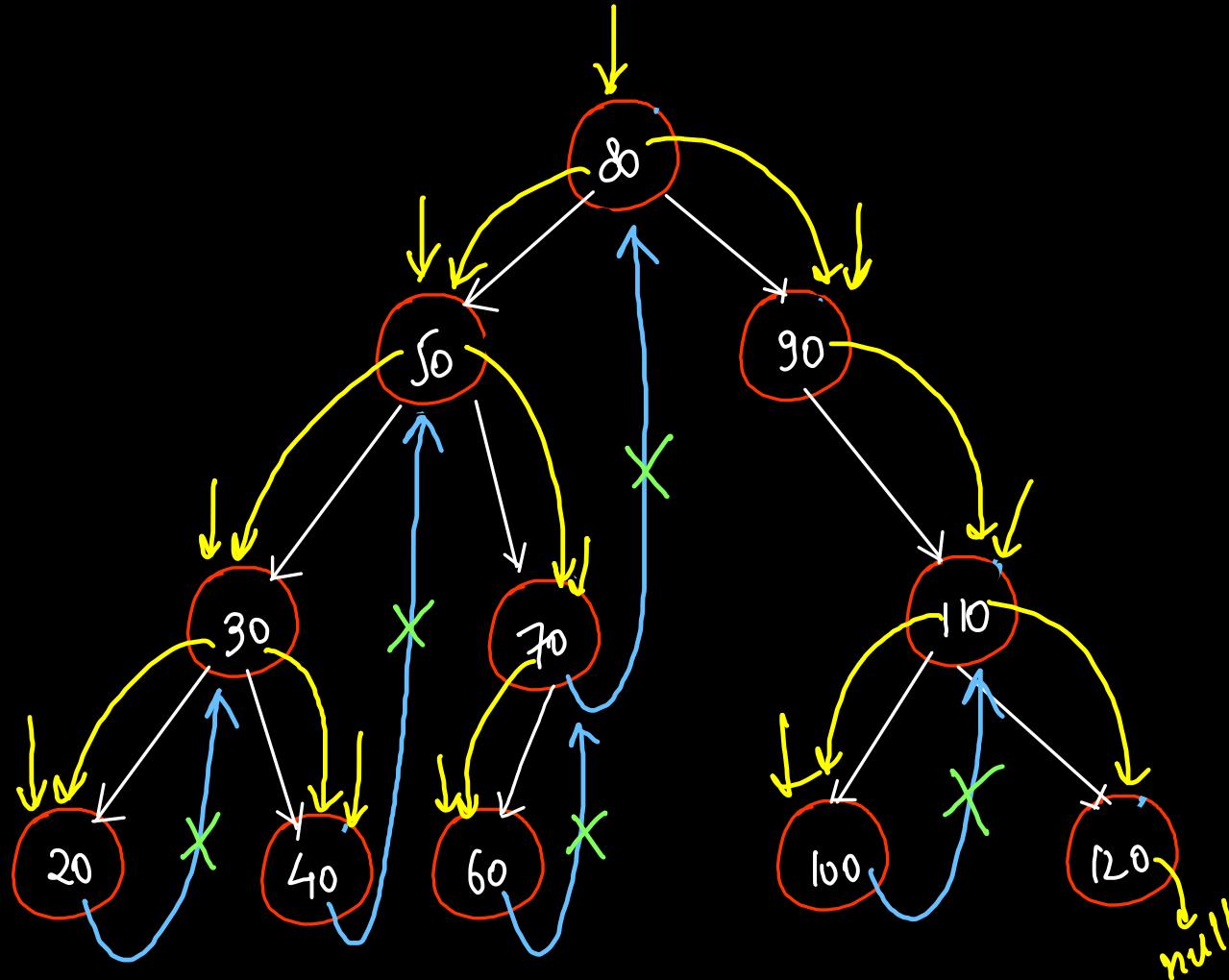
- preorder {
- ① thread → create
 - ② left subtree
- inorder {
- ③ thread → delete
 - ④ right subtree

preorder:

80 50 30 20 40 70 60

inorder:

20 30 40 50 60 70 80



preorder:

80 50 30 20 40 70 60 90 110 100 120

inorder:

