

# CS 223 Computer Architecture and Organization

## Control Transfer Instructions



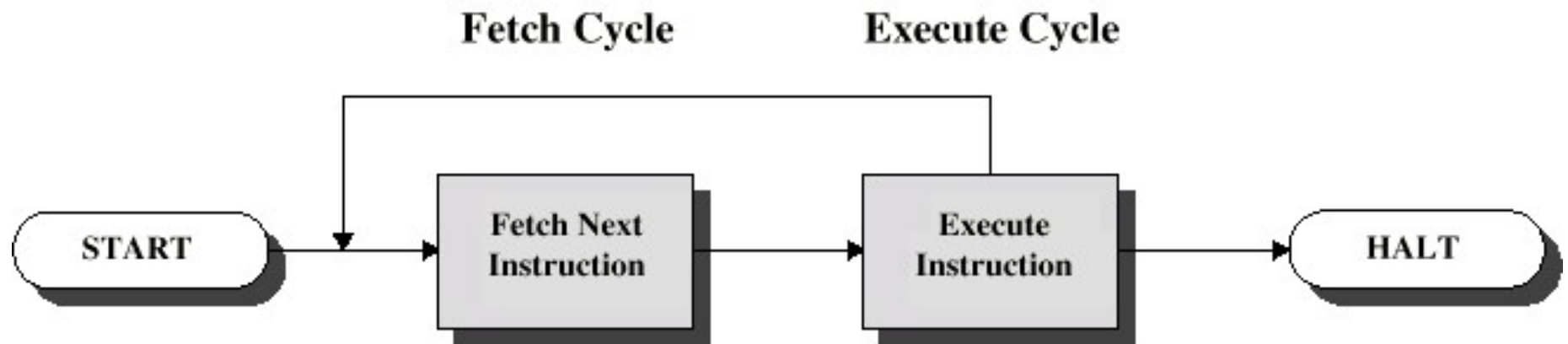
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# Instruction Cycle

- Two steps:
  - Fetch
  - Execute



# Conditional Instructions

- Unconditional Branch (BR)

- Format: 

OPCODE	Target Address
--------	----------------

Fetch Phase:

t1: MAR  $\leftarrow$  PC, Read

t2: MBR  $\leftarrow$  Memory

PC  $\leftarrow$  PC + 1

t3: IR  $\leftarrow$  MBR

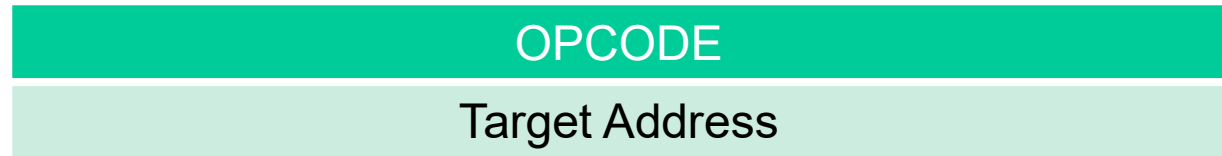
Execute Phase:

t4: PC  $\leftarrow$  IR<sub>Address</sub>

# Conditional Instructions

- Unconditional Branch

- Format:



Fetch Phase:

t1: MAR  $\leftarrow$  PC, Read

t2: MBR  $\leftarrow$  Memory

PC  $\leftarrow$  PC + 1

t3: IR  $\leftarrow$  MBR

Execute Phase:

# Conditional Instructions

- Conditional Branch (BRZ: Branch on Zero)
- Format: 

