

1108 Invalid IP Address

```
1
2 func defangIPaddr(address string) string {
3     res := strings.Replace(address, ".", "[.]", -1)
4     return res
5 }
6
```

1135 Connecting Cities With Minimum Cost

执行结果： **通过** [显示详情 >](#)

 

执行用时： **184 ms** ，在所有 C++ 提交中击败了 **97.73%** 的用户

内存消耗： **46.6 MB** ，在所有 C++ 提交中击败了 **44.66%** 的用户

炫耀一下：

```
1  /*
2     Kruskal 算法
3
4  */
5  class WeightedUnionFind{
6  public:
7      vector<int> sz;
8      vector<int> id;
9
10     WeightedUnionFind(int N) : sz(vector<int>(N, 1)) , id(vector<int>(N, 0)){
11         for(int i = 0; i < N; i++)
12             id[i] = i;
13     }
14
15     void uni(int p, int q){
```

```

16         int pRoot = find(p);
17         int qRoot = find(q);
18
19         if(pRoot == qRoot)
20             return;
21
22         if(sz[pRoot] >= sz[qRoot]){
23             sz[pRoot] += sz[qRoot];
24             id[qRoot] = id[pRoot];
25         }else{
26             sz[qRoot] += sz[pRoot];
27             id[pRoot] = id[qRoot];
28         }
29     }
30
31     int find(int p){
32         while(p != id[p]){
33             id[p] = id[id[p]];
34             p = id[p];
35         }
36         return p;
37     }
38
39     bool connected(int p, int q){
40         return find(p) == find(q);
41     }
42 };
43
44
45
46 class Edge{
47 public:
48     int val;
49     int thisEnd;
50     int anotherEnd;
51
52     Edge(int v, int t, int a) : val(v), thisEnd(t), anotherEnd(a){}
53 };
54
55 class Compare{
56 public:
57     bool operator()(Edge e1, Edge e2){
58         return e1.val > e2.val;
59     }
60 };
61
62
63 class Solution {
64 public:

```

```

65     int minimumCost(int n, vector<vector<int>>& connections) {
66         priority_queue<Edge, vector<Edge>, Compare> pq;
67
68         WeightedUnionFind wuf(n + 1);
69         for(vector<int>& edge : connections){
70             Edge myEdge(edge[2], edge[1], edge[0]);
71             pq.push(myEdge);
72         }
73
74         queue<Edge> myQueue;
75
76         int money = 0;
77         while(!pq.empty() && myQueue.size() != n - 1){
78             Edge edge = pq.top();
79             pq.pop();
80
81             int thisEnd = edge.thisEnd;
82             int thatEnd = edge.anotherEnd;
83             int val = edge.val;
84
85             if(wuf.connected(thisEnd, thatEnd))
86                 continue;
87
88             wuf.uni(thatEnd, thisEnd);
89
90             myQueue.push(edge);
91             money += val;
92         }
93
94         return myQueue.size() == n - 1? money : -1;
95     }
96 };
97

```

1123 Lowest Common Ancestor of Deepest Leaves

执行用时: **8 ms** , 在所有 Go 提交中击败了 **80.00%** 的用户

内存消耗: **5.2 MB** , 在所有 Go 提交中击败了 **82.22%** 的用户

```
1  /*
2     朴素的想法, 求出左右的最大深度, 如果一致, 那就是它了!
3     but still duplicated process, needs to improve on that
4  */
5  func lcaDeepestLeaves(root *TreeNode) *TreeNode {
6      if root == nil || (root.Left == nil && root.Right == nil){
7          return root
8      }
9
10     left := getMaxDepth(root.Left)
11     right := getMaxDepth(root.Right)
12
13     if left == right{
14         return root
15     }else{
16         if left > right{
17             return lcaDeepestLeaves(root.Left)
18         }else{
19             return lcaDeepestLeaves(root.Right)
20         }
21     }
22 }
23
24 func getMaxDepth(root *TreeNode) int{
25     if root == nil{
26         return 0
27     }
28
29     return 1 + max(getMaxDepth(root.Left), getMaxDepth(root.Right))
30 }
31
32 func max(a int, b int) int{
33     if a < b{
34         return b
35     }
36
37     return a
38 }
```

1143 Longest Common Subsequence

执行结果： **通过** [显示详情 >](#)



执行用时： **28 ms** ，在所有 C++ 提交中击败了 **59.93%** 的用户

内存消耗： **12.9 MB** ，在所有 C++ 提交中击败了 **7.67%** 的用户

炫耀一下：



```
1 class Solution {
2 public:
3     int longestCommonSubsequence(string text1, string text2) {
4         int len1 = text1.size();
5         int len2 = text2.size();
6
7         vector<vector<int>> > dp(len1 + 1, vector<int>(len2 + 1, 0));
8         for(int i = 1; i <= len1; i++){
9             for(int j = 1; j <= len2 ; j++){
10                 if(text1[i - 1] == text2[j - 1]){
11                     dp[i][j] = 1 + dp[i - 1][j - 1];
12                 }else{
13                     dp[i][j] = max(dp[i - 1][j], dp[i][j - 1]);
14                 }
15             }
16         }
17
18         return dp[len1][len2];
19     }
20
21     int max(int i ,int j){
22         return i >= j ? i : j;
23     }
24 };
```