20 30 40 50 60 70 80 90 100 110 120

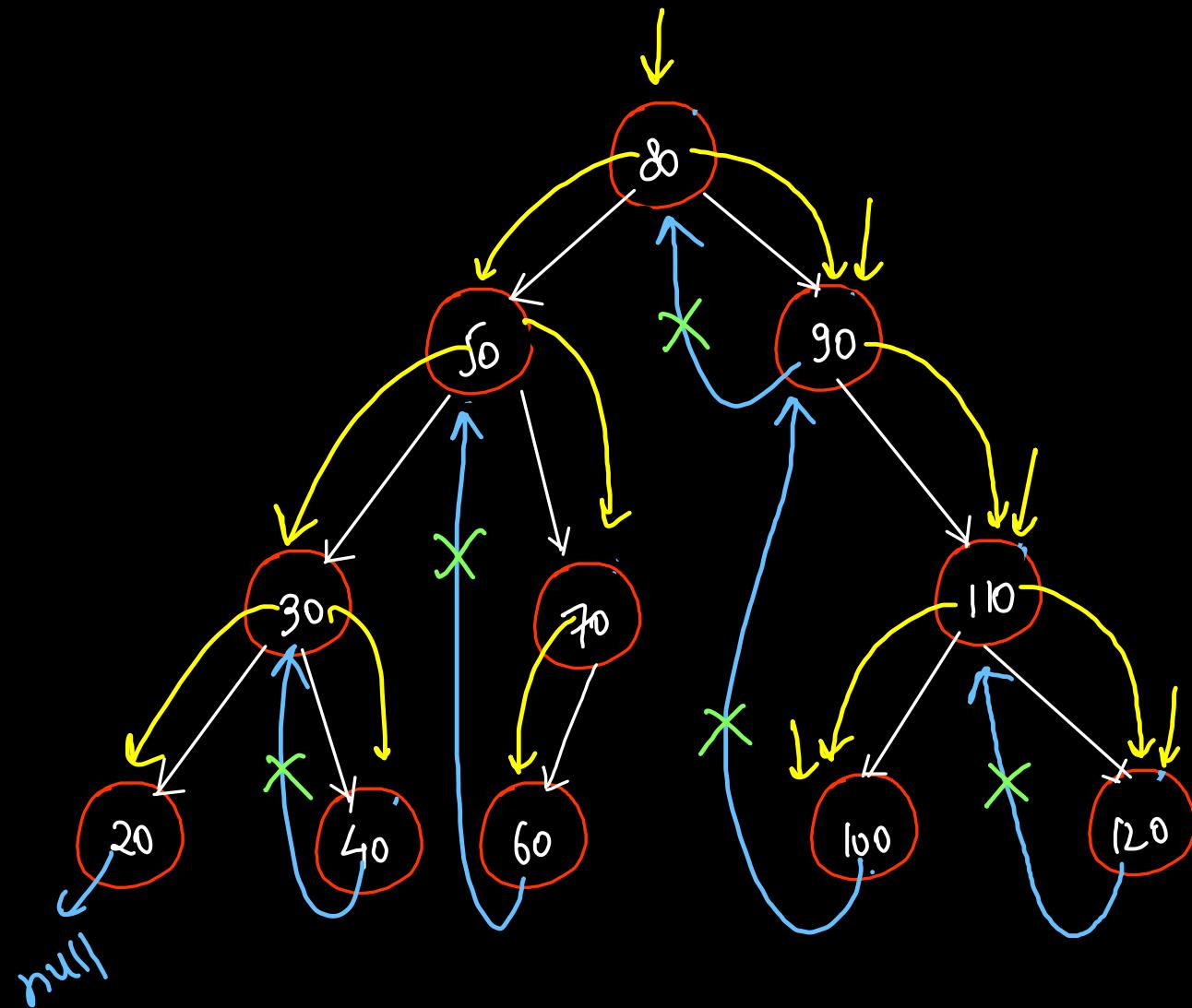
```
public TreeNode floor(TreeNode root){
    TreeNode ans = root.left;
    while(ans.right != null && ans.right != root)
        ans = ans.right;
    return ans;
}
```

```
List<Integer> preorder = new ArrayList<>();
List<Integer> inorder = new ArrayList<>();
```

```
while(root != null){
    if(root.left == null){
        preorder.add(root.val);
        inorder.add(root.val);
        root = root.right;
        continue;
    }
}
```