OPCODE	Target Address
--------	----------------

Fetch Phase:

t1: MAR  $\leftarrow$  PC, Read

t2: MBR  $\leftarrow$  Memory

PC  $\leftarrow$  PC + 1

t3: IR  $\leftarrow$  MBR

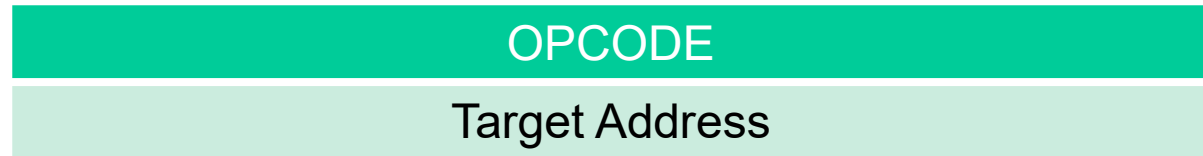
Execute Phase:

t4: If (Z = 1) PC  $\leftarrow$  IR<sub>Address</sub>

# Conditional Instructions

- Conditional Branch (BRZ: Branch on Zero)

- Format:



Fetch Phase:

t1: MAR  $\leftarrow$  PC, Read

t2: MBR  $\leftarrow$  Memory

PC  $\leftarrow$  PC + 1

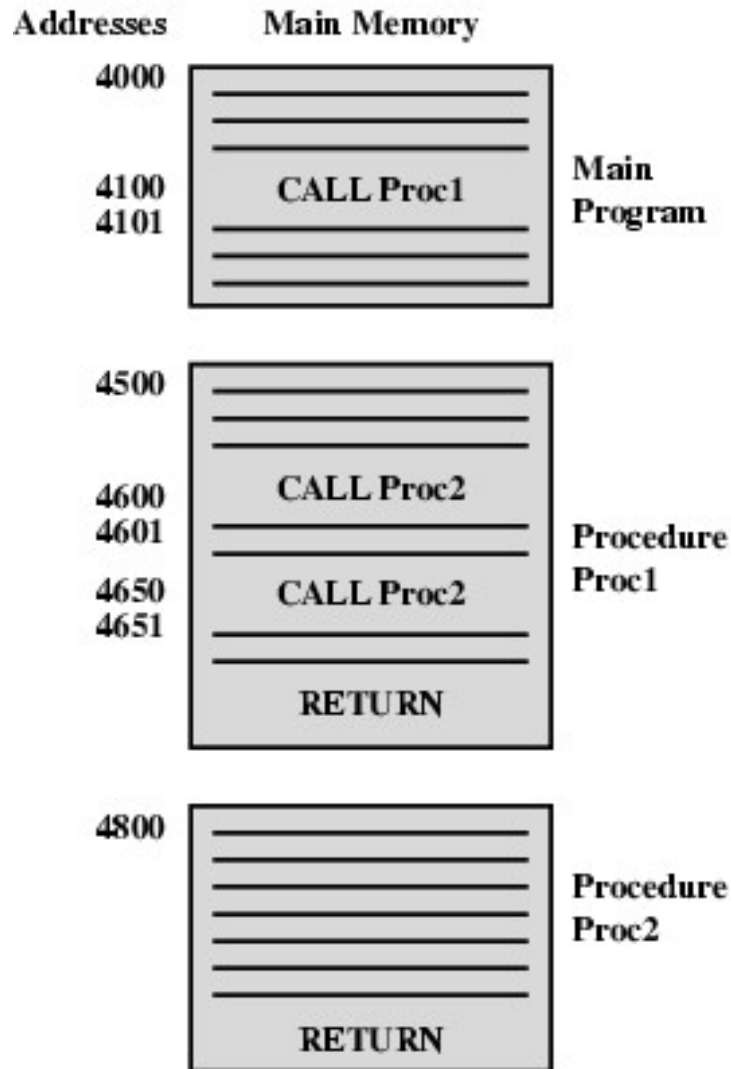
t3: IR  $\leftarrow$  MBR

Execute Phase:

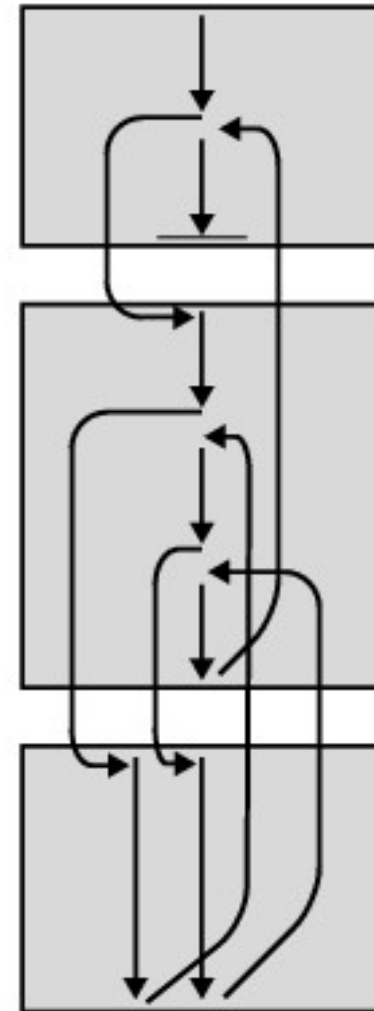
# Subroutine/Procedure/Function

- Independent unit of code to perform a subtask of the main task.
- Used in modular programming
- How to provide facility for procedure call
- Macros in Programming languages like C
- In case of Interrupt, device service routine is executed

# Nested Procedure Calls



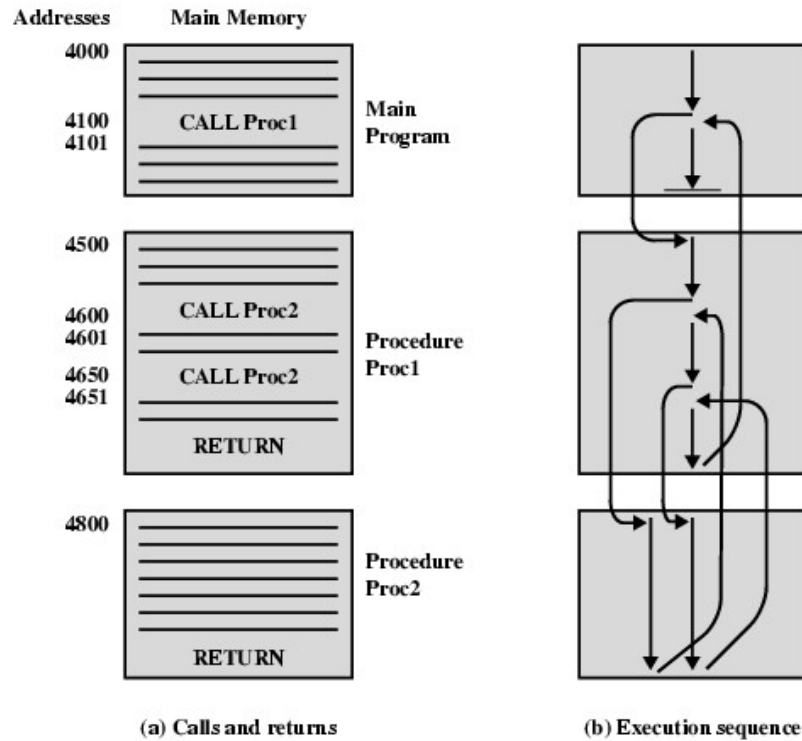
(a) Calls and returns



(b) Execution sequence



# Nested Procedure Calls



# Procedure Call

- Tasks to be performed before procedure CALL
  - Retain the current status of the processor
  - After returning from procedure/interrupt routine, we must restart the execution from the point where we have stopped.
- Current status of the processor
  - Program Counter
  - Program Status Word (PSW)
- How to Retain these information
- Any other information need to be saved?