1153 Strubg Transforms Into Another String

```

1  /*
2      本题目的思想就是， 类似常规映射的思想
3      但是 如果 string 2 中有 26个字母，那么一定无法映射
4      因为如下原因
5
6      abcdefghigklmnopqrstuvwxyz
7      bcdefghigklmnopqrstuvwxyz
8
9      对于上面的变换，我们一定要找到一个 中间 的temp 字母，否则就会一起变换
10     比如 abc
11         bca
12
13     如果说 a-> b
14     bbc
15     bca
16
17     那么就screw 的
18
19
20     注意一定得是 ch2 的 set
21     因为比如这个例子
22     "abcdefghijklmnopqrstuvwxyz"
23     "bcdefghijklmnopqrstuvwxxz"
24
25     实际上是可以被转换的，比如
26     先让 y 变成 z 然后 a b c 通过 y 的过渡， 变成 对应的字母
27     具体变换过程是
28     abc.....xyz
29     ->
30     abc.....xzz
31     ->
32     ybc.....xzz
33     yac.....xzz
34     bac.....xzz
35     byc.....xzz
36     bya.....xzz
37     bca.....xzz -> 最终的结果
38
39  */
40  class Solution {
41  public:
42      bool canConvert(string str1, string str2) {
43          int size1 = str1.size();
44          int size2 = str2.size();
45
46          if(str1 == str2)
47              return true;
48          if(size1 != size2)
49              return false;

```

```
50
51     unordered_map<char, char> map;
52     unordered_set<char> set;
53     for(int i = 0; i < size1; i++){
54         char ch1 = str1[i];
55         char ch2 = str2[i];
56
57         if(map.count(ch1) != 0){
58             if(map[ch1] != ch2)
59                 return false;
60         }else{
61             map.insert({ch1, ch2});
62         }
63
64         set.insert(ch2);
65     }
66
67     return set.size() < 26;
68 }
69 };
```

1162 As far From Land As Possible

1162. As Far from Land as Possible

难度 中等

200



Given an $n \times n$ grid containing only values 0 and 1, where 0 represents water and 1 represents land, find a water cell such that its distance to the nearest land cell is maximized, and return the distance. If no land or water exists in the grid, return -1.

The distance used in this problem is the Manhattan distance: the distance between two cells (x_0, y_0) and (x_1, y_1) is $|x_0 - x_1| + |y_0 - y_1|$.

Example 1:

1	0	1
0	0	0
1	0	1

Input: grid = [[1,0,1],[0,0,0],[1,0,1]]

Output: 2

Explanation: The cell (1, 1) is as far as possible from all the land with distance 2.

Example 2:

1	0	0
0	0	0
0	0	0

执行用时: 18 ms , 在所有 Java 提交中击败了 49.89% 的用户

内存消耗: 39.6 MB , 在所有 Java 提交中击败了 42.19% 的用户

炫耀一下:



```
1 int[][] dir = {{1, 0}, {0, 1}, {-1, 0}, {0, -1}};
2 int row;
3 int col;
4 public int maxDistance(int[][] grid) {
5     row = grid.length;
6     col = grid[0].length;
7     int ones = 0;
```



```

8      int zeros = 0;
9
10     Deque<int[]> queue = new ArrayDeque<>();
11     for(int i = 0; i < row; i++){
12         for(int j = 0; j < col; j++){
13             if(grid[i][j] == 1) {
14                 queue.add(new int[]{i, j});
15                 ones++;
16             }
17         }
18     }
19
20     zeros = row * col - ones;
21     int maxDistance = 0;
22
23     while(!queue.isEmpty()){
24         int size= queue.size();
25         for(int i = 0; i < size; i++){
26             int[] cur = queue.removeFirst();
27
28             for(int k = 0; k < 4; k++){
29                 int newX = cur[0] + dir[k][0];
30                 int newY = cur[1] + dir[k][1];
31
32                 if(isInRange(newX, newY) && grid[newX][newY] == 0){
33                     grid[newX][newY] = grid[cur[0]][cur[1]] + 10;
34                     zeros--;
35
36                     if(zeros == 0)
37                         break ;
38
39                     queue.addLast(new int[]{newX, newY});
40                 }
41             }
42
43             if(zeros == 0)
44                 break;
45         }
46
47         if(zeros == 0)
48             break;
49     }
50
51     for(int i = 0; i < row; i++){
52         for(int j =0; j < col; j++){
53             if(grid[i][j] > 1)
54                 maxDistance = Math.max(maxDistance, grid[i][j] / 10);
55         }
56     }

```

```

57
58     return maxDistance == 0 ? -1 : maxDistance;
59 }
60
61 private boolean isInRange(int i, int j){
62     return i >= 0 && j >= 0 && i < row && j < col;
63 }

```

```

1     public int maxDistance(int[][] grid) {
2         int row = grid.length;
3         int col = grid[0].length;
4
5         List<int[]> list = new ArrayList<>();
6
7         for(int i = 0; i < row; i++) {
8             for (int j = 0; j < col; j++) {
9                 if(grid[i][j] == 1)
10                     list.add(new int[]{i, j});
11             }
12         }
13
14
15         int maxDistance = 0;
16         for(int i = 0; i < row; i++) {
17             for (int j = 0; j < col; j++) {
18                 if (grid[i][j] == 0) {
19                     int distance = Integer.MAX_VALUE;
20                     for (int[] arr : list) {
21                         distance = Math.min(distance, Math.abs(arr[0] - i) +
Math.abs(arr[1] - j));
22                     }
23
24                     if(distance != Integer.MAX_VALUE)
25                         maxDistance = Math.max(maxDistance, distance);
26                 }
27             }
28         }
29
30         return maxDistance == 0 ? -1 : maxDistance;
31     }
32

```