Time = O(n)
Space = O(1)
Constant

```
TreeNode floor = floor(root);
if(floor.right == null){
    floor.right = root; // thread creation
    preorder.add(root.val);
    root = root.left;
} else {
    floor.right = null; // thread deletion
    inorder.add(root.val);
    root = root.right;
}
```



hint : Use left threaded binary tree

Morris Postorder

left → right → root
↓ reverse

root → right → left

80 90 110 120 100
50 70 60 30 40 20

{
 ↓ ↓ ↓ ↓ ↓
 10, 20, 30, 40 } n=4
 0 1 2 3

hasNext()

boolean

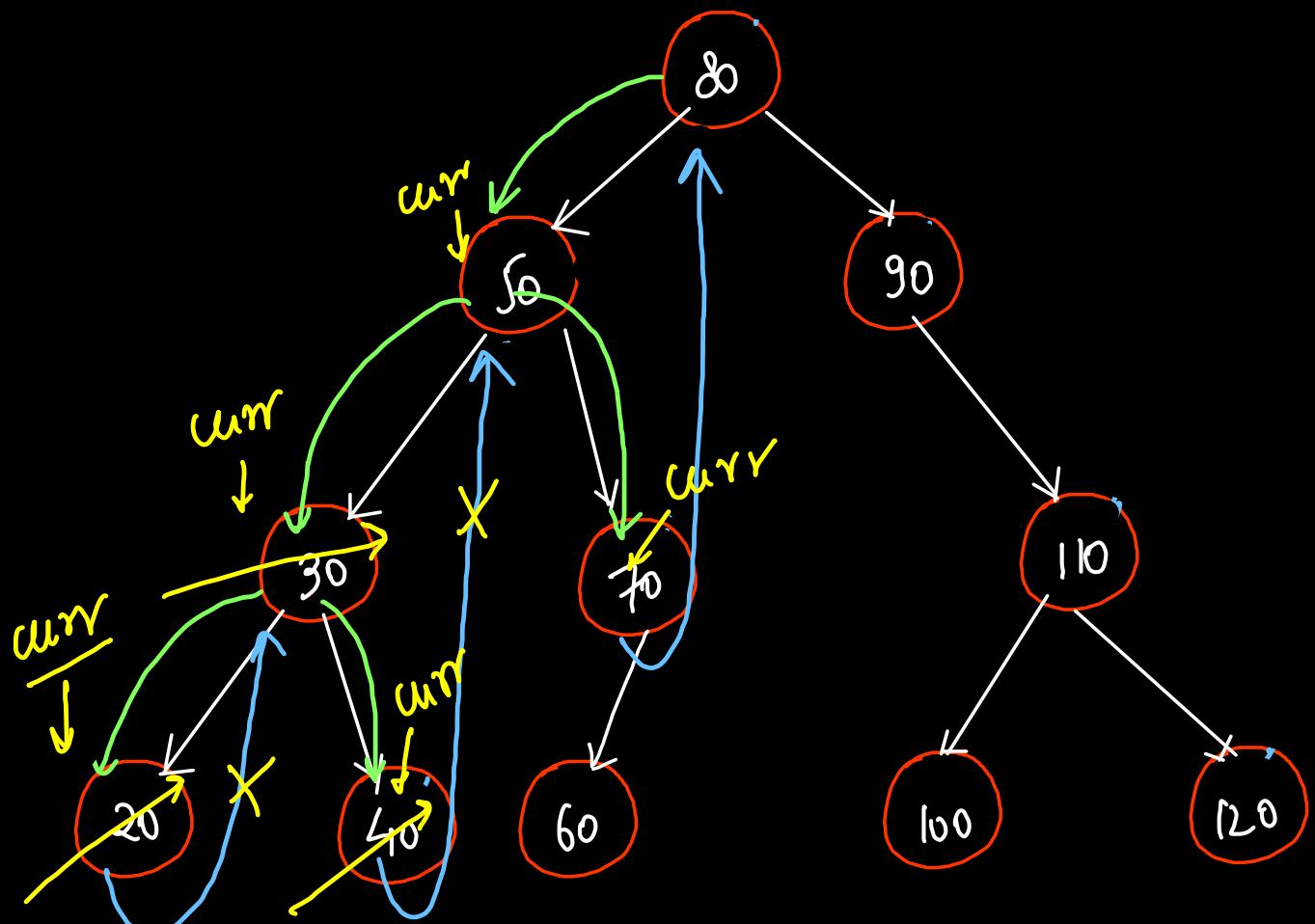
↓

if (ptr == n)
 return false
else return true;