# Modification in Organization

- Store the relevant information in main memory
  - Implement a stack in MM (Control Stack)
- Need to keep the address where to store
  - Use of a register, SP: Stack Pointer
  - To keep the address of the Top of the Stack
- After completion of the procedure, restore the information from stack

# Instructions

- PUSH R
  - source is the register R
- POP R
  - destination is the register R
- CALL address
  - starting address of the procedure
- RETURN
- Four different ways for implementation

# PUSH (Execute)

- PUSH Ri
  - $MAR \leftarrow SP$
  - $MDR \leftarrow Ri$
  - Write
  - $SP \leftarrow SP - 1$

# POP (Execute)

- POP Ri
  - $SP \leftarrow SP + 1$
  - $MAR \leftarrow SP$
  - Read
  - $Ri \leftarrow MDR$

# CALL (Execute)

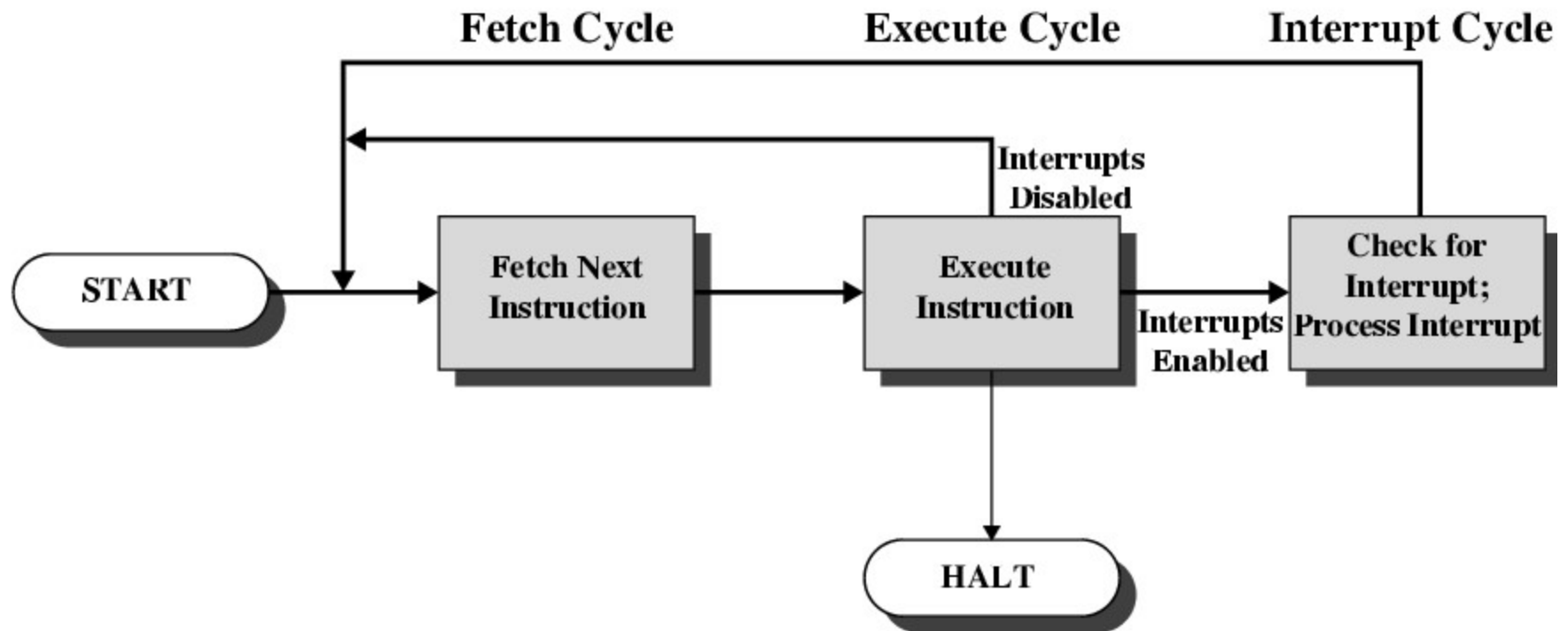
- CALL
  - $MAR \leftarrow SP$
  - $MDR \leftarrow PC$
  - Write
  - $SP \leftarrow SP - 1$
  - $MAR \leftarrow SP$
  - $MDR \leftarrow PSW$
  - Write
  - $SP \leftarrow SP - 1$
  - $PC \leftarrow IR_{\text{address}}$