next()

returns curr value

& goes to next value



Constructor

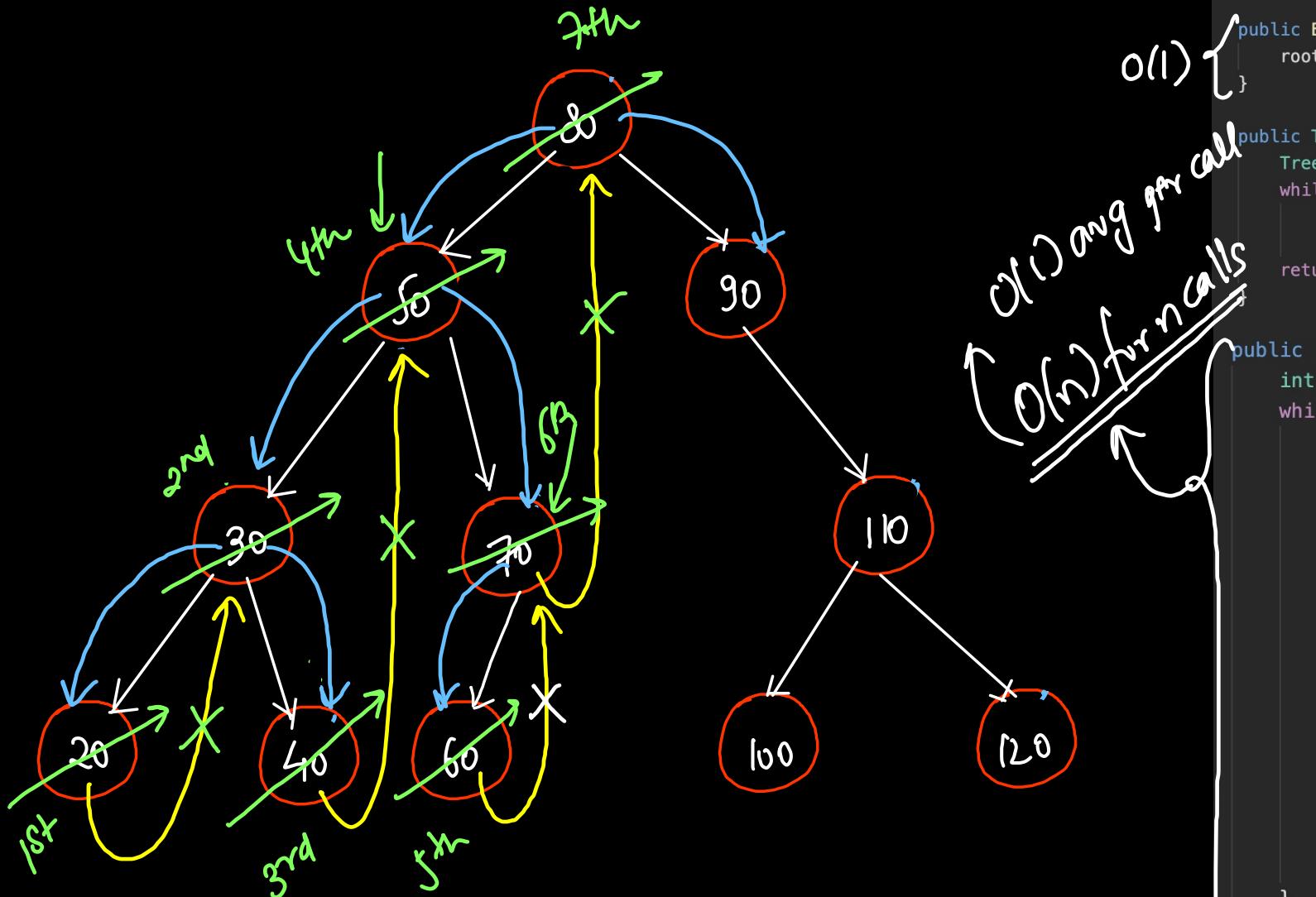
\rightarrow $\text{root} == \text{smallest node}$