# RETURN (Execute)

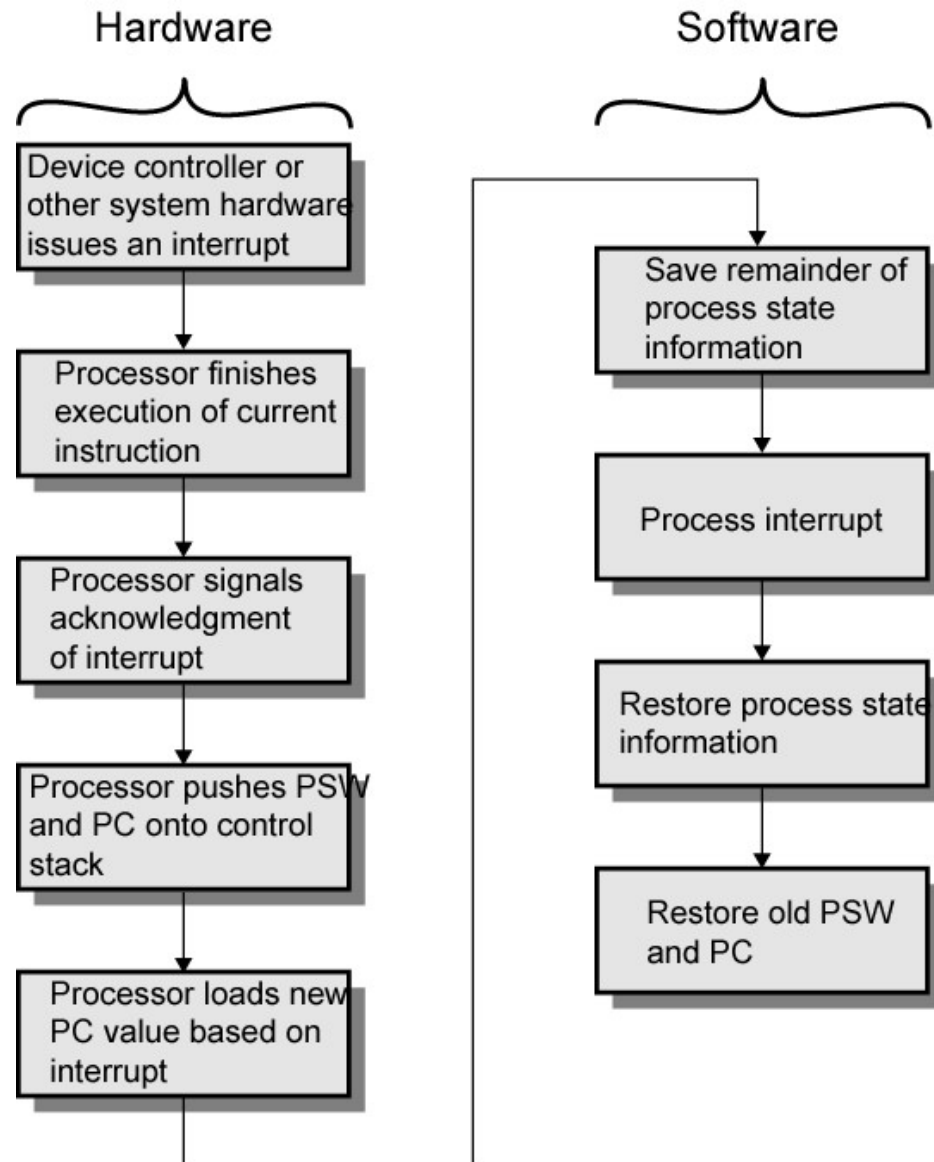
- RETURN
  - $SP \leftarrow SP + 1$
  - $MAR \leftarrow SP$
  - Read
  - $PSW \leftarrow MDR$
  - $SP \leftarrow SP + 1$
  - $MAR \leftarrow SP$
  - Read
  - $PC \leftarrow MDR$