$\text{next}() \rightarrow 20$

$\text{next}() \rightarrow 30$

$\text{next}() \rightarrow 40$

$\text{next}() \rightarrow 50$



```

TreeNode root;

public BSTIterator(TreeNode curr) {
    root = curr;
}

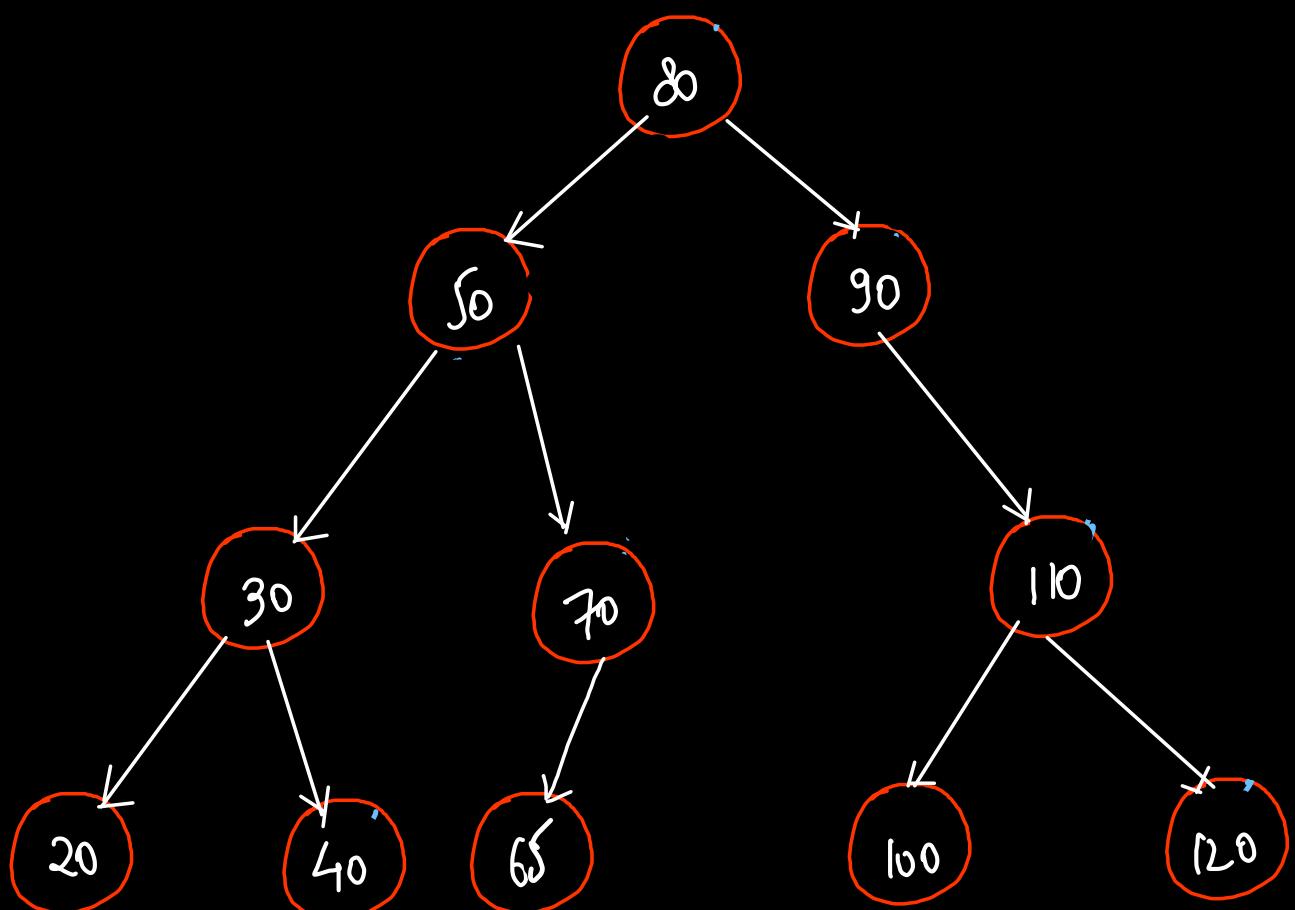
public TreeNode floor(TreeNode root) {
    TreeNode ans = root.left;
    while (ans.right != null && ans.right != root)
        ans = ans.right;
    return ans;
}

public int next() {
    int ans = -1;
    while(root != null){
        if(root.left == null){
            ans = root.val;
            root = root.right;
            return ans;
        }
        TreeNode floor = floor(root);
        if(floor.right == null){
            floor.right = root; // thread creation
            root = root.left;
        } else {
            floor.right = null; // thread deletion
            ans = root.val;
            root = root.right;
            break;
        }
    }
    return ans;
}

public boolean hasNext() {
    if(root == null) return false;
    return true;
}