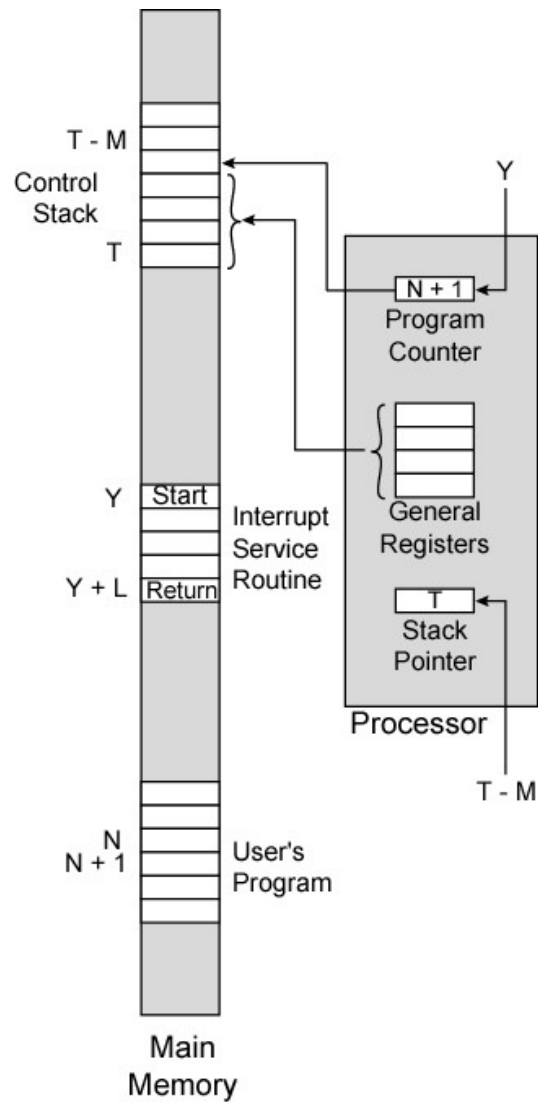
# Instruction Cycle with Interrupts



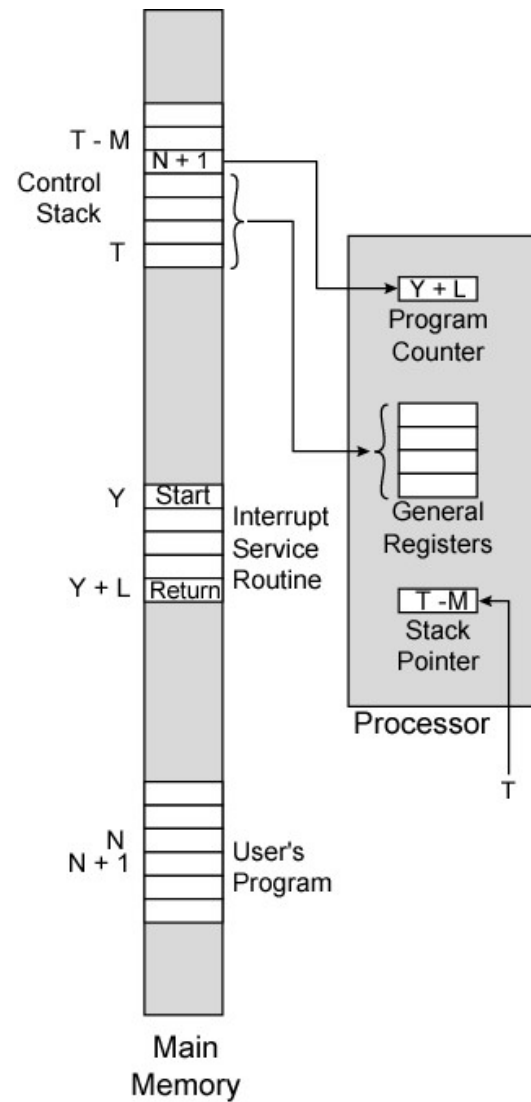
# Simple Interrupt Processing



# Memory and Registers for an Interrupt



(a) Interrupt occurs after instruction at location N



(b) Return from interrupt