```

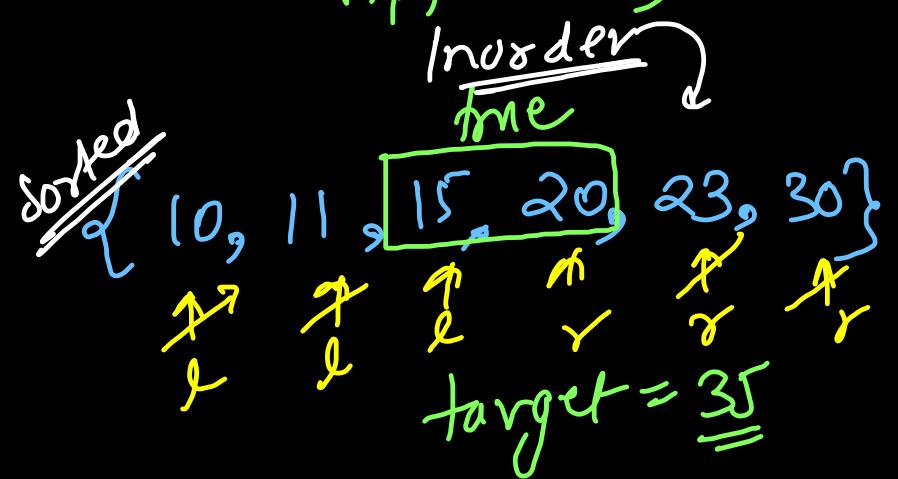
$O(1)$



Target Sum Pair

pair target = 135

Approach 1)



$$10 + 30 = 40 \Rightarrow \gamma^-$$

$$10 + 23 = 33 \Rightarrow \delta++$$

$$11 + 23 = 34 \Rightarrow \delta++$$

$$15 + 23 = 38 \Rightarrow \gamma^-$$

```

static class ForwardIterator{
    Stack<TreeNode> stk;

    public ForwardIterator(TreeNode root) {
        stk = new Stack<>();
        inorderSucc(root);
    }

    public void inorderSucc(TreeNode curr){
        while(curr != null){
            stk.push(curr);
            curr = curr.left;
        }
    }

    public int peek(){
        if(hasNext() == false) return 0;
        return stk.peek().val;
    }

    public int next() {
        if(hasNext() == false) return 0;
        TreeNode curr = stk.pop();
        inorderSucc(curr.right);
        return curr.val;
    }

    public boolean hasNext() {
        return (stk.size() > 0);
    }
}

```

```

static class BackwardIterator{
    Stack<TreeNode> stk;

    public BackwardIterator(TreeNode root) {
        stk = new Stack<>();
        inorderPred(root);
    }

    public void inorderPred(TreeNode curr){
        while(curr != null){
            stk.push(curr);
            curr = curr.right;
        }
    }

    public int peek(){
        if(hasPrev() == false) return 0;
        return stk.peek().val;
    }

    public int prev() {
        if(hasPrev() == false) return 0;
        TreeNode curr = stk.pop();
        inorderPred(curr.left);
        return curr.val;
    }

    public boolean hasPrev() {
        return (stk.size() > 0);
    }
}

```

Approach 2)

LC 653)

Target Sum Pair

in BST

```

public boolean findTarget(TreeNode root, int target) {
    if(root == null || (root.left == null & root.right == null)) return false;

    ForwardIterator left = new ForwardIterator(root);
    BackwardIterator right = new BackwardIterator(root);

    while(left.hasNext() == true && right.hasPrev() == true && left.peek() < right.peek()){
        if(left.peek() + right.peek() == target) return true;
        if(left.peek() + right.peek() < target) left.next();
        else right.prev();
    }

    return false;
}

```